

Constructions With Monoidal Categories

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01UF This chapter contains some material on constructions with monoidal categories.

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01UG 13.I Moduli Categories of Monoidal Structures

01UH 13.I.I The Moduli Category of Monoidal Structures on a Category

Let C be a category.

01UJ **Definition 13.I.I.I.1.** The **moduli category of monoidal structures on C** is the category $\mathcal{M}_{E_1}(C)$ defined by

$$\mathcal{M}_{E_1}(C) \stackrel{\text{def}}{=} \text{pt} \times_{\text{Cats}} \text{MonCats},$$

$$\begin{array}{ccc} \mathcal{M}_{E_1}(C) & \longrightarrow & \text{MonCats} \\ \downarrow & \lrcorner & \downarrow \tilde{\omega} \\ \text{pt} & \xrightarrow{[C]} & \text{Cats.} \end{array}$$

01UK **Remark 13.I.I.I.2.** In detail, **the moduli category of monoidal structures on C** is the category $\mathcal{M}_{E_1}(C)$ where:

- *Objects.* The objects of $\mathcal{M}_{E_1}(C)$ are monoidal categories $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ whose underlying category is C .
- *Morphisms.* A morphism from $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ to $(C, \boxtimes_C, 1'_C, \alpha^{C'}, \lambda^{C'}, \rho^{C'})$ is a strong monoidal functor structure

$$\begin{aligned} \text{id}_C^\otimes: A \boxtimes_C B &\xrightarrow{\sim} A \otimes_C B, \\ \text{id}_{1|C}^\otimes: 1'_C &\xrightarrow{\sim} 1_C \end{aligned}$$

on the identity functor $\text{id}_C: C \rightarrow C$ of C .

- *Identities.* For each $M \stackrel{\text{def}}{=} (C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C) \in \text{Obj}(\mathcal{M}_{E_1}(C))$, the unit map

$$\mathbb{1}_{M,M}^{\mathcal{M}_{E_1}(C)}: \text{pt} \rightarrow \text{Hom}_{\mathcal{M}_{E_1}(C)}(M, M)$$

of $\mathcal{M}_{E_1}(C)$ at M is defined by

$$\text{id}_M^{\mathcal{M}_{E_1}(C)} \stackrel{\text{def}}{=} (\text{id}_C^\otimes, \text{id}_{1|C}^\otimes),$$

where $(\text{id}_C^\otimes, \text{id}_{1|C}^\otimes)$ is the identity monoidal functor of C of ??.

- *Composition.* For each $M, N, P \in \text{Obj}(\mathcal{M}_{\mathbb{E}_1}(C))$, the composition map

$$\circ_{M,N,P}^{\mathcal{M}_{\mathbb{E}_1}(C)} : \text{Hom}_{\mathcal{M}_{\mathbb{E}_1}(C)}(N, P) \times \text{Hom}_{\mathcal{M}_{\mathbb{E}_1}(C)}(M, N) \rightarrow \text{Hom}_{\mathcal{M}_{\mathbb{E}_1}(C)}(M, P)$$

of $\mathcal{M}_{\mathbb{E}_1}(C)$ at (M, N, P) is defined by

$$\left(\text{id}_C^{\otimes, '}, \text{id}_{1|C}^{\otimes, '}, \right) \circ_{M,N,P}^{\mathcal{M}_{\mathbb{E}_1}(C)} \left(\text{id}_C^{\otimes}, \text{id}_{1|C}^{\otimes} \right) \stackrel{\text{def}}{=} \left(\text{id}_C^{\otimes, '}, \text{id}_C^{\otimes}, \text{id}_{1|C}^{\otimes, '}, \text{id}_{1|C}^{\otimes} \right).$$

01UL Remark 13.1.1.3. In particular, a morphism in $\mathcal{M}_{\mathbb{E}_1}(C)$ from $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ to $(C, \boxtimes_C, 1'_C, \alpha^{C, '}, \lambda^{C, '}, \rho^{C, '})$ satisfies the following conditions:

- 01UM** 1. *Naturality.* For each pair $f: A \rightarrow X$ and $g: B \rightarrow Y$ of morphisms of C , the diagram

$$\begin{array}{ccc} A \boxtimes_C B & \xrightarrow{f \boxtimes_C g} & X \boxtimes_C Y \\ \text{id}_{A,B}^{\otimes} \downarrow & & \downarrow \text{id}_{X,Y}^{\otimes} \\ A \otimes_C B & \xrightarrow{f \otimes_C g} & X \otimes_C Y \end{array}$$

commutes.

- 01UN** 2. *Monoidality.* For each $A, B, C \in \text{Obj}(C)$, the diagram

$$\begin{array}{ccc} & (A \boxtimes_C B) \boxtimes_C C & \\ \text{id}_{A,B}^{\otimes} \boxtimes \text{id}_C \swarrow & & \searrow \alpha_{A,B,C}^{C, '}, \\ (A \otimes_C B) \boxtimes_C C & & A \boxtimes_C (B \boxtimes_C C) \\ \text{id}_{A \otimes_C B, C}^{\otimes} \downarrow & & \downarrow \text{id}_A \boxtimes \text{id}_{B,C}^{\otimes} \\ (A \otimes_C B) \otimes_C C & & A \boxtimes_C (B \otimes_C C) \\ \alpha_{A,B,C}^C \swarrow & & \swarrow \text{id}_{A,B \otimes_C C}^{\otimes} \\ & A \otimes_C (B \otimes_C C) & \end{array}$$

commutes.

01UP 3. *Left Monoidal Unity*. For each $A \in \text{Obj}(C)$, the diagram

$$\begin{array}{ccc}
 & 1_C \boxtimes_C A & \xrightarrow{\text{id}_{1'_C}^\otimes} 1_C \otimes_C A \\
 \text{id}_1^\otimes \boxtimes_C \text{id}_A \nearrow & & \searrow \lambda_A^C \\
 1'_C \boxtimes_C A & \xrightarrow{\lambda_A^{C'}} & A
 \end{array}$$

commutes.

01UQ 4. *Right Monoidal Unity*. For each $A \in \text{Obj}(C)$, the diagram

$$\begin{array}{ccc}
 & A \boxtimes_C 1_C & \xrightarrow{\text{id}_{A,1'_C}^\otimes} A \otimes_C 1_C \\
 \text{id}_A \boxtimes_C \text{id}_{1'}^\otimes \nearrow & & \searrow \rho_A^C \\
 A \boxtimes_C 1'_C & \xrightarrow{\rho_A^{C'}} & A
 \end{array}$$

commutes.

01UR **Proposition 13.1.1.4.** Let C be a category.

01US I. *Extra Monoidality Conditions*. Let $(\text{id}_C^\otimes, \text{id}_{1_C}^\otimes)$ be a morphism of $\mathcal{M}_{\mathbb{E}_1}(C)$ from $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ to $(C, \boxtimes_C, 1'_C, \alpha^{C'}, \lambda^{C'}, \rho^{C'})$.

01UT (a) The diagram

$$\begin{array}{ccc}
 (A \boxtimes_C B) \boxtimes_C C & \xrightarrow{\text{id}_{A,B}^\otimes \boxtimes_C \text{id}_C} & (A \otimes_C B) \boxtimes_C C \\
 \text{id}_{A \boxtimes_C B, C}^\otimes \downarrow & & \downarrow \text{id}_{A \otimes_C B, C}^\otimes \\
 (A \boxtimes_C B) \otimes_C C & \xrightarrow{\text{id}_{A,B}^\otimes \otimes_C \text{id}_C} & (A \otimes_C B) \otimes_C C
 \end{array}$$

commutes.

01UU (b) The diagram

$$\begin{array}{ccc}
 A \boxtimes_C (B \boxtimes_C C) & \xrightarrow{\text{id}_A \boxtimes_C \text{id}_{B,C}^\otimes} & A \boxtimes_C (B \otimes_C C) \\
 \text{id}_{A, B \boxtimes_C C}^\otimes \downarrow & & \downarrow \text{id}_{A, B \otimes_C C}^\otimes \\
 A \otimes_C (B \boxtimes_C C) & \xrightarrow{\text{id}_A \otimes_C \text{id}_{B,C}^\otimes} & A \otimes_C (B \otimes_C C)
 \end{array}$$

commutes.

01WB 2. *Extra Monoidal Unity Constraints.* Let $(\text{id}_C^\otimes, \text{id}_{1|C}^\otimes)$ be a morphism of $\mathcal{M}_{\mathbb{E}_1}(C)$ from $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ to $(C, \boxtimes_C, 1'_C, \alpha^{C'}, \lambda^{C'}, \rho^{C'})$.

01WC (a) The diagram

$$\begin{array}{ccc} 1_C \otimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \boxtimes_C 1'_C \\ \lambda_{1'_C}^C \downarrow & & \downarrow \rho_{1_C}^{C'} \\ 1'_C & \xrightarrow{\text{id}_1^\otimes} & 1_C \end{array}$$

commutes.

01WD (b) The diagram

$$\begin{array}{ccc} 1'_C \otimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1'_C \boxtimes_C 1_C \\ \rho_{1'_C}^C \downarrow & & \downarrow \lambda_{1_C}^{C'} \\ 1'_C & \xrightarrow{\text{id}_1^\otimes} & 1_C \end{array}$$

commutes.

01WE (c) The diagram

$$\begin{array}{ccc} 1'_C \boxtimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1'_C \otimes_C 1_C \\ \lambda_{1_C}^{C'} \downarrow & & \downarrow \rho_{1'_C}^C \\ 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1'_C \end{array}$$

commutes.

01WF (d) The diagram

$$\begin{array}{ccc}
 1_C \boxtimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^\otimes} & 1_C \otimes_C 1'_C \\
 \rho_{1_C}^{C, '}\downarrow & & \downarrow \lambda_{1'_C}^C \\
 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1'_C
 \end{array}$$

commutes.

01UV 3. *Mixed Associators.* Let $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ and $(C, \boxtimes_C, 1'_C, \alpha^{C, '}, \lambda^{C, '}, \rho^{C, '})$ be monoidal structures on C and let

$$\text{id}_{-1, -2}^\otimes: -1 \boxtimes_C -2 \rightarrow -1 \otimes_C -2$$

be a natural transformation.

01UW (a) If there exists a natural transformation

$$\alpha_{A, B, C}^\otimes: (A \otimes_C B) \boxtimes_C C \rightarrow A \otimes_C (B \boxtimes_C C)$$

making the diagrams

$$\begin{array}{ccc}
 (A \otimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A, B, C}^\otimes} & A \otimes_C (B \boxtimes_C C) \\
 \text{id}_{A \otimes_C B, C}^\otimes \downarrow & & \downarrow \text{id}_A \otimes \text{id}_{B, C}^\otimes \\
 (A \otimes_C B) \otimes_C C & \xrightarrow{\alpha_{A, B, C}^C} & A \otimes_C (B \otimes_C C)
 \end{array}$$

and

$$\begin{array}{ccc}
 (A \boxtimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A, B, C}^{C, '}} & A \boxtimes_C (B \boxtimes_C C) \\
 \text{id}_{A, B}^\otimes \boxtimes \text{id}_C \downarrow & & \downarrow \text{id}_{A, B \boxtimes_C C} \\
 (A \otimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A, B, C}^\otimes} & A \otimes_C (B \boxtimes_C C)
 \end{array}$$

commute, then the natural transformation id^\otimes satisfies the monoidality condition of **Item 2** of **Definition 13.1.1.1.3**.

01UX (b) If there exists a natural transformation

$$\alpha_{A,B,C}^{\boxtimes} : (A \boxtimes_C B) \otimes_C C \rightarrow A \boxtimes_C (B \otimes_C C)$$

making the diagrams

$$\begin{array}{ccc} (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes}} & A \boxtimes_C (B \otimes_C C) \\ \text{id}_{A,B}^{\otimes} \otimes \text{id}_C \downarrow & & \downarrow \text{id}_{A,B \otimes C}^{\otimes} \\ (A \otimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^C} & A \otimes_C (B \otimes_C C) \end{array}$$

and

$$\begin{array}{ccc} (A \boxtimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A,B,C}^{C,\boxtimes}} & A \boxtimes_C (B \boxtimes_C C) \\ \text{id}_{A \boxtimes B,C}^{\otimes} \downarrow & & \downarrow \text{id}_A \boxtimes \text{id}_{B,C}^{\otimes} \\ (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes}} & A \boxtimes_C (B \otimes_C C) \end{array}$$

commute, then the natural transformation id^{\otimes} satisfies the monoidality condition of **Item 2** of **Definition 13.1.1.1.3**.

01UY (c) If there exists a natural transformation

$$\alpha_{A,B,C}^{\boxtimes,\otimes} : (A \boxtimes_C B) \otimes_C C \rightarrow A \otimes_C (B \boxtimes_C C)$$

making the diagrams

$$\begin{array}{ccc} (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes,\otimes}} & A \otimes_C (B \boxtimes_C C) \\ \text{id}_{A,B}^{\otimes} \otimes \text{id}_C \downarrow & & \downarrow \text{id}_A \otimes \text{id}_{B,C}^{\otimes} \\ (A \otimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^C} & A \otimes_C (B \otimes_C C) \end{array}$$

and

$$\begin{array}{ccc}
 (A \boxtimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A,B,C}^{C'}} & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow \text{id}_{A \boxtimes_C B, C}^{\otimes} & & \downarrow \text{id}_{A, B \boxtimes_C C}^{\otimes} \\
 (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes, \otimes}} & A \otimes_C (B \boxtimes_C C)
 \end{array}$$

commute, then the natural transformation id^{\otimes} satisfies the monoidality condition of **Item 2** of **Definition 13.1.1.1.3**.

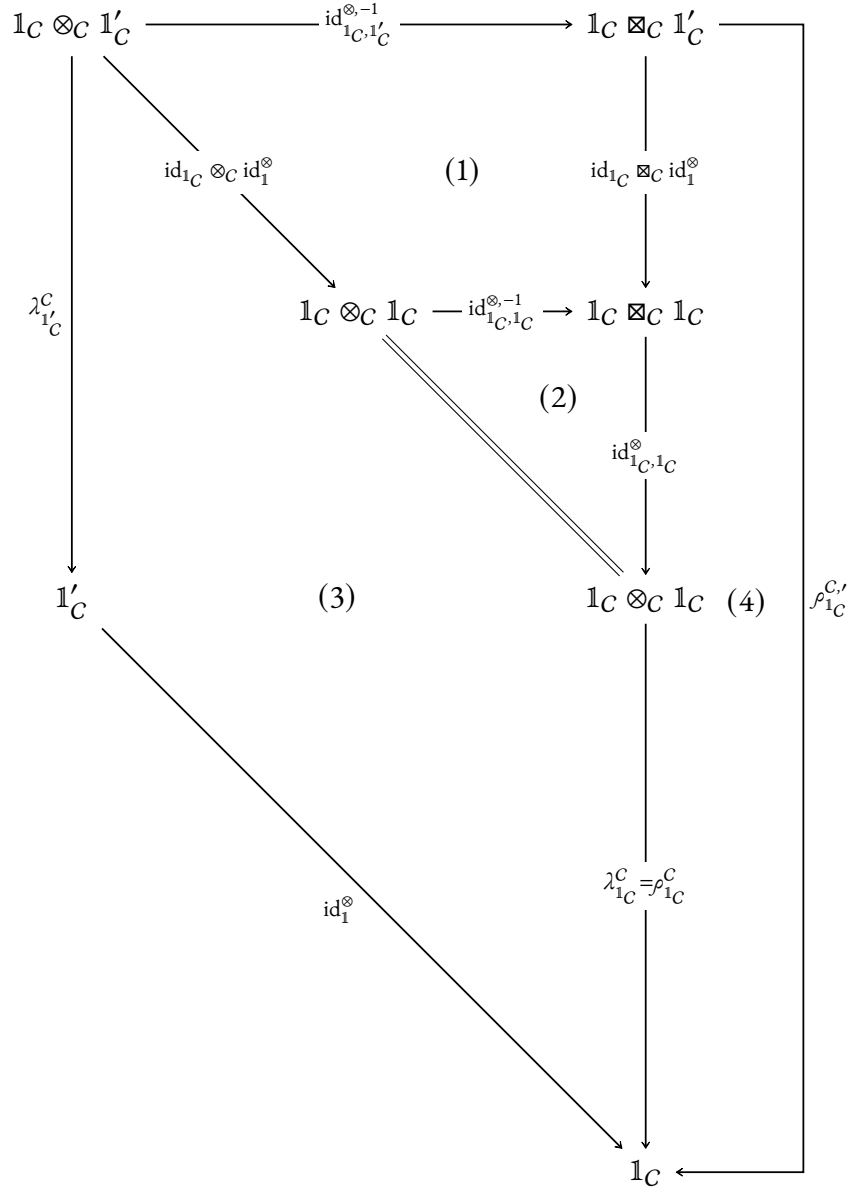
Proof. **Item 1, Extra Monoidality Conditions:** We claim that **Items 1a** and **1b** are indeed true:

1. *Proof of Item 1a:* This follows from the naturality of id^{\otimes} with respect to the morphisms $\text{id}_{A,B}^{\otimes}$ and id_C .
2. *Proof of Item 1b:* This follows from the naturality of id^{\otimes} with respect to the morphisms id_A and $\text{id}_{B,C}^{\otimes}$.

This finishes the proof.

Item 2, Extra Monoidal Unity Constraints: We claim that **Items 2a** and **2b** are indeed true:

I. *Proof of Item 1a:* Indeed, consider the diagram



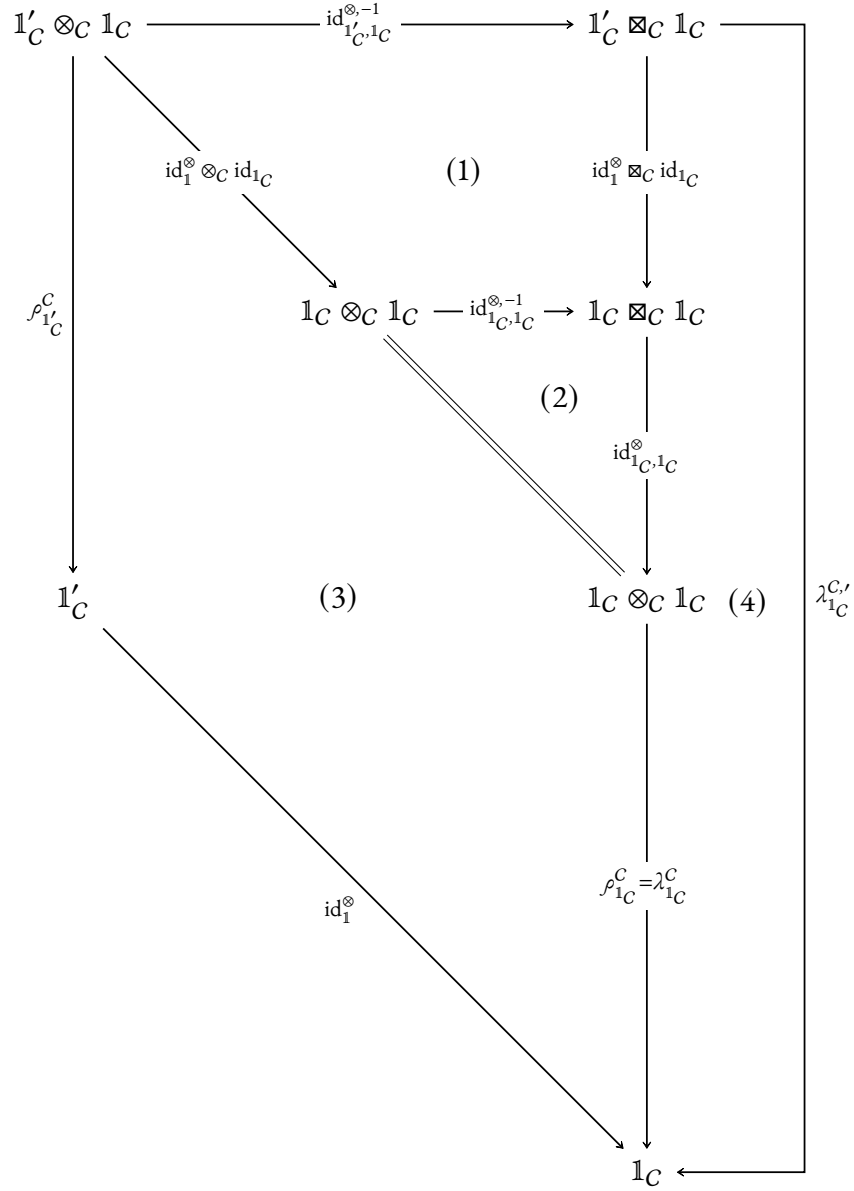
whose boundary diagram is the diagram whose commutativity we wish to prove. Since:

- Subdiagram (1) commutes by the naturality of $\text{id}_C^{\otimes, -1}$;

- Subdiagram (2) commutes trivially;
- Subdiagram (3) commutes by the naturality of λ^C , where the equality $\rho_{1_C}^C = \lambda_{1_C}^C$ comes from ??;
- Subdiagram (4) commutes by the right monoidal unity of $(\mathrm{id}_C, \mathrm{id}_C^\otimes, \mathrm{id}_{C|1}^\otimes)$;

so does the boundary diagram, and we are done.

2. *Proof of Item 1b:* Indeed, consider the diagram



whose boundary diagram is the diagram whose commutativity we wish to prove. Since:

- Subdiagram (1) commutes by the naturality of $\text{id}_C^{\otimes, -1}$;

- Subdiagram (2) commutes trivially;
- Subdiagram (3) commutes by the naturality of ρ^C , where the equality $\rho_{1_C}^C = \lambda_{1_C}^C$ comes from ??;
- Subdiagram (4) commutes by the left monoidal unity of $(\text{id}_C, \text{id}_C^\otimes, \text{id}_{C|1}^\otimes)$;

so does the boundary diagram, and we are done.

3. *Proof of Item 2c:* Indeed, consider the diagram

$$\begin{array}{ccccc}
 1'_C \otimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1'_C \boxtimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^\otimes} & 1'_C \otimes_C 1_C \\
 \rho_{1'_C}^C \downarrow & & \downarrow \lambda_{1_C}^{C, \prime} & & \downarrow \rho_{1'_C}^C \\
 1'_C & \xrightarrow{\text{id}_1^\otimes} & 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1'_C.
 \end{array}
 \quad \begin{array}{c} (1) \end{array} \quad \begin{array}{c} (\dagger)$$

Since:

- The boundary diagram commutes trivially;
- Subdiagram (1) commutes by Item 1b;

it follows that the diagram

$$\begin{array}{ccccc}
 1'_C \otimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1'_C \boxtimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^\otimes} & 1'_C \otimes_C 1_C \\
 & & \downarrow \lambda_{1_C}^{C, \prime} & & \downarrow \rho_{1'_C}^C \\
 & & 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1'_C
 \end{array}
 \quad \begin{array}{c} (\dagger)$$

commutes. But since $\text{id}_{1_C, 1'_C}^{\otimes, -1}$ is an isomorphism, it follows that the diagram (\dagger) also commutes, and we are done.

4. *Proof of Item 2d:* Indeed, consider the diagram

$$\begin{array}{ccccc}
 1_C \otimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \boxtimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \otimes_C 1'_C \\
 \lambda_{1'_C}^C \downarrow & & \downarrow \rho_{1_C}^{C, \prime} & & \downarrow \lambda_{1'_C}^C \\
 1'_C & \xrightarrow{\text{id}_1^{\otimes}} & 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1_C
 \end{array}
 \quad \begin{array}{c} (1) \end{array} \quad \begin{array}{c} (\dagger) \end{array}$$

Since:

- The boundary diagram commutes trivially;
- Subdiagram (1) commutes by **Item 1a**;

it follows that the diagram

$$\begin{array}{ccccc}
 1_C \otimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \boxtimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \otimes_C 1'_C \\
 & & \downarrow \rho_{1_C}^{C, \prime} & & \downarrow \lambda_{1'_C}^C \\
 & & 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1_C
 \end{array}
 \quad \begin{array}{c} (\dagger) \end{array}$$

commutes. But since $\text{id}_1^{\otimes, -1}$ is an isomorphism, it follows that the diagram (\dagger) also commutes, and we are done.

This finishes the proof.

Item 3, Mixed Associators: We claim that **Items 3a** to **3c** are indeed true:

01UZ 1. *Proof of Item 3a:* We may partition the monoidality diagram for id^{\otimes} of

Item 2 of Definition 13.1.1.3 as follows:

$$\begin{array}{ccccc}
 & (A \boxtimes_C B) \boxtimes_C C & & & \\
 \text{id}_{A,B}^\otimes \boxtimes \text{id}_C & \swarrow & \downarrow & \searrow & \alpha_{A,B,C}^C \\
 (A \otimes_C B) \boxtimes_C C & & \text{id}_{A \otimes_C B, C}^\otimes & & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow & (1) & \downarrow & (2) & \downarrow \\
 \text{id}_{A \otimes_C B, C}^\otimes & & (A \boxtimes_C B) \otimes_C C & & \text{id}_A \boxtimes \text{id}_{B,C}^\otimes \\
 \downarrow & \swarrow & \downarrow & \searrow & \downarrow \\
 (A \otimes_C B) \otimes_C C & & \text{id}_{A,B}^\otimes \otimes \text{id}_C & & A \boxtimes_C (B \otimes_C C) \\
 & (3) & \swarrow & \searrow & \\
 & & \text{id}_{A,B}^\otimes \otimes \text{id}_C & & \alpha_{A,B,C}^\otimes \\
 & & \downarrow & & \downarrow \\
 & & (A \otimes_C B) \otimes_C C & & A \boxtimes_C (B \otimes_C C) \\
 & & \searrow & \swarrow & \\
 & & \alpha_{A,B,C}^{C,\prime} & & \text{id}_{A, B \otimes_C C}^\otimes \\
 & & A \otimes_C (B \otimes_C C) & &
 \end{array}$$

Since:

- Subdiagram (1) commutes by Item 1a of Item 1.
- Subdiagram (2) commutes by assumption.
- Subdiagram (3) commutes by assumption.

it follows that the boundary diagram also commutes, i.e. id^\otimes satisfies the monoidality condition of Item 2 of Definition 13.1.1.3.

01V0 2. *Proof of Item 3b:* We may partition the monoidality diagram for id^\otimes of

Item 2 of Definition 13.1.1.3 as follows:

$$\begin{array}{ccccc}
 & (A \boxtimes_C B) \boxtimes_C C & & & \\
 \text{id}_{A,B}^\otimes \boxtimes \text{id}_C & \swarrow & & \searrow & \alpha_{A,B,C}^C \\
 (A \otimes_C B) \boxtimes_C C & & (1) & & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow \text{id}_{A \otimes_C B, C}^\otimes & \searrow \alpha_{A,B,C}^\boxtimes & & \swarrow \text{id}_{A, B \otimes_C C}^\otimes & \downarrow \text{id}_A \boxtimes \text{id}_{B,C}^\otimes \\
 & A \otimes_C (B \boxtimes_C C) & & & \\
 & (2) & & (3) & \\
 & \downarrow \text{id}_A \otimes \text{id}_{B,C}^\otimes & & & \\
 (A \otimes_C B) \otimes_C C & & A \boxtimes_C (B \otimes_C C) & & \\
 \searrow \alpha_{A,B,C}^{C,\prime} & \downarrow & \swarrow \text{id}_{A, B \otimes_C C}^\otimes & & \\
 & A \otimes_C (B \otimes_C C) & & &
 \end{array}$$

Since:

- Subdiagram (1) commutes by assumption.
- Subdiagram (2) commutes by assumption.
- Subdiagram (3) commutes by Item 1b of Item 1.

it follows that the boundary diagram also commutes, i.e. id^\otimes satisfies the monoidality condition of Item 2 of Definition 13.1.1.3.

01V1 3. *Proof of Item 3c:* We may partition the monoidality diagram for id^\otimes of

Item 2 of Definition 13.1.1.1.3 as follows:

$$\begin{array}{ccc}
 & (A \boxtimes_C B) \boxtimes_C C & \\
 \text{id}_{A,B}^\otimes \boxtimes_C \text{id}_C \swarrow & & \searrow \alpha_{A,B,C}^C \\
 (A \otimes_C B) \boxtimes_C C & & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow \text{id}_{A \otimes_C B, C}^\otimes & \searrow \alpha_{A,B,C}^{\boxtimes, \otimes} & \downarrow \text{id}_A \boxtimes_C \text{id}_{B,C}^\otimes \\
 (A \otimes_C B) \otimes_C C & & A \boxtimes_C (B \otimes_C C) \\
 \searrow \alpha_{A,B,C}^{C, \prime} & & \swarrow \text{id}_{A, B \otimes_C C}^\otimes \\
 & A \otimes_C (B \otimes_C C) &
 \end{array}
 \quad (1) \quad (2)$$

Since subdiagrams (1) and (2) commute by assumption, it follows that the boundary diagram also commutes, i.e. id^\otimes satisfies the monoidality condition of Item 2 of Definition 13.1.1.1.3.

This finishes the proof. \square

01V2 13.1.2 The Moduli Category of Braided Monoidal Structures on a Category

01V3 13.1.3 The Moduli Category of Symmetric Monoidal Structures on a Category

01V4 13.2 Moduli Categories of Closed Monoidal Structures

01V5 13.3 Moduli Categories of Refinements of Monoidal Structures

01V6 13.3.1 The Moduli Category of Braided Refinements of a Monoidal Structure

Appendices

A Other Chapters

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1. Introduction
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7. Tensor Products of Pointed Sets

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8. Relations
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11. Categories
12. Presheaves and the Yoneda Lemma

Monoidal Categories

13. Constructions With Monoidal Categories

Bicategories

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Extra Part

15. Notes