

Computer Arithmetic: Part 3 (Division Algorithms)



Division

- ◆ Implemented on computer systems by repeatedly subtracting the divisor from the dividend
- ◆ Counting the number of times that the divisor can be subtracted from the dividend before the dividend becomes smaller than the divisor
- ◆ Example: Division of 21 by 7
 - ◆ Subtract repeatedly from 21, getting 14, 7, and 0 as intermediate results
 - ◆ The quotient, 3, is the number of subtractions that had to be performed before the intermediate result became less than the dividend

TERMS to know

◆ Terms	:	Dividend	:	7
		Divisor	:	÷ 3
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		Quotient	:	2
		Remainder	:	1

Example and Analysis

◇ 0111) 10101 (011

000	

1010	
111	

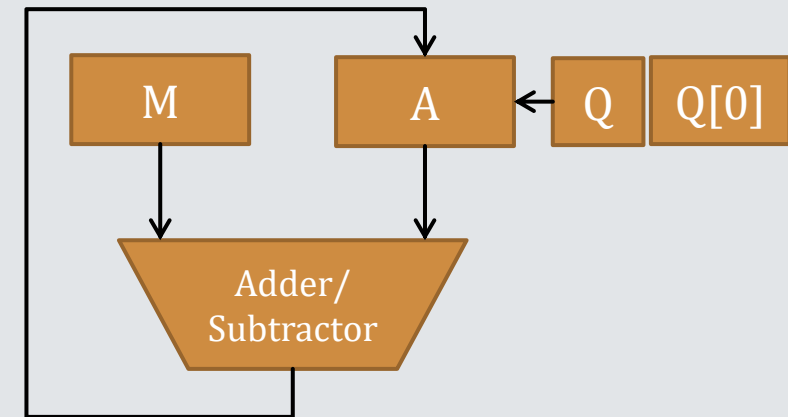
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111	
111	

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- ◇ If we divide a large number by a small number then we have to performed many subtraction
- ◇ 2^{24} divided by 2 is 2^{23} , which implies 2^{23} subtractions!

Division using restoring algorithm

- ◆ Restoring algorithm is performed on fixed point fractional numbers
- ◆ Here we are going to performed restoring algorithm of an unsigned integer
- ◆ Initially
 - ◆ M = Divisor
 - ◆ A (Accumulator) = 0
 - ◆ Q = Dividend
- ◆ After restoring division
 - ◆ Q = Quotient
 - ◆ A = Remainder



