Optimizing Network Traffic through Dynamic Load Balancing

Problem Statement:

In networking, the efficient distribution of network traffic is crucial for maintaining optimal performance, minimizing latency, and maximizing resource utilization. Unevenly distributed traffic patterns between a source and destination can lead to suboptimal network conditions, resulting in potential bottlenecks and degraded user experience.

The challenge addressed by this project is to design and implement a dynamic load-balancing solution that intelligently selects the most efficient path between a source and destination based on real-time traffic conditions. The primary issues to be tackled include:

- 1. **Traffic Imbalance:** Unequal distribution of network traffic across available paths can lead to congestion on certain routes, causing delays and reducing overall network efficiency.
- 2. **Resource Underutilization:** Inefficient load balancing can result in underutilization of available network resources, limiting the scalability and performance of the overall system.
- 3. **Dynamic Traffic Variations:** Networks experience dynamic changes in traffic patterns due to varying user demands, application requirements, and network topology alterations. A static load-balancing approach may fail to adapt to these fluctuations.

This project aims to address these challenges by implementing a load-balancing algorithm that dynamically evaluates and selects the shortest path with the minimum traffic between a source and destination. Through this approach, the load balancer seeks to optimize network performance, reduce latency, and enhance overall resource utilization in response to changing network conditions.

Approach:

1. Specify the source node and the destination node for which load balancing is to be performed.

```
Enter Host 1
2
Enter Host 2
3
```

- 2. Examine all the paths between the source node and the destination node.
- 3. Utilize Dijkstra's algorithm to determine the shortest paths between the source and destination. This ensures the exclusion of longer paths, limiting the computation to a select few.

```
All Paths
[2, 10, 1]
[2, 21, 1]
```

4. Retrieve real-time traffic information for the identified shortest paths.

```
Cost Computation...

TxRate...

1172

cost

0

Cost Computation...

TxRate...

1170

cost

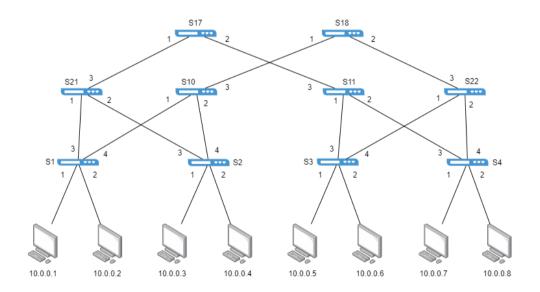
2

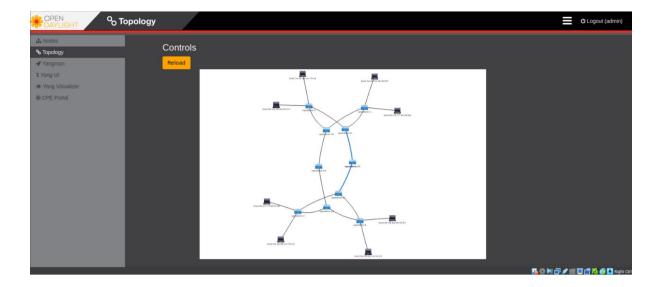
Final Link Cost
{'2::10::1': 4, '2::21::1': 2}
```

5. Select the path with the least traffic as the best path and insert flow rules into switches along the chosen path.

```
Shortest Path: 2::21::1
```

Testing Topology:





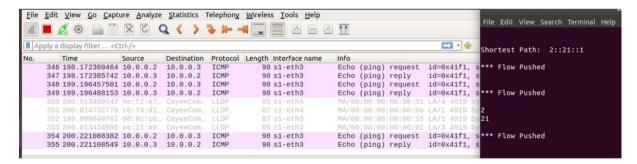
Discussion

When the shortest path(Minimum traffic path) changes, we can see a change of the interface name(port number) of the switch.

Ex:

Ping 10.0.0.2 to 10.0.0.3

When the shortest path (minimum traffic path) is 2::21::1 (s2-s21-s1) we can see the s1 interface is eth3.



Then the shortest path (minimum traffic path) is changed to 2::10::1 (s2-s10-s1). We can see the s1 interface is eth4.

