Constructions With Sets

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This chapter develops some material relating to constructions with sets with an eye towards its categorical and higher-categorical counterparts to be introduced later in this work. Of particular interest are perhaps the following:

- 1. Explicit descriptions of the major types of co/limits in Sets, including in particular explicit descriptions of pushouts and coequalisers (see Definitions 4.2.4.1.1, 4.2.4.1.3, 4.2.5.1.1 and 4.2.5.1.3).
- 2. A discussion of powersets as decategorifications of categories of presheaves, including in particular results such as:
 - (a) A discussion of the internal Hom of a powerset (Section 4.4.7).
 - (b) A o-categorical version of the Yoneda lemma (Presheaves and the Yoneda Lemma, Definition 12.1.5.1.1), which we term the Yoneda lemma for sets (Definition 4.5.5.1.1).
 - (c) A characterisation of powersets as free cocompletions (Section 4.4.5), mimicking the corresponding statement for categories of presheaves (??).
 - (d) A characterisation of powersets as free completions (Section 4.4.6), mimicking the corresponding statement for categories of copresheaves (??).
 - (e) A (-1)-categorical version of un/straightening (Item 2 of Definition 4.5.1.1.4 and Definition 4.5.1.1.5).
 - (f) A o-categorical form of Isbell duality internal to powersets (Section 4.4.8).

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3. A lengthy discussion of the adjoint triple

$$f_! \dashv f^{-1} \dashv f_* \colon \mathcal{P}(A) \xrightarrow{\rightleftarrows} \mathcal{P}(B)$$

of functors (i.e. morphisms of posets) between $\mathcal{P}(A)$ and $\mathcal{P}(B)$ induced by a map of sets $f: A \to B$, including in particular:

- (a) How f^{-1} can be described as a precomposition while $f_!$ and f_* can be described as Kan extensions (Definitions 4.6.1.1.4, 4.6.2.1.2 and 4.6.3.1.4).
- (b) An extensive list of the properties of $f_!$, f^{-1} , and f_* (Definitions 4.6.1.1.5, 4.6.1.1.6, 4.6.2.1.3, 4.6.2.1.4, 4.6.3.1.7 and 4.6.3.1.8).
- (c) How the functors $f_!$, f^{-1} , f_* , along with the functors

$$-_{1} \cap -_{2} \colon \mathcal{P}(X) \times \mathcal{P}(X) \to \mathcal{P}(X),$$
$$[-_{1}, -_{2}]_{X} \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

may be viewed as a six-functor formalism with the empty set \emptyset as the dualising object (Section 4.6.4).

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4.1 Limits of Sets

4.1.1 The Terminal Set

Definition 4.1.1.1. The **terminal set** is the terminal object of Sets as in Limits and Colimits, ??.

Construction 4.1.1.1.2. Concretely, the terminal set is the pair $(pt, \{!_A\}_{A \in Obj(Sets)})$ consisting of:

- 1. *The Limit*. The punctual set pt $\stackrel{\text{def}}{=} \{ \star \}$.
- 2. The Cone. The collection of maps

$$\{!_A : A \to \mathsf{pt}\}_{A \in \mathsf{Obj}(\mathsf{Sets})}$$

defined by

$$!_A(a) \stackrel{\text{def}}{=} \star$$

for each $a \in A$ and each $A \in Obj(Sets)$.

Proof. We claim that pt is the terminal object of Sets. Indeed, suppose we have a diagram of the form

$$A$$
 pt

in Sets. Then there exists a unique map $\phi\colon A\to \operatorname{pt}$ making the diagram

$$A - \frac{\phi}{\exists !} \rightarrow pt$$

commute, namely $!_A$.

4.1.2 Products of Families of Sets

Let $\{A_i\}_{i\in I}$ be a family of sets.

Definition 4.1.2.1.1. The **product**¹ of $\{A_i\}_{i\in I}$ is the product of $\{A_i\}_{i\in I}$ in Sets as in Limits and Colimits, ??.

¹Further Terminology: Also called the **Cartesian product of** $\{A_i\}_{i\in I}$.

Construction 4.1.2.1.2. Concretely, the product of $\{A_i\}_{i\in I}$ is the pair $\left(\prod_{i\in I}A_i, \left\{\operatorname{pr}_i\right\}_{i\in I}\right)$ consisting of:

1. *The Limit*. The set $\prod_{i \in I} A_i$ defined by

$$\prod_{i \in I} A_i \stackrel{\text{def}}{=} \left\{ f \in \operatorname{Sets} \left(I, \bigcup_{i \in I} A_i \right) \middle| \begin{array}{l} \text{for each } i \in I, \text{we} \\ \text{have } f(i) \in A_i \end{array} \right\}.$$

2. The Cone. The collection

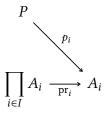
$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

of maps given by

$$\operatorname{pr}_{i}(f) \stackrel{\text{def}}{=} f(i)$$

for each $f \in \prod_{i \in I} A_i$ and each $i \in I$.

Proof. We claim that $\prod_{i \in I} A_i$ is the categorical product of $\{A_i\}_{i \in I}$ in Sets. Indeed, suppose we have, for each $i \in I$, a diagram of the form



in Sets. Then there exists a unique map $\phi\colon P o\prod_{i\in I}A_i$ making the diagram

$$P$$

$$\phi \mid \exists ! \qquad p_i$$

$$\prod_{i \in I} A_i \xrightarrow{\operatorname{pr}_i} A_i$$

commute, being uniquely determined by the condition $\operatorname{pr}_i \circ \phi = p_i$ for each $i \in I$ via

$$\phi(x) = (p_i(x))_{i \in I}$$

for each $x \in P$.

Remark 4.1.2.1.3. Less formally, we may think of Cartesian products and projection maps as follows:

- 1. We think of $\prod_{i \in I} A_i$ as the set whose elements are I-indexed collections $(a_i)_{i \in I}$ with $a_i \in A_i$ for each $i \in I$.
- 2. We view the projection maps

$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

as being given by

$$\operatorname{pr}_i\left(\left(a_j\right)_{j\in I}\right)\stackrel{\text{def}}{=}a_i$$

for each $(a_j)_{i \in I} \in \prod_{i \in I} A_i$ and each $i \in I$.

Proposition 4.1.2.1.4. Let $\{A_i\}_{i\in I}$ be a family of sets.

1. Functoriality. The assignment $\{A_i\}_{i\in I} \mapsto \prod_{i\in I} A_i$ defines a functor

$$\prod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

• *Action on Objects*. For each $(A_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$, we have

$$\left[\prod_{i\in I}\right]((A_i)_{i\in I})\stackrel{\text{def}}{=}\prod_{i\in I}A_i$$

• Action on Morphisms. For each $(A_i)_{i \in I}$, $(B_i)_{i \in I} \in \mathsf{Obj}(\mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}))$, the action on Hom-sets

$$\left(\prod_{i\in I}\right)_{(A_i)_{i\in I},(B_i)_{i\in I}}\colon\operatorname{Nat}((A_i)_{i\in I},(B_i)_{i\in I})\to\operatorname{Sets}\!\left(\prod_{i\in I}A_i,\prod_{i\in I}B_i\right)$$

of $\prod_{i \in I}$ at $((A_i)_{i \in I}, (B_i)_{i \in I})$ is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in Nat $((A_i)_{i \in I}, (B_i)_{i \in I})$ to the map of sets

$$\prod_{i \in I} f_i \colon \prod_{i \in I} A_i \to \prod_{i \in I} B_i$$

defined by

$$\left[\prod_{i\in I} f_i\right] ((a_i)_{i\in I}) \stackrel{\text{def}}{=} (f_i(a_i))_{i\in I}$$

for each $(a_i)_{i \in I} \in \prod_{i \in I} A_i$.

Proof. Item 1, *Functoriality*: This follows from Limits and Colimits, ?? of ??. □

4.1.3 Binary Products of Sets

Let *A* and *B* be sets.

Definition 4.1.3.1.1. The **product of** A **and** B^2 is the product of A and B in Sets as in Limits and Colimits, ??.

Construction 4.1.3.1.2. Concretely, the product of A and B is the pair $(A \times B, \{pr_1, pr_2\})$ consisting of:

1. *The Limit*. The set $A \times B$ defined by

$$A \times B \stackrel{\text{def}}{=} \prod_{z \in \{A,B\}} z$$

$$\stackrel{\text{def}}{=} \{ f \in \text{Sets}(\{0,1\}, A \cup B) \mid \text{we have } f(0) \in A \text{ and } f(1) \in B \}$$

$$\cong \{ \{ \{a\}, \{a,b\} \} \in \mathcal{P}(\mathcal{P}(A \cup B)) \mid \text{we have } a \in A \text{ and } b \in B \}$$

$$\cong \begin{cases} \text{ordered pairs } (a,b) \text{ with } \\ a \in A \text{ and } b \in B \end{cases}.$$

2. The Cone. The maps

$$\operatorname{pr}_1 : A \times B \to A,$$

 $\operatorname{pr}_2 : A \times B \to B$

²Further Terminology: Also called the **Cartesian product of** A **and** B.

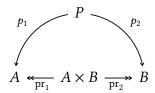
defined by

$$\operatorname{pr}_{1}(a, b) \stackrel{\text{def}}{=} a,$$

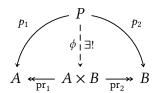
 $\operatorname{pr}_{2}(a, b) \stackrel{\text{def}}{=} b$

for each $(a, b) \in A \times B$.

Proof. We claim that $A \times B$ is the categorical product of A and B in the category of sets. Indeed, suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: P \to A \times B$ making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$

 $\operatorname{pr}_2 \circ \phi = p_2$

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each $x \in P$.

Proposition 4.1.3.1.3. Let *A*, *B*, *C*, and *X* be sets.

1. Functoriality. The assignments $A, B, (A, B) \mapsto A \times B$ define functors

$$A \times -:$$
 Sets \rightarrow Sets,
 $- \times B:$ Sets \rightarrow Sets,
 $-_1 \times -_2:$ Sets \times Sets \rightarrow Sets,

where -1×-2 is the functor where

• *Action on Objects.* For each $(A, B) \in Obj(Sets \times Sets)$, we have

$$[-1 \times -2](A, B) \stackrel{\text{def}}{=} A \times B.$$

• *Action on Morphisms*. For each $(A, B), (X, Y) \in Obj(Sets)$, the action on Hom-sets

$$\times_{(A,B),(X,Y)}$$
: Sets $(A,X) \times$ Sets $(B,Y) \rightarrow$ Sets $(A \times B, X \times Y)$

of \times at ((A, B), (X, Y)) is defined by sending (f, q) to the function

$$f \times q \colon A \times B \to X \times Y$$

defined by

$$[f \times g](a,b) \stackrel{\text{def}}{=} (f(a),g(b))$$

for each $(a, b) \in A \times B$.

and where $A \times -$ and $- \times B$ are the partial functors of $-_1 \times -_2$ at $A, B \in \text{Obj}(\mathsf{Sets})$.

2. Adjointness I. We have adjunctions

$$(A \times - + \operatorname{Sets}(A, -))$$
: Sets $\xrightarrow{A \times -}$ Sets, $\xrightarrow{\operatorname{Sets}(A, -)}$ Sets, $(- \times B + \operatorname{Sets}(B, -))$: Sets $\xrightarrow{\operatorname{Sets}(B, -)}$

witnessed by bijections

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$

 $Sets(A \times B, C) \cong Sets(B, Sets(A, C)),$

natural in $A, B, C \in Obj(Sets)$.

3. Adjointness II. We have an adjunction

$$(\Delta_{\mathsf{Sets}} \dashv -_1 \times -_2)$$
: Sets $\underbrace{\perp}_{-_1 \times -_2}$ Sets \times Sets,

witnessed by a bijection

$$Hom_{Sets \times Sets}((A, A), (B, C)) \cong Sets(A, B \times C),$$

natural in $A \in Obj(Sets)$ and in $(B, C) \in Obj(Sets \times Sets)$.

4. Associativity. We have an isomorphism of sets

$$\alpha_{ABC}^{\mathsf{Sets}} \colon (A \times B) \times C \xrightarrow{\sim} A \times (B \times C),$$

natural in $A, B, C \in Obj(Sets)$.

5. Unitality. We have isomorphisms of sets

$$\lambda_A^{\mathsf{Sets}} \colon \mathsf{pt} \times A \xrightarrow{\sim} A,$$

$$\rho_A^{\mathsf{Sets}} \colon A \times \mathsf{pt} \xrightarrow{\sim} A,$$

natural in $A \in Obj(Sets)$.

6. Commutativity. We have an isomorphism of sets

$$\sigma_{AB}^{\mathsf{Sets}} : A \times B \xrightarrow{\sim} B \times A,$$

natural in $A, B \in Obj(Sets)$.

7. Distributivity Over Coproducts. We have isomorphisms of sets

$$\delta_{\ell}^{\mathsf{Sets}} : A \times (B \ | \ C) \xrightarrow{\sim} (A \times B) \ | \ (A \times C),$$

$$\delta_r^{\mathsf{Sets}} : (A \coprod B) \times C \xrightarrow{\sim} (A \times C) \coprod (B \times C),$$

natural in $A, B, C \in Obj(Sets)$.

8. Annihilation With the Empty Set. We have isomorphisms of sets

$$\zeta_{\ell}^{\mathsf{Sets}} \colon \emptyset \times A \xrightarrow{\sim} \emptyset,$$

$$\zeta_r^{\mathsf{Sets}} \colon A \times \emptyset \xrightarrow{\sim} \emptyset,$$

natural in $A \in Obj(Sets)$.

9. Distributivity Over Unions. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \cup W) = (U \times V) \cup (U \times W),$$

$$(U \cup V) \times W = (U \times W) \cup (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

10. Distributivity Over Intersections. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \cap W) = (U \times V) \cap (U \times W),$$

$$(U \cap V) \times W = (U \times W) \cap (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

11. Distributivity Over Differences. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \setminus W) = (U \times V) \setminus (U \times W),$$

$$(U \setminus V) \times W = (U \times W) \setminus (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

12. Distributivity Over Symmetric Differences. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \triangle W) = (U \times V) \triangle (U \times W),$$

$$(U \triangle V) \times W = (U \times W) \triangle (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

13. Middle-Four Exchange with Respect to Intersections. The diagram

$$(\mathcal{P}(X) \times \mathcal{P}(X)) \times (\mathcal{P}(X) \times \mathcal{P}(X)) \xrightarrow{\cap \times \cap} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow^{\mathcal{P}_{X,X}^{\times} \times \mathcal{P}_{X,X}^{\times}} \qquad \qquad \downarrow^{\mathcal{P}_{X,X}^{\times}}$$

$$\mathcal{P}(X \times X) \times \mathcal{P}(X \times X) \xrightarrow{\cap} \mathcal{P}(X \times X)$$

commutes, i.e. we have

$$(U \times V) \cap (W \times T) = (U \cap V) \times (W \cap T).$$

for each $U, V, W, T \in \mathcal{P}(X)$.

- 14. Symmetric Monoidality. The 8-tuple (Sets, \times , pt, Sets $(-_1, -_2)$, α^{Sets} , λ^{Sets} , ρ^{Sets} , σ^{Sets}) is a closed symmetric monoidal category.
- 15. Symmetric Bimonoidality. The 18-tuple

$$\begin{split} & \Big(\mathsf{Sets}, \coprod, \times, \emptyset, \mathsf{pt}, \mathsf{Sets}(-_1, -_2), \alpha^{\mathsf{Sets}}, \lambda^{\mathsf{Sets}}, \rho^{\mathsf{Sets}}, \sigma^{\mathsf{Sets}}, \\ & \alpha^{\mathsf{Sets}, \coprod}, \lambda^{\mathsf{Sets}, \coprod}, \rho^{\mathsf{Sets}, \coprod}, \sigma^{\mathsf{Sets}, \coprod}, \delta^{\mathsf{Sets}}_{\ell}, \delta^{\mathsf{Sets}}_{r}, \zeta^{\mathsf{Sets}}_{\ell}, \zeta^{\mathsf{Sets}}_{r} \Big), \end{split}$$

is a symmetric closed bimonoidal category, where $\alpha^{\text{Sets},\coprod}$, $\lambda^{\text{Sets},\coprod}$, $\rho^{\text{Sets},\coprod}$, and $\sigma^{\text{Sets},\coprod}$ are the natural transformations from Items 3 to 5 of Definition 4.2.3.1.3.

Proof. Item 1, *Functoriality*: This follows from Limits and Colimits, ?? of ??. *Item* 2, *Adjointness*: We prove only that there's an adjunction $-\times B \dashv \mathsf{Sets}(B, -)$, witnessed by a bijection

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)).$$

natural in $B, C \in Obj(Sets)$, as the proof of the existence of the adjunction $A \times - \dashv Sets(A, -)$ follows almost exactly in the same way.

• Map I. We define a map

$$\Phi_{BC}$$
: Sets $(A \times B, C) \rightarrow \text{Sets}(A, \text{Sets}(B, C))$,

by sending a function

$$\xi \colon A \times B \to C$$

to the function

$$\xi^{\dagger} \colon A \longrightarrow \mathsf{Sets}(B,C),$$

$$a \mapsto \Big(\xi_a^{\dagger} \colon B \to C\Big),$$

where we define

$$\xi_a^{\dagger}(b) \stackrel{\text{def}}{=} \xi(a,b)$$

for each $b \in B$. In terms of the $[a \mapsto f(a)]$ notation of Sets, Definition 3.1.1.1.2, we have

$$\xi^{\dagger} \stackrel{\text{def}}{=} \llbracket a \mapsto \llbracket b \mapsto \xi(a,b) \rrbracket \rrbracket.$$

• Map II. We define a map

$$\Psi_{B,C}$$
: Sets $(A, \text{Sets}(B, C)), \rightarrow \text{Sets}(A \times B, C)$

given by sending a function

$$\xi: A \longrightarrow \mathsf{Sets}(B, C),$$

 $a \mapsto (\xi_a: B \to C),$

to the function

$$\xi^{\dagger}: A \times B \to C$$

defined by

$$\xi^{\dagger}(a,b) \stackrel{\text{def}}{=} \operatorname{ev}_{b}(\operatorname{ev}_{a}(\xi))$$

$$\stackrel{\text{def}}{=} \operatorname{ev}_{b}(\xi_{a})$$

$$\stackrel{\text{def}}{=} \xi_{a}(b)$$

for each $(a, b) \in A \times B$.

• Invertibility I. We claim that

$$\Psi_{AB} \circ \Phi_{AB} = \mathrm{id}_{\mathsf{Sets}(A \times BC)}$$
.

Indeed, given a function $\xi \colon A \times B \to C$, we have

$$\begin{split} \left[\Psi_{A,B} \circ \Phi_{A,B} \right] (\xi) &= \Psi_{A,B} \big(\Phi_{A,B} (\xi) \big) \\ &= \Psi_{A,B} \big(\Phi_{A,B} (\llbracket (a,b) \mapsto \xi(a,b) \rrbracket) \big) \\ &= \Psi_{A,B} (\llbracket a \mapsto \llbracket b \mapsto \xi(a,b) \rrbracket \rrbracket) \big) \\ &= \Psi_{A,B} (\llbracket a' \mapsto \llbracket b' \mapsto \xi(a',b') \rrbracket \rrbracket) \big) \\ &= \llbracket (a,b) \mapsto \operatorname{ev}_b (\operatorname{ev}_a (\llbracket a' \mapsto \llbracket b' \mapsto \xi(a',b') \rrbracket \rrbracket)) \rrbracket \\ &= \llbracket (a,b) \mapsto \operatorname{ev}_b (\llbracket b' \mapsto \xi(a,b') \rrbracket) \rrbracket \big) \\ &= \llbracket (a,b) \mapsto \xi(a,b) \rrbracket \\ &= \xi. \end{split}$$

• Invertibility II. We claim that

$$\Phi_{A,B} \circ \Psi_{A,B} = \mathrm{id}_{\mathsf{Sets}(A,\mathsf{Sets}(B,C))}$$
.

Indeed, given a function

$$\xi : A \longrightarrow \mathsf{Sets}(B, C),$$

 $a \mapsto (\xi_a : B \to C),$

we have

$$\begin{split} \left[\Phi_{A,B} \circ \Psi_{A,B} \right] (\xi) &\stackrel{\text{def}}{=} \Phi_{A,B} \big(\Psi_{A,B} (\xi) \big) \\ &\stackrel{\text{def}}{=} \Phi_{A,B} \big(\left[(a,b) \mapsto \xi_a(b) \right] \big) \\ &\stackrel{\text{def}}{=} \Phi_{A,B} \big(\left[(a',b') \mapsto \xi_{a'}(b') \right] \big) \\ &\stackrel{\text{def}}{=} \left[a \mapsto \left[b \mapsto \text{ev}_{(a,b)} \big(\left[(a',b') \mapsto \xi_{a'}(b') \right] \big) \right] \right] \\ &\stackrel{\text{def}}{=} \left[a \mapsto \left[b \mapsto \xi_a(b) \right] \right] \\ &\stackrel{\text{def}}{=} \left[a \mapsto \xi_a \right] \\ &\stackrel{\text{def}}{=} \xi. \end{split}$$

• *Naturality for* Φ , *Part I.* We need to show that, given a function $g: B \to B'$, the diagram

$$\begin{aligned} \mathsf{Sets}(A \times B', C) & \xrightarrow{\Phi_{B',C}} & \mathsf{Sets}(A, \mathsf{Sets}(B', C)), \\ & \mathsf{id}_A \times g^* \middle\downarrow & & & & & & & & \\ & & \mathsf{id}_A \times g^* \middle\downarrow & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & \\ & & & \\ & & \\ & & & \\ & &$$

commutes. Indeed, given a function

$$\xi \colon A \times B' \to C$$

we have

$$[\Phi_{B,C} \circ (\mathrm{id}_A \times g^*)](\xi) = \Phi_{B,C}([\mathrm{id}_A \times g^*](\xi))$$
$$= \Phi_{B,C}(\xi(-1, g(-2)))$$

$$= [\xi(-_{1}, g(-_{2}))]^{\dagger}$$

$$= \xi_{-_{1}}^{\dagger}(g(-_{2}))$$

$$= (g^{*})_{!}(\xi^{\dagger})$$

$$= (g^{*})_{!}(\Phi_{B',C}(\xi))$$

$$= [(g^{*})_{!} \circ \Phi_{B',C}](\xi).$$

Alternatively, using the $[\![a\mapsto f(a)]\!]$ notation of Sets, Definition 3.1.1.1.2, we have

$$\begin{split} \left[\Phi_{B,C} \circ (\mathrm{id}_A \times g^*) \right] (\xi) &= \Phi_{B,C} ([\mathrm{id}_A \times g^*] (\xi)) \\ &= \Phi_{B,C} ([\mathrm{id}_A \times g^*] ([\![(a,b') \mapsto \xi(a,b')]\!])) \\ &= \Phi_{B,C} ([\![(a,b) \mapsto \xi(a,g(b))]\!]) \\ &= [\![a \mapsto [\![b \mapsto \xi(a,g(b))]\!]] \\ &= [\![a \mapsto g^* ([\![b' \mapsto \xi(a,b')]\!])] \\ &= (g^*)_! ([\![a \mapsto [\![b' \mapsto \xi(a,b')]\!]])) \\ &= (g^*)_! (\Phi_{B',C} ([\![(a,b') \mapsto \xi(a,b')]\!])) \\ &= (g^*)_! (\Phi_{B',C} (\xi)) \\ &= [(g^*)_! \circ \Phi_{B',C}] (\xi). \end{split}$$

• *Naturality for* Φ , *Part II.* We need to show that, given a function $h: C \to C'$, the diagram

$$\begin{split} \mathsf{Sets}(A \times B, C) & \xrightarrow{\Phi_{B,C}} & \mathsf{Sets}(A, \mathsf{Sets}(B, C)), \\ h_! & & \downarrow^{(h_!)_!} \\ \mathsf{Sets}(A \times B, C') & \xrightarrow{\Phi_{B,C'}} & \mathsf{Sets}(A, \mathsf{Sets}(B, C')) \end{split}$$

commutes. Indeed, given a function

$$\xi: A \times B \to C$$

we have

$$[\Phi_{B,C} \circ h_!](\xi) = \Phi_{B,C}(h_!(\xi))$$

$$= \Phi_{B,C}(h_!(\llbracket (a,b) \mapsto \xi(a,b) \rrbracket))$$

$$= \Phi_{B,C}(\llbracket (a,b) \mapsto h(\xi(a,b)) \rrbracket)$$

$$= \llbracket a \mapsto \llbracket b \mapsto h(\xi(a,b)) \rrbracket \rrbracket$$

$$= \llbracket a \mapsto h_!(\llbracket b \mapsto \xi(a,b) \rrbracket \rrbracket)$$

$$= (h_!)_!(\llbracket a \mapsto \llbracket b \mapsto \xi(a,b) \rrbracket \rrbracket)$$

$$= (h_!)_!(\Phi_{B,C}(\llbracket (a,b) \mapsto \xi(a,b) \rrbracket))$$

$$= (h_!)_!(\Phi_{B,C}(\xi))$$

$$= [(h_!)_! \circ \Phi_{B,C}](\xi).$$

• Naturality for Ψ . Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Ψ is also natural in each argument.

This finishes the proof.

Item 3, Adjointness II: This follows from the universal property of the product. *Item 4, Associativity*: This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.1.4.1.1.

Item 5, Unitality: This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.1.5.1.1 and 5.1.6.1.1.

Item 6, Commutativity: This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.1.7.1.1.

Item 7, Distributivity Over Coproducts: This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.3.1.1.1 and 5.3.2.1.1.

Item 8, *Annihilation With the Empty Set*: This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.3.3.1.1 and 5.3.4.1.1.

Item 9, Distributivity Over Unions: See [Pro25c].

Item 10, *Distributivity Over Intersections*: See [Pro25d, Corollary 1].

Item 11, Distributivity Over Differences: See [Pro25a].

Item 12, *Distributivity Over Symmetric Differences*: See [Pro25b].

Item 13, Middle-Four Exchange With Respect to Intersections: See [Pro25d, Corollary 1].

Item 14, Symmetric Monoidality: This is a repetition of Monoidal Structures on the Category of Sets, Definition 5.1.9.1.1, and is proved there.

Item 15, Symmetric Bimonoidality: This is a repetition of Monoidal Structures on the Category of Sets, Definition 5.3.5.1.1, and is proved there. □

Remark 4.1.3.1.4. As shown in Item 1 of Definition 4.1.3.1.3, the Cartesian product of sets defines a functor

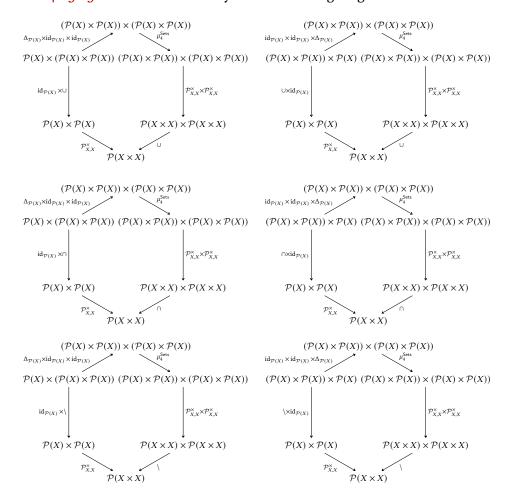
$$-1 \times -2 : \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}$$
.

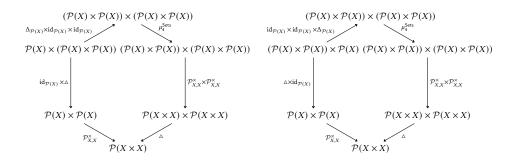
This functor is the $(k, \ell) = (-1, -1)$ case of a family of functors

$$\otimes_{k,\ell} \colon \mathsf{Mon}_{\mathbb{E}_k}(\mathsf{Sets}) \times \mathsf{Mon}_{\mathbb{E}_\ell}(\mathsf{Sets}) \to \mathsf{Mon}_{\mathbb{E}_{k+\ell}}(\mathsf{Sets})$$

of tensor products of \mathbb{E}_k -monoid objects on Sets with \mathbb{E}_ℓ -monoid objects on Sets; see $\ref{eq:tensor}$?.

Remark 4.1.3.1.5. We may state the equalities in Items 9 to 12 of Definition 4.1.3.1.3 as the commutativity of the following diagrams:





4.1.4 Pullbacks

Let A, B, and C be sets and let $f: A \to C$ and $g: B \to C$ be functions.

Definition 4.1.4.1.1. The **pullback of** A **and** B **over** C **along** f **and** g^3 is the pullback of A and B over C along f and g in Sets as in Limits and Colimits, ??.

Construction 4.1.4.1.2. Concretely, the pullback of A and B over C along f and g is the pair $(A \times_C B, \{pr_1, pr_2\})$ consisting of:

1. *The Limit*. The set $A \times_C B$ defined by

$$A \times_C B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid f(a) = g(b)\}.$$

2. The Cone. The maps⁴

$$\operatorname{pr}_1 : A \times_C B \to A,$$

 $\operatorname{pr}_2 : A \times_C B \to B$

defined by

$$\operatorname{pr}_{1}(a,b) \stackrel{\text{def}}{=} a,$$

 $\operatorname{pr}_{2}(a,b) \stackrel{\text{def}}{=} b$

for each $(a, b) \in A \times_C B$.

Proof. We claim that $A \times_C B$ is the categorical pullback of A and B over C with

 $^{^3}$ Further Terminology: Also called the **fibre product of** A **and** B **over** C **along** f **and** g. 4 Further Notation: Also written $\operatorname{pr}_1^{A \times_C B}$ and $\operatorname{pr}_2^{A \times_C B}$.

respect to (f,g) in Sets. First we need to check that the relevant pullback diagram commutes, i.e. that we have

$$f \circ \operatorname{pr}_1 = g \circ \operatorname{pr}_2, \qquad A \times_C B \xrightarrow{\operatorname{pr}_2} B$$

$$\downarrow^{g}$$

$$A \xrightarrow{f} C.$$

Indeed, given $(a, b) \in A \times_C B$, we have

$$[f \circ \operatorname{pr}_1](a,b) = f(\operatorname{pr}_1(a,b))$$

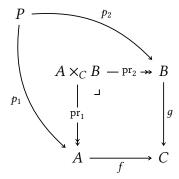
$$= f(a)$$

$$= g(b)$$

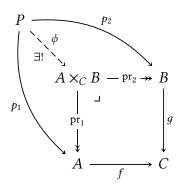
$$= g(\operatorname{pr}_2(a,b))$$

$$= [g \circ \operatorname{pr}_2](a,b),$$

where f(a) = g(b) since $(a,b) \in A \times_C B$. Next, we prove that $A \times_C B$ satisfies the universal property of the pullback. Suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: P \to A \times_C B$ making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$

 $\operatorname{pr}_2 \circ \phi = p_2$

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each $x \in P$, where we note that $(p_1(x), p_2(x)) \in A \times B$ indeed lies in $A \times_C B$ by the condition

$$f \circ p_1 = g \circ p_2$$
,

which gives

$$f(p_1(x)) = g(p_2(x))$$

for each $x \in P$, so that $(p_1(x), p_2(x)) \in A \times_C B$.

Remark 4.1.4.1.3. It is common practice to write $A \times_C B$ for the pullback of A and B over C along f and g, omitting the maps f and g from the notation and instead leaving them implicit, to be understood from the context.

However, the set $A \times_C B$ depends very much on the maps f and g, and sometimes it is necessary or useful to note this dependence explicitly. In such situations, we will write $A \times_{f,C,g} B$ or $A \times_C^{f,g} B$ for $A \times_C B$.

Example 4.1.4.1.4. Here are some examples of pullbacks of sets.

1. *Unions via Intersections*. Let *X* be a set. We have

$$A \cap B \cong A \times_{A \cup B} B,$$
 $A \cap B \xrightarrow{\longrightarrow} B$ $\downarrow \iota_B$ $\downarrow \iota_A \xrightarrow{\iota_A} A \cup B$

for each $A, B \in \mathcal{P}(X)$.

Proof. Item 1, Unions via Intersections: Indeed, we have

$$A \times_{A \cup B} B \cong \{(x, y) \in A \times B \mid x = y\}$$

\(\approx A \cap B.\)

This finishes the proof.

Proposition 4.1.4.1.5. Let *A*, *B*, *C*, and *X* be sets.

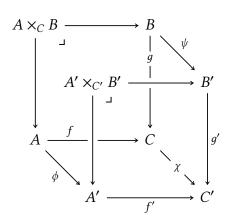
1. Functoriality. The assignment $(A,B,C,f,g)\mapsto A\times_{f,C,g}B$ defines a functor

$$-_1 \times_{-_3} -_1 : \operatorname{\mathsf{Fun}}(\mathcal{P}, \operatorname{\mathsf{Sets}}) \to \operatorname{\mathsf{Sets}},$$

where \mathcal{P} is the category that looks like this:

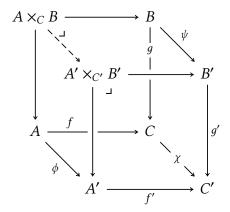


In particular, the action on morphisms of $-1 \times_{-3} -1$ is given by sending a morphism



in Fun(
$$\mathcal{P}$$
, Sets) to the map $\xi \colon A \times_C B \xrightarrow{\exists !} A' \times_{C'} B'$ given by
$$\xi(a,b) \stackrel{\text{def}}{=} (\phi(a), \psi(b))$$

for each $(a, b) \in A \times_C B$, which is the unique map making the diagram



commute.

2. Adjointness I. We have adjunctions

$$(A \times_X - + \mathbf{Sets}_{/X}(A, -)) : \operatorname{Sets}_{/X} \underbrace{\xrightarrow{A \times_X -}}_{\mathbf{Sets}_{/X}(A, -)} \operatorname{Sets}_{/X},$$

$$(- \times_X B + \mathbf{Sets}_{/X}(B, -)) : \operatorname{Sets}_{/X} \underbrace{\xrightarrow{A \times_X -}}_{\mathbf{Sets}_{/X}(B, -)} \operatorname{Sets}_{/X},$$

witnessed by bijections

$$\mathsf{Sets}_{/X}(A \times_X B, C) \cong \mathsf{Sets}_{/X}(A, \mathbf{Sets}_{/X}(B, C)),$$

 $\mathsf{Sets}_{/X}(A \times_X B, C) \cong \mathsf{Sets}_{/X}(B, \mathbf{Sets}_{/X}(A, C)),$

natural in (A, ϕ_A) , (B, ϕ_B) , $(C, \phi_C) \in \text{Obj}(\mathsf{Sets}_{/X})$, where $\mathsf{Sets}_{/X}(A, B)$ is the object of $\mathsf{Sets}_{/X}$ consisting of (see Fibred Sets, ??):

• *The Set.* The set **Sets** $_{/X}(A, B)$ defined by

$$\mathbf{Sets}_{/X}(A,B) \stackrel{\text{def}}{=} \coprod_{x \in X} \mathsf{Sets} \big(\phi_A^{-1}(x), \phi_Y^{-1}(x) \big)$$

• *The Map to X*. The map

$$\phi_{\mathsf{Sets}_{/X}(A,B)} \colon \mathsf{Sets}_{/X}(A,B) \to X$$

defined by

$$\phi_{\mathsf{Sets}_{/X}(A,B)}(x,f) \stackrel{\text{def}}{=} x$$

for each
$$(x, f) \in \mathbf{Sets}_{/X}(A, B)$$
.

3. Adjointness II. We have an adjunction

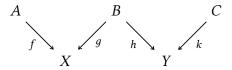
$$\left(\Delta_{\mathsf{Sets}_{/X}}\dashv -_1 \times -_2\right)$$
: $\mathsf{Sets}_{/X} \underbrace{\overset{\Delta_{\mathsf{Sets}_{/X}}}{\bot}}_{-_1 \times -_2} \mathsf{Sets}_{/X} \times \mathsf{Sets}_{/X}$,

witnessed by a bijection

$$\operatorname{Hom}_{\operatorname{\mathsf{Sets}}_{/X} \times \operatorname{\mathsf{Sets}}_{/X}}((A,A),(B,C)) \cong \operatorname{\mathsf{Sets}}_{/X}(A,B \times_X C),$$

natural in $A \in \text{Obj}(\mathsf{Sets}_{/X})$ and in $(B, C) \in \text{Obj}(\mathsf{Sets}_{/X} \times \mathsf{Sets}_{/X})$.

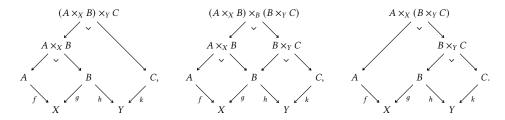
4. Associativity. Given a diagram



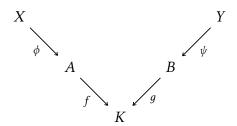
in Sets, we have isomorphisms of sets

$$(A \times_X B) \times_Y C \cong (A \times_X B) \times_B (B \times_Y C) \cong A \times_X (B \times_Y C),$$

where these pullbacks are built as in the diagrams



5. Interaction With Composition. Given a diagram



in Sets, we have isomorphisms of sets

$$\begin{split} X \times_K^{f \circ \phi, g \circ \psi} Y &\cong \left(X \times_A^{\phi, q_1} \left(A \times_K^{f, g} B \right) \right) \times_{A \times_K^{f, g} B}^{p_2, p_1} \left(\left(A \times_K^{f, g} B \right) \times_B^{q_2, \psi} Y \right) \\ &\cong X \times_A^{\phi, p} \left(\left(A \times_K^{f, g} B \right) \times_B^{q_2, \psi} Y \right) \\ &\cong \left(X \times_A^{\phi, q_1} \left(A \times_K^{f, g} B \right) \right) \times_B^{q, \psi} Y \end{split}$$

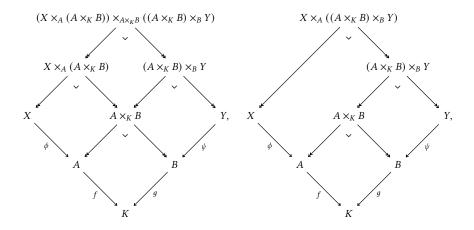
where

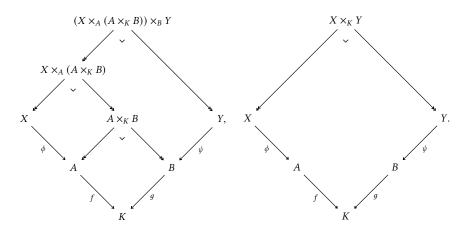
$$q_{1} = \operatorname{pr}_{1}^{A \times_{K}^{f,g} B}, \qquad q_{2} = \operatorname{pr}_{2}^{A \times_{K}^{f,g} B},$$

$$p_{1} = \operatorname{pr}_{1}^{\left(A \times_{K}^{f,g} B\right) \times_{Y}^{q_{2}, \psi}}, \qquad X \times_{A \times_{K}^{f,q_{1}} \left(A \times_{K}^{f,g} B\right)}^{\left(A \times_{K}^{f,g} B\right) \times_{B}^{q_{2}, \psi} Y},$$

$$p = q_{1} \circ \operatorname{pr}_{1}^{\left(A \times_{K}^{f,g} B\right) \times_{B}^{q_{2}, \psi} Y}, \qquad q = q_{2} \circ \operatorname{pr}_{2}^{X \times_{A}^{\phi, q_{1}} \left(A \times_{K}^{f,g} B\right)},$$

and where these pullbacks are built as in the following diagrams:





6. Unitality. We have isomorphisms of sets

natural in $(A, f) \in \text{Obj}(\mathsf{Sets}_{/X})$.

