

Lecture ÷ Hashing-2

Agenda

Pair sum = K

Count no. of pairs with sum = K

Subarray with sum = K

Distinct elements in every window.

Qn-1 Given $arr[n]$ and k . Check if there exists a pair (i, j) such that $arr[i] + arr[j] = k$ and $i \neq j$

8	9	1	-2	4	5	11	-6	4
---	---	---	----	---	---	----	----	---

$k = 6$ true $[1 + 5]$

$k = 22$ false

$k = 8$ true. $[4 + 4]$

3	5	1	2	1	2
---	---	---	---	---	---

$k = 7$ true $[5 + 2]$

$k = 10$ false.

Brute force:

TC: $O(n^2)$

SC: $O(1)$

Approach 1

0	1	2	3	4	5	6	7	8
8	9	2	-2	4	5	11	-6	4

target 1 $9-9=0$ $9-2=7$

$-2 + x = 9$
 $x = 11$ ✓

true.

hashset

8	9	2	-2
4	5	11	-6
4			

0	1	2	3	4	5	6	7	8
8	9	2	-2	4	5	11	-6	4

target -4 -5 $4-2=2$ ✓

true.
[Wrong]

hashset

8	9	2	-2
4	5	11	-6
4			

Approach 2

$k = 9$

0	1	2	3	4	5	6	7	8
8	9	2	-2	4	5	11	-6	4

target

↓ ↓ ↓ ↓ ↑ ↑

1 0 7 11 5 4 ✓

true.

hashset

8	9	2	-2
4			

$k = 4$

0	1	2	3	4	5	6	7	8
8	9	2	-2	4	5	11	-6	4

target

↑ ↑ ↑ ↑ ...

-4 -5 2

false

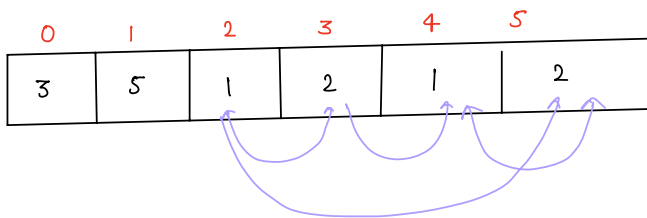
set

8	9	2
---	---	---

Code

```
boolean targetSum(int[] arr, int k) {  
    int n = arr.length;  
    Set<Integer> set;  
    for(i=0; i<n; i++) {  
        int target = k - arr[i];  
        if (set.contains(target)) {  
            return true;  
        }  
        set.add(arr[i]);  
    }  
    return false;  
}
```

TC: $O(n)$
SC: $O(n)$



$k=3$

ans = 4

Qu Count no. of pairs with $\text{sum} = k$.

Given $\text{arr}[n]$, count no. of pairs such that

$\text{arr}[i] + \text{arr}[j] = k$ and $i \neq j$

0	1	2	3	4	5	6	7
2	5	2	5	8	5	2	8

, $k=10$

pairs:- $[0, 4]$

$[0, 7]$

$[1, 5]$

$[1, 3]$

$[2, 4]$

$[2, 7]$

$[3, 5]$

$[4, 6]$

$[6, 7]$

ans = 9.

Approach

0	1	2	3	4	5	6	7
2	5	2	5	8	5	2	8

, k=10

target:

\uparrow 8
 \uparrow $\frac{10-5}{5}$
 \uparrow $\frac{10-2}{8}$
 \uparrow $\frac{10-5}{5}$
 \uparrow $\frac{10-8}{2}$
 \uparrow $\frac{10-5}{5}$
 \uparrow $\frac{10-2}{8}$
 \uparrow $\frac{10-8}{2}$

count:

+1 +2 +2 +1 +3

ans: 9

map

2 — ~~1~~ ~~2~~ 3
 5 — ~~1~~ ~~2~~ 3
 8 — ~~1~~ 2

Code

```
int countTargetSum(int[] arr, int k) {  
    n = arr.length;  
    Map<Integer, Integer> map;  
        array frequency.  
        elements  
    int cnt = 0;  
    for(i=0; i<n; i++) {  
        int target = k - arr[i];  
        if(map.containsKey(target)) {  
            cnt = cnt + map.get(target);  
        }  
        if(map.containsKey(arr[i])) {  
            int freq = map.get(arr[i]);  
            freq += 1;  
            map.put(arr[i], freq);  
        }  
        else {  
            map.put(arr[i], 1);  
        }  
    }  
    return cnt;  
}
```

TC: $O(n)$

SC: $O(n)$

Ques Given $arr[n]$, Check if there exists a subarray with $sum = k$

0	1	2	3	4	5	6	7	8
2	3	9	-4	1	5	6	2	5

$k = 11$ { 2, 3, 9, -4, 1 } { 5, 6 } true

$k = 10$ { 2, 3, 9, -4 } true.

