

Lecture: Quick sort and comparator

Agenda

- Sort 0 and 1.
- Quick sort
- Comparator problems.

Qn-1 Given an array $A[n]$ of 0's and 1's in random order.
Sort the given array.

ip:

0	1	0	0	1	1	0	1	0
---	---	---	---	---	---	---	---	---

op:

0	0	0	0	0	1	1	1	1
---	---	---	---	---	---	---	---	---

Brute force: Bubble sort — $O(n^2)$

Insertion " — $O(n^2)$

Array sort(arr) — $O(n \log n)$

Count sort: $O(n)$, $O(1)$
freq[2]

└─ Iterate array to build freq[]
└─ " freq[] for my answer.

Approach 2

0	1	0	0	1	1	0	1	0
---	---	---	---	---	---	---	---	---

count0 = 5

count1 = 4

0	0	0	0	0	1	1	1	1
---	---	---	---	---	---	---	---	---

Approach 3

TC: $O(n)$
SC: $O(1)$ } Expectation
Only 1 iteration

0	1	2	3	4	5	6	7	8
0	1 0	0	0	1 0	1	0 1	1	0 1
i	i	i	i	i	j	j	j	j

$$i = 0$$

$$j = 1$$

Algorithm

```
void sort01 (int[] A) {  
    i = 0;  
    j = A.length - 1;  
    while (i < j) {  
        if (arr[i] == 0) {  
            i++;  
        } else {  
            if (arr[j] == 0) {  
                swap(A, i, j);  
            }  
            j--;  
        }  
    }  
}
```

TC: $O(n)$

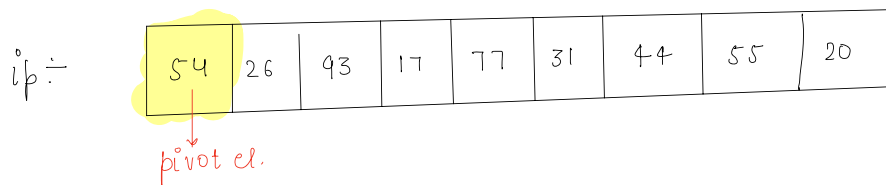
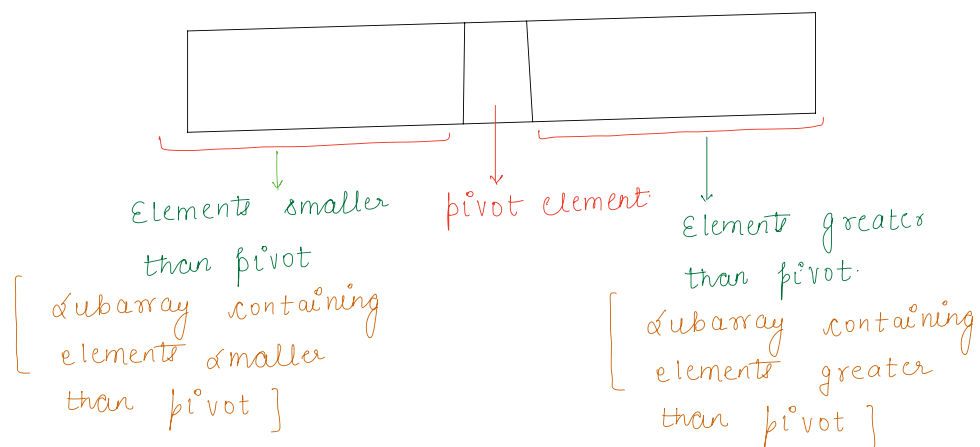
SC: $O(1)$

~~##~~ itr = 1 itr

Qu. Given $A[]$, consider first element as pivot, rearrange the elements such that for all i :-

- 1) if $A[i] < p$, then it should be present on left side.
- 2) if $A[i] > p$, then it should be present on right side.

Note: All elements are distinct.



