



## ■ Introduction aux bases de données relationnelles

- 7ème cours: **vues, autorisations et triggers en SQL**

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■ **Biblio:** chapitres 4 et 5 de *Database Systems Concepts* de Silberschatz et al, McGraw-Hill (6<sup>ème</sup> édition, 2010)

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# Intermediate and Advanced SQL

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# Chapter 4: Intermediate SQL

- **Views**
- Authorization
- Triggers (from Chap 5)



# Views: virtual relations

- In some cases, it is **not desirable for all users to see the entire logical model** (that is, all the actual relations stored in the database).
- Consider a person who needs to know an instructors name and department, but not the salary. This person should see a relation described, in SQL, by

```
select ID, name, dept_name  
from instructor
```

- A **view** provides a mechanism to **hide** certain data from the view of certain users.
- Views can also be used to **make it simpler to write queries**. Each time a view is used in a query, it is computed.



# Defining a View

- Syntax to create a new view called *v*

**create view *v* as** <query expression>

where <query expression> is any legal SQL expression.

- Once a view is defined, the **view name can be used to refer to the virtual relation** that the view generates.
- View definition is not the same as creating a new relation, by evaluating the query expression
  - The view definition **is substituted into queries using the view.**



# Examples: creating and using views

- Create a view of **instructors without their salary**  
**create view** *faculty* **as**  
    **select** *ID, name, dept\_name*  
    **from** *instructor*
- Query: find all instructors in the Biology department  
    **select** *name*  
    **from** *faculty*  
    **where** *dept\_name* = 'Biology'
- Create a **view of total salary per department. Specify attribute names in view definition.**  
    **create view** *departments\_total\_salary(dept\_name, total\_salary)* **as**  
        **select** *dept\_name, sum (salary)*  
        **from** *instructor*  
        **group by** *dept\_name*;

