

## Travelling Salesperson Problem (Using LCBB)

- Time complexity is  $O(n^2 \cdot 2^n)$  in case of dynamic programming.
- In Branch & Bound, the time complexity is same  $O(n^2 \cdot 2^n)$  in worst case, but using good bounding fns, will enable these branch & bound algs to solve some problem instances in much less time than required by the dynamic programming.
- Let  $G = (V, E)$  be a directed graph. Let  $c_{ij}$  be the cost of edge  $i \rightarrow j$  &  $c_{ij} = \infty$  if  $i, j \notin E$  & let  $|V| = n$ . Every tour starts & ends at vertex 1.
- To use LCBB to search the travelling salesperson state space tree, we need to define a cost fn  $C(\cdot)$  & two other fns  $\hat{C}(\cdot)$  &  $U(\cdot)$  such that  $\hat{C}(r) \leq C(r) \leq U(r)$  for all nodes  $r$ . The cost  $C(\cdot)$  such that the solution node with least  $C(\cdot)$  corresponds to a shortest tour in it.
- A better  $\hat{C}(\cdot)$  can be obtained by using the reduced cost matrix corresponding to  $A$ . A simple  $\hat{C}(\cdot)$  such that  $\hat{C}(A) \leq C(A) \leq U(A)$  is obtained by defining  $\hat{C}(A)$  to be the length of the path defined at node  $A$ .

## Procedure for solving TSP Problem :-

Step:-

Find out the reduced cost matrix from a given cost matrix (A matrix is to be reduced iff every row & column is reduced). This can be obtained as follows.

Find out (i) row-reduction

(ii), column-reduction.

(i) Row Reduction:- A row (column) is said to be reduced iff it contains at least one zero and all remaining entries are non-negative.

Take the min element from 1<sup>st</sup> row, subtract that element from 1<sup>st</sup> row (including that element), next take min element from 2<sup>nd</sup> row, subtract that element from 2<sup>nd</sup> row (including that element). Similarly apply same procedure for all rows. After applying row-reduction, we will get one resultant matrix. For this apply column reduction.

(ii) Column Reduction:-

Take min element from 1<sup>st</sup> column, subtract that element from 1<sup>st</sup> column. Similarly apply same process for remaining columns.

Now we should find row-wise reduction sum & column wise reduction sum

Row wise reduction sum:-

The sum of the elements which subtracted from rows.

Column wise reduction sum:-

The sum of the elements which subtracted from columns.

Cumulative reduction = Row wise reduction sum + Column wise reduction sum.

- Subtracting a constant  $t$  from every entry in one column or row of the cost matrix reduces the length of every tour by exactly  $t$ . A min-cost tour remains min-cost tour following this subtraction operation. If  $t$  is chosen to be the min entry in row  $i$  (col  $j$ ), then subtracting it from all entries in row  $i$  (col  $j$ ) introduces a zero into row  $i$  (col  $j$ ). Repeating this as often as needed, the cost matrix can be reduced. ~~The total amount~~

- The total amount (Cumulative reduction) subtracted from the cols and rows is a lower bound on the length of the min-cost tour, and can be used as a value for the root of the state space tree.

- Hence all tours in the original graph have a length atleast value of cumulative reduction.

Step2:- For root node, we take initial reduction as lower bound &  $\infty$  as upper bound.

- We can associate a reduced cost matrix with every node in the TSP state space tree. Let  $A$  be the reduced cost matrix for node  $R$ . Let  $S$  be the child of  $R$  such that the tree edge  $(R, S)$  corresponds to including edge  $\{i, j\}$  in the tour. If  $S$  is not a leaf, then the reduced cost matrix for  $S$  may be obtained as follows.

- (i) ~~change~~ If path  $\{i, j\}$  is considered, change all entries in row  $i$ , col  $j$  of  $A$  to  $\infty$ . This prevents the use of any more edges leaving vertex  $i$  or entering vertex  $j$ .
- (ii) Set  $A(j, 1)$  to  $\infty$ . This prevents the use of edge  $\{j, 1\}$ .
- (iii) Apply row reduction & col reduction except for rows & cols containing  $\infty$ . Also find cumulative reduction sum ( $r$ ).
- (iv) Let the resulting matrix  $\hat{A}(S)$ . Calculate

$$\hat{C}(S) = \hat{C}(R) + A(i, j) + r$$

$\hat{C}(S)$  is lower bound for node  $S$ .

