# Network Security and Cryptography, Exercises 1 (unassessed)

Decrypt the following ciphertext encrypted with the columnar transposition cipher (you need to find the key by brute-force search):

# AVUEVLETSEISBNACBOOLEOBTILBDLCOBOOE [1 point]

Ciphertext: AVUEVLETSEISBNACBOOLEOBTILBDLCOBOOE

Total No. of characters = 35 Factors of 35: 1, 5, 7, 35

Key is one of the factors, but 1 and 35 can not be the key because 1 as key will give 35 rows and 35 will give 1 row in both the case they do not make a sensible sentence

With 5 as a key:

35/5 = 7

That is 7 rows:

A	T	A	О	L
V	S	С	В	С
U	Е	В	T	О
Е	I	О	I	В
V	S	О	L	О
L	В	L	В	О
Е	N	Е	D	Е

This makes a gibberish sentence, hence cannot be the key.

With 7 as a key:

35/7 = 5

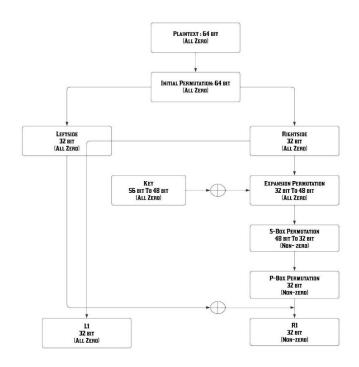
That is 5 rows:

A	L	I	С	Е	L	0
V	Е	S	В	0	В	В
U	T	В	0	В	D	О
Е	S	N	0	T	L	0
V	Е	A	L	I	С	Е

Sentence: ALICE LOVES BOB BUT BOB DOES NOT LOVE ALICE

Hence the key is 7

# What is the output of the first round of the DES algorithm when the plaintext and the key are both all zeros? [1 point]



## **S-box Permutation:**

	S1	S2	S3	S4	S5	S6	S7	S8
Input	000000	000000	000000	000000	000000	000000	000000	000000
(6 bit)								
Output	14	15	10	7	2	12	4	13
(Decimal)								
Output	1110	1111	1010	0111	0010	1100	0100	1101
(Binary)								
(4 bit)								

	P-Box Permutation														
16,	7,	20,	21,	29,	12,	28,	17,	1,	15,	23,	26,	5,	18,	31,	10
2,	8,	24,	14,	32,	27,	3,	9,	19,	13,	30,	6,	22,	11,	4,	2.5

11011000110110001101101110111100

XOR

11011000110110001101101110111100

# **Output of Round 1:**

Remember that it is desirable for good block ciphers that a change in one input bit affects many output bits, a property that is called *diffusion* or *avalanche effect*. We will try to get a feeling for the avalanche property of DES. Let x be all zeros (0x0000000000000000) and y be all zeros except 1 in the 13th bit (0x000800000000000). Let the key be all zeros. After just one round, how many bits in the block are different when x is the input, compared to when y is the input? What about after two rounds? Three? Four? (For this exercise, you might like to search for an implementation of DES on the web, and download it and modify it to output the answers.) [2 points]

