TECHNISCHE UNIVERSITÄT DRESDEN

Maik Thanh Nguyen

TITLE

CONFERENCE, DATE

Content

- Syntax for the basic modal language and semantics
  - Kripke frames
  - Topological space
  - Neighbourhood frames
- Multimodal logic and product of frames/spaces and logics
  - Multimodal logic and product of frames
  - Horizontal and Vertical topology/functions
  - Product of logics and the logic T
- Main result and ideas

Basic modal language, Kripke frames and models

- Basic modal langauge extends classical propositional logic. Formally:

Definition

Let Prop be a set of variable. Then a formula ϕ is defined as follows:

$$\phi ::= p \mid \bot \mid \neg\phi \mid \phi \vee \phi \mid \Box\phi$$

where $\Box$ is a modal operator and $p \in$ Prop

Basic modal language, Kripke frames and models

- Basic modal langauge extends classical propositional logic. Formally:

Definition

Let Prop be a set of variable. Then a formula ϕ is defined as follows:

$$\phi ::= p \mid \bot \mid \neg\phi \mid \phi \vee \phi \mid \Box\phi$$

where $\Box$ is a modal operator and $p \in \text{Prop}$

Definition

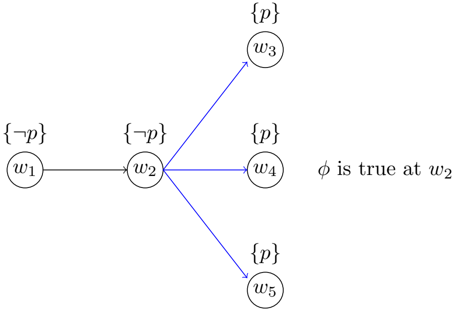
A frame $F = (W, R)$ is a pair where

- W is a non-empty set of worlds
- $R \subseteq W \times W$ is a binary relation

A model is a pair $M = (F, R)$ where V is a valuation and is of the form $V : Prop \rightarrow 2^W$

Example

- Let $\phi = \Box p$ and $M = (W, R, V)$ with $W = \{w_1, w_2, w_3, w_4, w_5\}$, $V(p) = \{w_3, w_4, w_5\}$ and $R =$



Kripke semantics

Definition

Let $M = (F, V)$ be a model and $w \in W$ a state in M. A formula being true at w is inductively defined as:

$$
\begin{array}{ll}
M, w \Vdash p & \text{iff } w \in V(p) \\
M, w \Vdash \bot & \text{never} \\
M, w \Vdash \neg\phi & \text{iff not } M, w \Vdash \phi \\
M, w \Vdash \phi \vee \psi & \text{iff } M, w \Vdash \phi \vee M, w \Vdash \psi \\
M, w \Vdash \Box\phi & \text{iff } \forall v \in W : wRv \to M, v \Vdash \phi \\
M, w \Vdash \Diamond\phi & \text{iff } \exists v \in W : wRv \wedge M, v \Vdash \phi
\end{array}
$$

Topological space

- deals with open sets, describes which points are "nearby"

Topological space

- deals with open sets, describes which points are "nearby"

Definition

A topological space is a pair (X, τ), where τ (called topology) is a collection of subsets of X (open sets) such that:

Topological space

- deals with open sets, describes which points are "nearby"

Definition

A topological space is a pair (X, τ), where τ (called topology) is a collection of subsets of X (open sets) such that:

- $\emptyset$ and X are open
- The union of arbitrary collection of open sets is open
- The intersection of finite collection of open sets is open

Topological space

- deals with open sets, describes which points are "nearby"

Definition

A topological space is a pair (X, τ), where τ (called topology) is a collection of subsets of X (open sets) such that:

- $\emptyset$ and X are open
- The union of arbitrary collection of open sets is open
- The intersection of finite collection of open sets is open

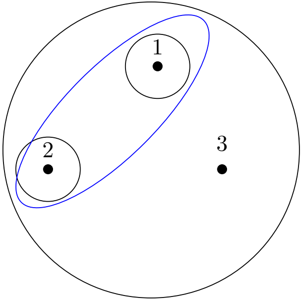
A topological model is a structure $M = (X, \tau, \upsilon)$ where (X, τ) is a topological space and υ a valuation of the form $\upsilon : \mathit{Prop} \rightarrow 2^X$

Example

- Let (X, τ) be a topological space with $X = \{1, 2, 3\}$, $\tau = \{\emptyset, \{1\}, \{2\}, \{1, 2\}, W\}$,$V(p) = \{1, 2\}$ and $\phi = \Box p$

