

# Constructions With Monoidal Categories

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This chapter contains some material on constructions with monoidal categories.

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## 13.1 Moduli Categories of Monoidal Structures

### 13.1.1 The Moduli Category of Monoidal Structures on a Category

Let  $\mathcal{C}$  be a category.

**Definition 13.1.1.1.** The **moduli category of monoidal structures on  $\mathcal{C}$**  is the category  $\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})$  defined by

$$\mathcal{M}_{\mathbb{E}_1}(\mathcal{C}) \stackrel{\text{def}}{=} \text{pt} \times_{\text{Cats}} \text{MonCats},$$

$$\begin{array}{ccc} \mathcal{M}_{\mathbb{E}_1}(\mathcal{C}) & \longrightarrow & \text{MonCats} \\ \downarrow & \lrcorner & \downarrow \text{忘} \\ \text{pt} & \xrightarrow{[\mathcal{C}]} & \text{Cats}. \end{array}$$

**Remark 13.1.1.2.** In detail, the **moduli category of monoidal structures on  $\mathcal{C}$**  is the category  $\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})$  where:

- *Objects.* The objects of  $\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})$  are monoidal categories  $(\mathcal{C}, \otimes_{\mathcal{C}}, \mathbb{1}_{\mathcal{C}}, \alpha^{\mathcal{C}}, \lambda^{\mathcal{C}}, \rho^{\mathcal{C}})$  whose underlying category is  $\mathcal{C}$ .
- *Morphisms.* A morphism from  $(\mathcal{C}, \otimes_{\mathcal{C}}, \mathbb{1}_{\mathcal{C}}, \alpha^{\mathcal{C}}, \lambda^{\mathcal{C}}, \rho^{\mathcal{C}})$  to  $(\mathcal{C}, \boxtimes_{\mathcal{C}}, \mathbb{1}'_{\mathcal{C}}, \alpha^{C'}, \lambda^{C'}, \rho^{C'})$  is a strong monoidal functor structure

$$\begin{aligned} \text{id}_{\mathcal{C}}^{\otimes} : A \otimes_{\mathcal{C}} B &\xrightarrow{\sim} A \boxtimes_{\mathcal{C}} B, \\ \text{id}_{\mathbb{1}_{\mathcal{C}}}^{\otimes} : \mathbb{1}'_{\mathcal{C}} &\xrightarrow{\sim} \mathbb{1}_{\mathcal{C}} \end{aligned}$$

on the identity functor  $\text{id}_{\mathcal{C}} : \mathcal{C} \rightarrow \mathcal{C}$  of  $\mathcal{C}$ .

- *Identities.* For each  $M \stackrel{\text{def}}{=} (\mathcal{C}, \otimes_{\mathcal{C}}, \mathbb{1}_{\mathcal{C}}, \alpha^{\mathcal{C}}, \lambda^{\mathcal{C}}, \rho^{\mathcal{C}}) \in \text{Obj}(\mathcal{M}_{\mathbb{E}_1}(\mathcal{C}))$ , the unit map

$$\mathbb{1}_{M,M}^{\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})} : \text{pt} \rightarrow \text{Hom}_{\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})}(M, M)$$

of  $\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})$  at  $M$  is defined by

$$\text{id}_M^{\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})} \stackrel{\text{def}}{=} (\text{id}_{\mathcal{C}}^{\otimes}, \text{id}_{\mathbb{1}_{\mathcal{C}}}^{\otimes}),$$

where  $(\text{id}_{\mathcal{C}}^{\otimes}, \text{id}_{\mathbb{1}_{\mathcal{C}}}^{\otimes})$  is the identity monoidal functor of  $\mathcal{C}$  of ??.