After partitioning:-

26	17	31	44	20	54	93	77	55
----	----	----	----	----	----	----	----	----

Approach

ip :-

54	26	93	17	77	31	44	55	20
----	----	----	----	----	----	----	----	----

↓ pivot i j

54	26	93	17	77	31	44	55	20
----	----	----	----	----	----	----	----	----

i j

26 < 54 i++
93 > 54 stop

54	26	93	17	77	31	44	55	20
----	----	----	----	----	----	----	----	----

i j

20 < 54. stop

54	26	20 93	17	77	31	44	55	20 93
----	----	--------------------------------	----	----	----	----	----	--------------------------------

i j j

swap(20, 93)

54	26	20 93	17	44 77	31	77 44	55	20 93
----	----	--------------------------------	----	--------------------------------	----	--------------------------------	----	--------------------------------

i i j i j j

77 > 54 stop
55 > 54.
44 < 54.

31 54	26	20 93	17	44 77	54 31	77 44	55	20 93
--------------------------------	----	--------------------------------	----	--------------------------------	--------------------------------	--------------------------------	----	--------------------------------

j i

Last step = swap(A, 0, j)

Algorithm

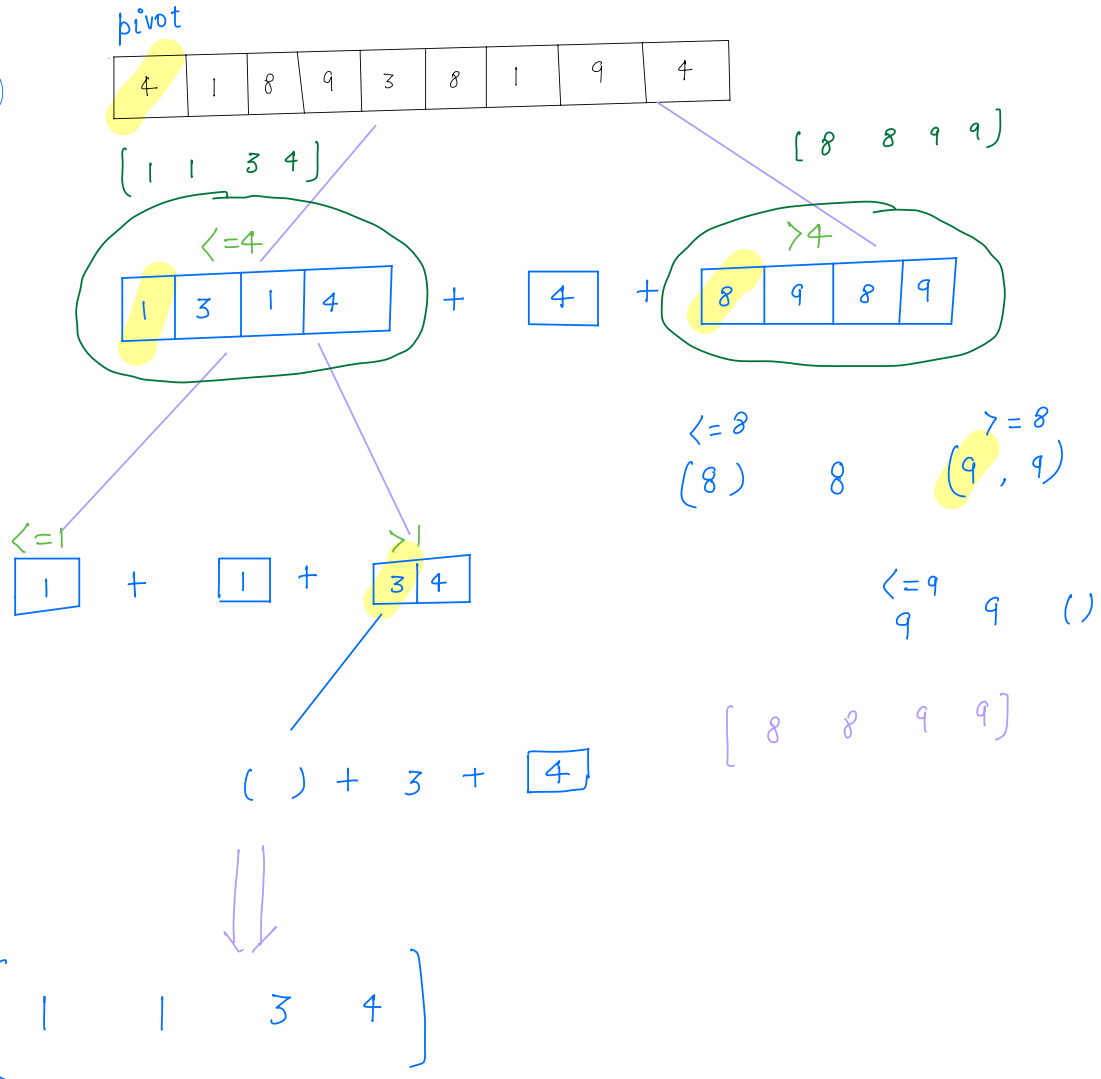
```
void partition(A[], 0start, A.length-1end) {  
    pivotvalue = A[start];  
    i = start+1;  
    j = end;  
    while(i <= j) {  
        if(A[i] <= pivotvalue) {  
            i++;  
        } else if(A[j] > pivotvalue) {  
            j--;  
        } else {  
            swap(A, i, j);  
            j--;  
            i++;  
        }  
    }  
    swap(A, start, j);  
}
```