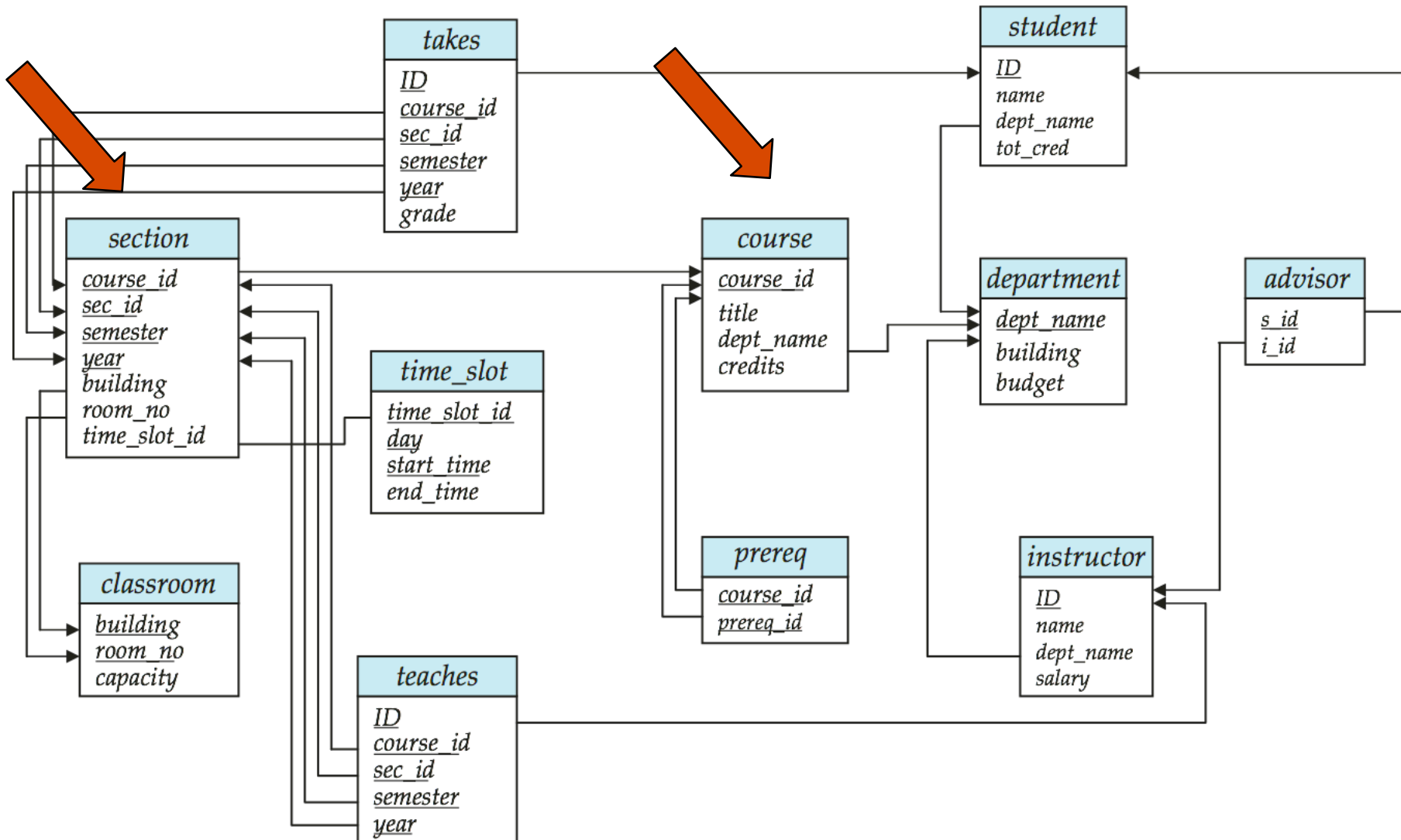


# Views Defined Using Other Views

- One view may be used in the expression defining another view,
- A view relation  $v_1$  is said to *depend directly* on a view relation  $v_2$  if  $v_2$  is used in the expression defining  $v_1$
- A view relation  $v_1$  is said to *depend on* view relation  $v_2$  if either  $v_1$  depends directly to  $v_2$  or there is a path of dependencies from  $v_1$  to  $v_2$
- A view relation  $v$  is said to be *recursive* if it depends on itself.



# Physics courses in Fall 2009







# Views Defined Using Other Views

- **create view** *physics\_fall\_2009* **as**  
    **select** *course.course\_id, sec\_id, building, room\_number*  
    **from** *course, section*  
    **where** *course.course\_id = section.course\_id*  
          **and** *course.dept\_name = 'Physics'*  
          **and** *section.semester = 'Fall'*  
          **and** *section.year = '2009'*;
- **create view** *physics\_fall\_2009\_watson* **as**  
    **select** *course\_id, room\_number*  
    **from** *physics\_fall\_2009*  
    **where** *building = 'Watson'*;



# View Expansion

- A view can't be computed and stored, because it may become out of date if the relations used to define it are modified.
- When a view appears in a query, it is replaced by the stored query expression.
- A way to define the meaning of views defined in terms of other views.
- Let view  $v_1$  be defined by an expression  $e_1$  that may itself contain uses of view relations.
- View expansion of an expression repeats the following replacement step:
  - repeat**
    - Find any view relation  $v_i$  in  $e_1$
    - Replace the view relation  $v_i$  by the expression defining  $v_i$
  - until** no more view relations are present in  $e_1$
- As long as the view definitions are not recursive, this loop will terminate.



# View Expansion

- Expand use of a view in a query/another view
- Our previous example:  
**create view** *physics\_fall\_2009\_watson* **as**  
**(select** *course\_id, room\_number*  
**from** (**select** *course.course\_id, building, room\_number*  
**from** *course, section*  
**where** *course.course\_id = section.course\_id*  
**and** *course.dept\_name = 'Physics'*  
**and** *section.semester = 'Fall'*  
**and** *section.year = '2009'*)  
**where** *building = 'Watson'*;



## Exo:

- Quelles sont des conditions de jointure que vous avez fréquemment du réécrire en TP sur la base INFOTOUR?
- Proposez des vues qui pourraient être utilisées pour simplifier des requêtes.



# Update of a View

- Keeping views up-to-date is a delicate issue!
- Add a new tuple to *faculty* view which we defined earlier

**insert into *faculty* values** ('30765', 'Green', 'Music');

- This insertion must be represented by the insertion of the following tuple into the *instructor* relation (with unknown salary!)

('30765', 'Green', 'Music', null)



# Some Updates cannot be Translated Uniquely

- **create view** *instructor\_info* **as**  
    **select** *ID, name, building*  
    **from** *instructor, department*  
    **where** *instructor.dept\_name= department.dept\_name;*
- **insert into** *instructor\_info* **values** ('69987', 'White', 'Taylor');
  - ▶ which department should we add, if multiple departments in Taylor?
  - ▶ what if no department is in Taylor?
  - ▶ ...



