Introduction aux Bases de Données Relationnelles

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Evaluation

N=(2*DS+TD+TP)/4

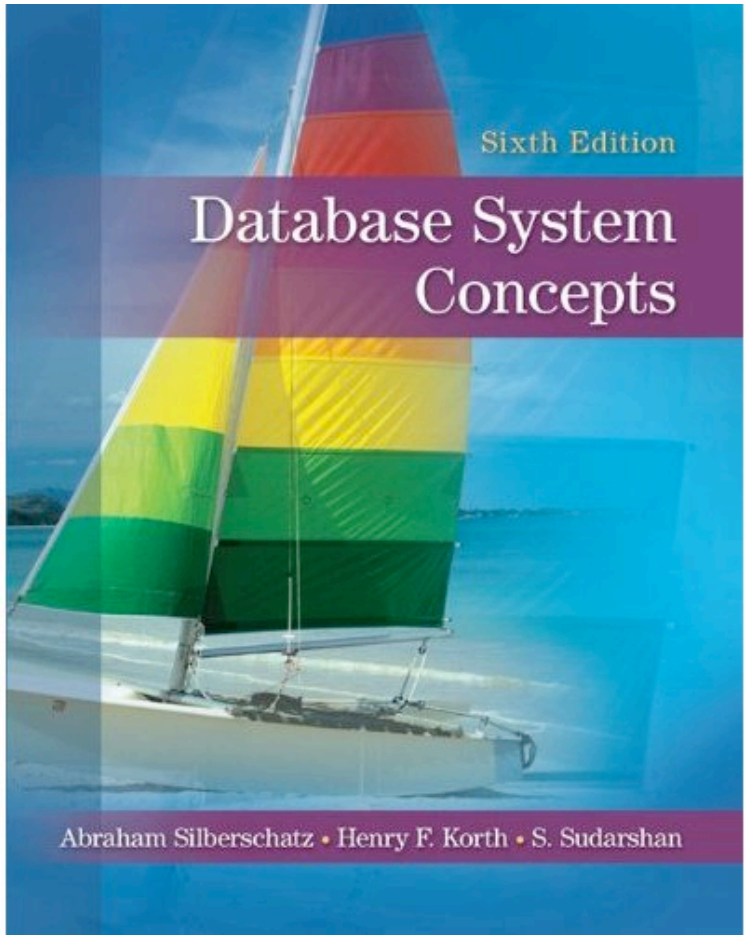
- **N**: note finale

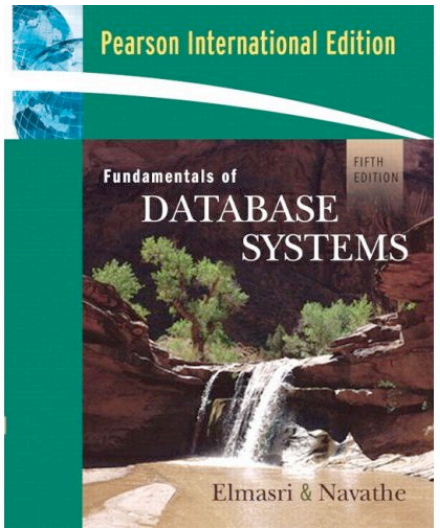
- **TD** : une note sur 20 de Travaux Dirigés, attribuée par l'enseignant de Travaux Dirigés: interrogations écrites, devoirs maison, participation, …

- **TP** : une note sur 20 de Travaux Pratiques: participation, TPs rendus et contrôle en fin de semestre

- **DS** : une note sur 20 pour l'examen de fin de semestre en décembre.

Bibliographie





www.db-book.com

Nous suivons le livre *Database System Concepts* de Silberschatz au moins pour 4 séances.

English Slides

- Textbooks provide slides for lectures
  www.db-book.com
- Please, systematically take notes of technical terms and their French translations
- University library has French editions of the books, both Silberschatz and Elmasri
- Lab manual is in French! ☺

Goals of today's lecture

- First overview of relational databases,
  - Of the topics we will cover this term,
  - of important topics, beyond this course.
- By the end of this week, you'll have seen the basics of SQL.

