# Eigensystems

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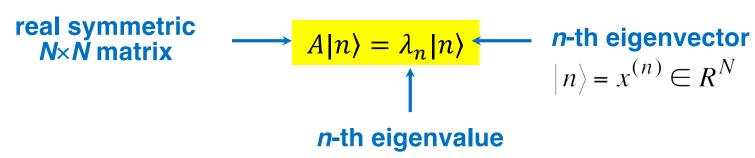
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- Matrix diagonalization methods in the context of quantum mechanics
- Matrix decompositions
- Vector space: projection & rotation



### **Eigenvalue Problem**

• Eigenvalue problem in N-dimensional vector space



or more explicitly

$$\sum_{j=1}^{N} A_{ij} x_j^{(n)} = \lambda_n x_i^{(n)}$$

*i*-th element of the *n*-th eigenvector

### **Orthonormal Basis**

- The basis set  $\{|n\rangle|n=1,...,N\}$  can be made orthonormal, i.e.,  $\langle m|n\rangle=\sum_{i=1}^N x_i^{(m)}x_i^{(n)}=\delta_{mn}$
- Orthogonal matrix:  $U = [x^{(1)} x^{(2)} ... x^{(N)}]$  or  $U_{in} \equiv x_i^{(n)}$  $U^T U = I \quad \because \sum_{i=1}^N x_i^{(m)} x_i^{(n)} = \sum_{i=1}^N U_{im}^{U_{mi}^T} U_{in} = (U^T U)_{mn} = \delta_{mn}$

#### (Proof: orthogonality)

For Hermitian matrix:

• 
$$\lambda_{m} \neq \lambda_{n}$$
  $\langle m|A|n \rangle = \lambda_{n} \langle m|n \rangle$   $\langle A^{\dagger} \rangle_{ij} = A_{ji}^{*} = A_{ij}$ 

-)  $\langle m|A|n \rangle = \lambda_{m} \langle m|n \rangle$  complex conjugate  $\langle n|A|m \rangle = \lambda_{m} \langle n|m \rangle$  complex conjugate  $\langle n|A|m \rangle = \langle n|A|n \rangle = \langle n|A^{\dagger}|n \rangle = \lambda_{n}^{*} \langle n|n \rangle$ 

$$\lambda_{n} \langle n|n \rangle = \langle n|A|n \rangle = \langle n|A^{\dagger}|n \rangle = \lambda_{n}^{*} \langle n|n \rangle$$

$$0 = (\lambda_{n} - \lambda_{n}^{*}) \langle n|n \rangle \Leftrightarrow \lambda_{n} = \lambda_{n}^{*}$$

- $\lambda_m = \lambda_n$  (degenerate): use Gram-Schmidt orthogonalization procedure
  - 1. Orthogonal projection:  $|n'\rangle \leftarrow |n\rangle |m\rangle\langle m|n\rangle = (1 |m\rangle\langle m|)|n\rangle$   $\langle m|n'\rangle = \langle m|n\rangle \overbrace{\langle m|m\rangle}^{1} \langle m|n\rangle = 0$   $|n'\rangle$
  - 2. Normalization:  $|n'\rangle \leftarrow |n'\rangle/\langle n'|n'\rangle^{1/2}$   $|m\rangle\langle m|n\rangle = |m\rangle\cos\theta$   $|m\rangle\langle m|n\rangle = |m\rangle\cos\theta$

 $|m\rangle\langle m|n\rangle = |m\rangle\cos\theta$ Directional cosine

# **Completeness**

• Arbitrary N-dimensional vector can be represented as a linear combination of (linearly independent) N vectors

$$|\psi\rangle = \sum_{n=1}^{N} |n\rangle\langle n|\psi\rangle$$
2D example (just Cartesian coordinates)  $|2\rangle\langle 2|\psi\rangle$ 

$$|1\rangle\langle 1|\psi\rangle$$

i.e., 
$$\sum_{n=1}^{N} |n\rangle\langle n| = 1$$
 or equivalently  $\sum_{n=1}^{N} x_i^{(n)} x_j^{(n)} = \delta_{ij}$ 

$$\psi_i = \sum_{n=1}^N x_i^{(n)} \sum_{j=1}^N x_j^{(n)} \psi_j = \sum_{j=1}^N \sum_{n=1}^N x_i^{(n)} x_j^{(n)} \psi_j$$

Orthogonal matrix

$$U^{T}U = UU^{T} = I$$

$$U^{T}U = U^{T} = I$$

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$$\delta_{ij} = \sum_{n=1}^{N} x_{i}^{(n)} x_{j}^{(n)} = \sum_{n=1}^{N} U_{in} \widetilde{U_{jn}^{T}} = (UU^{T})_{ij}$$