TC: $O(n)$

SC: $O(1)$

Quick sort

Dryrun ①

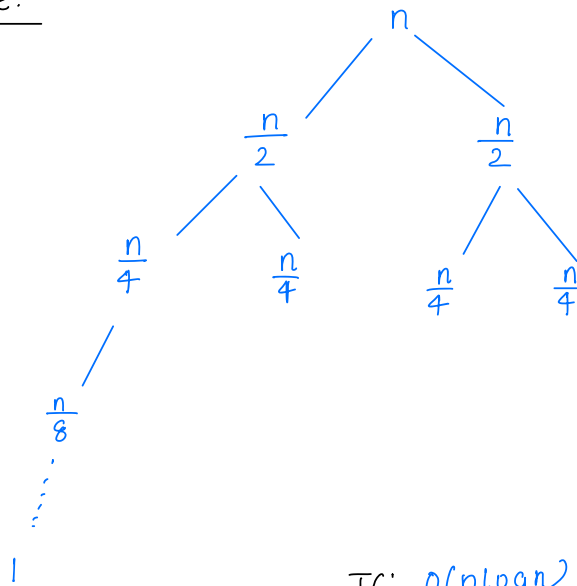


Algorithm

```
void quicksort(int[] A, int s, int e) {  
    if (s >= e) {  
        return;  
    }  
    int pivotIdx = partition(A, s, e);  
    quicksort(A, s, pivotIdx - 1);  
    quicksort(A, pivotIdx + 1, e);  
}
```

