

Q1 Introduction:

(1) Data refers to raw, unprocessed set of facts and figures that we collect.

② Information

→ Information answers questions like who, when, where & so forth.

(3) Knowledge:

Information can be synthesized based on business choices to obtain actionable insight. This is called knowledge.

Knowledge (synthesizing and decision making) e.g., Top performers

Information (processing data) e.g., A book

Data (collecting & organizing) (90 marks)

* Some features of the application in terms of data:

- Data Integrity: Data should be accurate, consistent and reliable.

(social media profile should contain a valid country name)

- Data Retrieval: Data should be available and usable.

(social media user should be able to access their data all the time)

- Data Security: Data should be protected from unauthorized access.

- Data Independence: A particular functionality should not depend on any specific application.

- Concurrent Access: Application should support multiple users and processes simultaneously without violating data Integrity.

* Flat Files :

Excel files
Word files
JSON
XML
etc.

Traditionally data ~~was~~ is stored in flat files

Date: / /

- * Can be accessed using a programming language

Limitations:

Data Redundancy

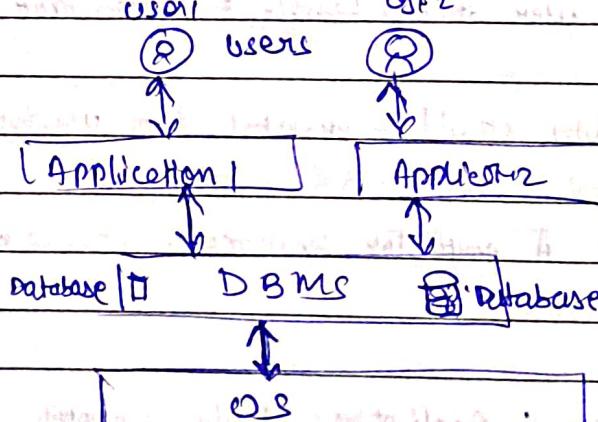
Loss of
data

Dependency of program

Scalability
challenges.

- * Database: Structured and organized collection of logically related data and description of raw data; description of these data, designed to manage, retrieve and update information as per the needs.

- * DBMS: Software Application for controlled access to the database.



Doubts \rightarrow Reliable
Consistent
(Integrity)

Date: / /

* which come first?

* Types of Database Systems:

(Attributes - values)

* Relational Model : stored data (Property)

in Relations i.e. Tables

* Attribute \rightarrow Column name
Employee (id, name, salary, bonus)

* Rows / Records / Tuples \rightarrow rows \rightarrow data

* Cardinality \rightarrow No. of rows.

* Degree \rightarrow No. of columns.

Represents a value of an attribute that is absent/unknown.

* NULL \rightarrow Null does not have any DataType



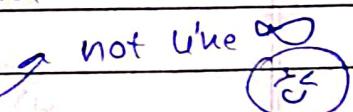
"Absence of Data" (Unknown)

* domain

Bonus can have values
 ≥ 0 but

↓ "Set of all possible values that an attribute can hold"

* Pure mathematics can't be applied



Only largest possible value of that data type,

example: Domain of Dept $\rightarrow \{ MCP, EEA, IUS \}$

Only

* Quiz :

Q1. Relational Database is a collection of:

Tables/ relations.

Q2. A relationship among a set of values in a table is represented by:

Row

represents the collection of interrelated data.

(Like the row represents the whole related data of
Employee 2 so that data is related to
Employee 2)

Q3. The set of permitted values
 ↓ Domain

Q4. Null (Unknown value)

Q5. Cardinality - 4

Q6. Degree - 3

↳ If customer can have multiple accounts
 and multiple shares.

Then

2	1173	}
2	1174	
3	1173	}

Both repeatly

but combination is unique

So here AcntNo, AcntNbr

both make the composite P.K.

* Data Integrity and Constraints

↳ Refers to accuracy, consistency, and reliability of data throughout its life cycle.

* Entity (Activity)
↳ one Tuple
represent
not
only.

* Can be ensured through CONSTRAINTS

which are used to restrict the data that can be entered.

→ Table on Next

Date: / /

Achieved through
Entity Integrity.

→ { } written in curly braces

* Candidate Key: → identified during the design of Database only

↳ based on Business logic and needs.

↳ "Minimal set of Attributes (columns) that can uniquely identify a Tuple (row) in a relation (table)."

Example: • EmployeeNo - ✓

• AddressNo - X

• Name, DateOfBirth - ✓ (for small group)

• Salary - X

• EmployeeNo, DateOfBirth - X because EmployeeNo itself is unique (see definition)

it can be a PRIMARY KEY

Example: Accounts in Bank AccountDetails(AccNo, CustNo, name, type)

1) If one customer can have only

one account and account cannot be shared?

Both {AccountNo} and {CustomerNo} possible CKs

2) If a customer can have multiple accounts

→ In this case {AccountNo}

only is the possible CK.

3) If each can only have 1 acc. Repeating Customer No
but can be Shared.

AccNo custNo

1173 2

Here accountNo will repeat.

← 1173 3

1173 4

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→ Here only {AccountNo} can be a possible CK.