- *Composition.* For each  $M, N, P \in \text{Obj}(\mathcal{M}_{\mathbb{E}_1}(\mathcal{C}))$ , the composition map

$$\circ_{M,N,P}^{\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})} : \text{Hom}_{\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})}(N, P) \times \text{Hom}_{\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})}(M, N) \rightarrow \text{Hom}_{\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})}(M, P)$$

of  $\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})$  at  $(M, N, P)$  is defined by

$$\left( \text{id}_{\mathcal{C}}^{\otimes, \prime}, \text{id}_{\mathbb{1}|\mathcal{C}}^{\otimes, \prime} \right) \circ_{M,N,P}^{\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})} \left( \text{id}_{\mathcal{C}}^{\otimes}, \text{id}_{\mathbb{1}|\mathcal{C}}^{\otimes} \right) \stackrel{\text{def}}{=} \left( \text{id}_{\mathcal{C}}^{\otimes, \prime} \circ \text{id}_{\mathcal{C}}^{\otimes}, \text{id}_{\mathbb{1}|\mathcal{C}}^{\otimes, \prime} \circ \text{id}_{\mathbb{1}|\mathcal{C}}^{\otimes} \right).$$

**Remark 13.1.1.1.3.** In particular, a morphism in  $\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})$  from  $(\mathcal{C}, \otimes_{\mathcal{C}}, \mathbb{1}_{\mathcal{C}}, \alpha^{\mathcal{C}}, \lambda^{\mathcal{C}}, \rho^{\mathcal{C}})$  to  $(\mathcal{C}, \boxtimes_{\mathcal{C}}, \mathbb{1}'_{\mathcal{C}}, \alpha^{C, \prime}, \lambda^{C, \prime}, \rho^{C, \prime})$  satisfies the following conditions:

1. *Naturality.* For each pair  $f: A \rightarrow X$  and  $g: B \rightarrow Y$  of morphisms of  $\mathcal{C}$ , the diagram

$$\begin{array}{ccc} A \boxtimes_{\mathcal{C}} B & \xrightarrow{f \boxtimes_{\mathcal{C}} g} & X \boxtimes_{\mathcal{C}} Y \\ \text{id}_{A,B}^{\otimes} \downarrow & & \downarrow \text{id}_{X,Y}^{\otimes} \\ A \otimes_{\mathcal{C}} B & \xrightarrow{f \otimes_{\mathcal{C}} g} & X \otimes_{\mathcal{C}} Y \end{array}$$

commutes.

2. *Monoidality.* For each  $A, B, C \in \text{Obj}(\mathcal{C})$ , the diagram

$$\begin{array}{ccc} & (A \boxtimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C & \\ \text{id}_{A,B}^{\otimes} \boxtimes_{\mathcal{C}} \text{id}_C \swarrow & & \searrow \alpha_{A,B,C}^{C, \prime} \\ (A \otimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C & & A \boxtimes_{\mathcal{C}} (B \boxtimes_{\mathcal{C}} C) \\ \text{id}_{A \otimes_{\mathcal{C}} B, C}^{\otimes} \downarrow & & \downarrow \text{id}_A \boxtimes_{\mathcal{C}} \text{id}_{B,C}^{\otimes} \\ (A \otimes_{\mathcal{C}} B) \otimes_{\mathcal{C}} C & & A \boxtimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C) \\ \alpha_{A,B,C}^{\mathcal{C}} \searrow & & \swarrow \text{id}_{A, B \otimes_{\mathcal{C}} C}^{\otimes} \\ & A \otimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C) & \end{array}$$

commutes.

3. *Left Monoidal Unity.* For each  $A \in \text{Obj}(\mathcal{C})$ , the diagram

$$\begin{array}{ccc}
 & \mathbb{1}_{\mathcal{C}} \boxtimes_{\mathcal{C}} A & \xrightarrow{\text{id}_{\mathbb{1}'_{\mathcal{C}}}^{\otimes} \text{id}_A} \mathbb{1}_{\mathcal{C}} \otimes_{\mathcal{C}} A \\
 \text{id}_{\mathbb{1}}^{\otimes} \boxtimes \text{id}_A \nearrow & & \searrow \lambda_A^{\mathcal{C}} \\
 \mathbb{1}'_{\mathcal{C}} \boxtimes_{\mathcal{C}} A & \xrightarrow{\lambda_A^{\mathcal{C},'}} & A
 \end{array}$$

commutes.

4. *Right Monoidal Unity.* For each  $A \in \text{Obj}(\mathcal{C})$ , the diagram

$$\begin{array}{ccc}
 & A \boxtimes_{\mathcal{C}} \mathbb{1}_{\mathcal{C}} & \xrightarrow{\text{id}_A^{\otimes} \text{id}_{\mathbb{1}'_{\mathcal{C}}}} A \otimes_{\mathcal{C}} \mathbb{1}_{\mathcal{C}} \\
 \text{id}_A \boxtimes \text{id}_{\mathbb{1}}^{\otimes} \nearrow & & \searrow \rho_A^{\mathcal{C}} \\
 A \boxtimes_{\mathcal{C}} \mathbb{1}'_{\mathcal{C}} & \xrightarrow{\rho_A^{\mathcal{C},'}} & A
 \end{array}$$

commutes.

