# **Constructions With Sets**

### The Clowder Project Authors

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This chapter develops some material relating to constructions with sets with an eye towards its categorical and higher-categorical counterparts to be introduced later in this work. Of particular interest are perhaps the following:

- 1. Explicit descriptions of the major types of co/limits in Sets, including in particular explicit descriptions of pushouts and coequalisers (see Definitions 4.2.4.1.1 and 4.2.5.1.1 and Remarks 4.2.4.1.4 and 4.2.5.1.4).
- 2. A discussion of powersets as decategorifications of categories of presheaves, including in particular results such as:
  - (a) A discussion of the internal Hom of a powerset (Section 4.4.7).
  - (b) A 0-categorical version of the Yoneda lemma (Presheaves and the Yoneda Lemma, Theorem 12.1.5.1.1), which we term the *Yoneda lemma for sets* (Proposition 4.5.5.1.1).
  - (c) A characterisation of powersets as free cocompletions (Section 4.4.5), mimicking the corresponding statement for categories of presheaves (??).
  - (d) A characterisation of powersets as free completions (Section 4.4.6), mimicking the corresponding statement for categories of copresheaves (??).
  - (e) A (-1)-categorical version of un/straightening (Item 2 of Proposition 4.5.1.1.4 and Remark 4.5.1.1.6).
  - (f) A 0-categorical form of Isbell duality internal to powersets (Section 4.4.8).
- 3. A lengthy discussion of the adjoint triple

$$f_! \dashv f^{-1} \dashv f_* \colon \mathcal{P}(A) \xrightarrow{\rightleftharpoons} \mathcal{P}(B)$$

of functors (i.e. morphisms of posets) between  $\mathcal{P}(A)$  and  $\mathcal{P}(B)$  induced by a map of sets  $f: A \to B$ , including in particular:

Contents 2

(a) How  $f^{-1}$  can be described as a precomposition while  $f_!$  and  $f_*$  can be described as Kan extensions (Remarks 4.6.1.1.4, 4.6.2.1.2 and 4.6.3.1.4).

- (b) An extensive list of the properties of  $f_1$ ,  $f^{-1}$ , and  $f_*$  (Propositions 4.6.1.1.5, 4.6.1.1.7, 4.6.2.1.3, 4.6.2.1.5, 4.6.3.1.7 and 4.6.3.1.9).
- (c) How the functors  $f_!$ ,  $f^{-1}$ ,  $f_*$ , along with the functors

$$-_{1} \cap -_{2} \colon \mathcal{P}(X) \times \mathcal{P}(X) \to \mathcal{P}(X),$$
$$[-_{1}, -_{2}]_{X} \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

may be viewed as a six-functor formalism with the empty set  $\emptyset$  as the dualising object (Section 4.6.4).

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# 4.1 Limits of Sets

# 4.1.1 The Terminal Set

# **DEFINITION 4.1.1.1.1 ► THE TERMINAL SET**

The **terminal set** is the terminal object of Sets as in Limits and Colimits, ??.

#### CONSTRUCTION 4.1.1.1.2 ► CONSTRUCTION OF THE TERMINAL SET

Concretely, the terminal set is the pair  $(pt, \{!_A\}_{A \in Obj(Sets)})$  consisting of:

- 1. *The Limit*. The punctual set pt  $\stackrel{\text{def}}{=} \{ \star \}$ .
- 2. The Cone. The collection of maps

$$\{!_A : A \to \mathsf{pt}\}_{A \in \mathsf{Obj}(\mathsf{Sets})}$$

defined by

$$!_A(a) \stackrel{\text{def}}{=} \star$$

for each  $a \in A$  and each  $A \in Obj(Sets)$ .

### PROOF 4.1.1.1.3 ► PROOF OF CONSTRUCTION 4.1.1.1.2

We claim that pt is the terminal object of Sets. Indeed, suppose we have a diagram of the form

in Sets. Then there exists a unique map  $\phi: A \to \operatorname{pt}$  making the diagram

$$A - \frac{\phi}{\exists 1} \rightarrow pt$$

commute, namely  $!_A$ .

### 4.1.2 Products of Families of Sets

Let  $\{A_i\}_{i\in I}$  be a family of sets.

### **DEFINITION 4.1.2.1.1** ► THE PRODUCT OF A FAMILY OF SETS

The **product**<sup>1</sup> **of**  $\{A_i\}_{i\in I}$  is the product of  $\{A_i\}_{i\in I}$  in Sets as in Limits and Colimits, **??**.

<sup>&</sup>lt;sup>1</sup>Further Terminology: Also called the **Cartesian product of**  $\{A_i\}_{i\in I}$ .

#### CONSTRUCTION 4.1.2.1.2 ► CONSTRUCTION OF THE PRODUCT OF A FAMILY OF SETS

Concretely, the product of  $\{A_i\}_{i\in I}$  is the pair  $\left(\prod_{i\in I}A_i, \left\{\operatorname{pr}_i\right\}_{i\in I}\right)$  consisting of:

1. The Limit. The set  $\prod_{i \in I} A_i$  defined by

$$\prod_{i \in I} A_i \stackrel{\text{def}}{=} \left\{ f \in \mathsf{Sets} \left( I, \bigcup_{i \in I} A_i \right) \middle| \begin{array}{l} \text{for each } i \in I, \text{ we} \\ \text{have } f(i) \in A_i \end{array} \right\}.$$

2. The Cone. The collection

$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

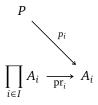
of maps given by

$$\operatorname{pr}_{i}(f) \stackrel{\text{def}}{=} f(i)$$

for each  $f \in \prod_{i \in I} A_i$  and each  $i \in I$ .

#### PROOF 4.1.2.1.3 ► PROOF OF CONSTRUCTION 4.1.2.1.2

We claim that  $\prod_{i \in I} A_i$  is the categorical product of  $\{A_i\}_{i \in I}$  in Sets. Indeed, suppose we have, for each  $i \in I$ , a diagram of the form



in Sets. Then there exists a unique map  $\phi \colon P \to \prod_{i \in I} A_i$  making the diagram

$$P$$

$$\phi \mid \exists ! \qquad p_i$$

$$\downarrow \qquad \qquad \downarrow$$

$$\prod_{i \in I} A_i \xrightarrow{pr_i} A_i$$

commute, being uniquely determined by the condition  $\operatorname{pr}_i \circ \phi = p_i$  for each  $i \in I$  via

$$\phi(x) = (p_i(x))_{i \in I}$$

for each  $x \in P$ .

#### REMARK 4.1.2.1.4 ► UNWINDING CONSTRUCTION 4.1.2.1.2

Less formally, we may think of Cartesian products and projection maps as follows:

- 1. We think of  $\prod_{i \in I} A_i$  as the set whose elements are *I*-indexed collections  $(a_i)_{i \in I}$  with  $a_i \in A_i$  for each  $i \in I$ .
- 2. We view the projection maps

$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

as being given by

$$\operatorname{pr}_i\left(\left(a_j\right)_{j\in I}\right)\stackrel{\text{def}}{=}a_i$$

for each  $(a_j)_{i \in I} \in \prod_{i \in I} A_i$  and each  $i \in I$ .

### PROPOSITION 4.1.2.1.5 ► PROPERTIES OF PRODUCTS OF FAMILIES OF SETS

Let  $\{A_i\}_{i\in I}$  be a family of sets.

1. Functoriality. The assignment  $\{A_i\}_{i\in I}\mapsto \prod_{i\in I}A_i$  defines a functor

$$\prod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

• *Action on Objects.* For each  $(A_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$ , we have

$$\left[\prod_{i\in I}\right]((A_i)_{i\in I})\stackrel{\text{def}}{=}\prod_{i\in I}A_i$$

• Action on Morphisms. For each  $(A_i)_{i \in I}, (B_i)_{i \in I} \in$ 

 $Obj(Fun(I_{disc}, Sets))$ , the action on Hom-sets

$$\left(\prod_{i\in I}\right)_{(A_i)_{i\in I},(B_i)_{i\in I}}: \operatorname{Nat}((A_i)_{i\in I},(B_i)_{i\in I}) \to \operatorname{Sets}\left(\prod_{i\in I}A_i,\prod_{i\in I}B_i\right)$$

of  $\prod_{i \in I}$  at  $((A_i)_{i \in I}, (B_i)_{i \in I})$  is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in  $Nat((A_i)_{i \in I}, (B_i)_{i \in I})$  to the map of sets

$$\prod_{i \in I} f_i \colon \prod_{i \in I} A_i \to \prod_{i \in I} B_i$$

defined by

$$\left[\prod_{i\in I} f_i\right] ((a_i)_{i\in I}) \stackrel{\text{def}}{=} (f_i(a_i))_{i\in I}$$

for each  $(a_i)_{i \in I} \in \prod_{i \in I} A_i$ .

#### PROOF 4.1.2.1.6 ► PROOF OF PROPOSITION 4.1.2.1.5

### Item 1: Functoriality

This follows from Limits and Colimits, ?? of ??.

# 4.1.3 Binary Products of Sets

Let A and B be sets.

### **DEFINITION 4.1.3.1.1** ► BINARY PRODUCTS OF SETS

The **product of** A and  $B^1$  is the product of A and B in Sets as in Limits and Colimits, ??.

<sup>1</sup>Further Terminology: Also called the **Cartesian product of** A **and** B.

#### CONSTRUCTION 4.1.3.1.2 ► CONSTRUCTION OF BINARY PRODUCTS OF SETS

Concretely, the product of *A* and *B* is the pair  $(A \times B, \{pr_1, pr_2\})$  consisting of:

1. The Limit. The set  $A \times B$  defined by

$$A \times B \stackrel{\text{def}}{=} \prod_{z \in \{A, B\}} z$$

$$\stackrel{\text{def}}{=} \{ f \in \text{Sets}(\{0, 1\}, A \cup B) \mid \text{we have } f(0) \in A \text{ and } f(1) \in B \}$$

$$\cong \{ \{ \{a\}, \{a, b\} \} \in \mathcal{P}(\mathcal{P}(A \cup B)) \mid \text{we have } a \in A \text{ and } b \in B \}$$

$$\cong \begin{cases} \text{ordered pairs } (a, b) \text{ with } \\ a \in A \text{ and } b \in B \end{cases}.$$

2. The Cone. The maps

$$\operatorname{pr}_1: A \times B \to A,$$
  
 $\operatorname{pr}_2: A \times B \to B$ 

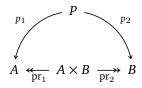
defined by

$$\operatorname{pr}_{1}(a,b) \stackrel{\text{def}}{=} a,$$
  
 $\operatorname{pr}_{2}(a,b) \stackrel{\text{def}}{=} b$ 

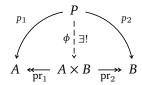
for each  $(a, b) \in A \times B$ .

### PROOF 4.1.3.1.3 ► PROOF OF CONSTRUCTION 4.1.3.1.2

We claim that  $A \times B$  is the categorical product of A and B in the category of sets. Indeed, suppose we have a diagram of the form



in Sets. Then there exists a unique map  $\phi: P \to A \times B$  making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$
  
 $\operatorname{pr}_2 \circ \phi = p_2$ 

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each  $x \in P$ .

### PROPOSITION 4.1.3.1.4 ► PROPERTIES OF PRODUCTS OF SETS

Let A, B, C, and X be sets.

1. Functoriality. The assignments  $A, B, (A, B) \mapsto A \times B$  define functors

$$A \times -:$$
 Sets  $\rightarrow$  Sets,  
 $- \times B:$  Sets  $\rightarrow$  Sets,  
 $-_1 \times -_2:$  Sets  $\times$  Sets  $\rightarrow$  Sets,

where  $-1 \times -2$  is the functor where

• *Action on Objects.* For each  $(A, B) \in \mathsf{Obj}(\mathsf{Sets} \times \mathsf{Sets})$ , we have

$$[-1 \times -2](A, B) \stackrel{\text{def}}{=} A \times B.$$

• *Action on Morphisms.* For each  $(A, B), (X, Y) \in Obj(Sets)$ , the action on Hom-sets

$$\times_{(A,B),(X,Y)}$$
:  $\mathsf{Sets}(A,X) \times \mathsf{Sets}(B,Y) \to \mathsf{Sets}(A \times B, X \times Y)$ 

of  $\times$  at ((A, B), (X, Y)) is defined by sending (f, g) to the function

$$f \times g \colon A \times B \to X \times Y$$

defined by

$$[f \times g](a,b) \stackrel{\text{def}}{=} (f(a),g(b))$$

for each  $(a, b) \in A \times B$ .

and where  $A \times -$  and  $- \times B$  are the partial functors of  $-_1 \times -_2$  at  $A, B \in \text{Obj}(\mathsf{Sets})$ .

2. Adjointness I. We have adjunctions

$$(A \times - \dashv \operatorname{Sets}(A, -))$$
: Sets  $\xrightarrow{A \times -}$  Sets,  $\subseteq$  Se

witnessed by bijections

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$
  
 $Sets(A \times B, C) \cong Sets(B, Sets(A, C)),$ 

natural in  $A, B, C \in Obj(Sets)$ .

3. Adjointness II. We have an adjunction

$$(\Delta_{\mathsf{Sets}} \dashv -_1 \times -_2)$$
: Sets  $\stackrel{\Delta_{\mathsf{Sets}}}{=}$  Sets  $\times$  Sets,

witnessed by a bijection

$$Hom_{Sets \times Sets}((A, A), (B, C)) \cong Sets(A, B \times C),$$

natural in  $A \in \text{Obj}(\mathsf{Sets})$  and in  $(B, C) \in \text{Obj}(\mathsf{Sets} \times \mathsf{Sets})$ .

4. Associativity. We have an isomorphism of sets

$$\alpha_{A.B.C}^{\mathsf{Sets}} : (A \times B) \times C \xrightarrow{\sim} A \times (B \times C),$$

natural in  $A, B, C \in Obj(Sets)$ .

5. Unitality. We have isomorphisms of sets

$$\lambda_A^{\mathsf{Sets}} \colon \mathsf{pt} \times A \xrightarrow{\sim} A,$$
  
 $\rho_A^{\mathsf{Sets}} \colon A \times \mathsf{pt} \xrightarrow{\sim} A,$ 

natural in  $A \in Obj(Sets)$ .

6. Commutativity. We have an isomorphism of sets

$$\sigma_{A,B}^{\mathsf{Sets}} \colon A \times B \xrightarrow{\sim} B \times A,$$

natural in  $A, B \in Obj(Sets)$ .

7. Distributivity Over Coproducts. We have isomorphisms of sets

$$\delta_{\ell}^{\mathsf{Sets}} : A \times (B \coprod C) \xrightarrow{\sim} (A \times B) \coprod (A \times C),$$
  
 $\delta_{r}^{\mathsf{Sets}} : (A \coprod B) \times C \xrightarrow{\sim} (A \times C) \coprod (B \times C),$ 

natural in  $A, B, C \in Obj(Sets)$ .

8. Annihilation With the Empty Set. We have isomorphisms of sets

$$\zeta_{\ell}^{\mathsf{Sets}} : \emptyset \times A \xrightarrow{\sim} \emptyset,$$

$$\zeta_{r}^{\mathsf{Sets}} : A \times \emptyset \xrightarrow{\sim} \emptyset,$$

natural in  $A \in Obj(Sets)$ .

9. *Distributivity Over Unions*. Let *X* be a set. For each  $U, V, W \in \mathcal{P}(X)$ , we have equalities

$$U \times (V \cup W) = (U \times V) \cup (U \times W),$$
  
$$(U \cup V) \times W = (U \times W) \cup (V \times W)$$

of subsets of  $\mathcal{P}(X \times X)$ .

10. Distributivity Over Intersections. Let X be a set. For each  $U, V, W \in \mathcal{P}(X)$ , we have equalities

$$U \times (V \cap W) = (U \times V) \cap (U \times W),$$
  
$$(U \cap V) \times W = (U \times W) \cap (V \times W)$$

of subsets of  $\mathcal{P}(X \times X)$ .

11. *Distributivity Over Differences*. Let *X* be a set. For each  $U, V, W \in \mathcal{P}(X)$ , we have equalities

$$U \times (V \setminus W) = (U \times V) \setminus (U \times W),$$
  
$$(U \setminus V) \times W = (U \times W) \setminus (V \times W)$$

of subsets of  $\mathcal{P}(X \times X)$ .

12. Distributivity Over Symmetric Differences. Let X be a set. For each  $U, V, W \in \mathcal{P}(X)$ , we have equalities

$$U \times (V \triangle W) = (U \times V) \triangle (U \times W),$$
  
$$(U \triangle V) \times W = (U \times W) \triangle (V \times W)$$

of subsets of  $\mathcal{P}(X \times X)$ .

13. Middle-Four Exchange with Respect to Intersections. The diagram

$$(\mathcal{P}(X) \times \mathcal{P}(X)) \times (\mathcal{P}(X) \times \mathcal{P}(X)) \xrightarrow{\cap \times \cap} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow^{\mathcal{P}_{X,X}^{\times}} \times \mathcal{P}_{X,X}^{\times}} \qquad \qquad \downarrow^{\mathcal{P}_{X,X}^{\times}}$$

$$\mathcal{P}(X \times X) \times \mathcal{P}(X \times X) \xrightarrow{\cap} \mathcal{P}(X \times X)$$

commutes, i.e. we have

$$(U \times V) \cap (W \times T) = (U \cap V) \times (W \cap T).$$

for each  $U, V, W, T \in \mathcal{P}(X)$ .

- 14. *Symmetric Monoidality*. The 8-tuple (Sets,  $\times$ , pt, Sets(-1, -2),  $\alpha^{Sets}$ ,  $\alpha^{Sets}$ ,  $\alpha^{Sets}$ ) is a closed symmetric monoidal category.
- 15. Symmetric Bimonoidality. The 18-tuple

$$\left(\mathsf{Sets}, \coprod, \times, \emptyset, \mathsf{pt}, \mathsf{Sets}(-_1, -_2), \alpha^{\mathsf{Sets}}, \lambda^{\mathsf{Sets}}, \rho^{\mathsf{Sets}}, \sigma^{\mathsf{Sets}}, \alpha^{\mathsf{Sets}}, \alpha^{\mathsf{Set$$

is a symmetric closed bimonoidal category, where  $\alpha^{Sets, \coprod}$ ,  $\lambda^{Sets, \coprod}$ ,  $\rho^{Sets, \coprod}$ , and  $\sigma^{Sets, \coprod}$  are the natural transformations from Items 3 to 5 of Proposition 4.2.3.1.4.

#### PROOF 4.1.3.1.5 ► PROOF OF PROPOSITION 4.1.3.1.4

#### Item 1: Functoriality

This follows from Limits and Colimits, ?? of ??.

### Item 2: Adjointness

We prove only that there's an adjunction  $- \times B \dashv \mathsf{Sets}(B, -)$ , witnessed by a bijection

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$

natural in  $B, C \in \text{Obj}(\mathsf{Sets})$ , as the proof of the existence of the adjunction  $A \times - \exists \mathsf{Sets}(A, -)$  follows almost exactly in the same way.

• Map I. We define a map

$$\Phi_{B,C} \colon \mathsf{Sets}(A \times B, C) \to \mathsf{Sets}(A, \mathsf{Sets}(B, C)),$$

by sending a function

$$\xi: A \times B \to C$$

to the function

$$\xi^{\dagger} \colon A \longrightarrow \mathsf{Sets}(B,C),$$

$$a \mapsto \Big(\xi_a^{\dagger} \colon B \to C\Big),$$

where we define

$$\xi_a^{\dagger}(b) \stackrel{\text{def}}{=} \xi(a,b)$$

for each  $b \in B$ . In terms of the  $[a \mapsto f(a)]$  notation of Sets, Notation 3.1.1.1.2, we have

$$\xi^{\dagger} \stackrel{\mathrm{def}}{=} \llbracket a \mapsto \llbracket b \mapsto \xi(a,b) \rrbracket \rrbracket.$$

• Map II. We define a map

$$\Psi_{B,C}$$
: Sets(A, Sets(B, C)),  $\rightarrow$  Sets(A  $\times$  B, C)

given by sending a function

$$\xi: A \longrightarrow \mathsf{Sets}(B, C),$$
  
 $a \mapsto (\xi_a: B \to C),$ 

to the function

$$\xi^{\dagger}: A \times B \to C$$

defined by

$$\xi^{\dagger}(a,b) \stackrel{\text{def}}{=} \operatorname{ev}_b(\operatorname{ev}_a(\xi))$$

$$\stackrel{\text{def}}{=} \operatorname{ev}_b(\xi_a)$$

$$\stackrel{\text{def}}{=} \xi_a(b)$$

for each  $(a, b) \in A \times B$ .

• Invertibility I. We claim that

$$\Psi_{A,B} \circ \Phi_{A,B} = \mathrm{id}_{\mathsf{Sets}(A \times B,C)}$$
.

Indeed, given a function  $\xi: A \times B \to C$ , we have

$$\begin{split} \left[ \Psi_{A,B} \circ \Phi_{A,B} \right] (\xi) &= \Psi_{A,B} \big( \Phi_{A,B} (\xi) \big) \\ &= \Psi_{A,B} \big( \Phi_{A,B} \big( \big[ \big[ (a,b) \mapsto \xi(a,b) \big] \big] \big) \big) \\ &= \Psi_{A,B} \big( \big[ a \mapsto \big[ b \mapsto \xi(a,b) \big] \big] \big] \big) \\ &= \Psi_{A,B} \big( \big[ a' \mapsto \big[ b' \mapsto \xi(a',b') \big] \big] \big] \big) \\ &= \big[ \big[ (a,b) \mapsto \operatorname{ev}_b \big( \operatorname{ev}_a \big( \big[ a' \mapsto \big[ b' \mapsto \xi(a',b') \big] \big] \big] \big) \big) \big] \\ &= \big[ \big[ (a,b) \mapsto \operatorname{ev}_b \big( \big[ \big[ b' \mapsto \xi(a,b') \big] \big] \big) \big] \big] \\ &= \big[ \big[ (a,b) \mapsto \xi(a,b) \big] \big] \\ &= \xi. \end{split}$$

• Invertibility II. We claim that

$$\Phi_{A,B} \circ \Psi_{A,B} = \mathrm{id}_{\mathsf{Sets}(A,\mathsf{Sets}(B,C))}$$
.

Indeed, given a function

$$\xi: A \longrightarrow \mathsf{Sets}(B, C),$$
  
 $a \mapsto (\xi_a: B \to C),$ 

we have

$$\left[\Phi_{A,B}\circ\Psi_{A,B}\right](\xi)\stackrel{\mathrm{def}}{=}\Phi_{A,B}\left(\Psi_{A,B}(\xi)\right)$$

$$\begin{split} &\overset{\mathrm{def}}{=} \Phi_{A,B}(\llbracket(a,b) \mapsto \xi_a(b)\rrbracket) \\ &\overset{\mathrm{def}}{=} \Phi_{A,B}(\llbracket(a',b') \mapsto \xi_{a'}(b')\rrbracket) \\ &\overset{\mathrm{def}}{=} \llbracket a \mapsto \llbracket b \mapsto \mathrm{ev}_{(a,b)}(\llbracket(a',b') \mapsto \xi_{a'}(b')\rrbracket)\rrbracket \rrbracket \\ &\overset{\mathrm{def}}{=} \llbracket a \mapsto \llbracket b \mapsto \xi_a(b)\rrbracket \rrbracket \end{bmatrix} \\ &\overset{\mathrm{def}}{=} \llbracket a \mapsto \xi_a \rrbracket \\ &\overset{\mathrm{def}}{=} \xi. \end{split}$$

• *Naturality for*  $\Phi$ , *Part I.* We need to show that, given a function  $g: B \to B'$ , the diagram

$$\begin{split} \mathsf{Sets}(A \times B', C) & \xrightarrow{\Phi_{B',C}} & \mathsf{Sets}(A, \mathsf{Sets}(B', C)), \\ & \mathsf{id}_A \times g^* \bigg| & & & & & & & & \\ & & \mathsf{id}_A \times g^* \bigg| & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & & \\ &$$

commutes. Indeed, given a function

$$\xi \colon A \times B' \to C$$

we have

$$\begin{split} \left[ \Phi_{B,C} \circ (\mathrm{id}_A \times g^*) \right] (\xi) &= \Phi_{B,C} ([\mathrm{id}_A \times g^*] (\xi)) \\ &= \Phi_{B,C} (\xi (-_1, g (-_2))) \\ &= [\xi (-_1, g (-_2))]^{\dagger} \\ &= \xi_{-_1}^{\dagger} (g (-_2)) \\ &= (g^*)_! (\xi^{\dagger}) \\ &= (g^*)_! (\Phi_{B',C} (\xi)) \\ &= [(g^*)_! \circ \Phi_{B',C}] (\xi). \end{split}$$

Alternatively, using the  $[a \mapsto f(a)]$  notation of Sets, Notation 3.1.1.1.2, we have

$$\begin{split} \left[\Phi_{B,C}\circ(\mathrm{id}_A\times g^*)\right](\xi) &= \Phi_{B,C}([\mathrm{id}_A\times g^*](\xi)) \\ &= \Phi_{B,C}([\mathrm{id}_A\times g^*](\left[\!\left[(a,b')\mapsto \xi(a,b')\right]\!\right])) \end{split}$$

$$= \Phi_{B,C}([(a,b) \mapsto \xi(a,g(b))])$$

$$= [a \mapsto [b \mapsto \xi(a,g(b))]]$$

$$= [a \mapsto g^*([b' \mapsto \xi(a,b')])]$$

$$= (g^*)_!([a \mapsto [b' \mapsto \xi(a,b')]])$$

$$= (g^*)_!(\Phi_{B',C}([(a,b') \mapsto \xi(a,b')]])$$

$$= (g^*)_!(\Phi_{B',C}(\xi))$$

$$= [(g^*)_! \circ \Phi_{B',C}](\xi).$$

• *Naturality for*  $\Phi$ , *Part II.* We need to show that, given a function  $h: C \to C'$ , the diagram

$$\mathsf{Sets}(A \times B, C) \xrightarrow{\Phi_{B,C}} \mathsf{Sets}(A, \mathsf{Sets}(B, C)),$$

$$\downarrow_{(h_!)_1} \qquad \qquad \downarrow_{(h_!)_1}$$

$$\mathsf{Sets}(A \times B, C') \xrightarrow{\Phi_{B,C'}} \mathsf{Sets}(A, \mathsf{Sets}(B, C'))$$

commutes. Indeed, given a function

$$\xi: A \times B \to C$$
,

we have

$$\begin{split} \big[ \Phi_{B,C} \circ h_! \big] (\xi) &= \Phi_{B,C} (h_!(\xi)) \\ &= \Phi_{B,C} (h_!( [\![ (a,b) \mapsto \xi(a,b) ]\!] )) \\ &= \Phi_{B,C} ( [\![ (a,b) \mapsto h(\xi(a,b)) ]\!] ) \\ &= [\![ a \mapsto [\![ b \mapsto h(\xi(a,b)) ]\!] ]\!] \\ &= [\![ a \mapsto h_!( [\![ b \mapsto \xi(a,b) ]\!] ]\!] ) \\ &= (h_!)_! ( [\![ a \mapsto [\![ b \mapsto \xi(a,b) ]\!] ]\!] ) \\ &= (h_!)_! ( \Phi_{B,C} ( [\![ (a,b) \mapsto \xi(a,b) ]\!] )) \\ &= (h_!)_! ( \Phi_{B,C} ( \xi) ) \\ &= [\![ (h_!)_! \circ \Phi_{B,C} ]\!] (\xi). \end{split}$$

• Naturality for  $\Psi$ . Since  $\Phi$  is natural in each argument and  $\Phi$  is a componentwise inverse to  $\Psi$  in each argument, it follows from

Categories, Item 2 of Proposition 11.9.7.1.2 that  $\Psi$  is also natural in each argument.

This finishes the proof.

Item 3: Adjointness II

This follows from the universal property of the product.

Item 4: Associativity

This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.1.4.1.1.

Item 5: Unitality

This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.1.5.1.1 and 5.1.6.1.1.

Item 6: Commutativity

This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.1.7.1.1.

Item 7: Distributivity Over Coproducts

This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.3.1.1.1 and 5.3.2.1.1.

Item 8: Annihilation With the Empty Set

This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.3.3.1.1 and 5.3.4.1.1.

Item 9: Distributivity Over Unions

See [Pro25c].

Item 10: Distributivity Over Intersections

See [Pro25d, Corollary 1].

Item 11: Distributivity Over Differences

See [Pro25a].

Item 12: Distributivity Over Symmetric Differences

See [Pro25b].

Item 13: Middle-Four Exchange With Respect to Intersections

See [Pro25d, Corollary 1].

Item 14: Symmetric Monoidality

This is a repetition of Monoidal Structures on the Category of Sets, Proposition 5.1.9.1.1, and is proved there.

#### Item 15: Symmetric Bimonoidality

This is a repetition of Monoidal Structures on the Category of Sets, Proposition 5.3.5.1.1, and is proved there.

#### REMARK 4.1.3.1.6 $\blacktriangleright$ The Cartesian Product of Sets as an $(\mathbb{E}_k, \mathbb{E}_\ell)$ -Tensor Product

As shown in Item 1 of Proposition 4.1.3.1.4, the Cartesian product of sets defines a functor

$$-1 \times -2$$
: Sets  $\times$  Sets  $\rightarrow$  Sets.

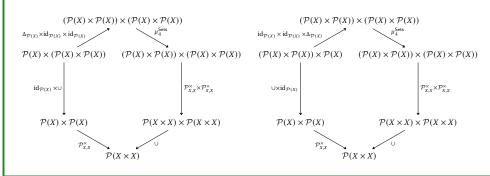
This functor is the  $(k, \ell) = (-1, -1)$  case of a family of functors

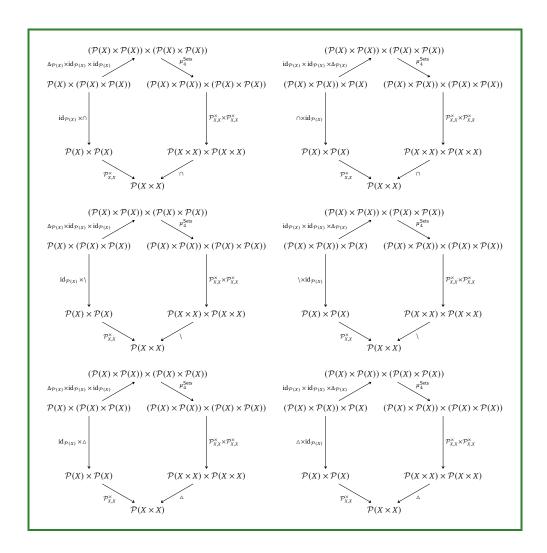
$$\otimes_{k,\ell} \colon \mathsf{Mon}_{\mathbb{E}_k}(\mathsf{Sets}) \times \mathsf{Mon}_{\mathbb{E}_\ell}(\mathsf{Sets}) \to \mathsf{Mon}_{\mathbb{E}_{k+\ell}}(\mathsf{Sets})$$

of tensor products of  $\mathbb{E}_k$ -monoid objects on Sets with  $\mathbb{E}_\ell$ -monoid objects on Sets; see ??.

#### REMARK 4.1.3.1.7 ► DIAGRAMS FOR ITEMS 9 TO 12 OF PROPOSITION 4.1.3.1.4

We may state the equalities in Items 9 to 12 of Proposition 4.1.3.1.4 as the commutativity of the following diagrams:





### 4.1.4 Pullbacks

Let A, B, and C be sets and let  $f: A \to C$  and  $g: B \to C$  be functions.

#### **DEFINITION 4.1.4.1.1** ► PULLBACKS OF SETS

The **pullback of** A **and** B **over** C **along** f **and**  $g^1$  is the pullback of A and B over C along f and g in Sets as in Limits and Colimits, ??.

<sup>&</sup>lt;sup>1</sup>Further Terminology: Also called the **fibre product of** A **and** B **over** C **along** f **and** g.

### CONSTRUCTION 4.1.4.1.2 ► CONSTRUCTION OF PULLBACKS OF SETS

Concretely, the pullback of A and B over C along f and g is the pair  $(A \times_C B, \{pr_1, pr_2\})$  consisting of:

1. *The Limit*. The set  $A \times_C B$  defined by

$$A \times_C B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid f(a) = g(b)\}.$$

2. The Cone. The maps<sup>1</sup>

$$\operatorname{pr}_1: A \times_C B \to A,$$
  
 $\operatorname{pr}_2: A \times_C B \to B$ 

defined by

$$\operatorname{pr}_{1}(a,b) \stackrel{\text{def}}{=} a,$$
  
 $\operatorname{pr}_{2}(a,b) \stackrel{\text{def}}{=} b$ 

for each  $(a, b) \in A \times_C B$ .

#### PROOF 4.1.4.1.3 ► PROOF OF CONSTRUCTION 4.1.4.1.2

We claim that  $A \times_C B$  is the categorical pullback of A and B over C with respect to (f,g) in Sets. First we need to check that the relevant pullback diagram commutes, i.e. that we have

$$f \circ \operatorname{pr}_1 = g \circ \operatorname{pr}_2, \qquad A \times_C B \xrightarrow{\operatorname{pr}_2} B$$

$$p_{r_1} \downarrow \qquad \qquad \downarrow g$$

$$A \xrightarrow{f} C.$$

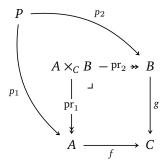
Indeed, given  $(a, b) \in A \times_C B$ , we have

$$[f \circ \operatorname{pr}_1](a, b) = f(\operatorname{pr}_1(a, b))$$
$$= f(a)$$
$$= g(b)$$

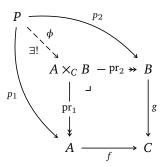
<sup>&</sup>lt;sup>1</sup>Further Notation: Also written  $pr_1^{A \times_C B}$  and  $pr_2^{A \times_C B}$ .

$$= g(\operatorname{pr}_{2}(a, b))$$
$$= [g \circ \operatorname{pr}_{2}](a, b),$$

where f(a) = g(b) since  $(a, b) \in A \times_C B$ . Next, we prove that  $A \times_C B$  satisfies the universal property of the pullback. Suppose we have a diagram of the form



in Sets. Then there exists a unique map  $\phi: P \to A \times_C B$  making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$
  
 $\operatorname{pr}_2 \circ \phi = p_2$ 

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each  $x \in P$ , where we note that  $(p_1(x), p_2(x)) \in A \times B$  indeed lies in  $A \times_C B$  by the condition

$$f \circ p_1 = g \circ p_2$$
,

which gives

$$f(p_1(x)) = g(p_2(x))$$

for each  $x \in P$ , so that  $(p_1(x), p_2(x)) \in A \times_C B$ .

#### REMARK 4.1.4.1.4 ▶ PULLBACKS OF SETS DEPEND ON THE MAPS

It is common practice to write  $A \times_C B$  for the pullback of A and B over C along f and g, omitting the maps f and g from the notation and instead leaving them implicit, to be understood from the context.

However, the set  $A \times_C B$  depends very much on the maps f and g, and sometimes it is necessary or useful to note this dependence explicitly. In such situations, we will write  $A \times_{f,C,g} B$  or  $A \times_C^{f,g} B$  for  $A \times_C B$ .

#### **EXAMPLE 4.1.4.1.5** ► **EXAMPLES OF PULLBACKS OF SETS**

Here are some examples of pullbacks of sets.

1. *Unions via Intersections*. Let *X* be a set. We have

$$A \cap B \cong A \times_{A \cup B} B, \qquad A \cap B \xrightarrow{J} B$$

$$A \cap B \cong A \times_{A \cup B} B, \qquad A \xrightarrow{\iota_A} A \cup B$$

for each  $A, B \in \mathcal{P}(X)$ .

#### PROOF 4.1.4.1.6 ► PROOF OF EXAMPLE 4.1.4.1.5

### Item 1: Unions via Intersections

Indeed, we have

$$A \times_{A \cup B} B \cong \{(x, y) \in A \times B \mid x = y\}$$
  
\approx A \cap B.

This finishes the proof.

#### PROPOSITION 4.1.4.1.7 ► PROPERTIES OF PULLBACKS OF SETS

Let A, B, C, and X be sets.

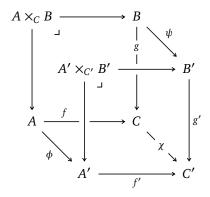
1. Functoriality. The assignment  $(A, B, C, f, g) \mapsto A \times_{f,C,g} B$  defines a functor

$$-1 \times_{-3} -1$$
: Fun( $\mathcal{P}$ , Sets)  $\rightarrow$  Sets,

where  $\mathcal{P}$  is the category that looks like this:



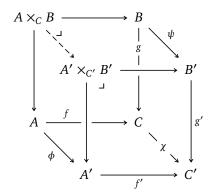
In particular, the action on morphisms of  $-1 \times_{-3} -1$  is given by sending a morphism



in  $\operatorname{Fun}(\mathcal{P},\operatorname{\mathsf{Sets}})$  to the map  $\xi\colon A\times_C B \xrightarrow{\exists !} A'\times_{C'} B'$  given by

$$\xi(a,b) \stackrel{\text{def}}{=} (\phi(a), \psi(b))$$

for each  $(a, b) \in A \times_C B$ , which is the unique map making the diagram



commute.

2. Adjointness I. We have adjunctions

$$(A \times_X - + \mathbf{Sets}_{/X}(A, -))$$
:  $Sets_{/X} \xrightarrow{A \times_X -} Sets_{/X}$ ,  $Sets_{/X}(A, -)$   
 $(- \times_X B + \mathbf{Sets}_{/X}(B, -))$ :  $Sets_{/X} \xrightarrow{A \times_X -} Sets_{/X}$ ,  $Sets_{/X}(B, -)$ 

witnessed by bijections

$$\mathsf{Sets}_{/X}(A \times_X B, C) \cong \mathsf{Sets}_{/X}(A, \mathbf{Sets}_{/X}(B, C)),$$
  
 $\mathsf{Sets}_{/X}(A \times_X B, C) \cong \mathsf{Sets}_{/X}(B, \mathbf{Sets}_{/X}(A, C)),$ 

natural in  $(A, \phi_A)$ ,  $(B, \phi_B)$ ,  $(C, \phi_C) \in \text{Obj}(\mathsf{Sets}_{/X})$ , where  $\mathsf{Sets}_{/X}(A, B)$  is the object of  $\mathsf{Sets}_{/X}$  consisting of (see Fibred Sets, ??):

• *The Set.* The set  $\mathbf{Sets}_{/X}(A, B)$  defined by

$$\mathbf{Sets}_{/X}(A,B) \stackrel{\text{def}}{=} \coprod_{x \in X} \mathsf{Sets}\Big(\phi_A^{-1}(x), \phi_Y^{-1}(x)\Big)$$

• *The Map to X*. The map

$$\phi_{\mathsf{Sets}_{/X}(A,B)} : \mathsf{Sets}_{/X}(A,B) \to X$$

defined by

$$\phi_{\mathsf{Sets}_{/X}(A,B)}(x,f) \stackrel{\mathrm{def}}{=} x$$

for each  $(x, f) \in \mathbf{Sets}_{/X}(A, B)$ .

3. Adjointness II. We have an adjunction

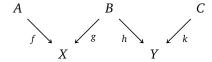
$$\left(\Delta_{\mathsf{Sets}_{/X}} \dashv -_1 \times -_2\right)$$
:  $\mathsf{Sets}_{/X} \underbrace{\perp}_{-_1 \times -_2} \mathsf{Sets}_{/X} \times \mathsf{Sets}_{/X}$ ,

witnessed by a bijection

$$\operatorname{\mathsf{Hom}}_{\mathsf{Sets}_{/X}\times\mathsf{Sets}_{/X}}((A,A),(B,C))\cong \operatorname{\mathsf{Sets}}_{/X}(A,B\times_XC),$$

natural in  $A \in \text{Obj}(\mathsf{Sets}_{/X})$  and in  $(B, C) \in \text{Obj}(\mathsf{Sets}_{/X} \times \mathsf{Sets}_{/X})$ .

