Qirby Optimisation

Extra Indexes

For the creation of extra indexes on the Z environment, multiple criteria were used to decide whether to create an index over that column, criteria such as:

- Usage of the column to merge tables (JOIN statements);
- Usage of the column to filter tables (WHERE statements);
- Usage of the column to aggregate tables (GROUP BY statements);

Furthermore, other criteria were used to decide the type of index, such as:

• Cardinality of the column;

A list of candidate columns for indexes was created that consisted on:

- ANO_LETIVO, PERIODO, CODIGO used to join tables on all queries that require to calculate number of hours and other information involving table XOCORRENCIAS and XTIPOSAULA;
- TIPO used extensively for aggregating data based on the type of occurrence and to filter data;
- ANO_LETIVO used vastly for filtering tables;
- CURSO used for filtering and aggregation on half the queries;
- CODIGO used for joining tables;
- ID used for joining tables;

To decide on the type of index to be created, we studied the statistics of the columns to help with the decision:

- CODIGO, ANO_LETIVO, PERIODO This composite is almost completely unique, and for that reason the BITMAP index wasn't even considered. This index will be created as a B-Tree index.
- TIPO Low cardinality column, only having 5 distinct values out of over 20k rows, and furthermore, the column is stable. For these reasons, it was opted to create a BITMAP index.
- ANO_LETIVO Based on the low cardinality of this column, the slow growth of unique values (only once a year) and the types of queries we have, it was decided to create a BITMAP index on this column, on both TIPOSAULA and OCORRENCIAS tables. However, when comparing execution plans, a BITMAP index on this column wouldn't be used as opposed to a B-Tree index, due to the need of performing range scans. B-tree index comes at the cost of using approximately 10x more space (0.5625MB vs 0.0625MB). Therefore, it was opted to use a B-tree index even with criteria pointing to a BITMAP index, as the former affected more queries than the latter.
- CURSO This column has a low cardinality (2:100), which is higher than the threshold considered (1:100). However, due to the stability of this table, as the table is never changed unless a restructure of the UCs is made, we opted to use a BITMAP index.
- CODIGO This column is a leading column on another index (primary key of OCORRENCIAS), therefore the creation of index in this column is redundant and wasn't performed;
- ID Column possesses a huge number of distinct values, therefore the only index that was considered
 was the B-Tree index.

Thus, the indexes created were:

• cap_idx - B-tree index on columns CODIGO, ANO_LETIVO, PERIODO of table ZTIPOSAULA

CREATE INDEX cap_idx ON ZTIPOSAULA(CODIGO, ANO_LETIVO, PERIODO);

• tipo_idx - Bitmap index on column TIPO of table ZTIPOSAULA

```
CREATE BITMAP INDEX tipo idx ON ZTIPOSAULA(TIPO);
```

• ano_tp_idx & ano_oc_idx - B-tree index on column ANO_LETIVO of table ZTIPOSAULA and on column ANO LETIVO of table ZOCORRENCIAS

```
CREATE INDEX ano_tp_idx ON ZTIPOSAULA(ANO_LETIVO);
CREATE INDEX ano_oc_idx ON ZOCORRENCIAS(ANO_LETIVO);
```

• curso_idx - Bitmap index on column CURSO of table ZUCS

```
CREATE BITMAP INDEX curso_idx ON ZUCS(CURSO);
```

• id_idx - B-tree index on column ID of table ZDSD

```
CREATE INDEX id_idx ON ZDSD(ID);
```

Query 1 - Select and Join

Show the codigo, designação, ano_letivo, inscritos, tipo, and turnos for the course 'Bases de Dados' of the program 275.

