

# Question 1: Theoretical Questions [12 points]

Are the following typing statements true or false? Explain why. [4x3 points]

- a.  $\{f : [T_2 \rightarrow T_3], g : [T_1 \rightarrow T_2], a : \text{Number}\} \vdash (f(g a)) : T_3$

The statement is True.

Explanation:

- The list of node in the AST are:

$$f : [T_2 \rightarrow T_3] \quad (\text{By } T_{\text{Env}})$$

$$g : [T_1 \rightarrow T_2] \quad (\text{By } T_{\text{Env}})$$

$$a : T_{\text{num}} \quad (\text{introduce new } T_{\text{var}})$$

$$(g a) : T_2 \quad (\text{By AppExp typing rule and equation } \{ T_{\text{num}} = T_1 \})$$

Since 'g' of type  $[T_1 \rightarrow T_2]$  and  $a : T_1$  (since  $g$  expects  $T_1$  as the first argument), therefore  $(g a)$  of type  $T_2$

$$(f(g a)) : T_3 \quad (\text{By AppExp typing rule})$$

Since  $(g a)$  of type  $T_2$  and 'f' maps from  $T_2$  to  $T_3$ , the result  $(f(g a))$  is of type  $T_3$

thus, the statement is true.

- b.  $\{f : [T_1 \rightarrow [T_2 \rightarrow \text{Boolean}]], x : T_1, y : T_2\} \vdash (f x y) : \text{Boolean}$

the statement is false.

Explanation:

- The list of node in the AST are:

$$f : [T_1 \rightarrow [T_2 \rightarrow \text{Boolean}]] \quad (\text{by } T_{\text{Env}})$$

$$x : T_1 \quad (\text{by } T_{\text{Env}})$$

$$y : T_2 \quad (\text{by } T_{\text{Env}})$$

$(f x y) : \text{Type mismatch (By AppExp typing rule)}$

This statement is false because  $(f x y)$  suggest applying 'f' to x and y simultaneously,

which would require 'f' to have type  $[T_1 \times T_2 \rightarrow \text{Boolean}]$ , to match the argument structure.

However, 'f' type is  $[T_1 \rightarrow [T_2 \rightarrow \text{Boolean}]]$  indicates that 'f' first takes an argument

of type ' $T_1$ ' and returns a function  $[T_2 \rightarrow \text{Boolean}]$ , not  $[T_1 \times T_2 \rightarrow \text{Boolean}]$ , therefore, the

statement is false due to the mismatch in argument structure.

c.  $\{f : [T_1 \times T_2 \rightarrow T_3], y : T_2\} \vdash (\lambda(x)(f x y)) : [T_1 \rightarrow T_3]$

The statement is True.

Explanation:

- The list of node in the AST are:

$f : [T_1 \times T_2 \rightarrow T_3]$  (By TEnv)

$y : T_2$  (By TEnv)

$x : T_1$  (Introduce new TVar)

$(f x y) : T_3$  (By AppExp typing rule and equation  $T_x = T_1$ )

Since 'f' is of type  $[T_1 \times T_2 \rightarrow T_3]$ , with 'x' of type  $T_1$  and 'y' of type  $T_2$  the expression  $(f x y)$  has type  $T_3$ .

$\lambda(x)(f x y) : [T_1 \rightarrow T_3]$  (By procExp typing rule)

the lambda takes an argument 'x' of type  $T_1$  and returns a result of type  $T_3$ , therefore,  $\lambda(x)(f x y)$  has type  $[T_1 \rightarrow T_3]$

(since 'f' expects a ' $T_1$ ' as the first argument)

Thus, The statement is true.

d.  $\{f : [T_2 \rightarrow T_1], x : T_1, y : T_3\} \vdash (f x) : T_1$

The statement is True.

Explanation:

- The list of node in the AST are:

$f : [T_2 \rightarrow T_1]$  (By TEnv)

$x : T_1$  (by TEnv)

$y : T_3$  (by TEnv)

$(f x) : T_1$

To type  $(f x)$ ,  $x$  should ideally match the expected type  $T_2$  of  $f$ , By substituting  $\{T_2 = T_1\}$ ,  $f$  type becomes  $[T_1 \rightarrow T_1]$ , Which aligns with  $x$ 's type,  $(f x)$  correctly returns type  $T_1$ , Therefore, the statement is true.

**2.1** We defined the type constructors `diff`, `inter`, and `union` according to set-theoretic principles. For each of the following TExps, write a simplified TExp denoting the same set of values. [3 pts]

- a.  $(\text{inter} \text{ number boolean}) \rightarrow \text{never}$
- b.  $(\text{inter} \text{ any string}) \rightarrow \text{String}$
- c.  $(\text{union} \text{ any never}) \rightarrow \text{any}$
- d.  $(\text{diff} (\text{union} \text{ number string}) \text{ string}) \rightarrow \text{number}$
- e.  $(\text{diff} \text{ string} (\text{union} \text{ number string})) \rightarrow \text{never}$
- f.  $(\text{inter} (\text{union} \text{ boolean number}) (\text{union} \text{ boolean} (\text{diff} \text{ string never}))) \rightarrow \text{boolean}$

**2.2** Complete the following code snippet (i.e. offer replacements to the **[bracketed gaps]**) such that it will pass type checking in L52 but not in L51. [8 pts]

```
; L52 return type is? boolean
;; L51 return type boolean
(define (isBoolean : (any -> is? boolean))
  (lambda ((x : any)) : is? boolean
    (boolean? x)))

;; Function to complete
(define (good_in_L52 : ((union number boolean) -> [a]))
  (lambda ((z : (union number boolean))) : [b]
    [c]
  ))
```

[a] → is? Boolean

[b] → is? Boolean

[c] → (isBoolean z)

**2.3** Complete the return type for f in L52. Explain. [7 pts]

Answer: (Union String boolean)

Explanation:

- the function 'f' takes an argument 'x' that can be either a number or a boolean, and returns either a 'string' "positive" or "negative" if x is a number, or returns 'x' if it's a boolean.
- the return type of 'f' is (Union String boolean), indicating it can return either a 'string' or a 'Boolean' based on the input type.