# Probability I Proof of Ergodic Theorem

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#### 1 Proof of Ergodic Theorem

**Note.** The ergodic theorem for finite state markov chain states the following. Let P be a the transition matrix. Then  $\lim_{n\to\infty}$  exists has constant columns and each row is made up of the unique stationary vectors of P. Furthermore, for any intial probability distribution v,  $\lim_{n\to\infty} P^n v = \pi$  the unique stationary distribution

*Proof.*  $\forall i, j \in \Omega$ , if  $x_n$  is irreducible, aperiodic, markov chain,

$$\lim_{n \to \infty} P[X_n = j \mid X_0 = i] = \pi_j.$$

where  $\pi$  is the stationary distribution vector. If  $x_0 \sim \pi$  i.e.  $P[x_0 = j] = \pi_j$ . What is the distribution of the following?

$$\begin{split} P[X_1 = k] &= \sum_{j \in \Omega} P[X_1 = k \mid X_0 = j] P[X_0 = j] \\ &= \sum_{j \in \Omega} P_{j,k} \pi \\ &= (\pi P)_k \cdot \pi_k. \end{split}$$

If the intial distribution is  $\pi$  then the distribution at all times is also  $\pi$ . Let's consider if  $v \neq \pi$ .

Let  $(X_n, Y_n)$  be two pariticles evolving with the same markov chain dynamics (i.e. same transition matrix) where  $X_0 = i$  and  $Y_0 \sim \pi$ . So  $X_n$  evolves in a complicated way but  $Y_n \sim \pi$ . At any time n, every state i chooses another state k with probability  $P_{i,k}$ , every particle that was at state i at times n moves to state k at time n+1 since the transition probability is the same for  $X_n$  and  $Y_n$ . This is known as a **coalescing walk**, when the two particles meet, they now have the same distribution at any time n+1 after.

Let  $T = \min\{n : X_n = Y_n\}$  at any time after  $T, X_n = Y_n$  since they follow the same set of instructions. Let's compute,

$$P[X_n = j \mid X_0 = i, T \le n] = P[Y_n = j \mid X_0 = i, T \le n] = \pi_j.$$

By law of total probability:

$$\begin{split} P[X_n = j \mid X_0 = i] &= P[X_n = j \mid n < T, x_0 = i] \cdot P[n < T \mid X_0 = i] + \\ P[X_n = j \mid n \geq T, X_0 = i] \cdot (1 - P[n < T \mid X_0 = i]) \\ &= P[X_n = j \mid n < T, X_0 = i] P[n < T \mid X_0 = i] + \pi_j (1 - P[n < T \mid X_0 = i]) \\ &= \pi_j + P[n < T \mid X_0 = i] (P[X_n = j \mid n < T, X_0 = i] - \pi_j) \\ P[X_n = j \mid X_0 = i] - \pi_j &= P[n < T \mid X_0 = i] (P[X_n = j \mid n < T, X_0 = i] - \pi_j) \\ \|P[X_n = j \mid X_0 = i] - \pi_j \| &= P[n < T \mid X_0 = i] \|(P[X_n = j \mid n < T, X_0 = i] - \pi_j) \| \\ &\leq P[n < T \mid X_0 = i] \end{split}$$

We want to show that  $P[n < T \mid X_0 = i] \to 0$  as  $n \to \infty$  since this implies that  $P[X_n = j \mid X_0 = i] = \pi_j$  and the random variable  $X_n$  approaches the stationary distribution despite the initial distribution.

Consider the  $|\Omega| \times |\Omega|$  joint state space  $(X_n, Y_n)$  that describes the markov chain dynamics of  $X_n, Y_n$  where

$$(i,j) = (X_n, Y_n).$$

Since P is irreducible and aperiodic, then we can reach any pair (i, j) from any other state (k, l). This means that P[ everything absorbed eventually] = 1. Equivalently,  $P[T < \infty] = 1$  and

$$\lim_{n \to \infty} P[n < T] = P[T = \infty] = 0.$$

and we conclude that  $||P[X_n = j \mid X_0 = i|| \to 0 \text{ as } n \to \infty.$  In fact,  $\exists p > 0 \text{ that } P[n < T] = P[n < T] \le (1-p)^n$ .

## 2 Mixing Times

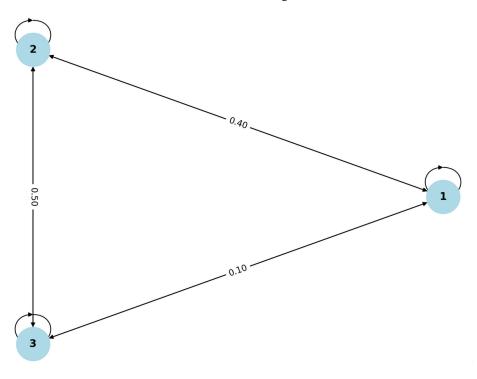
**Note.** How can we compare distributions? We know that eventually the distribution will converge to  $\pi$  but what is the distribution at n = 100, what about n = 1000? It will be almost  $\pi$  but not quite.