# Updatable views

- Most SQL implementations allow **updates, inserts and deletes only on simple views**
  - The **from** clause has only **one** database **relation**.
  - The **select** clause contains **only attribute** names of the relation, and does not have any expressions, aggregates, or **distinct** specification.
  - Any attribute not listed in the **select** clause can be set to null
  - The query does **not** have a **group** by or **having** clause.
- Beyond this: advanced topic!



# Some Updates can not be translated at all

- **create view** *history\_instructors* **as**  
**select** \*  
**from** *instructor*  
**where** *dept\_name*= 'History';
- **Insert** ('25566', 'Brown', 'Biology', 100000) into *history\_instructors*





# Exo

- Quelles des vues proposées sur INFOTOUR sont *simples*, et permettent donc des mises à jour sans problèmes?
- Quelles vues devraient être restreintes à l'utilisation dans des requêtes, sans permettre des mises à jour?



# Chapter 4: Intermediate SQL

- Join Expressions
- Integrity Constraints
- Views
- **Authorization**



# Authorization

Forms of authorization on parts of the database:

- **Read** - allows reading, but not modification of data.
- **Insert** - allows insertion of new data, but not modification of existing data.
- **Update** - allows modification, but not deletion of data.
- **Delete** - allows deletion of data.

Forms of authorization to modify the **database schema**

- **Index** - allows creation and deletion of indices.
- **Resources** - allows creation of new relations.
- **Alteration** - allows addition or deletion of attributes in a relation.
- **Drop** - allows deletion of relations.



# Authorization Specification in SQL

- The **grant** statement is used to confer authorization

**grant** <privilege list>

**on** <relation name or view name> **to** <user list>

- <user list> is:
  - a user-id
  - **public**, which allows all valid users the privilege granted
  - A role (more on this later)
- Granting a privilege on a view does not imply granting any privileges on the underlying relations.
- The grantor of the privilege must already hold the privilege on the specified item (or be the database administrator).



# Privileges in SQL

- **select:** allows read access to relation, or the ability to query using the view
  - Example: grant users  $U_1$ ,  $U_2$ , and  $U_3$  **select** authorization on the *branch* relation:

**grant select on *instructor* to  $U_1$ ,  $U_2$ ,  $U_3$**

- **insert:** the ability to insert tuples.
- **update:** the ability to update using the SQL update statement.
- **delete:** the ability to delete tuples.
- **all privileges:** used as a short form for all the allowable privileges.



# Revoking Authorization in SQL

- The **revoke** statement is used to revoke authorization.

**revoke** <privilege list>

**on** <relation name or view name> **from** <user list>

- Example:

**revoke select on** *branch* **from**  $U_1, U_2, U_3$

- <privilege-list> may be **all** to revoke all privileges the revokee may hold.
- If <revokee-list> includes **public**, all users lose the privilege except those granted it explicitly.
- If the same privilege was granted twice to the same user by different grantees, the user may retain the privilege after the revocation.
- All privileges that depend on the privilege being revoked are also revoked.



# Roles

- **create role** *instructor*;
- Privileges can be granted to roles:
  - **grant select on takes to** *instructor*;
  - This allows instructors to read all students' information on which courses they took, and grades they obtained.
- Roles can be granted to users, as well as to other roles
  - **create role** *student*
  - **grant instructor to** Amit;
  - **create role** *dean*;
  - **grant instructor to** *dean*;
  - **grant dean to** Satoshi;



# Authorization on Views

- **create view** *geo\_instructor* **as**  
(**select** \*  
**from** *instructor*  
**where** *dept\_name* = 'Geology');
- **grant select on** *geo\_instructor* **to** *staff*
- Suppose that a staff member issues
  - **select** \*  
**from** *geo\_instructor*;
- What if
  - staff does not have permissions on *instructor*?
  - creator of view did not have some permissions on *instructor*?