SQL Query

```
SELECT CODIGO, DESIGNACAO, ANO_LETIVO, INSCRITOS, TIPO, TURNOS FROM XUCS

INNER JOIN XOCORRENCIAS USING(CODIGO)

INNER JOIN XTIPOSAULA USING (ANO_LETIVO, PERIODO, CODIGO)

WHERE DESIGNACAO = 'Bases de Dados' AND CURSO = 275;
```

Result

		\$ ANO_LETIVO		∯ TIPO	
1 EIC3106	Bases de Dados	2003/2004	92	T	1
2 EIC3106	Bases de Dados	2003/2004	92	TP	4
3 EIC3106	Bases de Dados	2004/2005	114	T	1
4 EIC3106	Bases de Dados	2004/2005	114	TP	4
5 EIC3111	Bases de Dados	2005/2006	(null)	T	1
6 EIC3111	Bases de Dados	2005/2006	(null)	TP	6

Figure 1: Query 1 results

Execution Plan

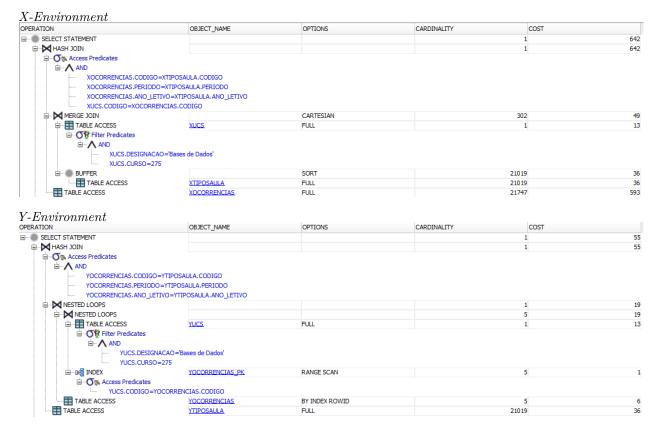
As can be seen in the execution plans below, the X environment presents the highest costs for the query, as expected. This is due to the non-existence of indexes as compared to Y (has indexes created from primary keys) and Z (has primary key indexes and extra indexes), which causes the query to perform full accesses to the tables. In the Y environment, this cost is optimised due to the existence of indexes created from the primary keys, reducing the cost of access on tables and also on the join of the tables (as can be seen on the join between YUCS and YOCORRENCIAS).

In the Z environment, the query is further optimised due to the creation of the index cap_idx on ZTIPOSAULA which permitted a more efficient join with the ZOCORRENCIAS table, cutting a lot of the cost from the Y

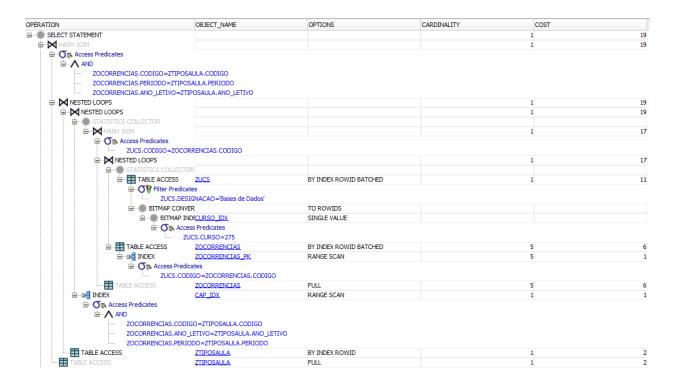
environment. The existence of the <code>curso_idx</code> index also helped when filtering the table <code>ZUCS</code> for the course <code>275</code>.

Cost optimisation in comparison to X environment

X	Y	Z
0%	-91.4%	-97%



Z-Environment



Query 2 - Aggregation

How many class hours of each type did the program 233 planned in year 2004/2005?