* PRIMARY KEY (UNIQUE, NOT NULL)

↓
"chosen specific candidate key" to uniquely identify a row.

Primary Key vs Candidate Key

* Cannot have null

* should have small/short size

* Composite Primary Key.

* Ensures Integrity

PR. → {EmployeeNo} by preventing duplication of data, it and

* Reference for relationships with other table

* UNIQUE KEY, can allow null values.

Set of one or more columns whose values is required to be unique

* foreign Key : → (can contain null values)

(Next Page
more detail)

In UNIQUE KEY

* Primary key of another (Parent) used in this table

table for their references.

(Referencing) $\xleftarrow{\text{Relationship}}$ Referenced
Child

* Double Classification : Both Candidate Key and Primary Key can be used for data integrity

by preventing
Duplication of data



1) 1 Table → only 1 P.K.

2) Candidate keys can be multiple

Employee

Date: / /

EmpNo	Name	Address	Salary	DOB
1	JKS	12345	2K	-
2	RKS	34567	2K	-
3	ABC	45628	20K	-
4	Xyz	56781	30K	-
5	Ashley	78190	35K	-
6	Anuradha	19235	35K	-
7	Dawood	23456	3K	-
8	KET	34901	3K	-

Tables

Candidate - Min set of columns

PK

UNIQUE NULL X

Candidate Key

{EmpNo}

{Address}

{Name, DOB} {Name, DOB}

Primary key based on

K C Manning

All CN \rightarrow PKAU PK \rightarrow PK

consistent over a period of time

Size should be small

attributes of Primary Key

EMPNO

Mandatory

Desirable

- Must uniquely identify a tuple.
- Should not chose with tiny size.
- Should have short size.
ex: numbers.

Composite Primary Key

When two or more columns together identify the unique row then it is referred as Composite Primary Key.

It is in Employee table

If we have chosen

(Name, DOB) as Primary Key

then it is Composite.

Candidate Keys me

Se hon-si

PRIMARY KEY

leni hai \rightarrow Depend on RAJDHANI
Aphne PC kartta.

* Qn 3 - Data Integrity

Employee

Date: / /

EmpNo	Name	Aadhar	Salary	DOB
1	J.S.	12345	8k	-
2	R.K.S.	34567	12k	2k
3	A.B.C.	45678	20k	-
4	X.Y.Z	56781	20k	-
5	Ashley	78190	35k	-
6	Anindita	19235	35k	1990
7	Ranvir	23456	3k	-
8	K.E.F.	345901	3k	1990

Candidate - Min set of columns

PK

UNIQUE

NULL X

Candidate Key

{EmpNo}

{Aadhar}

{Name, DOB}

consistent over a period of time

Size should be small

Primary key based on

* Criteria

All C.R. \rightarrow PK

AU PK \rightarrow PK

EMPNO attributes of Primary Key

Mandatory

Desirable

- Must uniquely identify a tuple.
- Should not choose with tiny

- NULL X
- Should have short size.

e.g. numbers.

* Composite Primary Key

When two or more columns together identify the unique row then it is referred as Composite Primary Key.

If we have chosen

(Name, DOB) as Primary Key

then it is Composite.

Composite Keys me

se kon si

PRIMARY KEY

leni hai \rightarrow

Depend ~~for~~ RAJENDRA
Apne PC karta.

Data Integrity and Constraints

ensured
through constraints.

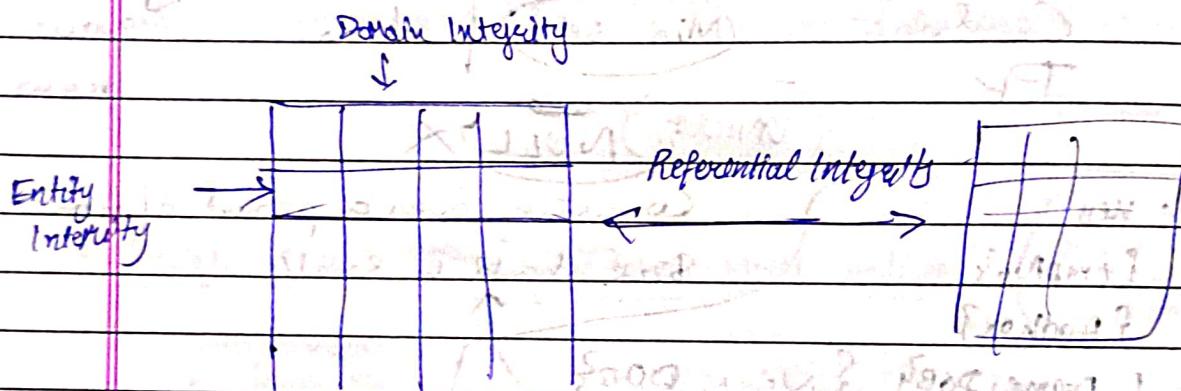
Date: / /

Integrity Type. Definition Enforced through

<u>Entity Integrity</u>	Row/Tuple (must be unique in itself) (prevents duplication)	PRIMARY KEY
-------------------------	---	-------------

<u>DOMAIN INTEGRITY</u>	column/attributes	Data Types & CHECK CONSTRAINT
-------------------------	-------------------	----------------------------------

<u>REFERENTIAL INTEGRITY</u>	How relations (or more)	FOREIGN KEY
----------------------------------	----------------------------	-------------



X FOREIGN KEY

UNIQUE key

When one table uses the primary key of other table
for reference.