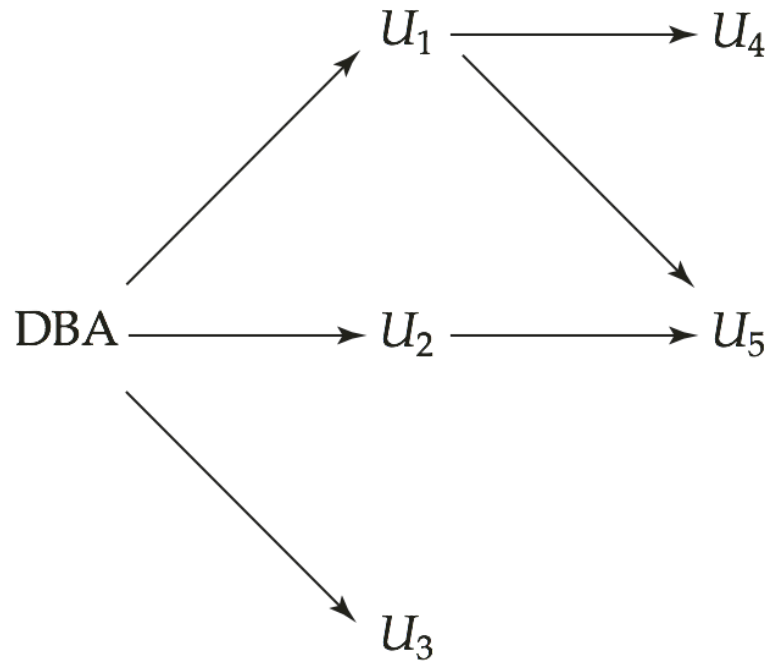
# Grant reference

- **references** privilege to **create foreign key**
  - **grant reference** (*dept\_name*) **on** *department* **to** Mariano;
  
- Why is this required?
  - User Michel creates a foreign key in a relation *r*, that references to *department.dept\_name*
  - User Michel inserts a tuple into *r*, for instance, about the Geology department (that didn't previously exist)
  - If this tuple is later deleted, the “mother” relation is modified. **This restricts the future activity of other users!**



# Transfer of privileges

- transfer of privileges
  - grant select on *department* to Amit **with grant option**;
  - revoke select on *department* from Amit, Satoshi **cascade**;
  - revoke select on *department* from Amit, Satoshi **restrict**;
- Etc.





# End of Chapter 4

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# Chapter 5: Advanced SQL - triggers

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# Triggers: event-condition-action

- Triggers allow to maintain database in a coherent state, beyond what key constraints allow.
  
- **Event**: database modification, such as insert, delete, update
- **Condition**: any true/false expression
- **Action**: Sequence of SQL statements that will automatically be executed.
  
- Triggers introduced to SQL standard in **SQL:1999**, but supported even earlier using non-standard syntax by most databases.
  - Here, we use this idealized syntax to explain the main principles.
  - Check the lab manual for Oracle Syntax!





# Trigger 1: ensure that time slot exists

- E.g. *time\_slot\_id* is not a primary key of *timeslot*, so we cannot create a foreign key constraint from *section* to *timeslot*.
- Alternative: use triggers on *section* and *timeslot* to enforce integrity constraints

