

Parallel Molecular Dynamics

Aiichiro Nakano

*Collaboratory for Advanced Computing & Simulations
Department of Computer Science
Department of Physics & Astronomy
Department of Chemical Engineering & Materials Science
Department of Quantitative & Computational Biology
University of Southern California*

Email: anakano@usc.edu

Objective: Operationally understand spatial decomposition (who does what) & message passing using a real-world application (pmd.c)



<https://aiichironakano.github.io/cs596/src/pmd>
<https://github.com/KenichiNomura/binary-LJ-pmd>



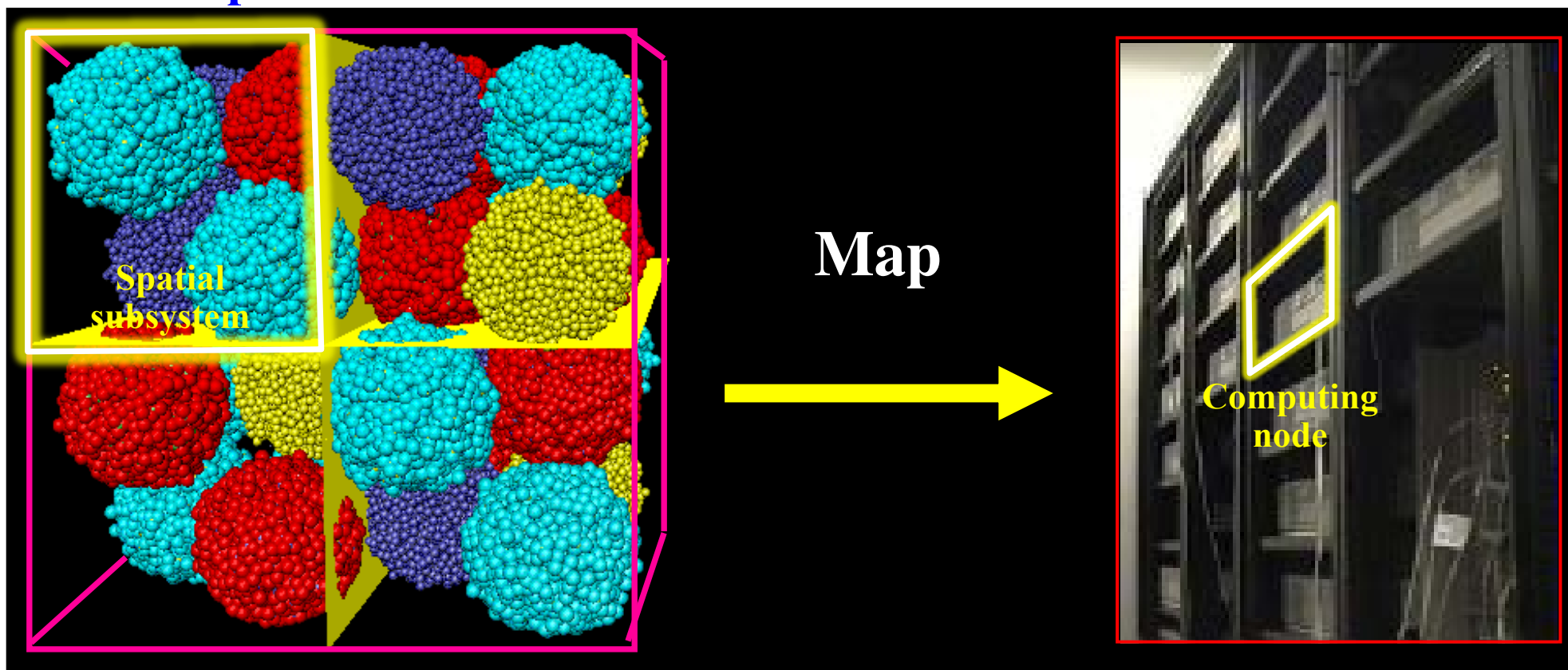
Parallel Molecular Dynamics

Spatial decomposition (short-ranged):

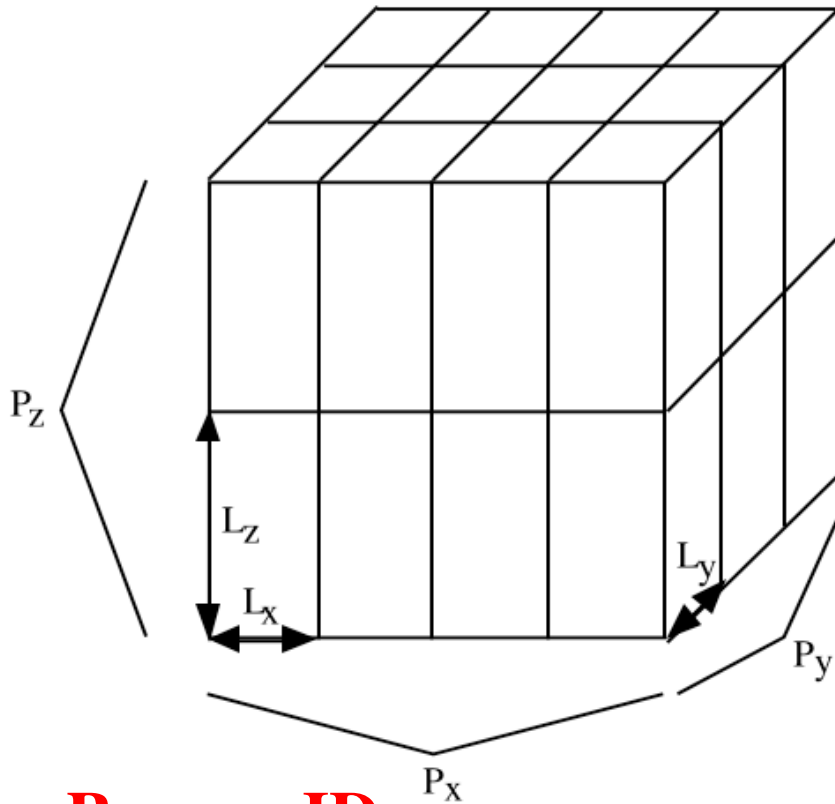
Will learn other decomposition schemes later:
<http://cacs.usc.edu/education/cs596/NT.pdf>

1. Divide the physical space into subspaces of equal volume
2. Assign each subspace to a computing node (more generally, to a process) in a parallel computer
or MPI rank
3. Each node computes forces on the atoms in its subspace & updates their positions & velocities

Who does what



Spatial Decomposition



- Process ID**

Vector

$$p_x = p / (P_y P_z)$$

$$p_y = (p / P_z) \bmod P_y$$

$$p_z = p \bmod P_z$$

Which 3D
subspace?

Scalar

$$p = p_x \times P_y P_z + p_y \times P_z + p_z$$

Rank

$$P_z - 1$$

5	11	17	23
3	9	15	21
1	7	13	19

$$P_y - 1$$

4	10	16	22
2	8	14	20
0	6	12	18

0 0 1 $P_x - 1$

P_z P_y P_x

$$nproc = vproc[0] \times vproc[1] \times vproc[2]$$

In pmd.h P_x P_y P_z

```
int vproc[3] = {1,1,2}, nproc = 2;
```

In pmd.c

```
MPI_Comm_rank(MPI_COMM_WORLD, &sid);
vid[0] = sid / (vproc[1] * vproc[2]);
vid[1] = (sid / vproc[2]) % vproc[1];
vid[2] = sid % vproc[2];
```

Neighbor Process ID

$$p'_\alpha(\kappa) = [p_\alpha + \delta_\alpha(\kappa) + P_\alpha] \bmod P_\alpha \quad (\kappa = 0, \dots, 5; \alpha = x, y, z)$$

$$p'(\kappa) = p'_x(\kappa) \times P_y P_z + p'_y(\kappa) \times P_z + p'_z(\kappa)$$

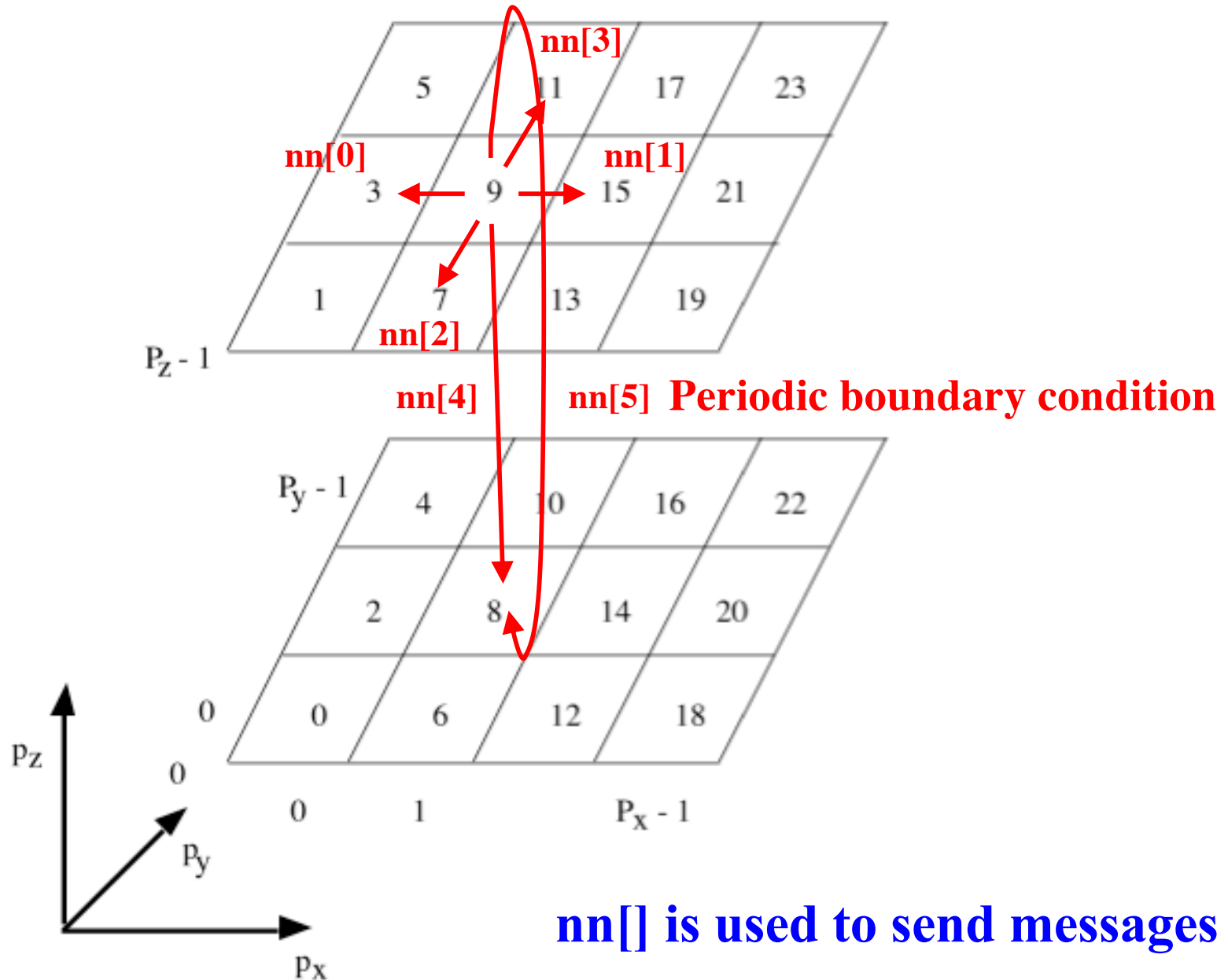
Neighbor ID, κ	$\vec{\delta} = (\delta_x, \delta_y, \delta_z)$	$\vec{\Delta} = (\Delta_x, \Delta_y, \Delta_z)$
0 (east)	(-1, 0, 0)	(-L _x , 0, 0)
1 (west)	(1, 0, 0)	(L _x , 0, 0)
2 (north)	(0, -1, 0)	(0, -L _y , 0)
3 (south)	(0, 1, 0)	(0, L _y , 0)
4 (up)	(0, 0, -1)	(0, 0, -L _z)
5 (down)	(0, 0, 1)	(0, 0, L _z)

