



# Using OpenMP 4.5 Target Offload for Programming Heterogeneous Systems

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# Outline

- Introduction
  - Concepts of Programming Heterogeneous Architectures
- OpenMP Heterogeneity Support Basics
  - Off-loading Work and Data to the GPU
  - Expressing Parallelism and Data Locality
- Using OpenMP 4.5 on Pleiades GPU Nodes
- Learning by Example
  - Laplace Kernel
- More OpenMP 4.5 Constructs and Clauses
- References

# Heterogeneous Systems



- A general purpose processor connected to an accelerator device
- An example is an Intel Xeon processor connected to a Nvidia GPU

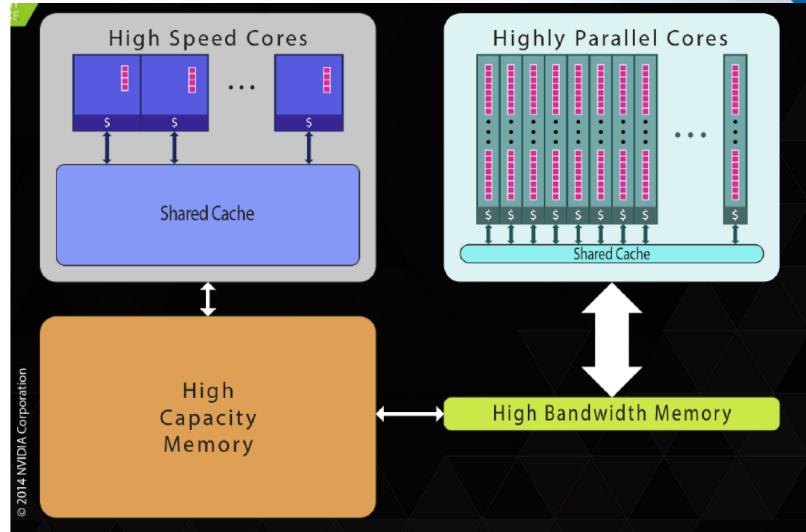


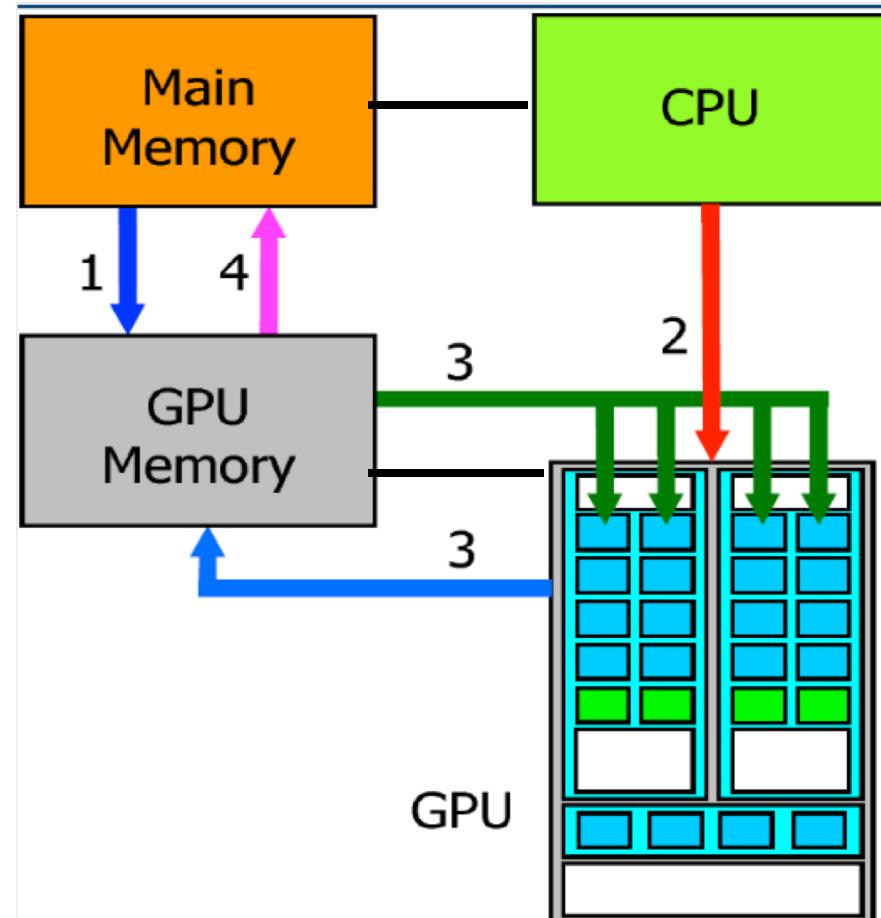
Image courtesy of Nvidia

- |   |  |
|---|--|
| <ul style="list-style-type: none"><li>• CPU<ul style="list-style-type: none"><li>- Fast clock (2.4-2.9 GHz on Pleiades)</li><li>- Multiple cores (16-40 on Pleiades)</li><li>- Complex cores<ul style="list-style-type: none"><li>o Large caches, complex branch prediction, OOO execution, multi-threading</li></ul></li><li>- Parallelism<ul style="list-style-type: none"><li>o Deep pipelines, multiple cores, vector units, SIMD execution</li></ul></li></ul></li></ul> | <ul style="list-style-type: none"><li>• Device<ul style="list-style-type: none"><li>- Slow clock (0.8-1.0 GHz)</li><li>- Thousands of cores<ul style="list-style-type: none"><li>o 2880 SP cores on Pleiades</li></ul></li><li>- Lightweight cores:<ul style="list-style-type: none"><li>o Small caches, little branch prediction, in-order execution, multi-threading</li></ul></li><li>- Parallelism: Theoretically enormous!<ul style="list-style-type: none"><li>o In practice limited; SIMD execution</li></ul></li></ul></li></ul> |
|---|--|

