

# 關於生命週期的一點事兒

**The relationship of Lifetimes and DataFlow**

Rnic / H.-S. Zheng

*Aug 17, 2019 @ COSCUP*

# Audience

- 讀過 Rust Book
- 想要了解編譯器怎麼看待 Lifetimes
- 對編譯器有那麼一點興趣
- ~~想要輕鬆駕馭 Rust's Lifetimes~~
- ~~想要快快樂樂寫 Rust~~

# Outline

- 1. Introduction
  - Example1
  - Basic Lifetimes Concepts
- 2. Borrow Checker
  - Collaborate with Data Flow
  - Example2
  - Datafrog (a datalog engine used in Polonius)

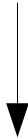
# Compilation of Rust

## Rust Code

```
let foo: T = Foo {};  
let bar: T = Bar {};  
  
let mut p = &foo;  
  
if cond {  
    println!("{}", *p);  
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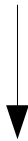


HIR

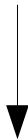
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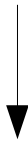
MIR



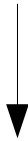
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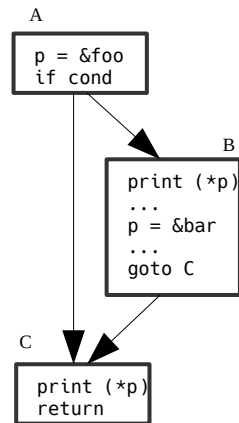
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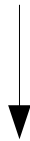
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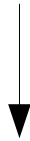
Control Flow Graph

# Compilation of Rust

Rust Code



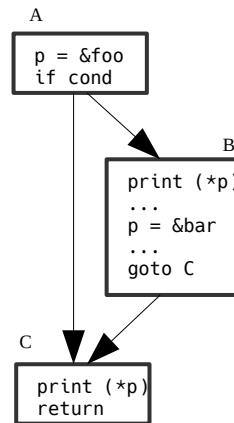
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Control Flow Graph



Borrow checker

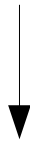
**Theorem:** Data Flow Analysis

**Tool:** Datafrog

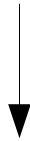


# Compilation of Rust

Rust Code



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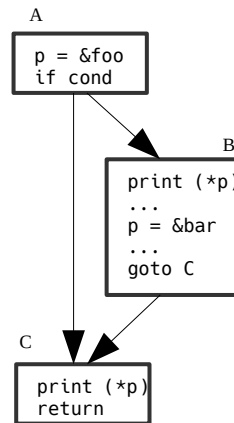


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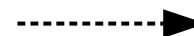


Control Flow Graph

Documents

**NLL RFC:**  
2094-nll

**Polonius:**  
an-alias-based-formulation  
-of-the-borrow-checker



Borrow checker

**Theorem:** Data Flow Analysis

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# Example

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fn list_not(mut head: Option<Box<ListNode>>) -> Option<Box<ListNode>>
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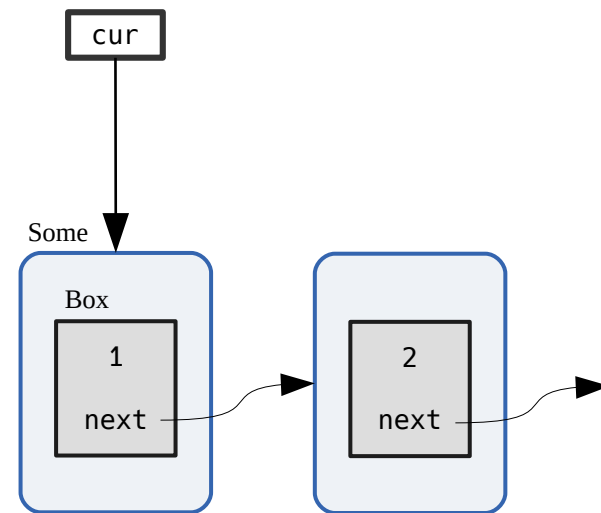
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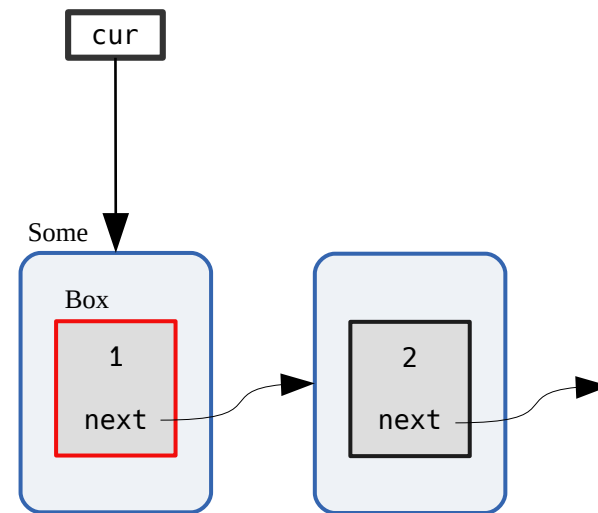


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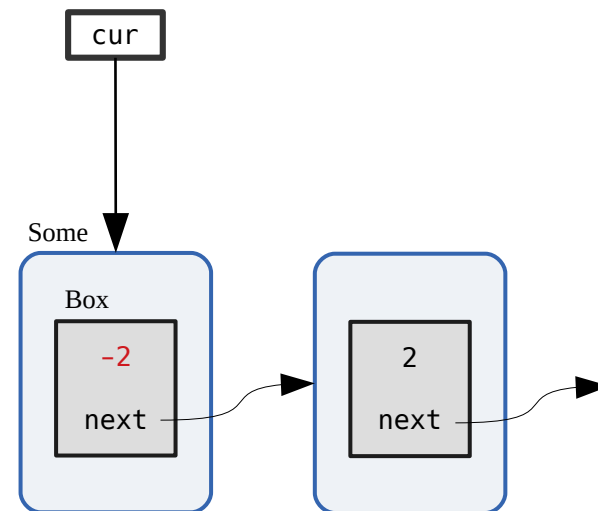


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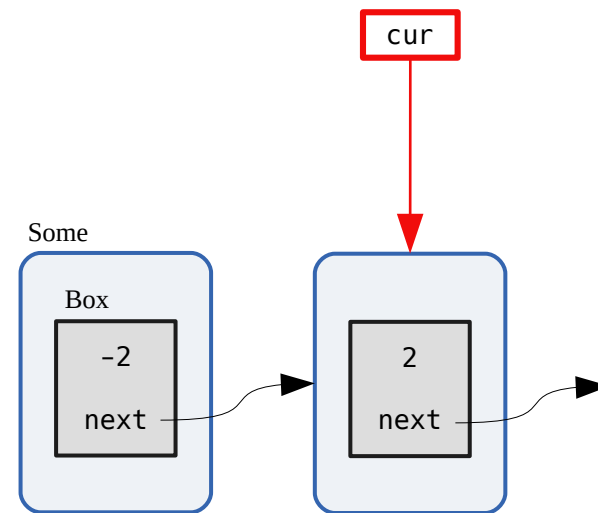
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# Using region

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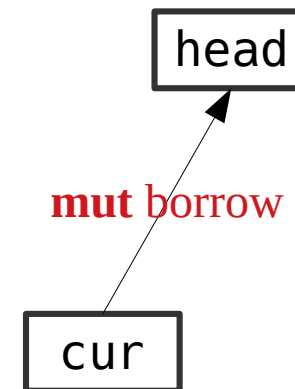
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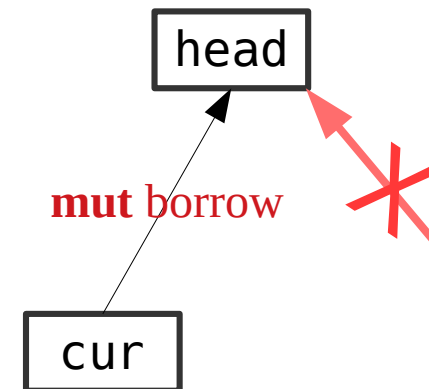
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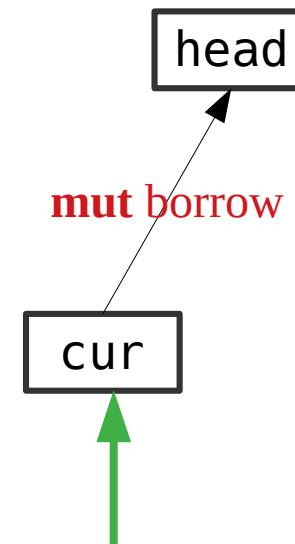
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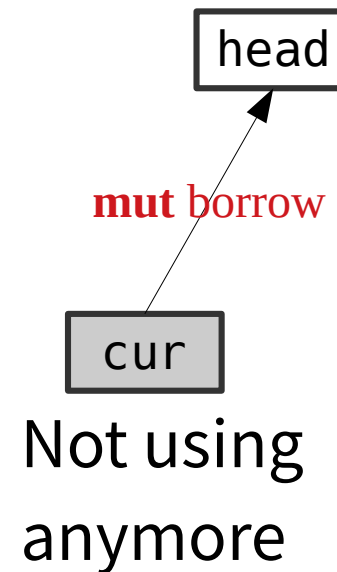
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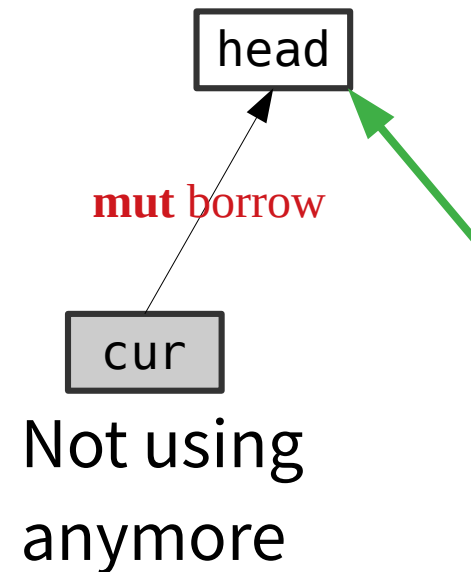
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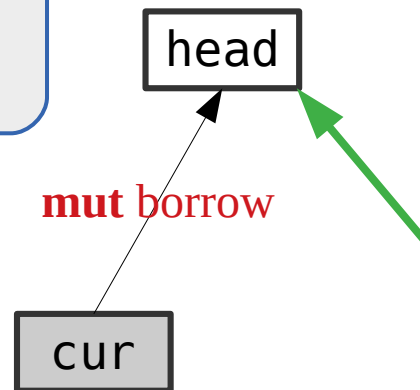
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```

``cur'` only used here