**Proposition 13.1.1.1.4.** Let  $\mathcal{C}$  be a category.

1. *Extra Monoidality Conditions.* Let  $(\text{id}_{\mathcal{C}}^{\otimes}, \text{id}_{\mathbb{1}_{\mathcal{C}}}^{\otimes})$  be a morphism of  $\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})$  from  $(\mathcal{C}, \otimes_{\mathcal{C}}, \mathbb{1}_{\mathcal{C}}, \alpha^{\mathcal{C}}, \lambda^{\mathcal{C}}, \rho^{\mathcal{C}})$  to  $(\mathcal{C}, \boxtimes_{\mathcal{C}}, \mathbb{1}'_{\mathcal{C}}, \alpha^{\mathcal{C},'}, \lambda^{\mathcal{C},'}, \rho^{\mathcal{C},'})$ .

(a) The diagram

$$\begin{array}{ccc}
 (A \boxtimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C & \xrightarrow{\text{id}_{A,B}^{\otimes} \boxtimes \text{id}_C} & (A \otimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C \\
 \text{id}_{A \boxtimes_{\mathcal{C}} B, C}^{\otimes} \downarrow & & \downarrow \text{id}_{A \otimes_{\mathcal{C}} B, C}^{\otimes} \\
 (A \boxtimes_{\mathcal{C}} B) \otimes_{\mathcal{C}} C & \xrightarrow{\text{id}_{A,B}^{\otimes} \otimes \text{id}_C} & (A \otimes_{\mathcal{C}} B) \otimes_{\mathcal{C}} C
 \end{array}$$

commutes.

(b) The diagram

$$\begin{array}{ccc}
 A \boxtimes_{\mathcal{C}} (B \boxtimes_{\mathcal{C}} C) & \xrightarrow{\text{id}_A \boxtimes \text{id}_{B,C}^{\otimes}} & A \boxtimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C) \\
 \text{id}_{A, B \boxtimes_{\mathcal{C}} C}^{\otimes} \downarrow & & \downarrow \text{id}_{A, B \otimes_{\mathcal{C}} C}^{\otimes} \\
 A \otimes_{\mathcal{C}} (B \boxtimes_{\mathcal{C}} C) & \xrightarrow{\text{id}_A \otimes \text{id}_{B,C}^{\otimes}} & A \otimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C)
 \end{array}$$

commutes.

2. *Extra Monoidal Unity Constraints.* Let  $(\text{id}_C^\otimes, \text{id}_{1|C}^\otimes)$  be a morphism of  $\mathcal{M}_{\mathbb{E}_1}(C)$  from  $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$  to  $(C, \boxtimes_C, 1'_C, \alpha^{C'}, \lambda^{C'}, \rho^{C'})$ .

(a) The diagram

$$\begin{array}{ccc} 1_C \otimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^\otimes, -1} & 1_C \boxtimes_C 1'_C \\ \lambda_{1'_C}^C \downarrow & & \downarrow \rho_{1_C}^{C, ' } \\ 1'_C & \xrightarrow{\text{id}_1^\otimes} & 1_C \end{array}$$

commutes.

(b) The diagram

$$\begin{array}{ccc} 1'_C \otimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^\otimes, -1} & 1'_C \boxtimes_C 1_C \\ \rho_{1'_C}^C \downarrow & & \downarrow \lambda_{1_C}^{C, ' } \\ 1'_C & \xrightarrow{\text{id}_1^\otimes} & 1_C \end{array}$$

commutes.

(c) The diagram

$$\begin{array}{ccc} 1'_C \boxtimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^\otimes} & 1'_C \otimes_C 1_C \\ \lambda_{1_C}^{C, ' } \downarrow & & \downarrow \rho_{1'_C}^C \\ 1_C & \xrightarrow{\text{id}_1^\otimes, -1} & 1'_C \end{array}$$

commutes.

