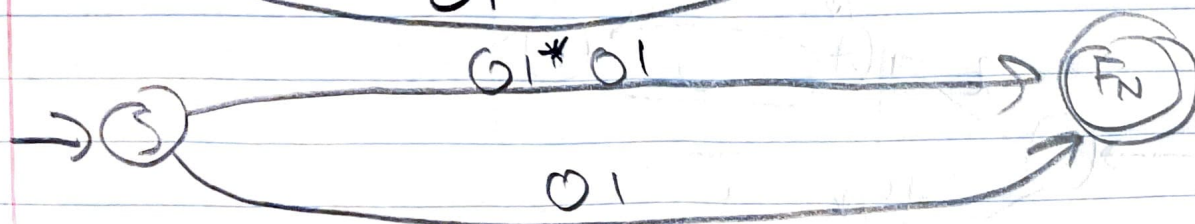
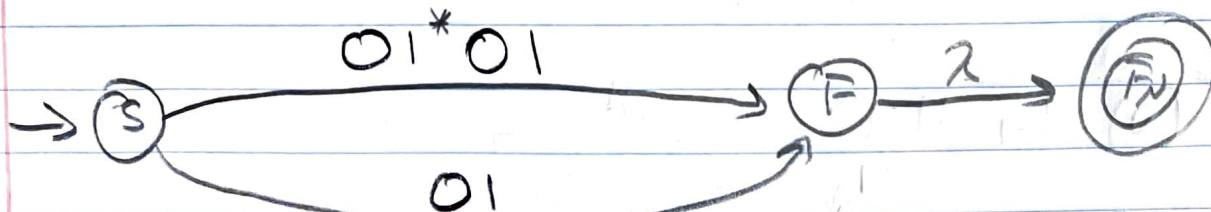
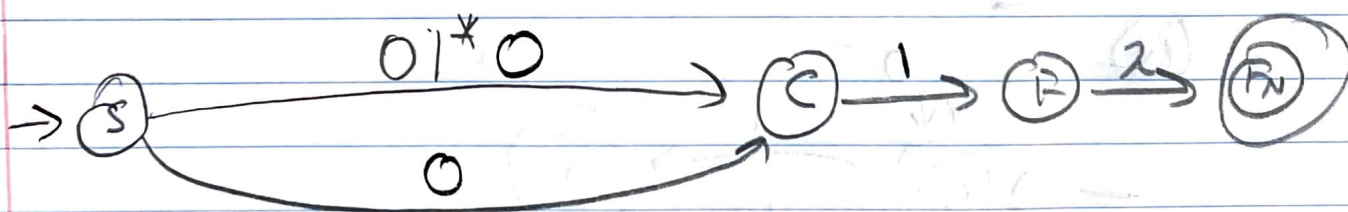
