# IIUM cat-us-trophy

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.vimrc		
set ai t	ts=4 sw=4 st=4 noet nu nohls	
syntax e	enable	
filetype	e plugin indent on	
map <f6></f6>	$>$ :w <cr>:!g++ % -g &amp;&amp; (ulimit -c unlimited; ./a.out &lt; <math>\sim</math>/input.txt)</cr>	<cr></cr>
map <f5></f5>	> <f6></f6>	
colo pab	blo	
map <f12< td=""><td>2&gt; :!gdb ./a.out -c core <cr></cr></td><td></td></f12<>	2> :!gdb ./a.out -c core <cr></cr>	
templa	ate.cpp	
#include	e <cstdio></cstdio>	
	e <sstream></sstream>	
	e <cstdlib></cstdlib>	
	e <cctype></cctype>	
#include	••	
	e <algorithm></algorithm>	
#include	•	
#include		
#include	-	
#include	e <list></list>	
#include	e <iostream></iostream>	
	e <string></string>	
	e <vector></vector>	
#include	e <cstring></cstring>	
#include	9	
#include	e <cassert></cassert>	
#include	e <climits></climits>	
using na	amespace std;	
#define	REP(i,n) for(int i=0, _e(n); i<_e; i++)	
	FOR(i,a,b) for(int i(a), _e(b); i<=_e; i++)	
	FORD(i,a,b) for(int i(a), _e(b); i>=_e; i)	
	<pre>FORIT(i, m) for (typeof((m).begin()) i=(m).begin(); i!=(m).end(</pre>	(); ++i)
	SET(t,v) memset((t), (v), sizeof(t))	•
	ALL(x) x.begin(), x.end()	
	UNIQUE(c) (c).resize( unique( ALL(c) ) - (c).begin() )	
#define	sz size()	
	pb push_back	
	VI vector <int></int>	

#define VS vector<string>

### **Combinatorics**

#### **Mathematical Sums**

$$\begin{array}{ll} \sum_{k=0}^n k = n(n+1)/2 & \sum_{k=0}^b k = (a+b)(b-a+1)/2 \\ \sum_{k=0}^n k^2 = n(n+1)(2n+1)/6 & \sum_{k=0}^n k^3 = n^2(n+1)^2/4 \\ \sum_{k=0}^n k^4 = (6n^5+15n^4+10n^3-n)/30 & \sum_{k=0}^n k^5 = (2n^6+6n^5+5n^4-n^2)/12 \\ \sum_{k=0}^n x^k = (x^{n+1}-1)/(x-1) & \sum_{k=0}^n kx^k = (x-(n+1)x^{n+1}+nx^{n+2})/(x-1)^2 \end{array}$$

#### Binomial coefficients

	0	1	2	3	4	5	6	7	8	9	10	11	12	$\binom{n}{k} = \frac{n!}{(n-k)!k!}$
0	1													$\binom{n}{-}$ $\binom{n-1}{-}$ $\perp$ $\binom{n-1}{-}$
1	1	1												$\binom{k}{k} = \binom{k}{k} + \binom{k-1}{k-1}$ $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$
2	1	2	1											$\binom{n}{k} = \frac{n}{n-k} \binom{n-1}{k}$ $\binom{n}{k} = \frac{n-k+1}{k} \binom{n}{k-1}$
3	1	3	3	1										$\binom{k}{k} \equiv \frac{1}{k} \binom{k-1}{n}$
4	1	4	6	4	1									$\binom{k}{k} - \frac{n-k+1}{n-k+1} \binom{k}{k}$
5	1	5	10	10	5	1								$\binom{n}{k+1} = \frac{n-k}{k+1} \binom{n}{k}$
6	1	6	15	20	15	6	1							
7	1	7	21	35	35	21	7	1						$\sum_{n=1}^{\infty} (n)$ and 1
8	1	8	28	56	70	56	28	8	1					$\sum_{k=1}^{n} k \binom{n}{k} = n2^{n-1}$ $\sum_{k=1}^{n} k^{2} \binom{n}{k} = (n+n^{2})2^{n-2}$
9	1	9	36	84	126	126	84	36	9	1				$\sum_{k=1}^{n} k^{2} \binom{n}{k} = (n+n^{2})2^{n-2}$
10	1	10	45	120	210	252	210	120	45	10	1			
11	1	11	55	165	330	462	462	330	165	55	11	1		$(m+n) \subseteq \sum^r (m) (n)$
12	1	12	66	220	495	792	924	792	495	220	66	12	1	$\binom{r}{r} = \sum_{k=0}^{\infty} \binom{k}{r-k}$
	0	1	2	3	4	5	6	7	8	9	10	11	12	$\binom{n}{k} = \prod_{i=1}^k rac{n-k+i}{i}$

