

Lab Report – Week 11

CS2023 Data Structures and Algorithms

Dept. of Computer Science and Engineering, University of Moratuwa

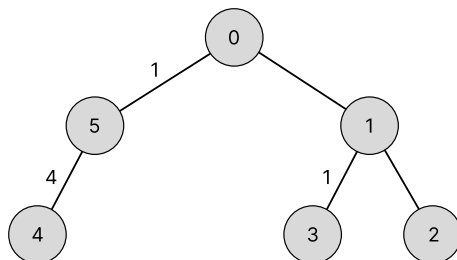
Name: Tharindu Perera

Index Number: 210472J

1.

	0	1	2	3	4	5
0	0	3	0	0	0	1
1	3	0	2	1	10	0
2	0	2	0	3	0	5
3	0	1	3	0	5	0
4	0	10	0	5	0	4
5	1	0	5	0	4	0

2.



3.

```
TERMINAL  PROBLEMS  2

PS C:\Users\thari\UoM-DSA-S2-Labs> cd .\Lab11\
PS C:\Users\thari\UoM-DSA-S2-Labs\Lab11> g++ -o bin\app.exe .\mst.cpp
PS C:\Users\thari\UoM-DSA-S2-Labs\Lab11> ./bin/app.exe
Edge  Weight
0 - 1  3
1 - 2  2
1 - 3  1
5 - 4  4
0 - 5  1
PS C:\Users\thari\UoM-DSA-S2-Labs\Lab11> 
```

4. Yes, MST in Question 2 and 3 are the same.

Edge weights have to be distinct.

5. The time complexity of Prim's algorithm, when using a priority queue to extract the minimal weighted edge, is either $O(E \log V)$ or $O((V + E) \log V)$. Here, V stands for the quantity of vertices, while E stands for the quantity of edges.

The sorting step in Kruskal's algorithm normally takes $O(E \log E)$ time, and the unionfind operations can be thought of as having an $O(E \log V)$ time complexity. So, either $O(E \log E)$ or $O(E \log V)$ can be used to express the overall time complexity.

Prim's algorithm typically performs better than Kruskal's approach for dense graphs, whereas Kruskal's technique performs better for sparse graphs.

GitHub Link : [Tharindu6516/UoM-DSA-S2-Labs \(github.com\)](https://github.com/Tharindu6516/UoM-DSA-S2-Labs)