2 3

3

5

## Contents IIUM cat-us-trophy

```
Combinatorics
 RMQ DP
 .vimrc
set ai ts=4 sw=4 st=4 noet nu nohls
svntax enable
filetype plugin indent on
map <F6> :w<CR>:!g++ % -g && (ulimit -c unlimited; ./a.out < ~/input.txt) <CR>
map <F5> <F6>
colo pablo
map <F12> :!gdb ./a.out -c core <CR>
template.cpp
#include<cstdio>
#include<sstream>
#include<cstdlib>
#include<cctype>
#include<cmath>
#include < algorithm >
#include<set>
#include<queue>
#include<stack>
#include<list>
#include<iostream>
#include<string>
#include<vector>
#include<cstring>
#include < map >
#include<cassert>
#include<climits>
using namespace std;
#define REP(i,n) for(int i=0, _e(n); i<_e; i++)
#define FOR(i,a,b) for(int i(a), _e(b); i <= _e; i++)
#define FORD(i,a,b) for(int i(a), _e(b); i \ge _e; i - -)
#define FORIT(i, m) for (__typeof((m).begin()) i=(m).begin(); i!=(m).end(); ++i)
#define SET(t,v) memset((t), (v), sizeof(t))
#define ALL(x) x.begin(), x.end()
#define UNIQUE(c) (c).resize( unique( ALL(c) ) - (c).begin() )
#define sz size()
#define pb push_back
#define VI vector<int>
#define VS vector<string>
```

### **Combinatorics**

#### **Mathematical Sums**

$$\begin{array}{ll} \sum_{k=0}^{n} k = n(n+1)/2 & \sum_{k=0}^{b} k = (a+b)(b-a+1)/2 \\ \sum_{k=0}^{n} k^2 = n(n+1)(2n+1)/6 & \sum_{k=0}^{n} k^3 = n^2(n+1)^2/4 \\ \sum_{k=0}^{n} k^4 = (6n^5+15n^4+10n^3-n)/30 & \sum_{k=0}^{n} k^5 = (2n^6+6n^5+5n^4-n^2)/12 \\ \sum_{k=0}^{n} x^k = (x^{n+1}-1)/(x-1) & \sum_{k=0}^{n} kx^k = (x-(n+1)x^{n+1}+nx^{n+2})/(x-1)^2 \end{array}$$

#### Binomial coefficients

	0	1	2	3	4	5	6	7	8	9	10	11	12	$\binom{n}{l} = \frac{n!}{(n-l)!l!}$
0	1													$\binom{k}{k} - \frac{\overline{(n-k)!k!}}{\binom{n}{k}} = \binom{n-1}{k} + \binom{n-1}{k-1}$
1	1	1												$\binom{k}{n} = \binom{k}{k} + \binom{k-1}{k}$
2	1	2	1											$\binom{k}{n} = \frac{n-k}{n-k+1} \binom{k}{n}$
3	1	3	3	1										$\binom{k}{k} - \binom{k}{k-1}$
4	1	4	6	4	1									$\begin{pmatrix} k \end{pmatrix} - \begin{pmatrix} n-k+1 \end{pmatrix} \begin{pmatrix} k \end{pmatrix}$
5	1	5	10	10	5	1								$\binom{n}{k+1} = \frac{n-k}{k+1} \binom{n}{k}$
6	1	6	15	20	15	6	1							
7	1	7	21	35	35	21	7	1						$\sum_{n=1}^{n} (n)$
8	1	8	28	56	70	56	28	8	1					$\sum_{k=1}^{n} k \binom{n}{k} = n2^{n-1}$ $\sum_{k=1}^{n} k^{2} \binom{n}{k} = (n+n^{2})2^{n-2}$
9	1	9	36	84	126	126	84	36	9	1				$\sum_{k=1}^{n} k^{2} \binom{n}{k} = (n+n^{2})2^{n-2}$
10	1	10	45	120	210	252	210	120	45	10	1			
11	1	11	55	165	330	462	462	330	165	55	11	1		$(m \perp n)$ $\sum_{r} r = (m) (n)$
12	1	12	66	220	495	792	924	792	495	220	66	12	1	$\binom{m+n}{r} = \sum_{k=0}^{r} \binom{m}{k} \binom{n}{r-k}$
	0	1	2	3	4	5	6	7	8	9	10	11	12	$\binom{r}{k} = \prod_{i=1}^{k} \frac{n-k+i}{i}$

Catalan numbers  $C_n = \frac{1}{n+1} {2n \choose n}$ .  $C_0 = 1$ ,  $C_n = \sum_{i=0}^{n-1} C_i C_{n-1-i}$ .  $C_{n+1} = C_n \frac{4n+2}{n+2}$ .  $C_0, C_1, \ldots = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900, \ldots$ 

 $C_n$  is the number of: properly nested sequences of n pairs of parentheses; rooted ordered binary trees with n+1 leaves; triangulations of a convex (n+2)-gon.

