

Types of Morphisms in Bicategories

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In this chapter, we study special kinds of morphisms in bicategories:

1. *Monomorphisms and Epimorphisms in Bicategories* ([Sections 14.1](#) and [14.2](#)). There is a large number of different notions capturing the idea of a “monomorphism” or of an “epimorphism” in a bicategory.

Arguably, the notion that best captures these concepts is that of a *pseudomononic morphism* ([Definition 14.1.10.1.1](#)) and of a *pseudoepic morphism* ([Definition 14.2.10.1.1](#)), although the other notions introduced in [Sections 14.1](#) and [14.2](#) are also interesting on their own.

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14.1 Monomorphisms in Bicategories

14.1.1 Representably Faithful Morphisms

Let C be a bicategory.

DEFINITION 14.1.1.1.1 ► REPRESENTABLY FAITHFUL MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **representably faithful**¹ if, for each $X \in \text{Obj}(C)$, the functor

$$f_*: \text{Hom}_C(X, A) \rightarrow \text{Hom}_C(X, B)$$

given by postcomposition by f is faithful.

¹*Further Terminology:* Also called simply a **faithful morphism**, based on [Item 1](#) of [Example 14.1.1.1.3](#).

REMARK 14.1.1.1.2 ► UNWINDING DEFINITION 14.1.1.1

In detail, f is representably faithful if, for all diagrams in C of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \downarrow \downarrow \downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if we have

$$\text{id}_f \star \alpha = \text{id}_f \star \beta,$$

then $\alpha = \beta$.

EXAMPLE 14.1.1.3 ► EXAMPLES OF REPRESENTABLY FAITHFUL MORPHISMS

Here are some examples of representably faithful morphisms.

1. *Representably Faithful Morphisms in \mathbf{Cats}_2 .* The representably faithful morphisms in \mathbf{Cats}_2 are precisely the faithful functors; see [Categories, Item 2](#) of [Proposition 11.6.1.1.2](#).
2. *Representably Faithful Morphisms in \mathbf{Rel} .* Every morphism of \mathbf{Rel} is representably faithful; see [Relations, Item 1](#) of [Proposition 8.5.11.1.1](#).

14.1.2 Representably Full Morphisms

Let C be a bicategory.

DEFINITION 14.1.2.1.1 ► REPRESENTABLY FULL MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **representably full**¹ if, for each $X \in \text{Obj}(C)$, the functor

$$f_*: \text{Hom}_C(X, A) \rightarrow \text{Hom}_C(X, B)$$

given by postcomposition by f is full.

¹*Further Terminology:* Also called simply a **full morphism**, based on [Item 1](#) of [Example 14.1.2.1.3](#).

REMARK 14.1.2.1.2 ► UNWINDING DEFINITION 14.1.2.1.1

In detail, f is representably full if, for each $X \in \text{Obj}(C)$ and each 2-morphism

$$\beta: f \circ \phi \Rightarrow f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of C , there exists a 2-morphism

$$\alpha: \phi \Rightarrow \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of C such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B = X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \text{id}_f \star \alpha.$$

EXAMPLE 14.1.2.1.3 ► EXAMPLES OF REPRESENTABLY FULL MORPHISMS

Here are some examples of representably full morphisms.

1. *Representably Full Morphisms in \mathbf{Cats}_2 .* The representably full morphisms in \mathbf{Cats}_2 are precisely the full functors; see [Categories](#), ?? of [Proposition 11.6.2.1.2](#).
2. *Representably Full Morphisms in \mathbf{Rel} .* The representably full morphisms in \mathbf{Rel} are characterised in [Relations](#), [Item 2](#) of [Proposition 8.5.11.1.1](#).

14.1.3 Representably Fully Faithful Morphisms

Let C be a bicategory.

DEFINITION 14.1.3.1.1 ► REPRESENTABLY FULLY FAITHFUL MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **representably fully faithful**¹ if the following equivalent conditions are satisfied:

1. The 1-morphism f is representably faithful ([Definition 14.1.1.1.1](#)) and representably full ([Definition 14.1.2.1.1](#)).
2. For each $X \in \text{Obj}(C)$, the functor

$$f_*: \text{Hom}_C(X, A) \rightarrow \text{Hom}_C(X, B)$$

given by postcomposition by f is fully faithful.