**create trigger** *timeslot\_check1*

**after insert on** *section*

**referencing new row as** *nrow*

**for each row**

**when**

*/\* time\_slot\_id not present in time\_slot \*/*

*(nrow.time\_slot\_id* **not in**

*(select time\_slot\_id from timeslot))*

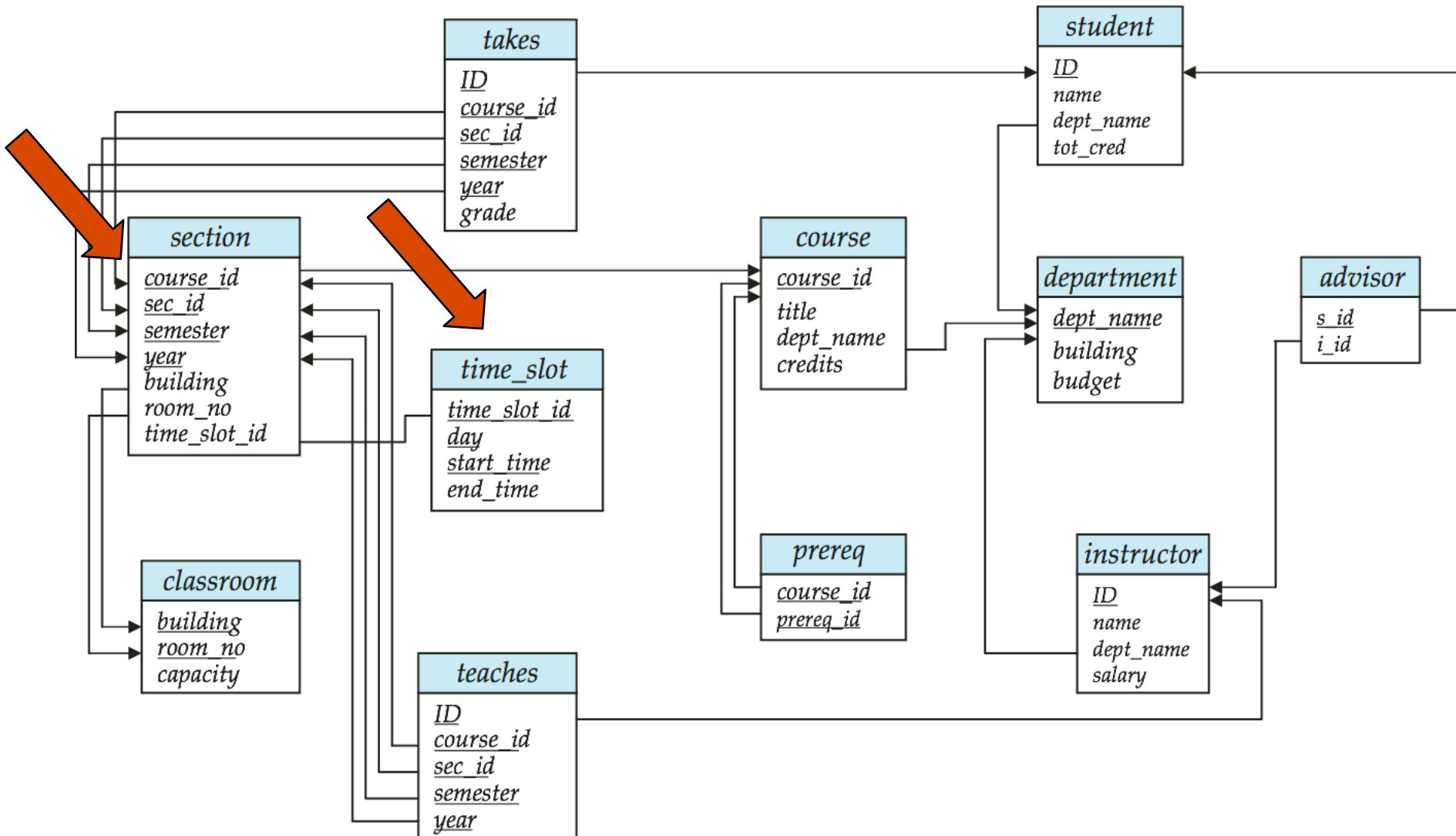
**begin**

**rollback**

**end;**



# University Database: deleting a time\_slot







# Trigger 2: don't delete a time\_slot that is still used!

```
create trigger timeslot_check2 after delete on time_slot
referencing old row as orow
for each row
when ( /* time_slot_id has just been deleted from time slot */
       orow.time_slot_id not in (
           select time_slot_id from time_slot)

       /* and time_slot_id still referenced from section*/
       and orow.time_slot_id in (
           select time_slot_id from section)
       )
begin
    rollback
end;
```