SQL Query

```
SELECT TIPO,
    SUM(COALESCE(N_AULAS, 1) * HORAS_TURNO * DECODE(PERIODO /*EXPR*/,
                    '1S', 6, /*SEARCH, RESULT*/
                    '2S', 6,
                    '1T', 3,
                   '2T', 3,
                   '3T', 3,
                   '4T', 3,
                   'T', 3,
                   'A', 12,
                   'B', 2,
                   6 /*DEFAULT*/) * 4) AS HOURS
FROM XOCORRENCIAS
    INNER JOIN XUCS USING (CODIGO)
    INNER JOIN XTIPOSAULA USING (CODIGO, ANO_LETIVO, PERIODO)
WHERE
    ANO_LETIVO = '2004/2005'
    AND CURSO = '233'
GROUP BY TIPO;
```

Result

Execution Plan

As can be seen in the execution plans below, the X environment presents the highest costs for the query, as expected. This is due to the non-existence of indexes as compared to Y (has indexes created from primary keys) and Z (has primary key indexes and extra indexes), which causes the query to perform full accesses to the tables. In the Y environment, this cost is optimised due to the existence of indexes created from the

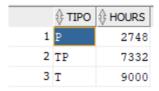


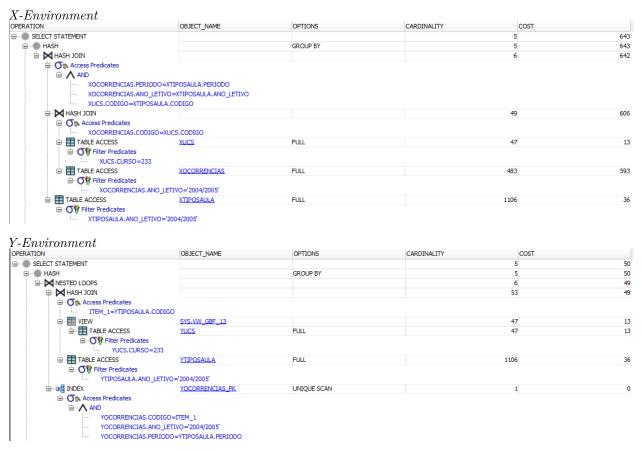
Figure 2: Query 2 result

primary keys, reducing the cost of access on tables and also on the join of the tables (as can be seen on the hash join operation between the tables).

The Z-environment is further optimised because of the creation of indexes cap_idx and ano_tp_idx which enable a more efficient join and filter of tables, respectively.

Cost optimisation in comparison to X environment

X	Y	Z
0%	-92.2%	-92.7%



Z-Environment



Query 3 - Negation

Which courses (show the code) did have occurrences planned but did not get service assigned in year 2003/2004?

```
Case A - Using NOT IN

SQL Query

SELECT CODIGO
FROM XUCS
INNER JOIN XOCORRENCIAS USING (CODIGO)
WHERE ANO_LETIVO = '2003/2004'
AND CODIGO NOT IN (
SELECT CODIGO
FROM XTIPOSAULA
INNER JOIN XDSD USING (ID)
WHERE ANO_LETIVO = '2003/2004'
);
```