- That column in child table (Reference table)
is called its FOREIGN KEY.

- FOREIGN key in the child table can contain NULL values.
- Self-Referencing foreign key
A foreign key of a table can reference its own table's PRIMARY KEY

CONNECT Problem

We can change the PRIMARY KEY/UNIQUE column
being referenced in the Parent table without any problem.

Can use only those computers that are available in the computer table.

uses A

Date: / /

Employee Table

id	ename	dept	compt
1001			
			Relationship
1002			

Computer Table

compid	make	model
1001		
1002		
1003		

child/Referencing Table

Parent/Master/Referenced Table

- Can be NULL
- Can be duplicates

* numerous rows in the child table can refer to the same row of the parent table, depending upon the type of the relationship.

Quiz - Data Integrity

- Relation b/w two tables → Foreign key.
- only 3 P.K. and although not desired but strng/char can be P.K.s.
- custId
RechargeId
(ध्यान से पढ़ो → Niche value table ka pata)
- Proprietor, Plan → Unique and Not Null

- FK → Niche

* Exercise - Data Integrity and Constraints

Table 1 : Trainer (TrainerId, TName, EmailId, ContactNo)

CRs: { TrainerId }, { EmailId }, { ContactNo } Date: / /

PK: { TrainerId }

FK: $\rightarrow \{ \}$, No Foreign Key

here it is null

Table 2 : Member (MemberId, Name)

CRs: { MemberId }

PK: { MemberId }

FK: $\rightarrow \{ \} \rightarrow \text{NO}$

Table 3 : WorkoutType (WorkoutTypeId, Level)

CRs: { WorkoutId }

PK: { WorkoutId }

FK: $\rightarrow \{ \}$

Table 4 : WorkoutSessions (TrainerId, WorkoutTypeId, MemberId, Tm)

CRs: { TrainerId, WorkoutTypeId, MemberId }

PK: { TrainerId, Same }

FK: { TrainerId }

{ WorkoutTypeId }

{ MemberId }

(Read the question carefully and get insight on what can be multiple times)

* Exercise #2: Keys (Bookstore)

• Book (BookId, ISBN, Status)

CR: { BookId } → multiply

PK: { BookId }

FK: { ISBN }

• Author (AuthorId, AuthorName)

- CR: { AuthorId }

- PK: { AuthorId }

- FKs: $\rightarrow \{ \}$

• BookDetails (ISBN, BookTitle, AuthorId)

- CR: { ISBN }

- PK: { ISBN }

- FKs: { BookTitle }

Ex partium

Book type Ni

book no rohi

- FK: { AuthorId }

Author
Author
Author
Author

- FK: { AuthorId }

key rohi bcoz can be multiple book with same title

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* Relationship b/w entities Date: / /
 * 1 Employee can have only 1 manager
 Cardinality 1 to 1
 Relationship 1 to 1

Junction Table

When 2 or more tables are connected with the help of the 3rd table

structure

N.V. Imp

* Relationships: associations of one entity with another entity through a Key

to link related data for efficient querying, and data integrity.

Each relationship has a name,

e.g. A computer is allocated to an employee.

* Mantra

Yaha homestha

* there can be more than 1 relationships b/w entities.

2-way relationship

e.g.

An employee works in a department,

An employee also manages the department

	Employee	works in	Department
What will be	Name		DeptId
be	Salary	manages	DeptName
xcept			Location

* A relationship can also exist between instances of the same entity

e.g. An employee reports to a manager

Employee	reports to
ID	
Name	
Salary	

Cardinality of Relationships:

Cardinality of relationship is the number of related records or entities that can be associated with each other record or entity in the other table.

* Employee and Computer:

- 1) How many computers can be allocated to an Employee?
- 2) Can computers be shared among employees.
- 3) Can employees exist without being allocated a computer.

If 0 or 1 computer can be allocated, 0 or 1 employee then the cardinality of relationship b/w these two entities is

1:1

Mantra

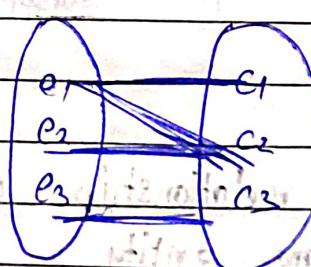
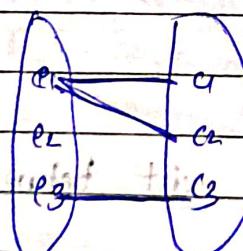
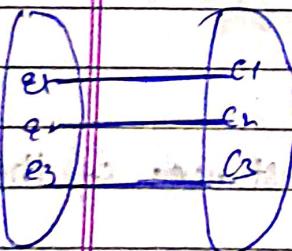
Hamesha ek side ka ek entity (now, ex computer) with or one relation durre ke elements se