(d) The diagram

$$\begin{array}{ccc} 1_C \boxtimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^\otimes} & 1_C \otimes_C 1'_C \\ \rho_{1_C}^{C, ' } \downarrow & & \downarrow \lambda_{1'_C}^C \\ 1_C & \xrightarrow{\text{id}_1^\otimes, -1} & 1'_C \end{array}$$

commutes.

3. *Mixed Associators.* Let  $(\mathcal{C}, \otimes_{\mathcal{C}}, 1_{\mathcal{C}}, \alpha^{\mathcal{C}}, \lambda^{\mathcal{C}}, \rho^{\mathcal{C}})$  and  $(\mathcal{C}, \boxtimes_{\mathcal{C}}, 1'_{\mathcal{C}}, \alpha^{\mathcal{C}'}, \lambda^{\mathcal{C}'}, \rho^{\mathcal{C}'})$  be monoidal structures on  $\mathcal{C}$  and let

$$\mathrm{id}_{-1,-2}^{\otimes}: -_1 \boxtimes_{\mathcal{C}} -_2 \rightarrow -_1 \otimes_{\mathcal{C}} -_2$$

be a natural transformation.

- (a) If there exists a natural transformation

$$\alpha_{A,B,C}^{\otimes}: (A \otimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C \rightarrow A \otimes_{\mathcal{C}} (B \boxtimes_{\mathcal{C}} C)$$

making the diagrams

$$\begin{array}{ccc} (A \otimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C & \xrightarrow{\alpha_{A,B,C}^{\otimes}} & A \otimes_{\mathcal{C}} (B \boxtimes_{\mathcal{C}} C) \\ \mathrm{id}_{A \otimes_{\mathcal{C}} B, C}^{\otimes} \downarrow & & \downarrow \mathrm{id}_A \otimes \mathrm{id}_{B,C}^{\otimes} \\ (A \otimes_{\mathcal{C}} B) \otimes_{\mathcal{C}} C & \xrightarrow{\alpha_{A,B,C}^{\mathcal{C}}} & A \otimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C) \end{array}$$

and

$$\begin{array}{ccc} (A \boxtimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C & \xrightarrow{\alpha_{A,B,C}^{\mathcal{C}'}} & A \boxtimes_{\mathcal{C}} (B \boxtimes_{\mathcal{C}} C) \\ \mathrm{id}_{A,B}^{\otimes} \boxtimes \mathrm{id}_C \downarrow & & \downarrow \mathrm{id}_{A,B \boxtimes_{\mathcal{C}} C} \\ (A \otimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C & \xrightarrow{\alpha_{A,B,C}^{\otimes}} & A \otimes_{\mathcal{C}} (B \boxtimes_{\mathcal{C}} C) \end{array}$$

commute, then the natural transformation  $\mathrm{id}^{\otimes}$  satisfies the monoidal-ity condition of [Item 2](#) of [Definition 13.1.1.1.3](#).

- (b) If there exists a natural transformation

$$\alpha_{A,B,C}^{\boxtimes}: (A \boxtimes_{\mathcal{C}} B) \otimes_{\mathcal{C}} C \rightarrow A \boxtimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C)$$

making the diagrams

$$\begin{array}{ccc} (A \boxtimes_{\mathcal{C}} B) \otimes_{\mathcal{C}} C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes}} & A \boxtimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C) \\ \mathrm{id}_{A,B}^{\otimes} \otimes \mathrm{id}_C \downarrow & & \downarrow \mathrm{id}_{A,B \otimes_{\mathcal{C}} C}^{\otimes} \\ (A \otimes_{\mathcal{C}} B) \otimes_{\mathcal{C}} C & \xrightarrow{\alpha_{A,B,C}^{\mathcal{C}}} & A \otimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C) \end{array}$$

and

$$\begin{array}{ccc}
 (A \boxtimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A,B,C}^{C'}} & A \boxtimes_C (B \boxtimes_C C) \\
 \text{id}_{A \boxtimes_C B, C}^\otimes \downarrow & & \downarrow \text{id}_A \boxtimes \text{id}_{B, C}^\otimes \\
 (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^\otimes} & A \boxtimes_C (B \otimes_C C)
 \end{array}$$

commute, then the natural transformation  $\text{id}^\otimes$  satisfies the monoidality condition of **Item 2** of **Definition 13.1.1.1.3**.

