

# Hardware Attacks

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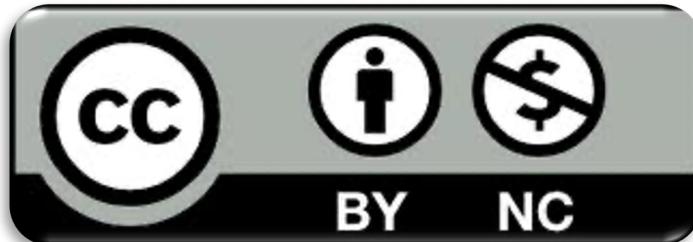
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# Acknowledgments

- The presentation includes material from
  - Giorgio DI NATALE
  - Nicolò MAUNERO
  - Gianluca ROASCIO

whose valuable contribution is here acknowledged and highly appreciated.

# Prerequisites

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## ➤ Lectures:

- *HS\_1.1 - The role of Hardware in Security*
- *HS\_1.2 - Hardware Vulnerabilities*

# Goal

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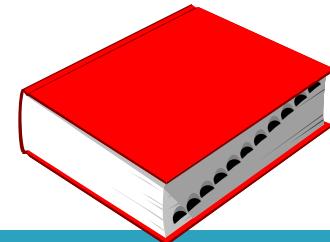
- Presenting the concept of *Hardware Attacks*
- Providing a comprehensive taxonomy and a detailed analysis of the most significant ones, clustered according to their
  - *goal* (stealing, corrupting or inhibiting)
  - *target* (information or property)
  - *modality* (invasive or non-invasive)
  - *domain* (logical or physical).

# Outline

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- Hardware Attacks
- Attack goals
- Attack targets
- Attack modalities
- Attack domains

# Hardware Attack



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- The act of exploiting a hardware vulnerability to perform an attack.

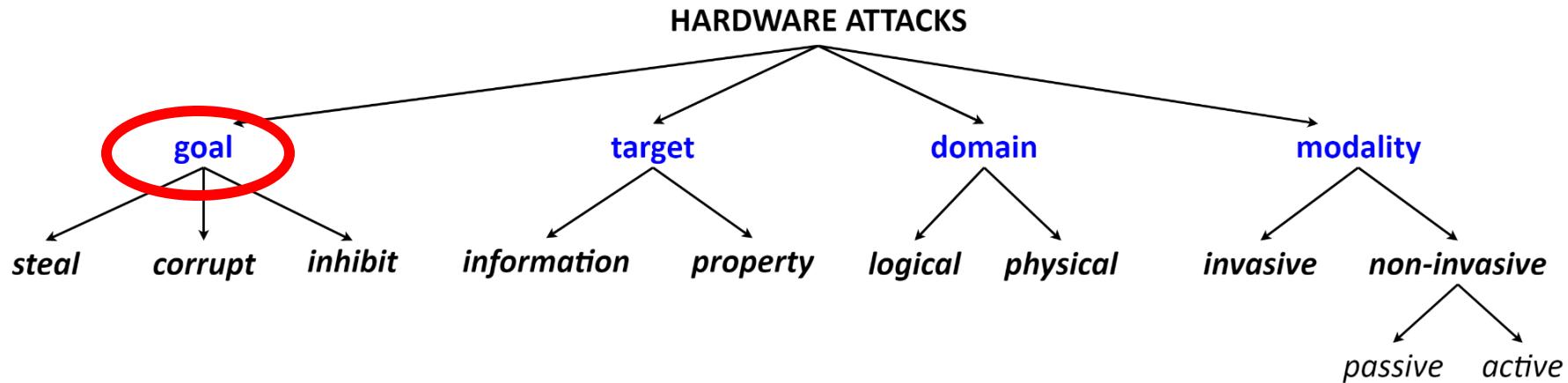
# Hardware Attacks

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- Can be clustered according to:
  - Their **goal** (*stealing, corrupting or inhibiting*)
  - Their **target** (*information or property*)
  - Their **modality** (*invasive or non-invasive*)
  - Their **domain** (*logical or physical*).

# Hardware Attacks Taxonomy

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# Hardware Attacks Goals

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- An attack always exploits a vulnerability, which always harms one of the three CIA properties

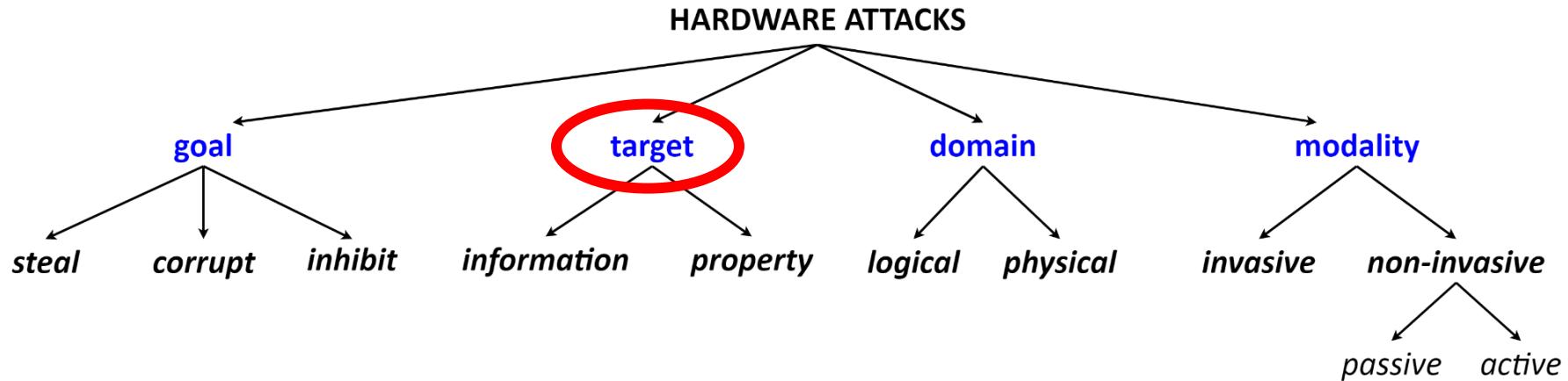
# Hardware Attacks Goals

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- An attack is then mounted to:
  - *steal* a target (e.g., a cryptographic key, a secret password, an intellectual property, a resource, etc.) → *Confidentiality* is broken
  - *corrupt* a target (e.g., a memory word, a permission file, a functionality, etc.) → *Integrity* is broken
  - *inhibit* a target (e.g., a service, a set of critical data, a defense mechanism, etc.) → *Availability* is broken

# Hardware Attacks Taxonomy

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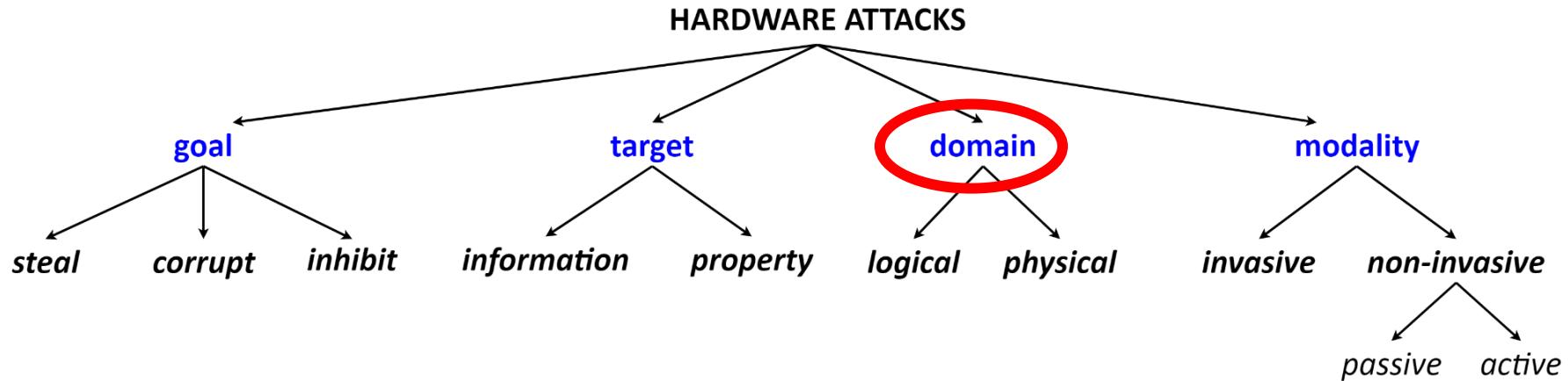
# Hardware Attacks Targets

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- The target of a hardware attack is always an asset treated or owned by a hardware component, which can be:
  - A piece of *information* treated by the hardware (a key, some credentials, a protected memory word or file, etc.)
  - A *property* owned by the hardware (its IP, one of its resource or functionalities, etc.)