Cardinality

One-to-one

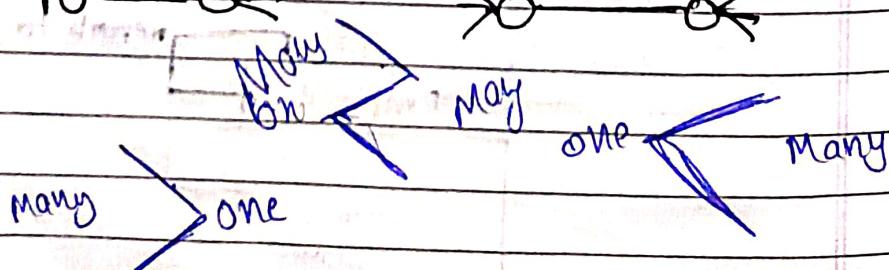
One-to-many

Many-to-many



1 to 1 or 1 to N

M:N



1
2
2

1 1 1

1 1 1

1 1 1

1 1 1

Date: / /

* Criss-Cross Notation is a visual representation used to depict the cardinality of relationship in an ER model.

* The notation comprises of 4 symbols and 3 of them needs to be used for representing each entity in a relationship.

1 1

Exactly one

1 0

zero or one

> 0

zero or more

X

one or more

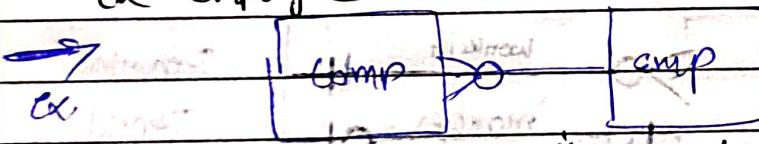
* Computer can be allocated to one and only one employee.

[comp] 1 [emp]

Ex computer staff. Ex employee no assign his staff.

* Employee can have zero or more computers.

Ex employee not both computers to staff ne.



Ex computer of

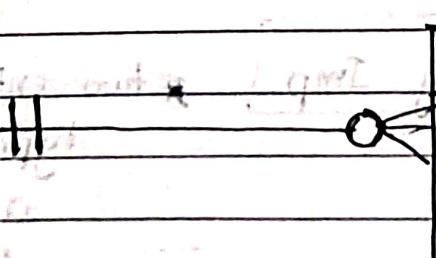
if ex employee
he 2 computer
to visit to
logistics

use no assign his

far use both no

nahi koo suda

Employee	
name	id
name	name
name	name
name	name



Computer

compld

Model

Model

V.V. Imp

- * Foreign Keys must be created in tables Date: 1/1
In order to establish the relationship b/w the entities.

Example:

Same as $O \leftarrow \rightarrow 1$

Computer	$O \leftarrow$	$\rightarrow 1$	Employee
compid			Id
Make			Name
Model			Salary

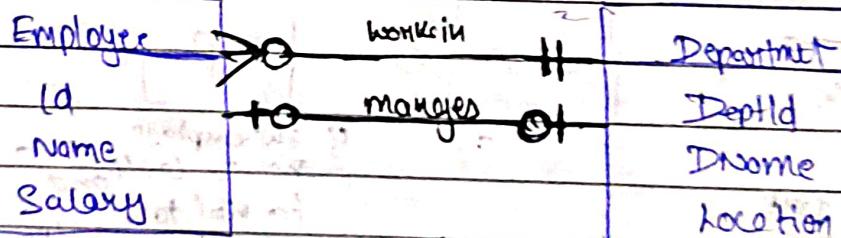
- Employee can have 0 or 1 computer
- Computer can be allocated to 0 or 1 Employee

FOREIGN KEYS

compid	Make	Model	Id	Name	salary	compid
1001	Dell	Vostro	1	Jones	72000	1001
102	Dell	Precision	2	Ethan	35000	NULL

Example:

* there is always 1 manager



Very - Very Imp

- * two entities can have various types of relationships and also different cardinalities for all those relations.

ER

- * Employee can work in exactly one department.

- * ER department can have one or more employees.

* An Employee can manage 0 or 1 Dept.

* except can be managed by 0 or 1 Emp.

* Foreign Key.

• here DeptId will go to Employee table as min

many-to-one
relation.

for works in relationship)

Employee			DeptId	Department		
id	ename	salony		dname	locatia	managid
1	James	72000	10	ETA	Mysore	1
2	Ethan	85000	20	TUS	Bangalore	NULL
3	Emily	25000	10			

* A (Ex)
department
can have many
Employees

(Primary keys of tables only)

→ 2 Relationships → matlab → 2 foreign keys

(2 tables me)

which one → where to place

V.V
Imp

- depends on →
- one-to-one (on one side)
 - one-to-many (on many side)
 - many-to-many. (separate table containing both table's primary keys)

if you want to think why
just take any example and

think / this will start

(repeatedly in the primary key)

