

Algorithm Theory - Assignment 1

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Problem 1

In this problem, the objective is to maximize the profit of selling products X and Y, while satisfying production constraints on these products. The profit can be computed as

$$\begin{aligned} \text{profit} &= \text{revenue} - \text{cost} \\ &= \text{revenue} - (\text{time}_{\text{machine}} \times \text{cost}_{\text{machine}} + \text{time}_{\text{craftsman}} \times \text{cost}_{\text{craftsman}}) \end{aligned}$$

We can compute profits for each of the products

$$\begin{aligned} \text{profit}(X) &= 200 - \left(\frac{15}{60} \times 100 + \frac{20}{60} \times 20 \right) = \frac{505}{3} \\ \text{profit}(Y) &= 300 - \left(\frac{20}{60} \times 100 + \frac{30}{60} \times 20 \right) = \frac{770}{3} \end{aligned}$$

And formulate the problem as a Linear Programming problem. Let n_X and n_Y be the number of products X and Y produced.

$$\begin{aligned} &\text{maximize } n_X \times \text{profit}(X) + n_Y \times \text{profit}(Y) \\ &\text{subject to } \begin{cases} n_X \times \text{time}_{\text{machine}}(X) + n_Y \times \text{time}_{\text{machine}}(Y) \leq 40 \times 60 \\ n_X \times \text{time}_{\text{craftsman}}(X) + n_Y \times \text{time}_{\text{craftsman}}(Y) \leq 35 \times 60 \\ n_X \geq 10 \\ n_X, n_Y \geq 0 \end{cases} \end{aligned}$$

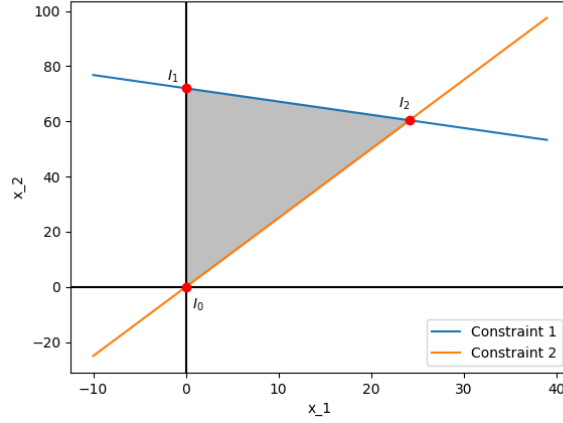
Which can be simplified as

$$\begin{aligned} &\text{maximize } n_X \times \frac{505}{3} + n_Y \times \frac{770}{3} \\ &\text{subject to } \begin{cases} 15 \times n_X + 20 \times n_Y \leq 2400 \\ 20 \times n_X + 30 \times n_Y \leq 2100 \\ n_X \geq 10 \\ n_Y \geq 0 \end{cases} \end{aligned}$$

Problem 2

a) Let n_A and n_B be the number of products A and B produced.

$$\begin{aligned} &\text{maximize } n_A \times 3 + n_B \times 5 \\ &\text{subject to } \begin{cases} 12 \times n_A + 25 \times n_B \leq 30 \times 60 \\ 2 \times n_B - 5 \times n_A \geq 0 \\ n_A, n_B \geq 0 \end{cases} \end{aligned}$$



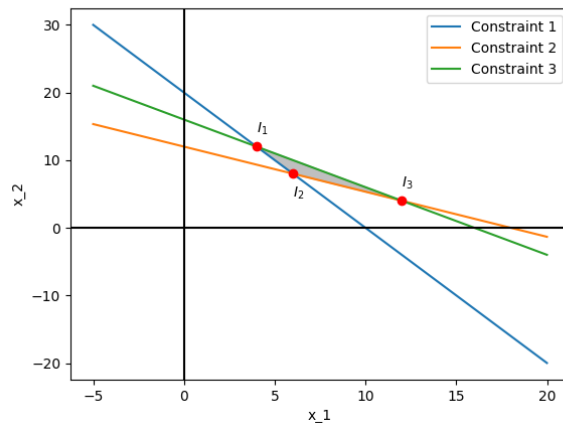
By graphing the feasible region, we see three candidate points. The first one $I_0 = (0, 0)$ can be trivially discarded as it leads to a profit of 0\$. Let's evaluate the objective function at the two other points. At $I_1 = (0, 72)$ the profit is 216\$. To compute the coordinates of the last point, we solve the following system. and at $I_2 = (24.2, 60.4)$ the profit is 374.5\$ which is the highest profit in this case. However, n_A and n_B are not integers. It is not specified in the exercise whether they should be or not, but as the problem is about a weekly production, we can assume that they do not need to be integer quantities as the weeks are continuous. If we wanted to restrict the problem to an integer solution, we should pick $n_A = 24$ and $n_B = 60$, leading to a profit of 372\$.