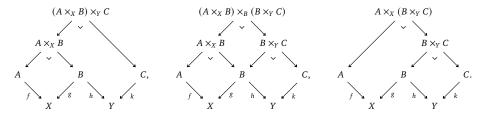
4. Associativity. Given a diagram



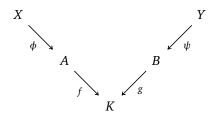
in Sets, we have isomorphisms of sets

$$(A \times_X B) \times_Y C \cong (A \times_X B) \times_B (B \times_Y C) \cong A \times_X (B \times_Y C),$$

where these pullbacks are built as in the diagrams



5. Interaction With Composition. Given a diagram



in Sets, we have isomorphisms of sets

$$\begin{split} X \times_K^{f \circ \phi, g \circ \psi} Y &\cong \left( X \times_A^{\phi, q_1} \left( A \times_K^{f, g} B \right) \right) \times_{A \times_K^{f, g} B}^{p_2, p_1} \left( \left( A \times_K^{f, g} B \right) \times_B^{q_2, \psi} Y \right) \\ &\cong X \times_A^{\phi, p} \left( \left( A \times_K^{f, g} B \right) \times_B^{q_2, \psi} Y \right) \\ &\cong \left( X \times_A^{\phi, q_1} \left( A \times_K^{f, g} B \right) \right) \times_B^{q, \psi} Y \end{split}$$

where

$$q_{1} = \operatorname{pr}_{1}^{A \times_{K}^{f,g} B}, \qquad q_{2} = \operatorname{pr}_{2}^{A \times_{K}^{f,g} B},$$

$$p_{1} = \operatorname{pr}_{1}^{\left(A \times_{K}^{f,g} B\right) \times_{Y}^{q_{2},\psi}}, \qquad X \times_{A \times_{K}^{f,g} B}^{\phi,q_{1}} \left(A \times_{K}^{f,g} B\right)},$$

$$p_{2} = \operatorname{pr}_{2}^{X \times_{K}^{\phi,q_{1}}} \left(A \times_{K}^{f,g} B\right),$$

$$p_{3} = \operatorname{pr}_{2}^{X \times_{K}^{\phi,q_{1}}} \left(A \times_{K}^{f,g} B\right),$$

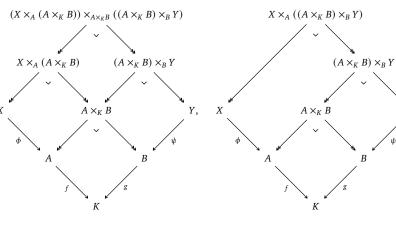
$$p_{4} = \operatorname{pr}_{2}^{X \times_{K}^{\phi,q_{1}}} \left(A \times_{K}^{f,g} B\right),$$

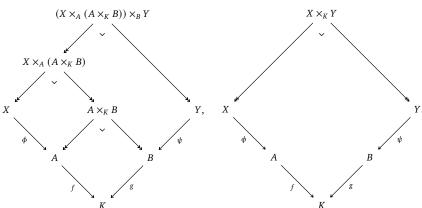
$$p_{5} = \operatorname{pr}_{2}^{X \times_{K}^{\phi,q_{1}}} \left(A \times_{K}^{f,g} B\right),$$

$$p_{7} = \operatorname{pr}_{2}^{X \times_{K}^{\phi,q_{1}}} \left(A \times_{K}^{f,g} B\right),$$

$$p_{8} = \operatorname{pr}_{2}^{X \times_{K}^{\phi,q_{1}}} \left(A \times_{K}^{f,g} B\right),$$

and where these pullbacks are built as in the following diagrams:





6. Unitality. We have isomorphisms of sets

natural in  $(A, f) \in \text{Obj}(\mathsf{Sets}_{/X})$ .

7. Commutativity. We have an isomorphism of sets

natural in (A, f),  $(B, g) \in Obj(Sets_{/X})$ .

8. *Distributivity Over Coproducts*. Let A, B, and C be sets and let  $\phi_A \colon A \to X$ ,  $\phi_B \colon B \to X$ , and  $\phi_C \colon C \to X$  be morphisms of sets. We have isomorphisms of sets

$$\delta_{\ell}^{\mathsf{Sets}_{/X}} : A \times_{X} (B \coprod C) \xrightarrow{\sim} (A \times_{X} B) \coprod (A \times_{X} C),$$
  
$$\delta_{r}^{\mathsf{Sets}_{/X}} : (A \coprod B) \times_{X} C \xrightarrow{\sim} (A \times_{X} C) \coprod (B \times_{X} C),$$

as in the diagrams

natural in  $A, B, C \in \text{Obj}(\mathsf{Sets}_{/X})$ .

9. Annihilation With the Empty Set. We have isomorphisms of sets

$$\emptyset \longrightarrow \emptyset \qquad \qquad \emptyset \longrightarrow A$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \zeta_{\ell}^{\mathsf{Sets}/X} : A \times_{X} \emptyset \xrightarrow{\sim} \emptyset, \qquad \qquad \downarrow \qquad \qquad \downarrow f$$

$$A \xrightarrow{f} X, \qquad \qquad \emptyset \longrightarrow X,$$

natural in  $(A, f) \in \text{Obj}(\mathsf{Sets}_{/X})$ .

10. Interaction With Products. We have an isomorphism of sets

$$A \times_{\text{pt}} B \cong A \times B,$$

$$A \times_{\text{pt}} B \cong A \times B,$$

$$A \xrightarrow{!_{A}} pt.$$

11. *Symmetric Monoidality*. The 8-tuple (Sets<sub>/X</sub>,  $\times_X$ , X, **Sets**<sub>/X</sub>,  $\alpha^{\text{Sets}_{/X}}$ ,  $\alpha^{\text{Sets}_{/X}}$ ,  $\sigma^{\text{Sets}_{/X}}$ ) is a symmetric closed monoidal category.

#### PROOF 4.1.4.1.8 ► PROOF OF PROPOSITION 4.1.4.1.7

#### Item 1: Functoriality

This is a special case of functoriality of co/limits, Limits and Colimits, ?? of ??, with the explicit expression for  $\xi$  following from the commutativity of the cube pullback diagram.

#### Item 2: Adjointness I

This is a repetition of Fibred Sets, ?? of ??, and is proved there.

#### Item 3: Adjointness II

This follows from the universal property of the product (pullbacks are products in  $\mathsf{Sets}_{/X}$ ).

### Item 4: Associativity

We have

$$(A \times_X B) \times_Y C \cong \{((a,b),c) \in (A \times_X B) \times C \mid h(b) = k(c)\}$$

$$\cong \{((a,b),c) \in (A \times B) \times C \mid f(a) = g(b) \text{ and } h(b) = k(c)\}$$

$$\cong \{(a,(b,c)) \in A \times (B \times C) \mid f(a) = g(b) \text{ and } h(b) = k(c)\}$$

$$\cong \{(a,(b,c)) \in A \times (B \times_Y C) \mid f(a) = g(b)\}$$

$$\cong A \times_X (B \times_Y C)$$

and

$$(A \times_X B) \times_B (B \times_Y C) \cong \{((a,b),(b',c)) \in (A \times_X B) \times (B \times_Y C) \mid b = b'\}$$

$$\cong \left\{ ((a,b),(b',c)) \in (A \times B) \times (B \times C) \mid f(a) = g(b), b = b', \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,(b,(b',c))) \in A \times (B \times (B \times C)) \mid f(a) = g(b), b = b', \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,((b,b'),c)) \in A \times ((B \times B) \times C) \mid f(a) = g(b), b = b', \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,((b,b'),c)) \in A \times ((B \times_B B) \times C) \mid f(a) = g(b) \text{ and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,(b,c)) \in A \times (B \times C) \mid f(a) = g(b) \text{ and } h(b) = k(c) \right\}$$

$$\cong A \times_X (B \times_Y C),$$

where we have used Item 6 for the isomorphism  $B \times_B B \cong B$ .

### Item 5: Interaction With Composition

By Item 4, it suffices to construct only the isomorphism

$$X \times_K^{f \circ \phi, g \circ \psi} Y \cong \left( X \times_A^{\phi, q_1} \left( A \times_K^{f, g} B \right) \right) \times_{A \times_{\nu}^{f, g} B}^{p_2, p_1} \left( \left( A \times_K^{f, g} B \right) \times_B^{q_2, \psi} Y \right).$$

We have

$$\left(X \times_A^{\phi,q_1} \left(A \times_K^{f,g} B\right)\right) \stackrel{\text{def}}{=} \left\{ (x,(a,b)) \in X \times \left(A \times_K^{f,g} B\right) \mid \phi(x) = q_1(a,b) \right\}$$

$$\stackrel{\text{def}}{=} \left\{ (x,(a,b)) \in X \times \left(A \times_K^{f,g} B\right) \mid \phi(x) = a \right\}$$

$$\cong \left\{ (x,(a,b)) \in X \times (A \times B) \mid \phi(x) = a \text{ and } f(a) = g(b) \right\},$$

$$\left( \left(A \times_K^{f,g} B\right) \times_B^{q_2,\psi} Y \right) \stackrel{\text{def}}{=} \left\{ ((a,b),y) \in \left(A \times_K^{f,g} B\right) \times Y \mid q_2(a,b) = \psi(y) \right\}$$

$$\stackrel{\text{def}}{=} \left\{ ((a,b),y) \in \left(A \times_K^{f,g} B\right) \times Y \mid b = \psi(y) \text{ and } f(a) = g(b) \right\},$$

$$\cong \left\{ ((a,b),y) \in (A \times B) \times Y \mid b = \psi(y) \text{ and } f(a) = g(b) \right\},$$

so writing

$$S = \left( X \times_A^{\phi, q_1} \left( A \times_K^{f, g} B \right) \right)$$
  
$$S' = \left( \left( A \times_K^{f, g} B \right) \times_B^{q_2, \psi} Y \right),$$

we have

$$\begin{split} S \times_{A \times_K^{f,g} B}^{p_2,p_1} S' &\stackrel{\text{def}}{=} \{ ((x,(a,b)), ((a',b'),y)) \in S \times S' \mid p_1(x,(a,b)) = p_2((a',b'),y) \} \\ &\stackrel{\text{def}}{=} \{ ((x,(a,b)), ((a',b'),y)) \in S \times S' \mid (a,b) = (a',b') \} \\ &\cong \{ ((x,a,b,y)) \in X \times A \times B \times Y \mid \phi(x) = a, \psi(y) = b, \text{ and } f(a) = g(b) \} \\ &\cong \{ ((x,a,b,y)) \in X \times A \times B \times Y \mid f(\phi(x)) = g(\psi(y)) \} \\ &\stackrel{\text{def}}{=} X \times_K Y. \end{split}$$

This finishes the proof.

#### Item 6: Unitality

We have

$$X \times_X A \cong \{(x, a) \in X \times A \mid f(a) = x\},\$$
  
$$A \times_X X \cong \{(a, x) \in X \times A \mid f(a) = x\},\$$

which are isomorphic to *A* via the maps  $(x, a) \mapsto a$  and  $(a, x) \mapsto a$ . The proof of the naturality of  $\lambda^{\text{Sets}/x}$  and  $\rho^{\text{Sets}/x}$  is omitted.

### Item 7: Commutativity

We have

$$A \times_C B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid f(a) = g(b)\}$$

$$= \{(a, b) \in A \times B \mid g(b) = f(a)\}$$

$$\cong \{(b, a) \in B \times A \mid g(b) = f(a)\}$$

$$\stackrel{\text{def}}{=} B \times_C A.$$

The proof of the naturality of  $\sigma^{\text{Sets}_{/X}}$  is omitted.

## Item 8: Distributivity Over Coproducts

We have

$$A \times_{X} (B \coprod C) \stackrel{\text{def}}{=} \left\{ (a, z) \in A \times (B \coprod C) \middle| \phi_{A}(a) = \phi_{B \coprod C}(z) \right\}$$

$$= \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (0, b) \text{ and } \phi_{A}(a) = \phi_{B \coprod C}(z) \right\}$$

$$\cup \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (1, c) \text{ and } \phi_{A}(a) = \phi_{B \coprod C}(z) \right\}$$

$$= \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (0, b) \text{ and } \phi_{A}(a) = \phi_{B}(b) \right\}$$

$$\cup \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (1, c) \text{ and } \phi_{A}(a) = \phi_{C}(c) \right\}$$

$$\cong \left\{ (a, b) \in A \times B \middle| \phi_{A}(a) = \phi_{B}(b) \right\}$$

$$\cup \left\{ (a, c) \in A \times C \middle| \phi_{A}(a) = \phi_{C}(c) \right\}$$

$$\stackrel{\text{def}}{=} (A \times_{X} B) \cup (A \times_{X} C)$$

$$\cong (A \times_{X} B) \coprod (A \times_{X} C),$$

with the construction of the isomorphism

$$\delta_r^{\mathsf{Sets}_{/X}} : (A \coprod B) \times_X C \xrightarrow{\sim} (A \times_X C) \coprod (B \times_X C)$$

being similar. The proof of the naturality of  $\delta_\ell^{\mathsf{Sets}_{/X}}$  and  $\delta_r^{\mathsf{Sets}_{/X}}$  is omitted.

# Item 9: Annihilation With the Empty Set

We have

$$A \times_X \emptyset \stackrel{\text{def}}{=} \{(a, b) \in A \times \emptyset \mid f(a) = g(b)\}$$
$$= \{k \in \emptyset \mid f(a) = g(b)\}$$
$$= \emptyset,$$

and similarly for  $\emptyset \times_X A$ , where we have used Item 8 of Proposition 4.1.3.1.4. The proof of the naturality of  $\zeta_\ell^{\mathsf{Sets}_{/X}}$  and  $\zeta_r^{\mathsf{Sets}_{/X}}$  is omitted.

### Item 10: Interaction With Products

We have

$$A \times_{\text{pt}} B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid !_A(a) = !_B(b)\}$$

$$\stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid \star = \star\}$$

$$= \{(a, b) \in A \times B\}$$

$$= A \times B.$$

### Item 11: Symmetric Monoidality

Omitted.

### 4.1.5 Equalisers

Let *A* and *B* be sets and let  $f, g: A \Rightarrow B$  be functions.

#### **DEFINITION 4.1.5.1.1** ► EQUALISERS OF SETS

The **equaliser of** f **and** g is the equaliser of f and g in Sets as in Limits and Colimits, ??.

### CONSTRUCTION 4.1.5.1.2 ► CONSTRUCTION OF EQUALISERS OF SETS

Concretely, the equaliser of f and g is the pair (Eq(f,g),eq(f,g)) consisting of:

1. *The Limit*. The set Eq(f,g) defined by

$$\operatorname{Eq}(f,g) \stackrel{\text{def}}{=} \{ a \in A \mid f(a) = g(a) \}.$$

2. The Cone. The inclusion map

$$eq(f,g): Eq(f,g) \hookrightarrow A.$$

#### PROOF 4.1.5.1.3 ► PROOF OF CONSTRUCTION 4.1.5.1.2

We claim that Eq(f,g) is the categorical equaliser of f and g in Sets. First we need to check that the relevant equaliser diagram commutes, i.e. that we have

$$f \circ eq(f,g) = g \circ eq(f,g),$$

which indeed holds by the definition of the set Eq(f,g). Next, we prove that Eq(f,g) satisfies the universal property of the equaliser. Suppose we have a diagram of the form

$$\operatorname{Eq}(f,g) \xrightarrow{\operatorname{eq}(f,g)} A \xrightarrow{f \atop g} B$$

in Sets. Then there exists a unique map  $\phi \colon E \to \operatorname{Eq}(f,g)$  making the diagram

$$Eq(f,g) \xrightarrow{eq(f,g)} A \xrightarrow{f} E$$

$$\downarrow f$$

$$\downarrow g$$

commute, being uniquely determined by the condition

$$eq(f,g) \circ \phi = e$$

via

$$\phi(x) = e(x)$$

for each  $x \in E$ , where we note that  $e(x) \in A$  indeed lies in Eq(f,g) by the condition

$$f \circ e = g \circ e$$
,

which gives

$$f(e(x)) = g(e(x))$$

for each  $x \in E$ , so that  $e(x) \in \text{Eq}(f, g)$ .

#### PROPOSITION 4.1.5.1.4 ▶ PROPERTIES OF EQUALISERS OF SETS

Let A, B, and C be sets.

1. Associativity. We have isomorphisms of sets<sup>1</sup>

$$\underbrace{\operatorname{Eq}(f \circ \operatorname{eq}(g,h), g \circ \operatorname{eq}(g,h))}_{=\operatorname{Eq}(f \circ \operatorname{eq}(g,h), h \circ \operatorname{eq}(g,h))} \cong \operatorname{Eq}(f,g,h) \cong \underbrace{\operatorname{Eq}(f \circ \operatorname{eq}(f,g), h \circ \operatorname{eq}(f,g))}_{=\operatorname{Eq}(g \circ \operatorname{eq}(f,g), h \circ \operatorname{eq}(f,g))},$$

where Eq(f, g, h) is the limit of the diagram

$$A \xrightarrow{\frac{f}{-g}} B$$

in Sets, being explicitly given by

$$Eq(f,g,h) \cong \{a \in A \mid f(a) = g(a) = h(a)\}.$$

2. Unitality. We have an isomorphism of sets

$$\operatorname{Eq}(f, f) \cong A$$
.

3. Commutativity. We have an isomorphism of sets

$$\operatorname{Eq}(f,g) \cong \operatorname{Eq}(g,f).$$

4. *Interaction With Composition*. Let

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C$$

be functions. We have an inclusion of sets

$$\operatorname{Eq}(h \circ f \circ \operatorname{eq}(f, g), k \circ g \circ \operatorname{eq}(f, g)) \subset \operatorname{Eq}(h \circ f, k \circ g),$$

where Eq( $h \circ f \circ \text{eq}(f, g), k \circ g \circ \text{eq}(f, g)$ ) is the equaliser of the composition

$$\operatorname{Eq}(f,g) \stackrel{\operatorname{eq}(f,g)}{\hookrightarrow} A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C.$$

<sup>&</sup>lt;sup>1</sup>That is, the following three ways of forming "the" equaliser of (f, g, h) agree:

(a) Take the equaliser of (f, g, h), i.e. the limit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets.

(b) First take the equaliser of f and g, forming a diagram

$$\operatorname{Eq}(f,g) \stackrel{\operatorname{eq}(f,g)}{\hookrightarrow} A \stackrel{f}{\underset{g}{\Longrightarrow}} B$$

and then take the equaliser of the composition

$$\operatorname{Eq}(f,g) \stackrel{\operatorname{eq}(f,g)}{\hookrightarrow} A \stackrel{f}{\underset{h}{\Longrightarrow}} B,$$

obtaining a subset

$$Eq(f \circ eq(f,g), h \circ eq(f,g)) = Eq(g \circ eq(f,g), h \circ eq(f,g))$$

of Eq(f, g).

(c) First take the equaliser of g and h, forming a diagram

$$\operatorname{Eq}(g,h) \stackrel{\operatorname{eq}(g,h)}{\hookrightarrow} A \stackrel{g}{\underset{h}{\Longrightarrow}} B$$

and then take the equaliser of the composition

$$\mathrm{Eq}(g,h) \overset{\mathrm{eq}(g,h)}{\hookrightarrow} A \overset{f}{\underset{g}{\Longrightarrow}} B,$$

obtaining a subset

$$Eq(f \circ eq(g, h), g \circ eq(g, h)) = Eq(f \circ eq(g, h), h \circ eq(g, h))$$

of Eq(g, h).

#### PROOF 4.1.5.1.5 ► PROOF OF PROPOSITION 4.1.5.1.4

### Item 1: Associativity

We first prove that Eq(f, g, h) is indeed given by

$$\operatorname{Eq}(f,g,h) \cong \{a \in A \mid f(a) = g(a) = h(a)\}.$$

Indeed, suppose we have a diagram of the form

$$\operatorname{Eq}(f,g,h) \xrightarrow{\operatorname{eq}(f,g,h)} A \xrightarrow{f \atop -g \xrightarrow{h}} B$$

in Sets. Then there exists a unique map  $\phi \colon E \to \mathrm{Eq}(f,g,h)$ , uniquely determined by the condition

$$eq(f,g) \circ \phi = e$$

being necessarily given by

$$\phi(x) = e(x)$$

for each  $x \in E$ , where we note that  $e(x) \in A$  indeed lies in Eq(f, g, h) by the condition

$$f \circ e = g \circ e = h \circ e$$
,

which gives

$$f(e(x)) = g(e(x)) = h(e(x))$$

for each  $x \in E$ , so that  $e(x) \in \text{Eq}(f, g, h)$ .

We now check the equalities

$$\operatorname{Eq}(f \circ \operatorname{eq}(g, h), g \circ \operatorname{eq}(g, h)) \cong \operatorname{Eq}(f, g, h) \cong \operatorname{Eq}(f \circ \operatorname{eq}(f, g), h \circ \operatorname{eq}(f, g)).$$

Indeed, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(g,h), g \circ \operatorname{eq}(g,h)) &\cong \{x \in \operatorname{Eq}(g,h) \mid [f \circ \operatorname{eq}(g,h)](a) = [g \circ \operatorname{eq}(g,h)](a)\} \\ &\cong \{x \in \operatorname{Eq}(g,h) \mid f(a) = g(a)\} \\ &\cong \{x \in A \mid f(a) = g(a) \text{ and } g(a) = h(a)\} \\ &\cong \{x \in A \mid f(a) = g(a) = h(a)\} \\ &\cong \operatorname{Eq}(f,g,h). \end{split}$$

Similarly, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(f,g), h \circ \operatorname{eq}(f,g)) &\cong \{x \in \operatorname{Eq}(f,g) \mid [f \circ \operatorname{eq}(f,g)](a) = [h \circ \operatorname{eq}(f,g)](a)\} \\ &\cong \{x \in \operatorname{Eq}(f,g) \mid f(a) = h(a)\} \\ &\cong \{x \in A \mid f(a) = h(a) \text{ and } f(a) = g(a)\} \\ &\cong \{x \in A \mid f(a) = g(a) = h(a)\} \end{split}$$

$$\cong$$
 Eq $(f, g, h)$ .

### Item 2: Unitality

Indeed, we have

$$\operatorname{Eq}(f, f) \stackrel{\text{def}}{=} \{ a \in A \mid f(a) = f(a) \}$$
$$= A.$$

### Item 3: Commutativity

Indeed, we have

$$\operatorname{Eq}(f,g) \stackrel{\text{def}}{=} \{ a \in A \mid f(a) = g(a) \}$$
$$= \{ a \in A \mid g(a) = f(a) \}$$
$$\stackrel{\text{def}}{=} \operatorname{Eq}(g,f).$$

### Item 4: Interaction With Composition

Indeed, we have

$$\begin{split} \operatorname{Eq}(h \circ f \circ \operatorname{eq}(f,g), k \circ g \circ \operatorname{eq}(f,g)) & \cong \{a \in \operatorname{Eq}(f,g) \mid h(f(a)) = k(g(a))\} \\ & \cong \{a \in A \mid f(a) = g(a) \text{ and } h(f(a)) = k(g(a))\}. \end{split}$$

and

$$Eq(h \circ f, k \circ g) \cong \{a \in A \mid h(f(a)) = k(g(a))\},\$$

and thus there's an inclusion from Eq $(h \circ f \circ eq(f,g), k \circ g \circ eq(f,g))$  to Eq $(h \circ f, k \circ g)$ .

#### 4.1.6 Inverse Limits

Let  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ :  $(I, \preceq)$   $\rightarrow$  Sets be an inverse system of sets.

#### **DEFINITION 4.1.6.1.1** ► INVERSE LIMITS OF SETS

The **inverse limit of**  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$  is the inverse limit of  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$  in Sets as in Limits and Colimits, ??.

#### CONSTRUCTION 4.1.6.1.2 ► CONSTRUCTION OF INVERSE LIMITS OF SETS

Concretely, the inverse limit of  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$  is the pair  $(\lim_{\alpha\in I}(X_{\alpha}), \{\operatorname{pr}_{\alpha}\}_{\alpha\in I})$  consisting of:

1. *The Limit*. The set  $\lim_{\alpha \in I} (X_{\alpha})$  defined by

$$\lim_{\substack{\longleftarrow \\ \alpha \in I}} (X_{\alpha}) \stackrel{\text{def}}{=} \left\{ (x_{\alpha})_{\alpha \in I} \in \prod_{\alpha \in I} X_{\alpha} \mid \text{for each } \alpha, \beta \in I, \text{ if } \alpha \preceq \beta, \right\}.$$

2. The Cone. The collection

$$\left\{ \operatorname{pr}_{\gamma} \colon \lim_{\substack{\longleftarrow \\ \alpha \in I}} (X_{\alpha}) \to X_{\gamma} \right\}_{\gamma \in I}$$

of maps of sets defined as the restriction of the maps

$$\left\{ \operatorname{pr}_{\gamma} : \prod_{\alpha \in I} X_{\alpha} \to X_{\gamma} \right\}_{\gamma \in I}$$

of Item 2 of Construction 4.1.2.1.2 to  $\lim_{\alpha \in I} (X_{\alpha})$  and hence given by

$$\operatorname{pr}_{\gamma}((x_{\alpha})_{\alpha\in I})\stackrel{\operatorname{def}}{=} x_{\gamma}$$

for each  $\gamma \in I$  and each  $(x_{\alpha})_{\alpha \in I} \in \lim_{\stackrel{\longleftarrow}{\alpha \in I}} (X_{\alpha})$ .

#### PROOF 4.1.6.1.3 ► PROOF OF CONSTRUCTION 4.1.6.1.2

We claim that  $\lim_{\leftarrow \alpha \in I} (X_{\alpha})$  is the limit of the inverse system of sets  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta \in I}$ . First we need to check that the limit diagram defined by

it commutes, i.e. that we have

$$f_{\alpha\beta} \circ \operatorname{pr}_{\alpha} = \operatorname{pr}_{\beta}, \qquad \lim_{\alpha \in I} (X_{\alpha})$$

$$X_{\alpha} \xrightarrow{f_{\alpha\beta}} X_{\beta}$$

for each  $\alpha, \beta \in I$  with  $\alpha \leq \beta$ . Indeed, given  $(x_{\gamma})_{\gamma \in I} \in \lim_{\leftarrow \gamma \in I} (X_{\gamma})$ , we have

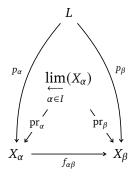
$$[f_{\alpha\beta} \circ \operatorname{pr}_{\alpha}] ((x_{\gamma})_{\gamma \in I}) \stackrel{\text{def}}{=} f_{\alpha\beta} (\operatorname{pr}_{\alpha} ((x_{\gamma})_{\gamma \in I}))$$

$$\stackrel{\text{def}}{=} f_{\alpha\beta} (x_{\alpha})$$

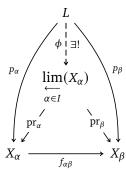
$$= x_{\beta}$$

$$\stackrel{\text{def}}{=} \operatorname{pr}_{\beta} ((x_{\gamma})_{\gamma \in I}),$$

where the third equality comes from the definition of  $\lim_{\leftarrow \alpha \in I} (X_{\alpha})$ . Next, we prove that  $\lim_{\leftarrow \alpha \in I} (X_{\alpha})$  satisfies the universal property of an inverse limit. Suppose that we have, for each  $\alpha, \beta \in I$  with  $\alpha \preceq \beta$ , a diagram of the form



in Sets. Then there indeed exists a unique map  $\phi: L \xrightarrow{\exists !} \varprojlim_{\alpha \in I} (X_{\alpha})$  making the diagram



commute, being uniquely determined by the family of conditions

$$\left\{p_{\alpha} = \operatorname{pr}_{\alpha} \circ \phi\right\}_{\alpha \in I}$$

via

$$\phi(\ell) = (p_{\alpha}(\ell))_{\alpha \in I}$$

for each  $\ell \in L$ , where we note that  $(p_{\alpha}(\ell))_{\alpha \in I} \in \prod_{\alpha \in I} X_{\alpha}$  indeed lies in  $\lim_{\leftarrow \alpha \in I} (X_{\alpha})$ , as we have

$$f_{\alpha\beta}(p_{\alpha}(\ell)) \stackrel{\text{def}}{=} [f_{\alpha\beta} \circ p_{\alpha}](\ell)$$

$$\stackrel{\text{def}}{=} p_{\beta}(\ell)$$

for each  $\beta \in I$  with  $\alpha \leq \beta$  by the commutativity of the diagram for  $(L, \{p_{\alpha}\}_{\alpha \in I})$ .

#### EXAMPLE 4.1.6.1.4 ► EXAMPLES OF INVERSE LIMITS OF SETS

Here are some examples of inverse limits of sets.

1. *The p-Adic Integers*. The ring of *p*-adic integers  $\mathbb{Z}_p$  of  $\ref{p-Adic}$  is the inverse limit

$$\mathbb{Z}_p \cong \lim_{\substack{\longleftarrow \\ n \in \mathbb{N}}} (\mathbb{Z}_{/p^n});$$

see ??.

2. Rings of Formal Power Series. The ring R[t] of formal power series in

a variable *t* is the inverse limit

$$R[[t]] \cong \lim_{\substack{t \in \mathbb{N} \\ n \in \mathbb{N}}} (R[t]/t^n R[t]);$$

see ??.

3. *Profinite Groups*. Profinite groups are inverse limits of finite groups; see ??.

# 4.2 Colimits of Sets

#### 4.2.1 The Initial Set

#### **DEFINITION 4.2.1.1.1** ► THE INITIAL SET

The **initial set** is the initial object of Sets as in Limits and Colimits, ??.

#### CONSTRUCTION 4.2.1.1.2 ► CONSTRUCTION OF THE INITIAL SET

Concretely, the initial set is the pair  $(\emptyset, \{\iota_A\}_{A \in \text{Obj}(\mathsf{Sets})})$  consisting of:

- 1. *The Colimit.* The empty set Ø of Definition 4.3.1.1.1.
- 2. The Cocone. The collection of maps

$$\{\iota_A \colon \emptyset \to A\}_{A \in \mathrm{Obj}(\mathsf{Sets})}$$

given by the inclusion maps from  $\emptyset$  to A.

#### PROOF 4.2.1.1.3 ► PROOF OF CONSTRUCTION 4.2.1.1.2

We claim that  $\emptyset$  is the initial object of Sets. Indeed, suppose we have a diagram of the form

$$\emptyset$$
 A

in Sets. Then there exists a unique map  $\phi: \emptyset \to A$  making the diagram

$$\emptyset - \frac{\phi}{\exists !} \rightarrow A$$

commute, namely the inclusion map  $\iota_A$ .

### 4.2.2 Coproducts of Families of Sets

Let  $\{A_i\}_{i\in I}$  be a family of sets.

### **DEFINITION 4.2.2.1.1** ► THE COPRODUCT OF A FAMILY OF SETS

The **coproduct of**  $\{A_i\}_{i\in I}^1$  is the coproduct of  $\{A_i\}_{i\in I}$  in Sets as in Limits and Colimits, ??.

#### CONSTRUCTION 4.2.2.1.2 ► CONSTRUCTION OF THE COPRODUCT OF A FAMILY OF SETS

Concretely, the disjoint union of  $\{A_i\}_{i\in I}$  is the pair  $\left(\coprod_{i\in I} A_i, \left\{\inf_i\right\}_{i\in I}\right)$  consisting of:

1. *The Colimit.* The set  $\coprod_{i \in I} A_i$  defined by

$$\coprod_{i \in I} A_i \stackrel{\text{def}}{=} \left\{ (i, x) \in I \times \left( \bigcup_{i \in I} A_i \right) \mid x \in A_i \right\}.$$

2. The Cocone. The collection

$$\left\{ \operatorname{inj}_i \colon A_i \to \coprod_{i \in I} A_i \right\}_{i \in I}$$

of maps given by

$$\operatorname{inj}_{i}(x) \stackrel{\text{def}}{=} (i, x)$$

for each  $x \in A_i$  and each  $i \in I$ .

<sup>&</sup>lt;sup>1</sup>Further Terminology: Also called the **disjoint union of the family**  $\{A_i\}_{i\in I}$ .

#### PROOF 4.2.2.1.3 ► PROOF OF CONSTRUCTION 4.2.2.1.2

We claim that  $\coprod_{i \in I} A_i$  is the categorical coproduct of  $\{A_i\}_{i \in I}$  in Sets. Indeed, suppose we have, for each  $i \in I$ , a diagram of the form

$$A_i \xrightarrow{\iota_i} A$$

$$A_i \xrightarrow{\text{inj}_i} \coprod_{i \in I} A$$

in Sets. Then there exists a unique map  $\phi: \coprod_{i \in I} A_i \to C$  making the diagram

$$A_{i} \xrightarrow{\lim_{i} \int_{i}^{i} \exists !} A_{i}$$

commute, being uniquely determined by the condition  $\phi \circ \text{inj}_i = \iota_i$  for each  $i \in I$  via

$$\phi((i,x)) = \iota_i(x)$$

for each  $(i, x) \in \coprod_{i \in I} A_i$ .

#### PROPOSITION 4.2.2.1.4 ► PROPERTIES OF COPRODUCTS OF FAMILIES OF SETS

Let  $\{A_i\}_{i\in I}$  be a family of sets.

1. Functoriality. The assignment  $\{A_i\}_{i\in I}\mapsto \coprod_{i\in I}A_i$  defines a functor

$$\coprod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

• *Action on Objects.* For each  $(A_i)_{i \in I} \in \mathsf{Obj}(\mathsf{Fun}(I_{\mathsf{disc}},\mathsf{Sets}))$ , we have

$$\left[ \bigsqcup_{i \in I} \right] ((A_i)_{i \in I}) \stackrel{\text{def}}{=} \bigsqcup_{i \in I} A_i$$

• Action on Morphisms. For each  $(A_i)_{i \in I}$ ,  $(B_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$ , the action on Hom-sets

$$\left(\bigsqcup_{i\in I}\right)_{(A_i)_{i\in I},(B_i)_{i\in I}}: \operatorname{Nat}((A_i)_{i\in I},(B_i)_{i\in I}) \to \operatorname{Sets}\left(\bigsqcup_{i\in I}A_i,\bigsqcup_{i\in I}B_i\right)$$

of  $\coprod_{i\in I}$  at  $((A_i)_{i\in I},(B_i)_{i\in I})$  is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in Nat $((A_i)_{i \in I}, (B_i)_{i \in I})$  to the map of sets

$$\coprod_{i \in I} f_i \colon \coprod_{i \in I} A_i \to \coprod_{i \in I} B_i$$

defined by

$$\left[ \bigsqcup_{i \in I} f_i \right] (i, a) \stackrel{\text{def}}{=} f_i(a)$$

for each  $(i, a) \in \coprod_{i \in I} A_i$ .

#### PROOF 4.2.2.1.5 ► PROOF OF PROPOSITION 4.2.2.1.4

#### Item 1: Functoriality

This follows from Limits and Colimits, ?? of ??.

### 4.2.3 Binary Coproducts

Let A and B be sets.

#### **DEFINITION 4.2.3.1.1** ► COPRODUCTS OF SETS

The **coproduct of** A **and**  $B^1$  is the coproduct of A and B in Sets as in Limits and Colimits,  $\ref{eq:coproduct}$ ?

<sup>&</sup>lt;sup>1</sup>Further Terminology: Also called the **disjoint union of** A **and** B.

#### CONSTRUCTION 4.2.3.1.2 ► CONSTRUCTION OF COPRODUCTS OF SETS

Concretely, the coproduct of *A* and *B* is the pair  $(A \coprod B, \{inj_1, inj_2\})$  consisting of:

1. *The Colimit.* The set  $A \coprod B$  defined by

$$A \coprod B \stackrel{\text{def}}{=} \coprod_{z \in \{A, B\}} z$$

$$\stackrel{\text{def}}{=} \{(0, a) \in S \mid a \in A\} \cup \{(1, b) \in S \mid b \in B\},$$

where  $S = \{0, 1\} \times (A \cup B)$ .

2. The Cocone. The maps

$$inj_1: A \to A \coprod B,$$
  
 $inj_2: B \to A \coprod B,$ 

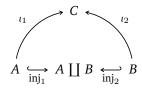
given by

$$\operatorname{inj}_{1}(a) \stackrel{\text{def}}{=} (0, a),$$
 $\operatorname{inj}_{2}(b) \stackrel{\text{def}}{=} (1, b),$ 

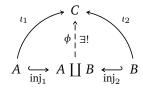
for each  $a \in A$  and each  $b \in B$ .

#### PROOF 4.2.3.1.3 ► PROOF OF CONSTRUCTION 4.2.3.1.2

We claim that  $A \coprod B$  is the categorical coproduct of A and B in Sets. Indeed, suppose we have a diagram of the form



in Sets. Then there exists a unique map  $\phi: A \coprod B \to C$  making the diagram



commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_A = \iota_A,$$
$$\phi \circ \operatorname{inj}_B = \iota_B$$

via

$$\phi(x) = \begin{cases} \iota_A(a) & \text{if } x = (0, a), \\ \iota_B(b) & \text{if } x = (1, b) \end{cases}$$

for each  $x \in A \coprod B$ .

#### PROPOSITION 4.2.3.1.4 ► PROPERTIES OF COPRODUCTS OF SETS

Let A, B, C, and X be sets.

1. Functoriality. The assignment  $A, B, (A, B) \mapsto A \coprod B$  defines functors

$$\begin{array}{cccc} A \coprod -\colon & \mathsf{Sets} & \to \mathsf{Sets}, \\ - \coprod B\colon & \mathsf{Sets} & \to \mathsf{Sets}, \\ -_1 \coprod -_2\colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}, \end{array}$$

where  $-1 \coprod -2$  is the functor where

• *Action on Objects.* For each  $(A, B) \in Obj(Sets \times Sets)$ , we have

$$[-_1 \coprod -_2](A, B) \stackrel{\text{def}}{=} A \coprod B.$$

• *Action on Morphisms*. For each  $(A, B), (X, Y) \in Obj(Sets)$ , the action on Hom-sets

$$\coprod_{(A,B),(X,Y)}$$
: Sets $(A,X) \times$  Sets $(B,Y) \rightarrow$  Sets $(A \coprod B,X \coprod Y)$  of  $\coprod$  at  $((A,B),(X,Y))$  is defined by sending  $(f,g)$  to the function

$$f \mid \mid g: A \mid \mid B \rightarrow X \mid \mid Y$$

defined by

$$[f \coprod g](x) \stackrel{\text{def}}{=} \begin{cases} (0, f(a)) & \text{if } x = (0, a), \\ (1, g(b)) & \text{if } x = (1, b), \end{cases}$$

for each  $x \in A \mid A \mid B$ .

and where  $A \coprod -$  and  $- \coprod B$  are the partial functors of  $-_1 \coprod -_2$  at  $A, B \in \text{Obj}(\mathsf{Sets})$ .

2. Adjointness. We have an adjunction

$$(-_1 \coprod -_2 \dashv \Delta_{\mathsf{Sets}})$$
: Sets  $\times$  Sets  $\underbrace{\bot}_{\Delta_{\mathsf{Sets}}}$  Sets,

witnessed by a bijection

$$\mathsf{Sets}(A \coprod B, C), \cong \mathsf{Hom}_{\mathsf{Sets} \times \mathsf{Sets}}((A, B), (C, C))$$

natural in  $(A, B) \in \text{Obj}(\mathsf{Sets} \times \mathsf{Sets})$  and in  $C \in \text{Obj}(\mathsf{Sets})$ .

3. Associativity. We have an isomorphism of sets

$$\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod} : (X \coprod Y) \coprod Z \xrightarrow{\sim} X \coprod (Y \coprod Z),$$

natural in  $X, Y, Z \in Obj(Sets)$ .

