

1 TO C

1 T 1. $\Sigma = \{0, 1, a\}$ and Σ^*

symbol \rightarrow Letter & digits

Alphabet \rightarrow Finite set of symbols

$$\Sigma^* = \{0, 1, a\}^*$$

string \rightarrow Finite sequence of symbols

finite & countable

$$\Sigma = \{0, 1\}$$

$$w_1 = 010, \quad w_2 = 110 \dots$$

\emptyset is empty string

1 T 2

$$\{\emptyset, \{a\}\} \subseteq \mathcal{P}$$

language \rightarrow finite collection of strings (Finite/Infinite)

By $\Sigma = \{a, b\}$

$L_1 = \text{set of all strings over } \Sigma \text{ of length 2}$
 $= \{aa, ab, ba, bb\}$

ϵ^* means $\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots$

| T4 Into to ~~DFA~~ DFA

Finite Automata Automata

FA with output

Minsky Mealy

FA without output

DFA

NFA

GNFA

Ex

$$\Sigma = \{a, b\}$$

L_1 = set of all strings start with an 'a'.

$$(Q, \Sigma, \delta, q_0, F)$$

Focal P.

Q = set of all states



Σ = input alphabets

q_0 = start state

$$\{a, b\} = \Sigma$$

F = set of final states

$$\{q_1\} = F$$

δ = transition function

$$A = \delta$$

$$\delta \leftarrow Q \times \Sigma \rightarrow Q$$

$$\delta = Q \times \Sigma \rightarrow Q$$

A state gets input and ~~other goes to~~
which state determined by transition
state.

| TS Formal definition of DFA |

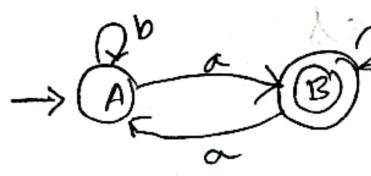
DFA has 5 tuples, $Q, \Sigma, q_0, F, \delta$

is it same with AFD or not

Ans answer $\{a, b\} = \Sigma$ sent more zero into

'a' file note write

Ex



$$Q = \{ A, B \}$$

$$\Sigma = \{ a, b \}$$

$$q_0 = A$$

$$\delta: Q \times \Sigma \rightarrow Q$$

A	a	b
B		A
B	A	B

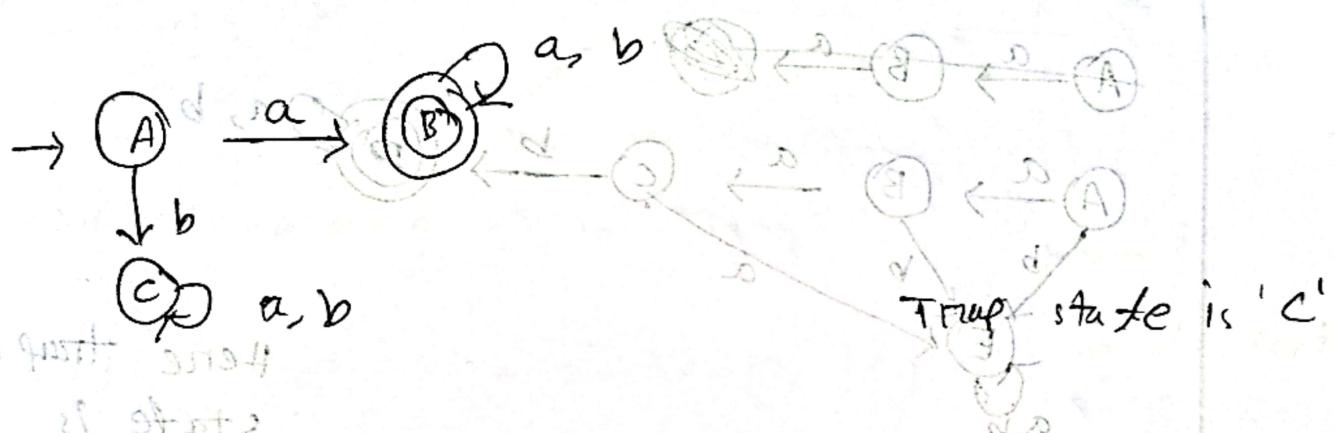
| T6, Ex 1 | Exercise on DFA | PF |

Ex Construct a DFA which accepts set of all strings over the $\Sigma = \{a, b\}$ where each string starts with 'a'

on

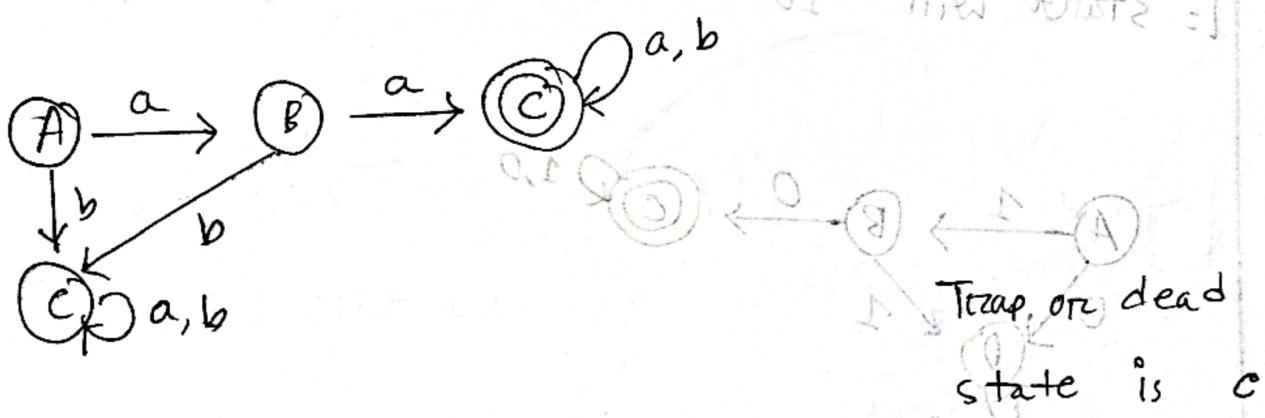
Q4 PT 1
 $L(M) = \{ w \mid w \text{ starts with an 'a'} \}$

Sol $\Sigma = \{ a, b \}$



| T7

$L = \text{set of all strings that start with 'aa'}$

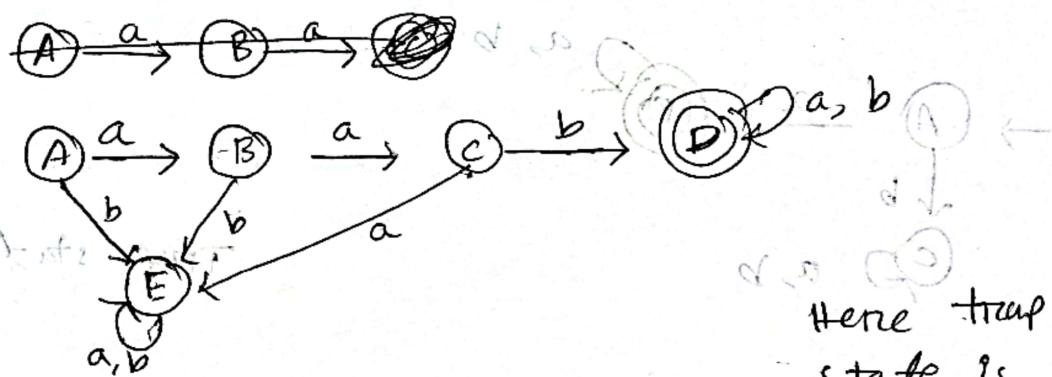


| T8

DFA

Ex. Find a DFA which accepts L = {aab}*

L = set of all strings over $\Sigma = \{a, b\}$ where each string starts with "aab"



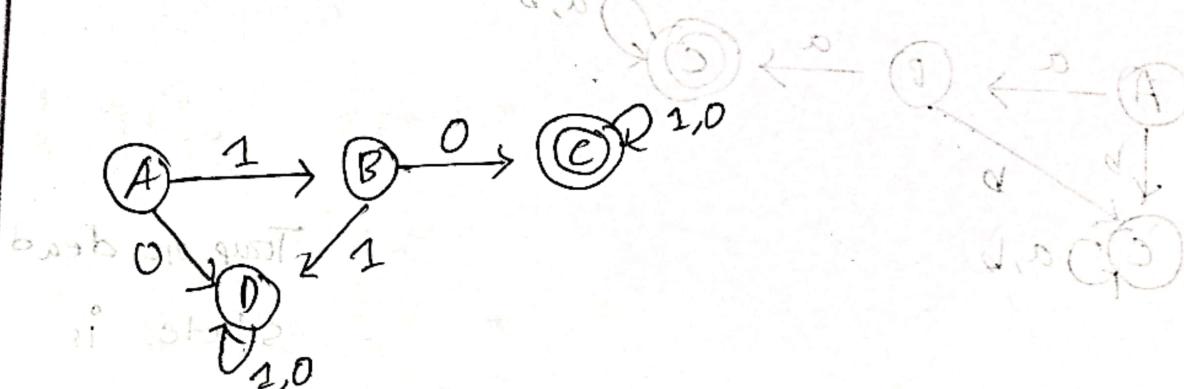
Here trap or dead state is E.