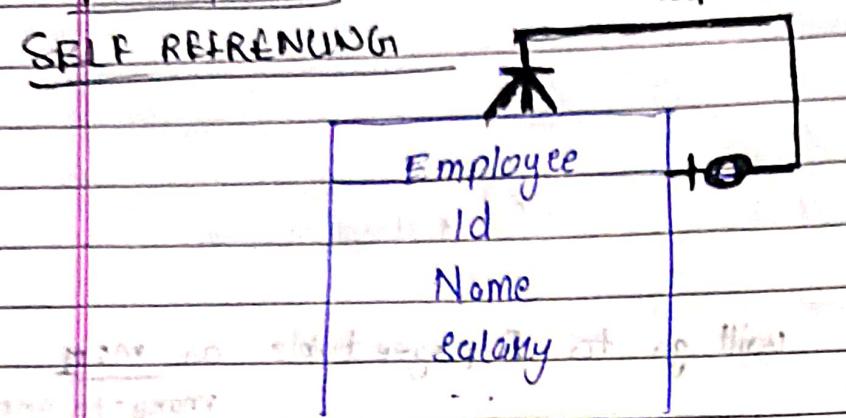
if you do like of
what said here!!

(and it's primary
key is their
combination.)

* Tricky Example

Date: / /

Example 3 SELF REFERENCE



* Every employee reports to 0 or 1 manager.

* Every Manager(employee) can have many Reportees (Employees)

* Foreign Key in same Table only

Id	Ename	Salary	*ManagerId	→ Foreign Key
1	James	72000	1	
2	Ethan	85000	NULL	

* The relation in which the foreign key will be created depends on Cardinality of the relationship

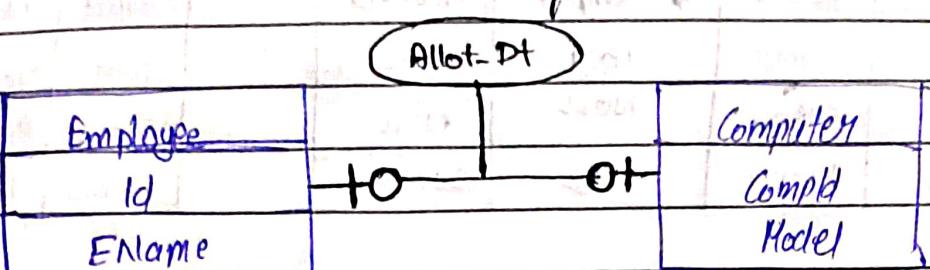
* one-to-one (1:1)

One is associated with one or one.
one or one with another is one ~~is~~

Date: / /

example:

Relationship Attribute



* in case of one-to-one (1:1)

- Put the other's table Primary Key in one Allot-DT to other other table as its Foreign Key (either table).

(ii) * Allot-DT - date on which employee was allocated a computer

↓
It is neither property of employee table nor computer table.

↓
It belongs to the relationship.

so, represented differently.

* This relationship can be mapped to Relations in either way:

1) use Compld attribute of Computer table as Foreign key in the emp table.

(Allot-DT along with the foreign key only)

(in the table where the foreign key is there)

2) Add Primary key of emp table EmpId as the (1:1) Foreign key in the Computer table with Allot-DT.

* Employee can have multiple computers.

* Computer can be given to 0 or 1 employee.

In 1:N relationship, the foreign key and relationship attributes are always added to (N) many side of the relationship.

the these attributes (Primary Key of EmpId as foreign key in Computer Table) and the Alt-Id also in Computer table along with the FOREIGN key.

The relationship is represented by a completely new table that has a composite Primary Key

↓

Primary of 1 , Primary of 2 }

That new table / design requires

2 - Foreign Keys on the new table linking the primary keys of each of the table.

Attribute A lot table goes with Foreign Key also

So it also go into new table.

Employee

computer

Id	Enam	CmpId	Model	Id	CmpId	Alt
1	Jans	1001	Lestra	1	1001	2014 -
2	Etha	1002	:	3	1002	:
3	Emij	1003	,	1	1003	
				3	1004	

V. Imp
try to read this

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Summary:

Entities, Relationships

* join relationship

* using foreign key

* there can be multiple relationship between two tables.

Types of Relationships

One-to-one relationship and One-to-many relationship

One-to-many relationship

one-to-many relationship exist in two ways

left with left join we have to join another table

right with right join it has more columns

* Introduction



Types of Database Systems:

Poubit

single 1 and

no 0 → ?

→ Hierarchical (Trees)



- each node can have only 1 parent

- but multiple children (each can have)

→ Network (Graph)



- each Node Connected to various

- Multiple parents and Children for each Node.

→ Relational → one or more tables

* Tables are related to each other through parent child relationship

, NOSQL database uses key-value pairs,

(Not only SQL) graph, or document data structures to store data.

