项目测试报告

任务三

一、项目任务一测试

1、准备工作

minic 的正则表达式:

sample.tny 文件:

```
//Sample program
//In MiniC language-computes factorial
//}
int x; // an integer
int fact;
double y;
int arr[100];
void test(int a, int b)
 a = 1;
 b = 2;
 return ;
void main(void)
 x = 2;
 y = 3;
 if (x > 1) //don't compute if x \le 0
      do
          fact = fact * x;
      while (fact >= 1);
 else
   X = 1;
  return; //return void
```

输入正则表达式:

■ 项目1 author: 李达良 — □	×
项目一: 正则表达式转lex	
姓名:李达良 班级:计科1班 学号: 20203231004 请输入正则表达式,具体输入规则请点击右边 "查看规则" 按钮	
+\ - \^\ % < <= > = = != = ; , \(\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \	;
功能选择:輸入正则表达式后,请先点击开始分析,再点击其他按钮查看结果,注意生成的NFA和DFA图不含关键词。 开始分析 NFA DFA DFA最小化 词法分析程序 查看lex文件	
	^
<	>

2、NFA



3、DFA图



4、DFA 图最小化



5、生成词法程序

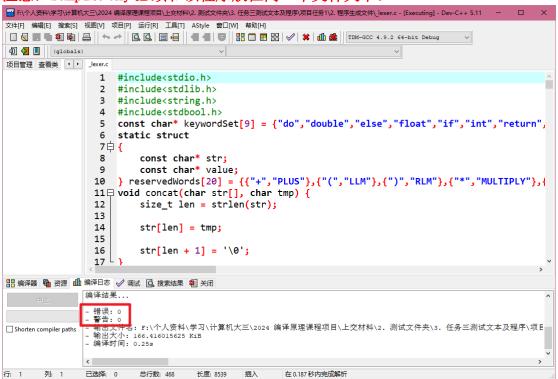


代码太长,具体可查看:

2. 测试文件夹\1. 任务一测试文本及程序\2. 程序生成文件中的_lexer.c文件

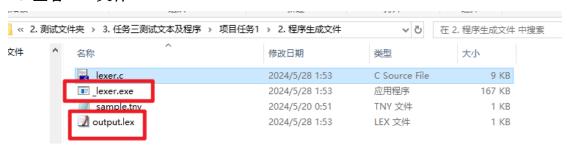
6、编译 lexer.c 文件并运行

注意: sample. tny 必须和该程序放在同一个文件夹下。



上图可知:编译成功 并成功生成 lex 文件:

7、 查看 1ex 文件





具体 1ex 如下:

```
int:int
ID:x
SEMI::
int:int
ID:fact
SEMI:;
double:double
ID:y
SEMI:;
int:int
ID:arr
LMM:
NUMBER: 100
RMM:]
SEMI::
void:void
ID:test
LLM:(
int:int
ID:a
DOU:,
int:int
```

```
ID:b
RLM:)
LBM: {
ID:a
ASSIGN:=
NUMBER: 1
SEMI:;
ID:b
ASSIGN:=
NUMBER: 2
SEMI:;
return:return
SEMI:;
RBM:}
void:void
ID:main
LLM:(
void:void
RLM:)
LBM: {
ID:x
ASSIGN:=
NUMBER: 2
SEMI:;
ID:y
ASSIGN:=
NUMBER: 3
SEMI:;
if:if
LLM:(
ID:x
RT:>
NUMBER: 1
RLM:)
do:do
ID:fact
ASSIGN:=
ID:fact
MULTIPLY:*
ID:x
SEMI:;
while:while
LLM:(
ID:fact
```

RTEQ:>=
NUMBER:1
RLM:)
SEMI:;
else:else
ID:x
ASSIGN:=
NUMBER:1
SEMI:;
return:return
SEMI:;
RBM:}
EOF:EOF