**Derangements**. Number of permutations of n = 0, 1, 2, ... elements without fixed points is 1, 0, 1, 2, 9, 44, 265, 1854, 14833, ... Recurrence:  $D_n = (n-1)(D_{n-1} + D_{n-2}) = nD_{n-1} + (-1)^n$ . Corollary: number of permutations with exactly k fixed points is  $\binom{n}{k}D_{n-k}$ .

Stirling numbers of  $1^{st}$  kind .  $s_{n,k}$  is  $(-1)^{n-k}$  times the number of permutations of n elements with exactly k permutation cycles.  $\begin{bmatrix} n \\ k \end{bmatrix} = |s_{n,k}| = |s_{n-1,k-1}| + (n-1)|s_{n-1,k}|$  s(0,0) = 1 and s(n,0) = s(0,n) = 0.

Stirling numbers of  $2^{nd}$  kind.  $S_{n,k}$  is the number of ways to partition a set of n elements into exactly k non-empty subsets.  $\binom{n}{k} = S_{n,k} = S_{n-1,k-1} + kS_{n-1,k}$ .  $S_{n,1} = S_{n,n} = 1$ .

**Bell numbers** .  $B_n$  is the number of partitions of n elements.  $B_0, \ldots = 1, 1, 2, 5, 15, 52, 203, 877, \ldots$   $B_{n+1} = \sum_{k=0}^{n} {n \choose k} B_k = \sum_{k=1}^{n+1} S_{n,k}$ . Bell triangle:  $B_r = a_{r,1} = a_{r-1,r-1}, \ a_{r,c} = a_{r-1,c-1} + a_{r,c-1}$ .

**Eulerian numbers** .  $E(n,k) = {n \choose k}$  is the number of permutations with exactly k descents  $(i: \pi_i < \pi_{i+1})$  / ascents  $(\pi_i > \pi_{i+1})$  / excedances  $(\pi_i > i)$  / k+1 weak excedances  $(\pi_i \ge i)$ .

Formula: E(n,m) = (m+1)E(n-1,m) + (n-m)E(n-1,m-1). E(n,0) = E(n,n-1) = 1.  $E(n,m) = \sum_{k=0}^{m} (-1)^k \binom{n+1}{k} (m+1-k)^n$ .

```
Double factorial. Permutations of the multiset \{1, 1, 2, 3, \dots n, n\} such that for each k, all the num-
bers between two occurrences of k in the permutation are greater than k. (2n-1)!! = \prod_{k=1}^{n} (2k-1).
Eulerian numbers of 2^{nd} kind. Related to Double factorial, number of all such permutations
that have exactly m ascents. \left\langle \left\langle {n \atop m} \right\rangle \right\rangle = (2n-m-1)\left\langle \left\langle {n-1 \atop m-1} \right\rangle \right\rangle + (m+1)\left\langle \left\langle {n-1 \atop m} \right\rangle \right\rangle. \left\langle \left\langle {n \atop 0} \right\rangle \right\rangle = 1
Multinomial theorem . (a_1 + \cdots + a_k)^n = \sum \binom{n}{n_1, \dots, n_k} a_1^{n_1} \dots a_k^{n_k}, where n_i \ge 0 and \sum n_i = n.
\binom{n}{n_1,\dots,n_k} = M(n_1,\dots,n_k) = \frac{n!}{n_1!\dots n_k!} M(a,\dots,b,c,\dots) = M(a+\dots+b,c,\dots)M(a,\dots,b)
RMQ DP
int make_dp(int n) { // N log N
     REP(i,n) H[i][0]=i;
     for(int l=0,k; (k=1<<1) < n; l++) for(int i=0;i+k< n;i++)
          H[i][1+1] = A[H[i][1]] > A[H[i+k][1]] ? H[i+k][1] : H[i][1];
} // query log N almost O(1)
int query_dp(int a, int b) {
     for(int l=0;;l++) if (a+(1<<l+1) > b) {
          int o2 = H[b-(1<<1)+1][1];
          return A[H[a][1]] < A[o2] ? H[a][1]:o2;
}
     }
String algorithms (SuffixA, Aho-Corasick, KMP)
Suffix arrays
const int N = 100 * 1000 + 10;
char str[N]; bool bh[N], b2h[N];
int rank[N], pos[N], cnt[N], next[N], lcp[N];
bool smaller(int a, int b) { return str[a] < str[b];}</pre>
void suffix_array(int n) {
     REP(i,n)pos[i]=i, b2h[i]=false;
     sort(pos,pos+n,smaller);
     REP(i,n) bh[i]=!i||str[pos[i]] != str[pos[i-1]];
     for(int h=1;h< n;h*=2) {
          int buckets=0;
          for(int i=0,j; i<n; i=j) {
               j=i+1;
               while(j<n && !bh[j])j++;
               next[i]=j;
               buckets++;
          }
          if(buckets==n)break;
          for(int i=0;i<n;i=next[i]) {</pre>
               cnt[i] = 0;
               FOR(j, i, next[i]-1) rank[pos[j]]=i;
          cnt[rank[n-h]]++;
          b2h[rank[n-h]]=true;
          for(int i=0;i<n;i=next[i]) {</pre>
               FOR(j, i, next[i]-1) {
                     int s = pos[j]-h;
                     if(s>=0){
                          rank[s] = rank[s] + cnt[rank[s]]++;
```