#### Code:

```
def hex2bin(s):
  mp = \{'0': "0000",
      '1': "0001",
      '2': "0010",
      '3': "0011",
      '4': "0100",
      '5': "0101",
      '6': "0110",
      '7': "0111",
      '8': "1000",
      '9': "1001",
      'A': "1010",
      'B': "1011",
      'C': "1100",
      'D': "1101",
      'E': "1110",
      'F': "1111"}
  bin = ""
  for i in range(len(s)):
    bin = bin + mp[s[i]]
  return bin
# Binary to hexadecimal conversion
def bin2hex(s):
  mp = {"0000": '0',}
      "0001": '1',
      "0010": '2'.
      "0011": '3',
      "0100": '4',
      "0101": '5',
      "0110": '6',
      "0111": '7',
      "1000": '8'.
      "1001": '9',
      "1010": 'A',
      "1011": 'B',
      "1100": 'C'.
      "1101": 'D'.
      "1110": 'E',
      "1111": 'F'}
  hex = ""
  for i in range(0, len(s), 4):
     ch = ""
     ch = ch + s[i]
     ch = ch + s[i + 1]
     ch = ch + s[i + 2]
     ch = ch + s[i + 3]
     hex = hex + mp[ch]
```

```
return hex
# Binary to decimal conversion
def bin2dec(binary):
  binary1 = binary
  decimal, i, n = 0, 0, 0
  while(binary != 0):
     dec = binary % 10
     decimal = decimal + dec * pow(2, i)
     binary = binary//10
     i += 1
  return decimal
# Decimal to binary conversion
def dec2bin(num):
  res = bin(num).replace("0b", "")
  if(len(res) \% 4 != 0):
     div = len(res) / 4
     div = int(div)
     counter = (4 * (div + 1)) - len(res)
     for i in range(0, counter):
       res = '0' + res
  return res
# Permute function to rearrange the bits
def permute(k, arr, n):
  permutation = ""
  for i in range(0, n):
     permutation = permutation + k[arr[i] - 1]
  return permutation
# shifting the bits towards left by nth shifts
def shift left(k, nth shifts):
  for i in range(nth_shifts):
     for j in range(1, len(k)):
       s = s + k[j]
     s = s + k[0]
     \mathbf{k} = \mathbf{s}
     s = ""
  return k
# calculating xow of two strings of binary number a and b
def xor(a, b):
  ans = ""
  for i in range(len(a)):
     if a[i] == b[i]:
       ans = ans + "0"
     else:
       ans = ans + "1"
```

# Table of Position of 64 bits at initial level: Initial Permutation Table

initial\_perm = [58, 50, 42, 34, 26, 18, 10, 2, 60, 52, 44, 36, 28, 20, 12, 4,

return ans

```
62, 54, 46, 38, 30, 22, 14, 6,
           64, 56, 48, 40, 32, 24, 16, 8,
           57, 49, 41, 33, 25, 17, 9, 1,
           59, 51, 43, 35, 27, 19, 11, 3,
           61, 53, 45, 37, 29, 21, 13, 5,
           63, 55, 47, 39, 31, 23, 15, 7]
# Expansion D-box Table
exp d = [32, 1, 2, 3, 4, 5, 4, 5,
      6, 7, 8, 9, 8, 9, 10, 11,
      12, 13, 12, 13, 14, 15, 16, 17,
      16, 17, 18, 19, 20, 21, 20, 21,
      22, 23, 24, 25, 24, 25, 26, 27,
      28, 29, 28, 29, 30, 31, 32, 1]
# Straight Permutation Table
per = [16, 7, 20, 21,
    29, 12, 28, 17,
    1, 15, 23, 26,
    5, 18, 31, 10,
    2, 8, 24, 14,
    32, 27, 3, 9,
     19, 13, 30, 6,
    22, 11, 4, 25]
# S-box Table
sbox = [[[14, 4, 13, 1, 2, 15, 11, 8, 3, 10, 6, 12, 5, 9, 0, 7],
      [0, 15, 7, 4, 14, 2, 13, 1, 10, 6, 12, 11, 9, 5, 3, 8],
      [4, 1, 14, 8, 13, 6, 2, 11, 15, 12, 9, 7, 3, 10, 5, 0],
      [15, 12, 8, 2, 4, 9, 1, 7, 5, 11, 3, 14, 10, 0, 6, 13]],
     [[15, 1, 8, 14, 6, 11, 3, 4, 9, 7, 2, 13, 12, 0, 5, 10],
      [3, 13, 4, 7, 15, 2, 8, 14, 12, 0, 1, 10, 6, 9, 11, 5],
      [0, 14, 7, 11, 10, 4, 13, 1, 5, 8, 12, 6, 9, 3, 2, 15],
      [13, 8, 10, 1, 3, 15, 4, 2, 11, 6, 7, 12, 0, 5, 14, 9]],
     [[10, 0, 9, 14, 6, 3, 15, 5, 1, 13, 12, 7, 11, 4, 2, 8],
      [13, 7, 0, 9, 3, 4, 6, 10, 2, 8, 5, 14, 12, 11, 15, 1],
      [13, 6, 4, 9, 8, 15, 3, 0, 11, 1, 2, 12, 5, 10, 14, 7],
      [1, 10, 13, 0, 6, 9, 8, 7, 4, 15, 14, 3, 11, 5, 2, 12]],
     [[7, 13, 14, 3, 0, 6, 9, 10, 1, 2, 8, 5, 11, 12, 4, 15],
      [13, 8, 11, 5, 6, 15, 0, 3, 4, 7, 2, 12, 1, 10, 14, 9],
      [10, 6, 9, 0, 12, 11, 7, 13, 15, 1, 3, 14, 5, 2, 8, 4],
      [3, 15, 0, 6, 10, 1, 13, 8, 9, 4, 5, 11, 12, 7, 2, 14]],
     [[2, 12, 4, 1, 7, 10, 11, 6, 8, 5, 3, 15, 13, 0, 14, 9],
      [14, 11, 2, 12, 4, 7, 13, 1, 5, 0, 15, 10, 3, 9, 8, 6],
      [4, 2, 1, 11, 10, 13, 7, 8, 15, 9, 12, 5, 6, 3, 0, 14],
      [11, 8, 12, 7, 1, 14, 2, 13, 6, 15, 0, 9, 10, 4, 5, 3]],
     [[12, 1, 10, 15, 9, 2, 6, 8, 0, 13, 3, 4, 14, 7, 5, 11],
      [10, 15, 4, 2, 7, 12, 9, 5, 6, 1, 13, 14, 0, 11, 3, 8],
      [9, 14, 15, 5, 2, 8, 12, 3, 7, 0, 4, 10, 1, 13, 11, 6],
      [4, 3, 2, 12, 9, 5, 15, 10, 11, 14, 1, 7, 6, 0, 8, 13]],
     [[4, 11, 2, 14, 15, 0, 8, 13, 3, 12, 9, 7, 5, 10, 6, 1],
      [13, 0, 11, 7, 4, 9, 1, 10, 14, 3, 5, 12, 2, 15, 8, 6],
      [1, 4, 11, 13, 12, 3, 7, 14, 10, 15, 6, 8, 0, 5, 9, 2],
      [6, 11, 13, 8, 1, 4, 10, 7, 9, 5, 0, 15, 14, 2, 3, 12]],
```

