### PRIORITY QUEUES & HEAPS

Heaps more fun than a barrel of monkeys

# PRIORITY QUEUE APPLICATIONS

Prioritizing data packets in routers

Tracking unexplored routes in path-finding

Bayesian spam filtering

Data compression

OS: load balancing, interrupt handling

### **PRIORITY QUEUES**

Abstract data type (ADT) that is like a queue, but each element has an associated priority number

Elements with smaller priority values are higher priority and are dequeued before lower priority elements (VIPs skip to the front of the queue)

How can you implement this ADT? (many ways)

### **PRIORITY QUEUES**

Almost always implemented with a heap

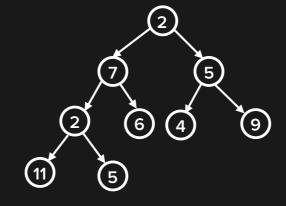
Elements are inserted in O(log n) time instead of O(n) time for a sorted array or linked list

Partial ordering happens with each insertion, so the cost of ordering is distributed across insertion and deletion instead of all at once

# **COMPLETE BINARY TREE**

Every level is completely filled except lowest level, and nodes on lowest level are as far left as possible

Height is always O(log<sub>2</sub> n)



# BINARY HEAP DEFINITION A complete binary tree Satisfies heap ordering property: min-heap - each node is less than or equal to its children (min value is root) max-heap - each node is greater than or equal to its children (max value is root) min-heap min-heap

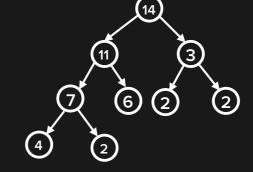
- 2 flavors of heap ordering --- just a different convention
- in heaps we use both, and care about being able to do both
- min-heap:
  - min at the root
  - most accessible
  - how a priority queue works!
  - smallest # comes out first

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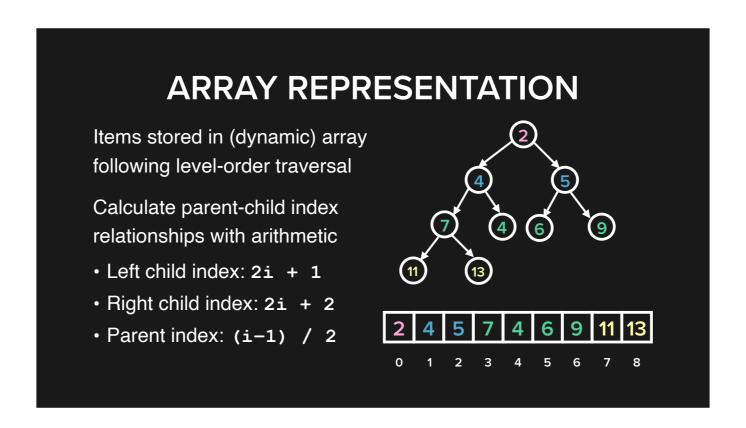


Heaps are *not* completely sorted (like a binary search tree)

Heaps are only "partially ordered"



max-heap

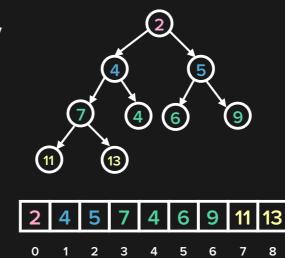


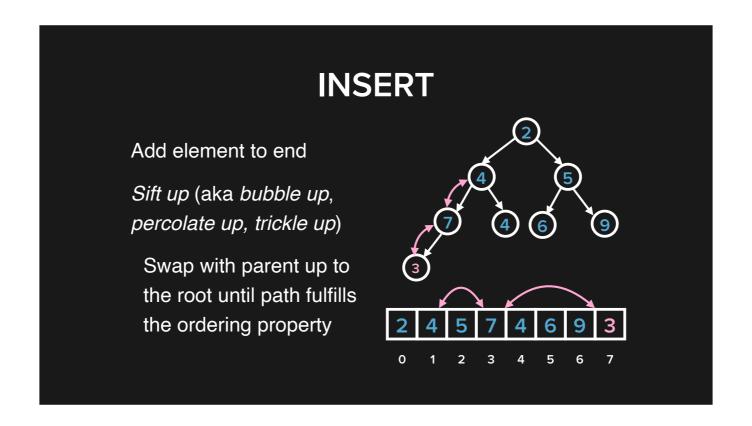
integer math always truncates a decimal value

### **ADVANTAGES**

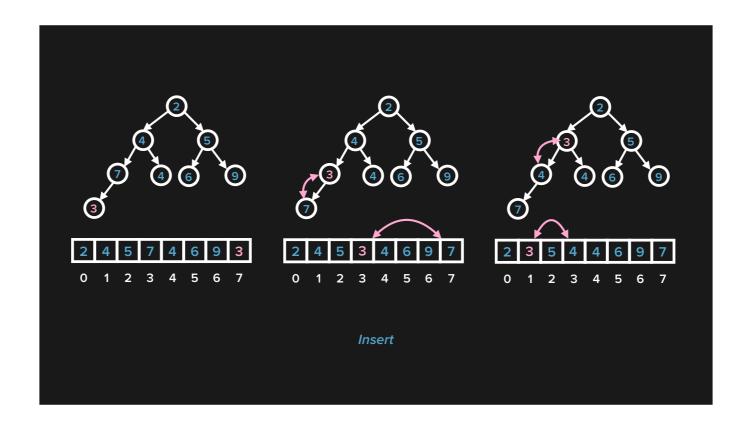
Uses less memory than binary tree represented with nodes (avoids node objects containing 3 pointers: data, left, right child)