¹*Further Terminology:* Also called simply a **fully faithful morphism**, based on [Item 1](#) of [Example 14.1.3.1.3](#).

REMARK 14.1.3.1.2 ► UNWINDING REPRESENTABLY FULLY FAITHFUL MORPHISMS

In detail, f is representably fully faithful if the conditions in [Remark 14.1.1.1.2](#) and [Remark 14.1.2.1.2](#) hold:

1. For all diagrams in C of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if we have

$$\mathrm{id}_f \star \alpha = \mathrm{id}_f \star \beta,$$

then $\alpha = \beta$.

2. For each $X \in \mathrm{Obj}(C)$ and each 2-morphism

$$\beta: f \circ \phi \Longrightarrow f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of C , there exists a 2-morphism

$$\alpha: \phi \Longrightarrow \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of C such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B = X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \mathrm{id}_f \star \alpha.$$

EXAMPLE 14.1.3.1.3 ► EXAMPLES OF REPRESENTABLY FULLY FAITHFUL MORPHISMS

Here are some examples of representably fully faithful morphisms.

1. *Representably Fully Faithful Morphisms in \mathbf{Cats}_2 .* The representably fully faithful morphisms in \mathbf{Cats}_2 are precisely the fully faithful functors; see [Categories](#), [Item 6](#) of [Proposition 11.6.3.1.2](#).
2. *Representably Fully Faithful Morphisms in \mathbf{Rel} .* The representably

fully faithful morphisms of **Rel** coincide (Relations, Item 3 of Proposition 8.5.11.1.1) with the representably full morphisms in **Rel**, which are characterised in Relations, Item 2 of Proposition 8.5.11.1.1.

14.1.4 Morphisms Representably Faithful on Cores

Let C be a bicategory.

DEFINITION 14.1.4.1.1 ► MORPHISMS REPRESENTABLY FAITHFUL ON CORES

A 1-morphism $f: A \rightarrow B$ of C is **representably faithful on cores** if, for each $X \in \text{Obj}(C)$, the functor

$$f_*: \text{Core}(\text{Hom}_C(X, A)) \rightarrow \text{Core}(\text{Hom}_C(X, B))$$

given by postcomposition by f is faithful.

REMARK 14.1.4.1.2 ► UNWINDING DEFINITION 14.1.4.1.1

In detail, f is representably faithful on cores if, for all diagrams in C of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \downarrow \downarrow \downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if α and β are 2-isomorphisms and we have

$$\text{id}_f \star \alpha = \text{id}_f \star \beta,$$

then $\alpha = \beta$.

14.1.5 Morphisms Representably Full on Cores

Let C be a bicategory.

DEFINITION 14.1.5.1.1 ► MORPHISMS REPRESENTABLY FULL ON CORES

A 1-morphism $f: A \rightarrow B$ of C is **representably full on cores** if, for each $X \in \text{Obj}(C)$, the functor

$$f_*: \text{Core}(\text{Hom}_C(X, A)) \rightarrow \text{Core}(\text{Hom}_C(X, B))$$

given by postcomposition by f is full.

REMARK 14.1.5.1.2 ► UNWINDING DEFINITION 14.1.5.1.1

In detail, f is representably full on cores if, for each $X \in \text{Obj}(C)$ and each 2-isomorphism

$$\beta: f \circ \phi \xRightarrow{\sim} f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of C , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of C such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B = X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \text{id}_f \star \alpha.$$

14.1.6 Morphisms Representably Fully Faithful on Cores

Let C be a bicategory.

DEFINITION 14.1.6.1.1 ► MORPHISMS REPRESENTABLY FULLY FAITHFUL ON CORES

A 1-morphism $f: A \rightarrow B$ of C is **representably fully faithful on cores** if the following equivalent conditions are satisfied:

1. The 1-morphism f is representably faithful on cores ([Definition 14.1.5.1.1](#)) and representably full on cores ([Definition 14.1.4.1.1](#)).
2. For each $X \in \text{Obj}(C)$, the functor

$$f_*: \text{Core}(\text{Hom}_C(X, A)) \rightarrow \text{Core}(\text{Hom}_C(X, B))$$

given by postcomposition by f is fully faithful.

