



# Graphs

## Graphs (Part 1)



1 → red  
0 → yellow

2 → white

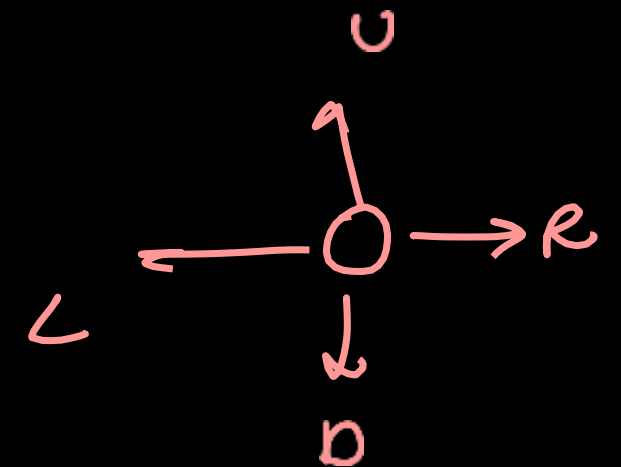
grid based

0	<u>1</u> 2	<u>1</u> 2	<u>1</u> 2
1	<u>1</u> 2	<u>1</u> 2	0
2	2 <u>1</u>	0	1
	0	1	2

