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IIUM cat-us-trophy Contents

```
Data structures (BIT, BIT2D, RMQ-DP, RMQ-segment tree, union-find, misof's tree) . . .
  Graph algorithms (LCA, SCC, BPM-konig, Bellman-Ford, NetFlow, MinCost MaxFlow) . . . . .
  15
  .vimrc
set ai ts=4 sw=4 st=4 noet nu nohls
syntax enable
filetype plugin indent on
map <F6> :w<CR>:!g++ % -g && (ulimit -c unlimited; ./a.out < ~/input.txt) <CR>
map <F5> <F6>
colo pablo
map <F12> :!gdb ./a.out -c core <CR>
template.cpp
#include<cstdio>
#include<sstream>
#include<cstdlib>
#include<cctype>
#include<cmath>
#include < algorithm >
#include<set>
#include<queue>
#include<stack>
#include<list>
#include<iostream>
#include<string>
#include<vector>
#include<cstring>
#include<map>
#include<cassert>
#include<climits>
using namespace std;
#define REP(i,n) for(int i=0, _e(n); i<_e; i++)
#define FOR(i,a,b) for(int i(a), _e(b); i <= _e; i++)
#define FORD(i,a,b) for(int i(a), _e(b); i>=_e; i--)
#define FORIT(i, m) for (__typeof((m).begin()) i=(m).begin(); i!=(m).end(); ++i)
#define SET(t,v) memset((t), (v), sizeof(t))
#define ALL(x) x.begin(), x.end()
#define UNIQUE(c) (c).resize( unique( ALL(c) ) - (c).begin() )
#define sz size()
#define pb push_back
```

Combinatorics

Mathematical Sums

$$\begin{array}{ll} \sum_{k=0}^{n} k = n(n+1)/2 & \sum_{k=0}^{b} k = (a+b)(b-a+1)/2 \\ \sum_{k=0}^{n} k^2 = n(n+1)(2n+1)/6 & \sum_{k=0}^{n} k^3 = n^2(n+1)^2/4 \\ \sum_{k=0}^{n} k^4 = (6n^5+15n^4+10n^3-n)/30 & \sum_{k=0}^{n} k^5 = (2n^6+6n^5+5n^4-n^2)/12 \\ \sum_{k=0}^{n} x^k = (x^{n+1}-1)/(x-1) & \sum_{k=0}^{n} kx^k = (x-(n+1)x^{n+1}+nx^{n+2})/(x-1)^2 \end{array}$$

Binomial coefficients

	0	1	2	3	4	5	6	7	8	9	10	11	12	$\binom{n}{l} = \frac{n!}{(l-1)! l!}$
0	1													$\binom{k}{k} = \frac{(n-k)!k!}{\binom{n}{k}} + \binom{n-1}{k-1}$
1	1	1												(n) n (n-1)
2	1	2	1											
3	1	3	3	1										
4	1	4	6	4	1									$\binom{k}{k} - \frac{1}{n-k+1} \binom{k}{k}$
5	1	5	10	10	5	1								$\binom{n}{k+1} = \frac{n-k}{k+1} \binom{n}{k}$
6	1	6	15	20	15	6	1							
7	1	7	21	35	35	21	7	1						$\sum_{n=1}^{n} (n)$ and 1
8	1	8	28	56	70	56	28	8	1					$\sum_{k=1}^{n} k \binom{n}{k} = n2^{n-1}$ $\sum_{k=1}^{n} k^2 \binom{n}{k} = (n+n^2)2^{n-2}$
9	1	9	36	84	126	126	84	36	9	1				$\sum_{k=1}^{n} k^{2} \binom{n}{k} = (n+n^{2}) 2^{n-1}$
10	1	10	45	120	210	252	210	120	45	10	1			
11	1	11	55	165	330	462	462	330	165	55	11	1		$\binom{m+n}{r} = \sum_{i=1}^{r} \binom{m}{i} \binom{n}{i}$
12	1	12	66	220	495	792	924	792	495	220	66	12	1	$\begin{pmatrix} r \end{pmatrix} - \angle k = 0 \begin{pmatrix} k \end{pmatrix} \begin{pmatrix} r - k \end{pmatrix}$
	0	1	2	3	4	5	6	7	8	9	10	11	12	$\binom{n}{k} = \prod_{i=1}^k \frac{n-k+i}{i}$

Catalan numbers $C_n = \frac{1}{n+1} {2n \choose n}$. $C_0 = 1$, $C_n = \sum_{i=0}^{n-1} C_i C_{n-1-i}$. $C_{n+1} = C_n \frac{4n+2}{n+2}$. $C_0, C_1, \ldots = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900, \ldots$

 C_n is the number of: properly nested sequences of n pairs of parentheses; rooted ordered binary trees with n+1 leaves; triangulations of a convex (n+2)-gon.