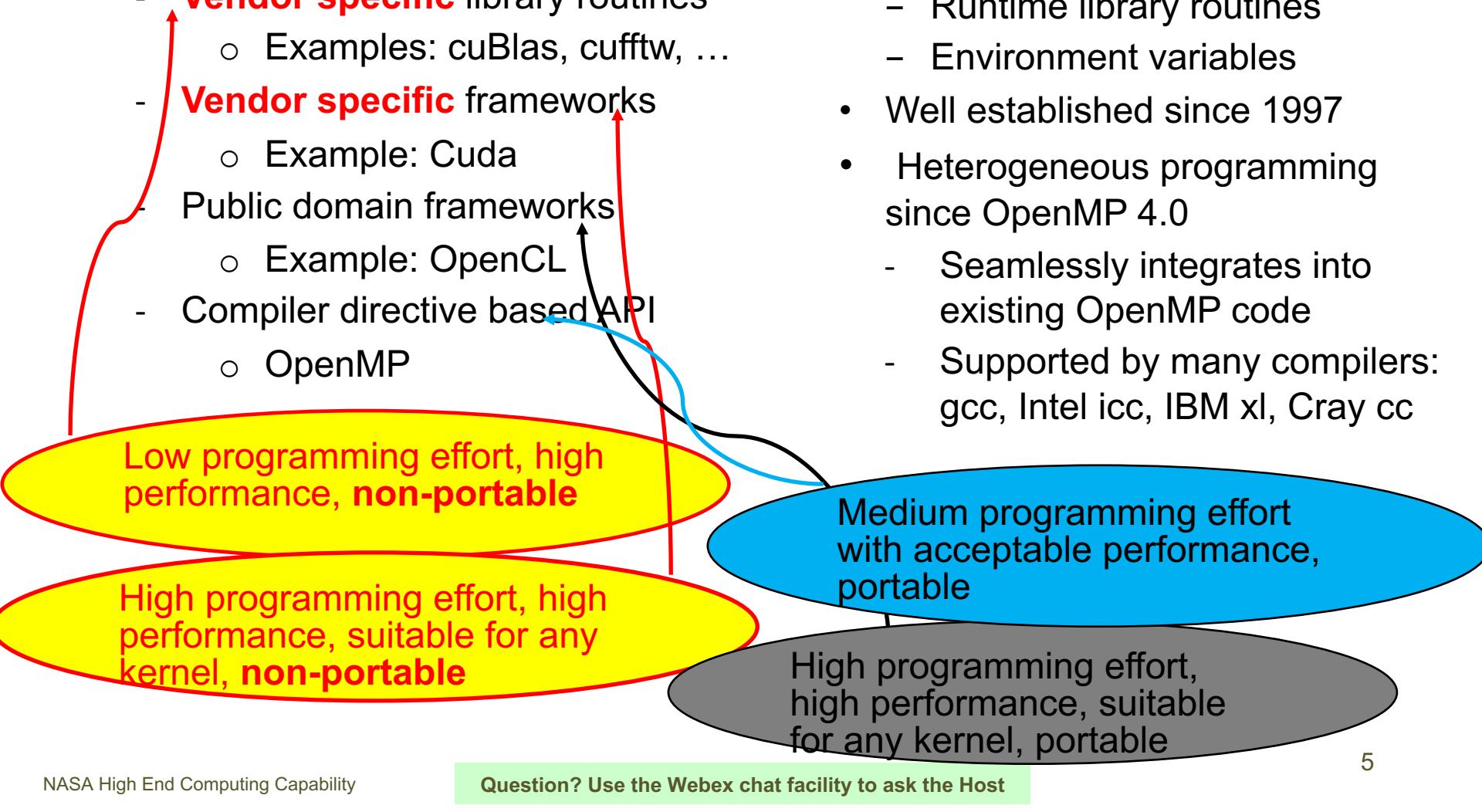
# Programming Heterogeneous Systems

- Necessary steps to be taken
  - Identification and offloading of compute kernels from host to device
  - Expressing parallelism within the kernel
  - Manage data transfer between CPU and Device
- Execution flow
  1. Data copy from main to device memory
  2. CPU initiates kernel for execution on the device
  3. Device executes the kernel using device memory
  4. Data copy from device to main memory

Example:  
Device is a GPU



# Why Use OpenMP?

- Methods to program heterogeneous systems:
    - **Vendor specific** library routines
      - Examples: cuBlas, cufftw, ...
    - **Vendor specific** frameworks
      - Example: Cuda
    - Public domain frameworks
      - Example: OpenCL
    - Compiler directive based API
      - OpenMP
  - OpenMP:
    - Compiler directives
    - Runtime library routines
    - Environment variables
  - Well established since 1997
  - Heterogeneous programming since OpenMP 4.0
  - Seamlessly integrates into existing OpenMP code
  - Supported by many compilers: gcc, Intel icc, IBM xl, Cray cc
- 
- The diagram illustrates the relationship between different programming paradigms and their characteristics. It features four colored ovals (yellow, blue, red, grey) arranged in a cross pattern, each containing a statement. Red arrows point from the first three ovals to the corresponding sections in the list above. A black arrow points from the fourth oval to the 'Compiler directive based API' section.
- Yellow oval (top-left):** Low programming effort, high performance, **non-portable**
  - Blue oval (top-right):** Medium programming effort with acceptable performance, portable
  - Red oval (bottom-left):** High programming effort, high performance, suitable for any kernel, **non-portable**
  - Grey oval (bottom-right):** High programming effort, high performance, suitable for any kernel, portable

# OpenMP Directive Syntax



- Compiler Directive
  - Programmer inserted hint/command for the compiler
- Directive Syntax
  - Fortran
    - Mostly paired with a matching **end** directive surrounding a structured code block

```
!$omp directive [clause [,] [clause] ...]
    code
```

```
!$omp end directive
```

- C
  - No **end** directive needed as the structured block is bracketed

```
#pragma omp directive [clause [,] [clause] ...]
{
    code
}
```

# OpenMP Heterogeneity Basic Support

**#pragma omp target or  
!\$omp target/end target**

- A device data environment is created for the marked region
- The marked code region is mapped to the device and executed.

**#pragma omp target teams or  
!\$omp teams/end teams**

- A **league** of thread teams is created
- The master thread of each team executes the code region

**#pragma omp target teams distribute**

- Worksharing construct: Share work across the teams

**#pragma omp parallel for or !\$omp parallel do**

- Worksharing across threads within a team

We have seen those before....