$\hat{C}(R)$  is lower bound value of  $S$ 's parent.

$A(i, j)$  is entry in reduced cost matrix  $A$  at node  $R$ .

$r$  = cumulative reduction sum.

For upper bound  $U_R$ , we can use

$$U(R) = \infty \text{ for all nodes } R.$$

$\hat{C}(S)$  is called ranking for repeat step 2 until all nodes are visited.

Example :-

Given Cost Matrix is

R <sub>1</sub>	$\infty$	20	30	10	11
R <sub>2</sub>	15	$\infty$	16	4	2
R <sub>3</sub>	3	5	$\infty$	2	4
R <sub>4</sub>	19	6	18	$\infty$	3
R <sub>5</sub>	16	4	7	16	$\infty$
	c <sub>1</sub>	c <sub>2</sub>	c <sub>3</sub>	c <sub>4</sub>	c <sub>5</sub>

Step 1 :-

(i) Perform row reduction R<sub>1</sub>-10, R<sub>2</sub>-2, R<sub>3</sub>-2,  
R<sub>4</sub>-3, R<sub>5</sub>-4

Then the resultant matrix is

$\infty$	10	20	0	1
13	$\infty$	14	2	0
1	3	$\infty$	0	2
16	3	15	$\infty$	0
12	0	3	12	$\infty$
c <sub>1</sub>	c <sub>2</sub>	c <sub>3</sub>	c <sub>4</sub>	c <sub>5</sub>

(ii) Perform col reduction c<sub>1</sub>-1, c<sub>3</sub>-3.

$\infty$	10	17	0	1
12	$\infty$	11	2	0
0	3	$\infty$	0	2
15	3	12	$\infty$	0
11	0	0	12	$\infty$

Row wise reduction sum = 10+2+2+3+4 = 21

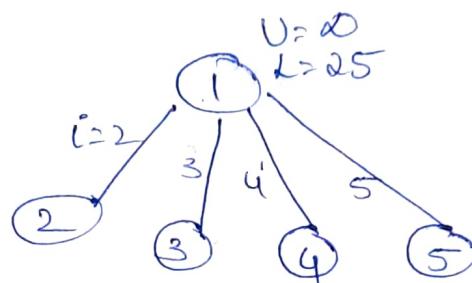
col wise reduction sum = 1+3 = 4

cumulative reduction sum = 21+4 = 25

Hence Reduced cost matrix indicated by:-

$$A = \begin{bmatrix} \infty & 10 & 17 & 0 & 1 \\ 12 & \infty & 11 & 2 & 0 \\ 0 & 3 & \infty & 0 & 2 \\ 15 & 3 & 12 & \infty & 0 \\ 11 & 0 & 0 & 12 & \infty \end{bmatrix} \rightarrow \text{Reduced cost matrix at node 1.}$$

Starting vertex node is 1, from node 1 he can visit next 2, 3, 4 or 5 vertices. This can be shown as



The variable ' $i$ ' indicates next node that sales person will visit.

For starting node, i.e. node 1 has lower bound as cumulative sum  $i=25$ .  
Upper bound is  $\infty$ .

$i=2$  means, next he visits node 2

$i=3$  means, next he visits node 3. & so on

Step 2:- Apply step 2, for reduced cost matrix A.

Consider path  $(\underset{i}{1}, \underset{j}{2})$

(i) Change all entries of 1<sup>st</sup> row & 2<sup>nd</sup> col of A to  $\infty$ . Then matrix is

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & 11 & 2 & 0 \\ 0 & \infty & \infty & 0 & 2 \\ 15 & \infty & 12 & \infty & 0 \\ 11 & \infty & 0 & 12 & \infty \end{bmatrix}$$

(ii) Set  $A(2,1)$  to  $\infty$ , resultant matrix is

$$\begin{array}{c} R_1 \\ R_2 \\ R_3 \\ R_4 \\ R_5 \end{array} \left[ \begin{array}{ccccc} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & 11 & 2 & 0 \\ 0 & \infty & \infty & 0 & 2 \\ 15 & \infty & 12 & \infty & 0 \\ 11 & \infty & 0 & 12 & \infty \end{array} \right] \rightarrow \text{reduced cost-matrix at node 2}$$

(iii) Apply row & col reduction, except for rows & cols containing  $\infty$ .