4. Unitality. We have isomorphisms of sets

$$\lambda_X^{\mathsf{Sets}, \coprod} : \emptyset \coprod X \xrightarrow{\sim} X,$$
$$\rho_X^{\mathsf{Sets}, \coprod} : X \coprod \emptyset \xrightarrow{\sim} X,$$

natural in  $X \in Obj(Sets)$ .

5. Commutativity. We have an isomorphism of sets

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod}: X\coprod Y\stackrel{\sim}{\dashrightarrow} Y\coprod X,$$

natural in  $X, Y \in Obj(Sets)$ .

6. *Symmetric Monoidality*. The 7-tuple (Sets,  $\coprod$ ,  $\emptyset$ ,  $\alpha_{\coprod}^{Sets}$ ,  $\lambda_{\coprod}^{Sets}$ ,  $\rho_{\coprod}^{Sets}$ ,  $\sigma_{\coprod}^{Sets}$ ) is a symmetric monoidal category.

#### PROOF 4.2.3.1.5 ► PROOF OF PROPOSITION 4.2.3.1.4

### Item 1: Functoriality

This follows from Limits and Colimits, ?? of ??.

### Item 2: Adjointness

This follows from the universal property of the coproduct.

#### Item 3: Associativity

This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.2.3.1.1.

### Item 4: Unitality

This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.2.4.1.1 and 5.2.5.1.1.

#### Item 5: Commutativity

This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.2.6.1.1.

### Item 6: Symmetric Monoidality

This is a repetition of Monoidal Structures on the Category of Sets, Proposition 5.2.7.1.1, and is proved there.

### 4.2.4 Pushouts

Let A, B, and C be sets and let  $f: C \to A$  and  $g: C \to B$  be functions.

### DEFINITION 4.2.4.1.1 ► PUSHOUTS OF SETS

The **pushout of** A **and** B **over** C **along** f **and**  $g^1$  is the pushout of A and B over C along f and g in Sets as in Limits and Colimits, ??.

#### CONSTRUCTION 4.2.4.1.2 ➤ CONSTRUCTION OF PUSHOUTS OF SETS

Concretely, the pushout of *A* and *B* over *C* along *f* and *g* is the pair  $(A \coprod_C B, \{\text{inj}_1, \text{inj}_2\})$  consisting of:

1. *The Colimit.* The set  $A \coprod_C B$  defined by

$$A \coprod_C B \stackrel{\mathrm{def}}{=} A \coprod B/\sim_C,$$

<sup>&</sup>lt;sup>1</sup>Further Terminology: Also called the **fibre coproduct of** A **and** B **over** C **along** f **and** g.

where  $\sim_C$  is the equivalence relation on  $A \coprod B$  generated by  $(0, f(c)) \sim_C (1, g(c))$ .

2. *The Cocone*. The maps

$$inj_1: A \to A \coprod_C B,$$
  
 $inj_2: B \to A \coprod_C B$ 

given by

$$inj_1(a) \stackrel{\text{def}}{=} [(0, a)] 
inj_2(b) \stackrel{\text{def}}{=} [(1, b)]$$

for each  $a \in A$  and each  $b \in B$ .

#### PROOF 4.2.4.1.3 ► PROOF OF CONSTRUCTION 4.2.4.1.2

We claim that  $A \coprod_C B$  is the categorical pushout of A and B over C with respect to (f,g) in Sets. First we need to check that the relevant pushout diagram commutes, i.e. that we have

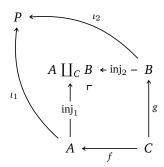
$$\begin{array}{ccc}
A \coprod_{C} B & \stackrel{\operatorname{inj}_{2}}{\longleftarrow} B \\
\operatorname{inj}_{1} \circ f & = \operatorname{inj}_{2} \circ g, & & \inf_{1} & f \\
A & \longleftarrow_{f} & C.
\end{array}$$

Indeed, given  $c \in C$ , we have

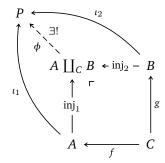
$$\begin{aligned} \left[ \inf_{1} \circ f \right](c) &= \inf_{1} (f(c)) \\ &= \left[ (0, f(c)) \right] \\ &= \left[ (1, g(c)) \right] \\ &= \inf_{2} (g(c)) \\ &= \left[ \inf_{2} \circ g \right](c), \end{aligned}$$

where [(0, f(c))] = [(1, g(c))] by the definition of the relation  $\sim$  on  $A \coprod B$ . Next, we prove that  $A \coprod {}_{C}B$  satisfies the universal property of the pushout.

Suppose we have a diagram of the form



in Sets. Then there exists a unique map  $\phi: A \coprod_C B \to P$  making the diagram



commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_1 = \iota_1,$$
  
$$\phi \circ \operatorname{inj}_2 = \iota_2$$

via

$$\phi(x) = \begin{cases} \iota_1(a) & \text{if } x = [(0, a)], \\ \iota_2(b) & \text{if } x = [(1, b)] \end{cases}$$

for each  $x \in A \coprod_C B$ , where the well-definedness of  $\phi$  is guaranteed by the equality  $\iota_1 \circ f = \iota_2 \circ g$  and the definition of the relation  $\sim$  on  $A \coprod B$  as follows:

1. Case 1: Suppose we have x = [(0, a)] = [(0, a')] for some  $a, a' \in A$ . Then, by Remark 4.2.4.1.4, we have a sequence

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0,a').$$

2. Case 2: Suppose we have x = [(1,b)] = [(1,b')] for some  $b,b' \in B$ . Then, by Remark 4.2.4.1.4, we have a sequence

$$(1,b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b').$$

3. Case 3: Suppose we have x = [(0, a)] = [(1, b)] for some  $a \in A$  and  $b \in B$ . Then, by Remark 4.2.4.1.4, we have a sequence

$$(0, a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1, b).$$

In all these cases, we declare  $x \sim' y$  iff there exists some  $c \in C$  such that x = (0, f(c)) and y = (1, g(c)) or x = (1, g(c)) and y = (0, f(c)). Then, the equality  $\iota_1 \circ f = \iota_2 \circ g$  gives

$$\begin{aligned} \phi([x]) &= \phi([(0, f(c))]) \\ &\stackrel{\text{def}}{=} \iota_1(f(c)) \\ &= \iota_2(g(c)) \\ &\stackrel{\text{def}}{=} \phi([(1, g(c))]) \\ &= \phi([y]), \end{aligned}$$

with the case where x = (1, g(c)) and y = (0, f(c)) similarly giving  $\phi([x]) = \phi([y])$ . Thus, if  $x \sim' y$ , then  $\phi([x]) = \phi([y])$ . Applying this equality pairwise to the sequences

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0,a'),$$
  
 $(1,b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b'),$   
 $(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b)$ 

gives

$$\begin{split} \phi([(0,a)]) &= \phi([(0,a')]), \\ \phi([(1,b)]) &= \phi([(1,b')]), \\ \phi([(0,a)]) &= \phi([(1,b)]), \end{split}$$

showing  $\phi$  to be well-defined.

#### REMARK 4.2.4.1.4 ► Unwinding Definition 4.2.4.1.1

In detail, by Conditions on Relations, Construction 10.5.2.1.2, the relation  $\sim$  of Definition 4.2.4.1.1 is given by declaring  $a \sim b$  iff one of the following conditions is satisfied:

- 1. We have  $a, b \in A$  and a = b.
- 2. We have  $a, b \in B$  and a = b.
- 3. There exist  $x_1, \ldots, x_n \in A \coprod B$  such that  $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$ , where we declare  $x \sim' y$  if one of the following conditions is satisfied:
  - (a) There exists  $c \in C$  such that x = (0, f(c)) and y = (1, g(c)).
  - (b) There exists  $c \in C$  such that x = (1, g(c)) and y = (0, f(c)).

In other words, there exist  $x_1, \ldots, x_n \in A \coprod B$  satisfying the following conditions:

- (c) There exists  $c_0 \in C$  satisfying one of the following conditions:
  - i. We have  $a = f(c_0)$  and  $x_1 = g(c_0)$ .
  - ii. We have  $a = g(c_0)$  and  $x_1 = f(c_0)$ .
- (d) For each  $1 \le i \le n-1$ , there exists  $c_i \in C$  satisfying one of the following conditions:
  - i. We have  $x_i = f(c_i)$  and  $x_{i+1} = g(c_i)$ .
  - ii. We have  $x_i = g(c_i)$  and  $x_{i+1} = f(c_i)$ .
- (e) There exists  $c_n \in C$  satisfying one of the following conditions:
  - i. We have  $x_n = f(c_n)$  and  $b = g(c_n)$ .
  - ii. We have  $x_n = g(c_n)$  and  $b = f(c_n)$ .

#### REMARK 4.2.4.1.5 ▶ PUSHOUTS OF SETS DEPEND ON THE MAPS

It is common practice to write  $A \coprod_C B$  for the pushout of A and B over C along f and g, omitting the maps f and g from the notation and instead leaving them implicit, to be understood from the context.

However, the set  $A \coprod_{C} B$  depends very much on the maps f and g, and sometimes it is necessary or useful to note this dependence explicitly. In such situations, we will write  $A \coprod_{f,C,g} B$  or  $A \coprod_{C} G B$  for  $A \coprod_{C} B B$ .

#### **EXAMPLE 4.2.4.1.6** ► **EXAMPLES OF PUSHOUTS OF SETS**

Here are some examples of pushouts of sets.

- 1. Wedge Sums of Pointed Sets. The wedge sum of two pointed sets of Pointed Sets, Definition 6.3.3.1.1 is an example of a pushout of sets.
- 2. *Intersections via Unions*. Let *X* be a set. We have

$$A \cup B \cong A \coprod_{A \cap B} B, \qquad A \longleftarrow B$$

$$A \longleftarrow A \longrightarrow A \cap B$$

for each  $A, B \in \mathcal{P}(X)$ .

#### PROOF 4.2.4.1.7 ► PROOF OF EXAMPLE 4.2.4.1.6

### Item 1: Wedge Sums of Pointed Sets

This follows by definition, as the wedge sum of two pointed sets is defined as a pushout.

### Item 2: Intersections via Unions

Indeed,  $A \coprod_{A \cap B} B$  is the quotient of  $A \coprod B$  by the equivalence relation obtained by declaring  $(0, a) \sim (1, b)$  iff  $a = b \in A \cap B$ , which is in bijection with  $A \cup B$  via the map with  $[(0, a)] \mapsto a$  and  $[(1, b)] \mapsto b$ .

#### PROPOSITION 4.2.4.1.8 ► PROPERTIES OF PUSHOUTS OF SETS

Let *A*, *B*, *C*, and *X* be sets.

1. Functoriality. The assignment  $(A, B, C, f, g) \mapsto A \coprod_{f,C,g} B$  defines a functor

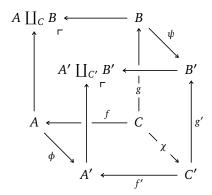
$$-_1 \coprod_{-_3} -_1 \colon \mathsf{Fun}(\mathcal{P},\mathsf{Sets}) \to \mathsf{Sets},$$

where  $\mathcal{P}$  is the category that looks like this:



In particular, the action on morphisms of  $-1 \coprod_{-3} -1$  is given by sending

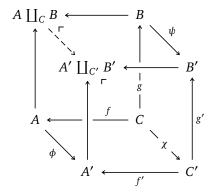
a morphism



in Fun( $\mathcal{P}$ , Sets) to the map  $\xi: A \coprod_C B \xrightarrow{\exists!} A' \coprod_{C'} B'$  given by

$$\xi(x) \stackrel{\text{def}}{=} \begin{cases} \phi(a) & \text{if } x = [(0, a)], \\ \psi(b) & \text{if } x = [(1, b)] \end{cases}$$

for each  $x \in A \coprod_C B$ , which is the unique map making the diagram



commute.

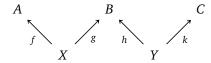
2. Adjointness. We have an adjunction

$$\left(-1 \coprod_{X} -_2 \dashv \Delta_{\mathsf{Sets}_{X/}}\right) \colon \quad \mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/} \underbrace{\bot_{\Delta_{\mathsf{Sets}_{X/}}}}_{\mathsf{Sets}_{X/}} \mathsf{Sets}_{X/},$$

witnessed by a bijection

$$\mathsf{Sets}_{X/}(A \coprod_X B, C), \cong \mathsf{Hom}_{\mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/}}((A, B), (C, C))$$
  
natural in  $(A, B) \in \mathsf{Obj}(\mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/})$  and in  $C \in \mathsf{Obj}(\mathsf{Sets}_{X/})$ .

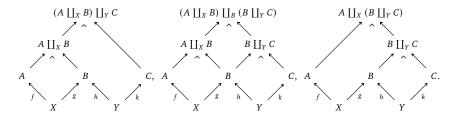
3. Associativity. Given a diagram



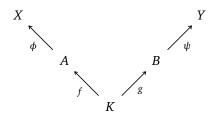
in Sets, we have isomorphisms of sets

$$(A \coprod_X B) \coprod_Y C \cong (A \coprod_X B) \coprod_B (B \coprod_Y C) \cong A \coprod_X (B \coprod_Y C)$$

where these pullbacks are built as in the diagrams



4. Interaction With Composition. Given a diagram



in Sets, we have isomorphisms of sets

$$\begin{split} X \coprod_{K}^{\phi \circ f, \psi \circ g} Y &\cong \left( X \coprod_{A}^{\phi, j_{1}} \left( A \coprod_{K}^{f, g} B \right) \right) \coprod_{A \coprod_{K}^{f, g} B} \left( \left( A \coprod_{K}^{f, g} B \right) \coprod_{B}^{j_{2}, \psi} Y \right) \\ &\cong X \coprod_{A}^{\phi, i} \left( \left( A \coprod_{K}^{f, g} B \right) \coprod_{B}^{j_{2}, \psi} Y \right) \\ &\cong \left( X \coprod_{A}^{\phi, i_{1}} \left( A \coprod_{K}^{f, g} B \right) \right) \coprod_{B}^{j, \psi} Y \end{split}$$

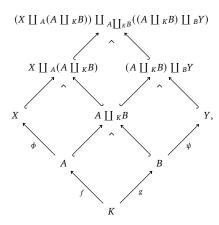
where

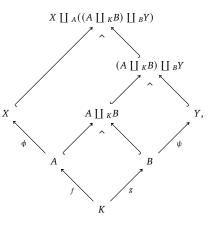
$$j_{1} = \operatorname{inj}_{1}^{A \times_{K}^{f,g} B}, \qquad j_{2} = \operatorname{inj}_{2}^{A \times_{K}^{f,g} B},$$

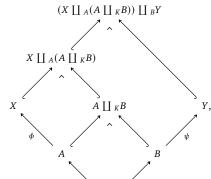
$$i_{1} = \operatorname{inj}_{1}^{\left(A \times_{K}^{f,g} B\right) \times_{Y}^{q_{2},\psi}}, \qquad i_{2} = \operatorname{inj}_{2}^{X \times_{K}^{f,g} B} \left(A \times_{K}^{f,g} B\right),$$

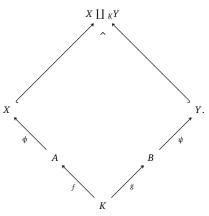
$$i = j_{1} \circ \operatorname{inj}_{1}^{\left(A \times_{K}^{f,g} B\right) \times_{B}^{q_{2},\psi} Y}, \qquad j = j_{2} \circ \operatorname{inj}_{2}^{X \times_{K}^{f,q_{1}} \left(A \times_{K}^{f,g} B\right)},$$

# and where these pullbacks are built as in the diagrams









5. Unitality. We have isomorphisms of sets

$$\begin{array}{ccc}
A & = & & A \\
f & & & \uparrow \\
X & = & & X
\end{array}$$

$$\lambda_A^{\mathsf{Sets}_{X/}} : X \coprod_X A \xrightarrow{\sim} A,$$
 $\rho_A^{\mathsf{Sets}_{X/}} : A \coprod_X X \xrightarrow{\sim} A,$ 

$$\begin{array}{cccc}
A & \stackrel{f}{\longleftarrow} & X \\
\parallel & & \parallel \\
X & \stackrel{f}{\longleftarrow} & X,
\end{array}$$

natural in  $(A, f) \in \text{Obj}(\mathsf{Sets}_{X/})$ .

6. Commutativity. We have an isomorphism of sets

$$A \coprod_{X} B \longleftarrow B$$

$$\uparrow \qquad \qquad \uparrow g \qquad \sigma_{A}^{\mathsf{Sets}_{X/}} : A \coprod_{X} B \xrightarrow{\sim} B \coprod_{X} A \qquad \qquad \uparrow \qquad \qquad \uparrow f$$

$$A \longleftarrow_{f} X, \qquad \qquad B \longleftarrow_{g} X.$$

natural in  $(A, f), (B, g) \in Obj(Sets_{X/})$ .

7. Interaction With Coproducts. We have

$$A \coprod_{\emptyset} B \cong A \coprod_{B}, \qquad \uparrow \qquad \uparrow \iota_{B}$$

$$A \longleftarrow_{\iota_{A}} \emptyset.$$

8. *Symmetric Monoidality*. The triple  $(Sets_{X/}, \coprod_X, X)$  is a symmetric monoidal category.

### PROOF 4.2.4.1.9 ► PROOF OF PROPOSITION 4.2.4.1.8

#### Item 1: Functoriality

This is a special case of functoriality of co/limits, Limits and Colimits, ?? of ??, with the explicit expression for  $\xi$  following from the commutativity of the cube pushout diagram.

#### Item 2: : Adjointness

This follows from the universal property of the coproduct (pushouts are coproducts in  $\mathsf{Sets}_{X/}$ ).

### Item 3: Associativity

Omitted.

#### Item 4: Interaction With Composition

Omitted.

### Item 5: Unitality

Omitted.

Item 6: Commutativity
Omitted.

Item 7: Interaction With Coproducts
Omitted.

Item 8: Symmetric Monoidality
Omitted.

### 4.2.5 Coequalisers

Let *A* and *B* be sets and let  $f, g: A \Rightarrow B$  be functions.

#### **DEFINITION 4.2.5.1.1** ► COEQUALISERS OF SETS

The **coequaliser of** f **and** g is the coequaliser of f and g in Sets as in Limits and Colimits, ??.

#### CONSTRUCTION 4.2.5.1.2 ► CONSTRUCTION OF COEQUALISERS OF SETS

Concretely, the coequaliser of f and g is the pair (CoEq(f, g), coeq(f, g)) consisting of:

1. *The Colimit.* The set CoEq(f, g) defined by

$$CoEq(f, g) \stackrel{\text{def}}{=} B/\sim$$
,

where  $\sim$  is the equivalence relation on *B* generated by  $f(a) \sim g(a)$ .

2. The Cocone. The map

$$coeq(f,g): B \rightarrow CoEq(f,g)$$

given by the quotient map  $\pi: B \twoheadrightarrow B/\sim$  with respect to the equivalence relation generated by  $f(a) \sim g(a)$ .

#### PROOF 4.2.5.1.3 ► PROOF OF CONSTRUCTION 4.2.5.1.2

We claim that CoEq(f,g) is the categorical coequaliser of f and g in Sets. First we need to check that the relevant coequaliser diagram commutes, i.e. that we have

$$coeq(f, g) \circ f = coeq(f, g) \circ g.$$

Indeed, we have

$$[\operatorname{coeq}(f,g) \circ f](a) \stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](f(a))$$

$$\stackrel{\text{def}}{=} [f(a)]$$

$$= [g(a)]$$

$$\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](g(a))$$

$$\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g) \circ g](a)$$

for each  $a \in A$ . Next, we prove that CoEq(f,g) satisfies the universal property of the coequaliser. Suppose we have a diagram of the form

$$A \xrightarrow{f \atop g} B \xrightarrow{\operatorname{coeq}(f,g)} \operatorname{CoEq}(f,g)$$

in Sets. Then, since c(f(a)) = c(g(a)) for each  $a \in A$ , it follows from Conditions on Relations, Items 4 and 5 of Proposition 10.6.2.1.3 that there exists a unique map  $CoEq(f,g) \xrightarrow{\exists !} C$  making the diagram

$$A \xrightarrow{f} B \xrightarrow{\operatorname{coeq}(f,g)} \operatorname{CoEq}(f,g)$$

$$\downarrow g$$

$$\downarrow \exists !$$

$$C$$

commute.

#### REMARK 4.2.5.1.4 ► Unwinding Definition 4.2.5.1.1

In detail, by Conditions on Relations, Construction 10.5.2.1.2, the relation  $\sim$  of Definition 4.2.5.1.1 is given by declaring  $a \sim b$  iff one of the following conditions is satisfied:

1. We have a = b;

- 2. There exist  $x_1, \ldots, x_n \in B$  such that  $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$ , where we declare  $x \sim' y$  if one of the following conditions is satisfied:
  - (a) There exists  $z \in A$  such that x = f(z) and y = g(z).
  - (b) There exists  $z \in A$  such that x = g(z) and y = f(z).

In other words, there exist  $x_1, ..., x_n \in B$  satisfying the following conditions:

- (a) There exists  $z_0 \in A$  satisfying one of the following conditions:
  - i. We have  $a = f(z_0)$  and  $x_1 = g(z_0)$ .
  - ii. We have  $a = g(z_0)$  and  $x_1 = f(z_0)$ .
- (b) For each  $1 \le i \le n-1$ , there exists  $z_i \in A$  satisfying one of the following conditions:
  - i. We have  $x_i = f(z_i)$  and  $x_{i+1} = g(z_i)$ .
  - ii. We have  $x_i = g(z_i)$  and  $x_{i+1} = f(z_i)$ .
- (c) There exists  $z_n \in A$  satisfying one of the following conditions:
  - i. We have  $x_n = f(z_n)$  and  $b = g(z_n)$ .
  - ii. We have  $x_n = g(z_n)$  and  $b = f(z_n)$ .

#### **EXAMPLE 4.2.5.1.5** ► **EXAMPLES OF COEQUALISERS OF SETS**

Here are some examples of coequalisers of sets.

1. *Quotients by Equivalence Relations*. Let *R* be an equivalence relation on a set *X*. We have a bijection of sets

$$X/\sim_R \cong \mathsf{CoEq}\bigg(R \hookrightarrow X \times X \overset{\mathsf{pr}_1}{\underset{\mathsf{pr}_2}{\Longrightarrow}} X\bigg).$$

#### PROOF 4.2.5.1.6 ► PROOF OF EXAMPLE 4.2.5.1.5

### Item 1: Quotients by Equivalence Relations

See [Pro25z].

#### PROPOSITION 4.2.5.1.7 ▶ PROPERTIES OF COEQUALISERS OF SETS

Let A, B, and C be sets.

1. Associativity. We have isomorphisms of sets<sup>1</sup>

$$\underbrace{\mathsf{CoEq}(\mathsf{coeq}(f,g) \circ f, \mathsf{coeq}(f,g) \circ h)}_{=\mathsf{CoEq}(\mathsf{coeq}(f,g) \circ g, \mathsf{coeq}(f,g) \circ h)} \cong \mathsf{CoEq}(f,g,h) \cong \underbrace{\mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ g)}_{=\mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ h)}$$

where CoEq(f, g, h) is the colimit of the diagram

$$A \xrightarrow{f \atop -g \xrightarrow{h}} B$$

in Sets.

2. Unitality. We have an isomorphism of sets

$$CoEq(f, f) \cong B$$
.

3. Commutativity. We have an isomorphism of sets

$$CoEq(f, g) \cong CoEq(g, f)$$
.

4. Interaction With Composition. Let

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C$$

be functions. We have a surjection

$$CoEq(h \circ f, k \circ g) \twoheadrightarrow CoEq(coeq(h, k) \circ h \circ f, coeq(h, k) \circ k \circ g)$$

exhibiting  $CoEq(coeq(h, k) \circ h \circ f, coeq(h, k) \circ k \circ g)$  as a quotient of  $CoEq(h \circ f, k \circ g)$  by the relation generated by declaring  $h(y) \sim k(y)$  for each  $y \in B$ .

(a) Take the coequaliser of (f, g, h), i.e. the colimit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets.

<sup>&</sup>lt;sup>1</sup>That is, the following three ways of forming "the" coequaliser of (f, g, h) agree:

(b) First take the coequaliser of f and g, forming a diagram

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{\operatorname{coeq}(f,g)}{\Longrightarrow} \operatorname{CoEq}(f,g)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\underset{h}{\Longrightarrow}} B \stackrel{\text{coeq}(f,g)}{\Longrightarrow} \text{CoEq}(f,g),$$

obtaining a quotient

$$\label{eq:coeq} {\sf CoEq}({\sf coeq}(f,g)\circ f, {\sf coeq}(f,g)\circ h) = {\sf CoEq}({\sf coeq}(f,g)\circ g, {\sf coeq}(f,g)\circ h)$$
 of  ${\sf CoEq}(f,g)$ 

(c) First take the coequaliser of g and h, forming a diagram

$$A \stackrel{g}{\underset{h}{\Longrightarrow}} B \stackrel{\operatorname{coeq}(g,h)}{\twoheadrightarrow} \operatorname{CoEq}(g,h)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{\operatorname{coeq}(g,h)}{\Longrightarrow} \operatorname{CoEq}(g,h),$$

obtaining a quotient

$$\label{eq:coeq} \begin{split} \mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ g) &= \mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ h) \\ \mathsf{of} \ \mathsf{CoEq}(g,h). \end{split}$$

### PROOF 4.2.5.1.8 ► PROOF OF PROPOSITION 4.2.5.1.7

Item 1: Associativity

Omitted.

Item 2: Unitality

Omitted.

Item 3: Commutativity

Omitted.

Item 4: Interaction With Composition

Omitted.

#### 4.2.6 Direct Colimits

Let  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ :  $(I, \preceq) \to \pi$  be a direct system of sets.

#### **DEFINITION 4.2.6.1.1** ► **DIRECT COLIMITS OF SETS**

The **direct colimit of**  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$  is the direct colimit of  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$  in Sets as in Limits and Colimits, ??.

#### CONSTRUCTION 4.2.6.1.2 ► CONSTRUCTION OF DIRECT COLIMITS OF SETS

Concretely, the direct colimit of  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$  is the pair  $\left( \underset{\alpha\in I}{\operatorname{colim}}(X_{\alpha}), \underset{\alpha\in I}{\longleftrightarrow} \right)$ 

1. *The Colimit*. The set  $\underset{\longrightarrow}{\operatorname{colim}}(X_{\alpha})$  defined by

$$\operatorname{colim}_{\underset{\alpha\in I}{\longrightarrow}}(X_{\alpha})\stackrel{\mathrm{def}}{=}\left(\bigsqcup_{\alpha\in I}X_{\alpha}\right)\bigg/\sim,$$

where  $\sim$  is the equivalence relation on  $\coprod_{\alpha \in I} X_{\alpha}$  generated by declaring  $(\alpha, x) \sim (\beta, y)$  iff there exists some  $\gamma \in I$  satisfying the following conditions:

- (a) We have  $\alpha \leq \gamma$ .
- (b) We have  $\beta \leq \gamma$ .
- (c) We have  $f_{\alpha \gamma}(x) = f_{\beta \gamma}(y)$ .
- 2. The Cocone. The collection

$$\left\{\operatorname{inj}_{\gamma}: X_{\gamma} \to \operatorname{colim}_{\alpha \in I}(X_{\alpha})\right\}_{\gamma \in I}$$

of maps of sets defined by

$$\operatorname{inj}_{\gamma}(x) \stackrel{\text{def}}{=} [(\gamma, x)]$$

for each  $\gamma \in I$  and each  $x \in X_{\gamma}$ .

#### PROOF 4.2.6.1.3 ► PROOF OF CONSTRUCTION 4.2.6.1.2

We will prove Construction 4.2.6.1.2 below in a bit, but first we need a lemma (which is interesting in its own right).

### LEMMA 4.2.6.1.4 $\blacktriangleright$ Identification of x with $f_{\alpha\beta}(x)$ in Direct Colimits

For each  $\alpha, \beta \in I$  and each  $x \in X_{\alpha}$ , if  $\alpha \leq \beta$ , then we have

$$(\alpha, x) \sim (\beta, f_{\alpha\beta}(x))$$

in  $\operatornamewithlimits{colim}_{\underset{\alpha\in I}{\longrightarrow}}(X_{\alpha})$ .

#### PROOF 4.2.6.1.5 ▶ PROOF OF LEMMA 4.2.6.1.4

Taking  $\gamma = \beta$ , we have  $f_{\alpha\gamma} = f_{\alpha\beta}$ , we have  $f_{\beta\gamma} = f_{\beta\beta} \stackrel{\text{def}}{=} \mathrm{id}_{X_{\beta}}$ , and we have

$$f_{\alpha\beta}(x) = f_{\beta\beta} (f_{\alpha\beta}(x))$$

$$\stackrel{\text{def}}{=} id_{X_{\beta}} (f_{\alpha\beta}(x)),$$

$$= f_{\alpha\beta}(x).$$

As a result, since  $\alpha \leq \beta$  and  $\beta \leq \beta$  as well, Items 1a to 1c of Construction 4.2.6.1.2 are met. Thus we have  $(\alpha, x) \sim (\beta, f_{\alpha\beta}(x))$ .

We can now prove Construction 4.2.6.1.2:

### PROOF 4.2.6.1.6 ► PROOF OF CONSTRUCTION 4.2.6.1.2

We claim that  $\operatornamewithlimits{colim}_{\alpha \in I}(X_{\alpha})$  is the colimit of the direct system of sets  $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta \in I}$ .

### Commutativity of the Colimit Diagram

First, we need to check that the colimit diagram defined by  $\operatornamewithlimits{colim}_{\alpha \in I}(X_{\alpha})$ 

commutes, i.e. that we have

for each  $\alpha, \beta \in I$  with  $\alpha \leq \beta$ . Indeed, given  $x \in X_{\alpha}$ , we have

$$[\inf_{\beta} \circ f_{\alpha\beta}](x) \stackrel{\text{def}}{=} \inf_{\beta} (f_{\alpha\beta}(x))$$

$$\stackrel{\text{def}}{=} [(\beta, f_{\alpha\beta}(x))]$$

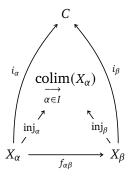
$$= [(\alpha, x)]$$

$$\stackrel{\text{def}}{=} \inf_{\alpha} (x),$$

where we have used Lemma 4.2.6.1.4 for the third equality.

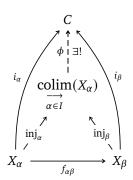
### Proof of the Universal Property of the Colimit

Next, we prove that  $\operatornamewithlimits{colim}_{\alpha \in I}(X_{\alpha})$  as constructed in Construction 4.2.6.1.2 satisfies the universal property of a direct colimit. Suppose that we have, for each  $\alpha, \beta \in I$  with  $\alpha \leq \beta$ , a diagram of the form



in Sets. We claim that there exists a unique map  $\phi$ :  $\underset{\longrightarrow}{\operatorname{colim}}(X_{\alpha}) \xrightarrow{\exists !} C$  making

the diagram



commute. To this end, first consider the diagram

$$\coprod_{\alpha \in I} X_{\alpha} \xrightarrow{\operatorname{pr}} \operatorname{colim}_{\alpha \in I} (X_{\alpha})$$

$$\coprod_{\alpha \in I} i_{\alpha} \qquad C.$$

**Lemma.** If  $(\alpha, x) \sim (\beta, y)$ , then we have

$$\left[ \bigsqcup_{\alpha \in I} i_{\alpha} \right](x) = \left[ \bigsqcup_{\alpha \in I} i_{\alpha} \right](y).$$

*Proof.* Indeed, if  $(\alpha, x) \sim (\beta, y)$ , then there exists some  $\gamma \in I$  satisfying the following conditions:

- 1. We have  $\alpha \leq \gamma$ .
- 2. We have  $\beta \leq \gamma$ .
- 3. We have  $f_{\alpha\gamma}(x) = f_{\beta\gamma}(y)$ .

We then have

$$\left[ \bigsqcup_{\alpha \in I} i_{\alpha} \right](x) \stackrel{\text{def}}{=} i_{\alpha}(x)$$

$$\stackrel{\text{def}}{=} \left[ i_{\gamma} \circ f_{\alpha \gamma} \right](x)$$

$$\stackrel{\text{def}}{=} i_{\gamma} (f_{\alpha \gamma}(x))$$

$$= i_{\gamma} (f_{\beta \gamma}(x))$$

$$\stackrel{\text{def}}{=} [i_{\gamma} \circ f_{\beta \gamma}](x)$$

$$= i_{\beta}(y)$$

$$\stackrel{\text{def}}{=} \left[ \coprod_{\alpha \in I} i_{\alpha} \right](y).$$

This finishes the proof of the lemma. Continuing, by Conditions on Relations, ?? of Proposition 10.6.2.1.3, there then exists a map  $\phi: \underset{\alpha \in I}{\operatorname{colim}}(X_{\alpha}) \xrightarrow{\exists !} C$  making the diagram

$$\coprod_{\alpha \in I} X_{\alpha} \xrightarrow{\operatorname{pr}} \operatorname{colim}_{\alpha \in I} (X_{\alpha})$$

$$\coprod_{\alpha \in I} i_{\alpha} \qquad \qquad \downarrow \phi$$

$$C$$

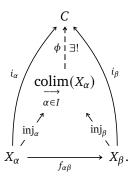
commute. In particular, this implies that the diagram

$$X_{\alpha} \xrightarrow{\inf_{\alpha}} \underset{\alpha \in I}{\underbrace{\operatorname{colim}}}(X_{\alpha})$$

$$\downarrow i_{\alpha} \qquad \qquad \downarrow \phi$$

$$C$$

also commutes, and thus so does the diagram



This finishes the proof.<sup>1</sup>



<sup>1</sup>Incidentally, the conditions

$$\left\{i_{\alpha} = \phi \circ \operatorname{inj}_{\alpha}\right\}_{\alpha \in I}$$

show that  $\phi$  must be given by

$$\phi([(\alpha, x)]) = (i_{\alpha}(x))_{\alpha \in I}$$

for each  $[(\alpha,x)] \in \underset{\alpha \in I}{\operatorname{colim}}(X_\alpha)$ , although we would need to show that this assignment is well-defined were we to prove Construction 4.2.6.1.2 in this way. Instead, invoking Conditions on Relations, ?? of Proposition 10.6.2.1.3 gave us a way to avoid having to prove this, leading to a cleaner alternative proof.

### **EXAMPLE 4.2.6.1.7** ► **EXAMPLES OF DIRECT COLIMITS OF SETS**

Here are some examples of direct colimits of sets.

1. *The Prüfer Group*. The Prüfer group  $\mathbb{Z}(p^{\infty})$  is defined as the direct colimit

$$\mathbb{Z}(p^{\infty}) \stackrel{\text{def}}{=} \underset{n \in \mathbb{N}}{\operatorname{colim}} (\mathbb{Z}_{/p^n});$$

see ??.

# 4.3 Operations With Sets

# 4.3.1 The Empty Set

## DEFINITION 4.3.1.1.1 ► THE EMPTY SET

The **empty set** is the set Ø defined by

$$\emptyset \stackrel{\text{def}}{=} \{ x \in X \mid x \neq x \},\$$

where X is the set in the set existence axiom, ?? of ??.

## 4.3.2 Singleton Sets

Let *X* be a set.

### **DEFINITION 4.3.2.1.1** ► SINGLETON SETS

The **singleton set containing** X is the set  $\{X\}$  defined by

$$\{X\} \stackrel{\mathrm{def}}{=} \{X, X\},\,$$

where  $\{X, X\}$  is the pairing of X with itself of Definition 4.3.3.1.1.

### 4.3.3 Pairings of Sets

Let *X* and *Y* be sets.

#### **DEFINITION 4.3.3.1.1** ► PAIRINGS OF SETS

The **pairing of** X **and** Y is the set  $\{X,Y\}$  defined by

$${X,Y} \stackrel{\text{def}}{=} {x \in A \mid x = X \text{ or } x = Y},$$

where *A* is the set in the axiom of pairing, ?? of ??.

#### 4.3.4 Ordered Pairs

Let *A* and *B* be sets.

#### **DEFINITION 4.3.4.1.1** ► ORDERED PAIRS

The **ordered pair associated to** A **and** B is the set (A, B) defined by

$$(A, B) \stackrel{\text{def}}{=} \{ \{A\}, \{A, B\} \}.$$

### PROPOSITION 4.3.4.1.2 ► PROPERTIES OF ORDERED PAIRS

Let A and B be sets.

- 1. *Uniqueness*. Let *A*, *B*, *C*, and *D* be sets. The following conditions are equivalent:
  - (a) We have (A, B) = (C, D).
  - (b) We have A = C and B = D.

#### PROOF 4.3.4.1.3 ► PROOF OF PROPOSITION 4.3.4.1.2

### Item 1: Uniqueness

See [Cie97, Theorem 1.2.3].

### 4.3.5 Sets of Maps

Let *A* and *B* be sets.

### **DEFINITION 4.3.5.1.1** ► **SETS OF MAPS**

The **set of maps from** A **to**  $B^1$  is the set  $Sets(A, B)^2$  whose elements are the functions from A to B.

### PROPOSITION 4.3.5.1.2 ► PROPERTIES OF SETS OF MAPS

Let A and B be sets.

1. *Functoriality*. The assignments  $X, Y, (X, Y) \mapsto \text{Hom}_{\mathsf{Sets}}(X, Y)$  define functors

$$\begin{array}{lll} \mathsf{Sets}(X,-)\colon & \mathsf{Sets} & \to \mathsf{Sets}, \\ \mathsf{Sets}(-,Y)\colon & \mathsf{Sets}^\mathsf{op} & \to \mathsf{Sets}, \\ \mathsf{Sets}(-_1,-_2)\colon & \mathsf{Sets}^\mathsf{op} \times \mathsf{Sets} \to \mathsf{Sets}. \end{array}$$

2. Adjointness. We have adjunctions

$$(A \times - \dashv \mathsf{Sets}(A, -))$$
: Sets  $\underbrace{ \xrightarrow{A \times -}}_{\mathsf{Sets}(A, -)} \mathsf{Sets},$ 

$$(-\times B + \mathsf{Sets}(B, -))$$
: Sets  $\underbrace{-\times B}_{\mathsf{Sets}(B, -)}$  Sets,

witnessed by bijections

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$
  
 $Sets(A \times B, C) \cong Sets(B, Sets(A, C)),$ 

natural in  $A, B, C \in Obj(Sets)$ .

<sup>&</sup>lt;sup>1</sup>Further Terminology: Also called the **Hom set from** *A* **to** *B*.

<sup>&</sup>lt;sup>2</sup>Further Notation: Also written  $Hom_{Sets}(A, B)$ .

3. Maps From the Punctual Set. We have a bijection

$$\mathsf{Sets}(\mathsf{pt},A)\cong A,$$

natural in  $A \in Obj(Sets)$ .

4. Maps to the Punctual Set. We have a bijection

$$Sets(A, pt) \cong pt$$
,

natural in  $A \in Obj(Sets)$ .

### PROOF 4.3.5.1.3 ► PROOF OF PROPOSITION 4.3.5.1.2

### Item 1: Functoriality

This follows from Categories, Items 2 and 5 of Proposition 11.1.4.1.2.

### Item 2: Adjointness

This is a repetition of Item 2 of Proposition 4.1.3.1.4 and is proved there.