| T9

DFA

L = start will '10'

$$\Sigma = \{0, 1\}$$



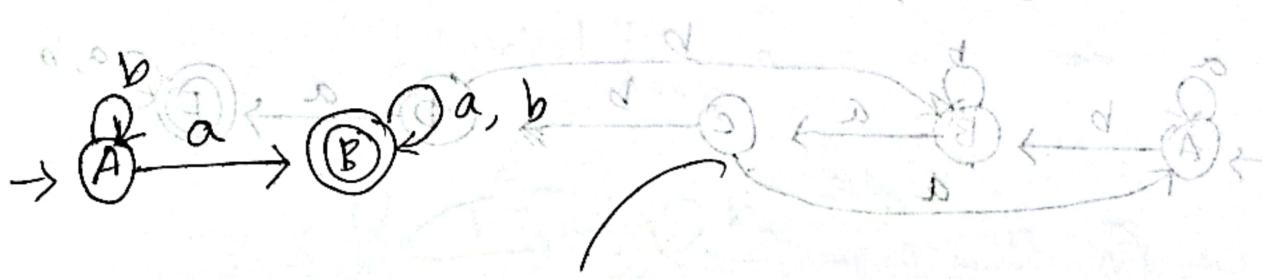
| T 10

DFA

AUT

2*+1

L_2 : set of all strings over $\Sigma = \{a, b\}$ in which every string contains 'a'



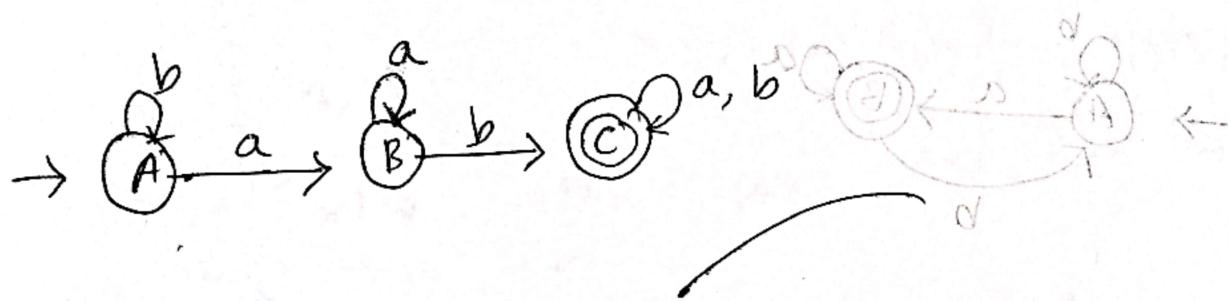
| T 11

DFA

AUT

ACT

L : set of all strings over $\{a, b\}$ which contains "ab" ('a' until now $w/w\}$) $\in \Sigma^*$