Result

138 unique values

		⊕ CODIGO	∯ CODIGO		⊕ CODIGO
1	MEA412	26 MEAM1312	51 MEM179	76 MDI1106	101 MPFCA203
2	MTM115	27 MEAM1310	52 MEM180	77 MDI1107	102 MPFCA204
3	MTM114	28 MEMT135	53 MEM181	78 MDI1205	103 MPFCA205
4	MEB205	29 MEMT100	54 MEM182	79 MDI1206	104 MPFCA206
5	EMM528	30 MPPAU2220	55 MEM183	80 MDI1207	105 EC5200
6	MEM1205	31 MPPAU2218	56 MEM184	81 MDI1208	106 EEC5022
7	GEI512	32 MPPAU2216	57 MEM187	82 MDI1209	107 EEC2207
8	MEMT107	33 MPPAU2215	58 MEM188	83 MDI1108	108 EIC4220
9	MEMT105	34 MPPAU1114	59 MEM189	84 MEB204	109 EIC4221
10	MEMT102	35 MPPAU1113	60 MEM191	85 MEMT120	110 EIC4222
11	MEB105	36 MPPAU1112	61 MVC1211	86 MPPAU1115	111 EIC4223
12	MTM111	37 EC5280	62 MEA112	87 MPPAU2217	112 EIC4224
13	MTM110	38 EEC5272	63 MEA219	88 MPPAU2219	113 EIC5122
14	MTM104	39 MEM5000	64 MEA215	89 MMCCE1220	114 EIC5123
15	EQ407	40 EQ418	65 MEA216	90 MPFCA100	115 EIC5124
16	EQ308	41 EQ411	66 MEA217	91 MPFCA101	116 EIC5125
17	MEAM5000	42 MTM108	67 MEA319	92 MPFCA102	117 EIC5126
18	MDI1204	43 MEMT131	68 MEA320	93 MPFCA103	118 EIC5127
19	MDI1105	44 EIC3209	69 MEA414	94 MPFCA104	119 EIC5129
20	MDI1103	45 MEEC1053	70 MEA415	95 MPFCA105	120 EIC4225
21	MDI1100	46 MEM157	71 MEST210	96 MPFCA106	121 CI014
22	MFAMF1108	47 MEM158	72 MEMT110	97 MPFCA107	122 CI018
23	MEMT2000	48 MEM163	73 MEMT106	98 MPFCA200	123 CI019
24	MEMT1000	49 MEM175	74 EI1107	99 MPFCA201	124 CI023
25	MEAM1314	50 MEM175	75 EC5287	100 MPFCA202	125 CI003

	⊕ CODIGO
115	EIC5124
	EIC5125
	EIC5126
	EIC5127
	EIC5129
	EIC4225
	CI014
	CI018
	CI019
	CI023
125	CI003
126	CI038
127	CI002
128	CI004
129	CI007
130	CI020
131	CI016
132	CI013
133	CI017
134	CI008
135	CI009
136	CI011
137	CI025
138	CI027
139	CI037

Execution Plan

As can be seen in the execution plans below, the X environment presents the highest costs for the query, as expected. This is due to the non-existence of indexes as compared to Y (has indexes created from primary keys) and Z (has primary key indexes and extra indexes), which causes the query to perform full accesses to the tables. In the Y environment, this cost is optimised due to the existence of indexes created from the primary keys, reducing the cost of access to tables and also on the join of the tables (mainly seen on the fast full scan performed on the YOCORRENCIAS instead of a full scan).

On Z-environment, the query is further optimised due to the creation of the index <code>ano_tp_idx</code> that improved the cost-effectiveness of the filtering of <code>ZTIPOSAULA</code> table. Another minor optimisation was due to the creation of the index <code>id_idx</code> which reduced the cost of the fast full scan on YDSD table, however this reduction was minor as the previous environment already had used the index created by the primary key to optimise this scan.

Cost optimisation in comparison to X environment

X	Y	Z
0%	-87.3%	-88.6%

X-Environment



Case B - Using OUTER JOIN and IS NULL

SQL Query

```
SELECT XUCS.CODIGO AS CODIGO

FROM XUCS

INNER JOIN XOCORRENCIAS ON XUCS.CODIGO = XOCORRENCIAS.CODIGO

LEFT OUTER JOIN (

SELECT CODIGO

FROM XTIPOSAULA

INNER JOIN XDSD USING (ID)

WHERE ANO_LETIVO = '2003/2004'

) temp ON XUCS.CODIGO = temp.CODIGO

WHERE ANO_LETIVO = '2003/2004' AND temp.CODIGO IS NULL;
```