Derangements. Number of permutations of $n=0,1,2,\ldots$ elements without fixed points is $1,0,1,2,9,44,265,1854,14833,\ldots$ Recurrence: $D_n=(n-1)(D_{n-1}+D_{n-2})=nD_{n-1}+(-1)^n$. Corollary: number of permutations with exactly k fixed points is $\binom{n}{k}D_{n-k}$.

Stirling numbers of 1^{st} kind . $s_{n,k}$ is $(-1)^{n-k}$ times the number of permutations of n elements with exactly k permutation cycles. $\begin{bmatrix} n \\ k \end{bmatrix} = |s_{n,k}| = |s_{n-1,k-1}| + (n-1)|s_{n-1,k}|$ s(0,0) = 1 and s(n,0) = s(0,n) = 0.

Stirling numbers of 2^{nd} kind. $S_{n,k}$ is the number of ways to partition a set of n elements into exactly k non-empty subsets. $\binom{n}{k} = S_{n,k} = S_{n-1,k-1} + kS_{n-1,k}$. $S_{n,1} = S_{n,n} = 1$.

Bell numbers . B_n is the number of partitions of n elements. $B_0, \ldots = 1, 1, 2, 5, 15, 52, 203, 877, \ldots$ $B_{n+1} = \sum_{k=0}^{n} \binom{n}{k} B_k = \sum_{k=1}^{n+1} S_{n,k}$. Bell triangle: $B_r = a_{r,1} = a_{r-1,r-1}, \ a_{r,c} = a_{r-1,c-1} + a_{r,c-1}$.

Eulerian numbers . $E(n,k) = {n \choose k}$ is the number of permutations with exactly k descents $(i : \pi_i < \pi_{i+1}) / \text{ascents } (\pi_i > \pi_{i+1}) / \text{excedances } (\pi_i > i) / k + 1 \text{ weak excedances } (\pi_i \ge i)$.

```
Formula: E(n,m) = (m+1)E(n-1,m) + (n-m)E(n-1,m-1). E(n,0) = E(n,n-1) = 1.
E(n,m) = \sum_{k=0}^{m} (-1)^k \binom{n+1}{k} (m+1-k)^n.
Double factorial. Permutations of the multiset \{1, 1, 2, 3, \dots n, n\} such that for each k, all the num-
bers between two occurrences of k in the permutation are greater than k. (2n-1)!! = \prod_{k=1}^{n} (2k-1).
Eulerian numbers of 2^{nd} kind. Related to Double factorial, number of all such permutations
that have exactly m ascents. \left\langle \left\langle {n \atop m} \right\rangle \right\rangle = (2n-m-1)\left\langle \left\langle {n-1 \atop m-1} \right\rangle \right\rangle + (m+1)\left\langle \left\langle {n-1 \atop m} \right\rangle \right\rangle. \left\langle \left\langle {n \atop 0} \right\rangle \right\rangle = 1
Multinomial theorem (a_1 + \dots + a_k)^n = \sum_{n_1,\dots,n_k} \binom{n}{n_1,\dots,n_k} a_1^{n_1} \dots a_k^{n_k}, where n_i \ge 0 and \sum_{n_i = n} n_i = n. \binom{n}{n_1,\dots,n_k} = M(n_1,\dots,n_k) = \frac{n!}{n_1!\dots n_k!}. M(a,\dots,b,c,\dots) = M(a+\dots+b,c,\dots)M(a,\dots,b)
Data structures (BIT, BIT2D, RMQ-DP, RMQ-segment tree, union-find,
misof's tree)
BIT - Binary indexed trees
int bit[M],n;
void update(int x, int v) { while( x \le n ) { bit[x] += v; x += x & -x; } }
int sum(int x) { int ret=0; while(x>0){ ret += bit[x]; x == x \& -x; } return ret; }
BIT 2D
int bit[M][M], n;
int sum( int x, int y ){
     int ret = 0;
     while (x > 0)
           int yy = y; while(yy > 0) ret += bit[x][yy], yy -= yy & -yy;
           x = (x \& -x);
     }
     return ret ;
}
void update(int x , int y , int val){
     int y1;
     while (x \le n){
           y1 = y;
           while (y1 \le n) \{ bit[x][y1] += val; y1 += (y1 & -y1); \}
           x += (x \& -x);
     }
}
RMQ DP
int make_dp(int n) { // N log N
     REP(i,n) H[i][0]=i;
     for(int l=0,k; (k=1<<1) < n; l++) for(int i=0;i+k< n;i++)
           H[i][1+1] = A[H[i][1]] > A[H[i+k][1]] ? H[i+k][1] : H[i][1];
} // query log N almost O(1)
int query_dp(int a, int b) {
     for(int l=0;;l++) if (a+(1<<l+1) > b) {
           int o2 = H[b-(1<<1)+1][1];
```