## Item 3: Maps From the Punctual Set

The bijection

$$\Phi_A \colon \mathsf{Sets}(\mathsf{pt},A) \xrightarrow{\sim} A$$

is given by

$$\Phi_A(f) \stackrel{\text{def}}{=} f(\star)$$

for each  $f \in Sets(pt, A)$ , admitting an inverse

$$\Phi_A^{-1} : A \xrightarrow{\sim} \mathsf{Sets}(\mathsf{pt}, A)$$

given by

$$\Phi_A^{-1}(a) \stackrel{\mathrm{def}}{=} \llbracket \star \mapsto a \rrbracket$$

for each  $a \in A$ . Indeed, we have

$$\begin{split} \left[\Phi_A^{-1} \circ \Phi_A\right](f) &\stackrel{\text{def}}{=} \Phi_A^{-1}(\Phi_A(f)) \\ &\stackrel{\text{def}}{=} \Phi_A^{-1}(f(\star)) \\ &\stackrel{\text{def}}{=} \left[\!\!\left[\star \mapsto f(\star)\right]\!\!\right] \\ &\stackrel{\text{def}}{=} f \end{split}$$

$$\stackrel{\text{def}}{=} \left[ \mathrm{id}_{\mathsf{Sets}(\mathsf{pt},A)} \right] (f)$$

for each  $f \in Sets(pt, A)$  and

$$\begin{split} \left[\Phi_{A} \circ \Phi_{A}^{-1}\right](a) &\stackrel{\text{def}}{=} \Phi_{A} \left(\Phi_{A}^{-1}(a)\right) \\ &\stackrel{\text{def}}{=} \Phi_{A} (\llbracket \star \mapsto a \rrbracket) \\ &\stackrel{\text{def}}{=} \operatorname{ev}_{\star} (\llbracket \star \mapsto a \rrbracket) \\ &\stackrel{\text{def}}{=} a \\ &\stackrel{\text{def}}{=} [\operatorname{id}_{A}](a) \end{split}$$

for each  $a \in A$ , and thus we have

$$\Phi_A^{-1} \circ \Phi_A = \mathrm{id}_{\mathsf{Sets}(\mathsf{pt},A)}$$
  
 $\Phi_A \circ \Phi_A^{-1} = \mathrm{id}_A$ .

To prove naturality, we need to show that the diagram

$$\begin{array}{ccc}
\mathsf{Sets}(\mathsf{pt},A) & \xrightarrow{f_!} & \mathsf{Sets}(\mathsf{pt},B) \\
& & \downarrow & & \downarrow \\
\Phi_A & & \downarrow & & \downarrow \\
A & & & & B
\end{array}$$

commutes. Indeed, we have

$$[f \circ \Phi_A](\phi) \stackrel{\text{def}}{=} f(\Phi_A(\phi))$$

$$\stackrel{\text{def}}{=} f(\phi(\star))$$

$$\stackrel{\text{def}}{=} [f \circ \phi](\star)$$

$$\stackrel{\text{def}}{=} \Phi_B(f \circ \phi)$$

$$\stackrel{\text{def}}{=} \Phi_B(f_!(\phi))$$

$$\stackrel{\text{def}}{=} [\Phi_B \circ f_!](\phi)$$

for each  $\phi \in Sets(pt, A)$ . This finishes the proof.

## Item 4: Maps to the Punctual Set

This follows from the universal property of pt as the terminal set, Definition 4.1.1.1.

### 4.3.6 Unions of Families of Subsets

Let *X* be a set and let  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

#### **DEFINITION 4.3.6.1.1** ► Unions of Families of Subsets

The **union of**  $\mathcal{U}$  is the set  $\bigcup_{U \in \mathcal{U}} U$  defined by

$$\bigcup_{U \in \mathcal{U}} U \stackrel{\text{def}}{=} \left\{ x \in X \;\middle|\; \text{there exists some } U \in \mathcal{U} \right\}.$$

#### PROPOSITION 4.3.6.1.2 ➤ PROPERTIES OF UNIONS OF FAMILIES OF SUBSETS

Let *X* be a set.

1. Functoriality. The assignment  $\mathcal{U} \mapsto \bigcup_{U \in \mathcal{U}} U$  defines a functor

$$\bigcup : (\mathcal{P}(\mathcal{P}(X)), \subset) \to (\mathcal{P}(X), \subset).$$

In particular, for each  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ , the following condition is satisfied:

$$(\star) \text{ If } \mathcal{U} \subset \mathcal{V} \text{, then } \bigcup_{U \in \mathcal{U}} U \subset \bigcup_{V \in \mathcal{V}} V.$$

2. Associativity. The diagram

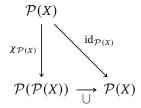
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(\mathcal{P}(X))) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \bigcup} & \mathcal{P}(\mathcal{P}(X)) \\
\cup \star \mathrm{id}_{\mathcal{P}(X)} & & & & & & & & & \\
\mathcal{P}(\mathcal{P}(X)) & & & & & & & & & & \\
\end{array}$$

commutes, i.e. we have

$$\bigcup_{U \in \bigcup_{A \in A} A} U = \bigcup_{A \in \mathcal{A}} \left( \bigcup_{U \in A} U \right)$$

for each  $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$ .

3. Left Unitality. The diagram

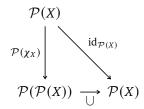


commutes, i.e. we have

$$\bigcup_{V \in \{U\}} V = U$$

for each  $U \in \mathcal{P}(X)$ .

4. Right Unitality. The diagram



commutes, i.e. we have

$$\bigcup_{\{u\}\in\chi_X(U)}\{u\}=U$$

for each  $U \in \mathcal{P}(X)$ .

5. Interaction With Unions I. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\cup} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\cup} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{W\in\mathcal{U}\cup\mathcal{V}}W=\left(\bigcup_{U\in\mathcal{U}}U\right)\cup\left(\bigcup_{V\in\mathcal{V}}V\right)$$

for each  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

6. Interaction With Unions II. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcup_{V \in \mathcal{V}} V\right) = \bigcup_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcup_{U \in \mathcal{U}} U\right) \cup V = \bigcup_{U \in \mathcal{U}} (U \cup V)$$

for each  $U, V \in \mathcal{P}(\mathcal{P}(X))$  and each  $U, V \in \mathcal{P}(X)$ .

7. Interaction With Intersections I. We have a natural transformation

with components

$$\bigcup_{W\in\mathcal{U}\cap\mathcal{V}}W\subset\left(\bigcup_{U\in\mathcal{U}}U\right)\cap\left(\bigcup_{V\in\mathcal{V}}V\right)$$

for each  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

8. Interaction With Intersections II. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcup_{V \in \mathcal{V}} V\right) = \bigcup_{V \in \mathcal{V}} (U \cup V),$$

$$\left(\bigcup_{U\in\mathcal{U}}U\right)\cup V=\bigcup_{U\in\mathcal{U}}(U\cup V)$$

for each  $U, V \in \mathcal{P}(\mathcal{P}(X))$  and each  $U, V \in \mathcal{P}(X)$ .

9. Interaction With Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\setminus} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\setminus} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in \mathcal{U} \setminus \mathcal{V}} W \neq \left(\bigcup_{U \in \mathcal{U}} U\right) \setminus \left(\bigcup_{V \in \mathcal{V}} V\right)$$

in general, where  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

10. Interaction With Complements I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\operatorname{op}} \xrightarrow{(-)^{c}} \mathcal{P}(\mathcal{P}(X))$$

$$\cup^{\operatorname{op}} \bigvee \bigvee \bigcup \bigcup$$

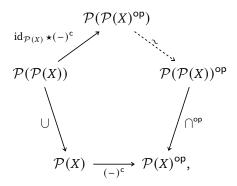
$$\mathcal{P}(X)^{\operatorname{op}} \xrightarrow{(-)^{c}} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{U \in \mathcal{U}^{\mathsf{c}}} U \neq \bigcup_{U \in \mathcal{U}} U^{\mathsf{c}}$$

in general, where  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

### 11. Interaction With Complements II. The diagram

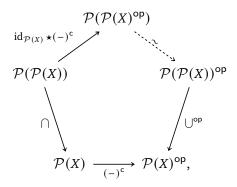


commutes, i.e. we have

$$\left(\bigcup_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

12. Interaction With Complements III. The diagram



commutes, i.e. we have

$$\left(\bigcap_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

13. Interaction With Symmetric Differences. The diagram

$$\begin{array}{cccc}
\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) & \stackrel{\triangle}{\longrightarrow} & \mathcal{P}(\mathcal{P}(X)) \\
 & & & & \downarrow & & \downarrow & \\
\mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{\triangle} & \mathcal{P}(X), & & & \\
\end{array}$$

does not commute in general, i.e. we may have

$$\bigcup_{W\in\mathcal{U}\triangle\mathcal{V}}W\neq\left(\bigcup_{U\in\mathcal{U}}U\right)\triangle\left(\bigcup_{V\in\mathcal{V}}V\right)$$

in general, where  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

14. Interaction With Internal Homs I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-_2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\downarrow^{\mathsf{op}} \times \cup^{\mathsf{op}} \qquad \qquad \downarrow \cup$$

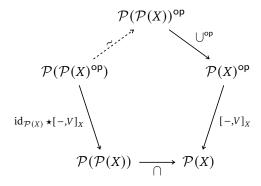
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-_2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcup_{U \in \mathcal{U}} U, \bigcup_{V \in \mathcal{V}} V\right]_X$$

in general, where  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

15. Interaction With Internal Homs II. The diagram



commutes, i.e. we have

$$\left[\bigcup_{U\in\mathcal{U}}U,V\right]_X=\bigcap_{U\in\mathcal{U}}[U,V]_X$$

for each  $U \in \mathcal{P}(\mathcal{P}(X))$  and each  $V \in \mathcal{P}(X)$ .

16. Interaction With Internal Homs III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcup} & \mathcal{P}(X) \\
id_{\mathcal{P}(X)} \star [U,-]_X & & & \downarrow [U,-]_X \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcup} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U, \bigcup_{V \in \mathcal{V}} V\right]_X = \bigcup_{V \in \mathcal{V}} [U, V]_X$$

for each  $U \in \mathcal{P}(X)$  and each  $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

17. *Interaction With Direct Images.* Let  $f: X \to Y$  be a map of sets. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_1)_1} & \mathcal{P}(\mathcal{P}(Y)) \\
\downarrow & & \downarrow \cup \\
\mathcal{P}(X) & \xrightarrow{f_1} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$\bigcup_{U\in\mathcal{U}}f_!(U)=\bigcup_{V\in f_!(\mathcal{U})}V$$

for each  $\mathcal{U} \in \mathcal{P}(X)$ , where  $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$ .

18. *Interaction With Inverse Images*. Let  $f: X \to Y$  be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \qquad \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{V \in \mathcal{V}} f^{-1}(V) = \bigcup_{U \in f^{-1}(\mathcal{U})} U$$

for each  $V \in \mathcal{P}(Y)$ , where  $f^{-1}(V) \stackrel{\text{def}}{=} (f^{-1})^{-1}(V)$ .

19. Interaction With Codirect Images. Let  $f: X \to Y$  be a map of sets. The diagram

$$\begin{array}{ccc} \mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_*)_*} & \mathcal{P}(\mathcal{P}(Y)) \\ & & & \downarrow \cup \\ & & & \downarrow \cup \\ & \mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y) \end{array}$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_*(U) = \bigcup_{V \in f_*(\mathcal{U})} V$$

for each  $\mathcal{U} \in \mathcal{P}(X)$ , where  $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$ .

20. Interaction With Intersections of Families I. The diagram

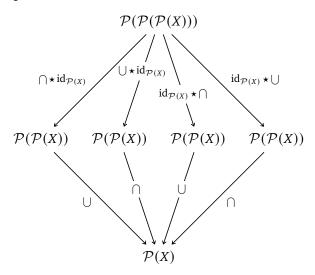
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \cap} & \mathcal{P}(\mathcal{P}(X)) \\
\downarrow \star \mathrm{id}_{\mathcal{P}(X)} & & & \downarrow \cap \\
\mathcal{P}(X) & \xrightarrow{\qquad \cap} & X
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \bigcup_{A \in A} A} U = \bigcap_{A \in \mathcal{A}} \left( \bigcap_{U \in A} U \right)$$

for each  $A \in \mathcal{P}(\mathcal{P}(X))$ .

21. *Interaction With Intersections of Families II.* Let *X* be a set and consider the compositions

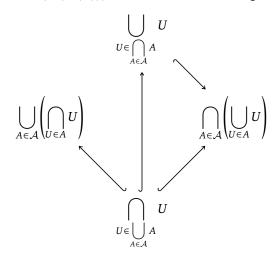


given by

$$A \mapsto \bigcup_{U \in \bigcap_{A \in \mathcal{A}} A} U, \quad A \mapsto \bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U,$$

$$A \mapsto \bigcup_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U\right), \quad A \mapsto \bigcap_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U\right)$$

for each  $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$ . We have the following inclusions:



All other possible inclusions fail to hold in general.

#### PROOF 4.3.6.1.3 ► PROOF OF PROPOSITION 4.3.6.1.2

### Item 1: Functoriality

Since  $\mathcal{P}(X)$  is posetal, it suffices to prove the condition  $(\star)$ . So let  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$  with  $\mathcal{U} \subset \mathcal{V}$ . We claim that

$$\bigcup_{U\in\mathcal{U}}U\subset\bigcup_{V\in\mathcal{V}}V.$$

Indeed, given  $x \in \bigcup_{U \in \mathcal{U}} U$ , there exists some  $U \in \mathcal{U}$  such that  $x \in U$ , but since  $\mathcal{U} \subset \mathcal{V}$ , we have  $U \in \mathcal{V}$  as well, and thus  $x \in \bigcup_{V \in \mathcal{V}} V$ , which gives our desired inclusion.

# Item 2: Associativity

We have

$$\bigcup_{U \in \bigcup_{A \in \mathcal{A}} A} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } U \in \bigcup_{A \in \mathcal{A}} A \right\}$$

$$= \left\{ x \in X \middle| \text{ there exists some } A \in \mathcal{A} \right\}$$

$$= \left\{ x \in X \middle| \text{ there exists some } A \in \mathcal{A} \right\}$$

$$= \left\{ x \in X \middle| \text{ there exists some } A \in \mathcal{A} \right\}$$

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$$= \left\{ x \in X \middle| \text{ there exists some } A \in \mathcal{A} \right\}$$

$$= \left\{ x \in X \middle| \text{ there exists some } A \in \mathcal{A} \right\}$$

$$= \left\{ x \in X \middle| \text{ there exists some } A \in \mathcal{A} \right\}$$

This finishes the proof.

### Item 3: Left Unitality

We have

$$\bigcup_{V \in \{U\}} V \stackrel{\text{def}}{=} \left\{ x \in X \;\middle|\; \text{there exists some } V \in \{U\} \right\}$$
 such that we have  $x \in U$ 

$$= \{x \in X \mid x \in U\}$$
$$= U.$$

This finishes the proof.

### Item 4: Right Unitality

We have

$$\bigcup_{\{u\} \in \chi_X(U)} \{u\} \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } \{u\} \in \chi_X(U) \right\}$$
 such that we have  $x \in \{u\}$ 

$$= \left\{ x \in X \middle| \text{ there exists some } \{u\} \in \chi_X(U) \right\}$$
 such that we have  $x = u$ 

$$= \left\{ x \in X \middle| \text{ there exists some } u \in U \right\}$$
 such that we have  $x = u$ 

$$= \left\{ x \in X \middle| \text{ there exists some } u \in U \right\}$$
 such that we have  $x = u$ 

$$= \left\{ x \in X \middle| x \in U \right\}$$

$$= U.$$

This finishes the proof.

# Item 5: Interaction With Unions I

We have

$$\bigcup_{W \in \mathcal{U} \cup \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{U} \cup \mathcal{V} \right\} \\
= \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{U} \text{ or some} \right\} \\
\stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{U} \text{ or some} \right\} \\
\downarrow \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{U} \right\} \\
\downarrow \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{V} \right\} \\
\downarrow \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{V} \right\} \\
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\downarrow \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{V} \right\} \\
\downarrow \left\{ x \in X \middle| \text{ there exists$$

This finishes the proof.

#### Item 6: Interaction With Unions II

Omitted.

### Item 7: Interaction With Intersections I

We have

$$\bigcup_{W \in \mathcal{U} \cap \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{c} \text{there exists some } W \in \mathcal{U} \cap \mathcal{V} \\ \text{such that we have } x \in W \end{array} \right\}$$

$$\subset \left\{ x \in X \middle| \begin{array}{c} \text{there exists some } U \in \mathcal{U} \text{ and some } V \in \mathcal{V} \\ \text{such that we have } x \in U \text{ and } x \in V \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{c} \text{there exists some } U \in \mathcal{U} \\ \text{such that we have } x \in U \end{array} \right\}$$

$$\cup \left\{ x \in X \middle| \begin{array}{c} \text{there exists some } V \in \mathcal{V} \\ \text{such that we have } x \in V \right\}$$

$$\stackrel{\text{def}}{=} \left( \bigcup_{U \in \mathcal{U}} U \right) \cap \left( \bigcup_{V \in \mathcal{V}} V \right).$$

This finishes the proof.

# Item 8: Interaction With Intersections II

Omitted.

### Item 9: Interaction With Differences

Let  $X = \{0, 1\}$ , let  $\mathcal{U} = \{\{0, 1\}\}$ , and let  $\mathcal{V} = \{\{0\}\}$ . We have

$$\bigcup_{W \in \mathcal{U} \setminus \mathcal{V}} U = \bigcup_{W \in \{\{0,1\}\}} W$$
$$= \{0,1\},$$

whereas

$$\left(\bigcup_{U \in \mathcal{U}} U\right) \setminus \left(\bigcup_{V \in \mathcal{V}} V\right) = \{0, 1\} \setminus \{0\}$$
$$= \{1\}.$$

Thus we have

$$\bigcup_{W\in\mathcal{U}\backslash\mathcal{V}}W=\left\{ 0,1\right\} \neq\left\{ 1\right\} =\left(\bigcup_{U\in\mathcal{U}}U\right)\backslash\left(\bigcup_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

# Item 10: Interaction With Complements I

Let  $X = \{0, 1\}$  and let  $\mathcal{U} = \{0\}$ . We have

$$\bigcup_{U \in \mathcal{U}^{c}} U = \bigcup_{U \in \{\emptyset, \{1\}, \{0,1\}\}} U$$
$$= \{0, 1\},$$

whereas

$$\bigcup_{U \in \mathcal{U}} U^{c} = \{0\}^{c}$$
$$= \{1\}.$$

Thus we have

$$\bigcup_{U\in\mathcal{U}^{\mathsf{c}}}U=\left\{ 0,1\right\} \neq\left\{ 1\right\} =\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}.$$

This finishes the proof.

### Item 11: Interaction With Complements II

Omitted.

### Item 12: Interaction With Complements III

Omitted.

# Item 13: Interaction With Symmetric Differences

Let  $X = \{0, 1\}$ , let  $\mathcal{U} = \{\{0, 1\}\}\$ , and let  $\mathcal{V} = \{\{0\}, \{0, 1\}\}\$ . We have

$$\bigcup_{W \in \mathcal{U} \triangle \mathcal{V}} W = \bigcup_{W \in \{\{0\}\}} W$$
$$= \{0\},$$

whereas

$$\left(\bigcup_{U\in\mathcal{U}}U\right)\triangle\left(\bigcup_{V\in\mathcal{V}}V\right)=\left\{0,1\right\}\triangle\left\{0,1\right\}$$

$$=\emptyset$$
,

Thus we have

$$\bigcup_{W\in\mathcal{U}\triangle\mathcal{V}}W=\left\{0\right\}\neq\emptyset=\left(\bigcup_{U\in\mathcal{U}}U\right)\triangle\left(\bigcup_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

Item 14: Interaction With Internal Homs I

This is a repetition of Item 7 of Proposition 4.4.7.1.4 and is proved there.

Item 15: Interaction With Internal Homs II

This is a repetition of Item 8 of Proposition 4.4.7.1.4 and is proved there.

Item 16: Interaction With Internal Homs III

This is a repetition of Item 9 of Proposition 4.4.7.1.4 and is proved there.

Item 17: Interaction With Direct Images

This is a repetition of Item 3 of Proposition 4.6.1.1.5 and is proved there.

Item 18: Interaction With Inverse Images

This is a repetition of Item 3 of Proposition 4.6.2.1.3 and is proved there.

Item 19: Interaction With Codirect Images

This is a repetition of Item 3 of Proposition 4.6.3.1.7 and is proved there.

Item 20: Interaction With Intersections of Families I

We have

$$\bigcup_{U \in \bigcup_{A \in \mathcal{A}} A} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{c} \text{for each } U \in \bigcup_{A \in \mathcal{A}} A, \\ \text{we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{c} \text{for each } A \in \mathcal{A} \text{ and each } \\ U \in A, \text{ we have } x \in U \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{A \in \mathcal{A}} \left( \bigcap_{U \in A} U \right).$$

This finishes the proof.

### Item 21: Interaction With Intersections of Families II

Omitted.

#### 4.3.7 Intersections of Families of Subsets

Let *X* be a set and let  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

#### **DEFINITION 4.3.7.1.1** ► Intersections of Families of Subsets

The **intersection of**  $\mathcal{U}$  is the set  $\bigcap_{U \in \mathcal{U}} U$  defined by

$$\bigcap_{U \in \mathcal{U}} U \stackrel{\text{def}}{=} \left\{ x \in X \, \middle| \, \text{for each } U \in \mathcal{U}, \right\}.$$
we have  $x \in U$ 

#### PROPOSITION 4.3.7.1.2 ► PROPERTIES OF INTERSECTIONS OF FAMILIES OF SUBSETS

Let *X* be a set.

1. Functoriality. The assignment  $\mathcal{U} \mapsto \bigcap_{U \in \mathcal{U}} U$  defines a functor

$$\bigcap : (\mathcal{P}(\mathcal{P}(X)), \supset) \to (\mathcal{P}(X), \subset).$$

In particular, for each  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ , the following condition is satisfied:

(\*) If 
$$\mathcal{U} \subset \mathcal{V}$$
, then  $\bigcap_{V \in \mathcal{V}} V \subset \bigcap_{U \in \mathcal{U}} U$ .

2. Oplax Associativity. We have a natural transformation

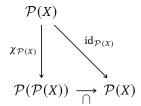
$$\begin{array}{ccc} \mathcal{P}(\mathcal{P}(\mathcal{P}(X))) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \, \cap} & \mathcal{P}(\mathcal{P}(X)) \\ \cap \star \mathrm{id}_{\mathcal{P}(X)} & & & & & & & & & & \\ \mathcal{P}(\mathcal{P}(X)) & & & & & & & & & & \\ \end{array}$$

with components

$$\bigcap_{A\in\mathcal{A}}\left(\bigcap_{U\in A}U\right)\subset\bigcap_{U\in\bigcap_{A\in\mathcal{A}}A}U$$

for each  $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$ .

3. Left Unitality. The diagram



commutes, i.e. we have

$$\bigcap_{V \in \{IJ\}} V = U$$

for each  $U \in \mathcal{P}(X)$ .

4. Oplax Right Unitality. The diagram

$$\begin{array}{c|c}
\mathcal{P}(X) & \operatorname{id}_{\mathcal{P}(X)} \\
 & \times & \\
\mathcal{P}(\mathcal{P}(X)) & \longrightarrow & \mathcal{P}(X)
\end{array}$$

does not commute in general, i.e. we may have

$$\bigcap_{\{x\}\in\chi_X(U)}\{x\}\neq U$$

in general, where  $U \in \mathcal{P}(X)$ . However, when U is nonempty, we have

$$\bigcap_{\{x\}\in\chi_X(U)}\{x\}\subset U.$$

5. Interaction With Unions I. The diagram

commutes, i.e. we have

$$\bigcap_{W \in \mathcal{U} \cup \mathcal{V}} W = \left(\bigcap_{U \in \mathcal{U}} U\right) \cap \left(\bigcap_{V \in \mathcal{V}} V\right)$$

for each  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

6. Interaction With Unions II. The diagram

commute, i.e. we have

$$U \cup \left(\bigcap_{V \in \mathcal{V}} V\right) = \bigcap_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cup V = \bigcap_{U \in \mathcal{U}} (U \cup V)$$

for each  $U, V \in \mathcal{P}(\mathcal{P}(X))$  and each  $U, V \in \mathcal{P}(X)$ .

7. Interaction With Intersections I. We have a natural transformation

with components

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\cap\left(\bigcap_{V\in\mathcal{V}}V\right)\subset\bigcap_{W\in\mathcal{U}\cap\mathcal{V}}W$$

for each  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

8. Interaction With Intersections II. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcap_{V \in \mathcal{V}} V\right) = \bigcap_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cup V = \bigcap_{U \in \mathcal{U}} (U \cup V)$$

for each  $U, V \in \mathcal{P}(\mathcal{P}(X))$  and each  $U, V \in \mathcal{P}(X)$ .

9. Interaction With Differences. The diagram

$$\begin{array}{ccc} \mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) & \stackrel{\backslash}{\longrightarrow} & \mathcal{P}(\mathcal{P}(X)) \\ & & & \swarrow & & \downarrow \cap \\ & \mathcal{P}(X) \times \mathcal{P}(X) & \stackrel{\backslash}{\longrightarrow} & \mathcal{P}(X), \end{array}$$

does not commute in general, i.e. we may have

$$\bigcap_{W\in\mathcal{U}\backslash\mathcal{V}}W\neq\left(\bigcap_{U\in\mathcal{U}}U\right)\backslash\left(\bigcap_{V\in\mathcal{V}}V\right)$$

in general, where  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

10. Interaction With Complements I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\text{op}} \xrightarrow{(-)^{c}} \mathcal{P}(\mathcal{P}(X))$$

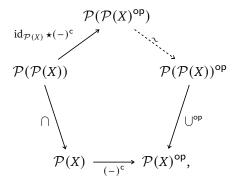
$$\uparrow^{\text{op}} \qquad \qquad \downarrow \\
\mathcal{P}(X)^{\text{op}} \xrightarrow{(-)^{c}} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W\in\mathcal{U}^\mathsf{c}}W\neq\bigcap_{U\in\mathcal{U}}U^\mathsf{c}$$

in general, where  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

11. Interaction With Complements II. The diagram

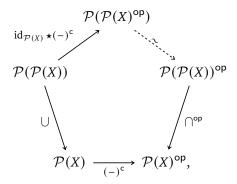


commutes, i.e. we have

$$\left(\bigcap_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

12. Interaction With Complements III. The diagram



commutes, i.e. we have

$$\left(\bigcup_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

13. Interaction With Symmetric Differences. The diagram

does not commute in general, i.e. we may have

$$\bigcap_{W \in \mathcal{U} \triangle \mathcal{V}} W \neq \left(\bigcap_{U \in \mathcal{U}} U\right) \triangle \left(\bigcap_{V \in \mathcal{V}} V\right)$$

in general, where  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

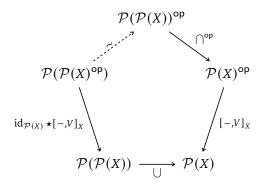
14. Interaction With Internal Homs I. The diagram

does not commute in general, i.e. we may have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcap_{U \in \mathcal{U}} U, \bigcap_{V \in \mathcal{V}} V\right]_X$$

in general, where  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

15. Interaction With Internal Homs II. The diagram



commutes, i.e. we have

$$\left[\bigcap_{U\in\mathcal{U}}U,V\right]_X=\bigcup_{U\in\mathcal{U}}[U,V]_X$$

for each  $U \in \mathcal{P}(\mathcal{P}(X))$  and each  $V \in \mathcal{P}(X)$ .

16. Interaction With Internal Homs III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcap}{\longrightarrow} & \mathcal{P}(X) \\
id_{\mathcal{P}(X)} \star [U,-]_X & & \downarrow [U,-]_X \\
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcap}{\longrightarrow} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U,\bigcap_{V\in\mathcal{V}}V\right]_X=\bigcap_{V\in\mathcal{V}}[U,V]_X$$

for each  $U \in \mathcal{P}(X)$  and each  $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

17. *Interaction With Direct Images.* Let  $f: X \to Y$  be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_!)_!} \mathcal{P}(\mathcal{P}(Y))$$

$$\cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcap_{U\in\mathcal{U}}f_!(U)=\bigcap_{V\in f_!(\mathcal{U})}V$$

for each  $\mathcal{U} \in \mathcal{P}(X)$ , where  $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$ .

18. *Interaction With Inverse Images*. Let  $f: X \to Y$  be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \bigcup_{f^{-1}} \bigcup_{f^{-1}} \bigcap_{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcap_{V \in \mathcal{V}} f^{-1}(V) = \bigcap_{U \in f^{-1}(\mathcal{U})} U$$

for each  $V \in \mathcal{P}(Y)$ , where  $f^{-1}(V) \stackrel{\text{def}}{=} (f^{-1})^{-1}(V)$ .

19. Interaction With Codirect Images. Let  $f: X \to Y$  be a map of sets. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_*)_*} & \mathcal{P}(\mathcal{P}(Y)) \\
& & \downarrow & & \downarrow & \\
\mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \mathcal{U}} f_*(U) = \bigcap_{V \in f_*(\mathcal{U})} V$$

for each  $\mathcal{U} \in \mathcal{P}(X)$ , where  $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$ .

20. Interaction With Unions of Families I. The diagram

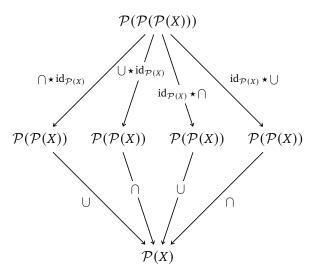
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \cap} & \mathcal{P}(\mathcal{P}(X)) \\
\downarrow \star \mathrm{id}_{\mathcal{P}(X)} & & & \downarrow \cap \\
\mathcal{P}(X) & \xrightarrow{\qquad \cap} & X
\end{array}$$

commutes, i.e. we have

$$\bigcup_{U \in \bigcup_{A \in A} A} U = \bigcap_{A \in \mathcal{A}} \left( \bigcap_{U \in A} U \right)$$

for each  $A \in \mathcal{P}(\mathcal{P}(X))$ .

21. *Interaction With Unions of Families II.* Let *X* be a set and consider the compositions

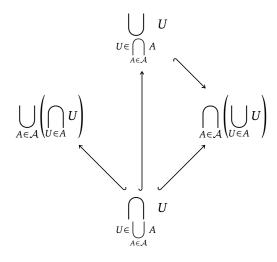


given by

$$A \mapsto \bigcup_{U \in \bigcap_{A \in \mathcal{A}} A} U, \quad A \mapsto \bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U,$$

$$A \mapsto \bigcup_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U\right), \quad A \mapsto \bigcap_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U\right)$$

for each  $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$ . We have the following inclusions:



All other possible inclusions fail to hold in general.

#### PROOF 4.3.7.1.3 ► PROOF OF PROPOSITION 4.3.6.1.2

# Item 1: Functoriality

Since  $\mathcal{P}(X)$  is posetal, it suffices to prove the condition  $(\star)$ . So let  $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$  with  $\mathcal{U} \subset \mathcal{V}$ . We claim that

$$\bigcap_{V\in\mathcal{V}}V\subset\bigcap_{U\in\mathcal{U}}U.$$

Indeed, if  $x \in \bigcap_{V \in \mathcal{V}} V$ , then  $x \in V$  for all  $V \in \mathcal{V}$ . But since  $\mathcal{U} \subset \mathcal{V}$ , it follows that  $x \in U$  for all  $U \in \mathcal{U}$  as well. Thus  $x \in \bigcap_{U \in \mathcal{U}} U$ , which gives our desired inclusion.

### Item 2: Oplax Associativity

We have

$$\bigcap_{A \in \mathcal{A}} \left( \bigcap_{U \in A} U \right) \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ for each } A \in \mathcal{A}, \\ \text{we have } x \in \bigcap_{U \in A} U \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ for each } A \in \mathcal{A} \text{ and each } \\ U \in A, \text{ we have } x \in U \right\}$$

$$= \left\{ x \in X \middle| \text{ for each } U \in \bigcup_{A \in \mathcal{A}} A, \\ \text{ we have } x \in U \right\}$$

$$\subset \left\{ x \in X \middle| \text{ for each } U \in \bigcap_{A \in \mathcal{A}} A, \\ \text{ we have } x \in U \right\}$$

$$\stackrel{\text{def}}{=} \bigcup_{A \in \mathcal{A}} U.$$

$$U \in \bigcap_{A \in \mathcal{A}} A$$

Since  $\mathcal{P}(X)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

### Item 3: Left Unitality

We have

$$\bigcap_{V \in \{U\}} V \stackrel{\text{def}}{=} \left\{ x \in X \mid \text{for each } V \in \{U\}, \right\}$$

$$= \left\{ x \in X \mid x \in U \right\}$$

$$= U.$$

This finishes the proof.

# Item 4: Oplax Right Unitality

If  $U = \emptyset$ , then we have

$$\bigcap_{\{u\} \in \chi_X(U)} \{u\} = \bigcap_{\{u\} \in \emptyset} \{u\}$$
$$= X,$$

so  $\bigcap_{\{u\} \in \chi_X(U)} \{u\} = X \neq \emptyset = U$ . When U is nonempty, we have two cases:

1. If *U* is a singleton, say  $U = \{u\}$ , we have

$$\bigcap_{\{u\}\in\chi_X(U)}\{u\}=\{u\}$$

$$\stackrel{\text{def}}{=}U.$$

2. If *U* contains at least two elements, we have

$$\bigcap_{\{u\} \in \chi_X(U)} \{u\} = \emptyset$$

$$\subset U.$$

This finishes the proof.

# Item 5: Interaction With Unions I

We have

$$\bigcap_{W \in \mathcal{U} \cup \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cup \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$
$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \text{ and each} \\ W \in \mathcal{V}, \text{ we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ for each } W \in \mathcal{U}, \right\} \\
\text{we have } x \in W \right\} \\
\cap \left\{ x \in X \middle| \text{ for each } W \in \mathcal{V}, \right\} \\
\stackrel{\text{def}}{=} \left( \bigcap_{W \in \mathcal{U}} W \right) \cap \left( \bigcap_{W \in \mathcal{V}} W \right) \\
= \left( \bigcap_{U \in \mathcal{U}} U \right) \cap \left( \bigcap_{V \in \mathcal{V}} V \right).$$

This finishes the proof.

# Item 6: Interaction With Unions II

Omitted.

### Item 7: Interaction With Intersections 1

We have

$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cap \left(\bigcap_{V \in \mathcal{V}} V\right) \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \mathcal{U}, \\ \text{we have } x \in U \end{array} \right\}$$

$$\cup \left\{ x \in X \middle| \begin{array}{l} \text{for each } V \in \mathcal{V}, \\ \text{we have } x \in V \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cap \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\subset \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cup \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{W \in \mathcal{U} \cap \mathcal{V}} W.$$

Since  $\mathcal{P}(X)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

### Item 8: Interaction With Intersections II

Omitted.

#### Item 9: Interaction With Differences

Let  $X = \{0, 1\}$ , let  $\mathcal{U} = \{\{0\}, \{0, 1\}\}\$ , and let  $\mathcal{V} = \{\{0\}\}\$ . We have

$$\bigcap_{W \in \mathcal{U} \setminus \mathcal{V}} U = \bigcap_{W \in \{\{0,1\}\}} W$$
$$= \{0,1\},$$

whereas

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\setminus\left(\bigcap_{V\in\mathcal{V}}V\right)=\left\{0\right\}\setminus\left\{0\right\}$$
$$=\emptyset.$$

Thus we have

$$\bigcap_{W\in\mathcal{U}\backslash\mathcal{V}}W=\left\{0,1\right\}\neq\emptyset=\left(\bigcap_{U\in\mathcal{U}}U\right)\backslash\left(\bigcap_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

# Item 10: Interaction With Complements I

Let  $X = \{0, 1\}$  and let  $\mathcal{U} = \{\{0\}\}$ . We have

$$\bigcap_{W \in \mathcal{U}^{c}} U = \bigcap_{W \in \{\emptyset, \{1\}, \{0,1\}\}} W$$
$$= \emptyset,$$

whereas

$$\bigcap_{U \in \mathcal{U}} U^{c} = \{0\}^{c}$$
$$= \{1\}.$$

Thus we have

$$\bigcap_{W\in\mathcal{U}^\mathsf{c}}U=\emptyset\neq\{1\}=\bigcap_{U\in\mathcal{U}}U^\mathsf{c}.$$

This finishes the proof.

# Item 11: Interaction With Complements II

This is a repetition of Item 12 of Proposition 4.3.6.1.2 and is proved there.

### Item 12: Interaction With Complements III

This is a repetition of Item 11 of Proposition 4.3.6.1.2 and is proved there.

#### Item 13: Interaction With Symmetric Differences

Let  $X = \{0, 1\}$ , let  $\mathcal{U} = \{\{0, 1\}\}\$ , and let  $\mathcal{V} = \{\{0\}, \{0, 1\}\}\$ . We have

$$\bigcap_{W \in \mathcal{U} \triangle \mathcal{V}} W = \bigcap_{W \in \{\{0\}\}} W$$
$$= \{0\},$$

whereas

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\triangle\left(\bigcap_{V\in\mathcal{V}}V\right)=\left\{0,1\right\}\triangle\left\{0\right\}$$
$$=\emptyset,$$

Thus we have

$$\bigcap_{W \in \mathcal{U} \triangle \mathcal{V}} W = \{0\} \neq \emptyset = \left(\bigcap_{U \in \mathcal{U}} U\right) \triangle \left(\bigcap_{V \in \mathcal{V}} V\right).$$

This finishes the proof.

### Item 14: Interaction With Internal Homs I

This is a repetition of Item 10 of Proposition 4.4.7.1.4 and is proved there.

### Item 15: Interaction With Internal Homs II

This is a repetition of Item 11 of Proposition 4.4.7.1.4 and is proved there.

### Item 16: Interaction With Internal Homs III

This is a repetition of Item 12 of Proposition 4.4.7.1.4 and is proved there.

### Item 17: Interaction With Direct Images

This is a repetition of Item 4 of Proposition 4.6.1.1.5 and is proved there.

#### Item 18: Interaction With Inverse Images

This is a repetition of Item 4 of Proposition 4.6.2.1.3 and is proved there.

### Item 19: Interaction With Codirect Images

This is a repetition of Item 4 of Proposition 4.6.3.1.7 and is proved there.

Item 20: Interaction With Unions of Families I

This is a repetition of Item 20 of Proposition 4.3.6.1.2 and is proved there.

# Item 21: Interaction With Unions of Families II

This is a repetition of Item 21 of Proposition 4.3.6.1.2 and is proved there.

# 4.3.8 Binary Unions

Let *X* be a set and let  $U, V \in \mathcal{P}(X)$ .

#### **DEFINITION 4.3.8.1.1** ► BINARY UNIONS

The **union of** *U* **and** *V* is the set  $U \cup V$  defined by

$$U \cup V \stackrel{\text{def}}{=} \bigcup_{z \in \{U, V\}} z$$

$$\stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in V\}.$$

#### PROPOSITION 4.3.8.1.2 ► PROPERTIES OF BINARY UNIONS

Let *X* be a set.

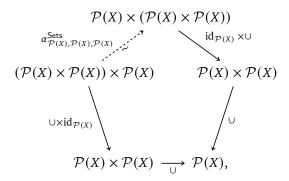
1. Functoriality. The assignments  $U, V, (U, V) \mapsto U \cup V$  define functors

$$\begin{array}{ll} U \cup -\colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ - \cup V \colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ -_1 \cup -_2 \colon (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset). \end{array}$$

In particular, the following statements hold for each  $U, V, A, B \in \mathcal{P}(X)$ :

- (a) If  $U \subset A$ , then  $U \cup V \subset A \cup V$ .
- (b) If  $V \subset B$ , then  $U \cup V \subset U \cup B$ .
- (c) If  $U \subset A$  and  $V \subset B$ , then  $U \cup V \subset A \cup B$ .