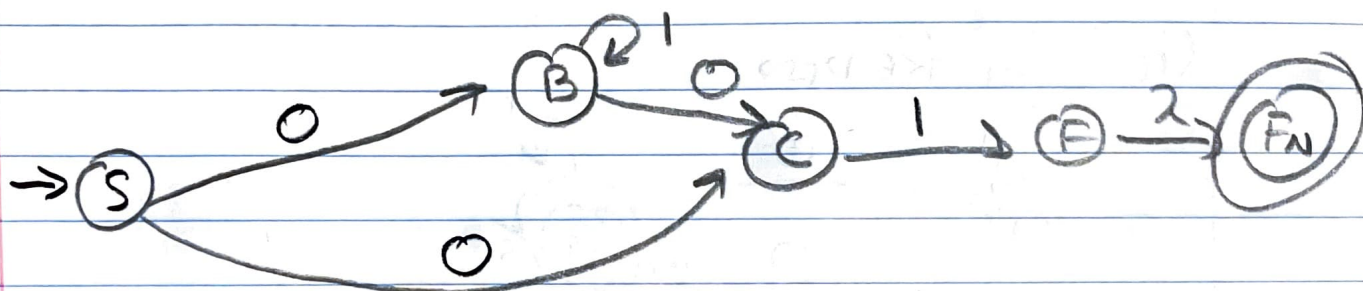
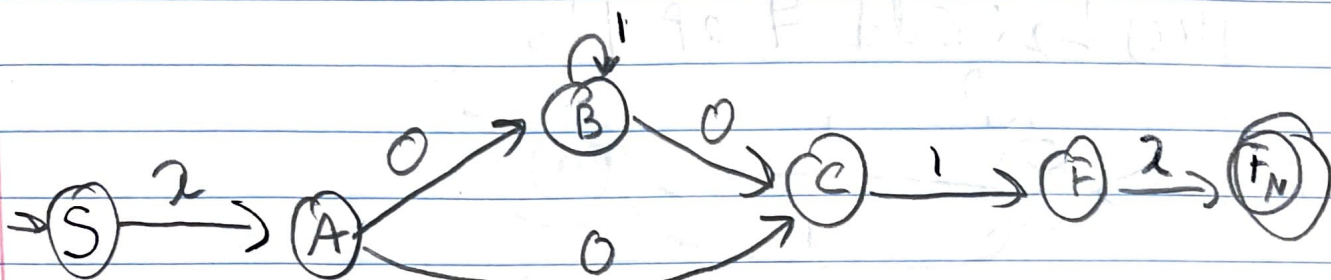
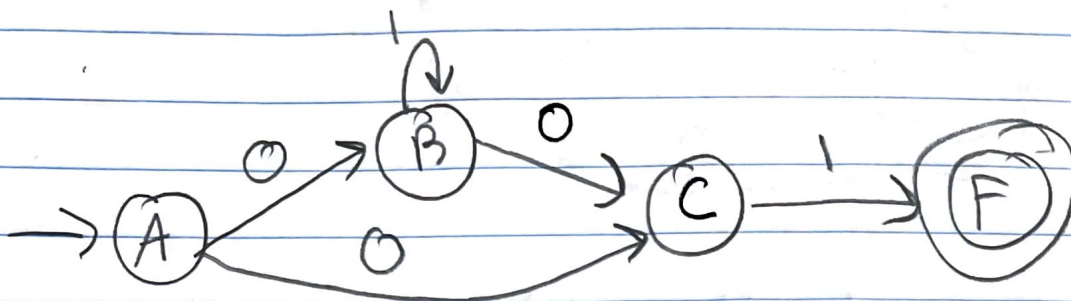


