

Constructions With Monoidal Categories

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01UF This chapter contains some material on constructions with monoidal categories.

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01UG 13.I Moduli Categories of Monoidal Structures

01UH 13.I.I The Moduli Category of Monoidal Structures on a Category

Let C be a category.

01UJ

DEFINITION 13.1.1.1.1 ► THE MODULI CATEGORY OF MONOIDAL STRUCTURES ON A CATEGORY

The **moduli category of monoidal structures on C** is the category $\mathcal{M}_{\mathbb{E}_1}(C)$ defined by

$$\mathcal{M}_{\mathbb{E}_1}(C) \stackrel{\text{def}}{=} \text{pt} \times_{\text{Cats}} \text{MonCats},$$

$$\begin{array}{ccc} \mathcal{M}_{\mathbb{E}_1}(C) & \longrightarrow & \text{MonCats} \\ \downarrow & \lrcorner & \downarrow \tilde{\omega} \\ \text{pt} & \xrightarrow{[C]} & \text{Cats.} \end{array}$$

01UK

REMARK 13.1.1.1.2 ► UNWINDING DEFINITION 13.1.1.1.1, I

In detail, **the moduli category of monoidal structures on C** is the category $\mathcal{M}_{\mathbb{E}_1}(C)$ where:

- *Objects.* The objects of $\mathcal{M}_{\mathbb{E}_1}(C)$ are monoidal categories $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ whose underlying category is C .
- *Morphisms.* A morphism from $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ to $(C, \boxtimes_C, 1'_C, \alpha^{C'}, \lambda^{C'}, \rho^{C'})$ is a strong monoidal functor structure

$$\text{id}_C^\otimes: A \boxtimes_C B \xrightarrow{\sim} A \otimes_C B,$$

$$\text{id}_{1|C}^\otimes: 1'_C \xrightarrow{\sim} 1_C$$

on the identity functor $\text{id}_C: C \rightarrow C$ of C .

- *Identities.* For each $M \stackrel{\text{def}}{=} (C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C) \in \text{Obj}(\mathcal{M}_{\mathbb{E}_1}(C))$, the unit map

$$\mathbb{1}_{M,M}^{\mathcal{M}_{\mathbb{E}_1}(C)}: \text{pt} \rightarrow \text{Hom}_{\mathcal{M}_{\mathbb{E}_1}(C)}(M, M)$$

of $\mathcal{M}_{\mathbb{E}_1}(C)$ at M is defined by

$$\mathrm{id}_M^{\mathcal{M}_{\mathbb{E}_1}(C)} \stackrel{\mathrm{def}}{=} \left(\mathrm{id}_C^{\otimes}, \mathrm{id}_{1|C}^{\otimes} \right),$$

where $\left(\mathrm{id}_C^{\otimes}, \mathrm{id}_{1|C}^{\otimes} \right)$ is the identity monoidal functor of C of ??.

- *Composition.* For each $M, N, P \in \mathrm{Obj}(\mathcal{M}_{\mathbb{E}_1}(C))$, the composition map

$$\circ_{M,N,P}^{\mathcal{M}_{\mathbb{E}_1}(C)} : \mathrm{Hom}_{\mathcal{M}_{\mathbb{E}_1}(C)}(N, P) \times \mathrm{Hom}_{\mathcal{M}_{\mathbb{E}_1}(C)}(M, N) \rightarrow \mathrm{Hom}_{\mathcal{M}_{\mathbb{E}_1}(C)}(M, P)$$

of $\mathcal{M}_{\mathbb{E}_1}(C)$ at (M, N, P) is defined by

$$\left(\mathrm{id}_C^{\otimes, \prime}, \mathrm{id}_{1|C}^{\otimes, \prime} \right) \circ_{M,N,P}^{\mathcal{M}_{\mathbb{E}_1}(C)} \left(\mathrm{id}_C^{\otimes}, \mathrm{id}_{1|C}^{\otimes} \right) \stackrel{\mathrm{def}}{=} \left(\mathrm{id}_C^{\otimes, \prime} \circ \mathrm{id}_C^{\otimes}, \mathrm{id}_{1|C}^{\otimes, \prime} \circ \mathrm{id}_{1|C}^{\otimes} \right).$$

