

PROJECT 2

Graph Algorithms and Related Data Structures

Algorithms and Data Structures

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UNIVERSITY OF NORTH CAROLINA AT CHARLOTTE

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PROBLEM 1:

SINGLE-SOURCE SHORTEST PATH ALGORITHM

Assume that G is a weighted graph. The length of a path P is the sum of the weights of the edges of P . If P consists of e_0, e_1, \dots, e_{k-1} , then length of P , denoted as $w(P)$, is defined as $w(P) = \sum_{i=0}^{k-1} w(e_i)$

The distance of a vertex v from a vertex s is the length of a shortest path between s and v , denoted as $d(s, v)$. If $d(s, v) = +\infty$ if no path exists.

DIJKSTRA'S ALGORITHM

Dijkstra's algorithm is a Greedy method of computing the shortest distances of all the vertices from a given source vertex s .

Assumption: Graphs are connected. Edges are undirected or directed, edge weights are nonnegative, i.e., $w(e) \geq 0$

Edge Relaxation:

```
RELAX (u, v, w)
if v.d > u.d + w(u, v)
    v.d = u.d + w(u, v)
    v. $\pi$  = u
```

Steps for finding shortest path using Dijkstra's algorithm:

- For all four inputs, created a graph with an adj list for all edges and their weight depending on whether the graph is directed or undirected.
- Created a shortest-path tree set visited that keeps track of the vertices in the shortest path tree.
- All vertices in the input graph were given a distance value. In the beginning, we set all distance values to infinite and set the distance value for the source vertex to 0 to be chosen first.
- Created a data structure for the parent that keeps track of the path taken and displays it in the output.
- Until all vertices have been visited, a vertex with the shortest distance has been chosen and added to visited.
- The distance value of all the adjacent vertices has been updated.

Algorithm:

DIJKSTRA-SHORT (G, w, s)

1 INITIALIZE (G, s)

```

2  S ←  $\phi$ 
3  Q ← V
4  while Q  $\neq \phi$ 
5      do u ← EXTRACT-MIN (Q)
6      S ← S  $\cup \{u\}$ 
7      for each vertex v  $\in$  Adj [u]
8          do RELAX (u, v, w)

```

Complexity Analysis of Dijkstra's algorithm:

When implemented with queues (which is a preferable approach), Dijkstra's algorithm has the following execution times:

Worst case time complexity: $O(E + V \log V)$
 Average case time complexity: $O(E + V \log V)$
 Best case time complexity: $O(E + V \log V)$
 Space complexity: $O(V)$

Where E is the number of edges and V is the number of vertices in the graph.

We have used python's default dictionary to build an adjacency list to represent the graph. Time Complexity of this implementation is $O(V^2)$. As the relaxation of vertex happens for the minimum picked vertex weight. And have not used a priority Queue so the final complexity for my code would be $O(E + V^2)$. V^2 for the Vertex relaxation and it happens for every edge.

Data Structures Used:

Graph, Adjacency List, python List, Python Default dictionary are the data Structures used in this program to find the shortest path.

PROBLEM 2:

MINIMUM SPANNING TREE

A spanning tree is a subset of Graph G, that covers all of the vertices with the minimum possible number of edges. Hence, there are no cycles in a spanning tree, and it cannot be disconnected. We can conclude from this definition that every connected and undirected Graph G contains at least one spanning tree. A disconnected graph does not have any spanning tree, as it cannot be spanned to all its vertices. Spanning trees have $n-1$ edges. Where, n is the number of vertices.

KRUSKAL'S ALGORITHM

Kruskal's algorithm is the greedy method to find the minimum spanning tree for the given graph. Kruskal's algorithm treats each node as a separate tree, connecting them only if they have the lowest cost compared to all other possibilities.

Steps to create MST using Kruskal's algorithm:

- It begins with a forest, with each vertex representing a tree (a single node tree).
- It finds a safe edge to add to the growing forest by finding the edge (u, v) with the least weight among all the edges that connect any two trees in the forest.
- Kruskal's algorithm qualifies as a greedy algorithm because at each step it adds to the forest an edge of least possible weight.
- At the end of the algorithm: We are left with one cloud that encompasses the MST A tree T which is our MST

Algorithm:

KRUSKAL-MST (G, w)

```
1  A =  $\phi$ 
2  for each vertex  $v \in G.V$ 
3      MAKE-SET( $v$ )
4  sort the edges of  $G.E$  into nondecreasing order by weight  $w$ 
5  for each edge  $(u, v) \in G.E$ , in nondecreasing order by weight
6      if FIND-SET( $u$ )  $\neq$  FIND-SET( $v$ )
7          A = A  $\cup$   $\{(u, v)\}$ 
8          UNION ( $u, v$ )
9  return A
```

Complexity Analysis of Kruskal's algorithm:

$O(E \log E)$ or $O(E \log V)$. Sorting of edges takes $O(E \log E)$ time. After sorting, we iterate through all edges and apply find-union algorithm. The find and union operations can take at most $O(\log V)$ time. So overall complexity is $O(E \log E + E \log V)$ time. The value of E can be at most $O(V^2)$, so $O(\log V)$ are $O(\log E)$ same. Therefore, overall time complexity is $O(E \log E)$ or $O(E \log V)$. Here E is number of edges in the graph and V is the number of vertices.

Data Structures used:

Graph, python List, Python default dictionary are the data Structures used in this program to find the MST using Kruskal's algorithm.

PROBLEM 3:

STRONGLY CONNECTED COMPONENTS:

Algorithm:

- Run DFS on D and every time DFS finishes expanding a vertex v , add v to a stack S .
- Let D^R be the graph D with the direction of all the edges reversed.
- While the stack S is non-empty, pop a vertex v from S and perform DFS on D^R starting at v .
- Let U be the set of all the vertices visited by $\text{DFS}(v)$. Then U is the strongly connected component containing v .
- Remove U from S and repeat.

The above algorithm is DFS based. It does DFS two times. DFS of a graph produces a single tree if all vertices are reachable from the DFS starting point. Otherwise DFS produces a forest. So DFS of a graph with only one SCC always produces a tree. The important point to note is DFS may produce a tree or a forest when there are more than one SCCs depending upon the chosen starting point.

Sequence of picking vertices as starting points of DFS:

There is no direct way to get this sequence. However, if we do a DFS of graph and store vertices according to their finish times, we ensure that the finish time of a vertex that connects to other SCCs (other than its own SCC) will always be greater than the finish time of vertices in the other SCC.

Time Complexity:

The above algorithm calls DFS finds reverse of the graph and again calls DFS. DFS takes $O(V+E)$ for a graph represented using an adjacency list. Reversing a graph also takes $O(V+E)$ time. For reversing the graph, we simply traverse all adjacency lists.

Data Structures used:

Graph, Adjacency List, python List, Python Default dictionary are the data Structures used in this program to find the SCCs.

Code:

PROBLEM 1:

Single-source Shortest Path Algorithm

```
: from collections import defaultdict
import sys
import time
Dijkstra_times = []

class Graph:
    def __init__(self, directed = False, vertices = None):
        self.directed = directed
        self.graph = defaultdict(list)
        self.vertices = vertices

    def addEdge(self, frm, to, w):
        self.graph[frm].append([to, w])

        if self.directed is False:
            self.graph[to].append([frm, w])
        elif self.directed is True:
            self.graph[to] = self.graph[to]

    def getEdges(self):
        print(self.graph.keys())
        print(self.graph.values())

    def find_min_dist(self, distance, visit):
        min_dist = float('inf')
        index = -1
        for key in self.graph.keys():
            if visit[key] is False and distance[key] < min_dist:
                min_dist = distance[key]
                index = key
        return index

    def path(self, root, k):
        if root[k] is None:
            return
        self.path(root, root[k])
        print(chr(k+65), end=" ")

    def solution(self, distance, root, source):
        print('{} \t\t {} \t {}'.format('Vertex', 'Path Cost', 'Path'))

        for key in self.graph.keys():
            if key == source:
                continue
            if distance[key] == float("inf"):
                continue
            print('{} -> {} \t\t {} \t\t {}'.format(chr(source+65), chr(key+65), distance[key], chr(source+65)), end=' ')
            tmp = []
            self.path(root, key)
            print()

    def dijkstra(self, source):
        visited = {i: False for i in self.graph}
        distance = {i: float('inf') for i in self.graph}
        root = {i: None for i in self.graph}

        distance[source] = 0

        # find shortest path for all vertices
        for i in range(len(self.graph) - 1):
            start = self.find_min_dist(distance, visited)
            visited[start] = True
            for vertex, w in self.graph[start]:

                if visited[vertex] is False and distance[start] + w < distance[vertex]:
                    distance[vertex] = distance[start] + w
                    root[vertex] = start
        return root, distance
```

```

if __name__ == '__main__':
    directed = False
    n_inputs = 0

    while n_inputs < 4:
        filepath = "graph" + str(n_inputs) + ".txt"
        inputFile = open(filepath, 'r')
        n_line = 0
        source = sys.maxsize
        print('\033[1m+Single source Shortest paths for+"\033[0m'+ " {} is: ".format(filepath))
        print("=====")

        for line in inputFile.readlines():
            tmp = line.split()
            if n_line == 0:
                print("Number of Vertices provided: {}".format(int(tmp[0])))
                print("Number of Edges provided: {}".format(int(tmp[1])))
                d_ud = tmp[2]
                if d_ud == "D":
                    directed = True
                    print("The Graph is: DIRECTED")
                else:
                    directed = False
                    print("The Graph is: UN DIRECTED")
                G = Graph(directed)
            elif len(tmp) == 1:
                source = ord(tmp[0])-65
            else:
                G.addEdge(ord(tmp[0])-65, ord(tmp[1])-65, int(tmp[2]))

            n_line += 1

        print("Start or Source provided: {}".format(chr(source+65)))
        st = time.time()
        root, distance = G.dijkstra(source)