# Hardware Attacks Taxonomy

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# Attack domains

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## Logical

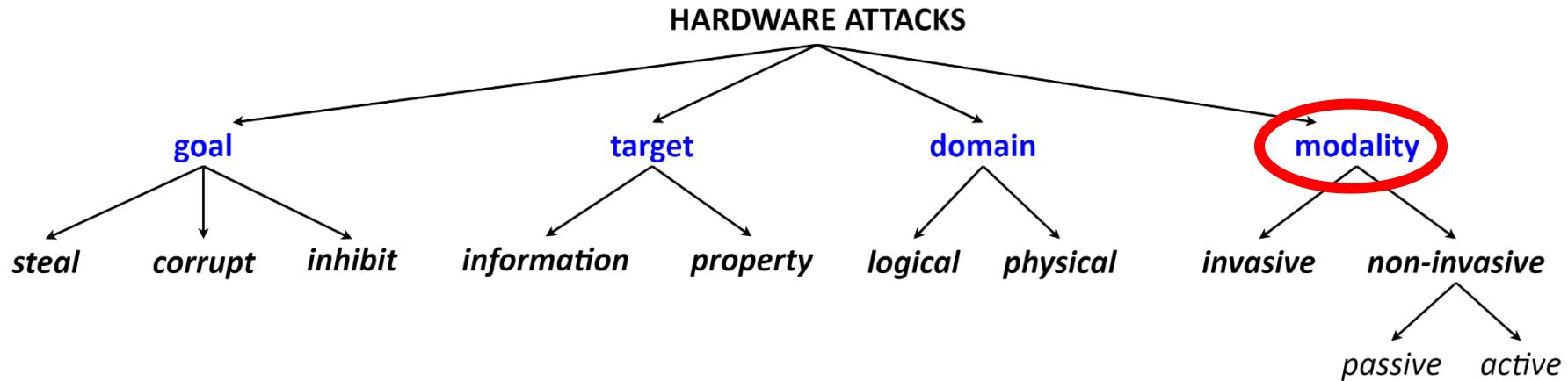
- Carried out exploiting the software run by the hardware
  - E.g., privilege escalation via software exploiting a hardware vulnerability

## Physical

- Carried out directly exploiting a hardware vulnerability

# Hardware Attacks Taxonomy

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# Hardware Attacks Modalities

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## Non-invasive

- Carried out without any physical contact with the device under attack.
- Further clustered in:
  - *Passive* non-invasive attacks
  - *Active* non-invasive attacks

## Invasive

- Require actions against the attacked hardware (e.g., desoldering, depackaging, disconnections) in order to allow a physical intrusions to easily access some internal components.

# Different “investments”

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## Non-invasive

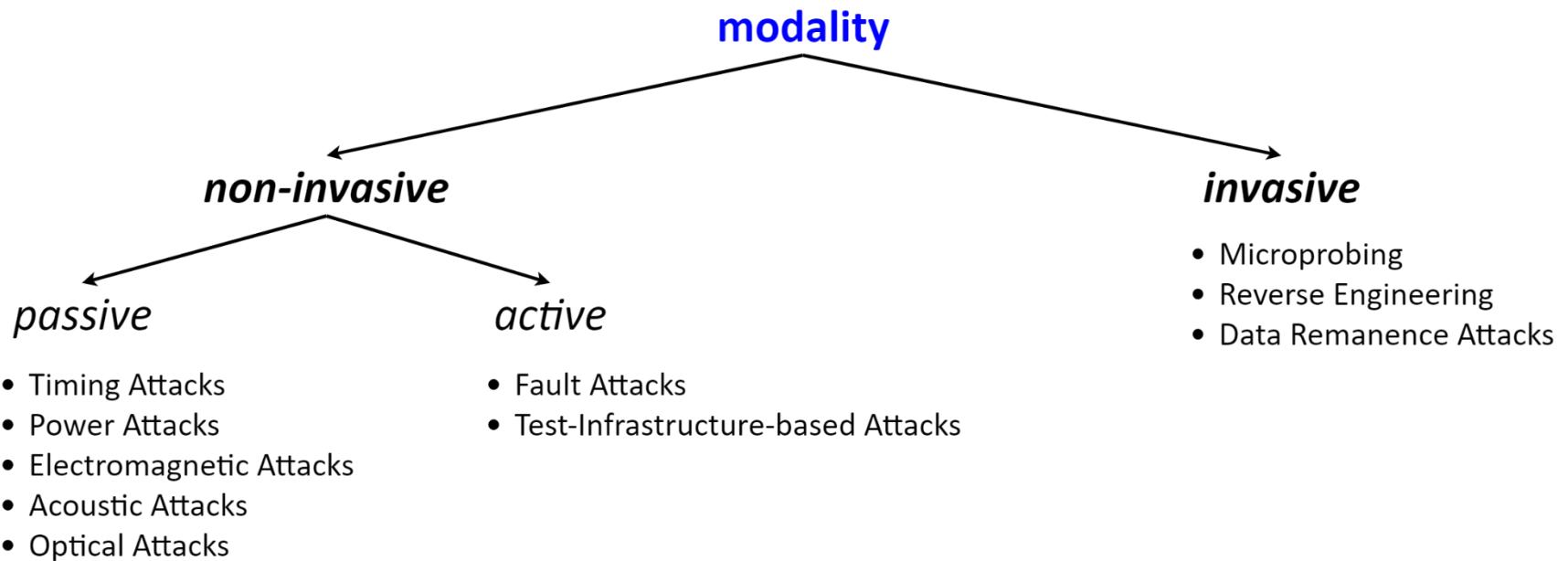
- They typically require moderate investments in terms of both lab equipment and effort in designing an attack on the target device
- Therefor the cost per device attacked is low

## Invasive

- They typically require:
  - a relatively high capital investment for lab equipment
  - a moderate investment of effort for each individual chip attacked.

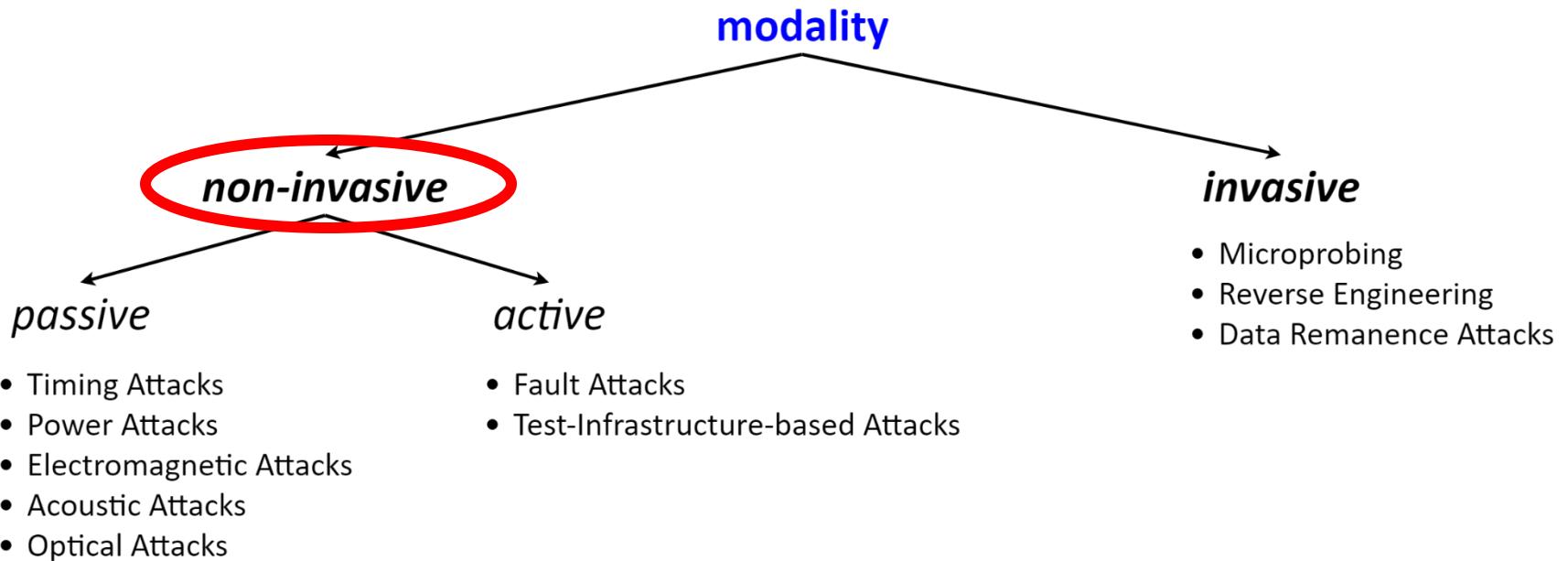
# Hardware Attacks Modalities

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# Hardware Attacks Modalities

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# Non-invasive Hardware Attacks

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## Passive

- Carried out by analyzing and measuring one (or more) physical dynamic entities of the attacked hardware.