```
while let Some(nodeBox) = cur.as_mut() {
    nodeBox.val = !nodeBox.val;
    cur = &mut nodeBox.next;
}
```

head



Not using  
anymore

# Borrow

$P_0$  : **let**  $r$ :  $\&T$

$P_1$  :  $r = \&x$ ;

Borrow Exp

# Borrow

$P_0$  : **let**  $r$  :  $\&'0$   $T$

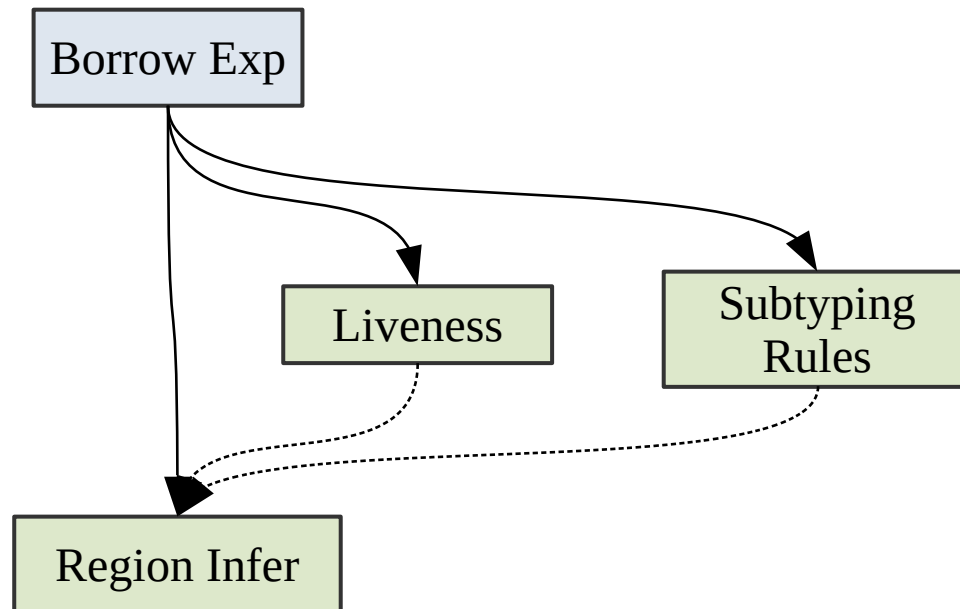
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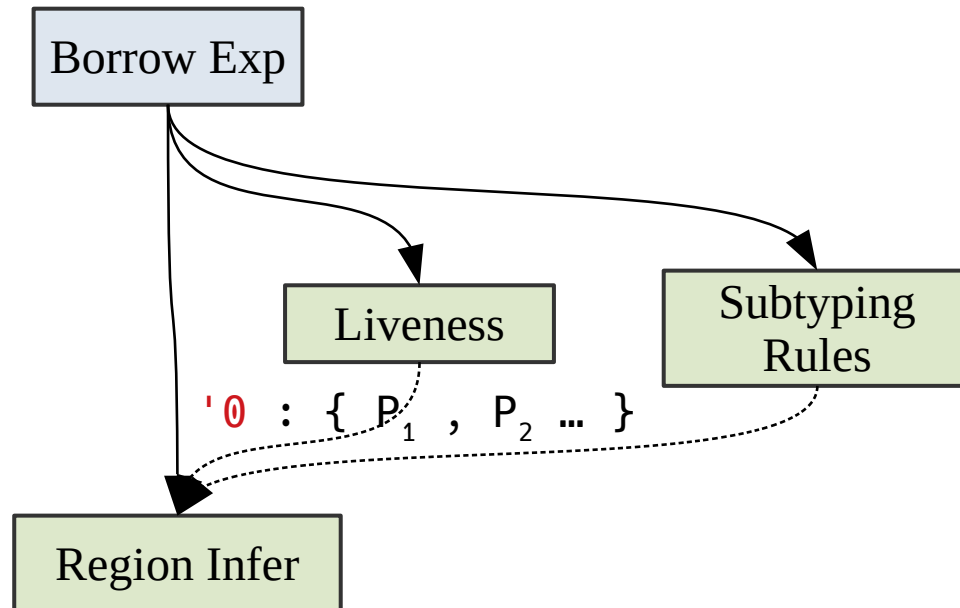




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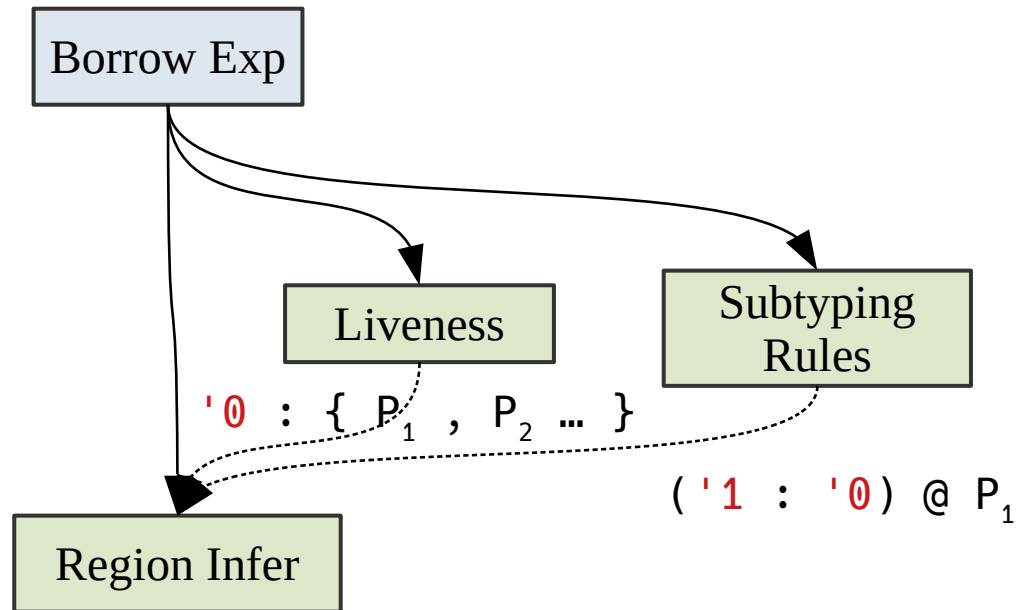
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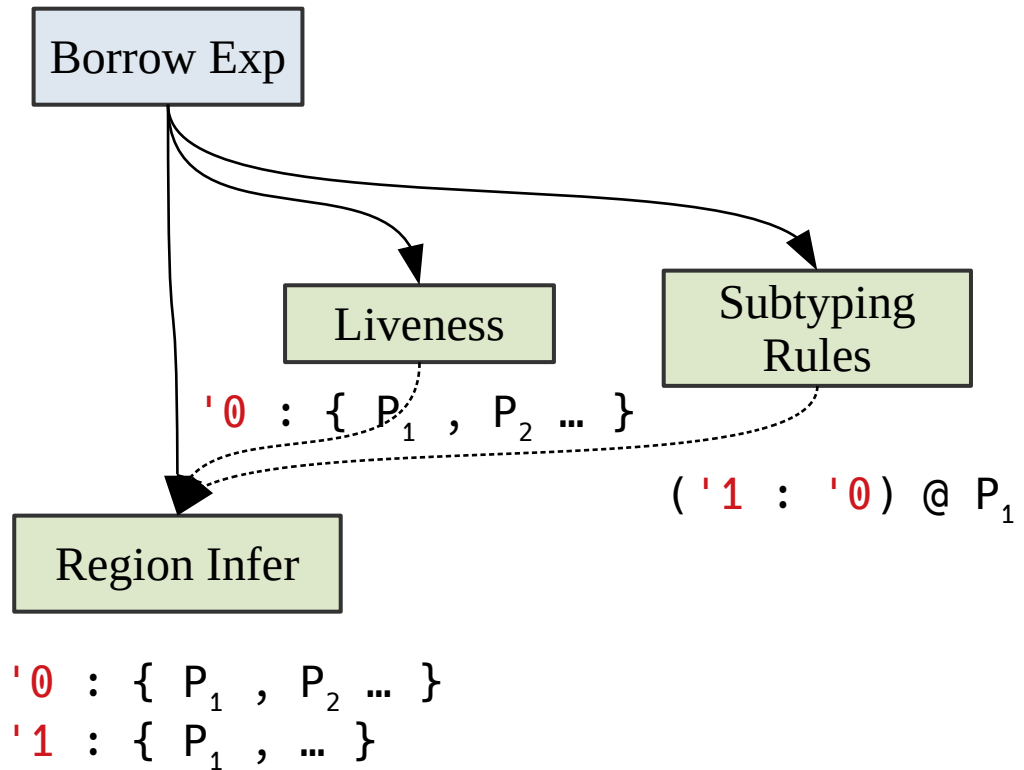
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## Borrow

 $P_0 : \text{let } r : \&'0 \text{ } T$ 
 $P_1 : r = \&'1 \text{ } x;$ 


# Borrow

Each Borrow expression will  
corresponding to each Loan