}

}

return A[H[a][1]] < A[o2] ? H[a][1]:o2;

```
const int M = 100005;
int n, in[M], f[M], st[M], en[M];
struct data { int 1, r, ans, next_1, next_r; };
data d[4*M]; // which is the range? :S
int nd;
int build( int 1, int r, int id ) {
    d[ id ].l = l, d[ id ].r = r;
    if( l == r ) d[ id ].ans = f[1];
    else {
        int bar = (r-1)/2+1;
        d[id].next_1 = ++nd, d[id].next_r = ++nd;
        int left = build( 1, bar, d[id].next_l );
        int right = build( bar+1, r, d[id].next_r );
        d[id].ans = max( left, right );
    return d[ id ].ans;
}
int query( int 1, int r, int id = 0 ) {
    if (1 > r) return 0;
    if( d[id].l == l && d[id].r == r ) return d[id].ans;
    else {
        int bar = (d[id].r-d[id].1) / 2 + d[id].1;
        int left = 0, right = 0;
        if( 1 <= bar ) {
            if( r <= bar ) left = query( 1, r, d[id].next_1 );</pre>
            else {
                left = query( 1, bar, d[id].next_l );
                right = query( bar+1, r, d[id].next_r );
        }else right = query( l, r, d[id].next_r );
        return max( left, right );
    }
}
Union find - rank by number of elements
struct unionfind {
    int p[MAX], r[MAX]; // r contains the population
    unionfind() { REP(i,MAX) p[i] = i, r[i] = 1; }
    int find( int x ) { if( p[x] == x ) return x; else return p[x] = find( p[x] ); }
    void Union(int x, int y) {
        int px = find(x), py = find(y);
        if( px == py ) return; //already joined
        if( r[px] < r[py] ) p[px] = py, r[py] += r[px];
        else p[ py ] = px, r[px] += r[py];
    }
};
misof's tree
int tree[17][65536];
void insert(int x) { for (int i=0; i<17; i++) { tree[i][x]++; x/=2; } }
void erase(int x) { for (int i=0; i<17; i++) { tree[i][x]--; x/=2; } }
```

```
int kThElement(int k) { int a=0, b=16;
    while (b--) { a*=2; if (tree[b][a]<k) k-=tree[b][a++]; }
    return a; }
Number theory
Congruence ax \equiv b \pmod{n}
int congruence( int a, int b, int n ) { // finds ax = b(mod n)
    int d = gcd(a, n);
    if( b \% d != 0) return 1<<30; // no solution
    pii ans = egcd( a, n );
    int ret = ans.x * (b/d + OLL), mul = n/d;
    ret %= mul;
    if ( ret < 0 ) ret += mul;
    return ret;
}
Extended GCD
pii egcd( LL a, LL b ) { // returns x,y | ax + by = gcd(a,b)
    if( b == 0 ) return pii( 1, 0 );
    else {
        pii d = egcd( b, a % b );
        return pii( d.y, d.x - d.y * ( a / b ) );
    }
}
GCD
LL gcd( LL a, LL b ) { return !b ? a : gcd( b, a%b ); }
General EGCD
template < class T > inline T euclid(T a,T b,T &X,T &Y)
{
    if(a<0) { T d=euclid(-a,b,X,Y); X=-X; return d; }
    if(b<0) { T d=euclid(a,-b,X,Y); Y=-Y; return d; }
    if(b==0) { X=1; Y=0; return a; }
    else { T = (a/b)*Y; t=X; X=Y; Y=t-(a/b)*Y; t=X; Y=t
}
int X[110], Y[110]; LL v[110];
void gen_euclid(int n)
{
    int g = a[0];
    FOR(i,1,n) g = euclid(g, a[i], X[i], Y[i]);
    LL mult = 1;
    FORD(i,n,1) v[i] = (mult * Y[i]) % m, mult = (mult * X[i]) % m;
    v[0] = mult;
}
```

```
int phi (int n) {
    int ret = n;
    for (int i=2; i*i <= n; ++i) if (n\%i == 0) {
        while(n \% i == 0) n /= i;
        ret -= ret / i;
    if (n > 1) ret -= ret / n;
    return ret;
}
Power iterative - expmod
LL power( LL a, LL b, LL mod ) {
    LL x = 1, y = a;
    while(b) {
        if( b\&1 ) x *= y, x %= mod;
        y *= y, y \%= mod, b/=2;
    return x%mod;
}
Rabin-Miller
bool Miller(LL p, LL s, int a){
    if(p==a) return 1;
    LL mod=expmod(a,s,p);
    for(;s-p+1 && mod-1 && mod-p+1;s*=2) mod=mulmod(mod,mod,p);
    return mod==p-1 || s%2;
bool isprime(LL n) {
    if (n<2) return 0; if (n\%2==0) return n==2;
    LL s=n-1;
    while(s\%2==0) s/=2;
    return Miller(n,s,2) && Miller(n,s,7) && Miller(n,s,61);
} // for 341*10^12 primes <= 17
Sieve prime
const int MAX = 100000000;
int p[MAX/64 + 2], np = 0;
#define on(x) ( p[x/64] & (1 << (x%64)/2) ) )
#define turn(x) p[x/64] = (1 << (x\%64)/2)
void sieve() {
    for(int i=3;i*i<MAX; i+=2) { if(!on(i)) {
            int ii = i*i, i2 = i+i;
            for(int j=ii; j<MAX; j+=i2) turn(j); }}}</pre>
inline bool prime(int num){return num>1 && (num==2 || ((num&1) && !on(num) ));}
Sieve totient
FOR(i,1,M) f[i] = i;
FOR(n,2,M) if ( f[n] == n ) for (int k=n; k<=M; k+=n) f[k] *= n-1, f[k] /= n;
```

