- 1. (9 pts) Give context-free grammars that generate the following languages.
 - a. $L_1 = \{ w \in \{0, 1\}^* \mid w \text{ contains at least three 1s } \}$
 - b. $L_2 = \{ w \in \{0, 1\}^* | w = w^R \text{ and } |w| \text{ is even } \}$
 - c. $L_3 = \{ a^i b^j c^k | i, j, k \ge 0, \text{ and } i = j \text{ or } i = k \}$
- 2. (5 pts) Consider the set of terminals $T = \{a, b, (,), +, *, \varnothing\}$. Construct a context-free grammar G=(V,T,S,P) that generates all strings in T^* that are regular expressions over $\Sigma = \{a, b\}$. Use the grammar to derive the regular expression $(a+b)^*$.

Recall the rules of regular expressions.

- 1. \emptyset and each member of Σ is a regular expression.
- 2. If r_1 and r_2 are regular expressions, then so is r_1 r_2
- 3. If r_1 and r_2 are regular expressions, then so is $r_1 + r_2$.
- 4. If r_1 is a regular expression, then so is r_1^* .
- 5. If r_1 is a regular expression, then so is (r_1) .
- 6. Nothing else is a regular expression.
- 3. (9 pts) Consider the following grammar G = ({S, A}, {a, b}, S, P} where P is defined below

$$S \rightarrow SS \mid AAA \mid \lambda$$

$$A \rightarrow aA \mid Aa \mid b$$

- a. Describe the language generated by this grammar.
- b. Give a left-most derivation for the terminal string abbaba.
- c. Show that the grammar is ambiguous by exhibiting two distinct derivation trees for some terminal string.
- d. If this language is regular, give a regular grammar generating it. If the language is not regular, prove that it is not.
- 4. (3 pts) Find an s-grammar for L = $\{a^nb^{n+1}: n \ge 1\}$.
- 5. (4 pts) Let L = $\{a^nb^n : n \ge 0\}$
 - a. Show that L² is a context-free language.
 - b. Show that L* is a context-free language.