[[13, 2, 8, 4, 6, 15, 11, 1, 10, 9, 3, 14, 5, 0, 12, 7],

```
[1, 15, 13, 8, 10, 3, 7, 4, 12, 5, 6, 11, 0, 14, 9, 2],
     [7, 11, 4, 1, 9, 12, 14, 2, 0, 6, 10, 13, 15, 3, 5, 8],
     [2, 1, 14, 7, 4, 10, 8, 13, 15, 12, 9, 0, 3, 5, 6, 11]]]
# Final Permutation Table
final_perm = [40, 8, 48, 16, 56, 24, 64, 32,
        39, 7, 47, 15, 55, 23, 63, 31,
        38, 6, 46, 14, 54, 22, 62, 30,
        37, 5, 45, 13, 53, 21, 61, 29,
        36, 4, 44, 12, 52, 20, 60, 28,
        35, 3, 43, 11, 51, 19, 59, 27,
        34, 2, 42, 10, 50, 18, 58, 26,
        33, 1, 41, 9, 49, 17, 57, 25]
def encrypt(pt, rkb, rk):
  pt = hex2bin(pt)
  # Initial Permutation
  pt = permute(pt, initial perm, 64)
  print("After initial permutation", bin2hex(pt))
  # Splitting
  left = pt[0:32]
  right = pt[32:64]
  for i in range(0, 16):
     # Expansion D-box: Expanding the 32 bits data into 48 bits
     right expanded = permute(right, exp d, 48)
     # XOR RoundKey[i] and right expanded
     xor x = xor(right expanded, rkb[i])
     # S-boxex: substituting the value from s-box table by calculating row and column
     sbox str = ""
     for j in range(0, 8):
       row = bin2dec(int(xor_x[j * 6] + xor_x[j * 6 + 5]))
       col = bin2dec(
         int(xor x[j*6+1] + xor x[j*6+2] + xor x[j*6+3] + xor x[j*6+4]))
       val = sbox[j][row][col]
       sbox str = sbox str + dec2bin(val)
     # Straight D-box: After substituting rearranging the bits
     sbox_str = permute(sbox_str, per, 32)
     # XOR left and sbox str
     result = xor(left, sbox str)
    left = result
     # Swapper
     if(i != 15):
       left, right = right, left
     print("Round", i + 1, "", bin2hex(left),
         " ", bin2hex(right), " ", rk[i])
  # Combination
  combine = left + right
  # Final permutation: final rearranging of bits to get cipher text
  cipher_text = permute(combine, final_perm, 64)
  return cipher_text
pt = "00000000000000000"
key = "00000000000000000"
```

```
# Key generation
# --hex to binary
key = hex2bin(key)
# --parity bit drop table
keyp = [57, 49, 41, 33, 25, 17, 9,
     1, 58, 50, 42, 34, 26, 18,
     10, 2, 59, 51, 43, 35, 27,
     19, 11, 3, 60, 52, 44, 36,
     63, 55, 47, 39, 31, 23, 15,
     7, 62, 54, 46, 38, 30, 22,
     14, 6, 61, 53, 45, 37, 29,
     21, 13, 5, 28, 20, 12, 4
# getting 56 bit key from 64 bit using the parity bits
key = permute(key, keyp, 56)
# Number of bit shifts
shift table = [1, 1, 2, 2,
         2, 2, 2, 2,
         1, 2, 2, 2,
         2, 2, 2, 1]
# Key- Compression Table: Compression of key from 56 bits to 48 bits
key comp = [14, 17, 11, 24, 1, 5,
       3, 28, 15, 6, 21, 10,
       23, 19, 12, 4, 26, 8,
       16, 7, 27, 20, 13, 2,
       41, 52, 31, 37, 47, 55,
       30, 40, 51, 45, 33, 48,
       44, 49, 39, 56, 34, 53,
       46, 42, 50, 36, 29, 32]
# Splitting
left = key[0:28] # rkb for RoundKeys in binary
right = key[28:56] # rk for RoundKeys in hexadecimal
rkb = []
rk = []
for i in range(0, 16):
  # Shifting the bits by nth shifts by checking from shift table
  left = shift_left(left, shift_table[i])
  right = shift_left(right, shift_table[i])
  # Combination of left and right string
  combine_str = left + right
  # Compression of key from 56 to 48 bits
  round_key = permute(combine_str, key_comp, 48)
  rkb.append(round_key)
  rk.append(bin2hex(round_key))
print("Encryption")
cipher_text = bin2hex(encrypt(pt, rkb, rk))
print("Cipher Text : ", cipher_text)
print("Decryption")
rkb rev = rkb[::-1]
rk rev = rk[::-1]
text = bin2hex(encrypt(cipher text, rkb rev, rk rev))
```

