

CS121 Data Structures

Positional Lists

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Lists

Indices are not good for describing positions within a linked list because we would need to traverse the list incrementally from its beginning or end, counting elements along the way.

Our goal is to design an ADT that provides a way to refer to elements anywhere in a sequence, and to perform arbitrary insertions and deletions.

A **position** formalizes the intuitive notion of the 'location' of an element relative to others in the list

The Position Abstract Data Type

A position supports the following single method:

`getElement()`: Returns (a reference to) the element stored at this position

A position acts as a marker within a positional list

A position p , which is associated with some element e in a list L , does not change, even if the index of e changes in L due to insertions or deletions elsewhere in the list

The position p does not change if we replace the element e stored at p with another element

A position becomes invalid if that position (and its element) are explicitly removed from the list

The Positional List Abstract Data Type

A **positional list** is a collection of positions, each storing an element

The positional list ADT supports **accessor** methods:

`first()`: Returns the position of the first element of L (or `null` if empty)

`last()`: Returns the position of the last element of L (or `null` if empty)

`before(p)`: Returns the position of L immediately before position p (or `null` if p is the first position)

`after(p)`: Returns the position of L immediately after position p (or `null` if p is the last position)

`isEmpty()`: Returns `true` if list L does not contain any elements

`size()`: Returns the number of elements in list L

An error occurs if a position p , sent as a parameter to a method, is not a valid position for the list

The Positional List Abstract Data Type (cont'd)

The positional list ADT supports **update** methods:

`addFirst(e)`: Inserts a new element e at the front of the list, returning the position of the new element

`addLast(e)`: Inserts a new element e at the back of the list, returning the position of the new element

`addBefore(p, e)`: Inserts a new element e in the list, just before position p , returning the position of the new element

`addAfter(p, e)`: Inserts a new element e in the list, just after position p , returning the position of the new element

`set(p, e)`: Replaces the element at position p with element e , returning the element formerly at position p

`remove(p)`: Removes and returns the element at position p in the list, invalidating the position

For a nonempty list, `addFirst(e)` and `addBefore(first(), e)`, `addLast(e)` and `addAfter(last(), e)` are pairwise equivalent

Traversing the List

`first()` and `last()` methods return the associated *positions*, not *elements* (in contrast to the Deque ADT)

First element accessed by: `first().getElement()`

Traverse and print each element of a list, named `guests`, that stores string elements:

```
1 Position<String> cursor = guests.first( );
2 while (cursor != null) {
3     System.out.println(cursor.getElement( ));
4     cursor = guests.after(cursor);           // advance to the next position (if any)
5 }
```

Using the convention of returning `null` references, the above code fragment works correctly on an empty `guests` list

Example

Method	Return Value	List Contents
<code>addLast(8)</code>	p	$(8p)$
<code>first()</code>	p	$(8p)$
<code>addAfter(p, 5)</code>	q	$(8p, 5q)$
<code>before(q)</code>	p	$(8p, 5q)$
<code>addBefore(q, 3)</code>	r	$(8p, 3r, 5q)$
<code>r.getElement()</code>	3	$(8p, 3r, 5q)$
<code>after(p)</code>	r	$(8p, 3r, 5q)$
<code>before(p)</code>	null	$(8p, 3r, 5q)$
<code>addFirst(9)</code>	s	$(9s, 8p, 3r, 5q)$
<code>remove(last())</code>	5	$(9s, 8p, 3r)$
<code>set(p, 7)</code>	8	$(9s, 7p, 3r)$
<code>remove(q)</code>	"error"	$(9s, 7p, 3r)$

The Position API

```
1 public interface Position<E> {  
2     /**  
3      * Returns the element stored at this position.  
4      *  
5      * @return the stored element  
6      * @throws IllegalStateException if position no longer valid  
7      */  
8     E getElement( ) throws IllegalStateException;  
9 }
```