REMARK 14.1.6.1.2 ► UNWINDING DEFINITION 14.1.6.1.1

In detail, f is representably fully faithful on cores if the conditions in [Remark 14.1.4.1.2](#) and [Remark 14.1.5.1.2](#) hold:

1. For all diagrams in C of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if α and β are 2-isomorphisms and we have

$$\mathrm{id}_f \star \alpha = \mathrm{id}_f \star \beta,$$

then $\alpha = \beta$.

2. For each $X \in \mathrm{Obj}(C)$ and each 2-isomorphism

$$\beta: f \circ \phi \xRightarrow{\sim} f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of C , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of C such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B = X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \mathrm{id}_f \star \alpha.$$

14.1.7 Representably Essentially Injective Morphisms

Let C be a bicategory.

DEFINITION 14.1.7.1.1 ► REPRESENTABLY ESSENTIALLY INJECTIVE MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **representably essentially injective** if, for each $X \in \text{Obj}(C)$, the functor

$$f_*: \text{Hom}_C(X, A) \rightarrow \text{Hom}_C(X, B)$$

given by postcomposition by f is essentially injective.

REMARK 14.1.7.1.2 ► UNWINDING DEFINITION 14.1.7.1.1

In detail, f is representably essentially injective if, for each pair of morphisms $\phi, \psi: X \rightrightarrows A$ of C , the following condition is satisfied:

(★) If $f \circ \phi \cong f \circ \psi$, then $\phi \cong \psi$.

14.1.8 Representably Conservative Morphisms

Let C be a bicategory.

DEFINITION 14.1.8.1.1 ► REPRESENTABLY CONSERVATIVE MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **representably conservative** if, for each $X \in \text{Obj}(C)$, the functor

$$f_*: \text{Hom}_C(X, A) \rightarrow \text{Hom}_C(X, B)$$

given by postcomposition by f is conservative.

REMARK 14.1.8.1.2 ► UNWINDING DEFINITION 14.1.8.1.1

In detail, f is representably conservative if, for each pair of morphisms $\phi, \psi: X \rightrightarrows A$ and each 2-morphism

$$\alpha: \phi \Rightarrow \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of C , if the 2-morphism

$$\mathrm{id}_f \star \alpha: f \circ \phi \Longrightarrow f \circ \psi, \quad \begin{array}{ccc} & f \circ \phi & \\ \nearrow & \parallel & \searrow \\ X & \mathrm{id}_f \star \alpha & B \\ \searrow & \Downarrow & \nearrow \\ & f \circ \psi & \end{array}$$

is a 2-isomorphism, then so is α .

14.1.9 Strict Monomorphisms

Let C be a bicategory.

DEFINITION 14.1.9.1.1 ► STRICT MONOMORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is a **strict monomorphism** if, for each $X \in \mathrm{Obj}(C)$, the functor

$$f_*: \mathrm{Hom}_C(X, A) \rightarrow \mathrm{Hom}_C(X, B)$$

given by postcomposition by f is injective on objects, i.e. its action on objects

$$f_*: \mathrm{Obj}(\mathrm{Hom}_C(X, A)) \rightarrow \mathrm{Obj}(\mathrm{Hom}_C(X, B))$$

is injective.

REMARK 14.1.9.1.2 ► UNWINDING DEFINITION 14.1.9.1.1

In detail, f is a strict monomorphism in C if, for each diagram in C of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if $f \circ \phi = f \circ \psi$, then $\phi = \psi$.

EXAMPLE 14.1.9.1.3 ► EXAMPLES OF STRICT MONOMORPHISMS

Here are some examples of strict monomorphisms.

1. *Strict Monomorphisms in Cats_2* . The strict monomorphisms in Cats_2 are precisely the functors which are injective on objects and injective on morphisms; see [Categories, Item 1](#) of [Proposition 11.7.2.1.2](#).