### 2. Associativity. The diagram

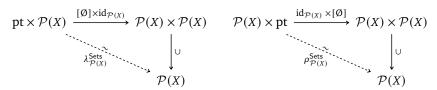


commutes, i.e. we have an equality of sets

$$(U \cup V) \cup W = U \cup (V \cup W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

3. Unitality. The diagrams



commute, i.e. we have equalities of sets

$$\emptyset \cup U = U$$
,  $U \cup \emptyset = U$ 

for each  $U \in \mathcal{P}(X)$ .

4. Commutativity. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\sigma_{\mathcal{P}(X),\mathcal{P}(X)}^{\mathsf{Sets}}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \cup \qquad \qquad \downarrow \cup$$

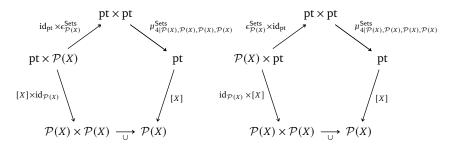
$$\mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cup V = V \cup U$$

for each  $U, V \in \mathcal{P}(X)$ .

### 5. Annihilation With X. The diagrams

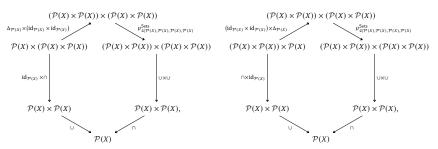


commute, i.e. we have equalities of sets

$$U \cup X = X,$$
$$X \cup V = X$$

for each  $U, V \in \mathcal{P}(X)$ .

# 6. Distributivity of Unions Over Intersections. The diagrams

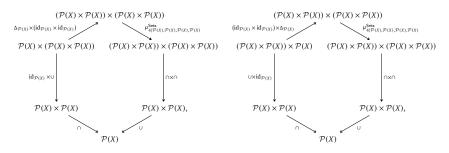


commute, i.e. we have equalities of sets

$$U \cup (V \cap W) = (U \cup V) \cap (U \cup W),$$
  
$$(U \cap V) \cup W = (U \cup W) \cap (V \cup W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

### 7. Distributivity of Intersections Over Unions. The diagrams

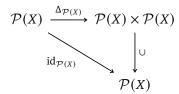


commute, i.e. we have equalities of sets

$$U \cap (V \cup W) = (U \cap V) \cup (U \cap W),$$
$$(U \cup V) \cap W = (U \cap W) \cup (V \cap W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

8. Idempotency. The diagram



commutes, i.e. we have an equality of sets

$$U \cup U = U$$

for each  $U \in \mathcal{P}(X)$ .

9. Via Intersections and Symmetric Differences. The diagram

$$(\mathcal{P}(X) \times \mathcal{P}(X)) \times (\mathcal{P}(X) \times \mathcal{P}(X)) \xrightarrow{\Delta \times \cap} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\Delta_{\mathcal{P}(X) \times \mathcal{P}(X)} / \Delta$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \longrightarrow \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cup V = (U \triangle V) \triangle (U \cap V)$$

for each  $U, V \in \mathcal{P}(X)$ .

10. Interaction With Characteristic Functions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each  $U, V \in \mathcal{P}(X)$ .

11. Interaction With Characteristic Functions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each  $U, V \in \mathcal{P}(X)$ .

12. *Interaction With Direct Images.* Let  $f: X \to Y$  be a function. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_! \times f_!} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\
\downarrow & & \downarrow & & \downarrow \\
\mathcal{P}(X) & \xrightarrow{f_!} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$

for each  $U, V \in \mathcal{P}(X)$ .

13. *Interaction With Inverse Images.* Let  $f: X \to Y$  be a function. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$

for each  $U, V \in \mathcal{P}(Y)$ .

14. *Interaction With Codirect Images.* Let  $f: X \to Y$  be a function. We have a natural transformation

with components

$$f_*(U) \cup f_*(V) \subset f_*(U \cup V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

15. *Interaction With Powersets and Semirings.* The quintuple  $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$  is an idempotent commutative semiring.

### PROOF 4.3.8.1.3 ► PROOF OF PROPOSITION 4.3.8.1.2

Item 1: Functoriality

See [Pro25an].

Item 2: Associativity

See [Pro25ba].

Item 3: Unitality

This follows from [Pro25bd] and Item 4.

Item 4: Commutativity

See [Pro25bb].

Item 5: Annihilation With *X* 

We have

$$U \cup X \stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in X\}$$
$$= \{x \in X \mid x \in X\},$$
$$= X$$

and

$$X \cup V \stackrel{\text{def}}{=} \{x \in X \mid x \in X \text{ or } x \in V\}$$
$$= \{x \in X \mid x \in X\}$$
$$= X.$$

This finishes the proof.

Item 6: Distributivity of Unions Over Intersections

See [Pro25az].

Item 7: Distributivity of Intersections Over Unions

See [Pro25aj].

Item 8: Idempotency

See [Pro25am].

Item 9: Via Intersections and Symmetric Differences

See [Pro25ay].

Item 10: Interaction With Characteristic Functions I

See [Pro25h].

Item 11: Interaction With Characteristic Functions II

See [Pro25h].

Item 12: Interaction With Direct Images

See [Pro25p].

Item 13: Interaction With Inverse Images

See [Pro25y].

Item 14: Interaction With Codirect Images

This is a repetition of Item 5 of Proposition 4.6.3.1.7 and is proved there.

Item 15: Interaction With Powersets and Semirings

This follows from Items 2 to 4 and 8 of this proposition and Items 3 to 6 and 8 of Proposition 4.3.9.1.2.

### 4.3.9 Binary Intersections

Let *X* be a set and let  $U, V \in \mathcal{P}(X)$ .

#### **DEFINITION 4.3.9.1.1** ► BINARY INTERSECTIONS

The **intersection of** *U* **and** *V* is the set  $U \cap V$  defined by

$$U \cap V \stackrel{\text{def}}{=} \bigcap_{z \in \{U, V\}} z$$

$$\stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in V\}.$$

#### PROPOSITION 4.3.9.1.2 ► PROPERTIES OF BINARY INTERSECTIONS

Let *X* be a set.

1. Functoriality. The assignments  $U, V, (U, V) \mapsto U \cap V$  define functors

$$\begin{array}{ll} U \cap -\colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ - \cap V \colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ -_1 \cap -_2 \colon (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset). \end{array}$$

In particular, the following statements hold for each  $U, V, A, B \in \mathcal{P}(X)$ :

- (a) If  $U \subset A$ , then  $U \cap V \subset A \cap V$ .
- (b) If  $V \subset B$ , then  $U \cap V \subset U \cap B$ .

- (c) If  $U \subset A$  and  $V \subset B$ , then  $U \cap V \subset A \cap B$ .
- 2. Adjointness. We have adjunctions

$$(U \cap - \dashv [U, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{U \cap -}{\downarrow}} \mathcal{P}(X),$$

$$(- \cap V \dashv [V, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{- \cap V}{\downarrow}} \mathcal{P}(X),$$

witnessed by bijections

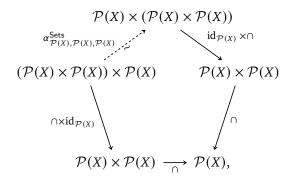
$$\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, [V, W]_X),$$
  
 $\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(V, [U, W]_X),$ 

natural in  $U, V, W \in \mathcal{P}(X)$ , where

$$[-1,-2]_X \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

is the bifunctor of Section 4.4.7. In particular, the following statements hold for each  $U, V, W \in \mathcal{P}(X)$ :

- (a) The following conditions are equivalent:
  - i. We have  $U \cap V \subset W$ .
  - ii. We have  $U \subset [V, W]_X$ .
- (b) The following conditions are equivalent:
  - i. We have  $U \cap V \subset W$ .
  - ii. We have  $V \subset [U, W]_X$ .
- 3. Associativity. The diagram

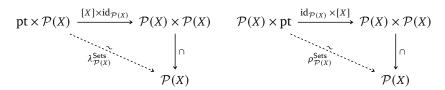


commutes, i.e. we have an equality of sets

$$(U \cap V) \cap W = U \cap (V \cap W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

4. Unitality. The diagrams



commute, i.e. we have equalities of sets

$$X \cap U = U,$$
$$U \cap X = U$$

for each  $U \in \mathcal{P}(X)$ .

5. Commutativity. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\sigma_{\mathcal{P}(X),\mathcal{P}(X)}^{\mathsf{Sets}}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\uparrow \cap \qquad \qquad \downarrow \cap$$

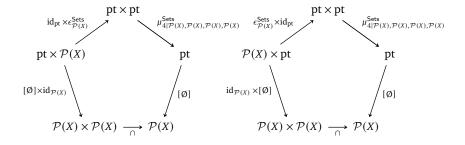
$$\mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cap V = V \cap U$$

for each  $U, V \in \mathcal{P}(X)$ .

6. Annihilation With the Empty Set. The diagrams

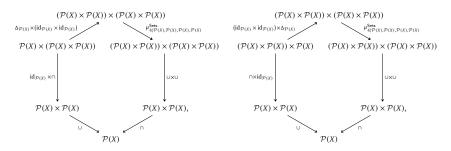


commute, i.e. we have equalities of sets

$$\emptyset \cap X = \emptyset$$
,  $X \cap \emptyset = \emptyset$ 

for each  $U \in \mathcal{P}(X)$ .

### 7. Distributivity of Unions Over Intersections. The diagrams

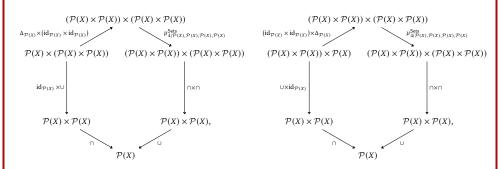


commute, i.e. we have equalities of sets

$$U \cup (V \cap W) = (U \cup V) \cap (U \cup W),$$
  
$$(U \cap V) \cup W = (U \cup W) \cap (V \cup W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

### 8. Distributivity of Intersections Over Unions. The diagrams



commute, i.e. we have equalities of sets

$$U \cap (V \cup W) = (U \cap V) \cup (U \cap W),$$
  
$$(U \cup V) \cap W = (U \cap W) \cup (V \cap W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

9. Idempotency. The diagram

$$\mathcal{P}(X) \xrightarrow{\Delta_{\mathcal{P}(X)}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \cap$$

$$\mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cap U = U$$

for each  $U \in \mathcal{P}(X)$ .

10. Interaction With Characteristic Functions I. We have

$$\chi_{U\cap V}=\chi_U\chi_V$$

for each  $U, V \in \mathcal{P}(X)$ .

11. Interaction With Characteristic Functions II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each  $U, V \in \mathcal{P}(X)$ .

12. *Interaction With Direct Images.* Let  $f: X \to Y$  be a function. We have a natural transformation

with components

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

13. Interaction With Inverse Images. Let  $f: X \to Y$  be a function. The diagram

$$\begin{array}{ccc} \mathcal{P}(Y) \times \mathcal{P}(Y) & \xrightarrow{f^{-1} \times f^{-1}} & \mathcal{P}(X) \times \mathcal{P}(X) \\ & & & & \downarrow \cap \\ & & & & \downarrow \cap \\ & & & \mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X) \end{array}$$

commutes, i.e. we have

$$f^{-1}(U \cap V) = f^{-1}(U) \cap f^{-1}(V)$$

for each  $U, V \in \mathcal{P}(Y)$ .

14. *Interaction With Codirect Images.* Let  $f: X \to Y$  be a function. The diagram

$$\begin{array}{ccccc} \mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_* \times f_*} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\ & & & & \downarrow \cap \\ & & & \downarrow \cap \\ & \mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y) \end{array}$$

commutes, i.e. we have

$$f_*(U) \cap f_*(V) = f_*(U \cap V)$$

for each  $U, V \in \mathcal{P}(X)$ .

- 15. *Interaction With Powersets and Monoids With Zero.* The quadruple  $((\mathcal{P}(X), \emptyset), \cap, X)$  is a commutative monoid with zero.
- 16. *Interaction With Powersets and Semirings.* The quintuple  $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$  is an idempotent commutative semiring.

# PROOF 4.3.9.1.3 ► PROOF OF PROPOSITION 4.3.9.1.2 Item 1: Functoriality See [Pro25al]. Item 2: Adjointness See [MSE 267469]. Item 3: Associativity See [Pro25r]. Item 4: Unitality This follows from [Pro25v] and Item 5. Item 5: Commutativity See [Pro25s]. Item 6: Annihilation With the Empty Set This follows from [Pro25t] and Item 5. Item 7: Distributivity of Unions Over Intersections See [Pro25az]. Item 8: Distributivity of Intersections Over Unions See [Pro25aj]. Item 9: Idempotency See [Pro25ak]. Item 10: Interaction With Characteristic Functions I See [Pro25e]. Item 11: Interaction With Characteristic Functions II See [Pro25e]. Item 12: Interaction With Direct Images See [Pro25n]. Item 13: Interaction With Inverse Images See [Pro25w]. Item 14: Interaction With Codirect Images

This is a repetition of Item 6 of Proposition 4.6.3.1.7 and is proved there.

### Item 15: Interaction With Powersets and Monoids With Zero

This follows from Items 3 to 6.

### Item 16: Interaction With Powersets and Semirings

This follows from Items 2 to 4 and 8 and Items 3 to 6 and 8 of Proposition 4.3.9.1.2.

### 4.3.10 Differences

Let *X* and *Y* be sets.

### **DEFINITION 4.3.10.1.1** ► **DIFFERENCES**

The **difference of** X **and** Y is the set  $X \setminus Y$  defined by

$$X \setminus Y \stackrel{\mathrm{def}}{=} \{ a \in X \mid a \notin Y \}.$$

### PROPOSITION 4.3.10.1.2 ► PROPERTIES OF DIFFERENCES

Let *X* be a set.

1. Functoriality. The assignments  $U, V, (U, V) \mapsto U \cap V$  define functors

$$\begin{array}{ccc} U \setminus -\colon & (\mathcal{P}(X), \supset) & \to (\mathcal{P}(X), \subset), \\ - \setminus V \colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ -_1 \setminus -_2 \colon (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \supset) \to (\mathcal{P}(X), \subset). \end{array}$$

In particular, the following statements hold for each  $U, V, A, B \in \mathcal{P}(X)$ :

- (a) If  $U \subset A$ , then  $U \setminus V \subset A \setminus V$ .
- (b) If  $V \subset B$ , then  $U \setminus B \subset U \setminus V$ .
- (c) If  $U \subset A$  and  $V \subset B$ , then  $U \setminus B \subset A \setminus V$ .
- 2. De Morgan's Laws. We have equalities of sets

$$X \setminus (U \cup V) = (X \setminus U) \cap (X \setminus V),$$
  
$$X \setminus (U \cap V) = (X \setminus U) \cup (X \setminus V)$$

for each  $U, V \in \mathcal{P}(X)$ .

3. Interaction With Unions I. We have equalities of sets

$$U \setminus (V \cup W) = (U \setminus V) \cap (U \setminus W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

4. Interaction With Unions II. We have equalities of sets

$$(U \setminus V) \cup W = (U \cup W) \setminus (V \setminus W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

5. Interaction With Unions III. We have equalities of sets

$$U \setminus (V \cup W) = (U \cup W) \setminus (V \cup W)$$
$$= (U \setminus V) \setminus W$$
$$= (U \setminus W) \setminus V$$

for each  $U, V, W \in \mathcal{P}(X)$ .

6. Interaction With Unions IV. We have equalities of sets

$$(U \cup V) \setminus W = (U \setminus W) \cup (V \setminus W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

7. Interaction With Intersections. We have equalities of sets

$$(U \setminus V) \cap W = (U \cap W) \setminus V$$
$$= U \cap (W \setminus V)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

8. *Interaction With Complements*. We have an equality of sets

$$U\setminus V=U\cap V^\mathsf{c}$$

for each  $U, V \in \mathcal{P}(X)$ .

9. Interaction With Symmetric Differences. We have an equality of sets

$$U \setminus V = U \triangle (U \cap V)$$

for each  $U, V \in \mathcal{P}(X)$ .

10. Triple Differences. We have

$$U \setminus (V \setminus W) = (U \cap W) \cup (U \setminus V)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

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11. Left Annihilation. We have

$$\emptyset \setminus U = \emptyset$$

for each  $U \in \mathcal{P}(X)$ .

12. Right Unitality. We have

$$U \setminus \emptyset = U$$

for each  $U \in \mathcal{P}(X)$ .

13. Right Annihilation. We have

$$U \setminus X = U$$

for each  $U \in \mathcal{P}(X)$ .

14. *Invertibility*. We have

$$U \setminus U = \emptyset$$

for each  $U \in \mathcal{P}(X)$ .

- 15. *Interaction With Containment*. The following conditions are equivalent:
  - (a) We have  $V \setminus U \subset W$ .
  - (b) We have  $V \setminus W \subset U$ .
- 16. Interaction With Characteristic Functions. We have

$$\chi_{U\setminus V} = \chi_U - \chi_{U\cap V}$$

for each  $U, V \in \mathcal{P}(X)$ .

17. Interaction With Direct Images. We have a natural transformation

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

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18. Interaction With Inverse Images. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\mathsf{op},-1} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad$$

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commutes, i.e. we have

$$f^{-1}(U \setminus V) = f^{-1}(U) \setminus f^{-1}(V)$$

for each  $U, V \in \mathcal{P}(X)$ .

19. Interaction With Codirect Images. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_{!}^{\mathsf{op}} \times f_{!}} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad$$

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

### PROOF 4.3.10.1.3 ► PROOF OF PROPOSITION 4.3.10.1.2

### Item 1: Functoriality

See [Pro25ad] and [Pro25ah].

Item 2: De Morgan's Laws

See [Pro25k].

Item 3: Interaction With Unions I

See [Pro251].

Item 4: Interaction With Unions II Omitted. Item 5: Interaction With Unions III See [Pro25ai]. Item 6: Interaction With Unions IV See [Pro25ac]. Item 7: Interaction With Intersections See [Pro25u]. Item 8: Interaction With Complements See [Pro25aa]. Item 9: Interaction With Symmetric Differences See [Pro25ab]. Item 10: Triple Differences See [Pro25ag]. Item 11: Left Annihilation Omitted. Item 12: Right Unitality See [Pro25ae]. Item 13: Right Annihilation Omitted. Item 14: Invertibility See [Pro25af]. Item 15: Interaction With Containment Omitted. Item 16: Interaction With Characteristic Functions See [Pro25f]. Item 17: Interaction With Direct Images See [Pro25o].

Item 18: Interaction With Inverse Images

See [Pro25x].

## 4.3.11 Complements

Let *X* be a set and let  $U \in \mathcal{P}(X)$ .

### **DEFINITION 4.3.11.1.1** ► COMPLEMENTS

The **complement of** U is the set  $U^{c}$  defined by

$$U^{c} \stackrel{\text{def}}{=} X \setminus U$$

$$\stackrel{\text{def}}{=} \{ a \in X \mid a \notin U \}.$$

### PROPOSITION 4.3.11.1.2 ► PROPERTIES OF COMPLEMENTS

Let *X* be a set.

1. Functoriality. The assignment  $U \mapsto U^c$  defines a functor

$$(-)^{c} : \mathcal{P}(X)^{op} \to \mathcal{P}(X).$$

In particular, the following statements hold for each  $U, V \in \mathcal{P}(X)$ :

- ( $\star$ ) If  $U \subset V$ , then  $V^{c} \subset U^{c}$ .
- 2. *De Morgan's Laws*. The diagrams

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)^{\mathsf{op}} \xrightarrow{\cup^{\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}} \qquad \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)^{\mathsf{op}} \xrightarrow{\cap^{\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}}$$

$$(-)^{\mathsf{c}} \times (-)^{\mathsf{c}} \downarrow \qquad \qquad \downarrow (-)^{\mathsf{c}} \qquad \qquad \downarrow (-)^{\mathsf{c}} \downarrow \qquad \qquad \downarrow (-)^{\mathsf{c}}$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\qquad \cap} \mathcal{P}(X) \qquad \qquad \mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\qquad \cup} \mathcal{P}(X)$$

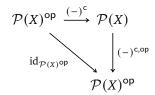
commute, i.e. we have equalities of sets

$$(U \cup V)^{c} = U^{c} \cap V^{c},$$
  

$$(U \cap V)^{c} = U^{c} \cup V^{c}$$

for each  $U, V \in \mathcal{P}(X)$ .

3. Involutority. The diagram



commutes, i.e. we have

$$(U^{\mathsf{c}})^{\mathsf{c}} = U$$

for each  $U \in \mathcal{P}(X)$ .

4. Interaction With Characteristic Functions. We have

$$\chi_{U^c} \equiv 1 - \chi_U \pmod{2}$$

for each  $U \in \mathcal{P}(X)$ .

5. Interaction With Direct Images. Let  $f: X \to Y$  be a function. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{f_*^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
\xrightarrow{(-)^{\mathsf{c}}} \qquad \qquad \downarrow^{(-)^{\mathsf{c}}} \\
\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_!(U^{\mathsf{c}}) = f_*(U)^{\mathsf{c}}$$

for each  $U \in \mathcal{P}(X)$ .

6. *Interaction With Inverse Images.* Let  $f: X \to Y$  be a function. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \xrightarrow{f^{-1},\mathsf{op}} \mathcal{P}(X)^{\mathsf{op}} \\
\xrightarrow{(-)^{\mathsf{c}}} \qquad \qquad \downarrow^{(-)^{\mathsf{c}}} \\
\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U^{c}) = f^{-1}(U)^{c}$$

for each  $U \in \mathcal{P}(X)$ .

7. *Interaction With Codirect Images.* Let  $f: X \to Y$  be a function. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{f_!^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} 
\xrightarrow[f_*]{} \mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U^{\mathsf{c}}) = f_!(U)^{\mathsf{c}}$$

for each  $U \in \mathcal{P}(X)$ .

# PROOF 4.3.11.1.3 ► PROOF OF PROPOSITION 4.3.11.1.2

Item 1: Functoriality

This follows from Item 1 of Proposition 4.3.10.1.2.

Item 2: De Morgan's Laws

See [Pro25k].

Item 3: Involutority

See [Pro25i].

Item 4: Interaction With Characteristic Functions

Omitted.

Item 5: Interaction With Direct Images

This is a repetition of Item 8 of Proposition 4.6.1.1.5 and is proved there.

Item 6: Interaction With Inverse Images

This is a repetition of Item 8 of Proposition 4.6.2.1.3 and is proved there.

Item 7: Interaction With Codirect Images

This is a repetition of Item 7 of Proposition 4.6.3.1.7 and is proved there.

### 4.3.12 Symmetric Differences

Let *X* be a set and let  $U, V \in \mathcal{P}(X)$ .

### **DEFINITION 4.3.12.1.1** ► SYMMETRIC DIFFERENCES

The **symmetric difference of** *U* **and** *V* is the set  $U \triangle V$  defined by  $^{\mathbf{1}}$ 

$$U \triangle V \stackrel{\text{def}}{=} (U \setminus V) \cup (V \setminus U).$$

 $^1{\it Illustration}:$ 



### PROPOSITION 4.3.12.1.2 ► PROPERTIES OF SYMMETRIC DIFFERENCES

Let *X* be a set.

1. Lack of Functoriality. The assignment  $(U,V) \mapsto U \triangle V$  does not in general define functors

$$\begin{array}{ll} U \bigtriangleup - \colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ - \bigtriangleup V \colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ -_1 \bigtriangleup -_2 \colon (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset). \end{array}$$

2. Via Unions and Intersections. We have

$$U \triangle V = (U \cup V) \setminus (U \cap V)$$

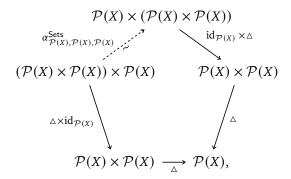
for each  $U, V \in \mathcal{P}(X)$ , as in the Venn diagram

$$\boxed{\bigcup_{U \triangle V}} = \boxed{\bigcup_{U \cup V}} \setminus \boxed{\bigcup_{U \cap V}}.$$

3. Symmetric Differences of Disjoint Sets. If U and V are disjoint, then we have

$$U \triangle V = U \cup V$$
.

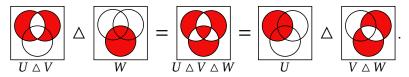
### 4. Associativity. The diagram



commutes, i.e. we have

$$(U \triangle V) \triangle W = U \triangle (V \triangle W)$$

for each  $U, V, W \in \mathcal{P}(X)$ , as in the Venn diagram



### 5. Unitality. The diagrams

$$\operatorname{pt} \times \mathcal{P}(X) \xrightarrow{[\emptyset] \times \operatorname{id}_{\mathcal{P}(X)}} \mathcal{P}(X) \times \mathcal{P}(X) \qquad \mathcal{P}(X) \times \operatorname{pt} \xrightarrow{\operatorname{id}_{\mathcal{P}(X)} \times [\emptyset]} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow^{\Delta} \qquad \qquad \downarrow^{\Delta} \qquad \qquad \downarrow^{\Delta}$$

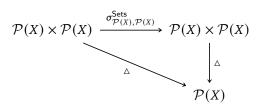
$$\mathcal{P}(X) \qquad \qquad \mathcal{P}(X)$$

commute, i.e. we have

$$U \triangle \emptyset = U,$$
  
$$\emptyset \triangle U = U$$

for each  $U \in \mathcal{P}(X)$ .

### 6. Commutativity. The diagram



commutes, i.e. we have

$$U \mathbin{\vartriangle} V = V \mathbin{\vartriangle} U$$

for each  $U, V \in \mathcal{P}(X)$ .

7. Invertibility. We have

$$U \triangle U = \emptyset$$

for each  $U \in \mathcal{P}(X)$ .

8. Interaction With Unions. We have

$$(U \triangle V) \cup (V \triangle T) = (U \cup V \cup W) \setminus (U \cap V \cap W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

9. Interaction With Complements I. We have

$$U \triangle U^{c} = X$$

for each  $U \in \mathcal{P}(X)$ .

10. Interaction With Complements II. We have

$$U \triangle X = U^{c},$$
  
 $X \triangle U = U^{c}$ 

for each  $U \in \mathcal{P}(X)$ .

11. Interaction With Complements III. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\Delta} \mathcal{P}(X) 
\xrightarrow{(-)^{c} \times (-)^{c}} \downarrow \qquad \qquad \downarrow^{(-)^{c}} 
\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\Delta} \mathcal{P}(X)$$

commutes, i.e. we have

$$U^{c} \triangle V^{c} = U \triangle V$$

for each  $U, V \in \mathcal{P}(X)$ .

12. "Transitivity". We have

$$(U \triangle V) \triangle (V \triangle W) = U \triangle W$$

for each  $U, V, W \in \mathcal{P}(X)$ .

13. The Triangle Inequality for Symmetric Differences. We have

$$U \triangle W \subset U \triangle V \cup V \triangle W$$

for each  $U, V, W \in \mathcal{P}(X)$ .

14. Distributivity Over Intersections. We have

$$U \cap (V \triangle W) = (U \cap V) \triangle (U \cap W),$$
  
$$(U \triangle V) \cap W = (U \cap W) \triangle (V \cap W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

15. Interaction With Characteristic Functions. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \pmod{2}$$

for each  $U, V \in \mathcal{P}(X)$ .

16. Bijectivity. Given  $U, V \in \mathcal{P}(X)$ , the maps

$$U \triangle -: \mathcal{P}(X) \to \mathcal{P}(X),$$
  
 $- \triangle V: \mathcal{P}(X) \to \mathcal{P}(X)$ 

are bijections with inverses given by

$$(U \triangle -)^{-1} = - \cup (U \cap -),$$
  
 $(- \triangle V)^{-1} = - \cup (V \cap -).$ 

Moreover, the map

$$\mathcal{P}(X) \longrightarrow \mathcal{P}(X)$$

$$C \longmapsto C \triangle (U \triangle V)$$

is a bijection of  $\mathcal{P}(X)$  onto itself sending U to V and V to U.

- 17. *Interaction With Powersets and Groups.* Let *X* be a set.
  - (a) The quadruple  $(\mathcal{P}(X), \triangle, \emptyset, \mathrm{id}_{\mathcal{P}(X)})$  is an abelian group.<sup>1</sup>
  - (b) Every element of  $\mathcal{P}(X)$  has order 2 with respect to  $\triangle$ , and thus  $\mathcal{P}(X)$  is a *Boolean group* (i.e. an abelian 2-group).
- 18. *Interaction With Powersets and Vector Spaces I.* The pair  $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$  consisting of
  - The group  $\mathcal{P}(X)$  of Item 17;
  - The map  $\alpha_{\mathcal{P}(X)} : \mathbb{F}_2 \times \mathcal{P}(X) \to \mathcal{P}(X)$  defined by

$$0 \cdot U \stackrel{\text{def}}{=} \emptyset,$$
$$1 \cdot U \stackrel{\text{def}}{=} U;$$

is an  $\mathbb{F}_2$ -vector space.

- 19. *Interaction With Powersets and Vector Spaces II.* If *X* is finite, then:
  - (a) The set of singletons sets on the elements of X forms a basis for the  $\mathbb{F}_2$ -vector space  $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$  of Item 18.
  - (b) We have

$$\dim(\mathcal{P}(X)) = \#X.$$

- 20. *Interaction With Powersets and Rings*. The quintuple  $(\mathcal{P}(X), \triangle, \cap, \emptyset, X)$  is a commutative ring.<sup>2</sup>
- 21. Interaction With Direct Images. We have a natural transformation

with components

$$f_!(U) \triangle f_!(V) \subset f_!(U \triangle V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

22. Interaction With Inverse Images. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\mathsf{op},-1} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

i.e. we have

$$f^{-1}(U) \triangle f^{-1}(V) = f^{-1}(U \triangle V)$$

for each  $U, V \in \mathcal{P}(Y)$ .

23. Interaction With Codirect Images. We have a natural transformation

with components

$$f_*(U \triangle V) \subset f_*(U) \triangle f_*(V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

i. When  $X = \emptyset$ , we have an isomorphism of groups between  $\mathcal{P}(\emptyset)$  and the trivial group:

$$(\mathcal{P}(\emptyset), \triangle, \emptyset, id_{\mathcal{P}(\emptyset)}) \cong pt.$$

ii. When X = pt, we have an isomorphism of groups between  $\mathcal{P}(pt)$  and  $\mathbb{Z}_{/2}$ :

$$(\mathcal{P}(\mathsf{pt}), \triangle, \emptyset, \mathrm{id}_{\mathcal{P}(\mathsf{pt})}) \cong \mathbb{Z}_{/2}.$$

iii. When  $X = \{0, 1\}$ , we have an isomorphism of groups between  $\mathcal{P}(\{0, 1\})$  and  $\mathbb{Z}_{/2} \times \mathbb{Z}_{/2}$ :

$$(\mathcal{P}(\{0,1\}), \triangle, \emptyset, id_{\mathcal{P}(\{0,1\})}) \cong \mathbb{Z}_{/2} \times \mathbb{Z}_{/2}.$$

Warning: The analogous statement replacing intersections by unions (i.e. that the quintuple  $(\mathcal{P}(X), \Delta, \cup, \emptyset, X)$  is a ring) is false, however. See [Pro25aw] for a proof.

<sup>&</sup>lt;sup>1</sup>Here are some examples:

### PROOF 4.3.12.1.3 ► PROOF OF PROPOSITION 4.3.12.1.2

Item 1: Lack of Functoriality

Omitted.

Item 2: Via Unions and Intersections

See [Pro25m].

Item 3: Symmetric Differences of Disjoint Sets

Since *U* and *V* are disjoint, we have  $U \cap V = \emptyset$ , and therefore we have

$$U \triangle V = (U \cup V) \setminus (U \cap V)$$
$$= (U \cup V) \setminus \emptyset$$
$$= U \cup V,$$

where we've used Item 2 and Item 12 of Proposition 4.3.10.1.2.

Item 4: Associativity

See [Pro25ao].

Item 5: Unitality

This follows from Item 6 and [Pro25at].

Item 6: Commutativity

See [Pro25ap].

Item 7: Invertibility

See [Pro25av].

Item 8: Interaction With Unions

See [Pro25bc].

Item 9: Interaction With Complements I

See [Pro25as].

Item 10: Interaction With Complements II

This follows from Item 6 and [Pro25ax].

Item 11: Interaction With Complements III

See [Pro25aq].

Item 12: "Transitivity"

We have

```
(U \triangle V) \triangle (V \triangle W) = U \triangle (V \triangle (V \triangle W))  (by Item 4)
= U \triangle ((V \triangle V) \triangle W) (by Item 4)
= U \triangle (\emptyset \triangle W) (by Item 7)
= U \triangle W. (by Item 5)
```

This finishes the proof.

Item 13: The Triangle Inequality for Symmetric Differences

This follows from Items 2 and 12.

Item 14: Distributivity Over Intersections

See [Pro25q].

Item 15: Interaction With Characteristic Functions

See [Pro25g].

Item 16: Bijectivity

Omitted.

Item 17: Interaction With Powersets and Groups

Item 17a follows from Items 4 to 7, while Item 17b follows from Item 7.1

Item 18: Interaction With Powersets and Vector Spaces I

See [MSE 2719059].

Item 19: Interaction With Powersets and Vector Spaces II

See [MSE 2719059].

Item 20: Interaction With Powersets and Rings

This follows from Items 6 and 15 of Proposition 4.3.9.1.2 and Items 14 and 17.<sup>2</sup>

Item 21: Interaction With Direct Images

This is a repetition of Item 9 of Proposition 4.6.1.1.5 and is proved there.

Item 22: Interaction With Inverse Images

This is a repetition of Item 9 of Proposition 4.6.2.1.3 and is proved there.

Item 23: Interaction With Codirect Images

This is a repetition of Item 8 of Proposition 4.6.3.1.7 and is proved there.

<sup>&</sup>lt;sup>1</sup>Reference: [Pro25ar].

<sup>&</sup>lt;sup>2</sup>Reference: [Pro25au].

### 4.4 Powersets

### 4.4.1 Foundations

Let *X* be a set.

### **DEFINITION 4.4.1.1.1** ▶ POWERSETS

The **powerset of** X is the set  $\mathcal{P}(X)$  defined by

$$\mathcal{P}(X) \stackrel{\text{def}}{=} \{ U \in P \mid U \subset X \},\,$$

where P is the set in the axiom of powerset, ?? of ??.

### REMARK 4.4.1.1.2 ▶ POWERSETS AS DECATEGORIFICATIONS OF CO/PRESHEAF CATEGORIES

Under the analogy that  $\{t, f\}$  should be the (-1)-categorical analogue of Sets, we may view the powerset of a set as a decategorification of the category of presheaves of a category (or of the category of copresheaves):

• The powerset of a set *X* is equivalently (Item 2 of Proposition 4.5.1.1.4) the set

$$Sets(X, \{t, f\})$$

of functions from X to the set  $\{t, f\}$  of classical truth values.

• The category of presheaves on a category *C* is the category

$$\operatorname{\mathsf{Fun}}(\mathcal{C}^{\operatorname{\mathsf{op}}},\operatorname{\mathsf{Sets}})$$

of functors from  $C^{op}$  to the category Sets of sets.

### NOTATION 4.4.1.1.3 ► FURTHER NOTATION FOR POWERSETS

Let *X* be a set.

- 1. We write  $\mathcal{P}_0(X)$  for the set of nonempty subsets of X.
- 2. We write  $\mathcal{P}_{fin}(X)$  for the set of finite subsets of X.

### PROPOSITION 4.4.1.1.4 ► ELEMENTARY PROPERTIES OF POWERSETS

Let *X* be a set.

- 1. *Co/Completeness*. The (posetal) category (associated to)  $(\mathcal{P}(X), \subset)$  is complete and cocomplete:
  - (a) *Products*. The products in  $\mathcal{P}(X)$  are given by intersection of subsets.

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- (b) *Coproducts*. The coproducts in  $\mathcal{P}(X)$  are given by union of subsets.
- (c) *Co/Equalisers*. Being a posetal category,  $\mathcal{P}(X)$  only has at most one morphisms between any two objects, so co/equalisers are trivial.
- 2. *Cartesian Closedness*. The category  $\mathcal{P}(X)$  is Cartesian closed.
- 3. Powersets as Sets of Relations. We have bijections

$$\mathcal{P}(X) \cong \text{Rel}(\text{pt}, X),$$
  
 $\mathcal{P}(X) \cong \text{Rel}(X, \text{pt}),$ 

natural in  $X \in Obi(Sets)$ .

4. Interaction With Products I. The map

$$\mathcal{P}(X) \times \mathcal{P}(Y) \longrightarrow \mathcal{P}(X \coprod Y)$$
$$(U, V) \longmapsto U \cup V$$

is an isomorphism of sets, natural in  $X,Y \in \text{Obj}(\mathsf{Sets})$  with respect to each of the functor structures  $\mathcal{P}_!, \mathcal{P}^{-1}$ , and  $\mathcal{P}_*$  on  $\mathcal{P}$  of Proposition 4.4.2.1.1. Moreover, this makes each of  $\mathcal{P}_!, \mathcal{P}^{-1}$ , and  $\mathcal{P}_*$  into a symmetric monoidal functor.

5. Interaction With Products II. The map

$$\mathcal{P}(X) \times \mathcal{P}(Y) \longrightarrow \mathcal{P}(X \coprod Y)$$
$$(U, V) \longmapsto U \boxtimes_{X \times Y} V,$$

where<sup>1</sup>

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \{(u, v) \in X \times Y \mid u \in U \text{ and } v \in V\}$$

is an inclusion of sets, natural in  $X,Y \in \text{Obj}(\mathsf{Sets})$  with respect to each of the functor structures  $\mathcal{P}_!$ ,  $\mathcal{P}^{-1}$ , and  $\mathcal{P}_*$  on  $\mathcal{P}$  of Proposition 4.4.2.1.1. Moreover, this makes each of  $\mathcal{P}_!$ ,  $\mathcal{P}^{-1}$ , and  $\mathcal{P}_*$  into a symmetric monoidal functor.

### 6. Interaction With Products III. We have an isomorphism

$$\mathcal{P}(X) \otimes \mathcal{P}(Y) \cong \mathcal{P}(X \times Y),$$

natural in  $X, Y \in \text{Obj}(\mathsf{Sets})$  with respect to each of the functor structures  $\mathcal{P}_!$ ,  $\mathcal{P}^{-1}$ , and  $\mathcal{P}_*$  on  $\mathcal{P}$  of Proposition 4.4.2.1.1, where  $\otimes$  denotes the tensor product of suplattices of ??. Moreover, this makes each of  $\mathcal{P}_!$ ,  $\mathcal{P}^{-1}$ , and  $\mathcal{P}_*$  into a symmetric monoidal functor.

### PROOF 4.4.1.1.5 ► PROOF OF PROPOSITION 4.4.1.1.4

### Item 1: Co/Completeness

Omitted.

### Item 2: Cartesian Closedness

See Section 4.4.7.

### Item 3: Powersets as Sets of Relations

Indeed, we have

$$Rel(pt, X) \stackrel{\text{def}}{=} \mathcal{P}(pt \times X)$$
$$\cong \mathcal{P}(X)$$

and

$$Rel(X, pt) \stackrel{\text{def}}{=} \mathcal{P}(X \times pt)$$

$$\cong \mathcal{P}(X),$$

where we have used Item 5 of Proposition 4.1.3.1.4.

### Item 4: Interaction With Products I

The inverse of the map in the statement is the map

$$\Phi \colon \mathcal{P}(X \coprod Y) \to \mathcal{P}(X) \times \mathcal{P}(Y)$$

defined by

$$\Phi(S) \stackrel{\text{def}}{=} (S_X, S_Y)$$

<sup>&</sup>lt;sup>1</sup>The set  $U \boxtimes_{X \times Y} V$  is usually denoted simply  $U \times V$ . Here we denote it in this somewhat weird way to highlight the similarity to external tensor products in six-functor formalisms (see also Section 4.6.4).

for each  $S \in \mathcal{P}(X \coprod Y)$ , where

$$S_X \stackrel{\text{def}}{=} \{ x \in X \mid (0, x) \in S \}$$
  
 $S_Y \stackrel{\text{def}}{=} \{ y \in Y \mid (1, y) \in S \}.$ 

The rest of the proof is omitted.

