Kempe Compiler & Language Manual

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Introduction

Kempe is a stack-based language, and kc is a toy compiler for x86_64.

Installing kc

```
First, install cabal and GHC. Then: cabal install kempe
```

This provides kc, the Kempe compiler.

kc requires NASM, an x86_64 assembler.

Editor Integration

```
A vim plugin is available.
```

To install with vim-plug:

```
Plug 'vmchale/kempe' , { 'rtp' : 'vim' }
```

Kempe Language

Types

Kempe has a stack-based type system. So if you see a type signature:

```
next : Word -- Word Word
```

that means that the stack must have a Word on it for next to be invoked, and that it will have two Words on the stack after it is invoked.

Polymorphism

Kempe allows polymorphic functions. So we can define:

```
id : a -- a =: []
```

The Kempe typechecker basically works though unification is slow.

Literals

Integer literals have type -- Int.

Positive literals followed by a u have type -- Word, e.g. 1u.

Negative integer literals are indicated by an underscore, $_$, i.e. $_1$ has type -- Int.

Builtins

The Kempe compiler has a few builtin functions that you can use for arithmetic and for shuffling data around. Many of them are familiar to stack-based programmers:

```
    swap : a b -- b a
    drop : a --
For arithmetic:
    + : Int Int -- Int
    * : Int Int -- Int
    - : Int Int -- Int
    / : Int Int -- Int
    / : Int Int -- Int
    / : Int Int -- Int
    % : Int Int -- Int
    % : Int Int8 -- Int
    << : Int Int8 -- Int
    xori : Int Int -- Int
    * * : Word Word -- Word
    *~ : Word Word -- Word</pre>
```

• dup : a -- a a

<< : Word Int8 -- Wordxoru : Word Word -- Wordpopcount : Word -- Int

/~: Word Word -- Word
 %~: Word Word -- Word
 >>~: Word Int8 -- Word

= : Int Int -- Bool
> : Int Int -- Bool
< : Int Int -- Bool
!= : Int Int -- Bool

!= : Int Int -- Bool<= : Int Int -- Bool>= : Int Int -- Bool

& : Bool Bool -- Bool|| : Bool Bool -- Boolxor : Bool Bool -- Bool

• ~ : Int -- Int

There is one higher-order construct, dip, which we illustrate by example:

If Blocks

If-blocks are atoms which contain two blocks of atoms on each arm. If the next item on the stack is **True**, the first will be executed, otherwise the second.

Sum Types

Kempe supports sum types, for instance:

```
type Either a b { Left a | Right b }
Note that empty sum types such as
type Void {}
are not really supported.
```

Pattern Matching

Sum types are taken apart with pattern matching, viz.

Note that pattern matches in Kempe must be exhaustive.

FFI

Kempe can call into C functions. Suppose we have

```
int rand(void);
```

Then we can declare this as:

```
rand : -- Int
=: $cfun"rand"
```

And rand will be available as a Kempe function.

Recursion

kc optimizes tail recursion.

Non-Features

Kempe is missing a good many features, such as:

- Floats
- Dynamically sized data types
- Strings
- Recursive data types
- Pointers
- Operator overloading

Programming in Kempe

Invoking the Compiler

kc cannot be used to produce executables. Rather, the Kempe compiler will produce .o files which contain functions.

Kempe functions can be exported with a C ABI:

%foreign cabi fac

This would be called with a C wrapper like so:

```
#include <stdio.h>
extern int fac(int);
int main(int argc, char *argv[]) {
```

```
printf("%d", fac(3));
}
```

Unlike the frontend and type checker, the backend is dodgy.

Internals

Kempe maintains its own stack and stores the pointer in rbp.

Kempe procedures do not require any registers to be preserved across function calls.

C Calls

When exporting to C, kc generates code that initializes the Kempe data pointer (rbx). Thus, one should avoid calling into Kempe code too often!

Note that the Kempe data pointer is static, so calling different Kempe functions in different threads will fail unpredictably.

Kempe ABI

Sum types have a guaranteed representation so that they can be used from other languages.

Consider:

```
type Param a b c
{ C a b b
| D a b c
}
```

Kempe types always have the same size; a value constructed with $\tt C$ will occupy the same number of bytes on the stack as a value constructed with $\tt D$.

So, for instance

```
mkD : Int8 Int Int8 -- (((Param Int8) Int) Int8)
=: [ D ]
```

will pad the value with 7 bytes, as a (((Param Int8) Int) Int8) constructed with C would be 7 bytes bigger.

Examples

Splitmix Pseudorandom Number Generator

The generator in question comes from a recent paper.

Implementation turns out to be quite nice thanks to Kempe's multiple return values:

%foreign kabi next

Note that 30i8 is an Int8 literal; shifts take an Int8 as the exponent.

Compare this C implementation:

```
#include <stdint.h>

// modified to have ""multiple return"" with destination-passing style
uint64_t next(uint64_t x, uint64_t* y) {
    uint64_t z = (x += 0x9e3779b97f4a7c15);
    z = (z ^ (z >> 30)) * 0xbf58476d1ce4e5b9;
    z = (z ^ (z >> 27)) * 0x94d049bb133111eb;
    *y = x;
    return z ^ (z >> 31);
}

GCD

gcd : Int Int -- Int
    =: [ dup 0 =
        if( drop
            , dup dip(%) swap gcd )
```

Mutual Recursion

kc supports mutual recursion:

```
not : Bool -- Bool
   =: [
    { case
       | True -> False
        | _ -> True
   }
]
odd : Int -- Bool
   =: [ dup 0 =
           if( drop False
             , - 1 even )
       ]
even : Int -- Bool
     =: [ dup 0 =
           if( drop True
             , - 1 odd )
       ]
```