The Positional List API

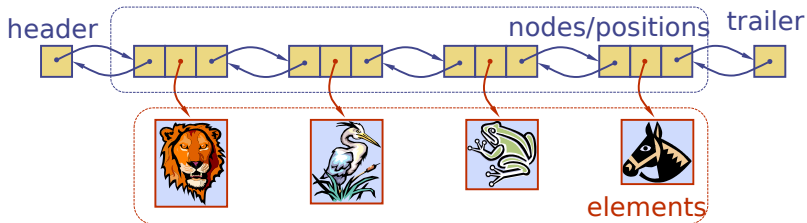
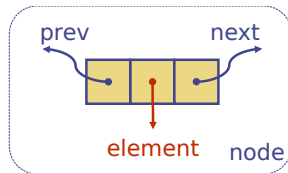
```
1  /** An interface for positional lists. */
2  public interface PositionalList<E> {
3
4      /** Returns the number of elements in the list. */
5      int size( );
6
7      /** Tests whether the list is empty. */
8      boolean isEmpty( );
9
10     /** Returns the first Position in the list (or null, if empty). */
11     Position<E> first( );
12
13     /** Returns the last Position in the list (or null, if empty). */
14     Position<E> last( );
15
16     /** Returns the Position immediately before Position p (or null, if p is first). */
17     Position<E> before(Position<E> p) throws IllegalArgumentException;
18
19     /** Returns the Position immediately after Position p (or null, if p is last). */
20     Position<E> after(Position<E> p) throws IllegalArgumentException;
21
```

The Positional List API (cont'd)

```
21
22  /** Inserts element e at the front of the list and returns its new Position. */
23  Position<E> addFirst(E e);
24
25  /** Inserts element e at the back of the list and returns its new Position. */
26  Position<E> addLast(E e);
27
28  /** Inserts element e immediately before Position p and returns its new Position. */
29  Position<E> addBefore(Position<E> p, E e) throws IllegalArgumentException;
30
31  /** Inserts element e immediately after Position p and returns its new Position. */
32  Position<E> addAfter(Position<E> p, E e) throws IllegalArgumentException;
33
34  /** Replaces the element stored at Position p and returns the replaced element. */
35  E set(Position<E> p, E e) throws IllegalArgumentException;
36
37  /** Removes the element stored at Position p and returns it (invalidating p). */
38  E remove(Position<E> p) throws IllegalArgumentException;
39  }
```

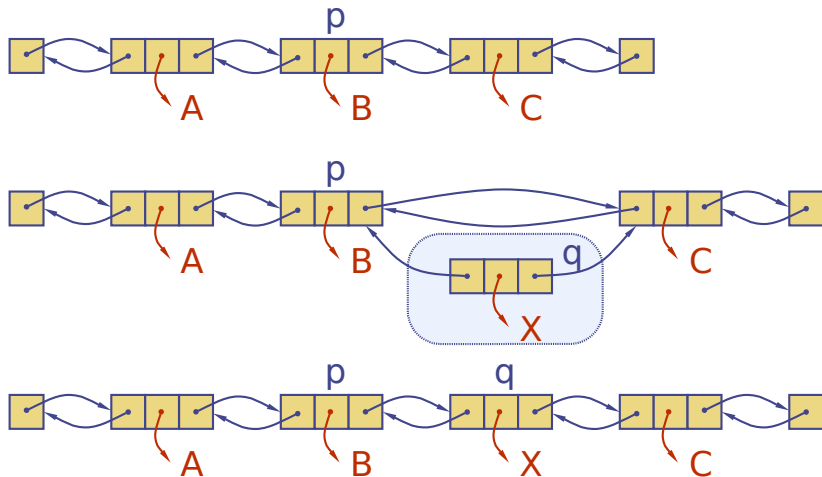
Doubly Linked List Implementation

A simple way of implementing the positional list ADT uses a doubly linked list, where the nodes are the positions