- L_x, L_y & L_z are the box lengths *per process* in the x, y & z directions
- Atom coordinates are in the range $[0, L_\alpha]$ ($\alpha = x, y, z$) in each process

In `pmd.c`

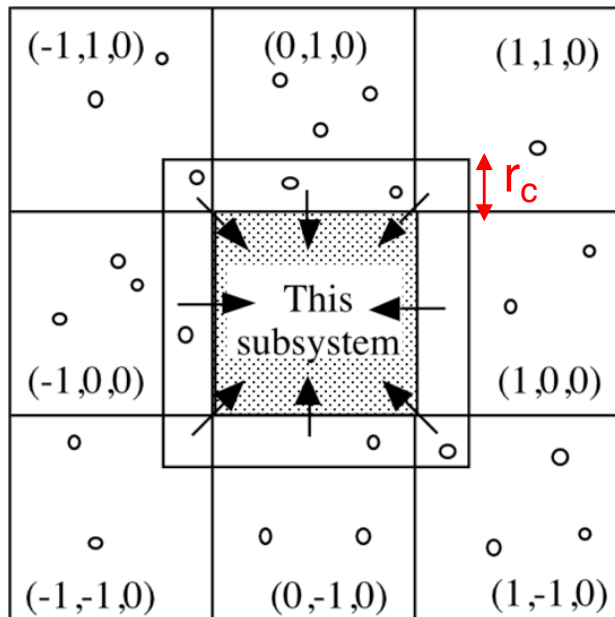
```
int iv[6][3]={{-1,0,0}, {1,0,0}, {0,-1,0}, {0,1,0}, {0,0,-1}, {0,0,1}};
...
for (ku=0; ku<6; ku++) {
    for (a=0; a<3; a++)
        k1[a] = (vid[a]+iv[ku][a]+vproc[a])%vproc[a]; Wrap around
    nn[ku] = k1[0]*vproc[1]*vproc[2]+k1[1]*vproc[2]+k1[2]; destination rank
    for (a=0; a<3; a++) sv[ku][a] = al[a]*iv[ku][a]; coordinate shift for
} self-centric parallelization
```

Neighbor Process ID Example

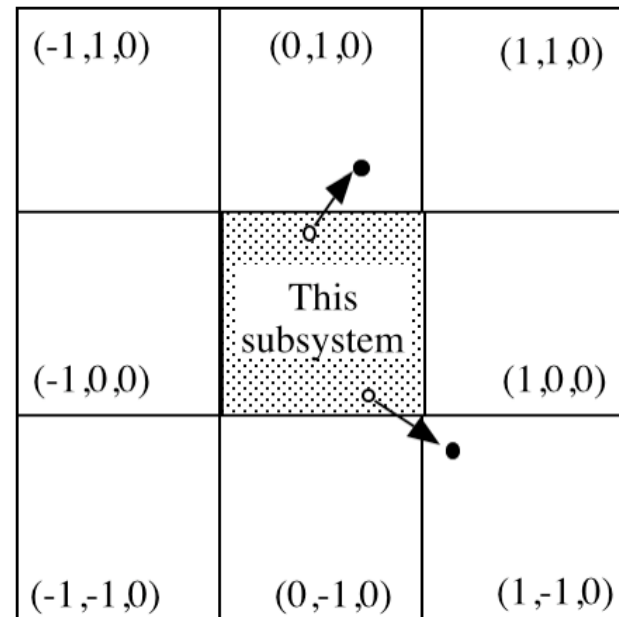


Parallel MD Concepts

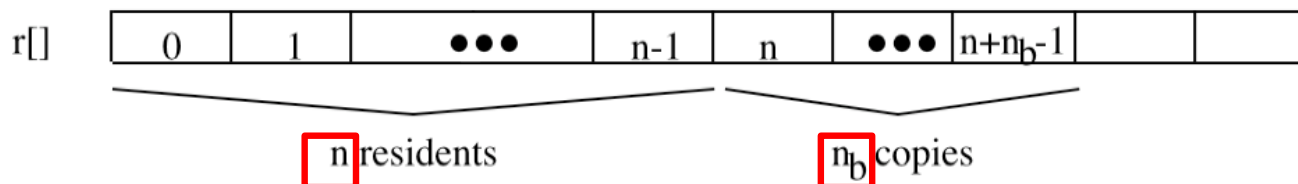
Atom caching



Atom migration



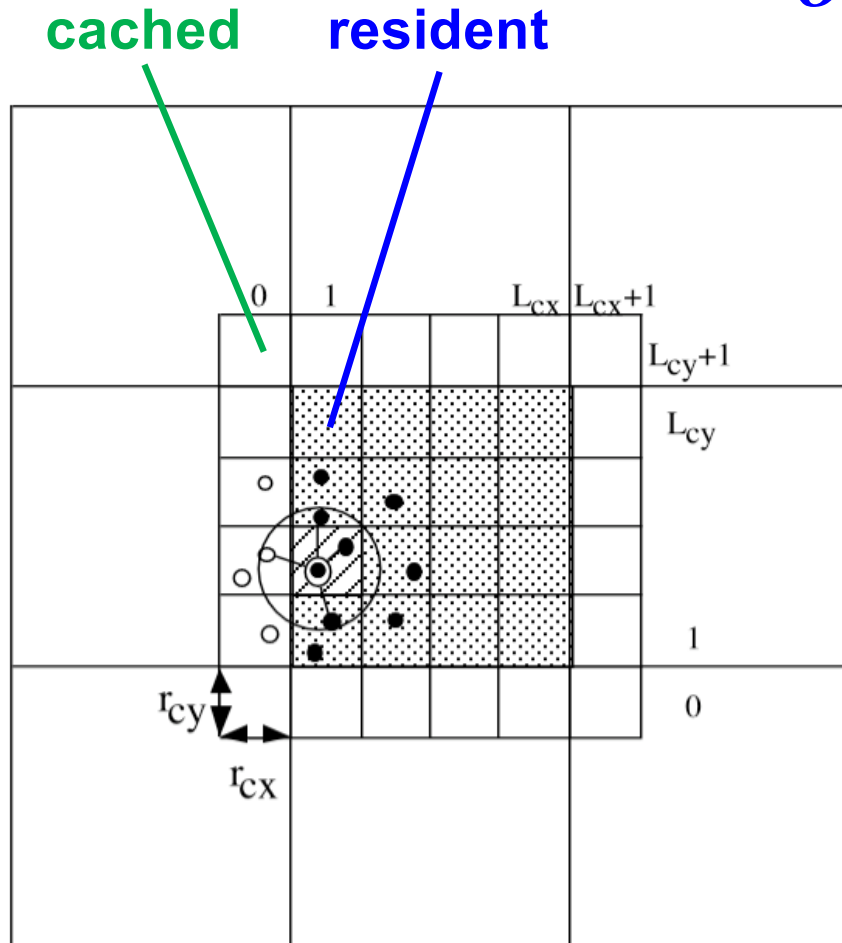
1. First half kick to obtain $v_i(t+Dt/2)$
2. Update atomic coordinates to obtain $r_i(t+Dt)$
3. **atom_move():** Migrate the moved-out atoms to the neighbor processes
4. **atom_copy():** Copy the surface atoms within distance r_c from the neighbors
5. **compute_accel():** Compute new accelerations, $a_i(t+Dt)$, including the contributions from the cached atoms
6. Second half kick to obtain $v_i(t+Dt)$



Linked-List Cell Method

Search for pairs only within the nearest neighbor cells:

$$O(N^2) \rightarrow O(N)$$



- **Cell size**

$$L_{c\alpha} = \lfloor L_{\alpha} / r_c \rfloor$$

$$r_{c\alpha} = L_{\alpha} / L_{c\alpha} \quad (\alpha = x, y, z)$$

- **Cell index**

$$c = c_x(L_{cy}+2)(L_{cz}+2) + c_y(L_{cz}+2) + c_z$$

$$c_x = c / [(L_{cy}+2)(L_{cz}+2)]$$

$$c_y = [c / (L_{cz}+2)] \bmod (L_{cy}+2)$$

$$c_z = c \bmod (L_{cz}+2)$$

- **Atom \rightarrow cell mapping**

$$c_{\alpha} = \lfloor (r_{\alpha} + r_{c\alpha}) / r_{c\alpha} \rfloor \quad (\alpha = x, y, z)$$

