(Jake Tantorski

Assignment 2

Section 1

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1.
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\begin{split} & \text{int } i=1; \ /\!/1 \\ & \text{while } (i <= n) \ /\!/ \\ & \text{some } O(1) \text{ time statements} \quad /\!/1 \\ & i=i*2; \ /\!/2 \\ & \text{end while} \\ \\ & \text{Iteration } 1: i=1 \\ & \text{Iteration } 2: i=2 \\ & \text{Iteration } 3: i=4 \\ & \text{Iteration } k: i=2^{(k-1)}=n \\ & k-1=\log(n) \\ & K=\log(n)+1 \end{split}
```

2.

Outer loop: $n/2^k-1=1$

Time: O(logn)

Using algebra I multiplied each side by 2^k-1.

$$N = 2^k-1$$

Logged each side

$$Log(n) = k-1$$

Moved the one over

$$Log(n)+1=k$$

Middle Loop:

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2^k = 1
Logged each side
n = log(n)
Move the one over
n = \log(n)
Inisde Loop:
n/2 = 1
Time: (\log(n)+1)*(\log(n))*(n/2) = (n\log(n)n)/2 + (n\log(n))/2 = O(n\log n)
3.
Outside Loop: n times
Middle Loop: 2^{(k-1)} - 1 = n - > (k-1) = \log(n) + 1 - > k = \log(n)
Inside Loop: j times = n
Time: n*log(n)*n = n^2*log(n) = O(n^2)
4.
j = 1, i = 0; //2
while (i \le n) {
       i = i + j; 1/2 * Sqrt(2n)
       j++; //2 * Sqrt(2n)
}
Time: 2+4Sqrt(2n) = O(sqrt(n))
Iteration k: (k*(k+1))/2
n = (k*(k+1))/2
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2n = k*(k+1)
k = k+1 assuming n is large
2n = k^2
k = Sqrt(2n)
5.
2010
log(log(n))
log(n)
Sum of n where i = 1 (1/i)
sqrt(n)
n
Sum of n where i = 1 (1)
nlog(n)
Sum of n where i = 1 (i)
n^2
n^4
2^n
e^n
n!
```

 n^n

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6.
squareroot(x)
Input X
              //2
Lower = 1
              //2
Upper = x
              //2
Ans = 0
              //2
while(lower<=upper) //logn
       Mid = (lower+upper)/2
       if(mid*mid == x)
                             //2
               Return mid
       Else if(mid*mid \leq x) //2
              lower = mid + 1
               ans = mid
       Else
              Upper = mid - 1
       End if
```

End while

Return ans

Time: O(log(n)) because it only uses binary search