CS582 Distributed Systems

Quiz 7 Solution (20 marks)

- 1. In Dynamo, N denotes the number of replicas, R denotes the read quorum, and W denotes the write quorum. For a fixed value N, R and W can be varied to achieve different levels of consistency, availability, performance, and durability suitable for a particular environment. For the following scenarios (assuming N = 5), suggest the most optimal value of R and W, while justifying your choice to receive full credit.
 - a) For a GET request, read the value from the closest replica to minimize delay, making sure the value is not stale.

Solution: $\begin{array}{l} R=1 \\ W=5 \\ Since \ we \ can \ only \ read \ from \ the \ closest \ replica, \ R=1, \ and \ to \ ensure \ that \ this \ replica \ never \ returns \ a \ stale \end{array}$

value, we must write to all replicas, so W = 5 (1 mark)

b) Speed up PUT requests as much as possible while making sure GET requests never return stale values.

To speed up PUT requests as much as possible, we can write to just one replica, so W=1, but for non-stale values to be returned, we must read from all replicas, so R=5 (1 mark)

c) Minimize read and write latency while achieving high durability, ensuring data is not lost and stale values aren't read even if two replicas fail.

Solution: R = 3 (0.5 marks) W = 3 (0.5 marks)

To maintain the durability of writes, we must replicate writes to at least more than two out of five replicas, so W = 3. Moreover, we should read from at least three replicas to ensure both read and write quorums intersect in order to avoid stale reads, so R = 3 (1 mark)

2. In Memcache, the concept of leasing is used to tackle the thundering herds problem. Consider the following scenario involving Temucache, a Memcache clone with different but self-explanatory APIs, and three clients C1, C2, and C3: At T = 0, the database has the key-value pair [X, 30], while the Temcache server does not have the key X.

T = 1 C1: temu_get(X)
T = 2 C2: temu_get(X)
T = 3 C1: db_get(X)
T = 4 C3: db_put(X, 60)
T = 5 C3: temu_delete(X)
T = 6 C1: temu_put(X, 30)
T = 7 C2: temu_get(X)
T = 8 C2: db_get(X)
T = 9 C3: temu_get(X)
T = 10 C2: temu_put(X, 60)
T = 11 C1: temu_get(X)

At the following points, is any value returned? If yes, state the value. Otherwise, explain why no value is returned.

a) T = 3

Solution: 30 (2 marks)

b) T = 7

Solution:

No value is returned as it is a cache miss.

(0.5 marks)

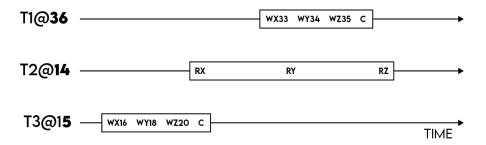
Because at T=4, C3 put the key-value pair [X,60] into the database, thus deleting [X,30] from the Temucache server and invalidating C1's lease. (1 mark)

Hence, no key-value pair is inserted into the Memcache server at T=6, leading to a cache miss when C2 tries to read X at T=7. (1 mark)

c) T = 11



3. Consider the following scenario involving Snapshot Isolation where the initial values of X, Y, and Z are 6, 7, and 8 respectively (written and committed by a much older transaction T_0).



a) Assuming clocks are synchronized, and without accounting for Replica Safe Time, what value of X, Y, and Z are read by transaction T_2 ?

b) Maintaining the same assumption as in (a), if T_2 has the timestamp 44 instead of 14, what values of X, Y, and Z are read by T_2 ?

Solution: X = 33 Y = 34	(0.5 marks) (0.5 marks)
Z = 35	$(0.5 \mathrm{marks})$

c) In Spanner, R/O transactions and R/W transactions are timestamped differently. If R/W transactions are timestamped like R/O transactions due to a glitch, what values of X, Y, and Z will T_2 read? Does this violate serializability and linearizability? Explain.

Solution:	
X = 6	(0.5 marks)
Y = 18	(0.5 marks)
Z = 20	(0.5 marks)
	,
Serializability will be violated	(0.5 marks)
Since T2 is reading partially from T0 and T3, we can not order them in any way	(1 mark)
Linearizability will be violated	$(0.5 \mathrm{marks})$
Since there is no serial order, let alone a real-time order	(1 mark)