... Column-aligned eigenvectors,  $U = [x^{(1)} x^{(2)} ... x^{(N)}]$ , can be made an orthogonal matrix

# **Orthogonal Transformation**

$$\sum_{i=1}^{N} x_i^{(m)} \times \left( \sum_{j=1}^{N} A_{ij} x_j^{(n)} = \lambda_n x_i^{(n)} \right)$$

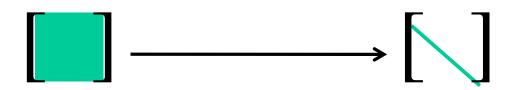
$$\sum_{i=1}^{N} \sum_{j=1}^{N} \underbrace{x_i^{(m)}}_{i} A_{ij} \underbrace{x_j^{(n)}}_{i} = \lambda_n \sum_{i=1}^{N} x_i^{(m)} x_i^{(n)} = \underbrace{\lambda_n \delta_{mn}}_{i}$$
orthogonality

**Matrix eigenvalue problem = find an orthogonal transformation matrix** 

Spectral  
decomposition
$$U^TAU = \Lambda$$
  
 $\Lambda_{mn} = \lambda_m \delta_{mn}$ 

**Grand strategy:** Nudge the matrix A towards diagonal form by a sequence of orthogonal transformations (successive elimination of off-diagonal elements)

$$A \to P_1^T A P_1 \to \overbrace{P_2^T P_1^T}^T A \overbrace{P_1 P_2}^U \to \cdots$$
$$U = P_1 P_2 \cdots$$



### **Rotation**

- General real symmetric 2×2 matrix:  $H = \begin{bmatrix} \varepsilon_1 & \delta \\ \delta & \varepsilon_2 \end{bmatrix}$
- General orthonormal matrices:  $|u\rangle = \begin{bmatrix} \cos\theta \\ \sin\theta \end{bmatrix} = \cos\theta |1\rangle + \sin\theta |2\rangle; |v\rangle = \begin{bmatrix} -\sin\theta \\ \cos\theta \end{bmatrix}$

• Eigenvalue solution

$$U = \begin{bmatrix} u & v \end{bmatrix} = \begin{bmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{bmatrix}$$

$$\begin{bmatrix} \lambda - \varepsilon_{1} & -\delta \\ -\delta & \lambda - \varepsilon_{2} \end{bmatrix} \begin{bmatrix} \cos \theta \\ \sin \theta \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix} \quad \det(\lambda I - H) = \begin{vmatrix} \lambda - \varepsilon_{1} & -\delta \\ -\delta & \lambda - \varepsilon_{2} \end{vmatrix} = (\lambda - \varepsilon_{1})(\lambda - \varepsilon_{2}) - \delta^{2} = 0$$

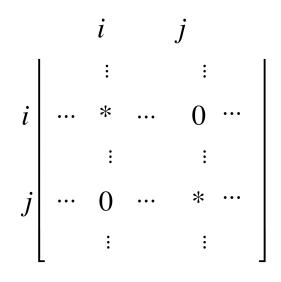
$$\lambda_{\pm} = \frac{\varepsilon_{1} + \varepsilon_{2} \pm \sqrt{(\varepsilon_{1} - \varepsilon_{2})^{2} + 4\delta^{2}}}{2}$$

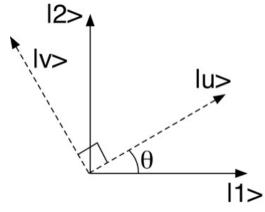
$$\theta = \tan^{-1} \left( \frac{-\varepsilon_{1} + \varepsilon_{2} + \sqrt{(\varepsilon_{1} - \varepsilon_{2})^{2} + 4\delta^{2}}}{2\delta} \right) \xrightarrow{\delta \to 0} \frac{\delta}{\varepsilon_{1} - \varepsilon_{2}}$$

$$\text{for } \lambda_{+} \text{ and } \varepsilon_{1} > \varepsilon_{2}$$

### **Jacobi Transformation**

• Successive 2D rotations to eliminate off-diagonal (i,j)–(j,i) pairs



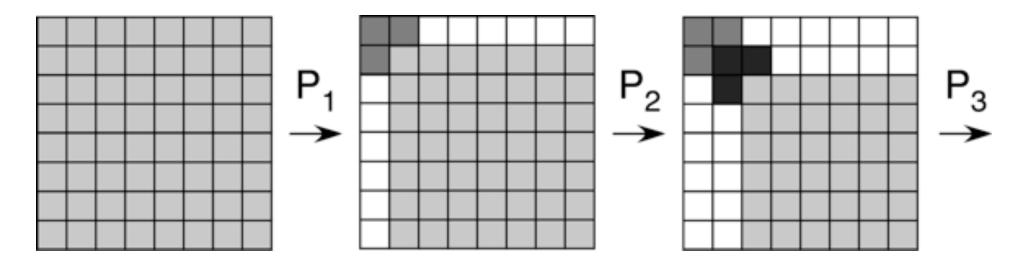




Carl Jacobi (1804-1851)

