IIUM cat-us-trophy

Contents

.vimrc

```
set ai ts=4 sw=4 st=4 noet nu nohls
syntax enable
filetype plugin indent on
map <F6> :w<CR>:!g++ % -g && (ulimit -c unlimited; ./a.out < ~/input.txt) <CR>
map <F5> <F6>
colo pablo
map <F12> :!gdb ./a.out -c core <CR>
```

template.cpp

```
#include<cstdio>
#include<sstream>
#include<cstdlib>
#include<cctype>
#include<cmath>
#include<algorithm>
#include<set>
#include<queue>
#include<stack>
#include<list>
#include<iostream>
#include<string>
#include<vector>
#include<cstring>
#include<map>
#include<cassert>
#include<climits>
using namespace std;

#define REP(i,n) for(int i=0, _e(n); i<_e; i++)
#define FOR(i,a,b) for(int i(a), _e(b); i<=_e; i++)
#define FORD(i,a,b) for(int i(a), _e(b); i>=_e; i--)
#define FORIT(i, m) for (__typeof((m).begin()) i=(m).begin(); i!=(m).end(); ++i)
#define SET(t,v) memset((t), (v), sizeof(t))
#define ALL(x) x.begin(), x.end()
#define UNIQUE(c) (c).resize( unique( ALL(c) ) - (c).begin() )

#define sz size()
#define pb push_back
#define VI vector<int>
```

```
#define VS vector<string>

typedef long long LL;
typedef long double LD;
typedef pair<int,int> pii;

#define D(x) if(1) cout << __LINE__ <<" "<< #x " = " << (x) << endl;
#define D2(x,y) if(1) cout << __LINE__ <<" "<< #x " = " << (x) \
     <<", " << #y " = " << (y) << endl;
```

Combinatorics

Mathematical Sums

$\sum_{k=0}^{n} k = n(n+1)/2$ $\quad\sum_{k=a}^{b} k = (a+b)(b-a+1)/2$
$\sum_{k=0}^{n} k^2 = n(n+1)(2n+1)/6$ $\quad\sum_{k=0}^{n} k^3 = n^2(n+1)^2/4$
$\sum_{k=0}^{n} k^4 = (6n^5+15n^4+10n^3-n)/30$ $\quad\sum_{k=0}^{n} k^5 = (2n^6+6n^5+5n^4-n^2)/12$
$\sum_{k=0}^{n} x^k = (x^{n+1}-1)/(x-1)$ $\quad\sum_{k=0}^{n} kx^k = (x-(n+1)x^{n+1}+nx^{n+2})/(x-1)^2$

Binomial coefficients

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	1												
1	1	1											
2	1	2	1										
3	1	3	3	1									
4	1	4	6	4	1								
5	1	5	10	10	5	1							
6	1	6	15	20	15	6	1						
7	1	7	21	35	35	21	7	1					
8	1	8	28	56	70	56	28	8	1				
9	1	9	36	84	126	126	84	36	9	1			
10	1	10	45	120	210	252	210	120	45	10	1		
11	1	11	55	165	330	462	462	330	165	55	11	1	
12	1	12	66	220	495	792	924	792	495	220	66	12	1
	0	1	2	3	4	5	6	7	8	9	10	11	12

$\binom{n}{k} = \frac{n!}{(n-k)!k!}$
$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$
$\binom{n}{k} = \frac{n}{n-k}\binom{n-1}{k}$
$\binom{n}{k} = \frac{n-k+1}{k}\binom{n}{k-1}$
$\binom{n+1}{k} = \frac{n+1}{n-k+1}\binom{n}{k}$
$\binom{n}{k+1} = \frac{n-k}{k+1}\binom{n}{k}$

$\sum_{k=1}^{n} k\binom{n}{k} = n2^{n-1}$
$\sum_{k=1}^{n} k^2\binom{n}{k} = (n+n^2)2^{n-2}$

$\binom{m+n}{r} = \sum_{k=0}^{r}\binom{m}{k}\binom{n}{r-k}$
$\binom{n}{k} = \prod_{i=1}^{k}\frac{n-k+i}{i}$

Catalan numbers $C_n = \frac{1}{n+1}\binom{2n}{n}$. $\quad C_0 = 1$, $C_n = \sum_{i=0}^{n-1} C_i C_{n-1-i}$. $\quad C_{n+1} = C_n\frac{4n+2}{n+2}$.
$C_0, C_1, \ldots = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900, \ldots$
C_n is the number of: properly nested sequences of n pairs of parentheses; rooted ordered binary trees with $n+1$ leaves; triangulations of a convex $(n+2)$-gon.