Only change from serial lmd.c in green:
Augmented cells to include cached atoms

List Construction Algorithm

```

/* Reset the headers, head */
for (c=0; c<lcxyz2; c++) head[c] = EMPTY;
/* Scan atoms to construct headers, head, & linked lists, lscl */
for (i=0; i<n+nb; i++) { Consider  $n_b$  cached atoms
    /* Vector cell index to which this atom belongs */
    for (a=0; a<3; a++) mc[a] = (r[i][a]+rc[a])/rc[a]; Position offset by one cell
    /* Translate the vector cell index, mc, to a scalar cell index */
    c = mc[0]*lcyz2+mc[1]*lc2[2]+mc[2];
    /* Link to the previous occupant (or EMPTY if you're the 1st) */
    lscl[i] = head[c];
    /* The last one goes to the header */
    head[c] = i;
}

```

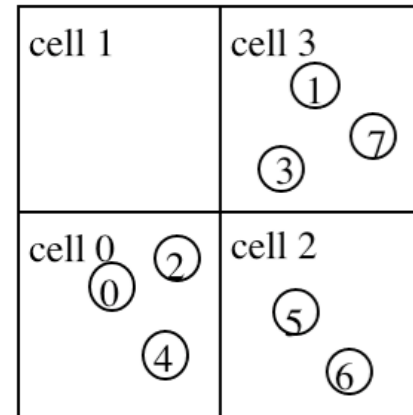
In the above:

$lcyz2 = lc2[1]*lc2[2]$

where

$lc2[a] = lc[a]+2$ ($a = 0, 1, 2$)

$lcxyz2 = lcyz2*lc2[0]$



	0	1	2	3															
head	4	E	6	7															
	0	1	2	3	4	5	6	7											
lscl	E	E	0	1	2	E	5	3											

Change from serial lmd.c in green



Interaction Computation

```

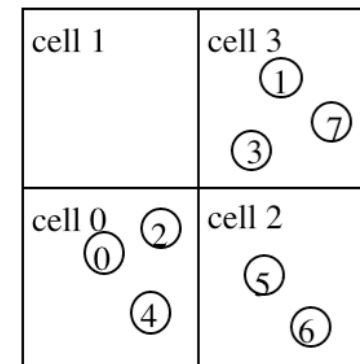
/* Scan inner cells (resident) */
for (mc[0]=1; mc[0]<=lc[0]; (mc[0])++)
for (mc[1]=1; mc[1]<=lc[1]; (mc[1])++)
for (mc[2]=1; mc[2]<=lc[2]; (mc[2])++) {
    /* Calculate a scalar cell index */
    c = mc[0]*lcyz2+mc[1]*lc2[2]+mc[2];
    /* Scan the neighbor cells (including itself) of cell c (resident + cached) */
    for (mc1[0]=mc[0]-1; mc1[0]<=mc[0]+1; (mc1[0])++)
    for (mc1[1]=mc[1]-1; mc1[1]<=mc[1]+1; (mc1[1])++)
    for (mc1[2]=mc[2]-1; mc1[2]<=mc[2]+1; (mc1[2])++) {
        /* Calculate the scalar cell index of the neighbor cell */
        c1 = mc1[0]*lcyz2+mc1[1]*lc2[2]+mc1[2];
        /* Scan atom i in cell c */
        i = head[c];
        while (i != EMPTY) {
            /* Scan atom j in cell c1 */
            j = head[c1];
            while (j != EMPTY) {
                ...
                if (i<j && rij<rc2) Process pair (i, j)
                ...
                j = lscl[j];
            }
            i = lscl[i];
        }
    }
}

```

Change from serial lmd.c in green

Who does what: Each rank computes forces on the **resident** atoms in its subspace & updates their positions & velocities

Resident atoms may interact with cached atoms (cf. slide 7)



	0	1	2	3				
head	4	E	6	7				
	0	1	2	3	4	5	6	7
lscl	E	E	0	1	2	E	5	3

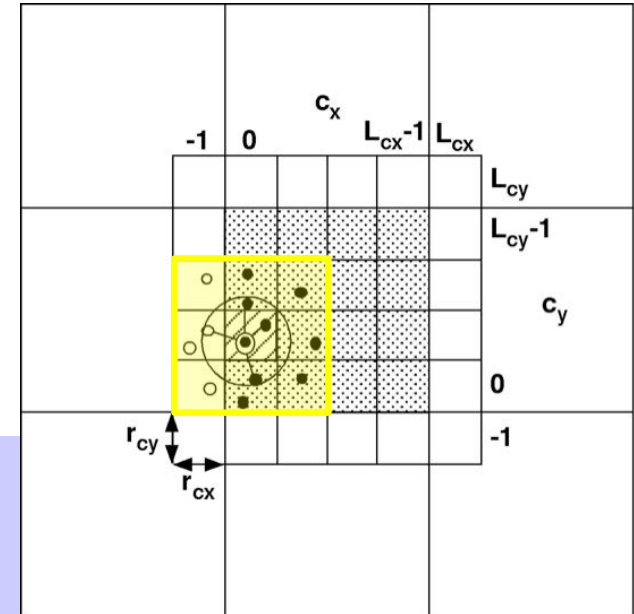


Parallel Interaction Computation

SPMD: Who does what?

Each process computes:

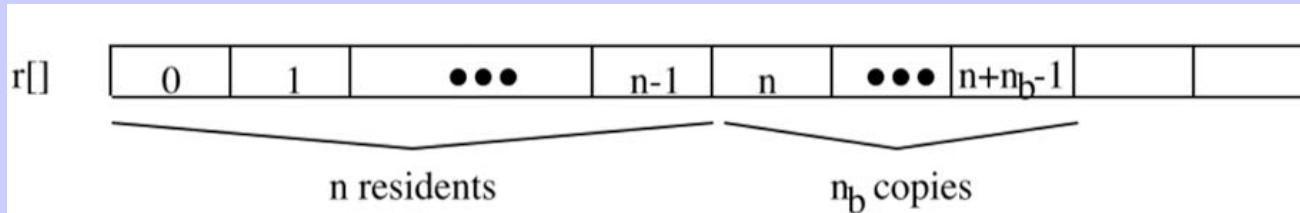
- 1. The forces on its resident atoms**
- 2. The potential energy between resident pairs & 1/2 of that between resident-cached pairs**



```

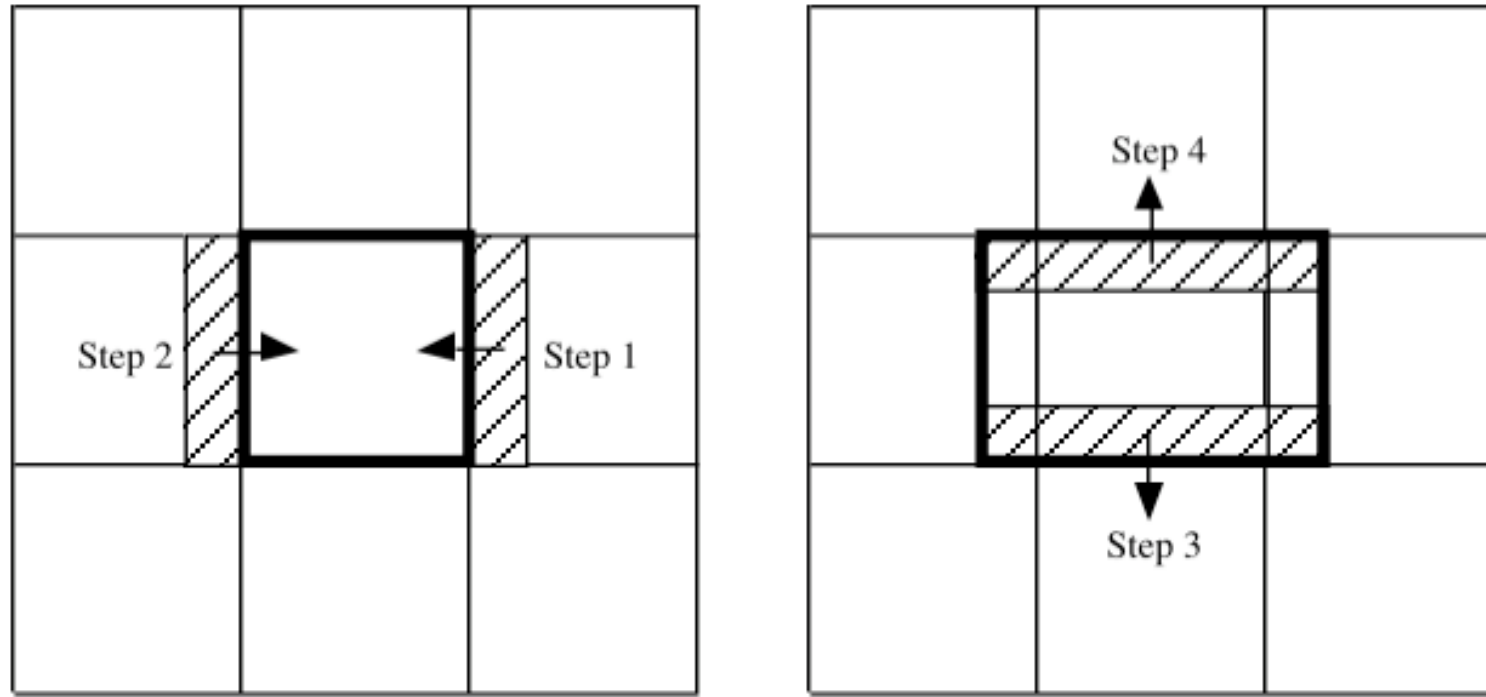
for resident cells, c {
  for neighbor (resident or cached) cells, c1 {
    scan atom i in cell c using c's linked list {
      scan atom j in cell c1 using c1's linked list {
        ...
        if (i < j && r_{ij} < r_c^2) {
          compute pair force aij & potential u(r_{ij})
          bintra = j < n; /* j is resident? */
          ai += aij; if (bintra) aj -= aij;
          if (bintra) lpe += u(r_{ij}); else lpe += u(r_{ij})/2;
        }
      }
    }
  }
}

MPI_Allreduce(&lpe, &potEnergy, ..., MPI_SUM, ...);
  
