哈尔滨工业大学(深圳)

《密码学基础》实验报告

RSA 密码算法实验

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一、实验步骤

请说明你的字符分组方式,以及关键的算法例如扩展欧几里德,素数检测,快速幂等函数的实现。

字符分组方式: 两个字符的 ascii 码拼接为一个 6 位的十进制数字作为一个分组

. 具体来说,对于加密,先从明文取两个字符,利用 ord 取 ascii 码,再通过 ord0 * 1000 + ord1 实现拼接:

```
m_group = ord(pair[0]) * 1000 + ord(pair[1])
```

对于孤立的单个字符,使用\0 补全:

```
m group = ord(plaintext[i]) * 1000
```

对于解密,也有类似操作,不过使用 chr 将 ascii 转为字符

```
decrypted_pair = mod_pow(num, d, n)
dechar_1 = chr(decrypted_pair // 1000)
dechar_2 = chr(decrypted_pair % 1000)
```

关键算法:

欧几里得算法求 gcd: 根据辗转相除法的定义, 最终 b==0 时 a 就是最后的 gcd, 也就是最后一步前的 b 必能整除 a(最极端情况下 b=1)

```
def gcd(a, b):
    while b:
        a, b = b, a % b
    return a
```

素数检测:使用 Miller-Rabin 算法检测素数,思路就是先做初始检查 (小边界和提高效率),之后分解 n-1=2^r * d 的形式,进行 k=10 轮平 方测试,每一轮测试中,依次检验平方模是否为 n - 1,如果没有(for 循环正常结束),则调用 for-else 语句的 else 返回合数的判断,否则算 暂时通过,进行下一轮的测试;

```
return False
if n <= 3:
   return True
if n % 2 == 0:
   return False
r, d = 0, n - 1
while d % 2 == 0:
   d //= 2
for in range(k):
   x = pow(a, d, n)
   if x == 1 \text{ or } x == n - 1:
       continue
   for in range(r - 1):
       if x == n - 1:
           break
    else:
       return False
return True
```

扩展欧几里德算法求逆: 思 n 路是每一轮的辗转相除需要得到 q 整除结果,然后利用 x(n)=x(n-2)-qx(n-1) 迭代即可,不过需要注意初始设置 x 的值即可(此处假定 a 是更大的那个数,故 x0=0, x1=1,其中 x0 比 x1 更晚更新)

```
def mod_inv(a, m):
    m0, x0, x1 = m, 0, 1
    if m == 1:
        return 0
    while a > 1:
        q = a // m
        m, a = a % m, m
```

```
x0, x1 = x1 - q * x0, x0
if x1 < 0:
    x1 += m0
return x1</pre>
```

快速幂函数: 从末位开始解析,每解析一个 bit 就让 b 平方模一次,同时让 exp 右移 1 位,实现快速幂计算

```
def mod_pow(b, exp, mod):
    res = 1
    b = b % mod # b < mod

while exp > 0:
    if (exp % 2) == 1:
        res = (res * b) % mod

    exp = exp // 2
    b = (b * b) % mod

return res
```

算法总体框架:

- 1. 生成 p、q, 计算 n=pq、phi(n)=(p-1)*(q-1)
- 2. 生成公钥 e, mod inv 计算模 phi(n)逆元私钥 d
- 3. 对每个明文分组,加密
- 4. 解密,解分组
- 5. 输出结果

以下是实现: 生成密钥

```
def generate_keys():
    p = gen_prime(64)
    q = gen_prime(64)
    n = p * q
    phi_n = (p - 1) * (q - 1)
    e = gen_prime(64)
    d = mod_inv(e, phi_n)

return {
        'p': p,
        'q': q,
        'n': n,
        'e': e,
        'd': d,
        'phi(n)': phi_n
}
```