## Output 1:

# Encryption

Round 1 00000000 D8D8DBBC 000000000000

Round 2 D8D8DBBC E73AED4F 000000000000

Round 3 E73AED4F 5BFA6BA6 0000000000000

Round 4 5BFA6BA6 5E45E257 0000000000000 ...

# Output 2:

# Encryption

Round 1 00000200 D8F8DBB4 000000000000

Round 2 D8F8DBB4 E528E26B 0000000000000

Round 3 E528E26B 87D70236 000000000000

Round 4 87D70236 3D14F53C 000000000000 ...

#### Hence,

After 1 round 8 bits are different

After 2 round 16 bits are different

After 3 round 48 bits are different

After 4 round 56 bits are different

Consider AES with 128-bit keys. Assume that the principal key k is all-zeros. Then the initial round key  $(k_0)$  is also all-zeros. What is the first round subkey  $(k_1)$  and the second round subkey  $(k_2)$ ? (Again, for this exercise and the following one, you might like to use a computer for help.) [2 points]

#### Code:

```
from Crypto.Cipher import AES
from Crypto.Util.Padding import pad, unpad
def xor bytes(a, b):
  return bytes(x \wedge y for x, y in zip(a, b))
def rot word(word):
  return word[1:] + word[:1]
def sub word(word):
  s box = [
              2 3 4 5 6 7 8 9 A B C D E F
    # 0
     0x63, 0x7c, 0x77, 0x7b, 0xf2, 0x6b, 0x6f, 0xc5, 0x30, 0x01, 0x67, 0x2b, 0xfe, 0xd7, 0xab, 0x76,
     0xca, 0x82, 0xc9, 0x7d, 0xfa, 0x59, 0x47, 0xf0, 0xad, 0xd4, 0xa2, 0xaf, 0x9c, 0xa4, 0x72, 0xc0,
     0xb7, 0xfd, 0x93, 0x26, 0x36, 0x3f, 0xf7, 0xcc, 0x34, 0xa5, 0xe5, 0xf1, 0x71, 0xd8, 0x31, 0x15,
     0x04, 0xc7, 0x23, 0xc3, 0x18, 0x96, 0x05, 0x9a, 0x07, 0x12, 0x80, 0xe2, 0xeb, 0x27, 0xb2, 0x75,
     0x09, 0x83, 0x2c, 0x1a, 0x1b, 0x6e, 0x5a, 0xa0, 0x52, 0x3b, 0xd6, 0xb3, 0x29, 0xe3, 0x2f, 0x84,
     0x53, 0xd1, 0x00, 0xed, 0x20, 0xfc, 0xb1, 0x5b, 0x6a, 0xcb, 0xbe, 0x39, 0x4a, 0x4c, 0x58, 0xcf,
     0xd0, 0xef, 0xaa, 0xfb, 0x43, 0x4d, 0x33, 0x85, 0x45, 0xf9, 0x02, 0x7f, 0x50, 0x3c, 0x9f, 0xa8,
     0x51, 0xa3, 0x40, 0x8f, 0x92, 0x9d, 0x38, 0xf5, 0xbc, 0xb6, 0xda, 0x21, 0x10, 0xff, 0xf3, 0xd2,
     0xcd, 0x0c, 0x13, 0xec, 0x5f, 0x97, 0x44, 0x17, 0xc4, 0xa7, 0x7e, 0x3d, 0x64, 0x5d, 0x19, 0x73,
     0x60, 0x81, 0x4f, 0xdc, 0x22, 0x2a, 0x90, 0x88, 0x46, 0xee, 0xb8, 0x14, 0xde, 0x5e, 0x0b, 0xdb,
     0xe0, 0x32, 0x3a, 0x0a, 0x49, 0x06, 0x24, 0x5c, 0xc2, 0xd3, 0xac, 0x62, 0x91, 0x95, 0xe4, 0x79,
     0xe7, 0xc8, 0x37, 0x6d, 0x8d, 0xd5, 0x4e, 0xa9, 0x6c, 0x56, 0xf4, 0xea, 0x65, 0x7a, 0xae, 0x08,
     0xba, 0x78, 0x25, 0x2e, 0x1c, 0xa6, 0xb4, 0xc6, 0xe8, 0xdd, 0x74, 0x1f, 0x4b, 0xbd, 0x8b, 0x8a,
     0x70, 0x3e, 0xb5, 0x66, 0x48, 0x03, 0xf6, 0x0e, 0x61, 0x35, 0x57, 0xb9, 0x86, 0xc1, 0x1d, 0x9e,
     0xe1, 0xf8, 0x98, 0x11, 0x69, 0xd9, 0x8e, 0x94, 0x9b, 0x1e, 0x87, 0xe9, 0xce, 0x55, 0x28, 0xdf,
     0x8c, 0xa1, 0x89, 0x0d, 0xbf, 0xe6, 0x42, 0x68, 0x41, 0x99, 0x2d, 0x0f, 0xb0, 0x54, 0xbb, 0x16
  return bytes([s box[b] for b in word])
def key expansion(key):
  RCON = [
     0x01, 0x02, 0x04, 0x08, 0x10, 0x20, 0x40, 0x80, 0x1B, 0x36, 0x6C, 0xD8, 0xAB, 0x4D, 0x9A
  expanded keys = [key]
  for i in range(1, 11):
     temp = expanded keys[-1][-4:]
     temp = sub word(rot word(temp))
     temp = xor_bytes(temp, bytes([RCON[i-1], 0, 0, 0]))
     new key = xor bytes(expanded keys[-1][:4], temp)
     for i in range(1, 4):
       new key += xor bytes(new key[-4:], expanded keys[-1][j * 4:(j + 1) * 4])
     expanded keys.append(new key)
  return expanded keys[1], expanded keys[2]
# Initial key is all-zeros for 128-bit AES key
initial key = bytes(16)
k1, k2 = key expansion(initial key)
print(f''k1: \{k1.hex()\}'')
print(f"k2: {k2.hex()}")
```