Row reduction:  $R_2 - 0, R_3 - 0, R_4 - 0, R_5 - 0$ .

row reduction sum =  $0+0+0+0=0$ .

Col reduction:  $-C_1 - 0, C_3 - 0, C_4 - 0, C_5 - 0$ .

Col reduction sum =  $0+0+0+0=0$ .

Cumulative reduction ( $r$ ) =  $0+0=0$ .

$$(iv) \hat{C}(S) = \hat{C}(R) + A(1,2) + r$$

$$\hat{C}(S) = 25 + 10 + 0 = 35$$

Consider the paths  $\begin{pmatrix} 1 \\ 3 \end{pmatrix}$  :-

change all entries of 1<sup>st</sup> row, 3<sup>rd</sup> col of A to  $\infty$ . set  $(3,1)$  to  $\infty$ . Then resultant matrix is

$$\left[ \begin{array}{ccccc} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & \infty & 2 & 0 \\ \infty & 3 & \infty & 0 & 2 \\ 15 & 3 & \infty & \infty & 0 \\ 11 & 0 & \infty & 12 & \infty \end{array} \right]$$

Apply row wise reduction i.e.  $R_2 \rightarrow 0$ ,  
 $R_4 \rightarrow 0$ ,  $R_5 \rightarrow 0$

$\therefore$  row reduction sum is = 0.

Apply col wise reduction  $C_1 \rightarrow 11$ ,  $C_2 \rightarrow 0$ ,  
 $C_4 \rightarrow 0$ ,  $C_5 \rightarrow 0$ . resultant matrix,

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 1 & \infty & \infty & 2 & 0 \\ \infty & 3 & \infty & 0 & 2 \\ 4 & 3 & \infty & \infty & 0 \\ 0 & 0 & \infty & 12 & \infty \end{bmatrix} \rightarrow \text{reduced cost matrix at node 3.}$$

Cumulative reduction ( $\gamma$ ) =  $0 + 11 = 11$ .

$$\begin{aligned} C(S) &= C(R) + A(1,3) + \gamma \\ &= 25 + 17 + 11 \\ &= \underline{\underline{53}} \end{aligned}$$

Consider path (1,4) :-

Change all entries of 1<sup>st</sup> row, 4<sup>th</sup> col of A to  $\infty$ . Set  $A(4,1) \rightarrow \infty$ . Then the resultant matrix is

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & 11 & \infty & 0 \\ 0 & 3 & \infty & \infty & 2 \\ \infty & 3 & 12 & \infty & 0 \\ 11 & 0 & 0 & \infty & \infty \end{bmatrix}$$

Apply row wise reduction,  $R_2 \rightarrow 0$ ,  $R_3 \rightarrow 0$ ,  $R_4 \rightarrow 0$ ,  $R_5 \rightarrow 0$ .

$\therefore$  row reduction sum = 0.

Apply col wise reduction  $C_1 \rightarrow 0$ ,  $C_2 \rightarrow 0$ ,  $C_3 \rightarrow 0$ ,  $C_5 \rightarrow 0$ .

$\therefore$  col reduction sum = 0.

*(Ans)* Cumulative reduction sum =  $0+0=0$ .

∴ Reduced matrix is

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & 11 & \infty & 0 \\ 0 & 3 & \infty & \infty & 2 \\ \infty & 3 & 12 & \infty & 0 \\ 11 & 0 & 0 & \infty & \infty \end{bmatrix} \rightarrow \text{Reduced cost matrix at node 4.}$$

$$\begin{aligned} \hat{C}(S) &= \hat{C}(R) + A(1,4) + r \\ &= 25 + 0 + 0 = 25 \end{aligned}$$

Consider the path  $(1,5)$  :-

Change all entries of 1<sup>st</sup> row, 5<sup>th</sup> col of  $A$  to  $\infty$ . Set  $A(5,1)$  to  $0$ .

