

WarmHoles

1. Problem Statement

You are given a directed graph representing star systems connected by wormholes. Each wormhole goes one way, and when you travel through it, you arrive t years in the future (if $t > 0$) or t years in the past (if $t < 0$).

The travel time through a wormhole is negligible.

You start from star system 0 (our Solar System). It is guaranteed that every star system is reachable from system 0 through some sequence of wormholes.

The goal is to determine whether there exists any cycle reachable from star system 0 such that the total time change along the cycle is negative.

If such a cycle exists, then by looping through it repeatedly, one can go back in time indefinitely — making it possible to reach the Big Bang.

For each test case, print:

- "possible" — if there exists a reachable negative time cycle
- "not possible" — otherwise

2. Hint

This problem can be transformed into checking whether a negative-weight cycle exists in a directed graph that is reachable from the starting node (0).

Interpret each time constraint t as the weight of an edge between two nodes.

Which well-known shortest-path algorithm is specifically designed to detect negative cycles that can be reached from a given source node?

3. Solution Approach

The problem is a direct application of the Bellman–Ford algorithm, which detects negative cycles in a weighted directed graph.

Key observations:

- Each wormhole is a directed edge with weight = time shift.
- You start from node 0.
- You're looking for any negative cycle reachable from node 0.
- Bellman–Ford relaxes edges $(N - 1)$ times to find shortest paths.
- A further relaxation on the N th pass indicates a negative cycle.

Steps:

1. Initialize distances with $\text{dist}[0] = 0$ and all others as INF.
2. Relax all edges for $N - 1$ iterations.
3. If no update occurs in an iteration, you can stop early (optimization).
4. On the Nth iteration, if any edge can still be relaxed, a negative cycle exists.

Output:

- If negative cycle detected \rightarrow "possible"
- Otherwise \rightarrow "not possible"

4.Pseudocode:

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FUNCTION BELLMAN_FORD(G, N, M, E)

// 1. Initialization
FOR each vertex v IN V:
    dist[v] = INF // Initialize all distances to infinity
dist[0] = 0      // Set distance for the source vertex (System 0/Earth) to zero

// 2. Relaxation Passes (Run N-1 times)
// N is the number of vertices. A path can have at most N-1 edges.
FOR i FROM 1 TO N - 1:
    relaxed_in_pass = FALSE
    FOR each edge (u, v, t) IN E (where t is the time/weight):
        IF dist[u] is NOT INF:
            IF dist[u] + t < dist[v]:
                dist[v] = dist[u] + t
                relaxed_in_pass = TRUE

    // Optimization: If no distance changed in a pass, paths are stable.
    IF relaxed_in_pass is FALSE:
        BREAK

// 3. Final Check for Negative Cycle (The N-th Pass)
FOR each edge (u, v, t) IN E:
    IF dist[u] is NOT INF:
        IF dist[u] + t < dist[v]:
            // If we can still relax an edge, a negative cycle exists.
            RETURN 1 // Possible

// If no negative cycle was found after the N-th check.
RETURN 0 // Not Possible

```

5. Time Complexity Analysis :

The time complexity of the Bellman-Ford algorithm is determined by the number of times we iterate through the edges.

1. Initialization: Setting all N distances takes $O(N)$ time.

2. Relaxation Passes: This is the main performance factor.

The outer loop runs a maximum of $N-1$ times.

The inner loop iterates over all M edges in the graph. O

Total time for N-1 passes: $O((N-1) \cdot M)$, which simplifies to $O(N \cdot M)$.

3. Final Check: The cycle check iterates over all M edges one last time. This takes $O(M)$ time.

Implementation Link:

https://github.com/TasnimaSultana/Algo_Lab_Final/blob/main/Bellman%20Ford/Warmhole/UVA558_Warmholesproblem.cpp