### **Householder Transformation**

• Eliminate an entire row (but the first 2 elements) at a time



• The outcome is a tridiagonal matrix



Alston Householder (1904-1993)

# **Projection Matrix**

• Let an N-dimensional vector  $v \in \mathbb{R}^N$ ) & the projection matrix

$$P = I - \frac{2vv^{T}}{v^{T}v} = I - \frac{2|v\rangle\langle v|}{\langle v|v\rangle}$$

then P is symmetric & orthonormal, i.e.,

$$P^TP = PP^T = I$$

#### (Proof)

$$P_{ij} = \delta_{ij} - \frac{2v_i v_j}{\sum_{k=1}^{N} v_k^2} \qquad \text{symmetric w.r.t. } i \leftrightarrow j$$

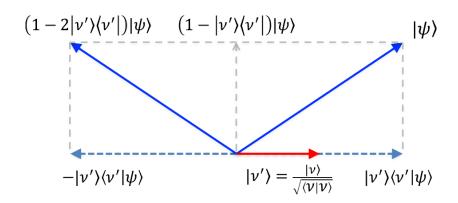
$$PP = \left(I - \frac{2vv^T}{v^T v}\right) \left(I - \frac{2vv^T}{v^T v}\right)$$

$$= I - \frac{4vv^T}{v^T v} + \frac{4vv^T vv^T}{v^T vv^T v} \qquad \text{Mirror imag}$$

$$= I - \frac{4vv^T}{v^Tv} + \frac{4vv^T}{v^Tv}$$

$$= I$$

#### **Mirror image: reflect twice = do nothing**



### **Householder Matrix**

• For  $x \in \mathbb{R}^N$ , let  $v = x \mp ||x||_2 e_1$  where

$$e_1 = \begin{bmatrix} 1 \\ 0 \\ \vdots \end{bmatrix}$$
 & the vector 2-norm is  $||x||_2 = \sqrt{x^T x} = \sqrt{\sum_{i=1}^N x_i^2}$ 

then the Householder matrix below, when multiplied, eliminates all the elements of x but one:

$$Px = \left(I - \frac{2vv^T}{v^Tv}\right)x = \mp \|x\|_2 e_1$$

(Proof)

$$v^T v = (x^T \pm ||x||_2 e_1^T) (x \pm ||x||_2 e_1) = ||x||_2^2 \pm 2||x||_2 x_1 + ||x||_2^2 = 2||x||_2 (||x||_2 \pm x_1)$$

$$Px = x - \frac{2vv^{T}}{2\|x\|_{2}(\|x\|_{2} \pm x_{1})}x$$

$$= x - \frac{(x \pm \|x\|_{2}e_{1})(x^{T} \pm \|x\|_{2}e_{1}^{T})x}{\|x\|_{2}(\|x\|_{2} \pm x_{1})}$$

$$= x - \frac{(x \pm \|x\|_{2}e_{1})\|x\|_{2}(\|x\|_{2} \pm x_{1})}{\|x\|_{2}(\|x\|_{2} \pm x_{1})}$$

$$= x - x \mp \|x\|_{2}e_{1} = \mp \|x\|_{2}e_{1}$$

# **Tridiagonalization**

Householder matrix can be used for tridiagonalization: Let

$$A = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1N} \\ a_{21} & & & \\ \vdots & & & \\ a_{N1} & & & \end{bmatrix} = \begin{bmatrix} a_{11} & A_{12} = A_{21}^T \\ & & & \\ A_{21} & & & \\ & & & \end{bmatrix}$$

$$v \in \mathbb{R}^{N-1} = A_{21} + \operatorname{sign}(a_{21}) ||A_{21}||_2 e_1$$