(c) If there exists a natural transformation

$$\alpha_{A,B,C}^{\boxtimes, \otimes}: (A \boxtimes_C B) \otimes_C C \rightarrow A \otimes_C (B \boxtimes_C C)$$

making the diagrams

$$\begin{array}{ccc}
 (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes, \otimes}} & A \otimes_C (B \boxtimes_C C) \\
 \text{id}_{A, B}^\otimes \otimes \text{id}_C \downarrow & & \downarrow \text{id}_A \otimes \text{id}_{B, C}^\otimes \\
 (A \otimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^C} & A \otimes_C (B \otimes_C C)
 \end{array}$$

and

$$\begin{array}{ccc}
 (A \boxtimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A,B,C}^{C'}} & A \boxtimes_C (B \boxtimes_C C) \\
 \text{id}_{A \boxtimes_C B, C}^\otimes \downarrow & & \downarrow \text{id}_{A, B \boxtimes_C C}^\otimes \\
 (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes, \otimes}} & A \otimes_C (B \boxtimes_C C)
 \end{array}$$

commute, then the natural transformation  $\text{id}^\otimes$  satisfies the monoidality condition of **Item 2** of **Definition 13.1.1.1.3**.

*Proof. **Item 1, Extra Monoidality Conditions:*** We claim that **Items 1a** and **1b** are indeed true:

1. *Proof of **Item 1a:*** This follows from the naturality of  $\text{id}^\otimes$  with respect to the morphisms  $\text{id}_{A, B}^\otimes$  and  $\text{id}_C$ .

2. *Proof of Item 1b:* This follows from the naturality of  $\text{id}^\otimes$  with respect to the morphisms  $\text{id}_A$  and  $\text{id}_{B,C}^\otimes$ .

This finishes the proof.

*Item 2, Extra Monoidal Unity Constraints:* We claim that **Items 2a** and **2b** are indeed true:

1. *Proof of Item 1a:* Indeed, consider the diagram

$$\begin{array}{ccccc}
 1_C \otimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \boxtimes_C 1'_C & & \\
 \downarrow \lambda_{1'_C}^C & \searrow \text{id}_{1_C} \otimes \text{id}_1^\otimes & \downarrow \text{id}_{1_C} \boxtimes \text{id}_1^\otimes & & \\
 & & 1_C \otimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1_C}^{\otimes, -1}} & 1_C \boxtimes_C 1_C \\
 & & \downarrow & & \downarrow \text{id}_{1_C, 1_C}^\otimes \\
 & & 1_C \otimes_C 1_C & \xrightarrow{\lambda_{1_C}^C = \rho_{1_C}^C} & 1_C \\
 & \searrow \text{id}_1^\otimes & & & \uparrow \rho_{1_C}^{C, '}} \\
 & & 1_C & & 
 \end{array}$$

(1) (2) (3) (4)