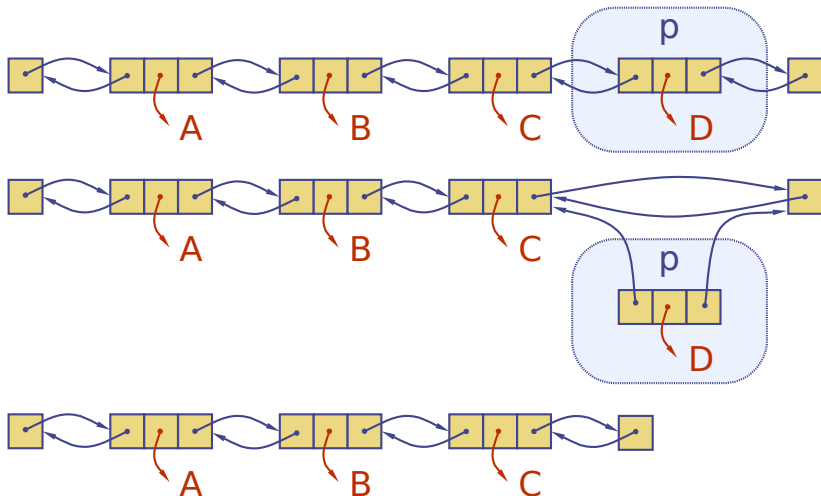
Linked Positional List: Element Insertion

Insert a new node q between p and its successor



Linked Positional List: Element Removal

Remove a node p from a doubly-linked list



Linked Positional List Implementation I

```
1  /** Implementation of a positional list stored as a doubly linked list. */
2  public class LinkedPositionalList<E> implements PositionalList<E> {
3      //----- nested Node class -----
4      private static class Node<E> implements Position<E> {
5          private E element;                // reference to the element stored at this node
6          private Node<E> prev;              // reference to the previous node in the list
7          private Node<E> next;              // reference to the subsequent node in the list
8          public Node(E e, Node<E> p, Node<E> n) {
9              element = e;
10             prev = p;
11             next = n;
12         }
13         public E getElement() throws IllegalStateException {
14             if (next == null)                // convention for defunct node
15                 throw new IllegalStateException("Position no longer valid");
16             return element;
17         }
18         public Node<E> getPrev() { return prev; }
19         public Node<E> getNext() { return next; }
20         public void setElement(E e) { element = e; }
21         public void setPrev(Node<E> p) { prev = p; }
22         public void setNext(Node<E> n) { next = n; }
23     } //----- end of nested Node class -----
24 }
```

Linked Positional List Implementation II

```
25 // instance variables of the LinkedPositionalList
26 private Node<E> header; // header sentinel
27 private Node<E> trailer; // trailer sentinel
28 private int size = 0; // number of elements in the list
29
30 /** Constructs a new empty list. */
31 public LinkedPositionalList() {
32     header = new Node<>(null, null, null); // create header
33     trailer = new Node<>(null, header, null); // trailer is preceded by header
34     header.setNext(trailer); // header is followed by trailer
35 }
36
37 // private utilities
38 /** Validates the position and returns it as a node. */
39 private Node<E> validate(Position<E> p) throws IllegalArgumentException {
40     if (!(p instanceof Node)) throw new IllegalArgumentException("Invalid p");
41     Node<E> node = (Node<E>) p; // safe cast
42     if (node.getNext() == null) // convention for defunct node
43         throw new IllegalArgumentException("p is no longer in the list");
44     return node;
45 }
46
47 /** Returns the given node as a Position (or null, if it is a sentinel). */
48 private Position<E> position(Node<E> node) {
49     if (node == header || node == trailer)
50         return null; // do not expose user to the sentinels
51     return node;
52 }
53
```

Linked Positional List Implementation III

```
54 // public accessor methods
55 /** Returns the number of elements in the linked list. */
56 public int size() { return size; }
57
58 /** Tests whether the linked list is empty. */
59 public boolean isEmpty() { return size == 0; }
60
61 /** Returns the first Position in the linked list (or null, if empty). */
62 public Position<E> first() {
63     return position(header.getNext());
64 }
65
66 /** Returns the last Position in the linked list (or null, if empty). */
67 public Position<E> last() {
68     return position(trailer.getPrev());
69 }
70
71 /** Returns the Position immediately before Position p (or null, if p is first). */
72 public Position<E> before(Position<E> p) throws IllegalArgumentException {
73     Node<E> node = validate(p);
74     return position(node.getPrev());
75 }
76
77 /** Returns the Position immediately after Position p (or null, if p is last). */
78 public Position<E> after(Position<E> p) throws IllegalArgumentException {
79     Node<E> node = validate(p);
80     return position(node.getNext());
81 }
82
```