## Active

- Require specific actions on the device, aimed at forcing the system into abnormal states, in which the attacker's goal is easier to reach.

# Non-invasive Hardware Attacks

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## Passive

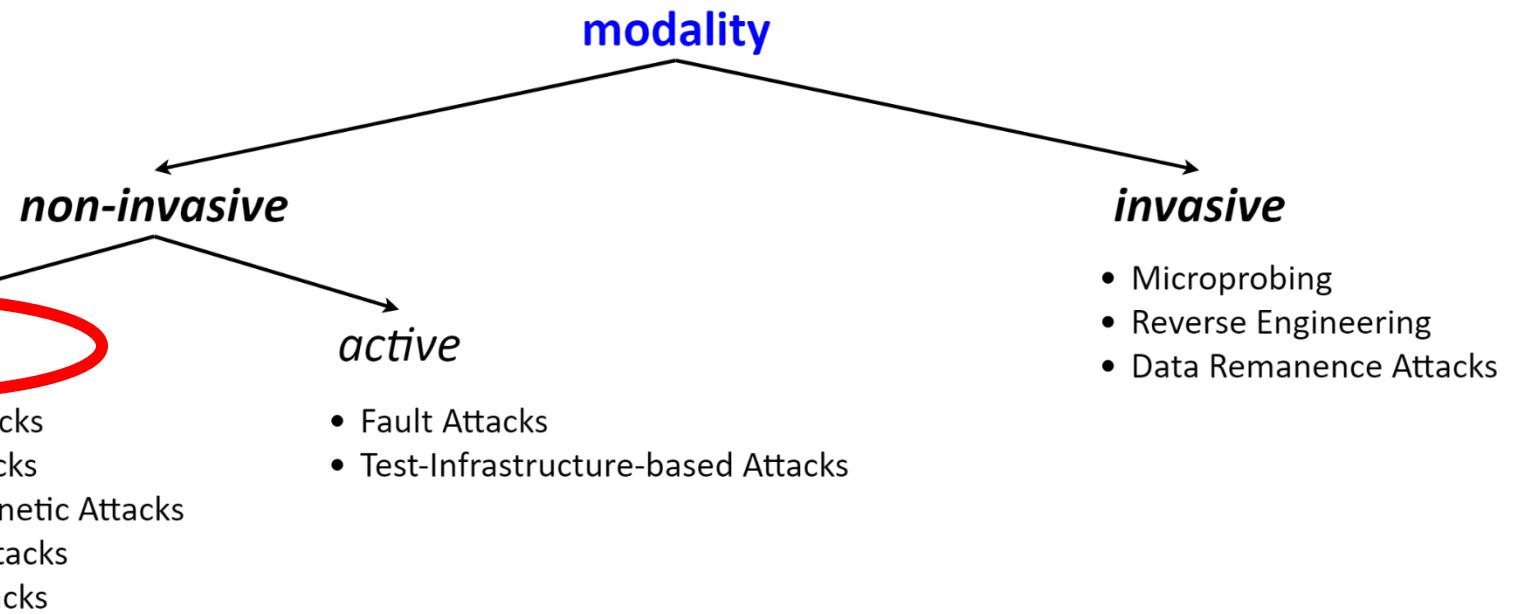
- The device works within its specifications

## Active

- The goal is reached by exploiting abnormal behaviors of the device
  - E.g., power glitches, laser pulses

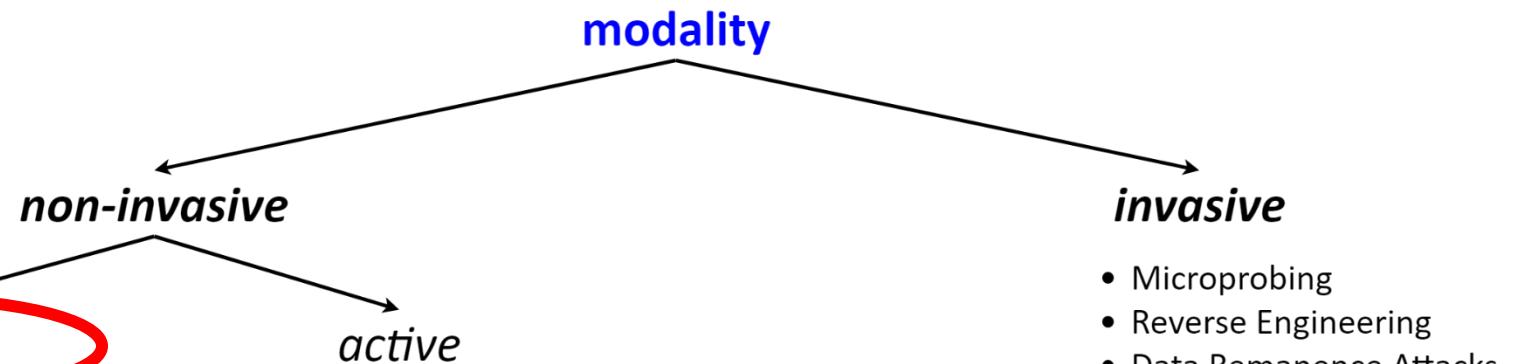
# Hardware Attacks Modalities

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# Hardware Attacks Modalities

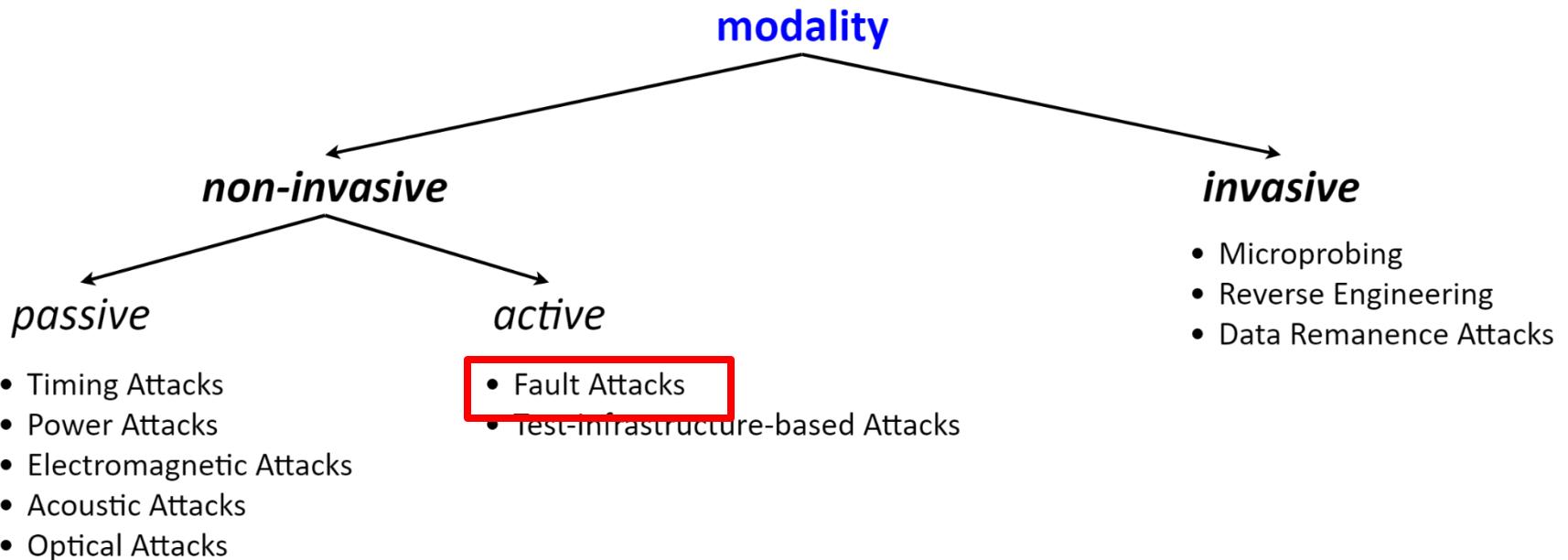
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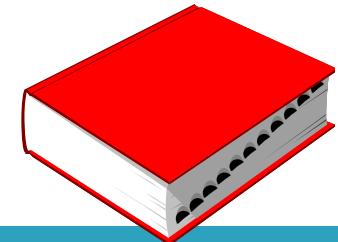
Deeply analyzed in the lectures:  
*HS\_3.1 - Side Channel Attacks*

# Hardware Attacks Modalities

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# Fault Attacks



- Consist in the injection of deliberate (malicious) faults into the target device, aimed at bringing it into a set of states from which private internal information items (e.g., a key) can be fraudulently extracted.