- b) By doubling the production capacity without modifying the other constraints, the resulting profit is doubled. The company should then pay less than 374.5\$ for renting an extra machine in order to stay profitable.

Problem 3

- a) We want to minimize the total cost of transportation, while respecting constraints on the flights. This problem can be formulated as the following linear programming problem. Let n_A and n_B be the number of flights flown with aircrafts A and B respectively.

$$\begin{aligned} & \text{minimize } n_A \times 10000 + n_B \times 12000 \\ & \text{subject to } \begin{cases} 30 \times n_A + 15 \times n_B \geq 300 \\ 500 \times n_A - 750 \times n_B \geq 9000 \\ n_A + n_B \leq 16 \\ n_A, n_B \geq 0 \end{cases} \end{aligned}$$



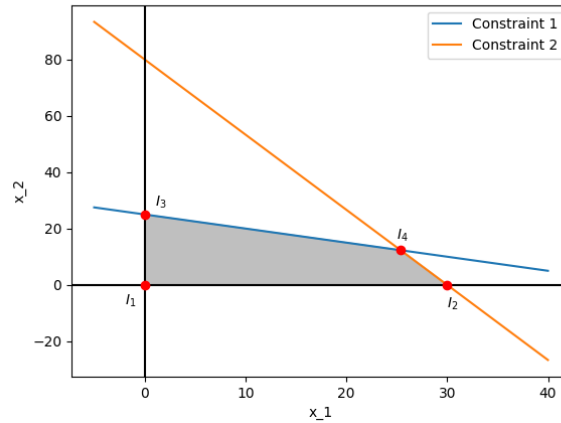
With the graph method, we can identify the three points delimiting the feasible region. We first determine their coordinates, and then evaluate the objective function at these points.

$$\begin{aligned}
 I_1 : \begin{cases} 30x + 15y = 300 \\ x + y = 16 \end{cases} &\Rightarrow \begin{cases} x = 4 \\ y = 12 \end{cases} \Rightarrow \text{cost} = 184000 \$ \\
 I_2 : \begin{cases} 30x + 15y = 300 \\ 500x + 750y = 9000 \end{cases} &\Rightarrow \begin{cases} x = 6 \\ y = 8 \end{cases} \Rightarrow \text{cost} = 156000 \$ \\
 I_3 : \begin{cases} x + y = 16 \\ 500x + 750y = 9000 \end{cases} &\Rightarrow \begin{cases} x = 12 \\ y = 4 \end{cases} \Rightarrow \text{cost} = 168000 \$
 \end{aligned}$$

The lowest cost is achieved by using 6 flights with A and 8 flights with B, for a total cost of 156000\$.

Problem 4

a) Graph method



$$\begin{aligned}
 I_1 : \begin{cases} x_1 = 0 \\ x_2 = 0 \end{cases} &\Rightarrow \text{objective} = 0 \\
 I_2 : \begin{cases} x_1 = 30 \\ x_2 = 0 \end{cases} &\Rightarrow \text{objective} = 90 \\
 I_3 : \begin{cases} x_1 = 0 \\ x_2 = 25 \end{cases} &\Rightarrow \text{objective} = 125 \\
 I_4 : \begin{cases} x_1 + 2x_2 = 50 \\ 8x_1 + 3x_2 = 240 \end{cases} &\Rightarrow \begin{cases} x_1 = \frac{330}{13} \\ x_2 = \frac{160}{13} \end{cases} \Rightarrow \text{objective} = \frac{1790}{13} \approx 137.7
 \end{aligned}$$

b) Simplex algorithm

We first convert the linear problem to its the slack form by introducing two slack variables s_1 and s_2 .

$$\begin{aligned} & \text{maximize } z = 3x_1 + 5x_2 \\ & \text{subject to } \begin{cases} s_1 = 50 - x_1 - 2x_2 \\ s_2 = 240 - 8x_1 - 3x_2 \\ x_1, x_2, s_1, s_2 \geq 0 \end{cases} \end{aligned}$$

