

1.

a.

if L is regular language, so  $L = \{0^n : n = 2^k \text{ for some } k > 1\}$

m be the constant in the pumping lemma, so  $\omega = 0^{2^m} \in L$ ,  $|\omega| = 2^m \geq m$ ,

so all possible x,y,z  $\omega = xyz$ ,  $|xy| \leq m$ ,  $|y| \geq 1$ .

Case:

$$x = 0^r, y = 0^s, z = 0^{2^m - r - s}, r + s \leq m, s \geq 1$$

$$\text{Let } i=2, \text{ so } \omega = xy^2z = 0^{2^m + s}$$

$$\text{Because } 2^m < 2^m + s \leq 2^m + m < 2^m + 2^m = 2^{m+1}$$

$$2^m + s \neq 2^K, \text{ so } i=2, \omega = 0^{2^m + s} \notin L, \text{ so L is not regular language.}$$

b.

if L is regular language, so  $L = \{abb, aab, aaab, \dots\}$

because  $n(a) \neq n(b)$ , so m be the constant in the pumping lemma,

$$\omega = a^m b^{m-s},$$

Case when  $\omega = aaabbbb$ ,  $x=aa, y=a, z=bbbb$ ,

$$i=2, \omega = xy^2z = (aa)(a)^2(bbbb) = aaaabbbb,$$

Then  $n(a)=n(b)=4$ ;

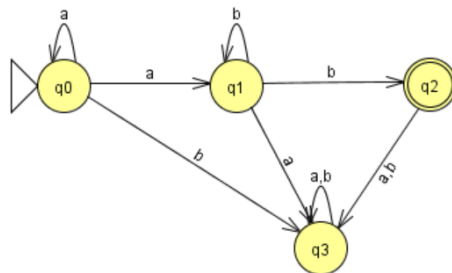
So L is not regular language.

2.

a.

let  $L1 = \{a^n b^n : n > 0\}$ ,  $L2 = \{a^n b^m : n > 0, m > 0\}$ ;

then because L2 is regular language



$L_2$  is regular by the  $L_2 = aa^*bb^*$ ;

So  $L = L_1 \cup L_2 = \{a^n b^n : n > 0\} \cup \{a^n b^m : n > 0, m > 0\}$  is regular language;

b.

because  $L = \{a^n b^m : n \leq m \leq 2n\}$ , if  $L$  is regular language

let  $m$  be as in pumping length

$w = a^m b^m$ , then let  $w = xyz$  be in pumping lemma,

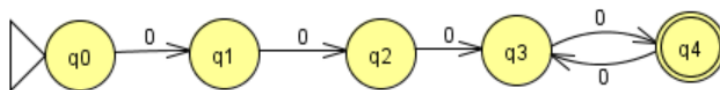
When  $i=2$ ,  $y=a^k$ ,  $w = xy^2z = a^{m+k}a^m$

Because  $m+k > m$ , so  $L$  is not regular language.

c.

$L$  is regular language, so  $L = \{0^n : n = 2k \text{ for some } k > 1\}$

The regular expression is  $L = 000(00)^*0$ ;

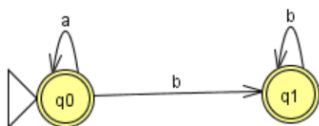


3.

As when  $L_1 = \{a^n b^m : n = m\}$ ,  $L_1$  is not a regular language

$L_2 = \{a^n b^m : n \neq m\}$ ,  $L_2$  is not a regular language

But when  $L_1 \cup L_2$ , the expression is  $L = a^*b^*$ ;



This is regular language;

4.

The definition of symmetric difference of two sets that:

$$S_1 \ominus S_2 = (S_1 \cap \overline{S_2}) \cup (S_2 \cap \overline{S_1})$$

Because of since regular sets are closed under union, intersection, and complement.