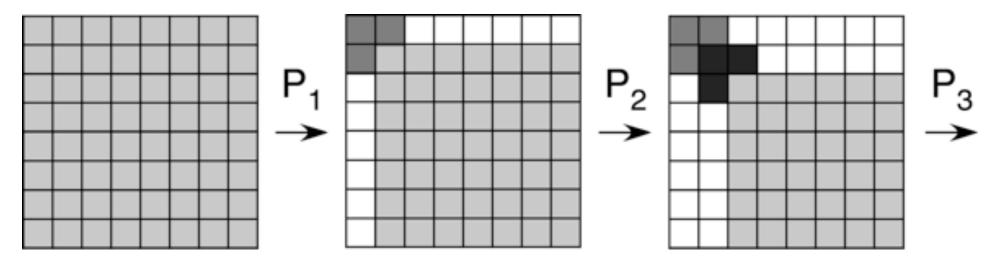
then

$$\overline{P}A_{21} = \left(I_{N-1} - \frac{2vv^{T}}{v^{T}v}\right)A_{21} = -\operatorname{sign}(a_{21})||A_{21}||_{2}e_{1} = ke_{1}$$

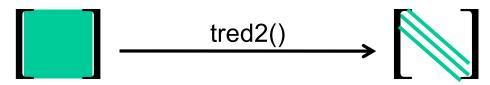
$$PAP = \begin{bmatrix} 1 & 0 & \cdots & 0 \\ 0 & & & \\ & \overline{P} & & \\ 0 & & & \overline{P}A_{21} \end{bmatrix} \begin{bmatrix} a_{11} & A_{21}^{T} & & \\ a_{21} & A_{22} & & \\ 0 & & & \overline{P}A_{22} \end{bmatrix} \begin{bmatrix} 1 & 0 & \cdots & 0 \\ 0 & & & \\ & \overline{P} & & & \\ 0 & & & \overline{P}A_{22} \overline{P} \end{bmatrix} = \begin{bmatrix} a_{11} & k & 0 & \cdots & 0 \\ k & & & & \overline{P}A_{22} \overline{P} \\ \vdots & & & & \\ 0 & & & & \overline{P}A_{22} \overline{P} \end{bmatrix}$$

### **Householder Transformation**

• After (N-2) such transformations, all the off-diagonal elements but the diagonal & upper/lower sub-diagonal elements are eliminated



• The outcome is a tridiagonal matrix (done in tred2() in *Numerical Recipes*)



# **QR** Decomposition

- Used for the diagonalization of a tridiagonal matrix
- Let A = QR, where Q is orthogonal & R is upper-triangular,  $R_{ij} = 0$  for i > j
- QR decomposition by Householder transformation

$$A = \begin{bmatrix} a_{11} \\ \vdots \\ a_{N1} \end{bmatrix} = \begin{bmatrix} A_1 & A_2 \end{bmatrix} \quad v \quad (\in R^N) = A_1 + \text{sign}(a_{11}) ||A_1||_2 e_1$$

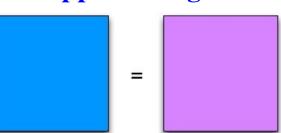
$$PA_1 = \left( I_N - \frac{2vv^T}{v^T v} \right) A_1 = -\text{sign}(a_{11}) ||A_1||_2 e_1 = ke_1$$

$$PA = \begin{bmatrix} PA_1 & PA_2 & \\ PA_1 & PA_2 & \\ 0$$

• After (N-1) transformations, the matrix is upper-triangular

$$P_{N-1} \cdots P_2 P_1 A = R$$

$$A = P_1^{-1} P_2^{-1} \cdots P_{N-1}^{-1} R \equiv QR$$





# Orthogonal Transformation by QR

$$A = QR \quad A' = RQ$$

$$R = Q^{-1}A = Q^{T}A$$

$$A \to A' = Q^{T}AQ$$

#### (QR algorithm)

$$\begin{cases} 1. Q_s R_s \leftarrow A_s \\ 2. A_{s+1} \leftarrow R_s Q_s \end{cases} \quad s = 1, 2, \dots$$

#### (Theorem)

- 1.  $\lim_{s\to\infty}A_s$  is upper-triangular
- 2. The eigenvalues appear on its diagonal
- tqli() in *Numerical Recipes* uses QL algorithm instead to obtain lowertriangular matrix
- Fast O(N) operations per iteration for a tridiagonal matrix
- tqli() diagonalizes a tridiagonal matrix by a sequence of rotations to eliminate subdiagonal elements, in addition to eigenvalue-shift to accelerate the convergence

# Top 10 algorithms in history *IEEE CiSE*, Jan/Feb ('00)

- Metropolis Algorithm for Monte Carlo
- Simplex Method for Linear Programming
- Krylov Subspace Iteration Methods
- The Decompositional Approach to Matrix Computations
- The Fortran Optimizing Compiler
- QR Algorithm for Computing Eigenvalues
- Quicksort Algorithm for Sorting
- Fast Fourier Transform
- Integer Relation Detection
- Fast Multipole Method