```

```

#print(parent, dist)
print("-----")
print("Dijkstra's Algorithm results:")
print("-----")
G.solution(distance, root, source)
n_inputs += 1
print("-----")
print('\033[1m+This input's execution time is :'+ '\033[0m', (time.time() - st)*1000, "\n")
print("-----\n")

```

The Results for the above graphs are: -

Single source Shortest paths for graph0.txt is:

=====

Number of Vertices provided: 9

Number of Edges provided: 18

The Graph is: UN DIRECTED

Start or Source provided: A

Dijkstra's Algorithm results:

Vertex	Path Cost	Path
A -> C	9	A C
A -> D	13	A D
A -> B	22	A B
A -> H	56	A B H
A -> F	54	A C F
A -> E	46	A D E
A -> I	53	A D I
A -> G	69	A D E G

This input's execution time is : 0.9992122650146484

Single source Shortest paths for graph1.txt is:
=====

Number of Vertices provided: 10
Number of Edges provided: 14
The Graph is: UN DIRECTED
Start or Source provided: A

Dijkstra's Algorithm results:

Vertex	Path Cost	Path
A -> B	3	A B
A -> C	1	A C
A -> D	5	A C D
A -> E	7	A C D E
A -> F	9	A C D E F
A -> G	13	A C D E F G
A -> H	10	A C D E H
A -> I	13	A C D E F J I
A -> J	11	A C D E F J

This input's execution time is : 0.9989738464355469

Single source Shortest paths for graph2.txt is:
=====

Number of Vertices provided: 10
Number of Edges provided: 20
The Graph is: DIRECTED
Start or Source provided: A

Dijkstra's Algorithm results:

Vertex	Path Cost	Path
A -> B	4	A B
A -> I	15	A B I
A -> D	13	A B D
A -> C	20	A B D C
A -> E	15	A B D E
A -> F	19	A B D F
A -> G	22	A B D E G
A -> H	19	A B D E H
A -> J	3	A J

This input's execution time is : 0.9922981262207031

Single source Shortest paths for graph3.txt is:

=====

Number of Vertices provided: 10

Number of Edges provided: 20

The Graph is: DIRECTED

Start or Source provided: A

Dijkstra's Algorithm results:

Vertex	Path Cost	Path
A -> B	4	A B
A -> C	15	A B C
A -> D	13	A B D
A -> E	15	A B D E
A -> F	19	A B D F
A -> G	22	A B D E G
A -> H	19	A B D E H
A -> I	29	A B D E H I
A -> J	3	A J

This input's execution time is : 1.9979476928710938

Graph Directed Graph's and Un-Directed Graph's Runtimes:

Graph0: 0.9992122650146484

Graph1: 0.9989738464355469

Graph2: 0.9922981262207031

Graph3: 1.9979476928710938

Runtimes in the above table are in Micro seconds

PROBLEM 2:

Minimum Spanning Tree Algorithm

```
class Graph:

    def __init__(self, vertices):
        self.vertices = vertices
        self.graph = []

    def addEdge(self, root, vertex, cost):
        self.graph.append([root, vertex, cost])

    def get_rootnode(self, root_node, v):
        if root_node[v] == v:
            return v
        return self.get_rootnode(root_node, root_node[v])

    def union_of_sets(self, root_node, rank, x, y):
        root_x = self.get_rootnode(root_node, x)
        root_y = self.get_rootnode(root_node, y)

        if rank[root_x] < rank[root_y]:
            root_node[root_x] = root_y
        elif rank[root_x] > rank[root_y]:
            root_node[root_y] = root_x
        else:
            root_node[root_y] = root_x
            rank[root_x] += 1

    def Kruskal(self):
        result = []
        idx_e = 0
        idx_r = 0
        self.graph = sorted(self.graph, key=lambda item: item[2])
        parent = []
        rank = []
        for node in range(self.vertices):
            parent.append(node)
            rank.append(0)
        while idx_e < self.vertices - 1:
            root, vertex, cost = self.graph[idx_r]
            idx_r += 1
            s1 = self.get_rootnode(parent, root)
            s2 = self.get_rootnode(parent, vertex)
            if s1 != s2:
                idx_e += 1
                result.append([root, vertex, cost])
                self.union_of_sets(parent, rank, s1, s2)

        print('\033[1m' + "Edge Selected \t\t Weight" + '\033[0m')
        print()
        total_cost=0
        for x, y, z in result:
            print(chr(x + 65), "->", chr(y + 65), "\t\t\t ", z)
            total_cost += z
        print("*****")
        print('\033[1m'+"The total Cost for Minimum Spanning Tree is", total_cost,'\033[0m')
```

```

if __name__ == '__main__':
    n_inputs = 0
    print()
    while (n_inputs < 4):
        print("=====")
        filepath = "graph" + str(n_inputs) + ".txt"
        inputFile = open(filepath, "r")
        print('\033[1m'+ "The Minimum Spanning Tree using Kruskal's Algorithm for" + '\033[0m' + " {} :".format(filepath))
        print("=====")

        n_line = 0
        for line in inputFile.readlines():
            tmp = line.split()
            if n_line == 0:
                no_of_vertices = int(tmp[0])
                graph = Graph(no_of_vertices)
            elif len(tmp) == 1:
                pass
            else:
                graph.addEdge(ord(tmp[0]) - 65, ord(tmp[1]) - 65, int(tmp[2]))
            n_line = n_line + 1

        print("-----")
        print("Kruskal's Algorithm Results")
        print("-----")
        st = time.time()
        graph.Kruskal()
        n_inputs += 1
        print("-----")
        print('\033[1m'+ "This input's execution time is : " + '\033[0m', (time.time() - st)*1000)
        print("-----\n")

```

The Results for the above graphs are: -

```

=====
The Minimum Spanning Tree using Kruskal's Algorithm for graph0.txt :
=====
-----
Kruskal's Algorithm Results
-----


| <b>Edge Selected</b> | <b>Weight</b> |
|----------------------|---------------|
| C -> D               | 4             |
| A -> C               | 9             |
| E -> F               | 18            |
| H -> I               | 19            |
| G -> I               | 20            |
| A -> B               | 22            |
| E -> G               | 23            |
| D -> E               | 33            |


-----
The total Cost for Minimum Spanning Tree is 148
-----
This input's execution time is : 1.992940902709961
-----

```

=====

The Minimum Spanning Tree using Kruskal's Algorithm for graph1.txt :

=====

Kruskal's Algorithm Results

Edge Selected	Weight
A -> C	1
F -> H	1
B -> C	2
D -> E	2
E -> F	2
I -> J	2
J -> F	2
B -> D	3
G -> H	3

The total Cost for Minimum Spanning Tree is 18

This input's execution time is : 1.0082721710205078

=====

The Minimum Spanning Tree using Kruskal's Algorithm for graph2.txt :

=====

Kruskal's Algorithm Results

Edge Selected	Weight
F -> C	1
I -> J	1
D -> E	2
H -> F	2
H -> J	2
A -> J	3
A -> B	4
E -> H	4
E -> G	7

The total Cost for Minimum Spanning Tree is 26

This input's execution time is : 0.0

=====

The Minimum Spanning Tree using Kruskal's Algorithm for graph3.txt :

=====

Kruskal's Algorithm Results

Edge Selected	Weight
---------------	--------

F -> C	1
I -> J	1
D -> E	2
H -> F	2
H -> J	2
A -> J	3
A -> B	4
E -> H	4
E -> G	7

The total Cost for Minimum Spanning Tree is 26

This input's execution time is : 0.0

PROBLEM 3:

Finding Strongly Connected Components

```
from collections import defaultdict
import sys
import time

class Graph:

    def __init__(self,vertices):
        self.V = vertices
        self.graph = defaultdict(list)
        self.scc_count = 0

    # adds edge to the graph
    def addEdge(self,root,vertex):
        self.graph[root].append(vertex)

    # helper function for DFS
    def helper_DFS(self,vertex,visited, st):
        visited[vertex]= True
        #print(vertex, end=" ")
        st.append(vertex)

        for i in self.graph[vertex]:
            if visited[i]==False:
                self.helper_DFS(i,visited, st)
                self.scc_count += 1
        self.scc_count += 1
        return st, self.scc_count
```

```

# helper function that marks current node as visited and traverse all the adjacent nodes to current node
def helper(self, vertex, visited, stack):
    visited[vertex] = True

    for node in self.graph[vertex]:
        if visited[node] == False:
            self.helper(node, visited, stack)
    stack = stack.append(vertex)

# Function that returns transposed graph
def get_transpose(self):
    gr = Graph(self.V)

    for i in self.graph:
        for j in self.graph[i]:
            gr.addEdge(j, i)
    return gr

# Function that finds strongly connected components in a directed graph
def strongly_connected_components(self):

    stack = []
    # first DFS all visited as False
    visited = [False] * (self.V)
    # fill vertices for first DFS
    for i in range(self.V):
        if visited[i] == False:
            self.helper(i, visited, stack)

    grph = self.get_transpose()

    # second DFS all visited as False
    visited = [False] * (self.V)

    # Process all vertices in order of Stack
    while stack:
        i = stack.pop()
        if visited[i] == False:
            tmp = []
            scc, sccCount = grph.helper_DFS(i, visited, tmp)
            print("Strongly Connected Components ", sccCount, " = ", scc)
            # print("\n")

if __name__ == '__main__':
    n_inputs = 0
    print()
    while (n_inputs < 4):
        print("=====")
        filepath = "graph" + str(n_inputs) + "_SCC.txt"
        inputFile = open(filepath, "r")
        print('\033[1m'+ "Strongly Connected Components in the given "+ '\033[0m'+ " {} :".format(filepath))
        print("=====")

        n_line = 0
        for line in inputFile.readlines():
            tmp = line.split()
            if n_line == 0:
                no_of_vertices = int(tmp[0])
                graph = Graph(no_of_vertices)
            elif len(tmp) == 1:
                pass
            else:
                # graph.add_edge(tmp[0], tmp[1])
                graph.addEdge(ord(tmp[0]) - 65, ord(tmp[1]) - 65)
            n_line = n_line + 1

```

```

print("*****")
print("SCC Algorithm Results")
print("*****")
st = time.time()
graph.strongly_connected_components()
n_inputs += 1
print("-----")
print('\033[1m'+ "This input's execution time is : " + '\033[0m', (time.time() - st)*1000)
print("-----\n")

```

The Results for the above graphs are: -

=====

Strongly Connected Components in the given graph0_SCC.txt :

=====

SCC Algorithm Results

Strongly Connected Components 1 = [0, 2, 1, 4, 3, 5, 7, 6]
 Strongly Connected Components 2 = [9]
 Strongly Connected Components 3 = [8]

This input's execution time is : 0.0

=====

Strongly Connected Components in the given graph1_SCC.txt :

=====

SCC Algorithm Results

Strongly Connected Components 1 = [0, 2, 3, 1, 4, 5, 7]
 Strongly Connected Components 2 = [6]
 Strongly Connected Components 3 = [8]
 Strongly Connected Components 4 = [9]

This input's execution time is : 0.0

=====

Strongly Connected Components in the given graph2_SCC.txt :

=====

SCC Algorithm Results

Strongly Connected Components 1 = [6]
 Strongly Connected Components 2 = [0, 2, 1, 4, 3, 5, 7]
 Strongly Connected Components 3 = [8]
 Strongly Connected Components 4 = [9]

This input's execution time is : 0.0

=====

Strongly Connected Components in the given graph3_SCC.txt :

=====

```
-----  
SCC Algorithm Results  
-----  
Strongly Connected Components 1 = [0, 2, 1, 4, 3, 5, 7, 6]  
Strongly Connected Components 2 = [8]  
Strongly Connected Components 3 = [9]  
-----  
This input's execution time is : 0.0  
-----
```