Chapter 1: Introduction

Database Management System (DBMS)

- DBMS contains **information** about a particular enterprise
  - Collection of interrelated **data**
  - Set of **programs** to access the data
  - An environment that is both *convenient* and *efficient* to use
- Database **applications**:
  - Banking: transactions
  - Airlines: reservations, schedules
  - Universities: registration, grades
  - Sales: customers, products, purchases
  - Online retailers: order tracking, customized recommendations
  - Manufacturing: production, inventory, orders, supply chain
  - Human resources: employee records, salaries, tax deductions
- Databases **touch all aspects of our lives**

University Database Example

- **Our running example for this course:**
  - Add new students, instructors, and courses
  - Register students for courses, and generate class rosters
  - Assign grades to students, compute grade point averages (GPA) and generate transcripts

Purpose of Database Systems

- In the **early days**, database applications were built directly on top of **file systems**

- **Drawbacks** of using **file systems** to store data:
  - Data **redundancy** and **inconsistency**
    - Multiple file formats, duplication of information in different files
  - Difficulty in accessing data
    - Need to write a new program to carry out each new task
  - Data **isolation** — multiple files and formats
  - **Integrity** problems
    - Integrity constraints (e.g., account balance > 0) become "buried" in program code rather than being stated explicitly
    - Hard to add new constraints or change existing ones

Purpose of Database Systems (Cont.)

- Drawbacks of using file systems (cont.)
  - **Atomicity** of updates
    - Failures may leave database in an inconsistent state with partial updates carried out
    - Example: Transfer of funds from one account to another should either complete or not happen at all
  - **Concurrent** access by multiple users
    - Concurrent access needed for performance
    - Uncontrolled concurrent accesses can lead to inconsistencies
      - Example: Two people reading a balance (say 100) and updating it by withdrawing money (say 50 each) at the same time
  - **Security** problems
    - Hard to provide user access to some, but not all, data
- **Database systems offer solutions** to all the above problems

Levels of Abstraction in DBMS

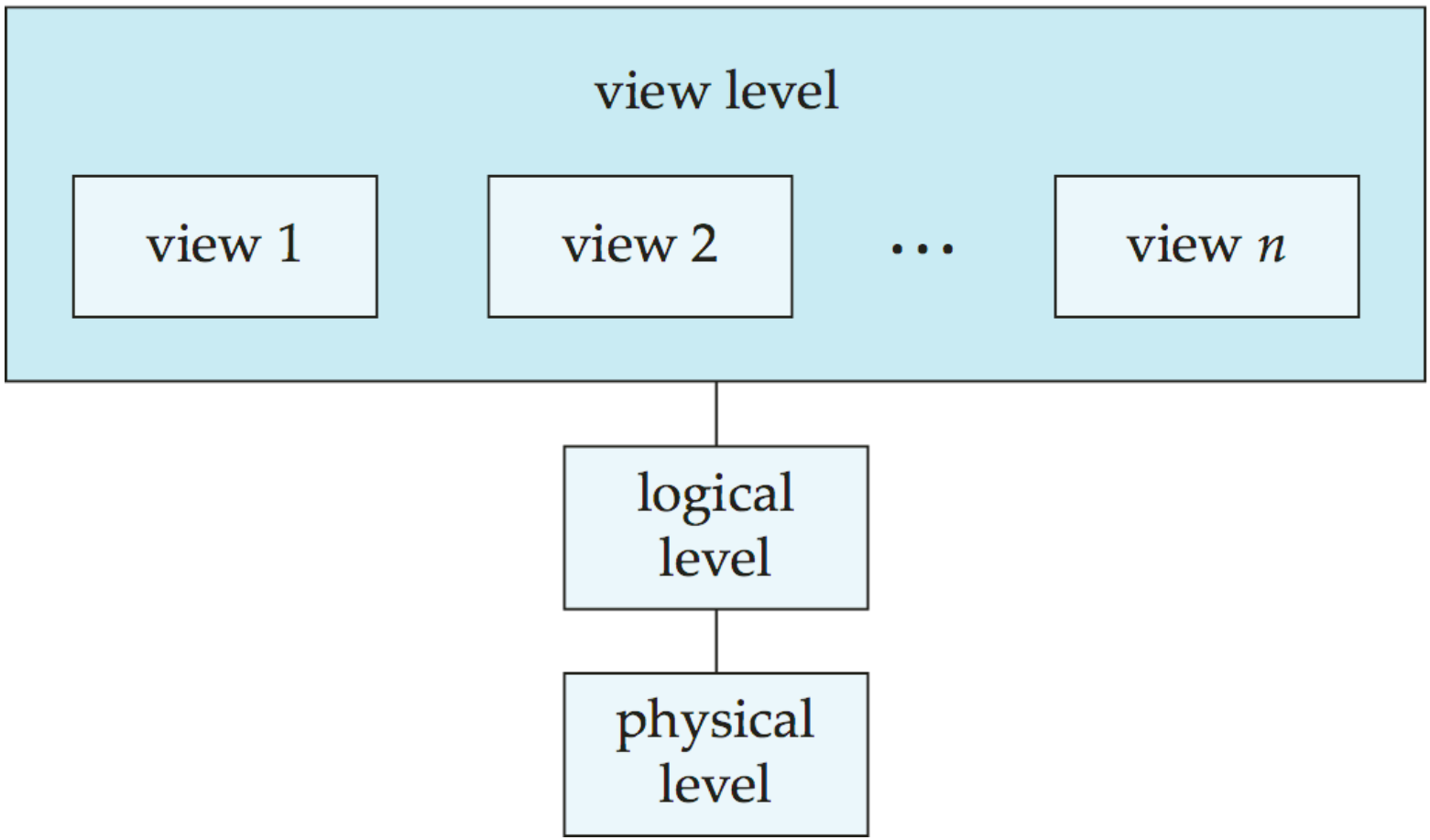
- **Physical level:** describes how a record (e.g., customer) is stored.
- **Logical level:** describes data stored in database, and the relationships among the data.

```
type instructor = record
        ID : string;
        name : string;
        dept_name : string;
        salary : integer;
    end;
```