A basic feasible solution is $(x_1, x_2, s_1, s_2) = (0, 0, 50, 240)$. However, this solution is not optimal as the objective function can be improved by picking a greater x_1 or x_2 . To do the first pivot, we choose x_2 as the entering variable because it has the highest coefficient in the objective function. To choose the leaving variable, we find the tightest constraint on x_2 . In the first constraint, x_2 is limited to 25, and in the second constraint it is limited to 80. Thus the leaving variable is s_1 , and we rewrite the problem with $x_2 = 25 - \frac{x_1}{2} - \frac{s_1}{2}$

$$\begin{aligned} & \text{maximize } z = 125 + \frac{x_1}{2} - \frac{5}{2}s_1 \\ & \text{subject to } s \begin{cases} x_2 = 25 - \frac{x_1}{2} - \frac{s_1}{2} \\ s_2 = 165 - \frac{13}{2}x_1 + \frac{3}{2}s_1 \\ x_1, x_2, s_1, s_2 \geq 0 \end{cases} \end{aligned}$$

The only remaining entering variable able to increase the objective function is x_1 . Constraint 1 limits it to 50 and constraint 2 limits it to $\frac{330}{13} \approx 25.4$. The second constraint being tighter than the first, we choose s_2 as leaving variable and rewrite the problem with $x_1 = \frac{330}{13} + \frac{3}{10}s_1 - \frac{2}{13}s_2$

$$\begin{aligned} & \text{maximize } z = \frac{1790}{13} - \frac{47}{20}s_1 - \frac{1}{26}s_2 \\ & \text{subject to } s \begin{cases} x_2 = \frac{160}{13} - \frac{8}{13}s_1 + \frac{1}{13}s_2 \\ x_1 = \frac{330}{13} + \frac{3}{10}s_1 - \frac{2}{13}s_2 \\ x_1, x_2, s_1, s_2 \geq 0 \end{cases} \end{aligned}$$

At this stage, the basic feasible solution of the objective function cannot be increased further because there is only subtractions of positive terms. Therefore, the optimal solution is $z = \frac{1790}{13} \approx 137.7$, obtained for the variables $x_1 = \frac{330}{13}$ and $x_2 = \frac{160}{13}$.

c) Programming method

Not wanting to run proprietary software, this problem was solved using GNU's GLPK. , the second is the resulting execution. This first image shows the program used.

```

from pulp import *

# problem definition
prob = LpProblem(sense=LpMaximize)

# variables
x_1 = LpVariable('x_1', lowBound=0)
x_2 = LpVariable('x_2', lowBound=0)

# objective function
prob += 3*x_1 + 5*x_2

# constraints
prob += x_1 + 2*x_2 <= 50
prob += 8*x_1 + 3*x_2 <= 240

GLPK().solve(prob)

# solution
for v in prob.variables():
    print(f'{v.name} = {v.varValue}')

print(f'objective = {value(prob.objective)}')

```

This second image shows the output of the execution.

```

$ python3 linprog.py
GLPSOL: GLPK LP/MIP Solver, v4.65
Parameter(s) specified in the command line:
--cpxlp /tmp/7ae8adc56e6b448a944b4eb315638c36-pulp.lp -o /tmp/7ae8adc56e6b448a944b4eb315638c36-pulp.sol
Reading problem data from '/tmp/7ae8adc56e6b448a944b4eb315638c36-pulp.lp'...
2 rows, 2 columns, 4 non-zeros
7 lines were read
GLPK Simplex Optimizer, v4.65
2 rows, 2 columns, 4 non-zeros
Preprocessing...
2 rows, 2 columns, 4 non-zeros
Scaling...
A: min|aij| = 1.000e+00 max|aij| = 8.000e+00 ratio = 8.000e+00
Problem data seem to be well scaled
Constructing initial basis...
Size of triangular part is 2
*   0: obj = -0.000000000e+00 inf = 0.000e+00 (2)
*   2: obj = -1.376923877e+02 inf = 0.000e+00 (0)
OPTIMAL LP SOLUTION FOUND
Time used: 0.0 secs
Memory used: 0.0 Mb (39693 bytes)
Writing basic solution to '/tmp/7ae8adc56e6b448a944b4eb315638c36-pulp.sol'...
x_1 = 25.3846
x_2 = 12.3077
Objective = 137.6923

```

We can see that with the three techniques, we get the same results. The optimal setting of variables to maximize the objective function is $x_1 = \frac{330}{13} \approx 25.4$ and $x_2 = \frac{160}{13} \approx 12.3$, for an objective value of $z = \frac{1790}{13} \approx 137.7$.