b2h[rank[s]]=true;

}

}

```
FOR(j, i, next[i]-1) {
                int s = pos[j]-h;
                if(s>=0 && b2h[rank[s]])
                     for(int k=rank[s]+1;!bh[k] && b2h[k]; k++) b2h[k]=false;
        }
            }
        REP(i,n) pos[rank[i]]=i, bh[i]|=b2h[i];
    REP(i,n) pos[rank[i]]=i;
}
void get_lcp(int n) {
    lcp[0]=0;
    int h=0;
    REP(i,n) if(rank[i]) {
        int j=pos[rank[i]-1];
        while(i+h < n \&\& j+h < n \&\& str[i+h] == str[j+h]) h++;
        lcp[rank[i]]=h;
        if(h)h--;
}
    }
Aho-Corasick
#define NC 26
#define NP 10005
#define M 100005
#define MM 500005
char a[M];
char b[NP][105];
int nb, cnt[NP], lenb[NP], alen;
int g[MM][NC], ng, f[MM], marked[MM];
int output[MM], pre[MM];
\#define\ init(x)\ \{REP(_i,NC)g[x][_i] = -1;\ f[x]=marked[x]=0;\ output[x]=pre[x]=-1;\ \}
void match() {
    ng = 0;
    init( 0 );
    // part 1 - building trie
    REP(i,nb) {
        cnt[i] = 0;
        int state = 0, j = 0;
         while(g[state][b[i][j]] != -1 \&\& j < lenb[i]) state = g[state][b[i][j]], j++; 
        while(j < lenb[i]) {
            g[state][ b[i][j] ] = ++ng;
            state = ng;
            init( ng );
            ++j;
        if( ng >= MM ) { cerr <<"i am dying"<<endl; while(1); // suicide }</pre>
        output[ state ] = i;
    // part 2 - building failure function
    queue < int > q;
    REP(i,NC) if(g[0][i] != -1) q.push(g[0][i]);
    while( !q.empty() ) {
        int r = q.front(); q.pop();
```

```
REP(i,NC) if (g[r][i] != -1) {
            int s = g[r][i];
           q.push(s);
            int state = f[r];
           while( g[state][i] == -1 && state ) state = f[state];
           f[s] = g[state][i] == -1 ? 0 : g[state][i];
        }
    }
    // final smash
    int state = 0;
    REP(i,alen) {
        while(g[state][a[i]] == -1) {
           state = f[state];
            if( !state ) break;
        }
        state = g[state][a[i]] == -1 ? 0 : g[state][a[i]];
        if( state && output[ state ] != -1 ) marked[ state ] ++;
    // counting
    REP(i,ng+1) if( i && marked[i] ) {
        int s = i;
        while(s != 0) cnt[output[s]] += marked[i], s = f[s];
    }
}
KMP
int f[ len ];
f[0] = f[1] = 0;
FOR(i,2,len) {
    int j = f[i-1];
    while( true ) {
        if(s[j] == s[i-1]) { f[i] = j + 1; break;}
        else if(!j) { f[i] = 0; break;}
        else j = f[j];
    }
}
i = j = 0;
while( true ) {
    if( i == len ) break;
    if(text[i] == s[j]) { i++, j++;}
        if( j == slen ) // match found
    else if( j > 0 ) j = f[j];
    else i++;
}
Graph algorithms (LCA, SCC, NetFlow, MinCost, BPM)
Tarjan's offline LCA
```

function TarjanOLCA(u)