2. *Strict Monomorphisms in Rel.* The strict monomorphisms in **Rel** are characterised in **Relations**, **Proposition 8.5.10.1.1**.

14.1.10 Pseudomonic Morphisms

Let C be a bicategory.

DEFINITION 14.1.10.1.1 ► PSEUDOMONIC MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **pseudomonic** if, for each $X \in \text{Obj}(C)$, the functor

$$f_*: \text{Hom}_C(X, A) \rightarrow \text{Hom}_C(X, B)$$

given by postcomposition by f is pseudomonic.

REMARK 14.1.10.1.2 ► UNWINDING DEFINITION 14.1.10.1.1

In detail, a 1-morphism $f: A \rightarrow B$ of C is pseudomonic if it satisfies the following conditions:

1. For all diagrams in C of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if we have

$$\text{id}_f \star \alpha = \text{id}_f \star \beta,$$

then $\alpha = \beta$.

2. For each $X \in \text{Obj}(C)$ and each 2-isomorphism

$$\beta: f \circ \phi \xrightarrow{\sim} f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of C , there exists a 2-isomorphism

$$\alpha: \phi \xrightarrow{\sim} \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of C such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B = X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \text{id}_f \star \alpha.$$

PROPOSITION 14.1.10.1.3 ► PROPERTIES OF PSEUDOMONIC MORPHISMS

Let $f: A \rightarrow B$ be a 1-morphism of C .

1. *Characterisations.* The following conditions are equivalent:

- (a) The morphism f is pseudomononic.
- (b) The morphism f is representably full on cores and representably faithful.
- (c) We have an isocomma square of the form

$$A \stackrel{\text{eq.}}{\cong} A \times_B A, \quad \begin{array}{ccc} A & \xrightarrow{\text{id}_A} & A \\ \text{id}_A \downarrow & \nearrow \text{dashed} & \downarrow F \\ A & \xrightarrow{F} & B \end{array}$$

in C up to equivalence.

2. *Interaction With Cotensors.* If C has cotensors with $\mathbb{1}$, then the following conditions are equivalent:

- (a) The morphism f is pseudomononic.
- (b) We have an isocomma square of the form

$$A \stackrel{\text{eq.}}{\cong} A \times_{\mathbb{1} \pitchfork F} B, \quad \begin{array}{ccc} A & \xrightarrow{\quad} & \mathbb{1} \pitchfork A \\ F \downarrow & \nearrow \text{dashed} & \downarrow \mathbb{1} \pitchfork F \\ B & \xrightarrow{\quad} & \mathbb{1} \pitchfork B \end{array}$$

in C up to equivalence.

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Item 1: Characterisations

Omitted.

Item 2: Interaction With Cotensors

Omitted.



14.2 Epimorphisms in Bicategories

14.2.1 Corepresentably Faithful Morphisms

Let C be a bicategory.

DEFINITION 14.2.1.1.1 ► COREPRESENTABLY FAITHFUL MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **corepresentably faithful** if, for each $X \in \text{Obj}(C)$, the functor

$$f^*: \text{Hom}_C(B, X) \rightarrow \text{Hom}_C(A, X)$$

given by precomposition by f is faithful.

REMARK 14.2.1.1.2 ► UNWINDING DEFINITION 14.2.1.1.1

In detail, f is corepresentably faithful if, for all diagrams in C of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \downarrow \downarrow \downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

EXAMPLE 14.2.1.1.3 ► EXAMPLES OF COREPRESENTABLY FAITHFUL MORPHISMS

Here are some examples of corepresentably faithful morphisms.

1. *Corepresentably Faithful Morphisms in Cats_2* . The corepresentably faithful morphisms in Cats_2 are characterised in [Categories, Item 5](#) of [Proposition 11.6.1.1.2](#).

2. *Corepresentably Faithful Morphisms in Rel.* Every morphism of **Rel** is corepresentably faithful; see [Relations](#), [Item 1](#) of [Proposition 8.5.13.1.1](#).

14.2.2 Corepresentably Full Morphisms

Let C be a bicategory.