String algorithms (SuffixA, Aho-Corasick, KMP)

Suffix arrays

```
const int N = 100 * 1000 + 10;
char str[N]; bool bh[N], b2h[N];
int rank[N], pos[N], cnt[N], next[N], lcp[N];
bool smaller(int a, int b) { return str[a] < str[b];}</pre>
void suffix_array(int n) {
    REP(i,n)pos[i]=i, b2h[i]=false;
    sort(pos,pos+n,smaller);
    REP(i,n) bh[i]=!i||str[pos[i]] != str[pos[i-1]];
    for(int h=1;h<n;h*=2) {
        int buckets=0;
        for(int i=0,j; i<n; i=j) {</pre>
            j=i+1;
            while(j<n && !bh[j])j++;
            next[i]=j;
            buckets++;
        }
        if(buckets==n)break;
        for(int i=0;i<n;i=next[i]) {</pre>
            cnt[i] = 0;
            FOR(j, i, next[i]-1) rank[pos[j]]=i;
        }
        cnt[rank[n-h]]++;
        b2h[rank[n-h]]=true;
        for(int i=0;i<n;i=next[i]) {</pre>
            FOR(j, i, next[i]-1) {
                 int s = pos[j]-h;
                 if(s>=0){
                     rank[s] = rank[s] + cnt[rank[s]]++;
                     b2h[rank[s]]=true;
            FOR(j, i, next[i]-1) {
                 int s = pos[j]-h;
                 if (s>=0 \&\& b2h[rank[s]])
                     for(int k=rank[s]+1;!bh[k] && b2h[k]; k++) b2h[k]=false;
        REP(i,n) pos[rank[i]]=i, bh[i]|=b2h[i];
    REP(i,n) pos[rank[i]]=i;
}
void get_lcp(int n) {
    lcp[0]=0;
    int h=0;
    REP(i,n) if(rank[i]) {
        int j=pos[rank[i]-1];
        while(i+h < n && j+h < n && str[i+h] == str[j+h]) h++;
        lcp[rank[i]]=h;
        if(h)h--;
}
    }
//slower version of SA, also works with get_lcp
```

```
struct data {
    int nr[2],
    bool operator<(const data &v)const{return nr[0] < v.nr[0] || nr[0] == v.nr[0] &&nr[1] < v.nr[
    bool operator == (const data &v)const{return nr[1] == v.nr[1] &&nr[0] == v.nr[0];}
} L[MAXN];
int P[MAXLG+2][MAXN], pos[MAXN], rank[MAXN];
int suffix_array(char *A, int N)
{
    int step,cnt;
    REP(i,N) P[0][i] = A[i];
    for(step=1,cnt=1;cnt/2<N;cnt*=2,step++) {</pre>
        REP(i,N) L[i]=(data){P[step-1][i],(i+cnt<N)?P[step-1][i+cnt]:-1,i};
        sort(L,L+N);
        REP(i,N) P[step][L[i].p] = i && L[i] == L[i-1]? P[step][L[i-1].p]:i;
    REP(i,N) rank[L[i].p]=i;