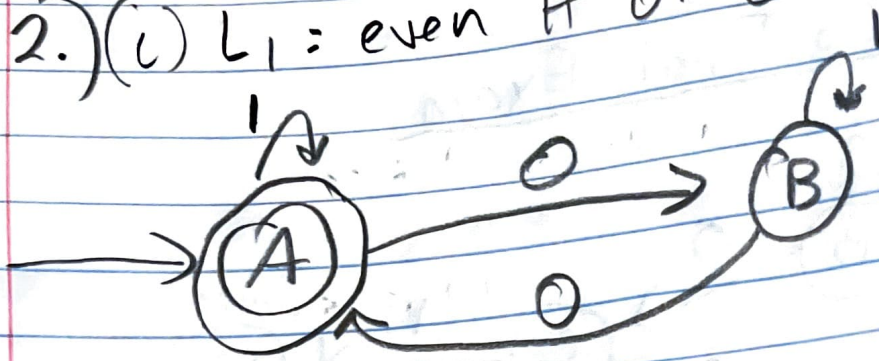
CS 3186 Final Exam

1.)

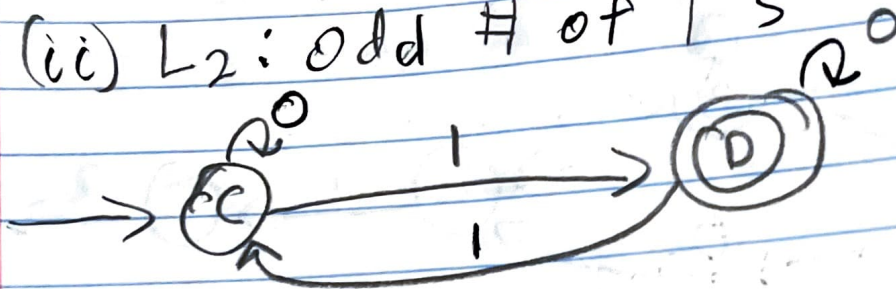


Regular expression: $(01^*0 + 0)1$

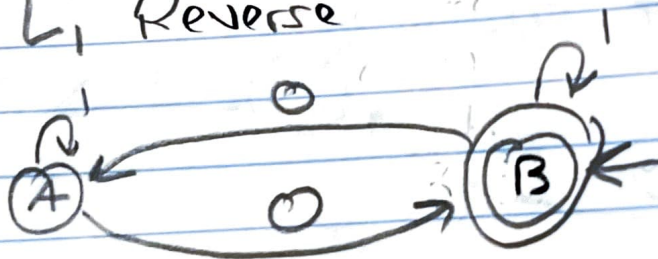
2.) (i) L_1 : even # of 0's



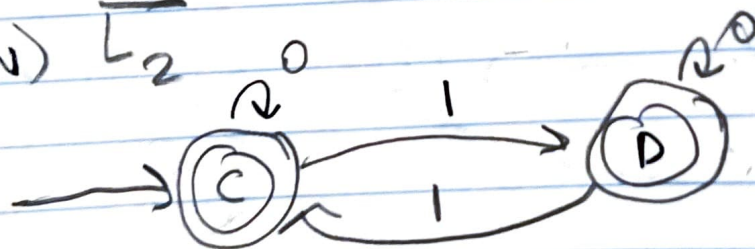
(ii) L_2 : odd # of 1's



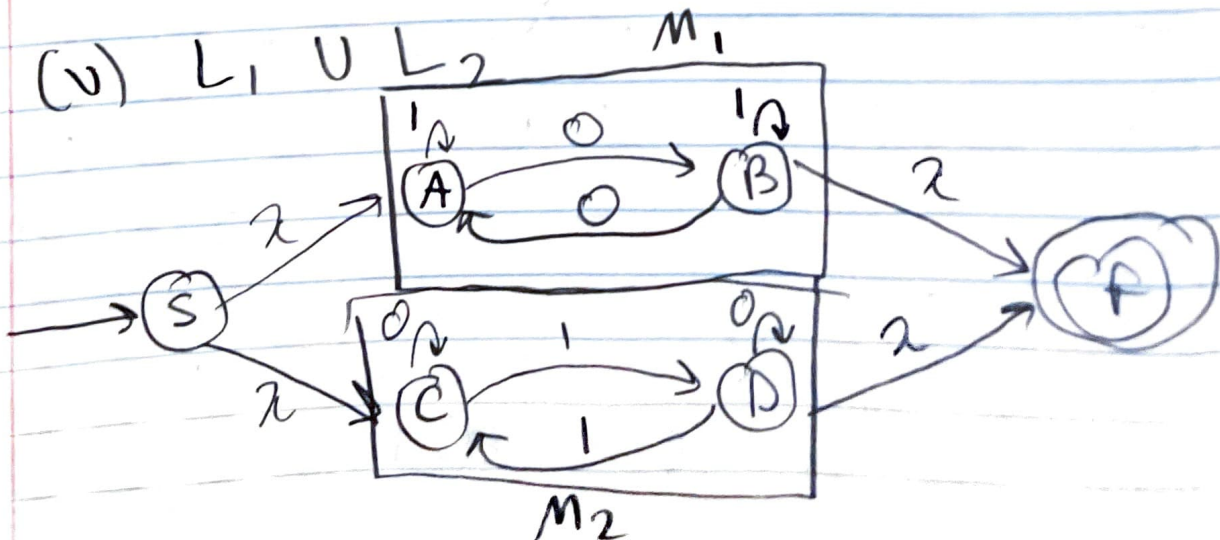
(iii.) L_1 Reverse



(iv) $\overline{L_2}$



(v) $L_1 \cup L_2$



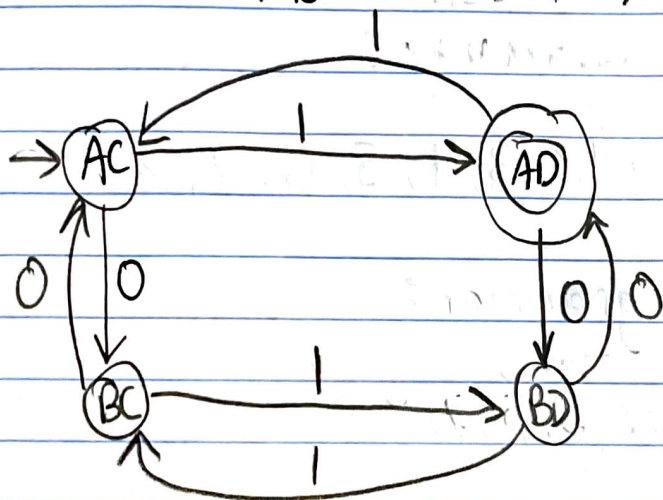
(vi) $L_1 \cap L_2$

$$\delta(AC, 0) = BC \quad \delta(AC, 1) = AD$$

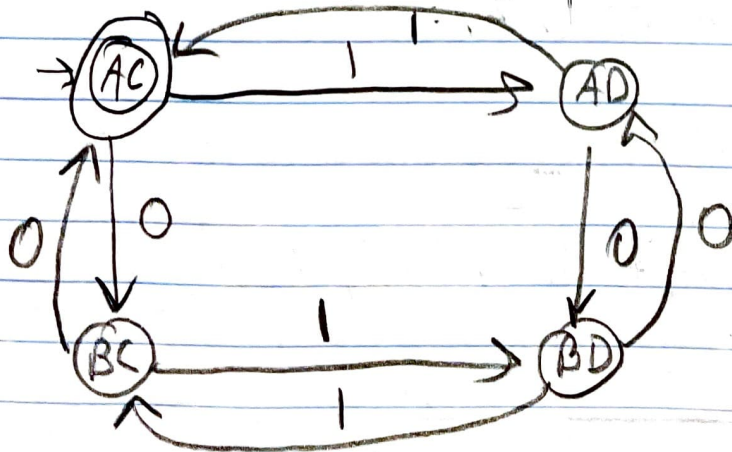
$$\delta(AD, 0) = BD \quad \delta(AD, 1) = AC$$

$$\delta(BC, 0) = AC \quad \delta(BC, 1) = BD$$

$$\delta(BD, 0) = AD \quad \delta(BD, 1) = BC$$



(vii) $L_1 - L_2 = L_1 \cap \overline{L_2}$



- 3.)
1. $S \rightarrow OA$
 2. $S \rightarrow IA$
 3. $S \rightarrow A\emptyset$
 4. $S \rightarrow AB$
 5. $S \rightarrow \emptyset$
 6. $S \rightarrow OA1$
 7. $A \rightarrow 1$
 8. $S \rightarrow \lambda$

(i) Context-free grammar:

Production rules: 1, 2, 3, 4, 5, 6, 7, 8

(ii) Left-linear grammar:

Production rules: 3, 5, 7, 8

(iii) Right Linear regular grammar:

Production rules: 1, 2, 5, 7, 8

(iv) Linear grammar:

Production rules: 1, 2, 3, 5, 7, 8

4.) Use pumping lemma to show $L = \{a^m b^m c^{m+1}\}$

Given L is an infinite language, assume L is context free

There exists a PDA with " n " of production rules & # of variables

Choose $w = a^n b^n c^{n+1}$

$$|w| = 3n+1 \geq n \text{ (as desired)}$$

$$w = a^n b^n c^{n+1} = uvxy z$$

Since $|vxy| \leq n$, leads to many cases

Case 1: vxy is within a 's

For $i=0$; you have less # of a 's than b 's
(Since $|vy| \geq 1$)

