

Heap size = N

For Parent i -

Left child = $2*i+1$ and Right child = $2*i+2$.

For Child i -

Parent = $\text{ceil}(i/2)-1$

Internal Node range = 0 to $(N/2)-1$ (Half Point)

Leaves range = $(N/2)$ to $(N-1)$ (Half Point)

Max-Heap = Root > Subtree Nodes

Min-Heap = Root < Subtree Nodes

Heapify Algorithm - (MAX HEAP) (LogN)

ith node to be rearranged to maintain heap property. (Except ith node, entire tree is a heap)

(Check with children and swap)

Heapify(ith node)

Check with children and swap till they reach stopping condition.

Stopping condition -

1. Reach Leaf.
2. Current Node > left child and right child.

```
Max_Heapify(A, i) {
    Left = 2*i+1;
    Right = 2*i+2;
    if(Left <= A.heap_size and A[Left] > A[i]) {
        Largest = Left;
    } else {
        Largest = i;
    }

    if(Right <= A.heap and A[Right] > A[Largest]) {
        Largest = Right;
    }

    if(Largest != i) {
        swap(A[Largest], A[i]);
        Max_Heapify(A, Largest);
    }
}
```

Build Heap - (MAX HEAP) $O(N)$

Heap size = N

Internal node range = 0 to $(N/2)-1$

Leaves range = $N/2$ to $N-1$

NOTE: (Leaf Node will always follow Heap property and No need to heapify)

(For applying heapify on ith node, All it's subtrees should be heapified)

Build Heap - Heapify All the internal nodes in Bottom Up (Right Left) Approach.

(Traverse over range of internal nodes and heapify those)

```
Build_MAX_Heap(A) {  
    for( $i = ((A.Length/2)-1)$  to  $0$ ) {  
        Max_Heapify(A,  $i$ );  
    }  
}
```

Extract Max - (Remove max element)

```
Heap_Extract_MAX(A) {  
    if( $A.heap\_size < 1$ ) {  
        print("Heap Underflow");  
    }  
    Max =  $A[0]$ ;  
     $A[0] = A[A.heap\_size-1]$ ;  
     $A.heap\_size--$ ;  
    Max_Heapify(A,  $0$ );  
    return max;  
}
```

Increase Key for MAX_HEAP - Percolate up.

Decrease Key for MAX_HEAP - Heapify. (Percolate down)

Stopping condition -

1. Parent > Child.
2. $I == \text{Root}$ (reached)

```

Max_Heap_Increase_Key(A, i, Key) {
    if(Key < A[i]) {
        print("Wrong operations");
    }
    A[i] = Key;
    while(i > 0 and A[i/2] < A[i]) {
        swap(A[i], A[i/2]);
        i = i/2;
    }
}

```

```

Max_Heap_Decrease_Key(A, i, Key) {
    if(Key > A[i]) {
        print("Wrong operations");
    }
    A[i] = Key;
    Max_Heapify(A, i);
}

```

Insert Element - In a complete binary tree, Insert elements at last level from left to right.

Steps -

- A. Insert node at the last position (Left to Right order)
- B. Percolate up till - (Same as Increase Key algorithm)
 - a. Parent > Current_Node
 - b. Current_Node is the root of the tree.

```

Heap_Insert_Element(A, val) {
    A.heap_size ++;
    A[A.heap_size-1] = val;
    i = A.heap_size-1;
    while(i > 0 and A[i/2] < A[i]) {
        swap(A[i/2], A[i]);
        i = i/2;
    }
}

```

Heap Sort - (Extract min/max) Heapify Algo. O(NlogN)
 (Given: Max Heap)
 (Sort: Sort elements in ASC order)

```

Heap_Sort(A) {
    Build_Heap(A);           //O(N)
    for(i = (A.heap_size-1 to 1)) { //O(N)
        swap(A[i], A[0]);
        Heapify(A[0], i, (Size)); //O(LogN)
    }
}

```

Heap Push / Pop O(LogN)

```

Push(A, element, max_size) {
    if(current_size > max_size) {
        print("Overflow");
    } else {
        A[current_size] = element;
        i = current_size;
        while(i > 0 and A[i/2] < A[i]) {
            swap(A[i], A[i/2]);
            i = i/2;
        }
        current_size++;
    }
}

```

```

Pop(A) {
    if(A.heap_size < 0) {
        print("Underflow");
    }
    Int max = A[0];
    A[0] = A[A.heap_size-1];
    A.heap_size--;
    Max_Heapify(A);
    return max;
}

```