## **Output:**

k1: 62636363626363636263636362636363 k2: 9b9898c9f9fbfbaa9b9898c9f9fbfbaa

Initial key is 128 bit which is all zeros. Initial sub key (k0) is 128 bit which is all zeros.

$$k1 = [w4, w5, w6, w7]$$

w' for k1:

w3 = 0x00000000

Rotword(w3) = 0x00000000

Subword(w3) = 0x63636363

XOR with Rcon[1]

Rcon[1] = 01000000

0x63636363 XOR 0x01000000

$$w' = 0x62636363$$

$$w4 = w' XOR w3 = 0x62636363$$

$$w5 = w4 \text{ XOR } w1 = 0x62636363$$

$$w6 = w5 \text{ XOR } w2 = 0x62636363$$

$$w7 = w6 \text{ XOR } w3 = 0x62636363$$

$$k1 = [w4, w5, w6, w7] = [0x62636363, 0x62636363, 0x62636363, 0x62636363]$$

$$k2 = [w8, w9, w10, w11]$$

w' for k2:

w7 = 0x62636363

Rotword(w3) = 0x63636362

Subword(w3) = 0xfbfbfbaa

XOR with Rcon[2]

Rcon[1] = 02000000

0xfbfbfbaa XOR 0x02000000

w' = 0xf9fbfbaa

$$w8 = w' XOR w7 = 0x9b9898c9$$

$$w9 = w4 \text{ XOR } w5 = 0 \text{xf9fbfbaa}$$

$$w10 = w5 \text{ XOR } w6 = 0x9b9898c9$$

$$w11 = w6 \text{ XOR } w7 = 0xf9fbfbaa$$

$$k1 = [w8, w9, w10, w11] = [0x9b9898c9, 0xf9fbfbaa, 0x9b9898c9, 0xf9fbfbaa]$$

Again, using AES-128, assume that the principal key is all-zeros, and that the plaintext is also all-zeros. What is the output of the first round, and what is the output of the second round? [2 points]

