

Name: Pradeep Rekapalli

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KSU10: 001100818

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NETID: PREKAPAL

CS 7530 Fall 2024

Assignment-2

Question-1

① Calculate $\phi(230)$.

Given to calculate $\phi(n)$ [Euler's function] which counts the number of integer less than or equal to n , that are relatively prime (co) to n .

To so in order to calculate $\phi(230)$ by this steps.

Step 1: we need find prime factorization of 230.

$$\begin{array}{r} 2 \overline{) 230} \\ 5 \overline{) 115} \\ 23 \end{array}$$

$$\text{So } 230 = 2 \times 5 \times 23$$

$p_1 \quad p_2 \quad p_3$

Step 2: Euler's function formula

$$\phi(n) = n \times \left(1 - \frac{1}{p_1}\right) \times \left(1 - \frac{1}{p_2}\right) \times \dots \times \left(1 - \frac{1}{p_k}\right)$$

So now we substitute this & find

$$\phi(230) = 230 \times \left(1 - \frac{1}{2}\right) \times \left(1 - \frac{1}{5}\right) \times \left(1 - \frac{1}{23}\right)$$

$$\phi(230) = 230 \times \frac{1}{2} \times \frac{4}{5} \times \frac{22}{23}$$

$$\text{So } \phi(230) = 88$$

② using Euler's theorem calculate

Ans

$$a^{\phi(n)} = 27^{88} \mod 230$$

$$n = 230$$

$$a = 27$$

before doing this but before we already calculate

$$\phi(230) = 88 \text{ so let}$$

$$27^{\phi(230)+1} = 27^{89} \mod 230$$

Equation

LHS

RHS

~~$$27^{88+1} = 27^{89} \mod 230$$~~

according to theorem it is $a^{\phi(n)+1} = 1 \mod n$.

Here 27, 230 are coprime.

according to theorem $a^{\phi(n)} = 1 \mod 230$

① $\leftarrow 27^{88} = 1 \mod 230$

Substitute it in equation:

$$27^{89} \bmod 230$$

$$27 \times 27^{88} \bmod 230 \rightarrow \text{from (1)}$$

$$27 \times (1) \text{ so}$$

$$27^{89} \bmod 230 = [27]$$

Section 2 (Question)

(a) Ans

Given to encrypt the plaintext "cryptography" using the

Vigenere Cipher Basics.

So we need to shift by the corresponding number in the key stream.

So for example a key shift of 9 shifts the first letter of the plaintext by 9 positions in the alphabet.

So

Given:

* plaintext = "cryptography".

Key stream = 9, 11, 15, 2, 8, 6, 1, 21, 16, 18, 13, 7

So first we need to convert alphabets we need to know its positions when

So $A=0, B=1, \dots, Z=25$ so

$C=2, X=17, Y=24, P=15, T=19, O=14, G=6, R=17, A=0,$

$P=15, H=7, Y=24.$

$\left[\begin{array}{l} \text{if we get } > 26 \text{ we do mod} \\ \text{if we get } < 0 \text{ we add } 26 \end{array} \right]$

So Now

$$C \rightarrow 2+9 = 11 \rightarrow L$$

$$X \rightarrow 17+11 = 28 \bmod 26 = 2 \rightarrow C$$

$$Y \rightarrow 24+15 = 39 \bmod 26 = 13 \rightarrow N$$

$$P \rightarrow 15+2 = 17 \rightarrow R$$

$$T \rightarrow 19+8 = 27 \bmod 26 = 1 \rightarrow B$$

$$O \rightarrow 14+6 = 20 \bmod 26 = 4 \rightarrow E$$

$$G = 6+1 = 7 \rightarrow H$$

$$R = 17+21 = 38 \bmod 26 = 12 \rightarrow M$$

$$A = 0+16 = 16 \rightarrow Q$$

$$P = 15+18 = 33 \bmod 26 = 7 \rightarrow H$$

$$P = 15+18 = 33 \bmod 26 = 7 \rightarrow H$$

$$H = 7+13 = 20 \rightarrow U$$

$$Y = 24+7 = 31 \bmod 26 = 5 \rightarrow F$$

So now ciphertext = "LCNRBEHMQHUF"