Catalan numbers  $C_n = \frac{1}{n+1} {2n \choose n}$ .  $C_0 = 1$ ,  $C_n = \sum_{i=0}^{n-1} C_i C_{n-1-i}$ .  $C_{n+1} = C_n \frac{4n+2}{n+2}$ .  $C_0, C_1, \ldots = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900, \ldots$ 

 $C_n$  is the number of: properly nested sequences of n pairs of parentheses; rooted ordered binary trees with n+1 leaves; triangulations of a convex (n+2)-gon.

**Derangements**. Number of permutations of  $n=0,1,2,\ldots$  elements without fixed points is  $1,0,1,2,9,44,265,1854,14833,\ldots$  Recurrence:  $D_n=(n-1)(D_{n-1}+D_{n-2})=nD_{n-1}+(-1)^n$ . Corollary: number of permutations with exactly k fixed points is  $\binom{n}{k}D_{n-k}$ .

Stirling numbers of  $1^{st}$  kind .  $s_{n,k}$  is  $(-1)^{n-k}$  times the number of permutations of n elements with exactly k permutation cycles.  $\binom{n}{k} = |s_{n,k}| = |s_{n-1,k-1}| + (n-1)|s_{n-1,k}|$  s(0,0) = 1 and s(n,0) = s(0,n) = 0.

Stirling numbers of  $2^{nd}$  kind.  $S_{n,k}$  is the number of ways to partition a set of n elements into exactly k non-empty subsets.  $\binom{n}{k} = S_{n,k} = S_{n-1,k-1} + kS_{n-1,k}$ .  $S_{n,1} = S_{n,n} = 1$ .

**Bell numbers**.  $B_n$  is the number of partitions of n elements.  $B_0, \ldots = 1, 1, 2, 5, 15, 52, 203, 877, \ldots$   $B_{n+1} = \sum_{k=0}^{n} \binom{n}{k} B_k = \sum_{k=1}^{n+1} S_{n,k}$ . Bell triangle:  $B_r = a_{r,1} = a_{r-1,r-1}, \ a_{r,c} = a_{r-1,c-1} + a_{r,c-1}$ .

**Eulerian numbers** .  $E(n,k) = {n \choose k}$  is the number of permutations with exactly k descents  $(i : \pi_i < \pi_{i+1}) / \text{ ascents } (\pi_i > \pi_{i+1}) / \text{ excedances } (\pi_i > i) / k + 1 \text{ weak excedances } (\pi_i \ge i).$ 

```
Formula: E(n,m) = (m+1)E(n-1,m) + (n-m)E(n-1,m-1). E(n,0) = E(n,n-1) = 1.
E(n,m) = \sum_{k=0}^{m} (-1)^k \binom{n+1}{k} (m+1-k)^n.
Double factorial. Permutations of the multiset \{1, 1, 2, 3, \dots n, n\} such that for each k, all the num-
bers between two occurrences of k in the permutation are greater than k. (2n-1)!! = \prod_{k=1}^{n} (2k-1).
Eulerian numbers of 2^{nd} kind. Related to Double factorial, number of all such permutations
that have exactly m ascents. \left\langle \left\langle {n \atop m} \right\rangle \right\rangle = (2n-m-1)\left\langle \left\langle {n-1 \atop m-1} \right\rangle \right\rangle + (m+1)\left\langle \left\langle {n-1 \atop m} \right\rangle \right\rangle. \left\langle \left\langle {n \atop 0} \right\rangle \right\rangle = 1
Multinomial theorem (a_1 + \dots + a_k)^n = \sum_{n_1,\dots,n_k} \binom{n}{n_1,\dots,n_k} a_1^{n_1} \dots a_k^{n_k}, where n_i \ge 0 and \sum_{n_i = n} n_i = n. \binom{n}{n_1,\dots,n_k} = M(n_1,\dots,n_k) = \frac{n!}{n_1!\dots n_k!}. M(a,\dots,b,c,\dots) = M(a+\dots+b,c,\dots)M(a,\dots,b)
Data structures
BIT - Binary indexed trees
int bit[M],n;
void update(int x, int v) { while( x \le n ) { bit[x] += v; x += x \& -x; } }
int sum(int x) { int ret=0; while(x>0){ ret += bit[x]; x == x \& -x; } return ret; }
BIT 2D
int bit[M][M], n;
int sum( int x, int y ){
     int ret = 0;
     while (x > 0)
           int yy = y; while(yy > 0) ret += bit[x][yy], yy -= yy & -yy;
           x = (x \& -x);
     return ret ;
void update(int x , int y , int val){
      int y1;
     while (x \le n){
           y1 = y;
           while (y1 \le n){ bit[x][y1] += val; y1 += (y1 & -y1); }
           x += (x \& -x);
     }
}
RMQ DP
int make_dp(int n) { // N log N
     REP(i,n) H[i][0]=i;
     for(int l=0,k; (k=1<<1) < n; l++) for(int i=0;i+k< n;i++)
           H[i][1+1] = A[H[i][1]] > A[H[i+k][1]] ? H[i+k][1] : H[i][1];
} // query log N almost O(1)
int query_dp(int a, int b) {
     for(int l=0;;l++) if (a+(1<<l+1) > b) {
           int o2 = H[b-(1<<1)+1][1];
           return A[H[a][1]] < A[o2] ? H[a][1]:o2;
}
     }
```