The resultant matrix is

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & 11 & 2 & \infty \\ 0 & 3 & \infty & 0 & \infty \\ 15 & 3 & 12 & \infty & \infty \\ \infty & 0 & 0 & 12 & \infty \end{bmatrix}$$

Apply row wise reduction:  $R_2 - 2, R_3 - 0,$   
 $R_4 - 3, R_5 - 0.$

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 10 & \infty & 9 & 0 & \infty \\ 0 & 3 & \infty & 0 & \infty \\ 12 & 0 & 9 & \infty & \infty \\ \infty & 0 & 0 & 12 & \infty \end{bmatrix}$$

row reduction sum =  $2+0+3+0=5$

Apply col reduction:  $C_1=0, C_2=0, C_3=0$

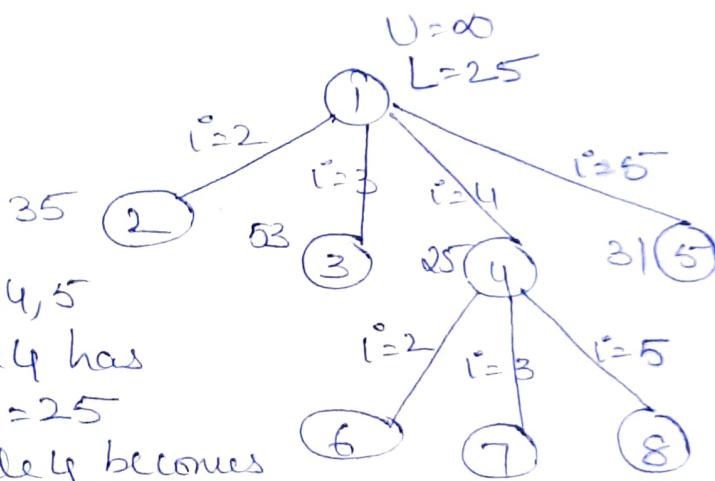
Col reduction sum =  $0+0+0+0=0$

Cumulative reduction sum =  $5+0=5$

$\infty$	$\infty$	$\infty$	$\infty$	$\infty$
10	$\infty$	9	0	$\infty$
0	3	$\infty$	0	$\infty$
12	0	9	$\infty$	$\infty$
$\infty$	0	0	12	$\infty$

- Reduced matrix  
at node 5

$$\begin{aligned} \hat{C}(S) &= \hat{C}(R) + A(1,5) + 8 \\ &= 25 + 1 + 5 \\ &= \underline{\underline{31}} \end{aligned}$$



Among 2, 3, 4, 5  
nodes, node 4 has  
 $\min \hat{C}(S) = 25$

Hence, node 4 becomes  
the next E-node.

So, we get nodes 6, 7, 8 live nodes as  
vertices 1, 4 are already considered. The options  
for node 4 are  $i=2, i=3, i=5$ . Now compute  
the cost of the nodes 6, 7, 8. whichever gets  
min cost that will become next E-node.

Now cost matrix  $A$  is changed to reduced cost matrix at node 4.

$$A = \begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & 11 & \infty & 0 \\ 0 & 3 & \infty & \infty & 2 \\ \infty & 3 & 12 & \infty & 0 \\ 11 & 0 & 0 & \infty & \infty \end{bmatrix}$$

- Reduced cost matrix at node 4. as we are exploring node 4.

Consider the path (4,2):

Change all entries of 4th row, 2nd col of  $A$  to  $\infty$ . Set  $A(2,1)$  to  $\infty$ .

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & 11 & \infty & 0 \\ 0 & \infty & \infty & \infty & 2 \\ \infty & \infty & \infty & \infty & \infty \\ 11 & \infty & 0 & \infty & \infty \end{bmatrix}$$

Apply row reduction

row wise reduction sum = 0.

Apply col wise reduction

col wise reduction sum = 0.

$$r = 0 + 0 = 0.$$

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & 11 & \infty & 0 \\ 0 & \infty & \infty & \infty & 2 \\ \infty & \infty & \infty & \infty & \infty \\ 11 & \infty & 0 & \infty & \infty \end{bmatrix}$$

- Reduced cost matrix at node 6.

$$\begin{aligned} \hat{C}(S) &= \hat{C}(R) + A(4,2) + r \\ &= 25 + 3 + 0 = \underline{\underline{28}} \end{aligned}$$

Consider the path (4,3) :-

change all entries of 6th row, 3rd col  
of A to  $\infty$ , set  $A(3,1)$  to  $\infty$ . Then the  
resultant matrix is

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & \infty & \infty & 0 \\ \infty & 3 & \infty & \infty & 2 \\ \infty & \infty & \infty & \infty & \infty \\ 11 & 0 & \infty & \infty & \infty \end{bmatrix}$$

Apply row wise reduction: R<sub>3-2</sub>

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & \infty & \infty & 0 \\ \infty & 1 & \infty & \infty & 0 \\ \infty & \infty & \infty & \infty & \infty \\ 11 & 0 & \infty & \infty & \infty \end{bmatrix}$$