"LCNRBEHMQHUF"

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(5)

⑬ Ans So ciphertext = "LCNRBEHMQHUF" → ①

Desired Plaintext = "arkennesawunw" → ②

So Now we need to do ① - ② to find key
if we get a

[If we get a number < 0 or ≥ 26 we do mod]

1. $L(1) - k(10) = 1 \rightarrow \text{key} = 1 \rightarrow \text{key}$

2. $C(2) - e(4) = -2 \equiv -2 \bmod 26 = 24$

3. $N(13) - n(13) = 0 \equiv 0$

4. $R(17) - n(13) = 4 \equiv 4$

5. $B(1) - e(4) = -3 \bmod 26 = 23$

6. $E(4) - S(18) = -14 \bmod 26 = 12$

7. $H(7) - a(0) = 7$

8. $M(12) - w(22) = -10 \equiv -10 \bmod 26 = 16$

9. $Q(16) - u(20) = -4 \bmod 26 = 22$

10. $H(7) - n(13) = -6 \bmod 26 = 20$

11. $u(20) - i(8) = 12$

12. $F(5) - v(21) = -16 \bmod 26 = 10$

So the key is

Key = 1, 24, 0, 4, 23, 12, 7, 16, 22, 20, 12, 10

Part 2 : Question 3:

Report:

As in question it is given to implement DES Function, and we need to give two inputs R_i (32 bits) and K_i (48bits) which produces a 32 bit output.

First lets understand DES (Data Encryption Standard). Here the F Function is main part.as f-
fucntion need to give two inputs R_i (32 bit) and K_i (48 bit) here we will do R_i to 48 bits by using
expansion which I will explain in steps. And then XOR it with K_i then finally permuate it to get 32
bit ouput.

So here I am dividing whole procedure into few sections for detailed explanation.I have written
sections based on the procedure.

1. Step:Expansion;
Here we need to expand R_i 32 bit to 48 bits by using expansion table
It generally expanded bots by repreating the certain bits from R_i
Here we have done in expand function
2. Step:XOR:
The 48 bit is then XOR by the 48 bit round key K_i .
3. Step:Subsititution (S box):
Here after xor oprtation the 48 bit ones are divided into 6 bit blocks
Each 6 bits are then converted by using S-box like in 6 bits , 1st and 6th bit are determine
the rows and inside 4 bits decide column
So that we can say that each Sbox output is 6 bit input to 4 bit putput
4. Step:Permutation (P box)
So total again sum 32 bits are formed as we are excluding 2 from 6 bits.and these are
permeated through but using permutation table
Here during the this procedure we reaarange ouputs for final output.This is key
procedure.

Additional :

Here I used `hex_to_bin` and `bin_to_hex` to convert the hexadecimal input to binary and vide versa
which I took reference from greeksforgreeks. And I have implemented code in C++.