- **View level:** application programs hide details of data types. Views can also hide information (such as an employee's salary) for security purposes.

View of Data

An architecture for a database system



Schemas and Instances

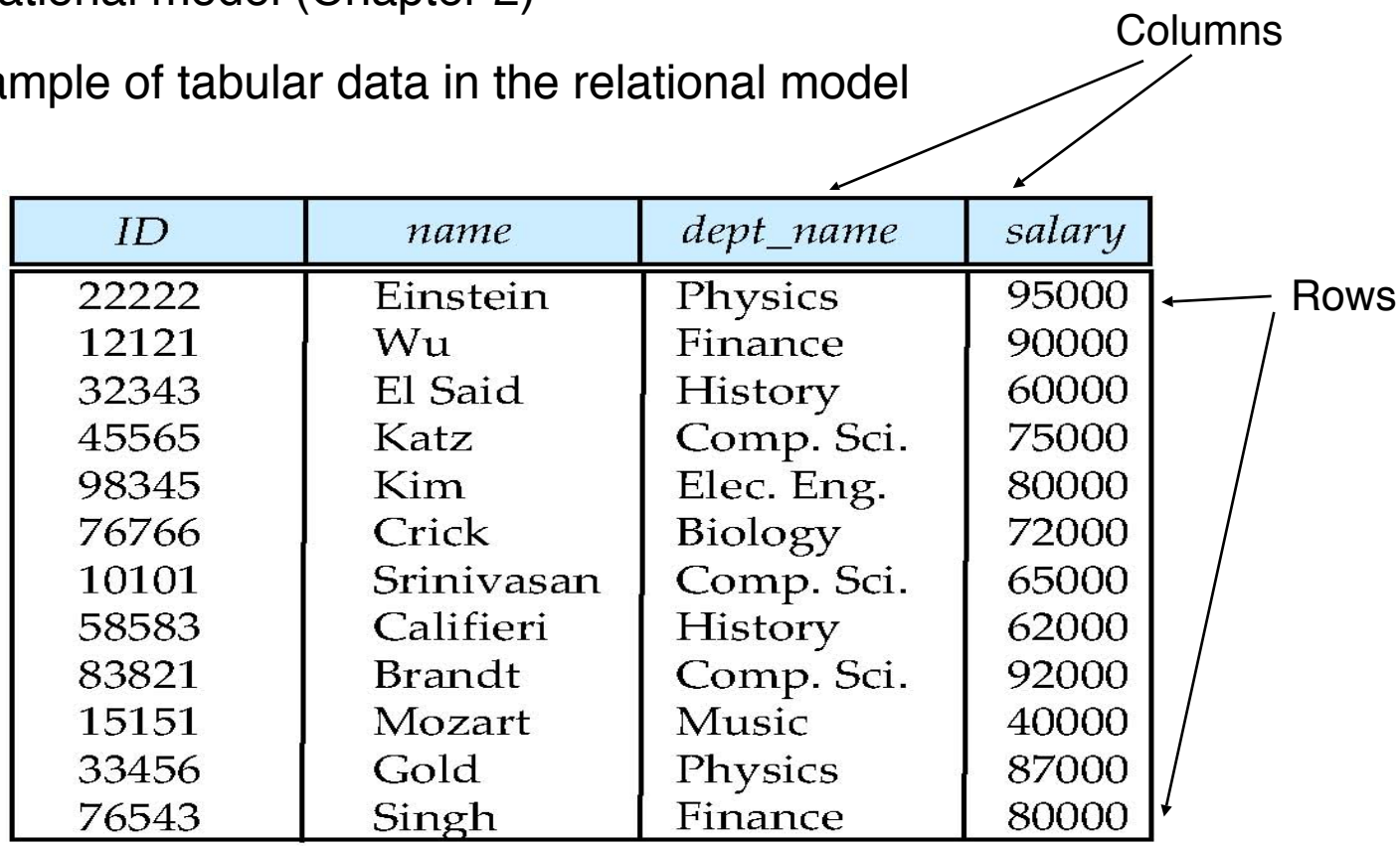
- Similar to types and variables in programming languages
- **Schema** – the logical structure of the database
  - Example: The database contains information about a set of instructors and courses and the relationship between them
  - Analogous to type information of a variable in a program
  - **Physical schema**: database design at the physical level
  - **Logical schema**: database design at the logical level
- **Instance** – the actual content of the database at a particular point in time
  - Analogous to the value of a variable
- **Physical Data Independence** – the ability to modify the physical schema without changing the logical schema
  - Applications depend on the logical schema
  - In general, the interfaces between the various levels and components should be well defined so that changes in some parts do not seriously