Derangements . Number of permutations of $n = 0, 1, 2, \ldots$ elements without fixed points is $1, 0, 1, 2, 9, 44, 265, 1854, 14833, \ldots$ Recurrence: $D_n = (n-1)(D_{n-1} + D_{n-2}) = nD_{n-1} + (-1)^n$. Corollary: number of permutations with exactly k fixed points is $\binom{n}{k}D_{n-k}$.

Stirling numbers of 1^{st} kind . $s_{n,k}$ is $(-1)^{n-k}$ times the number of permutations of n elements with exactly k permutation cycles. $\left[{n \atop k}\right] = |s_{n,k}| = |s_{n-1,k-1}| + (n-1)|s_{n-1,k}|$ $\quad s(0,0) = 1$ and $s(n,0) = s(0,n) = 0$.

Stirling numbers of 2^{nd} kind . $S_{n,k}$ is the number of ways to partition a set of n elements into exactly k non-empty subsets. $\left\{{n \atop k}\right\} = S_{n,k} = S_{n-1,k-1} + kS_{n-1,k}$. $S_{n,1} = S_{n,n} = 1$.

Bell numbers . B_n is the number of partitions of n elements. $B_0, \ldots = 1, 1, 2, 5, 15, 52, 203, 877, \ldots$
$B_{n+1} = \sum_{k=0}^{n}\binom{n}{k}B_k = \sum_{k=1}^{n+1} S_{n,k}$. Bell triangle: $B_r = a_{r,1} = a_{r-1,r-1}$, $a_{r,c} = a_{r-1,c-1} + a_{r,c-1}$.

Eulerian numbers . $E(n,k) = \left\langle{n \atop k}\right\rangle$ is the number of permutations with exactly k descents ($i : \pi_i < \pi_{i+1}$) / ascents ($\pi_i > \pi_{i+1}$) / excedances ($\pi_i > i$) / $k+1$ weak excedances ($\pi_i \geq i$).

Formula: $E(n,m) = (m+1)E(n-1,m) + (n-m)E(n-1,m-1)$. $E(n,0) = E(n,n-1) = 1$. $E(n,m) = \sum_{k=0}^{m}(-1)^k\binom{n+1}{k}(m+1-k)^n$.

Double factorial . Permutations of the multiset $\{1,1,2,3,\dots n,n\}$ such that for each k, all the numbers between two occurrences of k in the permutation are greater than k. $(2n-1)!! = \prod_{k=1}^{n}(2k-1)$.

Eulerian numbers of 2^{nd} kind . Related to Double factorial, number of all such permutations that have exactly m ascents. $\left\langle\!\left\langle {n \atop m} \right\rangle\!\right\rangle = (2n-m-1)\left\langle\!\left\langle {n-1 \atop m-1} \right\rangle\!\right\rangle + (m+1)\left\langle\!\left\langle {n-1 \atop m} \right\rangle\!\right\rangle$. $\left\langle\!\left\langle {n \atop 0} \right\rangle\!\right\rangle = 1$

Multinomial theorem . $(a_1+\dots+a_k)^n = \sum\binom{n}{n_1,\dots,n_k}a_1^{n_1}\dots a_k^{n_k}$, where $n_i \geq 0$ and $\sum n_i = n$. $\binom{n}{n_1,\dots,n_k} = M(n_1,\dots,n_k) = \frac{n!}{n_1!\dots n_k!}$. $M(a,\dots,b,c,\dots) = M(a+\dots+b,c,\dots)M(a,\dots,b)$

Data structures

BIT - Binary indexed trees

BIT 2D

```
int bit[M][M], n;
int sum( int x, int y ){
    int ret = 0;
    while( x > 0 ){
        int yy = y; while( yy > 0 ) ret += bit[x][yy], yy -= yy & -yy;
        x -= (x & -x);
    }
    return ret ;
}
void update(int x , int y , int val){
    int y1;
    while (x <= n){
        y1 = y;
        while (y1 <= n){ bit[x][y1] += val; y1 += (y1 & -y1); }
        x += (x & -x);
    }
}
```