    REP(i,N) pos[rank[i]]=i;
    return step-1;
}
Aho-Corasick
#define NC 26
#define NP 10005
#define M 100005
#define MM 500005
char a[M];
char b[NP][105];
int nb, cnt[NP], lenb[NP], alen;
int g[MM][NC], ng, f[MM], marked[MM];
int output[MM], pre[MM];
\#define\ init(x)\ \{REP(_i,NC)g[x][_i] = -1;\ f[x]=marked[x]=0;\ output[x]=pre[x]=-1;\ \}
void match() {
    ng = 0;
    init( 0 );
    // part 1 - building trie
    REP(i,nb) {
        cnt[i] = 0;
        int state = 0, j = 0;
        while (g[state][b[i][j]] != -1 \&\& j < lenb[i]) state = g[state][b[i][j]], j++;
        while(j < lenb[i]) {
            g[state][ b[i][j] ] = ++ng;
            state = ng;
            init( ng );
            ++j;
        }
        if( ng >= MM ) { cerr <<"i am dying"<<endl; while(1); // suicide }</pre>
        output[ state ] = i;
    }
    // part 2 - building failure function
    queue < int > q;
    REP(i,NC) if ( g[0][i] != -1 ) q.push( g[0][i] );
    while( !q.empty() ) {
```

```
int r = q.front(); q.pop();
        REP(i,NC) if(g[r][i] != -1) {
            int s = g[r][i];
            q.push(s);
            int state = f[r];
            while( g[state][i] == -1 && state ) state = f[state];
            f[s] = g[state][i] == -1 ? 0 : g[state][i];
        }
    }
    // final smash
    int state = 0;
    REP(i,alen) {
        while( g[state][a[i]] == -1 ) {
            state = f[state];
            if( !state ) break;
        }
        state = g[state][a[i]] == -1 ? 0 : g[state][a[i]];
        if( state && output[ state ] != -1 ) marked[ state ] ++;
    }
    // counting
    REP(i,ng+1) if( i && marked[i] ) {
        int s = i;
        while(s != 0) cnt[output[s]] += marked[i], s = f[s];
    }
}
KMP
int f[ len ];
f[0] = f[1] = 0;
FOR(i,2,len) {
    int j = f[i-1];
    while( true ) {
        if(s[j] == s[i-1]) { f[i] = j + 1; break;}
        }else if( !j ) { f[i] = 0; break;
        else j = f[j];
    }
}
i = j = 0;
while( true ) {
    if( i == len ) break;
    if(text[i] == s[j]) { i++, j++;}
        if( j == slen ) // match found
    else if(j > 0) j = f[j];
    else i++;
}
```

Graph algorithms (LCA, SCC, BPM-konig, Bellman-Ford, NetFlow, Min-Cost MaxFlow)