MakeSet(u); u.ancestor := u;
for each v in u.children do

```
TarjanOLCA(v); Union(u,v); Find(u).ancestor := u;
    u.colour := black;
    for each v such that \{u,v\} in P and v.color==black do
        print "LCA", u, v, Find(v).ancestor
Tarjan's Strong Connected Components
procedure tarjan(v)
  index = count; v.lowlink = count++; S.push(v);color[v] = 1;
  for all (v, v2) in E do
     if (!color[v2])
        tarjan(v2); v.lowlink = min(v.lowlink, v2.lowlink);
     else if (color[v2]==1)
        v.lowlink = min(v.lowlink, v2.lowlink);
  if (v.lowlink == index)
    do { v2 = S.top(); S.pop(); print v2; color[v2]=2; } while (v2 != v);
for all v in V do if(!color[v]) tarjan(v);
Bipartite matching
#define M 100
int grid[M][M];
int 1[M], r[M], seen[M];
int rows, cols;
bool dfs(int x) {
    if( seen[x] ) return 0;
    seen[x] = true;
    REP(i,cols) if( grid[x][i] ) {
        if( r[i] == -1 || dfs( r[i] ) ) {
            r[i] = x, l[x] = i;
            return true;
        }
    }
    return false;
int bpm() {
    SET( 1, -1 );
    SET(r, -1);
    int ret = 0;
    REP(i,rows) {
        SET( seen, 0 );
        if( dfs( i ) ) ret ++;
    return ret;
Network flow - Slow
#define M 750
int nr, nc, o = 355, source = 740, sink = 741;
vector<int> edge[M];
int cap[M][M];
bool vis[M];
```