Linked Positional List Implementation IV

```
83 // private utilities
84 /** Adds element e to the linked list between the given nodes. */
85 private Position<E> addBetween(E e, Node<E> pred, Node<E> succ) {
86     Node<E> newest = new Node<>(e, pred, succ); // create and link a new node
87     pred.setNext(newest);
88     succ.setPrev(newest);
89     size++;
90     return newest;
91 }
92
93 // public update methods
94 /** Inserts element e at the front of the linked list and returns its new Position. */
95 public Position<E> addFirst(E e) {
96     return addBetween(e, header, header.getNext()); // just after the header
97 }
98
99 /** Inserts element e at the back of the linked list and returns its new Position. */
100 public Position<E> addLast(E e) {
101     return addBetween(e, trailer.getPrev(), trailer); // just before the trailer
102 }
103
104 /** Inserts element e immediately before Position p, and returns its new Position. */
105 public Position<E> addBefore(Position<E> p, E e) throws IllegalArgumentException {
106     Node<E> node = validate(p);
107     return addBetween(e, node.getPrev(), node);
108 }
109
```

Linked Positional List Implementation V

```
110  /** Inserts element e immediately after Position p, and returns its new Position. */
111  public Position<E> addAfter(Position<E> p, E e) throws IllegalArgumentException {
112      Node<E> node = validate(p);
113      return addBetween(e, node, node.getNext());
114  }
115
116  /** Replaces the element stored at Position p and returns the replaced element. */
117  public E set(Position<E> p, E e) throws IllegalArgumentException {
118      Node<E> node = validate(p);
119      E answer = node.getElement();
120      node.setElement(e);
121      return answer;
122  }
123
124  /** Removes the element stored at Position p and returns it (invalidating p). */
125  public E remove(Position<E> p) throws IllegalArgumentException {
126      Node<E> node = validate(p);
127      Node<E> predecessor = node.getPrev();
128      Node<E> successor = node.getNext();
129      predecessor.setNext(successor);
130      successor.setPrev(predecessor);
131      size--;
132      E answer = node.getElement();
133      node.setElement(null);
134      node.setNext(null);
135      node.setPrev(null);
136      return answer;
137  }
138 }
```

*// help with garbage collection
// and convention for defunct node*

Linked Positional List: Analysis

For a positional list ADT implementation with a doubly linked list all operations run in worst-case constant time

This is in contrast to the `ArrayList` structure.

Which operations required linear time with an `ArrayList`?

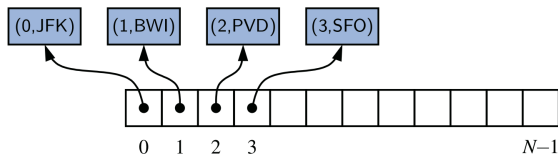
Method	Time
<code>size()</code>	$O(1)$
<code>isEmpty()</code>	$O(1)$
<code>first()</code> , <code>last()</code>	$O(1)$
<code>before(p)</code> , <code>after(p)</code>	$O(1)$
<code>addFirst(e)</code> , <code>addLast(e)</code>	$O(1)$
<code>addBefore(p, e)</code> , <code>addAfter(p, e)</code>	$O(1)$
<code>set(p, e)</code>	$O(1)$
<code>remove(p)</code>	$O(1)$

Space usage: $O(n)$, where n is the number of elements in the list

Implementing a Positional List with an Array

We store a new kind of position object in each cell of array A

A position p stores the element e as well as the current index i of that element within the list



We can determine the index currently associated with a position

And we can determine the position currently associated with a specific index

During insertions/deletions, we loop through the array to update the index variable stored in all positions in the list that are shifted

`addFirst`, `addBefore`, `addAfter`, `remove` take $O(n)$ time

All the other position-based methods take $O(1)$ time

Summary

Reading

Section 7.3 Positional Lists

Questions?