```



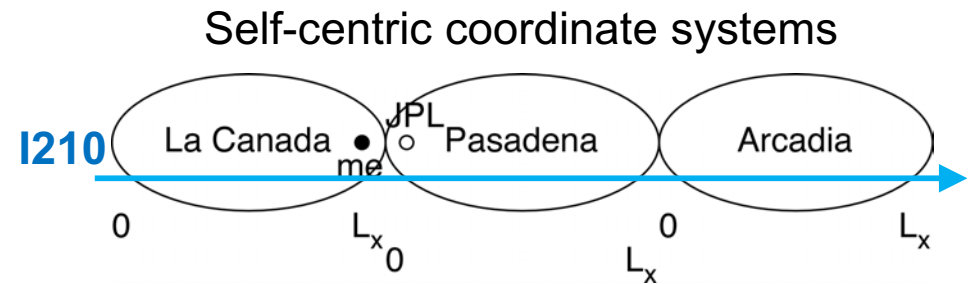
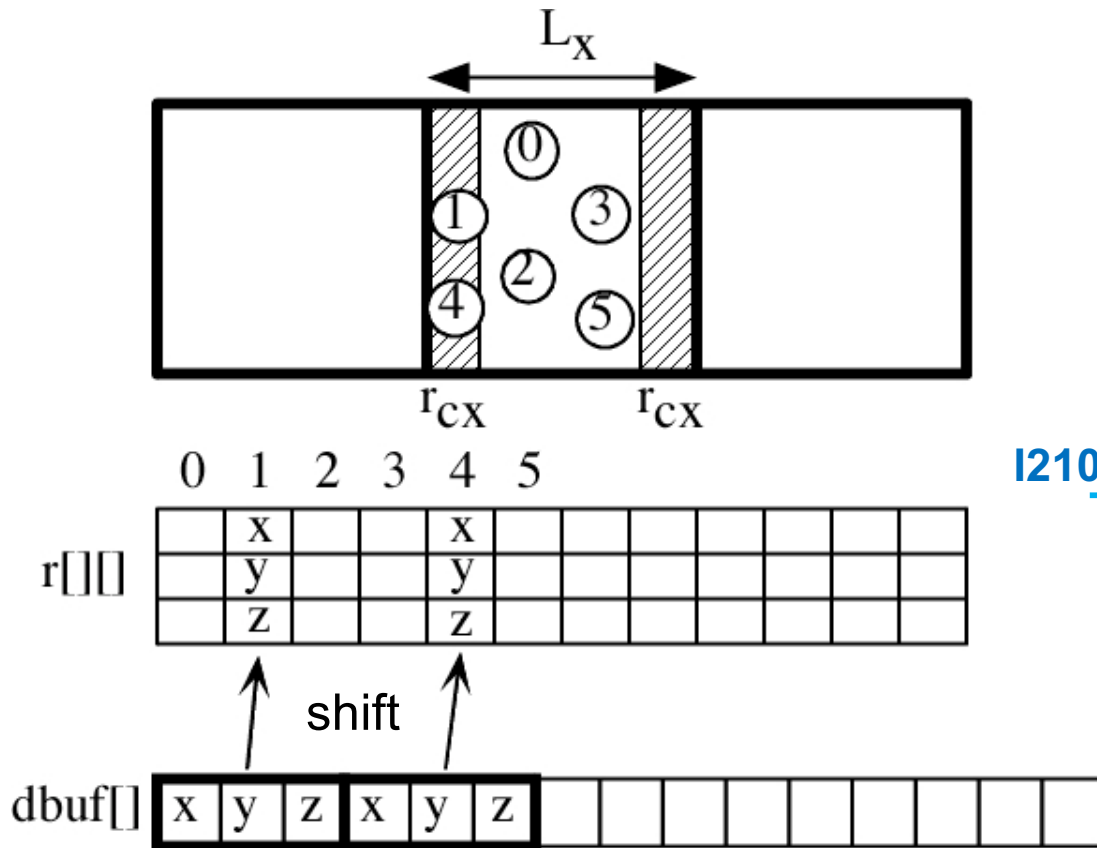
Atom Caching: `atom_copy()`

Caching from 26 neighbors in 6 steps



```
Reset the number of received cache atoms, nbnew = 0
for x, y, and z directions
  Make boundary-atom lists, lsb, for lower and higher directions including both
  resident, n, and cache, nbnew, atoms (within  $r_c$  from boundary)
  for lower and higher directions
    Send/receive boundary-atom coordinates to/from the neighbor
    Increment nbnew;
  endfor
endfor
nb = nbnew
```

Implementing Atom Caching



Copying condition

```

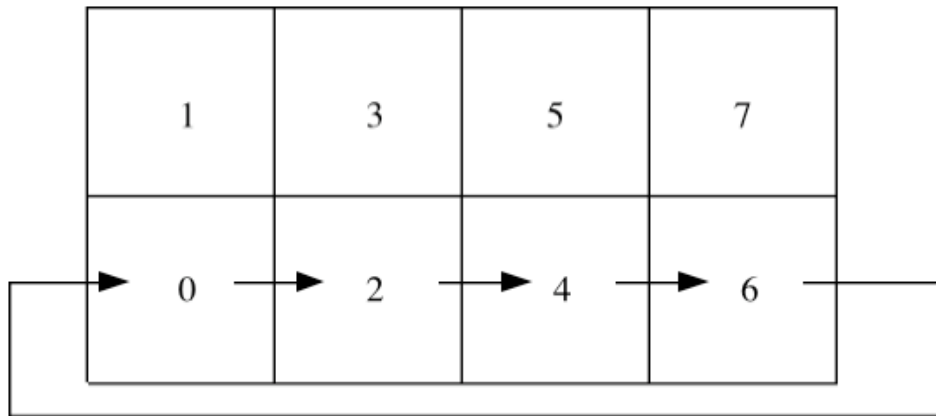
bbd(ri[], ku) {
  kd = ku / 2 (= 0|1|2) xlylz
  kdd = ku % 2 (= 0|1) lower|higher
  if (kdd == 0)
    return ri[kd] < RCUT
  else
    return al[kd] - RCUT < ri[kd]
}
    
```

3 phases of message passing

1. Message buffering: $dbuf \leftarrow r - sv$ (shift), gather
2. Message passing: $dbuf \leftarrow dbuf$
Send dbuf
Receive dbuf
3. Message storing: $r \leftarrow dbuf$, append after the residents

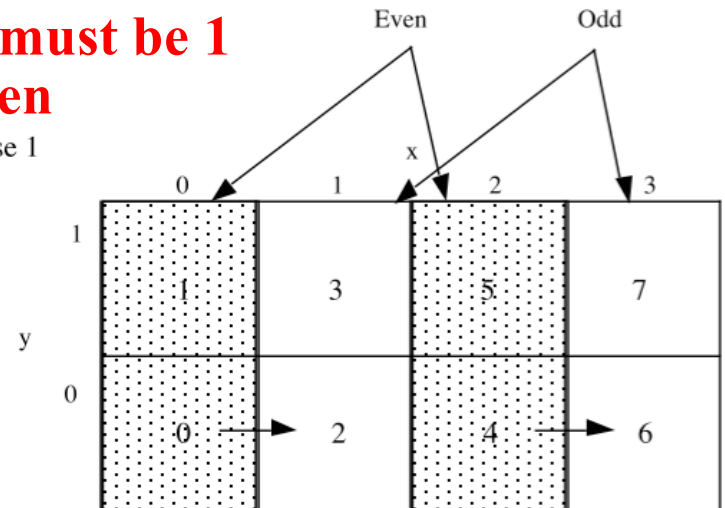
Deadlock Avoidance

Cyclic dependence

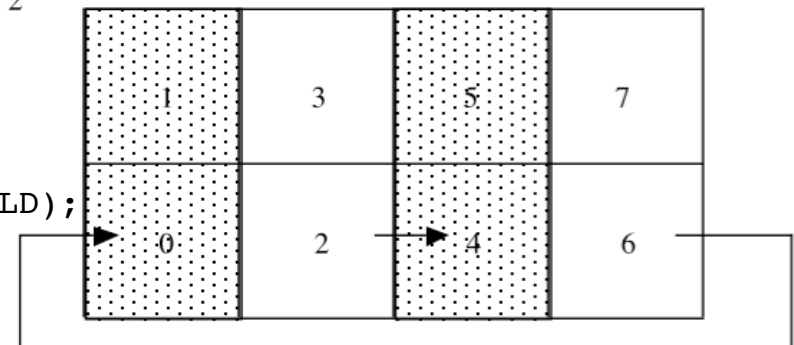


vproc[0|1|2] must be 1 or even

Phase 1



Phase 2



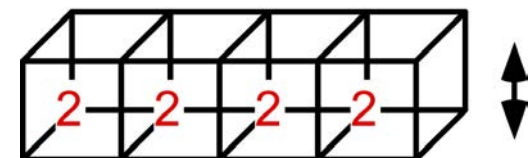
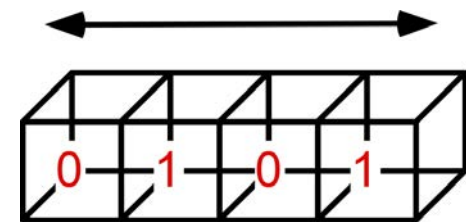
3-phase (deadlock-free) message passing

1. Message buffering: $\text{dbuf} \leftarrow r$, gather
2. Message passing: $\text{dbufr} \leftarrow \text{dbuf}$

```

/* Even node: send & recv, if not empty */
if (myparity[kd] == 0) {
    MPI_Send(dbuf, 3*nsd, MPI_DOUBLE, inode, 120, MPI_COMM_WORLD);
    MPI_Recv(dbufr, 3*nrc, MPI_DOUBLE, MPI_ANY_SOURCE, 120,
             MPI_COMM_WORLD, &status);
}
/* Odd node: recv & send, if not empty */
else if (myparity[kd] == 1) {
    MPI_Recv(dbufr, 3*nrc, MPI_DOUBLE, MPI_ANY_SOURCE, 120,
             MPI_COMM_WORLD, &status);
    MPI_Send(dbuf, 3*nsd, MPI_DOUBLE, inode, 120, MPI_COMM_WORLD);
}
/* Single layer: Exchange information with myself */
else
    for (i=0; i<3*nrc; i++) dbufr[i] = dbuf[i];

```
3. Message storing: $r \leftarrow \text{dbufr}$, append



ANL IBM SP1 User's Guide ('94)

11. Q: My parallel program runs on other parallel machines but seems to deadlock on the SP-1 when using EUI, EUI-H, or Chameleon.