For code I have took reference from and lectures provided

Reference : <https://github.com/alimasry/DES-Encryption/>

Code :

```
#include <iostream>
```

```
#include <string>
```

```
#include <bitset>
```

```
#include <sstream>
```

```
#include <iomanip>
```

```
using namespace std;
```

```
const int expansion_table[48] = {  
    32, 1, 2, 3, 4, 5, 4, 5, 6, 7, 8, 9,  
    8, 9, 10, 11, 12, 13, 12, 13, 14, 15, 16, 17,  
    16, 17, 18, 19, 20, 21, 20, 21, 22, 23, 24, 25,  
    24, 25, 26, 27, 28, 29, 28, 29, 30, 31, 32, 1  
};
```

```
const int permutation_table[32] = {  
    16, 7, 20, 21, 29, 12, 28, 17, 1, 15, 23, 26,  
    5, 18, 31, 10, 2, 8, 24, 14, 32, 27, 3, 9,  
    19, 13, 30, 6, 22, 11, 4, 25  
};
```

```
const int sbox[8][4][16] = {  
    {{14, 4, 13, 1, 2, 15, 11, 8, 3, 10, 6, 12, 5, 9, 0, 7},  
     {0, 15, 7, 4, 14, 2, 13, 1, 10, 6, 12, 11, 9, 5, 3, 8},  
     {4, 1, 14, 8, 13, 6, 2, 11, 15, 12, 9, 7, 3, 10, 5, 0},  
     {15, 12, 8, 2, 4, 9, 1, 7, 5, 11, 3, 14, 10, 0, 6, 13}},  
    // S2-S8 omitted for brevity. Include all 8 S-boxes in a real implementation.
```

```
};
```

```
string hex_to_bin(const string& hex) {  
    string bin;  
    for (char c : hex) {  
        int n = (c >= 'A') ? (c - 'A' + 10) : (c - '0');  
        bin += bitset<4>(n).to_string();  
    }  
    return bin;  
}
```

```
string bin_to_hex(const string& bin) {  
    stringstream ss;  
    for (size_t i = 0; i < bin.length(); i += 4) {  
        string chunk = bin.substr(i, 4);  
        int decimal = stoi(chunk, nullptr, 2);  
        ss << hex << uppercase << decimal;  
    }  
    return ss.str();  
}
```

```
string expand(const string& block) {  
    string expanded;  
    for (int i : expansion_table) {  
        expanded += block[i - 1];  
    }  
    return expanded;  
}
```

```
string xor_strings(const string& a, const string& b) {  
    string result;
```



```

for (size_t i = 0; i < a.length(); ++i) {
    result += (a[i] == b[i]) ? '0' : '1';
}
return result;
}

```

```

string apply_sbox(const string& block) {
    string output;
    for (int i = 0; i < 8; ++i) {
        int row = 2 * (block[i*6] - '0') + (block[i*6 + 5] - '0');
        int col = 8 * (block[i*6 + 1] - '0') + 4 * (block[i*6 + 2] - '0') +
            2 * (block[i*6 + 3] - '0') + (block[i*6 + 4] - '0');
        int val = sbox[i][row][col];
        output += bitset<4>(val).to_string();
    }
    return output;
}

```

```

string permute(const string& block) {
    string permuted;
    for (int i : permutation_table) {
        permuted += block[i - 1];
    }
    return permuted;
}

```

```

string f_function(const string& Ri, const string& Ki) {
    string expanded = expand(Ri);
    string xored = xor_strings(expanded, Ki);
    string substituted = apply_sbox(xored);
    string permuted = permute(substituted);
}

```

```
    return permuted;
}

int main() {
    // Test case 1

    string Ri_hex1 = "A1B2C3D4";
    string Ki_hex1 = "F0D532A490C6";

    string Ri_bin1 = hex_to_bin(Ri_hex1);
    string Ki_bin1 = hex_to_bin(Ki_hex1);

    string result_bin1 = f_function(Ri_bin1, Ki_bin1);
    string result_hex1 = bin_to_hex(result_bin1);

    cout << "Test Case 1:" << endl;
    cout << "Ri = " << Ri_hex1 << ", Ki = " << Ki_hex1 << endl;
    cout << "Output: " << result_hex1 << endl << endl;

    // Test case 2

    string Ri_hex2 = "3CF03C0F";
    string Ki_hex2 = "1E0F0380D293";

    string Ri_bin2 = hex_to_bin(Ri_hex2);
    string Ki_bin2 = hex_to_bin(Ki_hex2);

    string result_bin2 = f_function(Ri_bin2, Ki_bin2);
    string result_hex2 = bin_to_hex(result_bin2);

    cout << "Test Case 2:" << endl;
    cout << "Ri = " << Ri_hex2 << ", Ki = " << Ki_hex2 << endl;
    cout << "Output: " << result_hex2 << endl;
```

```
return 0;  
}
```

Output:

Output

Clear

```
/tmp/eSiVVJy20b.o  
Test Case 1:  
Ri = A1B2C3D4, Ki = F0D532A490C6  
Output: 00808002  
  
Test Case 2:  
Ri = 3CF03C0F, Ki = 1E0F0380D293  
Output: 00008000  
  
=== Code Execution Successful ===
```

Here is output, including test case. Professor, I am in hackthon at Morgan Stanley, so if anything needs to be done additionally, I missed anything. I would request for again submission after the deadline.

Thank you

Pradeep Rekapalli

KSU ID : 001100818

NET ID : prekapal