

# Week 01 - Lecture 2 Slides

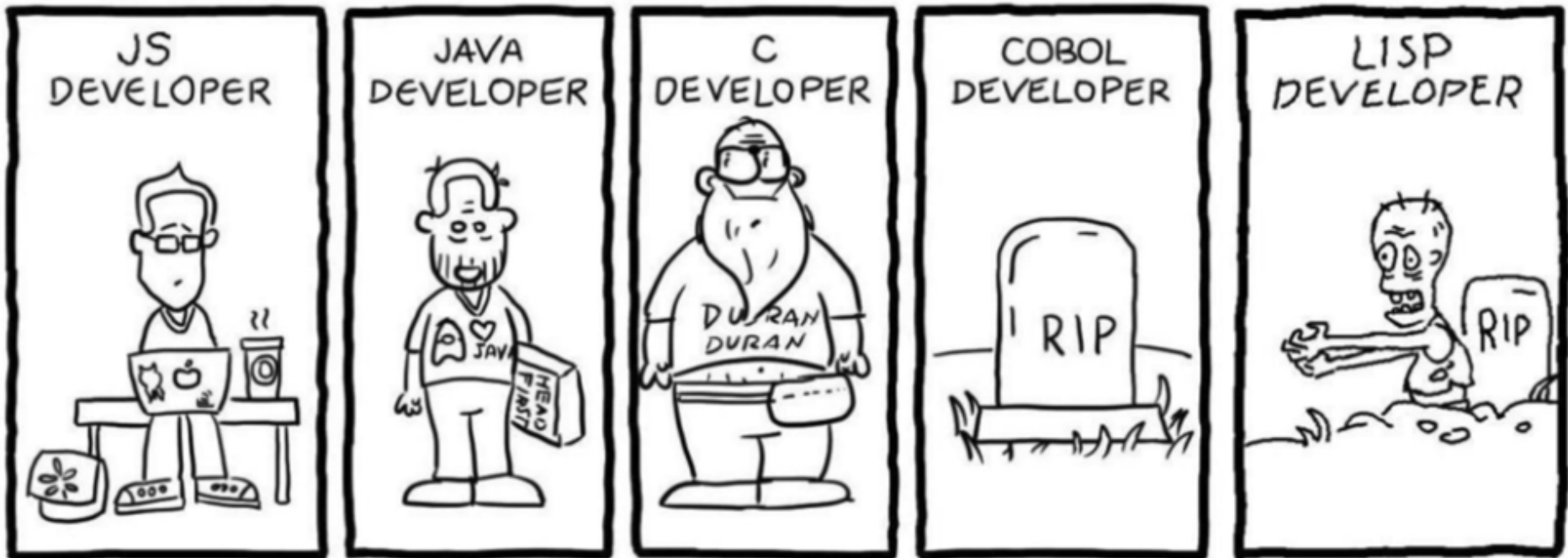
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## Lecture 2: A crash course on Lisp (cont.)



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Learning objectives:

By the end of this lecture you should be able to:

- Describe and write Lisp expressions that contain special forms
- Create programmer-defined functions
- Describe how functions are evaluated when applied to arguments

### **Lisp syntax and evaluation: special forms**

- In the previous lecture we learned about the simplest lisp expressions (aka *forms*):
  - constants
  - symbols
  - function applications
- How about ifs, loops, programmer-defined functions, etc.?

- There are 25 "special forms" in Lisp!
- What makes them "special"?
  - Each special form follows their own evaluation rule.

Let's first learn the five most used special forms.

