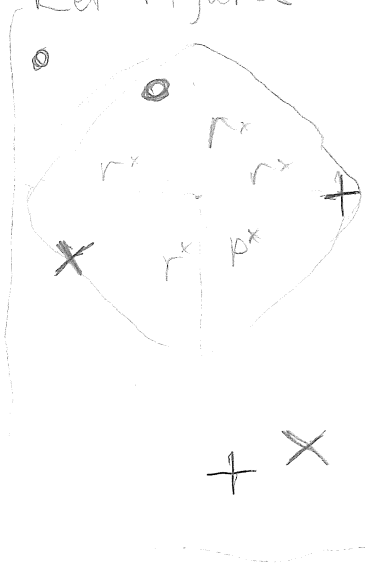


Let figures $+$, \times , \circ be different colors, " $G=X, C_2=+$, ..."



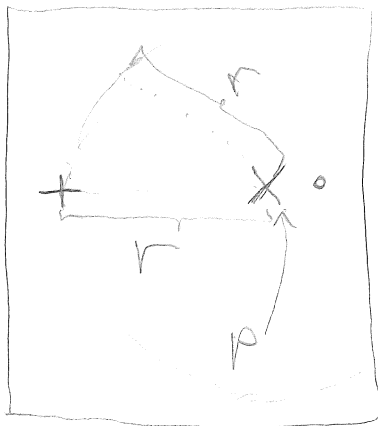
Just an example of how it can look like. (Not part of proof)

p^* is optimal point with radius r^*

Note, p^* need not be a point among the set of colored points as in the figure.

now fixate p to be the colored point that the greedy algorithm found gives minimum radius. r is

the best radius as described by greedy.



The approximation ratio, $\frac{r}{r^*}$, should be shown to be ≤ 2 .

(2/4)

Lets show $\frac{r}{r^*} \leq 2$.

let $p_1 \dots p_k$ be k color-diverse points
that the optimal solution covers.

p_i has color i ,

there might be multiple points
that the optimal solution covers
with the same color i , then

p_i can be any such point.

We also know that at least one
point with color i is covered

by the optimal solution,

since every solution must
do that.

(3/4)

the optimal solutionSince sol covered each p_i :

$$\forall i. \left| \overrightarrow{p^* p_i} \right| \leq r^*$$

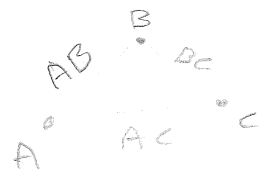
And now I claim that for any
fixated p_i (i constant):

$$\forall j. \left| \overrightarrow{p_i p_j} \right| \leq 2r^* \quad (\text{claim, to be shown})$$

Being on the plane means

triangle inequality holds. ($|\overrightarrow{AC}| \leq |\overrightarrow{AB}| + |\overrightarrow{BC}|$)

We can see p^*
as an intermediate point



$$\overrightarrow{p_i p_j} = \overrightarrow{p_i p^*} + \overrightarrow{p^* p_j} \quad \text{thus}$$

$$\forall j. \left| \overrightarrow{p_i p_j} \right| \stackrel{\text{T.I.}}{\leq} \left| \overrightarrow{p_i p^*} \right| + \left| \overrightarrow{p^* p_j} \right| \leq r^* + r^* = 2r^*$$

our claim is proven.

(4/4)

The greedy algorithm tries $p = p_1$, since in fact it tries with all colored points.

The attempt $p = p_1$

will not pick worse than picking

$p_2 \dots p_k$. And that have covered all colors.

More Formally, distance for p and picked point of color c is less or equal than $|\overrightarrow{pp_c}|$. (greedy assures this)

So that means

$$r \leq \max_j |\overrightarrow{pp_j}|$$

And since $\forall_j: |\overrightarrow{p_i p_j}| \leq 2r^*$ we get the desired

$$r \leq 2r^* \Rightarrow r/r^* \leq 2 \quad \square$$