

ASSIGNMENT 7

Aim:-

Insert the keys into a hash table of length m using open addressing using double hashing with $h(k)=1+(k \bmod (m-1))$.

Objective:-

To Insert the keys into a hash table of length m using open addressing using double hashing with $h(k)=1+(k \bmod (m-1))$.

Theory:-

Double hashing is a collision resolving technique in **Open Addressed Hash tables**. Double hashing uses the idea of applying a second hash function to key when a collision occurs.

Double hashing can be done using :

$(hash1(key) + i * hash2(key)) \% TABLE_SIZE$

Here $hash1()$ and $hash2()$ are hash functions and $TABLE_SIZE$ is size of hash table.

(We repeat by increasing i when collision occurs)

First hash function is typically $hash1(key) = key \% TABLE_SIZE$

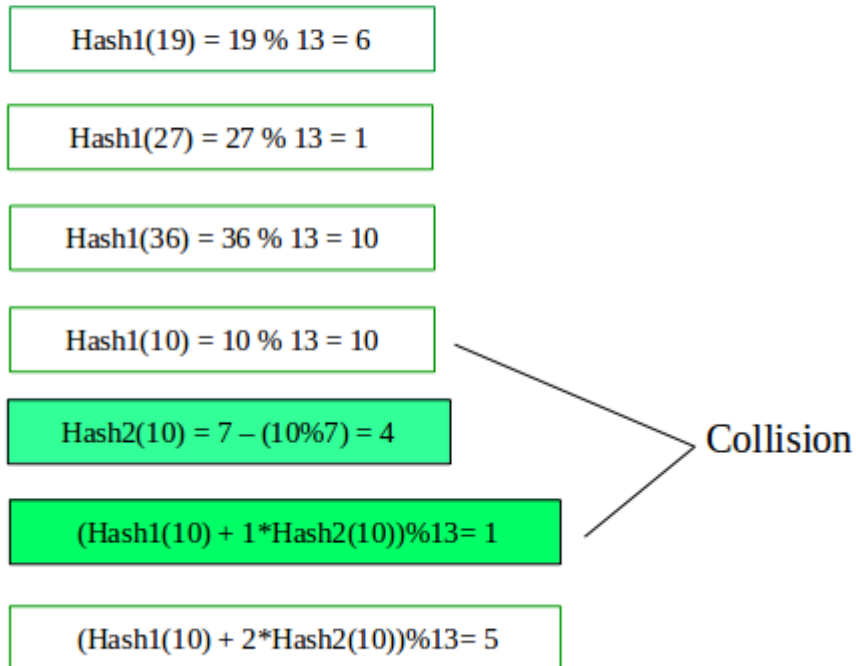
A popular second hash function is : **$hash2(key) = PRIME - (key \% PRIME)$** where $PRIME$ is a prime smaller than the $TABLE_SIZE$.

A good second Hash function is:

- It must never evaluate to zero
- Must make sure that all cells can be probed

Lets say, $\text{Hash1}(\text{key}) = \text{key} \% 13$

$\text{Hash2}(\text{key}) = 7 - (\text{key} \% 7)$



Algorithm:-

- Linear probing collision resolution leads to clusters in the table, because if two keys collide, the next position probed will be the same for both of them.
- The idea of double hashing: Make the offset to the next position probed depend on the key value, so it can be different for different keys
 - Need to introduce a second hash function $H_2(K)$, which is used as the offset in the probe sequence (think of linear probing as double hashing with $H_2(K) == 1$)

- For a hash table of size M , $H_2(K)$ should have values in the range 1 through $M-1$; if M is prime, one common choice is $H_2(K) = 1 + ((K/M) \bmod (M-1))$
- The insert algorithm for double hashing is then:
 1. Set $\text{indx} = H(K)$; $\text{offset} = H_2(K)$
 2. If table location indx already contains the key, no need to insert it. Done!
 3. Else if table location indx is empty, insert key there. Done!
 4. Else collision. Set $\text{indx} = (\text{indx} + \text{offset}) \bmod M$.
 5. If $\text{indx} == H(K)$, table is full! (Throw an exception, or enlarge table.) Else go to 2.
- With prime table size, double hashing works very well in practice