A: The following parallel program can deadlock on *any* system when the size of the message being sent is large enough:

```
send( to=partner, data, len, tag )
recv( from=partner, data, maxlen, tag )
```

where these are blocking send's and receives (`mp_bsend` in EUI/EUI-H and `PIbsend` in Chameleon). For many systems, deadlock does not occur until the message is very long (often 128 KBytes or more). For EUI, the size is (roughly) 128 bytes (*not* KBytes) and for EUI-H, the size is (again roughly) 4 KBytes. The limit for Chameleon is the same as the underlying transport layer (i.e., the EUI or EUI-H limits).

To fix this you have several choices:

pmd.c • Reorder your send and receive calls so that they are pair up. For example, if there are always an even number of processors, you could use

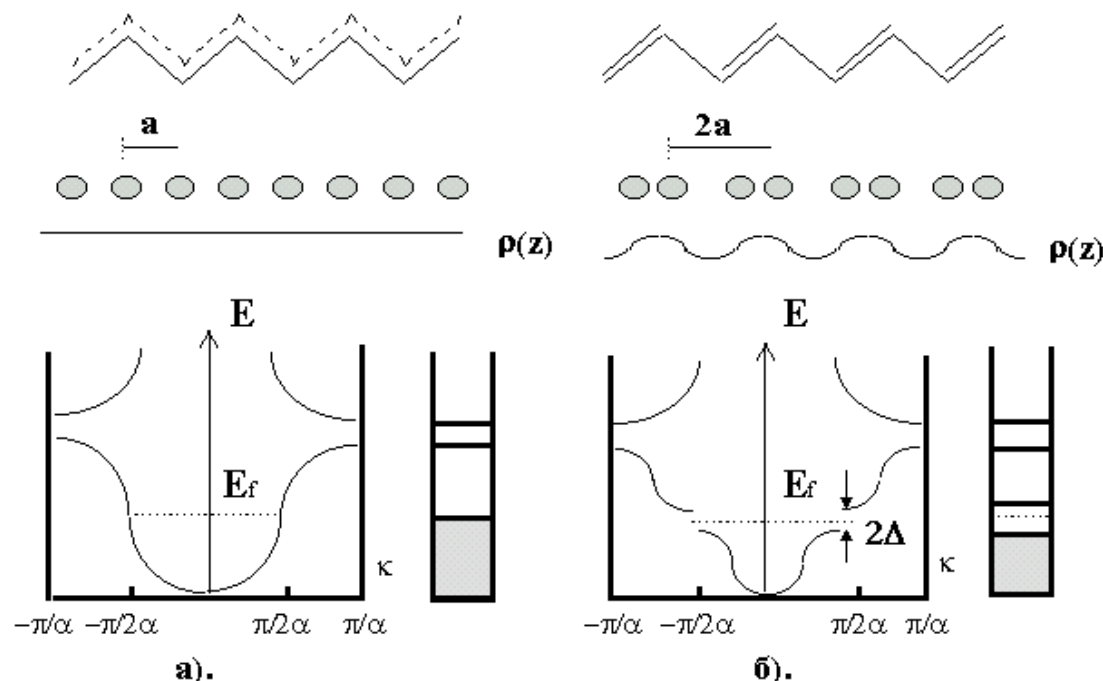
```
if (myid is even) {
    send( to=partner, data, len, tag )
    recv( from=partner, data, maxlen, tag )
}
else {
    recv( from=partner, data, maxlen, tag )
    send( to=partner, data, len, tag )
}
```

assignment • Use non-blocking sends and receives instead



```
MPI_Irecv();
MPI_Send();
MPI_Wait();
```

Digress: Polyacetylene & Peierls Distortion



Alan J. Heeger
Prize share: 1/3



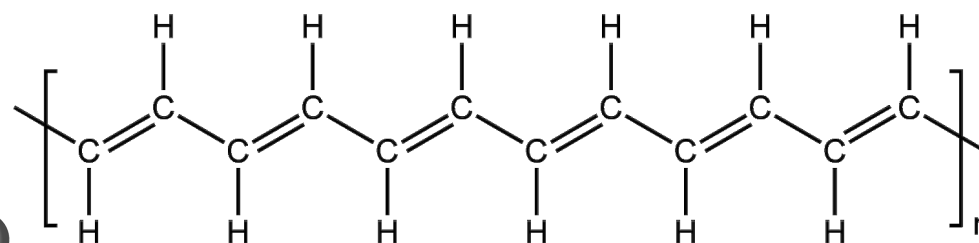
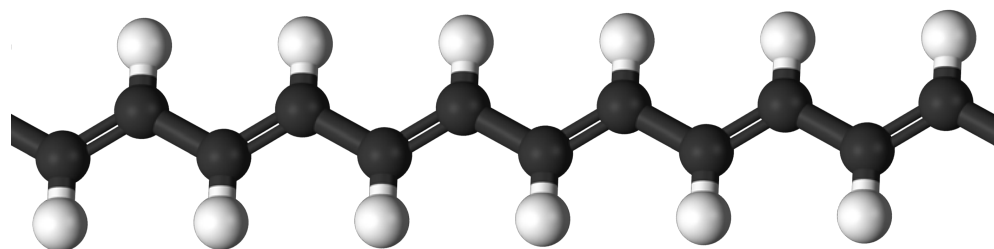
Alan G. MacDiarmid
Prize share: 1/3



Hideki Shirakawa
Prize share: 1/3

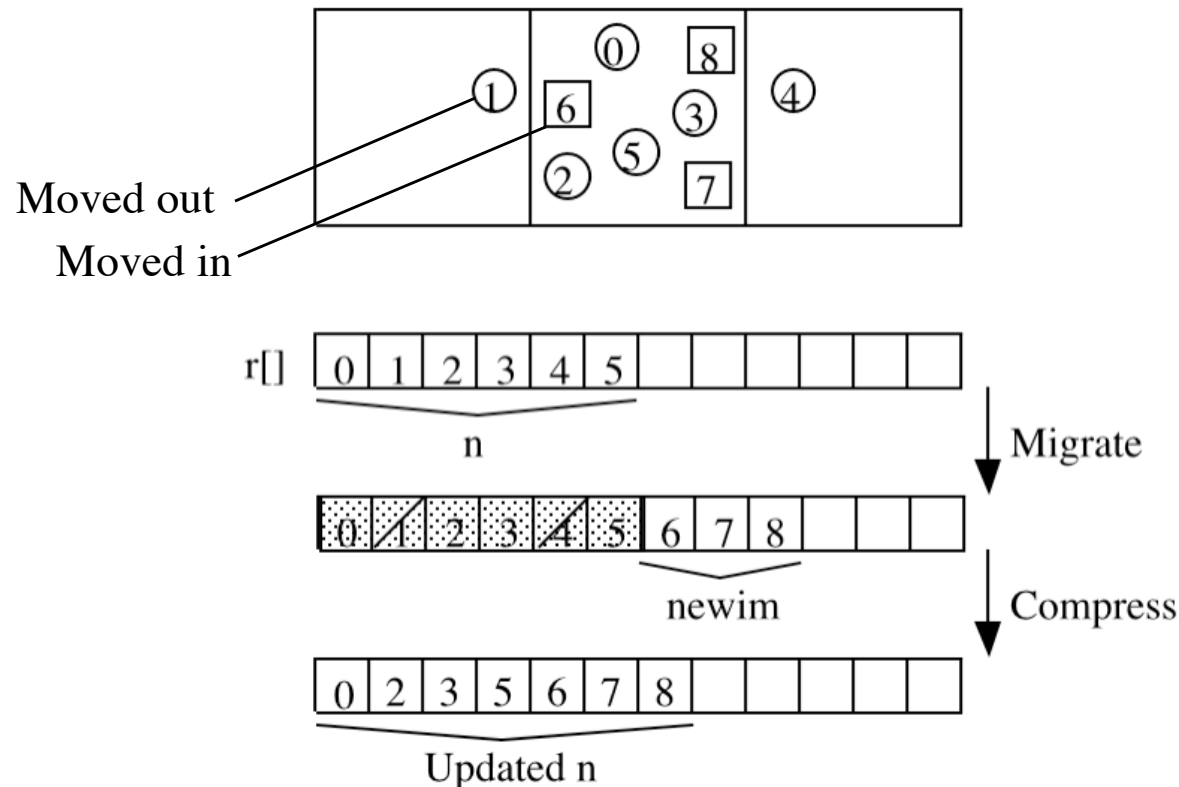
The Nobel Prize in Chemistry 2000 was awarded jointly to Alan J. Heeger, Alan G. MacDiarmid and Hideki Shirakawa "for the discovery and development of conductive polymers".

Fig. 1. Electron dispersion and a band pattern of one-dimensional molecular system: a). metallic and b). insulator state, ($\rho(z)$ —a electronic density, a —a lattice period).



