

Asymptotic Analysis

Stirling's formula: $n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$

First, simplifications of some functions:

$$1) \quad 2^{\log_2(\sqrt{n})} = \sqrt{n}$$

$$2) \quad \binom{n}{3} = \frac{n!}{3!(n-3)!} = \frac{(n-3)! \cdot (n-2) \cdot (n-1) \cdot n}{3!(n-3)!} = \\ = \frac{(n-2) \cdot (n-1) \cdot n}{3!} = \frac{(n-2)(n^2-n)}{3!} = \frac{n^3 - 3n^2 + 2n}{3!}$$

$$3) \quad \log_2(n^n) = n \log_2(n)$$

$$4) \quad (n+1)! \approx \sqrt{2\pi(n+1)} \left(\frac{n+1}{e}\right)^{n+1}$$

$$5) \quad \log_2(n!) \approx \log_2\left(\sqrt{2\pi n} \left(\frac{n}{e}\right)^n\right)$$

$$6) \quad \binom{n}{n/2} = \frac{n!}{(n/2)!(n-n/2)!} = \frac{n!}{((n/2)!)^2} \approx \frac{\sqrt{2\pi n} \left(\frac{n}{e}\right)^n}{\left(\sqrt{2\pi \frac{n}{2}} \left(\frac{n/2}{e}\right)^{\frac{n}{2}}\right)^2} = \\ = \frac{\sqrt{2\pi n} \left(\frac{n}{e}\right)^n}{\pi n \left(\frac{n}{2e}\right)^n} = \frac{\sqrt{2}}{\sqrt{\pi n} \left(\frac{1}{2}\right)^n} = \frac{(\sqrt{2})^n}{\sqrt{\pi n}}$$

$$7) \quad 8^{\log_2(n)} = 2^{3\log_2(n)} = 2^{\log_2(n^3)} = n^3$$

$$10 \log n \rightarrow 2^{\log \sqrt{n}} \rightarrow \sqrt{n} + 91^{2886} \log n \rightarrow \frac{n}{2} \rightarrow 5n \rightarrow \\ \rightarrow n \log n \rightarrow \log(n^n) \rightarrow \log(n!) \rightarrow \binom{n}{3} \rightarrow 8^{\log n} \rightarrow n + 5n^3 \rightarrow \\ \rightarrow \binom{n}{n/2} \rightarrow 2^{2n} \rightarrow 4^n \rightarrow 5n + n^n \rightarrow n! \rightarrow (n+1)!$$

1)

a) $O(n^3)$

b) $O(n^2)$

c) $O(n^3)$ tighter bound is $O(n^{2.25})$

d) $O(n^3)$

e) $O(n \log_2(n))$

f) $O(n \log^2_2(n))$

2)

a) $O(n)$

b) *undefined*:

Best case: the array will be sorted, best case is $O(n)$

“Expected case”: if we assume that all the permutations are equally likely, it might find the sorted array permutation in $n!$ trials

“Worst case” : infinite loop and stack overflow

c) *infinite loop*:

j in the second loop starts with 0 and is multiplied by 2 for each iteration, so it simply is evaluated to 0 infinitely

but, if j in the second for loop started with $j = 1$ instead of 0, the time complexity would be $O(n \log_2(n))$

d) $O(n \log_2(n))$

e) $O(n(\log_2(n))^2)$:

first iteration runs $\log n$, second runs $\log n$, third runs $n/2$

f) $O(n)$

GCD

```
public class GCD {
    public static int gcd(int a, int b) {
        if(b == 0) {
            return a;
        }
        return gcd(b, a % b);
    }

    public static void main(String[] args) {
        System.out.println(gcd(1071, 462));
    }
}
```

SumOfDigits

```
public class SumOfDigits {
    public static int sumOfDdigits(int n) {
        if (n < 10) {
            return n;
        }
        else {
            return n % 10 + sumOfDdigits(n / 10);
        }
    }

    public static void main(String[] args) {
        System.out.println(sumOfDdigits(1850374321));
    }
}
```

MazeSolver

```
public class MazeSolver {

    public static boolean pathExists(char[][] maze, int row, int col) {
        // Base cases
        if (row < 0 || row >= maze.length || col < 0 || col >=
maze[0].length) {
            return false; // Out of bounds
        }
        if (maze[row][col] == 'D') {
            return true; // Destination found
        }
        if (maze[row][col] != '.') {
            return false; // Blocked or already visited
        }

        // Mark the current cell as visited
        maze[row][col] = 'V';

        // Recursive calls in all four directions
        boolean pathExists = pathExists(maze, row - 1, col) || // Up
            pathExists(maze, row + 1, col) || // Down
            pathExists(maze, row, col - 1) || // Left
            pathExists(maze, row, col + 1); // Right

        // restore the current cell
        maze[row][col] = '.';

        return pathExists;
    }

    public static void main(String[] args) {
        char[][] maze = {{'.','X','X','X'},
            {'.','.', 'X','.'},
            {'X','.', 'X','X'},
            {'.','.', 'D'}};
        System.out.println(pathExists(maze, 0, 0));
    }
}
```

InfiniteExponent

```
public class InfiniteExponent {  
    public static double countInfiniteExponents(int i) {  
        if (i == 0) {  
            return 1;  
        }  
        else {  
            double x =  
Math.pow(Math.sqrt(2),countInfiniteExponents(i-1));  
            System.out.println("i is: " + i);  
            System.out.println("x is: " + x);  
            return x;  
        }  
    }  
  
    public static void main(String[] args) {  
        countInfiniteExponents(40);  
    }  
}
```