## The most common special forms

- $(\text{QUOTE } v)$ , also written as  $'v$ , where  $v$  is any lisp form. It means: return  $v$  as data, without evaluating it. (See [HyperSpec](#) doc)

```
(quote (+ 2 3))  ;==>  (+ 2 3)
'(+ 2 3)        ;==>  (+ 2 3)
(+ 2 3)         ;==>  5      ; Notice here that the form (+ 2 3) is not quoted
```

- $(\text{AND } a_1 \cdots a_n)$  and  $(\text{OR } a_1 \cdots a_n)$ : the forms  $a_i$  are evaluated from left to right
  - AND: As soon as any form  $a_i$  evaluates to nil, AND returns NIL without evaluating the remaining forms. If all forms but the last evaluate to true values, AND returns the result produced by evaluating the last form.
  - OR: As soon as any form  $a_i$  evaluates to true, OR returns  $a_i$ 's value without evaluating the remaining forms. If all forms but the last evaluate to NIL values, OR returns the result produced by evaluating the last form.
- The following forms are interpreted as **false** in a predicate test:  $()$ ,  $'()$ ,  $'\text{NIL}$ ,  $\text{NIL}$
- Any value other than NIL or  $()$  is interpreted as **true**.

```
(and 1 2 3) ;==> 3
(and 0 1 nil (+ 3 4) 3) ;==> nil
(and 0 1 2 (* 3 3)) ;==> 9
(or 10 (* 3 3) 2 1) ;==> 10
(or (> 5 6) (< 3 4) (= 1 2)) ;==> T
```

## Special forms (cont.): IF

$(\text{IF } \text{testForm } \text{thenForm } [\text{elseForm}])$

**Note:** in the [meta-notation](#) above,  $[abc]$  means:  $abc$  is optional.

- Is a special form for conditional branching. Let's read its description and see examples at [HyperSpec](#).

```
(if nil 4 (* 3 2)) ;==> 6

(if (and 0 1 nil (+ 3 4) 3) ; Notice: indentation enhances readability
    3
    10) ;==> 10

(if (or 0 (* 3 3) 2 1)
    (+ 5 5)
    (* 10 10)) ;==> 10
```

**Homework:** explore two other lisp conditional expressions: [WHEN](#) and [UNLESS](#)

## Special forms (cont.): LAMBDA

- Used for creating a nameless function
- (LAMBDA  $(x_1 \cdots x_n) e$ ), where each  $x_i$  is a distinct symbol and  $e$  is a form. It evaluates to the function of  $n$  arguments which, when given arguments  $v_1, \cdots v_n$ , returns as its value the result of evaluating  $e$  in a context where the variable  $x_i$  is bound to  $v_i$ .

```
(lambda (x) (* x x)) ;==> #<FUNCTION (LAMBDA (X)) {535F802B}> - this is a special form
```

```
((lambda (x) (* x x)) 4) ;==> 16 - this is a function application form
```

```
((lambda (x y) (+ (* x y) x)) 3 1) ;==> 6 - this is also function application form
```

## Special forms (cont.): LET

$$(\text{LET } \underbrace{((x_1 e_1) \cdots (x_n e_n))}_{\text{Bindings}} e)$$

- first, all  $e_i$  are evaluated at once and their values bound, in parallel, to their respective  $x_i$
- then  $e$  is evaluated

This is the usual way Lisp programmers give variables temporary local bindings (see [HyperSpec](#) doc)

```
(let ((x 4)) (* x x))  ;==> 16

(let ((x 3)                ; Notice: indentation enhances readability
      (y 1))
  (+ (* x y) x))  ;==> 6

(let ((x 3)
      (y (* x x)))  <=== ERROR!!! 3 has not been bound to x yet!
  (+ x y))
```

## Special forms (cont.): LET\* and RUTILS:WITH

$$(\text{LET}^* ((x_1 e_1) \cdots (x_n e_n)) e)$$
$$(\text{RUTILS:WITH} ((x_1 e_1) \cdots (x_n e_n)) e)$$

Unlike LET, in LET\* (and in WITH)  $e_1$  is evaluated and its value bound to  $x_1$ , then  $e_2$  is evaluated and its value bound to  $x_2$ , and so on.

In a LET, all  $e_i$  are evaluated at once and their values bound, in parallel, to their respective  $x_i$ .

```
CL-USER> (let* ((a 2)          ; Notice: indentation enhances readability
                 (b 4)
                 (c -1)
                 (delta (- (* b b) (* 4 a c)))
                 (res (sqrt delta)))
  (/ (+ (- b) res) (* 2 a)))
0.22474492
```

- In the textbook, you may see RTL:WITH instead of LET\*.
- WITH is a utility macro defined in the RUTILS package available in our lab machines.
- It offers more flexibility than LET\* for creating bindings with elaborate data structures.

