

H/W 3:

1.

a) $f(n) = 9n$ and $g(n) = 5n^3$.

$f \in O(g)$.

For O we need to check.

• $f \in O(g)$

• $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$

$= \frac{9n}{5n^3}$

$= \frac{9}{5n^2}$

$= 0 < \infty$

$\therefore f \in O(g)$ is true

$f \in \omega(g)$
 $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \neq \infty$
 \therefore relation not true

• $f \in o(g)$
 $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$ $\therefore f \in o(g)$ is true

• $f \in \Omega(g)$

$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 < 0$ Not true

$f \in \Omega$
 $\therefore f \in \Omega(g)$ is not true

• $f \in \Theta(g)$
 $f(n) \in O(g) \cap \Omega(g)$ gives us a relation which is false.

• $g \in O(f)$

$\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} = \frac{5n^3}{9n}$
 $= \frac{5}{9}n^2$
 $= \infty$

\therefore relation not true

• $g \in \Omega(f)$ is true

$\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} = \infty > 0$

\therefore relation is true

• $g \in \Theta(f)$ is false

$\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} < \infty$

is false \therefore relation is false

• $g \in o(f)$

$\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} \neq 0$

\therefore relation is false

• $g \in \omega(f)$

$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$

\therefore relation is true

$$b). f(n) = 9n^{0.8} + 2n^{0.3} + 14 \log n.$$

$$g(n) = \sqrt{n}.$$

$$\begin{aligned} i). \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} &= \frac{9n^{0.8} + 2n^{0.3} + 14 \log n}{n^{1/2}} \\ &= \lim_{n \rightarrow \infty} \left(9 \frac{n^{0.8}}{n^{0.5}} + 2 \frac{n^{0.3}}{n^{0.5}} + 14 \frac{\log n}{n^{0.5}} \right) \\ &= \lim_{n \rightarrow \infty} \left(9n^{0.3} + \frac{2}{n^{0.2}} + \frac{14 \log n}{n^{0.5}} \right) \\ &= \infty. \end{aligned}$$

$$\begin{aligned} ii). \lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} &= \frac{n^{0.5}}{9n^{0.8} + 2n^{0.3} + 14 \log n} \\ &= \frac{n^{0.5}}{9n^{0.8} + 2n^{0.3} + 14 \log n} \\ &= \frac{1}{9 + \frac{2}{n^{0.5}} + \frac{14 \log n}{n^{0.8}}} \\ &= 0. \end{aligned}$$

- $f \in o(g)$ is false because $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \neq 0$.
- $f \in O(g)$ is false because $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \text{ is not } < \infty$.
- $f \in \Omega(g)$ is true because $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} > 0$.
- $f \in \Theta(g)$ is false because $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \text{ is not } < \infty$.
- $f \in \omega(g)$ is true because $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$.
- $g \in o(f)$ is true because $\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} = 0$.
- $g \in O(f)$ is true because $\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} < \infty$.
- $g \in \Omega(f)$ is false because $\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} \text{ is not } > 0$.
- $g \in \Theta(f)$ is false because $\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} \text{ is not } > 0$.
- $g \in \omega(f)$ is false because $\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} \neq \infty$.

c) $f(n) = \frac{n^2}{\log n}$, $g(n) = n \log n$.

i). $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{\frac{n^2}{\log n}}{n \log n} = \frac{n}{(\log n)^2}$

$$= \lim_{n \rightarrow \infty} \frac{n}{(\log n)^2}$$

$$= \lim_{n \rightarrow \infty} \frac{dn}{\frac{d(\log n)^2}{dn} \times \frac{d \log n}{dn}}$$

$$= \lim_{n \rightarrow \infty} \frac{1}{2 \log n \times \frac{1}{n \ln 10}}$$

$$= \lim_{n \rightarrow \infty} \frac{n \ln 10}{2 \log n}$$

$$= \frac{\ln 10}{2} \lim_{n \rightarrow \infty} \frac{1}{\frac{1}{n \ln 10}}$$

$$= \frac{(\ln 10)^2}{2} \lim_{n \rightarrow \infty} \frac{n \ln 10}{1}$$

$$= \infty$$

ii) $\lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} = 0$.

- $f \in O(g)$ is false as $c_1 \neq 0$.
- $f \in O(g)$ is false as c_1 is not $< \infty$.
- $f \in \Omega(g)$ is true as c_1 is > 0 .
- $f \in \Theta(g)$ is false as c_1 is not $< \infty$.
- $f \in \omega(g)$ is true as $c_1 = \infty$.
- $g \in O(f)$ is true as $c_2 = 0$.
- $g \in \Omega(f)$ is false as c_2 is not > 0 .
- $g \in \Theta(f)$ is false as c_2 is not > 0 .
- $g \in \omega(f)$ is false as $c_2 \neq \infty$.

$$d) f(n) = (\log(3n))^3 = (\ln(3n))^3, \quad g(n) = 3 \log n = \frac{2 \ln n}{\ln 10}$$

$$b) \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{1}{(\ln 10)^3} \times (\ln 3n)^3$$

$$= \frac{1}{(\ln 10)^3} \lim_{n \rightarrow \infty} \frac{g \ln(n)}{(\ln 10)}$$

$$= \frac{1}{g(\ln 10)^3} \lim_{n \rightarrow \infty} \frac{\frac{d(\ln 3n)^3}{d \ln 3n} \times \frac{d \ln 3n}{d 3n} \times \frac{d 3n}{dn}}{\frac{d \ln(n)}{dn}}$$

$$= \frac{1}{g(\ln 10)^3}$$

$$= \frac{1}{g(\ln 10)^3} \lim_{n \rightarrow \infty} \frac{3(\ln 3n)^2 \times \frac{1}{3n} \times 3}{\frac{1}{n}}$$

$$= \infty$$

$$u) g(n) = \lim_{n \rightarrow \infty} \frac{g(n)}{f(n)} = 0$$

- $f \in O(g)$ is false as $d_i \neq 0$.
- $f \in O(g)$ is false as d_i is not $< \infty$.
- $f \in \Omega(g)$ is true as $d_i > 0$.
- $f \in \Theta(g)$ is false as d_i is not $< \infty$.
- $f \in \omega(g)$ is false as $d_i \neq \infty$.
- $g \in O(f)$ is true as $d_{ii} = 0$.
- $g \in O(f)$ is true as $d_{ii} < \infty$.
- $g \in \Omega(f)$ is false as d_{ii} is not > 0 .
- $g \in \Theta(f)$ is false as d_{ii} is not > 0 .
- $g \in \omega(f)$ is false as $d_{ii} \neq \infty$.

2a) is in the file Sorting.cpp

b) For this algorithm to be correct, it needs to pass 3 conditions: Initialization, Maintenance, and Termination.

- Initialization: \rightarrow The loop invariant must be true before 1st execution of the loop, which for selection sort is that the array is sorted for first i elements. In our case, the sorted array has no values in it, so it can be said that it is sorted.
- Maintenance: \rightarrow If invariant is true before iteration, it should be true after iteration, which is true in our case as i elements would be sorted after each iteration.
- Termination: \rightarrow When the loop is terminated, the invariant should show something useful. At the end of the loop, all n elements would be sorted.

c) Case A and B are shown in sorting up.

d) increased the sample size to 500 samples and the best, worst & average cases were stored in its respective txt files.

e) looking at the code, it seems ~~$O(n)$~~ both best & worst case ~~is~~ the ~~max~~ the complexity was ~~the~~ $\frac{n(n+1)}{2}$.

∴ complexity was n^2 ; but in graph - there were some variants.