Apply col wise reduction: C<sub>3-1</sub>

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 1 & \infty & \infty & \infty & 0 \\ \infty & 1 & \infty & \infty & 0 \\ \infty & \infty & \infty & \infty & \infty \\ 0 & 0 & \infty & \infty & \infty \end{bmatrix}$$

→ Reduced cost matrix  
at node 7

$$\gamma = 2 + 11 = \underline{\underline{13}}$$

$$\begin{aligned} \hat{C}(S) &= \hat{C}(R) + A(4,3) + \gamma \\ &= 25 + 12 + 13 \\ &= \underline{\underline{50}} \end{aligned}$$

consider the path (4,5):-

Change all entries of 4<sup>th</sup> row, 5<sup>th</sup> col of A to  $\infty$ . Set  $A(5,1)$  to  $\infty$ . The resultant matrix is

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & 11 & \infty & \infty \\ 0 & 3 & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & 0 & 0 & 0 & \infty \end{bmatrix}$$

Apply row wise reduction:  $R_2 - 11$

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 1 & \infty & 0 & \infty & \infty \\ 0 & 3 & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & 0 & 0 & \infty & \infty \end{bmatrix}$$

Apply col wise reduction:

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ 1 & \infty & 0 & \infty & \infty \\ 0 & 3 & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & 0 & 0 & \infty & \infty \end{bmatrix}$$

+ Reduced cost matrix at node 8

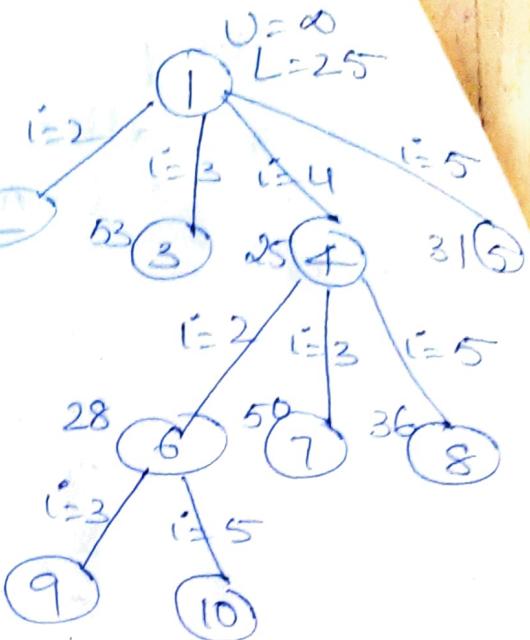
$$r = 11 + 0 = \underline{\underline{11}}$$

$$C(S) = E(R) + A(4,5) + r$$

$$= 25 + 0 + 11$$

$$= \underline{\underline{36}}$$

At this point-  
 node 6 has got min cost, hence it becomes next E-node & nodes 9 & 10 are generated with edge values  $i=3$  &  $i=5$  since the vertices  $i=1, i=4, i=2$  are already considered.



Now reduced cost matrix A is changed by reduced cost matrix at node 6.

$$A = \begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & 11 & \infty & 0 \\ 0 & \infty & \infty & \infty & 2 \\ \infty & \infty & \infty & \infty & \infty \\ 11 & \infty & 0 & \infty & \infty \end{bmatrix}$$

Consider the path (2,3) :-

change all entries of 2nd row, 3rd col of A to  $\infty$ . Set  $A(3,1)$  to  $\infty$ .

$$\begin{bmatrix} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & 2 \\ \infty & \infty & \infty & \infty & \infty \\ 11 & \infty & \infty & \infty & \infty \end{bmatrix}$$

Apply row wise reduction:- R<sub>3</sub>-2, R<sub>5</sub>-11  
 row wise reduction sum = 2+11=13.

$$\left[ \begin{array}{ccccc} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & 0 \\ \infty & \infty & \infty & \infty & \infty \\ 0 & \infty & \infty & \infty & \infty \end{array} \right]$$

Apply col wise reduction, ~~Get 000~~

$$\left[ \begin{array}{ccccc} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & 00 \\ \infty & \infty & \infty & \infty & 0 \\ \infty & \infty & \infty & \infty & \infty \\ 0 & \infty & \infty & \infty & \infty \end{array} \right]$$

- Reduced cost matrix at node 9.