**#pragma omp simd**

- Worksharing across vector length

**#pragma omp target map(*map-type*: *list*)**

- Map a variable to/from the device data environment



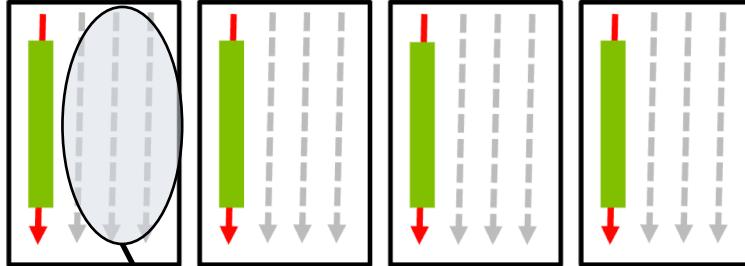
# Stencil Kernel in OpenMP 4.5

```
{
#pragma omp target teams  distribute num_teams(4)
    for( int j = 1; j < n-1; j++ ) {
        for( int i = 1; i < m-1; i++ ) {
            Anew[j][i] = 0.25 * ( A[j][i+1] + A[j][i-1]
                + A[j-1][i] + A[j+1][i]);
            error = fmax( error, fabs(Anew[j][i] - A[j][i]));
        }
    }
}
```

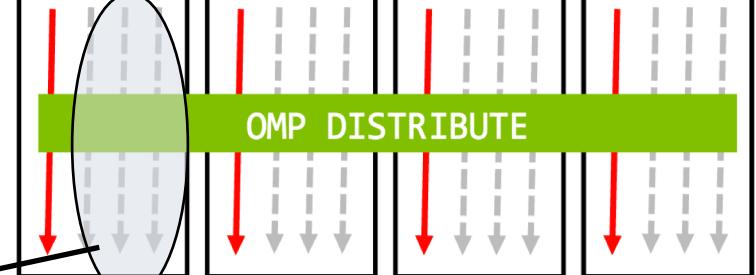
- League of teams is started via **target teams**
- All teams execute the target region