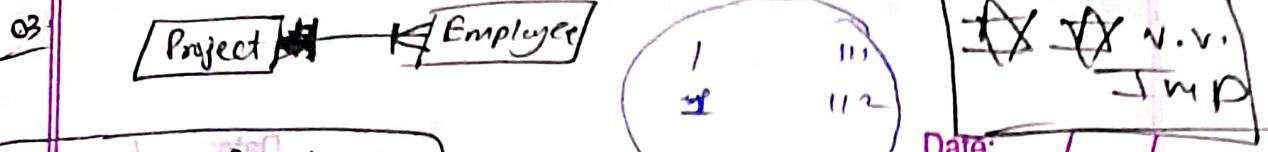
⇒ These databases aim for Simplicity of design

→ Horizontal scaling

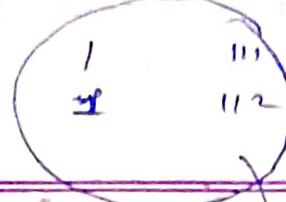
→ finer control over availability

Additional → OOP databases (DB4O)
(CopeDB)

Graph databases (Neo4J)



PRIMARY K for decision

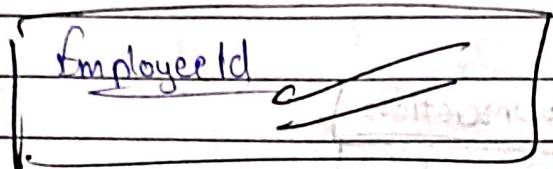


✓ ✓ v.v.
JMP

Date: / /

★ (one-to-many) (1:N) —> [Koi ek]

PRIMARY Key:



Many side
voali

Not Null (at least one)

(unique)

Q4

Project —> Employee [Many to Many]

Project Employee —> [Both composite]

1 1/18 Monday

11

1/2

One-to-one

Q5

Participant

Koi bhi le lo.
dono me se
' kisi ki bhi'

Affiliation

Course

Normalization:

#1 Introduction: (Retail Shop example for data Redundancy,

* See the case-study Inconsistencies and Anomalies

Summary * While designing database, relationship b/w different attributes and relation entities should be examined, to avoid

* Functional Dependency:



"relation b/w two attributes,"

- Inconsistencies
- Redundancy
- Anomalies.

→ whether one attribute can be uniquely determined from other attribute"

ex: Roll No.

must have only
one value
associated with
it in Name

In a relation R, two attributes/
Set of Attributes A and B,

B is functionally dependent on

A $(A \rightarrow B)$ if each value of A is associated with
exactly 1 value in B"

* (Black Copy: Last \Rightarrow) (Vishwadeep Groth)

(Determinant) (Dependent)

A \rightarrow B

For the Relation

retailoutletstock (retailoutletId, itemcode,
description, retailOutletLocation (street, city, zip),
email, retailUnitPrice, itemclass)

Several FD's are listed below:

{retailoutletId, itemcode} \rightarrow qtyAvailable
 {retailoutletId, itemcode} \rightarrow retailUnitPrice
 {retailoutletId} \rightarrow retailOutletLocation
retailUnitPrice \rightarrow itemclass
itemcode \rightarrow description

Closure of
{retailoutletId, itemcode}
gives all attributes
of the class.

{retailoutletId, itemcode}
is the Candidate Key

mention

(null cannot be inserted)

Se hui.

Date: / /

- Attributes part of the candidate key are called Prime attributes.
- Attributes that are not one Non-Prime attributes.

* Full Functional Dependency:

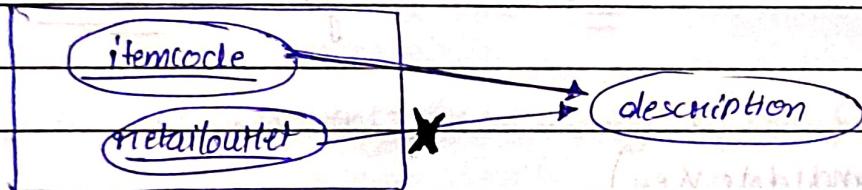
occurs when the dependent attribute is fully dependent on the whole Candidate Key and not a part of it.

"Any subset of Candidate Key should not determine the dependent".

* Partial Functional Dependency:

In the ~~retailoutlets~~ stock → description is unique for every ItemId, there is no need of retailoutletId

∴ description is Partially dependent on ~~retailoutletId, itemcode~~
attribute



"Any dependent attribute can be determined by a part of CK".

* Transitive Functional Dependency:

V.V-IMP occurs when any dependent is indirectly determined by the Candidate Key.

through non-key attributes

$$\text{CK} \xrightarrow{\text{non-key attr}} A \rightarrow B \rightarrow C$$

$$\text{CK} \xrightarrow{\text{non-key attr}} A \rightarrow C$$



Goto last
Page try
Anomalies topic

- ★ Full, Partial, Transitive Dependence Puchha toh?
- ↳ First find the Candidate Key

↳ if part of candidate key can determine any attribute

then Partial

Transitivity

if full is required → Full

if any type of transitivity is given then b/w CK and FDS

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$R(A, B, C, D, E, F)$

$(AB) \rightarrow CDEF$

$A \rightarrow CDF$

$B \rightarrow DEF$

Pehle Candidate Key
Nikalo

11.

Full, Partial, transitive

define Up Cyc.

f Candidate Key

↓

$\{AB\} \rightarrow CDEF$ or c are completely visible from A or B

(Answer) (Partial)

Q4.

Full (candidate key)

Se double key.