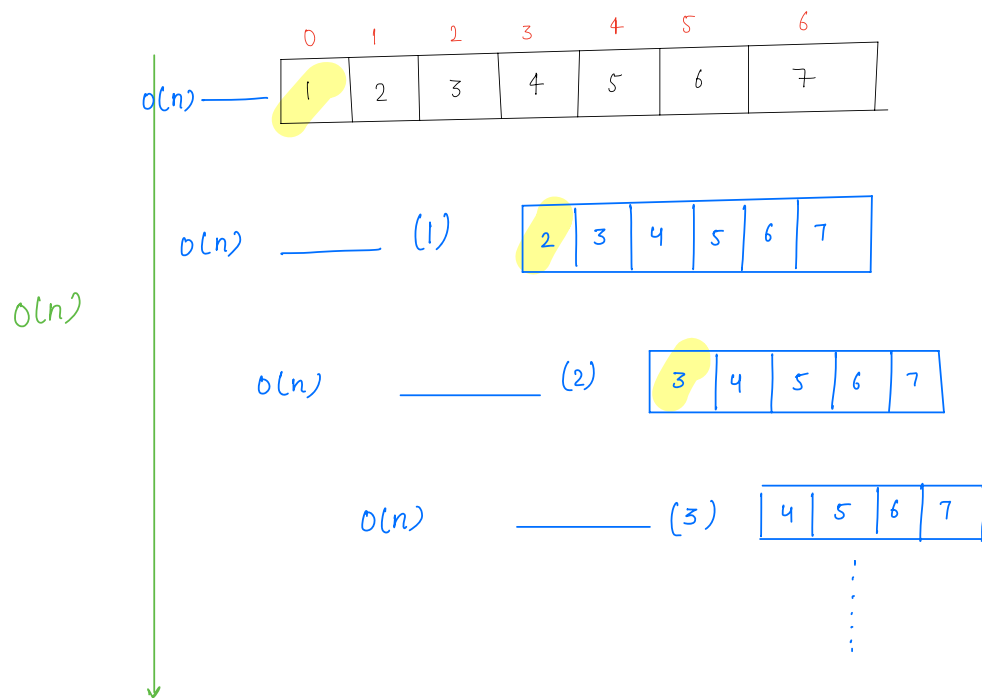
T.C and S.C analysis

Best case:



TC: $O(n \log n)$
SC: $O(\log n)$

Worst case time complexity



TC: $O(n^2)$ SC: $O(n)$

↓

sorted array.

Randomised Quick sort

→ Randomised quick sort helps us to get away with worst case T.C.

Why picking random el. helps?

→ Randomised Q.S helps us to get away with worst case T.C.

→ The odds/prob of choosing a random el as min is very
very very ^{ability} low.

Example:

n el, prob that a random el is min $\Rightarrow 1/n$

prob. that again next time, random el is min $\Rightarrow 1/n-1$

Then $\rightarrow \frac{1}{n-2}$

Then $\rightarrow \frac{1}{n-3}$

⋮

⋮

$\frac{1}{n} * \frac{1}{n-1} * \frac{1}{n-2} * \frac{1}{n-3} \dots$

$$\boxed{\frac{1}{n!}}$$

\rightarrow very very very $\dots \infty$ small.

Break: 8:23 - 8:33

Ques: Given $A[n]$, sort the array in ascending order on the basis of count of factors. If count of factors are equal, then sort the elements on basis of their magnitude.

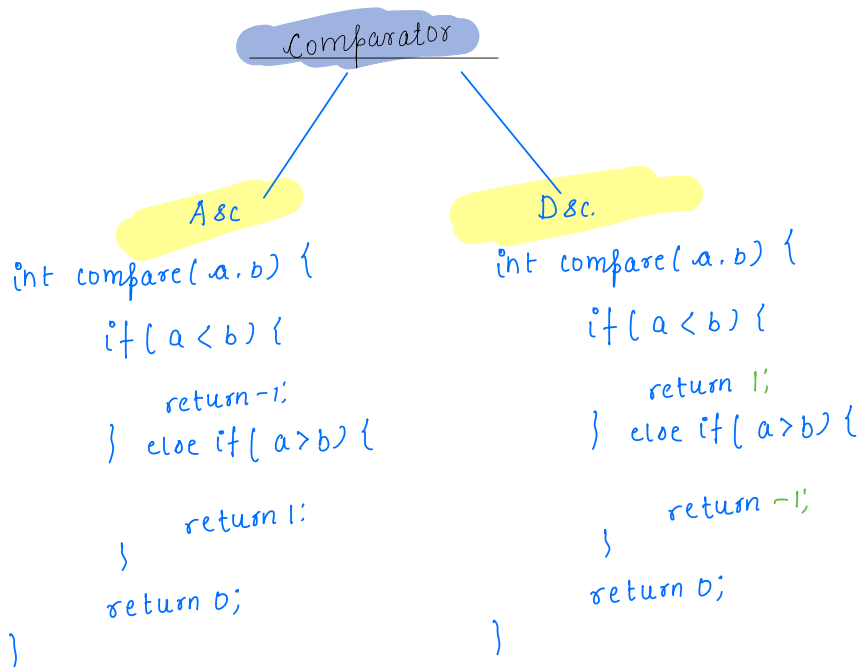
ip:-

	0	1	2	3	4
	9	3	10	6	4
factors	3	2	4	4	3

9: 1, 3, 9
 3: 1, 3
 10: 1, 2, 5, 10
 6: 1, 2, 3, 6
 4: 1, 2, 4

op:-

3	4	9	6	10
---	---	---	---	----



Code:

```
class FactorComparator implements Comparator<Integer> {
```

```
    @Override
```

```
    public int compare(int a, int b) {
```

```
        int fa = countFactors(a);
```

```
        int fb = countFactors(b);
```

```
        if (fa < fb) {
```

```
            return -1;
```

```
        } else if (fa > fb) {
```

```
            return 1;
```

```
        } else {
```

```
            if (a < b) {
```

```
                return -1;
```

```
            } else if (a > b) {
```

```
                return 1;
```

```
            }
```

```
            return 0;
```

```
        }
```

```
    }
```

```
}
```

```
void sortFactors(List<Integer> arr) {
```

```
    Collections.sort(arr, new FactorComparator());
```

```
}
```