### RMQ segment tree

```
const int M = 100005;
int n, in[M], f[M], st[M], en[M];
struct data { int 1, r, ans, next_1, next_r; };
data d[4*M]; // which is the range? :S
int nd;
int build( int 1, int r, int id ) {
    d[id].1 = 1, d[id].r = r;
    if( l == r ) d[ id ].ans = f[1];
    else {
        int bar = (r-1)/2+1;
        d[id].next_1 = ++nd, d[id].next_r = ++nd;
        int left = build( 1, bar, d[id].next_1 );
        int right = build( bar+1, r, d[id].next_r );
        d[id].ans = max( left, right );
    return d[ id ].ans;
}
int query( int 1, int r, int id = 0 ) {
    if (1 > r) return 0;
    if( d[id].1 == 1 && d[id].r == r ) return d[id].ans;
    else {
        int bar = (d[id].r-d[id].1) / 2 + d[id].1;
        int left = 0, right = 0;
        if( 1 <= bar ) {
            if( r <= bar ) left = query( 1, r, d[id].next_1 );</pre>
            else {
                left = query( 1, bar, d[id].next_l );
                right = query( bar+1, r, d[id].next_r );
        }else right = query( l, r, d[id].next_r );
        return max( left, right );
    }
}
String algorithms (SuffixA, Aho-Corasick, KMP)
Suffix arrays
const int N = 100 * 1000 + 10;
char str[N]; bool bh[N], b2h[N];
int rank[N], pos[N], cnt[N], next[N], lcp[N];
bool smaller(int a, int b) { return str[a] < str[b];}</pre>
void suffix_array(int n) {
    REP(i,n)pos[i]=i, b2h[i]=false;
    sort(pos,pos+n,smaller);
    REP(i,n) bh[i]=!i||str[pos[i]] != str[pos[i-1]];
    for(int h=1;h< n;h*=2) {
        int buckets=0;
        for(int i=0,j; i<n; i=j) {
            j=i+1;
            while(j<n && !bh[j])j++;
            next[i]=j;
```