DEFINITION 14.2.2.1.1 ► COREPRESENTABLY FULL MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **corepresentably full** if, for each $X \in \text{Obj}(C)$, the functor

$$f^*: \text{Hom}_C(B, X) \rightarrow \text{Hom}_C(A, X)$$

given by precomposition by f is full.

REMARK 14.2.2.1.2 ► UNWINDING DEFINITION 14.2.2.1.1

In detail, f is corepresentably full if, for each $X \in \text{Obj}(C)$ and each 2-morphism

$$\beta: \phi \circ f \Rightarrow \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of C , there exists a 2-morphism

$$\alpha: \phi \Rightarrow \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of C such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

EXAMPLE 14.2.2.1.3 ► EXAMPLES OF COREPRESENTABLY FULL MORPHISMS

Here are some examples of corepresentably full morphisms.

1. *Corepresentably Full Morphisms in \mathbf{Cats}_2* . The corepresentably full morphisms in \mathbf{Cats}_2 are characterised in [Categories, Item 7 of Proposition 11.6.2.1.2](#).
2. *Corepresentably Full Morphisms in \mathbf{Rel}* . The corepresentably full morphisms in \mathbf{Rel} are characterised in [Relations, Item 2 of Proposition 8.5.13.1.1](#).

14.2.3 Corepresentably Fully Faithful Morphisms

Let C be a bicategory.

DEFINITION 14.2.3.1.1 ► COREPRESENTABLY FULLY FAITHFUL MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **corepresentably fully faithful**¹ if the following equivalent conditions are satisfied:

1. The 1-morphism f is corepresentably full ([Definition 14.2.2.1.1](#)) and corepresentably faithful ([Definition 14.2.1.1.1](#)).
2. For each $X \in \mathbf{Obj}(C)$, the functor

$$f^*: \mathbf{Hom}_C(B, X) \rightarrow \mathbf{Hom}_C(A, X)$$

given by precomposition by f is fully faithful.

¹*Further Terminology:* Corepresentably fully faithful morphisms have also been called **lax epimorphisms** in the literature (e.g. in [\[Adá+01\]](#)), though we will always use the name “corepresentably fully faithful morphism” instead in this work.

REMARK 14.2.3.1.2 ► UNWINDING DEFINITION 14.2.3.1.1

In detail, f is corepresentably fully faithful if the conditions in [Remark 14.2.1.1.2](#) and [Remark 14.2.2.1.2](#) hold:

1. For all diagrams in C of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \downarrow \parallel \beta \\ \xrightarrow{\psi} \end{array} X,$$

if we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

2. For each $X \in \text{Obj}(C)$ and each 2-morphism

$$\beta: \phi \circ f \Rightarrow \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of C , there exists a 2-morphism

$$\alpha: \phi \Rightarrow \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of C such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

EXAMPLE 14.2.3.1.3 ► EXAMPLES OF COREPRESENTABLY FULLY FAITHFUL MORPHISMS

Here are some examples of corepresentably fully faithful morphisms.

1. *Corepresentably Fully Faithful Morphisms in \mathbf{Cats}_2* . The fully faithful epimorphisms in \mathbf{Cats}_2 are characterised in [Categories, Item 10 of Proposition 11.6.3.1.2](#).
2. *Corepresentably Fully Faithful Morphisms in \mathbf{Rel}* . The corepresentably fully faithful morphisms of \mathbf{Rel} coincide ([Relations, Item 3 of Proposition 8.5.13.1.1](#)) with the corepresentably full morphisms in \mathbf{Rel} , which are characterised in [Relations, Item 2 of Proposition 8.5.13.1.1](#).

14.2.4 Morphisms Corepresentably Faithful on Cores

Let C be a bicategory.

DEFINITION 14.2.4.1.1 ► MORPHISMS COREPRESENTABLY FAITHFUL ON CORES

A 1-morphism $f: A \rightarrow B$ of C is **corepresentably faithful on cores** if, for each $X \in \text{Obj}(C)$, the functor

$$f^*: \text{Core}(\text{Hom}_C(B, X)) \rightarrow \text{Core}(\text{Hom}_C(A, X))$$

given by precomposition by f is faithful.