Tarjan's offline LCA

function TarjanOLCA(u)

```
MakeSet(u); u.ancestor := u;
    for each v in u.children do
        TarjanOLCA(v); Union(u,v); Find(u).ancestor := u;
    u.colour := black;
    for each v such that {u,v} in P and v.color==black do
        print "LCA", u, v, Find(v).ancestor
Tarjan's Strong Connected Components
procedure tarjan(v)
  index = count; v.lowlink = count++; S.push(v);color[v] = 1;
 for all (v, v2) in E do
     if (!color[v2])
        tarjan(v2); v.lowlink = min(v.lowlink, v2.lowlink);
     else if (color[v2]==1)
        v.lowlink = min(v.lowlink, v2.lowlink);
  if (v.lowlink == index)
    do { v2 = S.top(); S.pop(); print v2; color[v2]=2; } while (v2 != v);
for all v in V do if(!color[v]) tarjan(v);
Bipartite matching with Konig
#define M 1010
int grid[M][M], 1[M], r[M], seen[M], rows, cols;
bool dfs(int x)
    if( seen[x] ) return false;
    seen[x] = true;
    Rep(i,cols) if(grid[x][i]) if(r[i] == -1 \mid \mid dfs(r[i]))
    {
        r[i] = x, l[x] = i;
        return true;
    return false;
}
int bpm() {
    SET( 1, -1 );
    SET(r, -1);
    int ret = 0;
    Rep(i,rows) {
        SET( seen, 0 );
        if( dfs( i ) ) ret ++;
    }
    return ret;
}
bool lT[M], rT[M];
void konigdfs(int x)
{
    if( !1T[x] ) return; 1T[x] = 0;
    Rep(i,cols) if(grid[x][i] && i != l[x])
    {
        rT[i] = true;
        if( r[i] != -1) konigdfs(r[i]);
```

```
}
}
int konig()
    SET(1T, 1); SET(rT, 0);
    Rep(i,rows) if(l[i] == -1) konigdfs(i);
}
Bellman-Ford
VI e[M], c[M];
int n, d[M], p[M];
int inf = 1 << 29;
int bford( int s, int f ) {
    REP(i,n) d[i] = i == s ? 0 : inf, p[i] = -1;
    REP(\_,n-1) {
        bool done = 1;
        REP(i,n) REP(j,e[i].sz) {
            int u = i, v = e[i][j], uv = c[i][j];
            if( d[u] + uv < d[v] ) d[v] = d[u] + uv, p[v] = u, done = 0;
        if( done ) break;
    REP(i,n) REP(j,e[i].sz) {
        int u = i, v = e[i][j], uv = c[i][j];
        if( d[u] + uv < d[v] ) return -33;
    if( d[f] == inf ) return -33;
    return d[f];
}
Network flow - Slow
#define M 750
int nr, nc, o = 355, source = 740, sink = 741;
vector<int> edge[M];
int cap[M][M];
bool vis[M];
void init() {
    REP(i,M) edge[i].clear();
    SET( cap, 0 );
}
void add( int a, int b, int c, int d ) {
    edge[a].pb(b), edge[b].pb(a);
    cap[a][b] += c, cap[b][a] += d;
}
int dfs( int src, int snk, int fl ) {
    if( vis[src] ) return 0;
    if( snk == src ) return fl;
    vis[src] = 1;
    REP(i,edge[src].sz) {
        int v = edge[src][i];
```

```
int x = min(fl, cap[src][v]);
        if(x > 0) {
            x = dfs(v, snk, x);
            if(!x) continue;
            cap[src][v] -= x;
            cap[v][src] += x;
            return x;
        }
    }
    return 0;
}
int flow( int src, int snk ) {
    int ret = 0;
    while(1) {
        SET( vis, 0 );
        int delta = dfs( src, snk, 1<<30 );</pre>
        if( !delta ) break;
        ret += delta;
    }
    return ret;
}
Network flow - Dinic fast
const int maxN = 5005;
const int maxE = 70000;
const int inf = 1000000005;
int nnode, nedge, src, snk;
int Q[ maxN ], pro[ maxN ], fin[ maxN ], dist[ maxN ];
int flow[ maxE ], cap[ maxE ], to[ maxE ], next[ maxE ];
void init( int _nnode, int _src, int _snk ) {
    nnode = _nnode, nedge = 0, src = _src, snk = _snk;
    FOR(i,1,nnode) fin[i] = -1;
void add( int a, int b, int c1, int c2 ) {
    to[nedge]=b, cap[nedge]=c1, flow[nedge]=0, next[nedge]=fin[a], fin[a]=nedge++;
    to[nedge] = a, cap[nedge] = c2, flow[nedge] = 0, next[nedge] = fin[b], fin[b] = nedge + +;
}
bool bfs() {
    SET( dist, -1);
    dist[src] = 0;
    int st = 0, en = 0;
    Q[en++] = src;
    while(st < en ) {
        int u = Q[st++];