```
Loan L0 {
  point: P1,
  path: x,
  kind: shared
  region: '1 {
    P1 ...
  }
}
```

$P_0 : \text{let } r : \&'0 \text{ T}$

$P_1 : r = \&'1 x;$

Borrow Exp

Liveness

Subtyping  
Rules

Region Infer

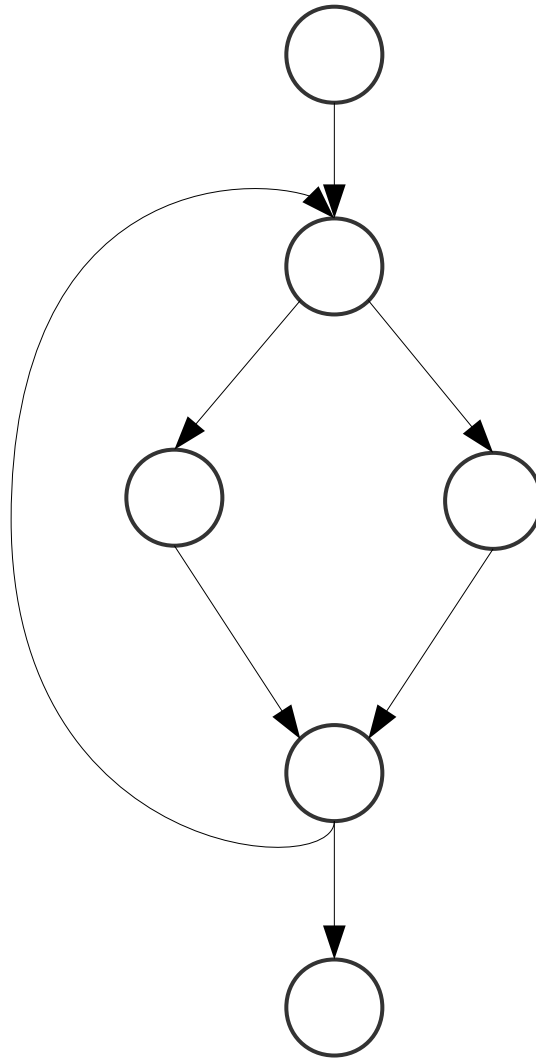
$'0 : \{ P_1, P_2, \dots \}$

$'1 : \{ P_1, \dots \}$

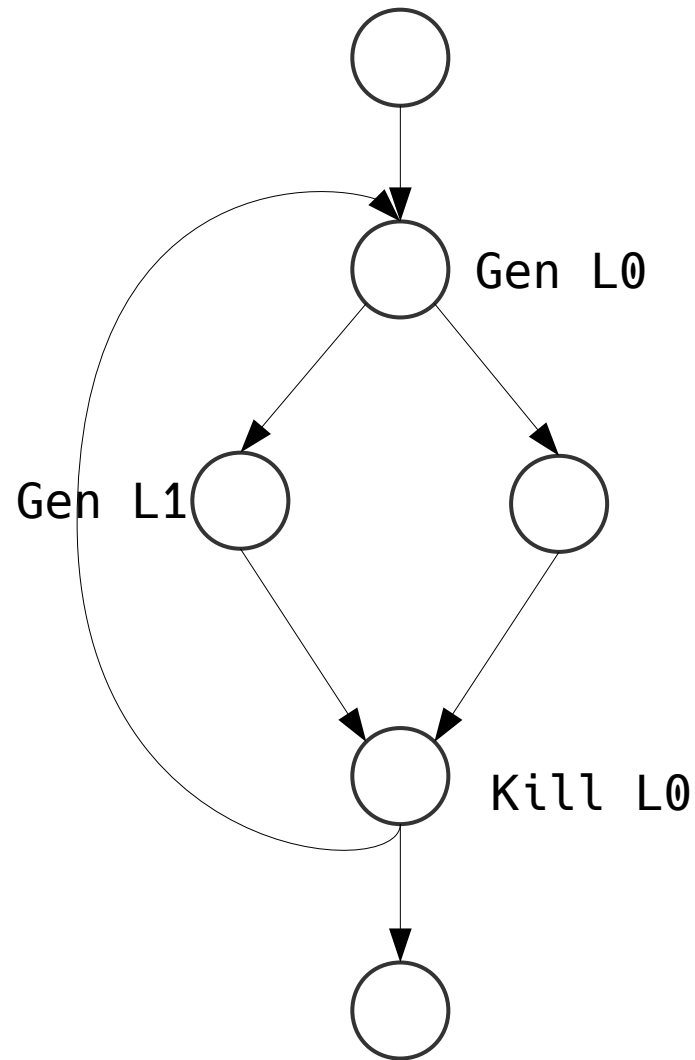
$'0 : \{ P_1, P_2, \dots \}$

$( '1 : '0 ) @ P_1$

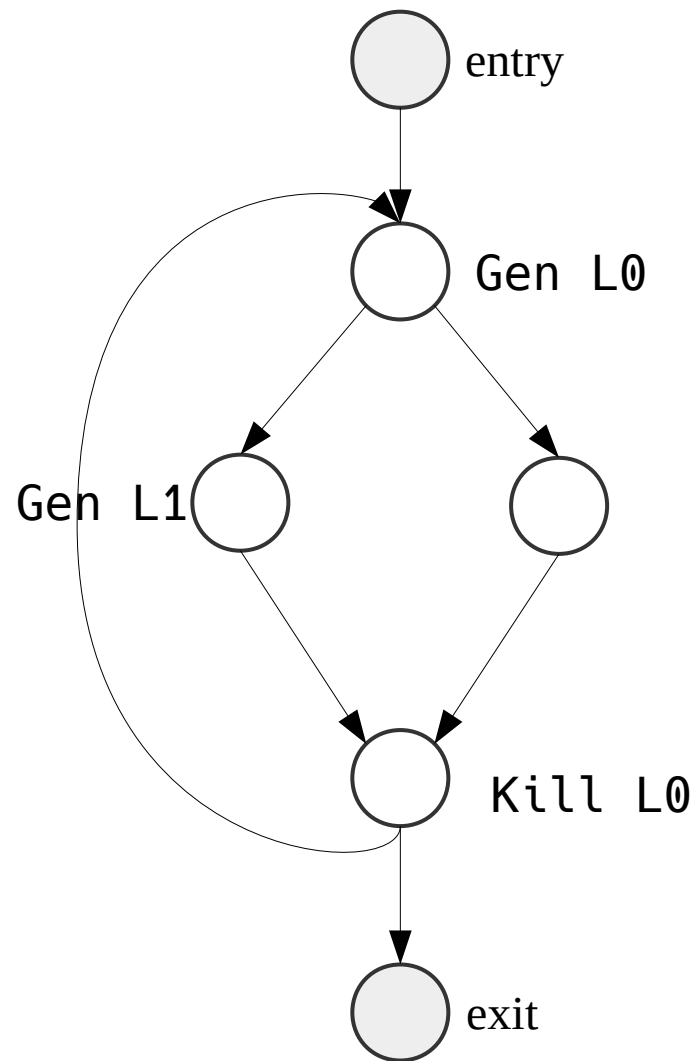
# The Data Flow of the Loan



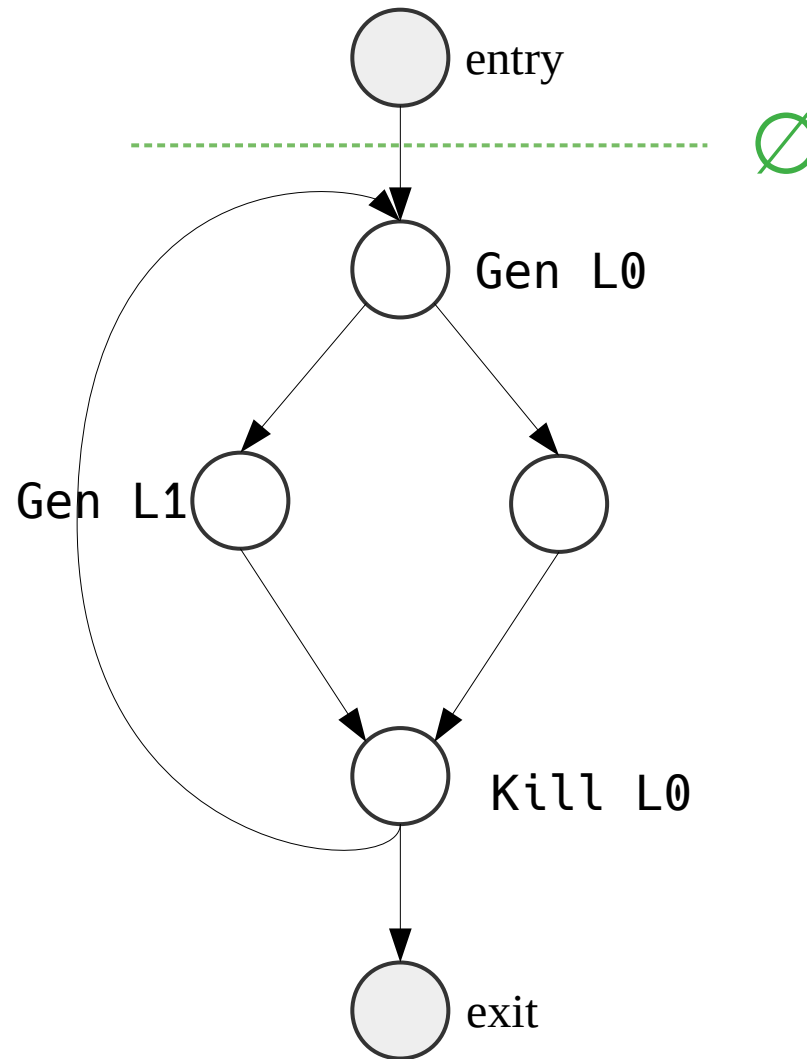
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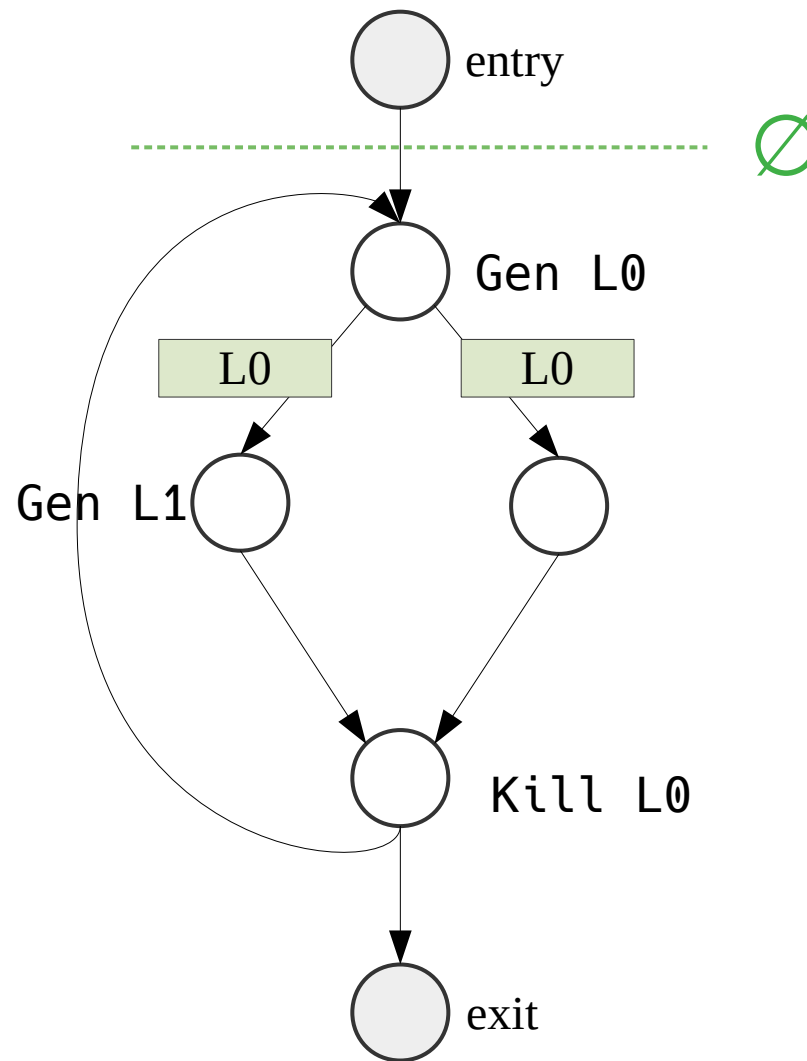


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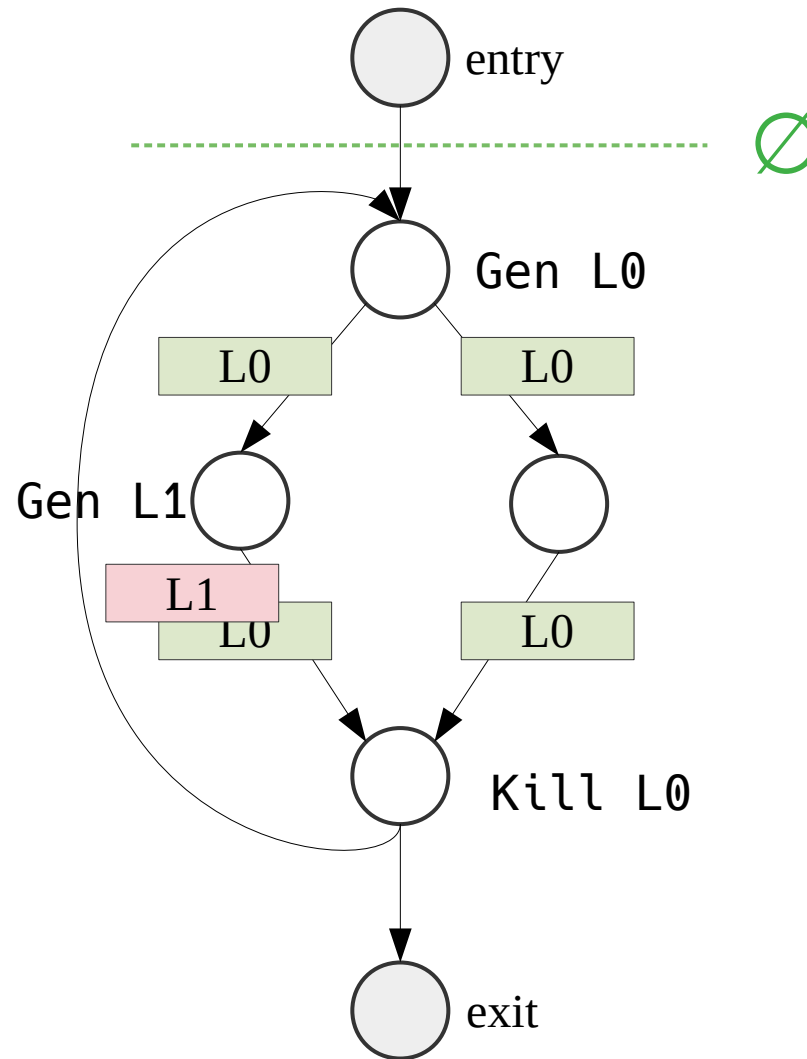




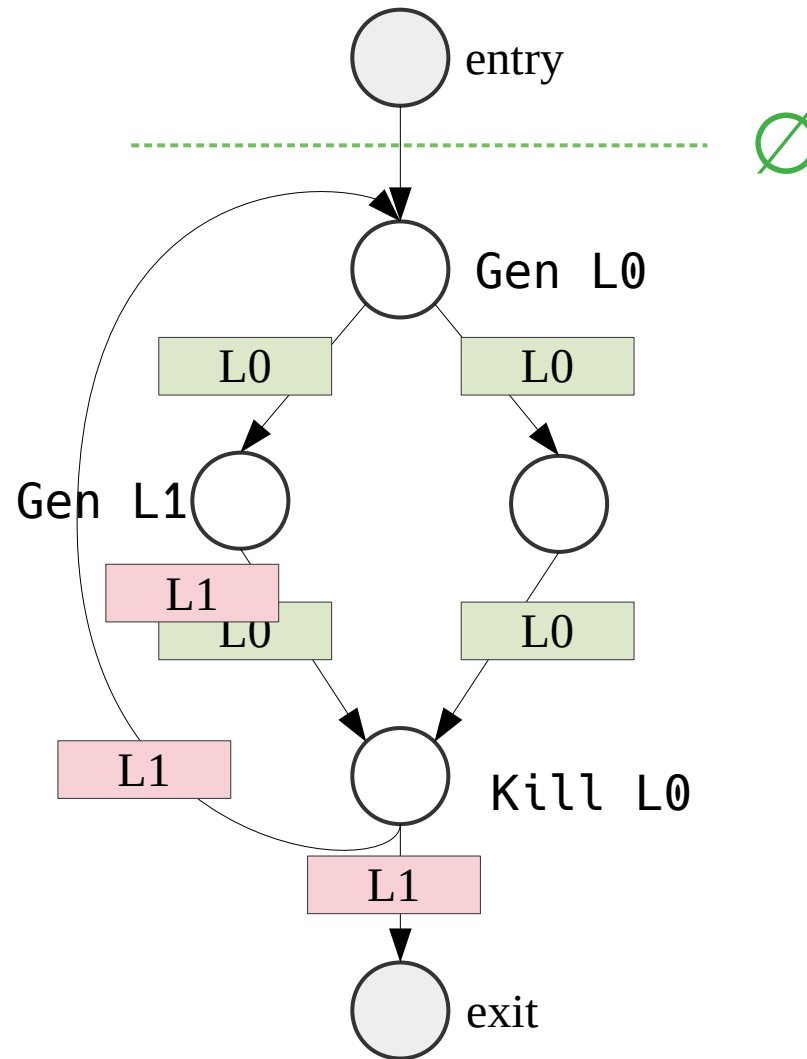
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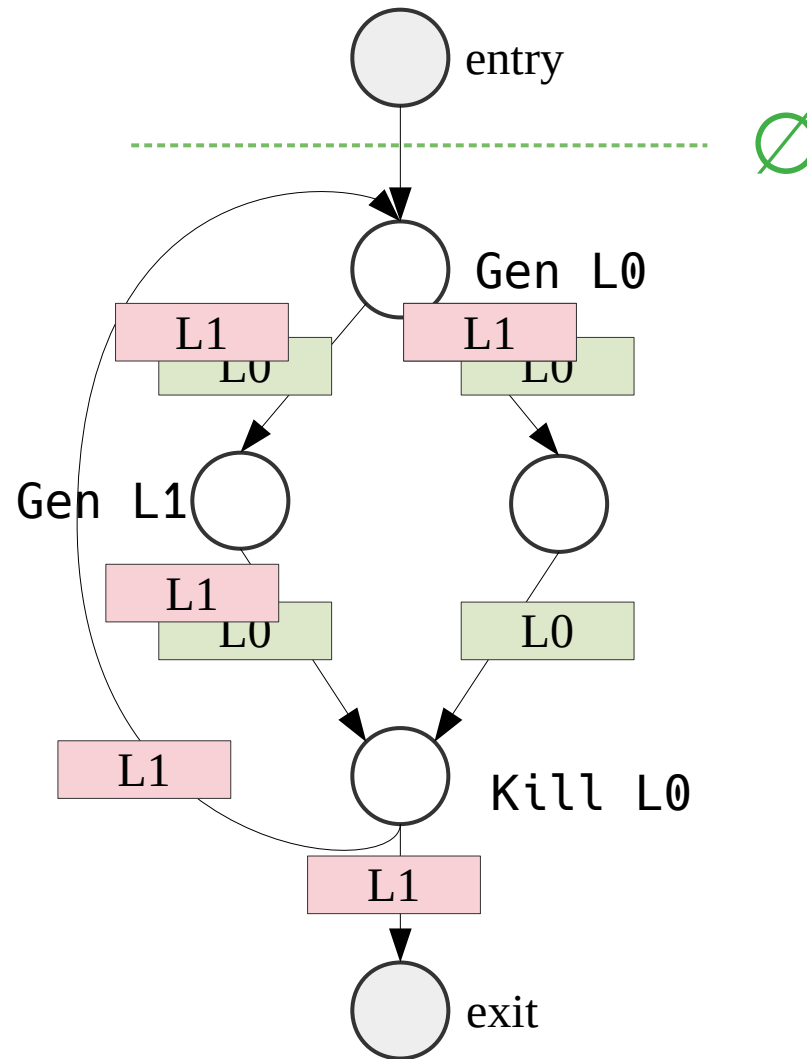
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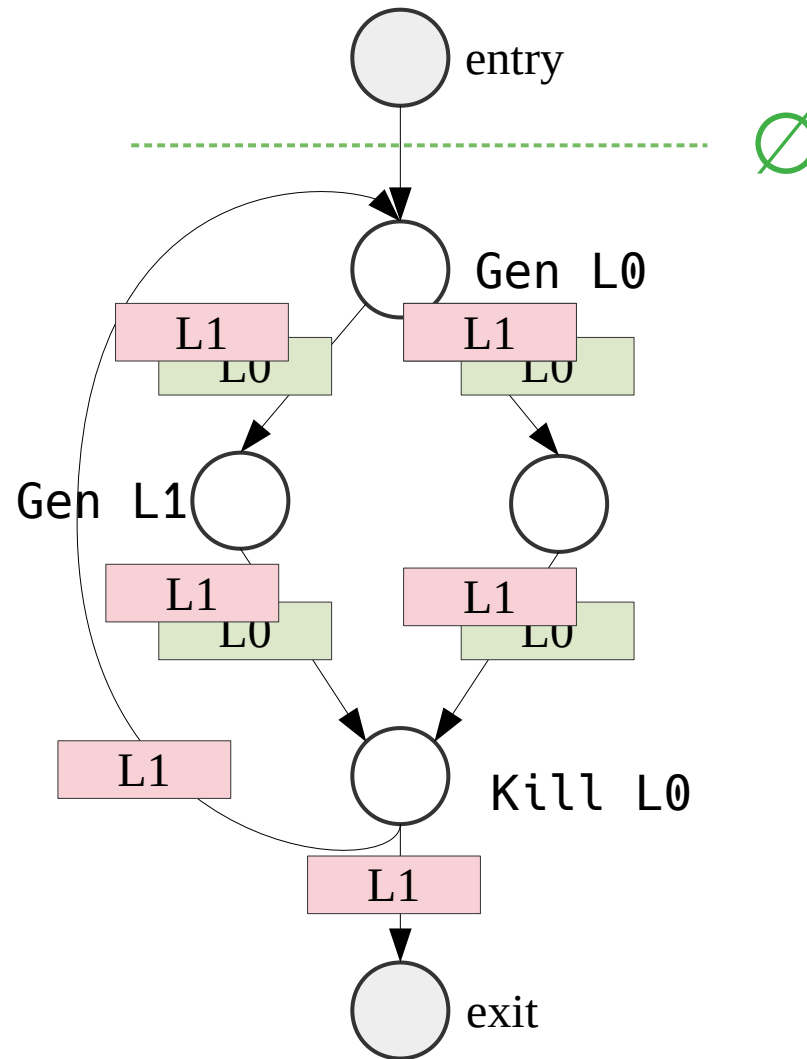