01UL

REMARK 13.1.1.1.3 ► UNWINDING DEFINITION 13.1.1.1, II

In particular, a morphism in $\mathcal{M}_{\mathbb{E}_1}(C)$ from $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ to $(C, \boxtimes_C, 1'_C, \alpha^{C, \prime}, \lambda^{C, \prime}, \rho^{C, \prime})$ satisfies the following conditions:

01UM

- I. *Naturality.* For each pair $f: A \rightarrow X$ and $g: B \rightarrow Y$ of morphisms of C , the diagram

$$\begin{array}{ccc} A \boxtimes_C B & \xrightarrow{f \boxtimes_C g} & X \boxtimes_C Y \\ \mathrm{id}_{A,B}^{\otimes} \downarrow & & \downarrow \mathrm{id}_{X,Y}^{\otimes} \\ A \otimes_C B & \xrightarrow{f \otimes_C g} & X \otimes_C Y \end{array}$$

commutes.

01UN

2. *Monoidality*. For each $A, B, C \in \text{Obj}(C)$, the diagram

$$\begin{array}{ccc}
 & (A \boxtimes_C B) \boxtimes_C C & \\
 \text{id}_{A,B}^\otimes \boxtimes \text{id}_C \swarrow & & \searrow \alpha_{A,B,C}^{C'} \\
 (A \otimes_C B) \boxtimes_C C & & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow \text{id}_{A \otimes_C B, C}^\otimes & & \downarrow \text{id}_A \boxtimes \text{id}_{B,C}^\otimes \\
 (A \otimes_C B) \otimes_C C & & A \boxtimes_C (B \otimes_C C) \\
 \searrow \alpha_{A,B,C}^C & & \swarrow \text{id}_{A,B \otimes_C C}^\otimes \\
 & A \otimes_C (B \otimes_C C) &
 \end{array}$$

commutes.

01UP

3. *Left Monoidal Unity*. For each $A \in \text{Obj}(C)$, the diagram

$$\begin{array}{ccccc}
 & & 1_C \boxtimes_C A & \xrightarrow{\text{id}_{1_C, A}^\otimes} & 1_C \otimes_C A \\
 \text{id}_{1_C}^\otimes \boxtimes \text{id}_A \nearrow & & & & \searrow \lambda_A^C \\
 1'_C \boxtimes_C A & \xrightarrow{\lambda_A^{C'}} & & & A
 \end{array}$$

commutes.

01UQ

4. *Right Monoidal Unity*. For each $A \in \text{Obj}(C)$, the diagram

$$\begin{array}{ccccc}
 & & A \boxtimes_C 1_C & \xrightarrow{\text{id}_{A, 1_C}^\otimes} & A \otimes_C 1_C \\
 \text{id}_A \boxtimes \text{id}_{1_C}^\otimes \nearrow & & & & \searrow \rho_A^C \\
 A \boxtimes_C 1'_C & \xrightarrow{\rho_A^{C'}} & & & A
 \end{array}$$

commutes.

PROPOSITION 13.1.1.4 ► PROPERTIES OF THE MODULI CATEGORY OF MONOIDAL STRUCTURES ON A CATEGORY

01UR

Let \mathcal{C} be a category.

01US

I. *Extra Monoidality Conditions.* Let $(\text{id}_{\mathcal{C}}^{\otimes}, \text{id}_{1|\mathcal{C}}^{\otimes})$ be a morphism of $\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})$ from $(\mathcal{C}, \otimes_{\mathcal{C}}, 1_{\mathcal{C}}, \alpha^{\mathcal{C}}, \lambda^{\mathcal{C}}, \rho^{\mathcal{C}})$ to $(\mathcal{C}, \boxtimes_{\mathcal{C}}, 1'_{\mathcal{C}}, \alpha^{C,'}, \lambda^{C,'}, \rho^{C,'})$.