7. Commutativity. We have an isomorphism of sets

natural in (A, f), $(B, g) \in Obj(Sets_{/X})$.

8. Distributivity Over Coproducts. Let A, B, and C be sets and let $\phi_A \colon A \to X$, $\phi_B \colon B \to X$, and $\phi_C \colon C \to X$ be morphisms of sets. We have isomorphisms of sets

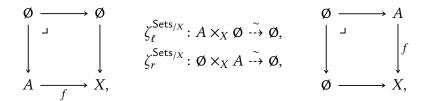
$$\delta_{\ell}^{\mathsf{Sets}_{/X}} : A \times_X (B \coprod C) \xrightarrow{\sim} (A \times_X B) \coprod (A \times_X C),$$

$$\delta_r^{\mathsf{Sets}_{/X}} : (A \coprod B) \times_X C \xrightarrow{\sim} (A \times_X C) \coprod (B \times_X C),$$

as in the diagrams

natural in $A, B, C \in Obj(Sets_{/X})$.

9. Annihilation With the Empty Set. We have isomorphisms of sets



natural in $(A, f) \in \text{Obj}(\mathsf{Sets}_{/X})$.

10. Interaction With Products. We have an isomorphism of sets

$$A \times_{\operatorname{pt}} B \cong A \times B, \qquad A \times_{\operatorname{pt}} B \cong A \times B, \qquad A \xrightarrow{!_{A}} \operatorname{pt}.$$

11. Symmetric Monoidality. The 8-tuple ($\mathsf{Sets}_{/X}$, x_X , X, $\mathsf{Sets}_{/X}$, $\alpha^{\mathsf{Sets}_{/X}}$, $\lambda^{\mathsf{Sets}_{/X}}$, $\rho^{\mathsf{Sets}_{/X}}$) is a symmetric closed monoidal category.

Proof. Item 1, *Functoriality*: This is a special case of functoriality of co/limits, Limits and Colimits, ?? of ??, with the explicit expression for ξ following from the commutativity of the cube pullback diagram.

Item 2, Adjointness I: This is a repetition of Fibred Sets, ?? of ??, and is proved there.

Item 3, Adjointness II: This follows from the universal property of the product (pullbacks are products in $Sets_{/X}$).

Item 4, Associativity: We have

$$(A \times_X B) \times_Y C \cong \{((a,b),c) \in (A \times_X B) \times C \mid h(b) = k(c)\}$$

$$\cong \{((a,b),c) \in (A \times B) \times C \mid f(a) = g(b) \text{ and } h(b) = k(c)\}$$

$$\cong \{(a,(b,c)) \in A \times (B \times C) \mid f(a) = g(b) \text{ and } h(b) = k(c)\}$$

$$\cong \{(a,(b,c)) \in A \times (B \times_Y C) \mid f(a) = g(b)\}$$

$$\cong A \times_X (B \times_Y C)$$

and

$$(A \times_X B) \times_B (B \times_Y C) \cong \{((a,b),(b',c)) \in (A \times_X B) \times (B \times_Y C) \mid b = b'\}$$

$$\cong \left\{ ((a,b),(b',c)) \in (A \times B) \times (B \times C) \mid f(a) = g(b), b = b', \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,(b,(b',c))) \in A \times (B \times (B \times C)) \mid f(a) = g(b), b = b', \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,((b,b'),c)) \in A \times ((B \times B) \times C) \mid f(a) = g(b), b = b', \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,((b,b'),c)) \in A \times ((B \times_B B) \times C) \mid f(a) = g(b) \text{ and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,(b,c)) \in A \times (B \times C) \mid f(a) = g(b) \text{ and } h(b) = k(c) \right\}$$

$$\cong A \times_X (B \times_Y C),$$

where we have used Item 6 for the isomorphism $B \times_B B \cong B$. *Item 5, Interaction With Composition*: By Item 4, it suffices to construct only the isomorphism

$$X\times_K^{f\circ\phi,g\circ\psi}Y\cong \left(X\times_A^{\phi,q_1}\left(A\times_K^{f,g}B\right)\right)\times_{A\times_V^{f,g}B}^{p_2,p_1}\left(\left(A\times_K^{f,g}B\right)\times_B^{q_2,\psi}Y\right).$$

We have

$$\left(X \times_A^{f,q_1} \left(A \times_K^{f,g} B \right) \right) \stackrel{\text{def}}{=} \left\{ (x, (a,b)) \in X \times \left(A \times_K^{f,g} B \right) \middle| \phi(x) = q_1(a,b) \right\}$$

$$\stackrel{\text{def}}{=} \left\{ (x, (a,b)) \in X \times \left(A \times_K^{f,g} B \right) \middle| \phi(x) = a \right\}$$

$$\cong \left\{ (x, (a,b)) \in X \times (A \times B) \middle| \phi(x) = a \text{ and } f(a) = g(b) \right\},$$

$$\left(\left(A \times_K^{f,g} B \right) \times_B^{q_2,\psi} Y \right) \stackrel{\text{def}}{=} \left\{ ((a,b),y) \in \left(A \times_K^{f,g} B \right) \times Y \middle| q_2(a,b) = \psi(y) \right\}$$

$$\stackrel{\text{def}}{=} \left\{ ((a,b),y) \in \left(A \times_K^{f,g} B \right) \times Y \middle| b = \psi(y) \right\}$$

$$\cong \{((a,b),y) \in (A \times B) \times Y \mid b = \psi(y) \text{ and } f(a) = g(b)\},$$

so writing

$$S = \left(X \times_A^{\phi, q_1} \left(A \times_K^{f, g} B\right)\right)$$
$$S' = \left(\left(A \times_K^{f, g} B\right) \times_B^{q_2, \psi} Y\right),$$

we have

$$\begin{split} S \times_{A \times_{K}^{f,g} B}^{p_{2},p_{1}} S' &\stackrel{\text{def}}{=} \{ ((x,(a,b)),((a',b'),y)) \in S \times S' \mid p_{1}(x,(a,b)) = p_{2}((a',b'),y) \} \\ &\stackrel{\text{def}}{=} \{ ((x,(a,b)),((a',b'),y)) \in S \times S' \mid (a,b) = (a',b') \} \\ &\cong \{ ((x,a,b,y)) \in X \times A \times B \times Y \mid \phi(x) = a, \psi(y) = b, \text{ and } f(a) = g(b) \} \\ &\stackrel{\text{def}}{=} \{ ((x,a,b,y)) \in X \times A \times B \times Y \mid f(\phi(x)) = g(\psi(y)) \} \\ &\stackrel{\text{def}}{=} X \times_{K} Y. \end{split}$$

This finishes the proof.

Item 6, Unitality: We have

$$X \times_X A \cong \{(x, a) \in X \times A \mid f(a) = x\},\$$
$$A \times_X X \cong \{(a, x) \in X \times A \mid f(a) = x\},\$$

which are isomorphic to A via the maps $(x, a) \mapsto a$ and $(a, x) \mapsto a$. The proof of the naturality of $\lambda^{\text{Sets}/X}$ and $\rho^{\text{Sets}/X}$ is omitted.

Item 7, Commutativity: We have

$$A \times_{C} B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid f(a) = g(b)\}$$

$$= \{(a, b) \in A \times B \mid g(b) = f(a)\}$$

$$\cong \{(b, a) \in B \times A \mid g(b) = f(a)\}$$

$$\stackrel{\text{def}}{=} B \times_{C} A.$$

The proof of the naturality of $\sigma^{\text{Sets}/X}$ is omitted. *Item 8, Distributivity Over Coproducts*: We have

$$A \times_{X} (B \coprod C) \stackrel{\text{def}}{=} \left\{ (a, z) \in A \times (B \coprod C) \middle| \phi_{A}(a) = \phi_{B \coprod C}(z) \right\}$$

$$= \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (0, b) \text{ and } \phi_{A}(a) = \phi_{B \coprod C}(z) \right\}$$

$$\cup \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (1, c) \text{ and } \phi_{A}(a) = \phi_{B \coprod C}(z) \right\}$$

$$= \{(a, z) \in A \times (B \coprod C) \mid z = (0, b) \text{ and } \phi_A(a) = \phi_B(b)\}$$

$$\cup \{(a, z) \in A \times (B \coprod C) \mid z = (1, c) \text{ and } \phi_A(a) = \phi_C(c)\}$$

$$\cong \{(a, b) \in A \times B \mid \phi_A(a) = \phi_B(b)\}$$

$$\cup \{(a, c) \in A \times C \mid \phi_A(a) = \phi_C(c)\}$$

$$\stackrel{\text{def}}{=} (A \times_X B) \cup (A \times_X C)$$

$$\cong (A \times_X B) \coprod (A \times_X C),$$

with the construction of the isomorphism

$$\delta_r^{\mathsf{Sets}_{/X}} : (A \coprod B) \times_X C \xrightarrow{\sim} (A \times_X C) \coprod (B \times_X C)$$

being similar. The proof of the naturality of $\delta_\ell^{\mathrm{Sets}/X}$ and $\delta_r^{\mathrm{Sets}/X}$ is omitted. *Item 9, Annihilation With the Empty Set*: We have

$$A \times_X \emptyset \stackrel{\text{def}}{=} \{ (a, b) \in A \times \emptyset \mid f(a) = g(b) \}$$
$$= \{ k \in \emptyset \mid f(a) = g(b) \}$$
$$= \emptyset.$$

and similarly for $\emptyset \times_X A$, where we have used Item 8 of Definition 4.1.3.1.3. The proof of the naturality of $\zeta_\ell^{\mathsf{Sets}_{/X}}$ and $\zeta_r^{\mathsf{Sets}_{/X}}$ is omitted. *Item 10, Interaction With Products*: We have

$$A \times_{\text{pt}} B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid !_A(a) = !_B(b)\}$$

$$\stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid \star = \star\}$$

$$= \{(a, b) \in A \times B\}$$

$$= A \times B.$$

Item 11, *Symmetric Monoidality*: Omitted.

4.1.5 Equalisers

Let *A* and *B* be sets and let $f, g: A \Rightarrow B$ be functions.

Definition 4.1.5.1.1. The **equaliser of** f **and** g is the equaliser of f and g in Sets as in Limits and Colimits, ??.

Construction 4.1.5.1.2. Concretely, the equaliser of f and g is the pair (Eq(f,g),eq(f,g)) consisting of:

1. *The Limit*. The set Eq(f, g) defined by

$$\operatorname{Eq}(f, q) \stackrel{\text{def}}{=} \{ a \in A \mid f(a) = q(a) \}.$$

2. *The Cone*. The inclusion map

$$eq(f, q) : Eq(f, q) \hookrightarrow A.$$

Proof. We claim that $\mathrm{Eq}(f,g)$ is the categorical equaliser of f and g in Sets. First we need to check that the relevant equaliser diagram commutes, i.e. that we have

$$f \circ \operatorname{eq}(f, q) = q \circ \operatorname{eq}(f, q),$$

which indeed holds by the definition of the set ${\rm Eq}(f,g)$. Next, we prove that ${\rm Eq}(f,g)$ satisfies the universal property of the equaliser. Suppose we have a diagram of the form

$$\operatorname{Eq}(f,g) \xrightarrow{\operatorname{eq}(f,g)} A \xrightarrow{f} B$$

$$E \xrightarrow{g} B$$

in Sets. Then there exists a unique map $\phi \colon E \to \text{Eq}(f, g)$ making the diagram

$$\begin{array}{ccc}
\operatorname{Eq}(f,g) & \xrightarrow{\operatorname{eq}(f,g)} & A & \xrightarrow{f} & B \\
\downarrow & & \downarrow & & \\
\downarrow & & & \\
E & & & \\
\end{array}$$

commute, being uniquely determined by the condition

$$eq(f, q) \circ \phi = e$$

via

$$\phi(x) = e(x)$$

for each $x \in E$, where we note that $e(x) \in A$ indeed lies in $\mathrm{Eq}(f,g)$ by the condition

$$f \circ e = g \circ e$$
,

which gives

$$f(e(x)) = g(e(x))$$

for each $x \in E$, so that $e(x) \in \text{Eq}(f, g)$.

Proposition 4.1.5.1.3. Let *A*, *B*, and *C* be sets.

1. Associativity. We have isomorphisms of sets⁵

$$\underbrace{\mathrm{Eq}(f \circ \mathrm{eq}(g,h), g \circ \mathrm{eq}(g,h))}_{=\mathrm{Eq}(f \circ \mathrm{eq}(g,h), h \circ \mathrm{eq}(g,h))} \cong \underbrace{\mathrm{Eq}(f,g,h)}_{=\mathrm{Eq}(g \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))} \underbrace{\mathrm{Eq}(f \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))}_{=\mathrm{Eq}(g \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))}$$

⁵That is, the following three ways of forming "the" equaliser of (f, g, h) agree:

1. Take the equaliser of (f, g, h), i.e. the limit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets.

2. First take the equaliser of f and g, forming a diagram

$$\operatorname{Eq}(f,g) \overset{\operatorname{eq}(f,g)}{\hookrightarrow} A \overset{f}{\underset{q}{\rightrightarrows}} B$$

and then take the equaliser of the composition

$$\operatorname{Eq}(f,g) \stackrel{\operatorname{eq}(f,g)}{\hookrightarrow} A \stackrel{f}{\underset{h}{\Longrightarrow}} B,$$

obtaining a subset

$$\mathrm{Eq}(f\circ\mathrm{eq}(f,g),h\circ\mathrm{eq}(f,g))=\mathrm{Eq}(g\circ\mathrm{eq}(f,g),h\circ\mathrm{eq}(f,g))$$
 of $\mathrm{Eq}(f,g).$

3. First take the equaliser of g and h, forming a diagram

$$\operatorname{Eq}(g,h) \overset{\operatorname{eq}(g,h)}{\hookrightarrow} A \overset{g}{\underset{h}{\Longrightarrow}} B$$

and then take the equaliser of the composition

$$\operatorname{Eq}(g,h) \overset{\operatorname{eq}(g,h)}{\hookrightarrow} A \overset{f}{\underset{g}{\Longrightarrow}} B,$$

obtaining a subset

$${\rm Eq}(f\circ {\rm eq}(g,h),g\circ {\rm eq}(g,h))={\rm Eq}(f\circ {\rm eq}(g,h),h\circ {\rm eq}(g,h))$$
 of ${\rm Eq}(g,h).$

where Eq(f, g, h) is the limit of the diagram

$$A \xrightarrow{f \atop h} B$$

in Sets, being explicitly given by

$$Eq(f, g, h) \cong \{a \in A \mid f(a) = g(a) = h(a)\}.$$

4. Unitality. We have an isomorphism of sets

$$\operatorname{Eq}(f, f) \cong A$$
.

5. Commutativity. We have an isomorphism of sets

$$\operatorname{Eq}(f,g) \cong \operatorname{Eq}(g,f)$$
.

6. Interaction With Composition. Let

$$A \stackrel{f}{\underset{q}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C$$

be functions. We have an inclusion of sets

$$\operatorname{Eq}(h \circ f \circ \operatorname{eq}(f, q), k \circ q \circ \operatorname{eq}(f, q)) \subset \operatorname{Eq}(h \circ f, k \circ q),$$

where ${\rm Eq}(h\circ f\circ {\rm eq}(f,g),k\circ g\circ {\rm eq}(f,g))$ is the equaliser of the composition

$$\operatorname{Eq}(f,g) \stackrel{\operatorname{eq}(f,g)}{\hookrightarrow} A \stackrel{f}{\underset{q}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C.$$

Proof. Item 1, Associativity: We first prove that Eq(f, g, h) is indeed given by

$$Eq(f, q, h) \cong \{a \in A \mid f(a) = q(a) = h(a)\}.$$

Indeed, suppose we have a diagram of the form

$$\operatorname{Eq}(f,g,h) \xrightarrow{\operatorname{eq}(f,g,h)} A \xrightarrow{f \atop h} B$$

in Sets. Then there exists a unique map $\phi\colon E\to \mathrm{Eq}(f,g,h),$ uniquely determined by the condition

$$eq(f, g) \circ \phi = e$$

being necessarily given by

$$\phi(x) = e(x)$$

for each $x \in E$, where we note that $e(x) \in A$ indeed lies in Eq(f, g, h) by the condition

$$f \circ e = g \circ e = h \circ e$$
,

which gives

$$f(e(x)) = g(e(x)) = h(e(x))$$

for each $x \in E$, so that $e(x) \in Eq(f, q, h)$.

We now check the equalities

$$\operatorname{Eq}(f \circ \operatorname{eq}(q, h), q \circ \operatorname{eq}(q, h)) \cong \operatorname{Eq}(f, q, h) \cong \operatorname{Eq}(f \circ \operatorname{eq}(f, q), h \circ \operatorname{eq}(f, q)).$$

Indeed, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(g,h), g \circ \operatorname{eq}(g,h)) &\cong \{x \in \operatorname{Eq}(g,h) \mid [f \circ \operatorname{eq}(g,h)](a) = [g \circ \operatorname{eq}(g,h)](a) \} \\ &\cong \{x \in \operatorname{Eq}(g,h) \mid f(a) = g(a) \} \\ &\cong \{x \in A \mid f(a) = g(a) \text{ and } g(a) = h(a) \} \\ &\cong \{x \in A \mid f(a) = g(a) = h(a) \} \\ &\cong \operatorname{Eq}(f,g,h). \end{split}$$

Similarly, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(f,g), h \circ \operatorname{eq}(f,g)) & \cong \{x \in \operatorname{Eq}(f,g) \mid [f \circ \operatorname{eq}(f,g)](a) = [h \circ \operatorname{eq}(f,g)](a) \} \\ & \cong \{x \in \operatorname{Eq}(f,g) \mid f(a) = h(a) \} \\ & \cong \{x \in A \mid f(a) = h(a) \text{ and } f(a) = g(a) \} \\ & \cong \{x \in A \mid f(a) = g(a) = h(a) \} \\ & \cong \operatorname{Eq}(f,g,h). \end{split}$$

Item 4, Unitality: Indeed, we have

$$\operatorname{Eq}(f, f) \stackrel{\text{def}}{=} \{ a \in A \mid f(a) = f(a) \}$$
$$= A.$$

Item 5, Commutativity: Indeed, we have

$$\operatorname{Eq}(f, q) \stackrel{\text{def}}{=} \{ a \in A \mid f(a) = g(a) \}$$

$$= \{a \in A \mid g(a) = f(a)\}\$$

$$\stackrel{\text{def}}{=} \text{Eq}(g, f).$$

Item 6, Interaction With Composition: Indeed, we have

$$\begin{aligned} \operatorname{Eq}(h \circ f \circ \operatorname{eq}(f, g), k \circ g \circ \operatorname{eq}(f, g)) & \cong \{a \in \operatorname{Eq}(f, g) \mid h(f(a)) = k(g(a))\} \\ & \cong \{a \in A \mid f(a) = g(a) \text{ and } h(f(a)) = k(g(a))\}. \end{aligned}$$

and

$$Eq(h \circ f, k \circ q) \cong \{a \in A \mid h(f(a)) = k(g(a))\},\$$

and thus there's an inclusion from $\operatorname{Eq}(h \circ f \circ \operatorname{eq}(f,g), k \circ g \circ \operatorname{eq}(f,g))$ to $\operatorname{Eq}(h \circ f, k \circ g)$.

4.1.6 Inverse Limits

Let $(X_{\alpha}, f_{\alpha\beta})_{\alpha \beta \in I} : (I, \preceq) \to \text{Sets be an inverse system of sets.}$

Definition 4.1.6.1.1. The **inverse limit of** $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the inverse limit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ in Sets as in Limits and Colimits, ??.

Construction 4.1.6.1.2. Concretely, the inverse limit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the pair $(\lim_{\longleftarrow} (X_{\alpha}), \{\operatorname{pr}_{\alpha}\}_{\alpha\in I})$ consisting of:

1. The Limit. The set $\lim_{\alpha \in I} (X_{\alpha})$ defined by

$$\lim_{\substack{\leftarrow \\ \alpha \in I}} (X_{\alpha}) \stackrel{\text{def}}{=} \left\{ (x_{\alpha})_{\alpha \in I} \in \prod_{\alpha \in I} X_{\alpha} \middle| \begin{array}{l} \text{for each } \alpha, \beta \in I, \text{ if } \alpha \preceq \beta, \\ \text{then we have } x_{\alpha} = f_{\alpha\beta}(x_{\beta}) \end{array} \right\}.$$

2. *The Cone*. The collection

$$\left\{ \operatorname{pr}_{\gamma} \colon \lim_{\stackrel{\longleftarrow}{\alpha \in I}} (X_{\alpha}) \to X_{\gamma} \right\}_{\gamma \in I}$$

of maps of sets defined as the restriction of the maps

$$\left\{ \operatorname{pr}_{\gamma} \colon \prod_{\alpha \in I} X_{\alpha} \to X_{\gamma} \right\}_{\gamma \in I}$$

4.1.6 Inverse Limits

of Item 2 of Definition 4.1.2.1.2 to $\lim_{\stackrel{\longleftarrow}{\alpha \in I}} (X_{\alpha})$ and hence given by

$$\operatorname{pr}_{V}((x_{\alpha})_{\alpha \in I}) \stackrel{\operatorname{def}}{=} x_{V}$$

for each $\gamma \in I$ and each $(x_{\alpha})_{\alpha \in I} \in \lim_{\alpha \in I} (X_{\alpha})$.

Proof. We claim that $\lim_{\alpha \to \alpha \in I} (X_{\alpha})$ is the limit of the inverse system of sets $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta \in I}$. First we need to check that the limit diagram defined by it commutes, i.e. that we have

$$f_{\alpha\beta} \circ \operatorname{pr}_{\alpha} = \operatorname{pr}_{\beta}, \qquad \varprojlim_{\alpha \in I} (X_{\alpha})$$

$$\operatorname{pr}_{\alpha} / \operatorname{pr}_{\beta}$$

$$X_{\alpha} \xrightarrow{f_{\alpha\beta}} X_{\beta}$$

for each $\alpha, \beta \in I$ with $\alpha \leq \beta$. Indeed, given $(x_{\gamma})_{\gamma \in I} \in \lim_{\leftarrow \gamma \in I} (X_{\gamma})$, we have

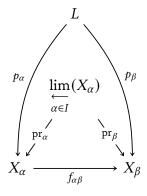
$$[f_{\alpha\beta} \circ \operatorname{pr}_{\alpha}] ((x_{\gamma})_{\gamma \in I}) \stackrel{\text{def}}{=} f_{\alpha\beta} (\operatorname{pr}_{\alpha} ((x_{\gamma})_{\gamma \in I}))$$

$$\stackrel{\text{def}}{=} f_{\alpha\beta} (x_{\alpha})$$

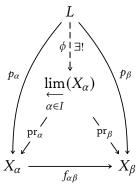
$$= x_{\beta}$$

$$\stackrel{\text{def}}{=} \operatorname{pr}_{\beta} ((x_{\gamma})_{\gamma \in I}),$$

where the third equality comes from the definition of $\lim_{\leftarrow \alpha \in I} (X_{\alpha})$. Next, we prove that $\lim_{\leftarrow \alpha \in I} (X_{\alpha})$ satisfies the universal property of an inverse limit. Suppose that we have, for each $\alpha, \beta \in I$ with $\alpha \preceq \beta$, a diagram of the form



in Sets. Then there indeed exists a unique map $\phi\colon L \xrightarrow{\exists !} \varprojlim_{\alpha \in I} (X_\alpha)$ making the diagram



commute, being uniquely determined by the family of conditions

$$\{p_{\alpha} = \operatorname{pr}_{\alpha} \circ \phi\}_{\alpha \in I}$$

via

$$\phi(\ell) = (p_{\alpha}(\ell))_{\alpha \in I}$$

for each $\ell \in L$, where we note that $(p_{\alpha}(\ell))_{\alpha \in I} \in \prod_{\alpha \in I} X_{\alpha}$ indeed lies in $\lim_{\epsilon \to \alpha \in I} (X_{\alpha})$, as we have

$$f_{lphaeta}(p_{lpha}(\ell))\stackrel{ ext{def}}{=} ig[f_{lphaeta}\circ p_{lpha}ig](\ell) \ \stackrel{ ext{def}}{=} p_{eta}(\ell)$$

for each $\beta \in I$ with $\alpha \leq \beta$ by the commutativity of the diagram for $(L, \{p_{\alpha}\}_{\alpha \in I})$.

Example 4.1.6.1.3. Here are some examples of inverse limits of sets.

1. *The p-Adic Integers*. The ring of *p*-adic integers \mathbb{Z}_p of $\ref{p-Adic}$ is the inverse limit

$$\mathbb{Z}_p \cong \lim_{\substack{\longleftarrow \\ n \in \mathbb{N}}} (\mathbb{Z}_{/p^n});$$

see??.

2. Rings of Formal Power Series. The ring R[t] of formal power series in a variable t is the inverse limit

$$R[[t]] \cong \lim_{\substack{\longleftarrow \\ n \in \mathbb{N}}} (R[t]/t^n R[t]);$$

see??.

3. *Profinite Groups*. Profinite groups are inverse limits of finite groups; see ??.

4.2 Colimits of Sets

4.2.1 The Initial Set

Definition 4.2.1.1.1. The **initial set** is the initial object of Sets as in Limits and Colimits, ??.

Construction 4.2.1.1.2. Concretely, the initial set is the pair $(\emptyset, \{\iota_A\}_{A \in \text{Obj}(\mathsf{Sets})})$ consisting of:

- 1. *The Colimit*. The empty set Ø of Definition 4.3.1.1.1.
- 2. The Cocone. The collection of maps

$$\{\iota_A\colon \emptyset\to A\}_{A\in \mathrm{Obj}(\mathsf{Sets})}$$

given by the inclusion maps from \emptyset to A.

Proof. We claim that Ø is the initial object of Sets. Indeed, suppose we have a diagram of the form

$$\emptyset$$
 A

in Sets. Then there exists a unique map $\phi \colon \emptyset \to A$ making the diagram

$$\emptyset - \frac{\phi}{\exists !} \rightarrow A$$

commute, namely the inclusion map ι_A .

4.2.2 Coproducts of Families of Sets

Let $\{A_i\}_{i\in I}$ be a family of sets.

Definition 4.2.2.1.1. The **coproduct of** $\{A_i\}_{i\in I}^6$ is the coproduct of $\{A_i\}_{i\in I}$ in Sets as in Limits and Colimits, ??.

⁶Further Terminology: Also called the **disjoint union of the family** $\{A_i\}_{i\in I}$.

Construction 4.2.2.1.2. Concretely, the disjoint union of $\{A_i\}_{i\in I}$ is the pair $\left(\coprod_{i\in I}A_i,\left\{\operatorname{inj}_i\right\}_{i\in I}\right)$ consisting of:

1. *The Colimit*. The set $\coprod_{i \in I} A_i$ defined by

$$\coprod_{i \in I} A_i \stackrel{\text{def}}{=} \left\{ (i, x) \in I \times \left(\bigcup_{i \in I} A_i \right) \middle| x \in A_i \right\}.$$

2. *The Cocone*. The collection

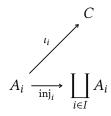
$$\left\{ \operatorname{inj}_i \colon A_i \to \coprod_{i \in I} A_i \right\}_{i \in I}$$

of maps given by

$$\operatorname{inj}_{i}(x) \stackrel{\text{def}}{=} (i, x)$$

for each $x \in A_i$ and each $i \in I$.

Proof. We claim that $\coprod_{i \in I} A_i$ is the categorical coproduct of $\{A_i\}_{i \in I}$ in Sets. Indeed, suppose we have, for each $i \in I$, a diagram of the form



in Sets. Then there exists a unique map $\phi\colon\coprod_{i\in I}A_i\to C$ making the diagram

$$A_{i} \xrightarrow[\text{inj}_{i}]{l_{i}} A_{i}$$

commute, being uniquely determined by the condition $\phi \circ \operatorname{inj}_i = \iota_i$ for each $i \in I$ via

$$\phi((i,x)) = \iota_i(x)$$

for each $(i, x) \in \coprod_{i \in I} A_i$.

Proposition 4.2.2.1.3. Let $\{A_i\}_{i\in I}$ be a family of sets.

1. Functoriality. The assignment $\{A_i\}_{i\in I}\mapsto \coprod_{i\in I}A_i$ defines a functor

$$\coprod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

• *Action on Objects*. For each $(A_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$, we have

$$\left[\bigsqcup_{i \in I} \right] ((A_i)_{i \in I}) \stackrel{\text{def}}{=} \bigsqcup_{i \in I} A_i$$

• Action on Morphisms. For each $(A_i)_{i \in I}$, $(B_i)_{i \in I} \in Obj(Fun(I_{disc}, Sets))$, the action on Hom-sets

$$\left(\bigsqcup_{i\in I}\right)_{(A_i)_{i\in I},(B_i)_{i\in I}}: \operatorname{Nat}((A_i)_{i\in I},(B_i)_{i\in I}) \to \operatorname{Sets}\left(\bigsqcup_{i\in I}A_i,\bigsqcup_{i\in I}B_i\right)$$

of $\coprod_{i\in I}$ at $((A_i)_{i\in I},(B_i)_{i\in I})$ is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in Nat $((A_i)_{i \in I}, (B_i)_{i \in I})$ to the map of sets

$$\coprod_{i \in I} f_i \colon \coprod_{i \in I} A_i \to \coprod_{i \in I} B_i$$

defined by

$$\left[\bigsqcup_{i \in I} f_i \right] (i, a) \stackrel{\text{def}}{=} f_i(a)$$

for each $(i, a) \in \coprod_{i \in I} A_i$.

Proof. Item 1, *Functoriality*: This follows from Limits and Colimits, ?? of ??. □

4.2.3 Binary Coproducts

Let *A* and *B* be sets.

Definition 4.2.3.1.1. The **coproduct of** A **and** B⁷ is the coproduct of A and B in Sets as in Limits and Colimits, ??.

Construction 4.2.3.1.2. Concretely, the coproduct of A and B is the pair $(A \coprod B, \{\text{inj}_1, \text{inj}_2\})$ consisting of:

1. *The Colimit*. The set $A \coprod B$ defined by

$$A \coprod B \stackrel{\text{def}}{=} \coprod_{z \in \{A, B\}} z$$

$$\stackrel{\text{def}}{=} \{(0, a) \in S \mid a \in A\} \cup \{(1, b) \in S \mid b \in B\},$$

where $S = \{0, 1\} \times (A \cup B)$.

2. The Cocone. The maps

$$\operatorname{inj}_1 : A \to A \coprod B,$$

 $\operatorname{inj}_2 : B \to A \coprod B,$

given by

$$\operatorname{inj}_{1}(a) \stackrel{\text{def}}{=} (0, a),$$

 $\operatorname{inj}_{2}(b) \stackrel{\text{def}}{=} (1, b),$

for each $a \in A$ and each $b \in B$.

Proof. We claim that $A \coprod B$ is the categorical coproduct of A and B in Sets. Indeed, suppose we have a diagram of the form

$$A \underset{\text{inj}_1}{\longleftrightarrow} A \coprod B \underset{\text{inj}_2}{\longleftrightarrow} B$$

in Sets. Then there exists a unique map $\phi: A \coprod B \to C$ making the diagram

$$A \xrightarrow{\text{inj}_{1}} A \coprod B \xrightarrow{\text{inj}_{2}} B$$

⁷Further Terminology: Also called the **disjoint union of** A **and** B.

commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_A = \iota_A,$$
 $\phi \circ \operatorname{inj}_B = \iota_B$

via

$$\phi(x) = \begin{cases} \iota_A(a) & \text{if } x = (0, a), \\ \iota_B(b) & \text{if } x = (1, b) \end{cases}$$

for each $x \in A \mid A \mid B$.

Proposition 4.2.3.1.3. Let *A*, *B*, *C*, and *X* be sets.

1. Functoriality. The assignment $A, B, (A, B) \mapsto A \coprod B$ defines functors

$$A \coprod -:$$
 Sets \rightarrow Sets,
 $- \coprod B:$ Sets \rightarrow Sets,
 $-_1 \coprod -_2:$ Sets \times Sets \rightarrow Sets,

where $-_1 \coprod -_2$ is the functor where

• *Action on Objects.* For each $(A, B) \in Obj(Sets \times Sets)$, we have

$$[-1][-2](A, B) \stackrel{\text{def}}{=} A [] B.$$

• Action on Morphisms. For each (A, B), $(X, Y) \in Obj(Sets)$, the action on Hom-sets

$$\coprod_{(A,B),(X,Y)} : \mathsf{Sets}(A,X) \times \mathsf{Sets}(B,Y) \to \mathsf{Sets}(A \coprod B,X \coprod Y)$$

of \coprod at ((A, B), (X, Y)) is defined by sending (f, g) to the function

$$f \coprod g: A \coprod B \to X \coprod Y$$

defined by

$$[f \coprod g](x) \stackrel{\text{def}}{=} \begin{cases} (0, f(a)) & \text{if } x = (0, a), \\ (1, g(b)) & \text{if } x = (1, b), \end{cases}$$

for each $x \in A \coprod B$.

and where $A \coprod -$ and $- \coprod B$ are the partial functors of $-_1 \coprod -_2$ at $A, B \in Obj(Sets)$.

2. Adjointness. We have an adjunction

$$(-_1 \coprod -_2 \dashv \Delta_{\mathsf{Sets}}) \colon \quad \mathsf{Sets} \times \mathsf{Sets} \underbrace{\stackrel{-_1 \coprod -_2}{\bot}}_{\Delta_{\mathsf{Sets}}} \mathsf{Sets},$$

witnessed by a bijection

$$Sets(A \coprod B, C), \cong Hom_{Sets \times Sets}((A, B), (C, C))$$

natural in $(A, B) \in Obj(Sets \times Sets)$ and in $C \in Obj(Sets)$.

3. Associativity. We have an isomorphism of sets

$$\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod} : (X \coprod Y) \coprod Z \xrightarrow{\sim} X \coprod (Y \coprod Z),$$

natural in $X, Y, Z \in Obj(Sets)$.

4. Unitality. We have isomorphisms of sets

$$\lambda_X^{\mathsf{Sets}, \coprod} \colon \emptyset \coprod X \xrightarrow{\sim} X,$$
$$\rho_X^{\mathsf{Sets}, \coprod} \colon X \coprod \emptyset \xrightarrow{\sim} X,$$

natural in $X \in Obj(Sets)$.

5. Commutativity. We have an isomorphism of sets

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod}: X\coprod Y\overset{\sim}{\dashrightarrow} Y\coprod X,$$

natural in $X, Y \in Obj(Sets)$.

6. *Symmetric Monoidality*. The 7-tuple (Sets, \coprod , \emptyset , α_{\coprod}^{Sets} , λ_{\coprod}^{Sets} , ρ_{\coprod}^{Sets} , σ_{\coprod}^{Sets}) is a symmetric monoidal category.

Proof. Item 1, Functoriality: This follows from Limits and Colimits, ?? of ??. *Item 2, Adjointness*: This follows from the universal property of the coproduct.

Item 3, Associativity: This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.2.3.1.1.

Item 4, Unitality: This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.2.4.1.1 and 5.2.5.1.1.

Item 5, Commutativity: This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.2.6.1.1.

Item 6, Symmetric Monoidality: This is a repetition of Monoidal Structures on the Category of Sets, Definition 5.2.7.1.1, and is proved there. □

4.2.4 Pushouts

Let A, B, and C be sets and let $f: C \to A$ and $g: C \to B$ be functions.

Definition 4.2.4.1.1. The **pushout of** A **and** B **over** C **along** f **and** g⁸ is the pushout of A and B over C along f and g in Sets as in Limits and Colimits, ??.

Construction 4.2.4.1.2. Concretely, the pushout of A and B over C along f and g is the pair $(A \coprod_C B, \{\text{inj}_1, \text{inj}_2\})$ consisting of:

1. *The Colimit*. The set $A \coprod_C B$ defined by

$$A \coprod_C B \stackrel{\text{def}}{=} A \coprod_B B/\sim_C$$
,

where \sim_C is the equivalence relation on $A \coprod B$ generated by $(0, f(c)) \sim_C (1, g(c))$.

2. The Cocone. The maps

$$inj_1: A \to A \coprod_C B,$$

 $inj_2: B \to A \coprod_C B$

given by

$$\operatorname{inj}_{1}(a) \stackrel{\text{def}}{=} [(0, a)]$$

 $\operatorname{inj}_{2}(b) \stackrel{\text{def}}{=} [(1, b)]$

for each $a \in A$ and each $b \in B$.

⁸ Further Terminology: Also called the **fibre coproduct of** A **and** B **over** C **along** f **and** g.

Proof. We claim that $A \coprod_C B$ is the categorical pushout of A and B over C with respect to (f,g) in Sets. First we need to check that the relevant pushout diagram commutes, i.e. that we have

$$\operatorname{inj}_{1} \circ f = \operatorname{inj}_{2} \circ g, \qquad A \coprod_{C} B \xleftarrow{\operatorname{inj}_{2}} B$$

$$\operatorname{inj}_{1} \downarrow \qquad \qquad \downarrow g$$

$$A \xleftarrow{f} C.$$

Indeed, given $c \in C$, we have

$$[\inf_{1} \circ f](c) = \inf_{1} (f(c))$$

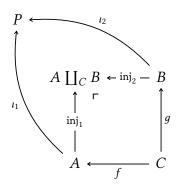
$$= [(0, f(c))]$$

$$= [(1, g(c))]$$

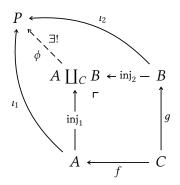
$$= \inf_{2} (g(c))$$

$$= [\inf_{2} \circ g](c),$$

where [(0, f(c))] = [(1, g(c))] by the definition of the relation \sim on $A \coprod B$. Next, we prove that $A \coprod {}_{C}B$ satisfies the universal property of the pushout. Suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: A \coprod_C B \to P$ making the diagram



commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_1 = \iota_1,$$

$$\phi \circ \operatorname{inj}_2 = \iota_2$$

via

$$\phi(x) = \begin{cases} \iota_1(a) & \text{if } x = [(0, a)], \\ \iota_2(b) & \text{if } x = [(1, b)] \end{cases}$$

for each $x \in A \coprod_C B$, where the well-definedness of ϕ is guaranteed by the equality $\iota_1 \circ f = \iota_2 \circ g$ and the definition of the relation \sim on $A \coprod B$ as follows:

1. Case 1: Suppose we have x = [(0, a)] = [(0, a')] for some $a, a' \in A$. Then, by Definition 4.2.4.1.3, we have a sequence

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0,a').$$

2. Case 2: Suppose we have x = [(1, b)] = [(1, b')] for some $b, b' \in B$. Then, by Definition 4.2.4.1.3, we have a sequence

$$(1,b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b').$$

3. Case 3: Suppose we have x = [(0, a)] = [(1, b)] for some $a \in A$ and $b \in B$. Then, by Definition 4.2.4.1.3, we have a sequence

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b).$$

In all these cases, we declare $x \sim' y$ iff there exists some $c \in C$ such that x = (0, f(c)) and y = (1, g(c)) or x = (1, g(c)) and y = (0, f(c)). Then, the equality $\iota_1 \circ f = \iota_2 \circ g$ gives

$$\phi([x]) = \phi([(0, f(c))])$$

$$\stackrel{\text{def}}{=} \iota_1(f(c))$$

$$= \iota_2(g(c))$$

$$\stackrel{\text{def}}{=} \phi([(1, g(c))])$$

$$= \phi([y]),$$

with the case where x=(1,g(c)) and y=(0,f(c)) similarly giving $\phi([x])=\phi([y])$. Thus, if $x\sim' y$, then $\phi([x])=\phi([y])$. Applying this equality pairwise to the sequences

$$(0, a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0, a'),$$

 $(1, b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1, b'),$
 $(0, a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1, b)$

gives

$$\phi([(0, a)]) = \phi([(0, a')]),$$

$$\phi([(1, b)]) = \phi([(1, b')]),$$

$$\phi([(0, a)]) = \phi([(1, b)]),$$

showing ϕ to be well-defined.

