OpenMP Target Offload for Heterogeneous Architectures

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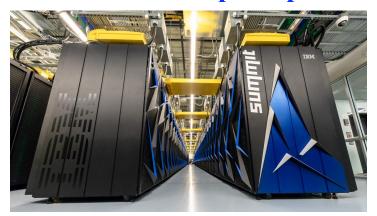
Goal: Unified open high-level programming of both CPU & GPU





Exaflop/s Supercomputing

Diverse exaflop/s supercomputing platforms



Summit (0.2 Exaflop/s)

IBM CPU/NVIDIA GPU



1 HPC and AI Optimized AMD EPYC CPU 4 Purpose Built AMD Radeon Instinct GPU



Frontier (1.7 Exaflop/s)

AMD CPU/AMD GPU



GPU Architecture

X^e arch-based "Ponte Vecchio" GPUTile-based, chiplets, HBM stack, Foveros 3D integration, 7nm

On-Node Interconnect

CPU-GPU: PCle GPU-GPU: X^e Link



Aurora (2.0 Exaflop/s)
Intel CPU/Intel GPU

• Need an *open* programming model for *heterogeneous* (*e.g.*, GPU-accelerated) clusters (note CUDA is a proprietary language by NVIDIA)

See https://extremecomputingtraining.anl.gov/agenda-2024/

Open Programming Models

OpenACC (Open Computing Language)
 Open standard for directive-based programming of heterogeneous devices

USING OPENMP THE NEXT STEP

Affinity, Accelerators, Tasking, and SIMD

Ruud van der Pas, Eric Stotzer,

https://www.openacc.org/

OpenMP 4.5/5
 Starting specification version 4.5,
 OpenMP allows offloading the execution of code & data to heterogeneous devices

https://www.openmp.org/specifications/

OpenMP Offload

- Latest version of OpenMP allows one to maintain one version of a code, which can run on either a general-purpose central processing unit (CPU) or an accelerator (e.g. graphics processing unit, GPU; tensor processing unit, TPU; digital signal processor, DSP; field-programmable gate array, FPGA)
- Objective is to execute parts of the program on a heterogeneous accelerator device (or target device), i.e., dedicated computer hardware outside CPU (which will be called host device) to execute certain functions faster than CPU
- In OpenMP, program execution begins on the host, which offloads the execution of parts of the code & data to accelerator

See "OpenMP 4.5 Target Offload" (NASA Ames) https://aiichironakano.github.io/cs653/OpenMP4.5 3-20-19.pdf

OpenMP Target Construct

Simple example

```
main() {
   float a[1000],b[1000],c,d;
   ...
   #pragma omp target map(a,b,c,d)
   {
      int i;
      #pragma omp parallel for
      for (i=0; i<N; i++)
        a[i] = b[i]*c+d;
   }
   ...
}</pre>
```

- When a host thread encounter the #pragma omp target directive, the target region specified by it will be executed by a new thread running on an accelerator, cf. CUDA GPU kernel
- Before the new thread starts executing the target region, the variable in the map() clause are mapped onto accelerator memory, which often is disjunct from the host memory, cf. cudaMemcpy()
- The offloaded code is usually a data-parallel structured block, which can be handled by multiple threads on accelerator using standard OpenMP constructs like #pragma omp parallel for

Computing the Value of π on GPU

omp_target_pi.c

```
#include <omp.h>
#include <stdio.h>
#define NBIN 1000000
#define NTRD 96
                                 The only addition for GPU offload
int main() {
  float step, sum=0.0, pi;
  step = 1.0/(float)NBIN;
                                          Thread reduction of sum
  #pragma omp target map(step,sum)
    # pragma omp parallel for reduction(+:sum) num threads(NTRD)
    for (long long i=0; i<NBIN; i++) {
      float x = (i+0.5)*step;
      sum += 4.0/(1.0+x*x);
                                            Specify # of GPU threads
                            This line is identical for CPU & GPU
 pi = sum*step;
  printf("PI = %f\n",pi);
                                             1. Black: original serial code
  return 0;
                                             2. Green: one-line multithreading
                                             3. Red: another line for GPU offload
```

GPU: Easy & Hard Ways

Serial: pi.c OpenMP: omp_target_pi.c

```
#include <stdio.h>
#define NBIN 100000000
int main() {
    double step, x, sum=0.0, pi;
    step = 1.0/NBIN;
    for (long long i=0; i<NBIN; i++) {
        x = (i+0.5)*step;
        sum += 4.0/(1.0+x*x);
    }
    pi = sum*step;
    printf("PI = %f\n",pi);
    return 0;
}</pre>
```

```
#include <omp.h>
#include <stdio.h>
#define NBIN 1000000
int main() {
  float step, sum=0.0, pi;
    step = 1.0/(float)NBIN;
    #pragma omp target map(step, sum)
    {
            # pragma omp parallel for reduction(+:sum)
            for (long long i=0; i<NBIN; i++) {
                float x = (i+0.5)*step;
                sum += 4.0/(1.0+x*x);
            }
            pi = sum*step;
            printf("PI = %f\n",pi);
            return 0;
}</pre>
```