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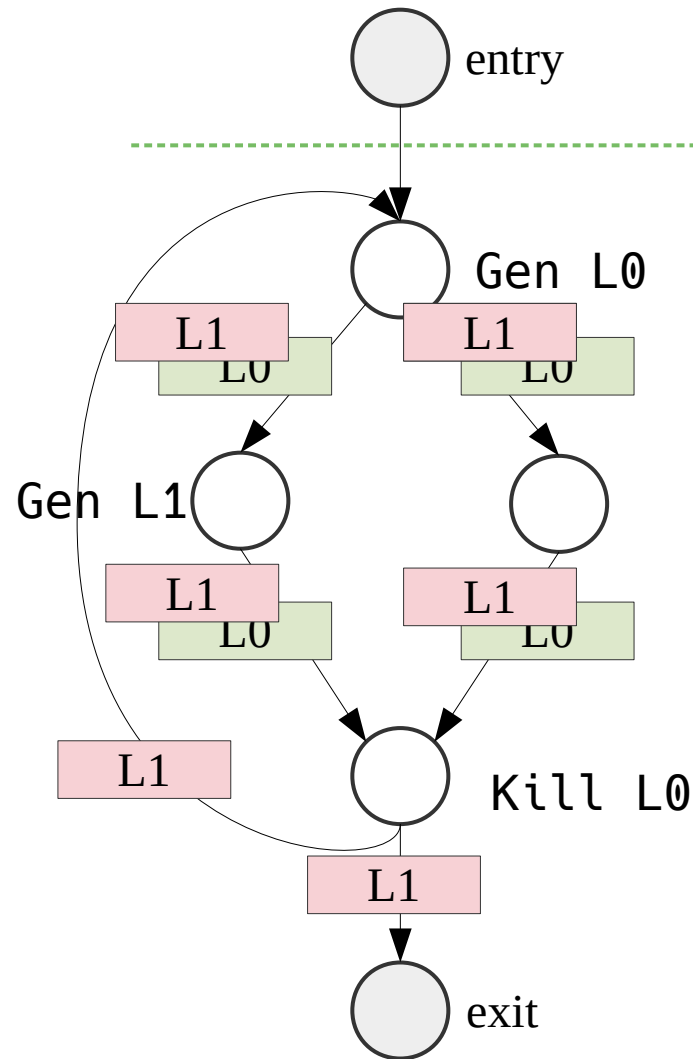


# The Data Flow of the Loan



# The Data Flow of the Loan

**Key:** which loan live at which points



When all the sets are stable, that's mean **the state is not changed anymore**, then the data flow computation is complete.

# When to Gen, Kill

## Gen Loan :

If it's a borrow expression, then gen a Loan

## Kill Loan :

- 1)  $LV = \text{Loan}_i . \text{path}$
- 2)  $\text{point} \notin \text{Loan}_i . \text{region}$

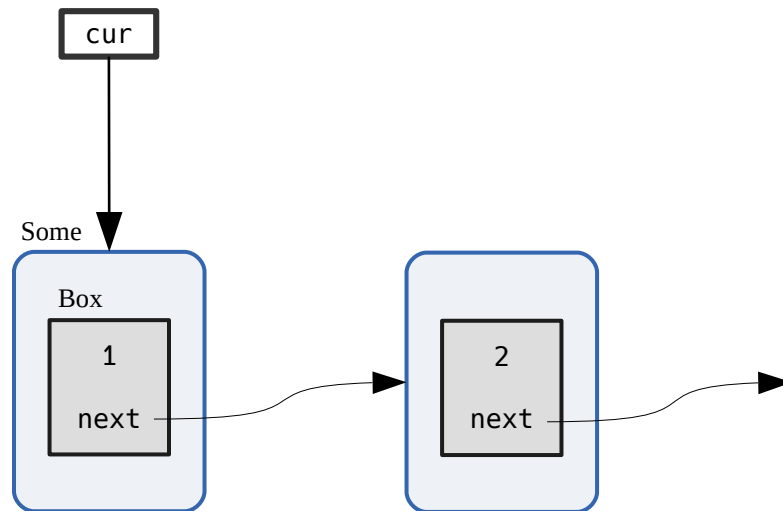
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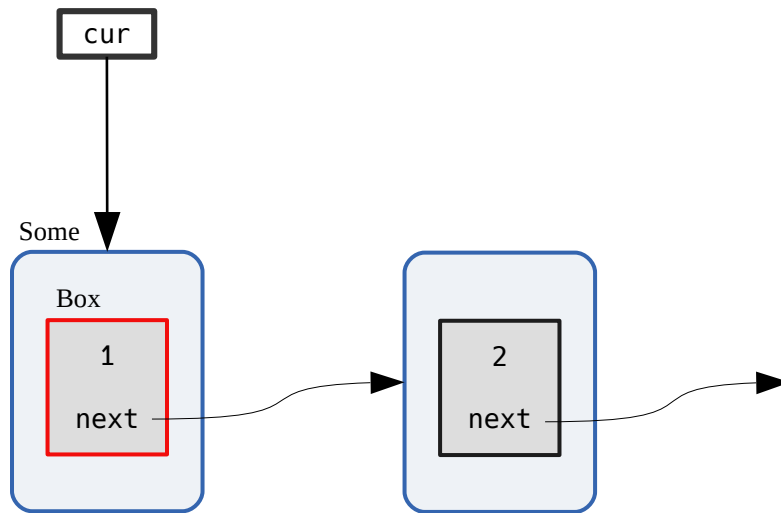
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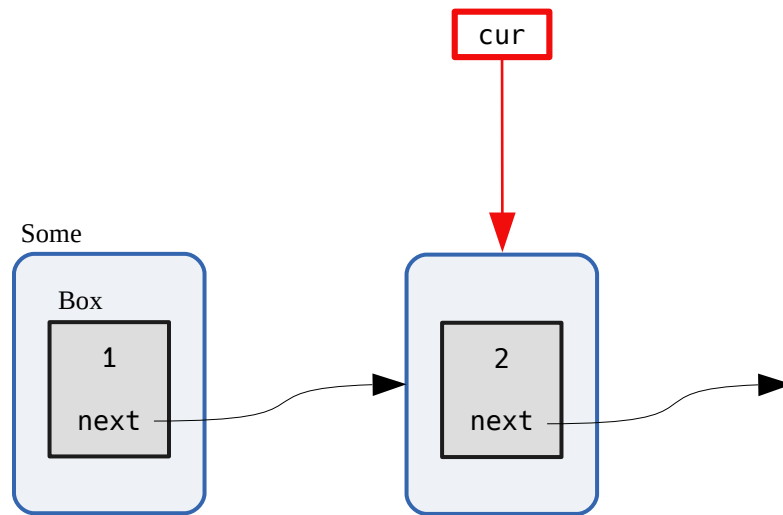
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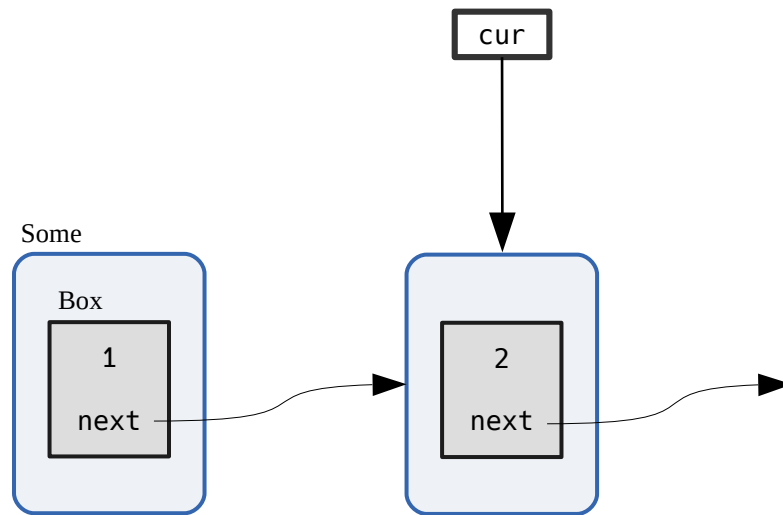