Remark 4.2.4.1.3. In detail, by Conditions on Relations, Definition 10.5.2.1.2, the relation \sim of Definition 4.2.4.1.1 is given by declaring $a \sim b$ iff one of the following conditions is satisfied:

- 1. We have $a, b \in A$ and a = b.
- 2. We have $a, b \in B$ and a = b.
- 3. There exist $x_1, \ldots, x_n \in A \coprod B$ such that $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$, where we declare $x \sim' y$ if one of the following conditions is satisfied:
 - (a) There exists $c \in C$ such that x = (0, f(c)) and y = (1, q(c)).

(b) There exists $c \in C$ such that x = (1, g(c)) and y = (0, f(c)).

In other words, there exist $x_1, \ldots, x_n \in A \coprod B$ satisfying the following conditions:

- (c) There exists $c_0 \in C$ satisfying one of the following conditions:
 - i. We have $a = f(c_0)$ and $x_1 = g(c_0)$.
 - ii. We have $a = q(c_0)$ and $x_1 = f(c_0)$.
- (d) For each $1 \le i \le n-1$, there exists $c_i \in C$ satisfying one of the following conditions:
 - i. We have $x_i = f(c_i)$ and $x_{i+1} = g(c_i)$.
 - ii. We have $x_i = q(c_i)$ and $x_{i+1} = f(c_i)$.
- (e) There exists $c_n \in C$ satisfying one of the following conditions:
 - i. We have $x_n = f(c_n)$ and $b = g(c_n)$.
 - ii. We have $x_n = g(c_n)$ and $b = f(c_n)$.

Remark 4.2.4.1.4. It is common practice to write $A \coprod_C B$ for the pushout of A and B over C along f and g, omitting the maps f and g from the notation and instead leaving them implicit, to be understood from the context.

However, the set $A \coprod_{C} B$ depends very much on the maps f and g, and sometimes it is necessary or useful to note this dependence explicitly. In such situations, we will write $A \coprod_{f,C,g} B$ or $A \coprod_{C} f B$ for $A \coprod_{C} B$.

Example 4.2.4.1.5. Here are some examples of pushouts of sets.

- 1. Wedge Sums of Pointed Sets. The wedge sum of two pointed sets of Pointed Sets, Definition 6.3.3.1.1 is an example of a pushout of sets.
- 2. *Intersections via Unions*. Let *X* be a set. We have

$$A \cup B \cong A \coprod_{A \cap B} B, \qquad \uparrow \qquad \uparrow \qquad \uparrow$$

$$A \longleftarrow A \cap B$$

for each $A, B \in \mathcal{P}(X)$.

Proof. Item 1, Wedge Sums of Pointed Sets: This follows by definition, as the wedge sum of two pointed sets is defined as a pushout.

Item 2, Intersections via Unions: Indeed, $A \coprod_{A \cap B} B$ is the quotient of $A \coprod B$ by the equivalence relation obtained by declaring $(0, a) \sim (1, b)$ iff $a = b \in A \cap B$, which is in bijection with $A \cup B$ via the map with $[(0, a)] \mapsto a$ and $[(1, b)] \mapsto b$.

Proposition 4.2.4.1.6. Let *A*, *B*, *C*, and *X* be sets.

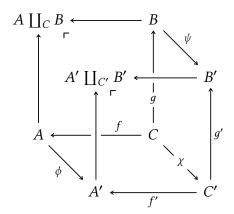
1. Functoriality. The assignment $(A,B,C,f,g)\mapsto A\coprod_{f,C,g} B$ defines a functor

$$-_1 \coprod_{-_3} -_1 : \mathsf{Fun}(\mathcal{P},\mathsf{Sets}) \to \mathsf{Sets},$$

where \mathcal{P} is the category that looks like this:



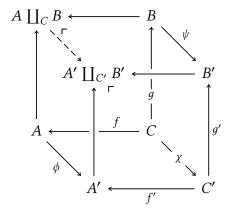
In particular, the action on morphisms of $-1 \coprod_{-3} -1$ is given by sending a morphism



in Fun(\mathcal{P} , Sets) to the map $\xi \colon A \coprod_C B \xrightarrow{\exists !} A' \coprod_{C'} B'$ given by

$$\xi(x) \stackrel{\text{def}}{=} \begin{cases} \phi(a) & \text{if } x = [(0, a)], \\ \psi(b) & \text{if } x = [(1, b)] \end{cases}$$

for each $x \in A \coprod_C B$, which is the unique map making the diagram



commute.

2. Adjointness. We have an adjunction

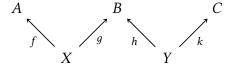
$$\left(-1 \coprod_{X} -_2 \dashv \Delta_{\mathsf{Sets}_{X/}}\right) \colon \mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/} \xrightarrow{-_1 \coprod_{X} -_2} \mathsf{Sets}_{X/},$$

witnessed by a bijection

$$\mathsf{Sets}_{X/}(A \coprod_X B, C), \cong \mathsf{Hom}_{\mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/}}((A, B), (C, C))$$

natural in $(A, B) \in \text{Obj}(\mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/})$ and in $C \in \text{Obj}(\mathsf{Sets}_{X/})$.

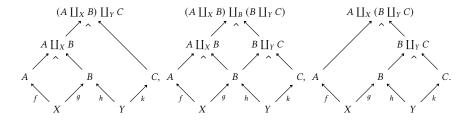
3. Associativity. Given a diagram



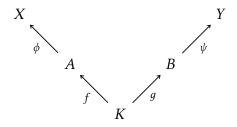
in Sets, we have isomorphisms of sets

$$(A \coprod_X B) \coprod_Y C \cong (A \coprod_X B) \coprod_B (B \coprod_Y C) \cong A \coprod_X (B \coprod_Y C)$$

where these pullbacks are built as in the diagrams



4. Interaction With Composition. Given a diagram



in Sets, we have isomorphisms of sets

$$\begin{split} X \coprod_K^{\phi \circ f, \psi \circ g} Y &\cong \left(X \coprod_A^{\phi, j_1} \left(A \coprod_K^{f, g} B \right) \right) \coprod_{A \coprod_K^{f, g} B}^{i_2, i_1} \left(\left(A \coprod_K^{f, g} B \right) \coprod_B^{j_2, \psi} Y \right) \\ &\cong X \coprod_A^{\phi, i} \left(\left(A \coprod_K^{f, g} B \right) \coprod_B^{j_2, \psi} Y \right) \\ &\cong \left(X \coprod_A^{\phi, i_1} \left(A \coprod_K^{f, g} B \right) \right) \coprod_B^{j, \psi} Y \end{split}$$

where

$$j_{1} = \operatorname{inj}_{1}^{A \times_{K}^{f,g} B}, \qquad j_{2} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B},$$

$$i_{1} = \operatorname{inj}_{1}^{\left(A \times_{K}^{f,g} B\right) \times_{Y}^{q_{2},\psi}}, \qquad i_{2} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B} \left(A \times_{K}^{f,g} B\right),$$

$$i_{2} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B} \left(A \times_{K}^{f,g} B\right),$$

$$i_{3} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B} \left(A \times_{K}^{f,g} B\right),$$

$$i_{4} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B} \left(A \times_{K}^{f,g} B\right),$$

$$i_{5} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B},$$

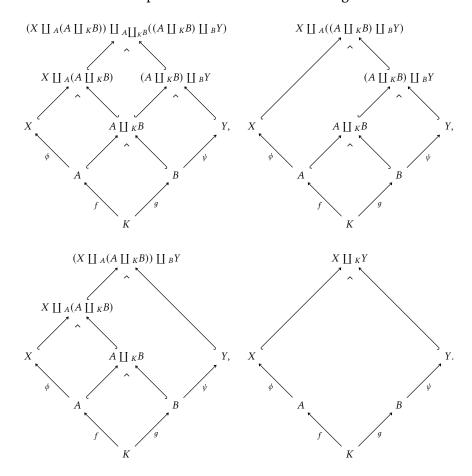
$$i_{6} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B},$$

$$i_{7} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B},$$

$$i_{8} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B},$$

$$i_{9} = \operatorname{inj}_{2}^{A \times_$$

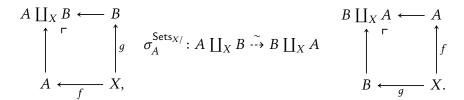
and where these pullbacks are built as in the diagrams



5. Unitality. We have isomorphisms of sets

 $\mathrm{natural}\,\mathrm{in}\;(A,f)\in\mathrm{Obj}\big(\mathsf{Sets}_{X/}\big).$

6. Commutativity. We have an isomorphism of sets



natural in (A, f), $(B, g) \in Obj(Sets_{X/})$.

7. Interaction With Coproducts. We have

$$A \coprod_{\emptyset} B \cong A \coprod B, \qquad \bigwedge^{\Gamma} \qquad \bigwedge^{I_{B}} l_{B}$$

$$A \longleftarrow_{I_{A}} \emptyset.$$

8. Symmetric Monoidality. The triple (Sets $_{X/}$, \coprod_{X} , X) is a symmetric monoidal category.

Proof. Item 1, *Functoriality*: This is a special case of functoriality of co/limits, Limits and Colimits, ?? of ??, with the explicit expression for ξ following from the commutativity of the cube pushout diagram.

Item 2, : Adjointness: This follows from the universal property of the coproduct (pushouts are coproducts in $\mathsf{Sets}_{X/}$).

Item 3, Associativity: Omitted.

Item 4, Interaction With Composition: Omitted.

Item 5, Unitality: Omitted.

Item 6, Commutativity: Omitted.

Item 7, Interaction With Coproducts: Omitted.

Item 8, Symmetric Monoidality: Omitted.

4.2.5 Coequalisers

Let A and B be sets and let $f, g: A \Rightarrow B$ be functions.

Definition 4.2.5.1.1. The **coequaliser of** f **and** g is the coequaliser of f and g in Sets as in Limits and Colimits, ??.

Construction 4.2.5.1.2. Concretely, the coequaliser of f and g is the pair (CoEq(f,g), coeq(f,g)) consisting of:

1. *The Colimit*. The set CoEq(f, g) defined by

$$CoEq(f, g) \stackrel{\text{def}}{=} B/\sim$$
,

where \sim is the equivalence relation on B generated by $f(a) \sim g(a)$.

2. The Cocone. The map

$$coeq(f,g): B \rightarrow CoEq(f,g)$$

given by the quotient map $\pi \colon B \twoheadrightarrow B/\sim$ with respect to the equivalence relation generated by $f(a) \sim g(a)$.

Proof. We claim that CoEq(f,g) is the categorical coequaliser of f and g in Sets. First we need to check that the relevant coequaliser diagram commutes, i.e. that we have

$$coeq(f, g) \circ f = coeq(f, g) \circ g.$$

Indeed, we have

$$[\operatorname{coeq}(f,g) \circ f](a) \stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](f(a))$$

$$\stackrel{\text{def}}{=} [f(a)]$$

$$= [g(a)]$$

$$\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](g(a))$$

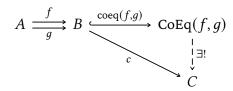
$$\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g) \circ g](a)$$

for each $a \in A$. Next, we prove that $\operatorname{CoEq}(f,g)$ satisfies the universal property of the coequaliser. Suppose we have a diagram of the form

$$A \xrightarrow{f} B \xrightarrow{\operatorname{coeq}(f,g)} \operatorname{CoEq}(f,g)$$

in Sets. Then, since c(f(a)) = c(g(a)) for each $a \in A$, it follows from Conditions on Relations, Items 4 and 5 of Definition 10.6.2.1.3 that there exists a

unique map $\operatorname{CoEq}(f,g) \stackrel{\exists !}{\longrightarrow} C$ making the diagram



commute.

Remark 4.2.5.1.3. In detail, by Conditions on Relations, Definition 10.5.2.1.2, the relation \sim of Definition 4.2.5.1.1 is given by declaring $a \sim b$ iff one of the following conditions is satisfied:

- 1. We have a = b;
- 2. There exist $x_1, \ldots, x_n \in B$ such that $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$, where we declare $x \sim' y$ if one of the following conditions is satisfied:
 - (a) There exists $z \in A$ such that x = f(z) and y = g(z).
 - (b) There exists $z \in A$ such that x = g(z) and y = f(z).

In other words, there exist $x_1, \ldots, x_n \in B$ satisfying the following conditions:

- (a) There exists $z_0 \in A$ satisfying one of the following conditions:
 - i. We have $a = f(z_0)$ and $x_1 = g(z_0)$.
 - ii. We have $a = g(z_0)$ and $x_1 = f(z_0)$.
- (b) For each $1 \le i \le n-1$, there exists $z_i \in A$ satisfying one of the following conditions:
 - i. We have $x_i = f(z_i)$ and $x_{i+1} = g(z_i)$.
 - ii. We have $x_i = g(z_i)$ and $x_{i+1} = f(z_i)$.
- (c) There exists $z_n \in A$ satisfying one of the following conditions:
 - i. We have $x_n = f(z_n)$ and $b = g(z_n)$.
 - ii. We have $x_n = g(z_n)$ and $b = f(z_n)$.

Example 4.2.5.1.4. Here are some examples of coequalisers of sets.

1. Quotients by Equivalence Relations. Let R be an equivalence relation on a set X. We have a bijection of sets

$$X/{\sim_R} \cong \operatorname{CoEq}\biggl(R \hookrightarrow X \times X \overset{\operatorname{pr}_1}{\underset{\operatorname{pr}_2}{\Longrightarrow}} X\biggr).$$

Proof. Item 1, Quotients by Equivalence Relations: See [Pro25z].

Proposition 4.2.5.1.5. Let *A*, *B*, and *C* be sets.

1. Associativity. We have isomorphisms of sets⁹

$$\underbrace{ \underbrace{ \text{CoEq}(\text{coeq}(f,g) \circ f, \text{coeq}(f,g) \circ h)}_{=\text{CoEq}(\text{coeq}(f,g) \circ g, \text{coeq}(f,g) \circ h)} \cong \underbrace{ \text{CoEq}(\text{coeq}(g,h) \circ f, \text{coeq}(g,h) \circ g, \text{coeq}(g,h) \circ h)}_{=\text{CoEq}(\text{coeq}(g,h) \circ f, \text{coeq}(g,h) \circ h)}$$

⁹That is, the following three ways of forming "the" coequaliser of (f, g, h) agree:

1. Take the coequaliser of (f, g, h), i.e. the colimit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets.

2. First take the coequaliser of f and g, forming a diagram

$$A \stackrel{f}{\underset{q}{\Longrightarrow}} B \stackrel{\operatorname{coeq}(f,g)}{\Longrightarrow} \operatorname{CoEq}(f,g)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\underset{h}{\Longrightarrow}} B \stackrel{\text{coeq}(f,g)}{\Longrightarrow} \text{CoEq}(f,g),$$

obtaining a quotient

$$\label{eq:coeq} \begin{aligned} \operatorname{CoEq}(\operatorname{coeq}(f,g) \circ f, \operatorname{coeq}(f,g) \circ h) &= \operatorname{CoEq}(\operatorname{coeq}(f,g) \circ g, \operatorname{coeq}(f,g) \circ h) \\ \\ \operatorname{of} \operatorname{CoEq}(f,g) &= \operatorname{CoEq}(f,g) \circ f, \operatorname{coeq}(f,g) \circ h \end{aligned}$$

3. First take the coequaliser of g and h, forming a diagram

$$A \stackrel{g}{\underset{h}{\Longrightarrow}} B \stackrel{\text{coeq}(g,h)}{\twoheadrightarrow} \text{CoEq}(g,h)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{\text{coeq}(g,h)}{\Longrightarrow} \text{CoEq}(g,h),$$

obtaining a quotient

$$\label{eq:coeq} \begin{split} \mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ g) &= \mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ h) \\ &\quad \mathsf{of}\, \mathsf{CoEq}(g,h). \end{split}$$

where CoEq(f, g, h) is the colimit of the diagram

$$A \xrightarrow{f \atop h} B$$

in Sets.

4. Unitality. We have an isomorphism of sets

$$CoEq(f, f) \cong B$$
.

5. Commutativity. We have an isomorphism of sets

$$CoEq(f,g) \cong CoEq(g,f)$$
.

6. Interaction With Composition. Let

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C$$

be functions. We have a surjection

 $CoEq(h \circ f, k \circ g) \twoheadrightarrow CoEq(coeq(h, k) \circ h \circ f, coeq(h, k) \circ k \circ g)$

exhibiting $\operatorname{CoEq}(\operatorname{coeq}(h,k) \circ h \circ f, \operatorname{coeq}(h,k) \circ k \circ g)$ as a quotient of $\operatorname{CoEq}(h \circ f, k \circ g)$ by the relation generated by declaring $h(y) \sim k(y)$ for each $y \in B$.

Proof. Item 1, Associativity: Omitted.

Item 4, Unitality: Omitted.

Item 5, Commutativity: Omitted.

Item 6, Interaction With Composition: Omitted.

4.2.6 Direct Colimits

Let $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$: $(I, \preceq) \to \mathbb{T}$ be a direct system of sets.

Definition 4.2.6.1.1. The **direct colimit of** $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the direct colimit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ in Sets as in Limits and Colimits, ??.

Construction 4.2.6.1.2. Concretely, the direct colimit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the pair $\left(\operatorname{colim}(X_{\alpha}), \left\{\operatorname{inj}_{\alpha}\right\}_{\alpha\in I}\right)$ consisting of:

1. The Colimit. The set $\operatornamewithlimits{colim}_{\substack{\alpha \in I \\ \alpha \in I}}(X_{\alpha})$ defined by

where \sim is the equivalence relation on $\coprod_{\alpha \in I} X_{\alpha}$ generated by declaring $(\alpha, x) \sim (\beta, y)$ iff there exists some $\gamma \in I$ satisfying the following conditions:

- (a) We have $\alpha \leq \gamma$.
- (b) We have $\beta \leq \gamma$.
- (c) We have $f_{\alpha\gamma}(x) = f_{\beta\gamma}(y)$.
- 2. The Cocone. The collection

$$\left\{\operatorname{inj}_{\gamma} \colon X_{\gamma} \to \operatorname{colim}_{\alpha \in I}(X_{\alpha})\right\}_{\gamma \in I}$$

of maps of sets defined by

$$\operatorname{inj}_{\gamma}(x) \stackrel{\text{def}}{=} [(\gamma, x)]$$

for each $y \in I$ and each $x \in X_y$.

Proof. We will prove Definition 4.2.6.1.2 below in a bit, but first we need a lemma (which is interesting in its own right).

Lemma 4.2.6.1.3. For each $\alpha, \beta \in I$ and each $x \in X_{\alpha}$, if $\alpha \leq \beta$, then we have

$$(\alpha, x) \sim (\beta, f_{\alpha\beta}(x))$$

in $\operatornamewithlimits{colim}_{\stackrel{\longrightarrow}{\alpha \in I}}(X_{\alpha})$.

Proof. Taking $\gamma = \beta$, we have $f_{\alpha\gamma} = f_{\alpha\beta}$, we have $f_{\beta\gamma} = f_{\beta\beta} \stackrel{\text{def}}{=} \mathrm{id}_{X_{\beta}}$, and we have

$$f_{\alpha\beta}(x) = f_{\beta\beta} (f_{\alpha\beta}(x))$$

$$\stackrel{\text{def}}{=} \mathrm{id}_{X_{\beta}} (f_{\alpha\beta}(x)),$$

$$= f_{\alpha\beta}(x).$$

As a result, since $\alpha \leq \beta$ and $\beta \leq \beta$ as well, Items 1a to 1c of Definition 4.2.6.1.2 are met. Thus we have $(\alpha, x) \sim (\beta, f_{\alpha\beta}(x))$.

We can now prove Definition 4.2.6.1.2:

Proof. We claim that $\operatornamewithlimits{colim}_{\alpha \in I}(X_\alpha)$ is the colimit of the direct system of sets $(X_\alpha, f_{\alpha\beta})_{\alpha,\beta \in I}$.

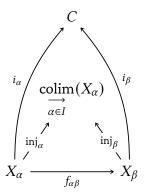
Commutativity of the Colimit Diagram: First, we need to check that the colimit diagram defined by colim (X_{α}) commutes, i.e. that we have

$$\operatorname{inj}_{\alpha} = \operatorname{inj}_{\beta} \circ f_{\alpha\beta}, \quad \begin{array}{c} \operatorname{colim}(X_{\alpha}) \\ \xrightarrow{\alpha \in I} \\ \operatorname{inj}_{\alpha} / \\ X_{\alpha} \xrightarrow{f_{\alpha\beta}} X_{\beta} \end{array}$$

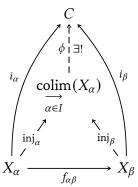
for each α , $\beta \in I$ with $\alpha \leq \beta$. Indeed, given $x \in X_{\alpha}$, we have

$$\left[\inf_{\beta} \circ f_{\alpha\beta} \right](x) \stackrel{\text{def}}{=} \inf_{\beta} \left(f_{\alpha\beta}(x) \right) \\
\stackrel{\text{def}}{=} \left[\left(\beta, f_{\alpha\beta}(x) \right) \right] \\
= \left[(\alpha, x) \right] \\
\stackrel{\text{def}}{=} \inf_{\alpha} (x),$$

where we have used Definition 4.2.6.1.3 for the third equality. Proof of the Universal Property of the Colimit: Next, we prove that colim (X_{α}) as constructed in Definition 4.2.6.1.2 satisfies the universal property of a direct colimit. Suppose that we have, for each $\alpha, \beta \in I$ with $\alpha \preceq \beta$, a diagram of the form



in Sets. We claim that there exists a unique map $\phi: \underset{\alpha \in I}{\operatorname{colim}}(X_{\alpha}) \xrightarrow{\exists !} C$ making the diagram



commute. To this end, first consider the diagram

$$\coprod_{\alpha \in I} X_{\alpha} \xrightarrow{\operatorname{pr}} \operatorname{colim}_{\alpha \in I} (X_{\alpha})$$

$$\coprod_{\alpha \in I} i_{\alpha}$$

$$C.$$

Lemma. If $(\alpha, x) \sim (\beta, y)$, then we have

$$\left[\coprod_{\alpha\in I}i_{\alpha}\right](x)=\left[\coprod_{\alpha\in I}i_{\alpha}\right](y).$$

Proof. Indeed, if $(\alpha, x) \sim (\beta, y)$, then there exists some $\gamma \in I$ satisfying the following conditions:

- 1. We have $\alpha \leq \gamma$.
- 2. We have $\beta \leq \gamma$.
- 3. We have $f_{\alpha y}(x) = f_{\beta y}(y)$.

We then have

$$\left[\bigsqcup_{\alpha \in I} i_{\alpha} \right](x) \stackrel{\text{def}}{=} i_{\alpha}(x)$$

$$\stackrel{\text{def}}{=} \left[i_{\gamma} \circ f_{\alpha \gamma} \right](x)$$

$$\stackrel{\text{def}}{=} i_{\gamma} (f_{\alpha \gamma}(x))$$

$$= i_{\gamma} (f_{\beta \gamma}(x))$$

$$\stackrel{\text{def}}{=} \left[i_{\gamma} \circ f_{\beta \gamma} \right](x)$$

$$= i_{\beta}(y)$$

$$\stackrel{\text{def}}{=} \left[\bigsqcup_{\alpha \in I} i_{\alpha} \right](y).$$

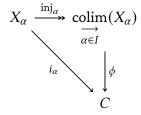
This finishes the proof of the lemma. Continuing, by Conditions on Relations, ?? of Definition 10.6.2.1.3, there then exists a map $\phi: \operatorname{colim}(X_\alpha) \stackrel{\exists !}{\longrightarrow} C$ making the diagram

$$\coprod_{\alpha \in I} X_{\alpha} \xrightarrow{\operatorname{pr}} \operatorname{colim}(X_{\alpha})$$

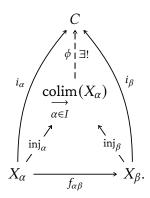
$$\coprod_{\alpha \in I} i_{\alpha} \qquad \downarrow \phi$$

$$C$$

commute. In particular, this implies that the diagram



also commutes, and thus so does the diagram



This finishes the proof.¹⁰

Example 4.2.6.1.4. Here are some examples of direct colimits of sets.

1. *The Prüfer Group.* The Prüfer group $\mathbb{Z}(p^{\infty})$ is defined as the direct colimit

$$\mathbb{Z}(p^{\infty}) \stackrel{\text{def}}{=} \underset{n \in \mathbb{N}}{\operatorname{colim}} (\mathbb{Z}_{/p^n});$$

see??.

4.3 Operations With Sets

4.3.1 The Empty Set

Definition 4.3.1.1.1. The **empty set** is the set \emptyset defined by

$$\emptyset \stackrel{\text{def}}{=} \{ x \in X \mid x \neq x \},\$$

$$\left\{i_{\alpha} = \phi \circ \operatorname{inj}_{\alpha}\right\}_{\alpha \in I}$$

show that ϕ must be given by

$$\phi([(\alpha, x)]) = (i_{\alpha}(x))_{\alpha \in I}$$

for each $[(\alpha,x)] \in \operatornamewithlimits{colim}_{\alpha \in I}(X_\alpha)$, although we would need to show that this assignment is well-defined were we to prove Definition 4.2.6.1.2 in this way. Instead, invoking Conditions

¹⁰Incidentally, the conditions

where X is the set in the set existence axiom, ?? of ??.

4.3.2 Singleton Sets

Let X be a set.

Definition 4.3.2.1.1. The **singleton set containing** X is the set $\{X\}$ defined by

$$\{X\} \stackrel{\text{def}}{=} \{X, X\},$$

where $\{X, X\}$ is the pairing of X with itself of Definition 4.3.3.1.1.

4.3.3 Pairings of Sets

Let *X* and *Y* be sets.

Definition 4.3.3.1.1. The **pairing of** X **and** Y is the set $\{X, Y\}$ defined by

$${X, Y} \stackrel{\text{def}}{=} {x \in A \mid x = X \text{ or } x = Y},$$

where A is the set in the axiom of pairing, ?? of ??.

4.3.4 Ordered Pairs

Let *A* and *B* be sets.

Definition 4.3.4.1.1. The **ordered pair associated to** A **and** B is the set (A, B) defined by

$$(A, B) \stackrel{\text{def}}{=} \{ \{A\}, \{A, B\} \}.$$

Proposition 4.3.4.1.2. Let *A* and *B* be sets.

- 1. *Uniqueness*. Let *A*, *B*, *C*, and *D* be sets. The following conditions are equivalent:
 - (a) We have (A, B) = (C, D).
 - (b) We have A = C and B = D.

Proof. Item 1, Uniqueness: See [Cie97, Theorem 1.2.3].

4.3.5 Sets of Maps

Let A and B be sets.

Definition 4.3.5.1.1. The **set of maps from** A **to** B^{11} is the set $Sets(A, B)^{12}$ whose elements are the functions from A to B.

Proposition 4.3.5.1.2. Let *A* and *B* be sets.

1. Functoriality. The assignments $X, Y, (X, Y) \mapsto \operatorname{Hom}_{\mathsf{Sets}}(X, Y)$ define functors

Sets
$$(X, -)$$
: Sets \rightarrow Sets,
Sets $(-, Y)$: Sets op \rightarrow Sets,
Sets $(-_1, -_2)$: Sets op \times Sets \rightarrow Sets.

2. Adjointness. We have adjunctions

$$(A \times - \dashv \operatorname{Sets}(A, -))$$
: Sets $\xrightarrow{A \times -}$ Sets, $(- \times B \dashv \operatorname{Sets}(B, -))$: Sets $\xrightarrow{- \times B}$ Sets, $\operatorname{Sets}(B, -)$

witnessed by bijections

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$

 $Sets(A \times B, C) \cong Sets(B, Sets(A, C)),$

natural in $A, B, C \in Obj(Sets)$.

3. *Maps From the Punctual Set*. We have a bijection

$$Sets(pt, A) \cong A$$
,

natural in $A \in Obj(Sets)$.

on Relations, ?? of Definition 10.6.2.1.3 gave us a way to avoid having to prove this, leading to a cleaner alternative proof.

¹¹Further Terminology: Also called the **Hom set from** A **to** B.

¹² Further Notation: Also written $Hom_{Sets}(A, B)$.

4. Maps to the Punctual Set. We have a bijection

$$Sets(A, pt) \cong pt$$
,

natural in $A \in Obj(Sets)$.

Proof. Item 1, Functoriality: This follows from Categories, Items 2 and 5 of Definition 11.1.4.1.2.

Item 2, Adjointness: This is a repetition of Item 2 of Definition 4.1.3.1.3 and is proved there.

Item 3, Maps From the Punctual Set: The bijection

$$\Phi_A \colon \mathsf{Sets}(\mathsf{pt},A) \xrightarrow{\sim} A$$

is given by

$$\Phi_A(f) \stackrel{\text{def}}{=} f(\star)$$

for each $f \in Sets(pt, A)$, admitting an inverse

$$\Phi_A^{-1} \colon A \xrightarrow{\sim} \mathsf{Sets}(\mathsf{pt}, A)$$

given by

$$\Phi_A^{-1}(a) \stackrel{\text{def}}{=} \llbracket \star \mapsto a \rrbracket$$

for each $a \in A$. Indeed, we have

$$\begin{split} \left[\Phi_A^{-1} \circ \Phi_A\right](f) &\stackrel{\text{def}}{=} \Phi_A^{-1}(\Phi_A(f)) \\ &\stackrel{\text{def}}{=} \Phi_A^{-1}(f(\star)) \\ &\stackrel{\text{def}}{=} \left[\!\!\left[\star \mapsto f(\star)\right]\!\!\right] \\ &\stackrel{\text{def}}{=} f \\ &\stackrel{\text{def}}{=} \left[id_{\mathsf{Sets}(\mathsf{pt},A)}\right](f) \end{split}$$

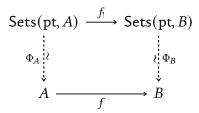
for each $f \in Sets(pt, A)$ and

$$\begin{split} \left[\Phi_{A} \circ \Phi_{A}^{-1}\right](a) &\stackrel{\text{def}}{=} \Phi_{A}\left(\Phi_{A}^{-1}(a)\right) \\ &\stackrel{\text{def}}{=} \Phi_{A}(\llbracket \star \mapsto a \rrbracket) \\ &\stackrel{\text{def}}{=} \operatorname{ev}_{\star}(\llbracket \star \mapsto a \rrbracket) \\ &\stackrel{\text{def}}{=} a \\ &\stackrel{\text{def}}{=} [\operatorname{id}_{A}](a) \end{split}$$

for each $a \in A$, and thus we have

$$\begin{split} & \Phi_A^{-1} \circ \Phi_A = \mathrm{id}_{\mathsf{Sets}(\mathsf{pt},A)} \\ & \Phi_A \circ \Phi_A^{-1} = \mathrm{id}_A \,. \end{split}$$

To prove naturality, we need to show that the diagram



commutes. Indeed, we have

$$[f \circ \Phi_A](\phi) \stackrel{\text{def}}{=} f(\Phi_A(\phi))$$

$$\stackrel{\text{def}}{=} f(\phi(\star))$$

$$\stackrel{\text{def}}{=} [f \circ \phi](\star)$$

$$\stackrel{\text{def}}{=} \Phi_B(f \circ \phi)$$

$$\stackrel{\text{def}}{=} \Phi_B(f_!(\phi))$$

$$\stackrel{\text{def}}{=} [\Phi_B \circ f_!](\phi)$$

for each $\phi \in Sets(pt, A)$. This finishes the proof.

Item 4, Maps to the Punctual Set: This follows from the universal property of pt as the terminal set, Definition 4.1.1.1.1. □

4.3.6 Unions of Families of Subsets

Let X be a set and let $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

Definition 4.3.6.1.1. The **union of** \mathcal{U} is the set $\bigcup_{U \in \mathcal{U}} U$ defined by

$$\bigcup_{U \in \mathcal{U}} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \mathcal{U} \\ \text{such that we have } x \in U \end{array} \right\}.$$

Proposition 4.3.6.1.2. Let X be a set.

1. Functoriality. The assignment $\mathcal{U}\mapsto \bigcup_{U\in\mathcal{U}} U$ defines a functor

$$[]: (\mathcal{P}(\mathcal{P}(X)), \subset) \to (\mathcal{P}(X), \subset).$$

In particular, for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$, the following condition is satisfied:

$$(\star) \ \text{ If } \mathcal{U} \subset \mathcal{V} \text{, then } \bigcup_{U \in \mathcal{U}} U \subset \bigcup_{V \in \mathcal{V}} V.$$

2. Associativity. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(\mathcal{P}(X))) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \bigcup} & \mathcal{P}(\mathcal{P}(X)) \\
\cup \star \mathrm{id}_{\mathcal{P}(X)} & & & \bigcup \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{} & & & \\
\end{array}$$

commutes, i.e. we have

$$\bigcup_{U \in \bigcup_{A \in A} A} U = \bigcup_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U \right)$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$.

3. Left Unitality. The diagram

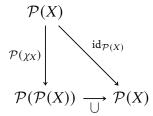
$$\begin{array}{c|c}
\mathcal{P}(X) & & \\
\downarrow^{\chi_{\mathcal{P}(X)}} & & & \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcup} \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\bigcup_{V \in \{U\}} V = U$$

for each $U \in \mathcal{P}(X)$.

4. Right Unitality. The diagram



commutes, i.e. we have

$$\bigcup_{\{u\}\in\chi_X(U)}\{u\}=U$$

for each $U \in \mathcal{P}(X)$.

5. Interaction With Unions I. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) & \stackrel{\cup}{\longrightarrow} \mathcal{P}(\mathcal{P}(X)) \\
\downarrow & & \downarrow \\
\mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{\bot} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\bigcup_{W \in \mathcal{U} \cup \mathcal{V}} W = \left(\bigcup_{U \in \mathcal{U}} U\right) \cup \left(\bigcup_{V \in \mathcal{V}} V\right)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

6. Interaction With Unions II. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcup_{V \in \mathcal{V}} V\right) = \bigcup_{V \in \mathcal{V}} (U \cup V),$$

$$\left(\bigcup_{U\in\mathcal{U}}U\right)\cup V=\bigcup_{U\in\mathcal{U}}(U\cup V)$$

for each $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

7. Interaction With Intersections I. We have a natural transformation

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\cap} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \qquad \qquad \bigcup \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\cap} \mathcal{P}(X),$$

with components

$$\bigcup_{W \in \mathcal{U} \cap \mathcal{V}} W \subset \left(\bigcup_{U \in \mathcal{U}} U\right) \cap \left(\bigcup_{V \in \mathcal{V}} V\right)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

8. Interaction With Intersections II. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcup_{V \in \mathcal{V}} V\right) = \bigcup_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcup_{U \in \mathcal{U}} U\right) \cup V = \bigcup_{U \in \mathcal{U}} (U \cup V)$$

for each $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

9. Interaction With Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\setminus} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\setminus} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in \mathcal{U} \setminus \mathcal{V}} W \neq \left(\bigcup_{U \in \mathcal{U}} U\right) \setminus \left(\bigcup_{V \in \mathcal{V}} V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

10. Interaction With Complements I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\text{op}} \xrightarrow{(-)^{c}} \mathcal{P}(\mathcal{P}(X))$$

$$\downarrow^{\text{op}} \qquad \qquad \downarrow \cup$$

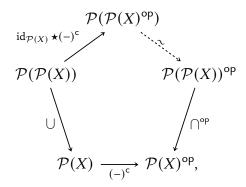
$$\mathcal{P}(X)^{\text{op}} \xrightarrow{(-)^{c}} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{U \in \mathcal{U}^{\mathsf{c}}} U \neq \bigcup_{U \in \mathcal{U}} U^{\mathsf{c}}$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

11. Interaction With Complements II. The diagram

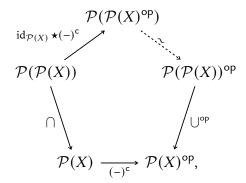


commutes, i.e. we have

$$\left(\bigcup_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

12. Interaction With Complements III. The diagram



commutes, i.e. we have

$$\left(\bigcap_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

13. Interaction With Symmetric Differences. The diagram

$$\begin{array}{c|c} \mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) & \stackrel{\triangle}{\longrightarrow} \mathcal{P}(\mathcal{P}(X)) \\ & & \swarrow & & \downarrow \cup \\ & \mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{\Delta} & \mathcal{P}(X), \end{array}$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in \mathcal{U} \triangle \mathcal{V}} W \neq \left(\bigcup_{U \in \mathcal{U}} U\right) \triangle \left(\bigcup_{V \in \mathcal{V}} V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

14. Interaction With Internal Homs I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\cup^{\mathsf{op}} \times \cup^{\mathsf{op}} \qquad \qquad \bigcup \cup$$

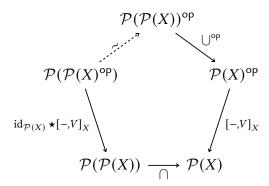
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcup_{U \in \mathcal{U}} U, \bigcup_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

15. Interaction With Internal Homs II. The diagram



commutes, i.e. we have

$$\left[\bigcup_{U \in \mathcal{U}} U, V\right]_X = \bigcap_{U \in \mathcal{U}} [U, V]_X$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

16. Interaction With Internal Homs III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcup} & \mathcal{P}(X) \\
id_{\mathcal{P}(X)} \star [U,-]_X & & & \downarrow [U,-]_X \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\downarrow} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U, \bigcup_{V \in \mathcal{V}} V\right]_X = \bigcup_{V \in \mathcal{V}} [U, V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

17. Interaction With Direct Images. Let $f: X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_!)_!} \mathcal{P}(\mathcal{P}(Y))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \bigcup$$

$$\mathcal{P}(X) \xrightarrow{f} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_!(U) = \bigcup_{V \in f_!(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

18. Interaction With Inverse Images. Let $f: X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{V\in\mathcal{V}}f^{-1}(V)=\bigcup_{U\in f^{-1}(\mathcal{U})}U$$

for each $\mathcal{V} \in \mathcal{P}(Y)$, where $f^{-1}(\mathcal{V}) \stackrel{\text{def}}{=} (f^{-1})^{-1}(\mathcal{V})$.

19. Interaction With Codirect Images. Let $f\colon X\to Y$ be a map of sets. The

diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \cup$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_*(U) = \bigcup_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

20. Interaction With Intersections of Families I. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \cap} & \mathcal{P}(\mathcal{P}(x)) \\
\downarrow \star \mathrm{id}_{\mathcal{P}(X)} & & & \downarrow \cap \\
\mathcal{P}(X) & \xrightarrow{\bigcap} & X
\end{array}$$

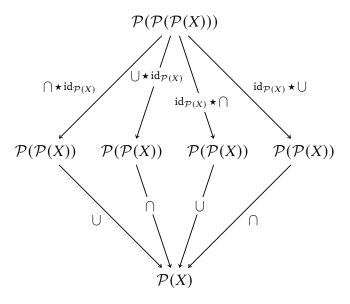
commutes, i.e. we have

$$\bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U = \bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right)$$

for each $A \in \mathcal{P}(\mathcal{P}(X))$.