Example

- Let (X, τ) be a topological space with $X = \{1, 2, 3\}$,
  $\tau = \{\emptyset, \{1\}, \{2\}, \{1, 2\}, W\}$,$V(p) = \{1, 2\}$ and $\phi = \Box p$



$M, 1 \Vdash \phi$
$M, 3 \nVdash \phi$

Topological semantics

Definition

Let $M = (X, \tau, \upsilon)$ be a topological model and $x \in X$ a point in M. A formula being true at x is inductively defined as:

$$
\begin{array}{ll}
M, x \vDash p & \text{iff } x \in v(p) \\
M, x \vDash \bot & \text{never} \\
M, x \vDash \neg\phi & \text{iff } M, x \nvDash \phi \\
M, x \vDash \phi \vee \psi & \text{iff } M, x \vDash \phi \text{ or } M, x \vDash \psi \\
M, x \vDash \Box\phi & \text{iff } \exists U \in \tau \text{ such that } x \in U \text{ and } \forall u \in U,\ M, u \vDash \phi \\
M, x \vDash \Diamond\phi & \text{iff } \forall U \in \tau : \text{if } x \in U \rightarrow \exists u \in U,\ M, u \vDash \phi
\end{array}
$$

Neighbourhood frames

- Generalize Kripke semantics
- Captures non-normal modal logics

Neighbourhood frames

- Generalize Kripke semantics
- Captures non-normal modal logics

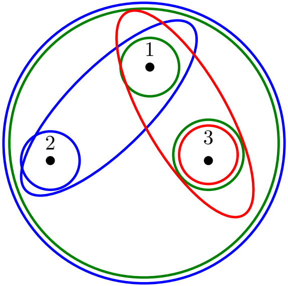
Definition

A neighbourhood frame is a pair (X, τ) where τ is a function $\tau : X \to 2^{2^X}$.
A neighbourhood model is a structure $M = (X, \tau, \upsilon)$, where υ is a valuation of the form $\upsilon : \mathit{Prop} \to 2^X$

Example

Assume $\phi = \Box p$. Let $W = \{1, 2, 3\}$, $V(p) = \{1, 2\}$ and

$$\tau(x) = \begin{cases} 1 \to \{\{1\}, \{3\}, W\} \\ 2 \to \{\{2\}, \{1, 2\}, W\} \\ 3 \to \{\{3\}, \{1, 3\}\} \end{cases}$$



True at all points

Neighbourhood semantics

Definition

Let $M = (X, \tau, \upsilon)$ be a neighbourhood model and $x \in X$ a point in M. A formula being true at x is inductively defined as:

$M, x \Vdash p$	iff $x \in V(p)$
$M, x \Vdash \bot$	never
$M, x \Vdash \neg\phi$	iff $M, x \nVdash \phi$
$M, x \Vdash \phi \vee \psi$	iff $M, x \vDash \phi \vee M, x \vDash \psi$
$M, x \Vdash \Box\phi$	iff $\exists V \in \tau(x) \forall y \in V : M, y \models \phi$

Modal logic and the logic T

Definition

A modal logic is a set of modal formulas containing all propositional tautologies, closed under Substitution ($\frac{\phi(p_i)}{\phi(\psi)}$), Modus Ponens ($\frac{\phi, \phi \rightarrow \psi}{\psi}$).

Modal logic and the logic T

Definition

A modal logic is a set of modal formulas containing all propositional tautologies, closed under Substitution ($\frac{\phi(p_i)}{\phi(\psi)}$), Modus Ponens ($\frac{\phi, \phi \rightarrow \psi}{\psi}$).

A modal logic is normal, if it contains $\Box(p \rightarrow q) \rightarrow (\Box p \rightarrow \Box q)$ (K) and is closed under Generalization ($\frac{\phi}{\Box \phi}$)

Definition

$T = K + \Box p \rightarrow p$

Multimodal logic and Fusion

- Simultaneously reasons about knowledge,time etc.
- Example: combines temporal and epistemic logic to reason "when does the agent know something"

Multimodal logic and Fusion

- Simultaneously reasons about knowledge,time etc.
- Example: combines temporal and epistemic logic to reason "when does the agent know something"