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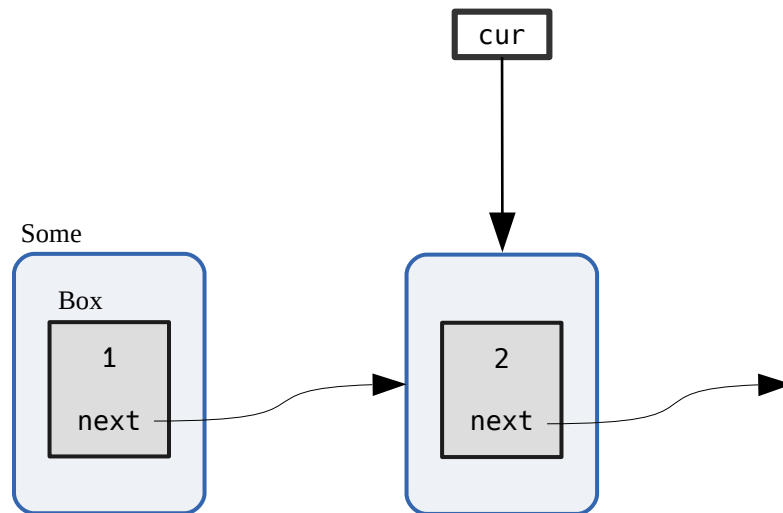
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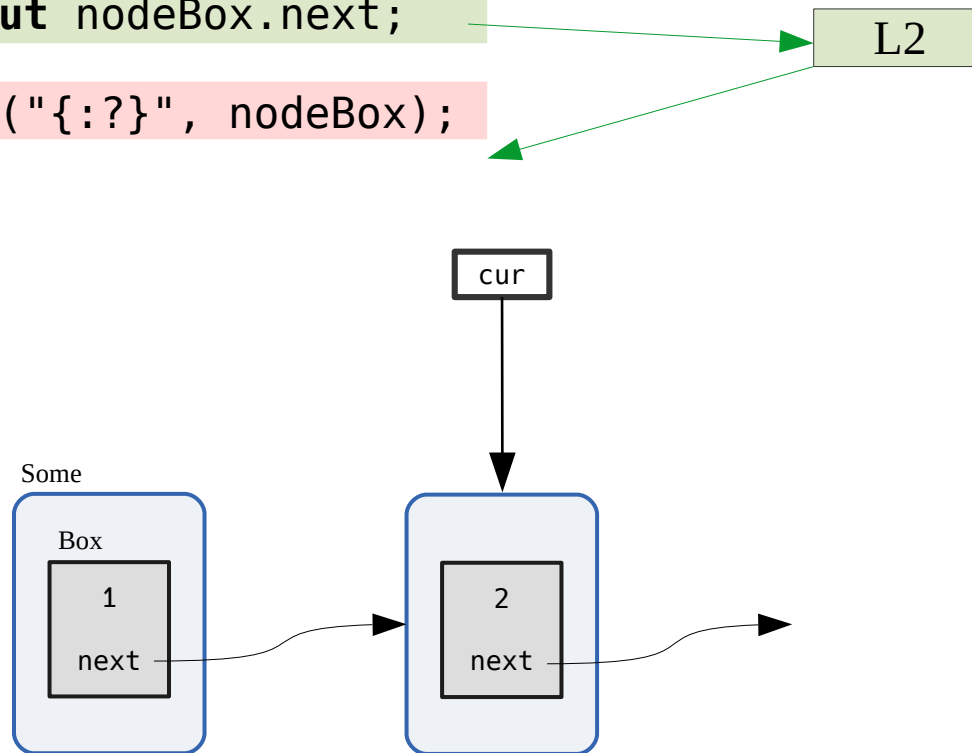
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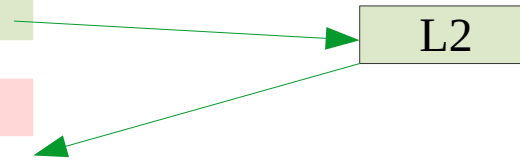
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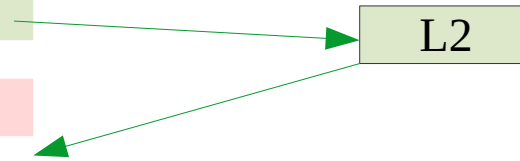
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Why **L2** live at this point ?

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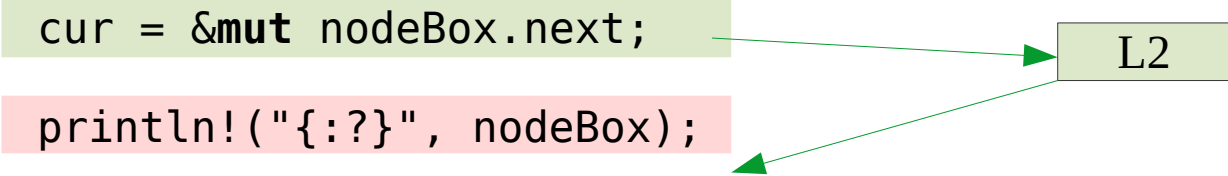
Why **`L2'** live at this point ?

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Why **`L2'** live at this point ?

1. no assignment to { nodeBox, nodeBox.next }
2. **`cur'** will be used later

It is rejected in the current borrow checker, but it is accepted by the Polonius borrow checker in the future.

## Example

## 參考題目