Data Models

- A collection of tools for describing
  - Data
  - Data relationships
  - Data semantics
  - Data constraints
- Relational model
- Entity-Relationship data model (mainly for database design)
- Object-based data models (Object-oriented and Object-relational)
- Semistructured data model (XML)
- Other older models:
  - Network model
  - Hierarchical model

Relational Model

- Relational model (Chapter 2)
- Example of tabular data in the relational model



ID	*name*	*dept_name*	*salary*
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
76766	Crick	Biology	72000
10101	Srinivasan	Comp. Sci.	65000
58583	Califieri	History	62000
83821	Brandt	Comp. Sci.	92000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
76543	Singh	Finance	80000

(a) The *instructor* table

A Sample Relational Database

ID	*name*	*dept_name*	*salary*
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
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15151	Mozart	Music	40000
33456	Gold	Physics	87000
76543	Singh	Finance	80000

(a) The *instructor* table

dept_name	*building*	*budget*
Comp. Sci.	Taylor	100000
Biology	Watson	90000
Elec. Eng.	Taylor	85000
Music	Packard	80000
Finance	Painter	120000
History	Painter	50000
Physics	Watson	70000

(b) The *department* table

Data Manipulation Language (DML)

- Language for accessing and manipulating the data organized by the appropriate data model
  - DML also known as query language
- Two classes of languages
  - **Procedural** – user specifies what data is required and how to get those data
  - **Declarative (nonprocedural)** – user specifies what data is required without specifying how to get those data
- SQL is the most widely used query language

Data Definition Language (DDL)

- Specification notation for **defining** the **database schema**

  Example: **create table** *instructor* (
  ID **char**(5),
  name **varchar**(20),
  dept_name **varchar**(20),
  salary **numeric**(8,2))

- DDL compiler generates a set of tables stored in a **data dictionary**

Define an example yourself!

- Following the example for instructor, define the database schema for **department!**

ID	*name*	*dept_name*	*salary*
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
76766	Crick	Biology	72000
10101	Srinivasan	Comp. Sci.	65000
58583	Califieri	History	62000
83821	Brandt	Comp. Sci.	92000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
76543	Singh	Finance	80000

(a) The *instructor* table

```
create table instructor (
        ID              char(5),
        name            varchar(20),
        dept_name       varchar(20),
        salary          numeric(8,2)
)
```

dept_name	*building*	*budget*
Comp. Sci.	Taylor	100000
Biology	Watson	90000
Elec. Eng.	Taylor	85000
Music	Packard	80000
Finance	Painter	120000
History	Painter	50000
Physics	Watson	70000

(b) The *department* table

Data Definition Language (DDL)

- **Data dictionary contains metadata** (i.e., data about data)
  - Database schema
  - **Integrity constraints**
    - Primary **key** (ID uniquely identifies instructors)
    - Referential integrity (**references** constraint in SQL)
      - e.g. *dept_name* value in any *instructor* tuple must appear in *department* relation
  - Authorization: who can access which data?

Data Manipulation Language: SQL

- **SQL**: widely used non-procedural language. Examples:
  - Find the name of the instructor with ID 22222

```
select  name
        from    instructor
        where   instructor.ID = '22222'
```

Query Language: SQL

- **SQL**: widely used non-procedural language. Examples:
  - Find the name of the instructor with ID 22222

```
select  name
        from    instructor
        where   instructor.ID = '22222'
```

  - **Find instructors from departments with a budget over 95000$**

```
select instructor.ID, department.dept_name

from instructor, department
    where instructor.dept_name= department.dept_name  and
    department.budget > 95000
```

- **Application programs** generally access databases through one of
  - Language extensions to allow embedded SQL

Database Design

The process of designing the general structure of the database:

- Logical Design – Deciding on the database schema. Database design requires that we find a "good" collection of relation schemas.
  - Business decision – What attributes should we record in the database?
  - Computer Science decision – What relation schemas should we have and how should the attributes be distributed among the various relation schemas?
- Physical Design – Deciding on the physical layout of the database

Good Database Design?