#### Code:

```
import numpy as np
# AES S-box used for SubWord operation in key expansion
sbox = np.array([
  [0x63, 0x7c, 0x77, 0x7b, 0xf2, 0x6b, 0x6f, 0xc5, 0x30, 0x01, 0x67, 0x2b, 0xfe, 0xd7, 0xab, 0x76],
  [0xca, 0x82, 0xc9, 0x7d, 0xfa, 0x59, 0x47, 0xf0, 0xad, 0xd4, 0xa2, 0xaf, 0x9c, 0xa4, 0x72, 0xc0],
  [0xb7, 0xfd, 0x93, 0x26, 0x36, 0x3f, 0xf7, 0xcc, 0x34, 0xa5, 0xe5, 0xf1, 0x71, 0xd8, 0x31, 0x15],
  [0x04, 0xc7, 0x23, 0xc3, 0x18, 0x96, 0x05, 0x9a, 0x07, 0x12, 0x80, 0xe2, 0xeb, 0x27, 0xb2, 0x75],
  [0x09, 0x83, 0x2c, 0x1a, 0x1b, 0x6e, 0x5a, 0xa0, 0x52, 0x3b, 0xd6, 0xb3, 0x29, 0xe3, 0x2f, 0x84],
  [0x53, 0xd1, 0x00, 0xed, 0x20, 0xfc, 0xb1, 0x5b, 0x6a, 0xcb, 0xbe, 0x39, 0x4a, 0x4c, 0x58, 0xcf],
  [0xd0, 0xef, 0xaa, 0xfb, 0x43, 0x4d, 0x33, 0x85, 0x45, 0xf9, 0x02, 0x7f, 0x50, 0x3c, 0x9f, 0xa8],
  [0x51, 0xa3, 0x40, 0x8f, 0x92, 0x9d, 0x38, 0xf5, 0xbc, 0xb6, 0xda, 0x21, 0x10, 0xff, 0xf3, 0xd2],
  [0xcd, 0x0c, 0x13, 0xec, 0x5f, 0x97, 0x44, 0x17, 0xc4, 0xa7, 0x7e, 0x3d, 0x64, 0x5d, 0x19, 0x73],
  [0x60, 0x81, 0x4f, 0xdc, 0x22, 0x2a, 0x90, 0x88, 0x46, 0xee, 0xb8, 0x14, 0xde, 0x5e, 0x0b, 0xdb],
  [0xe0, 0x32, 0x3a, 0x0a, 0x49, 0x06, 0x24, 0x5c, 0xc2, 0xd3, 0xac, 0x62, 0x91, 0x95, 0xe4, 0x79],
  [0xe7, 0xc8, 0x37, 0x6d, 0x8d, 0xd5, 0x4e, 0xa9, 0x6c, 0x56, 0xf4, 0xea, 0x65, 0x7a, 0xae, 0x08],
  [0xba, 0x78, 0x25, 0x2e, 0x1c, 0xa6, 0xb4, 0xc6, 0xe8, 0xdd, 0x74, 0x1f, 0x4b, 0xbd, 0x8b, 0x8a],
  [0x70, 0x3e, 0xb5, 0x66, 0x48, 0x03, 0xf6, 0x0e, 0x61, 0x35, 0x57, 0xb9, 0x86, 0xc1, 0x1d, 0x9e],
  [0xe1, 0xf8, 0x98, 0x11, 0x69, 0xd9, 0x8e, 0x94, 0x9b, 0x1e, 0x87, 0xe9, 0xce, 0x55, 0x28, 0xdf],
  [0x8c, 0xa1, 0x89, 0x0d, 0xbf, 0xe6, 0x42, 0x68, 0x41, 0x99, 0x2d, 0x0f, 0xb0, 0x54, 0xbb, 0x16]
1)
# Rcon used in key expansion
rcon = [0x01, 0x02, 0x04, 0x08, 0x10, 0x20, 0x40, 0x80, 0x1B, 0x36]
# Helper functions for key expansion
def sub_word(word):
  # Applies S-box to a word (4 bytes)
  return [sbox[byte \gg 4][byte & 0x0F] for byte in word]
def rot word(word):
  # Rotates a word (4 bytes)
  return word[1:] + word[:1]
def xor words(word1, word2):
  return [b1 \(^b2\) for b1, b2 in zip(word1, word2)]
# Key expansion function to generate round keys
def key expansion(key):
  key\_schedule = [list(key[i:i+4]) \text{ for i in range}(0, len(key), 4)]
  for i in range(4, 44):
     temp = key schedule[i - 1]
     if i % 4 == 0:
       temp = xor words(sub word(rot word(temp)), [rcon[(i/4) - 1], 0x00, 0x00, 0x00])
     key_schedule.append(xor_words(key_schedule[i - 4], temp))
  return key schedule
# Now, the full AES encryption process including rounds
# Helper functions for AES encryption steps
# ShiftRows step for AES (modifies state matrix row-wise)
def shift rows(state):
```

```
state = np.reshape(state, (4, 4)) # 4x4 matrix
  state = state.tolist() # Convert to list for manipulation