01UT

(a) The diagram

$$\begin{array}{ccc} (A \boxtimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C & \xrightarrow{\text{id}_{A,B}^{\otimes} \boxtimes_{\mathcal{C}} \text{id}_{\mathcal{C}}} & (A \otimes_{\mathcal{C}} B) \boxtimes_{\mathcal{C}} C \\ \text{id}_{A \boxtimes_{\mathcal{C}} B, C}^{\otimes} \downarrow & & \downarrow \text{id}_{A \otimes_{\mathcal{C}} B, C}^{\otimes} \\ (A \boxtimes_{\mathcal{C}} B) \otimes_{\mathcal{C}} C & \xrightarrow{\text{id}_{A,B}^{\otimes} \otimes_{\mathcal{C}} \text{id}_{\mathcal{C}}} & (A \otimes_{\mathcal{C}} B) \otimes_{\mathcal{C}} C \end{array}$$

commutes.

01UU

(b) The diagram

$$\begin{array}{ccc} A \boxtimes_{\mathcal{C}} (B \boxtimes_{\mathcal{C}} C) & \xrightarrow{\text{id}_A \boxtimes_{\mathcal{C}} \text{id}_{B,C}^{\otimes}} & A \boxtimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C) \\ \text{id}_{A, B \boxtimes_{\mathcal{C}} C}^{\otimes} \downarrow & & \downarrow \text{id}_{A, B \otimes_{\mathcal{C}} C}^{\otimes} \\ A \otimes_{\mathcal{C}} (B \boxtimes_{\mathcal{C}} C) & \xrightarrow{\text{id}_A \otimes_{\mathcal{C}} \text{id}_{B,C}^{\otimes}} & A \otimes_{\mathcal{C}} (B \otimes_{\mathcal{C}} C) \end{array}$$

commutes.

01WB

2. *Extra Monoidal Unity Constraints.* Let $(\text{id}_{\mathcal{C}}^{\otimes}, \text{id}_{1|\mathcal{C}}^{\otimes})$ be a morphism of $\mathcal{M}_{\mathbb{E}_1}(\mathcal{C})$ from $(\mathcal{C}, \otimes_{\mathcal{C}}, 1_{\mathcal{C}}, \alpha^{\mathcal{C}}, \lambda^{\mathcal{C}}, \rho^{\mathcal{C}})$ to $(\mathcal{C}, \boxtimes_{\mathcal{C}}, 1'_{\mathcal{C}}, \alpha^{C,'}, \lambda^{C,'}, \rho^{C,'})$.

01WC

(a) The diagram

$$\begin{array}{ccc}
 1_C \otimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \boxtimes_C 1'_C \\
 \lambda_{1'_C}^C \downarrow & & \downarrow \rho_{1_C}^{C, 'C} \\
 1'_C & \xrightarrow{\text{id}_1^{\otimes}} & 1_C
 \end{array}$$

commutes.

01WD

(b) The diagram

$$\begin{array}{ccc}
 1'_C \otimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1'_C \boxtimes_C 1_C \\
 \rho_{1'_C}^C \downarrow & & \downarrow \lambda_{1_C}^{C, 'C} \\
 1'_C & \xrightarrow{\text{id}_1^{\otimes}} & 1_C
 \end{array}$$

commutes.

01WE

(c) The diagram

$$\begin{array}{ccc}
 1'_C \boxtimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes}} & 1'_C \otimes_C 1_C \\
 \lambda_{1_C}^{C, 'C} \downarrow & & \downarrow \rho_{1'_C}^C \\
 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1'_C
 \end{array}$$

commutes.

01WF

(d) The diagram

$$\begin{array}{ccc}
 1_C \boxtimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes}} & 1_C \otimes_C 1'_C \\
 \rho_{1_C}^{C, 'C} \downarrow & & \downarrow \lambda_{1'_C}^C \\
 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1'_C
 \end{array}$$

commutes.

01UV

3. *Mixed Associators.* Let $(C, \otimes_C, 1_C, \alpha^C, \lambda^C, \rho^C)$ and $(C, \boxtimes_C, 1'_C, \alpha^{C'}, \lambda^{C'}, \rho^{C'})$ be monoidal structures on C and let

$$\text{id}_{-1, -2}^{\otimes}: -_1 \boxtimes_C -_2 \rightarrow -_1 \otimes_C -_2$$

be a natural transformation.