Nature's spontaneous even-odd
symmetry breaking

Atom Migration: `atom_move()`



Reset the number of received new immigrants, `newim = 0`

for x, y, and z directions

Make moving-atom lists, `mvque`, for lower and higher directions including both resident, `n`, and immigrant, `newim`, atoms but excluding those already moved out for lower and higher directions

Send/receive moving-atom coordinates to/from the neighbor

(When moving, `r[][0] ← MOVED_OUT = -1010`)

Increment `newim`

endfor

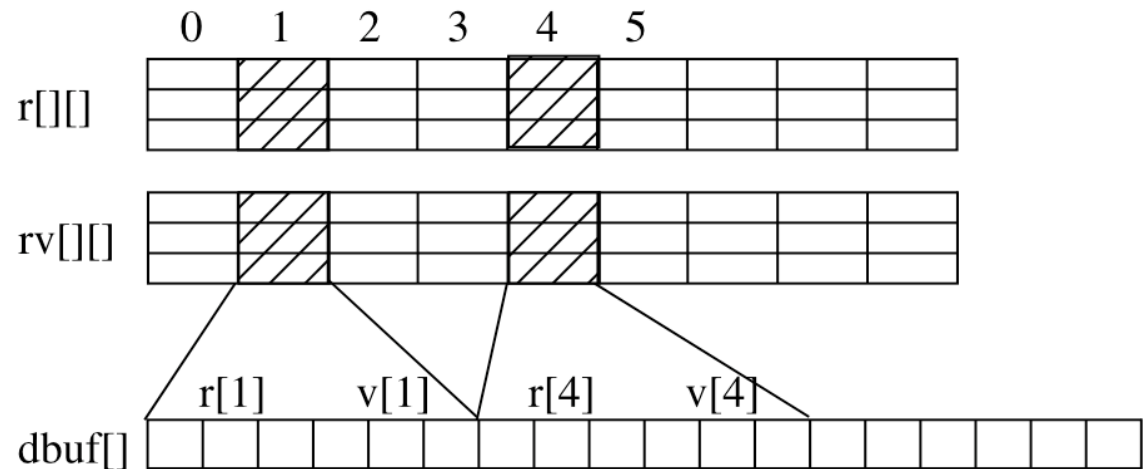
endfor

Compress the `r` array to eliminate the moved-out atoms

Implementing Atom Migration

Moving condition

```
bmv(ri[],ku) {  
    kd = ku / 2 (= 0|1|2)  
    kdd = ku % 2 (= 0|1)  
    if (kdd == 0)  
        return ri[kd] < 0.0  
    else  
        return al[kd] < ri[kd]  
}
```



3 phases of message passing

1. **Message buffering:** `dbuf` \leftarrow `r-sv` (shift) & `rv`, gather
Mark `MOVED_OUT` in `r`
2. **Message passing:** `dbufr` \leftarrow `dbuf`
Send `dbuf`
Receive `dbufr`
3. **Message storing:** `r` & `rv` \leftarrow `dbufr`, append after the residents

Bottom Line: Parallel MD

Parallel computing:

Specifies who does what — decomposition

Parallel molecular dynamics (spatial decomposition):

Who does what = each processor computes forces on only resident atoms in the subspace assigned to it & update their positions & velocities

Scalability Metrics for Parallel Molecular Dynamics

Aiichiro Nakano

*Collaboratory for Advanced Computing & Simulations
Department of Computer Science
Department of Physics & Astronomy
Department of Chemical Engineering & Materials Science
Department of Biological Sciences
University of Southern California*

Email: anakano@usc.edu

Objective: Consolidate your understanding of scalability analysis
(e.g., fixed-problem vs. isogranular scaling) using a real-world
example of pmd.c



Recap: Parallel Efficiency

Parallel computing = solving a big problem (W) in a short time (T) using many processors (P)

- **Execution time:** $T(W, P)$
 W : Workload
 P : Number of processors

- **Speed:** $S(W, P) = \frac{W}{T(W, P)}$

- **Speedup:** $S_P = \frac{S(W_P, P)}{S(W_1, 1)} = \frac{W_P T(W_1, 1)}{W_1 T(W_P, P)}$

- **Efficiency:** $E_P = \frac{S_P}{P} = \frac{W_P T(W_1, 1)}{P W_1 T(W_P, P)}$

How to scale W_P with P ?

Fixed Problem-Size (Strong) Scaling

Solve the same problem faster

$W_P = W$ —constant (strong scaling)

- **Speedup:** $S_P = \frac{W_P T(W_1, 1)}{W_1 T(W_P, P)} = \frac{T(W, 1)}{T(W, P)}$
- **Efficiency:** $E_P = \frac{T(W, 1)}{P T(W, P)}$
- **Amdahl's law:** f (= sequential fraction of the workload) limits the asymptotic speedup

$$T(W, P) = fT(W, 1) + \frac{(1-f)T(W, 1)}{P}$$
$$\therefore S_P = \frac{T(W, 1)}{T(W, P)} = \frac{1}{f + (1-f)/P}$$
$$\therefore S_P \rightarrow \frac{1}{f} \quad (P \rightarrow \infty)$$

Isogranular (Weak) Scaling

Solve a larger problem within the same time duration

$W_P = Pw$ (weak scaling)

w = constant workload per processor (granularity)

- **Speedup:**
$$S_P = \frac{S(P \bullet w, P)}{S(w, 1)} = \frac{P \bullet w / T(P \bullet w, P)}{w / T(w, 1)} = \frac{P \bullet T(w, 1)}{T(P \bullet w, P)}$$

- **Efficiency:**
$$E_P = \frac{S_P}{P} = \frac{T(w, 1)}{T(P \bullet w, P)}$$

Analysis of Parallel MD

- Parallel execution time:**

Workload \propto Number of atoms, N (linked-list cell algorithm)

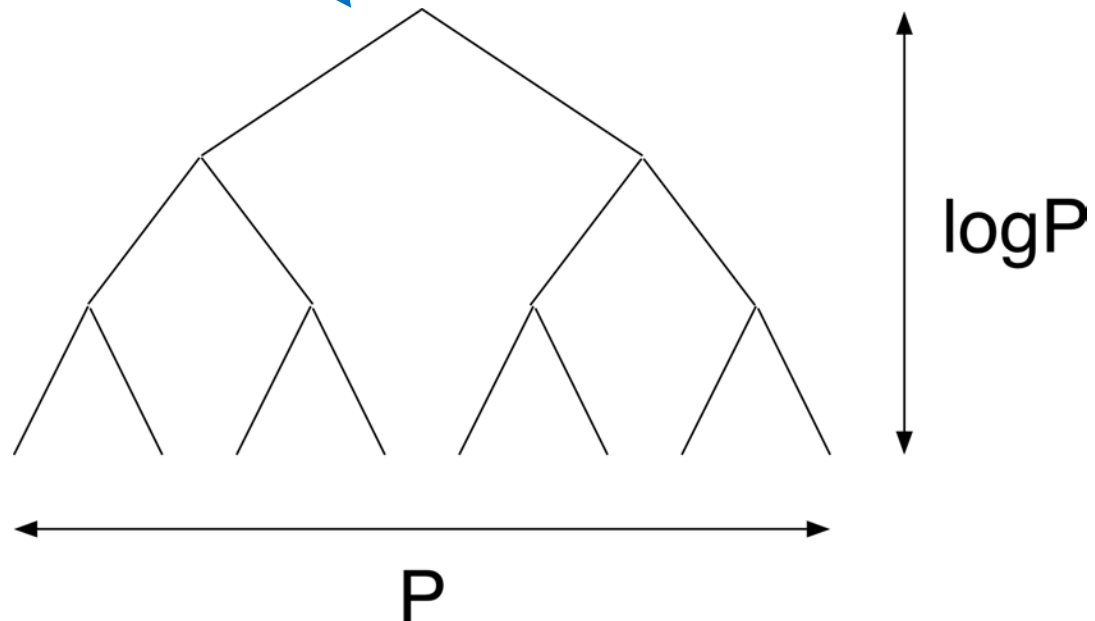
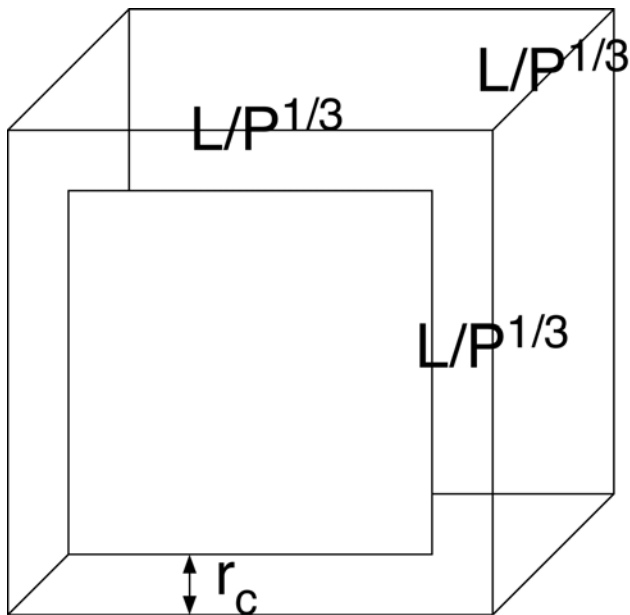
$$T(N, P) = T_{\text{comp}}(N, P) + T_{\text{comm}}(N, P) + T_{\text{global}}(P)$$

MPI_Allreduce()

$$= a \frac{N}{P} + b \left(\frac{N}{P} \right)^{2/3} + c \log P$$

$$\begin{aligned} & \text{facets } \tilde{6} \times \overbrace{\frac{L^2}{P^{2/3}} r_c}^{\text{cached volume}} \times \text{atom density } \tilde{\rho} \\ &= 6r_c \frac{N^{2/3} / \rho^{2/3}}{P^{2/3}} \rho \\ &= 6r_c \rho^{1/3} \left(\frac{N}{P} \right)^{2/3} \end{aligned}$$

$$\left(\because \frac{N}{L^3} = \rho \Rightarrow L^2 = \frac{N^{2/3}}{\rho^{2/3}} \right)$$



Fixed Problem-Size Scaling

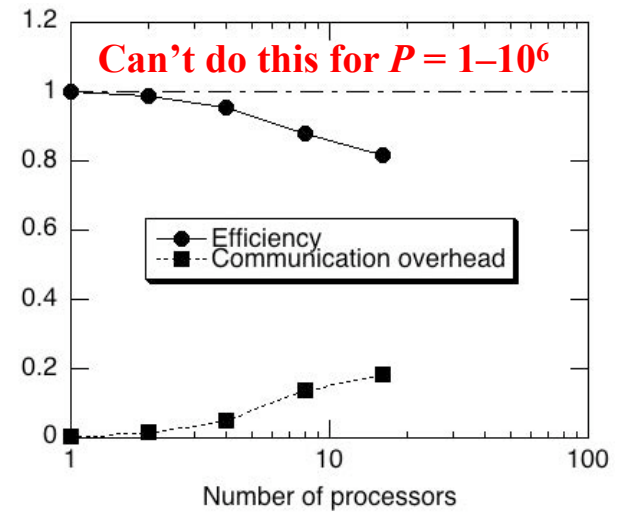
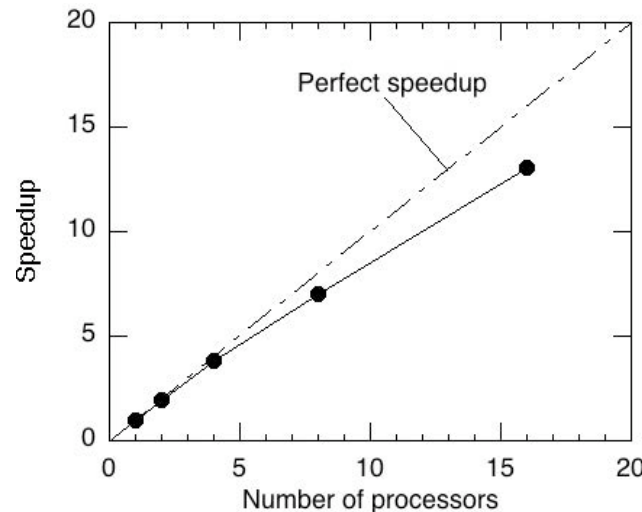
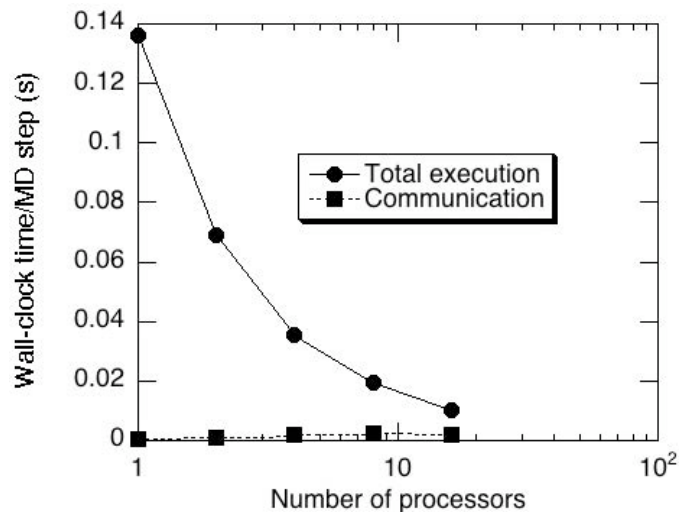
- **Speedup:**

$$S_P = \frac{T(N,1)}{T(N,P)} = \frac{aN}{a(N/P) + b(N/P)^{2/3} + c \log P}$$

$$= \frac{1}{1 + \frac{b}{a} \left(\frac{P}{N}\right)^{1/3} + \frac{c}{a} \frac{P \log P}{N}}$$

- **Efficiency:**

$$E_P = \frac{S_P}{P} = \frac{1}{1 + \frac{b}{a} \left(\frac{P}{N}\right)^{1/3} + \frac{c}{a} \frac{P \log P}{N}}$$

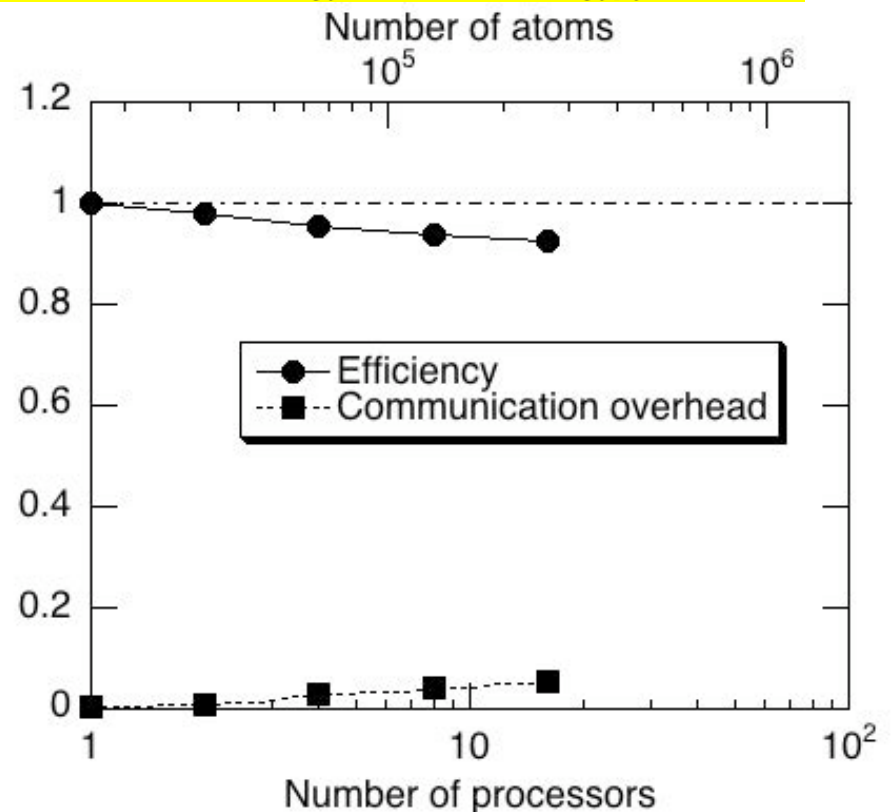
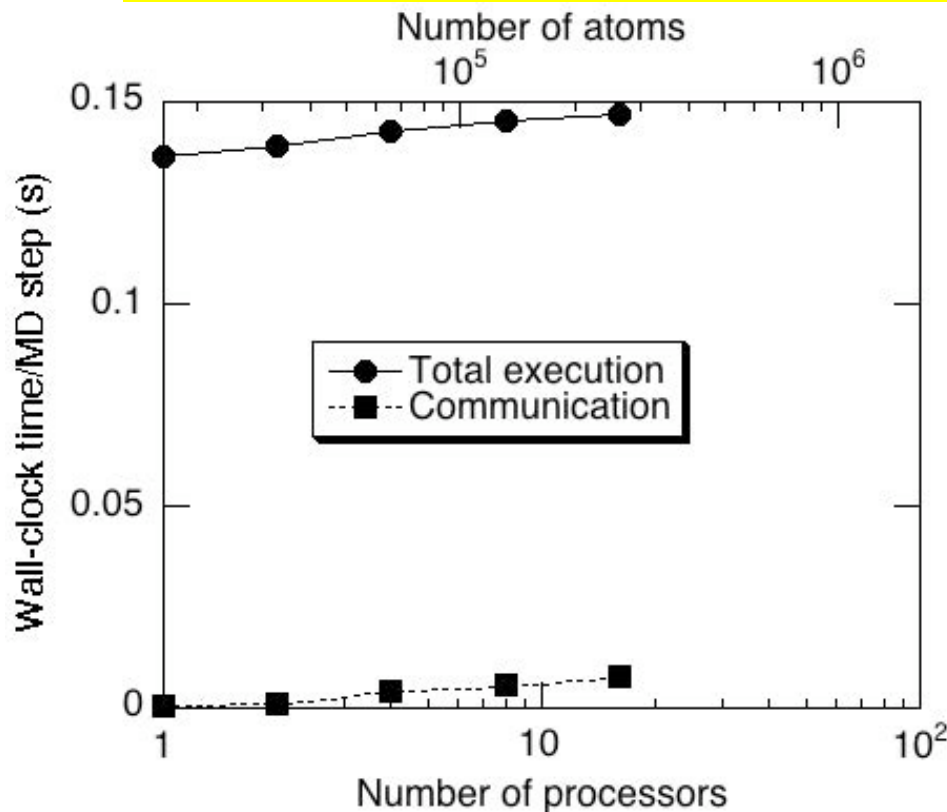


pmd.c: $N = 16,384$, on HPC (predecessor of CARC)

Isogranular Scaling of Parallel MD

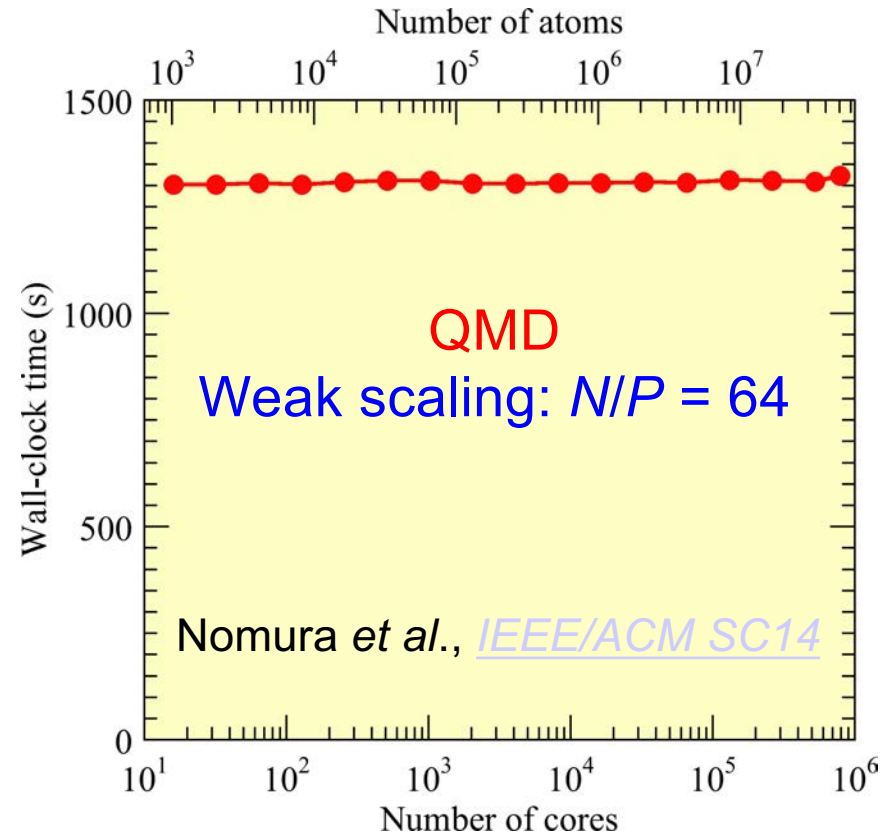
- $n = N/P = \text{constant}$: doable for arbitrarily large P
- Efficiency:

$$E_P = \frac{T(n,1)}{T(nP,P)} = \frac{an}{an + bn^{2/3} + c \log P} = \frac{1}{1 + \frac{b}{a}n^{-1/3} + \frac{c}{an} \log P}$$



pmd.c: $N/P = 16,384$, on HPC (predecessor of CARC)

High-End Parallel MD



- 4.9 trillion-atom space-time multiresolution MD (MRMD) of SiO_2
 - 8.5 billion-atom fast reactive force-field (F-ReaxFF) RMD of RDX
 - 39.8 trillion grid points (50.3 million-atom) DC-DFT QMD of SiC
- parallel efficiency 0.984 on 786,432 Blue Gene/Q cores

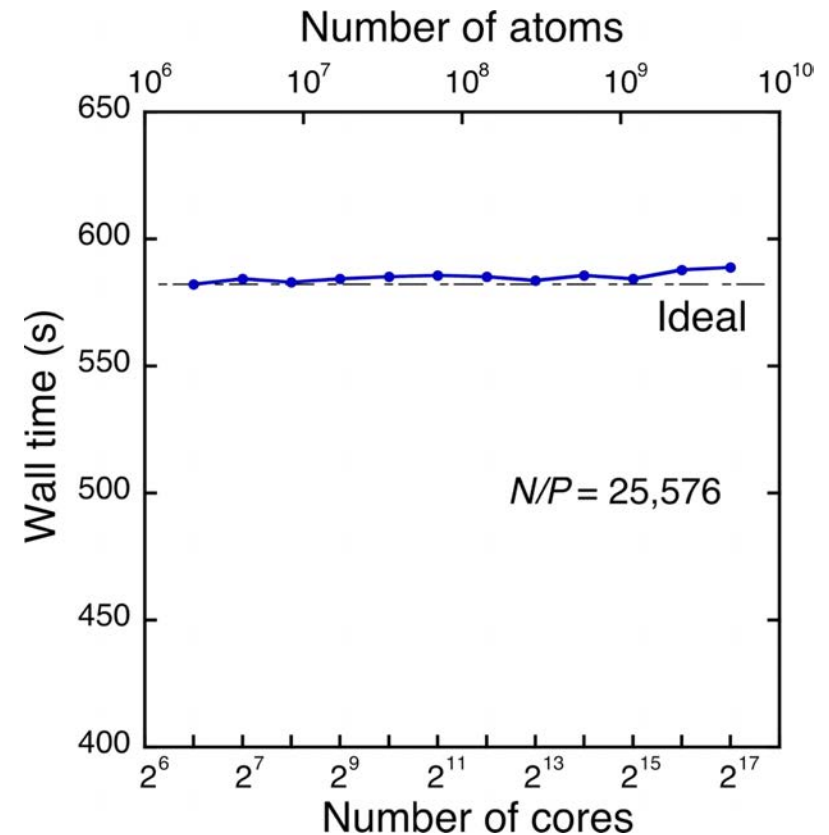
MD (molecular dynamics): MRMD

RMD (reactive molecular dynamics): F-ReaxFF

QMD (quantum molecular dynamics): DC-DFT

Portable Parallel Efficiency

- Weak-scaling parallel efficiency of 0.989 for a new generation of reactive molecular dynamics (RMD) on 131,072 Intel Knights Landing cores on Theta supercomputer at Argonne National Laboratory



K. Liu *et al.*, [Shift-collapse acceleration of generalized polarizable reactive molecular dynamics for machine learning-assisted computational synthesis of layered materials](#),
Proc. ScalA18 (IEEE, '18)

Quantum MD@Scale

Quantum dynamics at scale: ultrafast control of emergent functional materials

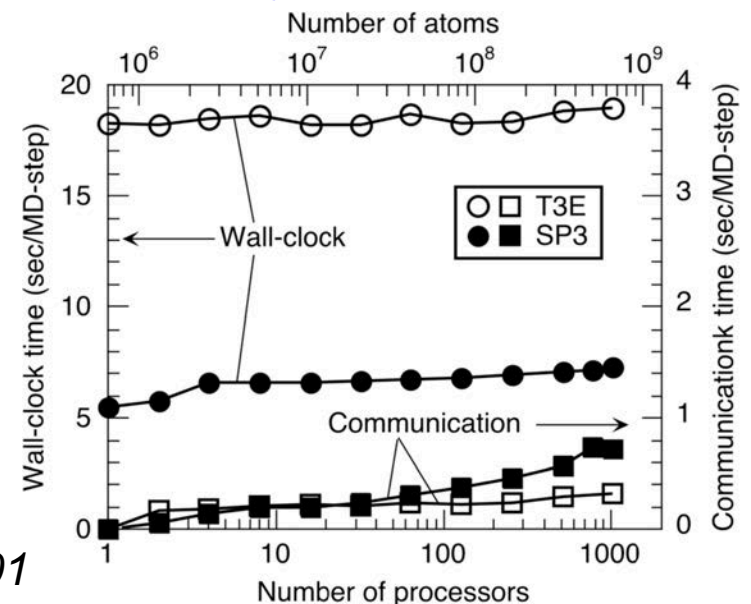
S. C. Tiwari, P. Sakdhnagool, R. K. Kalia, A. Krishnamoorthy, M. Kunaseth, A. Nakano, K. Nomura, P. Rajak, F. Shimojo, Y. Luo & P. Vashishta

Best Paper in *ACM HPC Asia 2020*



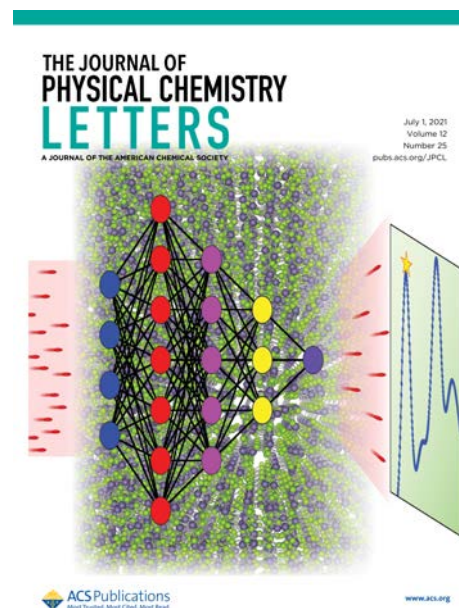
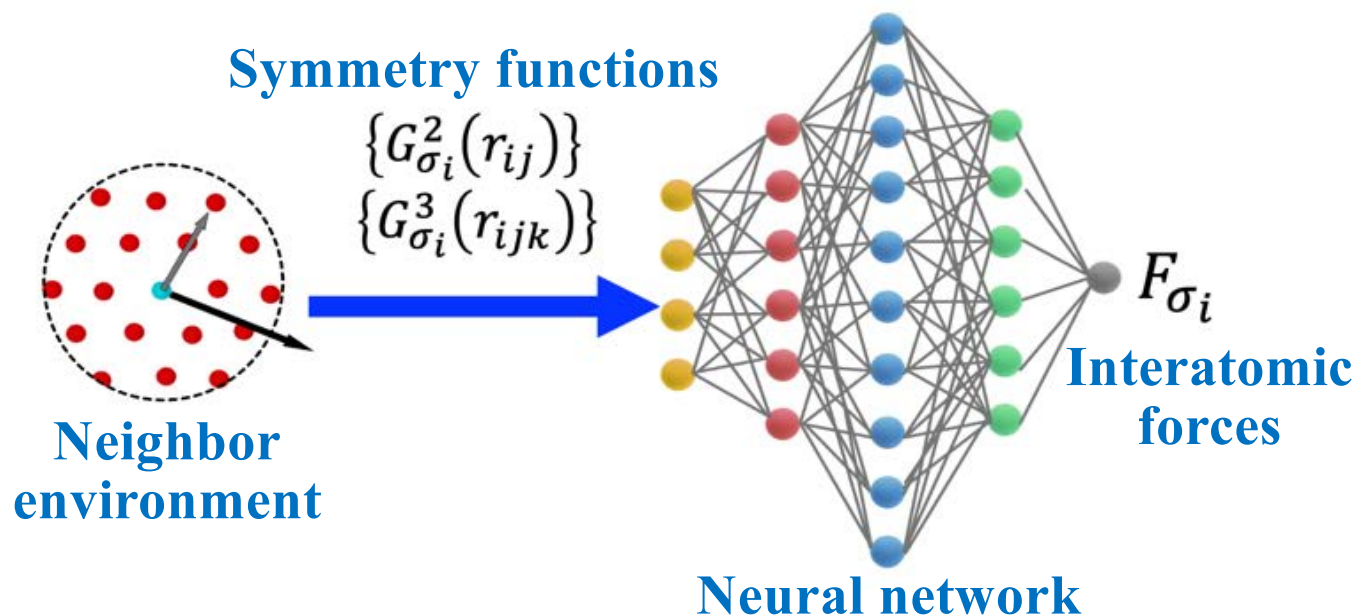
[Scalable atomistic simulation algorithms for materials research](#), A. Nakano *et al.*,
Best Paper, *IEEE/ACM Supercomputing 2001*, SC01

19 years since



Neural MD@Scale

- **Neural-network quantum molecular dynamics (NNQMD) could revolutionize atomistic modeling of materials, providing quantum-mechanical accuracy at a fraction of computational cost** [*Phys. Rev. Lett.* **126**, 216403 ('21); *J. Phys. Chem. Lett.* **12**, 6020 ('21)]



Neural network molecular dynamics at scale & Ex-NNQMD: extreme-scale neural network quantum molecular dynamics,

P. Rajak et al., *IEEE IPDPS ScaDL 20 & 21*

See also Pushing the limit of molecular dynamics with ab initio accuracy to 100 million atoms with machine learning

W. Jia et al., *ACM/IEEE Supercomputing, SC20*