对照 minic 源程序,可知解析完全正确。

测试结果

任务三-项目一的测试完全通过

二、项目任务二的测试

1、准备工作

minic 的文法:

```
program | definition-list | definition | variable-definition
function-definition | type-indicator | parameters | compound-stmt
parameter-list | parameter | local-definitions | statement-list
statement | expression-stmt | condition-stmt | dowhile-stmt | return-
stmt expression simple-expression variable additive-
expression | relop | term | addop | mulop | factor | call | arguments
argument-list
ID | SEMI | LMM | RMM | int | float | double | void | LLM | RLM | if
else | do | while | return | LTEQ | LT | RT | RTEQ | EQ | NE | PLUS
| MINUS | MULTIPLY | DIVIDE | MOD | NUMBER | DOU | LBM | RBM | ASSIGN
program -> definition-list
definition-list -> definition-list definition
definition-list -> definition
definition -> variable-definition
definition -> function-definition
variable-definition -> type-indicator ID SEMI
variable-definition -> type-indicator ID LMM NUMBER RMM SEMI
type-indicator -> int
type-indicator -> float
type-indicator -> double
type-indicator -> void
function-definition -> type-indicator ID LLM parameters RLM compound-
stmt
parameters -> parameter-list
parameters -> void
parameter-list -> parameter-list DOU parameter
parameter-list -> parameter
parameter -> type-indicator ID
parameter -> type-indicator ID LMM RMM
compound-stmt -> LBM local-definitions statement-list RBM
local-definitions -> local-definitions variable-definition
local-definitions -> @
statement-list -> statement-list statement
statement-list -> @
statement -> expression-stmt
statement -> compound-stmt
statement -> condition-stmt
statement -> dowhile-stmt
statement -> return-stmt
expression-stmt -> expression SEMI
expression-stmt -> SEMI
```

```
condition-stmt -> if LLM expression RLM statement
condition-stmt -> if LLM expression RLM statement else statement
dowhile-stmt -> do statement while LLM expression RLM SEMI
return-stmt -> return SEMI
return-stmt -> return expression SEMI
expression -> variable ASSIGN expression
expression -> simple-expression
variable -> ID
variable -> ID LMM expression RMM
simple-expression -> additive-expression relop additive-expression
simple-expression -> additive-expression
relop -> LTEQ
relop -> LT
relop -> RT
relop -> RTEQ
relop -> EQ
relop -> NE
additive-expression -> additive-expression addop term
additive-expression -> term
addop -> PLUS
addop -> MINUS
term -> term mulop factor
term -> factor
mulop -> MULTIPLY
mulop -> DIVIDE
mulop -> MOD
factor -> LLM expression RLM
factor -> variable
factor -> call
factor -> NUMBER
call -> ID LLM arguments RLM
arguments -> argument-list
arguments -> @
argument-list -> argument-list DOU expression
argument-list -> expression
```

语义函数:

```
0
0 1
0
0
0
0
1 0 -1
1 0 -1 2 -1 -1
```

```
0
0
0
0
1 0 -1 2 -1 3
0
0
0 - 1 1
0
1 0
1 0 -1 -1
-1 \ 0 \ 1 \ -1
0 1
-1
0 1
-1
0
0
0
0
0
0 - 1
-1
0 -1 1 -1 2
0 -1 1 -1 2 -1 3
0 1 -1 -1 2 -1 -1
0 -1
0 1 -1
1 0 2
0
0
0 -1 1 -1
1 0 2
0
0
0
0
0
0
0
0 1 2
0
0
0
```

```
0 1 2

0 0

0 0

0 -1 0 -1

0 0

0 0

0 -1 1 -1

0 -1

0 -1 1
```

输入文法:



2、求解 first 集合



项目过多, 在此不一一展示。

3、求解 follow 集合



项目过多, 在此不一一展示。

4、LR0

生成提示 LR(O)DFA图 开始生成 This is a second of the											
D:program->definition-list 1:definition-list >definition-list definition 2:definition-list-	1	状态 0	状态内文法 program	ASSIGN	DIVIDE	DOU	EQ	ID	LBM ^		
>definition 3:definition->variable-	2	1	definition-list								
definition 4:definition=>function=	3	2	program								
definition 5:variable-definition-	4	3	type-indicator								
>type-indicator ID SEMI 6:variable-definition-	5	4	type-indicator								
>type-indicator ID LMM NUMBER RMM SEMI 7:type-indicator->int	6	5	definition								
8: type=indicator=>float 9: type=indicator=>double	7	6	type-indicator								
10: type-indicator->void 11: function-definition-	Q <	7	function-					11	>		