}

}

```
void init() {
    REP(i,M) edge[i].clear();
    SET( cap, 0 );
void add( int a, int b, int c, int d ) {
    edge[a].pb(b), edge[b].pb(a);
    cap[a][b] += c, cap[b][a] += d;
}
int dfs( int src, int snk, int fl ) {
    if( vis[src] ) return 0;
    if( snk == src ) return fl;
    vis[src] = 1;
    REP(i,edge[src].sz) {
        int v = edge[src][i];
        int x = min(fl, cap[src][v]);
        if(x > 0)
            x = dfs(v, snk, x);
            if( !x ) continue;
            cap[src][v] -= x;
            cap[v][src] += x;
            return x;
        }
    }
    return 0;
}
int flow( int src, int snk ) {
    int ret = 0;
    while(1) {
        SET( vis, 0 );
        int delta = dfs( src, snk, 1<<30 );</pre>
        if( !delta ) break;
        ret += delta;
    }
    return ret;
}
Network flow - Dinic fast
const int maxN = 5005;
const int maxE = 70000;
const int inf = 1000000005;
int nnode, nedge, src, snk;
int Q[ maxN ], pro[ maxN ], fin[ maxN ], dist[ maxN ];
int flow[ maxE ], cap[ maxE ], to[ maxE ], next[ maxE ];
void init( int _nnode, int _src, int _snk ) {
    nnode = _nnode, nedge = 0, src = _src, snk = _snk;
    FOR(i,1,nnode) fin[i] = -1;
}
void add( int a, int b, int c1, int c2 ) {
    to[nedge]=b, cap[nedge]=c1, flow[nedge]=0, next[nedge]=fin[a], fin[a]=nedge++;
    to[nedge]=a, cap[nedge]=c2, flow[nedge]=0, next[nedge]=fin[b], fin[b]=nedge++;
}
```

```
bool bfs() {
    SET( dist, -1 );
    dist[src] = 0;
    int st = 0, en = 0;
    Q[en++] = src;
    while(st < en ) {
        int u = Q[st++];
        for(int e = fin[u]; e \ge 0; e = next[e]) {
            int v = to[e];
            if( flow[e] < cap[e] && dist[v] == -1 ) {
                dist[v] = dist[u] + 1;
                Q[en++] = v;
            }
        }
    }
    return dist[snk] != -1;
}
int dfs(int u, int fl) {
    if( u == snk ) return fl;
    for( int& e = pro[u]; e >= 0; e = next[e] ) {
        int v = to[e];
        if( flow[e] < cap[e] && dist[v] == dist[u]+1 ) {</pre>
            int x = dfs(v, min(cap[e] - flow[e], fl));
            if(x > 0) {
                flow[e] += x, flow[e^1] -= x;
                return x;
            }
        }
    }
    return 0;
}
LL dinic() {
    LL ret = 0;
    while( bfs() ) {
        FOR(i,1,nnode) pro[i] = fin[i];
        while(1) {
            int delta = dfs( src, inf );
            if( !delta ) break;
            ret += delta;
        }
    }
    return ret;
}
Min-Cost Max Flow
#define N 705
int n, nE;
int d[N], pre[N];
struct edge {
    int to, cost, cap;
    int back;
```

```
};
edge E[N*N];
vector< int > e[N];
int mincost( int s, int t, int lim ) {
    int flow = 0, ret = 0;
    while( flow < lim ) {</pre>
        SET( d, -1 ); SET( pre, -1 );
        d[s] = 0;
//
        cout <<"source "<< s <<" sink " << t << endl;</pre>
        // bellman ford
        int jump = n-1;
        bool done = 0;
        while (!done && --jump >= 0) {
            done = 1;
            REP(i,n) if (d[i] != -1) REP(j,e[i].sz) {
                 edge& x = E[e[i][j]];
                 int v = x.to;
                 if( x.cap > 0 && (d[v] == -1 || d[v] > d[i] + x.cost )) {
                     d[v] = d[i] + x.cost;
                     pre[v] = x.back;
                     done = 0;
//
                     cout << v << " " << d[v] << endl;
                 }
            }
            if( done ) break;
//
        cout << d[t] << endl;
        if( d[t] == -1 ) break;
//
        cout <<"found one path "<<endl;</pre>
        // traverse back
        int x = t, cflow = 1<<30;
        while(x != s) {
            edge& ed = E[ pre[x] ];
            cflow = min( cflow, E[ ed.back ].cap );
//
            cout << ed.to <<" to "<< x << endl;</pre>
            x = ed.to;
        }
        if( !cflow ) break;
        int take = min( lim - flow, cflow );
        ret += d[t] * take;
        flow += take;
//
        cout <<"taken flow "<< take <<" with cost "<< d[t] * take << endl << endl;
        x = t;
        while (x != s) {
            edge& back = E[ pre[x] ];
            edge& forw = E[ back.back ];
            back.cap += take;
            forw.cap -= take;
```

```
x = back.to;
        }
    }
// cout << "total flow " << flow << endl;</pre>
    if( flow < lim ) return -1;</pre>
    return ret;
}
// remember to add -cost in the opposite direction
void add( int u, int v, int uv, int vu, int fuv, int fvu ) {
    int a = nE, b = nE+1;
    nE += 2;
    E[ a ].to = v, E[ a ].cost = uv, E[ a ].cap = fuv, E[ a ].back = b;
    E[ b ].to = u, E[ b ].cost = vu, E[ b ].cap = fvu, E[ b ].back = a;
    e[ u ].pb( a ), e[ v ].pb( b );
}
Bipartite matching with Konig
#define M 100
int grid[M][M];
int 1[M], r[M], seen[M];
int rows, cols;
bool dfs(int x) {
    if( seen[x] ) return 0;
    seen[x] = true;
    REP(i,cols) if( grid[x][i] ) {
        if( r[i] == -1 || dfs( r[i] ) ) {
            r[i] = x, l[x] = i;
            return true;
        }
    }
    return false;
}
int bpm() {
    SET( 1, -1 );
    SET(r, -1);
    int ret = 0;
    REP(i,rows) {
        SET( seen, 0 );
        if( dfs( i ) ) ret ++;
    return ret;
```

}

# Appendices (ASCII table)

لـ	0	1	2	3	4	5	6	7	8	9	Α	В	C	D	Ε	ı F ı
0	NUL	SOH	STX	ETX	EOT	ENQ	ACK	BEL	BS	HT	LF	VT	FF	CR	S0	SI
1	DLE	DC1	DC2	DC3	DC4	NAK	SYN	ETB	CAN	EM	SUB	ESC	FS	GS	RS	US
2		!	- 11	#	\$	%	&	1	(	)	*	+	,	-	•	/
3	0	1	2	3	4	5	6	7	8	9		;	٧	=	^	?
4	0	Α	В	С	D	Е	F	G	Н	I	J	K	Г	М	N	0
5	Р	Q	R	S	T	U	V	W	X	Υ	Z	[	/	]	^	_
6	`	а	b	С	d	е	f	g	h	i	j	k	l	m	n	0
7	р	q	r	S	t	u	V	W	Х	у	Z	{		}	~	DEL