

# 实验一报告

## 一 移位密码

移位密码：将英文字母向前或向后移动一个固定位置。例如向后移动3个位置，即对字母表作置换（不分大小写）。

1	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
2																										
3	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C

设明文为：public keys,则经过以上置换就变成了：sxexlf nhbv。

如果将26个英文字母进行编码： $A \rightarrow 0, B \rightarrow 1, \dots, Z \rightarrow 25$ ，则以上加密过程可简单地写成：

- 明文： $m = m_1 m_2 \cdots m_n$ ,
- 密文： $c = c_1 c_2, \cdots c_n$ ，其中 $c_i \equiv (m_i + key) \pmod{26}, i = 1, 2, \dots, n$ .

### C++实现原理

```
1  #include <algorithm>
2  #include <cstdint>
3  #include <iostream>
4  #include <string>
5
6  #include "utils.hh"
7  int main(int argc, const char** argv)
8  {
9      std::cout << "Enter string: " << std::endl;
10     std::string input;
11     getline(std::cin, input);
12
13     const uint32_t distance = random_distance();
14
15     std::transform(input.begin(), input.end(), input.begin(), [distance](char
c) -> char {
16         if (c != ' ') {
17             return (char)((((std::tolower(c) - 'a' + distance) % 25 + 'a')));
18         } else {
19             return c;
```

```

20     }
21 });
22
23     std::cout << "Result: " << std::endl;
24     std::cout << input << std::endl;
25
26     return 0;
27 }

```

其工作原理可以概括如下：

1. 首先根据 Mersenne Twister 随机数种子均匀地生成一个随机密钥  
 $k \leftarrow \text{UniformRandom}(0, 25)$ :

```

1  uint32_t random_distance(void)
2  {
3      std::random_device rd;
4      std::mt19937 engine { rd() };
5      std::uniform_int_distribution<uint32_t> dist(0, 25);
6      return dist(engine);
7  }

```

2. 根据输入的字符串，先转换成小写字母，然后再根据公式进行代换：

$$c_i \equiv m_i + k \pmod{26}.$$

## 移位密码的攻击手段

很明显，这个密钥空间只有26，所以我们直接穷举即可，输出26个可能的字符串即可。这个过程和上述加密过程是对称的：

```

1     for (uint32_t i = 0; i < 26; i++) {
2         std::string guess;
3         std::transform(input.begin(), input.end(), std::back_inserter(guess),
4 [i](const char& c) -> char {
5             if (c == ' ') {
6                 return c;
7             } else {
8                 return (char)((((std::tolower(c) - 'a' + i) % 25 + 'a')));
9             }
10        });
11
12        std::cout << "Guess " << i << ": " << guess << std::endl;
13    }

```

## 二 单表置换密码

单表置换密码就是根据字母表的置换对明文进行变换的方法，例如，给定置换

1	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
2																										
3	H	K	W	T	X	Y	S	G	B	P	Q	E	J	A	Z	M	L	N	O	F	C	I	D	V	U	R

那么给定明文 `public keys`，则有密文：`mckebw qxuo`。单表置换实现的一个关键问题是关于置换表的构造。置换表的构造可以有各种不同的途径，主要考虑的是记忆的方便。如使用一个短语或句子，删去其中的重复部分，作为置换表的前面的部分，然后把没有用到的字母按字母表的顺序依次放入置换表中。

我们可以定义一个带密钥的伪随机置换函数  $\text{PRP}$  来做到这个置换密码， $\text{arr}' \leftarrow \text{PRP}_k(\text{arr})$ ，其中  $k \leftarrow \text{KeyGen}(1^\lambda)$  ( $\lambda$  为安全参数)。

### C++实现

```

1  #include <iostream>
2  #include <string>
3  #include <random>
4  #include <vector>
5
6  static std::string substitution_table = "abcdefghijklmnopqrstuvwxyz";
7
8  int main(int argc, const char** argv)
9  {

```

```

10     std::cout << "Enter string: " << std::endl;
11     std::string str;
12     getline(std::cin, str);
13
14     // Do the permutation.
15     std::random_device rd;
16     std::mt19937 engine { rd() };
17     std::shuffle(substitution_table.begin(), substitution_table.end(),
engine);
18
19     std::transform(str.begin(), str.end(), str.begin(), [](char c) -> char {
20         if (c != ' ') {
21             return substitution_table[std::tolower(c) - 'a'];
22         } else {
23             return c;
24         }
25     });
26
27     std::cout << "Result: " << std::endl << str << std::endl;
28     return 0;
29 }