| T 12

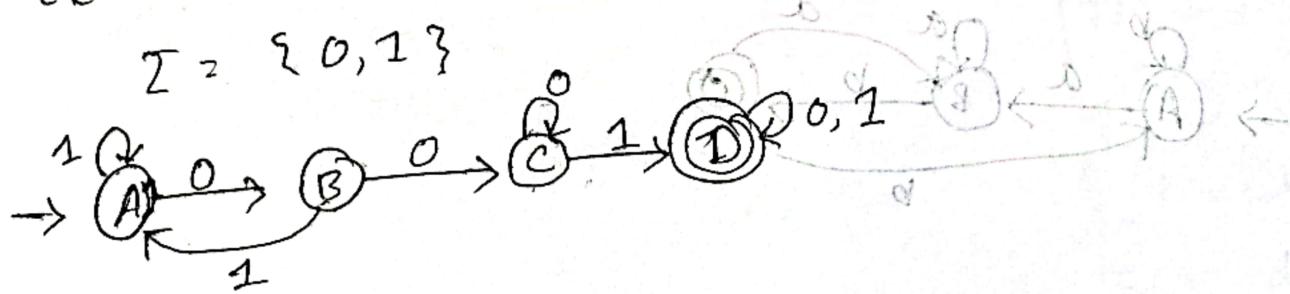
DFA

AUT

2*+1

$L(M) = \{ w | w \text{ contains the string } 001 \text{ as a substring} \}$

$$\Sigma = \{0, 1\}$$



T13

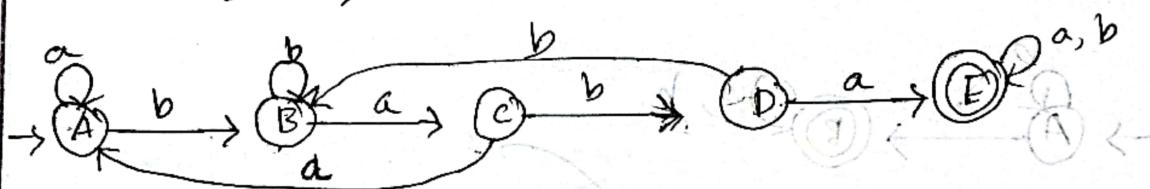
DFA

AFN

NFT

$L(M) = \{ w \mid w \text{ contains the substring } ba'b \}$

$$\Sigma = \{a, b\}$$



T14

DFA

AFN

NFT

$\Sigma = \{a, b\}$ no exists $w \in L(M)$ to a

$L(M) = \{ w \mid w \text{ ends with a } 'a' \}$



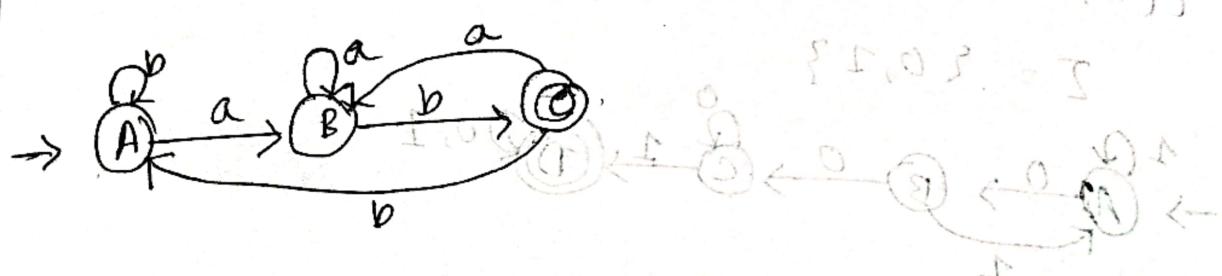
T145

DFA

AFN

NFT

$\Sigma = \{a, b\}$, $L(M) = \{ w \mid w \text{ end with } 'ab' \}$

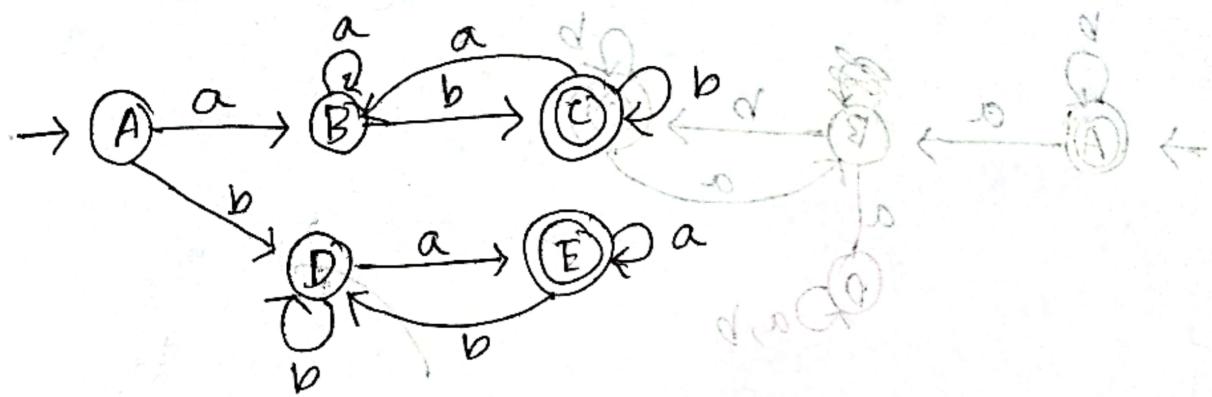


| T 16

DFA

A70

QDF

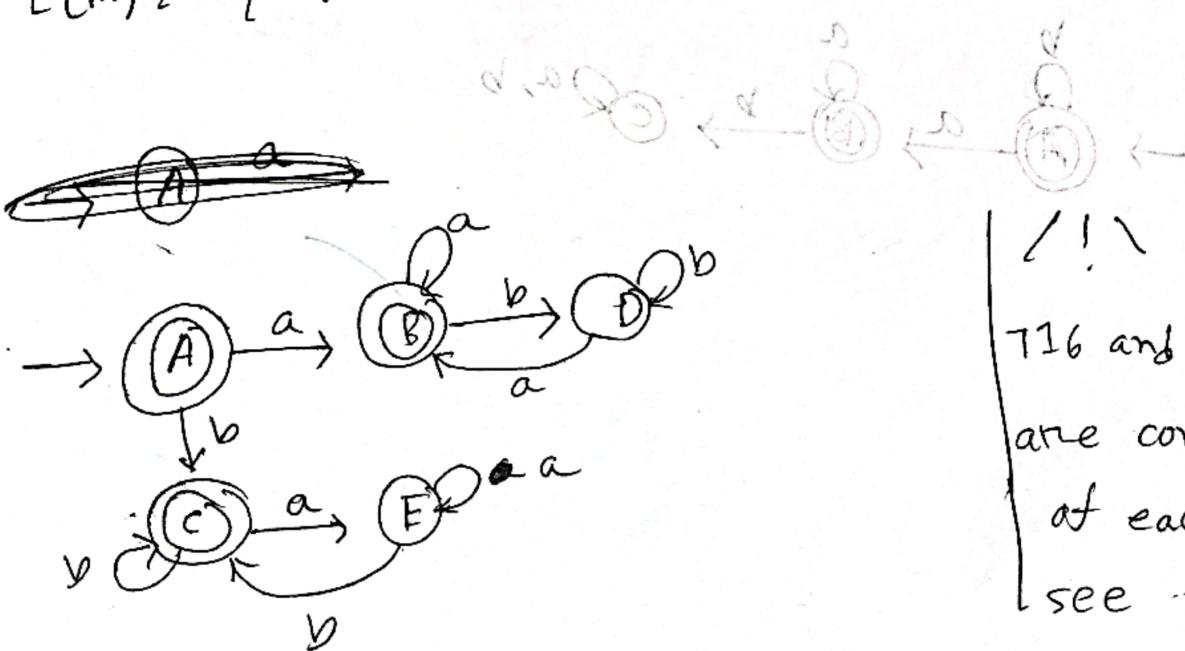
 $\Sigma = \{a, b\}$, $L(M) = \{w \mid w \text{ starts and ends with different symbols}\}$
 $L(M) = \{w \mid w \text{ starts and ends with different symbols}\}$


| T 17

DFA

A70

QDF

 $\Sigma = \{a, b\}$, $L(M) = \{w \mid w \text{ starts and ends with same symbol}\}$
 $L(M) = \{w \mid w \text{ starts and ends with same symbol}\}$


11

T16 and T17

are complement
of each other.
see deeply

T 28

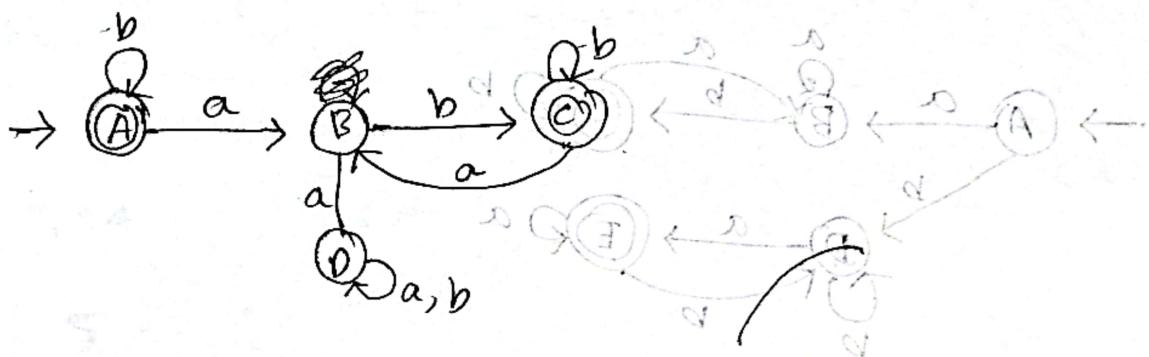
DFA

A/T

E/F

construct a DFA, $\Sigma = \{a, b\}$ and $A = \mathbb{Z}$

$L(M) = \{w \mid w \text{ every } a \text{ in } w \text{ is followed by one } 'b'\}$



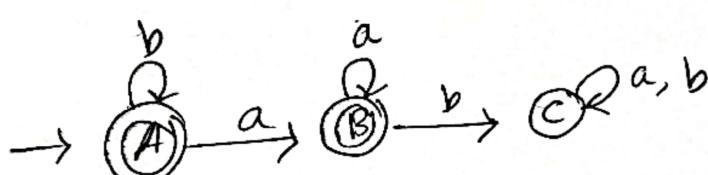
T 29

DFA

A/T

E/F

$L(M) = \{w \mid w \text{ every } a \text{ in } w \text{ is never followed by one } 'b'\}$



