Simple CPU

Description

The purpose of this learning activity is to test the final CPU after combining all the other components made before. More specifically, I was tasked to create a program to test all the functions of the CPU as a whole.

Schematic

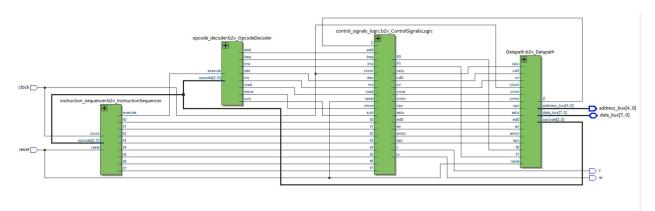


Figure 1: Simple CPU Schematic

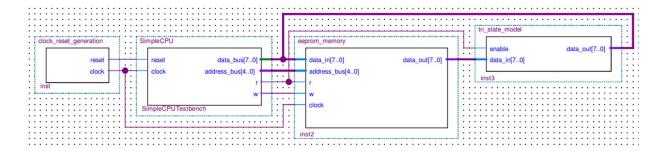


Figure 2: Simple Computer System

The schematic in Figure 1 illustrates the schematic for the simple CPU only using the components created from previous learning activities. The schematic in Figure 2 shows the entire computer system with the clock reset generation, CPU, memory, and tri-state model.

Program Explanation

The program used to verify the functionality of the simple CPU follows a simple logic.

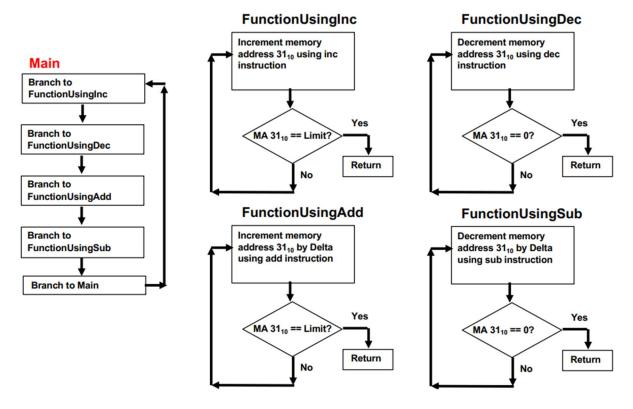


Figure 3: Flowchart of Program

Memory Address	Memory Address	Data	Memory File		Address					
(hex)	(decimal)	(Hex)	Contents	Instruction	(Decimal)	Description				
00	0	C5	00/C5;	bra	5	// (Main) unco	nditionally b	ranches to Fun	ctionUsingInc in	Memory Address = 5
01	1	CA	01/CA;	bra	10	//unconditionally branches to FunctionUsingDec in Memory Address = 10				
02	2	CD	02/CD;	bra	13	//unconditionally branches to FunctionUsingAdd in Memory Address = 13				
03	3	D4	03/D4;	bra	20	//unconditionally branches to FunctionUsingSub; MA = 19				
04	4	CO	04/C0;	bra	0	//unconditionally branches to Main; MA = 0				
05	5	9F	05/9F;	inc	31	//(FunctionUsingInc) Beginning of FunctionUsingInc; increment MA = 31				
06	6	1F	06/1F;	load	31	//load value in	MA = 31 into	D0		
07	7	7C	07/7C;	sub	28	//Subtract valu	ue in D0 by li	mit		
08	8	E1	08/E1;	beq	1	//invoke function again; runs only when MA = 31 is not equal to 0				
09	9	C5	09/C5;	bra	5	//unconditionally branch to the beginning of the function				
0A	10	BF	OA/BF;	dec	31	//decrement u	ising MA = 31	; (FunctionUsin	ngDec)	
0B	11	E2	0B/E2;	beq	2	//If MA = 31 is	equal to the	limit; yes => go	to invocation	
0C	12	CA	OC/CA;	bra	10	//invokes func	ction again if	MA = 31 is not	equal to 0	
0D	13	1F	0D/1F;	load	31	//(FunctionUsi	ingAdd) load	s the value in N	MA = 31 into D0	
0E	14	5D	0E/5D;	add	29	//adds delta to	D0			
OF	15	3F	0F/3F;	store	31	//stores the va	alue into MA	= 31		
10	16	1F	10/1F;	load	31	//load value in	MA = 31 into	D0		
11	17	7C	11/7C;	sub	28	//Subtract valu	ue in D0 by lin	nit		
12	18	E3	12/E3;	beq	3	//If MA = 31 is equal to the limit; yes => go to invocation				
13	19	CD	13/CD;	bra	13	//unconditionally branch to the beginning of the function				
14	20	1F	14/1F;	load	31	//(FunctionUsingSub) load MA=31 into D0 register				
15	21	7D	15/7D;	sub	29	//subtract the value by delta				
16	22	3F	16/3F;	store	31	//store the value into MA = 31				
17	23	1F	17/1F;	load	31	//load MA=31	back into D0			
18	24	7E	18/7E;	sub	30	//subtract from	n 0			
19	25	E4	19/E4;	beq	4	//conditional	tatement, if	value in MA=3	1 is 0 then go to i	nvocation
1A	26	D4	1A/D4;	bra	20	//unconditiona	ally branch to	the beginning	g of the function	
1 B	27	00	1B/00;	load	0					
1C	28	0A	1C/0A;	data	10	limit				
1 D	29	02	1D/02;	data	2	delta				
1E	30	00	1E/00;	data	0					
1F	31	00	1F/00;	data	0					