DPC++: pi.cpp

```
#include <CL/sycl.hpp>
#include <iostream>
#include <array>
using namespace cl::sycl;
#define NBIN 1000000 // # of bins for quadrature
#define NTRD 512
                    // # of threads
int main()
 float step = 1.0f/NBIN;
 std::array<float, NTRD> sum;
 for (int i=0; i<NTRD; ++i) sum[i] = 0.0f;
 queue q(gpu_selector{});
  std::cout << "Running on: " <<
 q.get device().get info<info::device::name>() << std::endl;</pre>
  range<1> sizeBuf{NTRD}:
   buffer<float, 1> sumBuf(sum.data(), sizeBuf);
   g.submit([&](handler &h){
     auto sumAccessor =
      sumBuf.get access<access::mode::read write>(h);
     h.parallel for(sizeBuf, [=](id<1> tid) {
        for (int i=tid; i<NBIN; i+=NTRD) {
         float x = (i+0.5f)*step;
         sumAccessor[tid] += 4.0f/(1.0f+x*x);
      }); // End parallel for
   }); // End queue submit
  float pi=0.0f;
  for (int i=0; i<NTRD; i++) // Inter-thread reduction
 pi *= step; // Multiply bin width to complete integration
 std::cout << "Pi = " << pi << std::endl;
 return 0;
```

CUDA: pi.cu

```
// Using CUDA device to calculate pi
#include <stdio.h>
#include <cuda.h>
#define NBIN 10000000 // Number of bins
#define NUM_BLOCK 13 // Number of thread blocks
#define NUM_THREAD 192 // Number of threads per block
float pi = 0;
// Kernel that executes on the CUDA device
__global__ void cal_pi(float *sum, int nbin, float step, int nthreads, int nblocks) {
  int i.
  float x:
  int idx = blockIdx.x*blockDim.x+threadIdx.x: // Seguential thread index across the blocks
  for (i=idx; i< nbin; i+=nthreads*nblocks) {
     x = (i+0.5)*step;
     sum[idx] += 4.0/(1.0+x*x);
// Main routine that executes on the host
int main(void) {
  dim3 dimGrid(NUM_BLOCK,1,1); // Grid dimensions
  dim3 dimBlock(NUM_THREAD,1,1); // Block dimensions
  float *sumHost, *sumDev; // Pointer to host & device arrays
  float step = 1.0/NBIN; // Step size
  size t size = NUM BLOCK*NUM THREAD*sizeof(float); //Array memory size
  sumHost = (float *)malloc(size); // Allocate array on host
  cudaMalloc((void **) &sumDev, size); // Allocate array on device
  // Initialize array in device to 0
  cudaMemset(sumDev, 0, size);
  // Do calculation on device
  cal_pi <<<dimGrid, dimBlock>>> (sumDev, NBIN, step, NUM_THREAD, NUM_BLOCK); // call CUDA kernel
  // Retrieve result from device and store it in host array
  cudaMemcpy(sumHost, sumDev, size, cudaMemcpyDeviceToHost);
  for(tid=0; tid<NUM THREAD*NUM BLOCK; tid++)
     pi += sumHost[tid];
  pi *= step;
  // Print results
  printf("PI = %f\n",pi);
  // Cleanup
  free(sumHost);
  cudaFree(sumDev);
  return 0:
```

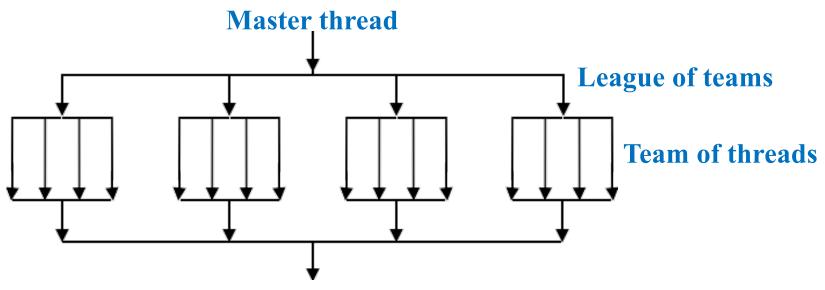
Hierarchical Parallelization

#pragma omp teams

Starts a league of multiple thread teams; teams construct must be nested immediately inside a target construct, so they are commonly used as #pragma omp target teams; num_teams clause can be used to specify the number of teams

(Example) #pragma omp target teams num_teams(13)

#pragma omp distribute
 Distribute the work across the teams



cf. CUDA grid of blocks & block of threads

Remember NVIDIA SM/SP & Intel slice/stack?