RMQ DP

```
int make_dp(int n) { // N log N
    REP(i,n) H[i][0]=i;
    for(int l=0,k; (k=1<<l) < n; l++) for(int i=0;i+k<n;i++)
        H[i][l+1] = A[H[i][l]] > A[H[i+k][l]] ? H[i+k][l] : H[i][l];
} // query log N almost O(1)
int query_dp(int a, int b) {
    for(int l=0;;l++)if(a+(1<<l+1) > b) {
        int o2 = H[b-(1<<l)+1][l];
        return A[H[a][l]]<A[o2] ? H[a][l]:o2;
}   }
```

RMQ segment tree

```
const int M = 100005;
int n, in[M], f[M], st[M], en[M];
struct data { int l, r, ans, next_l, next_r; };
data d[4*M]; // which is the range? :S
int nd;
int build( int l, int r, int id ) {
```

```
    d[ id ].l = l, d[ id ].r = r;
    if( l == r ) d[ id ].ans = f[l];
    else {
        int bar = ( r-l ) / 2 + l;
        d[id].next_l = ++nd, d[id].next_r = ++nd;
        int left = build( l, bar, d[id].next_l );
        int right = build( bar+1, r, d[id].next_r );
        d[id].ans = max( left, right );
    }
    return d[ id ].ans;
}
int query( int l, int r, int id = 0 ) {
    if( l > r ) return 0;
    if( d[id].l == l && d[id].r == r ) return d[id].ans;
    else {
        int bar = (d[id].r-d[id].l) / 2 + d[id].l;
        int left = 0, right = 0;
        if( l <= bar ) {
            if( r <= bar ) left = query( l, r, d[id].next_l );
            else {
                left = query( l, bar, d[id].next_l );
                right = query( bar+1, r, d[id].next_r );
            }
        }else right = query( l, r, d[id].next_r );
        return max( left, right );
    }
}
```

String algorithms (SuffixA, Aho-Corasick, KMP)

Suffix arrays

```
const int N = 100 * 1000 + 10;
char str[N]; bool bh[N], b2h[N];
int rank[N], pos[N], cnt[N], next[N], lcp[N];
bool smaller(int a, int b) { return str[a]<str[b];}
void suffix_array(int n) {
    REP(i,n)pos[i]=i, b2h[i]=false;
    sort(pos,pos+n,smaller);
    REP(i,n) bh[i]=!i||str[pos[i]] != str[pos[i-1]];
    for(int h=1;h<n;h*=2) {
        int buckets=0;
        for(int i=0,j; i<n; i=j) {
            j=i+1;
            while(j<n && !bh[j])j++;
            next[i]=j;
            buckets++;
        }
        if(buckets==n)break;
        for(int i=0;i<n;i=next[i]) {
            cnt[i] = 0;
            FOR(j, i, next[i]-1) rank[pos[j]]=i;
        }
```

```
        cnt[rank[n-h]]++;
        b2h[rank[n-h]]=true;
        for(int i=0;i<n;i=next[i]) {
            FOR(j, i, next[i]-1) {
                int s = pos[j]-h;
                if(s>=0){
                    rank[s] = rank[s] + cnt[rank[s]]++;
                    b2h[rank[s]]=true;
            }   }
            FOR(j, i, next[i]-1) {
                int s = pos[j]-h;
                if(s>=0 && b2h[rank[s]])
                    for(int k=rank[s]+1;!bh[k] && b2h[k]; k++) b2h[k]=false;
        }   }
        REP(i,n) pos[rank[i]]=i, bh[i]|=b2h[i];
    }
    REP(i,n) pos[rank[i]]=i;
}
void get_lcp(int n) {
    lcp[0]=0;
    int h=0;
    REP(i,n) if(rank[i]) {
        int j=pos[rank[i]-1];
        while(i+h<n && j+h<n && str[i+h] == str[j+h]) h++;
        lcp[rank[i]]=h;
        if(h)h--;
}   }
//slower version of SA, also works with get_lcp
struct data {
    int nr[2],  p;
    bool operator<(const data &v)const{return nr[0]<v.nr[0] || nr[0]==v.nr[0]&&nr[1]<v.nr[
    bool operator==(const data &v)const{return nr[1]==v.nr[1]&&nr[0]==v.nr[0];}
} L[MAXN];
int P[MAXLG+2][MAXN], pos[MAXN], rank[MAXN];
int suffix_array(char *A, int N)
{
    int step,cnt;
    REP(i,N) P[0][i] = A[i];
    for(step=1,cnt=1;cnt/2<N;cnt*=2,step++) {
        REP(i,N) L[i]=(data){P[step-1][i],(i+cnt<N)?P[step-1][i+cnt]:-1,i};
        sort(L,L+N);
        REP(i,N) P[step][L[i].p] = i && L[i]==L[i-1]? P[step][L[i-1].p]:i;
    }
    REP(i,N) rank[L[i].p]=i;
    REP(i,N) pos[rank[i]]=i;
    return step-1;
}
```