Q5

$R(A, B, CD)$

Candidate Key Ki

definition shi joh
Rome dinay me

$A \rightarrow BCD$

$A^+ = \{A, B, C, D\}$

$AD \rightarrow BC$

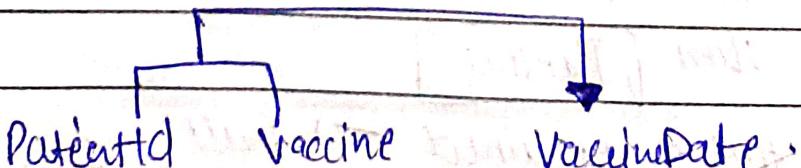
$CD \rightarrow AB$

$C^+ = \{C, A, B, D\}$

$AD^+ = \{A, B, C, D\}$

$CD^+ = \{A, B, C, D\}$

Q6 - Business requirements that hold true for F.U.D



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Normalization: for removing

- inconsistency
- Redundancy
- Anomalies.

Date: / /

Data Loss or Redundancy

"Decomposing larger tables into more organized smaller tables so as to minimize redundancy, inconsistency and anomalies."

* **Normal forms:** Set of rules based on which 3PLA in database is made.

1 NF
2 NF
3 NF

certain groups of attributes are analyzed.

* **1-NF:** (No relation functional Dependencies i.e no hi-ho in 1NF)

If a relation R is said to be in 1NF if and only if:

- ▷ All the attributes of R are atomic in nature.
- ▷ There should be no multi-valued attribute.

Atomic Attribute: → simple and cannot be decomposed.

ex: in RetailStockOutlet → there is a composite attribute → Location

ex: If I want a result based on city → it will be difficult.

so better if we have atomic attributes.

Solution → Split the composite attribute into multiple atomic attributes.

Multivalued Attribute: → use phone number

If result will be in 1NF

Nursing → **Solution 1** → Create separate columns (not good) →

Redundancy → **Solution 2** → have multiple rows for each value

of phone numbers (Primary key will repeat X)

Solution 3 → Create a new table

with one of the table and Multivalued attribute values.

Data Loss

* or
* Redundancy can be both table level + column level

CONSTRAINT ensures PRIMARY KEY (Candidate, Primary)

Date: / /

* Issues with bad database designs:

RetailOutletStock

CK → { RetailOutletId, ItemCode }

1. Insert Anomaly:

RetailOutletId	RetailOutletLocation	ItemCode	Description	Qty
R1004	Gear Street	NULL		

PK hi table me

Sub mukha hai
Inhone

* If a new retail outlet is to be added, but it does not offer any items as of now → still we need to add

↳ this cannot be done because { RetailOutletId, ItemCode } is PRIMARY KEY here → ItemCode cannot be NULL

" when some data in a relation cannot be added due to absence of other data,

2. Delete Anomaly:

RetailOutletId	RetailOutletLocation	ItemCode	Description	Qty
R1001	High Street	I1001	- - -	- - -
R1002	Saturn Street	I1005	Box	

* We want to drop the sale of item I1005,

If we try to delete it

↳

* we will also end up deleting the corresponding details of RetailOutlet

" When deleting some data results in unexpected deletion of other required data."

3. Update Anomaly:

RetailOutlet

* there was an item belonging to class C vehicle part price badhne pe vo class B me chala gaya.

Verin jo class C vela item the vo toh gaya no.

↳ So Data loss is there.

or modification of other data.

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4. When modification of some data results in unexpected deletion

Data loss
Redundancy.

galtat time pe

galtat data
Kahin mile or
Kahin nahi or
Date: / /

Anomaly: is an inconsistency or Irregularity in data that arise due to Insert, update or delete operation.

"Anomalies are the result of Poor Database Design or Redundancies"

Data Redundancy: when same data is present multiple times.

Redundant data will require more storage and lead to potential inconsistency.

(Kahin update ho gaya - Nahi nahi hoga)

Inconsistent data.)

- * INF \rightarrow Advantages \rightarrow queries effectively as it removes ambiguity.
- * Limitation (need for 2NF) \rightarrow There are Partial Functional Dependencies Line description can be Date: / / determined from itemcode only.

\rightarrow Hence we have redundant data in the table, which can be eliminated by having a separate table for item details.

- 2# 2NF (Agar Composite key hai hi naahi toh 2NF dono ka (Partial FD) check karne ka chahiye, nhi)
- A relation R is in 2NF if and only if:

- R is already in INF.
- There is no-partial dependency b/w non-key and key attributes.

