Note 1

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VC-Dimension. If $\exists x_1,...,x_n, s.t.\mathcal{F}$ shatters $x_1,...,x_n$ and $\forall x_1,...,x_{n+1}, \mathcal{F}$ can not shatter $x_1,...,x_{n+1}$. n is the **VC-dimension** of \mathcal{F}

Linear Classifier.

$$\max_{w,b,t} t$$

$$s.t.y_i(w^T x_i + b) \ge t$$

$$||w|| = 1$$

or

$$min_{w,b} \frac{1}{2} ||w||^2$$
$$s.t.y_i(w^T x_i + b) \ge 1$$

Appendix: Game Theory & Lagrange Duality. In two player zero-sum game, We have minmax = maxmin in a Nash equilibrium.¹

$$\max_{y \in A_2} \min_{x \in A_1} u_1(x, y) = \min_{x \in A_1} \max_{y \in A_2} u_1(x, y)$$

Problem 1. For a matrix M, $\min_i \max_j M_{ij} \ge \max_j \min_i M_{ij}$ which means that if two player choose pure strategy and to act sequentially, second one gets an advantage.

If two players can choose mixed strategy, a Nash equilibrium exists, so the equation above holds. **Saddle Point Theorem.** f(x,y) Fix y, $f(\cdot,y)$ is convex. Fix x, $f(x,\cdot)$ is concave, then

$$\max_{y} \min_{x} f(x, y) = \min_{x} \max_{y} f(x, y)$$

If $f(x), g_i(x)$ are convex and $h_i(x)$ is linear, the optimization problem

$$\begin{cases} minf(x) \\ s.t.g_i(x) \leq 0 \\ h_i(x) = 0 \end{cases}$$

$$\iff \begin{cases} min_x max_{\mu,\lambda} f(x) + \sum \mu_i h_i(x) + \sum \lambda_i g_i(x) \\ s.t.\lambda_i \geq 0 \end{cases}$$

$$\iff \begin{cases} max_{\mu,\lambda} min_x f(x) + \sum \mu_i h_i(x) + \sum \lambda_i g_i(x) \\ s.t.\lambda_i \geq 0 \end{cases}$$

Let $L(x,\mu,\lambda) = f(x) + \sum \mu_i h_i(x) + \sum \lambda_i g_i(x)$, solve $\frac{\partial L}{\partial x} = 0$ to get $x^* = \phi(\mu,\lambda)$. It suffices to solve $\max_{\mu,\lambda} f(\phi(\mu,\lambda)) + \sum \mu_i h_i(\phi(\mu,\lambda)) + \sum \lambda_i g_i(\phi(\mu,\lambda))$.

The linear classifier problem is
$$\begin{cases} \min_{w,b} \frac{1}{2} \|w\|^2 \\ s.t.y_i(w^T x_i + b) \geq 1 \end{cases}$$

¹A course in Game Theory, P25

$$\iff \begin{cases} \min_{w,b,\lambda_i} L(w,b) = \frac{1}{2} \|w\|^2 - \sum \lambda_i [y_i(w^T x_i + b) - 1] \\ \lambda_i \ge 0 \end{cases}$$

 $L(w,b) = \frac{1}{2} \|w\|^2 - \sum \lambda_i [y_i(w^T x_i + b) - 1], \ w^* = \sum \lambda_i y_i x_i, \ \sum \lambda_i y_i = 0$ **KKT-conditions.** (1) Stationary: $\nabla L(x,\lambda,\mu)|_{x^*,\lambda^*,\mu^*} = 0$ (2) Primal feasibility: $h_i(x^*) = 0, g_i(x^*) \leq 0$ (3) Dual Feasible $\lambda^* \geq 0$ (4) Complementary slackness $\lambda_i g_i(x^*) = 0$

Problem 2. KKT is necessary condition.

Problem 3. If $f(x), g_i(x)$ are convex and $h_i(x)$ is linear, KKT is sufficient condition.

By KKT condition, $\lambda_i^*[y_i(w^{*T}x_i+b^*)-1]=0$. $\lambda_i^*=0$ for all (x_i,y_i) that are not closest to the hyperspace.

 $\lambda_i^* \neq 0$ for all support vector. It is the Support Vector Machine (SVM). **Term Project.** Data space = [N], $\mathcal{F} \subseteq \{0,1\}^N$, VC dmension $VC(\mathcal{F}) = d$ if $\exists i_1,...,i_d$, \mathcal{F} 's projection onto $i_1,...,i_d$ contains $\{0,1\}^d$ and $\forall i_1,...,i_{d+1}$'s projection onto $i_1,...,i_{d+1} \neq \{0,1\}^{d+1}$

$$f \in \mathcal{F}, \exists X_f \subseteq X, f|_{X_f} \neq f'|_{X_f}, \forall f' \in \mathcal{F}$$

Teaching dimension of f is $TD(f,\mathcal{F}) = min|X_f|$, best case teaching dimension is $TD_{min}(\mathcal{F}) =$ $min_{f \in \mathcal{F}}TD(f, \mathcal{F}).$

Let $\mathcal{F}_0 = \mathcal{F}$ and assume $f_1, s.t.TD_{min}(\mathcal{F}) = TD(f_1, \mathcal{F}), \mathcal{F}_1 = \mathcal{F}_0 \setminus \{f_1\}$ and so on.

Define Recursive Teaching Dimension: $RTD(\mathcal{F}) = max_t TD_{min}(\mathcal{F}_t)$, then $RTD(\mathcal{F}) = O(VC(\mathcal{F})^2)^2$

²Quadratic Upper Bound for Recursive Teaching Dimension of Finite VC Classes