

What are Transaction Control Commands?

When you group SQL statements into a transaction, you need a way to **start**, **save**, **cancel**, or **complete** that transaction. These commands are like the buttons you press to control what happens to your group of SQL actions.

Let's break them down one by one in plain language:



BEGIN or START TRANSACTION

What it does:

This command marks the beginning of a transaction. It tells the database, "From this point forward, I'm grouping actions together."

How to explain to students:

"Imagine you're starting a recipe. You say, 'I'm beginning now, so if I mess up somewhere, I want to be able to throw everything away and start fresh.' That's what BEGIN does in SQL. It signals the start of something important that you may want to undo if needed."



COMMIT

What it does:

This command makes all the changes permanent. After a COMMIT, the database saves the results of everything done since BEGIN.

How to explain to students:

"Think of COMMIT as clicking Save on a document. After you commit, your work is stored — it's locked in. Even if the power goes out, the changes stay. You're saying, 'I'm done and happy with everything I've done."



ROLLBACK

What it does:

It cancels everything done in the transaction so far — nothing is saved. It returns the database to the way it was when you used BEGIN.

How to explain to students:

"This is like pressing Undo All or clicking 'Don't Save' and closing the file. If you made a mistake halfway, ROLLBACKlets you erase all the changes and start over."



SAVEPOINT

What it does:

This creates a **checkpoint** inside a transaction — a point you can return to if something goes wrong after it.

How to explain to students:

"Imagine you're writing an essay, and after finishing the introduction, you make a backup copy. That's a SAVEPOINT. If you mess up the next section, you don't have to throw everything away — you can go back to the backup."

You give it a name like SAVEPOINT intro done; so you can refer to it later.

◆ ROLLBACK TO SAVEPOINT

What it does:

It undoes part of the transaction — only the steps that happened after the savepoint.

How to explain to students:

"This is like saying, 'I liked what I had before this paragraph — let me go back to that version and fix it from there.' You don't erase everything, just the stuff you did after your last savepoint."

♦ RELEASE SAVEPOINT

What it does:

It **removes a savepoint** you no longer need. You can't roll back to it after it's released.

How to explain to students:

"Think of this like deleting that backup copy because you've moved on. It helps clean up memory and avoid confusion later on."

Wey Points to Emphasize to Students

- A transaction **starts with BEGIN**.
- Once you're sure everything is correct, you use **COMMIT**.
- If there's a mistake, use ROLLBACK to cancel.
- You can create **mini-backups** using SAVEPOINT.
- You can go back to a **safe point** using ROLLBACK TO SAVEPOINT.
- Once done, **clean up** with RELEASE SAVEPOINT.

4. Autocommit Behavior – Explained for Beginners



"When you write and run a single SQL statement — like an INSERT, UPDATE, or DELETE — many databases automatically save the changes right after that statement finishes. This is called autocommit.

It means:

- You don't have to say COMMIT the database does it for you behind the scenes.
- Every statement is treated like a tiny, automatic transaction."

Why This Matters

"This seems convenient, but it can be **dangerous** if you're running multiple steps that need to go together.

Let's say you want to do these in one group:

- 1. Subtract money from Account A
- 2. Add money to Account B

If autocommit is ON, and the first statement runs but the second fails, only half of your task happens. That's bad — your data is now **incomplete** and **inconsistent**.

You want **both** steps to succeed together or **none at all** — that's why you need manual transaction control."

Mow to Turn Off Autocommit

"In many systems, you can disable autocommit like this:

• In **MySQL**, run:

sql CopyEdit

SET autocommit = 0;

• In **PostgreSQL**, start a transaction manually:

sql CopyEdit

BEGIN;

•

- and it stays open until you COMMIT or ROLLBACK.
- In **programming languages** (like Python or Java), database drivers usually let you turn off autocommit in the connection settings.

Once autocommit is off:

- You control when to **commit**.
- You can **rollback** if there's a problem.
- You get the full power of transactions."

West Message for Students

"Autocommit is good for quick and simple changes. But for anything important or multi-step, **turn it off** so you can use proper transactions."

5. Isolation Levels – Explained for Beginners

What Are Isolation Levels?

"Imagine a busy kitchen where multiple chefs are preparing different meals. If they start using the same ingredients at the same time without rules, there'll be chaos — missing items, wrong orders, confusion.

The same happens in a **multi-user database**. If many people (or apps) are reading and writing at the same time, things can go wrong.

Isolation levels are the rules that control how transactions interact with each other. They decide how much one transaction can 'see' or be affected by others."

Why Isolation Matters

"Let's say:

- One transaction is reading the price of a product.
- Another transaction is updating that price at the same time.

Depending on the isolation level, the first transaction might:

- See the **old price**
- See the **new price**
- Or even see a price that's not finalized (a *dirty read*)!

That's why isolation levels exist: to **protect your data** when multiple things are happening at once."

The Four Isolation Levels

Explain them from weakest to strongest:

♦ 1. Read Uncommitted

- Most relaxed level.
- Transactions can see changes made by others even if those changes haven't been committed yet.
- This allows **dirty reads** risky because the other transaction might roll back.

"It's like reading someone's notes before they've finished writing. It might be wrong or get erased."

♦ 2. Read Committed

- Can only read data that's **already committed**.
- No dirty reads.
- But if another transaction updates the data in the meantime, you may see **different values if** you read twice.