Result

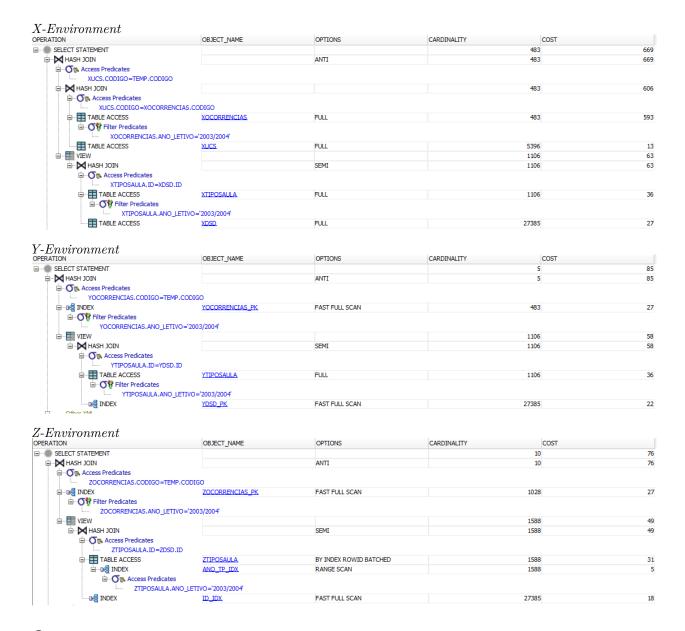
Same result as case A.

Execution Plan

Judging from the execution plan, the Oracle SQL environment optimised the query performed that resulted on the exact same execution plan as case A, and therefore all the conclusions explained above apply the same for this case.

Cost optimisation in comparison to X environment

X	Y	Z
0%	-87.3%	-88.6%



Query 4

Who is the professor with more class hours for each type of class, in the academic year 2003/2004? Show the number and name of the professor, the type of class and the total of class hours times the factor.

SQL Query

```
WITH AUX (NR, NOME, TIPO, HOURS) AS (
    SELECT NR, NOME, TIPO,
            SUM(HORAS * DECODE(PERIODO /*EXPR*/,
                     '1S', 6, /*SEARCH, RESULT*/
                     '2S', 6,
                     '1T', 3,
                     '2T', 3,
                     '3T', 3,
                     '4T', 3,
                     'T', 3,
                     'A', 12,
                     'B', 2,
                    6 /*DEFAULT*/) * 4) AS HOURS
        FROM XDOCENTES
            INNER JOIN XDSD USING(NR)
            INNER JOIN XTIPOSAULA USING(ID)
            INNER JOIN XOCORRENCIAS USING(CODIGO, ANO_LETIVO, PERIODO)
        WHERE ANO LETIVO = '2003/2004'
        GROUP BY (TIPO, NR, NOME)
)
SELECT NR, NOME, TIPO, HOURS
    FROM AUX
        INNER JOIN (
            SELECT TIPO, MAX(HOURS) AS HOURS
                FROM AUX
                GROUP BY TIPO
        ) USING (TIPO, HOURS);
```

Result

	∯ NR	♦ NOME	∜ TIPO	∯ HOURS
1	249564	Cecília do Carmo Ferreira da Silva	TP	624
2	210006	João Carlos Pascoal de Faria	OT	84
3	207638	Fernando Francisco Machado Veloso Gomes	T	500.04
4	208187	António Almerindo Pinheiro Vieira	P	720

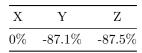
Figure 3: Query 4 result

Execution Plan

As can be seen in the execution plans below, the X environment presents the highest costs for the query, as expected. This is due to the nonexistance of indexes as compared to Y (has indexes created from primary keys) and Z (has primary key indexes and extra indexes), which causes the query to perform full accesses to the tables. In the Y environment, this cost is optimised due to the existence of indexes created from the primary keys, reducing the cost of access on tables and also on the join of the tables (as can be seen on the hash join operation between the tables).

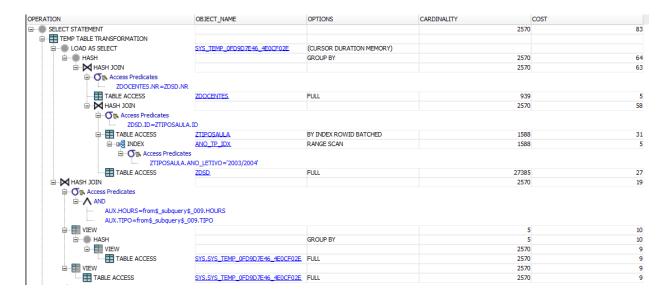
The Z-environment is slightly optimised thanks to the ano_tp_idx index, which allowed for a range scan instead of a full scan.