```

## 样例输出

```

1 Enter string:
2 jsdiajdisjaijdsaiisdaiasd
3
4 Guess 0: jsdiajdisjaijdsaiisdaiasd
5 Guess 1: ktejbkejtkbjketbjjtebjbte
6 Guess 2: lufkclfkulcklfuckkufckcuf
7 Guess 3: mvgldmgldvmdlmgvdlldvgldvg
8 Guess 4: nwhmenhmnemnhwemmhemewh
9 Guess 5: oxinfoinxofnoixfnxfnfxi
10 Guess 6: pyjogppjoypgopjygooyjgogyj
11 Guess 7: qakphqkpaqhpqkahppakhphak
12 Guess 8: rblqirlqbriqrlbiqqbliqibl
13 Guess 9: scmrjsmrscjrsmcjrrcmjrjcm
14 Guess 10: tdnsktnsdtkstndkssdnkskdn
15 Guess 11: ueotluoteultuoeltteoltleo
16 Guess 12: vfpumvpufvmuvpfmuufpmumfp
17 Guess 13: wgqvnwqvgwnvwqgnvvgqnvngq
18 Guess 14: xhrwoxrwhxowxrhowwhrowohr

```

```
19 Guess 15: yisxpysxiypxysipxxispxpis
20 Guess 16: ajtyqatyjaqyatjqyyjtqyqjt
21 Guess 17: bkuarbuakbrabukraakurarku
22 Guess 18: clvbscvblcsbcvlsbblvsbslv
23 Guess 19: dmwctdwcmtcdwmtccmwtctmw
24 Guess 20: enxduexdneudexnuddnxudunx
25 Guess 21: foyevfyeofvefyoveeoyvevoy
26 Guess 22: gpafwgafpgwfgapwffpawfwpaw
27 Guess 23: hqbgxhbgqhxghbqxggqbxgxqb
28 Guess 24: irchyichriyhicryhhrcyhyrc
29 Guess 25: jsdiajdisjaijdsaiisdaiasd
```

### 三 频率分析来破译单表置换密码

在单表置换密码中，由于置换表字母组合方式有 $26!$ 种，约为 $4.03 \times 10^{26}$ 种组合，所以采用穷举密钥的方法不是一种最有效的方法。对单表置换密码最有效的攻击方法是利用自然语言的使用频率：单字母、双字母组/三字母组、短语、词头/词尾等。我们考虑英文。

**短单词(small words):**在英文中只有很少几个非常短的单词。因此，如果在一个加密的文本中可以确定单词的范围，那么就能得出明显的结果。一个字母的单词只有a和I。如果不计单词的缩写，在从电子邮件中选取500k字节的样本中，只有两个字母的单词仅出现35次，而两个字母的所有组合为 $26 \times 26 = 676$ 种。而且，还是在那个样本中，只有三个字母的单词出现196次，而三个字母的所有组合为 $26 \times 26 \times 26 = 17576$ 种。

**常用单词(common words):**再次分析500k字节的样本，总共有5000多个不同的单词出现。在这里，9个最常用的单词出现的总次数占总单词数的21%，20个最常用的单词出现的总次数占总单词数的30%，104个最常用的单词占50%，247个最常用的单词占60%。样本中最常用的9个单词占总词数的百分比为：

```
1 the 4.65 to 3.02 of 2.61 I 2.2 a 1.95
2
3 and 1.82 is 1.68 that 1.62 in 1.57
```

**字母频率(character frequency):**在1M字节旧的电子文本中，对字母“A”到“Z”（忽略大小写）分别进行统计。发现近似频率（以百分比表示）：

```

1 e 11.67 t 9.53 o 8.22 i 7.81 a 7.73 n 6.71 s 6.55
2
3 r 5.97 h 4.52 l 4.3 d 3.24 u 3.21 c 3.06 m 2.8
4
5 p 2.34 y 2.22 f 2.14 g 2.00 w 1.69 b 1.58 v 1.03
6
7 k 0.79 x 0.30 j 0.23 q 0.12 z 0.09

```

从该表中可以看出，最常用的单字母英文是e和t，其他字母使用频率相对来说就小得多。这样，攻击一个单表置换密码，首先统计密文中最常出现的字母，并据此猜出两个最常用的字母，并根据英文统计的其他特征（如字母组合等）进行试译。

## 破译一段密文

```

1 SIC GCBSPNA XPMHACQ JB GPYXSMEPNXIY JR SINS MF SPNBRQJSSJBE JBFMPQNSJMB FPMQ N
  XMJBS N SM N XMJBS H HY QCNBR MF N XMRRJHAY JBRCGZPC GINBBCA JB RZGI N VNY
  SINS SIC MPJEJBNA QCRRNEC GNB MBAY HC PCGMTCPD HY SIC PJEISFZA PCGJXJCBSR SIC
  XNPSJGJXNBSR JB SIC SPNBRNGSJMB NPC NAJGC SIC MPJEJBNSMP MF SIC QCRRNEC HMH
  SIC PCGCJTCP NBD MRGNP N XMRRJHAC MXXMBCBS VIM VJRICR SM ENJB ZBNZSIMPJOCD
  GMBSPMA MF SIC QCRRNEC

```

1. 先编写一段朴素的C++程序来统计这个语言中每个字母的出现频率：

```

1 // Count occurrence.
2 uint32_t total = 0;
3 std::for_each(cipher.begin(), cipher.end(), [&freq, &total](const char&
  c) {
4     if (std::isalpha(c)) {
5         freq[c - 'a']++;
6         total++;
7     }
8 });
9 // Compute frequency.
10 std::vector<std::pair<char, double>> freq_cipher; // Frequency of the
  cipher.
11 char c = 'a';
12 std::transform(freq.begin(), freq.end(), std::back_inserter(freq_cipher),
  [total, &c](const uint32_t& count) {
13     std::cout << c << ", " << count * 1.0 / total << std::endl;
14     return std::make_pair(c++, count * 1.0 / total);
15 });

