# Implementation and Analysis of RSA Algorithm

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#### 1 Introduction to RSA Algorithm

RSA (Rivest-Shamir-Adleman) is a widely used public-key cryptosystem that facilitates secure data transmission. It relies on the mathematical complexity of prime factorization.

## 2 Implementation of RSA Algorithm in Python

The following Python code demonstrates the implementation of the RSA algorithm, including key generation, encryption, and decryption.

#### 2.1 Python Code for RSA Algorithm

Listing 1: RSA Algorithm in Python

```
1 import random
2 import sympy
  class RSA:
      """Function to initialize RSA and specify the key
      def __init__(self, key_size=1024):
          self.key_size = key_size
          self.public_key, self.private_key = self.
             generate_key_pair()
      """Function for generating a large prime number of
10
         specified bit size"""
      def generate_large_prime(self, bits):
11
          return sympy.randprime(2**(bits-1), 2**bits)
13
      """Function to compute the GCD using Euclidean
14
         algorithm"""
      def gcd(self, a, b):
15
          while b:
              a, b = b, a \% b
17
          return a
18
19
      """Function to compute modular inverse using
20
         Extended Euclidean Algorithm"""
      def mod_inverse(self, e, phi):
          def egcd(a, b):
22
```

```
if a == 0:
23
                   return b, 0, 1
24
              g, x, y = egcd(b \% a, a)
              return g, y - (b // a) * x, x
27
          g, x, _ = egcd(e, phi)
28
          if g != 1:
29
               raise ValueError("Modular inverse does not
30
                  exist")
          return x % phi
31
      """Generate RSA key pair"""
33
      def generate_key_pair(self):
34
          p = self.generate_large_prime(self.key_size //
35
          q = self.generate_large_prime(self.key_size //
          n = p * q
37
          phi = (p - 1) * (q - 1)
38
          e = 65537 # Common public exponent
          if self.gcd(e, phi) != 1:
              raise ValueError("e and phi(n) are not
                  coprime, regenerate primes.")
43
          d = self.mod_inverse(e, phi)
44
          public_key = (e, n)
          private_key = (d, n)
47
          return public_key, private_key
48
49
      """Encryption function to encrypt a plaintext
50
         message using the public key"""
      def encrypt(self, plaintext):
51
          e, n = self.public_key
52
          numeric_representation = [ord(char) for char in
             plaintext]
          encrypted_blocks = [pow(char, e, n) for char in
54
             numeric_representation]
          return encrypted_blocks
55
56
57
      """Decryption function to decrypt an encrypted
         message using the private key"""
```

```
def decrypt(self, encrypted_blocks):
          d, n = self.private_key
          decrypted_chars = [chr(pow(block, d, n)) for
             block in encrypted_blocks]
          return "".join(decrypted_chars)
61
63 # Initialize RSA with a 1024-bit key size
_{64} rsa = RSA(key_size=1024)
66 # Test Cases: Different Message Sizes
67 test_messages = [
      "T",
      "Tejas",
69
      "RSA Encryption by me!",
      "This is a test message to check RSA encryption
         using a medium length message",
      "I am Tejas Budhwal and now I am testing for a
72
         larger message and hence I will write a longer
         sentence here so that the test can successfully
         be conducted to show the results.",
73
for i, message in enumerate(test_messages, start=1):
      print(f"Test {i}: Original Message -> {message}")
      encrypted_message = rsa.encrypt(message)
      print(f"Encrypted Message: {encrypted_message}")
      decrypted_message = rsa.decrypt(encrypted_message)
      print(f"Decrypted Message: {decrypted_message}")
82
83
      assert message == decrypted_message, "Decryption
         failed!\n"
      print("Encryption and Decryption Successful!\n")
```

## 3 Explanation of the Code

- **Key Generation:** Generates large prime numbers p and q, computes  $n = p \times q$  and Euler's totient function  $\phi(n) = (p-1)(q-1)$ .
- **Public Key:** Composed of (e, n), where e is a common prime like 65537.
- Private Key: Computed as d, the modular inverse of e modulo  $\phi(n)$ .
- Encryption: Converts plaintext characters to their ASCII values and applies the encryption formula:

$$C = M^e \mod n$$

• Decryption: Recovers the original message using the formula:

$$M = C^d \mod n$$

## 4 Vulnerabilities of RSA and Mitigations

#### 4.1 Vulnerabilities

- 1. **Insufficient Key Length:** Keys smaller than 1024 bits are susceptible to brute-force attacks.
- 2. Common Public Exponent (e = 65537): Although 65537 is widely adopted for efficiency, poorly chosen exponents can weaken encryption strength.
- 3. Factorization Risks: If the modulus n can be factored into its prime components p and q, the private key can be easily derived.
- 4. **Timing Attacks:** By measuring the time taken for decryption, attackers may infer details about the private key.
- 5. Vulnerability to Chosen Ciphertext Attacks: Manipulating ciphertext strategically can help attackers uncover information about the original plaintext.
- 6. Weak Randomness in Prime Generation: Predictable patterns in generating prime numbers p and q can significantly compromise security.

#### 4.2 Mitigation Strategies

- Use Stronger Keys: Employ key sizes of at least 2048 bits to ensure resistance against brute-force attacks. Regularly update cryptographic standards as computational power advances.
- Secure Exponent Selection: While 65537 is secure due to its properties, avoid using extremely small or large exponents. Ensure the modulus is large enough to maintain security even with commonly used exponents.
- Generate Strong Primes: Use cryptographically secure random number generators (CSPRNGs) to create large, unpredictable prime numbers. Apply primality tests like the Miller-Rabin test to verify prime integrity.
- Implement Side-Channel Attack Protections: Design cryptographic operations to execute in constant time, regardless of input values, to prevent timing-based side-channel attacks.
- Padding Schemes: Use OAEP (Optimal Asymmetric Encryption Padding) to prevent chosen ciphertext attacks.
- Use Hybrid Encryption: Combine RSA with symmetric encryption (AES) to handle large data securely.

## 5 Analysis Based on Key Size and Input Size

- **Key Size:** Larger key sizes (2048 or 4096 bits) enhance security but increase computational time for encryption and decryption.
- Input Size: The plaintext message must be smaller than the modulus n. To encrypt larger messages, hybrid encryption (combining RSA with symmetric encryption like AES) is recommended.
- **Performance:** Encryption is faster with a small public exponent (e.g., 65537), while decryption is slower due to larger private exponents.

## 6 RSA Optimization Strategies

- Use Hybrid Encryption: Use RSA to encrypt symmetric keys (AES) instead of full messages.
- Increase Key Size Selectively: 2048-bit RSA is a good balance between security and speed.
- Optimize Prime Number Generation: Use efficient libraries like OpenSSL for prime generation.
- Parallel Processing: Leverage multi-core CPUs for encryption/decryption tasks.
- Use Fast Modular Exponentiation: Optimize mathematical operations using Montgomery Multiplication.

#### 7 Conclusion

RSA remains one of the most robust public-key cryptographic algorithms. However, its security depends heavily on proper implementation, key size, and mitigation of known vulnerabilities.