Teams for Computing π

Spatial decomposition via offset among teams & data privatization

```
#define NTMS 12
float sum_teams[NTMS];
for (int j=0; j<NTMS; j++) sum_teams[j] = 0.0;</pre>
                                                    - Data privatization among teams
#pragma omp target teams map(step,sum teams) num teams(NTMS)
  #pragma omp distribute ____
                                             New: teams & distribute constructs
  for (int j=0; j<NTMS; j++) {</pre>
    long long ibgn = NBIN/NTMS*j;
    long long iend = NBIN/NTMS*(j+1); NBIN/NTMS_XNTMS could be less than NBIN
    if (j == NTMS-1) iend = NBIN;
    # pragma omp parallel for reduction(+:sum teams[j]) num threads(NTRD)
    for (long long i=ibqn; i<iend; i++) {
       float x = (i+0.5)*step;
       sum teams[j] += 4.0/(1.0+x*x);
                Modified: offset & private accumulator
for (int j=0; j<NTMS; j++) sum += sum teams[j];</pre>
                                                              Master thread
                            League of teams
                        Team of threads
```

Using OpenMP Target on Discovery

• Necessary module

```
module purge
module load nvhpc
```

Compilation

```
nvc -o omp_target_pi omp_target_pi.c

NVIDIA C compiler supports newer C constructs
```

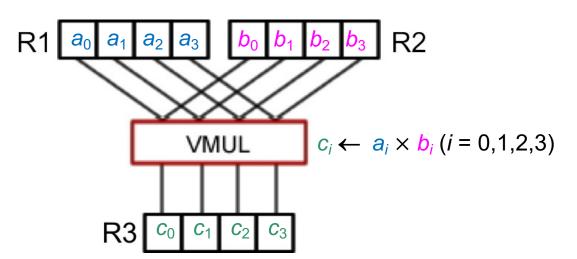
Execute on a GPU-accelerated node

```
[anakano@discovery]$ salloc --gres=gpu:1 --time=00:30:00
salloc: Nodes e16-03 are ready for job
[anakano@e16-03]$ ./omp_target_pi
PI = 3.141593
```

Single Instruction Multiple Data

- OpenMP 4.5 & later supports several other new features
- Single-instruction multiple-data (SIMD) parallelism: An arithmetic operation is operated on multiple operand-pairs stored in vector registers (each of which can hold multiple operands) using vector instructions
- OpenMP simd construct instructs the compiler to vectorize the loop

```
#pragma omp for simd
for (int i=0; i<n; i++)
    a[i] = b[i]+c[i];</pre>
```



Asynchronous Offload

```
main() {
  float a[1000],b[1000],c,d;
   ...
  #pragma omp target nowait map(a,b,c,d)
  {
    int i;
    #pragma omp parallel for
    for (i=0; i<N; i++)
        a[i] = b[i]*c+d;
  }
  func(b); // perform computation independent of device output
  #pragma omp taskwait
  func(a); // perform computation dependent on device output
}</pre>
```

- By default, the thread that encounters a device construct waits for the construct to complete before executing the next line
- When a nowait clause is added to the device construct, the encountering thread does not wait but instead continues executing the code passed the construct
- The taskwait constructs lets the original thread wait for the completion of the target task generated by it before continuing to the next line

Overlap CPU & GPU computations for high performance

Persistent GPU Memory Allocation

• Expensive CPU-GPU data transfer associated with the map clause can be minimized by making GPU memory allocation persistent

cf. 3-phase (host-to-device copy→kernel execution→device-to-host copy) performance bottleneck

```
// Pre-allocate GPU array
#pragma omp target enter data map(alloc:psi[0:Nmax])
                          Stand-alone directive to map variables to device memory
// Keep operating on device array only from device
#pragma omp target parallel for \
#pragma omp target parallel lor \
map(tofrom:psi[0:Nmax]) map(to:u[0:Nmax]) to: host-to-device
from: device-to-host
for (int i=0; i< Nmax; i++) psi[i] *= u[i];
                         Runtime system keeps track of CPU & GPU memory access
                         & avoids unnecessary CPU-GPU data transfer
// De-allocate GPU array
#pragma omp target exit data map(delete:psi[0:Nmax])
                           Stand-alone directive to unmap variables from device memory
```

Where to Go from Here

- Start developing your own OpenMP target offload codes for GPU acceleration
- Plenty of room for performance optimization *e.g.*, target region executed by different threads happens concurrently, used for , *e.g.*, overlapping computation & data transfer

```
#pragma omp parallel // Spawn multiple CPU threads
{
    #pragma omp target
    {// Different GPU threads perform computation or data transfer }
}
```

See OpenMP Offload Optimization (Ye Luo, Argonne National Lab.)

See also <u>Accelerating quantum light-matter dynamics on graphics processing units</u>, T. Razach *et al.*, *Proc. IEEE IPDPSW/PDSEC* ('24)