# Fault Attacks -- An example

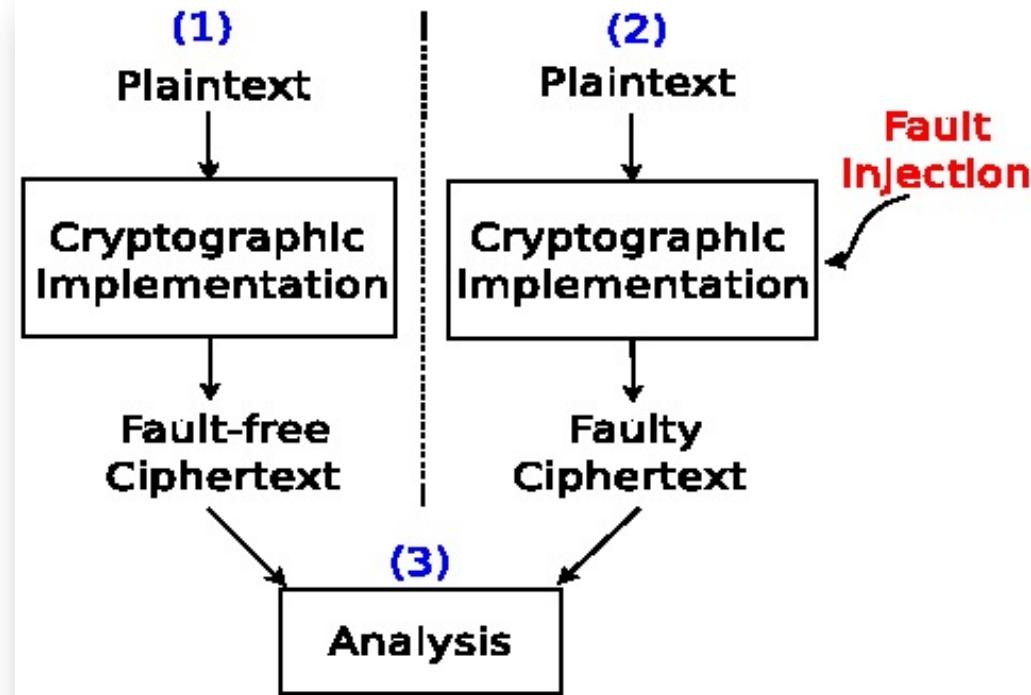
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- They exploit the possibility of inserting a fault in the process of algorithm execution in a way that it could help revealing the key.

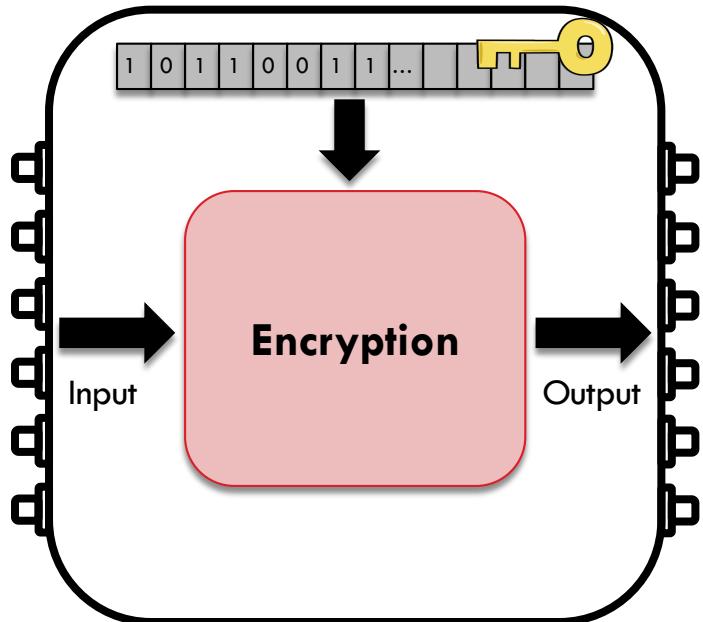
# Fault Attacks -- An example

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- Practically it changes the value of the key during the processing of the algorithm and compare the two outputs.



# Fault Attacks



Hypothesis: Injection forces a '0' on a single bit of the secret key

- 1)  $C_{OK} = E(P)$
- 2) Calculate  $C' = E(P)$ ,  
while injecting a fault
- 3) If  $C' = C_{OK}$  → target bit is '0'  
else → target bit is '1'

# Fault Injection Techniques

- To inject faults affecting critical paths:
  - Underpowering or overpowering
  - Altering the clock
  - Altering the temperature

# Fault Injection Techniques

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- To inject precise faults in space and time:
  - *Laser injections*

# Fault Injection Techniques

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- To inject precise faults in space and time:
  - *Laser injections*
- The most precise injection method
- Requires depackaging of the target device

# Fault Injection Techniques

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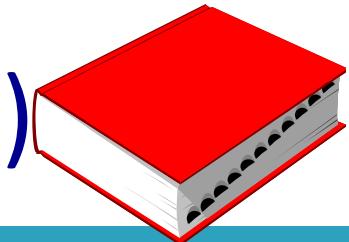
- To inject precise faults in space and time:
  - Laser injections
  - *Electromagnetic injections*

# Fault Injection Techniques

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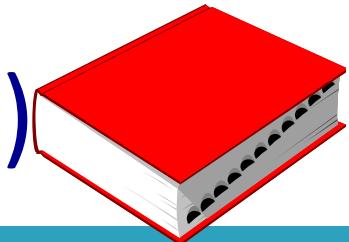
- To inject precise faults in space and time:
  - Laser injections
  - *Electromagnetic injections*
- *Harmonic signals* to target analog circuits
- *Pulse signals* to target digital circuits

# Differential Fault Analysis (DFA)



- A set of multiple injections with subsequent correlated analysis of the faulty behaviour of the target device

# Differential Fault Analysis (DFA)



- The attackers (cryptanalysts) focus on the fault influences of specific hardware implementations to collect some faulty outputs.
- Later, this information is compared and analyzed with correct results in order to harvest some partial or total compromise of the secret information.

# Injection goals

- Bypassing access/right control verification
  - To get access to the system
- Generating faulty encryptions/signatures
  - To be exploited by cryptanalysis in order to retrieve the secret key
- Denial of Service

# Countermeasures to Fault Attacks

- Several approaches have been proposed, including:
  - Preventing the attack
  - Detecting the fault injection
  - Detecting the fault effect (error)
  - De-synchronization
  - Robust package, “hardened” technologies

# Countermeasures to Fault Attacks

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# Fault detection

- Fault detectors:
  - Laser/light sensors
  - Bulk current sensors
  - Glitch detectors
- They can generate false positives:
  - Because of noise
  - Because the fault does not generate an error

# Countermeasures to Fault Attacks

- Several approaches have been proposed, including:
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# Important remark

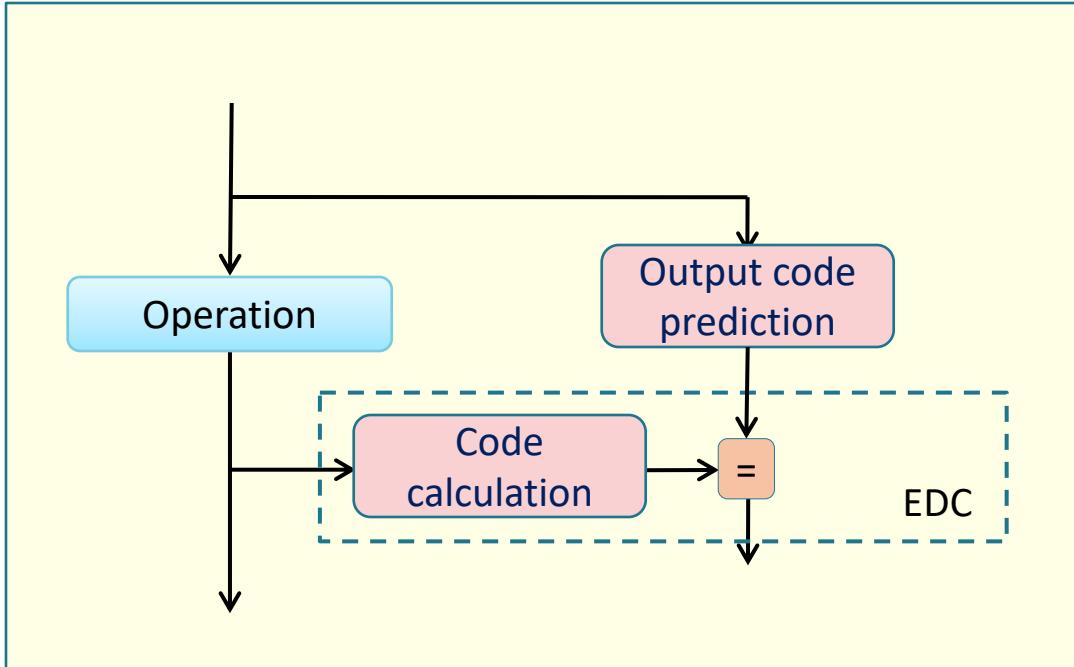
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- Safety-pushed solutions (e.g., all the solutions adopted to implement Fault Tolerant Architectures, aimed at “tolerating” i.e., detecting, correcting or simply making ineffective faults that could occur) are valuable countermeasures against Security-induced Fault attack !!