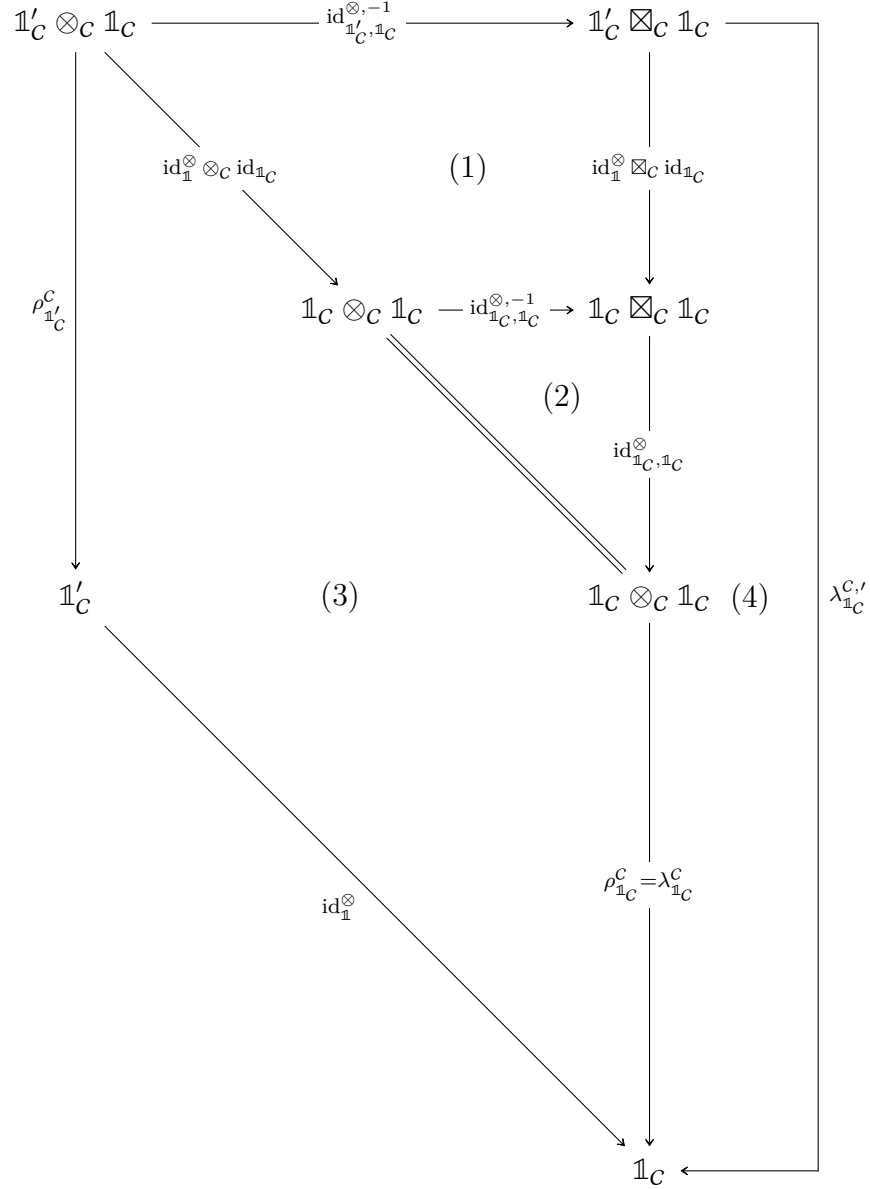


whose boundary diagram is the diagram whose commutativity we wish to prove. Since:

- Subdiagram (1) commutes by the naturality of  $\mathrm{id}_C^{\otimes, -1}$ ;
- Subdiagram (2) commutes trivially;
- Subdiagram (3) commutes by the naturality of  $\lambda^C$ , where the equality  $\rho_{1_C}^C = \lambda_{1_C}^C$  comes from ??;
- Subdiagram (4) commutes by the right monoidal unity of  $(\mathrm{id}_C, \mathrm{id}_C^{\otimes}, \mathrm{id}_{C|1}^{\otimes})$ ;

so does the boundary diagram, and we are done.

2. *Proof of Item 1b:* Indeed, consider the diagram



whose boundary diagram is the diagram whose commutativity we wish to prove. Since:

- Subdiagram (1) commutes by the naturality of  $\text{id}_C^{\otimes, -1}$ ;
- Subdiagram (2) commutes trivially;

- Subdiagram (3) commutes by the naturality of  $\rho^C$ , where the equality  $\rho_{\mathbb{1}_C}^C = \lambda_{\mathbb{1}_C}^C$  comes from ??;
- Subdiagram (4) commutes by the left monoidal unity of  $(\text{id}_C, \text{id}_C^\otimes, \text{id}_{C|\mathbb{1}}^\otimes)$ ;

so does the boundary diagram, and we are done.

3. *Proof of Item 2c:* Indeed, consider the diagram

$$\begin{array}{ccccc}
 \mathbb{1}'_C \otimes_C \mathbb{1}_C & \xrightarrow{\text{id}_{\mathbb{1}_C, \mathbb{1}'_C}^{\otimes, -1}} & \mathbb{1}'_C \boxtimes_C \mathbb{1}_C & \xrightarrow{\text{id}_{\mathbb{1}_C, \mathbb{1}'_C}^\otimes} & \mathbb{1}'_C \otimes_C \mathbb{1}_C \\
 \downarrow \rho_{\mathbb{1}'_C}^C & & \downarrow \lambda_{\mathbb{1}_C}^{C, '}& & \downarrow \rho_{\mathbb{1}'_C}^C \\
 \mathbb{1}'_C & \xrightarrow{\text{id}_{\mathbb{1}}^\otimes} & \mathbb{1}_C & \xrightarrow{\text{id}_{\mathbb{1}}^{\otimes, -1}} & \mathbb{1}'_C.
 \end{array}
 \quad \begin{array}{c} (1) \qquad \qquad (\dagger) \end{array}$$

