(I) Give a formal definition with any notations for the following:

### - DFA as a 5 tuple

- Deterministic Finite Accepter
- M = (Q,  $\Sigma$ ,  $\delta$ ,  $q_0 F$ ) where
  - Q = Set of states
  - $\Sigma$  = input alphabet
  - $-\Delta$  = transition function
  - q<sub>0</sub> = initial state
  - F = set of finite states  $F \subseteq Q$

### - Language accepted by automaton

- Let M be an Automaton
- L(M) is accepted by M if it contains all input strings accepted by M
- L(M) = {  $w \in \Sigma^*$ :  $\delta(q_0, w) \in F$  } for a DFA M = (Q,  $\Sigma$ ,  $\delta$ ,  $q_0 F$ )

#### - Regular language

- A language is regular if there is a DFA M such that L = L(M)

# (II) Given the alphabet as {0, 1}, write a DFA for the following three regular languages. (Give the complete description of the DFA, and also as a transition graph)

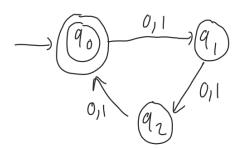
(i) L = {w | w is a string of even length}

$$\Xi = \{0,1\}$$

$$\longrightarrow (90) (91)$$

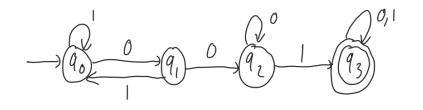
$\mathcal{W}$	[w]	8
人	0	90
0	1	9,
)	)	91
01	2	90
10	2	$q_o$
010	3	91

(ii) 
$$L = \{w \mid |w| \mod 3 = 0\}$$



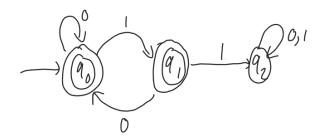
W	$ \omega $	8
$\sqrt{}$	0	90
O	1	91
[0	2	92
010	3	90
1010	4	9,

### (iii) L = {w | w contains the string 001 as a substring}



8	0	
90	91	90
$q_{l}$	92	90
92	92	93
93	93	93

### (iv) L = {w | w does not contain two consecutive 1's}



(III)Describe the extended transition function on a string recursively using transitions one symbol at a time for the automaton of following problems.

## (i) A string (of length >=4) that belongs to the language

i) 
$$L = \{ w \mid |w| \text{ is even } \}$$

0000
$$S^*(q_0,000),0)$$

$$S(S(S^*(q_0,00),0),0)$$

$$S(S(S(S(q_0,\lambda),0),0),0),0)$$

$$S(S(S(S(q_0,\lambda),0),0),0),0)$$

$$S(S(S(q_0,0),0),0)$$

$$S(S(q_0,0),0)$$

$$S(q_1,0) = q_0 \in F$$

ii) 
$$L = \{ w \mid |w| \% 2 = 0 \}$$
 $001100$ 
 $S * (q_0, 001100, 0)$ 
 $S(S(S(S), 0011), 0), 0)$ 
 $S(S(S(S), 0011), 0), 0)$ 
 $S(S(S(S(S), 001), 1), 0), 0)$ 
 $S(S(S(S(S), 000), 1), 1), 0), 0)$ 
 $S(S(S(S(S), 000), 0), 0)$ 
 $S(S(S(S), 000), 0)$ 
 $S(S(S(S), 000), 0)$ 
 $S(S(S(S), 000), 0)$ 
 $S(S(S(S), 000), 0)$ 

iii) W contains 001

0001

$$\delta^*(q_0,0001)$$
 $\delta(\delta^*(q_0,000),1)$ 
 $\delta(\delta(\delta^*(q_0,00),0),1)$ 
 $\delta(\delta(\delta(\delta(q_0,\lambda),0),0),0),1)$ 
 $\delta(\delta(\delta(\delta(q_0,\lambda),0),0),0),1)$ 
 $\delta(\delta(\delta(q_1,0),0),0),1)$ 
 $\delta(\delta(q_2,1)=q_3\in F$ 

iv) W does not contain 2 consecutive 1's 
$$1010$$
 $S * (90,1010)$ 
 $S(S * (90,101),0)$ 
 $S(S(S * (90,10),1),0)$ 
 $S(S(S(S * (90,1),0),1),0)$ 
 $S(S(S(S (90,1),0),1),0)$ 
 $S(S(S(90,1),0),1),0)$ 
 $S(S(S(90,1),0)$ 
 $S(S(90,1),0)$ 
 $S(90,1),0$ 

(II) A string (of length >=4) that does not belong to the language

i) 
$$00000$$
 $S^*(q_0,00000)$ 
 $S(S^*(q_0,0000),0)$ 
 $S(S(S(S^*(q_0,000),0),0)$ 
 $S(S(S(S(S^*(q_0,00),0),0),0),0)$ 
 $S(S(S(S(S(S(q_0,\lambda),0),0),0),0),0)$ 
 $S(S(S(S(S(q_0,\lambda),0),0),0),0)$ 
 $S(S(S(S(q_0,0),0),0),0)$ 
 $S(S(S(S(q_0,0),0),0),0)$ 
 $S(S(S(q_0,0),0),0)$ 
 $S(S(g(q_0,0),0),0)$ 
 $S(g(g,0)=g_1 \notin F$ 

ii) 
$$001/0$$
 $S^*(q_0,00110)$ 
 $S(S^*(q_0,0011),0)$ 
 $S(S(S^*(q_0,001),1),0)$ 
 $S(S(S(S(S^*(q_0,0),0),1),1),0)$ 
 $S(S(S(S(S(q_0,\lambda),0),0),1),1),0)$ 
 $S(S(S(S(S(q_0,\lambda),0),0),1),1),0)$ 
 $S(S(S(S(q_0,1),0),0)$ 
 $S(S(q_0,1),0)$ 
 $S(S(q_0,1),0)$ 
 $S(S(S^*(q_0,010),1)$ 
 $S(S(S^*(q_0,010),1))$ 
 $S(S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(q_0,\lambda),0),1),0),1)$ 
 $S(S(S(q_0,\lambda),0),1),0),1)$ 

iv) 
$$0110$$

$$S^*(q_0, 0110)$$

$$S(S^*(q_0, 011), 0)$$

$$S(S(S^*(q_0, 01), 1), 0)$$

$$S(S(S(S^*(q_0, 0), 1), 1), 0)$$

$$S(S(S(S(q_0, \lambda), 0), 1), 1), 0)$$

$$S(S(S(q_0, 1), 1), 0)$$

$$S(S(q_1, 1), 0)$$

$$S(S(q_1, 1), 0)$$

$$S(q_1, 0) = q_2 \notin F$$