Cost optimisation in comparison to X environment





 $Z ext{-}Environment$



Query 5

Compare the execution plans (just the environment Z) and the index sizes for the query giving the course code, the academic year, the period, and number of hours of the type 'OT' in the academic years of 2002/2003 and 2003/2004.

SQL Query

```
SELECT CODIGO, ANO LETIVO, PERIODO, SUM(COALESCE(N AULAS, 1) * HORAS TURNO * DECODE(PERIODO /*EXPR*/,
                                                      '1S', 6, /*SEARCH, RESULT*/
                                                      '2S', 6,
                                                      '1T', 3,
                                                      '2T', 3,
                                                      '3T', 3,
                                                      '4T', 3,
                                                      'T', 3,
                                                      'A', 12,
                                                      'B', 2,
                                                     6 /*DEFAULT*/) * 4) AS HOURS
    FROM XOCORRENCIAS
        INNER JOIN XTIPOSAULA USING (CODIGO, ANO_LETIVO, PERIODO)
    WHERE TIPO = 'OT' AND (
        ANO_LETIVO = '2002/2003' OR ANO_LETIVO = '2003/2004'
    )
    GROUP BY (CODIGO, ANO_LETIVO, PERIODO);
Result
```

\$\text{CODIGO} \$\tilde{\psi}\$ ANO_LETIVO \$\tilde{\psi}\$ PERIODO \$\tilde{\psi}\$ HOURS \$\\ 1 \text{EIC5202} 2002/2003 2S 12 \$\\ 2 \text{EIC5202} 2003/2004 2S 12

Figure 4: Query 5 result

Execution Plan

X-Environment



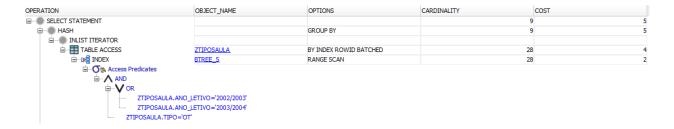
Case A - Using B-tree

SQL Query

```
CREATE INDEX BTREE 5 ON ZTIPOSAULA(ANO LETIVO, TIPO);
```

```
SELECT CODIGO, ANO_LETIVO, PERIODO, SUM(COALESCE(N_AULAS, 1) * HORAS_TURNO * DECODE(PERIODO /*EXPR*/,
                                                      '1S', 6, /*SEARCH, RESULT*/
                                                      '2S', 6,
                                                      '1T', 3,
                                                      '2T', 3,
                                                      '3T', 3,
                                                      '4T', 3,
                                                      'T', 3,
                                                      'A', 12,
                                                      'B', 2,
                                                     6 /*DEFAULT*/) * 4) AS HOURS
    FROM ZOCORRENCIAS
        INNER JOIN ZTIPOSAULA USING (CODIGO, ANO_LETIVO, PERIODO)
    WHERE TIPO = 'OT' AND (
        ANO_LETIVO = '2002/2003' OR ANO_LETIVO = '2003/2004'
    );
DROP INDEX BTREE_5;
```

Execution Plan
Z-Environment



Case B - Using Bitmap

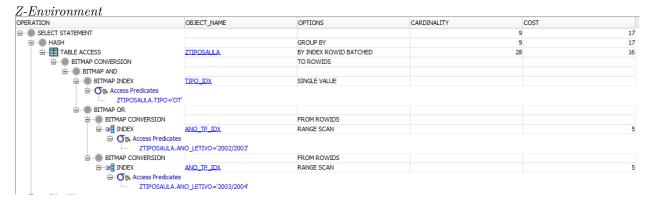
SQL Query

```
CREATE BITMAP INDEX BITMAP_5 ON ZTIPOSAULA(ANO_LETIVO, TIPO);
```

ANO_LETIVO = '2002/2003' OR ANO_LETIVO = '2003/2004'

