GTU Department of Computer Engineering CSE 222/505 - Spring 2022 Homework 4 Report

Atacan Başaran 200104004008

1. SYSTEM REQUIREMENTS

Question 1 -> Recursion method should find small string indice in bigger string according the occurrence

Question 2 -> Recursion method works as a binary search but it should not search an item, it should search for a range. Also after found the element it should look leftside and rightside

because there can be more element which in range

Question 3 -> Recursion method should traverse array and add them if totalsum equals to target that means it is a subarray. After found it, method should look until array ends.

2. CLASS DIAGRAM

<<utility>> Recursions

+ question1(smallerString : String, biggerString : String, wantedOccurence : int) : int

- question1Detailed(smallerString : String, biggerString : String, wantedOccurence : int, searchingPosition : int) : int

+ question2(sortedArray : int[], lowerBound : int, upperBound : int) : int

- question2Detailed(sortedArray : int[], lowerBound : int, upperBound : int, first : int, last : int) : int

+ question3(arr : int[], targetSum : int) : ArrayList<ArrayList<Integer>>

- question3Detailed(arr : int[], targetSum : int, tempSum : int, startingPosition : int, searchingPosition : int, subArrays : ArrayList<ArrayList<Integer>>) : ArrayList<ArrayList<Integer>>

<<utility>> DriverCode

+ main(args : String[]) : void

3. PROBLEM SOLUTION APPROACH

Question 1 -> I searched the small string in big string using indexof(String,int) which in String class.I searched by scrolling through the big string(searchingPosition) until the occurrence reached the 0 and when I found it I returned the index of it.

Question 2 -> I checked with binarysearch if the array contains elements in the range. If I find it, I also searched the left and right of the position of that element because there may be more elements in that range

Question 3-> I aimed to return the start and end positions of the subarrays. I create tempSum and add the elements to my tempSum by shifting the end position in the array. If it meets the subarray conditions, I add it in 2D Arraylist, if not, I try to find new subarrays by shifting the start position to the right.

Question 4 -> In time complexity analysis

4. TEST CASES

Testing each possibility of Question1 with 1 big string and 2 different small string

```
String bigString = "GTUdenemeGTUdenemeGTUdenemeGTUdeneme";
String smallString = "GTU";
String smallString2 = "deneme";
System.out.println("Big String: " + bigString);
System.out.println("Small string1: " + smallString);
System.out.println("Small string2: " + smallString2);
System.out.println("Search for smallstring1 in bigstring(occurrence 3): "+Recursions.question1(smallString, bigString,3));
System.out.println("Search for smallstring1 in bigstring(occurrence 50): "+Recursions.question1(smallString, bigString,50));
System.out.println("Search for Atacan in bigstring(occurrence 1): "+ Recursions.question1("Atacan", bigString,1));
System.out.println("Search for smallstring2 in bigstring(occurrence 1): "+Recursions.question1(smallString2, bigString,1));
System.out.println("Search for smallstring2 in bigstring(occurrence -5): "+Recursions.question1(smallString2, bigString,-5) + "\n\n\n");
```

Testing each possibility of Question2 with an array

Testing each possibility of Question2 with an array

5. RUNNING AND RESULTS

Q1)

```
Big String: GTUdenemeGTUdenemeGTUdenemeGTUdeneme
Small string1: GTU
Small string2: deneme
Search for smallstring1 in bigstring(occurrence 3): 18
Search for smallstring1 in bigstring(occurrence 50): -1
Search for Atacan in bigstring(occurrence 1): -1
Search for smallstring2 in bigstring(occurrence 1): 3
Search for smallstring2 in bigstring(occurrence -5): -1
```

Q2)

```
myarray for Question2: 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 Question 1 for myarray(Bound 4,10):7 Question 1 for myarray(Bound 15,25):6 Question 1 for myarray(Bound -5,0):0 Question 1 for myarray(Bound 25,50):0 Question 1 for empty array:0
```

Q3)

```
myarray for Question3: 1 5 2 9 3 4 3 2 5
Question 3 for myarray(Target 7): [[1, 2], [4, 5], [5, 6], [7, 8]]
Question 3 for myarray(Target 0): []
Question 3 for myarray(Target 5): [[1, 1], [6, 7], [8, 8]]
Question 3 for empty array(Target 5): []
```