- Work is distributed across the teams via **distribute**

OMP TEAMS



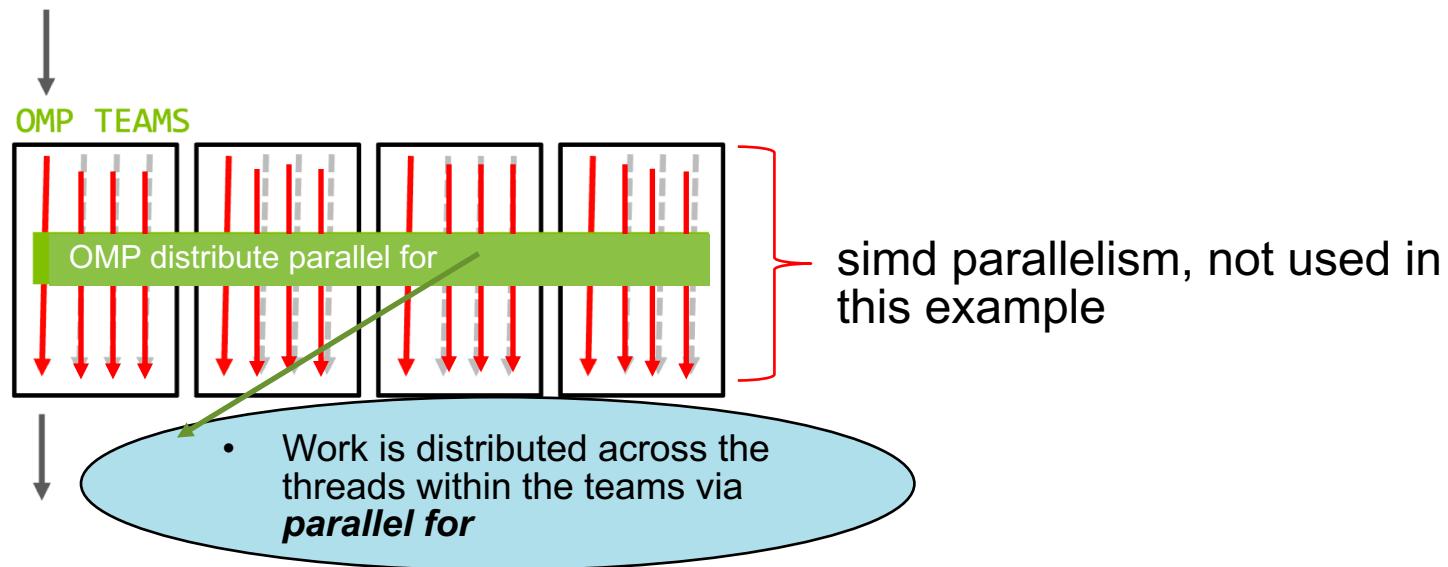
OMP TEAMS



More available hardware threads that are not used

# Stencil Kernel in OpenMP 4.5

```
{
#pragma omp target teams distribute num_teams(4) parallel for thread_limit(4)
for( int j = 1; j < n-1; j++) {
    for( int i = 1; i < m-1; i++ ) {
        Anew[j][i] = 0.25 * ( A[j][i+1] + A[j][i-1]
                               + A[j-1][i] + A[j+1][i]);
        error = fmax( error, fabs(Anew[j][i] - A[j][i]));
    }
}
```



For more details check out the presentation by Jeff Larkin, Nvidia:  
<http://on-demand.gputechconf.com/gtc/2016/presentation/s6510-jeff-larkin-targeting-gpus-openmp.pdf>

Image courtesy of Nvidia

# Structured Data Management



- The device data region:
  - A region of the program within which data is accessible to the device
  - It can be explicitly defined to reduce data copies
  - The **target data** construct is used to mark such regions

**#pragma omp target data map(*map-type*: *list*)**

- Example map types:
  - to (*list*)** 
    - Allocates memory on the device and copies data in when entering the region, the values are not copied back
  - from (*list*)**
    - Allocates memory on the device and copies the data to the host when exiting the region
  - alloc (*list*)**
    - Allocates memory on the device
    - If the data is already present on the device a reference counter is incremented

# Compiling and Running on Pleiades

- Pleiades GPU nodes:
  - 64 Sandy Bridge nodes with one Tesla K40 GPU 
- Compilation:
  - Intel icc/ifort does not generate code for Nvidia GPU execution ☹
  - PGI pgcc/pgf90 does not support the OpenMP target construct ☹
  - **Experimental gcc 8.1 with GPU support is available on Pleiades ☺**

```
module purge
module load cuda
module load /home1/gjost/public/modules/gcc8.1-module
