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    Developed at the Applied Logic, Programming Languages and Systems
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     (ALPS) Laboratory at UTD by Feliks Kluzniak.
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  %% This is an experimental version of a pure Prolog counterpart of verifier.tlp.
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%% The approach is to visit each node at most once, and rewrite the expression
%% to take into account the valuation of propositions in that node.
%% The cost is O( number of nodes ) * O( length of formula ).
%% We don't yet have a proof of correctness, but all the examples work.
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%% Written by Feliks Kluzniak at UTD (March 2009).
%% Last update: 24 April 2009.
:- [ 'operators.pl' ].
:- [ 'normalize.pl' ].
:- [ 'looping_prefix.pl' ].
:- [ 'consistency_checker.pl' ].
:- [ '../../general/higher_order.pl' ].
:- ensure_loaded( library( lists ) ). % Sicstus, reverse/2.
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```
%% Check whether the state satisfies the formula.
%% This is done by checking that it does not satisfy the formula's negation.
%% (We have to apply the conditional, because our tabling interpreter does not
%% support the cut, and we don't yet support negation for coinduction.)
check( State, Formula ) :-
        check_consistency,
            state( State )
            true
        ;
            write( '\"' ),
            write( State ),
            write( '\" is not a state' ),
            nl,
            fail
        write( 'Query for state ' ),
        write( State ),
        write(':'),
        write( Formula ),
        once( normalize( ~ Formula, NormalizedNegationOfFormula ) ),
        write( '(Negated and normalized: '),
        write( NormalizedNegationOfFormula ),
        write(')'),
        nl,
        (
            once( verify( State, NormalizedNegationOfFormula, [] ) )
        ->
            fail
        ;
            true
        ).
```

```
%% verify( + state, + formula, + path ) :
%% Verify whether the formula holds for this state (which we reached by this
%% path).
%% (The formula is our negated thesis, so we are looking for one path.)
verify( S, F, Path ) :-
       rewrite (F, S, NF),
       (
           NF = true
       ->
           show_path ( Path )
       ;
           NF = false
       ->
           fail
       ;
           (
               member( pair( S, F ), Path )
           ->
               (
                  disjunct( g _, F )
               ->
                  show_path( Path )
               ;
                  fail
               )
               once( strip_off_x( NF, NNF ) ),
               trans(S, NS),
               verify( NS, NNF, [ pair( S, F ) | Path ] )
           )
       ) .
show_path ( Path ) :-
       write( 'COUNTEREXAMPLE: ' ),
       reverse ( Path, RevPath ),
       map( first, RevPath, TruePath ),
       write( TruePath ),
       nl.
first( pair( State, _ ), State ).
%% disjunct( +- disjunct, + formula ):
%% Like member, only of an outermost disjunction rather than a list.
disjunct(A, A v _).
disjunct( A, _ v B ) :- disjunct( A, B ).
                ) .
disjunct (A, A
```

```
%% strip_off_x( + formula, - formula ):
%% Strip off the "x" operator from every disjunct, raise an alarm and abort if
%% there are disjuncts that are not so wrapped.

strip_off_x( x A v B, A v NB ) :- strip_off_x( B, NB ).

strip_off_x( x F, F ).

strip_off_x( F, _ ) :-
    F \= x F,
    write( 'Formula not in X : ' ),
    write( F ),
    nl,
    abort.
```

```
%% rewrite( + formula, + state, - new formula ):
%% The formula has been normalized, so that negations are applied only to
%% propositions.
rewrite(F, S, NF):-
       once( r(F, S, NF)).
r(A \lor B, S, NF) := r(A, S, NA), r(B, S, NB),
                      simplify ( NA v NB, NF ).
r(A ^B, S, NF) := r(A, S, NA), r(B, S, NB),
                      simplify( NA ^ NB, NF ).
r(xA,_,xA).
r(fA,S,NF):- r(A,S,NA),
                      simplify ( NA v x f A, NF ).
r(gA,S,NF):- r(A,S,NA),
                      simplify ( NA ^{\circ} x g A, NF ).
r(AuB, S, NF) := r(A, S, NA), r(B, S, NB),
                      simplify( NA ^{\circ} x (A u B), Conj ),
                      simplify( NB v Conj, NF ).
r(ArB, S, NF):-
                     r(A, S, NA), r(B, S, NB),
                      simplify( NB ^ NA, Conj ),
                      simplify (Conj v x (A r B), NF).
r( \tilde{P} , S, NF ) :- proposition( P ),
                      ( holds( S, P ) \rightarrow NF = false
                                        NF = true
                      ;
                      ) .
r(P
       , S, NF ) :- proposition(P),
                      ( holds( S, P ) \rightarrow NF = true
                                        NF = false
                      ) .
```

```
%% simplify( + formula, + state, - new formula ):
 %% The formula has now been rewritten: simplify the result.
simplify( F, NF ) :-
   once(s(F,NF)).
용
                        v_{-} , true ).
s( true
s( _
                           v true , true ).

      s(false
      v A
      , NA
      ):-
      s(A
      , NA).

      s(A
      v false, NA
      ):-
      s(A
      , NA).

      s(A v A v B
      , NF
      ):-
      s(A v B
      , NF).

      s(A v B)
      v C
      , NF
      ):-
      s(A v B v C, NF).

s(false ^{\prime} _ , false).
                           ^ false, false ).
s( _

      S( true
      ^ A
      , NA
      ) :- s( A
      , NA ).

      S( A
      ^ true
      , NA
      ) :- s( A
      , NA ).

      S( A
      ^ A
      A
      B
      , NF
      ) :- s( A
      B
      , NF ).

      S( (A
      B)
      ^ C
      , NF
      ) :- s( A
      B
      ^ C
      , NF ).

      S( X
      A
      ^ X
      B
      , X
      NF
      ) :- s( A
      B
      , NF
      ).

s(F
                                              , F ).
```