21. Interaction With Intersections of Families II. Let X be a set and consider

the compositions

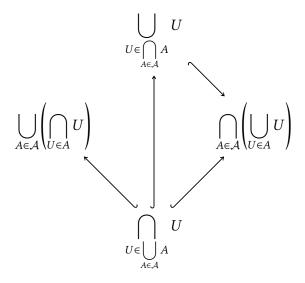


given by

$$\mathcal{A} \mapsto \bigcup_{U \in \bigcap_{A \in \mathcal{A}} A} U, \quad \mathcal{A} \mapsto \bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U,$$

$$\mathcal{A} \mapsto \bigcup_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U\right), \quad \mathcal{A} \mapsto \bigcap_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U\right)$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$. We have the following inclusions:



All other possible inclusions fail to hold in general.

Proof. Item 1, *Functoriality*: Since $\mathcal{P}(X)$ is posetal, it suffices to prove the condition (\star) . So let $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ with $\mathcal{U} \subset \mathcal{V}$. We claim that

$$\bigcup_{U\in\mathcal{U}}U\subset\bigcup_{V\in\mathcal{V}}V.$$

Indeed, given $x \in \bigcup_{U \in \mathcal{U}} U$, there exists some $U \in \mathcal{U}$ such that $x \in U$, but since $\mathcal{U} \subset \mathcal{V}$, we have $U \in \mathcal{V}$ as well, and thus $x \in \bigcup_{V \in \mathcal{V}} V$, which gives our desired inclusion.

Item 2, Associativity: We have

$$\bigcup_{U \in \bigcup_{A \in \mathcal{A}} A} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \bigcup_{A \in \mathcal{A}} A \\ \text{such that we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } A \in \mathcal{A} \\ \text{and some } U \in A \text{ such that} \\ \text{we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } A \in \mathcal{A} \\ \text{such that we have } x \in \bigcup_{U \in A} U \right\} \\ \stackrel{\text{def}}{=} \bigcup_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U \right). \end{array} \right.$$

This finishes the proof.

Item 3, Left Unitality: We have

$$\bigcup_{V \in \{U\}} V \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } V \in \{U\} \\ \text{such that we have } x \in U \end{array} \right\}$$
$$= \left\{ x \in X \mid x \in U \right\}$$
$$= U.$$

This finishes the proof.

Item 4, Right Unitality: We have

$$\bigcup_{\{u\} \in \chi_X(U)} \{u\} \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } \{u\} \in \chi_X(U) \right\}$$
 such that we have $x \in \{u\}$

$$= \left\{ x \in X \middle| \text{ there exists some } \{u\} \in \chi_X(U) \right\}$$
 such that we have $x = u$

$$= \left\{ x \in X \middle| \text{ there exists some } u \in U \right\}$$
 such that we have $x = u$

$$= \left\{ x \in X \middle| \text{ x } \in U \right\}$$

$$= U.$$

This finishes the proof.

Item 5, Interaction With Unions I: We have

$$\bigcup_{W \in \mathcal{U} \cup \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } W \in \mathcal{U} \cup \mathcal{V} \\ \text{such that we have } x \in W \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } W \in \mathcal{U} \text{ or some} \\ W \in \mathcal{V} \text{ such that we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } W \in \mathcal{U} \\ \text{such that we have } x \in W \end{array} \right\}$$

$$\cup \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } W \in \mathcal{V} \\ \text{such that we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left(\bigcup_{W \in \mathcal{U}} W \right) \cup \left(\bigcup_{W \in \mathcal{V}} W \right)$$

$$= \left(\bigcup_{U \in \mathcal{U}} U \right) \cup \left(\bigcup_{V \in \mathcal{V}} V \right).$$

This finishes the proof.

Item 6, Interaction With Unions II: Omitted.

Item 7, Interaction With Intersections I: We have

$$\bigcup_{W \in \mathcal{U} \cap \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } W \in \mathcal{U} \cap \mathcal{V} \\ \text{such that we have } x \in W \end{array} \right\}$$

$$\subset \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \mathcal{U} \text{ and some } V \in \mathcal{V} \\ \text{such that we have } x \in U \text{ and } x \in V \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \mathcal{U} \\ \text{such that we have } x \in U \end{array} \right\}$$

$$\cup \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } V \in \mathcal{V} \\ \text{such that we have } x \in V \right\}$$

$$\stackrel{\text{def}}{=} \left(\bigcup_{U \in \mathcal{U}} U \right) \cap \left(\bigcup_{V \in \mathcal{V}} V \right).$$

This finishes the proof.

Item 8, Interaction With Intersections II: Omitted.

Item 9, Interaction With Differences: Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}$, and let $\mathcal{V} = \{\{0\}\}$. We have

$$\bigcup_{W \in \mathcal{U} \setminus \mathcal{V}} U = \bigcup_{W \in \{\{0,1\}\}} W$$
$$= \{0,1\},$$

whereas

$$\left(\bigcup_{U \in \mathcal{U}} U\right) \setminus \left(\bigcup_{V \in \mathcal{V}} V\right) = \{0, 1\} \setminus \{0\}$$
$$= \{1\}.$$

Thus we have

$$\bigcup_{W\in\mathcal{U}\backslash\mathcal{V}}W=\left\{0,1\right\}\neq\left\{1\right\}=\left(\bigcup_{U\in\mathcal{U}}U\right)\backslash\left(\bigcup_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

Item 10, *Interaction With Complements I*: Let $X = \{0, 1\}$ and let $\mathcal{U} = \{0\}$. We have

$$\bigcup_{U \in \mathcal{U}^{c}} U = \bigcup_{U \in \{\emptyset, \{1\}, \{0,1\}\}} U$$
$$= \{0, 1\},$$

whereas

$$\bigcup_{U \in \mathcal{U}} U^{c} = \{0\}^{c}$$
$$= \{1\}.$$

Thus we have

$$\bigcup_{U\in\mathcal{U}^{\mathsf{c}}}U=\{0,1\}\neq\{1\}=\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}.$$

This finishes the proof.

Item 11, Interaction With Complements II: Omitted.

Item 12, Interaction With Complements III: Omitted.

Item 13, Interaction With Symmetric Differences: Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}$, and let $\mathcal{V} = \{\{0\}, \{0, 1\}\}$. We have

$$\bigcup_{W \in \mathcal{U} \triangle \mathcal{V}} W = \bigcup_{W \in \{\{0\}\}} W$$
$$= \{0\},$$

whereas

$$\left(\bigcup_{U\in\mathcal{U}}U\right)\triangle\left(\bigcup_{V\in\mathcal{V}}V\right)=\{0,1\}\triangle\{0,1\}$$
$$=\emptyset,$$

Thus we have

$$\bigcup_{W \in \mathcal{U} \triangle \mathcal{V}} W = \{0\} \neq \emptyset = \left(\bigcup_{U \in \mathcal{U}} U\right) \triangle \left(\bigcup_{V \in \mathcal{V}} V\right).$$

This finishes the proof.

Item 14, Interaction With Internal Homs I: This is a repetition of Item 7 of Definition 4.4.7.1.3 and is proved there.

Item 15, Interaction With Internal Homs II: This is a repetition of Item 8 of Definition 4.4.7.1.3 and is proved there.

Item 16, Interaction With Internal Homs III: This is a repetition of Item 9 of Definition 4.4.7.1.3 and is proved there.

Item 17, Interaction With Direct Images: This is a repetition of Item 3 of Definition 4.6.1.1.5 and is proved there.

Item 18, Interaction With Inverse Images: This is a repetition of Item 3 of Definition 4.6.2.1.3 and is proved there.

Item 19, Interaction With Codirect Images: This is a repetition of Item 3 of Definition 4.6.3.1.7 and is proved there.

Item 20, Interaction With Intersections of Families I: We have

$$\bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \bigcup_{A \in \mathcal{A}} A, \\ \text{we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } A \in \mathcal{A} \text{ and each} \\ U \in A, \text{ we have } x \in U \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right).$$

This finishes the proof.

Item 21, Interaction With Intersections of Families II: Omitted.

4.3.7 Intersections of Families of Subsets

Let X be a set and let $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

Definition 4.3.7.1.1. The intersection of $\mathcal U$ is the set $\bigcap_{U\in\mathcal U} U$ defined by

$$\bigcap_{U \in \mathcal{U}} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \mathcal{U}, \\ \text{we have } x \in U \end{array} \right\}.$$

Proposition 4.3.7.1.2. Let X be a set.

1. Functoriality. The assignment $\mathcal{U}\mapsto \bigcap_{U\in\mathcal{U}} U$ defines a functor

$$\bigcap : (\mathcal{P}(\mathcal{P}(X)), \supset) \to (\mathcal{P}(X), \subset).$$

In particular, for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$, the following condition is satisfied:

$$(\star)$$
 If $\mathcal{U} \subset \mathcal{V}$, then $\bigcap_{V \in \mathcal{V}} V \subset \bigcap_{U \in \mathcal{U}} U$.

2. Oplax Associativity. We have a natural transformation

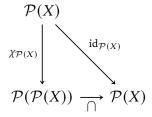
$$\begin{array}{c|c}
\mathcal{P}(\mathcal{P}(\mathcal{P}(X))) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \bigcap} & \mathcal{P}(\mathcal{P}(X)) \\
\cap \star \mathrm{id}_{\mathcal{P}(X)} & & & & & & & & & \\
\mathcal{P}(\mathcal{P}(X)) & & & & & & & & & & \\
\end{array}$$

with components

$$\bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right) \subset \bigcap_{U \in \bigcap_{A \in \mathcal{A}} A} U$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$.

3. Left Unitality. The diagram

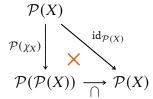


commutes, i.e. we have

$$\bigcap_{V\in\{U\}}V=U.$$

for each $U \in \mathcal{P}(X)$.

4. Oplax Right Unitality. The diagram



does not commute in general, i.e. we may have

$$\bigcap_{\{x\}\in\chi_X(U)}\{x\}\neq U$$

in general, where $U \in \mathcal{P}(X)$. However, when U is nonempty, we have

$$\bigcap_{\{x\}\in\chi_X(U)}\{x\}\subset U.$$

5. Interaction With Unions I. The diagram

commutes, i.e. we have

$$\bigcap_{W\in\mathcal{U}\cup\mathcal{V}}W=\left(\bigcap_{U\in\mathcal{U}}U\right)\cap\left(\bigcap_{V\in\mathcal{V}}V\right)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

6. Interaction With Unions II. The diagram

commute, i.e. we have

$$U \cup \left(\bigcap_{V \in \mathcal{V}} V\right) = \bigcap_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cup V = \bigcap_{U \in \mathcal{U}} (U \cup V)$$

for each $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

7. Interaction With Intersections I. We have a natural transformation

with components

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\cap\left(\bigcap_{V\in\mathcal{V}}V\right)\subset\bigcap_{W\in\mathcal{U}\cap\mathcal{V}}W$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

8. *Interaction With Intersections II*. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcap_{V \in \mathcal{V}} V\right) = \bigcap_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cup V = \bigcap_{U \in \mathcal{U}} (U \cup V)$$

for each $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

9. Interaction With Differences. The diagram

does not commute in general, i.e. we may have

$$\bigcap_{W \in \mathcal{U} \setminus \mathcal{V}} W \neq \left(\bigcap_{U \in \mathcal{U}} U\right) \setminus \left(\bigcap_{V \in \mathcal{V}} V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

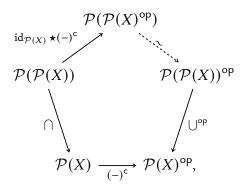
10. Interaction With Complements I. The diagram

does not commute in general, i.e. we may have

$$\bigcap_{W \in \mathcal{U}^{\mathsf{c}}} W \neq \bigcap_{U \in \mathcal{U}} U^{\mathsf{c}}$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

11. Interaction With Complements II. The diagram

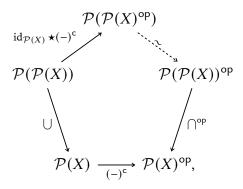


commutes, i.e. we have

$$\left(\bigcap_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

12. Interaction With Complements III. The diagram



commutes, i.e. we have

$$\left(\bigcup_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

13. Interaction With Symmetric Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\Delta} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \times \cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\wedge} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W \in \mathcal{U} \triangle \mathcal{V}} W \neq \left(\bigcap_{U \in \mathcal{U}} U\right) \triangle \left(\bigcap_{V \in \mathcal{V}} V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

14. Interaction With Internal Homs I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\operatorname{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\uparrow^{\operatorname{op}} \times \cap^{\operatorname{op}} \qquad \qquad \downarrow \cap$$

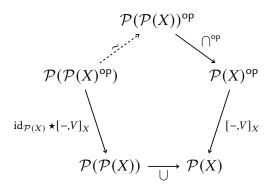
$$\mathcal{P}(X)^{\operatorname{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcap_{U \in \mathcal{U}} U, \bigcap_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

15. Interaction With Internal Homs II. The diagram



commutes, i.e. we have

$$\left[\bigcap_{U\in\mathcal{U}}U,V\right]_X=\bigcup_{V\in\mathcal{U}}[U,V]_X$$

for each $U \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

16. Interaction With Internal Homs III. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{\bigcap} \mathcal{P}(X)$$

$$id_{\mathcal{P}(X)} \star [U,-]_X \qquad \qquad \bigcup_{[U,-]_X} [U,-]_X$$

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{\bigcap} \mathcal{P}(X)$$

commutes, i.e. we have

$$\left[U, \bigcap_{V \in \mathcal{V}} V\right]_X = \bigcap_{V \in \mathcal{V}} [U, V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

17. Interaction With Direct Images. Let $f: X \to Y$ be a map of sets. The diagram

commutes, i.e. we have

$$\bigcap_{U \in \mathcal{U}} f_!(U) = \bigcap_{V \in f_!(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

18. Interaction With Inverse Images. Let $f: X \to Y$ be a map of sets. The

diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\uparrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcap_{V \in \mathcal{V}} f^{-1}(V) = \bigcap_{U \in f^{-1}(\mathcal{U})} U$$

for each $\mathcal{V} \in \mathcal{P}(Y)$, where $f^{-1}(\mathcal{V}) \stackrel{\text{def}}{=} (f^{-1})^{-1}(\mathcal{V})$.

19. Interaction With Codirect Images. Let $f: X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\cap \bigcup_{f_*} \qquad \bigcup_{f_*} \cap$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcap_{U \in \mathcal{U}} f_*(U) = \bigcap_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

20. Interaction With Unions of Families I. The diagram

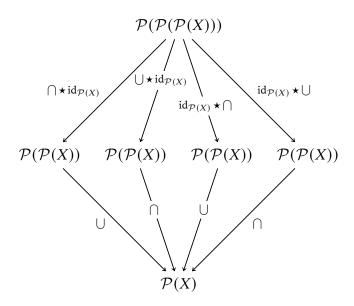
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\operatorname{id}_{\mathcal{P}(X)} \star \bigcap} & \mathcal{P}(\mathcal{P}(x)) \\
\downarrow \star \operatorname{id}_{\mathcal{P}(X)} & & & \downarrow \bigcap \\
\mathcal{P}(X) & \xrightarrow{\bigcap} & X
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U = \bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right)$$

for each $A \in \mathcal{P}(\mathcal{P}(X))$.

21. Interaction With Unions of Families II. Let X be a set and consider the compositions

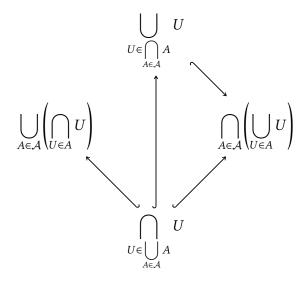


given by

$$A \mapsto \bigcup_{U \in \bigcap_{A \in \mathcal{A}} A} U, \quad A \mapsto \bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U,$$

$$A \mapsto \bigcup_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U\right), \quad A \mapsto \bigcap_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U\right)$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$. We have the following inclusions:



All other possible inclusions fail to hold in general.

Proof. Item 1, Functoriality: Since $\mathcal{P}(X)$ is posetal, it suffices to prove the condition (\star) . So let $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ with $\mathcal{U} \subset \mathcal{V}$. We claim that

$$\bigcap_{V\in\mathcal{V}}V\subset\bigcap_{U\in\mathcal{U}}U.$$

Indeed, if $x \in \bigcap_{V \in \mathcal{V}} V$, then $x \in V$ for all $V \in \mathcal{V}$. But since $\mathcal{U} \subset \mathcal{V}$, it follows that $x \in U$ for all $U \in \mathcal{U}$ as well. Thus $x \in \bigcap_{U \in \mathcal{U}} U$, which gives our desired inclusion.

Item 2, Oplax Associativity: We have

$$\bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right) \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{c} \text{for each } A \in \mathcal{A}, \\ \text{we have } x \in \bigcap_{U \in A} U \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{c} \text{for each } A \in \mathcal{A} \text{ and each} \\ U \in A, \text{ we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{c} \text{for each } U \in \bigcup_{A \in \mathcal{A}} A, \\ \text{we have } x \in U \end{array} \right\}$$
we have $x \in U$

$$\subset \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \bigcap_{A \in \mathcal{A}} A, \\ \text{we have } x \in U \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{A \in \mathcal{A}} U.$$

$$U \in \bigcap_{A \in \mathcal{A}} A$$

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 3, Left Unitality: We have

$$\bigcap_{V \in \{U\}} V \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } V \in \{U\}, \\ \text{we have } x \in U \end{array} \right\}$$
$$= \left\{ x \in X \middle| x \in U \right\}$$
$$= U.$$

This finishes the proof.

Item 4, Oplax Right Unitality: If $U = \emptyset$, then we have

$$\bigcap_{\{u\}\in\chi_X(U)} \{u\} = \bigcap_{\{u\}\in\emptyset} \{u\}$$
$$= X$$

so $\bigcap_{\{u\}\in\chi_X(U)}\{u\}=X\neq\emptyset=U.$ When U is nonempty, we have two cases:

1. If U is a singleton, say $U = \{u\}$, we have

$$\bigcap_{\{u\}\in\chi_X(U)}\{u\}=\{u\}$$

$$\stackrel{\text{def}}{=}U.$$

2. If U contains at least two elements, we have

$$\bigcap_{\{u\}\in\chi_X(U)}\{u\}=\emptyset$$

$$\subset U.$$

This finishes the proof.

Item 5, Interaction With Unions I: We have

$$\bigcap_{W \in \mathcal{U} \cup \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cup \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \text{ and each} \\ W \in \mathcal{V}, \text{ we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\cap \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left(\bigcap_{W \in \mathcal{U}} W \right) \cap \left(\bigcap_{W \in \mathcal{V}} W \right)$$

$$= \left(\bigcap_{U \in \mathcal{U}} U \right) \cap \left(\bigcap_{V \in \mathcal{V}} V \right).$$

This finishes the proof.

Item 6, Interaction With Unions II: Omitted.

Item 7, Interaction With Intersections I: We have

$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cap \left(\bigcap_{V \in \mathcal{V}} V\right) \stackrel{\text{def}}{=} \left\{x \in X \middle| \begin{array}{l} \text{for each } U \in \mathcal{U}, \\ \text{we have } x \in U \end{array}\right\}$$

$$\cup \left\{x \in X \middle| \begin{array}{l} \text{for each } V \in \mathcal{V}, \\ \text{we have } x \in V \end{array}\right\}$$

$$= \left\{x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cap \mathcal{V}, \\ \text{we have } x \in W \end{array}\right\}$$

$$\subset \left\{x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cup \mathcal{V}, \\ \text{we have } x \in W \end{array}\right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{W \in \mathcal{U} \cap \mathcal{V}} W.$$

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 8, Interaction With Intersections II: Omitted.

Item 9, Interaction With Differences: Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0\}, \{0, 1\}\}$, and let $\mathcal{V} = \{\{0\}\}$. We have

$$\bigcap_{W \in \mathcal{U} \setminus \mathcal{V}} U = \bigcap_{W \in \{\{0,1\}\}} W$$
$$= \{0,1\},$$

whereas

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\setminus\left(\bigcap_{V\in\mathcal{V}}V\right)=\{0\}\setminus\{0\}$$
$$=\emptyset.$$

Thus we have

$$\bigcap_{W\in\mathcal{U}\setminus\mathcal{V}}W=\left\{0,1\right\}\neq\emptyset=\left(\bigcap_{U\in\mathcal{U}}U\right)\setminus\left(\bigcap_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

Item 10, *Interaction With Complements I*: Let $X = \{0, 1\}$ and let $\mathcal{U} = \{\{0\}\}$. We have

$$\bigcap_{W \in \mathcal{U}^{c}} U = \bigcap_{W \in \{\emptyset, \{1\}, \{0,1\}\}} W$$
$$= \emptyset,$$

whereas

$$\bigcap_{U \in \mathcal{U}} U^{c} = \{0\}^{c}$$
$$= \{1\}.$$

Thus we have

$$\bigcap_{W\in\mathcal{U}^{\mathsf{c}}}U=\emptyset\neq\{1\}=\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}.$$

This finishes the proof.

Item 11, Interaction With Complements II: This is a repetition of Item 12 of Definition 4.3.6.1.2 and is proved there.

Item 12, Interaction With Complements III: This is a repetition of Item 11 of Definition 4.3.6.1.2 and is proved there.

Item 13, Interaction With Symmetric Differences: Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}$, and let $\mathcal{V} = \{\{0\}, \{0, 1\}\}$. We have

$$\bigcap_{W \in \mathcal{U} \triangle \mathcal{V}} W = \bigcap_{W \in \{\{0\}\}} W$$
$$= \{0\},$$

whereas

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\triangle\left(\bigcap_{V\in\mathcal{V}}V\right)=\left\{0,1\right\}\triangle\left\{0\right\}$$
$$=\emptyset.$$

Thus we have

$$\bigcap_{W \in \mathcal{U} \wedge \mathcal{V}} W = \{0\} \neq \emptyset = \left(\bigcap_{U \in \mathcal{U}} U\right) \triangle \left(\bigcap_{V \in \mathcal{V}} V\right).$$

This finishes the proof.

Item 14, Interaction With Internal Homs I: This is a repetition of Item 10 of Definition 4.4.7.1.3 and is proved there.

Item 15, Interaction With Internal Homs II: This is a repetition of Item 11 of Definition 4.4.7.1.3 and is proved there.

Item 16, Interaction With Internal Homs III: This is a repetition of Item 12 of Definition 4.4.7.1.3 and is proved there.

Item 17, Interaction With Direct Images: This is a repetition of Item 4 of Definition 4.6.1.1.5 and is proved there.

Item 18, Interaction With Inverse Images: This is a repetition of Item 4 of Definition 4.6.2.1.3 and is proved there.

Item 19, Interaction With Codirect Images: This is a repetition of Item 4 of Definition 4.6.3.1.7 and is proved there.

Item 20, *Interaction With Unions of Families I*: This is a repetition of Item 20 of Definition 4.3.6.1.2 and is proved there.

Item 21, Interaction With Unions of Families II: This is a repetition of Item 21 of Definition 4.3.6.1.2 and is proved there. □

4.3.8 Binary Unions

Let X be a set and let $U, V \in \mathcal{P}(X)$.

Definition 4.3.8.1.1. The **union of** U **and** V is the set $U \cup V$ defined by

$$U \cup V \stackrel{\text{def}}{=} \bigcup_{z \in \{U, V\}} z$$

$$\stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in V\}.$$

Proposition 4.3.8.1.2. Let *X* be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cup V$ define functors

$$U \cup -: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$- \cup V: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$-_1 \cup -_2: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset).$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $U \cup V \subset A \cup V$.
- (b) If $V \subset B$, then $U \cup V \subset U \cup B$.
- (c) If $U \subset A$ and $V \subset B$, then $U \cup V \subset A \cup B$.
- 2. Associativity. The diagram

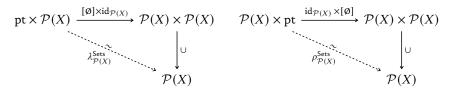
$$\begin{array}{c} \mathcal{P}(X)\times(\mathcal{P}(X)\times\mathcal{P}(X)) \\ \alpha^{\mathsf{Sets}}_{\mathcal{P}(X),\mathcal{P}(X),\mathcal{P}(X)} & \mathrm{id}_{\mathcal{P}(X)}\times\cup \\ (\mathcal{P}(X)\times\mathcal{P}(X))\times\mathcal{P}(X) & \mathcal{P}(X)\times\mathcal{P}(X) \\ & & & & & & & & \\ & & & & & & & \\ \mathcal{P}(X)\times\mathcal{P}(X) & \xrightarrow{} \mathcal{P}(X), \end{array}$$

commutes, i.e. we have an equality of sets

$$(U \cup V) \cup W = U \cup (V \cup W)$$

for each $U, V, W \in \mathcal{P}(X)$.

3. *Unitality*. The diagrams

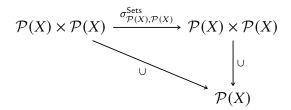


commute, i.e. we have equalities of sets

$$\emptyset \cup U = U$$
, $U \cup \emptyset = U$

for each $U \in \mathcal{P}(X)$.

4. Commutativity. The diagram

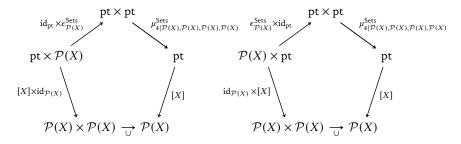


commutes, i.e. we have an equality of sets

$$U \cup V = V \cup U$$

for each $U, V \in \mathcal{P}(X)$.

5. Annihilation With X. The diagrams



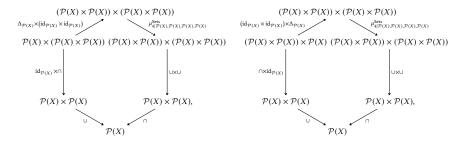
commute, i.e. we have equalities of sets

$$U \cup X = X$$
,

$$X \cup V = X$$

for each $U, V \in \mathcal{P}(X)$.

6. Distributivity of Unions Over Intersections. The diagrams



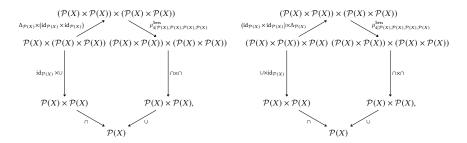
commute, i.e. we have equalities of sets

$$U \cup (V \cap W) = (U \cup V) \cap (U \cup W),$$

$$(U \cap V) \cup W = (U \cup W) \cap (V \cup W)$$

for each $U, V, W \in \mathcal{P}(X)$.

7. Distributivity of Intersections Over Unions. The diagrams



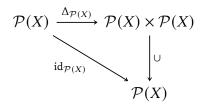
commute, i.e. we have equalities of sets

$$U \cap (V \cup W) = (U \cap V) \cup (U \cap W),$$

$$(U \cup V) \cap W = (U \cap W) \cup (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

8. Idempotency. The diagram



commutes, i.e. we have an equality of sets

$$U \cup U = U$$

for each $U \in \mathcal{P}(X)$.

9. Via Intersections and Symmetric Differences. The diagram

$$(\mathcal{P}(X) \times \mathcal{P}(X)) \times (\mathcal{P}(X) \times \mathcal{P}(X)) \xrightarrow{\triangle \times \cap} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\stackrel{\Delta_{\mathcal{P}(X) \times \mathcal{P}(X)}}{\longrightarrow} \mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\Delta} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cup V = (U \triangle V) \triangle (U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

10. Interaction With Characteristic Functions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

11. Interaction With Characteristic Functions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each $U, V \in \mathcal{P}(X)$.

12. Interaction With Direct Images. Let $f: X \to Y$ be a function. The diagram

$$P(X) \times P(X) \xrightarrow{f_1 \times f_1} P(Y) \times P(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \cup$$

$$P(X) \xrightarrow{f_1} P(Y)$$

commutes, i.e. we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$

for each $U, V \in \mathcal{P}(X)$.

13. Interaction With Inverse Images. Let $f\colon X\to Y$ be a function. The diagram

$$\begin{array}{cccc} \mathcal{P}(Y) \times \mathcal{P}(Y) & \xrightarrow{f^{-1} \times f^{-1}} & \mathcal{P}(X) \times \mathcal{P}(X) \\ & & & \downarrow \cup \\ & \mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X) \end{array}$$

commutes, i.e. we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

14. *Interaction With Codirect Images.* Let $f: X \to Y$ be a function. We have a natural transformation

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_* \times f_*} \mathcal{P}(Y) \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

with components

$$f_*(U) \cup f_*(V) \subset f_*(U \cup V)$$

indexed by $U, V \in \mathcal{P}(X)$.

15. *Interaction With Powersets and Semirings*. The quintuple $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$ is an idempotent commutative semiring.

Proof. Item 1, Functoriality: See [Pro25an].

Item 2, Associativity: See [Pro25ba].

Item 3, Unitality: This follows from [Pro25bd] and Item 4.

Item 4, Commutativity: See [Pro25bb].

Item 5, Annihilation With X: We have

$$U \cup X \stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in X\}$$
$$= \{x \in X \mid x \in X\},$$
$$= X$$

and

$$X \cup V \stackrel{\text{def}}{=} \{x \in X \mid x \in X \text{ or } x \in V\}$$
$$= \{x \in X \mid x \in X\}$$
$$= X.$$

This finishes the proof.

Item 6, Distributivity of Unions Over Intersections: See [Pro25az].

Item 7, *Distributivity of Intersections Over Unions*: See [Pro25aj].

Item 8, Idempotency: See [Pro25am].

Item 9, Via Intersections and Symmetric Differences: See [Pro25ay].

Item 10, *Interaction With Characteristic Functions I*: See [Pro25h].

Item 11, *Interaction With Characteristic Functions II*: See [Pro25h].

Item 12, *Interaction With Direct Images*: See [Pro25p].

Item 13, *Interaction With Inverse Images*: See [Pro25y].

Item 14, Interaction With Codirect Images: This is a repetition of Item 5 of Definition 4.6.3.1.7 and is proved there.

Item 15, Interaction With Powersets and Semirings: This follows from Items 2 to 4 and 8 of this proposition and Items 3 to 6 and 8 of Definition 4.3.9.1.2. □

4.3.9 Binary Intersections

Let X be a set and let $U, V \in \mathcal{P}(X)$.

Definition 4.3.9.1.1. The intersection of U and V is the set $U \cap V$ defined by

$$U \cap V \stackrel{\text{def}}{=} \bigcap_{z \in \{U, V\}} z$$

$$\stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in V\}.$$

Proposition 4.3.9.1.2. Let X be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cap V$ define functors

$$U \cap -: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$- \cap V: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$-_1 \cap -_2: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset).$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $U \cap V \subset A \cap V$.
- (b) If $V \subset B$, then $U \cap V \subset U \cap B$.
- (c) If $U \subset A$ and $V \subset B$, then $U \cap V \subset A \cap B$.
- 2. Adjointness. We have adjunctions

$$(U \cap - \dashv [U, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{U \cap -}{\bot}} \mathcal{P}(X),$$

$$(- \cap V \dashv [V, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{- \cap V}{\bot}} \mathcal{P}(X),$$

witnessed by bijections

$$\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, [V, W]_X),$$

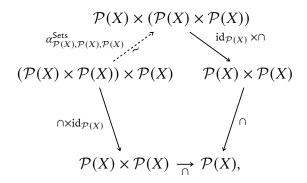
 $\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(V, [U, W]_X),$

natural in $U, V, W \in \mathcal{P}(X)$, where

$$[-1,-2]_X \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

is the bifunctor of Section 4.4.7. In particular, the following statements hold for each $U, V, W \in \mathcal{P}(X)$:

- (a) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $U \subset [V, W]_X$.
- (b) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $V \subset [U, W]_X$.
- 3. Associativity. The diagram

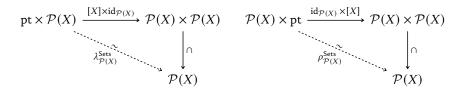


commutes, i.e. we have an equality of sets

$$(U \cap V) \cap W = U \cap (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

4. Unitality. The diagrams



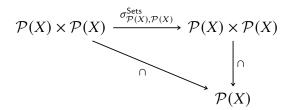
commute, i.e. we have equalities of sets

$$X \cap U = U,$$

$$U \cap X = U$$

for each $U \in \mathcal{P}(X)$.

5. Commutativity. The diagram

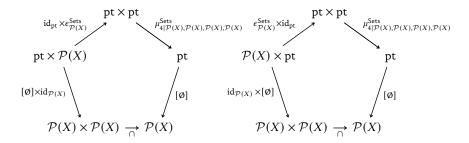


commutes, i.e. we have an equality of sets

$$U\cap V=V\cap U$$

for each $U, V \in \mathcal{P}(X)$.

6. Annihilation With the Empty Set. The diagrams

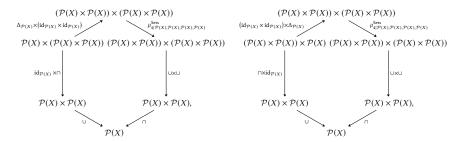


commute, i.e. we have equalities of sets

$$\emptyset \cap X = \emptyset$$
, $X \cap \emptyset = \emptyset$

for each $U \in \mathcal{P}(X)$.

7. Distributivity of Unions Over Intersections. The diagrams



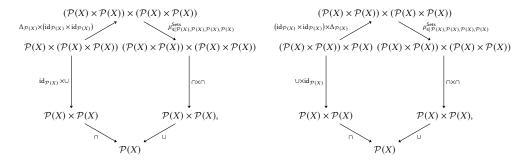
commute, i.e. we have equalities of sets

$$U \cup (V \cap W) = (U \cup V) \cap (U \cup W),$$

$$(U \cap V) \cup W = (U \cup W) \cap (V \cup W)$$

for each $U, V, W \in \mathcal{P}(X)$.

8. Distributivity of Intersections Over Unions. The diagrams



commute, i.e. we have equalities of sets

$$U \cap (V \cup W) = (U \cap V) \cup (U \cap W),$$

$$(U \cup V) \cap W = (U \cap W) \cup (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

9. Idempotency. The diagram

$$\mathcal{P}(X) \xrightarrow{\Delta_{\mathcal{P}(X)}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \cap$$

$$\mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cap U = U$$

for each $U \in \mathcal{P}(X)$.

10. Interaction With Characteristic Functions I. We have

$$\chi_{U\cap V} = \chi_U \chi_V$$

for each $U, V \in \mathcal{P}(X)$.

11. Interaction With Characteristic Functions II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

12. Interaction With Direct Images. Let $f: X \to Y$ be a function. We have a natural transformation

with components

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

13. Interaction With Inverse Images. Let $f\colon X\to Y$ be a function. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\uparrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U\cap V)=f^{-1}(U)\cap f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

14. Interaction With Codirect Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_* \times f_*} \mathcal{P}(Y) \times \mathcal{P}(Y)$$

$$\uparrow \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U) \cap f_*(V) = f_*(U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

- 15. *Interaction With Powersets and Monoids With Zero.* The quadruple $((\mathcal{P}(X), \emptyset), \cap, X)$ is a commutative monoid with zero.
- 16. *Interaction With Powersets and Semirings*. The quintuple $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$ is an idempotent commutative semiring.

Proof. Item 1, Functoriality: See [Pro25al].

Item 2, Adjointness: See [MSE 267469].

Item 3, *Associativity*: See [Pro25r].

Item 4, Unitality: This follows from [Pro25v] and Item 5.

Item 5, Commutativity: See [Pro25s].

Item 6, Annihilation With the Empty Set: This follows from [Pro25t] and Item 5.

Item 7, *Distributivity of Unions Over Intersections*: See [Pro25az].

Item 8, Distributivity of Intersections Over Unions: See [Pro25ai].

Item 9, Idempotency: See [Pro25ak].

Item 10, *Interaction With Characteristic Functions I*: See [Pro25e].

Item 11, *Interaction With Characteristic Functions II*: See [Pro25e].

Item 12, *Interaction With Direct Images*: See [Pro25n].

Item 13, *Interaction With Inverse Images*: See [Pro25w].

Item 14, Interaction With Codirect Images: This is a repetition of Item 6 of Definition 4.6.3.1.7 and is proved there.

Item 15, Interaction With Powersets and Monoids With Zero: This follows from Items 3 to 6.

Item 16, Interaction With Powersets and Semirings: This follows from Items 2 to 4 and 8 and Items 3 to 6 and 8 of Definition 4.3.9.1.2. □

4.3.10 Differences

Let *X* and *Y* be sets.

Definition 4.3.10.1.1. The **difference of** X **and** Y is the set $X \setminus Y$ defined by

$$X \setminus Y \stackrel{\text{def}}{=} \{ a \in X \mid a \notin Y \}.$$

Proposition 4.3.10.1.2. Let *X* be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cap V$ define functors

$$U \setminus -: \qquad (\mathcal{P}(X), \supset) \qquad \to (\mathcal{P}(X), \subset),$$

$$- \setminus V: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$-_1 \setminus -_2: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \supset) \to (\mathcal{P}(X), \subset).$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $U \setminus V \subset A \setminus V$.
- (b) If $V \subset B$, then $U \setminus B \subset U \setminus V$.
- (c) If $U \subset A$ and $V \subset B$, then $U \setminus B \subset A \setminus V$.
- 2. De Morgan's Laws. We have equalities of sets

$$X \setminus (U \cup V) = (X \setminus U) \cap (X \setminus V),$$

$$X \setminus (U \cap V) = (X \setminus U) \cup (X \setminus V)$$

for each $U, V \in \mathcal{P}(X)$.

3. Interaction With Unions I. We have equalities of sets

$$U \setminus (V \cup W) = (U \setminus V) \cap (U \setminus W)$$

for each $U, V, W \in \mathcal{P}(X)$.

4. Interaction With Unions II. We have equalities of sets

$$(U \setminus V) \cup W = (U \cup W) \setminus (V \setminus W)$$

for each $U, V, W \in \mathcal{P}(X)$.

5. Interaction With Unions III. We have equalities of sets

$$U \setminus (V \cup W) = (U \cup W) \setminus (V \cup W)$$
$$= (U \setminus V) \setminus W$$
$$= (U \setminus W) \setminus V$$

for each $U, V, W \in \mathcal{P}(X)$.

6. Interaction With Unions IV. We have equalities of sets

$$(U \cup V) \setminus W = (U \setminus W) \cup (V \setminus W)$$

for each $U, V, W \in \mathcal{P}(X)$.

7. Interaction With Intersections. We have equalities of sets

$$(U \setminus V) \cap W = (U \cap W) \setminus V$$
$$= U \cap (W \setminus V)$$

for each $U, V, W \in \mathcal{P}(X)$.

8. Interaction With Complements. We have an equality of sets

$$U \setminus V = U \cap V^{\mathsf{c}}$$

for each $U, V \in \mathcal{P}(X)$.

9. Interaction With Symmetric Differences. We have an equality of sets

$$U \setminus V = U \triangle (U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

10. Triple Differences. We have

$$U \setminus (V \setminus W) = (U \cap W) \cup (U \setminus V)$$

for each $U, V, W \in \mathcal{P}(X)$.