gcc --version
gcc or gfortran -fopenmp -foffload="-lm" test.c or test.f90
```

- Submit to GPU node

```
qsub -l select=1:ncpus=16:model=san_gpu -q k40
```
- Run the executable

```
./a.out
```





# Example: 2D Laplace Solver

```
#pragma omp target data map(to:Anew) map(A)
while ( error > tol && iter < iter_max ) {
    error = 0.0;

#pragma omp target teams distribute parallel for reduction(max:error)
map(error)
    for( int j = 1; j < n-1; j++) {
        for( int i = 1; i < m-1; i++ ) {
            Anew[j][i] = 0.25 * ( A[j][i+1] + A[j][i-1]
                                   + A[j-1][i] + A[j+1][i]);
            error = fmax( error, fabs(Anew[j][i] - A[j][i])); }
    }

#pragma omp target teams distribute parallel for
    for( int j = 1; j < n-1; j++) {
        for( int i = 1; i < m-1; i++ ) {
            A[j][i] = Anew[j][i]; } }

if(iter % 100 == 0) printf("%5d, %0.6f\n", iter, error);
iter++;} // end of while loop
```

total: 167.7766 seconds on Pleiades



# Example: 2D Laplace Solver Separated

```
#pragma omp target data map(to:Anew) map(A)
while ( error > tol && iter < iter_max ) {
    error = 0.0;

#pragma omp target teams distribute reduction(max:error) map(error)
    for( int j = 1; j < n-1; j++){
#pragma omp parallel for reduction(max:error)
        for( int i = 1; i < m-1; i++ ){
            Anew[j][i] = 0.25 * ( A[j][i+1] + A[j][i-1]
                + A[j-1][i] + A[j+1][i]);
            error = fmax( error, fabs(Anew[j][i] - A[j][i])); }
    }

#pragma omp target teams distribute
    for( int j = 1; j < n-1; j++){

#pragma omp parallel for
    for( int i = 1; i < m-1; i++ ){
        A[j][i] = Anew[j][i]; } }

if(iter % 100 == 0) printf("%5d, %0.6f\n", iter, error);
iter++; }
```

total: 155.659801 seconds on Pleiades

- Other possibilities, eg use collapse clause
- Depends on problem sizes and system