Case 2: vxy is within b 's

For $i=0$; you have less # of b 's than a 's
(Since $|vy| \geq 1$)

Case 3: vxy is within c 's

For $i=0$; you have same or less # of c 's than a 's and b 's (Since $|vy| \geq 1$)

Case 4: vxy spans a's and b's

For $i \geq 0$, you have at least 1 less than 2 a's or b's than c's. (Since $|vy| \geq 1$).

Case 5: vxy spans b's and c's

For $i \geq 0$, you have less # of b's or same # of c's than a's (Since $|vy| \geq 1$).

Example I used $i \geq 0$, $uxz \in L$

All the cases lead to a contradiction

Hence, our assumption that L is context free is not true. L is not a CFL.

$$5.) L = \{a^{2n}b^n \mid n \geq 1\} \cup \{b^{n+1}a^n \mid n \geq 1\}$$

(i) Algorithm:

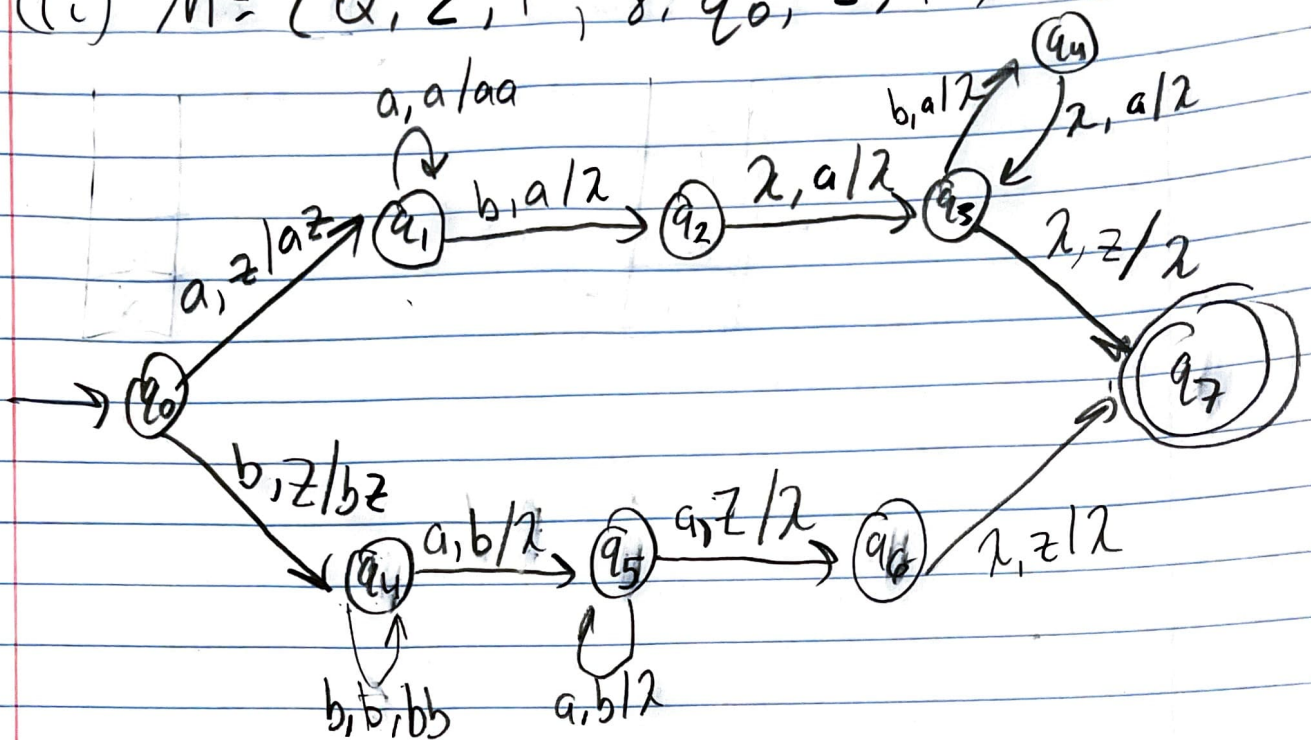
If the first letter is a, then:

When a, if top is ϵ or a, then push a
When b, if top is a, then pop two a's

If the first letter is b, then:

When b, if top is ϵ or b, push a b
When a, if top is b, pop, pop $(n+1)$ b's
When a, if top is ϵ , do nothing

(i) $M = (Q, \Sigma, \Gamma, \delta, q_0, z, F)$



Complete 7-tuple: $(\{q_0, q_1, q_2, q_3, q_4, q_5, q_6, q_7\}, \{a, b\}, \{z, a, b\}, \delta, q_0, z, \{q_7\})$

(iii.) Show $bbba$ leads to rejection

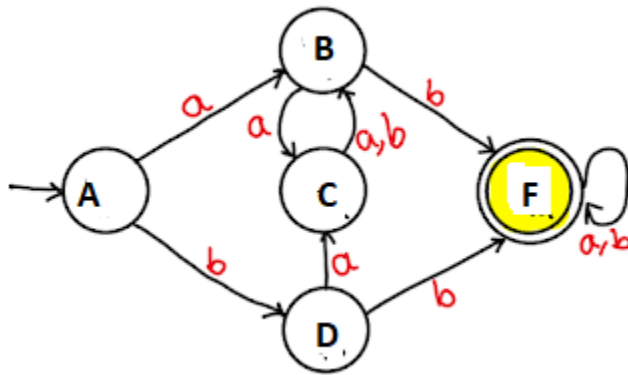
$\delta(q_0, bbb a, z) \vdash \delta(q_1, bbb a, bz) \vdash \delta(q_4, ba, bbz)$
 $\vdash \delta(q_4, a, bbbz) \vdash \delta(q_5, \lambda, bbbz)$

↓
 No transition
 defined leads to
 rejection.

Final

Name _____

1. (15 points) Minimize following DFA:



2. (30 points)

- (i) Give a DFA, M_1 , that accepts a Language L_1 that contains even number of 0's. (Hint: only 2 states)
- (ii) Give a DFA, M_2 , that accepts a Language L_2 that contains even number of 1's.
- (iii) Give acceptor for Reverse of L_1
- (iv) Give acceptor for complement of L_2
- (v) Give acceptor for L_1 union L_2
- (vi) Give acceptor for L_1 intersection L_2
- (vii) Give acceptor for $L_1 - L_2$

3. (15 points) Given the following grammar with production rules numbered from 1 to 8,

- { 1. $S \rightarrow 0A$,
- 2. $S \rightarrow 1A$,
- 3. $S \rightarrow A0$,
- 4. $S \rightarrow AB$,
- 5. $S \rightarrow 0$,
- 6. $S \rightarrow 0A1$,
- 7. $A \rightarrow 1$,
- 8. $S \rightarrow \lambda$ }

- (i) Indicate which of the rules satisfy the conditions for a context-free grammar
(Sample answer: Production rules 1,5,8)

- (ii) Indicate which of the rules satisfy the conditions for a left linear grammar
- (iii) Indicate which of the rules satisfy the conditions for a right linear regular grammar
- (iv) Indicate which of the rules satisfy the conditions for a linear grammar

4. (20 points) Use pumping lemma to show that the language $L = \{a^{m+1}b^m c^m\}$ is not context-free.
5. (20 points) Create a PDA that recognizes the following context free language with terminals $\{a, b\}$: $L = \{a^{n+1}b^n | n \geq 1\} \cup \{b^n a^{n+2} | n \geq 1\}$.
- (i) Describe your algorithm
 - (ii) Give the description as a complete 7-tuple with a transition diagram
 - (iii) Show configuration sequences on **aaabb** leading to acceptance.

