

5.1

a) In zip.

b) .txt in zip, did not make table

c) No, because of overflow that can occur ~~for~~ and rounding error occurring in closed form.

d) did not make plot.



52.

a)

$a \rightarrow n, b \rightarrow n$

Each position (ones, tens, hundreds, etc.) is multiplied with all the positions of the number we are multiplying with.

For example,  $a = 123, b = 456$ .

$$\begin{array}{r} 123 \\ \times 456 \\ \hline \end{array}$$

we would multiply 6 with 3, 2, 1  
5 with 3, 2, 1  
4 with 3, 2, 1

Then we add the values.

The multiplication is  $O(n^2)$  and addition is  $O(n)$ .

$$\therefore O(n^2) + O(n)$$

Hence, the overall time complexity is  $O(n^2)$ .

b)

First, consider  $x = 12$ , which can be split into  $x = 1 \cdot 10^{1/2} + 2$ .

consider

$$z = 1234$$

$$z = 12 \cdot 10^{4/2} + 34$$

$$= 12 \cdot 10^{4/2} + 3 \cdot 10^{3/2} + 4$$

So, our general formula would be

$$A = a \cdot 10^{n/2} + b, \text{ where } b \text{ would be recursively solved.}$$

So,  $B =$

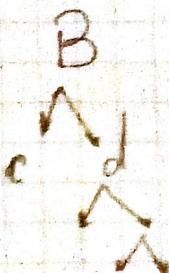
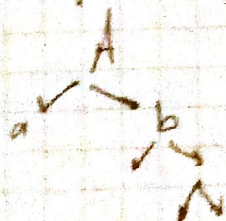
$$\text{let } B = c \cdot 10^{n/2} + d$$

$$A \times B = (a \cdot 10^{n/2} + b) (c \cdot 10^{n/2} + d)$$

$$= a \cdot c \cdot 10^n + ad \cdot 10^{n/2} + cb \cdot 10^{n/2} + b \cdot d$$

$$= (a \cdot c \cdot 10^n) + 10^{n/2} (a \cdot d + b \cdot c) + b \cdot d$$

Diagram.





Algor. thm:

multiply (A, B)

$n = \max(A, B)$

if ( $n == 1$ ):

return A\*B

else:

~~temp~~ temp = A

temp2 = B

$a = \text{temp} / (10^{n/2})$  // left side of A //

$b = \text{temp} \% (10^{n/2})$  // right side of A //

$c = \text{temp2} / (10^{n/2})$  // left side of B //

$d = \text{temp2} \% (10^{n/2})$  // right side of B //

~~W = multiply(a, c)~~

~~X = multiply(a, d)~~

~~Y = multiply(b, c)~~

~~Z = multiply(b, d)~~

~~Y = Y + multiply(a, d)~~

W = multiply(a, c)

X = multiply(b, d)

Y = multiply(a+b, c+d)

return (W \* (10^n) + ((10^(n/2)) \* (Y - W - X)) + X)

$$\begin{aligned} & ad + a(b+c)(b+d) \\ & ad + ab + ac + bc + bd \end{aligned}$$

$$\Rightarrow ac + bc = (a+b)(c+d) - ad - bd$$

c)

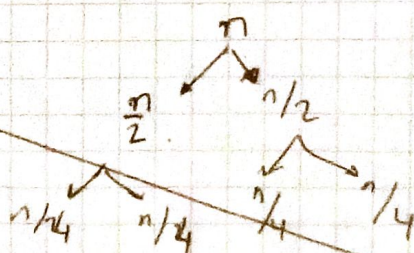
$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

divide  
(come from  
recursion)

conquer (comes from the  
addition)

⇒ 3 times because  
w, x, y.

d)



n  
n/2  
n/4

$$\begin{aligned} & \Rightarrow n + \frac{n}{2} + \frac{n}{4} + \frac{n}{8} \\ & \Rightarrow n \left( 1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} \right) \\ & \Rightarrow 2n \end{aligned}$$

$O(n)$



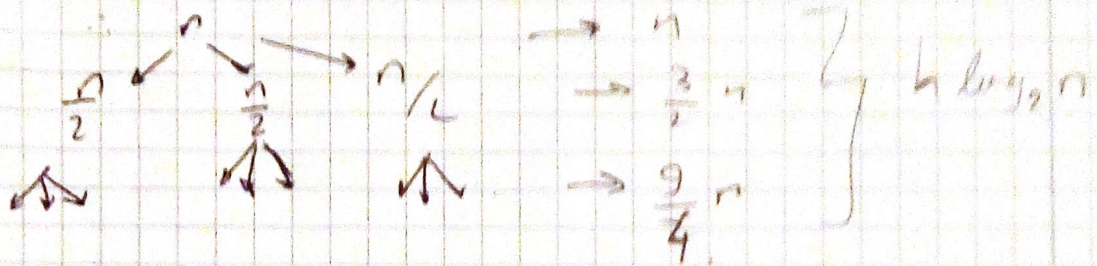
c)  $a = 3, b = 2$

$$n^{\log_3 3} = n^{1.58}$$

$$f(n) = n$$

$$= O(n^{1.58 - \epsilon}) \text{ for } \epsilon = 0.58$$

d)



$$\sum_{k=0}^h \left(\frac{3}{2}\right)^k \cdot n = n \sum_{k=0}^{\log_2 n} \left(\frac{3}{2}\right)^k$$

$$= n \left[ \frac{(1 - 3/2)^{\log_2 n}}{1 - 3/2} \right]$$

$$= n \left[ \frac{1 - n^{\log_2 3/2}}{-1/2} \right]$$

$$= 2n^{1.58} - 2n$$

$$\therefore O(n^{1.58})$$