  # Perform the row shifts
  state[1] = state[1][1:] + state[1][:1] # 1-byte shift to the left
  state[2] = state[2][2:] + state[2][:2] # 2-byte shift to the left
  state[3] = state[3][3:] + state[3][:3] # 3-byte shift to the left
  return np.reshape(state, (16,)) # Flatten the state back into a single array
# AddRoundKey step for AES (XOR with round key)
def add round key(state, round key):
  return [s ^ rk for s, rk in zip(state, round key)]
# SubBytes step for AES (substitute bytes using the S-box)
def sub bytes(state):
  return [sbox[byte \gg 4][byte & 0x0F] for byte in state]
# MixColumns step for AES (Galois field multiplication)
def mix columns(state):
  state = np.reshape(state, (4, 4)).T # Transpose to make it column-major
  for i in range(4):
     a = state[i, :]
     # Galois field multiplication and column mixing
     state[i, :] = [
       gf mult(a[0], 2) ^ gf mult(a[1], 3) ^ a[2] ^ a[3],
       a[0] ^ gf_{mult}(a[1], 2) ^ gf_{mult}(a[2], 3) ^ a[3],
       a[0] ^ a[1] ^ gf mult(a[2], 2) ^ gf mult(a[3], 3),
       gf mult(a[0], 3) ^ a[1] ^ a[2] ^ gf mult(a[3], 2)
  return np.reshape(state.T, (16,))
# Helper function for Galois field multiplication (used in MixColumns)
def gf mult(x, y):
  result = 0
  for i in range(8):
     if y & 1:
       result = x
     carry = x & 0x80
     x <<= 1
     if carry:
       x = 0x1b
     y >>= 1
  return result & 0xFF
# AES encryption round function
def aes round(state, round key, mix columns step=True):
  state = sub bytes(state)
                                # SubBytes
                                # ShiftRows
  state = shift rows(state)
  if mix columns step:
     state = mix_columns(state) # MixColumns (not in the last round)
  state = add_round_key(state, round_key) # AddRoundKey
  return state
# Initial all-zero plaintext and key
plaintext = [0x00] * 16
initial key = [0x00] * 16
# Generate the key schedule (44 words, or 11 round keys)
key schedule = key expansion(initial key)
# Initial AddRoundKey (round 0)
state = add round key(plaintext, initial key)
```

```
# First round
round_1_key = [byte for word in key_schedule[4:8] for byte in word]
state = aes_round(state, round_1_key)
# Store the output after the first round
first round output = state
# Second round
round 2 key = [byte for word in key schedule[8:12] for byte in word]
state = aes round(state, round 2 key, mix columns step=True)
# Store the output after the second round
second round output = state
# Formatting the output as hexadecimal strings
first_round_output_hex = ".join([f'{byte:02x}' for byte in first_round_output])
second_round_output_hex = ".join([f'{byte:02x}' for byte in second_round_output])
# Output the results
print("First round output:", first_round_output_hex)
print("Second round output:", second_round_output_hex)
```