## Exercise

What will the evaluation of each form return?

10 ;==> ?

(+ 5 3 4) ;==> ?

(- 9 1) ;==> ?

(+ (\* 2 4) (- 4 6)) ;==> ?

(lambda (x y) (+ x y)) ;==> ?

((lambda (x y) (+ x y)) (+ 2 1) 4) ;==> ?

```
(let ((a 4)
      (b 5)
      (c 30))
  (+ (* a b) c)) ;==> ?
```

## Solution

10 ;==> 10

(+ 5 3 4) ;==> 12

(- 9 1) ;==> 8

(+ (\* 2 4) (- 4 6)) ;==> 6

(lambda (x y) (+ x y)) ;==> #<FUNCTION (LAMBDA (X Y)) {2276001B}>

((lambda (x y) (+ x y)) (+ 2 1) 4) ;==> 7

```
(let ((a 4)
      (b 5)
      (c 30))
  (+ (* a b) c)) ;==> 50
```

## The most common special forms (cont.)

So far, we have identified in Lisp some elements that must appear in many programs:

- Numbers and arithmetic operations are primitive data and procedures

- Nesting of forms provides a means of combining operations
- LET expressions associate names with values providing a limited means of abstraction

**Procedure definitions** provide us a much more powerful abstraction technique.

## The DEFUN special form

Syntax: (DEFUN *functionName* ( $v_1 \cdots v_n$ )  $\underbrace{form^*}_{\text{Function body}}$  ), where:

- $v_i$  are the parameters of the function
- $form^*$  means zero or more forms (see [HyperSpec](#) doc)

The DEFUN special form is for DEfining a FUNction.

```
CL-USER> (defun average (x y)
           "Returns the average of x and y"
           (/ (+ x y) 2.0))
```

- The DEFUN operator creates in memory a function/procedure identified by its name and defined by its body.

```
(average (* 2 2) (/ 10 2)) ;==> 4.5
(describe 'average) ;==>
```

AVERAGE names a compiled function:

...

Documentation:

Returns the average of x and y

Source form:

```
(LAMBDA (X Y)
  "Returns the average of x and y"
  (BLOCK AVERAGE (/ (+ X Y) 2.0)))
```

## Example

Given that the built-in function ODDP tests if its argument is odd,

```
(oddp 3) ;==> T
```

```
(oddp 4)  ;==> NIL
```

write a function MAKE-EVEN that makes an odd number even by adding one to it. If the input to MAKE-EVEN is already even, it should be returned unchanged. For example:

```
(make-even 4)  ;==> 4
```

```
(make-even 5)  ;==> 6
```

---

Good styles

```
(defun make-even (x)
  (if (oddp x) (+ x 1)
      x))
;; Alternative style
(defun make-even (x)
  (if (oddp x)
      (+ x 1)
      x))
```

Bad style

```
(defun make-even (x)
  (if (oddp x)
    (+ x 1)
    x))
```

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For more information about Lisp programming style, read the [Formatting section of this web page](#).

## Quiz: defining functions



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<https://bit.ly/3Z3OAhT>



Write definitions for the following functions:

- CUBE: multiplies its input by itself three times.
- ADD1: adds one to its input
- ADD2: adds two to its input; it uses `add1` as a helper
- MAX2: returns the max if its two arguments, both are numbers
- ONEMOREP: returns `T` if its first input is exactly one greater than its second input; it uses `add1` as a helper

## Solution

```
(defun cube (x)
  (* x x x))
```

```
(defun add1 (x)
  (+ x 1))
```

```
(defun add2 (x)
  (add1 (add1 x)))
```

```
(defun max2 (x y)
  (if (> x y) x y))
```

```
(defun onemorep (x y)
  (= x (add1 y)))
```