01UW

(a) If there exists a natural transformation

$$\alpha_{A,B,C}^{\otimes}: (A \otimes_C B) \boxtimes_C C \rightarrow A \otimes_C (B \boxtimes_C C)$$

making the diagrams

$$\begin{array}{ccc} (A \otimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A,B,C}^{\otimes}} & A \otimes_C (B \boxtimes_C C) \\ \text{id}_{A \otimes_C B, C}^{\otimes} \downarrow & & \downarrow \text{id}_A \otimes \text{id}_{B,C}^{\otimes} \\ (A \otimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^C} & A \otimes_C (B \otimes_C C) \end{array}$$

and

$$\begin{array}{ccc} (A \boxtimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A,B,C}^{C'}} & A \boxtimes_C (B \boxtimes_C C) \\ \text{id}_{A \boxtimes_C B, C}^{\otimes} \boxtimes \text{id}_C \downarrow & & \downarrow \text{id}_{A, B \boxtimes_C C} \\ (A \otimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A,B,C}^{\otimes}} & A \otimes_C (B \boxtimes_C C) \end{array}$$

commute, then the natural transformation id^{\otimes} satisfies the monoidality condition of [Item 2](#) of [Remark 13.1.1.1.3](#).

01UX

(b) If there exists a natural transformation

$$\alpha_{A,B,C}^{\boxtimes}: (A \boxtimes_C B) \otimes_C C \rightarrow A \boxtimes_C (B \otimes_C C)$$

making the diagrams

$$\begin{array}{ccc} (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes}} & A \boxtimes_C (B \otimes_C C) \\ \text{id}_{A \boxtimes_C B, C}^{\otimes} \otimes \text{id}_C \downarrow & & \downarrow \text{id}_{A, B \otimes_C C}^{\otimes} \\ (A \otimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^C} & A \otimes_C (B \otimes_C C) \end{array}$$

and

$$\begin{array}{ccc}
 (A \boxtimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A,B,C}^{C'}} & A \boxtimes_C (B \boxtimes_C C) \\
 \text{id}_{A \boxtimes_C B, C}^\otimes \downarrow & & \downarrow \text{id}_A \boxtimes_C \text{id}_{B, C}^\otimes \\
 (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^\otimes} & A \boxtimes_C (B \otimes_C C)
 \end{array}$$

commute, then the natural transformation id^\otimes satisfies the monoidality condition of **Item 2** of **Remark 13.1.1.1.3**.

01UY

(c) If there exists a natural transformation

$$\alpha_{A,B,C}^{\boxtimes, \otimes}: (A \boxtimes_C B) \otimes_C C \rightarrow A \otimes_C (B \boxtimes_C C)$$

making the diagrams

$$\begin{array}{ccc}
 (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes, \otimes}} & A \otimes_C (B \boxtimes_C C) \\
 \text{id}_{A, B}^\otimes \otimes \text{id}_C \downarrow & & \downarrow \text{id}_A \otimes_C \text{id}_{B, C}^\otimes \\
 (A \otimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^C} & A \otimes_C (B \otimes_C C)
 \end{array}$$

and

$$\begin{array}{ccc}
 (A \boxtimes_C B) \boxtimes_C C & \xrightarrow{\alpha_{A,B,C}^{C'}} & A \boxtimes_C (B \boxtimes_C C) \\
 \text{id}_{A \boxtimes_C B, C}^\otimes \downarrow & & \downarrow \text{id}_{A, B \boxtimes_C C}^\otimes \\
 (A \boxtimes_C B) \otimes_C C & \xrightarrow{\alpha_{A,B,C}^{\boxtimes, \otimes}} & A \otimes_C (B \boxtimes_C C)
 \end{array}$$

commute, then the natural transformation id^\otimes satisfies the monoidality condition of **Item 2** of **Remark 13.1.1.1.3**.

PROOF 13.1.1.1.5 ► PROOF OF PROPOSITION 13.1.1.1.4**Item 1: Extra Monoidality Conditions**