REMARK 14.2.4.1.2 ► UNWINDING DEFINITION 14.2.4.1.1

In detail, f is corepresentably faithful on cores if, for all diagrams in C of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \downarrow \downarrow \downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if α and β are 2-isomorphisms and we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

14.2.5 Morphisms Corepresentably Full on Cores

Let C be a bicategory.

DEFINITION 14.2.5.1.1 ► MORPHISMS COREPRESENTABLY FULL ON CORES

A 1-morphism $f: A \rightarrow B$ of C is **corepresentably full on cores** if, for each $X \in \text{Obj}(C)$, the functor

$$f^*: \text{Core}(\text{Hom}_C(B, X)) \rightarrow \text{Core}(\text{Hom}_C(A, X))$$

given by precomposition by f is full.

REMARK 14.2.5.1.2 ► UNWINDING DEFINITION 14.2.5.1.1

In detail, f is corepresentably full on cores if, for each $X \in \text{Obj}(C)$ and each 2-isomorphism

$$\beta: \phi \circ f \xRightarrow{\sim} \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of C , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of C such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

14.2.6 Morphisms Corepresentably Fully Faithful on Cores

Let C be a bicategory.

DEFINITION 14.2.6.1.1 ► MORPHISMS COREPRESENTABLY FULLY FAITHFUL ON CORES

A 1-morphism $f: A \rightarrow B$ of C is **corepresentably fully faithful on cores** if the following equivalent conditions are satisfied:

1. The 1-morphism f is corepresentably full on cores (Definition 14.2.5.1.1) and corepresentably faithful on cores (Definition 14.2.1.1.1).
2. For each $X \in \text{Obj}(C)$, the functor

$$f^*: \text{Core}(\text{Hom}_C(B, X)) \rightarrow \text{Core}(\text{Hom}_C(A, X))$$

given by precomposition by f is fully faithful.

REMARK 14.2.6.1.2 ► UNWINDING DEFINITION 14.2.6.1.1

In detail, f is corepresentably fully faithful on cores if the conditions in Remark 14.2.4.1.2 and Remark 14.2.5.1.2 hold:

1. For all diagrams in C of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if α and β are 2-isomorphisms and we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

2. For each $X \in \text{Obj}(C)$ and each 2-isomorphism

$$\beta: \phi \circ f \xRightarrow{\sim} \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of C , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of C such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

14.2.7 Corepresentably Essentially Injective Morphisms

Let C be a bicategory.

Definition 14.5.11. *Corepresentably Essentially Injective Morphism*

A 1-morphism $f: A \rightarrow B$ of C is **corepresentably essentially injective** if, for each $X \in \text{Obj}(C)$, the functor

$$f^*: \text{Hom}_C(B, X) \rightarrow \text{Hom}_C(A, X)$$

given by precomposition by f is essentially injective.

REMARK 14.2.7.1.2 ► UNWINDING DEFINITION 14.2.7.1.1

In detail, f is corepresentably essentially injective if, for each pair of morphisms $\phi, \psi: B \rightrightarrows X$ of C , the following condition is satisfied:

(\star) If $\phi \circ f \cong \psi \circ f$, then $\phi \cong \psi$.

14.2.8 Corepresentably Conservative Morphisms

Let C be a bicategory.

DEFINITION 14.2.8.1.1 ► COREPRESENTABLY CONSERVATIVE MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **corepresentably conservative** if, for each $X \in \text{Obj}(C)$, the functor

$$f^*: \text{Hom}_C(B, X) \rightarrow \text{Hom}_C(A, X)$$

given by precomposition by f is conservative.

REMARK 14.2.8.1.2 ► UNWINDING DEFINITION 14.2.8.1.1

In detail, f is corepresentably conservative if, for each pair of morphisms $\phi, \psi: B \rightrightarrows X$ and each 2-morphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of C , if the 2-morphism

$$\alpha \star \text{id}_f: \phi \circ f \xRightarrow{\sim} \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \alpha \star \text{id}_f \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

is a 2-isomorphism, then so is α .