Aho-Corasick

```
#define NC 26
#define NP 10005
#define M 100005
```

```
#define MM 500005
char a[M];
char b[NP][105];
int nb, cnt[NP], lenb[NP], alen;
int g[MM][NC], ng, f[MM], marked[MM];
int output[MM], pre[MM];
#define init(x) {REP(_i,NC)g[x][_i] = -1; f[x]=marked[x]=0; output[x]=pre[x]=-1; }
void match() {
    ng = 0;
    init( 0 );
    // part 1 - building trie
    REP(i,nb) {
        cnt[i] = 0;
        int state = 0, j = 0;
        while(g[state][b[i][j]] != -1 && j < lenb[i]) state = g[state][b[i][j]], j++;
        while( j < lenb[i] ) {
            g[state][ b[i][j] ] = ++ng;
            state = ng;
            init( ng );
            ++j;
        }
        if( ng >= MM ) { cerr <<"i am dying"<<endl; while(1); // suicide }
        output[ state ] = i;
    }
    // part 2 - building failure function
    queue< int > q;
    REP(i,NC) if( g[0][i] != -1 ) q.push( g[0][i] );
    while( !q.empty() ) {
        int r = q.front(); q.pop();
        REP(i,NC) if( g[r][i] != -1 ) {
            int s = g[r][i];
            q.push( s );
            int state = f[r];
            while( g[state][i] == -1 && state ) state = f[state];
            f[s] = g[state][i] == -1 ? 0 : g[state][i];
        }
    }
    // final smash
    int state = 0;
    REP(i,alen) {
        while( g[state][a[i]] == -1 ) {
            state = f[state];
            if( !state ) break;
        }
        state = g[state][a[i]] == -1 ? 0 : g[state][a[i]];
        if( state && output[ state ] != -1 ) marked[ state ] ++;
    }
    // counting
    REP(i,ng+1) if( i && marked[i] ) {
        int s = i;
        while( s != 0 ) cnt[ output[s] ] += marked[i], s = f[s];
    }
```

```
}
```

KMP

```
int f[ len ];
f[0] = f[1] = 0;
FOR(i,2,len) {
    int j = f[i-1];
    while( true ) {
        if( s[j] == s[i-1] ) { f[i] = j + 1; break;
        }else if( !j ) { f[i] = 0; break;
        }else j = f[j];
    }
}
i = j = 0;
while( true ) {
    if( i == len ) break;
    if( text[i] == s[j] ) { i++, j++;
        if( j == slen ) // match found
    }else if( j > 0 ) j = f[j];
    else i++;
}
```

Graph algorithms (LCA, SCC, NetFlow, MinCost, BPM)

Tarjan's offline LCA

```
function TarjanOLCA(u)
    MakeSet(u); u.ancestor := u;
    for each v in u.children do
        TarjanOLCA(v); Union(u,v); Find(u).ancestor := u;
    u.colour := black;
    for each v such that {u,v} in P and v.color==black do
        print "LCA", u, v, Find(v).ancestor
```

Tarjan's Strong Connected Components

```
procedure tarjan(v)
  index = count; v.lowlink = count++; S.push(v);color[v] = 1;
  for all (v, v2) in E do
     if (!color[v2])
        tarjan(v2); v.lowlink = min(v.lowlink, v2.lowlink);
     else if (color[v2]==1)
        v.lowlink = min(v.lowlink, v2.lowlink);
  if (v.lowlink == index)
    do { v2 = S.top(); S.pop(); print v2; color[v2]=2; } while (v2 != v);
for all v in V do if(!color[v]) tarjan(v);
```

Bipartite matching with Konig

```
#define M 1010
int grid[M][M], l[M], r[M], seen[M], rows, cols;
bool dfs(int x)
{
```

```
    if( seen[x] ) return false;
    seen[x] = true;
    Rep(i,cols) if( grid[x][i] ) if( r[i] == -1 || dfs( r[i] ) )
    {
        r[i] = x, l[x] = i;
        return true;
    }
    return false;
}
int bpm() {
    SET( l, -1 );
    SET( r, -1 );
    int ret = 0;
    Rep(i,rows) {
        SET( seen, 0 );
        if( dfs( i ) ) ret ++;
    }
    return ret;
}
bool lT[M], rT[M];
void konigdfs(int x)
{
    if( !lT[x] ) return; lT[x] = 0;
    Rep(i,cols) if(grid[x][i] && i != l[x])
    {
        rT[i] = true;
        if( r[i] != -1) konigdfs(r[i]);
    }
}
int konig()
{
    SET(lT, 1); SET(rT, 0);
    Rep(i,rows) if(l[i] == -1) konigdfs(i);
}
```