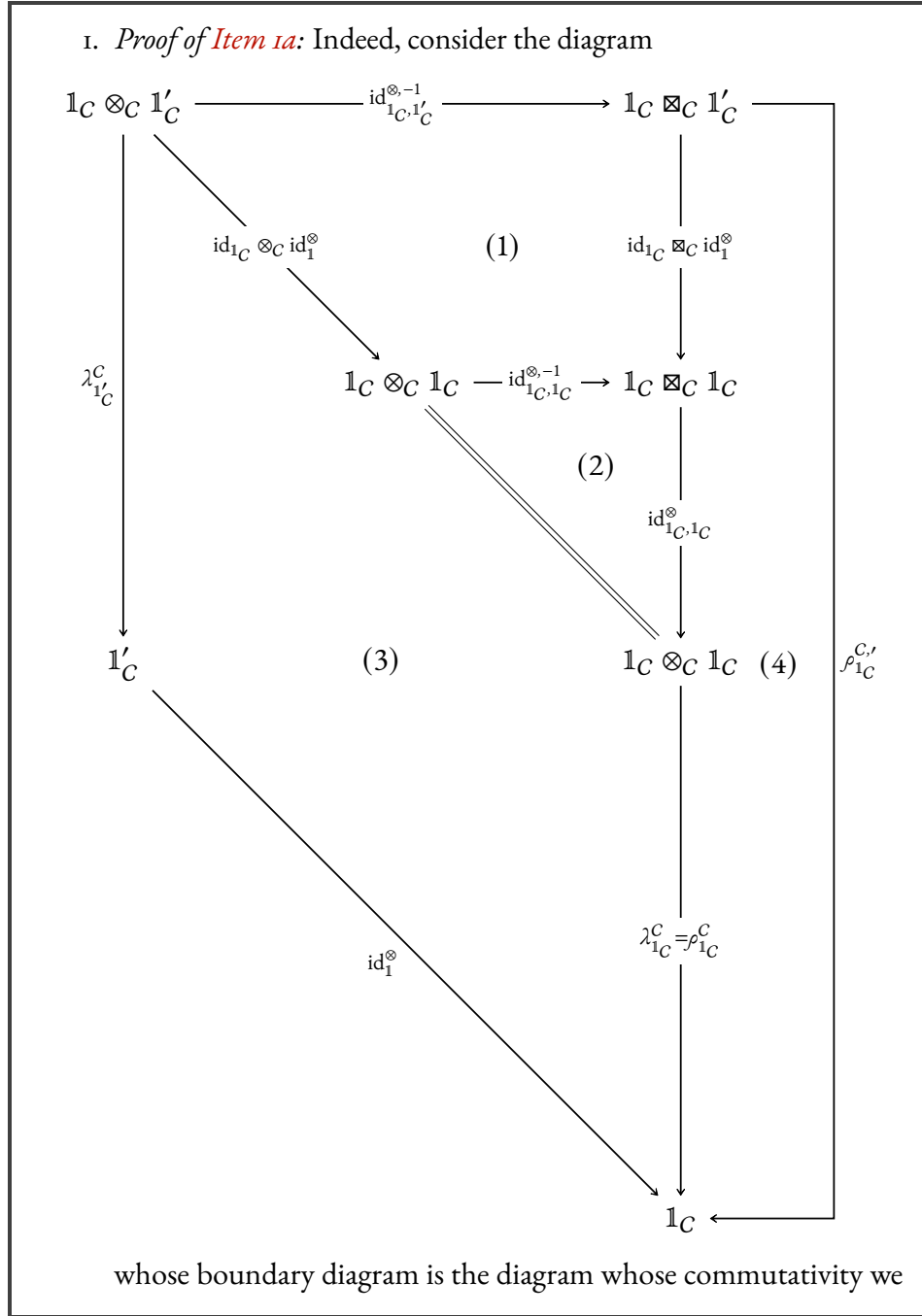
We claim that **Items 1a** and **1b** are indeed true:

1. *Proof of Item 1a:* This follows from the naturality of id^\otimes with respect to the morphisms $\mathrm{id}_{A,B}^\otimes$ and id_C .
2. *Proof of Item 1b:* This follows from the naturality of id^\otimes with respect to the morphisms id_A and $\mathrm{id}_{B,C}^\otimes$.

This finishes the proof.