Figure 4: Implementation in Excel

Figure 3 shows a simple flowchart of the logic of the program used to verify the functionality of the CPU and Figure 4 illustrates the actual implementation. In short, the program has four functions and two constants: delta and limit. In the case of my particular implementation, delta and limit were set to 2 and 10 respectively. The program first tests the *FunctionUsingInc* function where it increments the value in memory address 31 until it reaches the limit. When it reaches the limit, it decrements the value until it reaches 0 through *FunctionUsingDec*. Then the program increments the value in memory address 31 by delta using *FunctionUsingAdd* until reaches the limit and then decrements by delta using *FunctionUsingSub* until it reaches the value of 0.

Verification

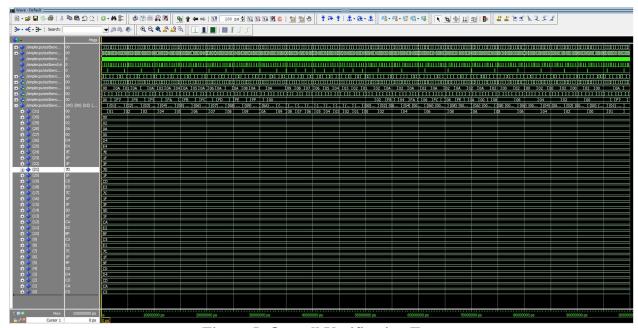


Figure 5: Overall Verification Test

Figure 5 shows the overall verification tests of the functionality of the program and the functionality of the CPU where memory address 31 is the variable that was transformed by the operations the most. To further illustrate the success of the functionality of the program and CPU, the verification section in this report is divided into four additional parts for each function.

<u>FunctionUsingInc</u>

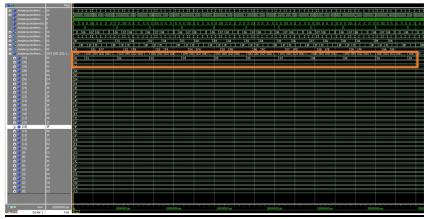


Figure 6: FunctionUsingInc Verification

Figure 6 shows the *FunctionUsingInc* works correctly because the value in memory address 31 increases by one every time. When the value in memory address 31 equals 10, it proceeds to begin decreasing showing the start of the *FunctionUsingDec* function.

<u>FunctionUsingDec</u>

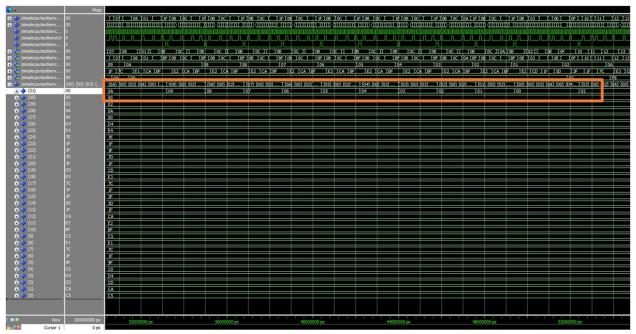


Figure 7: FunctionUsingDec Verification

Figure 7 shows that the *FunctionUsingDec* works correctly because when the value in memory address 31 is 10, it begins decreasing until it reaches 0. Then, the value starts transforming by the *FunctionUsingAdd* function.

FunctionUsingAdd

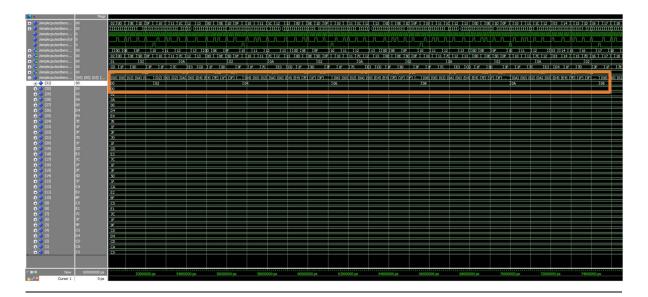


Figure 8: FunctionUsingAdd Verification

Figure 8 shows that the *FunctionUsingAdd* works correctly because when the value in memory address 31 is 0, it begins incrementing by delta, 2. Then, the value starts transforming by the *FunctionUsingSub* function when the value reaches 10.

<u>FunctionUsingSub</u>

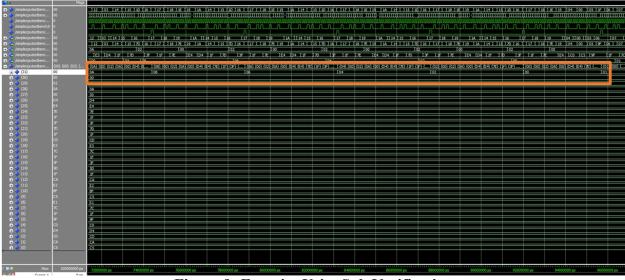


Figure 9: FunctionUsingSub Verification

Figure 8 shows that the *FunctionUsingSub* works correctly because when the value in memory address 31 is 10, it begins decrementing by delta, 2. Then, the value starts transforming by the *FunctionUsingInc* function when the value reaches 0 since the program will run as a continuous loop.