sr = 1  
 → source row

sc = 1  
 → source col

color = d

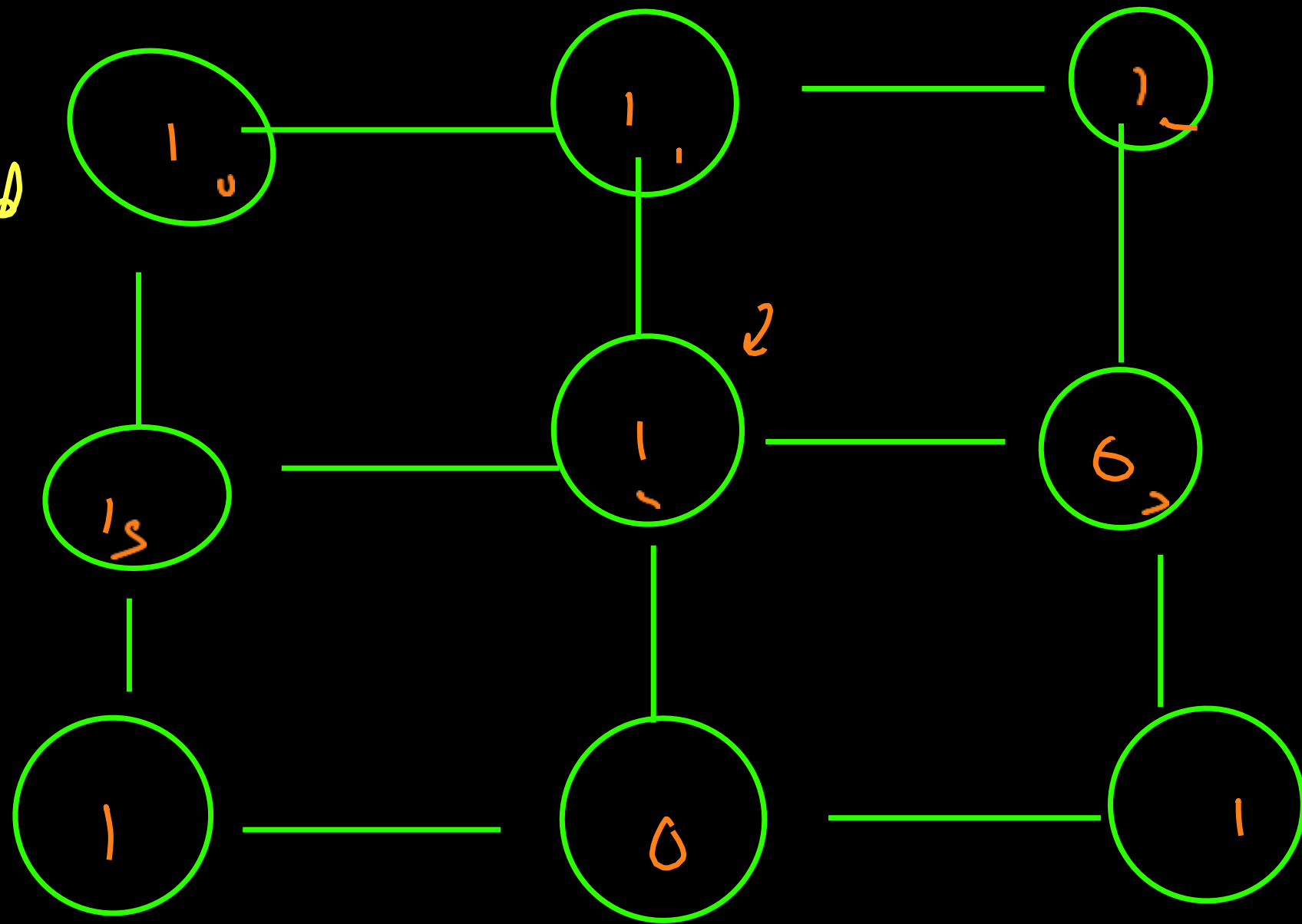


# Problem is giving an intuition of recursion

↳ dfs

↳ bfs

most of the  
grid based dfs/  
bfs can be solved  
without creating  
a graph of  
the grid.

