

xv6: Adding a system call

Background

RISC-V CPU modes

User mode

- Least privilege
- User processes run in this mode
- Uses paging with pagetable corresponding to running user process

Supervisor mode (or Kernel mode)

- Medium privilege
- Kernel (including kernel modules and device drivers), hypervisor run in this mode
- Uses paging with page table reserved for kernel

Machine mode

- Highest privilege
- Bootloader, firmware run in this mode
- Does not use paging

Life cycle of a system call: **sleep**

Suppose the following code is executed by a user program `test_sleep.c`:

```
#include "kernel/types.h"
#include "kernel/stat.h"
#include "user/user.h"

int main()
{
    int status = sleep(100);

    if (status == -1)
        printf("Error\n");
    else
        printf("Ok\n");

    return 0;
}
```

sleep is a system call. It will cause the following steps.

1. User function calls `sleep(100)`
 - It means sleep for 100 timer interrupts. Timer interrupts have an interval of 0.1 second (interval is defined in `kernel/start.c/timerinit`). So, this sleep takes about 10 seconds.
 - **sleep** system call takes an `int` as argument and returns an `int`, which is found from its declaration in `user/user.h`, `int sleep(int)`
2. The system call number of sleep is loaded into the register **a7** and a trap is generated.
 - The definition of the function `int sleep(int)` is provided in `user/usys.S`
 - **sleep** system call has system call number 13, which is defined as `SYS_sleep` in `kernel/syscall.h`
3. The trap generated in step 2, puts the processor into supervisor mode and calls the function `kernel/trampoline.S/uservec`. This function does some state management (saving state of registers and swapping pagetable) and finally jumps to `kernel/trap.c/usertrap`.
 - This function does not *call* `usertrap`. So, we will never get back to this function again while handling our `sleep(100)`.
4. The function `kernel/trap.c/usertrap` handles the trap based on its cause. As, **sleep** is a system call, it calls `kernel/syscall.c/syscall`.
 - The cause of a trap is given by `r_scause()` (supervisor trap cause). If it is 8, then the cause is system call. We will see other values in later assignments.

5. The function `kernel/syscall.c/syscall` is the single point of contact for handling system calls. It retrieves the system call number (13 in this case) from the saved (in step 3) value of register **a7**. Then it calls the handler function corresponding to the system call number. All system call handlers are saved in an array of function pointers, `kernel/syscall.c/syscalls`. In this array, the system call handler for **sleep** is `sys_sleep`.
6. The function `sys_sleep` is defined in `kernel/sysproc.c`. In this function, first the argument to **sleep** system call is retrieved. Then it does some work that actually makes the calling process sleep for the given amount of time. It returns -1, if there is any error, otherwise it returns 0.
7. Now, we are again in the `kernel/syscall.c/syscall`. The returned value of the system call handler (in this case, `sys_sleep`) is saved in the register **a0**.
8. Then it returns to `kernel/trap.c/usertrap` and calls `kernel/trap.c/usertrapret`.
9. In `kernel/trap.c/usertrapret`, it does some state management and calls `kernel/trampoline.S/userret`.
10. `kernel/trampoline.S/userret` does some state management (restoring state of registers and swapping pagetable) and returns to user mode.

System call related codes explanation:

`user/user.h`

- Contains declaration of system call related functions

`kernel/syscall.h`

- Contains definition of system call number for each system call

`user/usys.S`

- Generated by `user/usys.pl`
- Contains the definition of system call related functions
- Generates a trap with corresponding system call number that puts the CPU into supervisor mode

`kernel/trampoline.S`

`kernel/trampoline.S/uservec`

- Saves the state of the registers
- Switches the pagetable of the user process with the pagetable for kernel
- Calls `kernel/trap.c/usertrap`

`kernel/trampoline.S/userret`

- Restores the state of the registers
- Switches the pagetable for kernel with the pagetable of the user process
- Switches mode from supervisor to user

`kernel/trap.c`

`kernel/trap.c/usertrap`

- Handles an interrupt, exception, or system call from user space

- Calls `kernel/syscall.c/syscall` if the trap is generated by system call (supervisor trap cause = 8)
- Calls `kernel/trap.c/usertrapret` after the trap is handled

`kernel/trap.c/usertrapret`

- Returns to user space
- Calls `kernel/trampoline.S/uservec`

`kernel/syscall.c`

- Contains declarations of the system call handlers

`kernel/syscall.c/syscalls`

- An array of function pointers that points to the actual code for handling each system call
- All the system call handlers must have the same type. For xv6, the type is `uint64(void)`. So, if a system call has arguments, it retrieves them from saved registers using `kernel/syscall.c/argraw`. The returned value must be a `uint64`, which means a system call can either return an integer or a pointer.