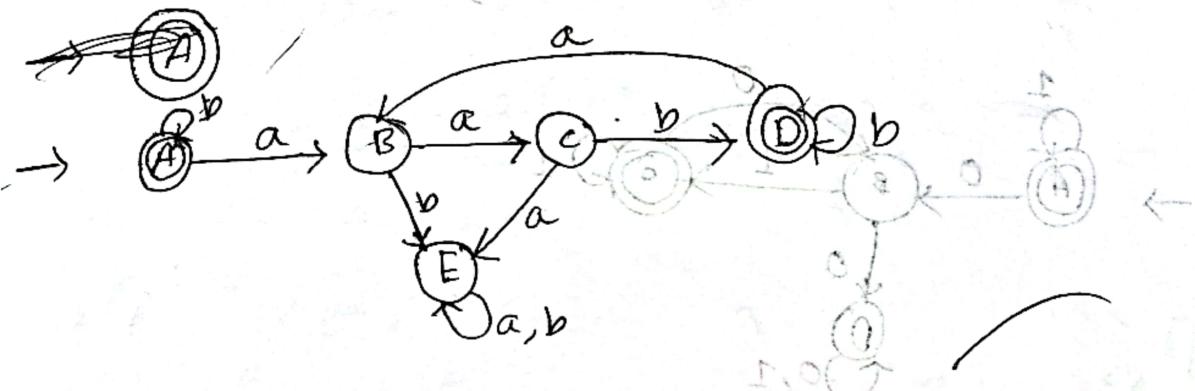
T20

DFA

ATQ

55

$L(M) = \{w \mid w \text{ - every 'a' in } w \text{ is followed by 'ab'}\}$

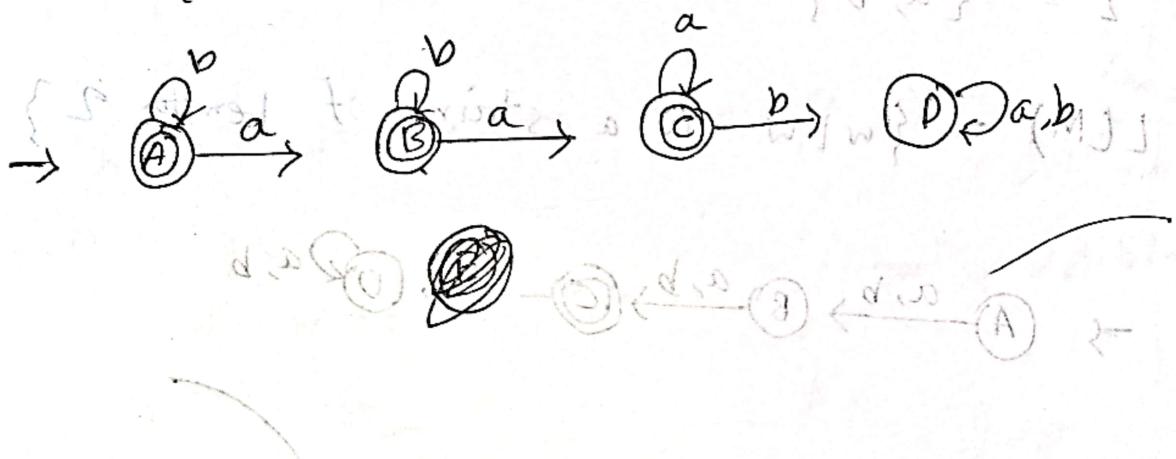


T21

DFA

ATQ

$L(M) = \{w \mid w \text{ every 'a' in } w \text{ is never followed by 'ab'}\}$



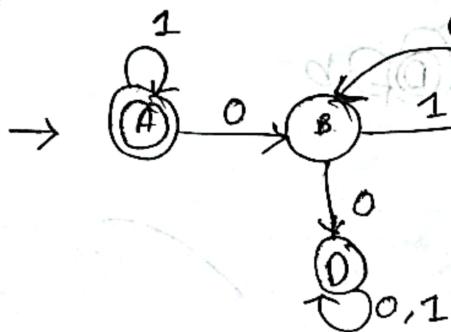
T 22

DFA

A/A

DST }

$L(M) = \{w \mid w \in \text{every '0' in } w \text{ is followed by at least one '1'}\}$



T 23

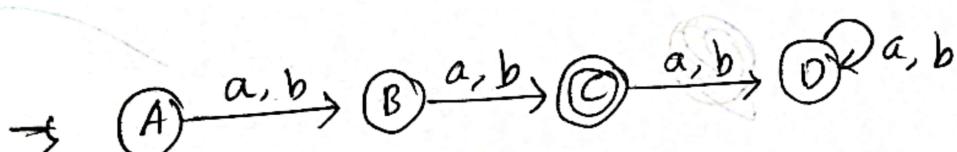
DFA

A/A

DST }

$$\Sigma = \{a, b\}$$

$L(M) = \{w \mid w \text{ is a string of length 2}\}$

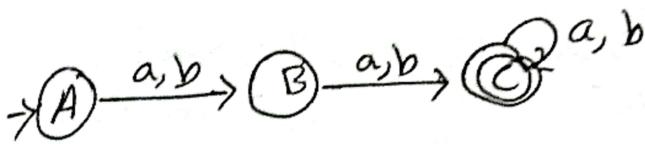


| T 24

DFA

$L(M) = \text{strings of length at least } 2$

$$\Sigma = \{a, b\}$$



| T 25

DFA

$L(M) = \text{strings of length at most } 2$. $\{a, b\}$



| T 26 Intco to minimal DFA (mDFA)

- Minimize as much as possible.



T-27

DFA

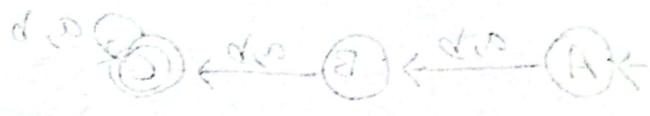
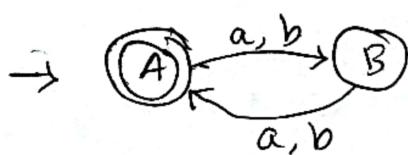
A.F.Q.

P.S.T.

Length of string is divisible by 2 i.e. (0)

$$\Sigma = \{a, b\}$$

$$\{0, 1\} \cdot \overline{1}$$



T-28

DFA

A.F.Q.

P.S.T.

Length of string is not divisible by 2



(A1)(A2)

T-29

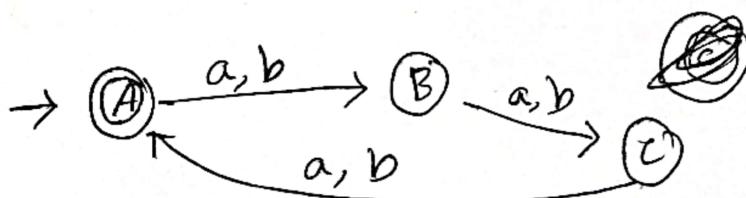
DFA

im

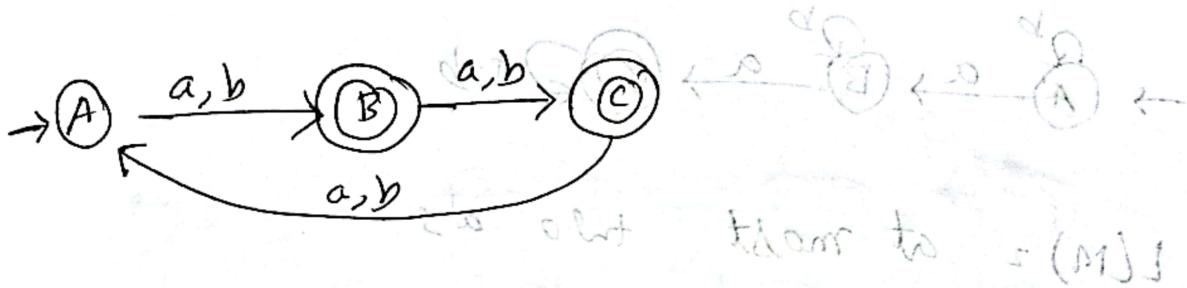
P.S.T.

Length of string divisible by 3

Length of string to be known by scanning it



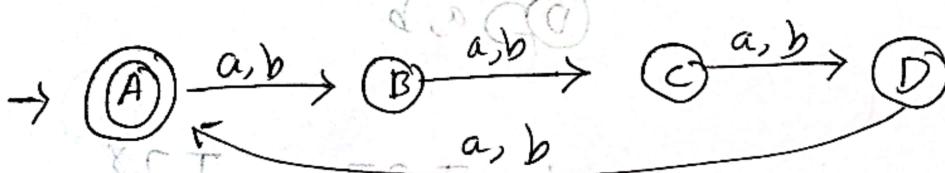
Not divisible by 3 part of string length



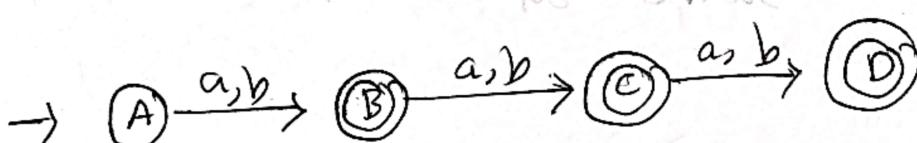
| T 30

DFA

string length \Rightarrow divisible by 4



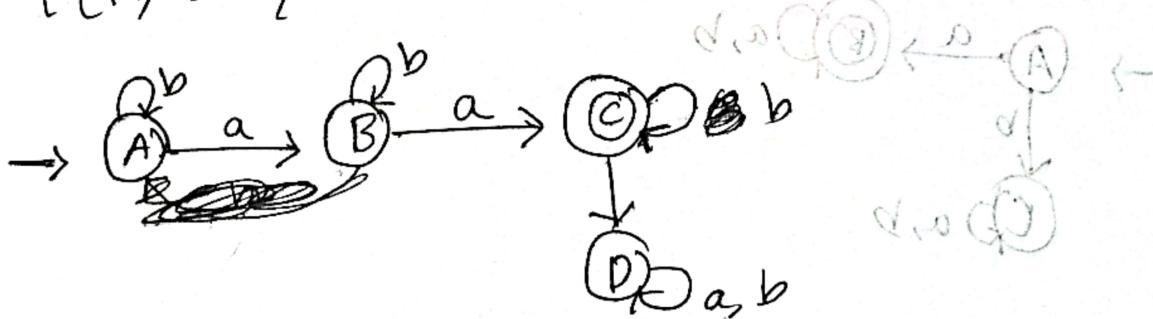
string length not divisible by 4



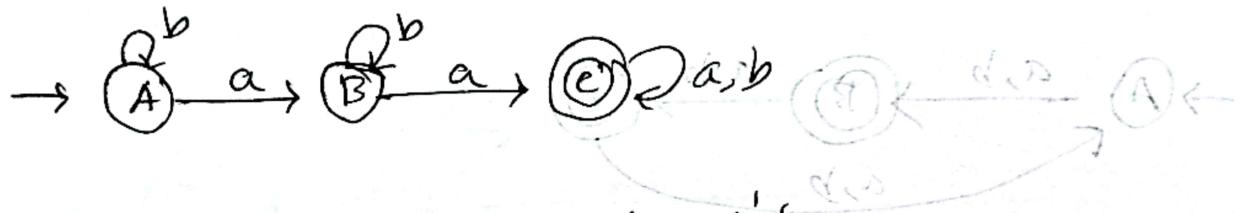
| T 31

DFA

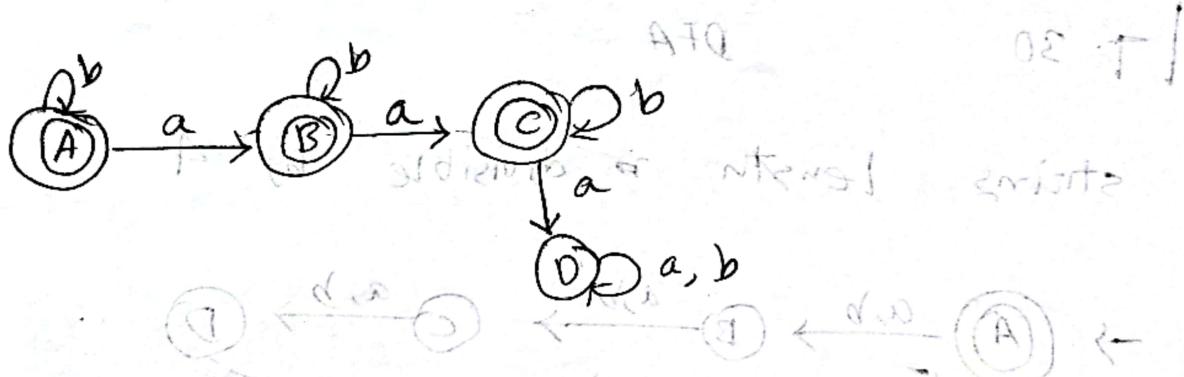
$L(M) = \{w \mid w \text{ has exactly two } a's\}$



$L(M)$ = at least two a's followed by b - 1



$L(M)$ = at most two a's

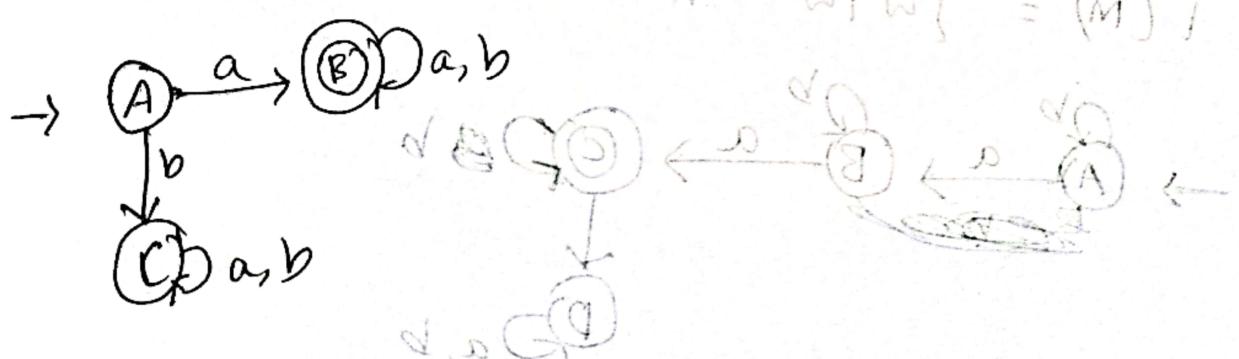


T 32 same as T 27, T 28

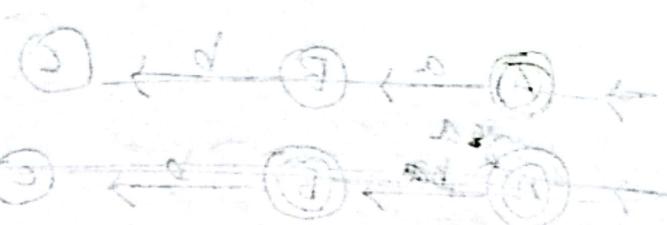
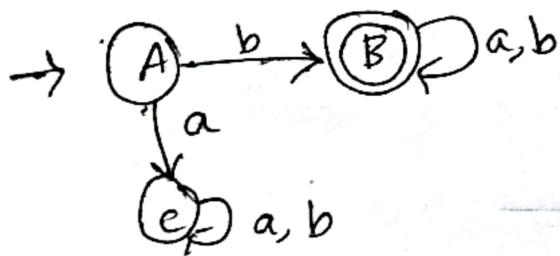
T 33 same as T 27, T 28

T 34

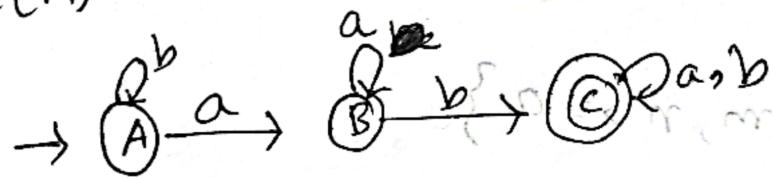
$L(M)$ = set of states with a
states with a = $\{w \mid w \in M\}$



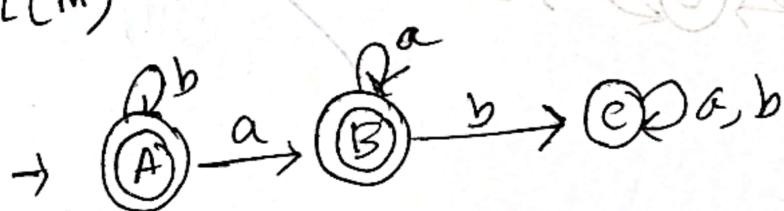
$L(M) =$ does not start with 'a'



$L(M) =$ contains substring 'ab'

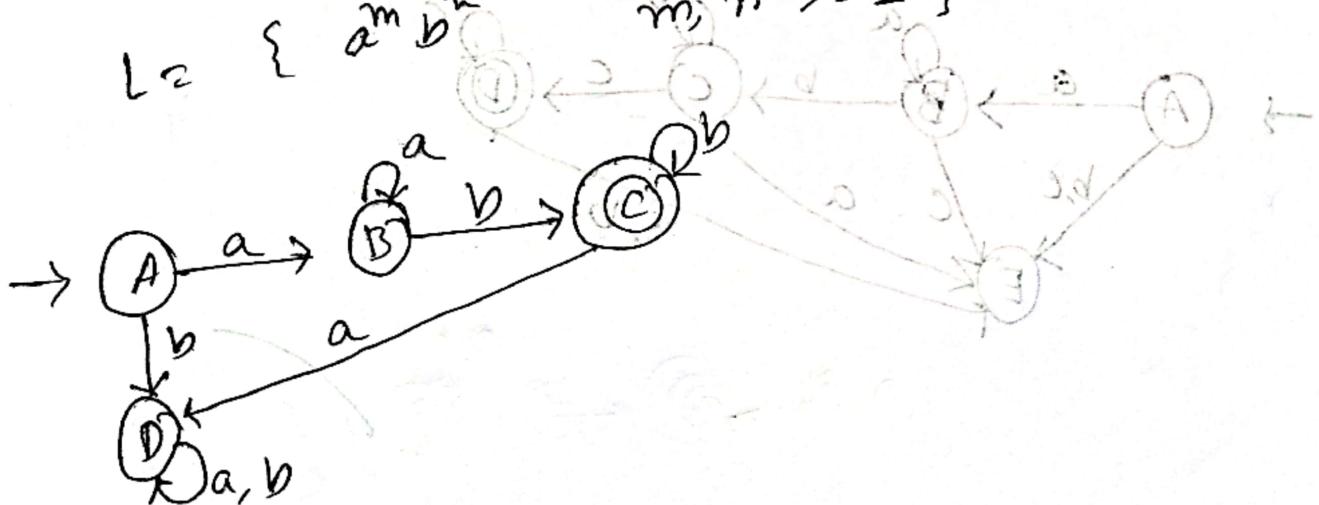


$L(M) =$ does not contain 'ab'