Since:

- The boundary diagram commutes trivially;
- Subdiagram (1) commutes by Item 1b;

it follows that the diagram

$$\begin{array}{ccc}
 \mathbb{1}'_C \otimes_C \mathbb{1}_C & \xrightarrow{\text{id}_{\mathbb{1}_C, \mathbb{1}'_C}^{\otimes, -1}} & \mathbb{1}'_C \boxtimes_C \mathbb{1}_C \xrightarrow{\text{id}_{\mathbb{1}_C, \mathbb{1}'_C}^\otimes} \mathbb{1}'_C \otimes_C \mathbb{1}_C \\
 & & \downarrow \lambda_{\mathbb{1}_C}^{C, '} \qquad \qquad \downarrow \rho_{\mathbb{1}'_C}^C \\
 & & \mathbb{1}_C \xrightarrow{\text{id}_{\mathbb{1}}^{\otimes, -1}} \mathbb{1}'_C
 \end{array}
 \quad (\dagger)$$

commutes. But since  $\text{id}_{\mathbb{1}_C, \mathbb{1}'_C}^{\otimes, -1}$  is an isomorphism, it follows that the diagram  $(\dagger)$  also commutes, and we are done.

4. *Proof of Item 2d:* Indeed, consider the diagram

$$\begin{array}{ccccc}
 \mathbb{1}_C \otimes_C \mathbb{1}'_C & \xrightarrow{\text{id}_{\mathbb{1}_C, \mathbb{1}'_C}^{\otimes, -1}} & \mathbb{1}_C \boxtimes_C \mathbb{1}'_C & \xrightarrow{\text{id}_{\mathbb{1}_C, \mathbb{1}'_C}^{\otimes, -1}} & \mathbb{1}_C \otimes_C \mathbb{1}'_C \\
 \downarrow \lambda_{\mathbb{1}'_C}^C & & \downarrow \rho_{\mathbb{1}_C}^{C, '}& & \downarrow \lambda_{\mathbb{1}'_C}^C \\
 \mathbb{1}'_C & \xrightarrow{\text{id}_{\mathbb{1}}^\otimes} & \mathbb{1}_C & \xrightarrow{\text{id}_{\mathbb{1}}^{\otimes, -1}} & \mathbb{1}_C
 \end{array}
 \quad \begin{array}{c} (1) \qquad \qquad (\dagger) \end{array}$$

Since:

- The boundary diagram commutes trivially;
- Subdiagram (1) commutes by **Item 1a**;

it follows that the diagram

$$\begin{array}{ccccc}
 1_C \otimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \boxtimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \otimes_C 1'_C \\
 & & \downarrow \rho_{1_C}^{C, '}& (\dagger) & \downarrow \lambda_{1'_C}^C \\
 & & 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1_C
 \end{array}$$

commutes. But since  $\text{id}_1^{\otimes, -1}$  is an isomorphism, it follows that the diagram  $(\dagger)$  also commutes, and we are done.

This finishes the proof.

**Item 3, Mixed Associators:** We claim that **Items 3a** to **3c** are indeed true:

1. *Proof of Item 3a:* We may partition the monoidality diagram for  $\text{id}^{\otimes}$  of **Item 2** of **Definition 13.1.1.3** as follows:

$$\begin{array}{ccccc}
 & & (A \boxtimes_C B) \boxtimes_C C & & \\
 & \swarrow \text{id}_{A,B}^{\otimes} \boxtimes \text{id}_C & \downarrow & \searrow \alpha_{A,B,C}^C & \\
 (A \otimes_C B) \boxtimes_C C & & \text{id}_{A \boxtimes_C B, C}^{\otimes} & & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow \text{id}_{A \otimes_C B, C}^{\otimes} & (1) & \downarrow & (2) & \downarrow \text{id}_A \boxtimes \text{id}_{B,C}^{\otimes} \\
 & & (A \boxtimes_C B) \otimes_C C & & \\
 & \swarrow \text{id}_{A,B}^{\otimes} \otimes \text{id}_C & & \searrow \alpha_{A,B,C}^{\otimes} & \\
 (A \otimes_C B) \otimes_C C & & (3) & & A \boxtimes_C (B \otimes_C C) \\
 & \swarrow \alpha_{A,B,C}^{C, '}& & \swarrow \text{id}_{A, B \otimes_C C}^{\otimes} & \\
 & & A \otimes_C (B \otimes_C C) & & 
 \end{array}$$

Since:

- Subdiagram (1) commutes by **Item 1a** of **Item 1**.