## **Output:**

First round output: 0100000001000000100000010000000 Second round output: c6e4e48b8587b9f7e7dac5b5bba687d6

Programming exercise.

Let MY24SHA a hash function which outputs the first 24 bits (6 nibbles) of SHA-1. For example, SHA-1 of "mark" is

f1b5a91d4d6ad523f2610114591c007e75d15084

so, the MY24SHA of "mark" is f1b5a9.

Find any collision for MY24SHA. Note: you should find two strings such that the unix command

echo -n str | sha1sum - | cut -c1-6

produces the same answer when *str* is replaced by each string. To enable me to verify your answer, please make sure the two strings are typable on a regular keyboard!

Hint: You should not write the code for SHA-1; you should use an existing library. [2 points]

#### Code:

```
import hashlib
def my24sha(input_string):
  sha1_hash = hashlib.sha1(input_string.encode()).hexdigest()
  return sha1_hash[:6]
collision\_found = False
attempts = 0
while not collision found:
  attempts += 1
  str1 = 'str' + str(attempts)
  str2 = 'str' + str(attempts + 1)
  if my24sha(str1) == my24sha(str2):
     collision found = True
print("Collision found!")
print("String 1: ", str1)
print("String 2: ", str2)
print("MY24SHA of both strings: ", my24sha(str1))
```

#### **Output:**

Collision found!

String 1: str6067189 String 2: str6067190

MY24SHA of both strings: 096b70