Code:-

```
#include <iostream>
using namespace std;
class dr
{
    int n=10;
    int arr[100][3];
    int c;
public:
    dr()
    {
        cout<<"Table of size "<<n<<" created\n";
        for(int i=0;i<n;i++)
        {
            arr[i][0]=0;
            arr[i][1]=-1;
            arr[i][2]=-1;
        }
        c=0;
    }
    void add(int,int);
    int find_key(int);
    void display();
    void update_val(int,int);
};
void dr::add(int key,int value)
{
    int new_hash_addr1,new_hash_addr2,main_hash_addr=-1,j=0;
    if(this->find_key(key)!=-1)
```


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```
{
    cout<<"Key already exists\n";
    return;
}
if(c==(n-1))
{
    cout<<"Table full, request denied\n";
}
new_hash_addr1=(key)%n;
new_hash_addr1=1+(key%(n-1));
if(arr[new_hash_addr1][1]==-1)
{
    arr[new_hash_addr1][0]=key;
    arr[new_hash_addr1][1]=value;
}
else if(arr[new_hash_addr2][1]==-1)
{
    arr[new_hash_addr2][0]=key;
    arr[new_hash_addr2][1]=value;
}
else
{
    while(arr[new_hash_addr2][2]!=-1)
    {
        main_hash_addr=new_hash_addr2;
        new_hash_addr2=arr[main_hash_addr][2];
    }
    main_hash_addr=new_hash_addr2;
    for(int i=0;i<n;i++)
    {
        new_hash_addr2=(main_hash_addr+i)%n;
        if(arr[new_hash_addr2][1]==-1)
        {
            arr[new_hash_addr2][0]=key;
            arr[new_hash_addr2][1]=value;
            arr[main_hash_addr][2]=new_hash_addr2;
            c++;
            break;
        }
    }
}
}
}
void dr::display()
{
    cout<<"Key\t\tValue\t\tChain\n";
    for(int i=0;i<n;i++)
    {
        cout<<arr[i][0]<<"\t\t"<<arr[i][1]<<"\t\t"<<arr[i][2]<<endl;
    }
}
int dr::find_key(int key)
{
    int search_addr=key%n,f=0;
    while(arr[search_addr][0]!=key && arr[search_addr][2]!=-1)
```

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```
{
    search_addr=arr[search_addr][2];
}
if(arr[search_addr][0]==key)
{
    return arr[search_addr][1];
}
else if(arr[search_addr][2]==-1)
{
    return -1;
}
}
int main()
{
    char r;
    do
    {
        char op;
        dr table;
        int c;
        do
        {
            cout<<"-----Menu-----\n";
            cout<<"1] Insert value\n2] Display\n";
            cout<<"\n";
            cout<<"Enter your choice: ";
            cin>>c;
            switch(c)
            {
                case 1: {
                    int key,val;
                    cout<<"Enter key: ";
                    cin>>key;
                    cout<<"Enter value: ";
                    cin>>val;
                    table.add(key,val);
                }
                break;
                case 2: table.display();
                break;
                default:cout<<"Invalid\n";
            }
            cout<<"\nDo you wish to go again? ";
            cin>>op;
        }while(op=='y' || op=='Y');
        cout << "Test pass?(y/n): " << endl;
        cin>>r;
    }while(r=='n' || r=='N');
    cout<<"*****\n";
    cout<<"*   Thank You!   *\n";
    cout<<"*****\n";
    return 0;
}
```

Output Screenshot:-

 "C:\Users\Dell\Downloads\main (4).exe"

```
Enter value: 2

Do you wish to go again? y
-----Menu-----
1] Insert value
2] Display

Enter your choice: 1
Enter key: 6
Enter value: 5

Do you wish to go again? y
-----Menu-----
1] Insert value
2] Display

Enter your choice: 2
Key          Value      Chain
0            -1         -1
0            -1         -1
1            2         -1
0            -1         -1
0            -1         -1
0            -1         -1
0            -1         -1
0            -1         -1
6            5         -1
0            -1         -1
0            -1         -1

Do you wish to go again?
```

Conclusion:-

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We Successfully implemented the Heap datastructure.