DROP INDEX BITMAP_5;

Execution Plan



Cost optimisation in comparison to X environment

GROUP BY (CODIGO, ANO_LETIVO, PERIODO);

X	Y	Z (A)	Z (B)
0%	-94.1%	-99.2%	-97.3%

The index sizes were obtained via this SQL statement:

B-tree allows for a more efficient way to fetch the results via a range scan on the index. The Bitmap suffers on the cost optimisation due to having to convert the existing indexes into Bitmap index to perform the operations, this could be avoided by making the indexes ano_tp_idx a bitmap index, however it would make other queries less cost-efficient.

From these, we can take the usual trade-off of cost vs space, as the B-tree index allows for a better cost optimisation (approximately 3x less cost than Bitmap index), however it comes with the cost of taking 10x more space to store the index.

Query 6

Select the programs (curso) that have classes with all the existing types.

SQL Query

```
SELECT CURSO
FROM XOCORRENCIAS

INNER JOIN XUCS USING (CODIGO)

INNER JOIN XTIPOSAULA USING (CODIGO, ANO_LETIVO, PERIODO)
GROUP BY CURSO
HAVING COUNT(DISTINCT TIPO) = 5;

-- without hard coded value on the number of distinct types
SELECT CURSO
FROM XOCORRENCIAS

INNER JOIN XUCS USING (CODIGO)

INNER JOIN XTIPOSAULA USING (CODIGO, ANO_LETIVO, PERIODO)
GROUP BY CURSO
HAVING COUNT(DISTINCT TIPO) = (

SELECT COUNT(DISTINCT TIPO)

FROM XTIPOSAULA);
```

Result

1	9461
2	4495
3	9508
4	2021

Figure 5: Query 6 result

Execution Plan

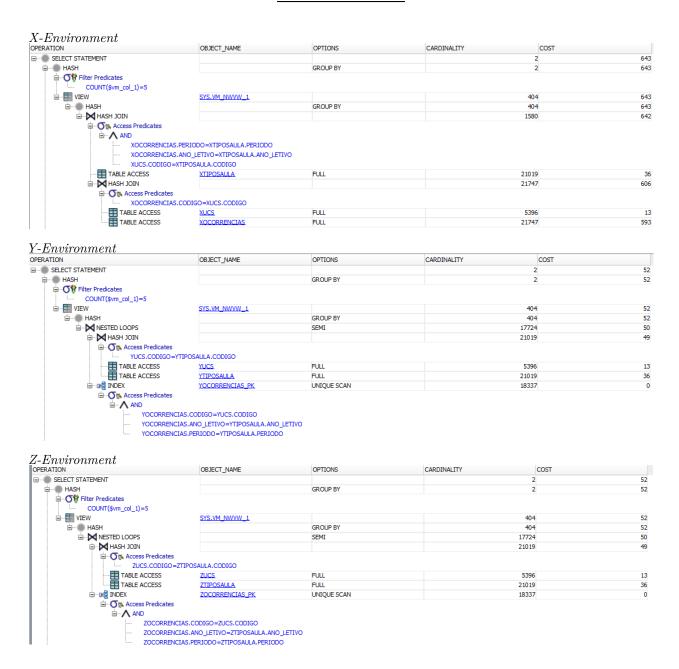
As can be seen in the execution plans below, the X environment presents the highest costs for the query, as expected. This is due to the non-existence of indexes as compared to Y (has indexes created from primary

keys) and Z (has primary key indexes and extra indexes), which causes the query to perform full accesses to the tables. In the Y environment, this cost is optimised due to the existence of indexes created from the primary keys, reducing the cost of access to tables and also on the join of the tables (as can be seen on the hash join operation between the tables).

However, our indexes didn't further optimise this query, staying with the same execution plan and the same cost for query.

Cost optimisation in comparison to X environment

X	Y	Z
0%	-91.9%	-91.9%



Overall System Cost Optimisation

X	Y	Z
0%	-90.7%	-92.8%