```

```

16 // Sort the frequency table.
17 std::sort(freq_cipher.begin(), freq_cipher.end(),
18     [](const std::pair<char, double>& lhs, const std::pair<char, double>&
19         rhs) -> bool {
19         return lhs.second > rhs.second;
20     });

```

经过上述的统计可以初步获取每个字母的出现频次排序为：

```

1 c, s, n, m, b, j, p, r, i, g, x, a, e, h, q, y, f, z, d, v, t, o, u, w, l,
  k

```

初步拟合的结果为：

```

1 the leatsou dsimuep na lsydticsodhy nr thot if tsoarpnttnac nafispotnia
  fsip o dinat o ti o dinat m my peoar if o dirrnmy narelgse lhoaau na
  rglh o boy thot the isncnaou perroce loa iauy me selivesew my the snchtfgu
  selndneatr the dostnlndoatr na the tsoaroltnia ose ounle the isncnaotis if
  the perroce mim the selenves oaw irlos o dirrnmue iddiaeat bhi bnrher ti
  cona gaagthisnkew liatsiu if the perroce

```

已经有一些词汇被正确识别了，接下来分析还没有分析正确的词汇。

2. 注意单个字母的单词只有a和l，因此单个字母的单词只能根据频率匹配这两个单词，故o匹配a，所以再次精细化得到结果：

```

1 the leatsau dsimuep na lsydticsadhy nr that if tsaarpnttnac nafispatnia
  fsip a dinat a ti a dinat m my peaar if a dirrnmy narelgse lhaaaau na
  rglh a bay that the isncnaau perrace laa iauy me selivesew my the snchtfgu
  selndneatr the dastnlndaatr na the tsaaraltnia ase aunle the isncnaatis if
  the perrace mim the selenves aaw irlas a dirrnmue iddiaeat bhi bnrher ti
  cana gaagthisnkew liatsiu if the perrace

```

3. 接下来分析两个字母的单词：

```

1 if:3 na:3 nr:1 ti:2 my:2 me:1

```

两个字母的单词也比较少，所以根据频率和逻辑可以推测

```
1 na -> in
2 if -> of
3 ti -> to
4 nr -> is
5 my -> by
```

可以再次精细化结果得到

```
1 the lentsau dsobuep in lsydticsadhy is that if tsanspittinc infospation
  fsop a doint a to a doint b by peans of a dossibuy inselgse lhanneu in
  sglh a bay that the osicinau pessace lan onuy be selovesew by the richtfgu
  selndients the dastilidants in the tsaasaltion ase auile the osiciaatos of
  the pessace bob the seleives aaw oslas a dossimue oddonent bhi bishes to
  cain gaagthosikew liatsou of the pessace
```

4. 此时很多文字已经可以辨认了，例

```
1 bishes -> wishes b -> w
2 infospation -> information s -> r, p -> m, c -> g
```

继续恢复：

```
1 the lentrau drobuep in lrydtogradhy is that if transmitting information
  from a doint a to a doint b by means of a dossibuy inselgre lhanneu in
  sglh a way that the originau message lan onuy be reloverew by the rightfgu
  relidients the dartilidants in the transaction are auile the originator of
  the message bob the releiver arw oslar a dossnmue oddonent who wishes to
  gain gnagthirikew lontrou of the message
```

5. 接下来可以继续认清单词：

```
1 onuy -> only u -> l
2 rightfgu -> rightful g -> u
3 dossibuy -> possibly d -> p
```

继续恢复：



- 1 the lentral problep in lryptography is that if transmitting information from a point a to a point b by means of a possibly inselgre channel in sglh a way that the original message lan only be reloverew by the rightful relipients the partilipants in the transaction are alile the originator of the message bob the releiver arw oslar a possimle opponent who wishes to gain gnagthorikew lontrol of the message

## 6. 基本上就已经清楚了：

- 1 lcryptography -> cryptography l -> c
- 2 problp -> problem p -> m
- 3 relocerew -> recovered w -> d

- 1 the central problem in cryptography is that if transmitting information from a point a to a point b by means of a possibly insecgre channel in sgch a way that the original message can only be recovered by the rightful recipients the participants in the transaction are alice the originator of the message bob the receiver and oscar a possimle opponent who wishes to gain gnagthoriked control of the message

## 7. 最后有

- 1 insecgre -> insecure g -> u
- 2 possimle -> possible m -> b
- 3 unauthoriked -> unauthorized k -> z

- 1 the central problem in cryptography is that if transmitting information from a point a to a point b by means of a possibly insecure channel in such a way that the original message can only be recovered by the rightful recipients the participants in the transaction are alice the originator of the message bob the receiver and oscar a possimle opponent who wishes to gain unauthorized control of the message.