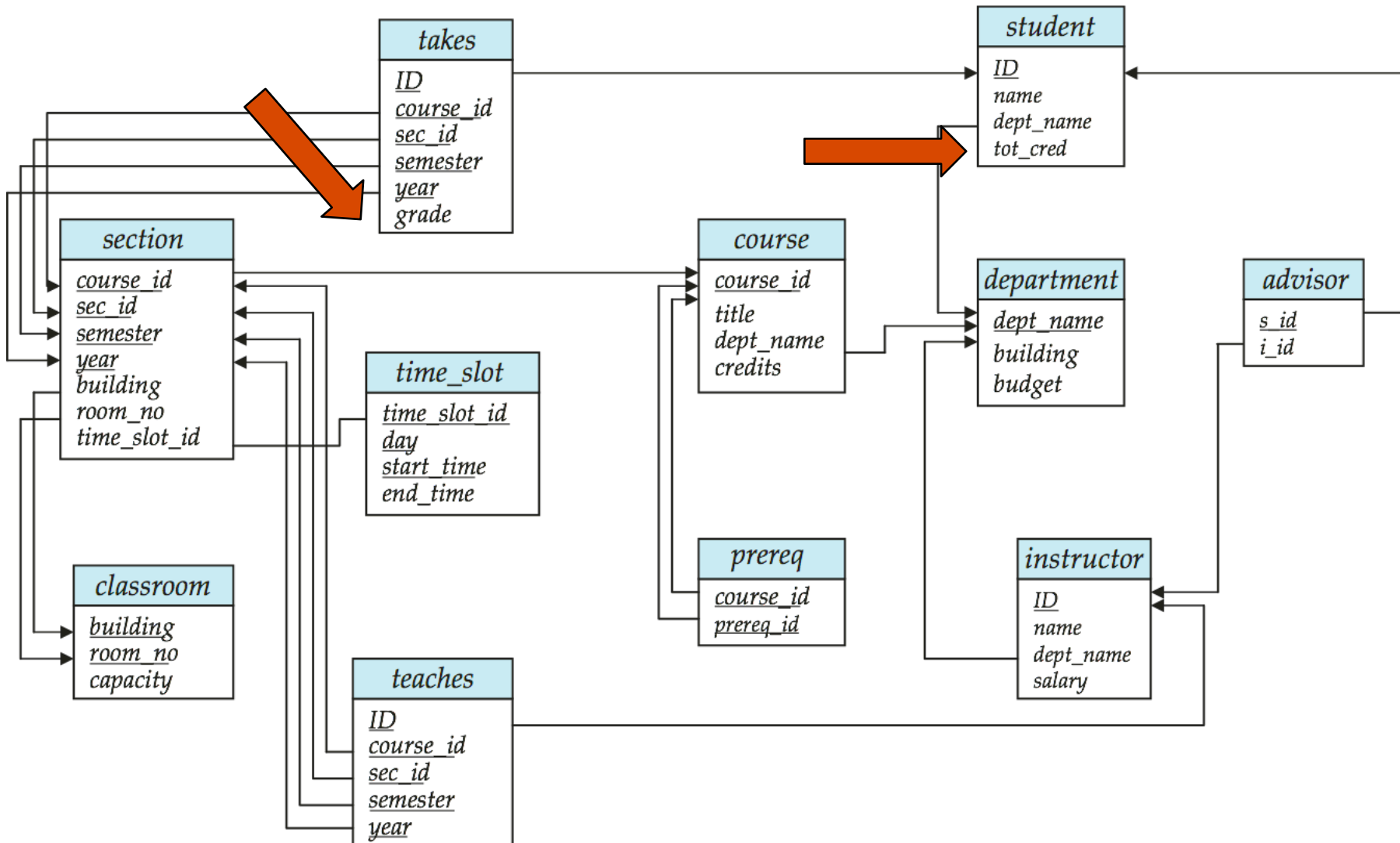
# Triggering Events and Actions in SQL

- Triggering **event** can be **insert**, **delete** or **update**
- Triggers on update can be restricted to specific attributes
  - **E.g., after update of** *takes* **on** *grade*
- Values of attributes **before** and **after** an update can be referenced
  - **referencing old row as** : for deletes and updates
  - **referencing new row as** : for inserts and updates
- Triggers can be activated before an event, which can serve as extra constraints. E.g. convert blank grades to null.

```
create trigger setnull_trigger before update of takes  
referencing new row as nrow  
for each row  
when (nrow.grade = ' ')  
begin atomic  
    set nrow.grade = null;  
end;
```



# At the end of the term, a student has passed an exam!





## Trigger 3: increase the lucky student's tot\_credit value

- **create trigger** *credits\_earned* **after update of** *takes* **on** (*grade*)  
**referencing new row as** *nrow*  
**referencing old row as** *orow*  
**for each row**  
**when** *nrow.grade*  $\neq$  'F' **and** *nrow.grade* **is not null**  
**and** (*orow.grade* = 'F' **or** *orow.grade* **is null**)  
**begin atomic**  
**update** *student*  
**set** *tot\_cred* = *tot\_cred* +  
**(select credits**  
**from course**  
**where course.course\_id =** *nrow.course\_id*)  
**where** *student.id* = *nrow.id*;  
**end;**



# When Not To Use Triggers

- Triggers were used earlier for tasks such as
  - maintaining summary data (e.g., total salary of each department)
  - Replicating databases by recording changes to special relations (called **change** or **delta** relations) and having a separate process that applies the changes over to a replica
- There are better ways of doing these now:
  - Databases today provide built in materialized view facilities to maintain summary data
  - Databases provide built-in support for replication
- Encapsulation facilities can be used instead of triggers in many cases
  - Define methods to update fields
  - Carry out actions as part of the update methods instead of through a trigger
- Risk of unintended execution of triggers, for example, when
  - loading data from a backup copy
  - replicating updates at a remote site
  - Trigger execution can be disabled before such actions.
- Other risks with triggers:



# Seen today

- Views
- Authorization
- Triggers