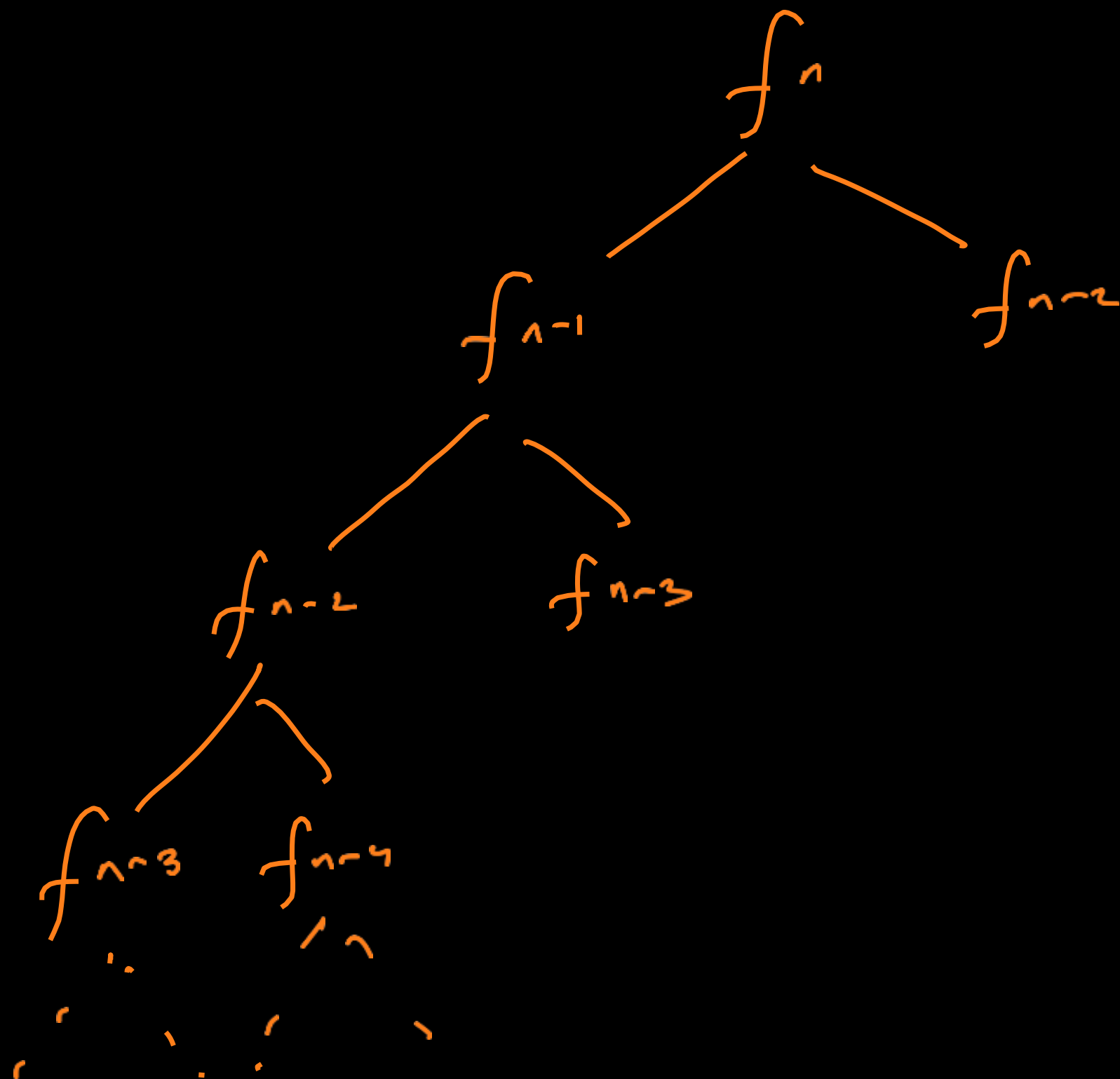


DFS

united color = 1

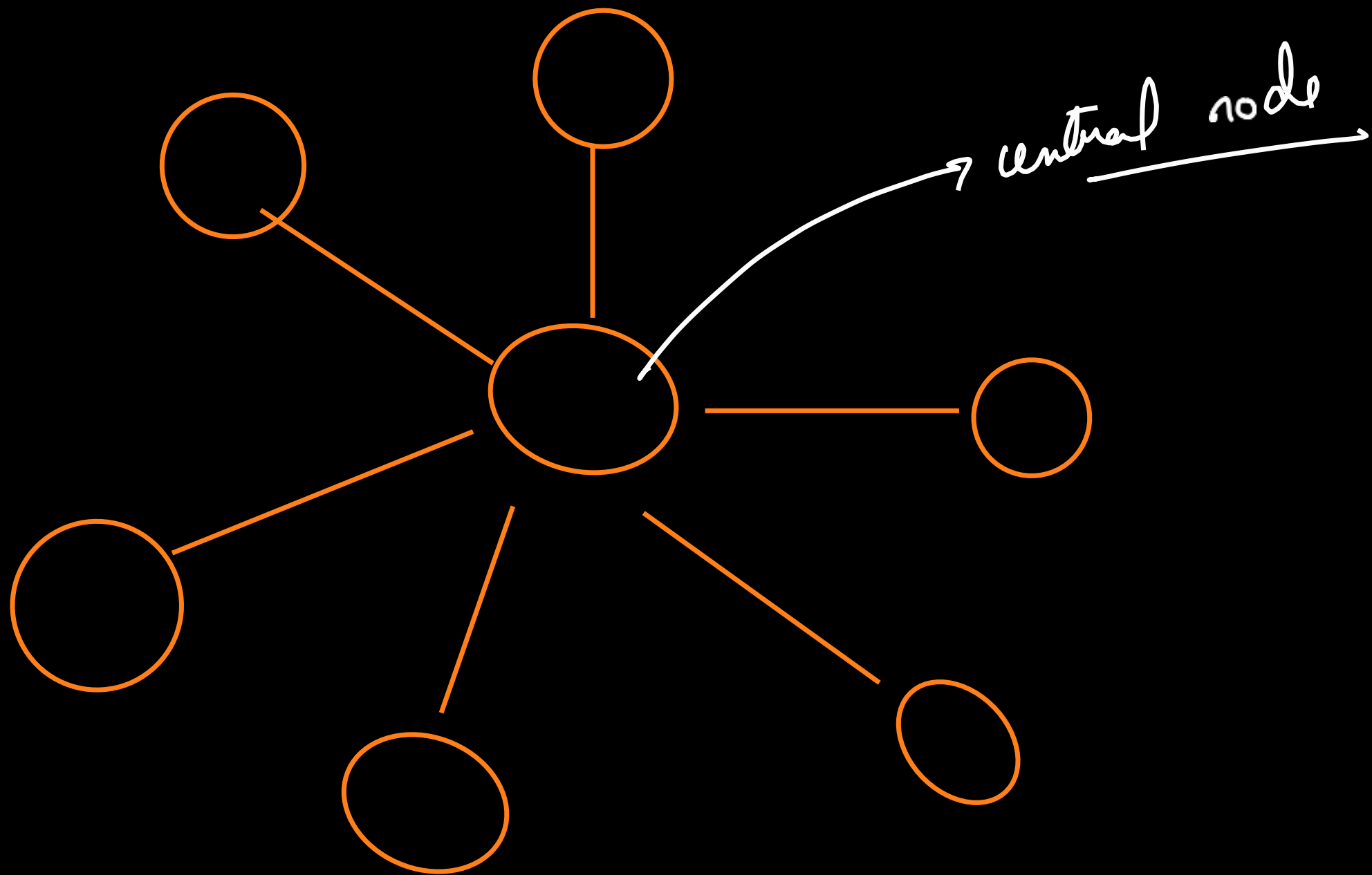
	0	1	2
0	2	2	2
1	2	2	0
2	2	0	1

$$f_n = f_{n-1} + f_{n-2}$$