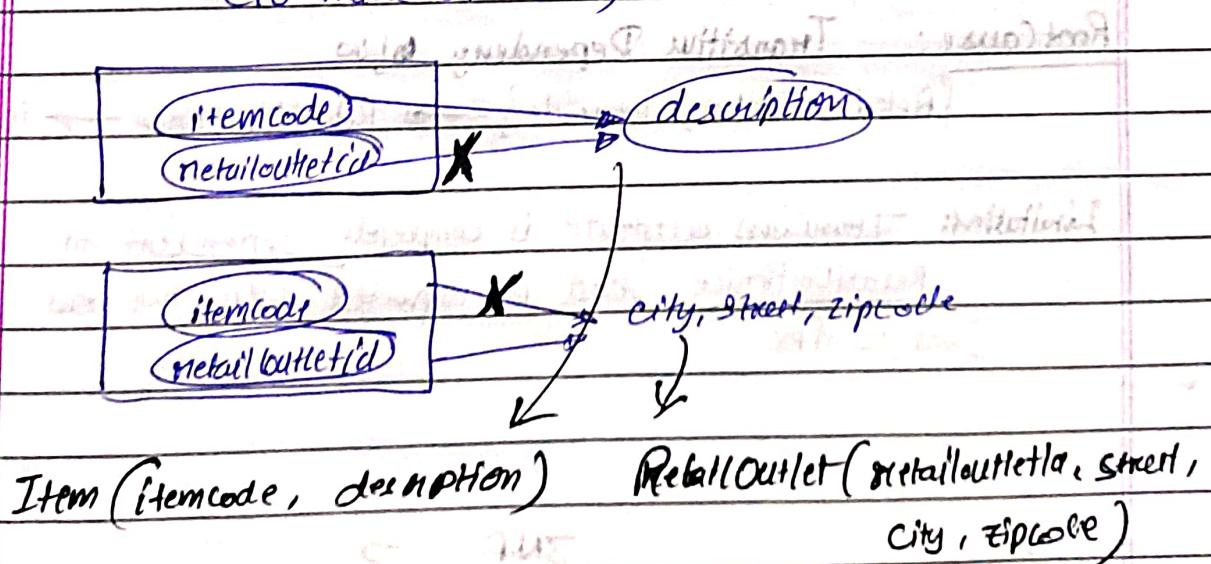
- * How to make a table 2NF?

\hookrightarrow Remove all such Partial Dependencies and decompose the table.

Montra * (Jitni Partial Dependencies Ohne tukde)

In RetailOutletStock, we have following partial dependencies:
which must be removed:

(to make it 2NF)



- * Remaining table as it is with the composite primary key

So tables after this:

Now for 3NF rules
Item, RetailOutletId, RetailerDetails, ItemClass, phonenumber Hajo nohi se example dit s.i.
Date: / /

* 3NF advantages: ensures data integrity, reduces the amount of data duplication.

Limitation → 3NF often leads to higher no. of smaller tables, can make queries more and more complex, may require combining multiple tables to retrieve data.

de-normalized

retailoutletstock { retailoutletid, itemcode, description, retail outlet locations, qtyavailable, retailunitprice, itemclass }

- composite attributes
- multi valued attributes

retailoutlet stock

{ retailoutletid, itemcode, description, street, city, zipcode, qtyavailable, retailunitprice, itemclass }

- partial dependency removed

Item (2NF)

(itemcode, description)

retailoutletstock (2NF)

(retailoutletid, itemcode, itemclass, retailunitprice, qtyavailable)

Transitive dependency here

(retailoutletid, itemcode) → itemclass → itemclass

remove

Item
(3NF)

retailoutletstock (3NF)

ItemClass
(3NF)

Retail Outlet
(3NF)

Normalization - Quiz:

~~CNF~~

Multi-valued attribute

A relation is in 2NF:

- It has a single-attribute Candidate Key(s).
- There is no partial dependency.
- It must be already in 1NF.

SQL Basics → DDL, DML, DCL, TPL

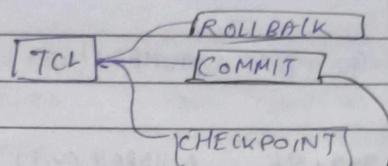
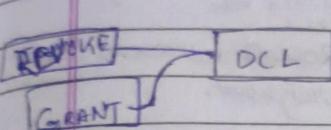
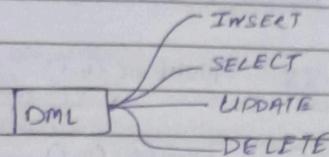
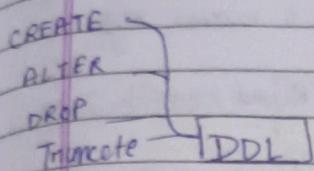
CQL (Structured Query Language) 1970s

↳ High Level language

↳ Platform independent and portable

↳ Standardized and fast.

Date: / /



Q6:

01. PCL TPL

02.

- * Q6: To save the changes permanently in the database.
↓
TCL → COMMIT is used.

Q6.

V.V-Imp



[SQL DataTypes]

#	Input	Used Data Type	Stored As	Reason
1.	134	INTEGER	134	int ← decimal
2.	234.6	INTEGER	234	Truncates value at decimal
3.	134	NUMERIC	134	Represented as it is.
4.	234.6	NUMERIC	935	[Round] towards nearest neighbor.
5.	134	NUMERIC(2)	Data Exception	Results in data loss so, it gives exception.
6.	23.6	NUMERIC(2)	24	[Round] towards nearest neighbor.
7.	23.5	NUMERIC(2)	23	if equidistant then Round towards
8.	23.4	NUMERIC(2)	23	③ ↳ Nearest neighbor. RAJDHANI