col wise reduction sum = 0.

$$r = 13 + 0 = \underline{\underline{13}}$$

$$\begin{aligned} C(S) &= C(R) + A(2,3) + r \\ &= 28 + 11 + 13 \\ &= \underline{\underline{52}} \end{aligned}$$

Consider the path (2,5) :-

change all entries of 2<sup>nd</sup> row, 5<sup>th</sup> col of A to  $\infty$ . Set  $A(5,1)$  to  $\infty$ . Then the resultant matrix is.

$$\left[ \begin{array}{ccccc} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ 0 & \infty & \infty & 0 & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & 0 & \infty & \infty \end{array} \right]$$

Row wise reduction sum = 0.

col wise reduction sum = 0.

$$r = 0 + 0 = \underline{\underline{0}}$$

$$\left[ \begin{array}{cccccc} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ 0 & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & 0 & \infty & \infty \end{array} \right] - \text{Reduced cost matrix at node}$$

$$\begin{aligned} C(S) &= C(R) + A(2,5) + r \\ &= 28 + 0 + 0 \\ &= \underline{\underline{28}} \end{aligned}$$

Next node to be explored is node 10.  
Since it has least cost.

Now reduced cost matrix  $A$  is change to reduced cost matrix at node 10.

$$A = \left[ \begin{array}{cccccc} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ 0 & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & 0 & \infty & \infty \end{array} \right]$$

Consider the path  $(5,3)$  :-

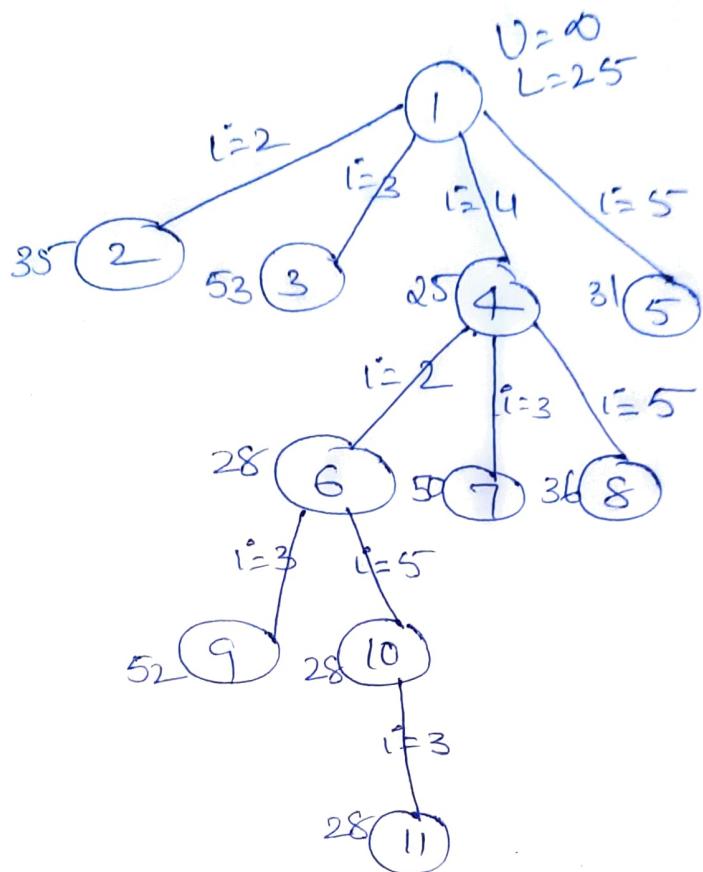
Change all entries of 5<sup>th</sup> row, 3<sup>rd</sup> col of  $A$  to  $\infty$ . Set  $A(3,1)$  to  $0$ .

$$\left[ \begin{array}{ccccc} \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ 0 & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & \infty & \infty & \infty \end{array} \right]$$

Apply row & col wise reduction

$$r = 0 + 0 = 0$$

$$\begin{aligned}
 C(S) &= C(R) + AC(5, 3) + r \\
 &= 28 + 0 + 0 = \underline{\underline{28}}
 \end{aligned}$$



Hence the paths for travelling salesperson problem is the paths from root to answer node i.e  $1 \rightarrow 4 \rightarrow 2 \rightarrow 5 \rightarrow 3 - 1$

From this paths we get min-cost

$$\text{as } 10 + 6 + 8 + 7 + 3 = \underline{\underline{28}}$$