`kernel/syscall.c/syscall`

- The main point of contact for handling system calls
- Calls the corresponding handler for the system call that generated the trap. The system call that generated the trap is found in the system call number which was provided in `user/usys.S`.
- Stores the returned value from the system call handler

`kernel/syscall.c/argraw`

- Retrieves the arguments to the system call as `uint64`
- Can retrieve at most 6 arguments, i.e., xv6 can only handle system call with at most 6 arguments

`kernel/syscall.c/argint`

- Retrieves the arguments to the system call and saves it at another variable of `int` type.

`kernel/syscall.c/argaddr`

- Retrieves the arguments to the system call and saves it at another variable which will be treated as an address.

`kernel/sysproc.c`

- Contains definition of the system call handlers related to process management

`kernel/sysfile.c`

- Contains definition of the system call handlers related to file management

Add a new system call

We will add a new system call to get an integer user id. We will call this system call **getuid**.

1. First we need to declare its signature in `user/user.h`.

```
// system calls
int fork(void);
...
int uptime(void);
int getuid(void); // add this line
```

2. Then we need to define the system call number for **getuid** in `kernel/syscall.h`.

```
// System call numbers
#define SYS_fork    1
...
#define SYS_close  21
#define SYS_getuid 22 // add this line
```

3. We need to have the definition of `int getuid(void)` in `user/usys.S`. This code is auto generated by `user/usys.pl`. So, we add an entry for **getuid** in `user/usys.pl`.

```
#!/usr/bin/perl -w
...
entry("uptime");
entry("getuid"); # add this line
```

4. We need to add a system call handler for **getuid**. For that we will add a function named `sys_getuid`. We need to add a record in the array `kernel/syscall.c/syscalls`. So, we make the following changes in `kernel/syscall.c`.

```
// Prototypes for the functions that handle system calls.
extern uint64 sys_fork(void);
...
extern uint64 sys_close(void);
extern uint64 sys_getuid(void); // add this line

// An array mapping syscall numbers from syscall.h
// to the function that handles the system call.
static uint64 (*syscalls[])(void) = {
[SYS_fork]    sys_fork,
...
[SYS_close]   sys_close,
[SYS_getuid]  sys_getuid, // add this line
};
```

5. We need to add the definition of the function `sys_getuid`. We can define it in both `kernel/sysproc.c` and `kernel/sysfile.c`. We will arbitrarily select `kernel/sysproc.c` this time. Add the following codes at the end of `kernel/sysproc.c`.

```
// return user id
uint64
sys_getuid(void)
{
    return getuid();
}
```

6. The function `sys_getuid` calls another function `getuid` that will be defined in `kernel/proc.c` (please note that this `getuid` function is different from that declared in `user/user.h`). Add the following codes at the end of `kernel/proc.c`. *[This step would not be required if `sys_getuid` did not call a newly defined function.]*

```
// define a global variable for user id
int uid = 123;

// getter for user id
int
getuid(void)
{
    return uid;
}
```

7. Now, to make `getuid` accessible from `kernel/sysproc.c`, add a line in `kernel/defs.h`. *[This step would not be required if `sys_getuid` did not call a newly defined function.]*

```
// proc.c
int          cpuid(void);
...
void         procdump(void);
int          getuid(); // add this line
```

8. Now, create the following user program `printuid.c` following the instructions in [xv6: Adding a user program](#) and run.

```
#include "kernel/types.h"
#include "kernel/stat.h"
#include "user/user.h"

int main()
{
    int uid = getuid();

    printf("%d\n", uid);

    return 0;
}
```

9. The output should look like this:

```
$ printuid
123
$
```

Practice

Add a system call to set the user id. Also, create a user program to test it.