Q Given `points[]`, where `points[i] = [xi, yi]` represents a point on the x-y plane and integer `B`, return `B` closest points to origin (0,0).

The distance b/w 2 points on x-y plane is the euclidean distance $\sqrt{(x_1 - x_2)^2 + (y_1 - y_2)^2}$.

You may return answer in any order.

Example 1:

`points[] = { [1, 3], [-2, 2] }`, `B = 1`

Distance b/w (1,3) and origin $\Rightarrow \sqrt{(1-0)^2 + (3-0)^2} = \sqrt{10}$

Distance b/w (-2,2) and origin $\Rightarrow \sqrt{(-2-0)^2 + (2-0)^2} = \sqrt{8}$

Ans = (-2,2).

Example 2: `points = { (3,3) (5,-1) (-2,4) }` `B = 2`.

Distance b/w (3,3) and origin $\Rightarrow \sqrt{(3-0)^2 + (3-0)^2} = \sqrt{18}$

Distance b/w (5,-1) and origin $\Rightarrow \sqrt{(5-0)^2 + (-1-0)^2} = \sqrt{26}$

Distance b/w (-2,4) and origin $\Rightarrow \sqrt{20}$

Ans = (3,3) and (-2,4)

Approach

points = { (3, 3) (5, -1) (-2, 4) } $\beta = 2$
↓ sort on euclidian distance
{ (3, 3) (-2, 4) (5, -1) }
Ans

```
int compare(Points a, Points b) {  
    double d1 = Math.sqrt(a.x * a.x + a.y * a.y);  
    double d2 = Math.sqrt(b.x * b.x + b.y * b.y);  
    if (d1 < d2) {  
        return -1;  
    } else if (d1 > d2) {  
        return 1;  
    }  
    return 0;  
}
```

Qn. Given a list of non-negative integer nums, arrange them such that they form largest number and return it. Since the result may be very large, so you need to return a string instead of an integer.

10	2
----	---

 \Rightarrow "210"

3	30	34	5	9
---	----	----	---	---

 \Rightarrow $\begin{matrix} 3430953 \\ 9534330 \end{matrix}$ ✓

Should we sort the numbers in descending order and append them?

3	30	34	5	9
---	----	----	---	---

↓ desc
34, 30, 9, 5, 4 \Rightarrow 3430954 (wrong)

34 < 30 \Rightarrow false

9 > 34 \Rightarrow false

"Ayush" > "Ayush" \Rightarrow false

"9" > "34" \Rightarrow true.

"34" > "30" \Rightarrow true.
↑↑ ↑↑

3 < 30 \Rightarrow false

"3" > "30" \Rightarrow

3	30	34	5	9
---	----	----	---	---

9534 3 30

9534 30 3

Approach

```
int compare(int a, int b) {  
    if (a < b) {  
        return -1;  
    } else if (a > b) {  
        return 1;  
    }  
    return 0;  
}
```

$a = '3'$ (string)
 $b = '30'$ (string)

$a+b$
330
✓
a

$b+a$
303
b

string a = "he"
string b = "llo"
print(a+b) \Rightarrow hello

String ab = String.valueOf(a) + String.valueOf(b);

String ba = String.valueOf(b) + String.valueOf(a);

int cmp = ab.compareTo(ba);

if (cmp < 0) {

// $ab < ba$

return 1;

}

else if (cmp > 0) {

// $ab > ba$

return -1;

}

return 0;

}

"Ay".compareTo("By")

↳ negative value

"Ay".compareTo("Ax")

↳ +ve value

"Ay".compareTo("Ay")

0

Thankyou 😊