

# Assignment 1: Promela

## Exercise 1.1: Verifying the log program

The commented code is attached.

## Exercise 1.2: Analysis of algorithm on slide 35

### What does the code do?

After passing the initial assert, the two processes will take turns running the critical section. However if one of the processes is preempted between setting `flag[_pid] = 0` and setting `flag[_pid] = 1`, then the other will be able to continuously run the critical section.

### How does it work?

The programs uses three components:

- **flag:** Each process has a flag. The process sets the flag when ready to run the critical section.
- **turn:** Is a variable. The turn variable allows the first process hitting the guard to pass, when the other process requests a turn by setting the variable to its own `_pid`. Without this variable, the system would be susceptible to a deadlock, if both processes sets their flags before passing the guard.
- **Guard:** The guard allows only one process to enter the critical section at a time. When hitting the guard a process will wait for the other process to complete the critical section, or the other to request a turn.

### How do you verify that?

To verify that the processes takes turns, a variable, `oldturn`, is introduced, holding the `_pid` of the last process to run the critical section. During the critical process, before setting `oldturn`, it is asserted, that `oldturn` holds the `_pid` of the other process. This solution has some special cases:

- The `oldturn` variable will not have been set on the first run of the critical section.
- If one process has not passed the “again” label, the other process will be able to continuously run the critical section.
- If one of the processes is preempted between setting `flag[_pid] = 0` and setting `flag[_pid] = 1`, then the other will be able to continuously run the critical section.