14.2.9 Strict Epimorphisms

Let C be a bicategory.

DEFINITION 14.2.9.1.1 ► STRICT EPIMORPHISMS

A 1-morphism $f: A \rightarrow B$ is a **strict epimorphism in C** if, for each $X \in \text{Obj}(C)$, the functor

$$f^*: \text{Hom}_C(B, X) \rightarrow \text{Hom}_C(A, X)$$

given by precomposition by f is injective on objects, i.e. its action on objects

$$f_*: \text{Obj}(\text{Hom}_C(B, X)) \rightarrow \text{Obj}(\text{Hom}_C(A, X))$$

is injective.

REMARK 14.2.9.1.2 ► UNWINDING DEFINITION 14.2.9.1.1

In detail, f is a strict epimorphism if, for each diagram in C of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \xrightarrow{\psi} \end{array} X,$$

if $\phi \circ f = \psi \circ f$, then $\phi = \psi$.

EXAMPLE 14.2.9.1.3 ► EXAMPLES OF STRICT EPIMORPHISMS

Here are some examples of strict epimorphisms.

1. *Strict Epimorphisms in Cats_2* . The strict epimorphisms in Cats_2 are characterised in [Categories, Item 1](#) of [Proposition 11.7.3.1.2](#).
2. *Strict Epimorphisms in \mathbf{Rel}* . The strict epimorphisms in \mathbf{Rel} are characterised in [Relations, Proposition 8.5.12.1.1](#).

14.2.10 Pseudoepic Morphisms

Let C be a bicategory.

DEFINITION 14.2.10.1.1 ► PSEUDOEPIC MORPHISMS

A 1-morphism $f: A \rightarrow B$ of C is **pseudoepic** if, for each $X \in \text{Obj}(C)$, the functor

$$f^*: \text{Hom}_C(B, X) \rightarrow \text{Hom}_C(A, X)$$

given by precomposition by f is pseudomonic.

REMARK 14.2.10.1.2 ► UNWINDING DEFINITION 14.2.10.1.1

In detail, a 1-morphism $f: A \rightarrow B$ of C is pseudoepic if it satisfies the following conditions:

1. For all diagrams in C of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

2. For each $X \in \text{Obj}(C)$ and each 2-isomorphism

$$\beta: \phi \circ f \xRightarrow{\sim} \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of C , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of C such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in C , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

Let $f: A \rightarrow B$ be a 1-morphism of C .

1. *Characterisations.* The following conditions are equivalent:

- (a) The morphism f is pseudoepic.
- (b) The morphism f is corepresentably full on cores and corepresentably faithful.
- (c) We have an isocomma square of the form

$$B \stackrel{\text{eq.}}{\cong} B \coprod_A B, \quad \begin{array}{ccc} B & \xleftarrow{\text{id}_B} & B \\ \text{id}_B \uparrow & \nearrow \text{dashed} & \uparrow F \\ B & \xleftarrow{F} & A \end{array}$$

in C up to equivalence.

PROOF 14.2.10.1.4 ► PROOF OF PROPOSITION 14.2.10.1.3

Item 1: Characterisations

Omitted.



Appendices

A Other Chapters

Preliminaries

- 1. Introduction
- 2. A Guide to the Literature

Sets

- 3. Sets
- 4. Constructions With Sets
- 5. Monoidal Structures on the Category of Sets

6. Pointed Sets

7. Tensor Products of Pointed Sets

Relations

- 8. Relations
- 9. Constructions With Relations
- 10. Conditions on Relations

Categories

- 11. Categories

12. Presheaves and the Yoneda Lemma
14. Types of Morphisms in Bicat-
egories

Monoidal Categories

13. Constructions With Monoidal Cat-
egories

Extra Part

Bicategories

15. Notes

References

- [Adá+01] Jiří Adámek, Robert El Bashir, Manuela Sobral, and Jiří Velebil. “On Functors Which Are Lax Epimorphisms”. In: *Theory Appl. Categ.* 8 (2001), pp. 509–521. issn: 1201-561X (cit. on p. 16).