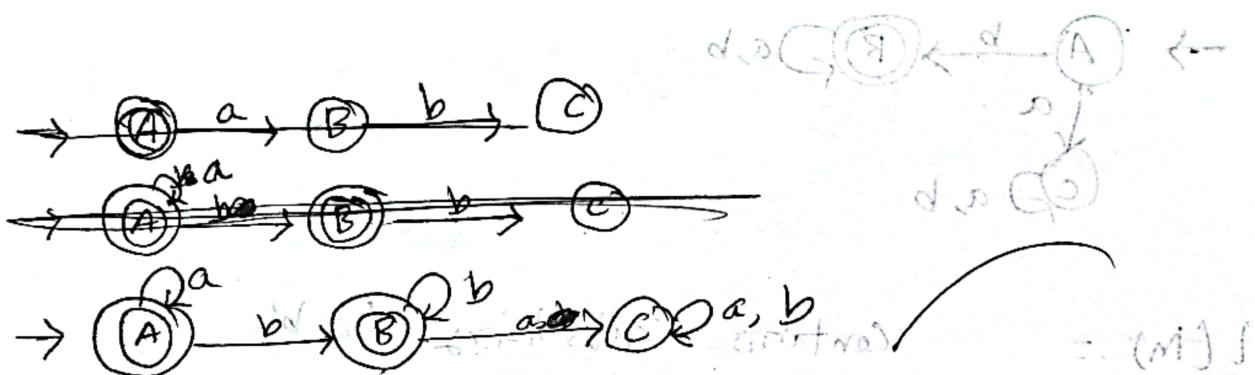
| T 35, -

$L = \{ a^m b^n \mid m, n \geq 1 \}$

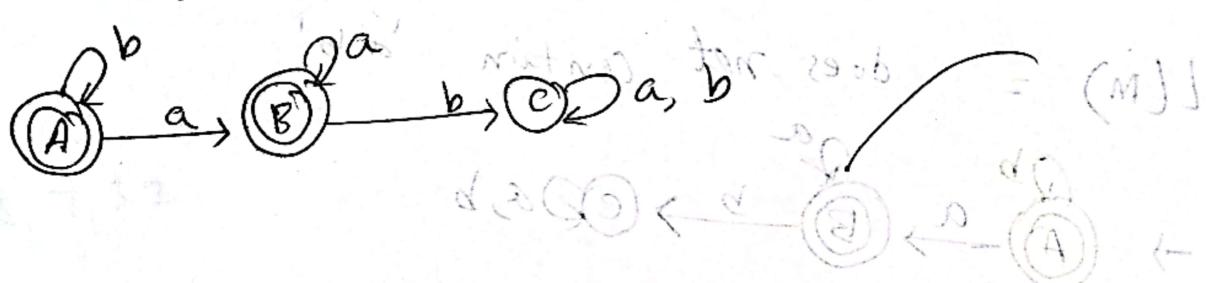


| L 36

$$L = \{ a^m b^n / m, n \geq 0 \} \rightarrow \text{RE} \rightarrow (M) \quad |$$

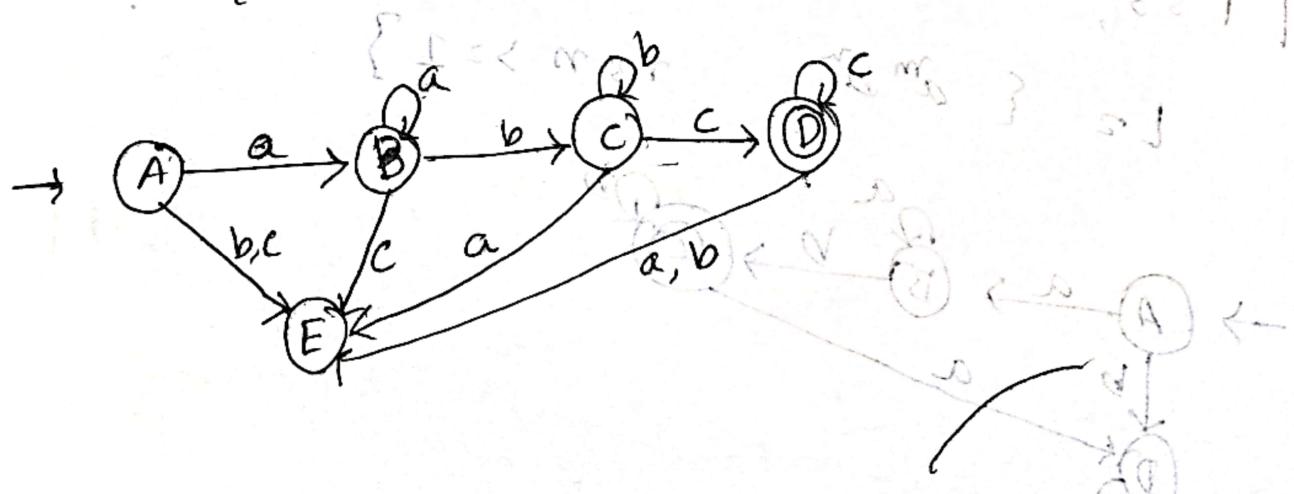


$$L = \{ b^m a^n / m, n \geq 0 \} \rightarrow \text{RE} \rightarrow (M) \quad |$$



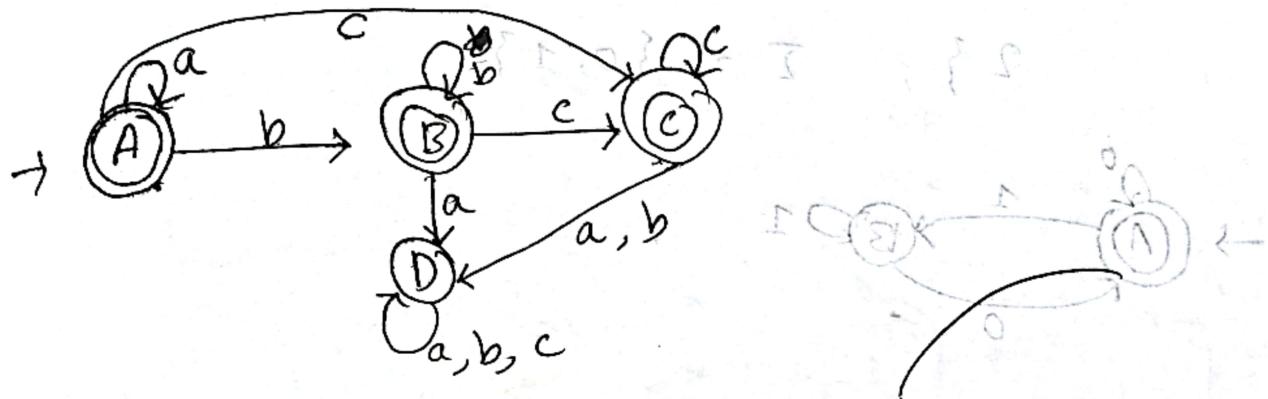
| L 37

$$L = \{ a^m b^n c^t / m, n, t \geq 1 \} \rightarrow \text{RE} \rightarrow (M) \quad |$$



| L 38

$L(m) = \{ a^m b^n c^k \mid a, b, c > 20 \}$

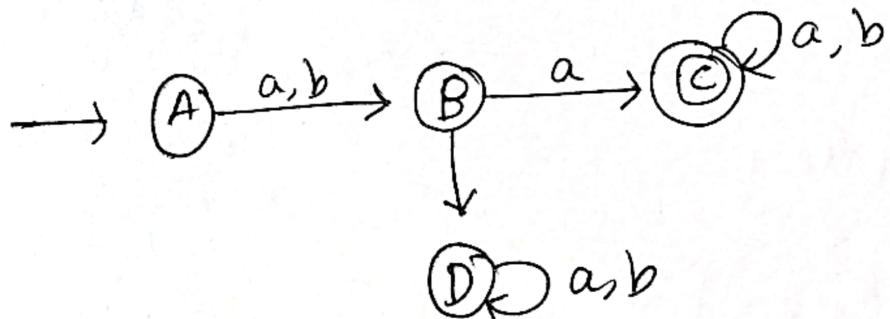


| L 39

$L(m) = \{ \text{2nd symbol must be } a \}$

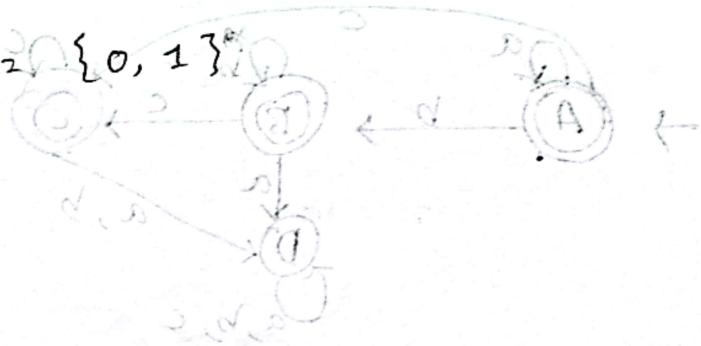
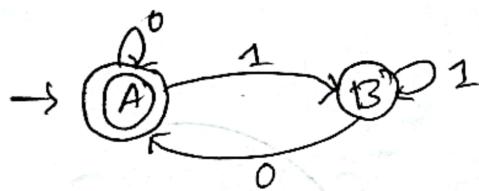
$$\sum_{\text{second symbol}} \{ a, b \}$$

if r = PAT

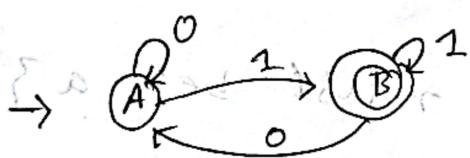


| T 40

$L = \{w \mid w \text{ binary number is divisible by } 2\}$, $\Sigma = \{0, 1\}^*$