Definition

Let L_1 and L_2 be modal logics with one modality $\Box$. The fusion is defined as:

$$L_1 \otimes L_2 = K_2 + L_{1(\Box \to \Box_1)} L_{2(\Box \to \Box_2)}$$

Product of frames

- Reasoning on horizontal and vertical relations in one dimension

Product of frames

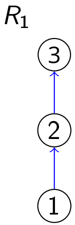
- Reasoning on horizontal and vertical relations in one dimension

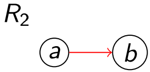
Definition

Let $F = (W, R_1)$ and $G = (V, R_2)$. We define the Kripke product on $W \times V$ as :

$$(w, v)R_1'(w', v') \text{ iff } wR_1 w' \text{ and } v = v' \text{ (horizontal)}$$
$$(w, v)R_2'(w', v') \text{ iff } w = w' \text{ and } vR_2 v' \text{(vertical)}$$

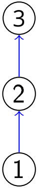
Example



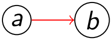


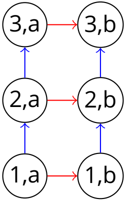
Example

R_1



R_2





Horizontal and Vertical topology

Definition

Let $\mathcal{X} = (X, \chi)$ and $\mathcal{Y} = (Y, \upsilon)$ be topological spaces and $N \subseteq X \times Y$

Horizontally open: N is horizontally open iff

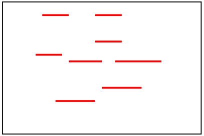
$$\forall (x, y) \in N \; \exists U \in \chi \text{ such that } x \in U \text{ and } U \times \{y\} \subseteq N.$$

Vertically open: N is vertically open iff

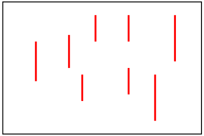
$$\forall (x, y) \in N \; \exists V \in \upsilon \text{ such that } y \in V \text{ and } \{x\} \times V \subseteq N.$$

τ_1 (horizontal topology) is the set of all horizontally open sets and τ_2 (vertical topology) the set of all vertically open sets.

Illustration, Standard product

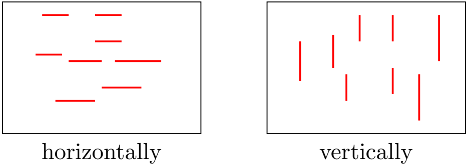


horizontally



vertically

Illustration, Standard product



- Reasoning in both directions simultaneously with the standard topology
- Basis of the standard topology is τ where

$$\tau = \{N \subseteq X \times Y \mid \exists U \in \chi \exists V \in \upsilon : N = U \times V\}$$

Horizontal, Vertical and Standard functions

Definition

Let $\mathcal{X} = (X, \tau_1)$ and $\mathcal{Y} = (Y, \tau_2)$ be two n-frames. We define the full product as

$$\mathcal{X} \times_n^+ \mathcal{Y} = (X \times Y, \tau_1', \tau_2', \tau) \text{ where}$$

$$\tau_1'(x, y) = \{U \subseteq \mathrm{X} \times \mathrm{Y} \mid \exists V \in \tau_1(x) : V \times \{y\} \subseteq U\}$$

$$\tau_2'(x, y) = \{U \subseteq \mathrm{X} \times \mathrm{Y} \mid \exists V \in \tau_2(y) : \{x\} \times V \subseteq U\}$$

$$\tau(x, y) = \{U \subseteq \mathrm{X} \times \mathrm{Y} \mid \exists W \in \tau_1(x) \, \exists V \in \tau_2(y) : W \times V \subseteq U\}$$

Product of logics

Definition

Let L_1 and L_2 be two unimodal logic. We define the full n-product of them as:

$$L_1 \times_n^+ L_2 = Log(\{\mathcal{X} \times_n^+ \mathcal{Y} \mid \mathcal{X} \Vdash L_1 \text{ and } \mathcal{Y} \Vdash L_2\})$$

where $Log(\mathcal{C}) = \{\phi \mid F \Vdash \phi \text{ for } F \in \mathcal{C}\}$

Main Research Question

Theorem

$$T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p = T \times_n^+ T$$

Main Research Question

Theorem

$$T \otimes T \otimes T + \Box p \to \Box_1 p \land \Box_2 p = T \times_n^+ T$$

- $T \otimes T \otimes T = K_3 + T_{\Box} + T_{(\Box \to \Box_1)} + T_{(\Box \to \Box_2)}$ is the logic with three modalities
- $\Box p \to \Box_1 p \land \Box_2 p$ is the interaction axiom

