Linear Algebra HW2

1

first three pivot are 2,7,6. 因為在做高斯消去時並不會影響到原來這3個 pivots

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syms a1 a2 a3 a4 a5 a6 a7 a8 a9 a10

A = [2    a1 a2 a3;...
        0    7    a4 a5;...
        0    0    6    a6;...
        a7 a8 a9 a10];

[n,~] = size(A);

for i=1:n-1
    m = A(i+1:n,i)/A(i,i);
    A(i+1:n,:) = A(i+1:n,:) - m*A(i,:);
end
```

after calculation

```
A =

[2, a1, a2,
[0, 7, a4,
[0, 0, 6,
[0, 0, 0, a10 - a5*(a8/7 - (a1*a7)/14) + a6*((a4*(a8/7 - (a1*a7)/14))/6 - a9/6 + (a2*a7)/12) - (a3*a7)/2]
```

只要A矩陣經計算後(如上圖),右下角為9就可以確定第4個pivot為9

2

- (a) **True**, Each column of AB is the product of A and a column of B
- (b) **False**, Each column of AB is a combination of the columns of A, 所以還是要看A的係數
- (c) False, $(AB)^2 = ABAB$

(d) **True**, if A is a lower triangle matrix, then $A_{ij} = 0 \ \forall i < j$. By the matrix operation, if both A and B are lower triangle matrices, then

 $(AB)_{ij} = 0 \ orall i < j$, hence AB is a lower triangle matrix.

3

For A_1 , 會將B的第一列乘4 ,第二列乘3 ,第三列乘2 ,第四列乘1 ,再將第一列到第四列顛倒。

$$(A_1)^{-1} = egin{bmatrix} 0 & 0 & 0 & rac{1}{4} \ 0 & 0 & rac{1}{3} & 0 \ 0 & rac{1}{4} & 0 & 0 \ 1 & 0 & 0 & 0 \end{bmatrix}$$

For A_2 , 會將B的第一列乘 $-\frac{1}{2}$ 加到第二列,乘4加到第四列,將第二列乘 $-\frac{2}{3}$ 加到第三列,將第三列乘 $-\frac{3}{4}$ 加到第四列

$$(A_2)^{-1} = egin{bmatrix} 1 & 0 & 0 & 0 \ rac{1}{2} & 1 & 0 & 0 \ rac{1}{3} & rac{2}{3} & 0 & 0 \ -rac{15}{4} & rac{1}{2} & rac{3}{4} & 1 \end{bmatrix}$$

4

以下是用高斯消去法的步驟

$$A = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 2 & 3 & 4 \end{bmatrix} \Rightarrow \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 2 & 3 & 4 \end{bmatrix}, P_1 = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$A' = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 2 & 3 & 4 \end{bmatrix} \Rightarrow \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 0 & 3 & 2 \end{bmatrix}, L_1 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -2 & 0 & 1 \end{bmatrix}$$

$$A' = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 0 & 3 & 2 \end{bmatrix} \Rightarrow \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & -1 \end{bmatrix}, L_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -3 & 1 \end{bmatrix}$$

(a) L和U分別為

$$L = (L_2 * L_1)^{-1}$$
 $L = egin{bmatrix} 1 & 0 & 0 \ 0 & 1 & 0 \ 2 & 3 & 1 \end{bmatrix}$
 $U = egin{bmatrix} 1 & 0 & 1 \ 0 & 1 & 1 \ 0 & 0 & -1 \end{bmatrix}$

(b)

$$L^{-1} = L_2 L_1 = egin{bmatrix} 1 & 0 & 0 \ 0 & 1 & 0 \ -2 & -3 & 1 \end{bmatrix}$$

5

(1)

$$A = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 1 & 4 & 1 \end{bmatrix} \begin{bmatrix} 5 & 7 & 8 \\ 0 & 2 & 3 \\ 0 & 0 & 6 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 0 & 4 & 1 \end{bmatrix} \begin{bmatrix} 5 & 7 & 8 \\ 0 & 2 & 3 \\ 0 & 0 & 6 \end{bmatrix}$$