**Definition 1.** Let  $\Omega$  be finite, and u, v be two rwo vectors. WE can denote the **variation distance** as

$$d_{tv}(u, v) = \frac{1}{2} \sum_{i \in \Omega} ||u_i - v_i||.$$

Consider the following example:

#### Markov Chain Diagram



The transition matrix P is

$$P = \begin{bmatrix} 0.5 & 0.3 & 0.2 \\ 0.4 & 0.4 & 0.2 \\ 0.1 & 0.5 & 0.4 \end{bmatrix}, \quad \pi^{(0)} = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$$

Lets consider some intial vector  $v = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$  We can compute the variation distance after two steps (how far our distribution v is from the stationary distribution) as follows.

1. Solve for  $\pi P = \pi$ , we get

$$\pi = \begin{bmatrix} 0.267 \\ 0.4 \\ 0.333 \end{bmatrix}$$

2. Compute  $P^2$  and  $vP^2 = v^2$ .

$$P^2 = P \times P = \begin{bmatrix} 0.5 & 0.3 & 0.2 \\ 0.4 & 0.4 & 0.2 \\ 0.1 & 0.5 & 0.4 \end{bmatrix} \times \begin{bmatrix} 0.5 & 0.3 & 0.2 \\ 0.4 & 0.4 & 0.2 \\ 0.1 & 0.5 & 0.4 \end{bmatrix} = \begin{bmatrix} 0.39 & 0.37 & 0.24 \\ 0.38 & 0.38 & 0.24 \\ 0.29 & 0.43 & 0.28 \end{bmatrix}$$

, so we get  $v^2 = vP^2 = \begin{bmatrix} 0.39 & 0.37 & 0.24 \end{bmatrix}$ 

3. We get that

$$d_{tv}(v^2, \pi) = \frac{1}{2} \sum_{i=1}^{3} \left| v_i^{(2)} - \pi_i \right| = \frac{1}{2} \left( |0.39 - 0.267| + |0.37 - 0.4| + |0.24 - 0.333| \right) = 0.123$$

Thus, the total variation distance after 2 steps is approximately 0.123.

**Definition 2.** Define  $DTM(n) = \text{distance to mixing time at } n = \sup_i d_{tv}(v_n^i, \pi)$ . If  $v_n$  is the distribution of  $v_n$  with any initial conditions,  $d_{tv}(v_n, \pi) \leq DTM(n)$ .

We define the mixing time with threshold  $\varepsilon$  as

$$\min\{n \mid DTM(n) \le \varepsilon\}.$$

i.e. the minimum number of steps to get  $\varepsilon$  close to the stationary distribution  $\pi$ . The markov chain after this timestep is consider **mixed** 

Let's continue our previous example. What is the mixing time with  $\varepsilon = .1$ ?

Step 0: 
$$v^{(0)} = \begin{bmatrix} 1 & 0 & 0 \end{bmatrix}$$
 
$$d_{\text{TV}}(v^{(0)}, \pi) = \frac{1}{2} \sum_{i=1}^{3} \left| \pi_i^{(0)} - \pi_i \right| = 0.733$$

Step 1: 
$$v^{(1)} = v^{(0)}P = \begin{bmatrix} 0.5 & 0.3 & 0.2 \end{bmatrix}$$
  
 $d_{\mathrm{TV}}(v^{(1)}, \pi) = 0.233$ 

Step 2: 
$$v^{(2)} = v^{(1)}P = \begin{bmatrix} 0.39 & 0.37 & 0.24 \end{bmatrix}$$
  
 $d_{\text{TV}}(v^{(2)}, \pi) = 0.123$ 

Step 3: 
$$v^{(3)} = v^{(2)}P = \begin{bmatrix} 0.35 & 0.39 & 0.26 \end{bmatrix}$$
  
 $d_{\text{TV}}(v^{(3)}, \pi) = 0.1$ 

Step 4: 
$$v^{(4)} = \pi^{(3)}P = \begin{bmatrix} 0.33 & 0.4 & 0.27 \end{bmatrix}$$
  
 $d_{\text{TV}}(v^{(4)}, \pi) = 0.0953$ 

Thus, the mixing time is 4.

#### **Theorem 1.** Perron-Frobenius Theorem

If P is the transition matrix of an irreducible markov chain then  $\lambda=1$  is always an eigenvalue with row-eigenvalue  $\pi$  of strictly positive entries. All other eigenvalues  $\{\lambda_i\}_{i=2}^n$  have  $|\lambda|<1$ . If  $\alpha$  is a row eigenvector of P with eigenvalue  $\neq 1$ , then  $\sum_{i=1}^n \alpha_i = 0$ .

Note. Eigenvalues can be complex.

**Corollary 1.** Any vector P with  $p_i \geq 0$  and  $\sum_{i=1}^n P_i = 1$  can be written as  $P = \pi + \sum_{i=1}^n c_i \vec{\alpha_i}$  hwere  $c_i \in \mathbb{R}$  and  $\alpha_i$  are row vectors of P with eigenvalue  $\neq 1$ .

### 3 Card Shuffling

Consider the n card deck is shuffled by moving the top card to a uniformly chosen random position.  $\Omega =$  arrangement of n cards. Clearly,  $|\Omega| = n!$ . Every row of  $P_{|\Omega| \times |\Omega|}$  would have n positive entries, with

probability to be  $\frac{1}{n}$  and all entries 0. It will be a very sparse matrix.

It turns out that this matrix is doubly stochastic! The stationary distribution is just the uniform distribution. By the ergodic theorem, this deck is well-shuffled, because  $n \to \infty$ ,  $n = \pi$  where  $\pi$  is the uniform distribution.