        for(int e = fin[u]; e \ge 0; e = next[e]) {
            int v = to[e];
            if( flow[e] < cap[e] && dist[v] == -1 ) {
                dist[v] = dist[u] + 1;
                Q[en++] = v;
            }
        }
    }
```

```
return dist[snk] != -1;
}
int dfs(int u, int fl) {
    if( u == snk ) return fl;
    for( int& e = pro[u]; e >= 0; e = next[e] ) {
        int v = to[e];
        if( flow[e] < cap[e] && dist[v] == dist[u]+1 ) {</pre>
            int x = dfs(v, min(cap[e] - flow[e], fl));
            if(x > 0) {
                flow[e] += x, flow[e^1] -= x;
                return x;
            }
        }
    }
    return 0;
}
LL dinic() {
    LL ret = 0;
    while( bfs() ) {
        FOR(i,1,nnode) pro[i] = fin[i];
        while(1) {
            int delta = dfs( src, inf );
            if( !delta ) break;
            ret += delta;
        }
    }
    return ret;
}
Min-Cost Max Flow
#define N 705
int n, nE;
int d[N], pre[N];
struct edge {
    int to, cost, cap;
    int back;
};
edge E[N*N];
vector< int > e[N];
int mincost( int s, int t, int lim ) {
    int flow = 0, ret = 0;
    while( flow < lim ) {</pre>
        SET( d, -1 ); SET( pre, -1 );
        d[s] = 0;
//
        cout <<"source "<< s <<" sink " << t << endl;</pre>
        // bellman ford
        int jump = n-1;
```

```
bool done = 0;
        while (!done && --jump >= 0) {
            done = 1;
            REP(i,n) if (d[i] != -1) REP(j,e[i].sz) {
                edge& x = E[e[i][j]];
                int v = x.to;
                if (x.cap > 0 && (d[v] == -1 || d[v] > d[i] + x.cost)) {
                     d[v] = d[i] + x.cost;
                     pre[v] = x.back;
                     done = 0;
//
                     cout << v << " " << d[v] << endl;
                }
            }
            if( done ) break;
//
        cout << d[t] << endl;
        if( d[t] == -1 ) break;
//
        cout <<"found one path "<<endl;</pre>
        // traverse back
        int x = t, cflow = 1 << 30;
        while (x != s) {
            edge& ed = E[pre[x]];
            cflow = min( cflow, E[ ed.back ].cap );
//
            cout << ed.to <<" to "<< x << endl;
            x = ed.to;
        }
        if( !cflow ) break;
        int take = min( lim - flow, cflow );
        ret += d[t] * take;
        flow += take;
//
        cout <<"taken flow "<< take <<" with cost "<< d[t] * take << endl << endl;</pre>
        x = t:
        while (x != s) {
            edge& back = E[ pre[x] ];
            edge& forw = E[ back.back ];
            back.cap += take;
            forw.cap -= take;
            x = back.to;
        }
   cout << "total flow " << flow << endl;</pre>
    if( flow < lim ) return -1;</pre>
    return ret;
}
// remember to add -cost in the opposite direction
void add( int u, int v, int uv, int vu, int fuv, int fvu ) {
    int a = nE, b = nE+1;
    nE += 2;
    E[ a ].to = v, E[ a ].cost = uv, E[ a ].cap = fuv, E[ a ].back = b;
    E[ b ].to = u, E[ b ].cost = vu, E[ b ].cap = fvu, E[ b ].back = a;
    e[u].pb(a), e[v].pb(b);
}
```

Misc (LIS)

Longest increasing subsequence

```
int n,total, nprob = 0;
vector< int > table;
while(scanf("%d", &total)==1){
   if( total == 0 ) break;
   table.clear();
   REP(kkk, total) {
    scanf("%d",&n);
   vector< int >::iterator i = lower_bound( table.begin(), table.end(), n );
   if( i== table.end() ) table.push_back( n );
   else *i <?= n;
}
printf("Set %d: %d\n",++nprob,table.size());
}</pre>
```

Geometry

Circle using three points Let A = (0,0) centers are $(C_y(B_x^2 + B_y^2) - B_y(C_x^2 + C_y^2))/D$ and $(B_x(C_x^2 + C_y^2) - C_x(B_x^2 + B_y^2))/D$ where $D = 2(B_xC_y - B_yC_x)$.

Appendices (ASCII table)