Network flow - Slow

```
#define M 750
int nr, nc, o = 355, source = 740, sink = 741;
vector<int> edge[M];
int cap[M][M];
bool vis[M];
void init() {
    REP(i,M) edge[i].clear();
    SET( cap, 0 );
}
void add( int a, int b, int c, int d ) {
    edge[a].pb(b), edge[b].pb(a);
    cap[a][b] += c, cap[b][a] += d;
}
int dfs( int src, int snk, int fl ) {
    if( vis[src] ) return 0;
    if( snk == src ) return fl;
```

```
    vis[src] = 1;

    REP(i,edge[src].sz) {
        int v = edge[src][i];
        int x = min( fl, cap[src][v] );
        if( x > 0 ) {
            x = dfs( v, snk, x );
            if( !x ) continue;
            cap[src][v] -= x;
            cap[v][src] += x;
            return x;
        }
    }
    return 0;
}
int flow( int src, int snk ) {
    int ret = 0;
    while( 1 ) {
        SET( vis, 0 );
        int delta = dfs( src, snk, 1<<30 );
        if( !delta ) break;
        ret += delta;
    }
    return ret;
}
```

Network flow - Dinic fast

```
const int maxN = 5005;
const int maxE = 70000;
const int inf = 1000000005;
int nnode, nedge, src, snk;
int Q[ maxN ], pro[ maxN ], fin[ maxN ], dist[ maxN ];
int flow[ maxE ], cap[ maxE ], to[ maxE ], next[ maxE ];
void init( int _nnode, int _src, int _snk ) {
    nnode = _nnode, nedge = 0, src = _src, snk = _snk;
    FOR(i,1,nnode) fin[i] = -1;
}
void add( int a, int b, int c1, int c2 ) {
    to[nedge]=b, cap[nedge]=c1, flow[nedge]=0, next[nedge]=fin[a], fin[a]=nedge++;
    to[nedge]=a, cap[nedge]=c2, flow[nedge]=0, next[nedge]=fin[b], fin[b]=nedge++;
}
bool bfs() {
    SET( dist, -1 );
    dist[src] = 0;
    int st = 0, en = 0;
    Q[en++] = src;
    while( st < en ) {
        int u = Q[ st++ ];
        for(int e = fin[u]; e >= 0; e = next[e] ) {
            int v = to[e];
            if( flow[e] < cap[e] && dist[v] == -1 ) {
                dist[v] = dist[u] + 1;
```

```
                Q[en++] = v;
            }
        }
    }
    return dist[snk] != -1;
}
int dfs(int u, int fl) {
    if( u == snk ) return fl;
    for( int& e = pro[u]; e >= 0; e = next[e] ) {
        int v = to[e];
        if( flow[e] < cap[e] && dist[v] == dist[u]+1 ) {
            int x = dfs( v, min( cap[e] - flow[e] , fl ) );
            if( x > 0 ) {
                flow[ e ] += x, flow[ e^1 ] -= x;
                return x;
            }
        }
    }
    return 0;
}
LL dinic() {
    LL ret = 0;
    while( bfs() ) {
        FOR(i,1,nnode) pro[i] = fin[i];
        while( 1 ) {
            int delta = dfs( src, inf );
            if( !delta ) break;
            ret += delta;
        }
    }
    return ret;
}
```