Item 2: Extra Monoidal Unity Constraints

We claim that **Items 2a** and **2b** are indeed true:

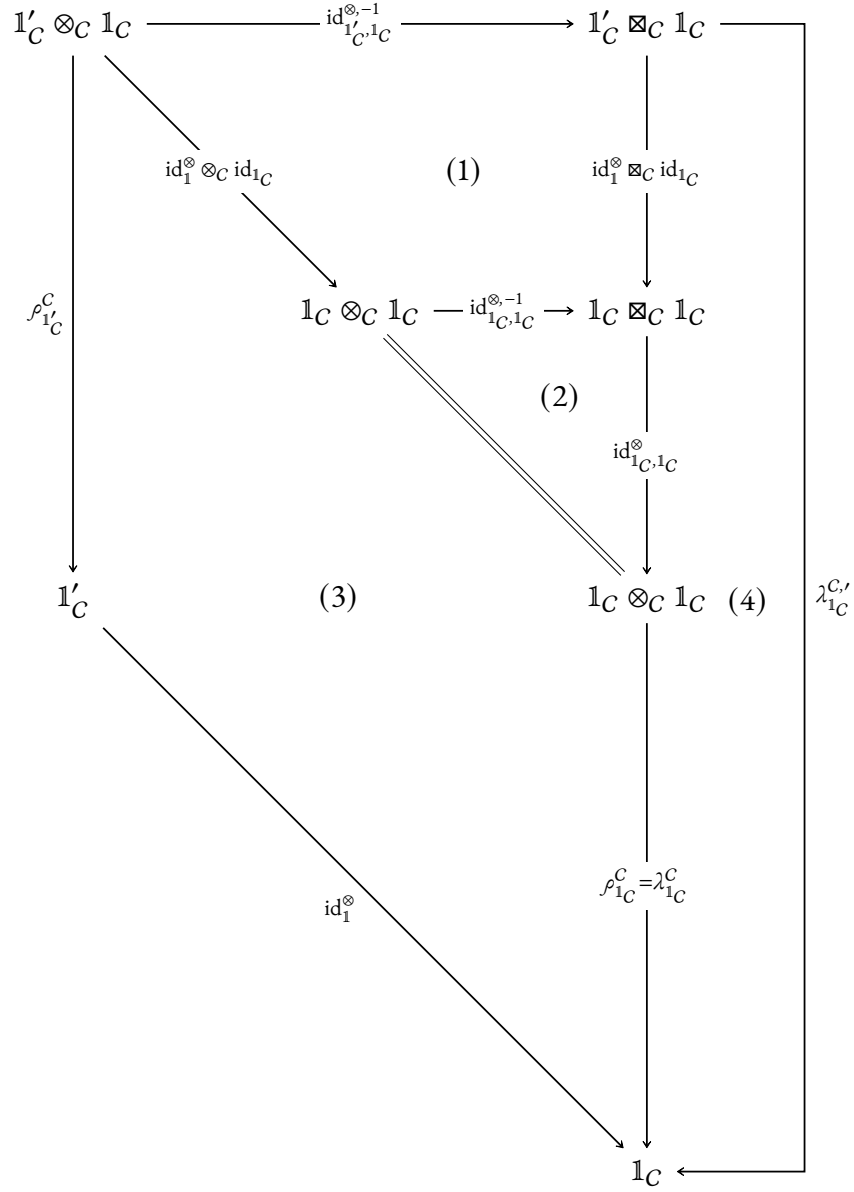


wish to prove. Since:

- Subdiagram (1) commutes by the naturality of $\mathrm{id}_C^{\otimes, -1}$;
- Subdiagram (2) commutes trivially;
- Subdiagram (3) commutes by the naturality of λ^C , where the equality $\rho_{1_C}^C = \lambda_{1_C}^C$ comes from ??;
- Subdiagram (4) commutes by the right monoidal unity of $(\mathrm{id}_C, \mathrm{id}_C^{\otimes}, \mathrm{id}_{C|1}^{\otimes})$;

so does the boundary diagram, and we are done.

2. *Proof of Item 1b:* Indeed, consider the diagram



whose boundary diagram is the diagram whose commutativity we

wish to prove. Since:

- Subdiagram (1) commutes by the naturality of $\text{id}_C^{\otimes, -1}$;
- Subdiagram (2) commutes trivially;
- Subdiagram (3) commutes by the naturality of ρ^C , where the equality $\rho_{1_C}^C = \lambda_{1_C}^C$ comes from ??;
- Subdiagram (4) commutes by the left monoidal unity of $(\text{id}_C, \text{id}_C^{\otimes}, \text{id}_{C|1}^{\otimes})$;

so does the boundary diagram, and we are done.