- Is there any problem with this design?

ID	*name*	*salary*	*dept_name*	*building*	*budget*
22222	Einstein	95000	Physics	Watson	70000
12121	Wu	90000	Finance	Painter	120000
32343	El Said	60000	History	Painter	50000
45565	Katz	75000	Comp. Sci.	Taylor	100000
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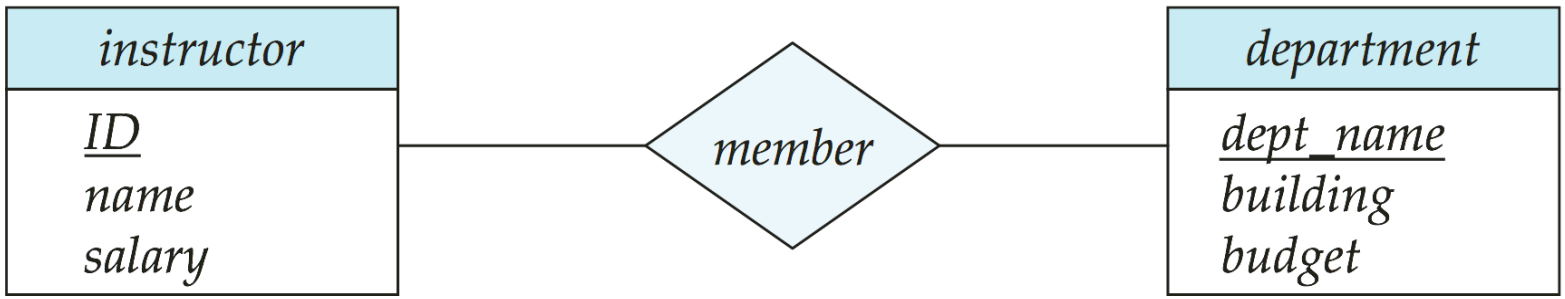
76543	Singh	80000	Finance	Painter	120000

Design Approaches

- Normalization Theory (Chapter 8)
  - Formalize what designs are bad, and test for them
- Entity Relationship Model (Chapter 7)
  - Represents an enterprise as a collection of *entities* and *relationships*
  - Represented diagrammatically by an *entity-relationship diagram*

The Entity-Relationship Model

- Models an enterprise as a collection of *entities* and *relationships*
  - Entity: a "thing" or "object" in the enterprise that is distinguishable from other objects
    - Described by a set of *attributes*
  - Relationship: an association among several entities
- Represented diagrammatically by an *entity-relationship diagram:*



What happened to dept_name of instructor?

Object-Relational Data Models

- Relational model: flat, “atomic” values
- Object Relational Data Models
  - Extend the relational data model by including object orientation and constructs to deal with added data types.
  - Allow attributes of tuples to have complex types, including non-atomic values such as nested relations.
  - Preserve relational foundations, in particular the declarative access to data, while extending modeling power.
  - Provide upward compatibility with existing relational languages.

XML: Extensible Markup Language

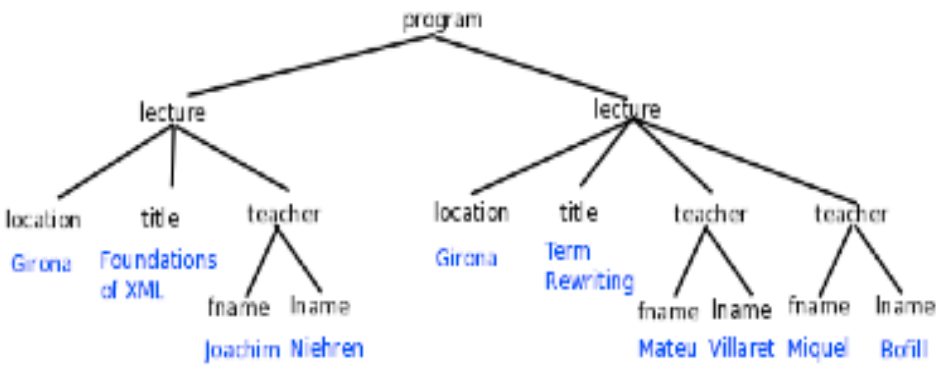
- Defined by the WWW Consortium (W3C)
- Originally intended as a document markup language not a database language
- The ability to specify new tags, and to create **nested tag structures (i.e. tree structures)** made XML a great way to exchange **data**, not just documents
- XML has become the **basis for all new generation data interchange formats**.
- A wide variety of tools is available for parsing, browsing and querying **XML** documents/data