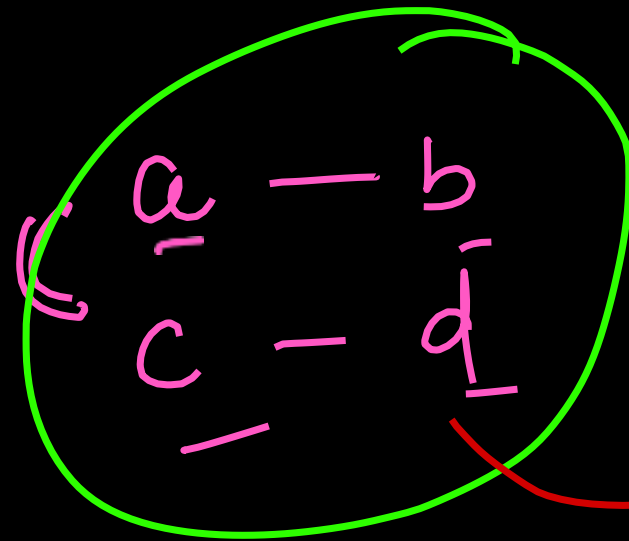
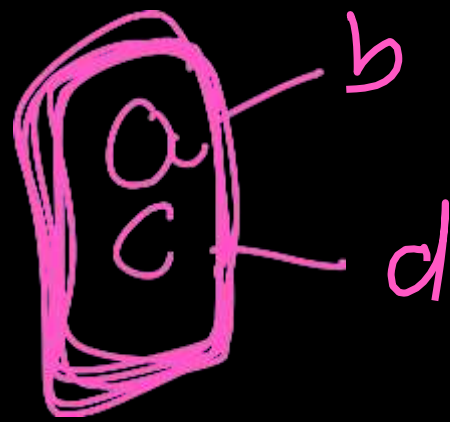
Generally a recursion  
sol<sup>n</sup> is more or  
less DfS