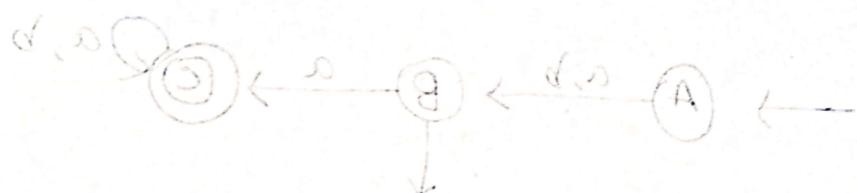
$L = \{w \mid w \text{ is not divisible by } 2\}$



| T 41 - T 49

~~skipped~~ Skipped Temporarily

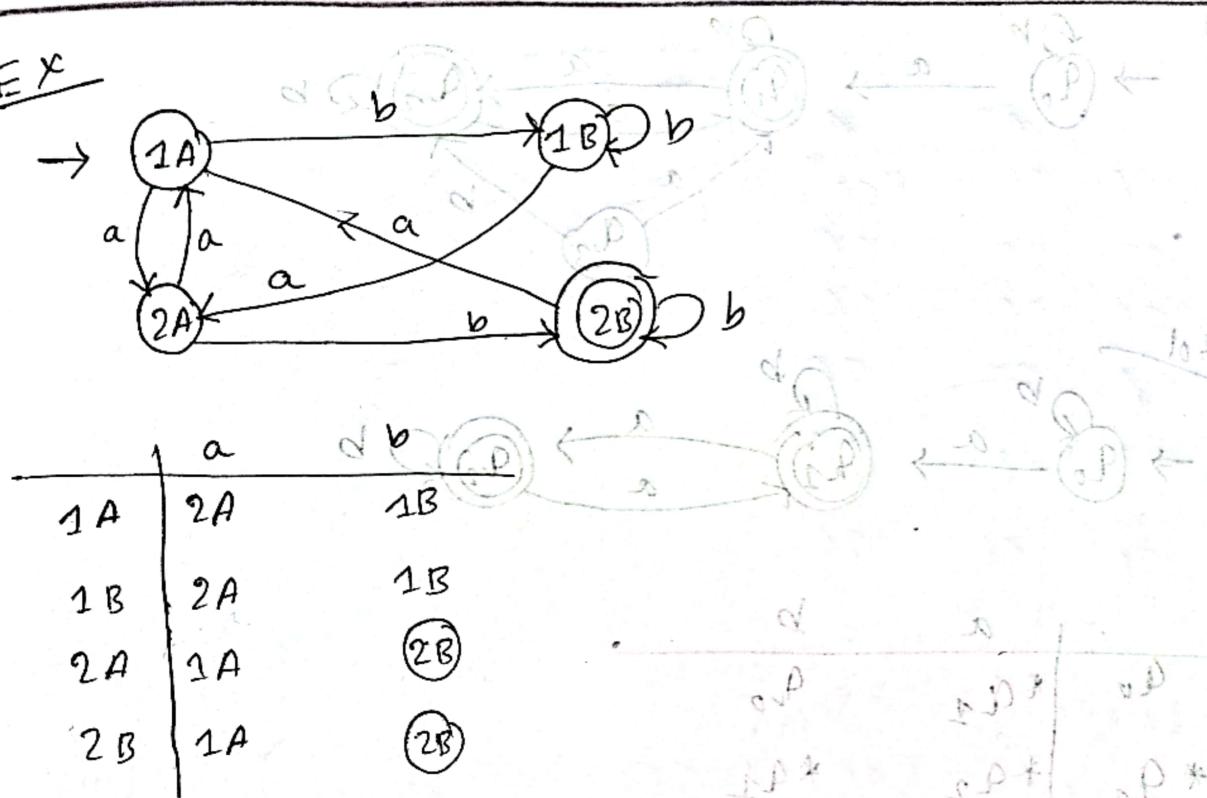
~~skip~~



~~do C @~~

1750 Minimization of DFA

Ex

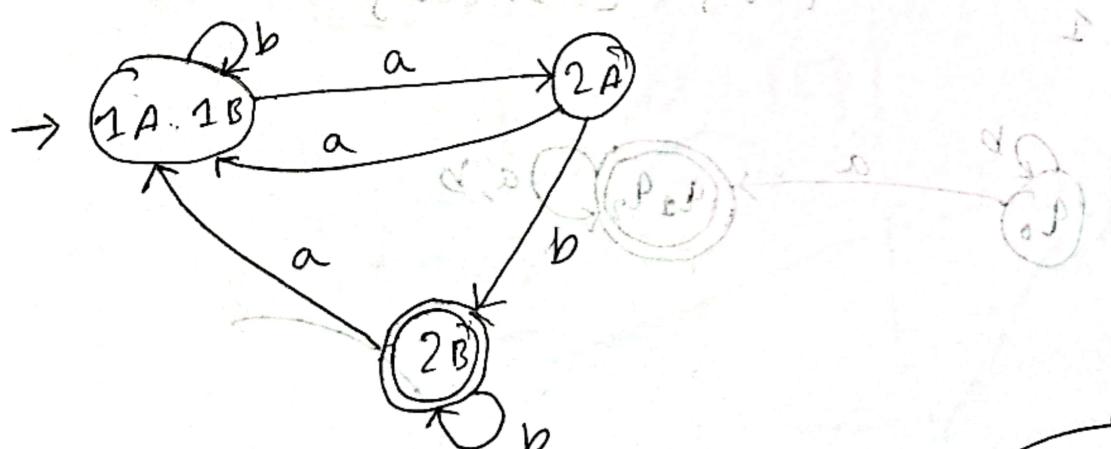


0 equivalent: $\{1A, 1B, 2A\}_{S^P} + \{2B\}_{S^P}$

1 - " : $\{1A, 1B\}, \{2A\}, \{2B\}$

$\{S^P, S^P\}, \{S^P, S^P\}$
 $\{1A, 1B\}, \{2A\}, \{2B\}$

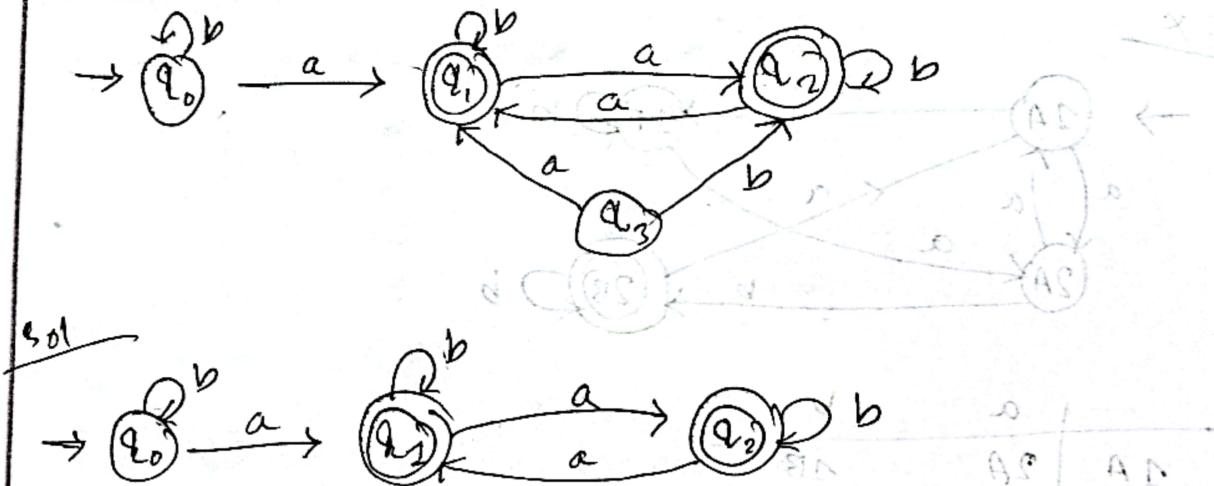
2 - " : $\{S^P, S^P\}, \{S^P, S^P\}$



T51

DFA minimization

part 1



	a	b
q_0	* q_1	q_0
* q_1	* q_2	* q_1
* q_2	* q_1 {AS} * q_2 {AF, AD}	{AF, AD}

	a	b
AS	AD	AS
AF	AF	AF
AD	AS	AS
FS	FS	FS

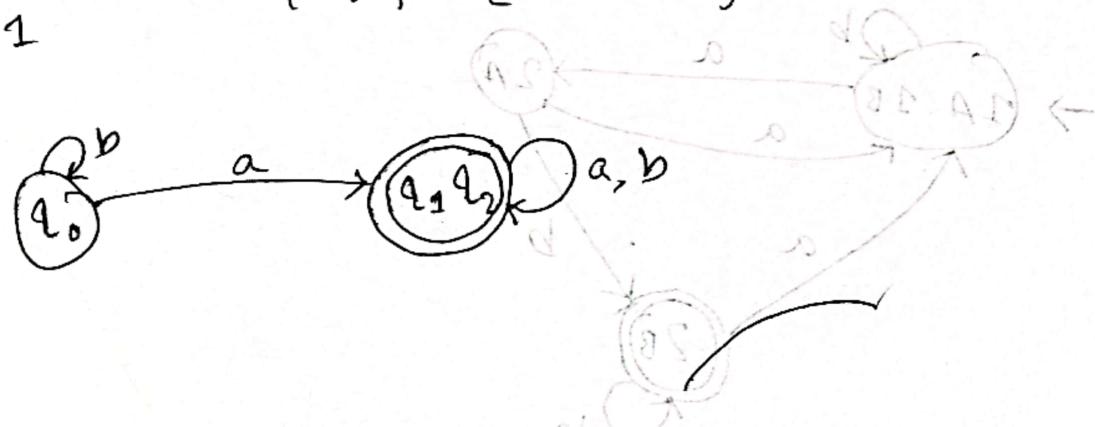
Equivalent sets: {AS}, {AF, AD}, {FS}

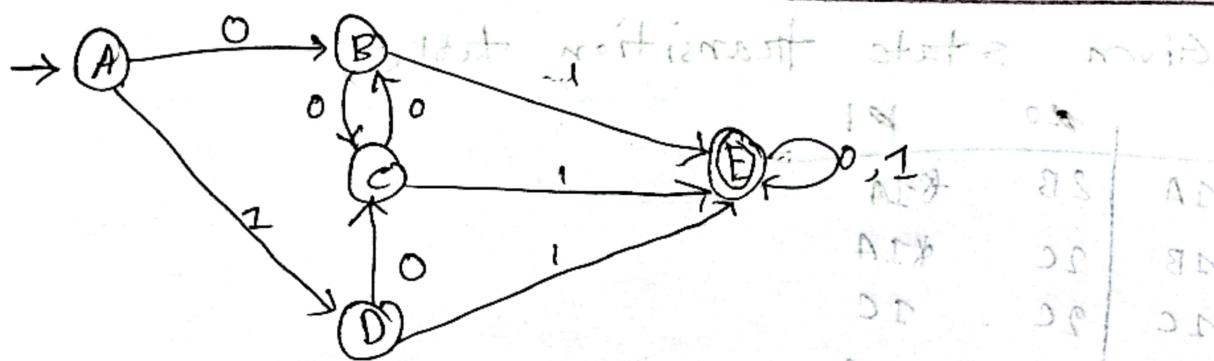
0 { q_0 } { q_1, q_2 }

{AS}, {FS}, {AF, AD}

{ q_0 } { q_1, q_2 }

1

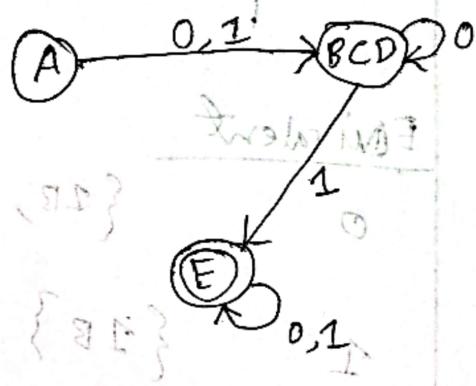