XML: Example

```
<program>
   <lecture>
      <location> Girona </location>
      <title> Foundations of XML </title>
      <teacher>
         <fname> Joachim </fname>
         <lname> Niehren </lname> </teacher> </lecture>
   <lecture>
      <location> Girona </location>
      <title> Term Rewriting </title>
      <teacher>
         <fname> Mateu </fname>
         <lname> Villaret </lname> </teacher>
      <teacher>
         <fname> Miquel </fname>
         <lname> Bofill </lname> <teacher> </lecture>
</program>
```

XML data markup

Same data as tree

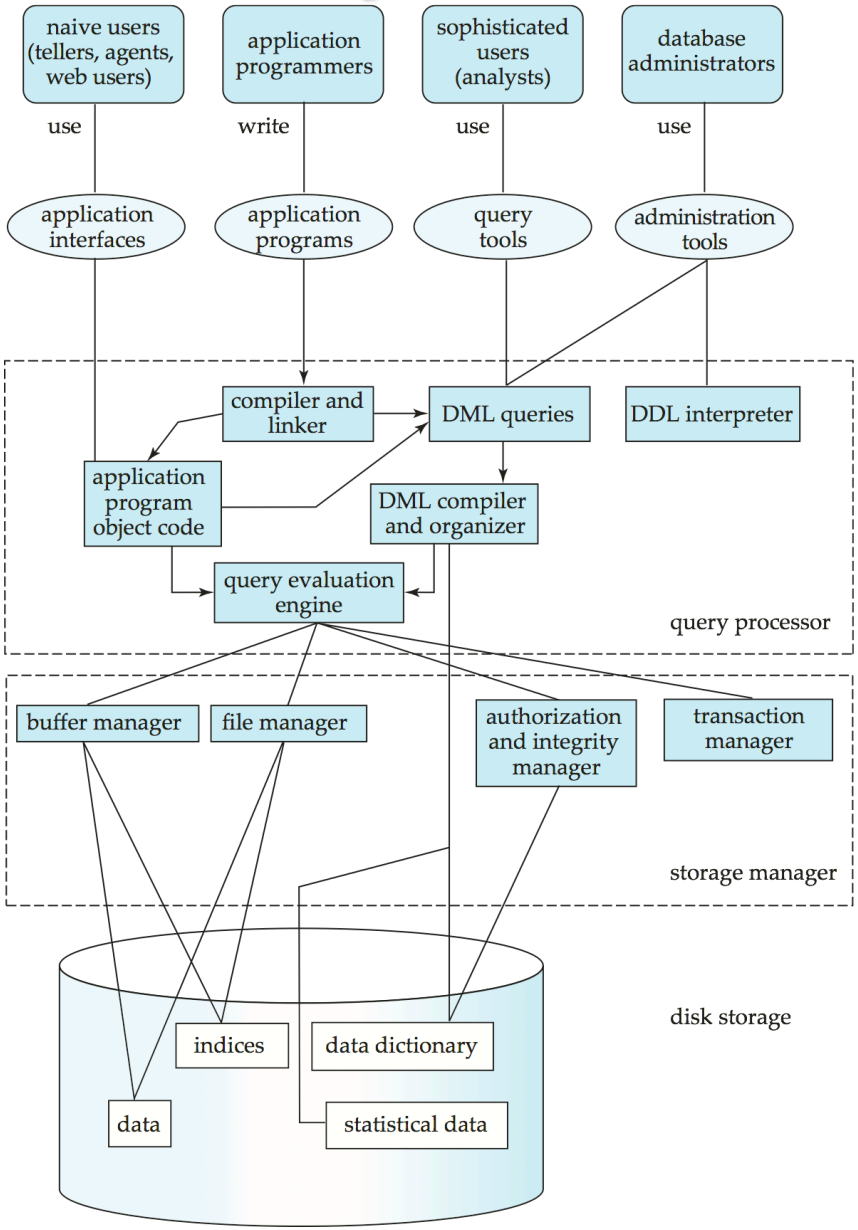


Same data as relations

lecture	location	title	teacher
$1	Girona	Foundations of XML	$3
$2	Girona	Term Rewriting	$4
$2	Girona	Term Rewriting	$5