"It's like checking a webpage twice and seeing different content each time."

♦ 3. Repeatable Read

- When you read data once, it **stays the same** for the whole transaction.
- Prevents **non-repeatable reads**.
- But **new rows** inserted by others may still be visible (phantom reads).

"It's like checking a document and freezing what you see — even if someone else edits it later."

• 4. Serializable

- Most strict and safest.
- Transactions are completely **isolated** as if they were running one at a time.

 Prevents all kinds of concurrency problems: dirty reads, non-repeatable reads, phantom reads.

"It's like you're the only person in the library — no one else can change anything while you're working."



"In most databases, you can set the isolation level at the beginning of a transaction:

sql
CopyEdit
SET TRANSACTION ISOLATION LEVEL READ COMMITTED;
BEGIN;
-- your SQL here
COMMIT;
You can also set it at the session level depending on your database."

Trade-offs: Performance vs. Safety

- **Lower isolation** (like Read Uncommitted) gives **better performance** faster and fewer locks.
- **Higher isolation** (like Serializable) gives **more safety** but can slow things down due to more locking and blocking.

"You have to **balance safety and speed**. For critical operations, use strict levels. For quick, non-critical reads, a lower level may be okay."

Key Message for Students

"Think of isolation levels as how *private* your transaction is.

The more private (isolated) it is, the safer it is — but also slower.

Choose the level based on how important accuracy is for your task."

6. Common Transaction Problems – SQL Server Focus

V Why Transaction Problems Happen

"When we don't carefully manage transactions — especially in systems with many users — problems can pop up. These problems don't always show up right away, but they can lead to **wrong results**, **inconsistent reports**, and even **data corruption** over time.

Let's go over the four major problems."

1. Dirty Reads

What is it?

A dirty read happens when a transaction reads data that was **changed by another transaction but not committed yet.** If that other transaction later rolls back, the data you read was never real - it was temporary and incorrect.

Example:

One transaction updates a customer balance. Another transaction reads the new balance before it's committed. Later, the update is rolled back — but the second transaction already used that false balance.

SQL Server Isolation Level to Prevent:

Use **Read Committed** or higher to block dirty reads.

(Default in SQL Server is READ COMMITTED, which does block dirty reads.)

♦ 2. Non-repeatable Reads

What is it?

When a transaction reads the same row twice and gets different values each time because another transaction changed the row in between.

Example:

You read a product's price. Another transaction changes it. You read it again and see a different price — in the same transaction.

SQL Server Isolation Level to Prevent:

Use Repeatable Read or higher.

This level locks rows for reading, so no other transaction can change them until your transaction ends.



3. Phantom Reads

What is it?

When a transaction reads a set of rows based on a condition, but later in the same transaction, it reads again and sees new rows that weren't there before — because another transaction inserted them.

Example:

You run a query to get all orders over \$100. Another user adds a new \$150 order. When you run the same query again, the result has changed.

SQL Server Isolation Level to Prevent:

Use Serializable.

It locks the entire range of rows, so others can't add new matching rows until your transaction is done.



4. Lost Updates

What is it?

Two transactions read the same data and update it, but one of the updates is overwritten by the other without warning.

Example:

Two people change the same product stock level at the same time. The last one to commit overwrites the changes made by the first.

SQL Server Fix:

This needs explicit row locking or handling with isolation level like Repeatable Read or **optimistic concurrency** (checking if data changed before updating).



Key Summary for Students

Problem	Fix by Using Isolation Level
Dirty Reads	Read Committed or higher
Non-repeatable Reads	Repeatable Read or higher
Phantom Reads	Serializable only
Lost Updates	Careful locking or versioning

7. Transactions and Error Handling – SQL Server (T-SQL)



Why Link Transactions with Error Handling?

"Even when we use transactions, things can go wrong:

- A statement might fail.
- A constraint might be violated.
- A deadlock might happen.

So we need to make sure that if something goes wrong inside a transaction, we don't accidentally commit broken or partial data."



How to Handle Errors in Transactions – T-SQL Style

In **SQL Server**, we use **TRY...CATCH** blocks in T-SQL to handle errors inside transactions.

Basic Structure:

```
sql
CopyEdit
BEGIN TRY
    BEGIN TRANSACTION;
    -- your SQL statements here
    COMMIT;
END TRY
BEGIN CATCH
    ROLLBACK;
    -- handle the error (optional logging, message, etc.)
    PRINT ERROR_MESSAGE();
END CATCH;
```

X Explain This Step-by-Step

- 1. **BEGIN TRANSACTION** Start the transaction.
- 2. Inside the TRY, run your SQL statements.
- 3. If everything works, COMMIT saves the changes.
- 4. If anything fails, SQL Server jumps to the CATCH block.
- 5. In the CATCH, you ROLLBACK to undo all changes.
- 6. You can also log the error or show a message.

Why This Matters

"If we don't catch errors and roll back, SQL Server might leave the transaction open, or worse, partially applied. That means **some changes happened, and some didn't** — exactly what we don't want."

Best Practices for Students

- Always pair transactions with error handling.
- Use **TRY...CATCH** when doing inserts, updates, or deletes in groups.

- Don't forget the ROLLBACK in the CATCH block.
- Optionally, use XACT_STATE () to check if the transaction is still valid before committing.