分组加密和解密:

```
def encrypt(plaintext, n, e):
   ciphertext = []
   start_time = time.time()
    for i in range(0, len(plaintext), 2):
        if i + 1 < len(plaintext):</pre>
            pair = plaintext[i:i + 2]
            m_group = ord(pair[0]) * 1000 + ord(pair[1])
            ciphertext.append(mod_pow(m_group, e, n))
        else:
           m_group = ord(plaintext[i]) * 1000
           ciphertext.append(mod_pow(m_group, e, n))
   end_time = time.time()
   encryption_time = end_time - start_time
   return ciphertext, encryption_time
def decrypt(ciphertext, n, d):
   decrypted = ''
   start_time = time.time()
    for num in ciphertext:
        decrypted_pair = mod_pow(num, d, n)
        dechar_1 = chr(decrypted_pair // 1000)
        dechar_2 = chr(decrypted_pair % 1000)
        decrypted += dechar_1 + dechar_2
   end_time = time.time()
    decryption_time = end_time - start_time
   return decrypted, decryption_time
```

结果输出部分:

```
with open('lab2-Plaintext.txt', 'r') as file:
   plaintext = file.read()
keys = generate_keys()
with open('public_key.txt', 'w') as pub_file:
   pub_file.write(f"n: {keys['n']}\n")
   pub_file.write(f"e: {keys['e']}\n")
with open('private_key.txt', 'w') as priv_file:
   priv_file.write(f"d: {keys['d']}\n")
ciphertext, encryption_time = encrypt(plaintext, keys['n'], keys['e'])
with open('ciphertext.txt', 'w') as file:
   for num in ciphertext:
       file.write(f"{num}\n") # 每个密文占—行
decrypted_text, decryption_time = decrypt(ciphertext, keys['n'], keys['d'])
  print("p:", keys['p'])
  print("q:", keys['q'])
  print("Public key written to 'public_key.txt'")
  print("Private key written to 'private_key.txt'")
  print("Ciphertext written to 'ciphertext.txt'")
  print("Decrypted text:", decrypted_text)
  print("Encryption time (seconds):", encryption_time)
  print("Decryption time (seconds):", decryption_time)
```

二、实验结果与分析

程序正确运行的结果截图,包括p, q, n,e, d, ϕ (n)这些参数的值,可输出加密解密时间作为时长参考。

```
[♠ Ionic from □ Terence_Lee][▼ 0.8835][● RAM: 9/1568][■ Thursday at 11:28:44 PM]

| [E:\Code\Github-Projects\HITsz-Cryptography-Project\code\Lab2]

| LA python \Less.py
| 10/1316218974401861
| q: 174855130480138978401861
| q: 17485513048013897840
| Public key written to 'public_key.txt'
| Private key written to 'private_key.txt'
| Ciphertext written to 'ciphertext.txt'
| Decrypted text: 2002 A.M. TURING AWARD. RSA, an acronym for Rivest, Shamir and Adleman, uses algorithmic number theory to provide an efficient realization
of a public-key cryptosystem, a concept first envisioned theoretically by Whitfield Diffie, Martin Hellman and Ralph Merkle. RSA is now the most widely used encryption method, with applications throughout the Internet to secure on-line transactions. It has also inspired breakthrough work in both theoretical to computer science and mathematics.
| Encryption time (seconds): 0.085001546952392578
| Decryption time (seconds): 0.08500154952392578
| Decryption time (seconds): 0.08502710975378418
| Decryption time (seconds): 0.08502710975378418
| Decryption time (seconds): 0.08502710975378418
```

可见,两次实验的 p, q 值均在 64bit 数字的范围,加密的结果是 Ciphertext,解密结果与原文相同,算法运行正常;加密解密的的时间都极短(解密 0.01s 左右)

其中, 压缩包内附有密文, 可查看

三、总结

- 1、实验过程中遇到的问题有哪些? 你是怎么解决的?
- 1) 一开始扩展欧几里得算法的迭代写反了,后续通过单元测试调初始值和数的顺序解决;
- 2) k 较低(一开始设置的是 3)存在解密崩掉成乱码的可能性,其原因可能是 p, q 不是素数但是经过了 is prime 的考验,后续增加 k 的轮次再也没出现过;
 - 2、关于本实验的意见或建议。

建议加入线上测试和验收平台,就像 MIT/Stanford 实验那样