### Item 5: Interaction With Products II

Omitted.

Item 6: Interaction With Products III

Omitted.

### 4.4.2 Functoriality of Powersets

### PROPOSITION 4.4.2.1.1 ► FUNCTORIALITY OF POWERSETS

Let *X* be a set.

1. Functoriality I. The assignment  $X \mapsto \mathcal{P}(X)$  defines a functor

$$\mathcal{P}_!$$
: Sets  $\to$  Sets,

where

• *Action on Objects*. For each  $A \in Obj(Sets)$ , we have

$$\mathcal{P}_{1}(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

• *Action on Morphisms*. For each  $A, B \in Obj(Sets)$ , the action on morphisms

$$\mathcal{P}_{*|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of  $\mathcal{P}_!$  at (A, B) is the map defined by sending a map of sets  $f \colon A \to B$  to the map

$$\mathcal{P}_!(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_!(f) \stackrel{\mathrm{def}}{=} f_!,$$

as in Definition 4.6.1.1.1.

2. Functoriality II. The assignment  $X \mapsto \mathcal{P}(X)$  defines a functor

$$\mathcal{P}^{-1}$$
: Sets<sup>op</sup>  $\rightarrow$  Sets.

where

• Action on Objects. For each  $A \in Obj(Sets)$ , we have

$$\mathcal{P}^{-1}(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

• *Action on Morphisms*. For each  $A, B \in Obj(Sets)$ , the action on morphisms

$$\mathcal{P}_{A,B}^{-1} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(B),\mathcal{P}(A))$$

of  $\mathcal{P}^{-1}$  at (A, B) is the map defined by sending a map of sets  $f: A \to B$  to the map

$$\mathcal{P}^{-1}(f) \colon \mathcal{P}(B) \to \mathcal{P}(A)$$

defined by

$$\mathcal{P}^{-1}(f) \stackrel{\text{def}}{=} f^{-1}$$
,

as in Definition 4.6.2.1.1.

3. Functoriality III. The assignment  $X \mapsto \mathcal{P}(X)$  defines a functor

$$\mathcal{P}_*$$
: Sets  $\rightarrow$  Sets,

where

• Action on Objects. For each  $A \in Obj(Sets)$ , we have

$$\mathcal{P}_*(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

• *Action on Morphisms*. For each  $A, B \in Obj(\mathsf{Sets})$ , the action on morphisms

$$\mathcal{P}_{!|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of  $\mathcal{P}_*$  at (A,B) is the map defined by sending a map of sets  $f\colon A\to B$  to the map

$$\mathcal{P}_*(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_*(f) \stackrel{\mathrm{def}}{=} f_*,$$

as in Definition 4.6.3.1.1.

# PROOF 4.4.2.1.2 ➤ PROOF OF PROPOSITION 4.4.2.1.1 Item 1: Functoriality I This follows from Items 3 and 4 of Proposition 4.6.1.1.7. Item 2: Functoriality II This follows from Items 3 and 4 of Proposition 4.6.2.1.5. Item 3: Functoriality III This follows from Items 3 and 4 of Proposition 4.6.3.1.9.

## 4.4.3 Adjointness of Powersets I

### PROPOSITION 4.4.3.1.1 ► ADJOINTNESS OF POWERSETS I

We have an adjunction

$$\left(\mathcal{P}^{-1} \dashv \mathcal{P}^{-1,\mathsf{op}}\right)$$
: Sets $\overset{\mathcal{P}^{-1}}{\underset{\mathcal{P}^{-1,\mathsf{op}}}{\smile}}$  Sets,

witnessed by a bijection

$$\underbrace{\mathsf{Sets}^\mathsf{op}(\mathcal{P}(X),Y)}_{\substack{\mathsf{def}\\ = \mathsf{Sets}(Y,\mathcal{P}(X))}} \cong \mathsf{Sets}(X,\mathcal{P}(Y)),$$

natural in  $X \in \text{Obj}(\mathsf{Sets})$  and  $Y \in \text{Obj}(\mathsf{Sets}^{\mathsf{op}})$ .

### PROOF 4.4.3.1.2 ► PROOF OF PROPOSITION 4.4.3.1.1

We have

where all bijections are natural in A and B.

<sup>1</sup>Here we are using Item 3 of Proposition 4.5.1.1.4.

### 4.4.4 Adjointness of Powersets II

### PROPOSITION 4.4.4.1.1 ► ADJOINTNESS OF POWERSETS II

We have an adjunction

$$(Gr \dashv \mathcal{P}_!)$$
: Sets  $\underbrace{\overset{Gr}{\downarrow}}_{\mathcal{P}_!}$  Rel,

witnessed by a bijection of sets

$$Rel(Gr(X), Y) \cong Sets(X, \mathcal{P}(Y))$$

natural in  $X \in \text{Obj}(\mathsf{Sets})$  and  $Y \in \text{Obj}(\mathsf{Rel})$ , where Gr is the graph functor of Relations, Item 1 of Proposition 8.2.2.1.2 and  $\mathcal{P}_!$  is the functor of Relations, Proposition 8.7.5.1.1.

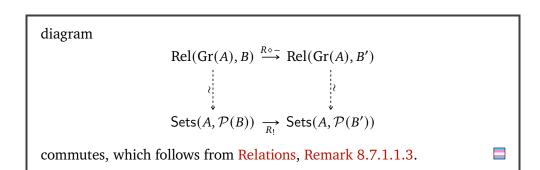
### PROOF 4.4.4.1.2 ► PROOF OF PROPOSITION 4.4.4.1.1

We have

```
Rel(Gr(A), B) \cong \mathcal{P}(A \times B)
\cong Sets(A \times B, \{t, f\}) \qquad \text{(by Item 2 of Proposition 4.5.1.1.4)}
\cong Sets(A, Sets(B, \{t, f\})) \qquad \text{(by Item 2 of Proposition 4.1.3.1.4)}
\cong Sets(A, \mathcal{P}(B)), \qquad \text{(by Item 2 of Proposition 4.5.1.1.4)}
```

where all bijections are natural in A, (where we are using Item 3 of Proposition 4.5.1.1.4). Explicitly, this isomorphism is given by sending a relation  $R: Gr(A) \to B$  to the map  $R^{\dagger}: A \to \mathcal{P}(B)$  sending a to the subset R(a) of B, as in Relations, Definition 8.1.1.1.1.

Naturality in *B* is then the statement that given a relation  $R: B \to B'$ , the



## 4.4.5 Powersets as Free Cocompletions

Let *X* be a set.

### PROPOSITION 4.4.5.1.1 ➤ POWERSETS AS FREE COCOMPLETIONS: UNIVERSAL PROPERTY

The pair  $(\mathcal{P}(X), \chi_{(-)})$  consisting of

- The powerset  $(\mathcal{P}(X), \subset)$  of X of Definition 4.4.1.1.1;
- The characteristic embedding  $\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$  of X into  $\mathcal{P}(X)$  of Definition 4.5.4.1.1;

satisfies the following universal property:

- $(\star)$  Given another pair (Y, f) consisting of
  - A suplattice (Y, ≤);
  - **–** A function f: X → Y;

there exists a unique morphism of suplattices

$$(\mathcal{P}(X),\subset) \xrightarrow{\exists !} (Y,\preceq)$$

making the diagram



commute.

This is a rephrasing of Proposition 4.4.5.1.3, which we prove below.<sup>1</sup>



<sup>1</sup>Here we only remark that the unique morphism of suplattices in the statement is given by the left Kan extension  $\text{Lan}_{\chi_X}(f)$  of f along  $\chi_X$ .

### PROPOSITION 4.4.5.1.3 ➤ POWERSETS AS FREE COCOMPLETIONS: ADJOINTNESS

We have an adjunction

$$(\mathcal{P} \dashv \overline{\mathbb{k}})$$
: Sets  $\underbrace{\overset{\mathcal{P}}{\vdash}}_{\overline{\mathbb{k}}}$  SupLat,

witnessed by a bijection

$$\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))\cong\mathsf{Sets}(X,Y),$$

natural in  $X \in \text{Obj}(\mathsf{Sets})$  and  $(Y, \preceq) \in \text{Obj}(\mathsf{SupLat})$ , where:

- The category SupLat is the category of suplattices of ??.
- The map

$$\chi_X^* \colon \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

witnessing the above bijection is defined by

$$\chi_X^*(f) \stackrel{\text{def}}{=} f \circ \chi_X,$$

i.e. by sending a morphism of suplattices  $f \colon \mathcal{P}(X) \to Y$  to the composition

$$X \stackrel{\chi_X}{\longleftrightarrow} \mathcal{P}(X) \stackrel{f}{\longrightarrow} Y.$$

The map

$$\operatorname{Lan}_{\chi_X} \colon \mathsf{Sets}(X,Y) \to \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$$

witnessing the above bijection is given by sending a function  $f: X \to Y$  to its left Kan extension along  $\chi_X$ ,

$$\operatorname{Lan}_{\chi_X}(f) \colon \mathcal{P}(X) \to Y, \qquad \chi_X = \left\{ \begin{array}{c} \mathcal{P}(X) \\ \operatorname{Lan}_{\chi_X}(f) \end{array} \right.$$

$$X \xrightarrow{f} Y.$$

Moreover, invoking the bijection  $\mathcal{P}(X) \cong \operatorname{Sets}(X, \{t, f\})$  of Item 2 of Proposition 4.5.1.1.4,  $\operatorname{Lan}_{\gamma_X}(f)$  can be explicitly computed by

$$\begin{aligned} \left[ \operatorname{Lan}_{\chi_{X}}(f) \right] (U) &= \int_{x \in X}^{x \in X} \chi_{\mathcal{D}(X)}(\chi_{x}, U) \odot f(x) \\ &= \int_{x \in X}^{x \in X} \chi_{U}(x) \odot f(x) \\ &= \bigvee_{x \in X} (\chi_{U}(x) \odot f(x)) \\ &= \left( \bigvee_{x \in U} (\chi_{U}(x) \odot f(x)) \right) \vee \left( \bigvee_{x \in U^{c}} (\chi_{U}(x) \odot f(x)) \right) \\ &= \left( \bigvee_{x \in U} f(x) \right) \vee \left( \bigvee_{x \in U^{c}} \varnothing_{Y} \right) \\ &= \bigvee_{x \in U} f(x) \end{aligned}$$

for each  $U \in \mathcal{P}(X)$ , where:

- We have used ?? for the first equality.
- We have used Proposition 4.5.5.1.1 for the second equality.
- We have used ?? for the third equality.
- **–** The symbol  $\lor$  denotes the join in  $(Y, \preceq)$ .
- The symbol  $\odot$  denotes the tensor of an element of Y by a truth value as in  $\ref{eq:total_point}$ . In particular, we have

true 
$$\odot f(x) \stackrel{\text{def}}{=} f(x)$$
,  
false  $\odot f(x) \stackrel{\text{def}}{=} \varnothing_Y$ ,

where  $\emptyset_Y$  is the bottom element of  $(Y, \preceq)$ .

In particular, when  $(Y, \preceq_Y) = (\mathcal{P}(B), \subset)$  for some set B, the Kan extension  $\text{Lan}_{\chi_X}(f)$  is given by

$$[\operatorname{Lan}_{\chi_X}(f)](U) = \bigvee_{x \in U} f(x)$$
$$= \bigcup_{x \in U} f(x)$$

for each  $U \in \mathcal{P}(X)$ .

### PROOF 4.4.5.1.4 ▶ PROOF OF PROPOSITION 4.4.5.1.3

### Map I

We define a map

$$\Phi_{X,Y} \colon \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

as in the statement, i.e. by

$$\Phi_{X,Y}(f) \stackrel{\text{def}}{=} f \circ \chi_X$$

for each  $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$ .

### Map II

We define a map

$$\Psi_{X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$$

as in the statement, i.e. by

$$\Psi_{X,Y}(f) \stackrel{\text{def}}{=} \operatorname{Lan}_{\chi_X}(f), \qquad X \xrightarrow{\chi_X} \downarrow \underset{f}{\downarrow} \operatorname{Lan}_{\chi_X}(f)$$

for each  $f \in Sets(X, Y)$ .

### Invertibility I

We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathrm{id}_{\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))}$$
.

We have

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) \stackrel{\text{def}}{=} \Psi_{X,Y} (\Phi_{X,Y}(f))$$

$$\stackrel{\text{def}}{=} \Psi_{X,Y}(f \circ \chi_X)$$

$$\stackrel{\text{def}}{=} \text{Lan}_{\chi_X}(f \circ \chi_X)$$

for each  $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$ . We now claim that

$$\operatorname{Lan}_{\chi_X}(f\circ\chi_X)=f$$

for each  $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$ . Indeed, we have

$$\begin{aligned} \left[ \operatorname{Lan}_{\chi_X} (f \circ \chi_X) \right] (U) &= \bigvee_{x \in U} f(\chi_X(x)) \\ &= f \left( \bigvee_{x \in U} \chi_X(x) \right) \\ &= f \left( \bigcup_{x \in U} \{x\} \right) \\ &= f(U) \end{aligned}$$

for each  $U \in \mathcal{P}(X)$ , where we have used that f is a morphism of suplattices and hence preserves joins for the second equality. This proves our claim. Since we have shown that

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) = f$$

for each  $f \in \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$ , it follows that  $\Psi_{X,Y} \circ \Phi_{X,Y}$  must be equal to the identity map  $\mathsf{id}_{\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))}$  of  $\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$ .

### Invertibility II

We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathsf{Sets}(X,Y)}$$
.

We have

$$\begin{aligned} \left[\Phi_{X,Y} \circ \Psi_{X,Y}\right](f) &\stackrel{\text{def}}{=} \Phi_{X,Y}\left(\Psi_{X,Y}(f)\right) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}\left(\text{Lan}_{\chi_X}(f)\right) \\ &\stackrel{\text{def}}{=} \text{Lan}_{\chi_X}(f) \circ \chi_X \end{aligned}$$

for each  $f \in Sets(X, Y)$ . We now claim that

$$\operatorname{Lan}_{\chi_X}(f) \circ \chi_X = f$$

for each  $f \in Sets(X, Y)$ . Indeed, we have

$$\left[\operatorname{Lan}_{\chi_X}(f)\circ\chi_X\right](x)=\bigvee_{y\in\{x\}}f(y)$$

$$= f(x)$$

for each  $x \in X$ . This proves our claim. Since we have shown that

$$[\Phi_{X,Y} \circ \Psi_{X,Y}](f) = f$$

for each  $f \in \mathsf{Sets}(X,Y)$ , it follows that  $\Phi_{X,Y} \circ \Psi_{X,Y}$  must be equal to the identity map  $\mathsf{id}_{\mathsf{Sets}(X,Y)}$  of  $\mathsf{Sets}(X,Y)$ .

### Naturality for Φ, Part I

We need to show that, given a function  $f: X \to X'$ , the diagram

$$\begin{array}{ccc} \mathsf{SupLat}((\mathcal{P}(X'),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X',Y}} & \mathsf{Sets}(X',Y) \\ & & & \downarrow f^* \\ & \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} & \mathsf{Sets}(X,Y) \end{array}$$

commutes. Indeed, we have

$$\begin{split} \left[\Phi_{X,Y} \circ \mathcal{P}_{!}(f)^{*}\right] &(\xi) \overset{\text{def}}{=} \Phi_{X,Y} \left(\mathcal{P}_{!}(f)^{*}(\xi)\right) \\ &\overset{\text{def}}{=} \Phi_{X,Y} (\xi \circ f_{!}) \\ &\overset{\text{def}}{=} (\xi \circ f_{!}) \circ \chi_{X} \\ &= \xi \circ (f_{!} \circ \chi_{X}) \\ &\overset{(\dagger)}{=} \xi \circ (\chi_{X'} \circ f) \\ &= (\xi \circ \chi_{X'}) \circ f \\ &\overset{\text{def}}{=} \Phi_{X',Y} (\xi) \circ f \\ &\overset{\text{def}}{=} f^{*} \left(\Phi_{X',Y}(\xi)\right) \\ &\overset{\text{def}}{=} \left[f^{*} \circ \Phi_{X',Y}\right] (\xi), \end{split}$$

for each  $\xi \in \text{SupLat}((\mathcal{P}(X'), \subset), (Y, \preceq))$ , where we have used Item 1 of Proposition 4.5.4.1.3 for the fifth equality above.

### Naturality for Φ, Part II

We need to show that, given a morphism of suplattices

$$g: (Y, \preceq_Y) \rightarrow (Y', \preceq_{Y'}),$$

the diagram

$$\begin{split} \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y) \\ & & & & \\ & & & & \\ \mathsf{g}_! & & & & \\ \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y',\preceq)) & \xrightarrow{\Phi_{X,Y'}} \mathsf{Sets}(X,Y') \end{split}$$

commutes. Indeed, we have

$$\begin{split} \left[\Phi_{X,Y'} \circ g_!\right] (\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g_!(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g \circ \xi) \\ &\stackrel{\text{def}}{=} (g \circ \xi) \circ \chi_X \\ &= g \circ (\xi \circ \chi_X) \\ &\stackrel{\text{def}}{=} g \circ \left(\Phi_{X,Y}(\xi)\right) \\ &\stackrel{\text{def}}{=} g_!\left(\Phi_{X,Y}(\xi)\right) \\ &\stackrel{\text{def}}{=} \left[g_! \circ \Phi_{X,Y}\right] (\xi). \end{split}$$

for each  $\xi \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$ .

### Naturality for Ψ

Since  $\Phi$  is natural in each argument and  $\Phi$  is a componentwise inverse to  $\Psi$  in each argument, it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that  $\Psi$  is also natural in each argument.

### WARNING 4.4.5.1.5 ► FREE COCOMPLETION IS NOT AN IDEMPOTENT OPERATION



Although the assignment  $X \mapsto \mathcal{P}(X)$  is called the *free cocompletion of X*, it is not an idempotent operation, i.e. we have  $\mathcal{P}(\mathcal{P}(X)) \neq \mathcal{P}(X)$ .

### 4.4.6 Powersets as Free Completions

Let *X* be a set.

### PROPOSITION 4.4.6.1.1 ▶ POWERSETS AS FREE COMPLETIONS: UNIVERSAL PROPERTY

The pair  $(\mathcal{P}(X), \chi_{(-)})$  consisting of

- The powerset of *X* together with reverse inclusion  $\mathcal{P}(X)^{op} = (\mathcal{P}(X), \supset)$  of Definition 4.4.1.1.1;
- The characteristic embedding  $\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$  of X into  $\mathcal{P}(X)$  of Definition 4.5.4.1.1;

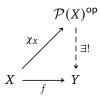
satisfies the following universal property:

- $(\star)$  Given another pair (Y, f) consisting of
  - An inflattice  $(Y, \preceq)$ ;
  - A function  $f: X \to Y$ ;

there exists a unique morphism of inflattices

$$(\mathcal{P}(X),\supset) \xrightarrow{\exists !} (Y,\preceq)$$

making the diagram



commute.

### PROOF 4.4.6.1.2 ► PROOF OF PROPOSITION 4.4.6.1.1

This is a rephrasing of Proposition 4.4.6.1.3, which we prove below.<sup>1</sup>

<sup>1</sup>Here we only remark that the unique morphism of inflattices in the statement is given by the right Kan extension  $\operatorname{Ran}_{\chi_X}(f)$  of f along  $\chi_X$ .

#### PROPOSITION 4.4.6.1.3 ▶ POWERSETS AS FREE COMPLETIONS: ADJOINTNESS

We have an adjunction

$$(\mathcal{P} + \overline{\mathbb{K}})$$
: Sets  $\underbrace{\perp}_{\overline{\mathbb{K}}}$  InfLat,

witnessed by a bijection

$$InfLat((\mathcal{P}(X),\supset),(Y,\preceq)) \cong Sets(X,Y),$$

natural in  $X \in \text{Obj}(\mathsf{Sets})$  and  $(Y, \preceq) \in \text{Obj}(\mathsf{InfLat})$ , where:

- The category InfLat is the category of inflattices of ??.
- The map

$$\chi_X^* : \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

witnessing the above bijection is defined by

$$\chi_X^*(f) \stackrel{\text{def}}{=} f \circ \chi_X,$$

i.e. by sending a morphism of inflattices  $f \colon \mathcal{P}(X)^{\mathsf{op}} \to Y$  to the composition

$$X \stackrel{\chi_X}{\hookrightarrow} \mathcal{P}(X)^{\mathsf{op}} \stackrel{f}{\longrightarrow} Y.$$

• The map

$$\operatorname{Ran}_{\gamma_X} : \operatorname{\mathsf{Sets}}(X,Y) \to \operatorname{\mathsf{InfLat}}((\mathcal{P}(X),\supset),(Y,\preceq))$$

witnessing the above bijection is given by sending a function  $f: X \to Y$  to its right Kan extension along  $\chi_X$ ,

$$\operatorname{Ran}_{\chi_X}(f) \colon \mathcal{P}(X)^{\operatorname{op}} \to Y, \qquad \begin{array}{c} \mathcal{P}(X)^{\operatorname{op}} \\ \chi_X / \text{ } & |_{\operatorname{Ran}_{\chi_X}(f)} \\ X / \text{ } & \downarrow \\ f / \end{array} Y.$$

Moreover, invoking the bijection  $\mathcal{P}(X) \cong \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$  of Item 2 of Proposition 4.5.1.1.4,  $\mathsf{Ran}_{\chi_X}(f)$  can be explicitly computed by

$$\left[\operatorname{Ran}_{\chi_X}(f)\right](U) = \int_{x \in X} \chi_{\mathcal{P}(X)^{\operatorname{op}}}(\chi_x, U) \, \, \text{th} \, \, f(x)$$

$$= \int_{x \in X} \chi_{\mathcal{D}(X)}(U, \chi_{x}) \, \pitchfork f(x)$$

$$= \int_{x \in X} \chi_{U}(x) \, \pitchfork f(x)$$

$$= \bigwedge_{x \in X} \chi_{U}(x) \, \pitchfork f(x)$$

$$= \left( \bigwedge_{x \in U} \chi_{U}(x) \, \pitchfork f(x) \right) \wedge \left( \bigwedge_{x \in U^{c}} \chi_{U}(x) \, \pitchfork f(x) \right)$$

$$= \left( \bigwedge_{x \in U} f(x) \right) \wedge \left( \bigwedge_{x \in U^{c}} \omega_{Y} \right)$$

$$= \left( \bigwedge_{x \in U} f(x) \right) \wedge \omega_{Y}$$

$$= \bigwedge_{x \in U} f(x)$$

for each  $U \in \mathcal{P}(X)$ , where:

- We have used ?? for the first equality.
- We have used Proposition 4.5.5.1.1 for the second equality.
- We have used ?? for the third equality.
- The symbol  $\land$  denotes the meet in  $(Y, \preceq)$ .
- The symbol  $\pitchfork$  denotes the cotensor of an element of *Y* by a truth value as in **??**. In particular, we have

true 
$$\pitchfork f(x) \stackrel{\text{def}}{=} f(x)$$
, false  $\pitchfork f(x) \stackrel{\text{def}}{=} \infty_V$ ,

where  $\infty_Y$  is the top element of  $(Y, \preceq)$ .

In particular, when  $(Y, \preceq_Y) = (\mathcal{P}(B), \subset)$  for some set B, the Kan extension  $\operatorname{Ran}_{\chi_X}(f)$  is given by

$$[\operatorname{Ran}_{\chi_X}(f)](U) = \bigwedge_{x \in U} f(x)$$
$$= \bigcap_{x \in U} f(x)$$

for each  $U \in \mathcal{P}(X)$ .

#### PROOF 4.4.6.1.4 ► PROOF OF PROPOSITION 4.4.5.1.3

### Map I

We define a map

$$\Phi_{X,Y} : \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

as in the statement, i.e. by

$$\Phi_{X,Y}(f) \stackrel{\text{def}}{=} f \circ \chi_X$$

for each  $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$ .

# Map II

We define a map

$$\Psi_{X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq))$$

as in the statement, i.e. by

$$\Psi_{X,Y}(f) \stackrel{\text{def}}{=} \operatorname{Ran}_{\chi_X}(f), \qquad \chi_X / \underset{f}{\swarrow} \operatorname{Ran}_{\chi_X}(f)$$

$$X \xrightarrow{f} Y,$$

for each  $f \in Sets(X, Y)$ .

# Invertibility I

We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathrm{id}_{\mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq))}$$
.

We have

$$\left[ \Psi_{X,Y} \circ \Phi_{X,Y} \right] (f) \stackrel{\text{def}}{=} \Psi_{X,Y} \left( \Phi_{X,Y} (f) \right)$$

$$\stackrel{\text{def}}{=} \Psi_{X,Y} (f \circ \chi_X)$$

$$\stackrel{\text{def}}{=} \operatorname{Ran}_{\chi_X} (f \circ \chi_X)$$

for each  $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$ . We now claim that

$$\operatorname{Ran}_{\chi_X}(f\circ\chi_X)=f$$

for each  $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$ . Indeed, we have

$$\begin{aligned} \left[ \operatorname{Ran}_{\chi_X} (f \circ \chi_X) \right] (U) &= \bigwedge_{x \in U} f(\chi_X(x)) \\ &= f \left( \bigwedge_{x \in U} \chi_X(x) \right) \\ &= f \left( \bigcup_{x \in U} \{x\} \right) \\ &= f(U) \end{aligned}$$

for each  $U \in \mathcal{P}(X)$ , where we have used that f is a morphism of inflattices and hence preserves meets in  $(\mathcal{P}(X), \supset)$  (i.e. joins in  $(\mathcal{P}(X), \subset)$ ) for the second equality. This proves our claim. Since we have shown that

$$\big[\Psi_{X,Y}\circ\Phi_{X,Y}\big](f)=f$$

for each  $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$ , it follows that  $\Psi_{X,Y} \circ \Phi_{X,Y}$  must be equal to the identity map  $id_{InfLat((\mathcal{P}(X), \supset), (Y, \preceq))}$  of  $InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$ .

### Invertibility II

We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathsf{Sets}(X,Y)} \; .$$

We have

$$\begin{split} \left[\Phi_{X,Y} \circ \Psi_{X,Y}\right](f) &\stackrel{\text{def}}{=} \Phi_{X,Y}\left(\Psi_{X,Y}(f)\right) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}\left(\text{Ran}_{\chi_X}(f)\right) \\ &\stackrel{\text{def}}{=} \text{Ran}_{\chi_X}(f) \circ \chi_X \end{split}$$

for each  $f \in Sets(X, Y)$ . We now claim that

$$\operatorname{Ran}_{\gamma_X}(f) \circ \chi_X = f$$

for each  $f \in Sets(X, Y)$ . Indeed, we have

$$[\operatorname{Ran}_{\chi_X}(f) \circ \chi_X](x) = \bigwedge_{y \in \{x\}} f(y)$$
$$= f(x)$$

for each  $x \in X$ . This proves our claim. Since we have shown that

$$\left[\Phi_{X,Y}\circ\Psi_{X,Y}\right](f)=f$$

for each  $f \in \mathsf{Sets}(X,Y)$ , it follows that  $\Phi_{X,Y} \circ \Psi_{X,Y}$  must be equal to the identity map  $\mathsf{id}_{\mathsf{Sets}(X,Y)}$  of  $\mathsf{Sets}(X,Y)$ .

# Naturality for Φ, Part I

We need to show that, given a function  $f: X \to X'$ , the diagram

$$\begin{array}{c|c} \mathsf{InfLat}((\mathcal{P}(X'),\supset),(Y,\preceq)) & \xrightarrow{\Phi_{X',Y}} & \mathsf{Sets}(X',Y) \\ & & & \downarrow f^* \\ \\ \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} & \mathsf{Sets}(X,Y) \end{array}$$

commutes. Indeed, we have

$$\begin{split} \left[\Phi_{X,Y} \circ \mathcal{P}_{!}(f)^{*}\right] &(\xi) \overset{\text{def}}{=} \Phi_{X,Y} \left(\mathcal{P}_{!}(f)^{*}(\xi)\right) \\ &\overset{\text{def}}{=} \Phi_{X,Y} (\xi \circ f_{!}) \\ &\overset{\text{def}}{=} (\xi \circ f_{!}) \circ \chi_{X} \\ &= \xi \circ (f_{!} \circ \chi_{X}) \\ &\overset{(\dagger)}{=} \xi \circ (\chi_{X'} \circ f) \\ &= (\xi \circ \chi_{X'}) \circ f \\ &\overset{\text{def}}{=} \Phi_{X',Y} (\xi) \circ f \\ &\overset{\text{def}}{=} f^{*} \left(\Phi_{X',Y}(\xi)\right) \\ &\overset{\text{def}}{=} \left[f^{*} \circ \Phi_{X',Y}\right] (\xi), \end{split}$$

for each  $\xi \in InfLat((\mathcal{P}(X'), \supset), (Y, \preceq))$ , where we have used Item 1 of Proposition 4.5.4.1.3 for the fifth equality above.

# Naturality for Φ, Part II

We need to show that, given a cocontinuous morphism of posets

$$g: (Y, \preceq_Y) \to (Y', \preceq_{Y'}),$$

the diagram

$$\mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y)$$
 
$$\downarrow^{g_!} \qquad \qquad \downarrow^{g_!}$$
 
$$\mathsf{InfLat}((\mathcal{P}(X),\supset),(Y',\preceq)) \xrightarrow{\Phi_{X,Y'}} \mathsf{Sets}(X,Y')$$

commutes. Indeed, we have

$$\begin{split} \left[\Phi_{X,Y'} \circ g_!\right] (\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g_!(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g \circ \xi) \\ &\stackrel{\text{def}}{=} (g \circ \xi) \circ \chi_X \\ &= g \circ (\xi \circ \chi_X) \\ &\stackrel{\text{def}}{=} g \circ (\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} g_!(\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} \left[g_! \circ \Phi_{X,Y}\right] (\xi). \end{split}$$

for each  $\xi \in \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq))$ .

## Naturality for Ψ

Since  $\Phi$  is natural in each argument and  $\Phi$  is a componentwise inverse to  $\Psi$  in each argument, it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that  $\Psi$  is also natural in each argument.

#### WARNING 4.4.6.1.5 ► FREE COMPLETION IS NOT AN IDEMPOTENT OPERATION



Although the assignment  $X \mapsto \mathcal{P}(X)^{\mathsf{op}}$  is called the *free completion of* X, it is not an idempotent operation, i.e. we have  $\mathcal{P}(\mathcal{P}(X)^{\mathsf{op}})^{\mathsf{op}} \neq \mathcal{P}(X)^{\mathsf{op}}$ .

#### 4.4.7 The Internal Hom of a Powerset

Let *X* be a set and let  $U, V \in \mathcal{P}(X)$ .

#### PROPOSITION 4.4.7.1.1 ► THE INTERNAL HOM OF A POWERSET

The **internal Hom of**  $\mathcal{P}(X)$  **from** U **to** V is the subset  $[U,V]_X^{-1}$  of X given by

$$[U, V]_X = U^{c} \cup V$$
$$= (U \setminus V)^{c}$$

where  $U^{c}$  is the complement of U of Definition 4.3.11.1.1.

<sup>1</sup>Further Notation: Also written  $\mathbf{Hom}_{\mathcal{P}(X)}(U,V)$ .

#### PROOF 4.4.7.1.2 ▶ PROOF OF PROPOSITION 4.4.7.1.1

## Proof of the Equality $U^{c} \cup V = (U \setminus V)^{c}$

We have

$$(U \setminus V)^{c} \stackrel{\text{def}}{=} X \setminus (U \setminus V)$$

$$= (X \cap V) \cup (X \setminus U)$$

$$= V \cup (X \setminus U)$$

$$\stackrel{\text{def}}{=} V \cup U^{c}$$

$$= U^{c} \cup V,$$

where we have used:

- 1. Item 10 of Proposition 4.3.10.1.2 for the second equality.
- 2. Item 4 of Proposition 4.3.9.1.2 for the third equality.
- 3. Item 4 of Proposition 4.3.8.1.2 for the last equality.

This finishes the proof.

#### Proof that $U^{c} \cup V$ Is Indeed the Internal Hom

This follows from Item 2 of Proposition 4.3.9.1.2.

### REMARK 4.4.7.1.3 $\blacktriangleright$ Intuition for the Internal Hom of $\mathcal{P}(X)$

Henning Makholm suggests the following heuristic intuition for the internal Hom of  $\mathcal{P}(X)$  from U to V ([MSE 267365]):

1. Since products in  $\mathcal{P}(X)$  are given by binary intersections (Item 1 of Proposition 4.4.1.1.4), the right adjoint  $\mathbf{Hom}_{\mathcal{P}(X)}(U, -)$  of  $U \cap -$  may be thought of as a function type [U, V].

- 2. Under the Curry–Howard correspondence (??), the function type [U,V] corresponds to implication  $U \Rightarrow V$ .
- 3. Implication  $U \Rightarrow V$  is logically equivalent to  $\neg U \lor V$ .
- 4. The expression  $\neg U \lor V$  then corresponds to the set  $U^c \cup V$  in  $\mathcal{P}(X)$ .
- 5. The set  $U^{c} \vee V$  turns out to indeed be the internal Hom of  $\mathcal{P}(X)$ .

#### PROPOSITION 4.4.7.1.4 ► PROPERTIES OF INTERNAL HOMS OF POWERSETS

Let *X* be a set.

1. Functoriality. The assignments  $U, V, (U, V) \mapsto \mathbf{Hom}_{\mathcal{P}(X)}$  define functors

$$\begin{split} [U,-]_X\colon & (\mathcal{P}(X),\supset) & \to (\mathcal{P}(X),\subset), \\ [-,V]_X\colon & (\mathcal{P}(X),\subset) & \to (\mathcal{P}(X),\subset), \\ [-_1,-_2]_X\colon & (\mathcal{P}(X)\times\mathcal{P}(X),\subset\times\supset) \to (\mathcal{P}(X),\subset). \end{split}$$

In particular, the following statements hold for each  $U, V, A, B \in \mathcal{P}(X)$ :

- (a) If  $U \subset A$ , then  $[A, V]_X \subset [U, V]_X$ .
- (b) If  $V \subset B$ , then  $[U,V]_X \subset [U,B]_X$ .
- (c) If  $U \subset A$  and  $V \subset B$ , then  $[A, V]_X \subset [U, B]_X$ .
- 2. Adjointness. We have adjunctions

$$(U \cap \neg \dashv [U, \neg]_X): \quad \mathcal{P}(X) \xrightarrow{\bot} \mathcal{P}(X),$$

$$(\neg \cap V \dashv [V, \neg]_X): \quad \mathcal{P}(X) \xrightarrow{\bot} \mathcal{P}(X),$$

witnessed by bijections

$$\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, [V, W]_X),$$
  
 $\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(V, [U, W]_X).$ 

In particular, the following statements hold for each  $U, V, W \in \mathcal{P}(X)$ :

(a) The following conditions are equivalent:

- i. We have  $U \cap V \subset W$ .
- ii. We have  $U \subset [V, W]_X$ .
- (b) The following conditions are equivalent:
  - i. We have  $U \cap V \subset W$ .
  - ii. We have  $V \subset [U, W]_X$ .
- 3. Interaction With the Empty Set I. We have

$$[U,\emptyset]_X=U^{\mathsf{c}},$$

$$[\emptyset,V]_X=X,$$

natural in  $U, V \in \mathcal{P}(X)$ .

4. Interaction With X. We have

$$[U,X]_X=X,$$

$$[X,V]_X = V,$$

natural in  $U, V \in \mathcal{P}(X)$ .

5. Interaction With the Empty Set II. The functor

$$D_X \colon \mathcal{P}(X)^{\mathsf{op}} \to \mathcal{P}(X)$$

defined by

$$D_X \stackrel{\text{def}}{=} [-, \emptyset]_X$$
$$= (-)^{\mathsf{c}}$$

is an involutory isomorphism of categories, making  $\emptyset$  into a dualising object for  $(\mathcal{P}(X), \cap, X, [-, -]_X)$  in the sense of **??**. In particular:

(a) The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{D_X} \mathcal{P}(X)$$

$$\mathsf{id}_{\mathcal{P}(X)^{\mathsf{op}}} \qquad \qquad \downarrow^{D_X}$$

$$\mathcal{P}(X)^{\mathsf{op}}$$

commutes, i.e. we have

$$\underbrace{D_X(D_X(U))}_{\stackrel{\text{def}}{=}[[U,\emptyset]_X,\emptyset]_X} = U$$

for each  $U \in \mathcal{P}(X)$ .

(b) The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)^{\mathsf{op}} \xrightarrow{\cap^{\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}}$$

$$\mathrm{id}_{\mathcal{P}(X)^{\mathsf{op}}} \times \mathcal{D}_{X} \xrightarrow{D_{X}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_{X}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

commutes, i.e. we have

$$\underbrace{D_X(U\cap D_X(V))}_{\stackrel{\mathrm{def}}{=}[U\cap[V,\emptyset]_X,\emptyset]_X}=[U,V]_X$$

for each  $U, V \in \mathcal{P}(X)$ .

- 6. *Interaction With the Empty Set III.* Let  $f: X \to Y$  be a function.
  - (a) Interaction With Direct Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_*^{\mathsf{op}}} & \mathcal{P}(Y)^{\mathsf{op}} \\
\downarrow^{D_X} & & \downarrow^{D_Y} \\
\mathcal{P}(X) & \xrightarrow{f_1} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_!(D_X(U)) = D_Y(f_*(U))$$

for each  $U \in \mathcal{P}(X)$ .

(b) Interaction With Inverse Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(Y)^{\mathsf{op}} & \xrightarrow{f^{-1},\mathsf{op}} & \mathcal{P}(X)^{\mathsf{op}} \\
\downarrow^{D_{Y}} & & \downarrow^{D_{X}} \\
\mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$f^{-1}(D_Y(U)) = D_X(f^{-1}(U))$$

for each  $U \in \mathcal{P}(X)$ .

(c) Interaction With Codirect Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_!^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
\downarrow^{D_X} & & \downarrow^{D_Y} \\
\mathcal{P}(X) & \xrightarrow{f_*} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(D_X(U)) = D_Y(f_!(U))$$

for each  $U \in \mathcal{P}(X)$ .

7. Interaction With Unions of Families of Subsets I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\cup^{\mathsf{op}} \times \cup^{\mathsf{op}} \bigvee \qquad \qquad \bigcup \cup$$

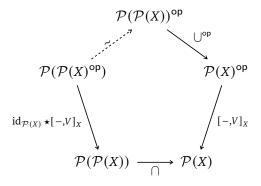
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcup_{U \in \mathcal{U}} U, \bigcup_{V \in \mathcal{V}} V\right]_X$$

in general, where  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

8. Interaction With Unions of Families of Subsets II. The diagram



commutes, i.e. we have

$$\left[\bigcup_{U\in\mathcal{U}}U,V\right]_X=\bigcap_{U\in\mathcal{U}}[U,V]_X$$

for each  $U \in \mathcal{P}(\mathcal{P}(X))$  and each  $V \in \mathcal{P}(X)$ .

9. Interaction With Unions of Families of Subsets III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcup}{\longrightarrow} & \mathcal{P}(X) \\
id_{\mathcal{P}(X)} \star [U,-]_X & & & \downarrow [U,-]_X \\
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcup}{\longrightarrow} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U, \bigcup_{V \in \mathcal{V}} V\right]_{X} = \bigcup_{V \in \mathcal{V}} [U, V]_{X}$$

for each  $U \in \mathcal{P}(X)$  and each  $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

10. Interaction With Intersections of Families of Subsets I. The diagram

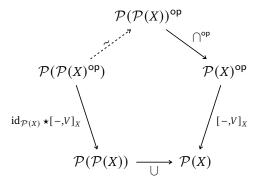
$$\mathcal{P}(\mathcal{P}(X))^{\text{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X)) \\
\cap^{\text{op}} \times \cap^{\text{op}} \downarrow \qquad \qquad \downarrow \cap \\
\mathcal{P}(X)^{\text{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcap_{U \in \mathcal{U}} U, \bigcap_{V \in \mathcal{V}} V\right]_X$$

in general, where  $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ .

11. Interaction With Intersections of Families of Subsets II. The diagram



commutes, i.e. we have

$$\left[\bigcap_{U\in\mathcal{U}}U,V\right]_X=\bigcup_{U\in\mathcal{U}}[U,V]_X$$

for each  $U \in \mathcal{P}(\mathcal{P}(X))$  and each  $V \in \mathcal{P}(X)$ .

12. Interaction With Intersections of Families of Subsets III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcap} & \mathcal{P}(X) \\
id_{\mathcal{P}(X)} \star [U,-]_X & & & \downarrow [U,-]_X \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcap} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U,\bigcap_{V\in\mathcal{V}}V\right]_{x}=\bigcap_{V\in\mathcal{V}}[U,V]_{X}$$

for each  $U \in \mathcal{P}(X)$  and each  $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ .

13. Interaction With Binary Unions. We have equalities of sets

$$[U \cap V, W]_X = [U, W]_X \cup [V, W]_X,$$
  
 $[U, V \cap W]_X = [U, V]_X \cap [U, W]_X$ 

for each  $U, V, W \in \mathcal{P}(X)$ .

14. Interaction With Binary Intersections. We have equalities of sets

$$[U \cup V, W]_X = [U, W]_X \cap [V, W]_X,$$
  
 $[U, V \cup W]_X = [U, V]_X \cup [U, W]_X$ 

for each  $U, V, W \in \mathcal{P}(X)$ .

15. Interaction With Differences. We have equalities of sets

$$[U \setminus V, W]_X = [U, W]_X \cup [V^c, W]_X$$
$$= [U, W]_X \cup [U, V]_X,$$
$$[U, V \setminus W]_X = [U, V]_X \setminus (U \cap W)$$

for each  $U, V, W \in \mathcal{P}(X)$ .

16. Interaction With Complements. We have equalities of sets

$$[U^{c}, V]_{X} = U \cup V,$$
  

$$[U, V^{c}]_{X} = U \cap V,$$
  

$$[U, V]_{X}^{c} = U \setminus V$$

for each  $U, V \in \mathcal{P}(X)$ .

17. Interaction With Characteristic Functions. We have

$$\chi_{[U,V]_{\mathcal{P}(X)}}(x) = \max(1 - \chi_U \pmod{2}, \chi_V)$$

for each  $U, V \in \mathcal{P}(X)$ .

18. *Interaction With Direct Images.* Let  $f: X \to Y$  be a function. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_*^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \\
[-1,-2]_X \downarrow \qquad \qquad \downarrow [-1,-2]_Y \\
\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have an equality of sets

$$f_!([U,V]_X) = [f_*(U), f_!(V)]_Y,$$

natural in  $U, V \in \mathcal{P}(X)$ .

19. *Interaction With Inverse Images.* Let  $f: X \to Y$  be a function. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1}, \mathsf{op} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) 
[-_1, -_2]_Y \qquad \qquad \qquad \downarrow [-_1, -_2]_X 
\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = \left[f^{-1}(U), f^{-1}(V)\right]_X,$$

natural in  $U, V \in \mathcal{P}(X)$ .

20. *Interaction With Codirect Images.* Let  $f: X \to Y$  be a function. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_*} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \\
[-1,-2]_X \downarrow \qquad \qquad \qquad \downarrow [-1,-2]_Y \\
\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

with components

$$[f_!(U), f_*(V)]_V \subset f_*([U, V]_X)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

## PROOF 4.4.7.1.5 ► PROOF OF PROPOSITION 4.4.7.1.4

### Item 1: Functoriality

Since  $\mathcal{P}(X)$  is posetal, it suffices to prove Items 1a to 1c.

1. Proof of Item 1a: We have

$$[A, V]_X \stackrel{\text{def}}{=} A^{\mathsf{c}} \cup V$$

$$\subset U^{\mathsf{c}} \cup V$$

$$\stackrel{\mathsf{def}}{=} [U, V]_X,$$

where we have used:

- (a) Item 1 of Proposition 4.3.11.1.2, which states that if  $U \subset A$ , then  $A^{c} \subset U^{c}$ .
- (b) Item 1a of Item 1 of Proposition 4.3.11.1.2, which states that if  $A^c \subset U^c$ , then  $A^c \cup K \subset U^c \cup K$  for any  $K \in \mathcal{P}(X)$ .
- 2. Proof of Item 1b: We have

$$[U, V]_X \stackrel{\text{def}}{=} U^c \cup V$$
$$\subset U^c \cup B$$
$$\stackrel{\text{def}}{=} [U, B]_X,$$

where we have used Item 1b of Item 1 of Proposition 4.3.11.1.2, which states that if  $V \subset B$ , then  $K \cup V \subset K \cup B$  for any  $K \in \mathcal{P}(X)$ .

3. Proof of Item 1c: We have

$$[A,V]_X \subset [U,V]_X$$
$$\subset [U,B]_X,$$

where we have used Items 1a and 1b.

This finishes the proof.

### Item 2: Adjointness

This is a repetition of Item 2 of Proposition 4.3.9.1.2 and is proved there.

#### Item 3: Interaction With the Empty Set I

We have

$$[U, \emptyset]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup \emptyset$$
$$= U^{\mathsf{c}},$$

where we have used Item 3 of Proposition 4.3.8.1.2, and we have

$$[\emptyset, V]_X \stackrel{\text{def}}{=} \emptyset^{\mathsf{c}} \cup V$$

$$\stackrel{\text{def}}{=} (X \setminus \emptyset) \cup V$$

$$= X \cup V$$

$$= X,$$

where we have used:

- 1. Item 12 of Proposition 4.3.10.1.2 for the first equality.
- 2. Item 5 of Proposition 4.3.8.1.2 for the last equality.

Since  $\mathcal{P}(X)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2).

# Item 4: Interaction With *X*

We have

$$[U,X]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup X$$
$$= X,$$

where we have used Item 5 of Proposition 4.3.8.1.2, and we have

$$[X, V]_X \stackrel{\text{def}}{=} X^{\mathsf{c}} \cup V$$

$$\stackrel{\text{def}}{=} (X \setminus X) \cup V$$

$$= \emptyset \cup V$$

$$= V,$$

where we have used Item 3 of Proposition 4.3.8.1.2 for the last equality. Since  $\mathcal{P}(X)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2).

# Item 5: Interaction With the Empty Set II

We have

$$D_X(D_X(U)) \stackrel{\text{def}}{=} [[U, \emptyset]_X, \emptyset]_X$$
$$= [U^c, \emptyset]_X$$
$$= (U^c)^c$$
$$= U,$$

where we have used:

- 1. Item 3 for the second and third equalities.
- 2. Item 3 of Proposition 4.3.11.1.2 for the fourth equality.

Since  $\mathcal{P}(X)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2), and thus we have

$$[[-,\emptyset]_X,\emptyset]_X \cong \mathrm{id}_{\mathcal{P}(X)}$$

This finishes the proof.

### Item 6: Interaction With the Empty Set III

Since  $D_X = (-)^c$ , this is essentially a repetition of the corresponding results for  $(-)^c$ , namely Items 5 to 7 of Proposition 4.3.11.1.2.

### Item 7: Interaction With Unions of Families of Subsets I

By Item 3 of Proposition 4.4.7.1.4, we have

$$[\mathcal{U}, \emptyset]_{\mathcal{P}(X)} = \mathcal{U}^{\mathsf{c}},$$
$$[\mathcal{U}, \emptyset]_{X} = \mathcal{U}^{\mathsf{c}}.$$

With this, the counterexample given in the proof of Item 10 of Proposition 4.3.6.1.2 then applies.

#### Item 8: Interaction With Unions of Families of Subsets II

We have

$$\left[\bigcup_{U \in \mathcal{U}} U, V\right]_{X}^{\text{def}} \left(\bigcup_{U \in \mathcal{U}} U\right)^{c} \cup V$$

$$= \left(\bigcap_{U \in \mathcal{U}} U^{c}\right) \cup V$$

$$= \bigcap_{U \in \mathcal{U}} (U^{c} \cup V)$$

$$\stackrel{\text{def}}{=} \bigcap_{U \in \mathcal{U}} [U, V]_{X},$$

where we have used:

1. Item 11 of Proposition 4.3.6.1.2 for the second equality.

### 2. Item 6 of Proposition 4.3.7.1.2 for the third equality.

This finishes the proof.

## Item 9: Interaction With Unions of Families of Subsets III

We have

$$\bigcup_{V \in \mathcal{V}} [U, V]_X \stackrel{\text{def}}{=} \bigcup_{V \in \mathcal{V}} (U^{\mathsf{c}} \cup V)$$

$$= U^{\mathsf{c}} \cup \left(\bigcup_{V \in \mathcal{V}} V\right)$$

$$\stackrel{\text{def}}{=} \left[U, \bigcup_{V \in \mathcal{V}} V\right]_X$$

where we have used Item 6. This finishes the proof.

# Item 10: Interaction With Intersections of Families of Subsets I

Let  $X = \{0, 1\}$ , let  $\mathcal{U} = \{\{0, 1\}\}\$ , and let  $\mathcal{V} = \{\{0\}, \{0, 1\}\}\$ . We have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W = \bigcap_{W \in \mathcal{P}(X)} W$$
$$= \{0, 1\},$$

whereas

$$\left[\bigcap_{U\in\mathcal{U}}U,\bigcap_{V\in\mathcal{V}}V\right]_X = [\{0,1\},\{0\}]$$
$$= \{0\},$$

Thus we have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W = \{0, 1\} \neq \{0\} = \left[\bigcap_{U \in \mathcal{U}} U, \bigcap_{V \in \mathcal{V}} V\right]_X.$$

This finishes the proof.

Item 11: Interaction With Intersections of Families of Subsets II

We have

$$\left[\bigcap_{U \in \mathcal{U}} U, V\right]_{X} \stackrel{\text{def}}{=} \left(\bigcap_{U \in \mathcal{U}} U\right)^{c} \cup V$$

$$= \left(\bigcup_{U \in \mathcal{U}} U^{c}\right) \cup V$$

$$= \bigcup_{U \in \mathcal{U}} (U^{c} \cup V)$$

$$\stackrel{\text{def}}{=} \bigcup_{U \in \mathcal{U}} [U, V]_{X},$$

where we have used:

- 1. Item 12 of Proposition 4.3.6.1.2 for the second equality.
- 2. Item 6 of Proposition 4.3.7.1.2 for the third equality.

This finishes the proof.

### Item 12: Interaction With Intersections of Families of Subsets III

We have

$$\bigcap_{V \in \mathcal{V}} [U, V]_X \stackrel{\text{def}}{=} \bigcap_{V \in \mathcal{V}} (U^c \cup V)$$

$$= U^c \cup \left(\bigcap_{V \in \mathcal{V}} V\right)$$

$$\stackrel{\text{def}}{=} \left[U, \bigcap_{V \in \mathcal{V}} V\right]_X.$$

where we have used Item 6. This finishes the proof.

# Item 13: Interaction With Binary Unions

We have

$$[U \cap V, W]_X \stackrel{\text{def}}{=} (U \cap V)^{c} \cup W$$
$$= (U^{c} \cup V^{c}) \cup W$$
$$= (U^{c} \cup V^{c}) \cup (W \cup W)$$

$$= (U^{c} \cup W) \cup (V^{c} \cup W)$$

$$\stackrel{\text{def}}{=} [U, W]_{X} \cup [V, W]_{X},$$

where we have used:

- 1. Item 2 of Proposition 4.3.11.1.2 for the second equality.
- 2. Item 8 of Proposition 4.3.8.1.2 for the third equality.
- 3. Several applications of Items 2 and 4 of Proposition 4.3.8.1.2 and for the fourth equality.

For the second equality in the statement, we have

$$\begin{split} [U, V \cap W]_X &\stackrel{\text{def}}{=} U^{\mathsf{c}} \cup (V \cap W) \\ &= (U^{\mathsf{c}} \cup V) \cap (U^{\mathsf{c}} \cap W) \\ &\stackrel{\text{def}}{=} [U, V]_X \cap [U, W]_X, \end{split}$$

where we have used Item 6 of Proposition 4.3.8.1.2 for the second equality.

## Item 14: Interaction With Binary Intersections

We have

$$\begin{aligned} [U \cup V, W]_X &\stackrel{\text{def}}{=} (U \cup V)^c \cup W \\ &= (U^c \cap V^c) \cup W \\ &= (U^c \cup W) \cap (V^c \cup W) \\ &\stackrel{\text{def}}{=} [U, W]_X \cap [V, W]_X, \end{aligned}$$

where we have used:

- 1. Item 2 of Proposition 4.3.11.1.2 for the second equality.
- 2. Item 6 of Proposition 4.3.8.1.2 for the third equality.

Now, for the second equality in the statement, we have

$$\begin{split} [U, V \cup W]_X &\stackrel{\text{def}}{=} U^\mathsf{c} \cup (V \cup W) \\ &= (U^\mathsf{c} \cup U^\mathsf{c}) \cup (V \cup W) \\ &= (U^\mathsf{c} \cup V) \cup (U^\mathsf{c} \cup W) \\ &\stackrel{\text{def}}{=} [U, V]_X \cup [U, W]_X, \end{split}$$

where we have used:

- 1. Item 8 of Proposition 4.3.8.1.2 for the second equality.
- 2. Several applications of Items 2 and 4 of Proposition 4.3.8.1.2 and for the third equality.

This finishes the proof.

# Item 15: Interaction With Differences

We have

$$\begin{split} [U \setminus V, W]_X &\stackrel{\text{def}}{=} (U \setminus V)^{\mathsf{c}} \cup W \\ &\stackrel{\text{def}}{=} (X \setminus (U \setminus V)) \cup W \\ &= ((X \cap V) \cup (X \setminus U)) \cup W \\ &= (V \cup (X \setminus U)) \cup W \\ &\stackrel{\text{def}}{=} (V \cup U^{\mathsf{c}}) \cup W \\ &= (V \cup (U^{\mathsf{c}} \cup U^{\mathsf{c}})) \cup W \\ &= (U^{\mathsf{c}} \cup W) \cup (U^{\mathsf{c}} \cup V) \\ &\stackrel{\text{def}}{=} [U, W]_X \cup [U, V]_X, \end{split}$$

where we have used:

- 1. Item 10 of Proposition 4.3.10.1.2 for the third equality.
- 2. Item 4 of Proposition 4.3.9.1.2 for the fourth equality.
- 3. Item 8 of Proposition 4.3.8.1.2 for the sixth equality.
- 4. Several applications of Items 2 and 4 of Proposition 4.3.8.1.2 and for the seventh equality.

We also have

$$\begin{split} [U \setminus V, W]_X &\stackrel{\text{def}}{=} (U \setminus V)^c \cup W \\ &\stackrel{\text{def}}{=} (X \setminus (U \setminus V)) \cup W \\ &= ((X \cap V) \cup (X \setminus U)) \cup W \\ &= (V \cup (X \setminus U)) \cup W \\ &\stackrel{\text{def}}{=} (V \cup U^c) \cup W \\ &= (V \cup U^c) \cup (W \cup W) \end{split}$$

$$= (U^{c} \cup W) \cup (V \cup W)$$
$$= (U^{c} \cup W) \cup ((V^{c})^{c} \cup W)$$
$$\stackrel{\text{def}}{=} [U, W]_{X} \cup [V^{c}, W]_{X},$$

where we have used:

- 1. Item 10 of Proposition 4.3.10.1.2 for the third equality.
- 2. Item 4 of Proposition 4.3.9.1.2 for the fourth equality.
- 3. Item 8 of Proposition 4.3.8.1.2 for the sixth equality.
- 4. Several applications of Items 2 and 4 of Proposition 4.3.8.1.2 and for the seventh equality.
- 5. Item 3 of Proposition 4.3.11.1.2 for the eighth equality.

Now, for the second equality in the statement, we have

$$\begin{split} [U,V\setminus W]_X &\stackrel{\mathrm{def}}{=} U^{\mathsf{c}} \cup (V\setminus W) \\ &= (V\setminus W) \cup U^{\mathsf{c}} \\ &= (V\cup U^{\mathsf{c}})\setminus (W\setminus U^{\mathsf{c}}) \\ &\stackrel{\mathrm{def}}{=} (V\cup U^{\mathsf{c}})\setminus (W\setminus (X\setminus U)) \\ &= (V\cup U^{\mathsf{c}})\setminus ((W\cap U)\cup (W\setminus X)) \\ &= (V\cup U^{\mathsf{c}})\setminus ((W\cap U)\cup \emptyset) \\ &= (V\cup U^{\mathsf{c}})\setminus (W\cap U) \\ &= (V\cup U^{\mathsf{c}})\setminus (U\cap W) \\ &\stackrel{\mathrm{def}}{=} [U,V]_X\setminus (U\cap W) \end{split}$$

where we have used:

- 1. Item 4 of Proposition 4.3.8.1.2 for the second equality.
- 2. Item 4 of Proposition 4.3.10.1.2 for the third equality.
- 3. Item 10 of Proposition 4.3.10.1.2 for the fifth equality.
- 4. Item 13 of Proposition 4.3.10.1.2 for the sixth equality.
- 5. Item 3 of Proposition 4.3.8.1.2 for the seventh equality.

6. Item 5 of Proposition 4.3.9.1.2 for the eighth equality.

This finishes the proof.

### Item 16: Interaction With Complements

We have

$$[U^{c}, V]_{X} \stackrel{\text{def}}{=} (U^{c})^{c} \cup V,$$
$$= U \cup V,$$

where we have used Item 3 of Proposition 4.3.11.1.2. We also have

$$[U, V^{c}]_{X} \stackrel{\text{def}}{=} U^{c} \cup V^{c}$$
$$= U \cap V$$

where we have used Item 2 of Proposition 4.3.11.1.2. Finally, we have

$$[U, V]_X^{c} = ((U \setminus V)^{c})^{c}$$
$$= U \setminus V,$$

where we have used Item 2 of Proposition 4.3.11.1.2.

#### Item 17: Interaction With Characteristic Functions

We have

$$\chi_{[U,V]_{\mathcal{P}(X)}}(x) \stackrel{\text{def}}{=} \chi_{U^{c} \cup V}(x)$$

$$= \max(\chi_{U^{c}}, \chi_{V})$$

$$= \max(1 - \chi_{U} \pmod{2}, \chi_{V}),$$

where we have used:

- 1. Item 10 of Proposition 4.3.8.1.2 for the second equality.
- 2. Item 4 of Proposition 4.3.11.1.2 for the third equality.

This finishes the proof.

### Item 18: Interaction With Direct Images

This is a repetition of Item 10 of Proposition 4.6.1.1.5 and is proved there.

### Item 19: Interaction With Inverse Images

This is a repetition of Item 10 of Proposition 4.6.2.1.3 and is proved there.

# Item 20: Interaction With Codirect Images

This is a repetition of Item 9 of Proposition 4.6.3.1.7 and is proved there.

## 4.4.8 Isbell Duality for Sets

Let *X* be a set.

#### **DEFINITION 4.4.8.1.1** ► THE ISBELL FUNCTION

The **Isbell function** of *X* is the map

$$I: \mathcal{P}(X) \to \mathsf{Sets}(X, \mathcal{P}(X))$$

defined by

$$\mathsf{I}(U) \stackrel{\mathrm{def}}{=} \llbracket x \mapsto [U, \{x\}]_X \rrbracket$$

for each  $U \in \mathcal{P}(X)$ .

#### REMARK 4.4.8.1.2 ► MOTIVATION FOR THE ISBELL FUNCTION

Recall from Remark 4.4.1.1.2 that we may view the powerset  $\mathcal{P}(X)$  of a set X as the decategorification of the category of presheaves  $\mathsf{PSh}(C)$  of a category C. Building upon this analogy, we want to mimic the definition of the Isbell Spec functor, which is given on objects by

$$Spec(\mathcal{F}) \stackrel{\text{def}}{=} Nat(\mathcal{F}, h_{(-)})$$

for each  $\mathcal{F} \in \text{Obj}(\mathsf{PSh}(C))$ . To this end, we could define

$$\mathsf{I}(U) \stackrel{\mathrm{def}}{=} \left[ U, \chi_{(-)} \right]_X,$$

replacing:

- The Yoneda embedding  $X \mapsto h_X$  of C into  $\mathsf{PSh}(C)$  with the characteristic embedding  $x \mapsto \chi_X$  of X into  $\mathcal{P}(X)$  of Definition 4.5.4.1.1.
- The internal Hom Nat of PSh(C) with the internal Hom  $[-,-]_X$  of  $\mathcal{P}(X)$  of Proposition 4.4.7.1.1.

However, since  $[U, \chi_x]_X$  is a subset of U instead of a truth value, we get a function

$$I: \mathcal{P}(X) \to \mathsf{Sets}(X, \mathcal{P}(X))$$

instead of a function

$$I: \mathcal{P}(X) \to \mathcal{P}(X).$$

This makes some of the properties involving I a bit more cumbersome to state, although we still have an analogue of Isbell duality in that  $I_! \circ I$  evaluates to  $id_{\mathcal{P}(X)}$  in the sense of Proposition 4.4.8.1.3.

\_

The diagram

$$\mathcal{P}(X) \xrightarrow{\mathsf{I}} \mathsf{Sets}(X, \mathcal{P}(X))$$

$$\Delta_{\Delta_{\mathrm{id}_{\mathcal{P}(X)}}} \qquad \qquad \mathsf{I}_{!}$$

$$\mathsf{Sets}(X, \mathsf{Sets}(X, \mathcal{P}(X)))$$

commutes, i.e. we have

$$I_!(I(U)) = \llbracket x \mapsto \llbracket y \mapsto U \rrbracket \rrbracket$$

for each  $U \in \mathcal{P}(X)$ .

### PROOF 4.4.8.1.4 ► PROOF OF PROPOSITION 4.4.8.1.3

We have

$$I_{!}(I(U)) \stackrel{\text{def}}{=} I_{!}(\llbracket x \mapsto U^{c} \cup \{x\} \rrbracket)$$

$$\stackrel{\text{def}}{=} \llbracket x \mapsto I(U^{c} \cup \{x\}) \rrbracket$$

$$\stackrel{\text{def}}{=} \llbracket x \mapsto \llbracket y \mapsto (U^{c} \cup \{x\})^{c} \cup \{x\} \rrbracket \rrbracket$$

$$= \llbracket x \mapsto \llbracket y \mapsto (U \cap (X \setminus \{x\})) \cup \{x\} \rrbracket \rrbracket$$

$$= \llbracket x \mapsto \llbracket y \mapsto (U \setminus \{x\}) \cup \{x\} \rrbracket \rrbracket$$

$$= \llbracket x \mapsto \llbracket y \mapsto U \rrbracket \rrbracket,$$

where we have used Item 2 of Proposition 4.3.11.1.2 for the fourth equality above.

### 4.5 Characteristic Functions

#### 4.5.1 The Characteristic Function of a Subset

Let *X* be a set and let  $U \in \mathcal{P}(X)$ .

#### **DEFINITION 4.5.1.1.1** ► THE CHARACTERISTIC FUNCTION OF A SUBSET

The **characteristic function of**  $U^1$  is the function  $\chi_U: X \to \{t, f\}^2$  defined by

$$\chi_U(x) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x \in U, \\ \text{false} & \text{if } x \notin U \end{cases}$$

for each  $x \in X$ .

<sup>1</sup>Further Terminology: Also called the **indicator function of** *U*.

# REMARK 4.5.1.1.2 ► CHARACTERISTIC FUNCTIONS OF SUBSETS AS DECATEGORIFICATIONS OF PRESHEAVES

Under the analogy that  $\{t,f\}$  should be the (-1)-categorical analogue of Sets, we may view a function

$$f: X \to \{\mathsf{t},\mathsf{f}\}$$

as a decategorification of presheaves and copresheaves

$$\mathcal{F}: C^{\mathsf{op}} \to \mathsf{Sets},$$
  
 $F: C \to \mathsf{Sets}.$ 

The characteristic functions  $\chi_U$  of the subsets of X are then the primordial examples of such functions (and, in fact, all of them).

#### NOTATION 4.5.1.1.3 ► FURTHER NOTATION FOR CHARACTERISTIC FUNCTIONS

We will often employ the bijection  $\{t, f\} \cong \{0, 1\}$  to make use of the arithmetical operations defined on  $\{0, 1\}$  when disucssing characteristic functions. Examples of this include Items 4 to 11 of Proposition 4.5.1.1.4 below.

#### PROPOSITION 4.5.1.1.4 ► PROPERTIES OF CHARACTERISTIC FUNCTIONS OF SUBSETS

Let *X* be a set.

1. Functionality. The assignment  $U \mapsto \chi_U$  defines a function

$$\chi_{(-)} : \mathcal{P}(X) \to \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\}).$$

- 2. *Bijectivity*. The function  $\chi_{(-)}$  from Item 1 is bijective.
- 3. Naturality. The collection

$$\left\{\chi_{(-)}\colon \mathcal{P}(X)\to \mathsf{Sets}(X,\{\mathsf{t},\mathsf{f}\})\right\}_{X\in \mathsf{Obj}(\mathsf{Sets})}$$

defines a natural isomorphism between  $\mathcal{P}^{-1}$  and  $\mathsf{Sets}(\mathsf{-}, \{\mathsf{t}, \mathsf{f}\})$ . In

<sup>&</sup>lt;sup>2</sup>Further Notation: Also written  $\chi_X(U, -)$  or  $\chi_X(-, U)$ .

particular, given a function  $f: X \to Y$ , the diagram

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

$$\chi_{(-)} \downarrow \chi \qquad \qquad \downarrow \chi_{(-)}$$

$$\mathsf{Sets}(Y, \{\mathsf{t}, \mathsf{f}\}) \xrightarrow{f^*} \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

commutes, i.e. we have

$$\chi_V \circ f = \chi_{f^{-1}(V)}$$

for each  $V \in \mathcal{P}(Y)$ .

4. Interaction With Unions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each  $U, V \in \mathcal{P}(X)$ .

5. Interaction With Unions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each  $U, V \in \mathcal{P}(X)$ .

6. Interaction With Intersections I. We have

$$\chi_{U\cap V} = \chi_U \chi_V$$

for each  $U, V \in \mathcal{P}(X)$ .

7. Interaction With Intersections II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each  $U, V \in \mathcal{P}(X)$ .

8. Interaction With Differences. We have

$$\chi_{U\setminus V} = \chi_U - \chi_{U\cap V}$$

for each  $U, V \in \mathcal{P}(X)$ .

9. Interaction With Complements. We have

$$\chi_{U^c} \equiv 1 - \chi_U \pmod{2}$$

for each  $U \in \mathcal{P}(X)$ .

10. Interaction With Symmetric Differences. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \pmod{2}$$

for each  $U, V \in \mathcal{P}(X)$ .

11. Interaction With Internal Homs. We have

$$\chi_{[U,V]_{\mathcal{P}(X)}} = \max(1 - \chi_U \pmod{2}, \chi_V)$$

for each  $U, V \in \mathcal{P}(X)$ .

### PROOF 4.5.1.1.5 ► PROOF OF PROPOSITION 4.5.1.1.4

### Item 1: Functionality

There is nothing to prove.

# Item 2: Bijectivity

We proceed in three steps:

1. The Inverse of  $\chi_{(-)}$ . The inverse of  $\chi_{(-)}$  is the map

$$\Phi \colon \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\}) \xrightarrow{\sim} \mathcal{P}(X),$$

defined by

$$\begin{split} \Phi(f) &\stackrel{\text{def}}{=} U_f \\ &\stackrel{\text{def}}{=} f^{-1}(\mathsf{true}) \\ &\stackrel{\text{def}}{=} \{x \in X \mid f(x) = \mathsf{true}\} \end{split}$$

for each  $f \in Sets(X, \{t, f\})$ .

2. Invertibility I. We have

$$\begin{split} \left[\Phi \circ \chi_{(-)}\right] &(U) \stackrel{\mathrm{def}}{=} \Phi(\chi_U) \\ &\stackrel{\mathrm{def}}{=} \chi_U^{-1}(\mathsf{true}) \\ &\stackrel{\mathrm{def}}{=} \left\{x \in X \mid \chi_U(x) = \mathsf{true}\right\} \\ &\stackrel{\mathrm{def}}{=} \left\{x \in X \mid x \in U\right\} \\ &= U \\ &\stackrel{\mathrm{def}}{=} \left[\mathrm{id}_{\mathcal{P}(X)}\right] &(U) \end{split}$$

for each  $U \in \mathcal{P}(X)$ . Thus, we have

$$\Phi \circ \chi_{(-)} = \mathrm{id}_{\mathcal{P}(X)}.$$

3. Invertibility II. We have

$$\begin{split} \left[\chi_{(-)} \circ \Phi\right] (U) &\stackrel{\text{def}}{=} \chi_{\Phi(f)} \\ &\stackrel{\text{def}}{=} \chi_{f^{-1}(\mathsf{true})} \\ &\stackrel{\text{def}}{=} \left[\!\!\left[x \mapsto \begin{cases} \mathsf{true} & \text{if } x \in f^{-1}(\mathsf{true}) \\ \mathsf{false} & \text{otherwise} \end{cases} \right] \\ &= \left[\!\!\left[x \mapsto f(x)\right]\!\!\right] \\ &= f \\ &\stackrel{\text{def}}{=} \left[\mathsf{id}_{\mathsf{Sets}(X,\{\mathsf{t},f\})}\right] (f) \end{split}$$

for each  $f \in Sets(X, \{t, f\})$ . Thus, we have

$$\chi_{(-)} \circ \Phi = \mathrm{id}_{\mathrm{Sets}(X,\{t,f\})}$$
.

This finishes the proof.

### Item 3: Naturality

We proceed in two steps:

1. *Naturality of*  $\chi_{(-)}$ . We have

$$\begin{split} [\chi_V \circ f](v) &\stackrel{\text{def}}{=} \chi_V(f(v)) \\ &= \begin{cases} \text{true} & \text{if } f(v) \in V, \\ \text{false} & \text{otherwise} \end{cases} \end{split}$$

$$= \begin{cases} \text{true} & \text{if } v \in f^{-1}(V), \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_{f^{-1}(V)}(v)$$

for each  $v \in V$ .

2. Naturality of  $\Phi$ . Since  $\chi_{(-)}$  is natural and a componentwise inverse to  $\Phi$ , it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that  $\Phi$  is also natural in each argument.

This finishes the proof.

#### Item 4: Interaction With Unions I

This is a repetition of Item 10 of Proposition 4.3.8.1.2 and is proved there.

#### Item 5: Interaction With Unions II

This is a repetition of Item 11 of Proposition 4.3.8.1.2 and is proved there.

### Item 6: Interaction With Intersections I

This is a repetition of Item 10 of Proposition 4.3.9.1.2 and is proved there.

#### Item 7: Interaction With Intersections II

This is a repetition of Item 11 of Proposition 4.3.9.1.2 and is proved there.

### Item 8: Interaction With Differences

This is a repetition of Item 16 of Proposition 4.3.10.1.2 and is proved there.

#### Item 9: Interaction With Complements

This is a repetition of Item 4 of Proposition 4.3.11.1.2 and is proved there.

### Item 10: Interaction With Symmetric Differences

This is a repetition of Item 15 of Proposition 4.3.12.1.2 and is proved there.

#### Item 11: Interaction With Internal Homs

This is a repetition of Item 17 of Proposition 4.4.7.1.4 and is proved there.

#### REMARK 4.5.1.1.6 ▶ POWERSETS AS SETS OF FUNCTIONS AND UN/STRAIGHTENING

The bijection

$$\mathcal{P}(X) \cong \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

of Item 2 of Proposition 4.5.1.1.4, which

• Takes a subset  $U \hookrightarrow X$  of X and straightens it to a function  $\chi_U \colon X \to X$ 

{true, false};

• Takes a function  $f: X \to \{\text{true}, \text{false}\}\$ and *unstraightens* it to a subset  $f^{-1}(\text{true}) \hookrightarrow X \text{ of } X;$ 

may be viewed as the (-1)-categorical version of the 0-categorical un/straightening isomorphism between indexed and fibred sets

$$\underbrace{\mathsf{FibSets}_X}_{\substack{\mathsf{def}\\ \mathsf{=}\mathsf{Sets}_{/X}}} \cong \underbrace{\mathsf{ISets}_X}_{\substack{\mathsf{def}\\ \mathsf{=}\mathsf{Fun}(X_{\mathsf{disc}},\mathsf{Sets})}}$$

of Un/Straightening for Indexed and Fibred Sets, ??. Here we view:

- Subsets  $U \hookrightarrow X$  as being analogous to X-fibred sets  $\phi_X \colon A \to X$ .
- Functions  $f: X \to \{t, f\}$  as being analogous to *X*-indexed sets  $A: X_{\sf disc} \to {\sf Sets}$ .

#### 4.5.2 The Characteristic Function of a Point

Let *X* be a set and let  $x \in X$ .

#### **DEFINITION 4.5.2.1.1** ► THE CHARACTERISTIC FUNCTION OF A POINT

The **characteristic function of** x is the function<sup>1</sup>

$$\chi_X : X \to \{\mathsf{t},\mathsf{f}\}$$

defined by

$$\chi_x \stackrel{\mathrm{def}}{=} \chi_{\{x\}}$$

i.e. by

$$\chi_x(y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each  $y \in X$ .

<sup>&</sup>lt;sup>1</sup>*Further Notation:* Also written  $\chi^x$ ,  $\chi_X(x, -)$ , or  $\chi_X(-, x)$ .

# REMARK 4.5.2.1.2 ➤ CHARACTERISTIC FUNCTIONS OF POINTS AS DECATEGORIFICATIONS OF REPRESENTABLE PRESHEAVES

Expanding upon Remark 4.5.1.1.2, we may think of the characteristic function

$$\chi_X \colon X \to \{\mathsf{t},\mathsf{f}\}$$

of an *element* x of X as a decategorification of the representable presheaf and of the representable copresheaf

$$h_X \colon C^{\mathsf{op}} \to \mathsf{Sets},$$
  
 $h^X \colon C \to \mathsf{Sets}$ 

associated of an object X of a category C.

#### 4.5.3 The Characteristic Relation of a Set

Let *X* be a set.

#### **DEFINITION 4.5.3.1.1** ► THE CHARACTERISTIC RELATION OF A SET

The **characteristic relation on**  $X^1$  is the relation<sup>2</sup>

$$\gamma_X(-1,-2): X \times X \to \{\mathsf{t},\mathsf{f}\}$$

on X defined by  $^3$ 

$$\chi_X(x,y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each  $x, y \in X$ .

# REMARK 4.5.3.1.2 ➤ THE CHARACTERISTIC RELATION OF A SET AS A DECATEGORIFICATION OF THE HOM PROFUNCTOR

Expanding upon Remarks 4.5.1.1.2 and 4.5.2.1.2, we may view the characteristic relation

$$\gamma_X(-1,-2): X \times X \to \{t,f\}$$

<sup>&</sup>lt;sup>1</sup>Further Terminology: Also called the **identity relation on** *X*.

<sup>&</sup>lt;sup>2</sup>Further Notation: Also written  $\chi_{-2}^{-1}$ , or  $\sim_{id}$  in the context of relations.

<sup>&</sup>lt;sup>3</sup>Under the bijection Sets( $X \times X$ , {t, f})  $\cong \mathcal{P}(X \times X)$  of Item 2 of Proposition 4.5.1.1.4, the relation  $\chi_X$  corresponds to the diagonal  $\Delta_X \subset X \times X$  of X.

of X as a decategorification of the Hom profunctor

$$\operatorname{Hom}_{\mathcal{C}}(-_1, -_2) \colon \mathcal{C}^{\mathsf{op}} \times \mathcal{C} \to \mathsf{Sets}$$

of a category C.

#### PROPOSITION 4.5.3.1.3 ► PROPERTIES OF CHARACTERISTIC RELATIONS

Let  $f: X \to Y$  be a function.

1. The Inclusion of Characteristic Relations Associated to a Function. Let  $f: A \to B$  be a function. We have an inclusion 1

$$\chi_B \circ (f \times f) \subset \chi_A, \qquad A \times A \xrightarrow{f \times f} B \times B$$

$$\chi_A \searrow \chi_B$$

$$\chi_A \downarrow \chi_B$$

$$\chi_A \downarrow \chi_B$$

$$\chi_A \downarrow \chi_B$$

<sup>1</sup>Note: This is the 0-categorical version of Categories, Definition 11.5.4.1.1.

#### PROOF 4.5.3.1.4 ► PROOF OF PROPOSITION 4.5.3.1.3

### Item 1: The Inclusion of Characteristic Relations Associated to a Function

The inclusion  $\chi_B(f(a), f(b)) \subset \chi_A(a, b)$  is equivalent to the statement "if a = b, then f(a) = f(b)", which is true.

# 4.5.4 The Characteristic Embedding of a Set

Let *X* be a set.

#### **DEFINITION 4.5.4.1.1** ► THE CHARACTERISTIC EMBEDDING OF A SET

The **characteristic embedding**<sup>1</sup> **of** *X* **into**  $\mathcal{P}(X)$  is the function

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

defined by<sup>2</sup>

$$\chi_{(-)}(x) \stackrel{\text{def}}{=} \chi_x$$
$$= \{x\}$$

for each  $x \in X$ .

<sup>1</sup>The name "characteristic *embedding*" is justified by Corollary 4.5.5.1.3, which gives an analogue of fully faithfulness for  $\chi_{(-)}$ .

<sup>2</sup>Here we are identifying  $\mathcal{P}(X)$  with Sets(X, {t, f}) as per Item 2 of Proposition 4.5.1.1.4.

# REMARK 4.5.4.1.2 ► THE CHARACTERISTIC EMBEDDING OF A SET AS A DECATEGORIFICATION OF THE YONEDA EMBEDDING

Expanding upon Remarks 4.5.1.1.2, 4.5.2.1.2 and 4.5.3.1.2, we may view the characteristic embedding

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

of X into  $\mathcal{P}(X)$  as a decategorification of the Yoneda embedding

$$\sharp : C^{\mathsf{op}} \hookrightarrow \mathsf{PSh}(C)$$

of a category C into PSh(C).

#### PROPOSITION 4.5.4.1.3 ► PROPERTIES OF CHARACTERISTIC EMBEDDINGS

Let  $f: X \to Y$  be a map of sets.

1. Interaction With Functions. We have

$$f_! \circ \chi_X = \chi_Y \circ f, \qquad \chi_X \qquad \qquad \downarrow^{f} \chi_Y$$

$$\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

# PROOF 4.5.4.1.4 ► PROOF OF PROPOSITION 4.5.4.1.3

## Item 1: Interaction With Functions

Indeed, we have

$$[f_! \circ \chi_X](x) \stackrel{\text{def}}{=} f_!(\chi_X(x))$$
$$\stackrel{\text{def}}{=} f_!(\{x\})$$
$$= \{f(x)\}$$

$$\stackrel{\text{def}}{=} \chi_{X'}(f(x))$$

$$\stackrel{\text{def}}{=} [\chi_{X'} \circ f](x),$$

for each  $x \in X$ , showing the desired equality.

#### 4.5.5 The Yoneda Lemma for Sets

Let X be a set and let  $U \subset X$  be a subset of X.