- Subdiagram (2) commutes by assumption.
- Subdiagram (3) commutes by assumption.

it follows that the boundary diagram also commutes, i.e.  $\text{id}^\otimes$  satisfies the monoidality condition of [Item 2](#) of [Definition 13.1.1.1.3](#).

2. *Proof of [Item 3b](#):* We may partition the monoidality diagram for  $\text{id}^\otimes$  of [Item 2](#) of [Definition 13.1.1.1.3](#) as follows:

$$\begin{array}{ccccc}
 & & (A \boxtimes_C B) \boxtimes_C C & & \\
 & \swarrow \text{id}_{A,B}^\otimes \boxtimes \text{id}_C & & \searrow \alpha_{A,B,C}^C & \\
 (A \otimes_C B) \boxtimes_C C & & (1) & & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow \text{id}_{A \otimes_C B, C}^\otimes & \swarrow \alpha_{A,B,C}^\boxtimes & & \nwarrow \text{id}_{A,B \boxtimes_C C}^\otimes & \downarrow \text{id}_A \boxtimes \text{id}_{B,C}^\otimes \\
 & A \otimes_C (B \boxtimes_C C) & & & \\
 & (2) & & (3) & \\
 & \downarrow \text{id}_A \otimes \text{id}_{B,C}^\otimes & & & \downarrow \\
 (A \otimes_C B) \otimes_C C & & A \boxtimes_C (B \otimes_C C) & & \\
 \swarrow \alpha_{A,B,C}^{C,\prime} & \downarrow & \swarrow \text{id}_{A,B \otimes_C C}^\otimes & & \\
 & A \otimes_C (B \otimes_C C) & & & 
 \end{array}$$

Since:

- Subdiagram (1) commutes by assumption.
- Subdiagram (2) commutes by assumption.
- Subdiagram (3) commutes by [Item 1b](#) of [Item 1](#).

it follows that the boundary diagram also commutes, i.e.  $\text{id}^\otimes$  satisfies the monoidality condition of [Item 2](#) of [Definition 13.1.1.1.3](#).

3. *Proof of [Item 3c](#):* We may partition the monoidality diagram for  $\text{id}^\otimes$

of **Item 2** of **Definition 13.1.1.1.3** as follows:

$$\begin{array}{ccc}
 & (A \boxtimes_C B) \boxtimes_C C & \\
 \text{id}_{A,B}^\otimes \boxtimes_C \text{id}_C \swarrow & & \searrow \alpha_{A,B,C}^C \\
 (A \otimes_C B) \boxtimes_C C & \xrightarrow{(1)} & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow \text{id}_{A \otimes_C B, C}^\otimes & \searrow \alpha_{A,B,C}^{\boxtimes, \otimes} & \downarrow \text{id}_A \boxtimes_C \text{id}_{B,C}^\otimes \\
 (A \otimes_C B) \otimes_C C & \xrightarrow{(2)} & A \boxtimes_C (B \otimes_C C) \\
 \searrow \alpha_{A,B,C}^{C, \prime} & & \swarrow \text{id}_{A, B \otimes_C C}^\otimes \\
 & A \otimes_C (B \otimes_C C) &
 \end{array}$$

Since subdiagrams (1) and (2) commute by assumption, it follows that the boundary diagram also commutes, i.e.  $\text{id}^\otimes$  satisfies the monoidality condition of **Item 2** of **Definition 13.1.1.1.3**.

This finishes the proof.  $\square$

### 13.1.2 The Moduli Category of Braided Monoidal Structures on a Category

### 13.1.3 The Moduli Category of Symmetric Monoidal Structures on a Category

## 13.2 Moduli Categories of Closed Monoidal Structures

## 13.3 Moduli Categories of Refinements of Monoidal Structures

### 13.3.1 The Moduli Category of Braided Refinements of a Monoidal Structure

# Appendices

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