Experiment 3.3

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1. Aim/Overview of the practical:

Knuth morris pratt to search pattern matching

2. Task to be done/which logistics used:

Knuth morris pratt to search pattern matching

3. Steps for experiment/practical/Code:

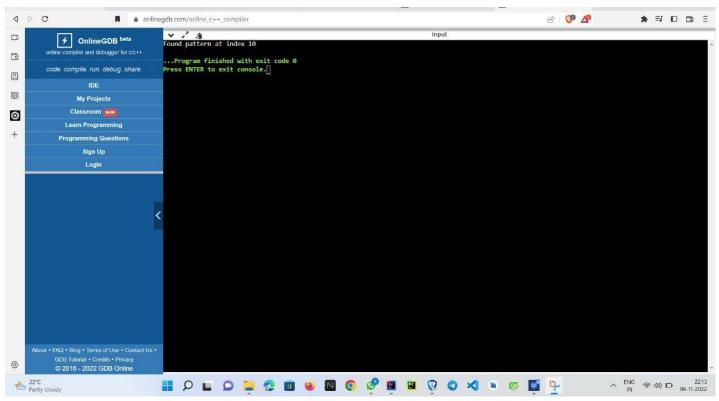
```
#include <bits/stdc++.h>
void computeLPSArray(char* pat, int M, int* lps);
// Prints occurrences of txt[] in pat[]
void KMPSearch(char* pat, char* txt)
{
        int M = strlen(pat);
        int N = strlen(txt);
        // create lps[] that will hold the longest prefix suffix
        // values for pattern
        int lps[M];
        // Preprocess the pattern (calculate lps[] array)
        computeLPSArray(pat, M, lps);
        int i = 0; // index for txt[]
        int j = 0; // index for pat[]
        while ((N - i) \ge (M - j)) {
                 if (pat[i] == txt[i]) {
                         j++;
                         i++;
                 if (j == M) {
                          printf("Found pattern at index %d ", i - j);
```

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```
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                          j = lps[j - 1];
                 // mismatch after j matches
                 else if (i \le N \&\& pat[j] != txt[i]) {
                         // Do not match lps[0..lps[j-1]] characters,
                          // they will match anyway
                          if (j != 0)
                                  j = lps[j - 1];
                          else
                                  i = i + 1;
}
// Fills lps[] for given pattern pat[0..M-1]
void computeLPSArray(char* pat, int M, int* lps)
{
        // length of the previous longest prefix suffix
        int len = 0;
        lps[0] = 0; // lps[0] is always 0
        // the loop calculates lps[i] for i = 1 to M-1
        int i = 1;
        while (i \le M) {
                 if(pat[i] == pat[len]) {
                          len++;
                          lps[i] = len;
                          i++;
                 else // (pat[i] != pat[len])
                          // This is tricky. Consider the example.
                          // AAACAAAA and i = 7. The idea is similar
                          // to search step.
                          if (len != 0) {
                                   len = lps[len - 1];
                                  // Also, note that we do not increment
                                  // i here
                          else // if (len == 0)
                                   lps[i] = 0;
                                   i++;
                 }
```

4. Output:



5. Learning Outcomes:

- a) Create a program keeping in mind the time complexity.
- b) Create a program keeping in mind the space complexity.



- c) Steps to make optimal algorithm.
- d) Learnt about how to implement 0-1 Knapsack problem using dynamic programming.