3. *Proof of Item 2c:* Indeed, consider the diagram

$$\begin{array}{ccccc}
 1'_C \otimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1'_C \boxtimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes}} & 1'_C \otimes_C 1_C \\
 \rho_{1'_C}^C \downarrow & & \downarrow \lambda_{1_C}^{C, '}& & \downarrow \rho_{1'_C}^C \\
 1'_C & \xrightarrow{\text{id}_1^{\otimes}} & 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1'_C.
 \end{array}
 \quad \begin{array}{c} (1) \end{array} \quad \begin{array}{c} (\dagger)$$

Since:

- The boundary diagram commutes trivially;
- Subdiagram (1) commutes by **Item 1b**;

it follows that the diagram

$$\begin{array}{ccccc}
 1'_C \otimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1'_C \boxtimes_C 1_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes}} & 1'_C \otimes_C 1_C \\
 & & \downarrow \lambda_{1_C}^{C, '}& & \downarrow \rho_{1'_C}^C \\
 & & 1_C & \xrightarrow{\text{id}_1^{\otimes, -1}} & 1'_C
 \end{array}
 \quad \begin{array}{c} (\dagger)$$

commutes. But since $\text{id}_{1_C, 1'_C}^{\otimes, -1}$ is an isomorphism, it follows that the diagram (\dagger) also commutes, and we are done.

4. *Proof of Item 2d:* Indeed, consider the diagram

$$\begin{array}{ccccc}
 1_C \otimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \boxtimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \otimes_C 1'_C \\
 \downarrow \lambda_{1'_C}^C & & \downarrow \rho_{1_C}^{C, \prime} & & \downarrow \lambda_{1'_C}^C \\
 1'_C & \xrightarrow{\text{id}_{1'_C}^{\otimes}} & 1_C & \xrightarrow{\text{id}_{1_C}^{\otimes, -1}} & 1_C
 \end{array}
 \quad (1) \qquad (\dagger)$$

Since:

- The boundary diagram commutes trivially;
- Subdiagram (1) commutes by **Item 1a**;

it follows that the diagram

$$\begin{array}{ccccc}
 1_C \otimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \boxtimes_C 1'_C & \xrightarrow{\text{id}_{1_C, 1'_C}^{\otimes, -1}} & 1_C \otimes_C 1'_C \\
 & & \downarrow \rho_{1_C}^{C, \prime} & & \downarrow \lambda_{1'_C}^C \\
 & & 1_C & \xrightarrow{\text{id}_{1_C}^{\otimes, -1}} & 1_C
 \end{array}
 \quad (\dagger)$$

commutes. But since $\text{id}_{1_C}^{\otimes, -1}$ is an isomorphism, it follows that the diagram (\dagger) also commutes, and we are done.

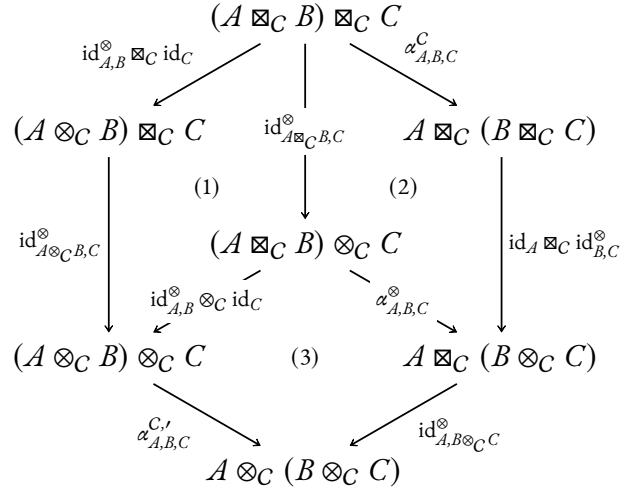
This finishes the proof.

Item 3: Mixed Associators

We claim that **Items 3a** to **3c** are indeed true:

01UZ

I. *Proof of Item 3a:* We may partition the monoidality diagram for id^\otimes of **Item 2** of **Remark 13.1.1.1.3** as follows:



Since:

- Subdiagram (1) commutes by **Item 1a** of **Item 1**.
- Subdiagram (2) commutes by assumption.
- Subdiagram (3) commutes by assumption.

it follows that the boundary diagram also commutes, i.e. id^\otimes satisfies the monoidality condition of **Item 2** of **Remark 13.1.1.1.3**.

