

# Homework 5

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## Problem 1

Let we build a CFG  $G = (V, T, P, S)$  for each lowercase letter is in  $T$ , and set of  $P$  has each element below:

$$S \longrightarrow aS_1dd$$

$$S_1 \longrightarrow aS_1dd \mid S_2$$

$$S_2 \longrightarrow bS_3 \mid S_4c$$

$$S_3 \longrightarrow bS_3 \mid bS_3c \mid \varepsilon$$

$$S_4 \longrightarrow S_4c \mid bS_4c \mid \varepsilon$$

## Problem 2

Let we build a CFG  $G = (P_t, A, P, E)$ . for  $P_t$  is the set of symbol which we would be used in  $P$ ;  $T$  is the set of symbol which would appear in arithmetic expressions for primary school students, so  $A = \{x, y, z, 0, 1, 2, +, -, *, \div, (, )\}$ . And we define some basic elements of  $P$  below:

$$V \longrightarrow x \mid y \mid z$$

$$C \longrightarrow 0 \mid 1 \mid 2$$

$$L \longrightarrow V \mid C$$

$$T_s \longrightarrow + \mid -$$

$$F_s \longrightarrow * \mid \div$$

And then, we add other complicated elements to  $P$  below:

$$E \longrightarrow T \mid F$$

$$T \longrightarrow TT_sT \mid TT_sF \mid FT_sT \mid FT_sF$$

$$F \longrightarrow (T)F_s(T) \mid (T)F_sF \mid FF_s(T) \mid FF_sF \mid L$$

Notably, we can generate arithmetic expressions without redundant parentheses by using this CFG.

### Problem 3

Let we write a CFG  $G = (V, \{a, b, c\}, P, S)$  to solve this problem.

In the beginning, we define some basic elements of  $P$  below:

$$W \longrightarrow a \mid b$$

$$T_2 \longrightarrow WT_2 \mid \varepsilon$$

$$T_3 \longrightarrow WT_3 \mid cT_3 \mid \varepsilon$$

And then, we consider the situation that each string could be legit without just two symbol of 'c' called  $G_1$ . So, the elements of  $P$  are inserted below:

$$S_1 \longrightarrow WS_1 \mid cF_1 \mid \varepsilon$$

$$F_1 \longrightarrow WF_1 \mid cF_2 \mid \varepsilon$$

$$F_2 \longrightarrow WF_2 \mid cF_3$$

$$F_3 \longrightarrow WF_3 \mid cF_3 \mid \varepsilon$$

Next, we consider the situation of each string with just two symbol of c. We suppose that a string  $w_1cw_2cw_3(w_i \in \{a, b\}^*)$ . Then we consider that  $G_2$  is allowable if  $\text{len}(w_1) \neq \text{len}(w_2)$  or  $\text{len}(w_2) \neq \text{len}(w_3)$ (the function  $\text{len}(w)$  will return the length of  $w$ ). So, we add other elements for  $G_2$  to  $P$  below:

$$S_2 \longrightarrow WPC_2T_2 \mid QWC_2T_2 \mid T_2cWP \mid T_2cQW$$

$$P \longrightarrow WP \mid WPW \mid c$$

$$Q \longrightarrow QW \mid WQW \mid c$$

However, we should also consider that it exists integer  $i$  such that the  $i$ -th position of  $w_1$  and  $w_2$  is different whether  $\text{len}(w_1) = \text{len}(w_2)$ , and it is similar to the situation of  $\text{len}(w_2) = \text{len}(w_3)$ . And we call it  $G_3$ , so we add other elements for  $G_3$  to  $P$  below:

$$S_3 \longrightarrow H_a b T_3 \mid H_b a T_3 \mid T_2 c H_a b T_3 \mid T_2 c H_b a T_3$$

$$H_a \longrightarrow WH_a W \mid a T_2 c$$

$$H_b \longrightarrow WH_b W \mid b T_2 c$$

Finally, we have comprehend that the union operation of CFG would also receive a CFG. So we let the following element in  $P$ :

$$S \longrightarrow S_1 \mid S_2 \mid S_3$$

So, we did it.