# Error detection

- Based on redundancy
  - Space redundancy
  - Information redundancy
  - Time redundancy
- Examples:
  - Performing twice the same operation and comparing the results
  - Performing the operation and its inverse, and checking

# Error Detection Codes



# Countermeasures to Fault Attacks

- Several approaches have been proposed, including:
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# De-synchronization

- A fault injection requires a precise timing to be effective
- Adding temporal randomness makes the timing of the fault harder to set
- How to:
  - Jittered clock
  - Dummy instructions
  - Randomized operation flow

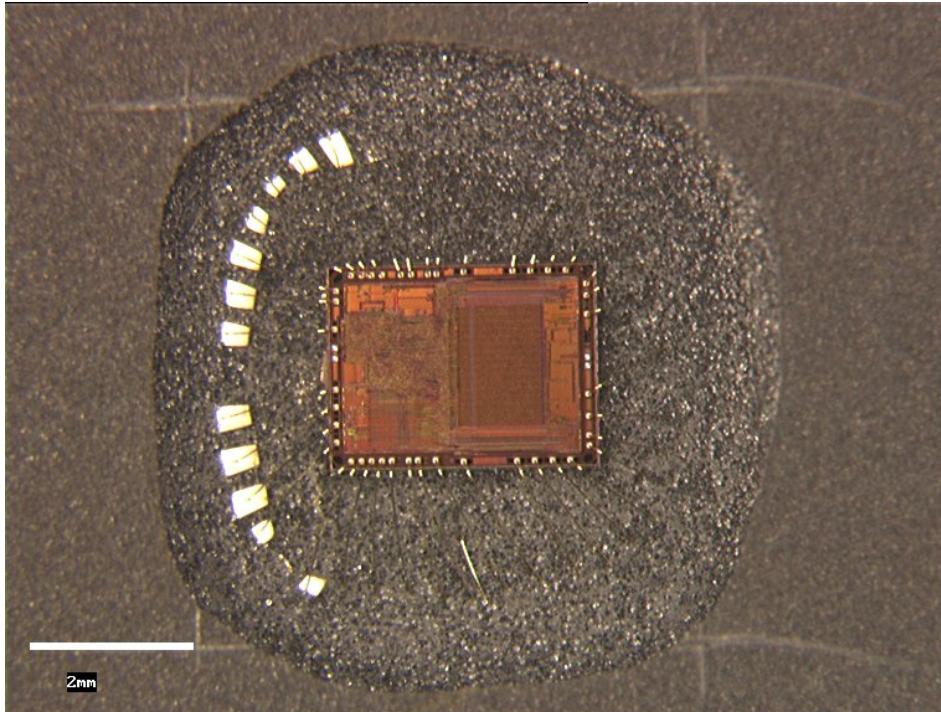
# Countermeasures to Fault Attacks

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# Hardened IC Package

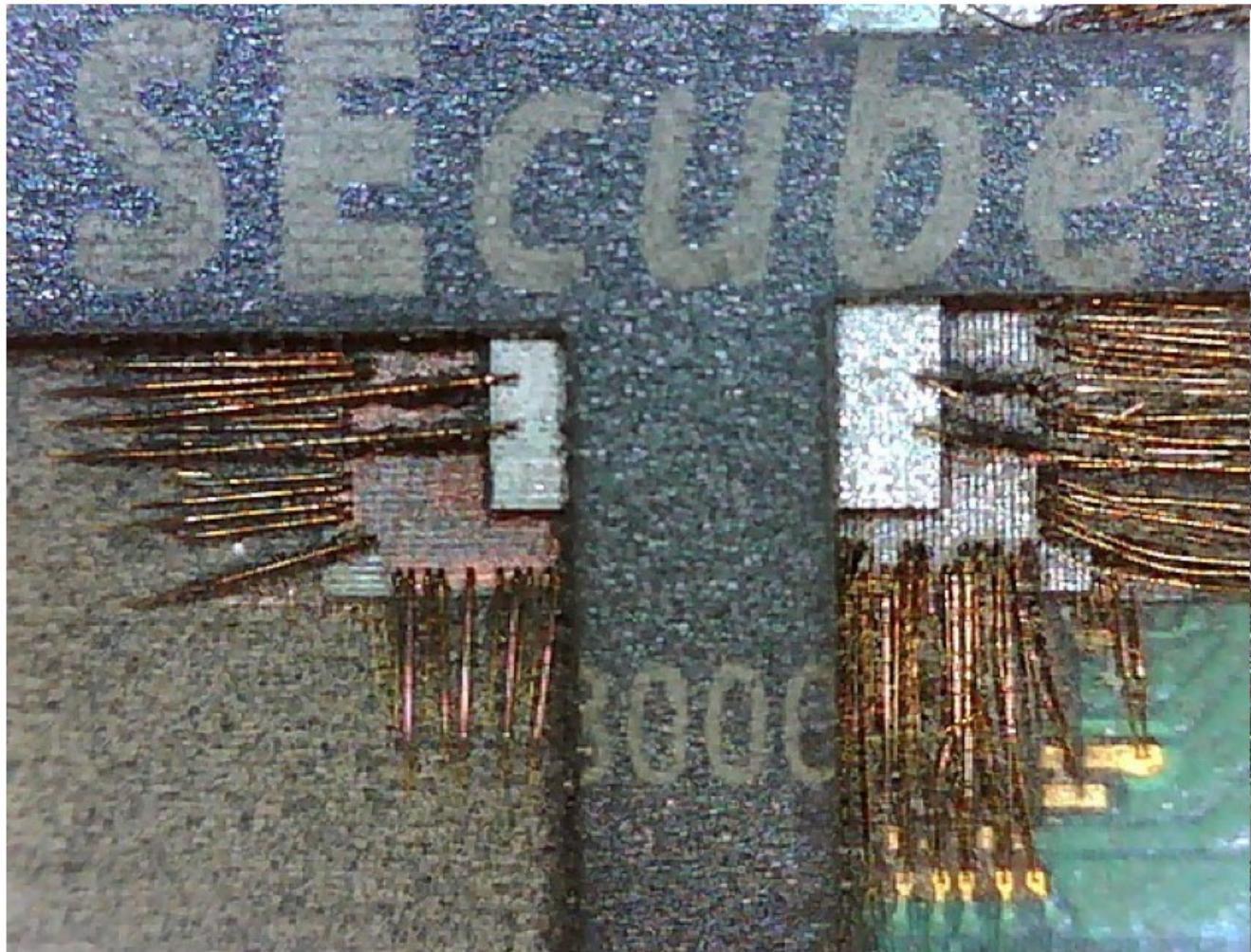
- Several fault injection techniques require the die of the IC to be exposed
  - FIB, laser, ...
- Hardening the package can be a countermeasure:
  - Smartcards are easy to open
  - Metallic package can be opened mechanically
  - Epoxy packages require chemical attacks
  - System-in-Package are very hard to open