SOL
①

Transition table

	0	1
A	B	D
B	C	*E
C	B	
D	{C, D}	{E}
*E	{E}	{A, B, C, D}

② Now new DFA (mDFA),

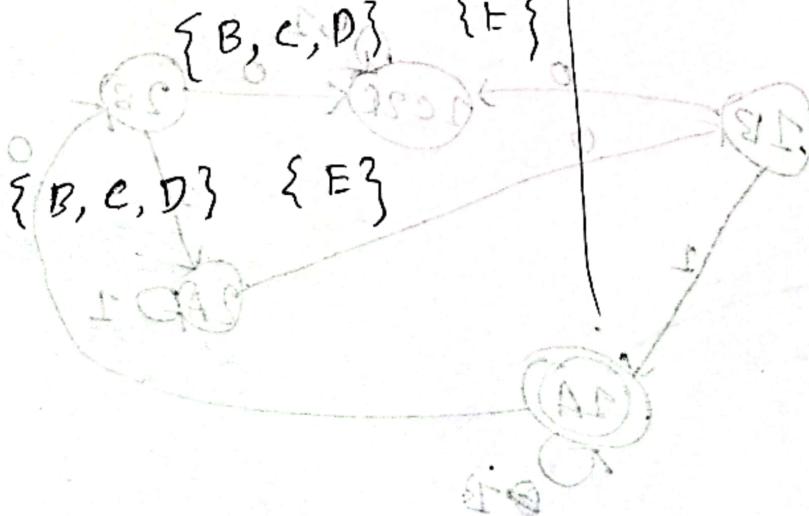


③ Equivalent

0 $\{A\} \{B, C, D\} \{E\}$

1 $\{A\}$

2 $\{A\} \{B, C, D\} \{E\}$



T 53

mDFA in reduced form

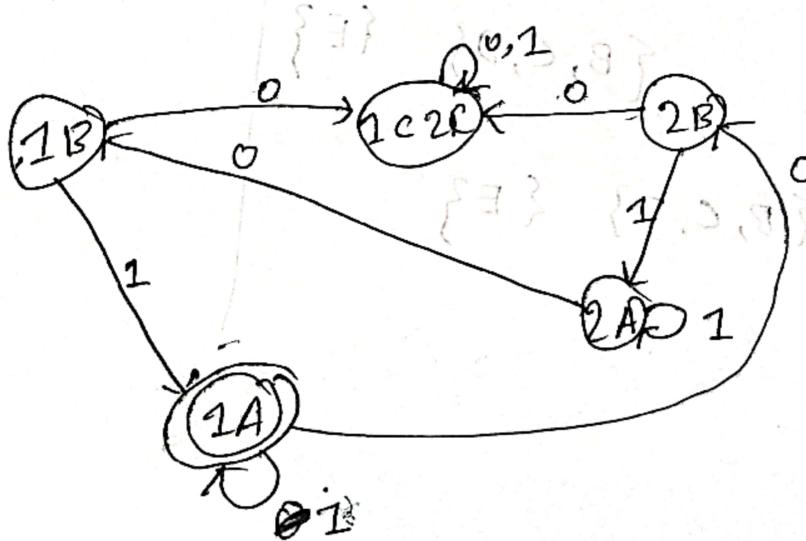
SFT

Given 3 state transition table,

	$\alpha 0$	$\alpha 1$
$*1A$	$2B$	$*1A$
$1B$	$2C$	$*1A$
$1C$	$2C$	$1C$
$2A$	$2B$	$2A$
$2B$	$1C$	$2A$
$2C$	$1C$	$2C$

Equivalent

0	$\{1B, 1C, 2A, 2B, 2C\}$	$\{1A\}$	A
1	$\{1B\}$	$\{1C, 2A, 2B, 2C\}$	$\{1A\}$
2	$\{1B\}$	$\{1C, 2B, 2C\}$	$\{2A\}$
3	$\{1B\}$	$\{1C, 2C\}$	$\{2B\}$
4	$\{1B\}$	$\{1C, 2C\}$	$\{2B\}$



T54

MDFA

(last ex)

Transition table

	a	b
1A	*2A	3B
* 2A	*2A	*2B
* 2B	*2B	2C
2C	2C	2C
3B	3B	3C
3C	3C	3C

Equivalent

0	$\{1A, 2C, 3B, 3C\}$	$\{2A, 2B\}$
1	$\{1A\}$	$\{2C, 3B, 3C\}$
2	$\{2A\}$	$\{2C, 3B, 3C\}$