11. Left Annihilation. We have

$$\emptyset \setminus U = \emptyset$$

for each $U \in \mathcal{P}(X)$.

12. Right Unitality. We have

$$U \setminus \emptyset = U$$

for each $U \in \mathcal{P}(X)$.

13. Right Annihilation. We have

$$U \setminus X = U$$

for each $U \in \mathcal{P}(X)$.

14. Invertibility. We have

$$U \setminus U = \emptyset$$

for each $U \in \mathcal{P}(X)$.

- 15. Interaction With Containment. The following conditions are equivalent:
 - (a) We have $V \setminus U \subset W$.
 - (b) We have $V \setminus W \subset U$.
- 16. Interaction With Characteristic Functions. We have

$$\chi_{U\setminus V} = \chi_U - \chi_{U\cap V}$$

for each $U, V \in \mathcal{P}(X)$.

17. Interaction With Direct Images. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad$$

with components

$$f_!(U)\setminus f_!(V)\subset f_!(U\setminus V)$$

indexed by $U, V \in \mathcal{P}(X)$.

18. Interaction With Inverse Images. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\mathsf{op},-1} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad$$

commutes, i.e. we have

$$f^{-1}(U \setminus V) = f^{-1}(U) \setminus f^{-1}(V)$$

for each $U, V \in \mathcal{P}(X)$.

19. Interaction With Codirect Images. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad$$

with components

$$f_i(U) \setminus f_i(V) \subset f_i(U \setminus V)$$

indexed by $U, V \in \mathcal{P}(X)$.

Proof. Item 1, Functoriality: See [Pro25ad] and [Pro25ah].

Item 2, De Morgan's Laws: See [Pro25k].

Item 3, *Interaction With Unions I*: See [Pro25l].

Item 4, Interaction With Unions II: Omitted.

Item 5, Interaction With Unions III: See [Pro25ai].

Item 6, Interaction With Unions IV: See [Pro25ac].

Item 7, *Interaction With Intersections*: See [Pro25u].

Item 8, Interaction With Complements: See [Pro25aa].

Item 9, Interaction With Symmetric Differences: See [Pro25ab].

Item 10, Triple Differences: See [Pro25ag].

Item 11, Left Annihilation: Omitted.

Item 12, Right Unitality: See [Pro25ae].

Item 13, Right Annihilation: Omitted.

Item 14, Invertibility: See [Pro25af].

Item 15, Interaction With Containment: Omitted.

Item 16, *Interaction With Characteristic Functions*: See [Pro25f].

Item 17, *Interaction With Direct Images*: See [Pro250].

Item 18, *Interaction With Inverse Images*: See [Pro25x].

4.3.11 Complements

Let X be a set and let $U \in \mathcal{P}(X)$.

Definition 4.3.11.1.1. The **complement of** U is the set U^{c} defined by

$$U^{c} \stackrel{\text{def}}{=} X \setminus U$$

$$\stackrel{\text{def}}{=} \{ a \in X \mid a \notin U \}.$$

Proposition 4.3.11.1.2. Let *X* be a set.

1. Functoriality. The assignment $U\mapsto U^{\mathsf{c}}$ defines a functor

$$(-)^{c} : \mathcal{P}(X)^{op} \to \mathcal{P}(X).$$

In particular, the following statements hold for each $U, V \in \mathcal{P}(X)$:

$$(\star)$$
 If $U \subset V$, then $V^{c} \subset U^{c}$.

2. De Morgan's Laws. The diagrams

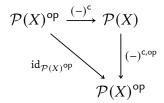
commute, i.e. we have equalities of sets

$$(U \cup V)^{c} = U^{c} \cap V^{c},$$

$$(U \cap V)^{c} = U^{c} \cup V^{c}$$

for each $U, V \in \mathcal{P}(X)$.

3. Involutority. The diagram



commutes, i.e. we have

$$(U^{\mathsf{c}})^{\mathsf{c}} = U$$

for each $U \in \mathcal{P}(X)$.

4. Interaction With Characteristic Functions. We have

$$\chi_{U^c} \equiv 1 - \chi_U \pmod{2}$$

for each $U \in \mathcal{P}(X)$.

5. Interaction With Direct Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{f_*^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}}$$

$$(-)^{\mathsf{c}} \qquad \qquad \downarrow (-)^{\mathsf{c}}$$

$$\mathcal{P}(X) \xrightarrow{f} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_!(U^{\mathsf{c}}) = f_*(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

6. Interaction With Inverse Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \xrightarrow{f^{-1,\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}}$$

$$\downarrow^{(-)^{\mathsf{c}}} \qquad \qquad \downarrow^{(-)^{\mathsf{c}}}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U^{\mathsf{c}}) = f^{-1}(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

7. Interaction With Codirect Images. Let $f\colon X\to Y$ be a function. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{f_!^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
\xrightarrow{(-)^{\mathsf{c}}} \qquad \qquad \downarrow^{(-)^{\mathsf{c}}} \\
\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U^{\mathsf{c}}) = f_!(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

Proof. Item 1, Functoriality: This follows from Item 1 of Definition 4.3.10.1.2.

Item 2, De Morgan's Laws: See [Pro25k].

Item 3, Involutority: See [Pro25i].

Item 4, Interaction With Characteristic Functions: Omitted.

Item 5, Interaction With Direct Images: This is a repetition of Item 8 of Defini-

tion 4.6.1.1.5 and is proved there.

Item 6, Interaction With Inverse Images: This is a repetition of Item 8 of Defini-

tion 4.6.2.1.3 and is proved there.

Item 7, Interaction With Codirect Images: This is a repetition of Item 7 of Definition 4.6.3.1.7 and is proved there.

4.3.12 Symmetric Differences

Let X be a set and let $U, V \in \mathcal{P}(X)$.

Definition 4.3.12.1.1. The **symmetric difference of** U **and** V is the set $U \triangle V$

defined by¹³

$$U \triangle V \stackrel{\text{def}}{=} (U \setminus V) \cup (V \setminus U).$$

Proposition 4.3.12.1.2. Let X be a set.

1. Lack of Functoriality. The assignment $(U, V) \mapsto U \triangle V$ does not in general define functors

$$U \triangle -: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$- \triangle V: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$-_{1} \triangle -_{2}: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset).$$

2. Via Unions and Intersections. We have

$$U \triangle V = (U \cup V) \setminus (U \cap V)$$

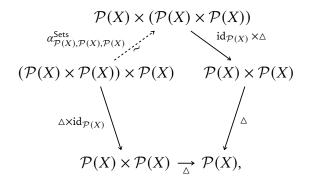
for each $U, V \in \mathcal{P}(X)$, as in the Venn diagram

$$\boxed{\bigcirc{U \triangle V}} = \boxed{\bigcirc{U \cup V}} \setminus \boxed{\bigcirc{U \cap V}}.$$

3. Symmetric Differences of Disjoint Sets. If U and V are disjoint, then we have

$$U \wedge V = U \cup V$$
.

4. Associativity. The diagram



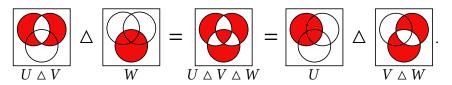
 $^{^{13}}$ Illustration:

$$\boxed{\bigcup_{U \, \triangle \, V}} = \boxed{\bigcup_{U \, \backslash \, V}} \cup \boxed{\bigcup_{V \, \backslash \, U}}$$

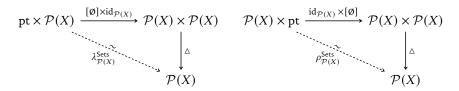
commutes, i.e. we have

$$(U \triangle V) \triangle W = U \triangle (V \triangle W)$$

for each $U, V, W \in \mathcal{P}(X)$, as in the Venn diagram



5. Unitality. The diagrams



commute, i.e. we have

$$U \triangle \emptyset = U,$$

 $\emptyset \triangle U = U$

for each $U \in \mathcal{P}(X)$.

6. Commutativity. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\sigma_{\mathcal{P}(X),\mathcal{P}(X)}^{\mathsf{Sets}}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow^{\triangle}$$

$$\mathcal{P}(X)$$

commutes, i.e. we have

$$U \triangle V = V \triangle U$$

for each $U, V \in \mathcal{P}(X)$.

7. Invertibility. We have

$$U \triangle U = \emptyset$$

for each $U \in \mathcal{P}(X)$.

8. Interaction With Unions. We have

$$(U \triangle V) \cup (V \triangle T) = (U \cup V \cup W) \setminus (U \cap V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

9. Interaction With Complements I. We have

$$U \triangle U^{c} = X$$

for each $U \in \mathcal{P}(X)$.

10. Interaction With Complements II. We have

$$U \triangle X = U^{\mathsf{c}},$$
$$X \triangle U = U^{\mathsf{c}}$$

$$X \triangle U = U$$

for each $U \in \mathcal{P}(X)$.

11. Interaction With Complements III. The diagram

commutes, i.e. we have

$$U^{\mathsf{c}} \triangle V^{\mathsf{c}} = U \triangle V$$

for each $U, V \in \mathcal{P}(X)$.

12. "Transitivity". We have

$$(U \triangle V) \triangle (V \triangle W) = U \triangle W$$

for each $U, V, W \in \mathcal{P}(X)$.

13. The Triangle Inequality for Symmetric Differences. We have

$$U \triangle W \subset U \triangle V \cup V \triangle W$$

for each $U, V, W \in \mathcal{P}(X)$.

14. Distributivity Over Intersections. We have

$$U \cap (V \triangle W) = (U \cap V) \triangle (U \cap W),$$

$$(U \triangle V) \cap W = (U \cap W) \triangle (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

15. Interaction With Characteristic Functions. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \pmod{2}$$

for each $U, V \in \mathcal{P}(X)$.

16. *Bijectivity*. Given $U, V \in \mathcal{P}(X)$, the maps

$$U \triangle -: \mathcal{P}(X) \to \mathcal{P}(X),$$

$$- \triangle V \colon \mathcal{P}(X) \to \mathcal{P}(X)$$

are bijections with inverses given by

$$(U \triangle -)^{-1} = - \cup (U \cap -),$$

$$(- \triangle V)^{-1} = - \cup (V \cap -).$$

Moreover, the map

$$\mathcal{P}(X) \longrightarrow \mathcal{P}(X)$$

$$C \longmapsto C \triangle (U \triangle V)$$

is a bijection of $\mathcal{P}(X)$ onto itself sending U to V and V to U.

- 17. Interaction With Powersets and Groups. Let X be a set.
 - (a) The quadruple $(\mathcal{P}(X), \triangle, \emptyset, \mathrm{id}_{\mathcal{P}(X)})$ is an abelian group.¹⁴
 - (b) Every element of $\mathcal{P}(X)$ has order 2 with respect to \triangle , and thus $\mathcal{P}(X)$ is a *Boolean group* (i.e. an abelian 2-group).
- 4. Interaction With Powersets and Vector Spaces I. The pair $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$ consisting of
 - The group $\mathcal{P}(X)$ of Item 17;
 - The map $\alpha_{\mathcal{P}(X)} \colon \mathbb{F}_2 \times \mathcal{P}(X) \to \mathcal{P}(X)$ defined by

$$0 \cdot U \stackrel{\text{def}}{=} \emptyset,$$

 $1 \cdot U \stackrel{\text{def}}{=} U;$

is an \mathbb{F}_2 -vector space.

- 5. *Interaction With Powersets and Vector Spaces II.* If *X* is finite, then:
 - (a) The set of singletons sets on the elements of X forms a basis for the \mathbb{F}_2 -vector space $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$ of Item 4.
 - (b) We have

$$\dim(\mathcal{P}(X)) = \#X.$$

6. *Interaction With Powersets and Rings.* The quintuple $(\mathcal{P}(X), \Delta, \cap, \emptyset, X)$

1. When $X = \emptyset$, we have an isomorphism of groups between $\mathcal{P}(\emptyset)$ and the trivial group:

$$(\mathcal{P}(\emptyset), \triangle, \emptyset, \mathrm{id}_{\mathcal{P}(\emptyset)}) \cong \mathrm{pt.}$$

2. When X = pt, we have an isomorphism of groups between $\mathcal{P}(pt)$ and $\mathbb{Z}_{/2}$:

$$(\mathcal{P}(\mathsf{pt}), \Delta, \emptyset, \mathsf{id}_{\mathcal{P}(\mathsf{pt})}) \cong \mathbb{Z}_{/2}.$$

3. When $X = \{0, 1\}$, we have an isomorphism of groups between $\mathcal{P}(\{0, 1\})$ and $\mathbb{Z}_{/2} \times \mathbb{Z}_{/2}$:

$$(\mathcal{P}(\{0,1\}), \Delta, \emptyset, \mathrm{id}_{\mathcal{P}(\{0,1\})}) \cong \mathbb{Z}_{/2} \times \mathbb{Z}_{/2}.$$

¹⁴Here are some examples:

is a commutative ring.15

7. Interaction With Direct Images. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

with components

$$f_!(U) \triangle f_!(V) \subset f_!(U \triangle V)$$

indexed by $U, V \in \mathcal{P}(X)$.

8. Interaction With Inverse Images. The diagram

$$\mathcal{P}(Y)^{\text{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\text{op},-1} \times f^{-1}} \mathcal{P}(X)^{\text{op}} \times \mathcal{P}(X)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

i.e. we have

$$f^{-1}(U) \triangle f^{-1}(V) = f^{-1}(U \triangle V)$$

for each $U, V \in \mathcal{P}(Y)$.

9. Interaction With Codirect Images. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_*^{\mathsf{op}} \times f_*} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

¹⁵ Warning: The analogous statement replacing intersections by unions (i.e. that the

with components

$$f_*(U \triangle V) \subset f_*(U) \triangle f_*(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

Proof. Item 1, Lack of Functoriality: Omitted.

Item 2, Via Unions and Intersections: See [Pro25m].

Item 3, Symmetric Differences of Disjoint Sets: Since U and V are disjoint, we have $U \cap V = G$ and therefore we have

have $U \cap V = \emptyset$, and therefore we have

$$U \triangle V = (U \cup V) \setminus (U \cap V)$$
$$= (U \cup V) \setminus \emptyset$$
$$= U \cup V,$$

where we've used Item 2 and Item 12 of Definition 4.3.10.1.2.

Item 4, Associativity: See [Pro25ao].

Item 5, Unitality: This follows from Item 6 and [Pro25at].

Item 6, *Commutativity*: See [Pro25ap].

Item 7, Invertibility: See [Pro25av].

Item 8, Interaction With Unions: See [Pro25bc].

Item 9, Interaction With Complements I: See [Pro25as].

Item 10, *Interaction With Complements II*: This follows from Item 6 and [Pro25ax].

Item 11, *Interaction With Complements III*: See [Pro25aq].

Item 12, "Transitivity": We have

$$(U \triangle V) \triangle (V \triangle W) = U \triangle (V \triangle (V \triangle W))$$
 (by Item 4)

$$= U \triangle ((V \triangle V) \triangle W)$$
 (by Item 4)

$$= U \triangle (\emptyset \triangle W)$$
 (by Item 7)

$$= U \triangle W.$$
 (by Item 5)

This finishes the proof.

Item 13, The Triangle Inequality for Symmetric Differences: This follows from

Items 2 and 12.

Item 14, *Distributivity Over Intersections*: See [Pro25q].

Item 15, Interaction With Characteristic Functions: See [Pro25g].

Item 16, Bijectivity: Omitted.

quintuple $(\mathcal{P}(X), \Delta, \cup, \emptyset, X)$ is a ring) is false, however. See [Pro25aw] for a proof.

Item 17, Interaction With Powersets and Groups: Item 17a follows from Items 4 to 7, while Item 3b follows from Item 7.¹⁶

Item 4, Interaction With Powersets and Vector Spaces I: See [MSE 2719059].

Item 5, Interaction With Powersets and Vector Spaces II: See [MSE 2719059].

Item 6, Interaction With Powersets and Rings: This follows from Items 6 and 15 of Definition 4.3.9.1.2 and Items 14 and 17.¹⁷

Item 7, Interaction With Direct Images: This is a repetition of Item 9 of Definition 4.6.1.1.5 and is proved there.

Item 8, Interaction With Inverse Images: This is a repetition of Item 9 of Definition 4.6.2.1.3 and is proved there.

Item 9, Interaction With Codirect Images: This is a repetition of Item 8 of Definition 4.6.3.1.7 and is proved there. □

4.4 Powersets

4.4.1 Foundations

Let *X* be a set.

Definition 4.4.1.1.1. The **powerset of** X is the set $\mathcal{P}(X)$ defined by

$$\mathcal{P}(X) \stackrel{\text{def}}{=} \{ U \in P \mid U \subset X \},\$$

where P is the set in the axiom of powerset, ?? of ??.

Remark 4.4.1.1.2. Under the analogy that $\{t, f\}$ should be the (-1)-categorical analogue of Sets, we may view the powerset of a set as a decategorification of the category of presheaves of a category (or of the category of copresheaves):

• The powerset of a set *X* is equivalently (Item 2 of Definition 4.5.1.1.4) the set

$$Sets(X, \{t, f\})$$

of functions from X to the set $\{t, f\}$ of classical truth values.

¹⁶Reference: [Pro25ar].

¹⁷Reference: [Pro25au].

• The category of presheaves on a category *C* is the category

$$\operatorname{Fun}(C^{\operatorname{op}},\operatorname{Sets})$$

of functors from C^{op} to the category Sets of sets.

Notation 4.4.1.1.3. Let X be a set.

- 1. We write $\mathcal{P}_0(X)$ for the set of nonempty subsets of X.
- 2. We write $\mathcal{P}_{fin}(X)$ for the set of finite subsets of X.

Proposition 4.4.1.1.4. Let X be a set.

- 1. *Co/Completeness*. The (posetal) category (associated to) $(\mathcal{P}(X), \subset)$ is complete and cocomplete:
 - (a) *Products*. The products in $\mathcal{P}(X)$ are given by intersection of subsets.
 - (b) *Coproducts*. The coproducts in $\mathcal{P}(X)$ are given by union of subsets.
 - (c) Co/Equalisers. Being a posetal category, $\mathcal{P}(X)$ only has at most one morphisms between any two objects, so co/equalisers are trivial.
- 2. *Cartesian Closedness*. The category $\mathcal{P}(X)$ is Cartesian closed.
- 3. Powersets as Sets of Relations. We have bijections

$$\mathcal{P}(X) \cong \text{Rel}(\text{pt}, X),$$

 $\mathcal{P}(X) \cong \text{Rel}(X, \text{pt}),$

natural in $X \in Obj(Sets)$.

4. Interaction With Products I. The map

$$\mathcal{P}(X) \times \mathcal{P}(Y) \longrightarrow \mathcal{P}(X \coprod Y)$$
$$(U, V) \longmapsto U \cup V$$

is an isomorphism of sets, natural in $X, Y \in \text{Obj}(\mathsf{Sets})$ with respect to each of the functor structures $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* on \mathcal{P} of Definition 4.4.2.1.1. Moreover, this makes each of $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* into a symmetric monoidal functor.

5. Interaction With Products II. The map

$$\mathcal{P}(X) \times \mathcal{P}(Y) \longrightarrow \mathcal{P}(X \coprod Y)$$
$$(U, V) \longmapsto U \boxtimes_{X \times Y} V,$$

where 18

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \{(u, v) \in X \times Y \mid u \in U \text{ and } v \in V\}$$

is an inclusion of sets, natural in $X, Y \in \text{Obj}(\mathsf{Sets})$ with respect to each of the functor structures $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* on \mathcal{P} of Definition 4.4.2.1.1. Moreover, this makes each of $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* into a symmetric monoidal functor.

6. Interaction With Products III. We have an isomorphism

$$\mathcal{P}(X) \otimes \mathcal{P}(Y) \cong \mathcal{P}(X \times Y),$$

natural in $X, Y \in \text{Obj}(\mathsf{Sets})$ with respect to each of the functor structures $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* on \mathcal{P} of Definition 4.4.2.1.1, where \otimes denotes the tensor product of suplattices of ??. Moreover, this makes each of $\mathcal{P}_!$, \mathcal{P}^{-1} , and \mathcal{P}_* into a symmetric monoidal functor.

Proof. Item 1, Co/Completeness: Omitted.

Item 2, Cartesian Closedness: See Section 4.4.7.

Item 3, Powersets as Sets of Relations: Indeed, we have

$$Rel(pt, X) \stackrel{\text{def}}{=} \mathcal{P}(pt \times X)$$
$$\cong \mathcal{P}(X)$$

and

$$\operatorname{Rel}(X,\operatorname{pt})\stackrel{\operatorname{def}}{=} \mathcal{P}(X \times \operatorname{pt})$$

 $\cong \mathcal{P}(X),$

where we have used Item 5 of Definition 4.1.3.1.3.

¹⁸The set $U \boxtimes_{X \times Y} V$ is usually denoted simply $U \times V$. Here we denote it in this somewhat weird way to highlight the similarity to external tensor products in six-functor formalisms

Item 4, Interaction With Products I: The inverse of the map in the statement is the map

$$\Phi \colon \mathcal{P}(X \mid \mid Y) \to \mathcal{P}(X) \times \mathcal{P}(Y)$$

defined by

$$\Phi(S) \stackrel{\text{def}}{=} (S_X, S_Y)$$

for each $S \in \mathcal{P}(X \coprod Y)$, where

$$S_X \stackrel{\text{def}}{=} \{ x \in X \mid (0, x) \in S \}$$

 $S_Y \stackrel{\text{def}}{=} \{ y \in Y \mid (1, y) \in S \}.$

The rest of the proof is omitted.

Item 5, Interaction With Products II: Omitted.

Item 6, Interaction With Products III: Omitted.

4.4.2 Functoriality of Powersets

Proposition 4.4.2.1.1. Let *X* be a set.

1. Functoriality I. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}_1 \colon \mathsf{Sets} \to \mathsf{Sets}$$

where

• Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}_!(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

 Action on Morphisms. For each A, B ∈ Obj(Sets), the action on morphisms

$$\mathcal{P}_{*|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of $\mathcal{P}_!$ at (A,B) is the map defined by sending a map of sets $f\colon A\to B$ to the map

$$\mathcal{P}_!(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_!(f) \stackrel{\text{def}}{=} f_!,$$

as in Definition 4.6.1.1.1.

2. Functoriality II. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}^{-1}$$
: Sets^{op} \rightarrow Sets,

where

• Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}^{-1}(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

• *Action on Morphisms*. For each $A, B \in \mathsf{Obj}(\mathsf{Sets})$, the action on morphisms

$$\mathcal{P}_{AB}^{-1} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(B),\mathcal{P}(A))$$

of \mathcal{P}^{-1} at (A, B) is the map defined by sending a map of sets $f: A \to B$ to the map

$$\mathcal{P}^{-1}(f) \colon \mathcal{P}(B) \to \mathcal{P}(A)$$

defined by

$$\mathcal{P}^{-1}(f) \stackrel{\text{def}}{=} f^{-1},$$

as in Definition 4.6.2.1.1.

3. Functoriality III. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}_* \colon \mathsf{Sets} \to \mathsf{Sets}$$
.

where

• Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}_*(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

• *Action on Morphisms*. For each $A, B \in Obj(Sets)$, the action on morphisms

$$\mathcal{P}_{!|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of \mathcal{P}_* at (A, B) is the map defined by by sending a map of sets $f: A \to B$ to the map

$$\mathcal{P}_*(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_*(f) \stackrel{\text{def}}{=} f_*,$$

as in Definition 4.6.3.1.1.

Proof. Item 1, Functoriality I: This follows from Items 3 and 4 of Definition 4.6.1.1.6. *Item 2, Functoriality II:* This follows from Items 3 and 4 of Definition 4.6.2.1.4. *Item 3, Functoriality III:* This follows from Items 3 and 4 of Definition 4.6.3.1.8.

4.4.3 Adjointness of Powersets I

Proposition 4.4.3.1.1. We have an adjunction

$$(\mathcal{P}^{-1} \dashv \mathcal{P}^{-1,op})$$
: Sets^{op} $\underbrace{\overset{\mathcal{P}^{-1}}{\downarrow}}_{\mathcal{P}^{-1,op}}$ Sets,

witnessed by a bijection

$$\underbrace{\mathsf{Sets}^{\mathsf{op}}(\mathcal{P}(X), Y)}_{\overset{\mathsf{def}}{=}\mathsf{Sets}(Y, \mathcal{P}(X))} \cong \mathsf{Sets}(X, \mathcal{P}(Y)),$$

natural in $X \in Obj(Sets)$ and $Y \in Obj(Sets^{op})$.

Proof. We have

$$Sets^{op}(\mathcal{P}(A), B) \stackrel{\text{def}}{=} Sets(B, \mathcal{P}(A))$$

$$\cong Sets(B, Sets(A, \{t, f\})) \quad \text{(by Item 2 of Definition 4.5.1.1.4)}$$

$$\cong Sets(A \times B, \{t, f\}) \quad \text{(by Item 2 of Definition 4.1.3.1.3)}$$

$$\cong Sets(A, Sets(B, \{t, f\})) \quad \text{(by Item 2 of Definition 4.1.3.1.3)}$$

$$\cong Sets(A, \mathcal{P}(B)), \quad \text{(by Item 2 of Definition 4.5.1.1.4)}$$

where all bijections are natural in A and B.¹⁹

¹⁹Here we are using Item 3 of Definition 4.5.1.1.4.

4.4.4 Adjointness of Powersets II

Proposition 4.4.4.1.1. We have an adjunction

$$(Gr \dashv \mathcal{P}_!)$$
: Sets $\underbrace{\overset{Gr}{\downarrow}}_{\mathcal{P}_l}$ Rel,

witnessed by a bijection of sets

$$Rel(Gr(X), Y) \cong Sets(X, \mathcal{P}(Y))$$

natural in $X \in \text{Obj}(\mathsf{Sets})$ and $Y \in \text{Obj}(\mathsf{Rel})$, where Gr is the graph functor of Relations, Item 1 of Definition 8.2.2.1.2 and $\mathcal{P}_!$ is the functor of Relations, Definition 8.7.5.1.1.

Proof. We have

where all bijections are natural in A, (where we are using Item 3 of Definition 4.5.1.1.4). Explicitly, this isomorphism is given by sending a relation $R: Gr(A) \to B$ to the map $R^{\dagger}: A \to \mathcal{P}(B)$ sending a to the subset R(a) of B, as in Relations, Definition 8.1.1.1.1.

Naturality in B is then the statement that given a relation $R \colon B \to B'$, the diagram

$$\operatorname{Rel}(\operatorname{Gr}(A),B) \xrightarrow{R \diamond -} \operatorname{Rel}(\operatorname{Gr}(A),B')$$

$$\downarrow \downarrow \qquad \qquad \downarrow \downarrow \downarrow$$

$$\operatorname{Sets}(A,\mathcal{P}(B)) \xrightarrow{R_!} \operatorname{Sets}(A,\mathcal{P}(B'))$$

commutes, which follows from Relations, Definition 8.7.1.1.3.

4.4.5 Powersets as Free Cocompletions

Let *X* be a set.

Proposition 4.4.5.1.1. The pair $(\mathcal{P}(X), \chi_{(-)})$ consisting of

- The powerset $(\mathcal{P}(X), \subset)$ of *X* of Definition 4.4.1.1.1;
- The characteristic embedding $\chi_{(-)} \colon X \hookrightarrow \mathcal{P}(X)$ of X into $\mathcal{P}(X)$ of Definition 4.5.4.1.1;

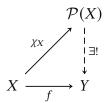
satisfies the following universal property:

- (\star) Given another pair (Y, f) consisting of
 - A suplattice (Y, \preceq) ;
 - A function $f: X \to Y$;

there exists a unique morphism of suplattices

$$(\mathcal{P}(X),\subset) \xrightarrow{\exists!} (Y,\preceq)$$

making the diagram



commute.

Proof. This is a rephrasing of Definition 4.4.5.1.2, which we prove below. \Box

Proposition 4.4.5.1.2. We have an adjunction

$$(\mathcal{P} \dashv \overline{\Xi})$$
: Sets $\underbrace{\overset{\mathcal{P}}{=}}$ SupLat,

²⁰Here we only remark that the unique morphism of suplattices in the statement is given

witnessed by a bijection

$$SupLat((\mathcal{P}(X), \subset), (Y, \preceq)) \cong Sets(X, Y),$$

natural in $X \in \text{Obj}(\mathsf{Sets})$ and $(Y, \preceq) \in \text{Obj}(\mathsf{SupLat})$, where:

- The category SupLat is the category of suplattices of ??.
- The map

$$\chi_X^* \colon \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

witnessing the above bijection is defined by

$$\chi_X^*(f) \stackrel{\text{def}}{=} f \circ \chi_X,$$

i.e. by sending a morphism of suplattices $f \colon \mathcal{P}(X) \to Y$ to the composition

$$X \stackrel{\chi_X}{\hookrightarrow} \mathcal{P}(X) \stackrel{f}{\longrightarrow} Y.$$

• The map

$$\operatorname{Lan}_{\chi_X} \colon \operatorname{\mathsf{Sets}}(X,Y) \to \operatorname{\mathsf{SupLat}}((\mathcal{P}(X),\subset),(Y,\preceq))$$

witnessing the above bijection is given by sending a function $f: X \to Y$ to its left Kan extension along χ_X ,

$$\operatorname{Lan}_{\chi_X}(f) \colon \mathcal{P}(X) \to Y, \qquad \begin{array}{c} \mathcal{P}(X) \\ \downarrow^{\chi_X} & \downarrow^{\downarrow} \operatorname{Lan}_{\chi_X}(f) \\ X & \xrightarrow{f} Y. \end{array}$$

Moreover, invoking the bijection $\mathcal{P}(X) \cong \operatorname{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ of Item 2 of Definition 4.5.1.1.4, $\operatorname{Lan}_{\chi_X}(f)$ can be explicitly computed by

$$\left[\operatorname{Lan}_{\chi_X}(f)\right](U) = \int^{x \in X} \chi_{\mathcal{P}(X)}(\chi_x, U) \odot f(x)$$

by the left Kan extension $\operatorname{Lan}_{\chi_X}(f)$ of f along χ_X .

$$= \int_{x \in X} \chi_{U}(x) \odot f(x)$$

$$= \bigvee_{x \in X} (\chi_{U}(x) \odot f(x))$$

$$= \left(\bigvee_{x \in U} (\chi_{U}(x) \odot f(x))\right) \vee \left(\bigvee_{x \in U^{c}} (\chi_{U}(x) \odot f(x))\right)$$

$$= \left(\bigvee_{x \in U} f(x)\right) \vee \left(\bigvee_{x \in U^{c}} \varnothing_{Y}\right)$$

$$= \bigvee_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$, where:

- We have used ?? for the first equality.
- We have used Definition 4.5.5.1.1 for the second equality.
- We have used?? for the third equality.
- − The symbol \bigvee denotes the join in (Y, \preceq) .
- The symbol \odot denotes the tensor of an element of Y by a truth value as in $\ref{eq:total_exp}$. In particular, we have

true
$$\odot f(x) \stackrel{\text{def}}{=} f(x)$$
,
false $\odot f(x) \stackrel{\text{def}}{=} \varnothing_Y$,

where \emptyset_Y is the bottom element of (Y, \preceq) .

In particular, when $(Y, \preceq_Y) = (\mathcal{P}(B), \subset)$ for some set B, the Kan extension $\operatorname{Lan}_{\chi_X}(f)$ is given by

$$[\operatorname{Lan}_{\chi_X}(f)](U) = \bigvee_{x \in U} f(x)$$
$$= \bigcup_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$.

Proof. Map I: We define a map

$$\Phi_{X,Y} \colon \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

as in the statement, i.e. by

$$\Phi_{X,Y}(f) \stackrel{\text{def}}{=} f \circ \chi_X$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$. *Map II*: We define a map

$$\Psi_{X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$$

as in the statement, i.e. by

$$\Psi_{X,Y}(f) \stackrel{\text{def}}{=} \operatorname{Lan}_{\chi_X}(f), \qquad X \xrightarrow{\chi_X} \downarrow \underset{f}{\downarrow} \operatorname{Lan}_{\chi_X}(f)$$

for each $f \in Sets(X, Y)$. *Invertibility I*: We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathrm{id}_{\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))}$$
.

We have

$$\begin{split} \big[\Psi_{X,Y} \circ \Phi_{X,Y}\big](f) &\stackrel{\text{def}}{=} \Psi_{X,Y} \big(\Phi_{X,Y}(f)\big) \\ &\stackrel{\text{def}}{=} \Psi_{X,Y}(f \circ \chi_X) \\ &\stackrel{\text{def}}{=} \operatorname{Lan}_{\chi_X}(f \circ \chi_X) \end{split}$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$. We now claim that

$$\mathrm{Lan}_{\chi_X}(f\circ\chi_X)=f$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq)).$ Indeed, we have

$$\left[\operatorname{Lan}_{\chi_X}(f\circ\chi_X)\right](U) = \bigvee_{x\in U} f(\chi_X(x))$$

$$= f\left(\bigvee_{x \in U} \chi_X(x)\right)$$
$$= f\left(\bigcup_{x \in U} \{x\}\right)$$
$$= f(U)$$

for each $U \in \mathcal{P}(X)$, where we have used that f is a morphism of suplattices and hence preserves joins for the second equality. This proves our claim. Since we have shown that

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) = f$$

for each $f \in \operatorname{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$, it follows that $\Psi_{X,Y} \circ \Phi_{X,Y}$ must be equal to the identity map $\operatorname{id}_{\operatorname{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))}$ of $\operatorname{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$. *Invertibility II*: We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathrm{Sets}(X,Y)}$$
.

We have

$$\begin{split} \big[\Phi_{X,Y} \circ \Psi_{X,Y} \big](f) &\stackrel{\text{def}}{=} \Phi_{X,Y} \big(\Psi_{X,Y}(f) \big) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y} \big(\text{Lan}_{\chi_X}(f) \big) \\ &\stackrel{\text{def}}{=} \text{Lan}_{\chi_X}(f) \circ \chi_X \end{split}$$

for each $f \in Sets(X, Y)$. We now claim that

$$Lan_{\chi_X}(f) \circ \chi_X = f$$

for each $f \in Sets(X, Y)$. Indeed, we have

$$\left[\operatorname{Lan}_{\chi_X}(f) \circ \chi_X \right](x) = \bigvee_{y \in \{x\}} f(y) \\
= f(x)$$

for each $x \in X$. This proves our claim. Since we have shown that

$$\big[\Phi_{X,Y}\circ\Psi_{X,Y}\big](f)=f$$

for each $f \in \mathsf{Sets}(X,Y)$, it follows that $\Phi_{X,Y} \circ \Psi_{X,Y}$ must be equal to the identity map $\mathsf{id}_{\mathsf{Sets}(X,Y)}$ of $\mathsf{Sets}(X,Y)$.

Naturality for Φ , Part I: We need to show that, given a function $f: X \to X'$, the diagram

$$\begin{split} \mathsf{SupLat}((\mathcal{P}(X'),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X',Y}} \mathsf{Sets}(X',Y) \\ & & \downarrow f^* \\ \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y) \end{split}$$

commutes. Indeed, we have

$$\begin{split} \left[\Phi_{X,Y} \circ \mathcal{P}_{!}(f)^{*}\right] (\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y} \left(\mathcal{P}_{!}(f)^{*}(\xi)\right) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y} (\xi \circ f_{!}) \\ &\stackrel{\text{def}}{=} (\xi \circ f_{!}) \circ \chi_{X} \\ &= \xi \circ (f_{!} \circ \chi_{X}) \\ &\stackrel{(\dagger)}{=} \xi \circ (\chi_{X'} \circ f) \\ &= (\xi \circ \chi_{X'}) \circ f \\ &\stackrel{\text{def}}{=} \Phi_{X',Y} (\xi) \circ f \\ &\stackrel{\text{def}}{=} f^{*} \left(\Phi_{X',Y} (\xi)\right) \\ &\stackrel{\text{def}}{=} \left[f^{*} \circ \Phi_{X',Y}\right] (\xi), \end{split}$$

for each $\xi \in \text{SupLat}((\mathcal{P}(X'), \subset), (Y, \preceq))$, where we have used Item 1 of Definition 4.5.4.1.3 for the fifth equality above.

Naturality for Φ , *Part II*: We need to show that, given a morphism of suplattices

$$g: (Y, \preceq_Y) \rightarrow (Y', \preceq_{Y'}),$$

the diagram

$$\begin{split} \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y) \\ & g_! & & \downarrow g_! \\ \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y',\preceq)) & \xrightarrow{\Phi_{X,Y'}} \mathsf{Sets}(X,Y') \end{split}$$

commutes. Indeed, we have

$$\left[\Phi_{X,Y'}\circ g_!\right](\xi)\stackrel{\text{def}}{=}\Phi_{X,Y'}(g_!(\xi))$$

$$\stackrel{\text{def}}{=} \Phi_{X,Y'}(g \circ \xi)$$

$$\stackrel{\text{def}}{=} (g \circ \xi) \circ \chi_X$$

$$= g \circ (\xi \circ \chi_X)$$

$$\stackrel{\text{def}}{=} g \circ (\Phi_{X,Y}(\xi))$$

$$\stackrel{\text{def}}{=} g_!(\Phi_{X,Y}(\xi))$$

$$\stackrel{\text{def}}{=} [g_! \circ \Phi_{X,Y}](\xi).$$

for each $\xi \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$.

Naturality for Ψ : Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Ψ is also natural in each argument.

Warning 4.4.5.1.3. Although the assignment $X \mapsto \mathcal{P}(X)$ is called the *free cocompletion of* X, it is not an idempotent operation, i.e. we have $\mathcal{P}(\mathcal{P}(X)) \neq \mathcal{P}(X)$.

4.4.6 Powersets as Free Completions

Let *X* be a set.

Proposition 4.4.6.1.1. The pair $(\mathcal{P}(X), \chi_{(-)})$ consisting of

- The powerset of X together with reverse inclusion $\mathcal{P}(X)^{\text{op}} = (\mathcal{P}(X), \supset)$ of Definition 4.4.1.1.1;
- The characteristic embedding $\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$ of X into $\mathcal{P}(X)$ of Definition 4.5.4.1.1;

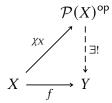
satisfies the following universal property:

- (\star) Given another pair (Y, f) consisting of
 - An inflattice (Y, \preceq) ;
 - A function $f: X \to Y$;

there exists a unique morphism of inflattices

$$(\mathcal{P}(X),\supset) \xrightarrow{\exists!} (Y,\preceq)$$

making the diagram



commute.

Proof. This is a rephrasing of Definition 4.4.6.1.2, which we prove below. \Box

Proposition 4.4.6.1.2. We have an adjunction

witnessed by a bijection

$$InfLat((\mathcal{P}(X),\supset),(Y,\preceq)) \cong Sets(X,Y),$$

natural in $X \in \text{Obj}(\mathsf{Sets})$ and $(Y, \preceq) \in \text{Obj}(\mathsf{InfLat})$, where:

- The category InfLat is the category of inflattices of ??.
- The map

$$\chi_X^* : \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

witnessing the above bijection is defined by

$$\chi_X^*(f) \stackrel{\text{def}}{=} f \circ \chi_X,$$

i.e. by sending a morphism of inflattices $f\colon \mathcal{P}(X)^{\mathrm{op}} \to Y$ to the composition

$$X \stackrel{\chi\chi}{\hookrightarrow} \mathcal{P}(X)^{\mathsf{op}} \stackrel{f}{\longrightarrow} Y.$$

²¹Here we only remark that the unique morphism of inflattices in the statement is given by the right Kan extension $\operatorname{Ran}_{\chi_X}(f)$ of f along χ_X .