# Example: Epoxy package opened with nitric acid



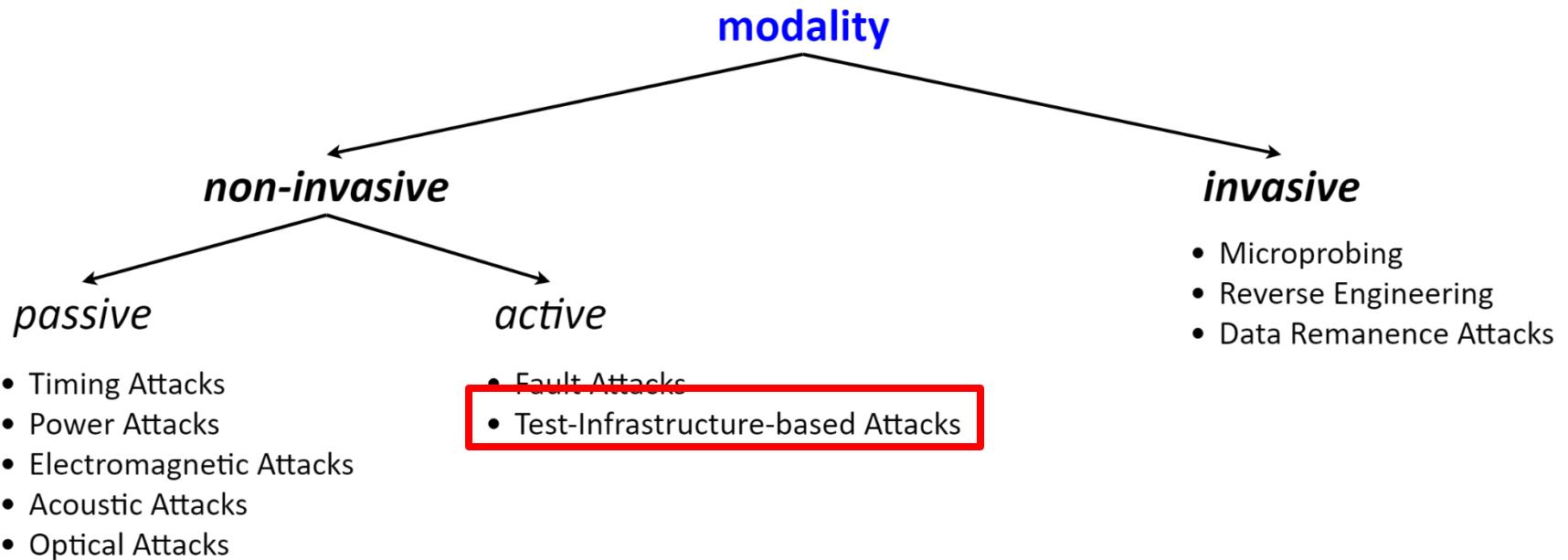
# 3D SiP

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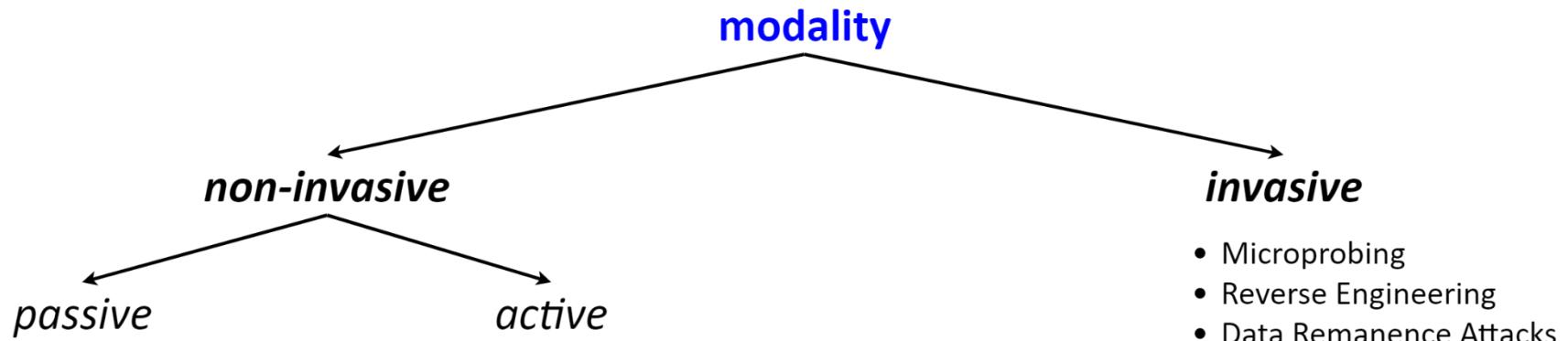
# Hardware Attacks Modalities

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# Hardware Attacks Modalities

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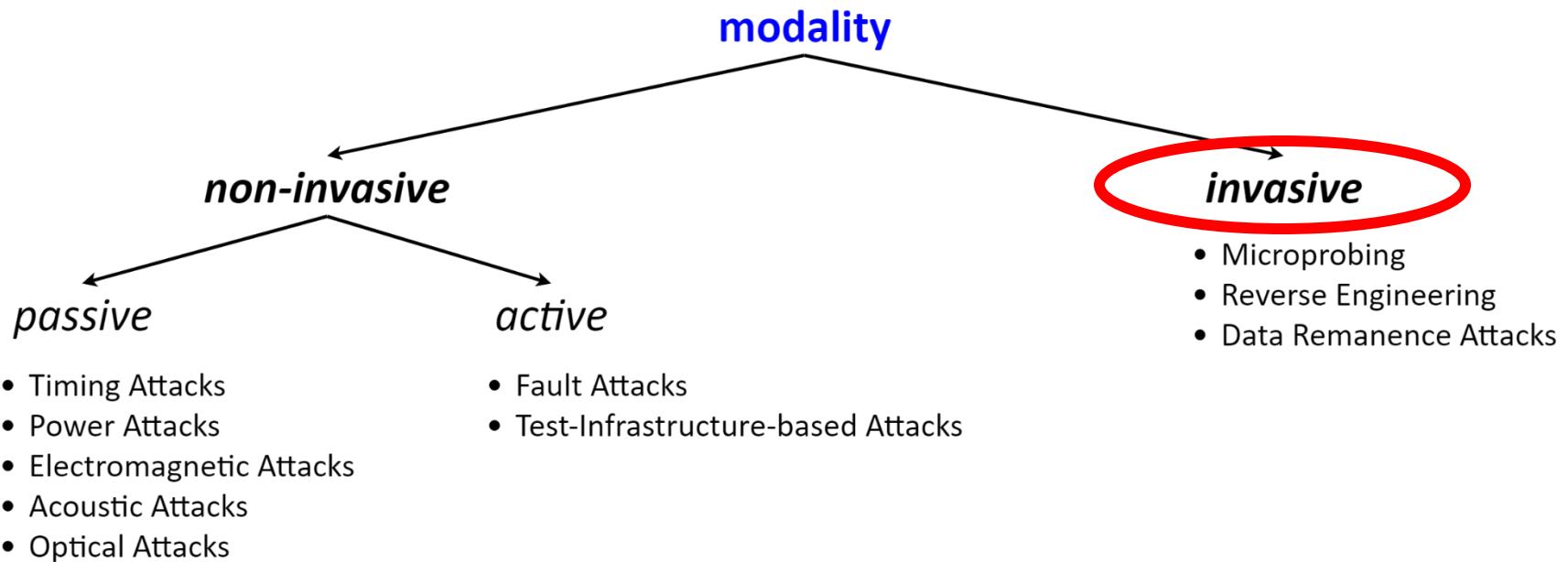


For further details, please refer to the lecture:

*HS\_2.1 - Vulnerabilities in test infrastructures*

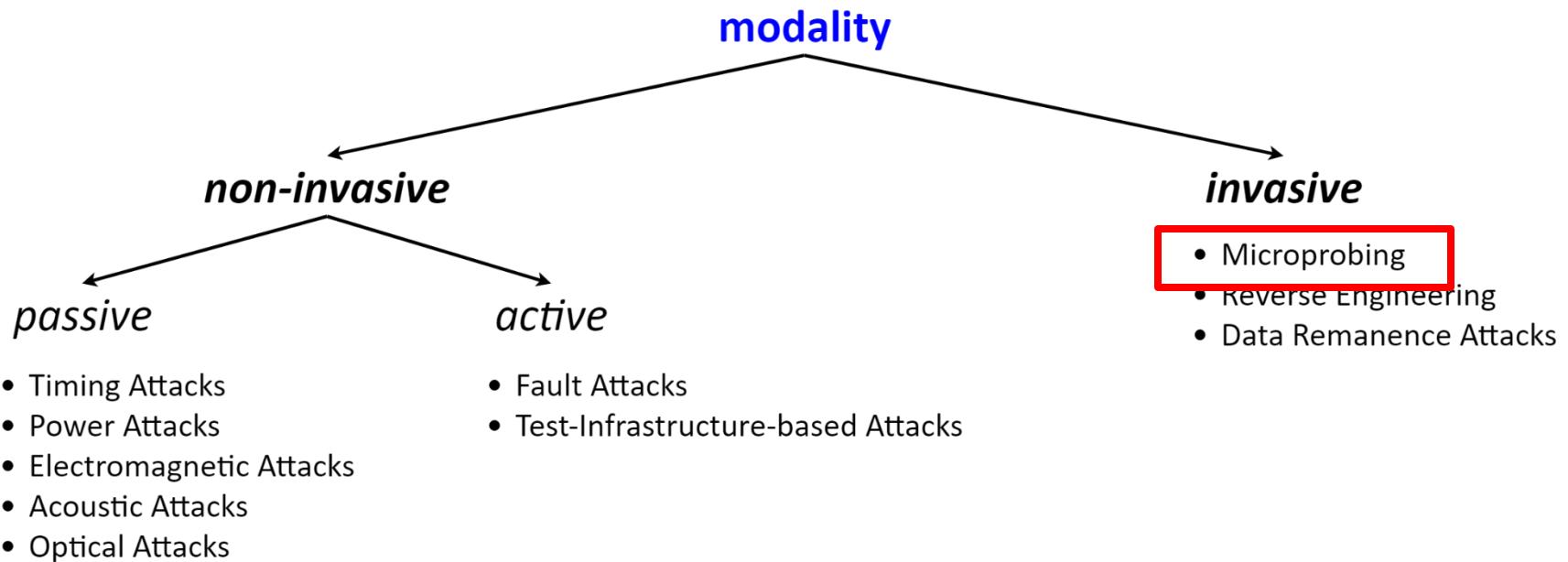
# Hardware Attacks Modalities

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# Hardware Attacks Modalities

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# Microprobing

## What

- Tries to extract information by measuring electrical quantities directly on the silicon die of the target device, once obtained physical access to it.

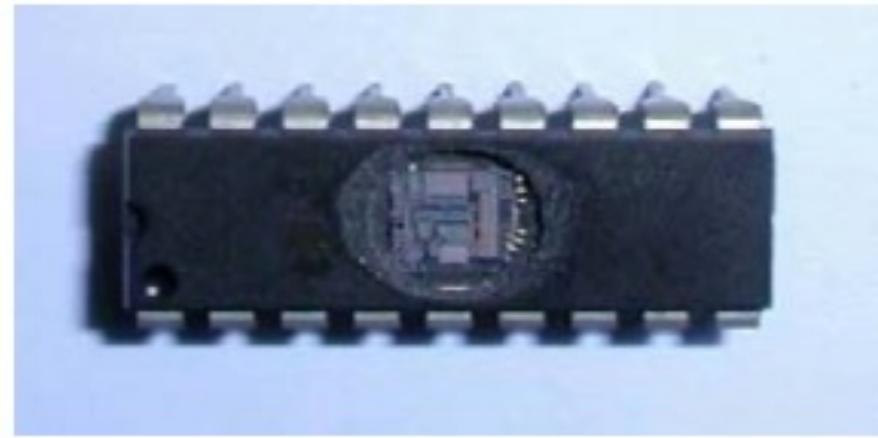
## How

- The die exposition is usually achieved by removing the plastic packages via chemical etching and/or by mechanical approaches.

# Example of depackaging

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- Microcontroller PIC16F84 original and depackaged

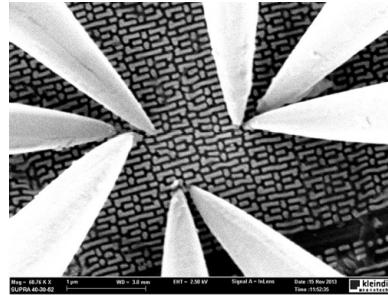
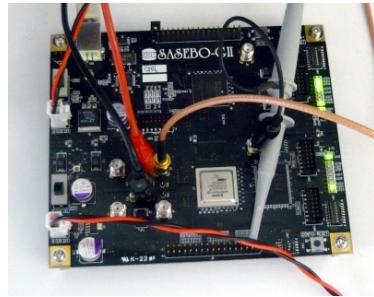
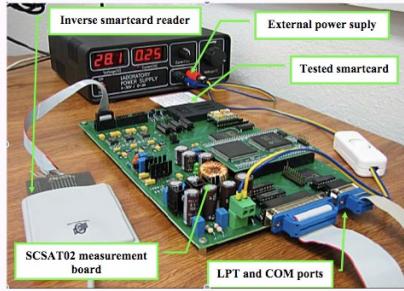


[B.S. Kaliski Jr. et al. (Eds.): CHES 2002, LNCS 2523, pp. 4, 2003]

# Implementation Attacks

- Microprobing-based attacks can, in turn, be
  - *Passive*: e.g., using microprobes to monitor a smartcard's bus while a program is executing:
  - *Active*: in this case signals may be also injected, the classic example being the use of a grounded microprobe needle on the clock line to the instruction latch to disable jump instructions.

# Implementation Attacks



# Example: Sat-TV Decoder

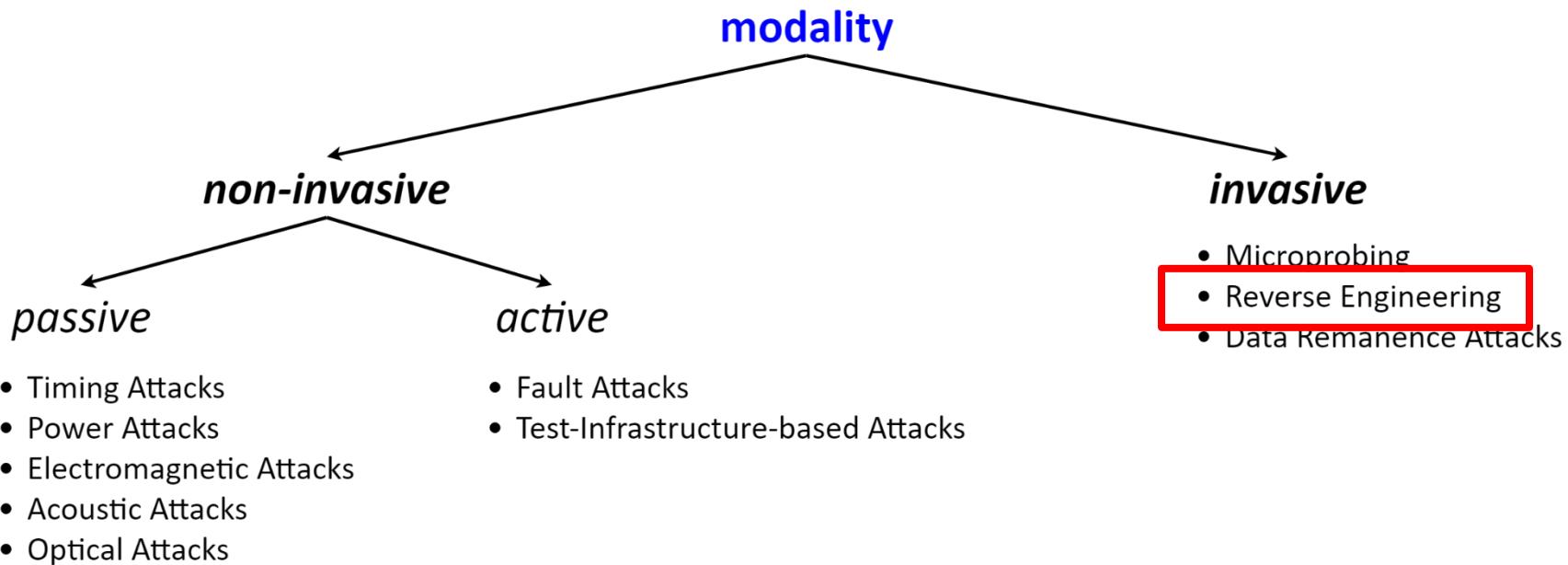
- Let's have a look at this video on YouTube:  
“How to Reverse-Engineer a Satellite TV Smart Card”

<https://www.youtube.com/watch?v=tnY7UVyaFiQ>



# Hardware Attacks Modalities

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# Reverse Engineering (RE) Attacks

- They aim at understanding the structure of a semiconductor device and its functions, i.e., at stealing the IP - *Intellectual Properties* - of the design