# Leetcode 1791 → Brentford → Calculate frequency of  
vertices in the edge list.  $O(V+E)$



↳ pick any 2 edges, my graph is a star

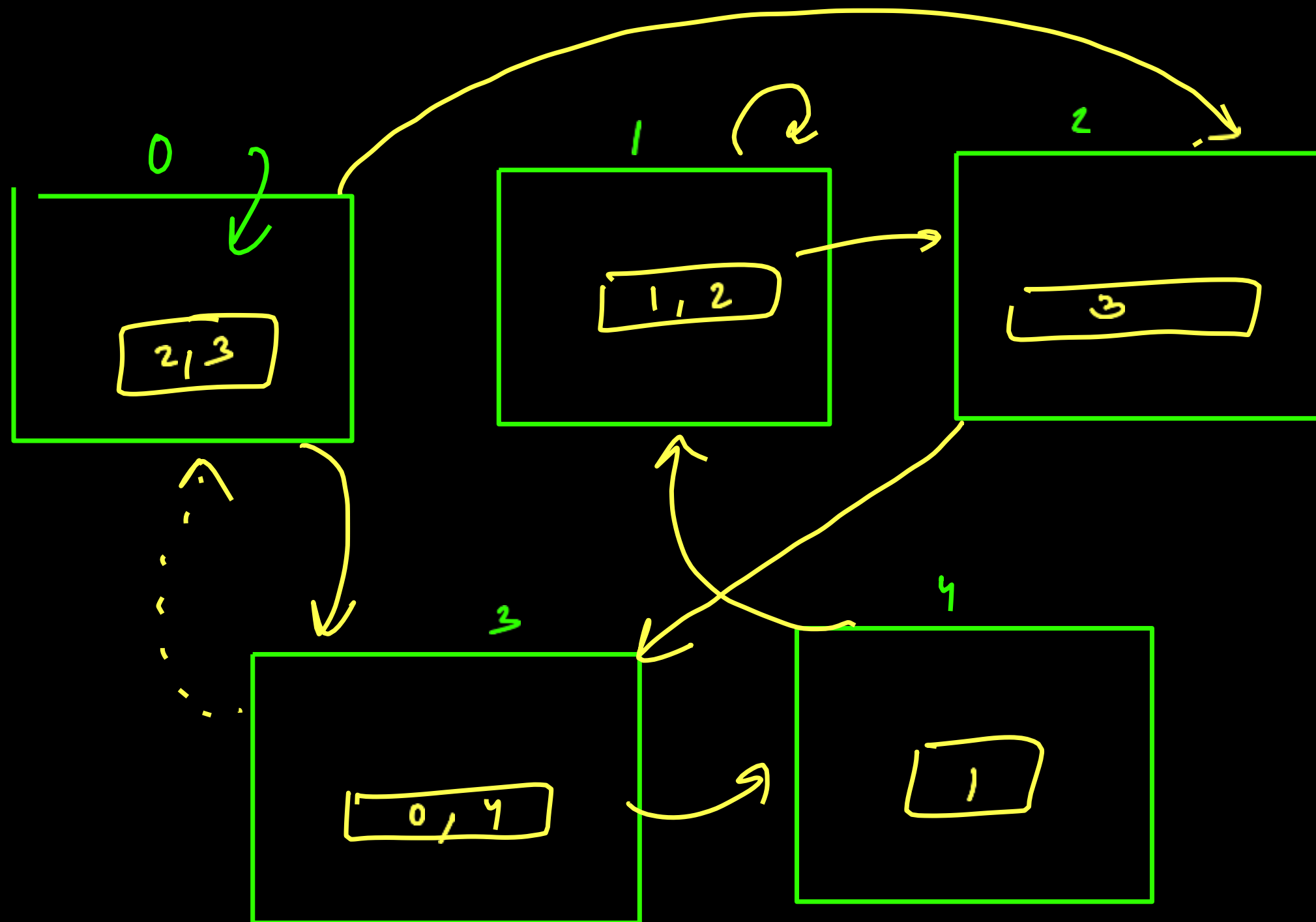
graph-



3 uniq

one of the nodes is central

$$(a == c \parallel b == c) \quad ? \quad c \circ d$$



the set of keys inside a room tell us the neighbours  
of the node (room).

