

# Homework 4

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## Problem 1

We consider a basic homomorphism  $h$ :

$$\begin{cases} h(0) = & 0 \\ h(\hat{0}) = & 001 \\ h(1) = & 1 \end{cases}$$

and we can get  $h^{-1}(R)$  which is similar of the answer.  $h^{-1}(R)$  is regular.

But, the element of regular set  $h^{-1}(R)$  should not own substring 001. Because this substring 001 should be  $\hat{0}$ . So, we minus these wrong answer. Therefore, we have

$$h^{-1}(R) - (\hat{0} + 0 + 1)^*001(\hat{0} + 0 + 1)^* = R_A$$

It's easy to realize that  $R_A$  is regular.

Finally, we define a simple homomorphism  $g$ :

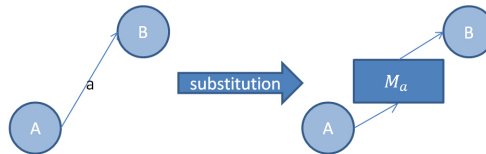
$$\begin{cases} g(0) = & 0 \\ g(\hat{0}) = & 0 \\ g(1) = & 1 \end{cases}$$

Then, the answer :  $R_B = g(R_A)$  is what we want. Obviously,  $R_B$  is regular.

## Problem 2

We define  $R$  is the initial regular sets. And  $R$  can be written to a  $\varepsilon$ -NFA  $M$  directly which means there are at least one state in  $M$  represent for each symbol in  $R$ .

And thinking about the substitution  $s(a) = R_a$ . We define  $M_a$  is the  $\varepsilon$ -NFA of  $R_a$ . So, if we change  $\delta(x, a)$  in  $M$  to pass  $M_a$ . Which means if we need via  $\delta(x, a)$ , we become to pass  $R_a$  after substitution. Here are a picture below to show this idea.



And we can also do similar operation of  $s(b) = R_b$ . So, the new FA named  $M_s$  is also a  $\varepsilon$ -NFA. And the regular sets  $R_s$  of  $M_s$  is also regular.

Therefore, we have been proved that regular sets are closed under substitution of regular sets.

### Problem 3

We use pumping lemma to proof they are not regular.

1.  $L = \{a^i b^j c^k \mid \text{either } i = j \text{ or } i = k \text{ or } j = k\}$

$\forall n \geq 1$ , we found  $w = a^n b^n c^{n-1}$ . And because of  $w = xyz \mid y \neq \varepsilon, |xy| \leq n$ . So,  $y$  would be  $a^+$ . We define that  $y = a^t (t > 0)$ .

So, there exists  $\forall k > 1$  to make  $xy^k z = a^{n+(k-1)t} b^n c^{n-1} \notin L$ .

2.  $L = \{(0+1)^n 1^n \mid n \geq 1\}$

$\forall n \geq 1$ , we found  $w = 1^n 0^n 1^{2n}$ . And because of  $w = xyz \mid y \neq \varepsilon, |xy| \leq n$ . So,  $y$  would be  $1^+$ . We define that  $y = 1^t (t > 0)$ .

So, there exists  $\forall k > 1$  to make  $xy^k z = 1^{n+(k-1)t} 0^n 1^{2n} \notin L$ .

(we can't maintain " $1^n$ " in  $k > 1$ )

3.  $L = \{(0)^i 1^j \mid i \text{ and } j \text{ are relatively prime}\}$

$\forall n \geq 1$ , we found  $w = 0^n 1^m$ . which  $m$  is the minimum prime number of  $m > n$ . And because of  $w = xyz \mid y \neq \varepsilon, |xy| \leq n$ . So,  $y$  would be  $0^+$ . We define that  $y = 0^t (t > 0)$ .

So, there exists  $k > 1$  to make  $xy^k z = 0^{n+(k-1)t} 1^m \notin L$ . Here is the reason:

Because  $m$  is prime number. So, we found that

$$(k-1)t \equiv m - n \pmod{m}$$

Obviously,  $k$  must be exist.