Allows sorting an array inplace (heapsort)





insert = enqueue in priority queue
just talking about # bc it's simple --- all that is important is that the objects are orderable / comparable
tree doesn't exist in memory, just a way to perceive / view it
insert on far left side of last level, next spot in reading order, must stay a complete tree
then we do a new kind of op called "sift up"



use the math trick to go to parent index

Left child index: 2i + 1 Right child index: 2i + 2 Parent index: (i-1) / 2

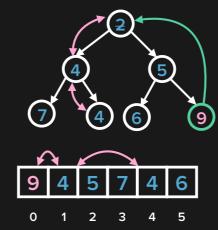
worst case: log n

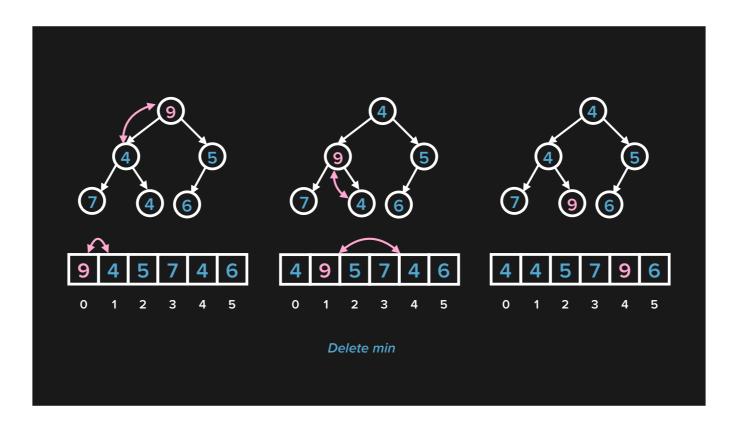
### **DELETE MIN / MAX**

Replace root with last element

Sift down (aka bubble down, percolate down, trickle down)

Swap with smaller child (min) or larger child (max) until trio fulfills the ordering property

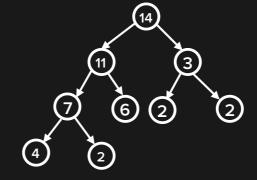




# **OTHER METHODS**

peek (aka find-min or findmax) returns the root value

size (aka count or length) returns number of elements



max-heap

### **HEAPIFY**

Input is an array (usually unsorted, unordered)

Output is an array that satisfies the binary heap ordering property

### **HEAPIFY**

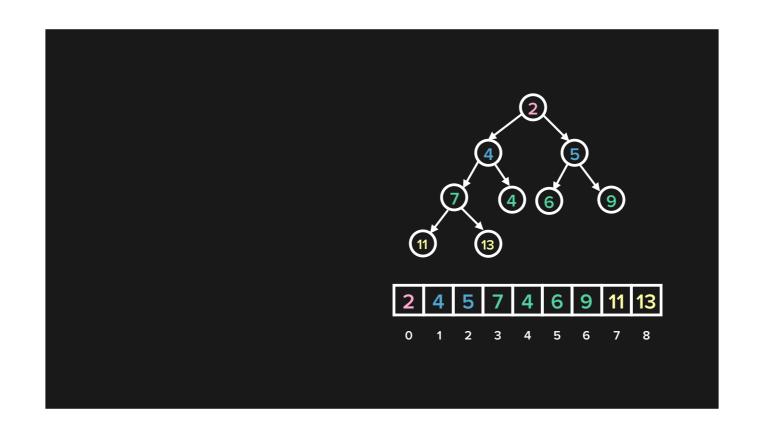
```
Start at last parent node

index = (count - 2) / 2

while index >= 0:

Sift down element at index

index -= 1
```



### **HEAPSORT**

```
heapify array
```

```
while (count > 0):
```

Grab min or max element (peek)

delete-min or delete-max element

# **HEAP RUNTIME**

Average Worst
Case Case

Space O(n) O(n)

Insert O(log n) O(log n)

Delete O(log n) O(log n)

### **HEAPSORT RUNTIME**

Average Worst
Case Case

Space O(n) O(n)

Heapify O(n log n) O(n log n)

Heapsort O(n log n) O(n log n)

### **RESOURCES**

http://www.cs.cmu.edu/~adamchik/15-121/lectures/Binary%20Heaps/heaps.html

http://en.wikipedia.org/wiki/Binary\_heap

http://en.wikipedia.org/wiki/Heap\_(data\_structure)

http://en.wikipedia.org/wiki/Priority\_queue