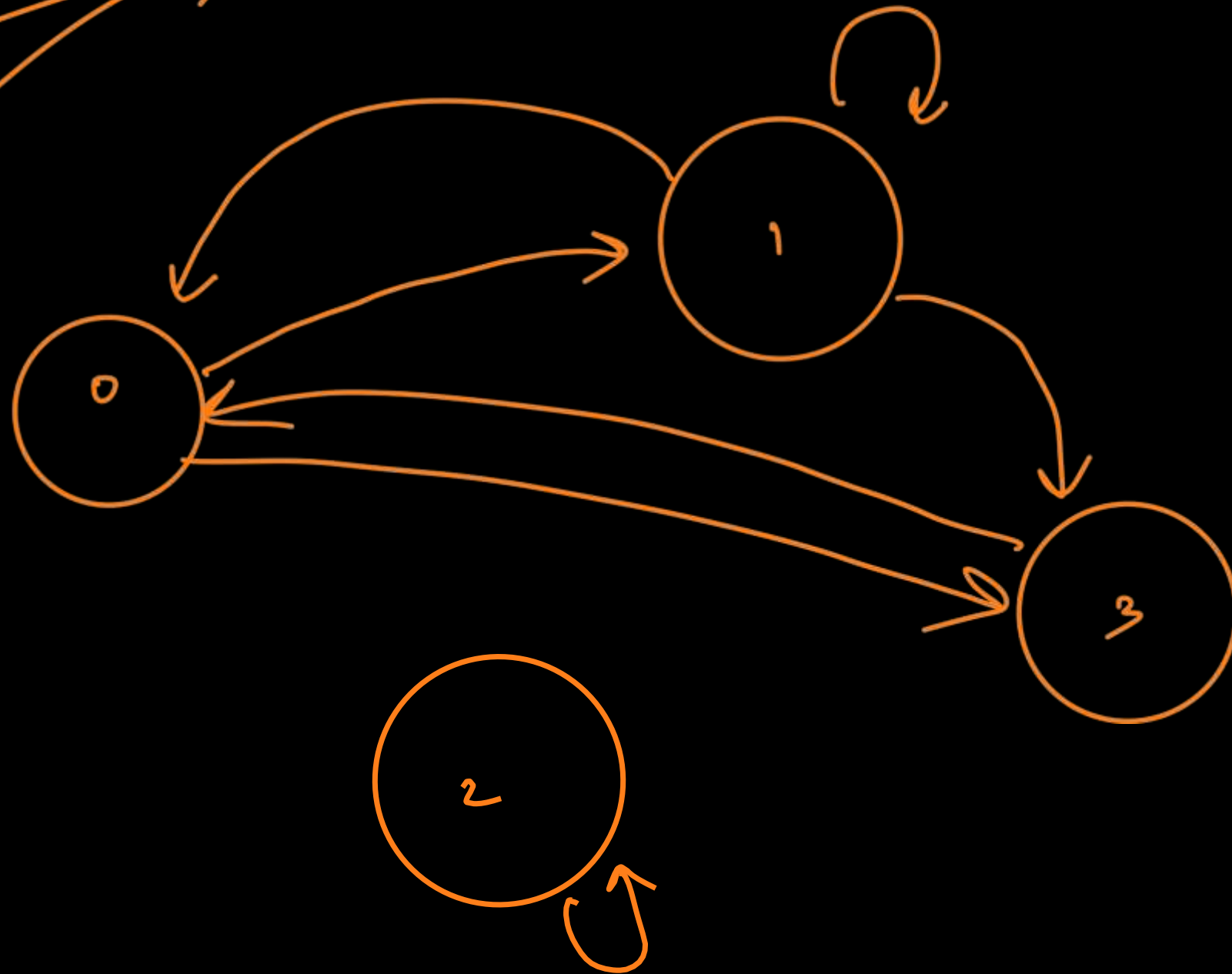


<sup>0</sup>      <sup>1</sup>      <sup>2</sup>      <sup>3</sup>  
[[1, 3]] [[3, 0, 1]] [[2]] [[0, 2]]