• The map

$$\operatorname{Ran}_{\gamma_X} : \operatorname{\mathsf{Sets}}(X,Y) \to \operatorname{\mathsf{InfLat}}((\mathcal{P}(X),\supset),(Y,\preceq))$$

witnessing the above bijection is given by sending a function $f: X \to Y$ to its right Kan extension along χ_X ,

$$\operatorname{Ran}_{\chi_X}(f) \colon \mathcal{P}(X)^{\operatorname{op}} \to Y, \qquad \chi_X / \underset{f}{ \downarrow_{\operatorname{Ran}_{\chi_X}(f)}} X \xrightarrow{f} Y.$$

Moreover, invoking the bijection $\mathcal{P}(X) \cong \operatorname{Sets}(X, \{t, f\})$ of Item 2 of Definition 4.5.1.1.4, $\operatorname{Ran}_{\chi_X}(f)$ can be explicitly computed by

$$[\operatorname{Ran}_{\chi_{X}}(f)](U) = \int_{x \in X} \chi_{\mathcal{P}(X)^{\operatorname{op}}}(\chi_{x}, U) \, \, \, \, \, f(x)$$

$$= \int_{x \in X} \chi_{\mathcal{P}(X)}(U, \chi_{x}) \, \, \, \, \, f(x)$$

$$= \int_{x \in X} \chi_{U}(x) \, \, \, \, \, f(x)$$

$$= \left(\bigwedge_{x \in U} \chi_{U}(x) \, \, \, \, \, \, f(x) \right) \wedge \left(\bigwedge_{x \in U^{c}} \chi_{U}(x) \, \, \, \, \, \, f(x) \right)$$

$$= \left(\bigwedge_{x \in U} f(x) \right) \wedge \left(\bigwedge_{x \in U^{c}} \infty_{Y} \right)$$

$$= \left(\bigwedge_{x \in U} f(x) \right) \wedge \infty_{Y}$$

$$= \bigwedge_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$, where:

- We have used ?? for the first equality.

- We have used Definition 4.5.5.1.1 for the second equality.
- We have used ?? for the third equality.
- − The symbol \land denotes the meet in (Y, \preceq) .
- The symbol \pitchfork denotes the cotensor of an element of Y by a truth value as in $\ref{eq:total_eq}$. In particular, we have

true
$$\pitchfork f(x) \stackrel{\text{def}}{=} f(x)$$
, false $\pitchfork f(x) \stackrel{\text{def}}{=} \infty_Y$,

where ∞_Y is the top element of (Y, \preceq) .

In particular, when $(Y, \preceq_Y) = (\mathcal{P}(B), \subset)$ for some set B, the Kan extension $\operatorname{Ran}_{\chi_X}(f)$ is given by

$$[\operatorname{Ran}_{\chi_X}(f)](U) = \bigwedge_{x \in U} f(x)$$
$$= \bigcap_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$.

Proof. Map I: We define a map

$$\Phi_{X,Y} \colon \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) \to \mathsf{Sets}(X,Y)$$

as in the statement, i.e. by

$$\Phi_{X,Y}(f) \stackrel{\text{def}}{=} f \circ \chi_X$$

 $\text{for each } f \in \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq)).$

Map II: We define a map

$$\Psi_{X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq))$$

as in the statement, i.e. by

$$\Psi_{X,Y}(f) \stackrel{\text{def}}{=} \operatorname{Ran}_{\chi_X}(f), \qquad \chi_X / \underset{f}{ \downarrow_{\operatorname{Ran}_{\chi_X}(f)}} X \xrightarrow{\chi_X} Y,$$

for each $f \in Sets(X, Y)$. *Invertibility I*: We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathrm{id}_{\mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq))}$$
.

We have

$$\begin{split} \big[\Psi_{X,Y} \circ \Phi_{X,Y}\big](f) &\stackrel{\text{def}}{=} \Psi_{X,Y} \big(\Phi_{X,Y}(f)\big) \\ &\stackrel{\text{def}}{=} \Psi_{X,Y}(f \circ \chi_X) \\ &\stackrel{\text{def}}{=} \operatorname{Ran}_{\chi_X}(f \circ \chi_X) \end{split}$$

for each $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$. We now claim that

$$\operatorname{Ran}_{\chi_X}(f\circ\chi_X)=f$$

for each $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$. Indeed, we have

$$\begin{aligned} \left[\operatorname{Ran}_{\chi_X} (f \circ \chi_X) \right] (U) &= \bigwedge_{x \in U} f(\chi_X(x)) \\ &= f \left(\bigwedge_{x \in U} \chi_X(x) \right) \\ &= f \left(\bigcup_{x \in U} \{x\} \right) \\ &= f(U) \end{aligned}$$

for each $U \in \mathcal{P}(X)$, where we have used that f is a morphism of inflattices and hence preserves meets in $(\mathcal{P}(X), \supset)$ (i.e. joins in $(\mathcal{P}(X), \subset)$) for the second equality. This proves our claim. Since we have shown that

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) = f$$

for each $f \in \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq))$, it follows that $\Psi_{X,Y} \circ \Phi_{X,Y}$ must be equal to the identity map $\mathsf{id}_{\mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq))}$ of $\mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq))$. *Invertibility II*: We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathsf{Sets}(X,Y)} .$$

We have

$$\begin{split} \big[\Phi_{X,Y} \circ \Psi_{X,Y} \big] (f) &\stackrel{\text{def}}{=} \Phi_{X,Y} \big(\Psi_{X,Y} (f) \big) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y} \big(\text{Ran}_{\chi_X} (f) \big) \\ &\stackrel{\text{def}}{=} \text{Ran}_{\chi_X} (f) \circ \chi_X \end{split}$$

for each $f \in Sets(X, Y)$. We now claim that

$$\operatorname{Ran}_{\chi_X}(f) \circ \chi_X = f$$

for each $f \in Sets(X, Y)$. Indeed, we have

$$[\operatorname{Ran}_{\chi_X}(f) \circ \chi_X](x) = \bigwedge_{y \in \{x\}} f(y)$$
$$= f(x)$$

for each $x \in X$. This proves our claim. Since we have shown that

$$[\Phi_{X,Y} \circ \Psi_{X,Y}](f) = f$$

for each $f \in \mathsf{Sets}(X,Y)$, it follows that $\Phi_{X,Y} \circ \Psi_{X,Y}$ must be equal to the identity map $\mathsf{id}_{\mathsf{Sets}(X,Y)}$ of $\mathsf{Sets}(X,Y)$.

Naturality for Φ , *Part I*: We need to show that, given a function $f: X \to X'$, the diagram

$$\mathsf{InfLat}((\mathcal{P}(X'),\supset),(Y,\preceq)) \xrightarrow{\Phi_{X',Y}} \mathsf{Sets}(X',Y)$$

$$\downarrow^{f^*} \qquad \qquad \downarrow^{f^*}$$

$$\mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y)$$

commutes. Indeed, we have

$$\begin{split} \left[\Phi_{X,Y} \circ \mathcal{P}_{!}(f)^{*}\right] (\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y} (\mathcal{P}_{!}(f)^{*}(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y} (\xi \circ f_{!}) \\ &\stackrel{\text{def}}{=} (\xi \circ f_{!}) \circ \chi_{X} \\ &= \xi \circ (f_{!} \circ \chi_{X}) \\ &\stackrel{(\dagger)}{=} \xi \circ (\chi_{X'} \circ f) \end{split}$$

$$= (\xi \circ \chi_{X'}) \circ f$$

$$\stackrel{\text{def}}{=} \Phi_{X',Y}(\xi) \circ f$$

$$\stackrel{\text{def}}{=} f^* (\Phi_{X',Y}(\xi))$$

$$\stackrel{\text{def}}{=} [f^* \circ \Phi_{X',Y}](\xi),$$

for each $\xi \in \text{InfLat}((\mathcal{P}(X'), \supset), (Y, \preceq))$, where we have used Item 1 of Definition 4.5.4.1.3 for the fifth equality above.

Naturality for Φ , Part II: We need to show that, given a cocontinuous morphism of posets

$$g: (Y, \preceq_Y) \rightarrow (Y', \preceq_{Y'}),$$

the diagram

$$\begin{split} \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y) \\ & g_! & & \downarrow g_! \\ & \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y',\preceq)) & \xrightarrow{\Phi_{X,Y'}} \mathsf{Sets}(X,Y') \end{split}$$

commutes. Indeed, we have

$$\begin{split} \left[\Phi_{X,Y'} \circ g_!\right](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g_!(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g \circ \xi) \\ &\stackrel{\text{def}}{=} (g \circ \xi) \circ \chi_X \\ &= g \circ (\xi \circ \chi_X) \\ &\stackrel{\text{def}}{=} g \circ \left(\Phi_{X,Y}(\xi)\right) \\ &\stackrel{\text{def}}{=} g_!(\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} \left[g_! \circ \Phi_{X,Y}\right](\xi). \end{split}$$

for each $\xi \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$.

Naturality for Ψ : Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Ψ is also natural in each argument.

Warning 4.4.6.1.3. Although the assignment $X \mapsto \mathcal{P}(X)^{\text{op}}$ is called the *free completion of X*, it is not an idempotent operation, i.e. we have $\mathcal{P}(\mathcal{P}(X)^{\text{op}})^{\text{op}} \neq \mathcal{P}(X)^{\text{op}}$.

4.4.7 The Internal Hom of a Powerset

Let X be a set and let $U, V \in \mathcal{P}(X)$.

Proposition 4.4.7.1.1. The **internal Hom of** $\mathcal{P}(X)$ **from** U **to** V is the subset $[U,V]_X^{22}$ of X given by

$$[U, V]_X = U^{c} \cup V$$
$$= (U \setminus V)^{c}$$

where U^{c} is the complement of U of Definition 4.3.11.1.

Proof. Proof of the Equality $U^c \cup V = (U \setminus V)^c$: We have

$$(U \setminus V)^{c} \stackrel{\text{def}}{=} X \setminus (U \setminus V)$$

$$= (X \cap V) \cup (X \setminus U)$$

$$= V \cup (X \setminus U)$$

$$\stackrel{\text{def}}{=} V \cup U^{c}$$

$$= U^{c} \cup V,$$

where we have used:

- 1. Item 10 of Definition 4.3.10.1.2 for the second equality.
- 2. Item 4 of Definition 4.3.9.1.2 for the third equality.
- 3. Item 4 of Definition 4.3.8.1.2 for the last equality.

This finishes the proof.

Proof that $U^c \cup V$ *Is Indeed the Internal Hom*: This follows from Item 2 of Definition 4.3.9.1.2.

Remark 4.4.7.1.2. Henning Makholm suggests the following heuristic intuition for the internal Hom of $\mathcal{P}(X)$ from U to V ([MSE 267365]):

1. Since products in $\mathcal{P}(X)$ are given by binary intersections (Item 1 of Definition 4.4.1.1.4), the right adjoint $\operatorname{Hom}_{\mathcal{P}(X)}(U,-)$ of $U\cap-$ may be thought of as a function type [U,V].

²² Further Notation: Also written $\mathbf{Hom}_{\mathcal{P}(X)}(U,V)$.

- 2. Under the Curry–Howard correspondence (??), the function type [U, V] corresponds to implication $U \Rightarrow V$.
- 3. Implication $U \Rightarrow V$ is logically equivalent to $\neg U \lor V$.
- 4. The expression $\neg U \lor V$ then corresponds to the set $U^{c} \cup V$ in $\mathcal{P}(X)$.
- 5. The set $U^{c} \vee V$ turns out to indeed be the internal Hom of $\mathcal{P}(X)$.

Proposition 4.4.7.1.3. Let X be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto \operatorname{Hom}_{\mathcal{P}(X)}$ define functors

$$\begin{split} [U,-]_X\colon & (\mathcal{P}(X),\supset) & \to (\mathcal{P}(X),\subset), \\ [-,V]_X\colon & (\mathcal{P}(X),\subset) & \to (\mathcal{P}(X),\subset), \\ [-_1,-_2]_X\colon (\mathcal{P}(X)\times\mathcal{P}(X),\subset\times\supset) & \to (\mathcal{P}(X),\subset). \end{split}$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $[A, V]_X \subset [U, V]_X$.
- (b) If $V \subset B$, then $[U, V]_X \subset [U, B]_X$.
- (c) If $U \subset A$ and $V \subset B$, then $[A, V]_X \subset [U, B]_X$.
- 2. Adjointness. We have adjunctions

$$(U \cap - + [U, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{U \cap -}{\bot}} \mathcal{P}(X),$$

$$(- \cap V + [V, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{- \cap V}{\bot}} \mathcal{P}(X),$$

witnessed by bijections

$$\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, [V, W]_X),$$

 $\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(V, [U, W]_X).$

In particular, the following statements hold for each $U, V, W \in \mathcal{P}(X)$:

- (a) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $U \subset [V, W]_X$.
- (b) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $V \subset [U, W]_X$.
- 3. Interaction With the Empty Set I. We have

$$[U,\emptyset]_X=U^{\mathsf{c}},$$

$$[\emptyset, V]_X = X,$$

natural in $U, V \in \mathcal{P}(X)$.

4. Interaction With X. We have

$$[U,X]_X = X,$$

$$[X, V]_X = V,$$

natural in $U, V \in \mathcal{P}(X)$.

5. Interaction With the Empty Set II. The functor

$$D_X \colon \mathcal{P}(X)^{\mathsf{op}} \to \mathcal{P}(X)$$

defined by

$$D_X \stackrel{\text{def}}{=} [-, \emptyset]_X$$
$$= (-)^{\mathsf{c}}$$

is an involutory isomorphism of categories, making Ø into a dualising object for $(\mathcal{P}(X), \cap, X, [-, -]_X)$ in the sense of **??**. In particular:

(a) The diagram

$$\mathcal{P}(X)^{\operatorname{op}} \xrightarrow{D_X} \mathcal{P}(X)$$

$$\operatorname{id}_{\mathcal{P}(X)^{\operatorname{op}}} \qquad \downarrow^{D_X}$$

$$\mathcal{P}(X)^{\operatorname{op}}$$

commutes, i.e. we have

$$\underbrace{D_X(D_X(U))}_{\stackrel{\text{def}}{=}[[U,\emptyset]_X,\emptyset]_X} = U$$

for each $U \in \mathcal{P}(X)$.

(b) The diagram

$$\mathcal{P}(X)^{\operatorname{op}} \times \mathcal{P}(X)^{\operatorname{op}} \xrightarrow{\cap^{\operatorname{op}}} \mathcal{P}(X)^{\operatorname{op}}$$

$$\operatorname{id}_{\mathcal{P}(X)^{\operatorname{op}} \times D_X} / \qquad D_X$$

$$\mathcal{P}(X)^{\operatorname{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X)$$

commutes, i.e. we have

$$\underbrace{D_X(U\cap D_X(V))}_{\stackrel{\text{def}}{=}[U\cap [V,\emptyset]_X,\emptyset]_X}=[U,V]_X$$

for each $U, V \in \mathcal{P}(X)$.

- 6. *Interaction With the Empty Set III.* Let $f: X \to Y$ be a function.
 - (a) Interaction With Direct Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\text{op}} & \xrightarrow{f_*^{\text{op}}} \mathcal{P}(Y)^{\text{op}} \\
\downarrow^{D_X} & & \downarrow^{D_Y} \\
\mathcal{P}(X) & \xrightarrow{f_!} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_!(D_X(U)) = D_Y(f_*(U))$$

for each $U \in \mathcal{P}(X)$.

(b) Interaction With Inverse Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(Y)^{\mathsf{op}} & \xrightarrow{f^{-1,\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}} \\
D_{Y} & & \downarrow D_{X} \\
\mathcal{P}(Y) & \xrightarrow{f^{-1}} \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$f^{-1}(D_Y(U)) = D_X(f^{-1}(U))$$

for each $U \in \mathcal{P}(X)$.

(c) Interaction With Codirect Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_!^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
D_X & & \downarrow D_Y \\
\mathcal{P}(X) & \xrightarrow{f_*} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(D_X(U)) = D_Y(f_!(U))$$

for each $U \in \mathcal{P}(X)$.

7. Interaction With Unions of Families of Subsets I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\operatorname{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\cup^{\operatorname{op}} \times \cup^{\operatorname{op}} \qquad \qquad \bigcup \cup$$

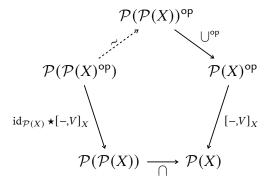
$$\mathcal{P}(X)^{\operatorname{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_{Y}} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcup_{U \in \mathcal{U}} U, \bigcup_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

8. Interaction With Unions of Families of Subsets II. The diagram



commutes, i.e. we have

$$\left[\bigcup_{U \in \mathcal{U}} U, V\right]_X = \bigcap_{U \in \mathcal{U}} [U, V]_X$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

9. Interaction With Unions of Families of Subsets III. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{\bigcup} \mathcal{P}(X)$$

$$\mathrm{id}_{\mathcal{P}(X)} \star [U,-]_X \downarrow \qquad \qquad \downarrow [U,-]_X$$

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{\prod} \mathcal{P}(X)$$

commutes, i.e. we have

$$\left[U, \bigcup_{V \in \mathcal{V}} V\right]_X = \bigcup_{V \in \mathcal{V}} [U, V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

10. Interaction With Intersections of Families of Subsets I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\operatorname{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\uparrow^{\operatorname{op}} \times \cap^{\operatorname{op}} \downarrow \qquad \qquad \downarrow \cap$$

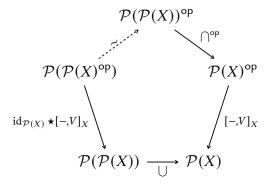
$$\mathcal{P}(X)^{\operatorname{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcap_{U \in \mathcal{U}} U, \bigcap_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

11. Interaction With Intersections of Families of Subsets II. The diagram



commutes, i.e. we have

$$\left[\bigcap_{U\in\mathcal{U}}U,V\right]_X=\bigcup_{U\in\mathcal{U}}[U,V]_X$$

for each $U \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

12. Interaction With Intersections of Families of Subsets III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcap} \mathcal{P}(X) \\
\downarrow^{[U,-]_X} & & \downarrow^{[U,-]_X} \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcap} \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U, \bigcap_{V \in \mathcal{V}} V\right]_X = \bigcap_{V \in \mathcal{V}} [U, V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

13. Interaction With Binary Unions. We have equalities of sets

$$[U \cap V, W]_X = [U, W]_X \cup [V, W]_X,$$

 $[U, V \cap W]_X = [U, V]_X \cap [U, W]_X$

for each $U, V, W \in \mathcal{P}(X)$.

14. Interaction With Binary Intersections. We have equalities of sets

$$[U \cup V, W]_X = [U, W]_X \cap [V, W]_X,$$

 $[U, V \cup W]_X = [U, V]_X \cup [U, W]_X$

for each $U, V, W \in \mathcal{P}(X)$.

15. Interaction With Differences. We have equalities of sets

$$[U \setminus V, W]_X = [U, W]_X \cup [V^{c}, W]_X$$
$$= [U, W]_X \cup [U, V]_X,$$
$$[U, V \setminus W]_X = [U, V]_X \setminus (U \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

16. Interaction With Complements. We have equalities of sets

$$[U^{c}, V]_{X} = U \cup V,$$

$$[U, V^{c}]_{X} = U \cap V,$$

$$[U, V]_{X}^{c} = U \setminus V$$

for each $U, V \in \mathcal{P}(X)$.

17. Interaction With Characteristic Functions. We have

$$\chi_{[U,V]_{\mathcal{P}(X)}}(x) = \max(1 - \chi_U \pmod{2}, \chi_V)$$
 for each $U, V \in \mathcal{P}(X)$.

18. *Interaction With Direct Images*. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(X)^{\text{op}} \times \mathcal{P}(X) \xrightarrow{f_*^{\text{op}} \times f_!} \mathcal{P}(Y)^{\text{op}} \times \mathcal{P}(Y) \\
 \downarrow [-_1, -_2]_X \downarrow \qquad \qquad \downarrow [-_1, -_2]_Y \\
 \mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have an equality of sets

$$f_!([U,V]_X) = [f_*(U), f_!(V)]_Y,$$

natural in $U, V \in \mathcal{P}(X)$.

19. Interaction With Inverse Images. Let $f\colon X\to Y$ be a function. The diagram

$$\mathcal{P}(Y)^{\text{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1,\text{op}} \times f^{-1}} \mathcal{P}(X)^{\text{op}} \times \mathcal{P}(X) \\
\downarrow [-1,-2]_{X} \\
\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_X,$$

natural in $U, V \in \mathcal{P}(X)$.

20. *Interaction With Codirect Images.* Let $f: X \to Y$ be a function. We have a natural transformation

with components

$$[f_{!}(U), f_{*}(V)]_{V} \subset f_{*}([U, V]_{X})$$

indexed by $U, V \in \mathcal{P}(X)$.

Proof. Item 1, Functoriality: Since $\mathcal{P}(X)$ is posetal, it suffices to prove Items 1a to 1c.

1. Proof of Item 1a: We have

$$[A, V]_X \stackrel{\text{def}}{=} A^{\mathsf{c}} \cup V$$

$$\subset U^{\mathsf{c}} \cup V$$

$$\stackrel{\text{def}}{=} [U, V]_X,$$

where we have used:

- (a) Item 1 of Definition 4.3.11.1.2, which states that if $U \subset A$, then $A^{c} \subset U^{c}$.
- (b) Item 1a of Item 1 of Definition 4.3.11.1.2, which states that if $A^c \subset U^c$, then $A^c \cup K \subset U^c \cup K$ for any $K \in \mathcal{P}(X)$.
- 2. Proof of Item 1b: We have

$$[U, V]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup V$$
$$\subset U^{\mathsf{c}} \cup B$$
$$\stackrel{\text{def}}{=} [U, B]_X,$$

where we have used Item 1b of Item 1 of Definition 4.3.11.1.2, which states that if $V \subset B$, then $K \cup V \subset K \cup B$ for any $K \in \mathcal{P}(X)$.

3. Proof of Item 1c: We have

$$[A, V]_X \subset [U, V]_X$$
$$\subset [U, B]_X,$$

where we have used Items 1a and 1b.

This finishes the proof.

Item 2, Adjointness: This is a repetition of Item 2 of Definition 4.3.9.1.2 and is proved there.

Item 3, Interaction With the Empty Set I: We have

$$[U, \emptyset]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup \emptyset$$
$$= U^{\mathsf{c}},$$

where we have used Item 3 of Definition 4.3.8.1.2, and we have

$$[\emptyset, V]_X \stackrel{\text{def}}{=} \emptyset^{\mathsf{c}} \cup V$$

$$\stackrel{\text{def}}{=} (X \setminus \emptyset) \cup V$$

$$= X \cup V$$

$$= X,$$

where we have used:

- 1. Item 12 of Definition 4.3.10.1.2 for the first equality.
- 2. Item 5 of Definition 4.3.8.1.2 for the last equality.

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2).

Item 4, Interaction With X: We have

$$[U,X]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup X$$
$$= X.$$

where we have used Item 5 of Definition 4.3.8.1.2, and we have

$$\begin{split} [X,V]_X &\stackrel{\text{def}}{=} X^\mathsf{c} \cup V \\ &\stackrel{\text{def}}{=} (X \setminus X) \cup V \\ &= \emptyset \cup V \\ &= V, \end{split}$$

where we have used Item 3 of Definition 4.3.8.1.2 for the last equality. Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). Item 5, Interaction With the Empty Set II: We have

$$D_X(D_X(U)) \stackrel{\text{def}}{=} [[U, \emptyset]_X, \emptyset]_X$$
$$= [U^c, \emptyset]_X$$
$$= (U^c)^c$$
$$= U,$$

where we have used:

- 1. Item 3 for the second and third equalities.
- 2. Item 3 of Definition 4.3.11.1.2 for the fourth equality.

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2), and thus we have

$$[[-,\emptyset]_X,\emptyset]_X \cong \mathrm{id}_{\mathcal{P}(X)}$$

This finishes the proof.

Item 6, Interaction With the Empty Set III: Since $D_X = (-)^c$, this is essentially a repetition of the corresponding results for $(-)^c$, namely Items 5 to 7 of Definition 4.3.11.1.2.

Item 7, Interaction With Unions of Families of Subsets I: By Item 3 of Definition 4.4.7.1.3, we have

$$[\mathcal{U}, \emptyset]_{\mathcal{P}(X)} = \mathcal{U}^{\mathsf{c}},$$

 $[\mathcal{U}, \emptyset]_X = \mathcal{U}^{\mathsf{c}}.$

With this, the counterexample given in the proof of Item 10 of Definition 4.3.6.1.2 then applies.

Item 8, Interaction With Unions of Families of Subsets II: We have

$$\left[\bigcup_{U \in \mathcal{U}} U, V\right]_{X} \stackrel{\text{def}}{=} \left(\bigcup_{U \in \mathcal{U}} U\right)^{c} \cup V$$

$$= \left(\bigcap_{U \in \mathcal{U}} U^{c}\right) \cup V$$

$$= \bigcap_{U \in \mathcal{U}} (U^{c} \cup V)$$

$$\stackrel{\text{def}}{=} \bigcap_{U \in \mathcal{U}} [U, V]_{X},$$

where we have used:

- 1. Item 11 of Definition 4.3.6.1.2 for the second equality.
- 2. Item 6 of Definition 4.3.7.1.2 for the third equality.

This finishes the proof.

Item 9, Interaction With Unions of Families of Subsets III: We have

$$\bigcup_{V \in \mathcal{V}} [U, V]_X \stackrel{\text{def}}{=} \bigcup_{V \in \mathcal{V}} (U^{\mathsf{c}} \cup V)$$

$$= U^{\mathsf{c}} \cup \left(\bigcup_{V \in \mathcal{V}} V\right)$$

$$\stackrel{\mathsf{def}}{=} \left[U, \bigcup_{V \in \mathcal{V}} V\right]_{V}.$$

where we have used Item 6. This finishes the proof.

Item 10, *Interaction With Intersections of Families of Subsets I*: Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}$, and let $\mathcal{V} = \{\{0\}, \{0, 1\}\}$. We have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W = \bigcap_{W \in \mathcal{P}(X)} W$$
$$= \{0, 1\},$$

whereas

$$\left[\bigcap_{U\in\mathcal{U}}U,\bigcap_{V\in\mathcal{V}}V\right]_X = \left[\{0,1\},\{0\}\right]$$
$$= \{0\},$$

Thus we have

$$\bigcap_{W\in [\mathcal{U},\mathcal{V}]_{\mathcal{P}(X)}}W=\left\{0,1\right\}\neq\left\{0\right\}=\left[\bigcap_{U\in\mathcal{U}}U,\bigcap_{V\in\mathcal{V}}V\right]_X.$$

This finishes the proof.

Item 11, Interaction With Intersections of Families of Subsets II: We have

$$\left[\bigcap_{U \in \mathcal{U}} U, V\right]_{X} \stackrel{\text{def}}{=} \left(\bigcap_{U \in \mathcal{U}} U\right)^{c} \cup V$$

$$= \left(\bigcup_{U \in \mathcal{U}} U^{c}\right) \cup V$$

$$= \bigcup_{U \in \mathcal{U}} (U^{c} \cup V)$$

$$\stackrel{\text{def}}{=} \bigcup_{U \in \mathcal{U}} [U, V]_{X},$$

where we have used:

- 1. Item 12 of Definition 4.3.6.1.2 for the second equality.
- 2. Item 6 of Definition 4.3.7.1.2 for the third equality.

This finishes the proof.

Item 12, Interaction With Intersections of Families of Subsets III: We have

$$\bigcap_{V \in \mathcal{V}} [U, V]_X \stackrel{\text{def}}{=} \bigcap_{V \in \mathcal{V}} (U^{\mathsf{c}} \cup V)$$

$$= U^{\mathsf{c}} \cup \left(\bigcap_{V \in \mathcal{V}} V\right)$$

$$\stackrel{\text{def}}{=} \left[U, \bigcap_{V \in \mathcal{V}} V\right]_Y$$

where we have used Item 6. This finishes the proof. *Item 13, Interaction With Binary Unions*: We have

$$\begin{split} [U \cap V, W]_X &\stackrel{\text{def}}{=} (U \cap V)^{\text{c}} \cup W \\ &= (U^{\text{c}} \cup V^{\text{c}}) \cup W \\ &= (U^{\text{c}} \cup V^{\text{c}}) \cup (W \cup W) \\ &= (U^{\text{c}} \cup W) \cup (V^{\text{c}} \cup W) \\ &\stackrel{\text{def}}{=} [U, W]_X \cup [V, W]_X, \end{split}$$

where we have used:

- 1. Item 2 of Definition 4.3.11.1.2 for the second equality.
- 2. Item 8 of Definition 4.3.8.1.2 for the third equality.
- 3. Several applications of Items 2 and 4 of Definition 4.3.8.1.2 and for the fourth equality.

For the second equality in the statement, we have

$$[U, V \cap W]_X \stackrel{\text{def}}{=} U^{c} \cup (V \cap W)$$
$$= (U^{c} \cup V) \cap (U^{c} \cap W)$$
$$\stackrel{\text{def}}{=} [U, V]_X \cap [U, W]_X,$$

where we have used Item 6 of Definition 4.3.8.1.2 for the second equality. *Item 14, Interaction With Binary Intersections*: We have

$$[U \cup V, W]_X \stackrel{\text{def}}{=} (U \cup V)^{c} \cup W$$
$$= (U^{c} \cap V^{c}) \cup W$$
$$= (U^{c} \cup W) \cap (V^{c} \cup W)$$
$$\stackrel{\text{def}}{=} [U, W]_X \cap [V, W]_X,$$

where we have used:

- 1. Item 2 of Definition 4.3.11.1.2 for the second equality.
- 2. Item 6 of Definition 4.3.8.1.2 for the third equality.

Now, for the second equality in the statement, we have

$$\begin{split} [U, V \cup W]_X &\stackrel{\text{def}}{=} U^{\mathsf{c}} \cup (V \cup W) \\ &= (U^{\mathsf{c}} \cup U^{\mathsf{c}}) \cup (V \cup W) \\ &= (U^{\mathsf{c}} \cup V) \cup (U^{\mathsf{c}} \cup W) \\ &\stackrel{\text{def}}{=} [U, V]_X \cup [U, W]_X, \end{split}$$

where we have used:

- 1. Item 8 of Definition 4.3.8.1.2 for the second equality.
- 2. Several applications of Items 2 and 4 of Definition 4.3.8.1.2 and for the third equality.

This finishes the proof.

Item 15, Interaction With Differences: We have

$$\begin{split} [U \setminus V, W]_X &\stackrel{\mathrm{def}}{=} (U \setminus V)^{\mathsf{c}} \cup W \\ &\stackrel{\mathrm{def}}{=} (X \setminus (U \setminus V)) \cup W \\ &= ((X \cap V) \cup (X \setminus U)) \cup W \\ &= (V \cup (X \setminus U)) \cup W \\ &\stackrel{\mathrm{def}}{=} (V \cup U^{\mathsf{c}}) \cup W \\ &= (V \cup (U^{\mathsf{c}} \cup U^{\mathsf{c}})) \cup W \\ &= (U^{\mathsf{c}} \cup W) \cup (U^{\mathsf{c}} \cup V) \\ &\stackrel{\mathrm{def}}{=} [U, W]_X \cup [U, V]_X, \end{split}$$

where we have used:

- 1. Item 10 of Definition 4.3.10.1.2 for the third equality.
- 2. Item 4 of Definition 4.3.9.1.2 for the fourth equality.
- 3. Item 8 of Definition 4.3.8.1.2 for the sixth equality.
- 4. Several applications of Items 2 and 4 of Definition 4.3.8.1.2 and for the seventh equality.

We also have

$$\begin{split} [U \setminus V, W]_X &\stackrel{\mathrm{def}}{=} (U \setminus V)^{\mathsf{c}} \cup W \\ &\stackrel{\mathrm{def}}{=} (X \setminus (U \setminus V)) \cup W \\ &= ((X \cap V) \cup (X \setminus U)) \cup W \\ &= (V \cup (X \setminus U)) \cup W \\ &\stackrel{\mathrm{def}}{=} (V \cup U^{\mathsf{c}}) \cup W \\ &= (V \cup U^{\mathsf{c}}) \cup (W \cup W) \\ &= (U^{\mathsf{c}} \cup W) \cup (V \cup W) \\ &= (U^{\mathsf{c}} \cup W) \cup ((V^{\mathsf{c}})^{\mathsf{c}} \cup W) \\ &\stackrel{\mathrm{def}}{=} [U, W]_X \cup [V^{\mathsf{c}}, W]_X, \end{split}$$

where we have used:

- 1. Item 10 of Definition 4.3.10.1.2 for the third equality.
- 2. Item 4 of Definition 4.3.9.1.2 for the fourth equality.
- 3. Item 8 of Definition 4.3.8.1.2 for the sixth equality.
- 4. Several applications of Items 2 and 4 of Definition 4.3.8.1.2 and for the seventh equality.
- 5. Item 3 of Definition 4.3.11.1.2 for the eighth equality.

Now, for the second equality in the statement, we have

$$\begin{split} [U, V \setminus W]_X &\stackrel{\text{def}}{=} U^{\mathsf{c}} \cup (V \setminus W) \\ &= (V \setminus W) \cup U^{\mathsf{c}} \\ &= (V \cup U^{\mathsf{c}}) \setminus (W \setminus U^{\mathsf{c}}) \end{split}$$

$$\stackrel{\text{def}}{=} (V \cup U^{c}) \setminus (W \setminus (X \setminus U))$$

$$= (V \cup U^{c}) \setminus ((W \cap U) \cup (W \setminus X))$$

$$= (V \cup U^{c}) \setminus ((W \cap U) \cup \emptyset)$$

$$= (V \cup U^{c}) \setminus (W \cap U)$$

$$= (V \cup U^{c}) \setminus (U \cap W)$$

$$\stackrel{\text{def}}{=} [U, V]_{Y} \setminus (U \cap W)$$

where we have used:

- 1. Item 4 of Definition 4.3.8.1.2 for the second equality.
- 2. Item 4 of Definition 4.3.10.1.2 for the third equality.
- 3. Item 10 of Definition 4.3.10.1.2 for the fifth equality.
- 4. Item 13 of Definition 4.3.10.1.2 for the sixth equality.
- 5. Item 3 of Definition 4.3.8.1.2 for the seventh equality.
- 6. Item 5 of Definition 4.3.9.1.2 for the eighth equality.

This finishes the proof.

Item 16, Interaction With Complements: We have

$$[U^{c}, V]_{X} \stackrel{\text{def}}{=} (U^{c})^{c} \cup V,$$
$$= U \cup V,$$

where we have used Item 3 of Definition 4.3.11.1.2. We also have

$$[U, V^{\mathsf{c}}]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup V^{\mathsf{c}}$$
$$= U \cap V$$

where we have used Item 2 of Definition 4.3.11.1.2. Finally, we have

$$[U, V]_X^{c} = ((U \setminus V)^{c})^{c}$$
$$= U \setminus V,$$

where we have used Item 2 of Definition 4.3.11.1.2.

Item 17, Interaction With Characteristic Functions: We have

$$\chi_{[U,V]_{\mathcal{P}(X)}}(x) \stackrel{\text{def}}{=} \chi_{U^{\mathsf{c}} \cup V}(x)$$

$$= \max(\chi_{U^{\mathsf{c}}}, \chi_{V})$$

$$= \max(1 - \chi_{U} \pmod{2}, \chi_{V}),$$

where we have used:

- 1. Item 10 of Definition 4.3.8.1.2 for the second equality.
- 2. Item 4 of Definition 4.3.11.1.2 for the third equality.

This finishes the proof.

Item 18, Interaction With Direct Images: This is a repetition of Item 10 of Definition 4.6.1.1.5 and is proved there.

Item 19, Interaction With Inverse Images: This is a repetition of Item 10 of Definition 4.6.2.1.3 and is proved there.

Item 20, *Interaction With Codirect Images*: This is a repetition of Item 9 of Definition 4.6.3.1.7 and is proved there. □

4.4.8 Isbell Duality for Sets

Let *X* be a set.

Definition 4.4.8.1.1. The **Isbell function** of X is the map

$$I: \mathcal{P}(X) \to \mathsf{Sets}(X, \mathcal{P}(X))$$

defined by

$$I(U) \stackrel{\text{def}}{=} \llbracket x \mapsto \llbracket U, \{x\} \rrbracket_X \rrbracket$$

for each $U \in \mathcal{P}(X)$.

Remark 4.4.8.1.2. Recall from Definition 4.4.1.1.2 that we may view the powerset $\mathcal{P}(X)$ of a set X as the decategorification of the category of presheaves $\mathsf{PSh}(C)$ of a category C. Building upon this analogy, we want to mimic the definition of the Isbell Spec functor, which is given on objects by

$$\mathsf{Spec}(\mathcal{F}) \stackrel{\mathsf{def}}{=} \mathsf{Nat}(\mathcal{F}, h_{(-)})$$

for each $\mathcal{F} \in \mathsf{Obj}(\mathsf{PSh}(C))$. To this end, we could define

$$\mathsf{I}(U) \stackrel{\mathrm{def}}{=} \left[U, \chi_{(-)} \right]_X,$$

replacing:

- The Yoneda embedding $X \mapsto h_X$ of C into PSh(C) with the characteristic embedding $x \mapsto \chi_x$ of X into $\mathcal{P}(X)$ of Definition 4.5.4.1.1.
- The internal Hom Nat of PSh(C) with the internal Hom $[-,-]_X$ of $\mathcal{P}(X)$ of Definition 4.4.7.1.1.

However, since $[U, \chi_x]_X$ is a subset of U instead of a truth value, we get a function

$$I: \mathcal{P}(X) \to \mathsf{Sets}(X, \mathcal{P}(X))$$

instead of a function

$$I: \mathcal{P}(X) \to \mathcal{P}(X).$$

This makes some of the properties involving I a bit more cumbersome to state, although we still have an analogue of Isbell duality in that $I_! \circ I$ evaluates to $id_{\mathcal{P}(X)}$ in the sense of Definition 4.4.8.1.3.

Proposition 4.4.8.1.3. The diagram

$$\mathcal{P}(X) \xrightarrow{\mathsf{I}} \mathsf{Sets}(X, \mathcal{P}(X))$$

$$\Delta_{\Delta_{\mathsf{id}_{\mathcal{P}}(X)}} \qquad \qquad \mathsf{I}_{!}$$

$$\mathsf{Sets}(X, \mathsf{Sets}(X, \mathcal{P}(X)))$$

commutes, i.e. we have

$$I_!(I(U)) = \llbracket x \mapsto \llbracket y \mapsto U \rrbracket \rrbracket$$

for each $U \in \mathcal{P}(X)$.

Proof. We have

$$I_{!}(I(U)) \stackrel{\text{def}}{=} I_{!}([\![x \mapsto U^{c} \cup \{x\}]\!])$$

$$\stackrel{\text{def}}{=} [\![x \mapsto I(U^{c} \cup \{x\})]\!]$$

$$\stackrel{\text{def}}{=} [\![x \mapsto [\![y \mapsto (U^{c} \cup \{x\})^{c} \cup \{x\}]\!]]\!]$$

$$= [\![x \mapsto [\![y \mapsto (U \cap (X \setminus \{x\})) \cup \{x\}]\!]]\!]$$

$$= [\![x \mapsto [\![y \mapsto (U \setminus \{x\}) \cup \{x\}]\!]]\!]$$

$$= [\![x \mapsto [\![y \mapsto U]\!]]\!],$$

where we have used Item 2 of Definition 4.3.11.1.2 for the fourth equality above.