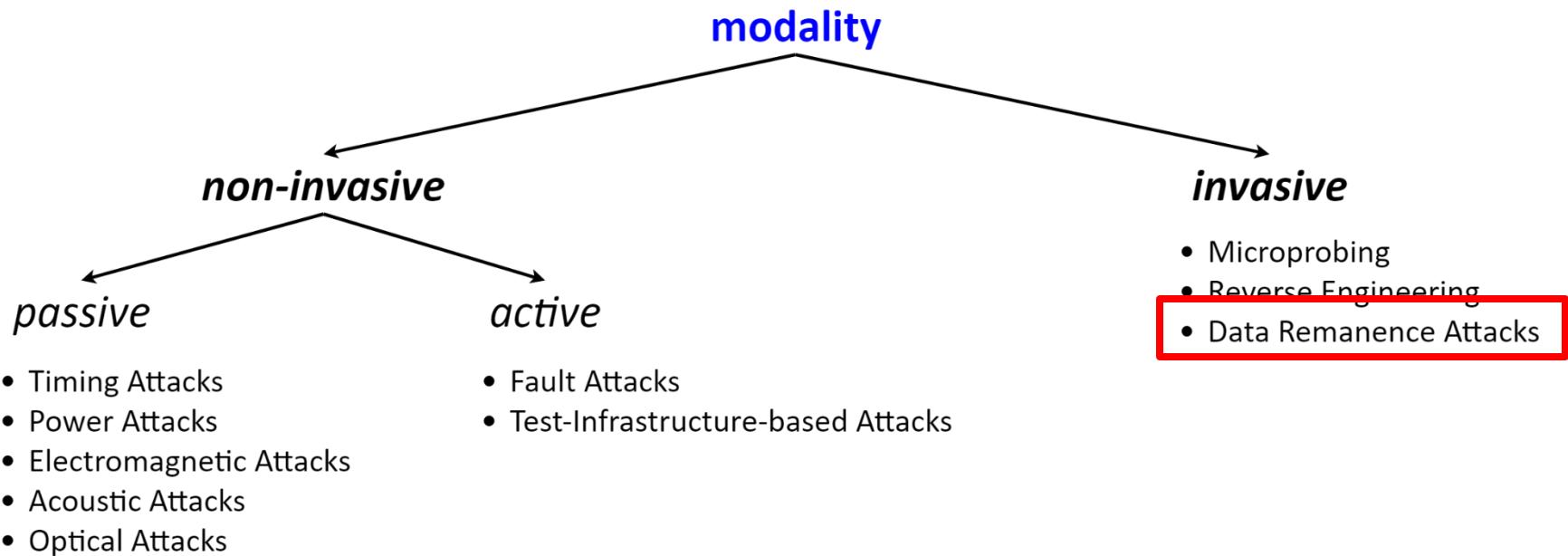
# Advanced tools for RE

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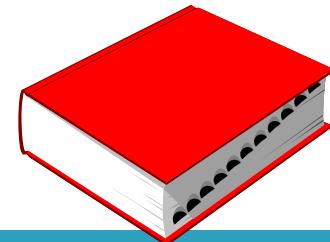
- A very interesting presentation on advanced tools for RE can be found here:
  - <https://media.hardware.io/integrated-circuit-offensive-security-olivier-thomas/>

# Hardware Attacks Modalities

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# Data remanence



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- Is the residual information remaining on storage media after clearing, i.e., after the foreseen actions to remove or erase the target data.

# Example of residues

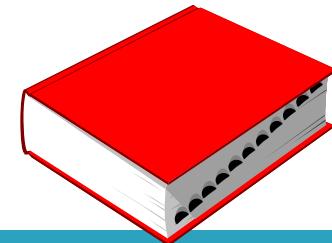
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- This residue may result from:
  - data being left intact by a nominal file deletion operation
  - by reformatting of storage media that does not remove data previously written to the media
  - through physical properties of the storage media that allow previously written data to be recovered

# Data Remanence Attacks

- An attacker with physical access to a computing system may just disconnect the memory chips and then retrieve their content at disconnection time by analyzing them off-line

# Cold boot attack



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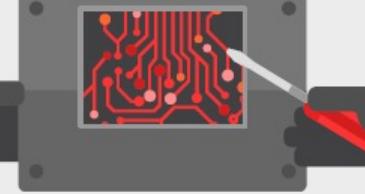
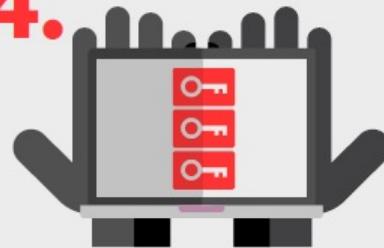
- A cold boot attack is a process for obtaining unauthorized access to a computer's encryption keys when the computer is left physically unattended.

# Consequences

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## ➤ Cold boot attack

Cold boot attacks can steal encryption keys from nearly any laptop

1.   
Attacker gets physical access to a company laptop
2.   
Attacker manipulates firmware settings
3.   
Attacker performs cold reboot into USB key
4.   
Attacker gets encryption keys from memory

# Cold boot attack

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- Contrary to popular assumption, DRAMs used in most modern computers retain their contents for several seconds after power is lost, even at room temperature and even if removed from a motherboard.
- Although DRAMs become less reliable when they are not refreshed, they are not immediately erased, and their contents persist sufficiently for malicious (or forensic) acquisition of usable full-system memory images.

[J. Alex Halderman et al:  
*"Lest We Remember: Cold Boot Attacks on Encryption Keys"*  
in Proc. 2008 USENIX Security Symposium]

- This video from the Center for Information Technology Policy demonstrates how a cold boot attack works:
  - <https://www.youtube.com/watch?v=JDaicPlgn9U>

# **Lest We Remember:**

## Cold Boot Attacks on Encryption Keys

[citp.princeton.edu/memory](http://citp.princeton.edu/memory)

# Important remark on volatility

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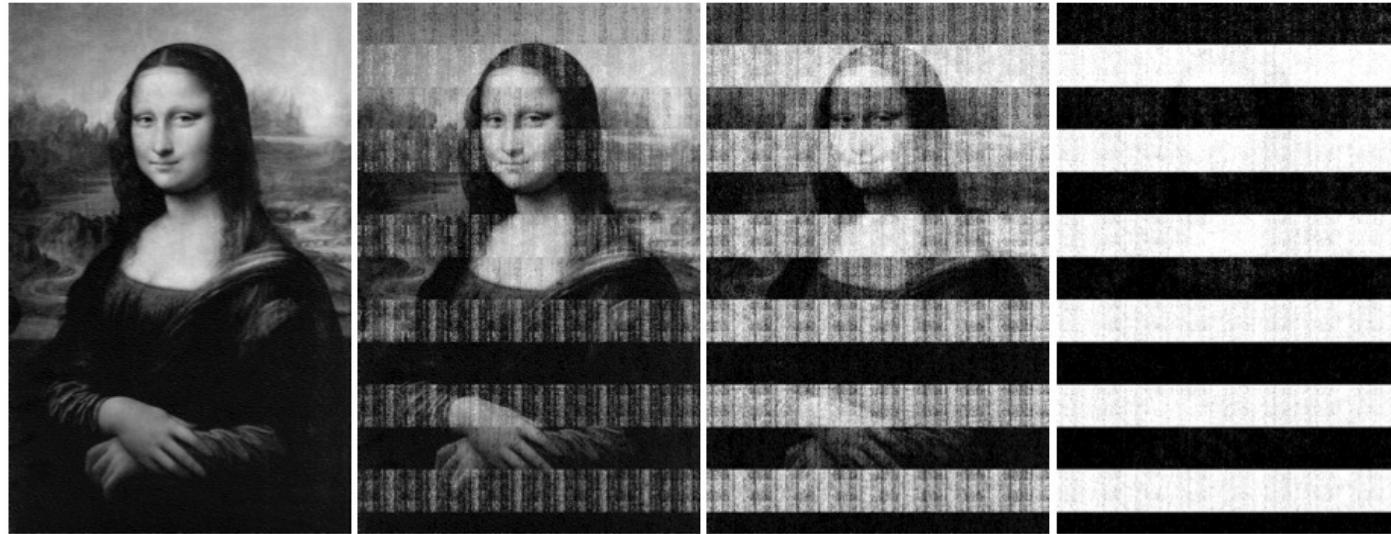


Figure 4: We loaded a bitmap image into memory on Machine A, then cut power for varying lengths of time. After 5 seconds (left), the image is indistinguishable from the original. It gradually becomes more degraded, as shown after 30 seconds, 60 seconds, and 5 minutes.

# Important remark on volatility

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Figure 5: Before powering off the computer, we spray an upside-down canister of multipurpose duster directly onto the memory chips, cooling them to  $-50^{\circ}\text{C}$ . At this temperature, the data will persist for several minutes after power loss with minimal error, even if we remove the DIMM from the computer.

# Conclusions

- Cryptography has 2,000+ years history and experience
- Hardware Security is still a young research field...

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