6. TIME COMPLEXITY ANALYSIS

Question 1

```
T(n) = T(n-1) + O(l*m)
n = occurrence
O(l*m) comes from indeox(String,int) method
l = small string size
m = big string size
```

Mathematical Induction

n = 0 -> (True) if indexfof method founds small string return index otherwise return -1 (Base case)

Let assume n<= k is True

n = k+1 -> (True) Method found k. occurrence and start searching again after new searchingPosition(k. occurrence + length of small string). We reach base case. if indexof method founds small string return index otherwise return -1

```
T(n) = O(n*l*m)
```

Question 2

```
private static int question2Detailed(int sortedArray[], int lowerBound, int upperBound, int first, int last) {
   if (first > last) // It is for unsuccessful search
        return 0;

   int middle = (first + last) / 2;

   if (sortedArray[middle] >= lowerBound && sortedArray[middle] <= upperBound) {
        // I include lower and upperbound to my output counter
        int found = 1;
        found += question2Detailed(sortedArray, lowerBound, upperBound, first, middle - 1);
        // It found the element inside the bounds but need to look leftside and rightside for all elements which inside the bounds
        found += question2Detailed(sortedArray, lowerBound, upperBound, middle + 1, last);

        return found;

   else if (sortedArray[middle] < lowerBound)
        return question2Detailed(sortedArray, lowerBound, upperBound, middle + 1, last);
        // Element smaller than we are looking for so search in rightside

else
        return question2Detailed(sortedArray, lowerBound, upperBound, first, middle - 1);
        // Element bigger than we are looking for so search in rightside
}
</pre>
```

```
T(n) = T(n/2) + \theta(1)

T(0) = 1

n = number of array elements

T(n) = T(n/2^2) + 2^* \theta(1)

T(n) = T(n/2^k) + 2^k^* \theta(1)

T(0) = 1, k = logn

T(n) = O(logn)
```

Question 3

```
private static Arraylist<Arraylist<Integer>> question3Detailed(int arr[], int targetSum, int tempSum,
        int startingPosition,
        int searchingPosition,
       ArrayList<ArrayList<Integer>> subArrays) {
    if (searchingPosition >= arr.length) // If searching position equals array size it means finding subarray is finished
       return subArrays;
    tempSum += arr[searchingPosition];
    if (tempSum == targetSum) { // It means it found a subarray add it into subArrays and continue to finding new subarray
       ArrayList<Integer> subArraysElement = new ArrayList<>();
       subArraysElement.add(startingPosition);
       subArraysElement.add(searchingPosition);
       subArrays.add(subArraysElement);
       question3Detailed(arr, targetSum, 0, startingPosition + 1, startingPosition + 1, subArrays);
    } else if (tempSum < targetSum) // If it could not find subarray yet continue to finding subarray
       question3Detailed(arr, targetSum, tempSum, startingPosition, searchingPosition + 1, subArrays);
    else if (tempSum > targetSum) // It means it could not be a subarray because exceeded the sum, find a new subarray
       question3Detailed(arr, targetSum, 0, startingPosition + 1, startingPosition + 1, subArrays);
    return subArrays;
```

```
T(n) = T(n-1) + O(n)

T(0) = 1

n = number of array elements

T(n) = T(n-2) + 2*O(n)

T(n) = T(n-k) + k*O(n)

T(0) = 1, k=n

T(n) = O(n)*n + 1

T(n) = O(n^2)
```

Question 4

foo method splits the integers on half and calling byself 3 times until splits to 1 digit and returning a value which represent with formula

```
\begin{split} T(n) &= 3^*T(n/2) + \theta(n) \\ \theta(n) \text{ for split\_integer} \\ n &= \text{digit number of element} \\ T(n) &= 3^k * T(n/2^k) + 3^k * \theta(n) \\ T(1) &= 1 \text{ , } k = logn \text{ (I couldn't type logn in base 2 in pdf editor but logn means in base 2 logn in below )} \\ T(n) &= n^log 3 + n^log 3 * \theta(n) \\ T(n) &= \theta(n^2) \end{split}
```

Mathematical Induction

n= 1 -> (True) return integer1*integer2 (base case)

Let assume n<=k is true

n = k+1 -> (True) Program did all n<=k foo method(each n calls 3 foo) in n=k+1 program doing extra mathematical equation (constant time) so n=k+1 is true