# Example: 2D Laplace Solver in Fortran

```
!$omp target data map(to:Anew) map(A)
do while ( error .gt. tol .and. iter .lt. iter_max )
    error=0.0_fp_kind
 !$omp target teams distribute parallel do reduction(max:error) map(error)
    do j=1,m-2
        do i=1,n-2
            Anew(i,j) = 0.25_fp_kind * ( A(i+1,j     ) + A(i-1,j     ) + &
                                         A(i     ,j-1) + A(i     ,j+1) )
            error = max( error, abs(Anew(i,j)-A(i,j)) )
        end do
    end do
 !$omp end target teams distribute parallel do
    if(mod(iter,100).eq.0) write(*,'(i5,f10.6)') iter, error
    iter = iter + 1
 !$omp target teams distribute parallel do
    do j=1,m-2
        do i=1,n-2
            A(i,j) = Anew(i,j)
        end do
    end do
 !$omp end target teams distribute parallel do
end do
 !$omp end target data
```



# Performance Challenges

- Function calls in inner loops:
  - Challenging for the compiler's dependence analysis
    - o Possibly use `!$omp declare target` to compile routine for the device (p. 20)
    - o manually in-line
- IF-Banches:
  - Can get expensive if threads take different execution paths
    - o Move your branches up to a level in your code where all threads go down the same code branch
    - o Avoid branches in inner loops
- Structures with complex and/or dynamic components:
  - Might prevent stride 1 memory access
  - Example: NPB FT uses complex data; performance increase via manual linearization
  - Unstructured data management for structures with dynamic components (next slide)
- Loop strides/memory layout
  - Inner loop should be long and move along the fastest dimension



# Unstructured Data Management: Enter/Exit Data Constructs

- Real life applications, just like real life, are not always a nicely structured sequence of code regions
  - Example: C++ Structures or Fortran user defined types with dynamic arrays
- Unstructured data directives
  - `!$omp target enter data map(map_type:list)`**
    - Allocate memory on the device for the remainder of the program or until explicitly deleted
    - Possible map types are `to`, `alloc`
  - `!$omp target exit data map(map_type:list)`**
    - Deallocate the memory on the device
    - Possible map types are `from`, `release`, or `delete`
- Multiple **`enter/exit data`** constructs, branched across different function calls are allowed



# Unstructured Data Management

```
module matrix_mod
    implicit none
    type matrix
        integer :: num_rows
        integer :: nnz
        integer, pointer :: row_offsets(:)
        integer, pointer :: cols(:)
        real(8), pointer :: coefs(:)
    end type matrix
    public :: allocate_3d
    ...
    subroutine allocate_3d (a, n)
        implicit none
        type(matrix) :: a
        allocate(a%row_offsets(num_rows+1))
        allocate(a%cols(nnz))
        allocate(a%coefs(nnz))
        arrow_offsets => a%row_offsets
        acols => a%cols
        acoefs => a%coefs
        !$omp target enter data map(to:a)
        !$omp target enter data map(to:acols(1:nnz)...
    end subroutine allocate_3d
end module matrix_mod
```

```
subroutine free_3d(a)
    implicit none
    type(matrix) :: a
    ...
    arow_offsets => a%row_offsets
    acols => a%cols
    acoefs => a%coefs
    !$omp exit data map(delete:acols(1:nnz),...)
    !$omp exit data map (delete: a)
    deallocate(arow_offsets)
    deallocate(acols)
    deallocate(acoefs)
end subroutine free_matrix
```

Call `allocate_3d` and `free_3d` as often as you need to, also across subroutine calls

# Routines and global variables

**!\$omp declare target**

- Specifies that variables, functions and subroutines are mapped to a device for the duration of the program

in LBM.cpp

```
#pragma omp target teams distribute parallel do
for (int index1 = 0; index1 < a_numPts; ++index1)
{
    Real feq[NUMV];
    getFEQ(feq, index1);
```

in LBM.H

```
#pragma omp declare target
inline void getFEQ(Real* const __restrict a_feq,
                   const int a_index1)
#pragma omp end declare target
```