```
fn do_something(mut head: Option<Box<ListNode>>)  
{
```

```
    let mut cur = &mut head;
```

**Leetcode : remove linked list elements**

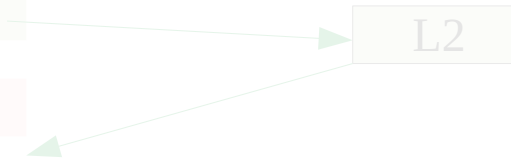
```
    if let Some(nodeBox) = cur.as_mut() {
```

```
        cur = &mut nodeBox.next;
```

```
        println!("{:?}", nodeBox);
```

```
    }
```

```
}
```



Why `L2` live at this point ?

1. no assignment to { nodeBox, nodeBox.next }
2. `cur` will be used later

It is rejected in the current borrow checker, but it is accepted by the Polonius borrow checker in the future.

# Datafrog

The tool used in Rust's new borrow checker called  
Polonius

# Idea

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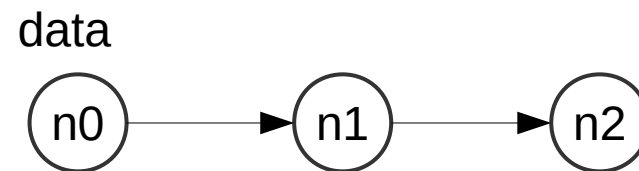
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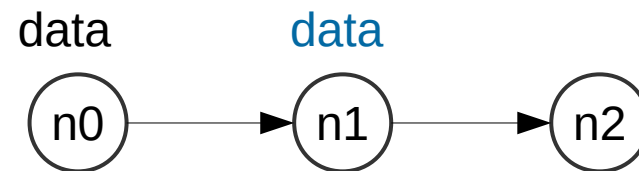
$$\frac{N(a, x) \cdot e(a, b)}{N(b, x)}$$



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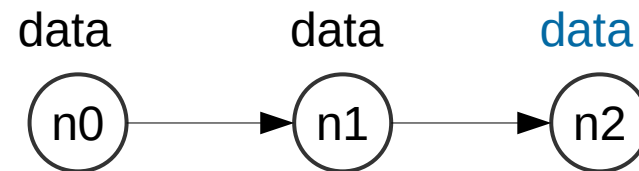




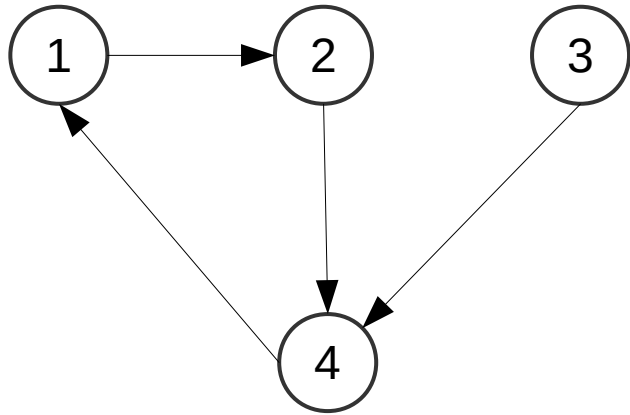
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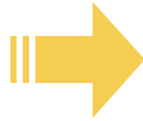
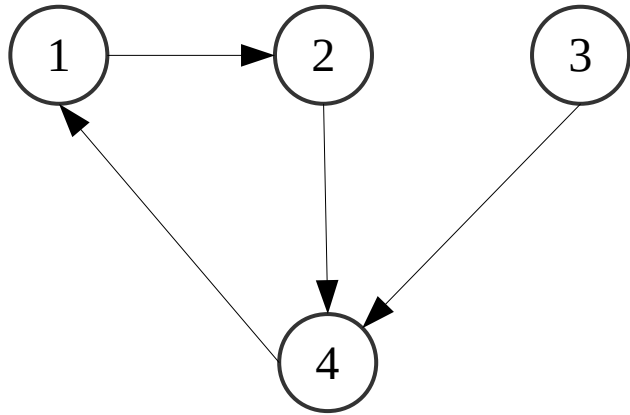
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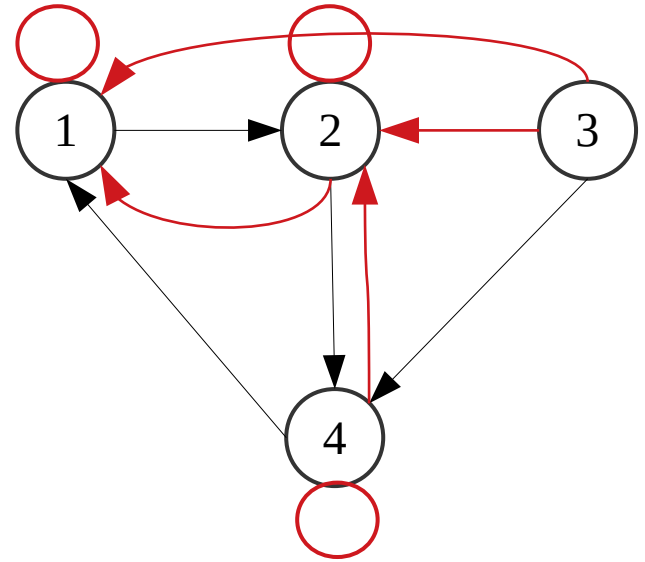
## Example . Transitive Closure



# Example . Transitive Closure



$$\frac{e(a,b) \quad e(b,c)}{e(a,c)}$$



# Implementation – Initial

```
// create a iteration context
let mut iteration = Iteration::new();

// create some variables for later use
let v_edges  = iteration.variable::<(u32, u32)>("edges");
let v_redges = iteration.variable::<(u32, u32)>("reverse edges");

// load the initial variables
v_edges.insert(edges.into());

// start iteration
while iteration.changed() {

    ...

}


let result = v_edges.complete();
```

# Implementation – Initial

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// start iteration
while iteration.changed() {
     Writing Rules here
}

let result = v_edges.complete();
```

# Implementation – Writing Rules

```
while iteration.changed() {  
    // reverse edges for mapping  
    v_redges.from_map(&v_edges, |&(a, b)| (b, a));  
  
    // e(a,c) <- e(a,b), e(b,c)  
    v_edges.from_join(&v_redges, &v_edges, |_b, &a, &c| (a, c));  
}
```

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$$\frac{e(a,b)}{e(b,c)} \frac{e(b,c)}{e(a,c)}$$