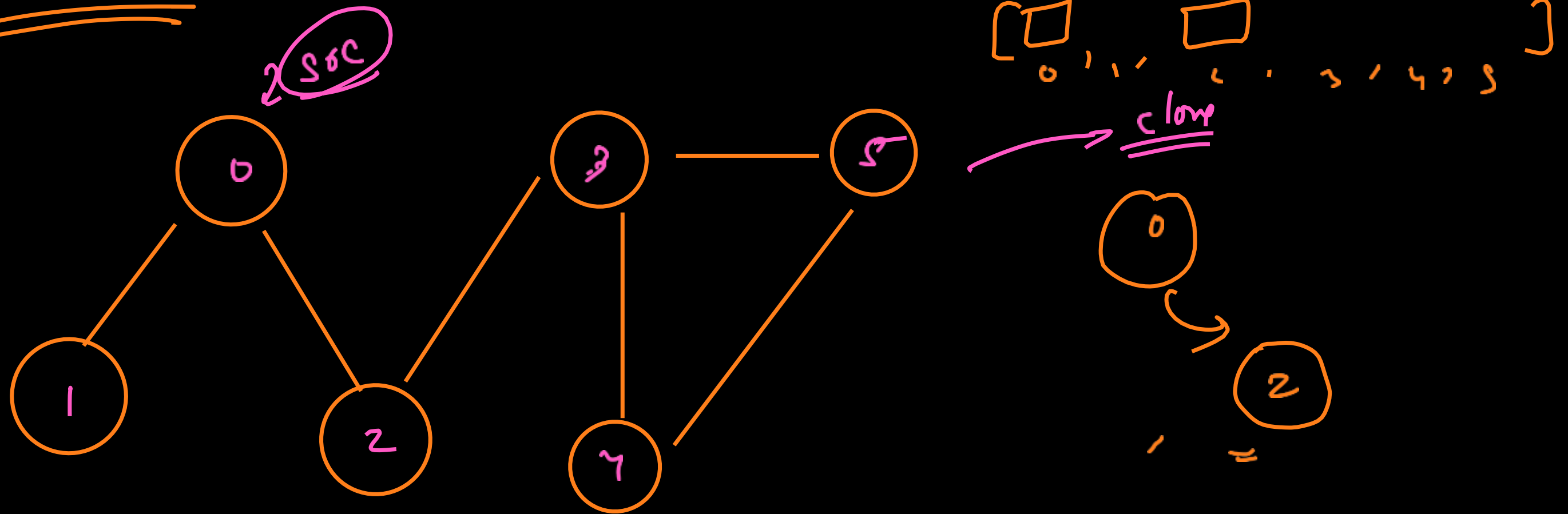
adj list

Read the graph

visited



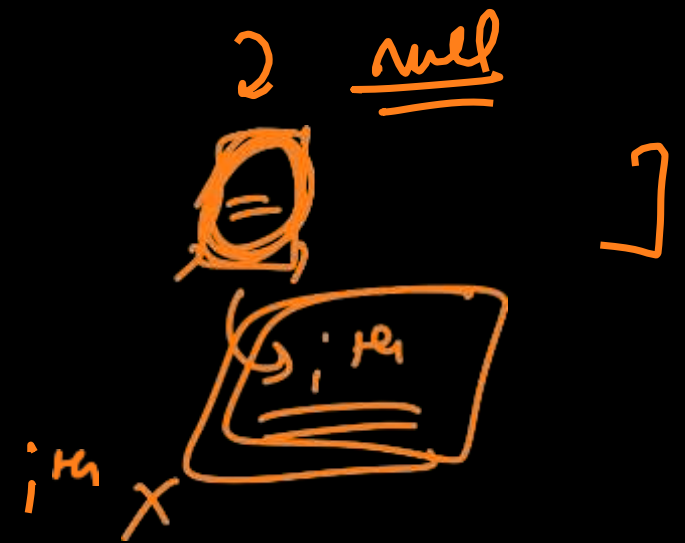
# Lecture 133

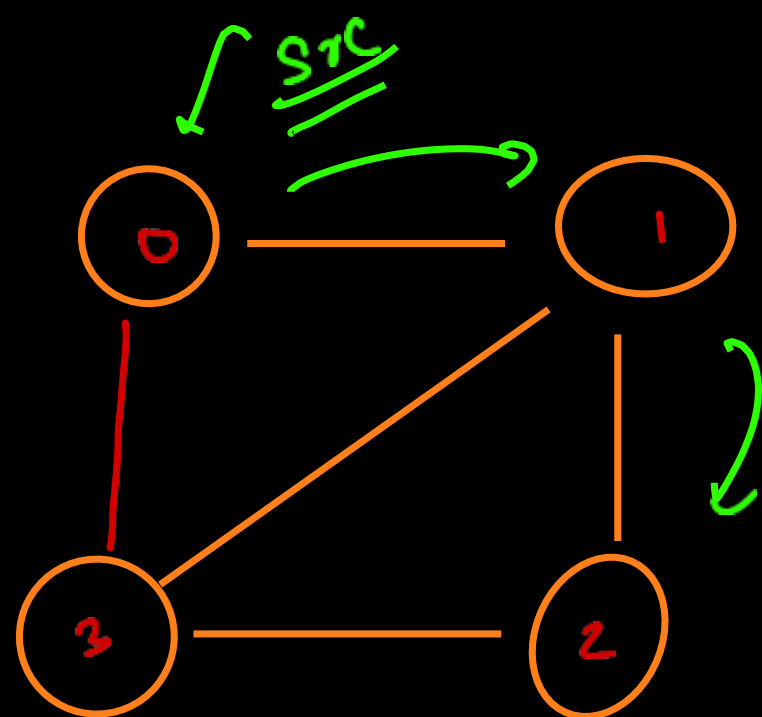


How to identify if we already created a node <?