Min-Cost Max Flow

```
#define N 705
int n, nE;
int d[N], pre[N];

struct edge {
    int to, cost, cap;
    int back;
};

edge E[N*N];
vector< int > e[N];

int mincost( int s, int t, int lim ) {

    int flow = 0, ret = 0;
    while( flow < lim ) {

        SET( d, -1 ); SET( pre, -1 );
```

```
        d[s] = 0;
//      cout <<"source "<< s <<" sink " << t << endl;
        // bellman ford
        int jump = n-1;
        bool done = 0;
        while( !done && --jump >= 0) {
            done = 1;
            REP(i,n) if( d[i] != -1 ) REP(j,e[i].sz) {
                edge& x = E[ e[i][j] ];
                int v = x.to;
                if( x.cap > 0 && ( d[v] == -1 || d[v] > d[i] + x.cost )) {
                    d[v] = d[i] + x.cost;
                    pre[v] = x.back;
                    done = 0;
//                  cout<<v<<" "<<d[v]<<endl;
                }
            }
            if( done ) break;
        }
//      cout << d[t] << endl;
        if( d[t] == -1 ) break;
//      cout <<"found one path "<<endl;
        // traverse back
        int x = t, cflow = 1<<30;
        while( x != s ) {
            edge& ed = E[ pre[x] ];
            cflow = min( cflow, E[ ed.back ].cap );
//          cout << ed.to <<" to "<< x << endl;
            x = ed.to;
        }
        if( !cflow ) break;
        int take = min( lim - flow, cflow );
        ret += d[t] * take;
        flow += take;
//      cout <<"taken flow "<< take <<" with cost "<< d[t] * take << endl << endl;
        x = t;
        while( x != s ) {
            edge& back = E[ pre[x] ];
            edge& forw = E[ back.back ];
            back.cap += take;
            forw.cap -= take;
            x = back.to;
        }
    }
//  cout << "total flow " << flow << endl;
    if( flow < lim ) return -1;
    return ret;
}
// remember to add -cost in the opposite direction
void add( int u, int v, int uv, int vu, int fuv, int fvu ) {
    int a = nE, b = nE+1;
    nE += 2;
```

```
    E[ a ].to = v, E[ a ].cost = uv, E[ a ].cap = fuv, E[ a ].back = b;
    E[ b ].to = u, E[ b ].cost = vu, E[ b ].cap = fvu, E[ b ].back = a;
    e[ u ].pb( a ), e[ v ].pb( b );
}
```

Geometry

Circle using three points Let $A = (0,0)$ centers are $(C_y(B_x^2 + B_y^2) - B_y(C_x^2 + C_y^2))/D$ and $(B_x(C_x^2 + C_y^2) - C_x(B_x^2 + B_y^2))/D$ where $D = 2(B_xC_y - B_yC_x)$.

Appendices (ASCII table)

32 0010 0000	33 0010 0001	34 0010 0010	35 0010 0011	36 0010 0100	37 0010 0101	38 0010 0110	39 0010 0111	40 0010 1000	41 0010 1001	42 0010 1010	43 0010 1011	44 0010 1100	45 0010 1101	46 0010 1110	47 0010 1111
SP	!	"	#	$	%	&	'	(	)	*	+	,	−	.	/
48 0011 0000	49 0011 0001	50 0011 0010	51 0011 0011	52 0011 0100	53 0011 0101	54 0011 0110	55 0011 0111	56 0011 1000	57 0011 1001	58 0011 1010	59 0011 1011	60 0011 1100	61 0011 1101	62 0011 1110	63 0011 1111
0	1	2	3	4	5	6	7	8	9	:	;	<	=	>	?
64 0100 0000	65 0100 0001	66 0100 0010	67 0100 0011	68 0100 0100	69 0100 0101	70 0100 0110	71 0100 0111	72 0100 1000	73 0100 1001	74 0100 1010	75 0100 1011	76 0100 1100	77 0100 1101	78 0100 1110	79 0100 1111
@	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
80 0101 0000	81 0101 0001	82 0101 0010	83 0101 0011	84 0101 0100	85 0101 0101	86 0101 0110	87 0101 0111	88 0101 1000	89 0101 1001	90 0101 1010	91 0101 1011	92 0101 1100	93 0101 1101	94 0101 1110	95 0101 1111
P	Q	R	S	T	U	V	W	X	Y	Z	[	\	]	^	_
96 0110 0000	97 0110 0001	98 0110 0010	99 0110 0011	100 0110 0100	101 0110 0101	102 0110 0110	103 0110 0111	104 0110 1000	105 0110 1001	106 0110 1010	107 0110 1011	108 0110 1100	109 0110 1101	110 0110 1110	111 0110 1111
`	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o
112 0111 0000	113 0111 0001	114 0111 0010	115 0111 0011	116 0111 0100	117 0111 0101	118 0111 0110	119 0111 0111	120 0111 1000	121 0111 1001	122 0111 1010	123 0111 1011	124 0111 1100	125 0111 1101	126 0111 1110	127 0111 1111
p	q	r	s	t	u	v	w	x	y	z	{	|	}	~	DEL