$$\frac{e(a,b)}{r(b,a)}$$

$$\frac{r(b,a)}{e(b,c)} \frac{e(b,c)}{e(a,c)}$$



# Implementation – Writing Rules

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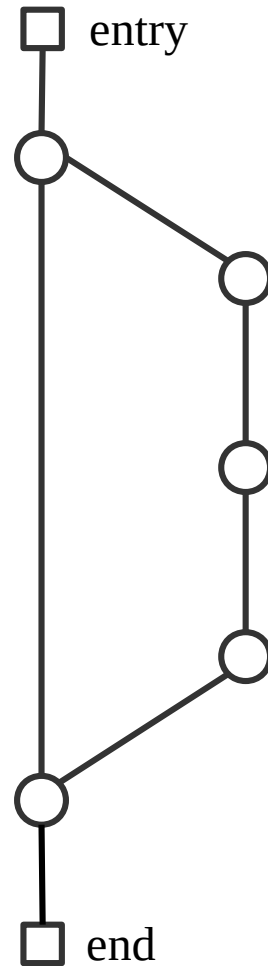
```

$$\frac{e(a,b)}{e(b,c)} \quad \Rightarrow \quad \frac{e(a,b)}{r(b,a)} \quad \frac{r(\textcolor{red}{b}, \textcolor{green}{a})}{e(\textcolor{red}{b}, \textcolor{blue}{c})} \quad \frac{e(a,c)}{e(a,c)}$$

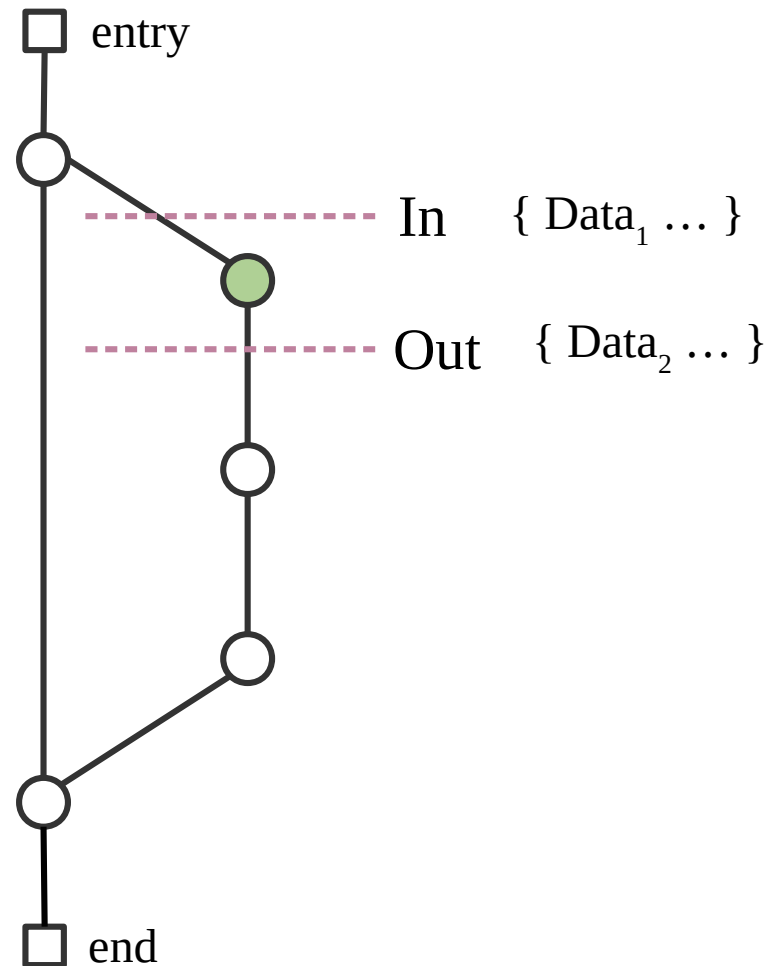
In Polonius

$$\begin{array}{c}
\text{Require}(R, B, P) \\
\text{cfg-edges}(P, Q) \\
!\text{killed}(B, P) \\
\text{region\_live\_at}(R, Q) \\
\hline
\text{Require}(R, B, Q)
\end{array}$$

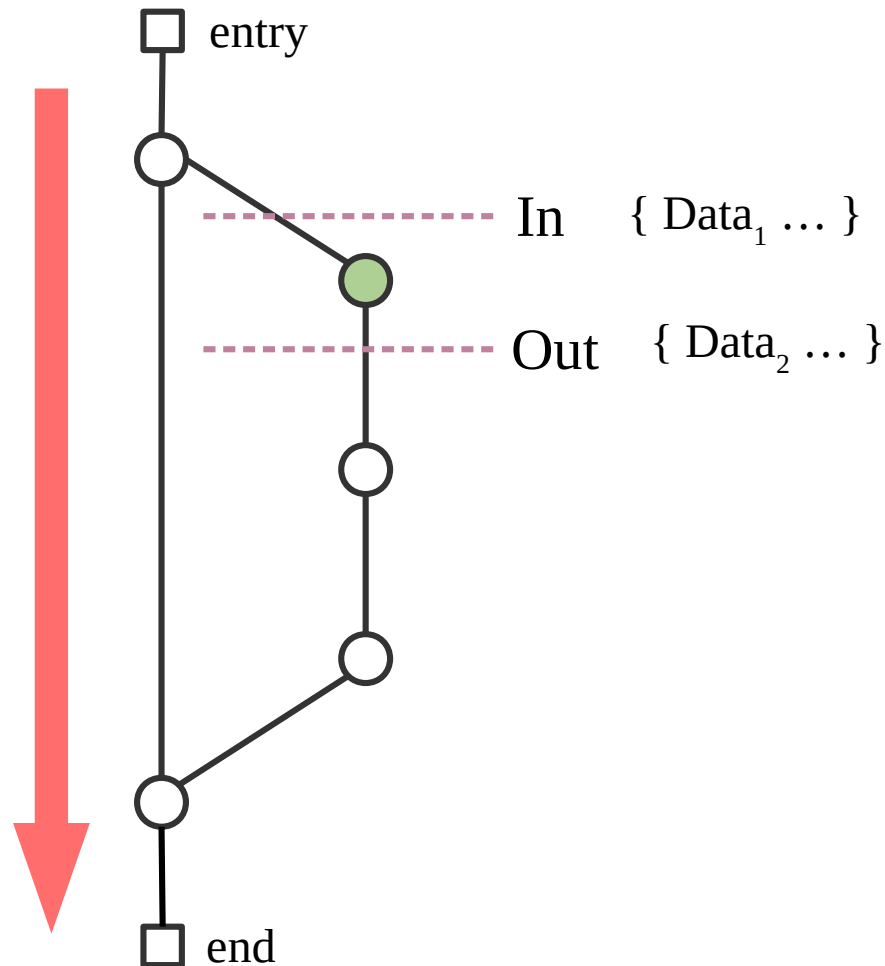
# Data Flow Concepts $\langle D, V, \wedge, F \rangle$



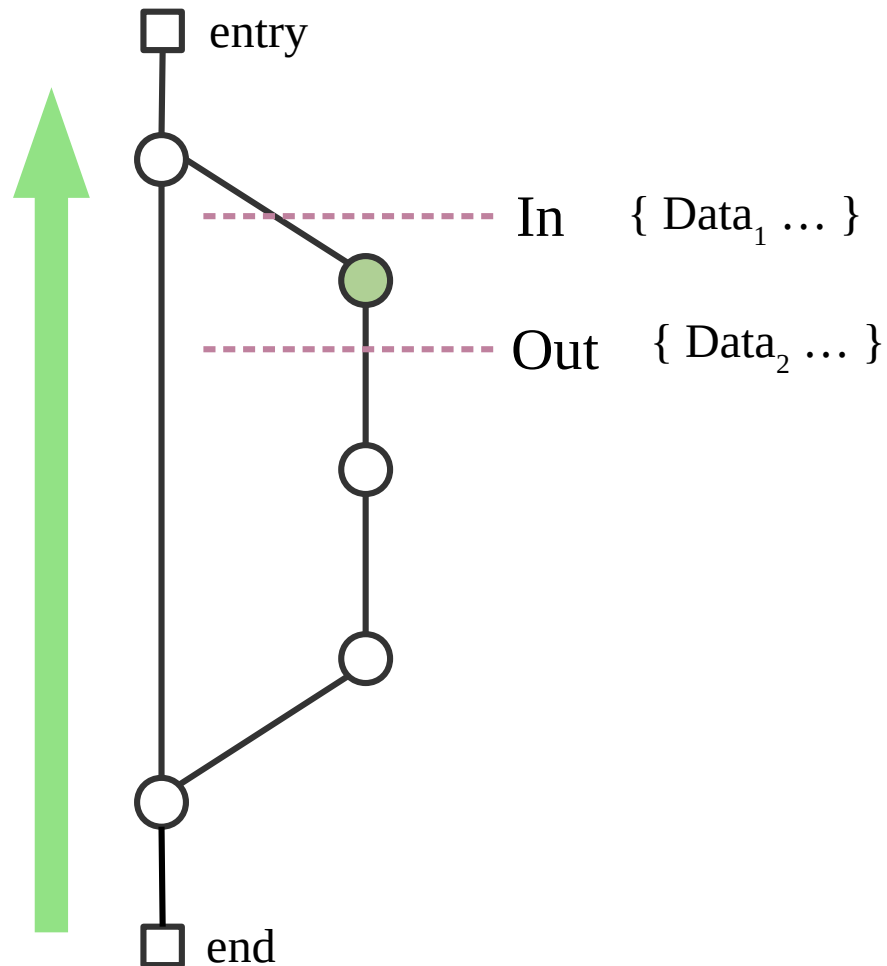
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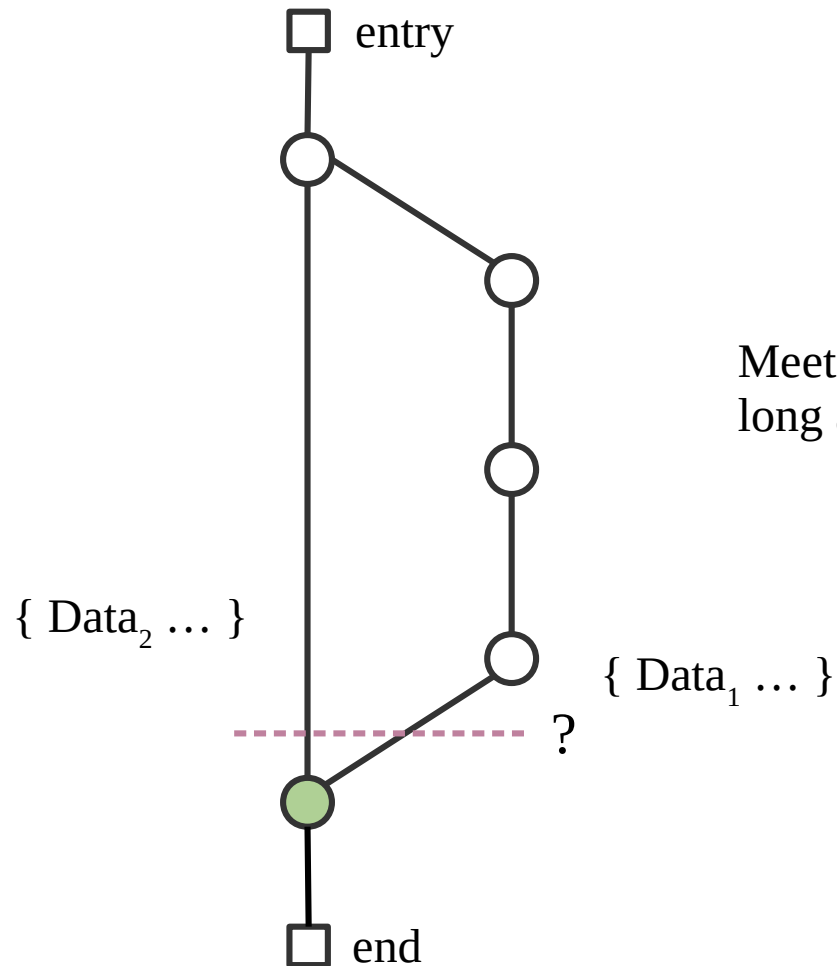