$$= L * U$$

將第2列乘4加到第3列,將第一列乘2加到第二列,乘1加到第三列

(2) no row exchange required

(3)

$$Lc = b \Rightarrow egin{bmatrix} 1 & 0 & 0 \ 2 & 1 & 0 \ 1 & 4 & 1 \end{bmatrix} egin{bmatrix} c_1 \ c_2 \ c_3 \end{bmatrix} = egin{bmatrix} 1 \ 2 \ 3 \ \end{bmatrix} \ c_1 = 1, c_2 = 0, c_3 = 2 \ Ux = c \Rightarrow egin{bmatrix} 5 & 7 & 8 \ 0 & 2 & 3 \ 0 & 0 & 6 \end{bmatrix} egin{bmatrix} x_1 \ x_2 \ x_3 \end{bmatrix} = egin{bmatrix} 1 \ 0 \ 2 \ \end{bmatrix} \ ANS: x = egin{bmatrix} 1/30 \ -1/2 \ 11/30 \end{bmatrix}$$

6

$$P = egin{bmatrix} 0 & 1 & 0 \ 1 & 0 & 0 \ 0 & 0 & 1 \end{bmatrix} \ PA = egin{bmatrix} 1 & 0 & 1 \ 0 & 1 & 1 \ 2 & 3 & 4 \end{bmatrix} \ from & the & result & in & (4) \ LDU = egin{bmatrix} 1 & 0 & 0 \ 0 & 1 & 0 \ 2 & 3 & 1 \end{bmatrix} egin{bmatrix} 1 & 0 & 0 \ 0 & 1 & 0 \ 0 & 0 & -1 \end{bmatrix} egin{bmatrix} 1 & 0 & 1 \ 0 & 1 & 1 \ 0 & 0 & 1 \end{bmatrix}$$

7

(1)Gaussian elimination and back substitution

left side: $(n^2+\ldots+1^2)-(n+\ldots+1)=\frac{n(n+1)(2n+1)}{6}=\frac{n^3-n}{3}$ right side : $(1+2+\ldots+n)+[(n-1)+(n-2)+\ldots 1]\approx n^2$ back substitution $\approx \frac{n^2}{2}$ total $\approx \frac{2n^3+3n^2}{6}$

(2)Gauss-Jordan Method (algorithm)

Reference: gauss-complexity.pdf (ryerson.ca)

For Each row i (R_i) from 1 to n

If any row j below row i has non zero entries to the right of the first non zero entry in row i

$$R_i \leftrightarrow R_j$$

 $R_i
ightarrow rac{1}{c} R_j$ where c = the first non-zero entry of row i

For each row j < i

 $R_i \leftrightarrow R_j - dR_i$ where d = the entry in row j which is directly below the pivot in row i

If any 0 rows have appeared exchange them to the bottom if the matrix.

next i

For each non zero row i (R_i) from n to 1

For each j < i

 $R_i \leftrightarrow R_j - dR_i$ where b = the value in row j directly above the pivot in row i.

the complexity is:

$$egin{aligned} \sum_{i=1}^n [(n+1) + (n-i)(n+1)] + \sum_{i=1}^n [(n-i)(n+1)] \ &= \sum_{i=1}^n \sum_{i=1}^n (2n-2i+1)(n+1) \ &= \sum_{i=1}^n 2n^2 + 3n^2 + 1 - 2(n+1)i \ &= 2n^3 + 3n^2 + n + n(n+1)^2 \ &= 3n^3 + 5n^2 + 2n \end{aligned}$$

結論

兩者比較起來Gaussian elimination and back substitution會比Gauss-Jordan Method好一些,但資料量很大時複雜度都是 $O(n^3)$

$$P^3 = I$$

將identity matrix做兩次permutation

$$P = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

$$P^4
eq I$$

將identity matrix做三次permutation

$$P = egin{bmatrix} 0 & 0 & 1 & 0 \ 1 & 0 & 0 & 0 \ 0 & 1 & 0 & 0 \ 0 & 0 & 0 & 1 \end{bmatrix}$$

9

$$PA = LDU$$
 $(PA)^{-1} = (LDU)^{-1}$ $A^{-1}P^{-1} = U^{-1}D^{-1}L^{-1}$ $A^{-1} = U^{-1}D^{-1}L^{-1}P$