# Synchronize host and device buffers

## ! \$omp target update clause

- Makes the variables on the device consistent with the original variables
- Possible motion clauses are **to** or **from**

```
#pragma omp declare target
void work_on_device_work(int* srcA, int* result);
#pragma omp end declare target

#pragma omp target data map(to:a[0:size])map(from:c[0:size])
{
    #pragma omp target
    {
        work_on_device_work(a, c);
    }
    /* Copy modified target data to the host */
    #pragma omp target update from(c[0:size])

    work_on_host(c);
    #pragma omp target
    {
        work_on_device_work(a, c);
    }
}
```

# Device Pointers



- Pointer variable on the host containing device address
- Device memory routines
  - Support of allocation and management of pointers in the data environment of the target device

The diagram illustrates a pointer variable `u1` pointing to a device address. A blue oval labeled "device address" contains the text "device address". Another blue oval labeled "u1 must appear in the clause" contains the text "u1 must appear in the clause". An arrow points from the text "u1 must appear in the clause" to the variable `u1` in the code snippet below.

```
u1 = (double*)omp_target_alloc(n3*n2*n1*sizeof(double),  
omp_get_default_device());  
...  
  
#pragma omp target map(tofrom: ou[0:n3*n2*n1]) is_device_ptr(u1)  
  
#pragma omp teams distribute  
for (i3 = 1; i3 < n3-1; i3++) {  
  
#pragma omp parallel for collapse(2)  
for (i2 = 1; i2 < n2-1; i2++) {  
    for (i1 = 0; i1 < n1; i1++) {  
        I3D(u1, n1, n2, i3, i2, i1) =  
            I3D(ou, n1, n2, i3, i2-1, i1)  
            + I3D(ou, n1, n2, i3, i2+1, i1)  
            + I3D(ou, n1, n2, i3-1, i2, i1)  
            + I3D(ou, n1, n2, i3+1, i2, i1); }}}
```

Note :

```
#define I3D(array,n1,n2,i3,i2,i1) (array[(i3)*n2*n1 + (i2)*n1 + (i1)])
```



# References

Using GPUs codes on Pleiades HECC KB

[https://www.nas.nasa.gov/hecc/support/kb/using-gpu-nodes\\_298.html](https://www.nas.nasa.gov/hecc/support/kb/using-gpu-nodes_298.html)

Jeff Larkin Presentation

(see slide) <https://www.nas.nasa.gov/publications/ams/2015/04-21-15.html>

- OpenMP Home Page and Books

<http://www.openmp.org/>

<http://www.openmp.org/tech/using-openmp-next-step/>

<https://mitpress.mit.edu/books/using-openmp>

# Fortran Target Routines

```

module mod_fib
contains
    subroutine fib(N)
        integer :: N
        !$omp declare target
        ! ...
    end subroutine
    end module
program my_fib
    !$omp target if (N > 1000)
        call fib(N)
    !$omp end target
end program

```

```

program my_fib
    Integer :: N=8
    !$omp declare target(fib)
    !$omp target
        call fib(N)
    !$omp end target
    end program
    subroutine fib(N)
        integer :: N
        !$omp declare target
            print*, "hello from fib"
            !...
    end subroutine

```

in a module

external



# Host/Device Asynchronous Execution

```
#pragma omp target map(a,b,c,d) nowait
{
    {
        #pragma parallel for
        for (i=0; i<N; i++) (
            a[i] = b[i] * c + d;
        )
    }
    F(b); //Execute in parallel with target task

    #pragma omp taskwait //wait for target task to finish
```