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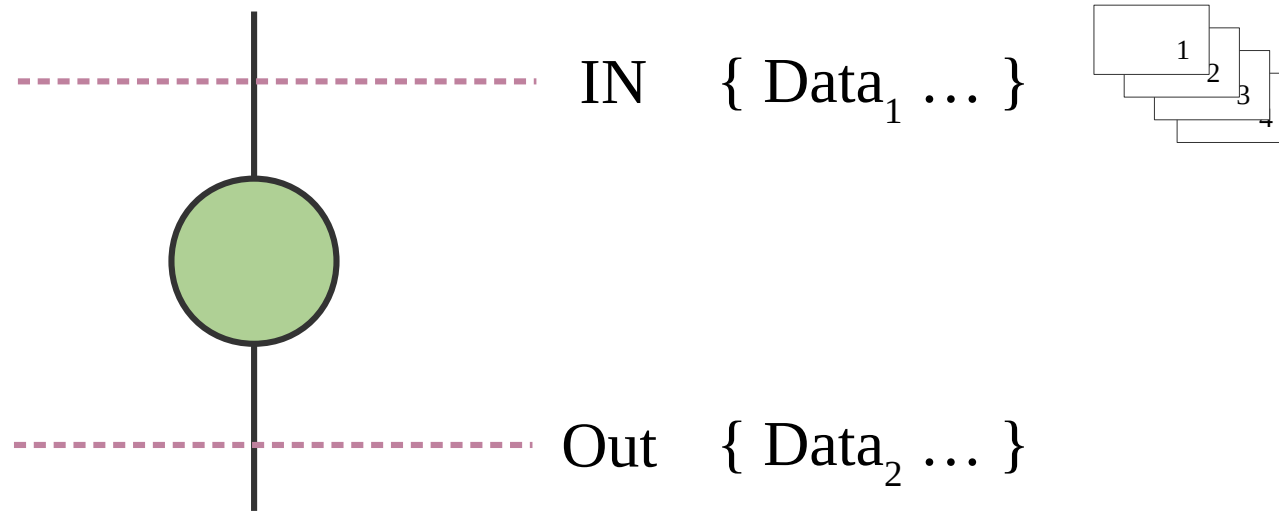
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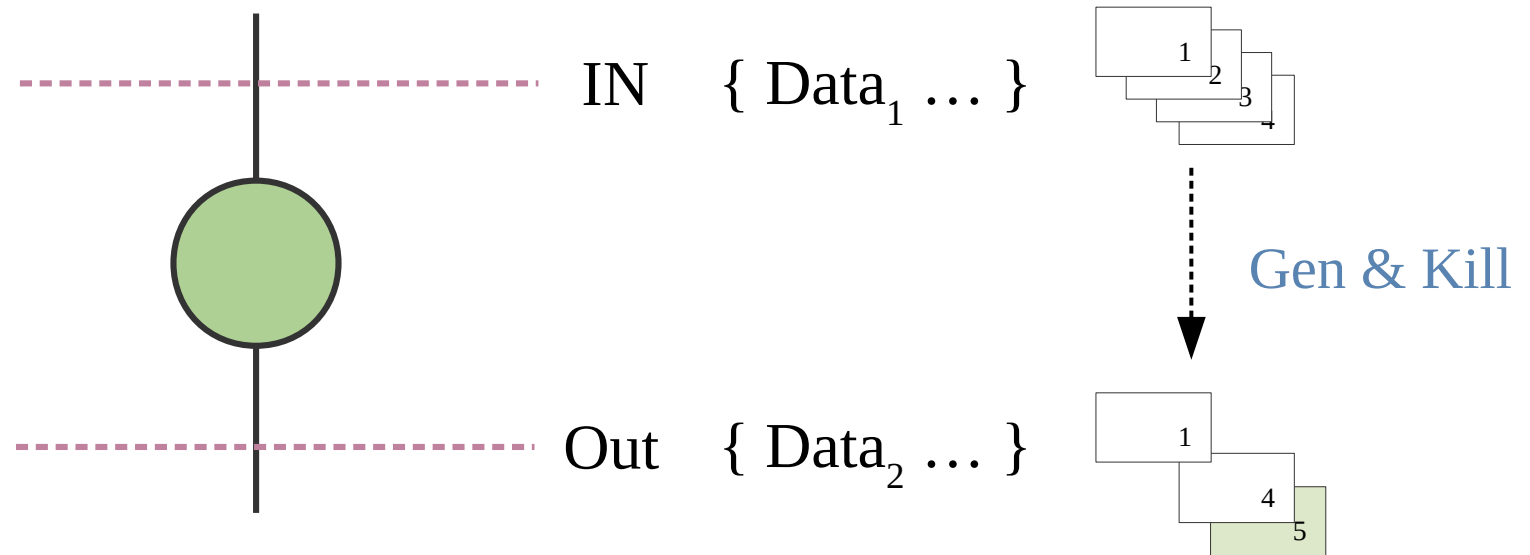
Meet operator can be union, intersection, as long as it can make  $(V, \wedge)$  semilattice.



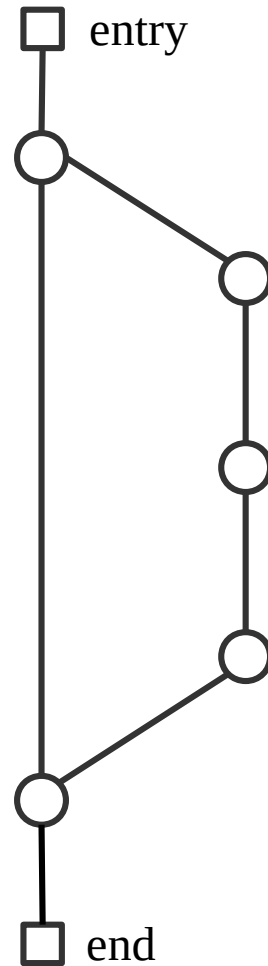
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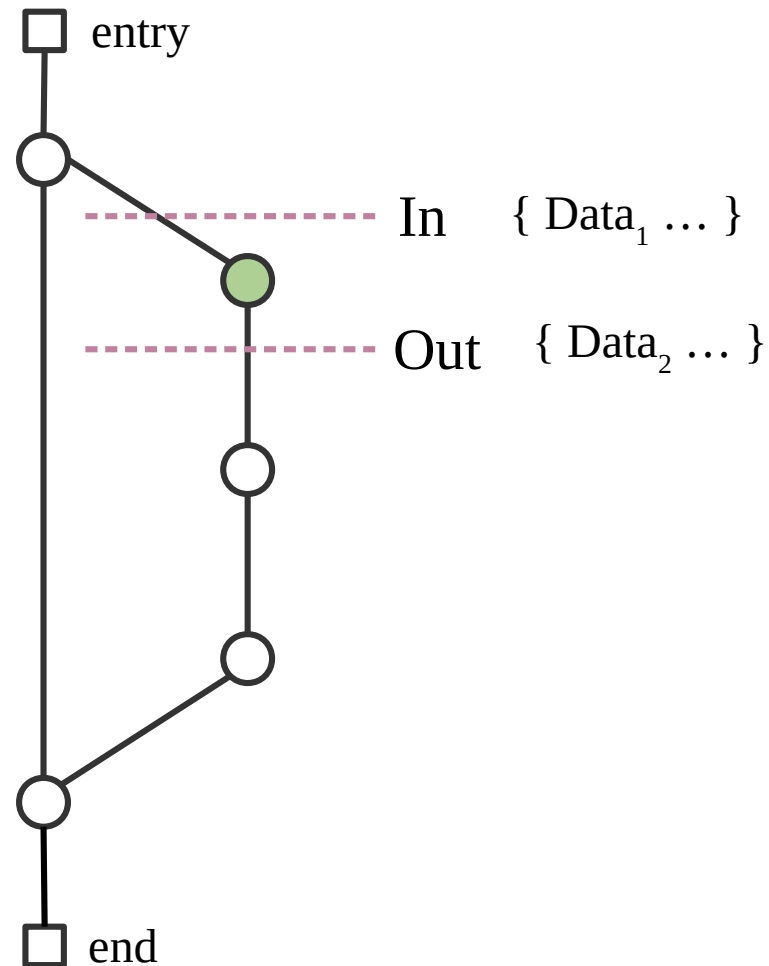
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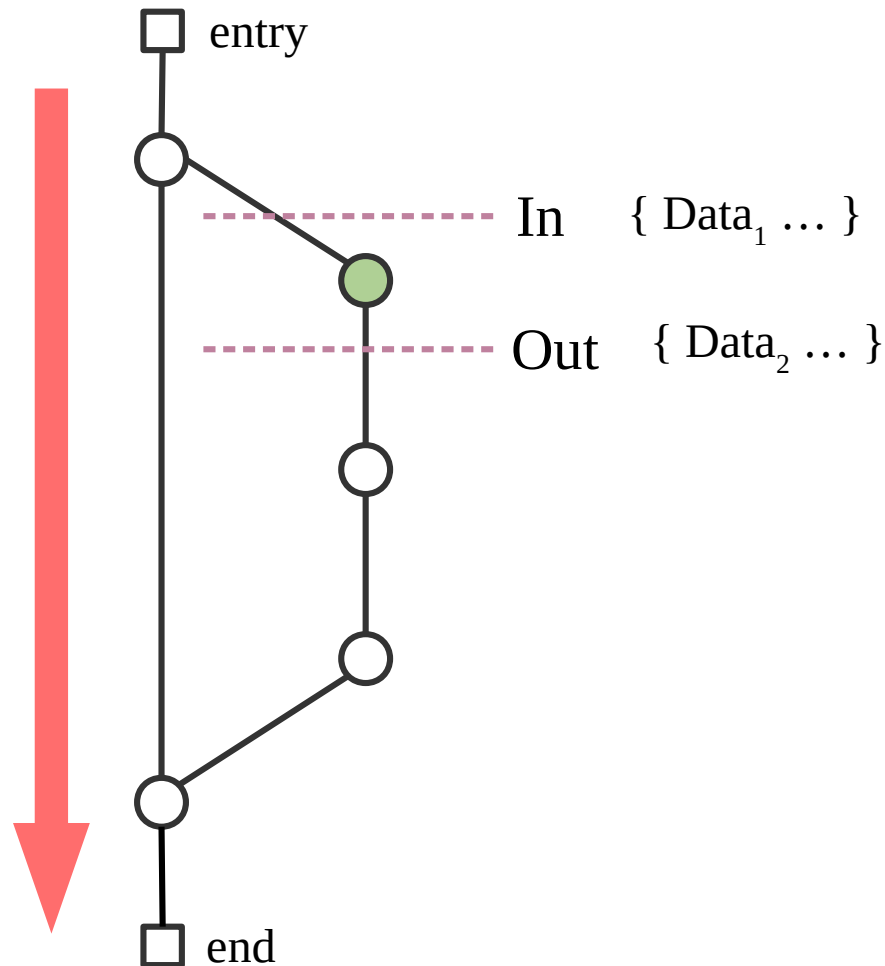
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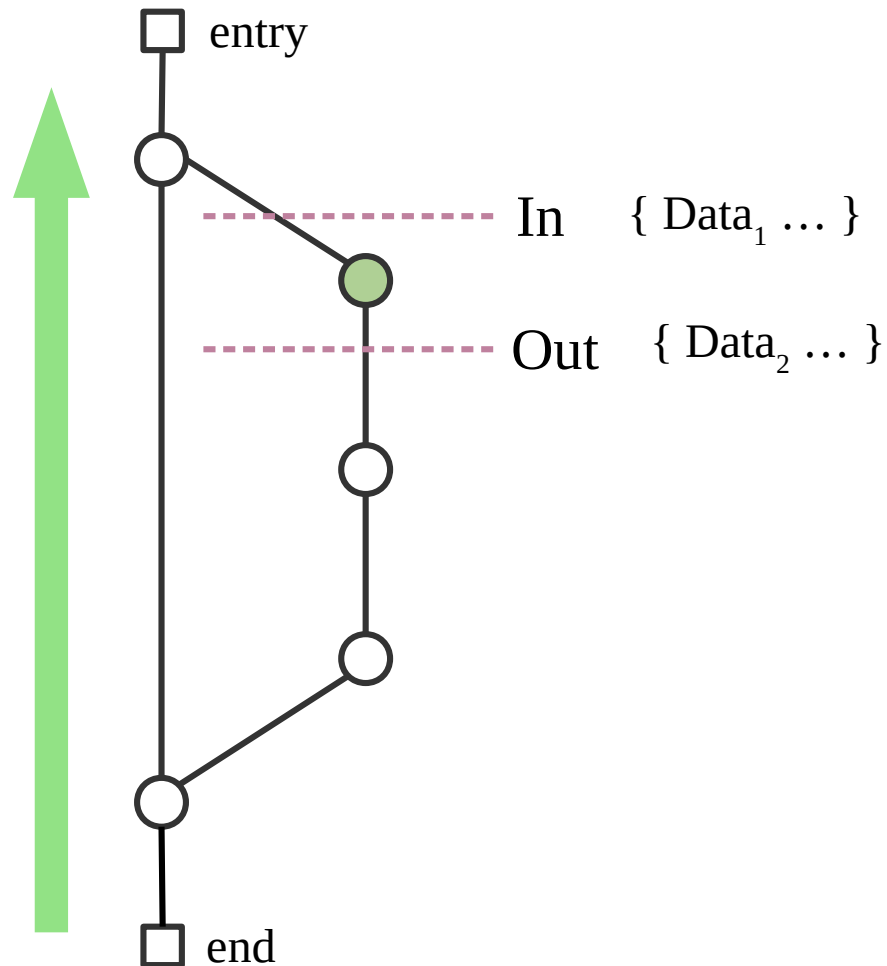
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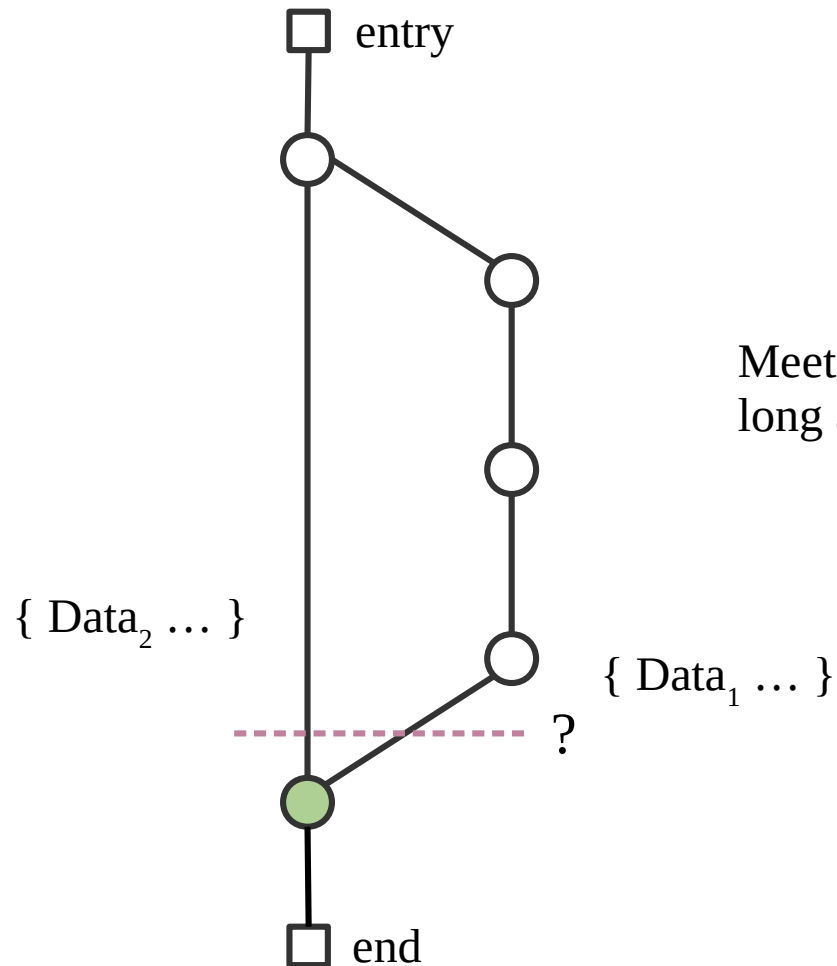
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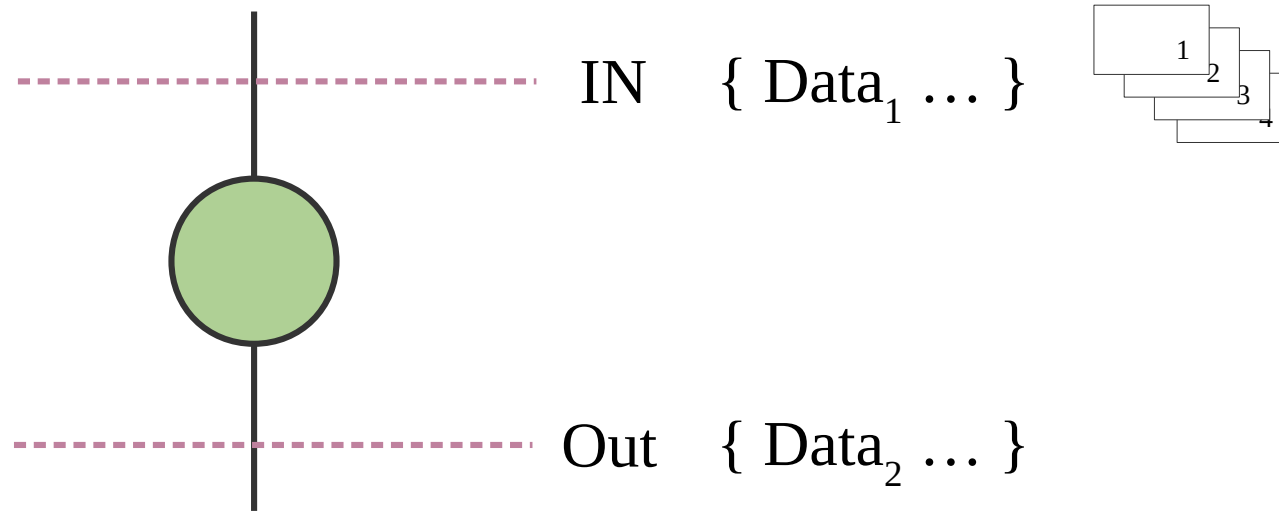


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