4.5 Characteristic Functions

4.5.1 The Characteristic Function of a Subset

Let X be a set and let $U \in \mathcal{P}(X)$.

Definition 4.5.1.1.1. The **characteristic function of** U^{23} is the function $\chi_U: X \to \{t, f\}^{24}$ defined by

$$\chi_U(x) \stackrel{\text{def}}{=} \begin{cases} \mathsf{true} & \text{if } x \in U, \\ \mathsf{false} & \text{if } x \notin U \end{cases}$$

for each $x \in X$.

Remark 4.5.1.1.2. Under the analogy that $\{t, f\}$ should be the (-1)-categorical analogue of Sets, we may view a function

$$f: X \to \{\mathsf{t}, \mathsf{f}\}$$

as a decategorification of presheaves and copresheaves

$$\mathcal{F} \colon C^{\mathsf{op}} \to \mathsf{Sets},$$

 $F \colon C \to \mathsf{Sets}.$

The characteristic functions χ_U of the subsets of X are then the primordial examples of such functions (and, in fact, all of them).

Notation 4.5.1.1.3. We will often employ the bijection $\{t, f\} \cong \{0, 1\}$ to make use of the arithmetical operations defined on $\{0, 1\}$ when disucssing characteristic functions.

Examples of this include Items 4 to 11 of Definition 4.5.1.1.4 below.

Proposition 4.5.1.1.4. Let *X* be a set.

1. Functionality. The assignment $U \mapsto \chi_U$ defines a function

$$\chi_{(-)} \colon \mathcal{P}(X) \to \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\}).$$

2. *Bijectivity*. The function $\chi_{(-)}$ from Item 1 is bijective.

²³ Further Terminology: Also called the **indicator function of** U.

²⁴ Further Notation: Also written $\chi_X(U, -)$ or $\chi_X(-, U)$.

3. Naturality. The collection

$$\left\{\chi_{(-)}\colon \mathcal{P}(X) \to \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})\right\}_{X \in \mathsf{Obi}(\mathsf{Sets})}$$

defines a natural isomorphism between \mathcal{P}^{-1} and $\mathsf{Sets}(-, \{\mathsf{t}, \mathsf{f}\})$. In particular, given a function $f \colon X \to Y$, the diagram

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

$$\chi_{(-)} \downarrow \chi \qquad \qquad \downarrow \chi_{(-)} \downarrow \chi_{(-)}$$

$$\mathsf{Sets}(Y, \{\mathsf{t}, \mathsf{f}\}) \xrightarrow{f^*} \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

commutes, i.e. we have

$$\chi_V \circ f = \chi_{f^{-1}(V)}$$

for each $V \in \mathcal{P}(Y)$.

4. Interaction With Unions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

5. Interaction With Unions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each $U, V \in \mathcal{P}(X)$.

6. Interaction With Intersections I. We have

$$\chi_{U\cap V} = \chi_U \chi_V$$

for each $U, V \in \mathcal{P}(X)$.

7. Interaction With Intersections II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

8. Interaction With Differences. We have

$$\chi_{U\setminus V}=\chi_U-\chi_{U\cap V}$$

for each $U, V \in \mathcal{P}(X)$.

9. Interaction With Complements. We have

$$\chi_{U^c} \equiv 1 - \chi_U \pmod{2}$$

for each $U \in \mathcal{P}(X)$.

10. Interaction With Symmetric Differences. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \pmod{2}$$

for each $U, V \in \mathcal{P}(X)$.

11. Interaction With Internal Homs. We have

$$\chi_{[U,V]_{\mathcal{P}(X)}} = \max(1 - \chi_U \pmod{2}, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

Proof. Item 1, Functionality: There is nothing to prove. *Item 2, Bijectivity*: We proceed in three steps:

1. The Inverse of $\chi_{(-)}$. The inverse of $\chi_{(-)}$ is the map

$$\Phi \colon \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\}) \xrightarrow{\sim} \mathcal{P}(X),$$

defined by

$$\begin{split} \Phi(f) &\stackrel{\text{def}}{=} U_f \\ &\stackrel{\text{def}}{=} f^{-1}(\mathsf{true}) \\ &\stackrel{\text{def}}{=} \{x \in X \,|\, f(x) = \mathsf{true}\} \end{split}$$

for each $f \in \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$.

2. Invertibility I. We have

$$\begin{split} \big[\Phi \circ \chi_{(-)}\big](U) &\stackrel{\text{def}}{=} \Phi(\chi_U) \\ &\stackrel{\text{def}}{=} \chi_U^{-1}(\mathsf{true}) \\ &\stackrel{\text{def}}{=} \{x \in X \mid \chi_U(x) = \mathsf{true}\} \\ &\stackrel{\text{def}}{=} \{x \in X \mid x \in U\} \\ &= U \\ &\stackrel{\text{def}}{=} \big[\mathsf{id}_{\mathcal{P}(X)}\big](U) \end{split}$$

for each $U \in \mathcal{P}(X)$. Thus, we have

$$\Phi \circ \chi_{(-)} = \mathrm{id}_{\mathcal{P}(X)} .$$

3. Invertibility II. We have

$$\begin{split} \left[\chi_{(-)} \circ \Phi\right](U) &\stackrel{\text{def}}{=} \chi_{\Phi(f)} \\ &\stackrel{\text{def}}{=} \chi_{f^{-1}(\mathsf{true})} \\ &\stackrel{\text{def}}{=} \left[\!\!\left[x \mapsto \begin{cases} \mathsf{true} & \text{if } x \in f^{-1}(\mathsf{true}) \\ \mathsf{false} & \text{otherwise} \end{cases} \right] \\ &= \left[\!\!\left[x \mapsto f(x) \right]\!\!\right] \\ &= f \\ &\stackrel{\text{def}}{=} \left[\mathrm{id}_{\mathsf{Sets}(X, \{\mathsf{t}, f\})} \right](f) \end{split}$$

for each $f \in Sets(X, \{t, f\})$. Thus, we have

$$\chi_{(-)} \circ \Phi = \mathrm{id}_{\mathrm{Sets}(X,\{\mathsf{t},\mathsf{f}\})}$$
.

This finishes the proof.

Item 3, Naturality: We proceed in two steps:

1. *Naturality of* $\chi_{(-)}$. We have

$$[\chi_V \circ f](v) \stackrel{\text{def}}{=} \chi_V(f(v))$$

$$= \begin{cases} \text{true} & \text{if } f(v) \in V, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } v \in f^{-1}(V), \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_{f^{-1}(V)}(v)$$

for each $v \in V$.

2. Naturality of Φ . Since $\chi_{(-)}$ is natural and a componentwise inverse to Φ , it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Φ is also natural in each argument.

This finishes the proof.

Item 4, Interaction With Unions I: This is a repetition of Item 10 of Definition 4.3.8.1.2 and is proved there.

Item 5, Interaction With Unions II: This is a repetition of Item 11 of Definition 4.3.8.1.2 and is proved there.

Item 6, Interaction With Intersections I: This is a repetition of Item 10 of Definition 4.3.9.1.2 and is proved there.

Item 7, Interaction With Intersections II: This is a repetition of Item 11 of Definition 4.3.9.1.2 and is proved there.

Item 8, Interaction With Differences: This is a repetition of Item 16 of Definition 4.3.10.1.2 and is proved there.

Item 9, Interaction With Complements: This is a repetition of Item 4 of Definition 4.3.11.1.2 and is proved there.

Item 10, *Interaction With Symmetric Differences*: This is a repetition of Item 15 of Definition 4.3.12.1.2 and is proved there.

Item 11, Interaction With Internal Homs: This is a repetition of Item 17 of Definition 4.4.7.1.3 and is proved there. □

Remark 4.5.1.1.5. The bijection

$$\mathcal{P}(X) \cong \operatorname{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

of Item 2 of Definition 4.5.1.1.4, which

- Takes a subset $U \hookrightarrow X$ of X and *straightens* it to a function $\chi_U \colon X \to \{\text{true}, \text{false}\};$
- Takes a function $f: X \to \{\text{true}, \text{false}\}\$ and *unstraightens* it to a subset $f^{-1}(\text{true}) \hookrightarrow X \text{ of } X;$

may be viewed as the (-1)-categorical version of the o-categorical un/s-traightening isomorphism between indexed and fibred sets

$$\underbrace{\mathsf{FibSets}_X}_{\substack{\underline{\mathsf{def}}\\ = \mathsf{Sets}_{/X}}} \cong \underbrace{\mathsf{ISets}_X}_{\substack{\underline{\mathsf{def}}\\ = \mathsf{Fun}(X_{\mathsf{disc}},\mathsf{Sets})}}$$

of Un/Straightening for Indexed and Fibred Sets, ??. Here we view:

- Subsets $U \hookrightarrow X$ as being analogous to X-fibred sets $\phi_X \colon A \to X$.
- Functions $f: X \to \{\mathsf{t}, \mathsf{f}\}$ as being analogous to X-indexed sets $A: X_{\mathsf{disc}} \to \mathsf{Sets}$.

4.5.2 The Characteristic Function of a Point

Let X be a set and let $x \in X$.

Definition 4.5.2.1.1. The **characteristic function of** x is the function x

$$\gamma_x \colon X \to \{\mathsf{t},\mathsf{f}\}$$

defined by

$$\chi_x \stackrel{\mathrm{def}}{=} \chi_{\{x\}},$$

i.e. by

$$\chi_x(y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each $y \in X$.

Remark 4.5.2.1.2. Expanding upon Definition 4.5.1.1.2, we may think of the characteristic function

$$\gamma_x \colon X \to \{\mathsf{t},\mathsf{f}\}$$

of an element x of X as a decategorification of the representable presheaf and of the representable copresheaf

$$h_X \colon C^{\mathsf{op}} \to \mathsf{Sets},$$

 $h^X \colon C \to \mathsf{Sets}$

associated of an *object* X of a category C.

²⁵ *Further Notation:* Also written χ^x , $\chi_X(x, -)$, or $\chi_X(-, x)$.

4.5.3 The Characteristic Relation of a Set

Let *X* be a set.

Definition 4.5.3.1.1. The characteristic relation on X^{26} is the relation 27

$$\gamma_X(-1,-2): X \times X \to \{\mathsf{t},\mathsf{f}\}$$

on X defined by 28

$$\chi_X(x,y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each $x, y \in X$.

Remark 4.5.3.1.2. Expanding upon Definitions 4.5.1.1.2 and 4.5.2.1.2, we may view the characteristic relation

$$\chi_X(-1,-2): X \times X \to \{\mathsf{t},\mathsf{f}\}$$

of X as a decategorification of the Hom profunctor

$$\operatorname{Hom}_{\mathcal{C}}(-1,-2)\colon \mathcal{C}^{\operatorname{op}}\times\mathcal{C}\to\operatorname{Sets}$$

of a category C.

Proposition 4.5.3.1.3. Let $f: X \to Y$ be a function.

1. The Inclusion of Characteristic Relations Associated to a Function. Let $f: A \to B$ be a function. We have an inclusion 29

$$\chi_{B} \circ (f \times f) \subset \chi_{A}, \qquad A \times A \xrightarrow{f \times f} B \times B$$

$$\chi_{A} \searrow \chi_{A} \searrow \chi_{B}$$

$$\{t, f\}.$$

Proof. Item 1, *The Inclusion of Characteristic Relations Associated to a Function*: The inclusion $\chi_B(f(a), f(b)) \subset \chi_A(a, b)$ is equivalent to the statement "if a = b, then f(a) = f(b)", which is true.

²⁶ Further Terminology: Also called the **identity relation on** X.

²⁷ Further Notation: Also written χ_{-2}^{-1} , or \sim_{id} in the context of relations.

²⁸Under the bijection Sets($X \times X$, $\{t, f\}$) $\cong \mathcal{P}(X \times X)$ of Item 2 of Definition 4.5.1.1.4, the relation χ_X corresponds to the diagonal $\Delta_X \subset X \times X$ of X.

²⁹ Note: This is the 0-categorical version of Categories, Definition 11.5.4.1.1.

4.5.4 The Characteristic Embedding of a Set

Let *X* be a set.

Definition 4.5.4.1.1. The characteristic embedding³⁰ of X into $\mathcal{P}(X)$ is the function

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

defined by³¹

$$\chi_{(-)}(x) \stackrel{\text{def}}{=} \chi_x$$
$$= \{x\}$$

for each $x \in X$.

Remark 4.5.4.1.2. Expanding upon Definitions 4.5.1.1.2, 4.5.2.1.2 and 4.5.3.1.2, we may view the characteristic embedding

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

of X into $\mathcal{P}(X)$ as a decategorification of the Yoneda embedding

$$\sharp: C^{\mathsf{op}} \hookrightarrow \mathsf{PSh}(C)$$

of a category C into PSh(C).

Proposition 4.5.4.1.3. Let $f: X \to Y$ be a map of sets.

1. Interaction With Functions. We have

$$f_! \circ \chi_X = \chi_Y \circ f, \qquad \chi_X \downarrow \qquad \downarrow_{\chi_Y} \downarrow_{\chi_Y} \\ \mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y).$$

³⁰The name "characteristic *embedding*" is justified by Definition 4.5.5.1.2, which gives an analogue of fully faithfulness for $\chi_{(-)}$.

³¹Here we are identifying $\mathcal{P}(X)$ with Sets $(X, \{t, f\})$ as per Item 2 of Definition 4.5.1.1.4.

Proof. Item 1, Interaction With Functions: Indeed, we have

$$[f! \circ \chi_X](x) \stackrel{\text{def}}{=} f!(\chi_X(x))$$

$$\stackrel{\text{def}}{=} f!(\{x\})$$

$$= \{f(x)\}$$

$$\stackrel{\text{def}}{=} \chi_{X'}(f(x))$$

$$\stackrel{\text{def}}{=} [\chi_{X'} \circ f](x),$$

for each $x \in X$, showing the desired equality.

4.5.5 The Yoneda Lemma for Sets

Let *X* be a set and let $U \subset X$ be a subset of *X*.

Proposition 4.5.5.1.1. We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_U) = \chi_U(x)$$

for each $x \in X$, giving an equality of functions

$$\chi_{\mathcal{P}(X)}(\chi_{(-)},\chi_U)=\chi_U,$$

where

$$\chi_{\mathcal{P}(X)}(U,V) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } U \subset V, \\ \text{false} & \text{otherwise.} \end{cases}$$

Proof. We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_U) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } \{x\} \subset U, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } x \in U \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_U(x).$$

This finishes the proof.

Corollary 4.5.5.1.2. The characteristic embedding is fully faithful, i.e., we have

$$\chi_{\mathcal{P}(X)}(\chi_x,\chi_y)\cong\chi_X(x,y)$$

for each $x, y \in X$.

Proof. We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_y) = \chi_y(x)$$

$$\stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x \in \{y\} \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } x = y \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_X(x, y).$$

where we have used Definition 4.5.5.1.1 for the first equality.

4.6 The Adjoint Triple $f_! \dashv f^{-1} \dashv f_*$

4.6.1 Direct Images

Let $f: X \to Y$ be a function.

Definition 4.6.1.1.1. The direct image function associated to f is the function f is the function

$$f_1 \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by³³

$$f(U) \stackrel{\text{def}}{=} \left\{ y \in Y \middle| \begin{array}{l} \text{there exists some } x \in U \\ \text{such that } y = f(x) \end{array} \right\}$$
$$= \left\{ f(x) \in Y \middle| x \in U \right\}$$

for each $U \in \mathcal{P}(X)$.

Notation 4.6.1.1.2. Sometimes one finds the notation

$$\exists_f \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

for $f_!$. This notation comes from the fact that the following statements are equivalent, where $y \in Y$ and $U \in \mathcal{P}(X)$:

³² Further Notation: Also written simply $f: \mathcal{P}(X) \to \mathcal{P}(Y)$.

³³Further Terminology: The set f(U) is called the **direct image of** U **by** f.

- We have $y \in \exists_f(U)$.
- There exists some $x \in U$ such that f(x) = y.

We will not make use of this notation elsewhere in Clowder.

Warning 4.6.1.1.3. Notation for direct images between powersets is tricky:

- 1. Direct images for powersets and presheaves are both adjoint to their corresponding inverse image functors. However, the direct image functor for powersets is a *left* adjoint, while the direct image functor for presheaves is a *right* adjoint:
 - (a) *Powersets*. Given a function $f: X \to Y$, we have an inverse image functor

$$f^{-1} \colon \mathcal{P}(Y) \to \mathcal{P}(X).$$

The *left* adjoint of this functor is the usual direct image, defined above in Definition 4.6.1.1.1.

(b) *Presheaves*. Given a morphism of topological spaces $f: X \to Y$, we have an inverse image functor

$$f^{-1} \colon \mathsf{PSh}(Y) \to \mathsf{PSh}(X).$$

The *right* adjoint of this functor is the direct image functor of presheaves, defined in **??**.

- 2. The presheaf direct image functor is denoted f_* , but the direct image functor for powersets is denoted $f_!$ (as it's a left adjoint).
- 3. Adding to the confusion, it's somewhat common for $f_! \colon \mathcal{P}(X) \to \mathcal{P}(Y)$ to be denoted f_* .

We chose to write $f_!$ for the direct image to keep the notation aligned with the following similar adjoint situations:

Situation	Adjoint String
Functoriality of Powersets	$(f_! \dashv f^{-1} \dashv f_*) \colon \mathcal{P}(X) \xrightarrow{\rightleftarrows} \mathcal{P}(Y)$
Functoriality of Presheaf Categories	$(f_! \dashv f^{-1} \dashv f_*) \colon PSh(X) \xrightarrow{\rightleftarrows} PSh(Y)$
Base Change	$(f_!\dashv f^*\dashv f_*)\colon C_{/X}\stackrel{\rightleftarrows}{\to} C_{/Y}$
Kan Extensions	$(F_! \dashv F^* \dashv F_*) \colon \operatorname{Fun}(\mathcal{C}, \mathcal{E}) \xrightarrow{\rightleftarrows} \operatorname{Fun}(\mathcal{D}, \mathcal{E})$

Remark 4.6.1.1.4. Identifying $\mathcal{P}(X)$ with $\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ via Item 2 of Definition 4.5.1.1.4, we see that the direct image function associated to f is equivalently the function

$$f_! \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by

$$f_{!}(\chi_{U}) \stackrel{\text{def}}{=} \operatorname{Lan}_{f}(\chi_{U})$$

$$= \operatorname{colim}\left(\left(f \stackrel{\rightarrow}{\times} (-_{1})\right) \stackrel{\operatorname{pr}}{\twoheadrightarrow} A \stackrel{\chi_{U}}{\longrightarrow} \{\mathsf{t}, \mathsf{f}\}\right)$$

$$= \operatorname{colim}_{x \in X} (\chi_{U}(x))$$

$$f(x) = -_{1}$$

$$= \bigvee_{\substack{x \in X \\ f(x) = -_{1}}} (\chi_{U}(x)),$$

where we have used ?? for the second equality. In other words, we have

$$[f!(\chi_U)](y) = \bigvee_{\substack{x \in X \\ f(x) = y}} (\chi_U(x))$$

$$= \begin{cases} \text{true} & \text{if there exists some } x \in X \text{ such} \\ & \text{that } f(x) = y \text{ and } x \in U, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if there exists some } x \in U \\ & \text{such that } f(x) = y, \\ \text{false} & \text{otherwise} \end{cases}$$

for each $y \in Y$.

Proposition 4.6.1.1.5. Let $f: X \to Y$ be a function.

1. Functoriality. The assignment $U \mapsto f_!(U)$ defines a functor

$$f_! : (\mathcal{P}(X), \subset) \to (\mathcal{P}(Y), \subset).$$

In particular, for each $U, V \in \mathcal{P}(X)$, the following condition is satisfied:

$$(\star)$$
 If $U \subset V$, then $f_!(U) \subset f_!(V)$.

2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$id_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad id_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$

 $f_! \circ f^{-1} \hookrightarrow id_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow id_{\mathcal{P}(X)},$

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*(f^{-1}(V)),$$

 $f_!(f^{-1}(V)) \subset V, \qquad f^{-1}(f_*(U)) \subset U$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

(b) Bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(Y)}(f_!(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f^{-1}(V)),$$

 $\operatorname{Hom}_{\mathcal{P}(X)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f_*(V)),$

natural in $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$ and (respectively) $V \in \mathcal{P}(X)$ and $U \in \mathcal{P}(Y)$. In particular:

i. The following conditions are equivalent:

A. We have
$$f_!(U) \subset V$$
.

B. We have
$$U \subset f^{-1}(V)$$
.

ii. The following conditions are equivalent:

A. We have
$$f^{-1}(U) \subset V$$
.

B. We have
$$U \subset f_*(V)$$
.

3. Interaction With Unions of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f)_!} \mathcal{P}(\mathcal{P}(Y))$$

$$\cup \bigcup_{f} \bigcup_{f} \bigcup_{f} \bigcup_{f} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_!(U) = \bigcup_{V \in f_!(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

4. Interaction With Intersections of Families of Subsets. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_{i})_{!}} & \mathcal{P}(\mathcal{P}(Y)) \\
& & & & \downarrow \\
& & & \downarrow \\
\mathcal{P}(X) & \xrightarrow{f_{i}} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U\in\mathcal{U}}f_!(U)=\bigcap_{V\in f_!(\mathcal{U})}V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

5. Interaction With Binary Unions. The diagram

$$\begin{array}{cccc} \mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_{1} \times f_{1}} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\ & & & \downarrow & & \downarrow \\ \mathcal{P}(X) & \xrightarrow{f_{1}} & \mathcal{P}(Y) & \end{array}$$

commutes, i.e. we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$

for each $U, V \in \mathcal{P}(X)$.

6. Interaction With Binary Intersections. We have a natural transformation

with components

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

7. Interaction With Differences. We have a natural transformation

$$\mathcal{P}(X)^{\text{op}} \times \mathcal{P}(X) \xrightarrow{f_{!}^{\text{op}} \times f_{!}} \mathcal{P}(Y)^{\text{op}} \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad$$

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by $U, V \in \mathcal{P}(X)$.

8. Interaction With Complements. The diagram

$$\mathcal{P}(X)^{\operatorname{op}} \xrightarrow{f_*^{\operatorname{op}}} \mathcal{P}(Y)^{\operatorname{op}}$$

$$\downarrow^{(-)^{\operatorname{c}}} \qquad \qquad \downarrow^{(-)^{\operatorname{c}}}$$

$$\mathcal{P}(X) \xrightarrow{f} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_!(U^\mathsf{c}) = f_*(U)^\mathsf{c}$$

for each $U \in \mathcal{P}(X)$.

9. Interaction With Symmetric Differences. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\stackrel{\triangle}{\longrightarrow} \qquad \qquad \qquad \downarrow_{\triangle}$$

$$\mathcal{P}(X) \xrightarrow{f} \qquad \qquad \mathcal{P}(Y)$$

with components

$$f_!(U) \triangle f_!(V) \subset f_!(U \triangle V)$$

indexed by $U, V \in \mathcal{P}(X)$.

10. Interaction With Internal Homs of Powersets. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_*^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \\
 \downarrow [-_1, -_2]_Y \\
 \mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have an equality of sets

$$f_!([U,V]_X) = [f_*(U), f_!(V)]_Y$$

natural in $U, V \in \mathcal{P}(X)$.

11. Preservation of Colimits. We have an equality of sets

$$f[\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f[(U_i),$$

natural in $\{U_i\}_{i\in I}\in\mathcal{P}(X)^{\times I}$. In particular, we have equalities

$$f_!(U) \cup f_!(V) = f_!(U \cup V),$$

 $f_!(\emptyset) = \emptyset,$

natural in $U, V \in \mathcal{P}(X)$.

12. Oplax Preservation of Limits. We have an inclusion of sets

$$f_! \left(\bigcap_{i \in I} U_i \right) \subset \bigcap_{i \in I} f_! (U_i),$$

natural in $\{U_i\}_{i\in I} \in \mathcal{P}(X)^{\times I}$. In particular, we have inclusions

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V),$$

 $f_!(X) \subset Y,$

natural in $U, V \in \mathcal{P}(X)$.

13. Symmetric Strict Monoidality With Respect to Unions. The direct image function of Item 1 has a symmetric strict monoidal structure

$$(f_!, f_!^{\otimes}, f_{!|1}^{\otimes}) \colon (\mathcal{P}(X), \cup, \emptyset) \to (\mathcal{P}(Y), \cup, \emptyset),$$

being equipped with equalities

$$f_{!|U,V}^{\otimes} : f_{!}(U) \cup f_{!}(V) \xrightarrow{=} f_{!}(U \cup V),$$

 $f_{!|\mathbb{1}}^{\otimes} : \emptyset \xrightarrow{=} \emptyset,$

natural in $U, V \in \mathcal{P}(X)$.

14. Symmetric Oplax Monoidality With Respect to Intersections. The direct image function of Item 1 has a symmetric oplax monoidal structure

$$(f_!, f_!^{\otimes}, f_{!|1}^{\otimes}) \colon (\mathcal{P}(X), \cap, X) \to (\mathcal{P}(Y), \cap, Y),$$

being equipped with inclusions

$$f_{!|U,V}^{\otimes} \colon f_{!}(U \cap V) \hookrightarrow f_{!}(U) \cap f_{!}(V),$$
$$f_{!|\mathfrak{A}}^{\otimes} \colon f_{!}(X) \hookrightarrow Y,$$

natural in $U, V \in \mathcal{P}(X)$.

15. *Interaction With Coproducts.* Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \coprod g)_!(U \coprod V) = f_!(U) \coprod g_!(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

16. *Interaction With Products.* Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \boxtimes_{X \times Y} q)_!(U \boxtimes_{X \times Y} V) = f_!(U) \boxtimes_{X' \times Y'} q_!(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

17. Relation to Codirect Images. We have

$$f_!(U) = f_*(U^c)^c$$

 $\stackrel{\text{def}}{=} Y \setminus f_*(X \setminus U)$

for each $U \in \mathcal{P}(X)$.

Proof. Item 1, Functoriality: Omitted.

Item 2, Triple Adjointness: This follows from Definition 4.6.1.1.4, Definition 4.6.2.1.2, Definition 4.6.3.1.4, and Kan Extensions, ?? of ??.

Item 3, Interaction With Unions of Families of Subsets: We have

$$\bigcup_{V \in f_{!}(\mathcal{U})} V = \bigcup_{V \in \{f_{!}(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcup_{U \in \mathcal{U}} f_{!}(U).$$

This finishes the proof.

Item 4, Interaction With Intersections of Families of Subsets: We have

$$\bigcap_{V \in f_{!}(\mathcal{U})} V = \bigcap_{V \in \{f_{!}(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcap_{U \in \mathcal{U}} f_{!}(U).$$

This finishes the proof.

Item 5, Interaction With Binary Unions: See [Pro25p].

Item 6, *Interaction With Binary Intersections*: See [Pro25n].

Item 7, Interaction With Differences: See [Pro250].

Item 8, Interaction With Complements: Applying Item 17 to $X \setminus U$, we have

$$f_!(U^c) = f_!(X \setminus U)$$

$$= Y \setminus f_*(X \setminus (X \setminus U))$$

$$= Y \setminus f_*(U)$$

$$= f_*(U)^c.$$

This finishes the proof.

Item 9, Interaction With Symmetric Differences: We have

$$f_{!}(U) \triangle f_{!}(V) = (f_{!}(U) \cup f_{!}(V)) \setminus (f_{!}(U) \cap f_{!}(V))$$

$$\subset (f_{!}(U) \cup f_{!}(V)) \setminus (f_{!}(U \cap V))$$

$$= (f_{!}(U \cup V)) \setminus (f_{!}(U \cap V))$$

$$\subset f_{!}((U \cup V) \setminus (U \cap V))$$

$$= f_{!}(U \triangle V),$$

where we have used:

- 1. Item 2 of Definition 4.3.12.1.2 for the first equality.
- 2. Item 6 of this proposition together with Item 1 of Definition 4.3.10.1.2 for the first inclusion.
- 3. Item 5 for the second equality.
- 4. Item 7 for the second inclusion.
- 5. Item 2 of Definition 4.3.12.1.2 for the tchird equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 10, Interaction With Internal Homs of Powersets: We have

$$f_!([U,V]_X) \stackrel{\text{def}}{=} f_!(U^c \cup V)$$
$$= f_!(U^c) \cup f_!(V)$$

$$= f_*(U)^{\mathsf{c}} \cup f_!(V)$$

$$\stackrel{\text{def}}{=} [f_*(U), f_!(V)]_{\gamma},$$

where we have used:

- 1. Item 5 for the second equality.
- 2. Item 17 for the third equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 11, Preservation of Colimits: This follows from Item 2 and ??, ?? of ??.34

Item 12, Oplax Preservation of Limits: The inclusion $f_!(X) \subset Y$ is automatic. See [Pro25n] for the other inclusions.

Item 13, Symmetric Strict Monoidality With Respect to Unions: This follows from Item 11.

Item 14, Symmetric Oplax Monoidality With Respect to Intersections: The inclusions in the statement follow from Item 12. Since $\mathcal{P}(Y)$ is posetal, the commutativity of the diagrams in the definition of a symmetric oplax monoidal functor is automatic (Categories, Item 4 of Definition 11.2.7.1.2).

Item 15, Interaction With Coproducts: Omitted.

Item 16, Interaction With Products: Omitted.

Item 17, Relation to Codirect Images: Applying Item 16 of Definition 4.6.3.1.7 to $X \setminus U$, we have

$$f_*(X \setminus U) = B \setminus f_!(X \setminus (X \setminus U))$$
$$= B \setminus f_!(U).$$

Taking complements, we then obtain

$$f_!(U) = B \setminus (B \setminus f_!(U)),$$

= $B \setminus f_*(X \setminus U),$

which finishes the proof.

Proposition 4.6.1.1.6. Let $f: X \to Y$ be a function.

³⁴Reference: [Pro25p].

1. Functionality I. The assignment $f \mapsto f_!$ defines a function

$$(-)_{*|X|Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(X),\mathcal{P}(Y)).$$

2. Functionality II. The assignment $f \mapsto f_!$ defines a function

$$(-)_{*|X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(X),\subset),(\mathcal{P}(Y),\subset)).$$

3. *Interaction With Identities.* For each $X \in Obj(Sets)$, we have

$$(\mathrm{id}_X)_! = \mathrm{id}_{\mathcal{P}(X)}$$
.

4. *Interaction With Composition*. For each pair of composable functions $f: X \to Y$ and $g: Y \to Z$, we have

$$(g \circ f)_{!} = g_{!} \circ f_{!},$$

$$\mathcal{P}(X) \xrightarrow{f_{!}} \mathcal{P}(Y)$$

$$\downarrow^{g_{!}}$$

$$\mathcal{P}(Z).$$

Proof. Item 1, *Functionality I*: There is nothing to prove.

Item 2, Functionality II: This follows from Item 1 of Definition 4.6.1.1.5.

Item 3, Interaction With Identities: This follows from Definition 4.6.1.1.4 and Kan Extensions, ?? of ??.

Item 4, Interaction With Composition: This follows from Definition 4.6.1.1.4 and Kan Extensions, ?? of ??. □

4.6.2 Inverse Images

Let $f: X \to Y$ be a function.

Definition 4.6.2.1.1. The inverse image function associated to f is the function 35

$$f^{-1} \colon \mathcal{P}(Y) \to \mathcal{P}(X)$$

³⁵ Further Notation: Also written $f^* : \mathcal{P}(Y) \to \mathcal{P}(X)$.

defined by³⁶

$$f^{-1}(V) \stackrel{\text{def}}{=} \{x \in X \mid \text{we have } f(x) \in V\}$$

for each $V \in \mathcal{P}(Y)$.

Remark 4.6.2.1.2. Identifying $\mathcal{P}(Y)$ with Sets $(Y, \{t, f\})$ via Item 2 of Definition 4.5.1.1.4, we see that the inverse image function associated to f is equivalently the function

$$f^* \colon \mathcal{P}(Y) \to \mathcal{P}(X)$$

defined by

$$f^*(\chi_V) \stackrel{\text{def}}{=} \chi_V \circ f$$

for each $\chi_V \in \mathcal{P}(Y)$, where $\chi_V \circ f$ is the composition

$$X \xrightarrow{f} Y \xrightarrow{\chi_V} \{\text{true, false}\}$$

in Sets.

Proposition 4.6.2.1.3. Let $f: X \to Y$ be a function.

1. Functoriality. The assignment $V\mapsto f^{-1}(V)$ defines a functor

$$f^{-1} \colon (\mathcal{P}(Y), \subset) \to (\mathcal{P}(X), \subset).$$

In particular, for each $U, V \in \mathcal{P}(Y)$, the following condition is satisfied:

$$(\star) \ \ \text{If} \ U \subset V \text{, then} \ f^{-1}(U) \subset f^{-1}(V).$$

2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

 $[\]overline{^{36}}$ Further Terminology: The set $f^{-1}(V)$ is called the **inverse image of** V by f.

(a) Units and counits of the form

$$id_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad id_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$

 $f_! \circ f^{-1} \hookrightarrow id_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow id_{\mathcal{P}(X)},$

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*(f^{-1}(V)),$$

 $f_!(f^{-1}(V)) \subset V, \qquad f^{-1}(f_*(U)) \subset U$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

(b) Bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(Y)}(f_!(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f^{-1}(V)),$$

$$\operatorname{Hom}_{\mathcal{P}(X)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f_*(V)),$$

natural in $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$ and (respectively) $V \in \mathcal{P}(X)$ and $U \in \mathcal{P}(Y)$. In particular:

- i. The following conditions are equivalent:
 - A. We have $f_!(U) \subset V$.
 - B. We have $U \subset f^{-1}(V)$.
- ii. The following conditions are equivalent:
 - A. We have $f^{-1}(U) \subset V$.
 - B. We have $U \subset f_*(V)$.
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\bigcup_{U} \qquad \qquad \bigcup_{U} \qquad \qquad \downarrow_{U}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{V\in\mathcal{V}}f^{-1}(V)=\bigcup_{U\in f^{-1}(\mathcal{U})}U$$

for each $\mathcal{V} \in \mathcal{P}(Y)$, where $f^{-1}(\mathcal{V}) \stackrel{\text{def}}{=} (f^{-1})^{-1}(\mathcal{V})$.

4. Interaction With Intersections of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcap_{V \in \mathcal{V}} f^{-1}(V) = \bigcap_{U \in f^{-1}(\mathcal{U})} U$$

for each $\mathcal{V} \in \mathcal{P}(Y)$, where $f^{-1}(\mathcal{V}) \stackrel{\text{def}}{=} (f^{-1})^{-1}(\mathcal{V})$.

5. Interaction With Binary Unions. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

6. Interaction With Binary Intersections. The diagram

$$\begin{array}{ccc} \mathcal{P}(Y) \times \mathcal{P}(Y) & \xrightarrow{f^{-1} \times f^{-1}} & \mathcal{P}(X) \times \mathcal{P}(X) \\ & & & \downarrow \cap \\ & & & \downarrow \cap \\ & & \mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X) \end{array}$$

commutes, i.e. we have

$$f^{-1}(U\cap V) = f^{-1}(U)\cap f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

7. Interaction With Differences. The diagram

$$\mathcal{P}(Y)^{\text{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\text{op},-1} \times f^{-1}} \mathcal{P}(X)^{\text{op}} \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad$$

commutes, i.e. we have

$$f^{-1}(U \setminus V) = f^{-1}(U) \setminus f^{-1}(V)$$

for each $U, V \in \mathcal{P}(X)$.

8. Interaction With Complements. The diagram

$$\mathcal{P}(Y)^{\text{op}} \xrightarrow{f^{-1,\text{op}}} \mathcal{P}(X)^{\text{op}} \\
\xrightarrow{(-)^{c}} \qquad \qquad \downarrow^{(-)^{c}} \\
\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U^{\mathsf{c}}) = f^{-1}(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

9. Interaction With Symmetric Differences. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\mathsf{op},-1} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

i.e. we have

$$f^{-1}(U) \vartriangle f^{-1}(V) = f^{-1}(U \vartriangle V)$$

for each $U, V \in \mathcal{P}(Y)$.

10. Interaction With Internal Homs of Powersets. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1,\mathsf{op}} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow [-1,-2]_{Y} \qquad \qquad \downarrow [-1,-2]_{X}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = \left[f^{-1}(U), f^{-1}(V)\right]_X,$$

natural in $U, V \in \mathcal{P}(X)$.

11. Preservation of Colimits. We have an equality of sets

$$f^{-1}\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f^{-1}(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(Y)^{\times I}$. In particular, we have equalities

$$f^{-1}(U) \cup f^{-1}(V) = f^{-1}(U \cup V),$$

 $f^{-1}(\emptyset) = \emptyset,$

natural in $U, V \in \mathcal{P}(Y)$.

12. Preservation of Limits. We have an equality of sets

$$f^{-1}\left(\bigcap_{i\in I}U_i\right)=\bigcap_{i\in I}f^{-1}(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(Y)^{\times I}$. In particular, we have equalities

$$f^{-1}(U) \cap f^{-1}(V) = f^{-1}(U \cap V),$$

 $f^{-1}(Y) = X,$

natural in $U, V \in \mathcal{P}(Y)$.

13. Symmetric Strict Monoidality With Respect to Unions. The inverse image function of Item 1 has a symmetric strict monoidal structure

$$\left(f^{-1}, f^{-1, \otimes}, f_{\mathbb{1}}^{-1, \otimes}\right) \colon (\mathcal{P}(Y), \cup, \emptyset) \to (\mathcal{P}(X), \cup, \emptyset),$$

being equipped with equalities

$$f_{U,V}^{-1,\otimes} \colon f^{-1}(U) \cup f^{-1}(V) \xrightarrow{=} f^{-1}(U \cup V),$$

$$f_{\parallel}^{-1,\otimes} \colon \emptyset \xrightarrow{=} f^{-1}(\emptyset),$$

natural in $U, V \in \mathcal{P}(Y)$.

14. Symmetric Strict Monoidality With Respect to Intersections. The inverse image function of Item 1 has a symmetric strict monoidal structure

$$\left(f^{-1}, f^{-1,\otimes}, f_{\mathbb{1}}^{-1,\otimes}\right) \colon (\mathcal{P}(Y), \cap, Y) \to (\mathcal{P}(X), \cap, X),$$

being equipped with equalities

$$f_{U,V}^{-1,\otimes} \colon f^{-1}(U) \cap f^{-1}(V) \xrightarrow{=} f^{-1}(U \cap V),$$

$$f_{1}^{-1,\otimes} \colon X \xrightarrow{=} f^{-1}(Y),$$

natural in $U, V \in \mathcal{P}(Y)$.

15. *Interaction With Coproducts.* Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \coprod g)^{-1}(U' \coprod V') = f^{-1}(U') \coprod g^{-1}(V')$$

for each $U' \in \mathcal{P}(X')$ and each $V' \in \mathcal{P}(Y')$.

16. Interaction With Products. Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \boxtimes_{X' \times Y'} g)^{-1}(U' \boxtimes_{X' \times Y'} V') = f^{-1}(U') \boxtimes_{X \times Y} g^{-1}(V')$$

for each $U' \in \mathcal{P}(X')$ and each $V' \in \mathcal{P}(Y')$.

Proof. Item 1, Functoriality: Omitted.

Item 2, Triple Adjointness: This follows from Definition 4.6.1.1.4, Definition 4.6.2.1.2,

Definition 4.6.3.1.4, and Kan Extensions, ?? of ??.

Item 3, Interaction With Unions of Families of Subsets: We have

$$\bigcup_{U \in f^{-1}(\mathcal{V})} U = \bigcup_{U \in \{f^{-1}(V) \in \mathcal{P}(X) \mid V \in \mathcal{V}\}} U$$
$$= \bigcup_{V \in \mathcal{V}} f^{-1}(V).$$

This finishes the proof.

Item 4, Interaction With Intersections of Families of Subsets: We have

$$\bigcap_{U \in f^{-1}(\mathcal{V})} U = \bigcap_{U \in \{f^{-1}(V) \in \mathcal{P}(X) \mid V \in \mathcal{V}\}} U$$
$$= \bigcap_{V \in \mathcal{V}} f^{-1}(V).$$

This finishes the proof.

Item 5, Interaction With Binary Unions: See [Pro25y].

Item 6, *Interaction With Binary Intersections*: See [Pro25w].

Item 7, Interaction With Differences: See [Pro25x].

Item 8, *Interaction With Complements*: See [Pro25]].

Item 9, Interaction With Symmetric Differences: We have

$$\begin{split} f^{-1}(U \triangle V) &= f^{-1}((U \cup V) \setminus (U \cap V)) \\ &= f^{-1}(U \cup V) \setminus f^{-1}(U \cap V) \\ &= f^{-1}(U) \cup f^{-1}(V) \setminus f^{-1}(U \cap V) \\ &= f^{-1}(U) \cup f^{-1}(V) \setminus f^{-1}(U) \cap f^{-1}(V) \\ &= f^{-1}(U) \triangle f^{-1}(V), \end{split}$$

where we have used:

- 1. Item 2 of Definition 4.3.12.1.2 for the first equality.
- 2. Item 7 for the second equality.
- 3. Item 5 for the third equality.

- 4. Item 6 for the fourth equality.
- 5. Item 2 of Definition 4.3.12.1.2 for the fifth equality.

This finishes the proof.

Item 10, Interaction With Internal Homs of Powersets: We have

$$f^{-1}([U, V]_Y) \stackrel{\text{def}}{=} f^{-1}(U^{c} \cup V)$$

$$= f^{-1}(U^{c}) \cup f^{-1}(V)$$

$$= f^{-1}(U)^{c} \cup f^{-1}(V)$$

$$\stackrel{\text{def}}{=} [f^{-1}(U), f^{-1}(V)]_X,$$

where we have used:

- 1. Item 8 for the second equality.
- 2. Item 5 for the third equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 11, Preservation of Colimits: This follows from Item 2 and??,?? of??.37

Item 12, Preservation of Limits: This follows from Item 2 and??,?? of??.38

Item 13, Symmetric Strict Monoidality With Respect to Unions: This follows from Item 11.

Item 14, Symmetric Strict Monoidality With Respect to Intersections: This follows from Item 12.

Item 15, Interaction With Coproducts: Omitted.

Item 16, Interaction With Products: Omitted.

Proposition 4.6.2.1.4. Let $f: X \to Y$ be a function.

1. Functionality I. The assignment $f\mapsto f^{-1}$ defines a function

$$(-)_{X,Y}^{-1} \colon \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(Y),\mathcal{P}(X)).$$

2. Functionality II. The assignment $f \mapsto f^{-1}$ defines a function

$$(-)_{X,Y}^{-1} \colon \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(Y),\subset),(\mathcal{P}(X),\subset)).$$

³⁷Reference: [Pro25y].

³⁸Reference: [Pro25w].

3. *Interaction With Identities.* For each $X \in Obj(Sets)$, we have

$$\mathrm{id}_X^{-1}=\mathrm{id}_{\mathcal{P}(X)}.$$

4. *Interaction With Composition*. For each pair of composable functions $f: X \to Y$ and $g: Y \to Z$, we have

$$(g \circ f)^{-1} = f^{-1} \circ g^{-1},$$

$$\mathcal{P}(Z) \xrightarrow{g^{-1}} \mathcal{P}(Y)$$

$$\downarrow^{f^{-1}}$$

$$\mathcal{P}(X).$$

Proof. Item 1, *Functionality I*: There is nothing to prove.

Item 2, Functionality II: This follows from Item 1 of Definition 4.6.2.1.3.

Item 3, Interaction With Identities: This follows from Definition 4.6.2.1.2 and Categories, Item 5 of Definition 11.1.4.1.2.

Item 4, Interaction With Composition: This follows from Definition 4.6.2.1.2 and Categories, Item 2 of Definition 11.1.4.1.2. □

4.6.3 Codirect Images

Let $f: X \to Y$ be a function.

Definition 4.6.3.1.1. The **codirect image function associated to** f is the function

$$f_* \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by^{39,40}

$$f_*(U) \stackrel{\text{def}}{=} \left\{ y \in Y \middle| \begin{array}{l} \text{for each } x \in X, \text{ if we have} \\ f(x) = y, \text{ then } x \in U \end{array} \right\}$$

$$f_*(U) = f_!(U^c)^c$$

 $\stackrel{\text{def}}{=} Y \setminus f_!(X \setminus U)$

see Item 16 of Definition 4.6.3.1.7.

³⁹ Further Terminology: The set $f_*(U)$ is called the **codirect image of** U **by** f.

⁴⁰We also have

$$= \{ y \in Y \mid \text{we have } f^{-1}(y) \subset U \}$$

for each $U \in \mathcal{P}(X)$.

Notation 4.6.3.1.2. Sometimes one finds the notation

$$\forall_f \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

for $f_!$. This notation comes from the fact that the following statements are equivalent, where $y \in Y$ and $U \in \mathcal{P}(X)$:

- We have $y \in \forall_f(U)$.
- For each $x \in X$, if y = f(x), then $x \in U$.

We will not make use of this notation elsewhere in Clowder.

Warning 4.6.3.1.3. See Definition 4.6.1.1.3.

Remark 4.6.3.1.4. Identifying $\mathcal{P}(X)$ with $\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ via Item 2 of Definition 4.5.1.1.4, we see that the codirect image function associated to f is equivalently the function

$$f_* \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by

$$\begin{split} f_*(\chi_U) &\stackrel{\text{def}}{=} \operatorname{Ran}_f(\chi_U) \\ &= \lim \left(\left(\underbrace{(-_1)}_{X \in X} \stackrel{\rightarrow}{\to} f \right) \stackrel{\operatorname{pr}}{\twoheadrightarrow} X \stackrel{\chi_U}{\longrightarrow} \left\{ \text{true, false} \right\} \right) \\ &= \lim_{\substack{x \in X \\ f(x) = -_1}} (\chi_U(x)) \\ &= \bigwedge_{\substack{x \in X \\ f(x) = -_1}} (\chi_U(x)). \end{split}$$

where we have used ?? for the second equality. In other words, we have

$$[f_*(\chi_U)](y) = \bigwedge_{\substack{x \in X \\ f(x) = y}} (\chi_U(x))$$

$$= \begin{cases} \text{true} & \text{if, for each } x \in X \text{ such that} \\ & f(x) = y, \text{ we have } x \in U, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } f^{-1}(y) \subset U \\ \text{false} & \text{otherwise} \end{cases}$$

for each $y \in Y$.

Definition 4.6.3.1.5. Let U be a subset of X.^{41,42}

1. The image part of the codirect image $f_*(U)$ of U is the set $f_{*,\mathrm{im}}(U)$ defined by

$$f_{*,\text{im}}(U) \stackrel{\text{def}}{=} f_*(U) \cap \text{Im}(f)$$

$$= \left\{ y \in Y \middle| \begin{array}{l} \text{we have } f^{-1}(y) \subset U \\ \text{and } f^{-1}(y) \neq \emptyset. \end{array} \right\}.$$

2. The complement part of the codirect image $f_*(U)$ of U is the set $f_{*,cp}(U)$ defined by

$$f_{*,cp}(U) \stackrel{\text{def}}{=} f_*(U) \cap (Y \setminus \text{Im}(f))$$

$$f_*(U) = f_{*,im}(U) \cup f_{*,cp}(U),$$

as

$$\begin{split} f_*(U) &= f_*(U) \cap Y \\ &= f_*(U) \cap (\operatorname{Im}(f) \cup (Y \setminus \operatorname{Im}(f))) \\ &= (f_*(U) \cap \operatorname{Im}(f)) \cup (f_*(U) \cap (Y \setminus \operatorname{Im}(f))) \\ &\stackrel{\text{def}}{=} f_{*,\operatorname{im}}(U) \cup f_{*,\operatorname{cp}}(U). \end{split}$$

⁴²In terms of the meet computation of $f_*(U)$ of Definition 4.6.3.1.4, namely

$$f_*(\chi_U) = \bigwedge_{\substack{x \in X \\ f(x) = -1}} (\chi_U(x)),$$

we see that $f_{*,\text{im}}$ corresponds to meets indexed over nonempty sets, while $f_{*,\text{cp}}$ corresponds to meets indexed over the empty set.

⁴¹Note that we have

$$= Y \setminus \operatorname{Im}(f)$$

$$= \left\{ y \in Y \middle| \text{ we have } f^{-1}(y) \subset U \right\}$$

$$= \left\{ y \in Y \middle| f^{-1}(y) = \emptyset \right\}.$$

Example 4.6.3.1.6. Here are some examples of codirect images.

1. *Multiplication by Two*. Consider the function $f: \mathbb{N} \to \mathbb{N}$ given by

$$f(n) \stackrel{\text{def}}{=} 2n$$

for each $n \in \mathbb{N}$. Since f is injective, we have

$$f_{*,im}(U) = f_!(U)$$

 $f_{*,cn}(U) = \{ \text{odd natural numbers} \}$

for any $U \subset \mathbb{N}$. In particular, we have

$$f_*(\{\text{even natural numbers}\}) = \mathbb{N}.$$

2. *Parabolas*. Consider the function $f: \mathbb{R} \to \mathbb{R}$ given by

$$f(x) \stackrel{\text{def}}{=} x^2$$

for each $x \in \mathbb{R}$. We have

$$f_{*,cp}(U) = \mathbb{R}_{<0}$$

for any $U \subset \mathbb{R}$. Moreover, since $f^{-1}(x) = \{-\sqrt{x}, \sqrt{x}\}$, we have e.g.:

$$f_{*,\text{im}}([0,1]) = \{0\},$$

$$f_{*,\text{im}}([-1,1]) = [0,1],$$

$$f_{*,\text{im}}([1,2]) = \emptyset,$$

$$f_{*,\text{im}}([-2,-1] \cup [1,2]) = [1,4].$$

3. Circles. Consider the function $f \colon \mathbb{R}^2 \to \mathbb{R}$ given by

$$f(x,y) \stackrel{\text{def}}{=} x^2 + y^2$$

for each $(x, y) \in \mathbb{R}^2$. We have

$$f_{*,cp}(U) = \mathbb{R}_{<0}$$

for any $U \subset \mathbb{R}^2$, and since

$$f^{-1}(r) = \begin{cases} \text{a circle of radius } r \text{ about the origin} & \text{if } r > 0, \\ \{(0,0)\} & \text{if } r = 0, \\ \emptyset & \text{if } r < 0, \end{cases}$$

we have e.g.:

$$f_{*,\text{im}}([-1,1] \times [-1,1]) = [0,1],$$

$$f_{*,\text{im}}(([-1,1] \times [-1,1]) \setminus [-1,1] \times \{0\}) = \emptyset.$$

Proposition 4.6.3.1.7. Let $f: X \to Y$ be a function.

1. Functoriality. The assignment $U \mapsto f_*(U)$ defines a functor

$$f_*: (\mathcal{P}(X), \subset) \to (\mathcal{P}(Y), \subset).$$

In particular, for each $U, V \in \mathcal{P}(X)$, the following condition is satisfied:

$$(\star)$$
 If $U \subset V$, then $f_*(U) \subset f_*(V)$.

2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$id_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad id_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$

 $f_! \circ f^{-1} \hookrightarrow id_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow id_{\mathcal{P}(X)},$

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*(f^{-1}(V)),$$

 $f_!(f^{-1}(V)) \subset V, \qquad f^{-1}(f_*(U)) \subset U$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

(b) Bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(Y)}(f_!(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f^{-1}(V)),$$

$$\operatorname{Hom}_{\mathcal{P}(X)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f_*(V)),$$

natural in $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$ and (respectively) $V \in \mathcal{P}(X)$ and $U \in \mathcal{P}(Y)$. In particular:

- i. The following conditions are equivalent:
 - A. We have $f_!(U) \subset V$.
 - B. We have $U \subset f^{-1}(V)$.
- ii. The following conditions are equivalent:
 - A. We have $f^{-1}(U) \subset V$.
 - B. We have $U \subset f_*(V)$.
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\bigcup_{f_*} \bigcup_{f_*} \bigcup_{f_*} \bigcup_{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_*(U) = \bigcup_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

4. Interaction With Intersections of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcap_{U \in \mathcal{U}} f_*(U) = \bigcap_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

5. Interaction With Binary Unions. Let $f: X \to Y$ be a function. We have a natural transformation

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_* \times f_*} \mathcal{P}(Y) \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow$$

with components

$$f_*(U) \cup f_*(V) \subset f_*(U \cup V)$$

indexed by $U, V \in \mathcal{P}(X)$.

6. Interaction With Binary Intersections. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_* \times f_*} \mathcal{P}(Y) \times \mathcal{P}(Y)$$

$$\uparrow \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U) \cap f_*(V) = f_*(U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

7. Interaction With Complements. The diagram

$$\mathcal{P}(X)^{\text{op}} \xrightarrow{f_!^{\text{op}}} \mathcal{P}(Y)^{\text{op}} \\
\xrightarrow{(-)^c} \qquad \qquad \downarrow^{(-)^c} \\
\mathcal{P}(X) \xrightarrow{f} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U^{\mathsf{c}}) = f_!(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

8. Interaction With Symmetric Differences. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_{*}^{\mathsf{op}} \times f_{*}} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(X) \xrightarrow{f_{*}} \mathcal{P}(Y)$$

with components

$$f_*(U \triangle V) \subset f_*(U) \triangle f_*(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

9. Interaction With Internal Homs of Powersets. We have a natural transformation

with components

$$[f_!(U), f_*(V)]_V \subset f_*([U, V]_X)$$

indexed by $U, V \in \mathcal{P}(X)$.

10. Lax Preservation of Colimits. We have an inclusion of sets

$$\bigcup_{i\in I} f_*(U_i) \subset f_*\left(\bigcup_{i\in I} U_i\right),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{\times I}$. In particular, we have inclusions

$$f_*(U) \cup f_*(V) \hookrightarrow f_*(U \cup V),$$

 $\emptyset \hookrightarrow f_*(\emptyset),$

natural in $U, V \in \mathcal{P}(X)$.

11. Preservation of Limits. We have an equality of sets

$$f_*\left(\bigcap_{i\in I}U_i\right)=\bigcap_{i\in I}f_*(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{\times I}$. In particular, we have equalities

$$f^{-1}(U \cap V) = f_*(U) \cap f^{-1}(V),$$

$$f_*(X) = Y,$$

natural in $U, V \in \mathcal{P}(X)$.

12. Symmetric Lax Monoidality With Respect to Unions. The codirect image function of Item 1 has a symmetric lax monoidal structure

$$(f_*, f_*^{\otimes}, f_{*|1}^{\otimes}) \colon (\mathcal{P}(X), \cup, \emptyset) \to (\mathcal{P}(Y), \cup, \emptyset),$$

being equipped with inclusions

$$f_{*|U,V}^{\otimes} \colon f_{*}(U) \cup f_{*}(V) \hookrightarrow f_{*}(U \cup V),$$
$$f_{*|\mathfrak{1}}^{\otimes} \colon \emptyset \hookrightarrow f_{*}(\emptyset),$$

natural in $U, V \in \mathcal{P}(X)$.

13. Symmetric Strict Monoidality With Respect to Intersections. The direct image function of Item 1 has a symmetric strict monoidal structure

$$\left(f_*, f_*^{\otimes}, f_{*|1}^{\otimes}\right) \colon (\mathcal{P}(X), \cap, X) \to (\mathcal{P}(Y), \cap, Y),$$

being equipped with equalities

$$f_{*|U,V}^{\otimes} \colon f_{*}(U \cap V) \xrightarrow{=} f_{*}(U) \cap f_{*}(V),$$
$$f_{*|\mathbb{1}}^{\otimes} \colon f_{*}(X) \xrightarrow{=} Y,$$

natural in $U, V \in \mathcal{P}(X)$.

14. Interaction With Coproducts. Let $f\colon X\to X'$ and $g\colon Y\to Y'$ be maps of sets. We have

$$(f \coprod g)_*(U \coprod V) = f_*(U) \coprod g_*(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

15. Interaction With Products. Let $f\colon X\to X'$ and $g\colon Y\to Y'$ be maps of sets. We have

$$(f \boxtimes_{X \times Y} g)_*(U \boxtimes_{X \times Y} V) = f_*(U) \boxtimes_{X' \times Y'} g_*(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

16. Relation to Direct Images. We have

$$f_*(U) = f_!(U^{\mathsf{c}})^{\mathsf{c}}$$
$$= Y \setminus f_!(X \setminus U)$$

for each $U \in \mathcal{P}(X)$.

17. *Interaction With Injections*. If f is injective, then we have

$$f_{*,\text{im}}(U) = f_!(U),$$

 $f_{*,\text{cp}}(U) = Y \setminus \text{Im}(f),$

and so

$$f_*(U) = f_{*,\text{im}}(U) \cup f_{*,\text{cp}}(U)$$
$$= f_!(U) \cup (Y \setminus \text{Im}(f))$$

for each $U \in \mathcal{P}(X)$.

18. Interaction With Surjections. If f is surjective, then we have

$$f_{*,\text{im}}(U) \subset f_!(U),$$

 $f_{*,\text{cp}}(U) = \emptyset,$

and so

$$f_*(U) \subset f_!(U)$$

for each $U \in \mathcal{P}(X)$.

Proof. Item 1, Functoriality: Omitted.

Item 2, Triple Adjointness: This follows from Definition 4.6.1.1.4, Definition 4.6.2.1.2,

Definition 4.6.3.1.4, and Kan Extensions, ?? of ??.

Item 3, Interaction With Unions of Families of Subsets: We have

$$\bigcup_{V \in f_*(\mathcal{U})} V = \bigcup_{V \in \{f_*(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcup_{U \in \mathcal{U}} f_*(U).$$

This finishes the proof.

Item 4, Interaction With Intersections of Families of Subsets: We have

$$\bigcap_{V \in f_*(\mathcal{U})} V = \bigcap_{V \in \{f_*(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcap_{U \in \mathcal{U}} f_*(U).$$

This finishes the proof.

Item 5, Interaction With Binary Unions: We have

$$f_*(U) \cup f_*(V) = f_!(U^c)^c \cup f_!(V^c)^c$$

$$= (f_!(U^c) \cap f_!(V^c))^c$$

$$\subset (f_!(U^c \cap V^c))^c$$

$$= f_!((U \cup V)^c)^c$$

$$= f_*(U \cup V),$$

where:

- 1. We have used Item 16 for the first equality.
- 2. We have used Item 2 of Definition 4.3.11.1.2 for the second equality.
- 3. We have used Item 6 of Definition 4.6.1.1.5 for the third equality.
- 4. We have used Item 2 of Definition 4.3.11.1.2 for the fourth equality.
- 5. We have used Item 16 for the last equality.

This finishes the proof.

Item 6, *Interaction With Binary Intersections*: This follows from Item 11.

Item 7, Interaction With Complements: Omitted.

Item 8, Interaction With Symmetric Differences: Omitted.

Item 9, Interaction With Internal Homs of Powersets: We have

$$[f_!(U), f^!(V)]_X \stackrel{\text{def}}{=} f_!(U)^{\mathsf{c}} \cup f_*(V)$$

$$= f_*(U^{\mathsf{c}}) \cup f_*(V)$$

$$\subset f_*(U^{\mathsf{c}} \cup V)$$

$$\stackrel{\text{def}}{=} f_*([U, V]_X),$$

where we have used:

- 1. Item 7 of Definition 4.6.3.1.7 for the second equality.
- 2. Item 5 of Definition 4.6.3.1.7 for the inclusion.

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 10, *Lax Preservation of Colimits*: Omitted.

Item 11, Preservation of Limits: This follows from Item 2 and ??, ?? of ??.

Item 12, Symmetric Lax Monoidality With Respect to Unions: This follows from Item 10.

Item 13, Symmetric Strict Monoidality With Respect to Intersections: This follows from Item 11.

Item 14, *Interaction With Coproducts*: Omitted.

Item 15, *Interaction With Products*: Omitted.

Item 16, *Relation to Direct Images*: We claim that $f_*(U) = Y \setminus f_!(X \setminus U)$.

• The First Implication. We claim that

$$f_*(U) \subset Y \setminus f_!(X \setminus U).$$

Let $y \in f_*(U)$. We need to show that $y \notin f_!(X \setminus U)$, i.e. that there is no $x \in X \setminus U$ such that f(x) = y.

This is indeed the case, as otherwise we would have $x \in f^{-1}(y)$ and $x \notin U$, contradicting $f^{-1}(y) \subset U$ (which holds since $y \in f_*(U)$).

Thus $y \in Y \setminus f_!(X \setminus U)$.

• The Second Implication. We claim that

$$Y \setminus f_!(X \setminus U) \subset f_*(U)$$
.

Let $y \in Y \setminus f_!(X \setminus U)$. We need to show that $y \in f_*(U)$, i.e. that $f^{-1}(y) \subset U$.

Since $y \notin f_!(X \setminus U)$, there exists no $x \in X \setminus U$ such that y = f(x), and hence $f^{-1}(y) \subset U$.

Thus $y \in f_*(U)$.

This finishes the proof of Item 16.

Item 17, Interaction With Injections: Omitted.

Item 18, Interaction With Surjections: Omitted.

Proposition 4.6.3.1.8. Let $f: X \to B$ be a function.

1. Functionality I. The assignment $f \mapsto f_*$ defines a function

$$(-)_{!|X,Y} : \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(X),\mathcal{P}(Y)).$$

2. Functionality II. The assignment $f \mapsto f_*$ defines a function

$$(-)_{1|X|Y}$$
: Sets $(X,Y) \to \mathsf{Pos}((\mathcal{P}(X),\subset),(\mathcal{P}(Y),\subset)).$

3. *Interaction With Identities.* For each $X \in Obj(Sets)$, we have

$$(\mathrm{id}_X)_* = \mathrm{id}_{\mathcal{D}(X)}$$
.

4. *Interaction With Composition*. For each pair of composable functions $f: X \to Y$ and $g: Y \to Z$, we have

$$(g \circ f)_* = g_* \circ f_*,$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

$$\downarrow^{g_*}$$

$$\mathcal{P}(Z).$$

Proof. Item 1, Functionality I: There is nothing to prove.

Item 2, Functionality II: This follows from Item 1 of Definition 4.6.3.1.7.

Item 3, Interaction With Identities: This follows from Definition 4.6.3.1.4 and Kan Extensions, ?? of ??.

Item 4, Interaction With Composition: This follows from Definition 4.6.3.1.4 and Kan Extensions, ?? of ??. □

4.6.4 A Six-Functor Formalism for Sets

Remark 4.6.4.1.1. The assignment $X \mapsto \mathcal{P}(X)$ together with the functors f_*, f^{-1} , and $f_!$ of Item 1 of Definition 4.6.1.1.5, Item 1 of Definition 4.6.2.1.3, and Item 1 of Definition 4.6.3.1.7, and the functors

$$-_{1} \cap -_{2} \colon \mathcal{P}(X) \times \mathcal{P}(X) \to \mathcal{P}(X),$$
$$[-_{1}, -_{2}]_{X} \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

of Item 1 of Definition 4.3.9.1.2 and Item 1 of Definition 4.4.7.1.3 satisfy several properties reminiscent of a six functor formalism in the sense of ??.

We collect these properties in Definition 4.6.4.1.2 below.⁴³

Proposition 4.6.4.1.2. Let *X* be a set.

1. The Beck-Chevalley Condition. Let

$$X \times_{Z} Y \xrightarrow{\operatorname{pr}_{2}} Y$$

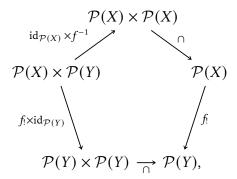
$$\downarrow^{\operatorname{pr}_{1}} \qquad \downarrow^{g}$$

$$X \xrightarrow{f} Z$$

be a pullback diagram in Sets. We have

⁴³ See also [nLa25].

2. The Projection Formula I. The diagram

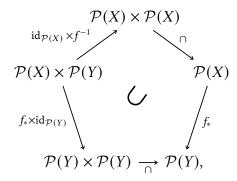


commutes, i.e. we have

$$f_!(U \cap f^{-1}(V)) = f_!(U) \cap V$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

3. The Projection Formula II. We have a natural transformation



with components

$$f_*(U) \cap V \subset f_*(U \cap f^{-1}(V))$$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

4. Strong Closed Monoidality. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1,\mathsf{op}} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \\
 \downarrow [-1,-2]_{Y} \qquad \qquad \downarrow [-1,-2]_{X} \\
 \mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_X,$$

natural in $U, V \in \mathcal{P}(X)$.

5. The External Tensor Product. We have an external tensor product

$$-_1 \boxtimes_{X \times Y} -_2 : \mathcal{P}(X) \times \mathcal{P}(Y) \to \mathcal{P}(X \times Y)$$

given by

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \operatorname{pr}_{1}^{-1}(U) \cap \operatorname{pr}_{2}^{-1}(V)$$
$$= \{(u, v) \in X \times Y \mid u \in U \text{ and } v \in V\}.$$

This is the same map as the one in Item 5 of Definition 4.4.1.1.4. Moreover, the following conditions are satisfied:

(a) Interaction With Direct Images. Let $f: X \to X'$ and $g: Y \to Y'$ be functions. The diagram

$$\begin{array}{c|c} \mathcal{P}(X) \times \mathcal{P}(Y) & \xrightarrow{f_! \times g_!} & \mathcal{P}(X') \times \mathcal{P}(Y') \\ & \boxtimes_{X \times Y} & & & & & & & \\ & & & & & & & \\ & \mathcal{P}(X \times Y) & \xrightarrow{f_! \times g_!} & \mathcal{P}(X' \times Y') & & & \end{array}$$

commutes, i.e. we have

$$[f_! \times g_!](U \boxtimes_{X \times Y} V) = f_!(U) \boxtimes_{X' \times Y'} g_!(V)$$

for each
$$(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$$
.

(b) Interaction With Inverse Images. Let $f: X \to X'$ and $g: Y \to Y'$ be functions. The diagram

commutes, i.e. we have

$$[f^{-1} \times g^{-1}](U \boxtimes_{X' \times Y'} V) = f^{-1}(U) \boxtimes_{X \times Y} g^{-1}(V)$$
 for each $(U, V) \in \mathcal{P}(X') \times \mathcal{P}(Y')$.

(c) *Interaction With Codirect Images*. Let $f: X \to X'$ and $g: Y \to Y'$ be functions. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X) \times \mathcal{P}(Y) & \xrightarrow{f_{*} \times g_{*}} \mathcal{P}(X') \times \mathcal{P}(Y') \\
\boxtimes_{X \times Y} & & & & & & & \\
\mathbb{E}_{X \times Y} & & & & & & \\
\mathcal{P}(X \times Y) & \xrightarrow{f_{*} \times g_{*}} \mathcal{P}(X' \times Y')
\end{array}$$

commutes, i.e. we have

$$[f_* \times g_*](U \boxtimes_{X \times Y} V) = f_*(U) \boxtimes_{X' \times Y'} g_*(V)$$

for each $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$.

(d) Interaction With Diagonals. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\boxtimes_{X \times X}} \mathcal{P}(X \times X)$$

$$\downarrow^{\Delta_X^{-1}}$$

$$\mathcal{P}(X),$$

i.e. we have

$$U\cap V=\Delta_X^{-1}(U\boxtimes_{X\times X}V)$$

for each $U, V \in \mathcal{P}(X)$.

6. The Dualisation Functor. We have a functor

$$D_X \colon \mathcal{P}(X)^{\mathsf{op}} \to \mathcal{P}(X)$$

given by

$$D_X(U) \stackrel{\text{def}}{=} [U, \emptyset]_X$$

for each $U \in \mathcal{P}(X)$, as in Item 5 of Definition 4.4.7.1.3, satisfying the following conditions:

(a) Duality. We have

$$D_X(D_X(U)) = U, \qquad D_X \xrightarrow{\mathrm{id}_{\mathcal{P}(X)}} \mathcal{P}(X)$$

$$\mathcal{P}(X).$$

(b) Duality. The diagram

$$\mathcal{P}(X)^{\text{op}} \times \mathcal{P}(X)^{\text{op}} \xrightarrow{\cap^{\text{op}}} \mathcal{P}(X)^{\text{op}}$$

$$id_{\mathcal{P}(X)^{\text{op}}} \times D_X \nearrow D_X$$

$$\mathcal{P}(X)^{\text{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X)$$

commutes, i.e. we have

$$\underbrace{D_X(U\cap D_X(V))}_{\stackrel{\mathrm{def}}{=}[U\cap [V,\emptyset]_X,\emptyset]_X}=[U,V]_X$$

for each $U, V \in \mathcal{P}(X)$.

(c) Interaction With Direct Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_*^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
\downarrow^{D_X} & & \downarrow^{D_Y} \\
\mathcal{P}(X) & \xrightarrow{f} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_!(D_X(U)) = D_Y(f_*(U))$$

for each $U \in \mathcal{P}(X)$.

(d) Interaction With Inverse Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(Y)^{\mathsf{op}} & \xrightarrow{f^{-1,\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}} \\
\downarrow^{D_{Y}} & & \downarrow^{D_{X}} \\
\mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$f^{-1}(D_Y(U)) = D_X(f^{-1}(U))$$

for each $U \in \mathcal{P}(X)$.

(e) Interaction With Codirect Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\text{op}} & \xrightarrow{f_!^{\text{op}}} \mathcal{P}(Y)^{\text{op}} \\
\downarrow^{D_X} & & \downarrow^{D_Y} \\
\mathcal{P}(X) & \xrightarrow{f_*} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(D_X(U)) = D_Y(f_!(U))$$

for each $U \in \mathcal{P}(X)$.

Proof. Item 1, The Beck-Chevalley Condition: We have

$$[g^{-1} \circ f_!](U) \stackrel{\text{def}}{=} g^{-1}(f_!(U))$$

$$\stackrel{\text{def}}{=} \{y \in Y \mid g(y) \in f_!(U)\}$$

$$= \left\{ y \in Y \mid \text{there exists some } x \in U \right\}$$

$$\text{such that } f(x) = g(y)$$

$$= \left\{ y \in Y \mid \text{there exists some } (x, y) \in \{(x, y) \in X \times_Z Y \mid x \in U\} \right\}$$

$$= \left\{ y \in Y \mid (x, y) \in \{(x, y) \in X \times_Z Y \mid x \in U\} \right\}$$

$$\text{such that } y = y$$

$$= \left\{ y \in Y \mid (x, y) \in \{(x, y) \in X \times_Z Y \mid x \in U\} \right\}$$

$$\text{such that pr}_2(x, y) = y$$

$$\stackrel{\text{def}}{=} (\text{pr}_2)_!(\{(x, y) \in X \times_Z Y \mid x \in U\})$$

$$= (\operatorname{pr}_2)_! (\{(x, y) \in X \times_Z Y \mid \operatorname{pr}_1(x, y) \in U\})$$

$$\stackrel{\text{def}}{=} (\operatorname{pr}_2)_! (\operatorname{pr}_1^{-1}(U))$$

$$\stackrel{\text{def}}{=} [(\operatorname{pr}_2)_! \circ \operatorname{pr}_1^{-1}](U)$$

for each $U \in \mathcal{P}(X)$. Therefore, we have

$$g^{-1} \circ f_! = (pr_2)_! \circ pr_1^{-1}$$
.

For the second equality, we have

$$[f^{-1} \circ g_!](U) \stackrel{\text{def}}{=} f^{-1}(g_!(U))$$

$$\stackrel{\text{def}}{=} \{x \in X \mid f(x) \in g_!(V)\}$$

$$= \left\{x \in X \mid \text{there exists some } y \in V \right\}$$

$$= \left\{x \in X \mid \text{there exists some } (x, y) \in \{(x, y) \in X \times_Z Y \mid y \in V\} \right\}$$

$$= \left\{x \in X \mid \text{there exists some } (x, y) \in \{(x, y) \in X \times_Z Y \mid y \in V\} \right\}$$

$$= \left\{x \in X \mid \text{there exists some } (x, y) \in \{(x, y) \in X \times_Z Y \mid y \in V\} \right\}$$

$$= \left\{x \in X \mid \text{there exists some } (x, y) \in \{(x, y) \in X \times_Z Y \mid y \in V\} \right\}$$

$$= \left\{x \in X \mid \text{there exists some } (x, y) \in \{(x, y) \in X \times_Z Y \mid y \in V\} \right\}$$

$$= \left\{\text{pr}_1\right\}_! (\{(x, y) \in X \times_Z Y \mid \text{pr}_2(x, y) \in V\})$$

$$= \left\{\text{pr}_1\right\}_! (\{(x, y) \in X \times_Z Y \mid \text{pr}_2(x, y) \in V\})$$

$$= \left\{\text{pr}_1\right\}_! (\text{pr}_2^{-1}(V))$$

$$= \left\{(\text{pr}_1\right)_! (\text{pr}_2^{-1}(V))$$

$$= \left\{(\text{pr}_1\right)_! (\text{pr}_2^{-1}(V)) \right\}$$

for each $V \in \mathcal{P}(Y)$. Therefore, we have

$$f^{-1} \circ g_! = (pr_1)_! \circ pr_2^{-1}$$
.

This finishes the proof.

Item 2, The Projection Formula I: We claim that

$$f_!(U) \cap V \subset f_!(U \cap f^{-1}(V)).$$

Indeed, we have

$$f_!(U) \cap V \subset f_!(U) \cap f_!(f^{-1}(V))$$
$$= f_!(U \cap f^{-1}(V)),$$

where we have used:

- 1. Item 2 of Definition 4.6.1.1.5 for the inclusion.
- 2. Item 6 of Definition 4.6.1.1.5 for the equality.

Conversely, we claim that

$$f_!(U \cap f^{-1}(V)) \subset f_!(U) \cap V.$$

Indeed:

- 1. Let $y \in f_!(U \cap f^{-1}(V))$.
- 2. Since $y \in f_!(U \cap f^{-1}(V))$, there exists some $x \in U \cap f^{-1}(V)$ such that f(x) = y.
- 3. Since $x \in U \cap f^{-1}(V)$, we have $x \in U$, and thus $f(x) \in f_!(U)$.
- 4. Since $x \in U \cap f^{-1}(V)$, we have $x \in f^{-1}(V)$, and thus $f(x) \in V$.
- 5. Since $f(x) \in f(U)$ and $f(x) \in V$, we have $f(x) \in f(U) \cap V$.
- 6. But y = f(x), so $y \in f(U) \cap V$.
- 7. Thus $f_!(U \cap f^{-1}(V)) \subset f_!(U) \cap V$.

This finishes the proof.

Item 3, The Projection Formula II: We have

$$f_*(U) \cap V \subset f_*(U) \cap f_*(f^{-1}(V))$$

= $f_*(U \cap f^{-1}(V)),$

where we have used:

1. Item 2 of Definition 4.6.3.1.7 for the inclusion.

2. Item 6 of Definition 4.6.3.1.7 for the equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2).

Item 4, Strong Closed Monoidality: This is a repetition of Item 19 of Definition 4.4.7.1.3 and is proved there.

Item 5, The External Tensor Product: We have

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \operatorname{pr}_{1}^{-1}(U) \cap \operatorname{pr}_{2}^{-1}(V)$$

$$\stackrel{\text{def}}{=} \left\{ (x, y) \in X \times Y \mid \operatorname{pr}_{1}(x, y) \in U \right\}$$

$$\cup \left\{ (x, y) \in X \times Y \mid \operatorname{pr}_{2}(x, y) \in V \right\}$$

$$= \left\{ (x, y) \in X \times Y \mid x \in U \right\}$$

$$\cup \left\{ (x, y) \in X \times Y \mid y \in V \right\}$$

$$\stackrel{\text{def}}{=} U \times V.$$

Next, we claim that Items 5a to 5d are indeed true:

- 1. *Proof of Item 5a*: This is a repetition of Item 16 of Definition 4.6.1.1.5 and is proved there.
- 2. *Proof of Item 5b*: This is a repetition of Item 16 of Definition 4.6.2.1.3 and is proved there.
- 3. *Proof of Item 5c*: This is a repetition of Item 15 of Definition 4.6.3.1.7 and is proved there.
- 4. Proof of Item 5d: We have

$$\begin{split} \Delta_X^{-1}(U \boxtimes_{X \times X} V) &\stackrel{\text{def}}{=} \{x \in X \mid (x, x) \in U \boxtimes_{X \times X} V\} \\ &= \{x \in X \mid (x, x) \in \{(u, v) \in X \times X \mid u \in U \text{ and } v \in V\}\} \\ &= U \cap V. \end{split}$$

This finishes the proof.

Item 6, The Dualisation Functor: This is a repetition of Items 5 and 6 of Definition 4.4.7.1.3 and is proved there. □

Appendices

A Other Chapters

Preliminaries

- 1. Introduction
- 2. A Guide to the Literature

Sets

- 3. Sets
- 4. Constructions With Sets
- 5. Monoidal Structures on the Category of Sets
- 6. Pointed Sets
- 7. Tensor Products of Pointed Sets

Relations

- 8. Relations
- 9. Constructions With Relations

10. Conditions on Relations

Categories

- 11. Categories
- 12. Presheaves and the Yoneda Lemma

Monoidal Categories

13. Constructions With Monoidal Categories

Bicategories

14. Types of Morphisms in Bicategories

Extra Part

15. Notes

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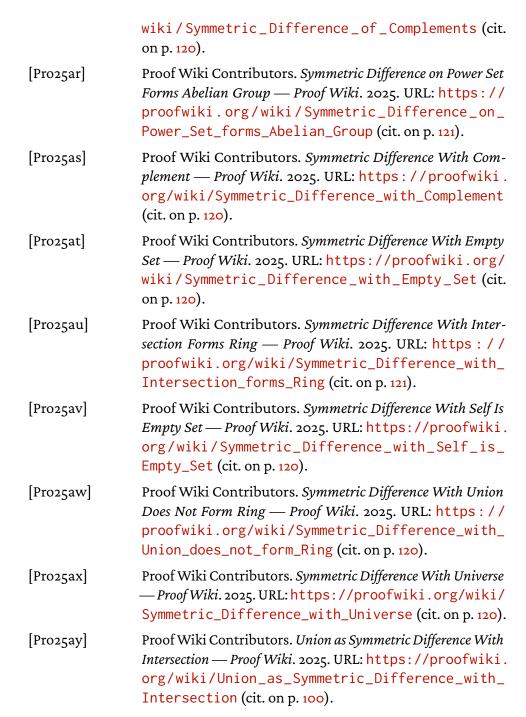
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