buckets++;

```
}
        if(buckets==n)break;
        for(int i=0;i<n;i=next[i]) {</pre>
            cnt[i] = 0;
            FOR(j, i, next[i]-1) rank[pos[j]]=i;
        }
        cnt[rank[n-h]]++;
        b2h[rank[n-h]]=true;
        for(int i=0;i<n;i=next[i]) {</pre>
            FOR(j, i, next[i]-1) {
                 int s = pos[j]-h;
                 if(s>=0)
                     rank[s] = rank[s] + cnt[rank[s]]++;
                     b2h[rank[s]]=true;
            FOR(j, i, next[i]-1) {
                 int s = pos[j]-h;
                 if (s>=0 \&\& b2h[rank[s]])
                     for(int k=rank[s]+1;!bh[k] && b2h[k]; k++) b2h[k]=false;
        REP(i,n) pos[rank[i]]=i, bh[i]|=b2h[i];
    REP(i,n) pos[rank[i]]=i;
}
void get_lcp(int n) {
    lcp[0]=0;
    int h=0;
    REP(i,n) if(rank[i]) {
        int j=pos[rank[i]-1];
        while(i+h < n \&\& j+h < n \&\& str[i+h] == str[j+h]) h++;
        lcp[rank[i]]=h;
        if(h)h--;
}
    }
Aho-Corasick
#define NC 26
#define NP 10005
#define M 100005
#define MM 500005
char a[M];
char b[NP][105];
int nb, cnt[NP], lenb[NP], alen;
int g[MM][NC], ng, f[MM], marked[MM];
int output[MM], pre[MM];
\#define\ init(x)\ \{REP(_i,NC)g[x][_i] = -1;\ f[x]=marked[x]=0;\ output[x]=pre[x]=-1;\ \}
void match() {
    ng = 0;
    init( 0 );
    // part 1 - building trie
    REP(i,nb) {
        cnt[i] = 0;
        int state = 0, j = 0;
```

```
while (g[state][b[i][j]] != -1 \&\& j < lenb[i]) state = g[state][b[i][j]], j++;
        while( j < lenb[i] ) {</pre>
            g[state][b[i][j]] = ++ng;
            state = ng;
            init( ng );
            ++j;
        }
        if( ng >= MM ) { cerr <<"i am dying"<<endl; while(1); // suicide }</pre>
        output[ state ] = i;
    // part 2 - building failure function
    queue< int > q;
    REP(i,NC) if ( g[0][i] != -1 ) q.push( g[0][i] );
    while( !q.empty() ) {
        int r = q.front(); q.pop();
        REP(i,NC) if (g[r][i] != -1) {
            int s = g[r][i];
            q.push(s);
            int state = f[r];
            while( g[state][i] == -1 && state ) state = f[state];
            f[s] = g[state][i] == -1 ? 0 : g[state][i];
        }
    }
    // final smash
    int state = 0;
    REP(i,alen) {
        while( g[state][a[i]] == -1 ) {
            state = f[state];
            if( !state ) break;
        }
        state = g[state][a[i]] == -1 ? 0 : g[state][a[i]];
        if( state && output[ state ] != -1 ) marked[ state ] ++;
    // counting
    REP(i,ng+1) if( i && marked[i] ) {
        int s = i;
        while(s != 0) cnt[output[s]] += marked[i], s = f[s];
    }
KMP
int f[ len ];
f[0] = f[1] = 0;
FOR(i,2,len) {
    int j = f[i-1];
    while(true) {
        if(s[j] == s[i-1]) { f[i] = j + 1; break;}
        }else if( !j ) { f[i] = 0; break;
        else j = f[j];
    }
i = j = 0;
```