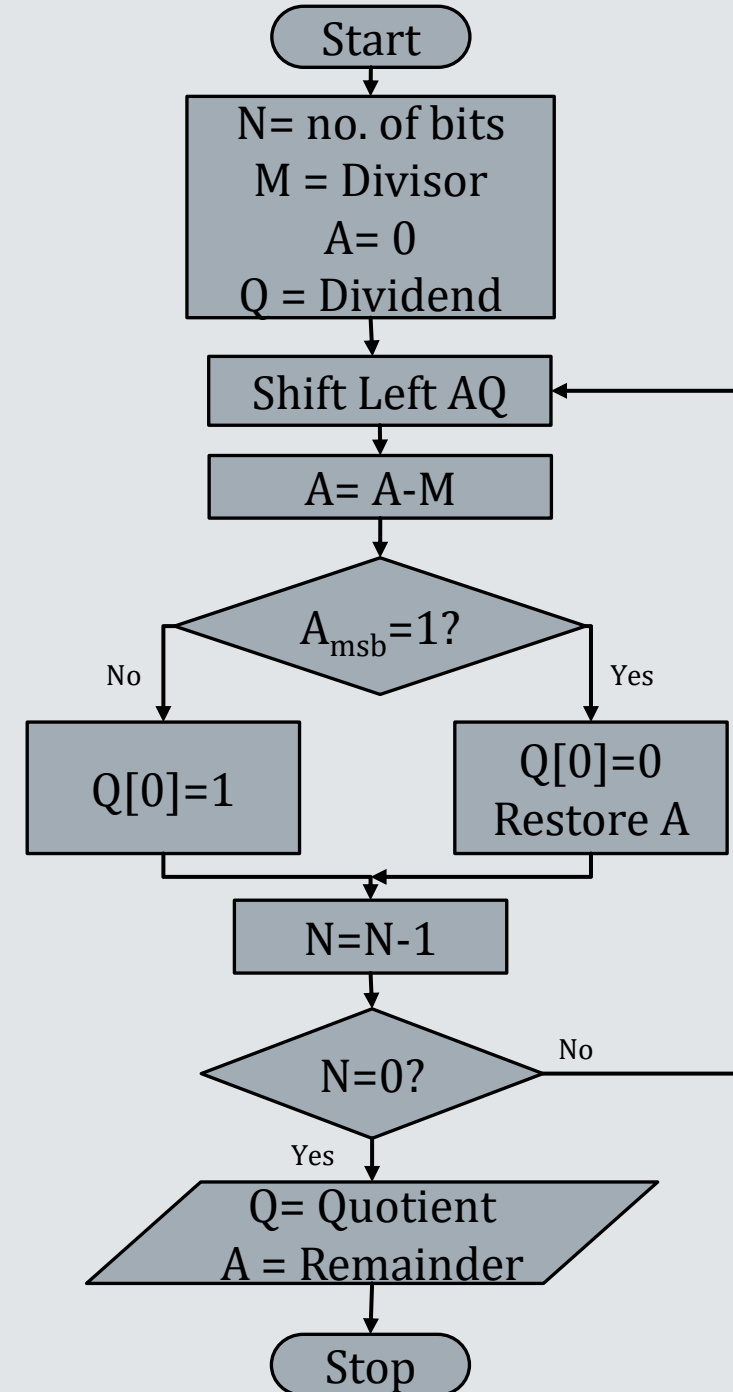
Example Restoring Algorithm

$11 \div 3$

N	A	Q	Action
4	00000 00001 11110 00001	1011 011_ 011_ 0110	Initial Configuration Shift Left AQ $A = A - M$ $Q[0] = 0$; Restore A (Since $A_{msb} = 1$)
3	00010 11111 00010	110_ 110_ 1100	Shift Left AQ $A = A - M$ $Q[0] = 0$; Restore A (Since $A_{msb} = 1$)
2	00101 00010 00010	100_ 100_ 1001	Shift Left AQ $A = A - M$ $Q[0] = 1$ (Since $A_{msb} = 0$)
1	00101 00010 00010	001_ 001_ 0011	Shift Left AQ $A = A - M$ $Q[0] = 1$ (Since $A_{msb} = 0$)

Remainder

Quotient



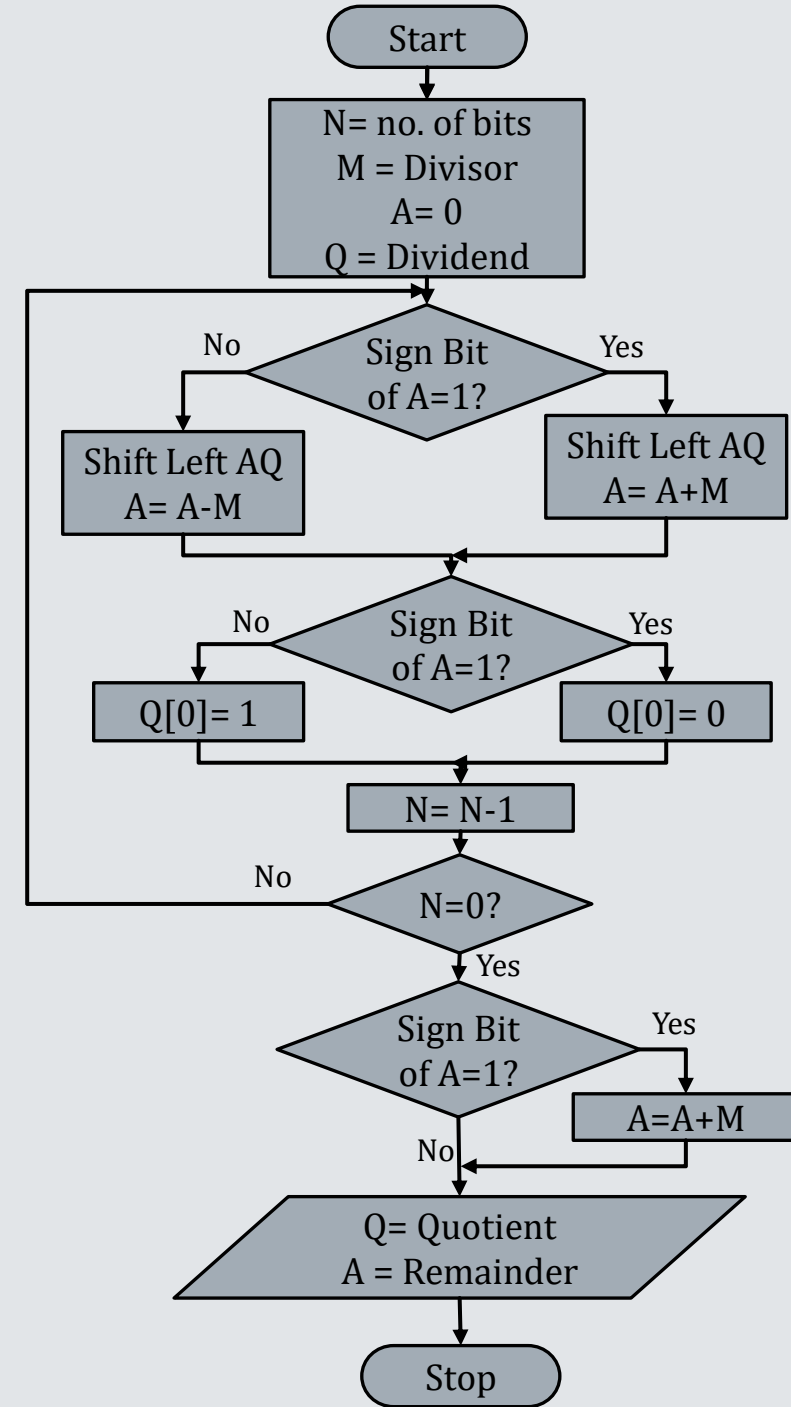
Non Restoring Algorithm

- ◆ Non-restoring division algorithm is more complex than the restoring division algorithm.
- ◆ However, it has less hardware complexity as compared to restoring algorithm.
- ◆ It has only addition/subtraction per quotient bit and no restoring is required.
- ◆ Task performed is almost half as compared to restoring division
- ◆ Faster than restoring division

Example: Non-Restoring Algorithm

$11 \div 3$

N	A	Q	Action
4	00000	1011	Initial configuration
	00001	011_	Shift Left AQ
	11110	011_	$A = A - M$
	11110	0110	$Q[0] = 0$
3	11100	110_	Shift Left AQ
	11111	110_	$A = A + M$
	11111	1100	$Q[0] = 0$
2	11111	100_	Shift Left AQ
	00010	100_	$A = A + M$
	00010	1001	$Q[0] = 1$
1	00101	001_	Shift Left AQ
	00010	001_	$A = A - M$
	00010	0011	$Q[0] = 1$
Remainder		Quotient	





Thank You