$k = 15$ { -4, 1, 5, 6, 2, 5 } true

0	1	2	3	4
5	10	20	100	105

$k = 110$ false.

Brute force

Go to all subarrays — $O(n^2)$
find sum — $O(n)$
compare.

TC: $O(n^3)$

SC: $O(1)$

Approach

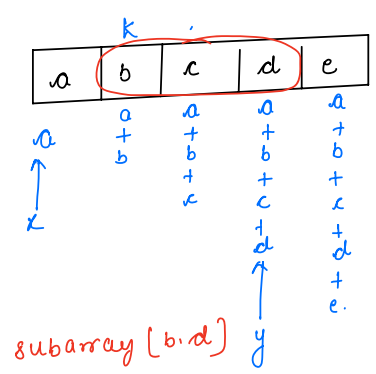
	0	1	2	3	4	5	6	7	8
	2	3	9	-4	1	5	6	2	5
pf[] =	↓	↓	↓	↓	↓	↓			
	2	5	14	10	11	16			
target:	2-11	5-11	14-11						
	= -9	= -6	= 3						
				↓	↓	↓			
				10-11	11-11	16-11			
				= -1	= 0	= 5			

subarray exists

set:

2	5	14
10	11	

k = 11



subarray [b,d]

$y - x = k$

$y - k = x$ search in set

Edge case

	0	1	2	3	4
	2	3	9	-4	1
pf[]	2	↓ 5	↓ 14	↓ 10	↓ 11
target	2-10 =-8	5-10 =-5	14-10 =4	10-10 =0	11-10 =1

k=10

false, (true),

set

2	5	14
10	11	

Resolve:- Add 0 at the start in set

Code

```
boolean targetSubArray (arr, k) {  
    n = arr.length;  
    set<Integer> set;  
    set.add(0);  
    int sum = 0;  
    for(int i=0; i<n; i++){  
        sum += arr[i];  
        int target = sum-k;  
        if(set.contains(sum-k)){  
            return true;  
        }  
        set.add(sum);  
    }  
    return false;  
}
```

TC: $O(n)$
SC: $O(n)$

Break: 845 - 8:55

Sliding window - Map

Sliding window suggests to insert all elements of first window in the set. Now slide the window, throw the previous element out and insert new adjacent element in the set.

0	1	2	3	4	5	6
1	2	10	6	4	3	2

size of window = 4

Window	Indices	All elements of window

Q. Given arr[n] and number k. find count of distinct elements in every window of size k in array.

0	1	2	3	4	5	6
1	2	1	3	4	2	3

k=4

Window	Indices	All elements of window	distinct
first	[0-3]	1, 2, 1, 3	3
2nd	[1-4]	2, 1, 3, 4	4
3rd	[2-5]	1, 3, 4, 2	4
4th	[3-6]	3, 4, 2, 3	3

0	1	2	3	4	5	6
1	2	1	3	4	2	3

k=4

1st window: $\langle 1, 2, \cancel{1}, 3 \rangle$ ans = 3

2nd window: $\langle 2, 3, 4 \rangle$ ans = 3 ~~Wrong~~ wrong.

Approach

0	1	2	3	4	5	6	
1	2	1	3	4	2	3	k=4

Window	Indices	Map	distinct elements
1st	[0-3] {1, 2, 1, 3}	1: 2 2: 1 3: 1	3 [map.size()]
2nd	[1-4] {2, 1, 3, 4}	1: 2 1 2: 1 3: 1 4: 1	4 [map.size()]
3rd	{2-5} {1, 3, 4, 2}	1: 1 2: 1 3: 1 4: 1 2: 1	4 [map.size()]
4th	[3, 6] {3, 4, 2, 3} remove=1 add=3	1: 1 0 2: 1 3: 1 2 4: 1 2: 1	3 [map.size()]

code: h/w

Thankyou 😊

8th Dec:

