Pointed Sets

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July 29, 2025

This chapter contains some foundational material on pointed sets.

Contents

6.1 Pointed Sets

6.1.1 Foundations

Definition 6.1.1.1.1. A **pointed set**[1] is equivalently:

- An $\mathbb{E}_0$-monoid in $(\mathsf{N}_\bullet(\mathsf{Sets}), \mathrm{pt})$.
- A pointed object in $(\mathsf{Sets}, \mathrm{pt})$.

Remark 6.1.1.1.2. In detail, a **pointed set** is a pair (X, x_0) consisting of:

- *The Underlying Set.* A set X, called the **underlying set of** (X, x_0).
- *The Basepoint.* A morphism

$$[x_0]\colon \mathrm{pt} \to X$$

  in Sets, determining an element $x_0 \in X$, called the **basepoint of** X.

Example 6.1.1.1.3. The 0**-sphere**[2] is the pointed set $(S^0, 0)$[3] consisting of:

- *The Underlying Set.* The set S^0 defined by

$$S^0 \overset{\mathrm{def}}{=} \{0, 1\}.$$

- *The Basepoint.* The element 0 of S^0.

Example 6.1.1.1.4. The **trivial pointed set** is the pointed set $(\mathrm{pt}, \star)$ consisting of:

- *The Underlying Set.* The punctual set $\mathrm{pt} \overset{\mathrm{def}}{=} \{\star\}$.
- *The Basepoint.* The element $\star$ of pt.

Example 6.1.1.1.5. The **standard pointed set with** $n+1$ **elements** is the pointed set $\langle n \rangle$ consisting of

- *The Underlying Set.* The set $\langle n \rangle$ defined by

$$\langle n \rangle \overset{\mathrm{def}}{=} \{*\} \cup \{1, \ldots, n\}.$$

- *The Basepoint.* The element $*$ of $\langle n \rangle$.

[1] *Further Terminology:* In the context of monoids with zero as models for $\mathbb{F}_1$-algebras, pointed sets are viewed as $\mathbb{F}_1$**-modules**.

[2] *Further Terminology:* In the context of monoids with zero as models for $\mathbb{F}_1$-algebras, the 0-sphere is viewed as the **underlying pointed set of the field with one element**.

[3] *Further Notation:* In the context of monoids with zero as models for $\mathbb{F}_1$-algebras, S^0 is also

6.1.2 Morphisms of Pointed Sets

Definition 6.1.2.1.1. A **morphism of pointed sets**[4,5] is equivalently:

- A morphism of $\mathbb{E}_0$-monoids in $(\mathrm{N}_\bullet(\mathsf{Sets}), \mathrm{pt})$.
- A morphism of pointed objects in $(\mathsf{Sets}, \mathrm{pt})$.

Remark 6.1.2.1.2. In detail, a **morphism of pointed sets** $f\colon (X, x_0) \to (Y, y_0)$ is a morphism of sets $f\colon X \to Y$ such that the diagram

$$\begin{tikzcd} & \mathrm{pt} \arrow[dl, "{[x_0]}"'] \arrow[dr, "{[y_0]}"] & \\ X \arrow[rr, "f"'] & & Y \end{tikzcd}$$

commutes, i.e. such that

$$f(x_0) = y_0.$$

6.1.3 The Category of Pointed Sets

Definition 6.1.3.1.1. The **category of pointed sets** is the category Sets_* defined equivalently as:

- The homotopy category of the ∞-category $\mathsf{Mon}_{\mathbb{E}_0}(\mathrm{N}_\bullet(\mathsf{Sets}), \mathrm{pt})$ of **??**, **??**.
- The category Sets_* of Constructions With Categories, **??**.

Remark 6.1.3.1.2. In detail, the **category of pointed sets** is the category Sets_* where:

- *Objects.* The objects of Sets_* are pointed sets.
- *Morphisms.* The morphisms of Sets_* are morphisms of pointed sets.
- *Identities.* For each $(X, x_0) \in \mathrm{Obj}(\mathsf{Sets}_*)$, the unit map

  $$\mathbb{1}^{\mathsf{Sets}_*}_{(X,x_0)} : \mathrm{pt} \to \mathsf{Sets}_*((X, x_0), (X, x_0))$$

  of Sets_* at (X, x_0) is defined by[6]

  $$\mathrm{id}^{\mathsf{Sets}_*}_{(X,x_0)} \overset{\text{def}}{=} \mathrm{id}_X .$$

denoted $(\mathbb{F}_1, 0)$.

[4] *Further Terminology:* Also called a **pointed function**.

[5] *Further Terminology:* In the context of monoids with zero as models for $\mathbb{F}_1$-algebras, morphisms of pointed sets are also called **morphism of $\mathbb{F}_1$-modules**.

[6] Note that id_X is indeed a morphism of pointed sets, as we have $\mathrm{id}_X(x_0) = x_0$.

- *Composition.* For each $(X,x_0),(Y,y_0),(Z,z_0) \in \mathrm{Obj}(\mathsf{Sets}_*)$, the composition map

$$\circ^{\mathsf{Sets}_*}_{(X,x_0),(Y,y_0),(Z,z_0)} \colon \mathsf{Sets}_*((Y,y_0),(Z,z_0)) \times \mathsf{Sets}_*((X,x_0),(Y,y_0)) \to \mathsf{Sets}_*((X,x_0),(Z,z_0))$$

  of Sets_* at $((X,x_0),(Y,y_0),(Z,z_0))$ is defined by[7]

$$g \circ^{\mathsf{Sets}_*}_{(X,x_0),(Y,y_0),(Z,z_0)} f \overset{\text{def}}{=} g \circ f.$$

6.1.4 Elementary Properties of Pointed Sets

Proposition 6.1.4.1.1. Let (X,x_0) be a pointed set.

1. *Completeness.* The category Sets_* of pointed sets and morphisms between them is complete, having in particular:
   (a) Products, described as in Definition 6.2.3.1.1.
   (b) Pullbacks, described as in Definition 6.2.4.1.1.
   (c) Equalisers, described as in Definition 6.2.5.1.1.

2. *Cocompleteness.* The category Sets_* of pointed sets and morphisms between them is cocomplete, having in particular:
   (a) Coproducts, described as in Definition 6.3.3.1.1.
   (b) Pushouts, described as in Definition 6.3.4.1.1;
   (c) Coequalisers, described as in Definition 6.3.5.1.1.

3. *Failure To Be Cartesian Closed.* The category Sets_* is not Cartesian closed.[8]

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[7]Note that the composition of two morphisms of pointed sets is indeed a morphism of pointed sets, as we have

$$\begin{aligned} g(f(x_0)) &= g(y_0) \\ &= z_0, \end{aligned}$$

$$\begin{tikzcd} & \mathrm{pt} \arrow[dl, "{[x_0]}"'] \arrow[d, "{[y_0]}"] \arrow[dr, "{[z_0]}"] & \\ X \arrow[r, "f"'] & Y \arrow[r, "g"'] & Z \end{tikzcd}$$

[8]The category Sets_* does admit a natural monoidal closed structure, however; see Tensor Products of Pointed Sets.

4. *Morphisms From the Monoidal Unit.* We have a bijection of sets[9]

$$\mathsf{Sets}_*(S^0, X) \cong X,$$

natural in $(X, x_0) \in \mathrm{Obj}(\mathsf{Sets}_*)$, internalising also to an isomorphism of pointed sets

$$\mathbf{Sets}_*(S^0, X) \cong (X, x_0),$$

again natural in $(X, x_0) \in \mathrm{Obj}(\mathsf{Sets}_*)$.

5. *Relation to Partial Functions.* We have an equivalence of categories[10]

$$\mathsf{Sets}_* \overset{\text{eq.}}{\cong} \mathsf{Sets}^{\text{part.}}$$

between the category of pointed sets and pointed functions between them and the category of sets and partial functions between them, where:

(a) *From Pointed Sets to Sets With Partial Functions.* The equivalence

$$\xi \colon \mathsf{Sets}_* \xrightarrow{\cong} \mathsf{Sets}^{\text{part.}}$$

sends:

i. A pointed set (X, x_0) to X.

ii. A pointed function

$$f \colon (X, x_0) \to (Y, y_0)$$

to the partial function

$$\xi_f \colon X \to Y$$

defined on $f^{-1}(Y \setminus y_0)$ and given by

$$\xi_f(x) \overset{\text{def}}{=} f(x)$$

for each $x \in f^{-1}(Y \setminus y_0)$.

(b) *From Sets With Partial Functions to Pointed Sets.* The equivalence

$$\xi^{-1} \colon \mathsf{Sets}^{\text{part.}} \xrightarrow{\cong} \mathsf{Sets}_*$$

sends:

[9]In other words, the forgetful functor

$$忘 \colon \mathsf{Sets}_* \to \mathsf{Sets}$$

defined on objects by sending a pointed set to its underlying set is corepresentable by S^0.

[10] *Warning:* This is not an isomorphism of categories, only an equivalence.

i. A set X is to the pointed set $(X, \star)$ with $\star$ an element that is not in X.

ii. A partial function

$$f : X \to Y$$

defined on $U \subset X$ to the pointed function

$$\xi_f^{-1} : (X, x_0) \to (Y, y_0)$$

defined by

$$\xi_f(x) \overset{\text{def}}{=} \begin{cases} f(x) & \text{if } x \in U, \\ y_0 & \text{otherwise.} \end{cases}$$

for each $x \in X$.

Proof. *Item 1, Completeness*: This follows from (the proofs) of Definitions 6.2.3.1.1, 6.2.4.1.1 and 6.2.5.1.1 and **??**.
Item 2, Cocompleteness: This follows from (the proofs) of Definitions 6.3.3.1.1, 6.3.4.1.1 and 6.3.5.1.1 and **??**.
Item 3, Failure To Be Cartesian Closed: See [`MSE 2855868`].
Item 4, Morphisms From the Monoidal Unit: Since a morphism from S^0 to a pointed set (X, x_0) sends $0 \in S^0$ to x_0 and then can send $1 \in S^0$ to any element of X, we obtain a bijection between pointed maps $S^0 \to X$ and the elements of X.

The isomorphism then

$$\mathsf{Sets}_*(S^0, X) \cong (X, x_0)$$

follows by noting that $\Delta_{x_0} : S^0 \to X$, the basepoint of $\mathsf{Sets}_*(S^0, X)$, corresponds to the pointed map $S^0 \to X$ picking the element x_0 of X, and thus we see that the bijection between pointed maps $S^0 \to X$ and elements of X is compatible with basepoints, lifting to an isomorphism of pointed sets.
Item 5, Relation to Partial Functions: See [`MSE 884460`]. □

6.1.5 Active and Inert Morphisms of Pointed Sets

Definition 6.1.5.1.1. Let $f : (X, x_0) \to (Y, y_0)$ be a morphism of pointed sets.

1. The morphism f is **active** if $f^{-1}(y_0) = x_0$.

2. The morphism f is **inert** if, for each $y \in Y$, the set $f^{-1}(y)$ has exactly one element.

Notation 6.1.5.1.2. We write $\mathsf{Sets}_*^{\mathrm{actv}}$ for the wide subcategory of Sets_* spanned by pointed sets and the active maps between them.

Example 6.1.5.1.3. Here are some examples of active and inert maps of pointed sets.

1. The map $\mu\colon \langle 2\rangle \to \langle 1\rangle$ given by

$$\begin{tikzcd}[row sep=small] 1 \arrow[r, mapsto] & 1 \\ 2 \arrow[ur, mapsto] & \\ * \arrow[r, mapsto] & * \end{tikzcd}$$

   is active but not inert.

2. The map $f\colon \langle 2\rangle \to \langle 2\rangle$ given by

$$\begin{tikzcd}[row sep=small] 1 \arrow[r, mapsto] & 1 \\ 2 \arrow[dr, mapsto] & 2 \\ * \arrow[r, mapsto] & * \end{tikzcd}$$

   is inert but not active.

3. The map $f\colon \langle 3\rangle \to \langle 1\rangle$ given by

$$\begin{tikzcd}[row sep=small] 1 \arrow[r, mapsto] & 1 \\ 2 \arrow[ur, mapsto] & \\ 3 \arrow[dr, mapsto] & \\ * \arrow[r, mapsto] & * \end{tikzcd}$$

   is neither inert nor active. However, it factors as $f = a \circ i$, where

$$\begin{aligned} i &\colon \langle 3\rangle \to \langle 2\rangle, \\ a &\colon \langle 2\rangle \to \langle 1\rangle \end{aligned}$$

   are the morphisms of pointed sets given by

$$\begin{tikzcd}[row sep=small] 1 \arrow[r, mapsto] & 1 & 1 \arrow[r, mapsto] & 1 \\ 2 \arrow[r, mapsto] & 2 & 2 \arrow[ur, mapsto] & \\ 3 \arrow[dr, mapsto] & & * \arrow[r, mapsto] & *, \\ * \arrow[r, mapsto] & * & & \end{tikzcd}$$

   with i being inert and a being active.

Proposition 6.1.5.1.4. Let (X, x_0) and (Y, y_0) be pointed sets.

1. *Active-Inert Factorisation.* Every morphism of pointed sets $f\colon (X, x_0) \to (Y, y_0)$ factors uniquely as

   $$f = a \circ i,$$

   where:

   (a) The map $i\colon (X, x_0) \to (K, k_0)$ is an inert morphism of pointed sets

   (b) The map $a\colon (K, k_0) \to (Y, y_0)$ is an active morphism of pointed sets.

   Moreover, this determines an orthogonal factorisation system in Sets_*.

Proof. *Item 1, Active-Inert Factorisation*: Let $f\colon X \to Y$ be a morphism of pointed sets. We can factor f as

$$X \xrightarrow{i} K \xrightarrow{a} Y,$$

where:

- K is the pointed set given by

  $$\begin{aligned} K &= \{x \in X \mid f(x) \neq y_0\} \cup \{x_0\} \\ &= (X \setminus f^{-1}(y_0)) \cup \{x_0\}; \end{aligned}$$

- $i\colon X \to K$ is the inert morphism of pointed sets given by

  $$i(x) \stackrel{\text{def}}{=} \begin{cases} x & \text{if } x \in K, \\ x_0 & \text{otherwise} \end{cases}$$

  for each $x \in X$;

- $a\colon K \to Y$ is the active morphism of pointed sets given by

  $$a(x) \stackrel{\text{def}}{=} f(x)$$

  for each $x \in K$.

Next, let

$$\begin{tikzcd} X \arrow[r, "i"] \arrow[d, "f"'] & Y \arrow[d, "g"] \\ A \arrow[r, "a"'] & B \end{tikzcd}$$

be a commutative diagram in Sets_*. Consider the morphism $\phi\colon Y \to A$ given by

$$\phi(y) = f(i^{-1}(y))$$

for each $y \in Y$ (which is well-defined since, as i is inert, $i^{-1}(y)$ is a singleton for all $y \in Y$). We claim that ϕ is the unique diagonal filler in the diagram

$$\begin{tikzcd} X \arrow[r, "i"] \arrow[d, "f"'] & Y \arrow[d, "g"] \arrow[dl, dashed, "\exists!"', "\phi"] \\ A \arrow[r, "a"'] & B. \end{tikzcd}$$

Indeed, this diagram commutes, as we have

$$\begin{aligned} [\phi \circ i](x) &\overset{\text{def}}{=} \phi(i(x)) \\ &\overset{\text{def}}{=} f(i^{-1}(i(x))) \\ &= f(x) \end{aligned}$$

for each $x \in X$ and

$$\begin{aligned} [a \circ \phi](y) &\overset{\text{def}}{=} a(\phi(y)) \\ &\overset{\text{def}}{=} a(f(i^{-1}(y))) \\ &\overset{\text{def}}{=} [a \circ f](i^{-1}(y)) \\ &= [g \circ i](i^{-1}(y)) \\ &\overset{\text{def}}{=} g(i(i^{-1}(y))) \\ &\overset{\text{def}}{=} g(y) \end{aligned}$$

for each $y \in Y$. Moreover, given another morphism ψ such that the diagram

$$\begin{tikzcd} X \arrow[r, "i"] \arrow[d, "f"'] & Y \arrow[d, "g"] \arrow[dl, dashed, "\psi"'] \\ A \arrow[r, "a"'] & B \end{tikzcd}$$

commutes, it follows that we must have $\psi = \phi$, since, given $y \in Y$, there exists a unique $x \in X$ such that $i(x) = y$, so we have

$$\begin{aligned} \psi(y) &= \psi(i(x)) \\ &= f(x) \\ &= f(i^{-1}(y)) \\ &\overset{\text{def}}{=} \phi(y). \end{aligned}$$

This finishes the proof. □

6.2 Limits of Pointed Sets

6.2.1 The Terminal Pointed Set

Definition 6.2.1.1.1. The **terminal pointed set** is the terminal object of Sets_* as in Limits and Colimits, **??**.

Construction 6.2.1.1.2. Concretely, the **terminal pointed set** is the pair $((\mathrm{pt}, \star), \{!_X\}_{(X,x_0)\in\mathrm{Obj}(\mathsf{Sets}_*)})$ consisting of:

- *The Limit.* The pointed set $(\mathrm{pt}, \star)$.
- *The Cone.* The collection of morphisms of pointed sets

  $$\{!_X \colon (X, x_0) \to (\mathrm{pt}, \star)\}_{(X,x_0)\in\mathrm{Obj}(\mathsf{Sets})}$$

  defined by

  $$!_X(x) \stackrel{\mathrm{def}}{=} \star$$

  for each $x \in X$ and each $(X, x_0) \in \mathrm{Obj}(\mathsf{Sets})$.

Proof. We claim that $(\mathrm{pt}, \star)$ is the terminal object of Sets_*. Indeed, suppose we have a diagram of the form

$$(X, x_0) \qquad (\mathrm{pt}, \star)$$

in Sets_*. Then there exists a unique morphism of pointed sets

$$\phi \colon (X, x_0) \to (\mathrm{pt}, \star)$$

making the diagram

$$(X, x_0) \underset{\exists!}{\overset{\phi}{\dashrightarrow}} (\mathrm{pt}, \star)$$

commute, namely $!_X$. □

6.2.2 Products of Families of Pointed Sets

Let $\{(X_i, x_0^i)\}_{i\in I}$ be a family of pointed sets.

Definition 6.2.2.1.1. The **product of** $\{(X_i, x_0^i)\}_{i\in I}$ is the product of $\{(X_i, x_0^i)\}_{i\in I}$ in Sets_* as in Limits and Colimits, **??**.

Construction 6.2.2.1.2. Concretely, the **product of** $\{(X_i, x_0^i)\}_{i\in I}$ is the pair $((\prod_{i\in I} X_i, (x_0^i)_{i\in I}), \{\mathrm{pr}_i\}_{i\in I})$ consisting of:

- *The Limit.* The pointed set $(\prod_{i\in I} X_i, (x_0^i)_{i\in I})$.
- *The Cone.* The collection

  $$\left\{\mathrm{pr}_i \colon (\prod_{i\in I} X_i, (x_0^i)_{i\in I}) \to (X_i, x_0^i)\right\}_{i\in I}$$

  of maps given by

  $$\mathrm{pr}_i((x_j)_{j\in I}) \stackrel{\text{def}}{=} x_i$$

  for each $(x_j)_{j\in I} \in \prod_{i\in I} X_i$ and each $i \in I$.

Proof. We claim that $(\prod_{i\in I} X_i, (x_0^i)_{i\in I})$ is the categorical product of $\left\{(X_i, x_0^i)\right\}_{i\in I}$ in Sets_*. Indeed, suppose we have, for each $i \in I$, a diagram of the form

$$\begin{tikzcd} & (P,*) \arrow[dr, "p_i"] & \\ (\prod_{i\in I} X_i, (x_0^i)_{i\in I}) \arrow[rr, "\mathrm{pr}_i"'] & & (X_i, x_0^i) \end{tikzcd}$$

in Sets_*. Then there exists a unique morphism of pointed sets

$$\phi \colon (P,*) \to (\prod_{i\in I} X_i, (x_0^i)_{i\in I})$$

making the diagram

$$\begin{tikzcd} (P,*) \arrow[d, dashed, "\phi"', "\exists!"] \arrow[dr, "p_i"] & \\ (\prod_{i\in I} X_i, (x_0^i)_{i\in I}) \arrow[r, "\mathrm{pr}_i"'] & (X_i, x_0^i) \end{tikzcd}$$

commute, being uniquely determined by the condition $\mathrm{pr}_i \circ \phi = p_i$ for each $i \in I$ via

$$\phi(x) = (p_i(x))_{i\in I}$$

for each $x \in P$. Note that this is indeed a morphism of pointed sets, as we have

$$\begin{aligned} \phi(*) &= (p_i(*))_{i\in I} \\ &= (x_0^i)_{i\in I}, \end{aligned}$$

where we have used that p_i is a morphism of pointed sets for each $i \in I$. $\square$

Proposition 6.2.2.1.3. Let $\{(X_i, x_0^i)\}_{i\in I}$ be a family of pointed sets.

1. *Functoriality.* The assignment $\{(X_i, x_0^i)\}_{i\in I} \mapsto (\prod_{i\in I} X_i, (x_0^i)_{i\in I})$ defines a functor
$$\prod_{i\in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}_*) \to \mathsf{Sets}_*.$$

Proof. *Item 1*, *Functoriality*: This follows from Limits and Colimits, **??** of **??**. □

6.2.3 Products

Let (X, x_0) and (Y, y_0) be pointed sets.

Definition 6.2.3.1.1. The **product of** (X, x_0) **and** (Y, y_0) is the product of (X, x_0) and (Y, y_0) in Sets_* as in Limits and Colimits, **??**.

Construction 6.2.3.1.2. Concretely, the **product of** (X, x_0) **and** (Y, y_0) is the pair consisting of:

- *The Limit.* The pointed set $(X \times Y, (x_0, y_0))$.
- *The Cone.* The morphisms of pointed sets
$$\mathrm{pr}_1 \colon (X \times Y, (x_0, y_0)) \to (X, x_0),$$
$$\mathrm{pr}_2 \colon (X \times Y, (x_0, y_0)) \to (Y, y_0)$$
defined by
$$\mathrm{pr}_1(x, y) \overset{\mathrm{def}}{=} x,$$
$$\mathrm{pr}_2(x, y) \overset{\mathrm{def}}{=} y$$
for each $(x, y) \in X \times Y$.

Proof. We claim that $(X \times Y, (x_0, y_0))$ is the categorical product of (X, x_0) and (Y, y_0) in Sets_*. Indeed, suppose we have a diagram of the form

$$\begin{tikzcd}
 & (P, *) \arrow[dl, bend right, "p_1"'] \arrow[dr, bend left, "p_2"] & \\
(X, x_0) & (X \times Y, (x_0, y_0)) \arrow[l, two heads, "\mathrm{pr}_1"] \arrow[r, two heads, "\mathrm{pr}_2"'] & (Y, y_0)
\end{tikzcd}$$

in Sets_*. Then there exists a unique morphism of pointed sets

$$\phi\colon (P,*) \to (X\times Y,(x_0,y_0))$$

making the diagram

$$\begin{tikzcd}
 & (P,*) \arrow[dl, bend right, "p_1"'] \arrow[d, dashed, "\phi", "\exists!"'] \arrow[dr, bend left, "p_2"] & \\
(X,x_0) & (X\times Y,(x_0,y_0)) \arrow[l, two heads, "\mathrm{pr}_1"] \arrow[r, two heads, "\mathrm{pr}_2"'] & (Y,y_0)
\end{tikzcd}$$

commute, being uniquely determined by the conditions

$$\begin{aligned}\mathrm{pr}_1\circ\phi &= p_1,\\ \mathrm{pr}_2\circ\phi &= p_2\end{aligned}$$

via

$$\phi(x) = (p_1(x),p_2(x))$$

for each $x\in P$. Note that this is indeed a morphism of pointed sets, as we have

$$\begin{aligned}\phi(*) &= (p_1(*),p_2(*))\\ &= (x_0,y_0),\end{aligned}$$

where we have used that p_1 and p_2 are morphisms of pointed sets. □

Proposition 6.2.3.1.3. Let (X,x_0), (Y,y_0), and (Z,z_0) be pointed sets.

1. *Functoriality.* The assignments

$$(X,x_0),(Y,y_0),((X,x_0),(Y,y_0)) \mapsto (X\times Y,(x_0,y_0))$$

   define functors

$$\begin{aligned} A\times -\colon &\quad \mathsf{Sets}_* \to \mathsf{Sets}_*,\\ -\times B\colon &\quad \mathsf{Sets}_* \to \mathsf{Sets}_*,\\ -_1\times -_2\colon &\quad \mathsf{Sets}_*\times\mathsf{Sets}_* \to \mathsf{Sets}_*,\end{aligned}$$

   defined in the same way as the functors of Constructions With Sets, Item 1 of Definition 4.1.3.1.3.

2. *Lack of Adjointness.* The functors $X \times -$ and $- \times Y$ do not admit right adjoints.

3. *Associativity.* We have an isomorphism of pointed sets

$$((X \times Y) \times Z, ((x_0, y_0), z_0)) \cong (X \times (Y \times Z), (x_0, (y_0, z_0)))$$

natural in $(X, x_0), (Y, y_0), (Z, z_0) \in \mathrm{Obj}(\mathsf{Sets}_*)$.

4. *Unitality.* We have isomorphisms of pointed sets

$$\begin{aligned}(\mathrm{pt}, \star) \times (X, x_0) &\cong (X, x_0), \\ (X, x_0) \times (\mathrm{pt}, \star) &\cong (X, x_0),\end{aligned}$$

natural in $(X, x_0) \in \mathrm{Obj}(\mathsf{Sets}_*)$.

5. *Commutativity.* We have an isomorphism of pointed sets

$$(X \times Y, (x_0, y_0)) \cong (Y \times X, (y_0, x_0)),$$

natural in $(X, x_0), (Y, y_0) \in \mathrm{Obj}(\mathsf{Sets}_*)$.

6. *Symmetric Monoidality.* The triple $(\mathsf{Sets}_*, \times, (\mathrm{pt}, \star))$ is a symmetric monoidal category.

Proof. *Item 1, Functoriality*: This is a special case of functoriality of limits, Limits and Colimits, **??** of **??**.
Item 2, Lack of Adjointness: See [MSE 2855868].
Item 3, Associativity: This follows from Constructions With Sets, Item 4 of Definition 4.1.3.1.3.
Item 4, Unitality: This follows from Constructions With Sets, Item 5 of Definition 4.1.3.1.3.
Item 5, Commutativity: This follows from Constructions With Sets, Item 6 of Definition 4.1.3.1.3.
Item 6, Symmetric Monoidality: This follows from Constructions With Sets, Item 14 of Definition 4.1.3.1.3. □

6.2.4 Pullbacks

Let (X, x_0), (Y, y_0), and (Z, z_0) be pointed sets and let $f : (X, x_0) \to (Z, z_0)$ and $g : (Y, y_0) \to (Z, z_0)$ be morphisms of pointed sets.

Definition 6.2.4.1.1. The **pullback of** (X, x_0) **and** (Y, y_0) **over** (Z, z_0) **along** (f, g) is the pullback of (X, x_0) and (Y, y_0) over (Z, z_0) along (f, g) in Sets_* as in Limits and Colimits, **??**.

Construction 6.2.4.1.2. Concretely, the **pullback of** (X, x_0) **and** (Y, y_0) **over** (Z, z_0) **along** (f, g) is the pair consisting of:

- *The Limit.* The pointed set $(X \times_Z Y, (x_0, y_0))$.
- *The Cone.* The morphisms of pointed sets

  $$\begin{aligned} \mathrm{pr}_1 \colon (X \times_Z Y, (x_0, y_0)) &\to (X, x_0), \\ \mathrm{pr}_2 \colon (X \times_Z Y, (x_0, y_0)) &\to (Y, y_0) \end{aligned}$$

  defined by

  $$\begin{aligned} \mathrm{pr}_1(x, y) &\stackrel{\text{def}}{=} x, \\ \mathrm{pr}_2(x, y) &\stackrel{\text{def}}{=} y \end{aligned}$$

  for each $(x, y) \in X \times_Z Y$.

Proof. We claim that $X \times_Z Y$ is the categorical pullback of (X, x_0) and (Y, y_0) over (Z, z_0) with respect to (f, g) in Sets_*. First we need to check that the relevant pullback diagram commutes, i.e. that we have

$$f \circ \mathrm{pr}_1 = g \circ \mathrm{pr}_2, \qquad \begin{tikzcd} (X \times_Z Y, (x_0, y_0)) \arrow[r, "\mathrm{pr}_2"] \arrow[d, "\mathrm{pr}_1"'] & (Y, y_0) \arrow[d, "g"] \\ (X, x_0) \arrow[r, "f"'] & (Z, z_0). \end{tikzcd}$$

Indeed, given $(x, y) \in X \times_Z Y$, we have

$$\begin{aligned} [f \circ \mathrm{pr}_1](x, y) &= f(\mathrm{pr}_1(x, y)) \\ &= f(x) \\ &= g(y) \\ &= g(\mathrm{pr}_2(x, y)) \\ &= [g \circ \mathrm{pr}_2](x, y), \end{aligned}$$

where $f(x) = g(y)$ since $(x, y) \in X \times_Z Y$. Next, we prove that $X \times_Z Y$ satisfies the universal property of the pullback. Suppose we have a diagram of the form

$$\begin{tikzcd} (P, *) \arrow[drr, bend left, "p_2"] \arrow[ddr, bend right, "p_1"'] & & \\ & (X \times_Z Y, (x_0, y_0)) \arrow[r, dashed, "\mathrm{pr}_2"] \arrow[d, dashed, "\mathrm{pr}_1"] \arrow[dr, phantom, "\lrcorner", very near start] & (Y, y_0) \arrow[d, "g"] \\ & (X, x_0) \arrow[r, "f"'] & (Z, z_0) \end{tikzcd}$$

in Sets_*. Then there exists a unique morphism of pointed sets

$$\phi\colon (P,*) \to (X\times_Z Y,(x_0,y_0))$$

making the diagram

$$\begin{tikzcd}
(P,*) \arrow[rd, dashed, "\phi", "\exists!"'] \arrow[rrd, bend left, "p_2"] \arrow[rdd, bend right, "p_1"'] & & \\
& (X\times_Z Y,(x_0,y_0)) \arrow[r, "\mathrm{pr}_2"] \arrow[d, "\mathrm{pr}_1"'] \arrow[rd, phantom, "\lrcorner", very near start] & (Y,y_0) \arrow[d, "g"] \\
& (X,x_0) \arrow[r, "f"'] & (Z,z_0)
\end{tikzcd}$$

commute, being uniquely determined by the conditions

$$\begin{aligned}
\mathrm{pr}_1\circ\phi &= p_1,\\
\mathrm{pr}_2\circ\phi &= p_2
\end{aligned}$$

via

$$\phi(x)=(p_1(x),p_2(x))$$

for each $x\in P$, where we note that $(p_1(x),p_2(x))\in X\times Y$ indeed lies in $X\times_Z Y$ by the condition

$$f\circ p_1=g\circ p_2,$$

which gives

$$f(p_1(x))=g(p_2(x))$$

for each $x\in P$, so that $(p_1(x),p_2(x))\in X\times_Z Y$. Lastly, we note that ϕ is indeed a morphism of pointed sets, as we have

$$\begin{aligned}
\phi(*) &= (p_1(*),p_2(*))\\
&= (x_0,y_0),
\end{aligned}$$

where we have used that p_1 and p_2 are morphisms of pointed sets. □

Proposition 6.2.4.1.3. Let (X,x_0), (Y,y_0), (Z,z_0), and (A,a_0) be pointed sets.

1. *Functoriality.* The assignment $(X,Y,Z,f,g)\mapsto X\times_{f,Z,g}Y$ defines a functor

$$-_1\times_{-_3}-_1\colon \mathsf{Fun}(\mathcal{P},\mathsf{Sets}_*)\to\mathsf{Sets}_*,$$

where $\mathcal{P}$ is the category that looks like this:

$$
\begin{tikzcd}
 & \bullet \arrow[d] \\
\bullet \arrow[r] & \bullet.
\end{tikzcd}
$$

In particular, the action on morphisms of $-_1 \times_{-_3} -_1$ is given by sending a morphism

$$
\begin{tikzcd}
X \times_Z Y \arrow[rr] \arrow[dd] \arrow[dr, phantom, very near start, "\lrcorner"] & & Y \arrow[dd, "g"] \arrow[dr, "\psi"] & \\
 & X' \times_{Z'} Y' \arrow[rr] \arrow[dd] \arrow[dr, phantom, very near start, "\lrcorner"] & & Y' \arrow[dd, "g'"] \\
X \arrow[rr, "f"] \arrow[dr, "\phi"'] & & Z \arrow[dr, "\chi"] & \\
 & X' \arrow[rr, "f'"'] & & Z'
\end{tikzcd}
$$

in $\mathsf{Fun}(\mathcal{P}, \mathsf{Sets}_*)$ to the morphism of pointed sets

$$\xi \colon (X \times_Z Y, (x_0, y_0)) \xrightarrow{\exists!} (X' \times_{Z'} Y', (x_0', y_0'))$$

given by

$$\xi(x, y) \stackrel{\text{def}}{=} (\phi(x), \psi(y))$$

for each $(x, y) \in X \times_Z Y$, which is the unique morphism of pointed sets making the diagram

$$
\begin{tikzcd}
X \times_Z Y \arrow[rr] \arrow[dd] \arrow[dr, dashed] & & Y \arrow[dd, "g"] \arrow[dr, "\psi"] & \\
 & X' \times_{Z'} Y' \arrow[rr] \arrow[dd] \arrow[dr, phantom, very near start, "\lrcorner"] & & Y' \arrow[dd, "g'"] \\
X \arrow[rr, "f"] \arrow[dr, "\phi"'] & & Z \arrow[dr, "\chi"] & \\
 & X' \arrow[rr, "f'"'] & & Z'
\end{tikzcd}
$$

commute.

2. *Associativity.* Given a diagram

$$\begin{tikzcd} X \arrow[dr, "f"'] & & Y \arrow[dl, "g"] \arrow[dr, "h"'] & & Z \arrow[dl, "k"] \\ & W & & V & \end{tikzcd}$$

in Sets_*, we have isomorphisms of pointed sets

$$(X \times_W Y) \times_V Z \cong (X \times_W Y) \times_Y (Y \times_V Z) \cong X \times_W (Y \times_V Z),$$

where these pullbacks are built as in the diagrams

$$\begin{tikzcd}[column sep=tiny] & & (X \times_W Y) \times_Y Z \arrow[dl] \arrow[ddrr] \arrow[d, phantom, "\lrcorner" rotate=-45] & & \\ & X \times_W Y \arrow[dl] \arrow[dr] \arrow[d, phantom, "\lrcorner" rotate=-45] & {} & & \\ X \arrow[dr, "f"'] & {} & Y \arrow[dl, "g"] \arrow[dr, "h"'] & & Z, \arrow[dl, "k"] \\ & W & & V & \end{tikzcd}
\quad
\begin{tikzcd}[column sep=tiny] & & (X \times_W Y) \times_Y (Y \times_V Z) \arrow[dl] \arrow[dr] \arrow[d, phantom, "\lrcorner" rotate=-45] & & \\ & X \times_W Y \arrow[dl] \arrow[dr] \arrow[d, phantom, "\lrcorner" rotate=-45] & {} & Y \times_V Z \arrow[dl] \arrow[dr] \arrow[d, phantom, "\lrcorner" rotate=-45] & \\ X \arrow[dr, "f"'] & {} & Y \arrow[dl, "g"] \arrow[dr, "h"'] & {} & Z, \arrow[dl, "k"] \\ & W & & V & \end{tikzcd}
\quad
\begin{tikzcd}[column sep=tiny] & & X \times_W (Y \times_V Z) \arrow[ddll] \arrow[dr] \arrow[d, phantom, "\lrcorner" rotate=-45] & & \\ & & {} & Y \times_V Z \arrow[dl] \arrow[dr] \arrow[d, phantom, "\lrcorner" rotate=-45] & \\ X \arrow[dr, "f"'] & & Y \arrow[dl, "g"] \arrow[dr, "h"'] & {} & Z. \arrow[dl, "k"] \\ & W & & V & \end{tikzcd}$$

3. *Unitality.* We have isomorphisms of pointed sets

$$\begin{tikzcd} A \arrow[r, equal] \arrow[d, "f"'] \arrow[dr, phantom, "\lrcorner", very near start] & A \arrow[d, "f"] \\ X \arrow[r, equal] & X \end{tikzcd}
\qquad
\begin{aligned} X \times_X A &\cong A, \\ A \times_X X &\cong A, \end{aligned}
\qquad
\begin{tikzcd} A \arrow[r, "f"] \arrow[d, equal] \arrow[dr, phantom, "\lrcorner", very near start] & X \arrow[d, equal] \\ X \arrow[r, "f"'] & X. \end{tikzcd}$$

4. *Commutativity.* We have an isomorphism of pointed sets

$$\begin{tikzcd} A \times_X B \arrow[r] \arrow[d] \arrow[dr, phantom, "\lrcorner", very near start] & B \arrow[d, "g"] \\ A \arrow[r, "f"'] & X, \end{tikzcd}
\qquad
A \times_X B \cong B \times_X A
\qquad
\begin{tikzcd} B \times_X A \arrow[r] \arrow[d] \arrow[dr, phantom, "\lrcorner", very near start] & A \arrow[d, "f"] \\ B \arrow[r, "g"'] & X. \end{tikzcd}$$

5. *Interaction With Products.* We have an isomorphism of pointed sets

$$X \times_{\mathrm{pt}} Y \cong X \times Y,
\qquad
\begin{tikzcd} X \times Y \arrow[r] \arrow[d] \arrow[dr, phantom, "\lrcorner", very near start] & Y \arrow[d, "!_Y"] \\ X \arrow[r, "!_X"'] & \mathrm{pt}. \end{tikzcd}$$

6. *Symmetric Monoidality*. The triple $(\mathsf{Sets}_*, \times_X, X)$ is a symmetric monoidal category.

Proof. *Item 1, Functoriality*: This is a special case of functoriality of co/limits, Limits and Colimits, **??** of **??**, with the explicit expression for ξ following from the commutativity of the cube pullback diagram.
Item 2, Associativity: This follows from Constructions With Sets, Item 4 of Definition 4.1.4.1.5.
Item 3, Unitality: This follows from Constructions With Sets, Item 6 of Definition 4.1.4.1.5.
Item 4, Commutativity: This follows from Constructions With Sets, Item 7 of Definition 4.1.4.1.5.
Item 5, Interaction With Products: This follows from Constructions With Sets, Item 10 of Definition 4.1.4.1.5.
Item 6, Symmetric Monoidality: This follows from Constructions With Sets, Item 11 of Definition 4.1.4.1.5. □

6.2.5 Equalisers

Let $f, g \colon (X, x_0) \rightrightarrows (Y, y_0)$ be morphisms of pointed sets.

Definition 6.2.5.1.1. The **equaliser of** (f, g) is the equaliser of f and g in Sets_* as in Limits and Colimits, **??**.

Construction 6.2.5.1.2. Concretely, the **equaliser of** (f, g) is the pair consisting of:

- *The Limit*. The pointed set $(\mathrm{Eq}(f, g), x_0)$.
- *The Cone*. The morphism of pointed sets

  $$\mathrm{eq}(f, g) \colon (\mathrm{Eq}(f, g), x_0) \hookrightarrow (X, x_0)$$

  given by the canonical inclusion $\mathrm{eq}(f, g) \hookrightarrow \mathrm{Eq}(f, g) \hookrightarrow X$.

Proof. We claim that $(\mathrm{Eq}(f, g), x_0)$ is the categorical equaliser of f and g in Sets_*. First we need to check that the relevant equaliser diagram commutes, i.e. that we have

$$f \circ \mathrm{eq}(f, g) = g \circ \mathrm{eq}(f, g),$$

which indeed holds by the definition of the set $\mathrm{Eq}(f, g)$. Next, we prove that $\mathrm{Eq}(f, g)$ satisfies the universal property of the equaliser. Suppose we have a diagram of the form

$$\begin{tikzcd}
(\mathrm{Eq}(f,g),x_0) \arrow[r, hook, "\mathrm{eq}(f,g)"] & (X,x_0) \arrow[r, shift left, "f"] \arrow[r, shift right, "g"'] & (Y,y_0) \\
(E,*) \arrow[ur, "e"'] & &
\end{tikzcd}$$

in Sets_*. Then there exists a unique morphism of pointed sets

$$\phi\colon (E,*) \to (\mathrm{Eq}(f,g),x_0)$$

making the diagram

$$\begin{tikzcd}
(\mathrm{Eq}(f,g),x_0) \arrow[r, hook, "\mathrm{eq}(f,g)"] & (X,x_0) \arrow[r, shift left, "f"] \arrow[r, shift right, "g"'] & (Y,y_0) \\
(E,*) \arrow[u, dashed, "\phi", "\exists!"'] \arrow[ur, "e"'] & &
\end{tikzcd}$$

commute, being uniquely determined by the condition

$$\mathrm{eq}(f,g) \circ \phi = e$$

via

$$\phi(x) = e(x)$$

for each $x \in E$, where we note that $e(x) \in A$ indeed lies in $\mathrm{Eq}(f,g)$ by the condition

$$f \circ e = g \circ e,$$

which gives

$$f(e(x)) = g(e(x))$$

for each $x \in E$, so that $e(x) \in \mathrm{Eq}(f,g)$. Lastly, we note that ϕ is indeed a morphism of pointed sets, as we have

$$\begin{aligned}
\phi(*) &= e(*) \\
&= x_0,
\end{aligned}$$

where we have used that e is a morphism of pointed sets. □

Proposition 6.2.5.1.3. Let (X,x_0) and (Y,y_0) be pointed sets and let $f,g,h\colon (X,x_0) \to (Y,y_0)$ be morphisms of pointed sets.

1. *Associativity.* We have isomorphisms of pointed sets

$$\underbrace{\mathrm{Eq}(f \circ \mathrm{eq}(g,h), g \circ \mathrm{eq}(g,h))}_{=\mathrm{Eq}(f \circ \mathrm{eq}(g,h), h \circ \mathrm{eq}(g,h))} \cong \mathrm{Eq}(f,g,h) \cong \underbrace{\mathrm{Eq}(f \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))}_{=\mathrm{Eq}(g \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))},$$

   where $\mathrm{Eq}(f,g,h)$ is the limit of the diagram

$$(X, x_0) \overset{f}{\underset{h}{\overset{g}{\rightrightarrows}}} (Y, y_0)$$

   in Sets_*, being explicitly given by

$$\mathrm{Eq}(f,g,h) \cong \{a \in A \mid f(a) = g(a) = h(a)\}.$$

2. *Unitality.* We have an isomorphism of pointed sets

$$\mathrm{Eq}(f,f) \cong X.$$

3. *Commutativity.* We have an isomorphism of pointed sets

$$\mathrm{Eq}(f,g) \cong \mathrm{Eq}(g,f).$$

Proof. *Item 1, Associativity*: This follows from Constructions With Sets, Item 1 of Definition 4.1.5.1.3.
Item 2, Unitality: This follows from Constructions With Sets, Item 4 of Definition 4.1.5.1.3.
Item 3, Commutativity: This follows from Constructions With Sets, Item 5 of Definition 4.1.5.1.3. □

6.3 Colimits of Pointed Sets

6.3.1 The Initial Pointed Set

Definition 6.3.1.1.1. The **initial pointed set** is the initial object of Sets_* as in Limits and Colimits, **??**.

Construction 6.3.1.1.2. Concretely, the **initial pointed set** is the pair $\left((\mathrm{pt}, \star), \{\iota_X\}_{(X,x_0)\in\mathrm{Obj}(\mathsf{Sets}_*)}\right)$ consisting of:

- *The Limit.* The pointed set $(\mathrm{pt}, \star)$.

- *The Cone.* The collection of morphisms of pointed sets

$$\{\iota_X\colon (\mathrm{pt},\star) \to (X,x_0)\}_{(X,x_0)\in\mathsf{Obj}(\mathsf{Sets})}$$

  defined by

$$\iota_X(\star) \overset{\text{def}}{=} x_0.$$

Proof. We claim that $(\mathrm{pt},\star)$ is the initial object of Sets_*. Indeed, suppose we have a diagram of the form

$$(\mathrm{pt},\star) \qquad (X,x_0)$$

in Sets_*. Then there exists a unique morphism of pointed sets

$$\phi\colon (\mathrm{pt},\star) \to (X,x_0)$$

making the diagram

$$(\mathrm{pt},\star) \xrightarrow[\exists!]{\phi} (X,x_0)$$

commute, namely ι_X. □

6.3.2 Coproducts of Families of Pointed Sets

Let $\left\{(X_i,x_0^i)\right\}_{i\in I}$ be a family of pointed sets.

Definition 6.3.2.1.1. The **coproduct of the family** $\left\{(X_i,x_0^i)\right\}_{i\in I}$[11] is the coproduct of $\left\{(X_i,x_0^i)\right\}_{i\in I}$ in Sets_* as in Limits and Colimits, **??**.

Construction 6.3.2.1.2. Concretely, the **coproduct of the family** $\left\{(X_i,x_0^i)\right\}_{i\in I}$ is the pair $((\bigvee_{i\in I}X_i,p_0),\left\{\mathrm{inj}_i\right\}_{i\in I})$ consisting of:

- *The Colimit.* The pointed set $(\bigvee_{i\in I}X_i,p_0)$ consisting of:
  - *The Underlying Set.* The set $\bigvee_{i\in I}X_i$ defined by

$$\bigvee_{i\in I} X_i \overset{\text{def}}{=} \Big(\coprod_{i\in I} X_i\Big)/\sim,$$

    where $\sim$ is the equivalence relation on $\coprod_{i\in I}X_i$ given by declaring

$$(i,x_0^i) \sim (j,x_0^j)$$

    for each $i,j\in I$.

[11] *Further Terminology:* Also called the **wedge sum of the family** $\left\{(X_i,x_0^i)\right\}_{i\in I}$.

- *The Basepoint.* The element p_0 of $\bigvee_{i \in I} X_i$ defined by

$$\begin{aligned} p_0 &\overset{\text{def}}{=} [(i, x_0^i)] \\ &= [(j, x_0^j)] \end{aligned}$$

for any $i, j \in I$.

- *The Cocone.* The collection

$$\left\{ \mathrm{inj}_i \colon (X_i, x_0^i) \to (\bigvee_{i \in I} X_i, p_0) \right\}_{i \in I}$$

of morphism of pointed sets given by

$$\mathrm{inj}_i(x) \overset{\text{def}}{=} (i, x)$$

for each $x \in X_i$ and each $i \in I$.

Proof. We claim that $(\bigvee_{i \in I} X_i, p_0)$ is the categorical coproduct of $\left\{(X_i, x_0^i)\right\}_{i \in I}$ in Sets_*. Indeed, suppose we have, for each $i \in I$, a diagram of the form

$$\begin{tikzcd} & (C, *) \\ (X_i, x_0^i) \arrow[ur, "\iota_i"] \arrow[r, "\mathrm{inj}_i"'] & (\bigvee_{i \in I} X_i, p_0) \end{tikzcd}$$

in Sets_*. Then there exists a unique morphism of pointed sets

$$\phi \colon (\bigvee_{i \in I} X_i, p_0) \to (C, *)$$

making the diagram

$$\begin{tikzcd} & (C, *) \\ (X_i, x_0^i) \arrow[ur, "\iota_i"] \arrow[r, "\mathrm{inj}_i"'] & (\bigvee_{i \in I} X_i, p_0) \arrow[u, dashed, "\phi", "\exists !"'] \end{tikzcd}$$

commute, being uniquely determined by the condition $\phi \circ \mathrm{inj}_i = \iota_i$ for each $i \in I$ via

$$\phi([(i, x)]) = \iota_i(x)$$

for each $[(i,x)] \in \bigvee_{i\in I} X_i$, where we note that ϕ is indeed a morphism of pointed sets, as we have

$$\begin{aligned}\phi(p_0) &= \iota_i([(i,x_0^i)])\\ &= *,\end{aligned}$$

as ι_i is a morphism of pointed sets. □

Proposition 6.3.2.1.3. Let $\{(X_i,x_0^i)\}_{i\in I}$ be a family of pointed sets.

1. *Functoriality.* The assignment $\{(X_i,x_0^i)\}_{i\in I} \mapsto (\bigvee_{i\in I} X_i, p_0)$ defines a functor

$$\bigvee_{i\in I}\colon \mathsf{Fun}(I_{\mathsf{disc}},\mathsf{Sets}_*) \longrightarrow \mathsf{Sets}_*.$$

Proof. Item 1, *Functoriality*: This follows from Limits and Colimits, **??** of **??**. □

6.3.3 Coproducts

Let (X,x_0) and (Y,y_0) be pointed sets.

Definition 6.3.3.1.1. The **coproduct of** (X,x_0) **and** (Y,y_0)[12] is the coproduct of (X,x_0) and (Y,y_0) in Sets_* as in Limits and Colimits, **??**.

Construction 6.3.3.1.2. Concretely, the **coproduct of** (X,x_0) **and** (Y,y_0), also called their **wedge sum**, is the pair consisting of:

- *The Colimit.* The pointed set $(X\vee Y,p_0)$ consisting of:
  - *The Underlying Set.* The set $X\vee Y$ defined by

$$\begin{aligned}(X\vee Y,p_0) &\overset{\text{def}}{=} (X,x_0)\textstyle\coprod(Y,y_0)\\ &\cong (X\textstyle\coprod_{\mathrm{pt}} Y,p_0)\\ &\cong (X\textstyle\coprod Y/\sim,p_0),\end{aligned} \qquad \begin{tikzcd} X\vee Y \arrow[dr, phantom, "\ulcorner", very near start] & Y \arrow[l] \\ X \arrow[u] & \mathrm{pt}, \arrow[l, hook', "{[x_0]}"] \arrow[u, hook, "{[y_0]}"'] \end{tikzcd}$$

    where $\sim$ is the equivalence relation on $X\coprod Y$ obtained by declaring $(0,x_0)\sim(1,y_0)$.
  - *The Basepoint.* The element p_0 of $X\vee Y$ defined by

$$\begin{aligned}p_0 &\overset{\text{def}}{=} [(0,x_0)]\\ &= [(1,y_0)].\end{aligned}$$

[12] *Further Terminology:* Also called the **wedge sum of** (X,x_0) **and** (Y,y_0).

- *The Cocone.* The morphisms of pointed sets

  $$\begin{aligned}\mathrm{inj}_1 &\colon (X, x_0) \to (X \vee Y, p_0),\\ \mathrm{inj}_2 &\colon (Y, y_0) \to (X \vee Y, p_0),\end{aligned}$$

  given by

  $$\begin{aligned}\mathrm{inj}_1(x) &\stackrel{\text{def}}{=} [(0, x)],\\ \mathrm{inj}_2(y) &\stackrel{\text{def}}{=} [(1, y)],\end{aligned}$$

  for each $x \in X$ and each $y \in Y$.

Proof. We claim that $(X \vee Y, p_0)$ is the categorical coproduct of (X, x_0) and (Y, y_0) in Sets_*. Indeed, suppose we have a diagram of the form

$$\begin{tikzcd}
 & (C, *) & \\
(X, x_0) \arrow[r, hook, "\mathrm{inj}_1"'] \arrow[ur, bend left, "\iota_1"] & (X \vee Y, p_0) & (Y, y_0) \arrow[l, hook', "\mathrm{inj}_2"] \arrow[ul, bend right, "\iota_2"']
\end{tikzcd}$$

in Sets. Then there exists a unique morphism of pointed sets

$$\phi \colon (X \vee Y, p_0) \to (C, *)$$

making the diagram

$$\begin{tikzcd}
 & (C, *) & \\
(X, x_0) \arrow[r, hook, "\mathrm{inj}_1"'] \arrow[ur, bend left, "\iota_1"] & (X \vee Y, p_0) \arrow[u, dashed, "\phi", "\exists!"'] & (Y, y_0) \arrow[l, hook', "\mathrm{inj}_2"] \arrow[ul, bend right, "\iota_2"']
\end{tikzcd}$$

commute, being uniquely determined by the conditions

$$\begin{aligned}\phi \circ \mathrm{inj}_X &= \iota_X,\\ \phi \circ \mathrm{inj}_Y &= \iota_Y\end{aligned}$$

via

$$\phi(z) = \begin{cases} \iota_X(x) & \text{if } z = [(0, x)] \text{ with } x \in X, \\ \iota_Y(y) & \text{if } z = [(1, y)] \text{ with } y \in Y \end{cases}$$

for each $z \in X \vee Y$, where we note that ϕ is indeed a morphism of pointed sets, as we have

$$\begin{aligned}\phi(p_0) &= \iota_X([(0, x_0)])\\ &= \iota_Y([(1, y_0)])\\ &= *,\end{aligned}$$

as ι_X and ι_Y are morphisms of pointed sets. □

Proposition 6.3.3.1.3. Let (X, x_0) and (Y, y_0) be pointed sets.

1. *Functoriality.* The assignments

$$(X, x_0), (Y, y_0), ((X, x_0), (Y, y_0)) \mapsto (X \vee Y, p_0)$$

   define functors

$$\begin{aligned}X \vee - \colon \mathsf{Sets}_* &\longrightarrow \mathsf{Sets}_*,\\ - \vee Y \colon \mathsf{Sets}_* &\longrightarrow \mathsf{Sets}_*,\\ -_1 \vee -_2 \colon \mathsf{Sets}_* \times \mathsf{Sets}_* &\longrightarrow \mathsf{Sets}_*.\end{aligned}$$

2. *Associativity.* We have an isomorphism of pointed sets

$$(X \vee Y) \vee Z \cong X \vee (Y \vee Z),$$

   natural in $(X, x_0), (Y, y_0), (Z, z_0) \in \mathsf{Sets}_*$.

3. *Unitality.* We have isomorphisms of pointed sets

$$\begin{aligned}(\mathrm{pt}, *) \vee (X, x_0) &\cong (X, x_0),\\ (X, x_0) \vee (\mathrm{pt}, *) &\cong (X, x_0),\end{aligned}$$

   natural in $(X, x_0) \in \mathsf{Sets}_*$.

4. *Commutativity.* We have an isomorphism of pointed sets

$$X \vee Y \cong Y \vee X,$$

   natural in $(X, x_0), (Y, y_0) \in \mathsf{Sets}_*$.

5. *Symmetric Monoidality.* The triple $(\mathsf{Sets}_*, \vee, \mathrm{pt})$ is a symmetric monoidal category.

6. *The Fold Map.* We have a natural transformation

$$\nabla\colon \vee \circ \Delta^{\mathsf{Cats}}_{\mathsf{Sets}_*} \Longrightarrow \mathrm{id}_{\mathsf{Sets}_*}, \qquad \begin{tikzcd} & \mathsf{Sets}_* \times \mathsf{Sets}_* \arrow[d, phantom, "\Downarrow\nabla"] \arrow[dr, "\vee"] & \\ \mathsf{Sets}_* \arrow[ur, "\Delta^{\mathsf{Cats}}_{\mathsf{Sets}_*}"] \arrow[rr, bend right, "\mathrm{id}_{\mathsf{Sets}_*}"'] & & \mathsf{Sets}_*, \end{tikzcd}$$

called the **fold map**, whose component

$$\nabla_X\colon X \vee X \to X$$

at X is given by

$$\nabla_X(p) \stackrel{\text{def}}{=} \begin{cases} x & \text{if } p = [(0,x)], \\ x & \text{if } p = [(1,x)] \end{cases}$$

for each $p \in X \vee X$.

Proof. *Item 1, Functoriality*: This follows from Limits and Colimits, **??** of **??**.
Item 2, Associativity: Omitted.
Item 3, Unitality: Omitted.
Item 4, Commutativity: Omitted.
Item 5, Symmetric Monoidality: Omitted.
Item 6, The Fold Map: Naturality for the transformation ∇ is the statement that, given a morphism of pointed sets $f\colon (X,x_0) \to (Y,y_0)$, we have

$$\nabla_Y \circ (f \vee f) = f \circ \nabla_X, \qquad \begin{tikzcd} X \vee X \arrow[r, "\nabla_X"] \arrow[d, "f \vee f"'] & X \arrow[d, "f"] \\ Y \vee Y \arrow[r, "\nabla_Y"'] & Y. \end{tikzcd}$$

Indeed, we have

$$\begin{aligned} [\nabla_Y \circ (f \vee f)]([(i,x)]) &= \nabla_Y([(i,f(x))]) \\ &= f(x) \\ &= f(\nabla_X([(i,x)])) \\ &= [f \circ \nabla_X]([(i,x)]) \end{aligned}$$

for each $[(i,x)] \in X \vee X$, and thus ∇ is indeed a natural transformation. □

6.3.4 Pushouts

Let (X, x_0), (Y, y_0), and (Z, z_0) be pointed sets and let $f\colon (Z, z_0) \to (X, x_0)$ and $g\colon (Z, z_0) \to (Y, y_0)$ be morphisms of pointed sets.

Definition 6.3.4.1.1. The **pushout of** (X, x_0) **and** (Y, y_0) **over** (Z, z_0) **along** (f, g) is the pushout of (X, x_0) and (Y, y_0) over (Z, z_0) along (f, g) in Sets_* as in Limits and Colimits, **??**.

Construction 6.3.4.1.2. Concretely, the **pushout of** (X, x_0) **and** (Y, y_0) **over** (Z, z_0) **along** (f, g) is the pair consisting of:

- *The Colimit.* The pointed set $(X \coprod_{f,Z,g} Y, p_0)$, where:
  - The set $X \coprod_{f,Z,g} Y$ is the pushout (of unpointed sets) of X and Y over Z with respect to f and g;
  - We have $p_0 = [x_0] = [y_0]$.
- *The Cocone.* The morphisms of pointed sets

$$\begin{aligned}\mathrm{inj}_1 &\colon (X, x_0) \to (X \textstyle\coprod_Z Y, p_0),\\ \mathrm{inj}_2 &\colon (Y, y_0) \to (X \textstyle\coprod_Z Y, p_0)\end{aligned}$$

  given by

$$\begin{aligned}\mathrm{inj}_1(x) &\stackrel{\text{def}}{=} [(0, x)]\\ \mathrm{inj}_2(y) &\stackrel{\text{def}}{=} [(1, y)]\end{aligned}$$

  for each $x \in X$ and each $y \in Y$.

Proof. Firstly, we note that indeed $[x_0] = [y_0]$, as we have

$$\begin{aligned}x_0 &= f(z_0),\\ y_0 &= g(z_0)\end{aligned}$$

since f and g are morphisms of pointed sets, with the relation $\sim$ on $X \coprod_Z Y$ then identifying $x_0 = f(z_0) \sim g(z_0) = y_0$.

We now claim that $(X \coprod_Z Y, p_0)$ is the categorical pushout of (X, x_0) and (Y, y_0) over (Z, z_0) with respect to (f, g) in Sets_*. First we need to check that the relevant pushout diagram commutes, i.e. that we have

$$\mathrm{inj}_1 \circ f = \mathrm{inj}_2 \circ g, \qquad \begin{tikzcd} (X \coprod_Z Y, p_0) & (Y, y_0) \arrow[l, "\mathrm{inj}_2"'] \\ (X, x_0) \arrow[u, "\mathrm{inj}_1"] & (Z, z_0). \arrow[l, "f"] \arrow[u, "g"'] \end{tikzcd}$$

Indeed, given $z \in Z$, we have

$$\begin{aligned}
[\mathrm{inj}_1 \circ f](z) &= \mathrm{inj}_1(f(z)) \\
&= [(0, f(z))] \\
&= [(1, g(z))] \\
&= \mathrm{inj}_2(g(z)) \\
&= [\mathrm{inj}_2 \circ g](z),
\end{aligned}$$

where $[(0, f(z))] = [(1, g(z))]$ by the definition of the relation $\sim$ on $X \coprod Y$ (the coproduct of unpointed sets of X and Y). Next, we prove that $X \coprod_Z Y$ satisfies the universal property of the pushout. Suppose we have a diagram of the form

$$\begin{tikzcd}
(P, *) & & \\
& (X \coprod_Z Y, p_0) & (Y, y_0) \arrow[l, "\mathrm{inj}_2"'] \arrow[ull, bend right, "\iota_2"'] \arrow[dl, phantom, "\ulcorner" very near end] \\
& (X, x_0) \arrow[u, "\mathrm{inj}_1"] \arrow[uul, bend left, "\iota_1"] & (Z, z_0) \arrow[l, "f"] \arrow[u, "g"']
\end{tikzcd}$$

in Sets_*. Then there exists a unique morphism of pointed sets

$$\phi \colon (X \coprod_Z Y, p_0) \to (P, *)$$

making the diagram

$$\begin{tikzcd}
(P, *) & & \\
& (X \coprod_Z Y, p_0) \arrow[ul, dashed, "\exists!"', "\phi"] & (Y, y_0) \arrow[l, "\mathrm{inj}_2"'] \arrow[ull, bend right, "\iota_2"'] \arrow[dl, phantom, "\ulcorner" very near end] \\
& (X, x_0) \arrow[u, "\mathrm{inj}_1"] \arrow[uul, bend left, "\iota_1"] & (Z, z_0) \arrow[l, "f"] \arrow[u, "g"']
\end{tikzcd}$$

commute, being uniquely determined by the conditions

$$\begin{aligned}
\phi \circ \mathrm{inj}_1 &= \iota_1, \\
\phi \circ \mathrm{inj}_2 &= \iota_2
\end{aligned}$$

via

$$\phi(p) = \begin{cases} \iota_1(x) & \text{if } x = [(0, x)], \\ \iota_2(y) & \text{if } x = [(1, y)] \end{cases}$$

for each $p \in X \coprod_Z Y$, where the well-definedness of ϕ is proven in the same way as in the proof of Constructions With Sets, Definition 4.2.4.1.1. Finally, we show that ϕ is indeed a morphism of pointed sets, as we have

$$\begin{aligned} \phi(p_0) &= \phi([(0, x_0)]) \\ &= \iota_1(x_0) \\ &= *, \end{aligned}$$

or alternatively

$$\begin{aligned} \phi(p_0) &= \phi([(1, y_0)]) \\ &= \iota_2(y_0) \\ &= *, \end{aligned}$$

where we use that ι_1 (resp. ι_2) is a morphism of pointed sets. □

Proposition 6.3.4.1.3. Let (X, x_0), (Y, y_0), (Z, z_0), and (A, a_0) be pointed sets.

1. *Functoriality.* The assignment $(X, Y, Z, f, g) \mapsto X \coprod_{f,Z,g} Y$ defines a functor

   $$-_1 \coprod_{-_3} -_1 \colon \mathsf{Fun}(\mathcal{P}, \mathsf{Sets}) \to \mathsf{Sets}_*,$$

   where $\mathcal{P}$ is the category that looks like this:

   $$\begin{tikzcd} & \bullet \\ \bullet & \bullet \arrow[l] \arrow[u] \end{tikzcd}.$$

   In particular, the action on morphisms of $-_1 \coprod_{-_3} -_1$ is given by sending a morphism

$$\begin{tikzcd} X \coprod_Z Y \arrow[dr, phantom, "\ulcorner"] & & Y \arrow[ll] \arrow[dr, "\psi"] & \\ & X' \coprod_{Z'} Y' \arrow[dr, phantom, "\ulcorner"] & & Y' \arrow[ll] \\ X \arrow[uu] \arrow[dr, "\phi"'] & & Z \arrow[ll, "f"'] \arrow[uu, "g"] \arrow[dr, "\chi"] & \\ & X' \arrow[uu] & & Z' \arrow[ll, "f'"] \arrow[uu, "g'"'] \end{tikzcd}$$

in $\mathsf{Fun}(\mathcal{P}, \mathsf{Sets}_*)$ to the morphism of pointed sets

$$\xi\colon (X \textstyle\coprod_Z Y, p_0) \xrightarrow{\exists!} (X' \textstyle\coprod_{Z'} Y', p_0')$$

given by

$$\xi(p) \overset{\text{def}}{=} \begin{cases} \phi(x) & \text{if } p = [(0,x)], \\ \psi(y) & \text{if } p = [(1,y)] \end{cases}$$

for each $p \in X \coprod_Z Y$, which is the unique morphism of pointed sets making the diagram

$$\begin{tikzcd}
X \coprod_Z Y \arrow[dr, dashed] & & Y \arrow[ll] \arrow[dr, "\psi"] & \\
& X' \coprod_{Z'} Y' & & Y' \arrow[ll] \\
X \arrow[uu] \arrow[dr, "\phi"'] & & Z \arrow[ll, "f"'] \arrow[uu, "g"] \arrow[dr, "\chi"] & \\
& X' \arrow[uu] & & Z' \arrow[ll, "f'"] \arrow[uu, "g'"']
\end{tikzcd}$$

commute.

2. *Associativity.* Given a diagram

$$\begin{tikzcd}
X & & Y & & Z \\
& W \arrow[ul, "f"] \arrow[ur, "g"'] & & V \arrow[ul, "h"] \arrow[ur, "k"'] &
\end{tikzcd}$$

in Sets, we have isomorphisms of pointed sets

$$(X \textstyle\coprod_W Y) \coprod_V Z \cong (X \coprod_W Y) \coprod_Y (Y \coprod_V Z) \cong X \coprod_W (Y \coprod_V Z),$$

where these pullbacks are built as in the diagrams

$$\begin{tikzcd}[column sep=tiny]
& & (X \coprod_W Y) \coprod_V Z & & \\
& X \coprod_W Y \arrow[ur] & & & \\
X \arrow[ur] & & Y \arrow[ul] & & Z, \arrow[uull] \\
& W \arrow[ul, "f"] \arrow[ur, "g"'] & & V \arrow[ul, "h"] \arrow[ur, "k"'] &
\end{tikzcd}
\quad
\begin{tikzcd}[column sep=tiny]
& & (X \coprod_W Y) \coprod_Y (Y \coprod_V Z) & & \\
& X \coprod_W Y \arrow[ur] & & Y \coprod_V Z \arrow[ul] & \\
X \arrow[ur] & & Y \arrow[ul] \arrow[ur] & & Z, \arrow[ul] \\
& W \arrow[ul, "f"] \arrow[ur, "g"'] & & V \arrow[ul, "h"] \arrow[ur, "k"'] &
\end{tikzcd}
\quad
\begin{tikzcd}[column sep=tiny]
& & X \coprod_W (Y \coprod_V Z) & & \\
& & & Y \coprod_V Z \arrow[ul] & \\
X \arrow[uurr] & & Y \arrow[ur] & & Z. \arrow[ul] \\
& W \arrow[ul, "f"] \arrow[ur, "g"'] & & V \arrow[ul, "h"] \arrow[ur, "k"'] &
\end{tikzcd}$$

3. *Unitality.* We have isomorphisms of sets

$$\begin{tikzcd} A \arrow[r, equal] & A \\ X \arrow[u, "f"] \arrow[r, equal] \arrow[ur, phantom, very near end, "\ulcorner"] & X \arrow[u, "f"'] \end{tikzcd} \qquad \begin{aligned} X \textstyle\coprod_X A &\cong A, \\ A \textstyle\coprod_X X &\cong A, \end{aligned} \qquad \begin{tikzcd} A & X \arrow[l, "f"'] \\ X \arrow[u, equal] \arrow[ur, phantom, very near end, "\ulcorner"] & X. \arrow[l, "f"] \arrow[u, equal] \end{tikzcd}$$

4. *Commutativity.* We have an isomorphism of sets

$$\begin{tikzcd} X \coprod_Z Y & Y \arrow[l] \\ X \arrow[u] \arrow[ur, phantom, very near end, "\ulcorner"] & Z, \arrow[l, "f"] \arrow[u, "g"'] \end{tikzcd} \qquad X \textstyle\coprod_Z Y \cong Y \textstyle\coprod_Z X \qquad \begin{tikzcd} Y \coprod_Z X & X \arrow[l] \\ Y \arrow[u] \arrow[ur, phantom, very near end, "\ulcorner"] & Z. \arrow[l, "g"] \arrow[u, "f"'] \end{tikzcd}$$

5. *Interaction With Coproducts.* We have

$$X \textstyle\coprod_{\mathrm{pt}} Y \cong X \vee Y, \qquad \begin{tikzcd} X \vee Y & Y \arrow[l] \\ X \arrow[u] \arrow[ur, phantom, very near end, "\ulcorner"] & \mathrm{pt}. \arrow[l, hook', "{[x_0]}"] \arrow[u, hook, "{[y_0]}"'] \end{tikzcd}$$

6. *Symmetric Monoidality.* The triple $(\mathsf{Sets}_*, \coprod_X, (X, x_0))$ is a symmetric monoidal category.

Proof. *Item 1, Functoriality*: This is a special case of functoriality of co/limits, Limits and Colimits, **??** of **??**, with the explicit expression for ξ following from the commutativity of the cube pushout diagram.
Item 2, Associativity: This follows from Constructions With Sets, Item 3 of Definition 4.2.4.1.6.
Item 3, Unitality: This follows from Constructions With Sets, Item 5 of Definition 4.2.4.1.6.
Item 4, Commutativity: This follows from Constructions With Sets, Item 6 of Definition 4.2.4.1.6.
Item 5, Interaction With Coproducts: Omitted.
Item 6, Symmetric Monoidality: Omitted. □

6.3.5 Coequalisers

Let $f, g\colon (X, x_0) \rightrightarrows (Y, y_0)$ be morphisms of pointed sets.

Definition 6.3.5.1.1. The **coequaliser of** (f, g) is the pointed set $(\mathrm{CoEq}(f, g), [y_0])$.

Construction 6.3.5.1.2. The **coequaliser of** (f, g) is the pair $((\mathrm{CoEq}(f, g), [y_0]), \mathrm{coeq}(f, g))$ consisting of:

- *The Colimit.* The pointed set $(\mathrm{CoEq}(f, g), [y_0])$, where $\mathrm{CoEq}(f, g)$ is the coequaliser of f and g as in Constructions With Sets, Definition 4.2.5.1.1.
- *The Cocone.* The map

$$\mathrm{coeq}(f, g)\colon Y \twoheadrightarrow (\mathrm{CoEq}(f, g), [y_0])$$

  given by the quotient map, as in Constructions With Sets, Item 2 of Definition 4.2.5.1.2.

Proof. We claim that $(\mathrm{CoEq}(f, g), [y_0])$ is the categorical coequaliser of f and g in Sets_*. First we need to check that the relevant coequaliser diagram commutes, i.e. that we have

$$\mathrm{coeq}(f, g) \circ f = \mathrm{coeq}(f, g) \circ g.$$

Indeed, we have

$$\begin{aligned}
[\mathrm{coeq}(f, g) \circ f](x) &\overset{\text{def}}{=} [\mathrm{coeq}(f, g)](f(x)) \\
&\overset{\text{def}}{=} [f(x)] \\
&= [g(x)] \\
&\overset{\text{def}}{=} [\mathrm{coeq}(f, g)](g(x)) \\
&\overset{\text{def}}{=} [\mathrm{coeq}(f, g) \circ g](x)
\end{aligned}$$

for each $x \in X$. Next, we prove that $\mathrm{CoEq}(f, g)$ satisfies the universal property of the coequaliser. Suppose we have a diagram of the form

$$\begin{tikzcd}
(X, x_0) \arrow[r, shift left, "f"] \arrow[r, shift right, "g"'] & (Y, y_0) \arrow[r, hook, "\mathrm{coeq}(f,g)"] \arrow[drr, "c"'] & (\mathrm{CoEq}(f, g), [y_0]) & \\
& & & (C, *)
\end{tikzcd}$$

in Sets. Then, since $c(f(a)) = c(g(a))$ for each $a \in A$, it follows from Conditions on Relations, Items 4 and 5 of Definition 10.6.2.1.3 that there exists a unique map $\phi \colon \mathrm{CoEq}(f,g) \xrightarrow{\exists!} C$ making the diagram

$$\begin{tikzcd}
(X,x_0) \arrow[r, shift left, "f"] \arrow[r, shift right, "g"'] & (Y,y_0) \arrow[r, hook, "\mathrm{coeq}(f,g)"] \arrow[drr, "c"'] & (\mathrm{CoEq}(f,g),[y_0]) \arrow[d, dashed, "\phi"', "\exists!"] \\
& & (C,*)
\end{tikzcd}$$

commute, where we note that ϕ is indeed a morphism of pointed sets since

$$\begin{aligned}
\phi([y_0]) &= [\phi \circ \mathrm{coeq}(f,g)]([y_0]) \\
&= c([y_0]) \\
&= *,
\end{aligned}$$

where we have used that c is a morphism of pointed sets. □

Proposition 6.3.5.1.3. Let (X,x_0) and (Y,y_0) be pointed sets and let $f,g,h \colon (X,x_0) \to (Y,y_0)$ be morphisms of pointed sets.

1. *Associativity.* We have isomorphisms of pointed sets

   $$\underbrace{\mathrm{CoEq}(\mathrm{coeq}(f,g)\circ f, \mathrm{coeq}(f,g)\circ h)}_{=\mathrm{CoEq}(\mathrm{coeq}(f,g)\circ g, \mathrm{coeq}(f,g)\circ h)} \cong \mathrm{CoEq}(f,g,h) \cong \underbrace{\mathrm{CoEq}(\mathrm{coeq}(g,h)\circ f, \mathrm{coeq}(g,h)\circ g)}_{=\mathrm{CoEq}(\mathrm{coeq}(g,h)\circ f, \mathrm{coeq}(g,h)\circ h)},$$

   where $\mathrm{CoEq}(f,g,h)$ is the colimit of the diagram

   $$\begin{tikzcd}
   (X,x_0) \arrow[r, shift left=2, "f"] \arrow[r, dashed, "g" description] \arrow[r, shift right=2, "h"'] & (Y,y_0)
   \end{tikzcd}$$

   in Sets_*.

2. *Unitality.* We have an isomorphism of pointed sets

   $$\mathrm{CoEq}(f,f) \cong B.$$

3. *Commutativity.* We have an isomorphism of pointed sets

   $$\mathrm{CoEq}(f,g) \cong \mathrm{CoEq}(g,f).$$

Proof. *Item 1*, *Associativity*: This follows from Constructions With Sets, Item 1 of Definition 4.2.5.1.5.
Item 2, *Unitality*: This follows from Constructions With Sets, Item 4 of Definition 4.2.5.1.5.
Item 3, *Commutativity*: This follows from Constructions With Sets, Item 5 of Definition 4.2.5.1.5. □

6.4 Constructions With Pointed Sets

6.4.1 Free Pointed Sets

Let X be a set.

Definition 6.4.1.1.1. The **free pointed set on** X is the pointed set X^+ consisting of:

- *The Underlying Set.* The set X^+ defined by[13]

$$\begin{aligned} X^+ &\overset{\text{def}}{=} X \coprod \mathrm{pt} \\ &\overset{\text{def}}{=} X \coprod \{\star\}. \end{aligned}$$

- *The Basepoint.* The element $\star$ of X^+.

Proposition 6.4.1.1.2. Let X be a set.

1. *Functoriality.* The assignment $X \mapsto X^+$ defines a functor

$$(-)^+ \colon \mathsf{Sets} \to \mathsf{Sets}_*,$$

   where:

   - *Action on Objects.* For each $X \in \mathrm{Obj}(\mathsf{Sets})$, we have

$$[(-)^+](X) \overset{\text{def}}{=} X^+,$$

     where X^+ is the pointed set of Definition 6.4.1.1.1.
   - *Action on Morphisms.* For each morphism $f \colon X \to Y$ of Sets, the image

$$f^+ \colon X^+ \to Y^+$$

     of f by $(-)^+$ is the map of pointed sets defined by

$$f^+(x) \overset{\text{def}}{=} \begin{cases} f(x) & \text{if } x \in X, \\ \star_Y & \text{if } x = \star_X. \end{cases}$$

2. *Adjointness.* We have an adjunction

$$\left((-)^+ \dashv 忘\right) \colon \quad \begin{tikzcd} \mathsf{Sets} \arrow[r, bend left, "(-)^+"] \arrow[r, phantom, "\perp"] & \mathsf{Sets}_*, \arrow[l, bend left, "忘"] \end{tikzcd}$$

[13] *Further Notation:* We sometimes write $\star_X$ for the basepoint of X^+ for clarity, specially when

witnessed by a bijection of sets

$$\mathsf{Sets}_*((X^+,\star_X),(Y,y_0)) \cong \mathsf{Sets}(X,Y),$$

natural in $X \in \mathrm{Obj}(\mathsf{Sets})$ and $(Y,y_0) \in \mathrm{Obj}(\mathsf{Sets}_*)$.

3. *Symmetric Strong Monoidality With Respect to Wedge Sums*. The free pointed set functor of Item 1 has a symmetric strong monoidal structure

$$\left((-)^+,(-)^{+,\coprod},(-)^{+,\coprod}_{\mathbb{1}}\right)\colon (\mathsf{Sets},\textstyle\coprod,\emptyset) \to (\mathsf{Sets}_*,\vee,\mathrm{pt}),$$

being equipped with isomorphisms of pointed sets

$$(-)^{+,\coprod}_{X,Y}\colon X^+ \vee Y^+ \overset{\sim}{\dashrightarrow} (X \textstyle\coprod Y)^+,$$
$$(-)^{+,\coprod}_{\mathbb{1}}\colon \mathrm{pt} \overset{\sim}{\dashrightarrow} \emptyset^+,$$

natural in $X,Y \in \mathrm{Obj}(\mathsf{Sets})$.

4. *Symmetric Strong Monoidality With Respect to Smash Products*. The free pointed set functor of Item 1 has a symmetric strong monoidal structure

$$\left((-)^+,(-)^+,(-)^+_{\mathbb{1}}\right)\colon (\mathsf{Sets},\times,\mathrm{pt}) \to (\mathsf{Sets}_*,\wedge,S^0),$$

being equipped with isomorphisms of pointed sets

$$(-)^+_{X,Y}\colon X^+ \wedge Y^+ \overset{\sim}{\dashrightarrow} (X \times Y)^+,$$
$$(-)^+_{\mathbb{1}}\colon S^0 \overset{\sim}{\dashrightarrow} \mathrm{pt}^+,$$

natural in $X,Y \in \mathrm{Obj}(\mathsf{Sets})$.

Proof. *Item 1, Functoriality*: We claim that $(-)^+$ is indeed a functor:

- *Preservation of Identities*. Let $X \in \mathrm{Obj}(\mathsf{Sets})$. We have

$$\mathrm{id}^+_X(x) \overset{\mathrm{def}}{=} \begin{cases} x & \text{if } x \in X, \\ \star_X & \text{if } x = \star_X, \end{cases}$$

for each $x \in X^+$, so $\mathrm{id}^+_X = \mathrm{id}_{X^+}$.

- *Preservation of Composition*. Given morphisms of sets

$$f\colon X \to Y,$$

$$g : Y \to Z,$$

we have

$$\begin{aligned}[g^+ \circ f^+](x) &\overset{\text{def}}{=} g^+(f^+(x)) \\ &\overset{\text{def}}{=} g^+(f(x)) \\ &\overset{\text{def}}{=} g(f(x)) \\ &\overset{\text{def}}{=} [g \circ f]^+(x)\end{aligned}$$

for each $x \in X$ and

$$\begin{aligned}[g^+ \circ f^+](\star_X) &\overset{\text{def}}{=} g^+(f^+(\star_X)) \\ &\overset{\text{def}}{=} g^+(\star_Y) \\ &\overset{\text{def}}{=} \star_Z \\ &\overset{\text{def}}{=} [g \circ f]^+(\star_X),\end{aligned}$$

so $(g \circ f)^+ = g^+ \circ f^+$.

This finishes the proof.
Item 2, *Adjointness*: We proceed in a few steps:

- *Map I*. We define a map

$$\Phi_{X,Y} : \mathsf{Sets}_*(X^+, Y) \to \mathsf{Sets}(X, Y)$$

by sending a morphism of pointed sets

$$\xi : (X^+, \star_X) \to (Y, y_0)$$

to the function

$$\xi^\dagger : X \to Y$$

given by

$$\xi^\dagger(x) \overset{\text{def}}{=} \xi(x)$$

for each $x \in X$.

- *Map II*. We define a map

$$\Psi_{X,Y} : \mathsf{Sets}(X, Y) \to \mathsf{Sets}_*(X^+, Y)$$

given by sending a function $\xi : X \to Y$ to the morphism of pointed sets

$$\xi^\dagger : (X^+, \star_X) \to (Y, y_0)$$

defined by

$$\xi^{\dagger}(x) \stackrel{\text{def}}{=} \begin{cases} \xi(x) & \text{if } x \in X, \\ y_0 & \text{if } x = \star_X \end{cases}$$

for each $x \in X^+$.

- *Invertibility I.* Given a morphism of pointed sets

$$\xi \colon (X^+, \star_X) \to (Y, y_0),$$

we have

$$\begin{aligned} [\Psi_{X,Y} \circ \Phi_{X,Y}](\xi) &\stackrel{\text{def}}{=} \Psi_{X,Y}(\Phi_{X,Y}(\xi)) \\ &= \Psi_{X,Y}(\xi^{\dagger}) \\ &\stackrel{\text{def}}{=} [\![x \mapsto \begin{cases} \xi^{\dagger}(x) & \text{if } x \in X \\ y_0 & \text{if } x = \star_X \end{cases}]\!] \\ &= [\![x \mapsto \begin{cases} \xi(x) & \text{if } x \in X \\ y_0 & \text{if } x = \star_X \end{cases}]\!] \\ &= \xi \\ &\stackrel{\text{def}}{=} [\mathrm{id}_{\mathsf{Sets}_*(X^+,Y)}](\xi). \end{aligned}$$

Therefore we have

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathrm{id}_{\mathsf{Sets}_*(X^+,Y)}.$$

- *Invertibility II.* Given a map of sets $\xi \colon X \to Y$, we have

$$\begin{aligned} [\Phi_{X,Y} \circ \Psi_{X,Y}](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y}(\Psi_{X,Y}(\xi)) \\ &= \Phi_{X,Y}(\xi^{\dagger}) \\ &= \Phi_{X,Y}([\![x \mapsto \begin{cases} \xi(x) & \text{if } x \in X \\ y_0 & \text{if } x = \star_X \end{cases}]\!]) \\ &= [\![x \mapsto \xi(x)]\!] \\ &= \xi \\ &\stackrel{\text{def}}{=} [\mathrm{id}_{\mathsf{Sets}(X,Y)}](\xi). \end{aligned}$$

Therefore we have

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathsf{Sets}(X,Y)}.$$

- *Naturality for* Φ*, Part I.* We need to show that, given a morphism of pointed sets
$$f\colon (X,x_0)\to(X',x'_0),$$
the diagram
$$\begin{tikzcd} \mathsf{Sets}_*(X'^{,+},Y) \arrow[r, "\Phi_{X',Y}"] \arrow[d, "f^*"'] & \mathsf{Sets}(X',Y) \arrow[d, "f^*"] \\ \mathsf{Sets}_*(X^+,Y) \arrow[r, "\Phi_{X,Y}"'] & \mathsf{Sets}(X,Y) \end{tikzcd}$$
commutes. Indeed, given a morphism of pointed sets $\xi\colon X'^{,+}\to Y$, we have
$$\begin{aligned} [\Phi_{X,Y}\circ f^*](\xi) &= \Phi_{X,Y}(f^*(\xi)) \\ &= \Phi_{X,Y}(\xi\circ f) \\ &= \xi\circ f \\ &= \Phi_{X',Y}(\xi)\circ f \\ &= f^*(\Phi_{X',Y}(\xi)) \\ &= f^*(\Phi_{X',Y}(\xi)) \\ &= [f^*\circ\Phi_{X',Y}](\xi). \end{aligned}$$
Therefore we have
$$\Phi_{X,Y}\circ f^* = f^*\circ\Phi_{X',Y}$$
and the naturality diagram for Φ above indeed commutes.

- *Naturality for* Φ*, Part II.* We need to show that, given a morphism of pointed sets
$$g\colon (Y,y_0)\to(Y',y'_0),$$
the diagram
$$\begin{tikzcd} \mathsf{Sets}_*(X^+,Y) \arrow[r, "\Phi_{X,Y}"] \arrow[d, "g_*"'] & \mathsf{Sets}(X,Y) \arrow[d, "g_*"] \\ \mathsf{Sets}_*(X^+,Y'), \arrow[r, "\Phi_{X,Y'}"'] & \mathsf{Sets}(X,Y') \end{tikzcd}$$
commutes. Indeed, given a morphism of pointed sets
$$\xi^\dagger\colon X^+\to Y,$$

we have

$$\begin{aligned}[\Phi_{X,Y'} \circ g_*](\xi) &= \Phi_{X,Y'}(g_*(\xi)) \\ &= \Phi_{X,Y'}(g \circ \xi) \\ &= g \circ \xi \\ &= g \circ \Phi_{X,Y'}(\xi) \\ &= g_*(\Phi_{X,Y'}(\xi)) \\ &= [g_* \circ \Phi_{X,Y'}](\xi).\end{aligned}$$

Therefore we have

$$\Phi_{X,Y'} \circ g_* = g_* \circ \Phi_{X,Y'}$$

and the naturality diagram for Φ above indeed commutes.

- *Naturality for Ψ.* Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Ψ is also natural in each argument.

This finishes the proof.

Item 3, Symmetric Strong Monoidality With Respect to Wedge Sums: We construct the strong monoidal structure on $(-)^+$ with respect to $\coprod$ and $\vee$ as follows:

- *The Strong Monoidality Constraints.* The isomorphism

$$(-)^{+,\coprod}_{X,Y} \colon X^+ \vee Y^+ \xrightarrow{\sim} (X \coprod Y)^+$$

is given by

$$(-)^{+,\coprod}_{X,Y}(z) = \begin{cases} x & \text{if } z = [(0,x)] \text{ with } x \in X, \\ y & \text{if } z = [(1,y)] \text{ with } y \in Y, \\ \star_{X \coprod Y} & \text{if } z = [(0,\star_X)], \\ \star_{X \coprod Y} & \text{if } z = [(1,\star_Y)] \end{cases}$$

for each $z \in X^+ \vee Y^+$, with inverse

$$(-)^{+,\coprod,-1}_{X,Y} \colon (X \coprod Y)^+ \xrightarrow{\sim} X^+ \vee Y^+$$

given by

$$(-)^{+,\coprod,-1}_{X,Y}(z) \overset{\text{def}}{=} \begin{cases} [(0,x)] & \text{if } z = [(0,x)], \\ [(1,y)] & \text{if } z = [(1,y)], \\ p_0 & \text{if } z = \star_{X \coprod Y} \end{cases}$$

for each $z \in (X \coprod Y)^+$.

- *The Strong Monoidal Unity Constraint.* The isomorphism

$$(-)^{+,\coprod,\mathbb{1}}_{X,Y} \colon \mathrm{pt} \overset{\sim}{\dashrightarrow} \emptyset^{+}$$

  is given by sending $\star_X$ to $\star_{\emptyset}$.

The verification that these isomorphisms satisfy the coherence conditions making the functor $(-)^{+}$ into a symmetric strong monoidal functor is omitted.
Item 4, Symmetric Strong Monoidality With Respect to Smash Products: We construct the strong monoidal structure on $(-)^{+}$ with respect to $\times$ and $\wedge$ as follows:

- *The Strong Monoidality Constraints.* The isomorphism

$$(-)^{+}_{X,Y} \colon X^{+} \wedge Y^{+} \overset{\sim}{\dashrightarrow} (X \times Y)^{+}$$

  is given by

$$(-)^{+}_{X,Y}(x \wedge y) = \begin{cases} (x, y) & \text{if } x \neq \star_X \text{ and } y \neq \star_Y \\ \star_{X \times Y} & \text{otherwise} \end{cases}$$

  for each $x \wedge y \in X^{+} \wedge Y^{+}$, with inverse

$$(-)^{+,-1}_{X,Y} \colon (X \times Y)^{+} \overset{\sim}{\dashrightarrow} X^{+} \wedge Y^{+}$$

  given by

$$(-)^{+,-1}_{X,Y}(z) \overset{\text{def}}{=} \begin{cases} x \wedge y & \text{if } z = (x, y) \text{ with } (x, y) \in X \times Y, \\ \star_X \wedge \star_Y & \text{if } z = \star_{X \times Y}, \end{cases}$$

  for each $z \in (X \times Y)^{+}$.

- *The Strong Monoidal Unity Constraint.* The isomorphism

$$(-)^{+,\mathbb{1}}_{X,Y} \colon S^0 \overset{\sim}{\dashrightarrow} \mathrm{pt}^{+}$$

  is given by sending 0 to $\star_{\mathrm{pt}}$ and 1 to $\star$, where $\mathrm{pt}^{+} = \{\star, \star_{\mathrm{pt}}\}$.

The verification that these isomorphisms satisfy the coherence conditions making the functor $(-)^{+}$ into a symmetric strong monoidal functor is omitted. □

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there are multiple free pointed sets involved in the current discussion.

6.4.2 Deleting Basepoints

Let (X, x_0) be a pointed set.

Definition 6.4.2.1.1. The **set with deleted basepoint associated to** X is the set X^- defined by

$$X^- \overset{\text{def}}{=} X \setminus \{x_0\}.$$

Proposition 6.4.2.1.2. Let (X, x_0) be a pointed set.

1. *Functoriality.* The assignment $(X, x_0) \mapsto X^-$ defines a functor

$$X^- \colon \mathsf{Sets}_*^{\mathrm{actv}} \to \mathsf{Sets},$$

where:

- *Action on Objects.* For each $X \in \mathrm{Obj}(\mathsf{Sets}_*^{\mathrm{actv}})$, we have

$$[(-)^-](X) \overset{\text{def}}{=} X^-,$$

where X^- is the set of Definition 6.4.2.1.1.

- *Action on Morphisms.* For each morphism $f \colon X \to Y$ of $\mathsf{Sets}_*^{\mathrm{actv}}$, the image

$$f^- \colon X^- \to Y^-$$

of f by $(-)^-$ is the map defined by

$$f^-(x) \overset{\text{def}}{=} f(x)$$

for each $x \in X^-$.

2. *Adjoint Equivalence.* We have an adjoint equivalence of categories

$$\big((-)^- \dashv (-)^+\big)\colon \quad \begin{tikzcd} \mathsf{Sets}_*^{\mathrm{actv}} \arrow[r, bend left, "(-)^-"] \arrow[r, phantom, "\perp_{\mathrm{eq}}"] & \mathsf{Sets}, \arrow[l, bend left, "(-)^+"] \end{tikzcd}$$

witnessed by a bijection of sets

$$\mathsf{Sets}(X^-, Y) \cong \mathsf{Sets}_*(X, Y^+),$$

natural in $X \in \mathrm{Obj}(\mathsf{Sets}_*)$ and $Y \in \mathrm{Obj}(\mathsf{Sets})$, and by isomorphisms

$$(X^-)^+ \cong X,$$
$$(Y^+)^- \cong Y,$$

once again natural in $X \in \mathrm{Obj}(\mathsf{Sets}_*)$ and $Y \in \mathrm{Obj}(\mathsf{Sets})$.

3. *Symmetric Strong Monoidality With Respect to Wedge Sums.* The functor of Item 1 has a symmetric strong monoidal structure

$$\left((-)^{-},(-)^{-,\vee},(-)^{-,\vee}_{\mathbb{1}}\right)\colon(\mathsf{Sets}^{\mathrm{actv}}_{*},\vee,\mathrm{pt}),\longrightarrow(\mathsf{Sets},\textstyle\coprod,\emptyset),$$

being equipped with isomorphisms of pointed sets

$$\begin{gathered}(-)^{-,\vee}_{X,Y}\colon X^{-}\textstyle\coprod Y^{-}\overset{\sim}{\dashrightarrow}(X\vee Y)^{-},\\(-)^{-,\vee}_{\mathbb{1}}\colon\emptyset\overset{\sim}{\dashrightarrow}\mathrm{pt}^{-},\end{gathered}$$

natural in $X,Y\in\mathrm{Obj}(\mathsf{Sets})$.

4. *Symmetric Strong Monoidality With Respect to Smash Products.* The free pointed set functor of Item 1 has a symmetric strong monoidal structure

$$\left((-)^{-},(-)^{-,\times},(-)^{-,\times}_{\mathbb{1}}\right)\colon(\mathsf{Sets}^{\mathrm{actv}}_{*},\wedge,S^{0}),\longrightarrow(\mathsf{Sets},\times,\mathrm{pt})$$

being equipped with isomorphisms of pointed sets

$$\begin{gathered}(-)^{-}_{X,Y}\colon X^{-}\times Y^{-}\overset{\sim}{\dashrightarrow}(X\wedge Y)^{-},\\(-)^{-}_{\mathbb{1}}\colon\mathrm{pt}\overset{\sim}{\dashrightarrow}(S^{0})^{-},\end{gathered}$$

natural in $X,Y\in\mathrm{Obj}(\mathsf{Sets})$.

Proof. *Item 1, Functoriality*: We claim that $(-)^{-}$ is indeed a functor:

- *Preservation of Identities.* Let $X\in\mathrm{Obj}(\mathsf{Sets})$. We have

$$\mathrm{id}^{-}_{X}(x)\overset{\mathrm{def}}{=}x$$

for each $x\in X^{-}$, so $\mathrm{id}^{-}_{X}=\mathrm{id}_{X^{-}}$.

- *Preservation of Composition.* Given morphisms of pointed sets

$$\begin{aligned}f\colon(X,x_0)&\longrightarrow(Y,y_0),\\g\colon(Y,y_0)&\longrightarrow(Z,z_0),\end{aligned}$$

we have

$$\begin{aligned}[g^{-}\circ f^{-}](x)&\overset{\mathrm{def}}{=}g^{-}(f^{-}(x))\\&\overset{\mathrm{def}}{=}g^{-}(f(x))\\&\overset{\mathrm{def}}{=}g(f(x))\\&\overset{\mathrm{def}}{=}[g\circ f]^{-}(x)\end{aligned}$$

for each $x\in X$, so $(g\circ f)^{-}=g^{-}\circ f^{-}$.

This finishes the proof.

Item 2, *Adjoint Equivalence*: We proceed in a few steps:

1. *Map I.* We define a map
$$\Phi_{X,Y} \colon \mathsf{Sets}(X^-, Y) \to \mathsf{Sets}_*^{\mathrm{actv}}(X, Y^+)$$
by sending a map $\xi \colon X^- \to Y$ to the active morphism of pointed sets
$$\xi^\dagger \colon X \to Y^+$$
given by
$$\xi^\dagger(x) \stackrel{\text{def}}{=} \begin{cases} \xi(x) & \text{if } x \in X^-, \\ \star_Y & \text{if } x = x_0, \end{cases}$$
for each $x \in X$, where this morphism is indeed active since $\xi(x) \in Y = Y^+ \setminus \{\star_Y\}$ for all $x \in X^-$.

2. *Map II.* We define a map
$$\Psi_{X,Y} \colon \mathsf{Sets}_*^{\mathrm{actv}}(X, Y^+) \to \mathsf{Sets}(X^-, Y)$$
given by sending an active morphism of pointed sets $\xi \colon X \to Y^+$ to the map
$$\xi^\dagger \colon X^- \to Y$$
defined by
$$\xi^\dagger(x) \stackrel{\text{def}}{=} \xi(x)$$
for each $x \in X^-$, which is indeed well-defined (in that $\xi(x) \in Y$ for all $x \in X^-$) since ξ is active.

3. *Invertibility I.* Given a map of sets $\xi \colon X^- \to Y$, we have
$$\begin{aligned} [\Psi_{X,Y} \circ \Phi_{X,Y}](\xi) &\stackrel{\text{def}}{=} \Psi_{X,Y}(\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} \Psi_{X,Y}\left(\left[\!\left[x \mapsto \begin{cases} \xi(x) & \text{if } x \in X^- \\ \star_Y & \text{if } x = x_0 \end{cases}\right]\!\right]\right) \\ &= [\![x \mapsto \xi(x)]\!] \\ &= \xi \\ &= [\mathrm{id}_{\mathsf{Sets}(X^-,Y)}](\xi). \end{aligned}$$
Therefore we have
$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathrm{id}_{\mathsf{Sets}(X^-,Y)}.$$

4. *Invertibility II.* Given a morphism of pointed sets

$$\xi\colon (X,x_0)\to(Y^+,\star_Y),$$

we have

$$\begin{aligned}
[\Phi_{X,Y}\circ\Psi_{X,Y}](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y}(\Psi_{X,Y}(\xi))\\
&= \Phi_{X,Y}([\![x\mapsto\xi(x)]\!])\\
&= [\![x\mapsto\begin{cases}\xi(x) & \text{if } x\in X^-\\ \star_Y & \text{if } x=x_0\end{cases}]\!]\\
&= \xi\\
&= [\mathrm{id}_{\mathsf{Sets}_*^{\mathrm{actv}}(X,Y^+)}](\xi).
\end{aligned}$$

Therefore we have

$$\Phi_{X,Y}\circ\Psi_{X,Y} = \mathrm{id}_{\mathsf{Sets}_*^{\mathrm{actv}}(X,Y^+)}.$$

5. *Naturality for Φ, Part I.* We need to show that, given a morphism of pointed sets

$$f\colon (X,x_0)\to(X',x_0'),$$

the diagram

$$\begin{tikzcd}
\mathsf{Sets}(X'^{,-},Y) \arrow[r, "\Phi_{X',Y}"] \arrow[d, "f^*"'] & \mathsf{Sets}_*^{\mathrm{actv}}(X',Y^+) \arrow[d, "f^*"]\\
\mathsf{Sets}_*(X^-,Y) \arrow[r, "\Phi_{X,Y}"'] & \mathsf{Sets}_*^{\mathrm{actv}}(X,Y^+)
\end{tikzcd}$$

commutes. Indeed, given a map of sets $\xi\colon X'\to Y$, we have

$$\begin{aligned}
[\Phi_{X,Y}\circ f^*](\xi) &= \Phi_{X,Y}(f^*(\xi))\\
&= \Phi_{X,Y}(\xi\circ f)\\
&= [\![x\mapsto\begin{cases}\xi(f(x)) & \text{if } f(x)\in X'^{,-}\\ \star_Y & \text{if } f(x)=x_0'\end{cases}]\!]\\
&= f^*([\![x'\mapsto\begin{cases}\xi(x') & \text{if } x'\in X'^{,-}\\ \star_Y & \text{if } x'=x_0'\end{cases}]\!])\\
&= f^*(\Phi_{X',Y}(\xi))\\
&= [f^*\circ\Phi_{X',Y}](\xi).
\end{aligned}$$

Therefore we have

$$\Phi_{X,Y} \circ f^* = f^* \circ \Phi_{X',Y},$$

and the naturality diagram for Φ above indeed commutes.

6. *Naturality for Φ, Part II.* We need to show that, given a morphism of pointed sets

$$g\colon (Y, y_0) \to (Y', y'_0),$$

the diagram

$$\begin{tikzcd}
\mathsf{Sets}(X^-, Y) \arrow[r, "\Phi_{X,Y}"] \arrow[d, "g_*"'] & \mathsf{Sets}_*^{\mathrm{actv}}(X, Y^+) \arrow[d, "g_*"] \\
\mathsf{Sets}(X^-, Y') \arrow[r, "\Phi_{X,Y'}"'] & \mathsf{Sets}_*^{\mathrm{actv}}(X, Y'^{,+})
\end{tikzcd}$$

commutes. Indeed, given a map of sets $\xi\colon X^- \to Y$, we have

$$\begin{aligned}
[\Phi_{X,Y'} \circ g_*](\xi) &= \Phi_{X,Y'}(g_*(\xi)) \\
&= \Phi_{X,Y'}(g \circ \xi) \\
&= [\![x \mapsto \begin{cases} g(\xi(x)) & \text{if } x \in X^- \\ \star_{Y'} & \text{if } x = x_0 \end{cases}]\!] \\
&= g_*([\![x \mapsto \begin{cases} \xi(x) & \text{if } x \in X^- \\ \star_{Y} & \text{if } x = x_0 \end{cases}]\!]) \\
&= g_*(\Phi_{X,Y'}(\xi)) \\
&= [g_* \circ \Phi_{X,Y'}](\xi).
\end{aligned}$$

Therefore we have

$$\Phi_{X,Y'} \circ g_* = g_* \circ \Phi_{X,Y'},$$

and the naturality diagram for Φ above indeed commutes.

7. *Naturality for Ψ.* Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Ψ is also natural in each argument.

8. *Fully Faithfulness of $(-)^-$.* We aim to show that the assignment $f \mapsto f^-$ sets up a bijection

$$(-)^-_{X,Y}\colon \mathsf{Sets}_*^{\mathrm{actv}}(X, Y) \overset{\sim}{\dashrightarrow} \mathsf{Sets}(X^-, Y^-).$$

Indeed, the inverse map

$$(-)^{-,-1}_{X,Y} \colon \mathsf{Sets}(X^-, Y^-) \overset{\sim}{\dashrightarrow} \mathsf{Sets}^{\mathrm{actv}}_*(X, Y)$$

is given by sending a map of sets $f \colon X^- \to Y^-$ to the active morphism of pointed sets $f^\dagger \colon X \to Y$ defined by

$$f^\dagger(x) \overset{\text{def}}{=} \begin{cases} f(x) & \text{if } x \in X^-, \\ y_0 & \text{if } x = x_0 \end{cases}$$

for each $x \in X$.

9. *Essential Surjectivity of* $(-)^-$. We need to show that, given an object $X \in \mathrm{Obj}(\mathsf{Sets})$, there exists some $X' \in \mathrm{Obj}(\mathsf{Sets}^{\mathrm{actv}}_*)$ such that $(X')^- \cong X$. Indeed, taking $X' = X^+$, we have

$$\begin{aligned} (X^+)^- &\overset{\text{def}}{=} (X \cup \{\star_X\})^- \\ &\overset{\text{def}}{=} (X \cup \{\star_X\}) \setminus \{\star_X\} \\ &= X, \end{aligned}$$

and thus we have in fact an *equality* $(X^+)^- = X$, showing $(-)^-$ to be essentially surjective.

10. *The Functor* $(-)^-$ *Is an Equivalence.* Since $(-)^-$ is fully faithful and essentially surjective, it is an equivalence by Categories, Item 1 of Definition 11.6.7.1.2.

This finishes the proof.

Item 3, Symmetric Strong Monoidality With Respect to Wedge Sums: We construct the strong monoidal structure on $(-)^-$ with respect to $\vee$ and $\coprod$ as follows:

- *The Strong Monoidality Constraints.* The isomorphism

$$(-)^{-,\vee}_{X,Y} \colon X^- \coprod Y^- \overset{\sim}{\dashrightarrow} (X \vee Y)^-$$

is given by

$$(-)^{-,\vee}_{X,Y}(z) = \begin{cases} [(0,x)] & \text{if } z = (0,x) \text{ with } x \in X, \\ [(1,y)] & \text{if } z = (1,y) \text{ with } y \in Y \end{cases}$$

for each $z \in X^- \coprod Y^-$, with inverse

$$(-)^{-,\vee,-1}_{X,Y} \colon (X \vee Y)^- \overset{\sim}{\dashrightarrow} X^- \coprod Y^-$$

given by

$$(-)^{-,\vee,-1}_{X,Y}(z) \overset{\text{def}}{=} \begin{cases} (0,x) & \text{if } z = [(0,x)], \\ (1,y) & \text{if } z = [(1,y)], \end{cases}$$

for each $z \in (X \vee Y)^{-}$.

- *The Strong Monoidal Unity Constraint.* The isomorphism

  $$(-)^{+,\vee,\mathbb{1}}_{X,Y} \colon \emptyset \overset{\sim}{\dashrightarrow} \mathrm{pt}^{-}$$

  is an equality.

The verification that these isomorphisms satisfy the coherence conditions making the functor $(-)^{-}$ into a symmetric strong monoidal functor is omitted.
Item 4, Symmetric Strong Monoidality With Respect to Smash Products: We construct the strong monoidal structure on $(-)^{+}$ with respect to $\wedge$ and $\times$ as follows:

- *The Strong Monoidality Constraints.* The isomorphism

  $$(-)^{-}_{X,Y} \colon X^{-} \times Y^{-} \overset{\sim}{\dashrightarrow} (X \wedge Y)^{-}$$

  is given by

  $$(-)^{-}_{X,Y}(x,y) = x \wedge y$$

  for each $(x,y) \in X^{-} \times Y^{-}$, with inverse

  $$(-)^{-,-1}_{X,Y} \colon (X \wedge Y)^{-} \overset{\sim}{\dashrightarrow} X^{-} \times Y^{-}$$

  given by

  $$(-)^{-,-1}_{X,Y}(x \wedge y) \overset{\text{def}}{=} (x,y)$$

  for each $x \wedge y \in (X \wedge Y)^{-}$.

- *The Strong Monoidal Unity Constraint.* The isomorphism

  $$(-)^{-,\mathbb{1}}_{X,Y} \colon \mathrm{pt} \overset{\sim}{\dashrightarrow} (S^{0})^{-}$$

  is given by sending $\star$ to 1.

The verification that these isomorphisms satisfy the coherence conditions making the functor $(-)^{+}$ into a symmetric strong monoidal functor is omitted. □

Appendices

A Other Chapters

References

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