vector < Node\* > nodeRegister

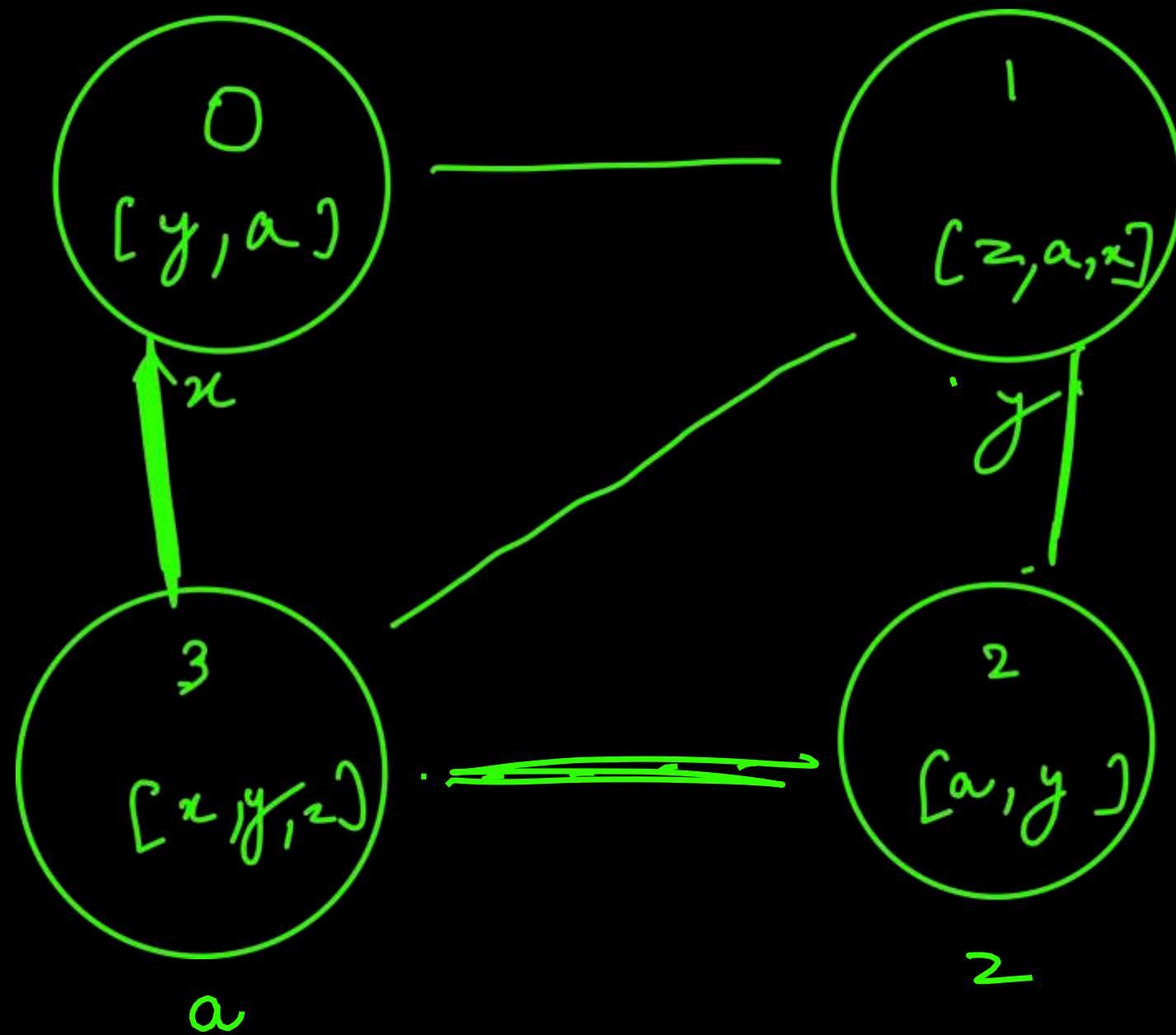
nodeRegister[i] → address of the  
newly created i<sup>th</sup> node





$[ \begin{matrix} x & y & z & a \\ 0 & 1 & 2 & 3 \end{matrix} ]$ 
  
 node key

dfs(node, clare)



dfs(node, clone)

for (neigh: node)

if (!nodeKeySet {neigh.val}) {

—— count

—— me

—— Rem

} else

—— neigh.all

}



▶ **THANK YOU** ◀