}

}

```
while( true ) {
    if( i == len ) break;
    if(text[i] == s[j]) { i++, j++;}
        if( j == slen ) // match found
    else if( j > 0 ) j = f[j];
    else i++;
}
Graph algorithms (LCA, SCC, NetFlow, MinCost, BPM)
Tarjan's offline LCA
function TarjanOLCA(u)
    MakeSet(u); u.ancestor := u;
    for each v in u.children do
        TarjanOLCA(v); Union(u,v); Find(u).ancestor := u;
    u.colour := black;
    for each v such that \{u,v\} in P and v.color==black do
        print "LCA", u, v, Find(v).ancestor
Tarjan's Strong Connected Components
procedure tarjan(v)
  index = count; v.lowlink = count++; S.push(v);color[v] = 1;
  for all (v, v2) in E do
     if (!color[v2])
        tarjan(v2); v.lowlink = min(v.lowlink, v2.lowlink);
     else if (color[v2]==1)
        v.lowlink = min(v.lowlink, v2.lowlink);
  if (v.lowlink == index)
    do { v2 = S.top(); S.pop(); print v2; color[v2]=2; } while (v2 != v);
for all v in V do if(!color[v]) tarjan(v);
Bipartite matching with Konig
#define M 1010
int grid[M][M], 1[M], r[M], seen[M], rows, cols;
bool dfs(int x)
{
    if( seen[x] ) return false;
    seen[x] = true;
    Rep(i,cols) if(grid[x][i]) if(r[i] == -1 \mid \mid dfs(r[i]))
    {
        r[i] = x, l[x] = i;
        return true;
    return false;
}
int bpm() {
    SET( 1, -1 );
    SET(r, -1);
    int ret = 0;
    Rep(i,rows) {
```

```
SET( seen, 0 );
        if( dfs( i ) ) ret ++;
    }
    return ret;
}
bool lT[M], rT[M];
void konigdfs(int x)
{
    if( !1T[x] ) return; 1T[x] = 0;
    Rep(i,cols) if(grid[x][i] && i != l[x])
    {
        rT[i] = true;
        if( r[i] != -1) konigdfs(r[i]);
    }
}
int konig()
{
    SET(1T, 1); SET(rT, 0);
    Rep(i,rows) if(l[i] == -1) konigdfs(i);
}
Network flow - Slow
#define M 750
int nr, nc, o = 355, source = 740, sink = 741;
vector<int> edge[M];
int cap[M][M];
bool vis[M];
void init() {
    REP(i,M) edge[i].clear();
    SET( cap, 0 );
}
void add( int a, int b, int c, int d ) {
    edge[a].pb(b), edge[b].pb(a);
    cap[a][b] += c, cap[b][a] += d;
}
int dfs( int src, int snk, int fl ) {
    if( vis[src] ) return 0;
    if( snk == src ) return fl;
    vis[src] = 1;
    REP(i,edge[src].sz) {
        int v = edge[src][i];
        int x = min(fl, cap[src][v]);
        if(x > 0) {
            x = dfs(v, snk, x);
            if( !x ) continue;
            cap[src][v] -= x;
            cap[v][src] += x;
            return x;
        }
    }
    return 0;
```

```
}
int flow( int src, int snk ) {
    int ret = 0;
    while(1) {
        SET( vis, 0 );
        int delta = dfs( src, snk, 1<<30 );</pre>
        if( !delta ) break;
        ret += delta;
    }
    return ret;
}
Network flow - Dinic fast
const int maxN = 5005;
const int maxE = 70000;
const int inf = 1000000005;
int nnode, nedge, src, snk;
int Q[ maxN ], pro[ maxN ], fin[ maxN ], dist[ maxN ];
int flow[ maxE ], cap[ maxE ], to[ maxE ], next[ maxE ];
void init( int _nnode, int _src, int _snk ) {
    nnode = _nnode, nedge = 0, src = _src, snk = _snk;
    FOR(i,1,nnode) fin[i] = -1;
}
void add( int a, int b, int c1, int c2 ) {
    to[nedge]=b, cap[nedge]=c1, flow[nedge]=0, next[nedge]=fin[a], fin[a]=nedge++;
    to[nedge]=a, cap[nedge]=c2, flow[nedge]=0, next[nedge]=fin[b], fin[b]=nedge++;
}
bool bfs() {
    SET( dist, -1 );
    dist[src] = 0;
    int st = 0, en = 0;
    Q[en++] = src;
    while(st < en ) {
        int u = Q[st++];
        for(int e = fin[u]; e \ge 0; e = next[e]) {
            int v = to[e];
            if( flow[e] < cap[e] && dist[v] == -1 ) {
                dist[v] = dist[u] + 1;
                Q[en++] = v;
            }
        }
    }
    return dist[snk] != -1;
}
int dfs(int u, int fl) {
    if( u == snk ) return fl;
    for( int& e = pro[u]; e >= 0; e = next[e] ) {
        int v = to[e];
        if( flow[e] < cap[e] && dist[v] == dist[u]+1 ) {
            int x = dfs(v, min(cap[e] - flow[e], fl));
            if(x > 0) {
                flow[ e ] += x, flow[ e^1 ] -= x;
```

```
return x;
            }
        }
    return 0;
}
LL dinic() {
    LL ret = 0;
    while( bfs() ) {
        FOR(i,1,nnode) pro[i] = fin[i];
        while( 1 ) {
            int delta = dfs( src, inf );
            if( !delta ) break;
            ret += delta;
        }
    }
    return ret;
}
Min-Cost Max Flow
#define N 705
int n, nE;
int d[N], pre[N];
struct edge {
    int to, cost, cap;
    int back;
};
edge E[N*N];
vector< int > e[N];
int mincost( int s, int t, int lim ) {
    int flow = 0, ret = 0;
    while( flow < lim ) {</pre>
        SET( d, -1 ); SET( pre, -1 );
        d[s] = 0;
//
        cout <<"source "<< s <<" sink " << t << endl;</pre>
        // bellman ford
        int jump = n-1;
        bool done = 0;
        while (!done && --jump >= 0) {
            done = 1;
            REP(i,n) if ( d[i] != -1 ) REP(j,e[i].sz) {
                edge& x = E[e[i][j]];
                int v = x.to;
                 if (x.cap > 0 \& (d[v] == -1 || d[v] > d[i] + x.cost)) {
                     d[v] = d[i] + x.cost;
                     pre[v] = x.back;
                     done = 0;
```

```
//
                     cout<<v<" "<<d[v]<<endl;
                 }
            if( done ) break;
        }
//
        cout << d[t] << endl;</pre>
        if( d[t] == -1 ) break;
//
        cout <<"found one path "<<endl;</pre>
        // traverse back
        int x = t, cflow = 1 << 30;
        while(x != s) {
            edge& ed = E[ pre[x] ];
            cflow = min( cflow, E[ ed.back ].cap );
//
            cout << ed.to <<" to "<< x << endl;</pre>
            x = ed.to;
        }
        if( !cflow ) break;
        int take = min( lim - flow, cflow );
        ret += d[t] * take;
        flow += take;
//
        cout <<"taken flow "<< take <<" with cost "<< d[t] * take << endl << endl;</pre>
        x = t;
        while (x != s) {
            edge& back = E[ pre[x] ];
            edge& forw = E[ back.back ];
            back.cap += take;
            forw.cap -= take;
            x = back.to;
        }
    }
// cout << "total flow " << flow << endl;</pre>
    if( flow < lim ) return -1;</pre>
    return ret;
}
// remember to add -cost in the opposite direction
void add( int u, int v, int uv, int vu, int fuv, int fvu ) {
    int a = nE, b = nE+1;
    nE += 2;
    E[ a ].to = v, E[ a ].cost = uv, E[ a ].cap = fuv, E[ a ].back = b;
    E[ b ].to = u, E[ b ].cost = vu, E[ b ].cap = fvu, E[ b ].back = a;
    e[u].pb(a), e[v].pb(b);
}
```