# Example: 2D Laplace Solver Collapse

```
#pragma omp target data map(to:Anew) map(A)
while ( error > tol && iter < iter_max ) {
    error = 0.0;

#pragma omp target teams distribute parallel for reduction(max:error)
map(error) collapse (2)
    for ( int j = 1; j < n-1; j++) {
        for( int i = 1; i < m-1; i++ ) {
            Anew[j][i] = 0.25 * ( A[j][i+1] + A[j][i-1]
                                   + A[j-1][i] + A[j+1][i]);
            error = fmax( error, fabs(Anew[j][i] - A[j][i])); }
    }

#pragma omp target teams distribute parallel for collapse(2)
    for( int j = 1; j < n-1; j++) {
        for( int i = 1; i < m-1; i++ ) {
            A[j][i] = Anew[j][i]; } }

if(iter % 100 == 0) printf("%5d, %0.6f\n", iter, error);
iter++; }
```



total: 167.67 seconds on Pleiades



# Get to know your GPU

```
PBS r313i1n2 51> module load cuda
```

```
PBS r313i1n2 52> nvidia-smi
```

```
Mon Mar 18 14:58:08 2019
```

```
+-----+  
| NVIDIA-SMI 384.81      Driver Version: 384.81 |  
+-----+  
| GPU Name Persistence-M| Bus-Id Disp.A | Volatile Uncorr. ECC |  
| Fan Temp Perf Pwr:Usage/Cap| Memory-Usage | GPU-Util Compute M. |  
+-----+-----+-----+-----+-----+-----+  
| 0 Tesla K40m     Off | 00000000:02:00.0 Off |          0 |  
| N/A 25C   P0    67W / 235W |    0MiB / 11439MiB |   95%  Default |  
+-----+-----+-----+-----+-----+-----+
```

```
+-----+  
| Processes:                               GPU Memory |  
| GPU PID Type Process name           Usage |  
+-----+-----+-----+-----+-----+  
| No running processes found             |  
+-----+-----+-----+-----+-----+
```



# nvprof textual Profiles

Sparse Matvec Better

Total Iterations:		100		Time (s):	16.21501	
==95387== Profiling result:						
Time(%)	Time	Calls	Avg	Min	Max	Name
97.02%	15.8332s	101	156.76ms	156.50ms	156.90ms	matvec_114_gpu
1.91%	311.27ms	302	1.0307ms	696.17us	1.0377ms	waxpby_60_gpu
0.89%	145.68ms	200	728.41us	620.77us	843.05us	dot_46_gpu
0.18%	29.020ms	200	145.10us	144.10us	146.27us	dot_46_gpu_red
0.00%	422.63us	200	2.1130us	2.0800us	2.7840us	[CUDA memcpy DtoH]
0.00%	186.50us	200	932ns	896ns	1.6320us	[CUDA memcpy HtoD]

Sparse Matvec Best

nvprof ./cg.exe

Total Iterations:		100		Time (s):	5.814235					
Profiling application: ./cg.x										
==64924== Profiling result:										
Time(%)	Time	Calls	Avg	Min	Max	Name				
87.90%	5.36086s	101	53.078ms	53.044ms	53.107ms	matvec_118_gpu				
5.07%	309.17ms	302	1.0237ms	680.83us	1.0302ms	waxpby_64_gpu				
4.17%	254.03ms	370	686.57us	896ns	1.7418ms	[CUDA memcpy HtoD]				
2.37%	144.67ms	200	723.35us	622.12us	831.88us	dot_50_gpu				
0.47%	28.874ms	200	144.37us	143.39us	145.63us	dot_50_gpu_red				
0.01%	641.54us	2	320.77us	320.55us	320.99us	initialize_vector_26_gpu				
0.01%	414.98us	200	2.0740us	2.0480us	2.8480us	[CUDA memcpy DtoH]				



# Unstructured Data Management

```
module matrix_mod
    implicit none
    type matrix
        integer :: num_rows
        integer :: nnz
        integer, pointer :: row_offsets(:)
        integer, pointer :: cols(:)
        real(8), pointer :: coefs(:)
    end type matrix
    public :: allocate_3d
    ...

```

**!\$omp target exit data map(delete: a)**

```
subroutine allocate_3d (a, n)
    implicit none
    type(matrix) :: a
    allocate(a%row_offsets(num_rows+1))
    allocate(a%cols(nnz))
    allocate(a%coefs(nnz))
    arrow_offsets => a%row_offsets
    acols => a%cols
    acoefs => a%coefs
```

**!\$omp target enter data map(to: a)**

**!\$omp target enter data map(to : a%row\_offsets,a%cols,a%coefs) ...**

```
end subroutine allocate_3d
end module matrix_mod
```

call allocate\_3d(a,n)  
call free\_3d\_(a,n)

Call as often as you need to,  
also across subroutine calls