teacher	fname	lname
$3	Joachim	Niehren
$4	Mateu	Villaret
$5	Miquel	Bofill

Database System Internals

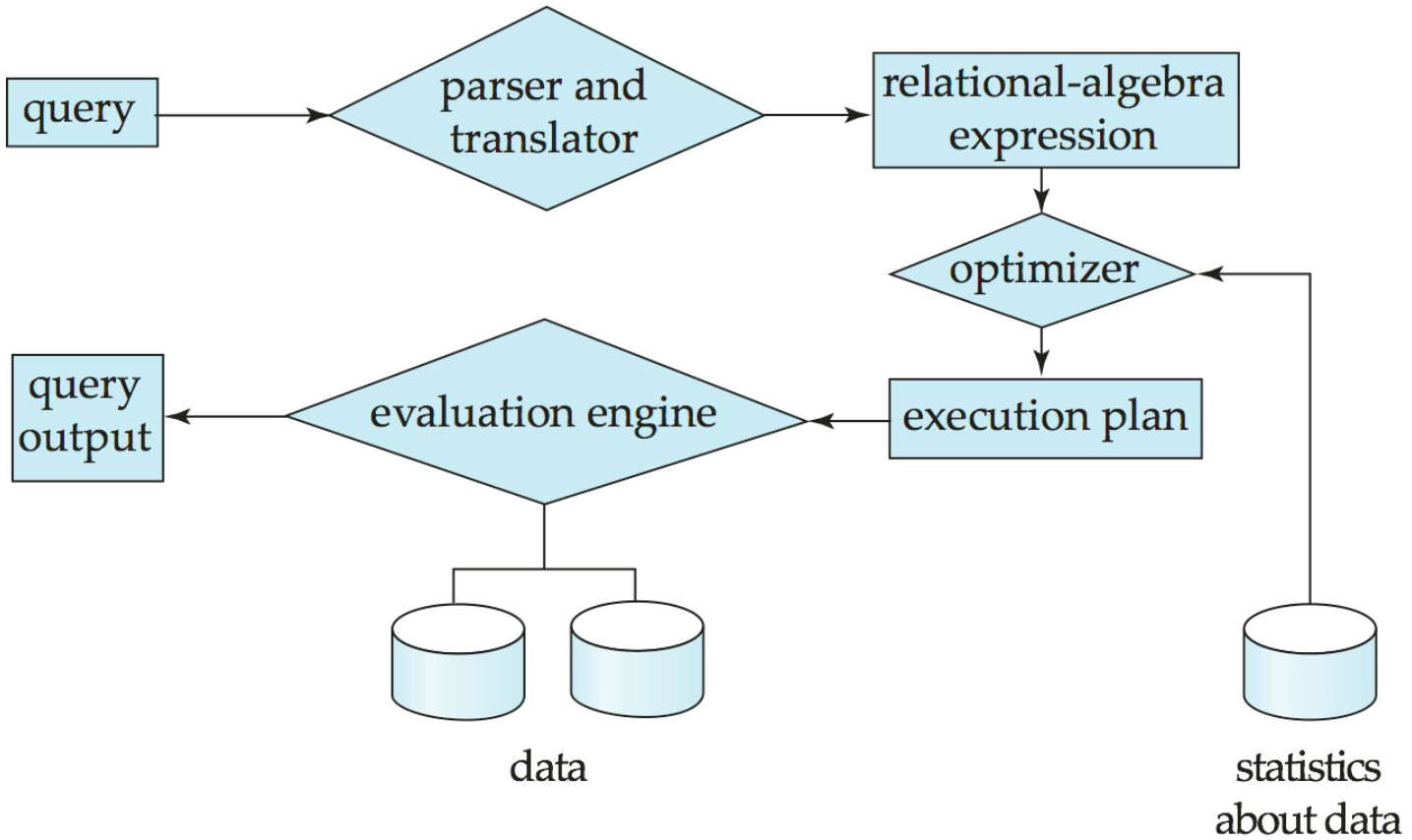


Storage Management

- **Storage manager** is a program module that provides the interface between the low-level data stored in the database and the application programs and queries submitted to the system.
- The storage manager is responsible to the following tasks:
  - Interaction with the file manager
  - Efficient storing, retrieving and updating of data
- Issues:
  - Storage access
  - File organization
  - **Indexing and hashing**

Query Processing

1. Parsing and translation
2. **Optimization**
3. Evaluation



Query Processing (Cont.)

- Alternative ways of evaluating a given SQL query
  - Equivalent expressions in relational algebra
  - Different algorithms for each operation
- Cost difference between a good and a bad way of evaluating a query can be enormous
- Need to estimate the cost of operations
  - Depends critically on statistical information about relations which the database must maintain
  - Need to estimate statistics for intermediate results to compute cost of complex expressions