项目太多,未能展示完全

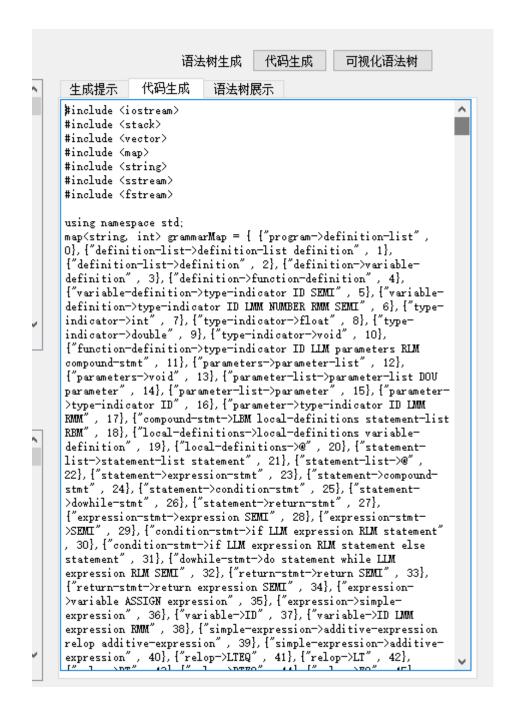
5、SLR1 表



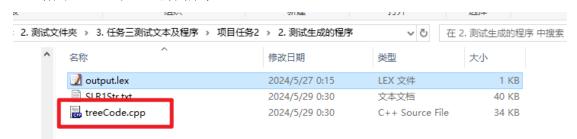
项目太多,未能展示完全 出现移进归约冲突,我们做只移进不规约的处理

6、语法树代码生成

注意: lex 所在路径即为代码生成路径

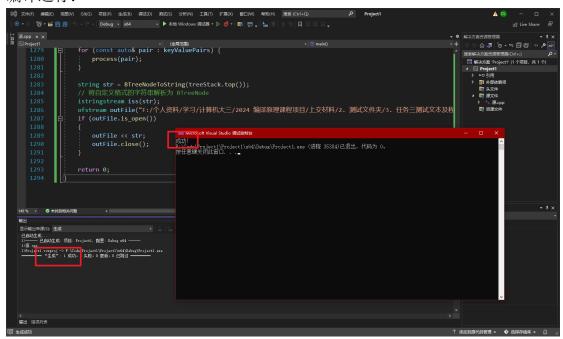


7、查看语法树生成代码文件 PS:请用 C++11 以上进行编译

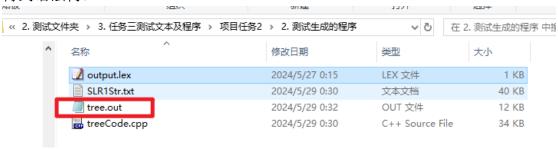


具体语法树代码太长,在此不进行展示,请打开 treeCode. cpp 查看

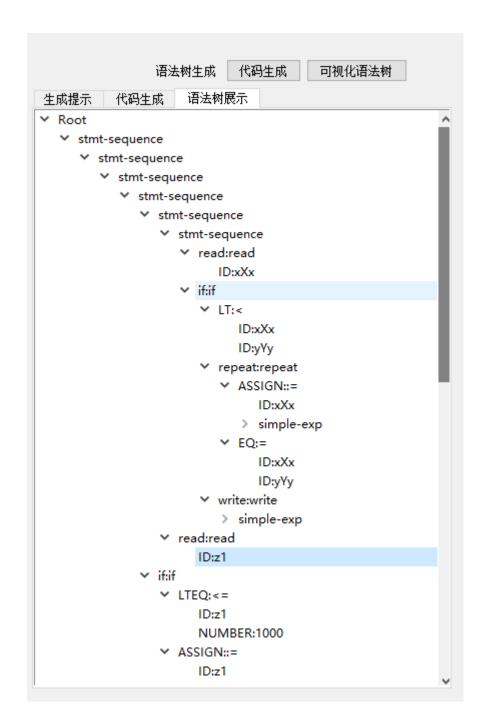
编译运行:



得到语法树:



8、可视化语法树



测试结果

任务三-项目二的测试完全通过