## Geometry

Circle using three points Let A = (0,0) centers are  $(C_y(B_x^2 + B_y^2) - B_y(C_x^2 + C_y^2))/D'$  and  $(B_x(C_x^2 + C_y^2) - C_x(B_x^2 + B_y^2))/D$  where  $D = 2(B_xC_y - B_yC_x)$ .

# Appendices (ASCII table)

32 0010 0000	33   0010   0001	34 0010 0010	35 0010 0011	36 0010 0100	37 0010 0101	38 0010 0110	39 0010 0111	40 0010	41 0010 1001	42 0010	43 0010	44 0010	45 0010 1101	46 0010	47 0010
SP	!	"	#	\$	%	&	,	(	)	*	+	,	_		/
48 0011 0000	49 0011	50 0011 0010	51 0011	52 0011 0100	53 0011	54 0011 0110	55 0011 0111	56 0011	57 0011 1001	58 0011 1010	59 0011 1011	60 0011	61 0011	62 0011	63 0011
0	1	2	3	4	5	6	7	8	9	:	;	<	=	>	?
64 0100 0000	65 0100 0001	66 0100	67 0100	68 0100 0100	69 0100	70 0100	71 0100	72 0100	73 0100	74 0100 1010	75 0100 1011	76 0100	77 0100 1101	78 0100 1110	79 0100
@	Α	В	С	D	Ε	F	G	Н	1	J	K	L	М	Ν	Ø
80 0101	81 0101	82 0101 0010	83 0101	84 0101	85 0101 0101	86 0101 0110	87 0101 0111	88 0101	89 0101	90 0101	91 0101	92 0101	93 0101	94 0101	95 0101 1111
Р	Q	R	S	Т	U	٧	W	Χ	Υ	Z	[	\	]	^	_
96 0110	97 0110	98 0110	99 0110	100 0110	101 0110	102 0110	103 0110	104 0110 1000	105 0110	106 0110	107 1011	108 1100	109 9110	$110^{\frac{0110}{1110}}$	111 0110
`	a	b	С	d	е	f	g	h	i	j	k	1	m	n	0
112 0111	113 0111	114 0111	115	116 0111	117 0111	$118_{0110}^{0111}$	119 0111	$120\frac{0111}{1000}$	121 0111	122 1010	123 1011	124 1100	125	126	127
р	q	r	S	t	u	٧	W	X	у	Z	{		}	~	DEL DEL