Transaction Management

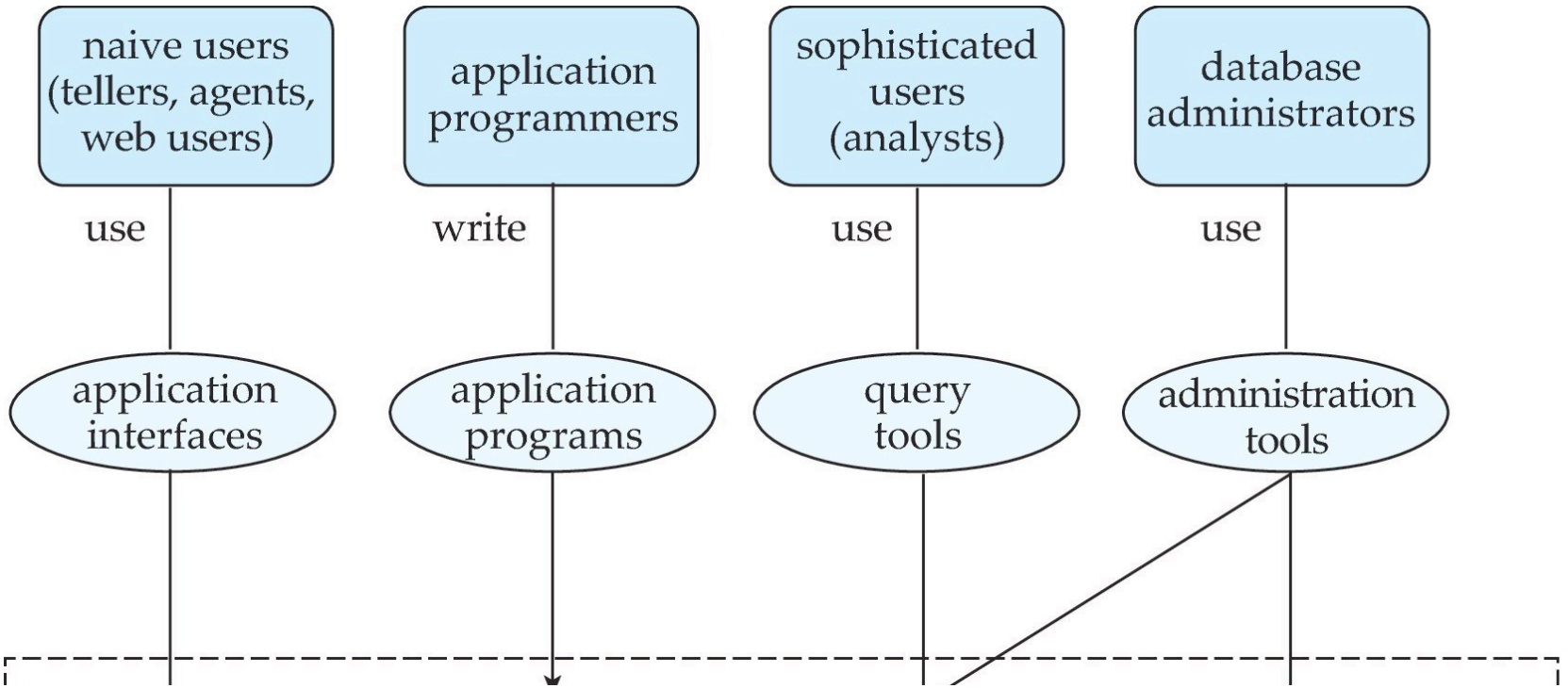
- What if the system fails?
- What if more than one user is concurrently updating the same data?
- A **transaction** is a collection of operations that performs a single logical function in a database application
- **Transaction-management component** ensures that the database remains in a consistent (correct) state despite system failures (e.g., power failures and operating system crashes) and transaction failures.
- **Concurrency-control manager** controls the interaction among the concurrent transactions, to ensure the consistency of the database.

Database Architecture

The architecture of a database systems is greatly influenced by the underlying computer system on which the database is running:

- Centralized
- Client-server
- Parallel (multi-processor)
- Distributed

Database Users and Administrators



History of Database Systems

- 1950s and early 1960s:
  - Data processing using magnetic tapes for storage
    - Tapes provided only sequential access
  - Punched cards for input
- Late 1960s and **1970s**:
  - Hard disks allowed direct access to data
  - Network and hierarchical data models in widespread use
  - Ted Codd defines the **relational data model**
    - Would win the **ACM Turing Award** for this work
    - IBM Research begins System R prototype
    - UC Berkeley begins Ingres prototype
  - High-performance (for the era) transaction processing

History (cont.)

- 1980s:
  - Research relational prototypes evolve into commercial systems
    - SQL becomes industrial standard
  - Parallel and distributed database systems
  - Object-oriented database systems
- 1990s:
  - Large decision support and data-mining applications
  - Large multi-terabyte data warehouses
  - Emergence of Web commerce
- Early 2000s:
  - **XML** and **XQuery** standards
  - Automated database administration
- Later 2000s:
  - Giant data storage systems

End of Lecture 1

Schema diagram for university database