Main Research Question

Theorem

$$T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p = T \times_n^+ T$$

- $T \otimes T \otimes T = K_3 + T_{\Box} + T_{(\Box \to \Box_1)} + T_{(\Box \to \Box_2)}$ is the logic with three modalities
- $\Box p \to \Box_1 p \wedge \Box_2 p$ is the interaction axiom
- Sketch the inclusion from right to left
- Proof is splitted into two parts (ideas are from [1], [2])

Sketch of the main ideas

- Pick $\mathcal{C} = \{F \mid F \Vdash T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p\}$

Sketch of the main ideas

- Pick $\mathcal{C} = \{F \mid F \Vdash T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p\}$

Claim

$Log(\mathcal{C}) = T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$

Sketch of the main ideas

- Pick $\mathcal{C} = \{F \mid F \Vdash T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p\}$

Claim

$Log(\mathcal{C}) = T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$

- Shown by Sahlqvist Theorem

Sketch of the main ideas

- Pick $\mathcal{C} = \{F \mid F \Vdash T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p\}$

Claim

$Log(\mathcal{C}) = T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$

- Shown by Sahlqvist Theorem

Claim

$T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$ *has the finite model property*

Sketch of the main ideas

- Pick $\mathcal{C} = \{F \mid F \Vdash T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p\}$

Claim

$Log(\mathcal{C}) = T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$

- Shown by Sahlqvist Theorem

Claim

$T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$ *has the finite model property*

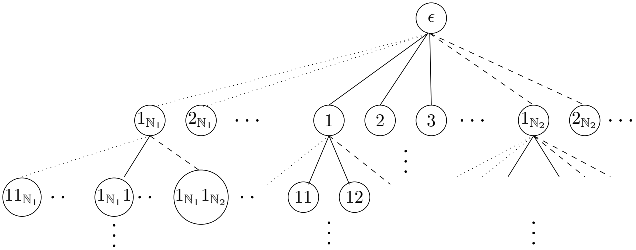
- Use Filtration Theorem

Sketch continue

- Construct an infinite branching and infinite depth tree with three reflexive relations ($T_{\omega,\omega,\omega[rn]}$)

Sketch continue

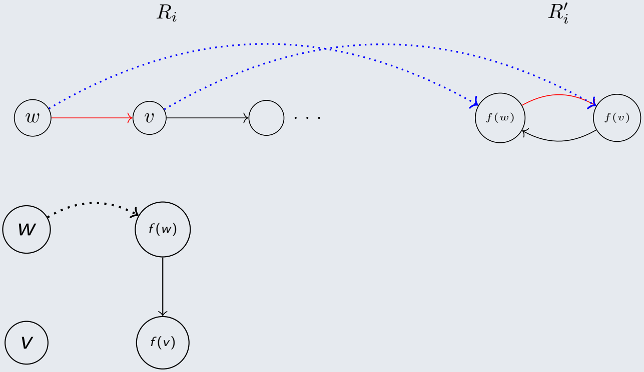
- Construct an infinite branching and infinite depth tree with three reflexive relations ($T_{\omega,\omega,\omega[rn]}$)



Sketch continue

Definition

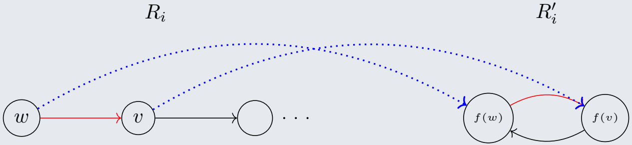
Let $F = (W, R_1, R_2, \ldots)$ and $F' = (W', R'_1, R'_2, \ldots)$ be two frames. A bounded morphism $f : W \to W'$ can be illustrated as:

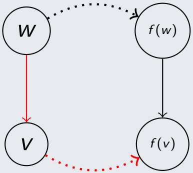


Sketch continue

Definition

Let $F = (W, R_1, R_2, ...)$ and $F' = (W', R'_1, R'_2, ...)$ be two frames. A bounded morphism $f : W \to W'$ can be illustrated as:



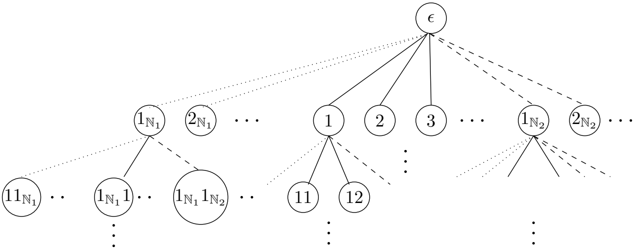


Sketch continue

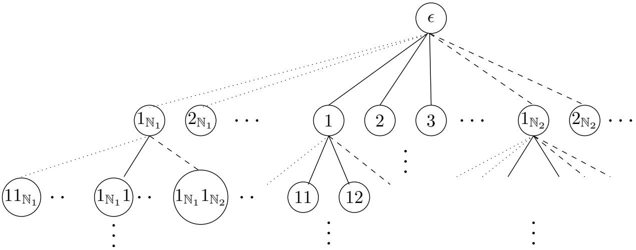
Claim

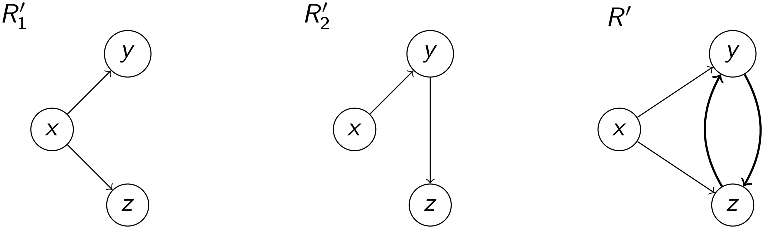
$Log(T_{\omega,\omega,\omega[rn]}) = T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$

Detailed sketch

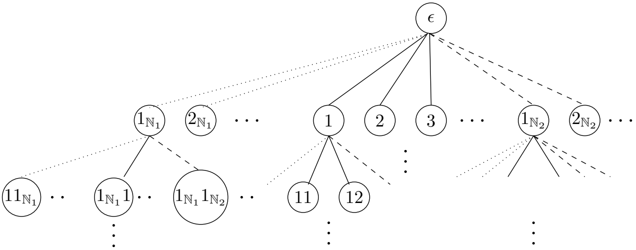


Detailed sketch

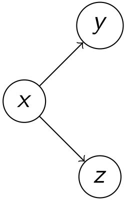




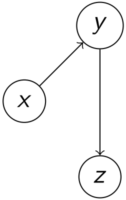
Detailed sketch



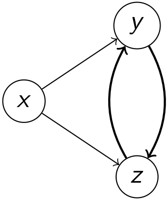
R'_1



R'_2



R'



Required components

$$Log(T_{\omega,\omega,\omega[rn]}) = T \otimes T \otimes T + \Box p \rightarrow \Box_1 p \land \Box_2 p$$

- Introduce $T_{\omega[rn]}$ a simpler version of $T_{\omega,\omega,\omega[rn]}$

Required components

$$Log(T_{\omega,\omega,\omega[rn]}) = T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$$

- Introduce $T_{\omega[rn]}$ a simpler version of $T_{\omega,\omega,\omega[rn]}$
- Construct $N_\omega(T_{\omega[rn]})$ the neighbourhood version of $T_{\omega[rn]}$
- Show $T = Log(N_\omega(T_{\omega[rn]}))$

Required components

$$Log(T_{\omega,\omega,\omega[rn]}) = T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$$

- Introduce $T_{\omega[rn]}$ a simpler version of $T_{\omega,\omega,\omega[rn]}$
- Construct $N_\omega(T_{\omega[rn]})$ the neighbourhood version of $T_{\omega[rn]}$
- Show $T = Log(N_\omega(T_{\omega[rn]}))$
- Construct a bounded morphism

$$N_\omega(T_{\omega[rn]}) \times_n^+ N_\omega(T_{\omega[rn]}) \to T_{\omega,\omega,\omega[rn]}$$

- With some further steps we can conclude

$$T \times_n^+ T \subseteq T \otimes T \otimes T + \Box p \to \Box_1 p \wedge \Box_2 p$$

Future works

- Many ways to continue the research
- Discover how it works for logic K
- Combine different logics for example $D \times_n^+ K$, $T \times_n^+ K, \ldots$
- For logic Λ with $T \subseteq \Lambda \subseteq S4$, does the following hold:

$$\Lambda \otimes \Lambda \otimes \Lambda + \Box p \to \Box_1 p \wedge \Box_2 p = \Lambda \times_n^+ \Lambda$$

Conclusion

References

[1] Johan van Benthem, Guram Bezhanishvili, Balder ten Cate, and Darko Sarenac.
Multimodal logics of products of topologies. *Studia Logica*, 84:369–392, 2006.

[2] Andrei Kudinov. Modal logic of some products of neighbourhood frames. In *Advances in Modal Logic, Volume 9*, pages 286–294, London, 2012. College Publications.