01V0

2. *Proof of Item 3b:* We may partition the monoidality diagram for id^\otimes

of **Item 2** of **Remark 13.1.1.1.3** as follows:

$$\begin{array}{ccccc}
 & & (A \boxtimes_C B) \boxtimes_C C & & \\
 & \swarrow \text{id}_{A,B}^\otimes \boxtimes \text{id}_C & & \searrow \alpha_{A,B,C}^C & \\
 (A \otimes_C B) \boxtimes_C C & & (1) & & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow \text{id}_{A \otimes_C B, C}^\otimes & \swarrow \alpha_{A,B,C}^\boxtimes & & \swarrow \text{id}_{A,B \boxtimes_C C}^\otimes & \downarrow \text{id}_A \boxtimes \text{id}_{B,C}^\otimes \\
 & A \otimes_C (B \boxtimes_C C) & & & \\
 & (2) & & (3) & \\
 (A \otimes_C B) \otimes_C C & \xrightarrow{\text{id}_A \otimes \text{id}_{B,C}^\otimes} & A \otimes_C (B \otimes_C C) & \xleftarrow{\text{id}_{A,B \otimes_C C}^\otimes} & A \boxtimes_C (B \otimes_C C) \\
 \searrow \alpha_{A,B,C}^{C'} & & \downarrow & & \swarrow \text{id}_{A,B \otimes_C C}^\otimes \\
 & A \otimes_C (B \otimes_C C) & & &
 \end{array}$$

Since:

- Subdiagram (1) commutes by assumption.
- Subdiagram (2) commutes by assumption.
- Subdiagram (3) commutes by **Item 1b** of **Item 1**.

it follows that the boundary diagram also commutes, i.e. id^\otimes satisfies the monoidality condition of **Item 2** of **Remark 13.1.1.1.3**.

01V1


3. *Proof of **Item 3c**:* We may partition the monoidality diagram for id^\otimes

of **Item 2** of **Remark 13.1.1.1.3** as follows:

$$\begin{array}{ccc}
 & (A \boxtimes_C B) \boxtimes_C C & \\
 \text{id}_{A,B}^\otimes \boxtimes \text{id}_C \swarrow & & \searrow \alpha_{A,B,C}^C \\
 (A \otimes_C B) \boxtimes_C C & & A \boxtimes_C (B \boxtimes_C C) \\
 \downarrow \text{id}_{A \otimes_C B, C}^\otimes & \searrow \alpha_{A,B,C}^{\boxtimes, \otimes} & \downarrow \text{id}_A \boxtimes \text{id}_{B,C}^\otimes \\
 (A \otimes_C B) \otimes_C C & & A \boxtimes_C (B \otimes_C C) \\
 \searrow \alpha_{A,B,C}^{C, \otimes} & & \swarrow \text{id}_{A, B \otimes_C C}^\otimes \\
 & A \otimes_C (B \otimes_C C) &
 \end{array}
 \quad (1)$$

(2)

Since subdiagrams (1) and (2) commute by assumption, it follows that the boundary diagram also commutes, i.e. id^\otimes satisfies the monoidality condition of **Item 2** of **Remark 13.1.1.1.3**.

This finishes the proof. 

- 01V2 13.1.2 The Moduli Category of Braided Monoidal Structures on a Category
- 01V3 13.1.3 The Moduli Category of Symmetric Monoidal Structures on a Category
- 01V4 13.2 Moduli Categories of Closed Monoidal Structures
- 01V5 13.3 Moduli Categories of Refinements of Monoidal Structures
- 01V6 13.3.1 The Moduli Category of Braided Refinements of a Monoidal Structure

Appendices

A Other Chapters

Preliminaries

- 1. Introduction
- 2. A Guide to the Literature

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- 3. Sets
- 4. Constructions With Sets
- 5. Monoidal Structures on the Category of Sets
- 6. Pointed Sets

7. Tensor Products of Pointed Sets

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- 8. Relations
- 9. Constructions With Relations
- 10. Conditions on Relations

Categories

- 11. Categories
- 12. Presheaves and the Yoneda Lemma

Monoidal Categories

13. Constructions With Monoidal Categories

Bicategories

Extra Part

14. Types of Morphisms in Bicate- 15. Notes