#### PROPOSITION 4.5.5.1.1 ► THE YONEDA LEMMA FOR SETS

We have

$$\chi_{\mathcal{P}(X)}(\chi_x,\chi_U)=\chi_U(x)$$

for each  $x \in X$ , giving an equality of functions

$$\chi_{\mathcal{P}(X)}(\chi_{(-)},\chi_U)=\chi_U,$$

where

$$\chi_{\mathcal{P}(X)}(U,V) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } U \subset V, \\ \text{false} & \text{otherwise.} \end{cases}$$

## PROOF 4.5.5.1.2 ► PROOF OF PROPOSITION 4.5.5.1.1

We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_U) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } \{x\} \subset U, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } x \in U \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_U(x).$$

This finishes the proof.

#### COROLLARY 4.5.5.1.3 ► THE CHARACTERISTIC EMBEDDING IS FULLY FAITHFUL

The characteristic embedding is fully faithful, i.e., we have

$$\chi_{\mathcal{P}(X)}(\chi_x,\chi_y)\cong\chi_X(x,y)$$

for each  $x, y \in X$ .

## PROOF 4.5.5.1.4 ► PROOF OF COROLLARY 4.5.5.1.3

We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_y) = \chi_y(x)$$

$$\stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x \in \{y\} \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } x = y \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_X(x, y).$$

where we have used Proposition 4.5.5.1.1 for the first equality.

# **4.6** The Adjoint Triple $f_! \dashv f^{-1} \dashv f_*$

## 4.6.1 Direct Images

Let  $f: X \to Y$  be a function.

#### **DEFINITION 4.6.1.1.1** ► **DIRECT IMAGES**

The **direct image function associated to** f is the function<sup>1</sup>

$$f_! \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by<sup>2</sup>

$$f_!(U) \stackrel{\text{def}}{=} \left\{ y \in Y \mid \text{there exists some } x \in U \\ \text{such that } y = f(x) \right\}$$

$$= \left\{ f(x) \in Y \mid x \in U \right\}$$

for each  $U \in \mathcal{P}(X)$ .

<sup>&</sup>lt;sup>1</sup>Further Notation: Also written simply  $f: \mathcal{P}(X) \to \mathcal{P}(Y)$ .

<sup>&</sup>lt;sup>2</sup>Further Terminology: The set f(U) is called the **direct image of** U **by** f.

## NOTATION 4.6.1.1.2 ► FURTHER NOTATION FOR DIRECT IMAGES

Sometimes one finds the notation

$$\exists_f \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

for  $f_!$ . This notation comes from the fact that the following statements are equivalent, where  $y \in Y$  and  $U \in \mathcal{P}(X)$ :

- We have  $y \in \exists_f(U)$ .
- There exists some  $x \in U$  such that f(x) = y.

We will not make use of this notation elsewhere in Clowder.

WARNING 4.6.1.1.3 ► NOTATION FOR DIRECT IMAGES IS CONFUSING

Notation for direct images between powersets is tricky:

- 1. Direct images for powersets and presheaves are both adjoint to their corresponding inverse image functors. However, the direct image functor for powersets is a *left* adjoint, while the direct image functor for presheaves is a *right* adjoint:
  - (a) *Powersets*. Given a function  $f: X \to Y$ , we have an inverse image functor

$$f^{-1} \colon \mathcal{P}(Y) \to \mathcal{P}(X).$$

The *left* adjoint of this functor is the usual direct image, defined above in Definition 4.6.1.1.1.

(b) *Presheaves*. Given a morphism of topological spaces  $f: X \to Y$ , we have an inverse image functor

$$f^{-1} \colon \mathsf{PSh}(Y) \to \mathsf{PSh}(X).$$

The *right* adjoint of this functor is the direct image functor of presheaves, defined in ??.

- 2. The presheaf direct image functor is denoted  $f_*$ , but the direct image functor for powersets is denoted  $f_!$  (as it's a left adjoint).
- 3. Adding to the confusion, it's somewhat common for  $f_!: \mathcal{P}(X) \to \mathcal{P}(Y)$  to be denoted  $f_*$ .

We chose to write  $f_!$  for the direct image to keep the notation aligned with the following similar adjoint situations:

SITUATION	Adjoint String
Functoriality of Powersets	$(f_! \dashv f^{-1} \dashv f_*) \colon \mathcal{P}(X) \xrightarrow{\rightleftarrows} \mathcal{P}(Y)$
Functoriality of Presheaf Categories	$(f_! \dashv f^{-1} \dashv f_*) \colon PSh(X) \xrightarrow{\rightleftarrows} PSh(Y)$
Base Change	$(f_!\dashv f^*\dashv f_*)\colon C_{/X}\overset{\rightleftarrows}{ o} C_{/Y}$
Kan Extensions	$(F_! \dashv F^* \dashv F_*) \colon \operatorname{Fun}(C, \mathcal{E}) \xrightarrow{\rightleftarrows} \operatorname{Fun}(\mathcal{D}, \mathcal{E})$



#### REMARK 4.6.1.1.4 ► Unwinding Definition 4.6.1.1.1

Identifying  $\mathcal{P}(X)$  with  $\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$  via Item 2 of Proposition 4.5.1.1.4, we see that the direct image function associated to f is equivalently the function

$$f_! \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by

$$f_{!}(\chi_{U}) \stackrel{\text{def}}{=} \operatorname{Lan}_{f}(\chi_{U})$$

$$= \operatorname{colim}\left(\left(f \stackrel{\rightarrow}{\times} (\underline{-_{1}})\right) \stackrel{\text{pr}}{\twoheadrightarrow} A \stackrel{\chi_{U}}{\longrightarrow} \{t, f\}\right)$$

$$= \operatorname{colim}_{\substack{x \in X \\ f(x) = -_{1}}} (\chi_{U}(x))$$

$$= \bigvee_{\substack{x \in X \\ f(x) = -_{1}}} (\chi_{U}(x)),$$

where we have used ?? for the second equality. In other words, we have

$$[f_!(\chi_U)](y) = \bigvee_{\substack{x \in X \\ f(x) = y}} (\chi_U(x))$$

$$= \begin{cases} \text{true} & \text{if there exists some } x \in X \text{ such that } f(x) = y \text{ and } x \in U, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if there exists some } x \in U \\ & \text{such that } f(x) = y, \\ \text{false} & \text{otherwise} \end{cases}$$

for each  $y \in Y$ .

#### PROPOSITION 4.6.1.1.5 ► PROPERTIES OF DIRECT IMAGES I

Let  $f: X \to Y$  be a function.

1. Functoriality. The assignment  $U \mapsto f_1(U)$  defines a functor

$$f_! : (\mathcal{P}(X), \subset) \to (\mathcal{P}(Y), \subset).$$

In particular, for each  $U, V \in \mathcal{P}(X)$ , the following condition is satisfied:

 $(\star)$  If  $U \subset V$ , then  $f_!(U) \subset f_!(V)$ .

2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$id_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad id_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$
  
 $f_! \circ f^{-1} \hookrightarrow id_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow id_{\mathcal{P}(X)},$ 

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*\Big(f^{-1}(V)\Big),$$
  
 $f_!\Big(f^{-1}(V)\Big) \subset V, \qquad f^{-1}(f_*(U)) \subset U$ 

indexed by  $U \in \mathcal{P}(X)$  and  $V \in \mathcal{P}(Y)$ .

(b) Bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(Y)}(f_!(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f^{-1}(V)),$$
  
 $\operatorname{Hom}_{\mathcal{P}(X)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f_*(V)),$ 

natural in  $U \in \mathcal{P}(X)$  and  $V \in \mathcal{P}(Y)$  and (respectively)  $V \in \mathcal{P}(X)$  and  $U \in \mathcal{P}(Y)$ . In particular:

- i. The following conditions are equivalent:
  - A. We have  $f_!(U) \subset V$ .
  - B. We have  $U \subset f^{-1}(V)$ .
- ii. The following conditions are equivalent:
  - A. We have  $f^{-1}(U) \subset V$ .
  - B. We have  $U \subset f_*(V)$ .
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\begin{array}{ccc} \mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_!)_1} & \mathcal{P}(\mathcal{P}(Y)) \\ & & & \downarrow \cup \\ & & & \downarrow \cup \\ & \mathcal{P}(X) & \xrightarrow{f_!} & \mathcal{P}(Y) \end{array}$$

commutes, i.e. we have

$$\bigcup_{U\in\mathcal{U}} f_!(U) = \bigcup_{V\in f_!(\mathcal{U})} V$$

for each  $\mathcal{U} \in \mathcal{P}(X)$ , where  $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$ .

4. Interaction With Intersections of Families of Subsets. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_!)_!} & \mathcal{P}(\mathcal{P}(Y)) \\
& & & & \downarrow \cap \\
& & & \downarrow \cap \\
& & & \mathcal{P}(X) & \xrightarrow{f_!} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U\in\mathcal{U}}f_!(U)=\bigcap_{V\in f_!(\mathcal{U})}V$$

for each  $\mathcal{U} \in \mathcal{P}(X)$ , where  $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$ .

5. Interaction With Binary Unions. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_! \times f_!} \mathcal{P}(Y) \times \mathcal{P}(Y)$$

$$\downarrow \bigcup_{\mathcal{P}(X) \xrightarrow{f_!}} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$

for each  $U, V \in \mathcal{P}(X)$ .

6. *Interaction With Binary Intersections*. We have a natural transformation

with components

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

7. Interaction With Differences. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad$$

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

8. Interaction With Complements. The diagram

$$\mathcal{P}(X)^{\text{op}} \xrightarrow{f_{*}^{\text{op}}} \mathcal{P}(Y)^{\text{op}} \\
\xrightarrow{(-)^{c}} \qquad \qquad \downarrow^{(-)^{c}} \\
\mathcal{P}(X) \xrightarrow{f_{1}} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_!(U^\mathsf{c}) = f_*(U)^\mathsf{c}$$

for each  $U \in \mathcal{P}(X)$ .

9. *Interaction With Symmetric Differences*. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\stackrel{\triangle}{\downarrow} \qquad \qquad \qquad \qquad \downarrow \triangle$$

$$\mathcal{P}(X) \xrightarrow{f_!} \qquad \qquad \mathcal{P}(Y)$$

with components

$$f_!(U) \triangle f_!(V) \subset f_!(U \triangle V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

10. Interaction With Internal Homs of Powersets. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_*^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) 
[-1,-2]_X \downarrow \qquad \qquad \downarrow [-1,-2]_Y 
\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have an equality of sets

$$f_!([U,V]_X) = [f_*(U), f_!(V)]_Y,$$

natural in  $U, V \in \mathcal{P}(X)$ .

11. Preservation of Colimits. We have an equality of sets

$$f_!\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f_!(U_i),$$

natural in  $\{U_i\}_{i\in I} \in \mathcal{P}(X)^{\times I}$ . In particular, we have equalities

$$f_!(U) \cup f_!(V) = f_!(U \cup V),$$
  
 $f_!(\emptyset) = \emptyset,$ 

natural in  $U, V \in \mathcal{P}(X)$ .

12. Oplax Preservation of Limits. We have an inclusion of sets

$$f!\left(\bigcap_{i\in I}U_i\right)\subset\bigcap_{i\in I}f_!(U_i),$$

natural in  $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{\times I}$ . In particular, we have inclusions

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V),$$
  
 $f_!(X) \subset Y,$ 

natural in  $U, V \in \mathcal{P}(X)$ .

13. Symmetric Strict Monoidality With Respect to Unions. The direct image function of Item 1 has a symmetric strict monoidal structure

$$(f_!, f_!^{\otimes}, f_{!|1}^{\otimes}) \colon (\mathcal{P}(X), \cup, \emptyset) \to (\mathcal{P}(Y), \cup, \emptyset),$$

being equipped with equalities

$$f_{!|U,V}^{\otimes} \colon f_{!}(U) \cup f_{!}(V) \xrightarrow{=} f_{!}(U \cup V),$$
  
 $f_{!|1}^{\otimes} \colon \emptyset \xrightarrow{=} \emptyset,$ 

natural in  $U, V \in \mathcal{P}(X)$ .

14. *Symmetric Oplax Monoidality With Respect to Intersections.* The direct image function of Item 1 has a symmetric oplax monoidal structure

$$(f_!, f_!^{\otimes}, f_{!|1}^{\otimes}) \colon (\mathcal{P}(X), \cap, X) \to (\mathcal{P}(Y), \cap, Y),$$

being equipped with inclusions

$$f_{!|U,V}^{\otimes} : f_{!}(U \cap V) \hookrightarrow f_{!}(U) \cap f_{!}(V),$$
  
 $f_{!|1}^{\otimes} : f_{!}(X) \hookrightarrow Y,$ 

natural in  $U, V \in \mathcal{P}(X)$ .

15. *Interaction With Coproducts.* Let  $f: X \to X'$  and  $g: Y \to Y'$  be maps of sets. We have

$$(f \mid \mid g)_!(U \mid \mid V) = f_!(U) \mid \mid g_!(V)$$

for each  $U \in \mathcal{P}(X)$  and each  $V \in \mathcal{P}(Y)$ .

16. *Interaction With Products.* Let  $f: X \to X'$  and  $g: Y \to Y'$  be maps of sets. We have

$$(f \boxtimes_{X \times Y} g)_!(U \boxtimes_{X \times Y} V) = f_!(U) \boxtimes_{X' \times Y'} g_!(V)$$

for each  $U \in \mathcal{P}(X)$  and each  $V \in \mathcal{P}(Y)$ .

17. Relation to Codirect Images. We have

$$f_!(U) = f_*(U^{\mathsf{c}})^{\mathsf{c}}$$

$$\stackrel{\text{def}}{=} Y \setminus f_*(X \setminus U)$$

for each  $U \in \mathcal{P}(X)$ .

#### PROOF 4.6.1.1.6 ► PROOF OF PROPOSITION 4.6.1.1.5

## Item 1: Functoriality

Omitted.

## Item 2: Triple Adjointness

This follows from Remark 4.6.1.1.4, Remark 4.6.2.1.2, Remark 4.6.3.1.4, and Kan Extensions, ?? of ??.

## Item 3: Interaction With Unions of Families of Subsets

We have

$$\bigcup_{V \in f_{!}(\mathcal{U})} V = \bigcup_{V \in \{f_{!}(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcup_{U \in \mathcal{U}} f_{!}(U).$$

This finishes the proof.

## Item 4: Interaction With Intersections of Families of Subsets

We have

$$\bigcap_{V \in f_!(\mathcal{U})} V = \bigcap_{V \in \{f_!(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcap_{U \in \mathcal{U}} f_!(U).$$

This finishes the proof.

Item 5: Interaction With Binary Unions

See [Pro25p].

Item 6: Interaction With Binary Intersections

See [Pro25n].

Item 7: Interaction With Differences

See [Pro25o].

Item 8: Interaction With Complements

Applying Item 17 to  $X \setminus U$ , we have

$$f_!(U^c) = f_!(X \setminus U)$$

$$= Y \setminus f_*(X \setminus (X \setminus U))$$
  
= Y \ f\_\*(U)  
= f\_\*(U)^c.

This finishes the proof.

## Item 9: Interaction With Symmetric Differences

We have

$$f_{!}(U) \triangle f_{!}(V) = (f_{!}(U) \cup f_{!}(V)) \setminus (f_{!}(U) \cap f_{!}(V))$$

$$\subset (f_{!}(U) \cup f_{!}(V)) \setminus (f_{!}(U \cap V))$$

$$= (f_{!}(U \cup V)) \setminus (f_{!}(U \cap V))$$

$$\subset f_{!}((U \cup V) \setminus (U \cap V))$$

$$= f_{!}(U \triangle V),$$

where we have used:

- 1. Item 2 of Proposition 4.3.12.1.2 for the first equality.
- 2. Item 6 of this proposition together with Item 1 of Proposition 4.3.10.1.2 for the first inclusion.
- 3. Item 5 for the second equality.
- 4. Item 7 for the second inclusion.
- 5. Item 2 of Proposition 4.3.12.1.2 for the tchird equality.

Since  $\mathcal{P}(Y)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

## Item 10: Interaction With Internal Homs of Powersets

We have

$$f_{!}([U,V]_{X}) \stackrel{\text{def}}{=} f_{!}(U^{c} \cup V)$$

$$= f_{!}(U^{c}) \cup f_{!}(V)$$

$$= f_{*}(U)^{c} \cup f_{!}(V)$$

$$\stackrel{\text{def}}{=} [f_{*}(U), f_{!}(V)]_{Y},$$

where we have used:

- 1. Item 5 for the second equality.
- 2. Item 17 for the third equality.

Since  $\mathcal{P}(Y)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

#### Item 11: Preservation of Colimits

This follows from Item 2 and ??, ?? of ??.<sup>1</sup>

#### Item 12: Oplax Preservation of Limits

The inclusion  $f_!(X) \subset Y$  is automatic. See [Pro25n] for the other inclusions.

Item 13: Symmetric Strict Monoidality With Respect to Unions

This follows from Item 11.

## Item 14: Symmetric Oplax Monoidality With Respect to Intersections

The inclusions in the statement follow from Item 12. Since  $\mathcal{P}(Y)$  is posetal, the commutativity of the diagrams in the definition of a symmetric oplax monoidal functor is automatic (Categories, Item 4 of Proposition 11.2.7.1.2).

## Item 15: Interaction With Coproducts

Omitted.

#### Item 16: Interaction With Products

Omitted.

## Item 17: Relation to Codirect Images

Applying Item 16 of Proposition 4.6.3.1.7 to  $X \setminus U$ , we have

$$f_*(X \setminus U) = B \setminus f_!(X \setminus (X \setminus U))$$
$$= B \setminus f_!(U).$$

Taking complements, we then obtain

$$f_!(U) = B \setminus (B \setminus f_!(U)),$$
  
=  $B \setminus f_*(X \setminus U),$ 

which finishes the proof.

<sup>1</sup>Reference: [Pro25p].

#### PROPOSITION 4.6.1.1.7 ▶ PROPERTIES OF DIRECT IMAGES II

Let  $f: X \to Y$  be a function.

1. Functionality I. The assignment  $f \mapsto f_!$  defines a function

$$(-)_{*|X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(X),\mathcal{P}(Y)).$$

2. Functionality II. The assignment  $f \mapsto f_!$  defines a function

$$(-)_{*|X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(X),\subset),(\mathcal{P}(Y),\subset)).$$

3. *Interaction With Identities.* For each  $X \in Obj(Sets)$ , we have

$$(\mathrm{id}_X)_! = \mathrm{id}_{\mathcal{P}(X)}$$
.

4. *Interaction With Composition*. For each pair of composable functions  $f: X \to Y$  and  $g: Y \to Z$ , we have

$$(g \circ f)_! = g_! \circ f_!, \qquad \begin{array}{c} \mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y) \\ & \downarrow g_! \\ & \mathcal{P}(Z). \end{array}$$

#### PROOF 4.6.1.1.8 ► PROOF OF PROPOSITION 4.6.1.1.7

## Item 1: Functionality I

There is nothing to prove.

## Item 2: Functionality II

This follows from Item 1 of Proposition 4.6.1.1.5.

## Item 3: Interaction With Identities

This follows from Remark 4.6.1.1.4 and Kan Extensions, ?? of ??.

## Item 4: Interaction With Composition

This follows from Remark 4.6.1.1.4 and Kan Extensions, ?? of ??.

## 4.6.2 Inverse Images

Let  $f: X \to Y$  be a function.

The **inverse image function associated to** f is the function<sup>1</sup>

$$f^{-1} \colon \mathcal{P}(Y) \to \mathcal{P}(X)$$

defined by<sup>2</sup>

$$f^{-1}(V) \stackrel{\text{def}}{=} \{x \in X \mid \text{we have } f(x) \in V\}$$

for each  $V \in \mathcal{P}(Y)$ .

## REMARK 4.6.2.1.2 ► UNWINDING DEFINITION 4.6.2.1.1

Identifying  $\mathcal{P}(Y)$  with  $\mathsf{Sets}(Y, \{\mathsf{t}, \mathsf{f}\})$  via Item 2 of Proposition 4.5.1.1.4, we see that the inverse image function associated to f is equivalently the function

$$f^* \colon \mathcal{P}(Y) \to \mathcal{P}(X)$$

defined by

$$f^*(\chi_V) \stackrel{\mathrm{def}}{=} \chi_V \circ f$$

for each  $\chi_V \in \mathcal{P}(Y)$ , where  $\chi_V \circ f$  is the composition

$$X \xrightarrow{f} Y \xrightarrow{\chi_V} \{\text{true, false}\}$$

in Sets.

## PROPOSITION 4.6.2.1.3 ► PROPERTIES OF INVERSE IMAGES I

Let  $f: X \to Y$  be a function.

1. Functoriality. The assignment  $V \mapsto f^{-1}(V)$  defines a functor

$$f^{-1} \colon (\mathcal{P}(Y), \subset) \to (\mathcal{P}(X), \subset).$$

In particular, for each  $U, V \in \mathcal{P}(Y)$ , the following condition is satisfied:

$$(\star)$$
 If  $U \subset V$ , then  $f^{-1}(U) \subset f^{-1}(V)$ .

<sup>&</sup>lt;sup>1</sup>Further Notation: Also written  $f^* : \mathcal{P}(Y) \to \mathcal{P}(X)$ .

<sup>&</sup>lt;sup>2</sup>Further Terminology: The set  $f^{-1}(V)$  is called the **inverse image of** V **by** f.

2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$id_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad id_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$
  
 $f_! \circ f^{-1} \hookrightarrow id_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow id_{\mathcal{P}(X)},$ 

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*\Big(f^{-1}(V)\Big),$$
  
 $f_!\Big(f^{-1}(V)\Big) \subset V, \qquad f^{-1}(f_*(U)) \subset U$ 

indexed by  $U \in \mathcal{P}(X)$  and  $V \in \mathcal{P}(Y)$ .

(b) Bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(Y)}(f_!(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f^{-1}(V)),$$
  
 $\operatorname{Hom}_{\mathcal{P}(X)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f_*(V)),$ 

natural in  $U \in \mathcal{P}(X)$  and  $V \in \mathcal{P}(Y)$  and (respectively)  $V \in \mathcal{P}(X)$  and  $U \in \mathcal{P}(Y)$ . In particular:

- i. The following conditions are equivalent:
  - A. We have  $f_!(U) \subset V$ .
  - B. We have  $U \subset f^{-1}(V)$ .
- ii. The following conditions are equivalent:
  - A. We have  $f^{-1}(U) \subset V$ .
  - B. We have  $U \subset f_*(V)$ .
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{V \in \mathcal{V}} f^{-1}(V) = \bigcup_{U \in f^{-1}(\mathcal{U})} U$$

for each  $V \in \mathcal{P}(Y)$ , where  $f^{-1}(V) \stackrel{\text{def}}{=} (f^{-1})^{-1}(V)$ .

4. Interaction With Intersections of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcap_{V \in \mathcal{V}} f^{-1}(V) = \bigcap_{U \in f^{-1}(\mathcal{U})} U$$

for each  $\mathcal{V} \in \mathcal{P}(Y)$ , where  $f^{-1}(\mathcal{V}) \stackrel{\text{def}}{=} (f^{-1})^{-1}(\mathcal{V})$ .

5. Interaction With Binary Unions. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$

for each  $U, V \in \mathcal{P}(Y)$ .

6. Interaction With Binary Intersections. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\uparrow \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U \cap V) = f^{-1}(U) \cap f^{-1}(V)$$

for each  $U, V \in \mathcal{P}(Y)$ .

7. Interaction With Differences. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\mathsf{op},-1} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad$$

commutes, i.e. we have

$$f^{-1}(U \setminus V) = f^{-1}(U) \setminus f^{-1}(V)$$

for each  $U, V \in \mathcal{P}(X)$ .

8. Interaction With Complements. The diagram

$$\mathcal{P}(Y)^{\text{op}} \xrightarrow{f^{-1,\text{op}}} \mathcal{P}(X)^{\text{op}}$$

$$(-)^{c} \downarrow \qquad \qquad \downarrow (-)^{c}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U^{\mathsf{c}}) = f^{-1}(U)^{\mathsf{c}}$$

for each  $U \in \mathcal{P}(X)$ .

9. Interaction With Symmetric Differences. The diagram

$$\begin{array}{ccc} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) & \xrightarrow{f^{\mathsf{op},-1} \times f^{-1}} & \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \\ & & & & & \downarrow \triangle \\ & & & & & \downarrow \triangle \\ & & & & & \mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X) \end{array}$$

i.e. we have

$$f^{-1}(U) \triangle f^{-1}(V) = f^{-1}(U \triangle V)$$

for each  $U, V \in \mathcal{P}(Y)$ .

10. Interaction With Internal Homs of Powersets. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1}, \mathsf{op} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \\
 \downarrow [-1, -2]_{X} \\
 \mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_Y$$

natural in  $U, V \in \mathcal{P}(X)$ .

11. Preservation of Colimits. We have an equality of sets

$$f^{-1}\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f^{-1}(U_i),$$

natural in  $\{U_i\}_{i\in I}\in \mathcal{P}(Y)^{\times I}$ . In particular, we have equalities

$$f^{-1}(U) \cup f^{-1}(V) = f^{-1}(U \cup V),$$
  
 $f^{-1}(\emptyset) = \emptyset,$ 

natural in  $U, V \in \mathcal{P}(Y)$ .

12. Preservation of Limits. We have an equality of sets

$$f^{-1}\left(\bigcap_{i\in I}U_i\right)=\bigcap_{i\in I}f^{-1}(U_i),$$

natural in  $\{U_i\}_{i\in I} \in \mathcal{P}(Y)^{\times I}$ . In particular, we have equalities

$$f^{-1}(U) \cap f^{-1}(V) = f^{-1}(U \cap V),$$
  
 $f^{-1}(Y) = X,$ 

natural in  $U, V \in \mathcal{P}(Y)$ .

13. *Symmetric Strict Monoidality With Respect to Unions.* The inverse image function of Item 1 has a symmetric strict monoidal structure

$$\left(f^{-1}, f^{-1, \otimes}, f_{1}^{-1, \otimes}\right) \colon (\mathcal{P}(Y), \cup, \emptyset) \to (\mathcal{P}(X), \cup, \emptyset),$$

being equipped with equalities

$$f_{U,V}^{-1,\otimes} \colon f^{-1}(U) \cup f^{-1}(V) \xrightarrow{=} f^{-1}(U \cup V),$$
$$f_{1}^{-1,\otimes} \colon \emptyset \xrightarrow{=} f^{-1}(\emptyset),$$

natural in  $U, V \in \mathcal{P}(Y)$ .

14. *Symmetric Strict Monoidality With Respect to Intersections*. The inverse image function of Item 1 has a symmetric strict monoidal structure

$$\left(f^{-1}, f^{-1, \otimes}, f_{\mathbb{1}}^{-1, \otimes}\right) \colon (\mathcal{P}(Y), \cap, Y) \to (\mathcal{P}(X), \cap, X),$$

being equipped with equalities

$$f_{U,V}^{-1,\otimes} \colon f^{-1}(U) \cap f^{-1}(V) \stackrel{=}{\to} f^{-1}(U \cap V),$$
  
 $f_{\mathbb{I}}^{-1,\otimes} \colon X \stackrel{=}{\to} f^{-1}(Y),$ 

natural in  $U, V \in \mathcal{P}(Y)$ .

15. *Interaction With Coproducts.* Let  $f: X \to X'$  and  $g: Y \to Y'$  be maps of sets. We have

$$(f \coprod g)^{-1}(U' \coprod V') = f^{-1}(U') \coprod g^{-1}(V')$$

for each  $U' \in \mathcal{P}(X')$  and each  $V' \in \mathcal{P}(Y')$ .

16. *Interaction With Products.* Let  $f: X \to X'$  and  $g: Y \to Y'$  be maps of sets. We have

$$(f \boxtimes_{X' \times Y'} g)^{-1}(U' \boxtimes_{X' \times Y'} V') = f^{-1}(U') \boxtimes_{X \times Y} g^{-1}(V')$$

for each  $U' \in \mathcal{P}(X')$  and each  $V' \in \mathcal{P}(Y')$ .

#### PROOF 4.6.2.1.4 ► PROOF OF PROPOSITION 4.6.2.1.3

## Item 1: Functoriality

Omitted.

## Item 2: Triple Adjointness

This follows from Remark 4.6.1.1.4, Remark 4.6.2.1.2, Remark 4.6.3.1.4, and Kan Extensions, ?? of ??.

## Item 3: Interaction With Unions of Families of Subsets

We have

$$\bigcup_{U \in f^{-1}(\mathcal{V})} U = \bigcup_{U \in \left\{ f^{-1}(V) \in \mathcal{P}(X) \mid V \in \mathcal{V} \right\}} U$$
$$= \bigcup_{V \in \mathcal{V}} f^{-1}(V).$$

This finishes the proof.

## Item 4: Interaction With Intersections of Families of Subsets

We have

$$\bigcap_{U \in f^{-1}(\mathcal{V})} U = \bigcap_{U \in \{f^{-1}(V) \in \mathcal{P}(X) \mid V \in \mathcal{V}\}} U$$
$$= \bigcap_{V \in \mathcal{V}} f^{-1}(V).$$

This finishes the proof.

Item 5: Interaction With Binary Unions

See [Pro25y].

Item 6: Interaction With Binary Intersections

See [Pro25w].

Item 7: Interaction With Differences

See [Pro25x].

Item 8: Interaction With Complements

See [Pro25j].

Item 9: Interaction With Symmetric Differences

We have

$$\begin{split} f^{-1}(U \triangle V) &= f^{-1}((U \cup V) \setminus (U \cap V)) \\ &= f^{-1}(U \cup V) \setminus f^{-1}(U \cap V) \\ &= f^{-1}(U) \cup f^{-1}(V) \setminus f^{-1}(U \cap V) \\ &= f^{-1}(U) \cup f^{-1}(V) \setminus f^{-1}(U) \cap f^{-1}(V) \\ &= f^{-1}(U) \triangle f^{-1}(V), \end{split}$$

where we have used:

- 1. Item 2 of Proposition 4.3.12.1.2 for the first equality.
- 2. Item 7 for the second equality.
- 3. Item 5 for the third equality.
- 4. Item 6 for the fourth equality.
- 5. Item 2 of Proposition 4.3.12.1.2 for the fifth equality.

This finishes the proof.

## Item 10: Interaction With Internal Homs of Powersets

We have

$$f^{-1}([U,V]_Y) \stackrel{\text{def}}{=} f^{-1}(U^c \cup V)$$

$$= f^{-1}(U^c) \cup f^{-1}(V)$$

$$= f^{-1}(U)^c \cup f^{-1}(V)$$

$$\stackrel{\text{def}}{=} [f^{-1}(U), f^{-1}(V)]_X,$$

where we have used:

- 1. Item 8 for the second equality.
- 2. Item 5 for the third equality.

Since  $\mathcal{P}(Y)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

Item 11: Preservation of Colimits

This follows from Item 2 and ??, ?? of ??.1

## Item 12: Preservation of Limits

This follows from Item 2 and ??, ?? of ??.<sup>2</sup>

Item 13: Symmetric Strict Monoidality With Respect to Unions

This follows from Item 11.

Item 14: Symmetric Strict Monoidality With Respect to Intersections

This follows from Item 12.

Item 15: Interaction With Coproducts

Omitted.

#### Item 16: Interaction With Products

Omitted.

<sup>1</sup>Reference: [Pro25y]. <sup>2</sup>Reference: [Pro25w].

#### PROPOSITION 4.6.2.1.5 ► PROPERTIES OF INVERSE IMAGES II

Let  $f: X \to Y$  be a function.

1. Functionality I. The assignment  $f \mapsto f^{-1}$  defines a function

$$(-)_{XY}^{-1} : \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(Y),\mathcal{P}(X)).$$

2. Functionality II. The assignment  $f \mapsto f^{-1}$  defines a function

$$(-)_{X,Y}^{-1} \colon \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(Y),\subset),(\mathcal{P}(X),\subset)).$$

3. *Interaction With Identities.* For each  $X \in \mathsf{Obj}(\mathsf{Sets})$ , we have

$$\mathrm{id}_X^{-1}=\mathrm{id}_{\mathcal{P}(X)}\;.$$

4. *Interaction With Composition*. For each pair of composable functions  $f: X \to Y$  and  $g: Y \to Z$ , we have

$$(g \circ f)^{-1} = f^{-1} \circ g^{-1},$$

$$\mathcal{P}(Z) \xrightarrow{g^{-1}} \mathcal{P}(Y)$$

$$\downarrow^{f^{-1}}$$

$$\mathcal{P}(X).$$

## Item 1: Functionality I

There is nothing to prove.

## Item 2: Functionality II

This follows from Item 1 of Proposition 4.6.2.1.3.

## Item 3: Interaction With Identities

This follows from Remark 4.6.2.1.2 and Categories, Item 5 of Proposition 11.1.4.1.2.

## Item 4: Interaction With Composition

This follows from Remark 4.6.2.1.2 and Categories, Item 2 of Proposition 11.1.4.1.2.

## 4.6.3 Codirect Images

Let  $f: X \to Y$  be a function.

#### **DEFINITION 4.6.3.1.1** ► CODIRECT IMAGES

The **codirect image function associated to** f is the function

$$f_* \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by<sup>1,2</sup>

$$f_*(U) \stackrel{\text{def}}{=} \left\{ y \in Y \middle| \begin{array}{l} \text{for each } x \in X, \text{ if we have} \\ f(x) = y, \text{ then } x \in U \end{array} \right\}$$
$$= \left\{ y \in Y \middle| \text{ we have } f^{-1}(y) \subset U \right\}$$

for each  $U \in \mathcal{P}(X)$ .

$$f_*(U) = f_!(U^{\mathsf{c}})^{\mathsf{c}}$$

$$\stackrel{\text{def}}{=} Y \setminus f_!(X \setminus U);$$

see Item 16 of Proposition 4.6.3.1.7.

<sup>&</sup>lt;sup>1</sup>Further Terminology: The set  $f_*(U)$  is called the **codirect image of** U **by** f.

<sup>&</sup>lt;sup>2</sup>We also have

#### NOTATION 4.6.3.1.2 ► FURTHER NOTATION FOR CODIRECT IMAGES

Sometimes one finds the notation

$$\forall_f \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

for  $f_!$ . This notation comes from the fact that the following statements are equivalent, where  $y \in Y$  and  $U \in \mathcal{P}(X)$ :

- We have  $y \in \forall_f(U)$ .
- For each  $x \in X$ , if y = f(x), then  $x \in U$ .

We will not make use of this notation elsewhere in Clowder.

#### WARNING 4.6.3.1.3 ► NOTATION FOR CODIRECT IMAGES IS CONFUSING



See Warning 4.6.1.1.3.

#### REMARK 4.6.3.1.4 ► UNWINDING DEFINITION 4.6.3.1.1

Identifying  $\mathcal{P}(X)$  with  $\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$  via Item 2 of Proposition 4.5.1.1.4, we see that the codirect image function associated to f is equivalently the function

$$f_* \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by

$$f_{*}(\chi_{U}) \stackrel{\text{def}}{=} \operatorname{Ran}_{f}(\chi_{U})$$

$$= \lim \left( \left( \underbrace{(-_{1})}_{\times} \stackrel{\longrightarrow}{\times} f \right) \stackrel{\operatorname{pr}}{\twoheadrightarrow} X \stackrel{\chi_{U}}{\longrightarrow} \left\{ \text{true, false} \right\} \right)$$

$$= \lim_{\substack{x \in X \\ f(x) = -_{1}}} (\chi_{U}(x))$$

$$= \bigwedge_{\substack{x \in X \\ f(x) = -_{1}}} (\chi_{U}(x)).$$

where we have used ?? for the second equality. In other words, we have

$$[f_*(\chi_U)](y) = \bigwedge_{\substack{x \in X \\ f(x) = y}} (\chi_U(x))$$

$$= \begin{cases} \mathsf{true} & \text{if, for each } x \in X \text{ such that} \\ & f(x) = y, \text{ we have } x \in U, \\ \mathsf{false} & \text{otherwise} \end{cases}$$
 
$$= \begin{cases} \mathsf{true} & \text{if } f^{-1}(y) \subset U \\ \mathsf{false} & \text{otherwise} \end{cases}$$

for each  $y \in Y$ .

#### **DEFINITION 4.6.3.1.5** $\blacktriangleright$ The Image and Complement Parts of $f_*$

Let *U* be a subset of X.<sup>1,2</sup>

1. The **image part of the codirect image**  $f_*(U)$  **of** U is the set  $f_{*,im}(U)$  defined by

$$f_{*,\text{im}}(U) \stackrel{\text{def}}{=} f_*(U) \cap \text{Im}(f)$$

$$= \left\{ y \in Y \middle| \begin{array}{l} \text{we have } f^{-1}(y) \subset U \\ \text{and } f^{-1}(y) \neq \emptyset. \end{array} \right\}.$$

2. The complement part of the codirect image  $f_*(U)$  of U is the set  $f_{*,cp}(U)$  defined by

$$f_{*,cp}(U) \stackrel{\text{def}}{=} f_*(U) \cap (Y \setminus \text{Im}(f))$$

$$= Y \setminus \text{Im}(f)$$

$$= \left\{ y \in Y \middle| \text{ we have } f^{-1}(y) \subset U \right\}$$

$$= \left\{ y \in Y \middle| f^{-1}(y) = \emptyset \right\}.$$

$$f_*(U) = f_{*,im}(U) \cup f_{*,cp}(U),$$

as

$$\begin{split} f_*(U) &= f_*(U) \cap Y \\ &= f_*(U) \cap (\operatorname{Im}(f) \cup (Y \setminus \operatorname{Im}(f))) \\ &= (f_*(U) \cap \operatorname{Im}(f)) \cup (f_*(U) \cap (Y \setminus \operatorname{Im}(f))) \\ &\stackrel{\text{def}}{=} f_{*,\operatorname{im}}(U) \cup f_{*,\operatorname{cp}}(U). \end{split}$$

<sup>&</sup>lt;sup>1</sup>Note that we have

<sup>2</sup>In terms of the meet computation of  $f_*(U)$  of Remark 4.6.3.1.4, namely

$$f_*(\chi_U) = \bigwedge_{\substack{x \in X \\ f(x) = -1}} (\chi_U(x)),$$

we see that  $f_{*,\text{im}}$  corresponds to meets indexed over nonempty sets, while  $f_{*,\text{cp}}$  corresponds to meets indexed over the empty set.

#### **EXAMPLE 4.6.3.1.6** ► **EXAMPLES OF CODIRECT IMAGES**

Here are some examples of codirect images.

1. *Multiplication by Two*. Consider the function  $f: \mathbb{N} \to \mathbb{N}$  given by

$$f(n) \stackrel{\text{def}}{=} 2n$$

for each  $n \in \mathbb{N}$ . Since f is injective, we have

$$f_{*,\text{im}}(U) = f_!(U)$$
  
 $f_{*,\text{cp}}(U) = \{\text{odd natural numbers}\}$ 

for any  $U \subset \mathbb{N}$ . In particular, we have

 $f_*(\{\text{even natural numbers}\}) = \mathbb{N}.$ 

2. *Parabolas*. Consider the function  $f: \mathbb{R} \to \mathbb{R}$  given by

$$f(x) \stackrel{\text{def}}{=} x^2$$

for each  $x \in \mathbb{R}$ . We have

$$f_{*,cp}(U) = \mathbb{R}_{<0}$$

for any  $U \subset \mathbb{R}$ . Moreover, since  $f^{-1}(x) = \{-\sqrt{x}, \sqrt{x}\}$ , we have e.g.:

$$f_{*,\text{im}}([0,1]) = \{0\},$$

$$f_{*,\text{im}}([-1,1]) = [0,1],$$

$$f_{*,\text{im}}([1,2]) = \emptyset,$$

$$f_{*,\text{im}}([-2,-1] \cup [1,2]) = [1,4].$$

3. *Circles*. Consider the function  $f: \mathbb{R}^2 \to \mathbb{R}$  given by

$$f(x,y) \stackrel{\text{def}}{=} x^2 + y^2$$

for each  $(x, y) \in \mathbb{R}^2$ . We have

$$f_{*,\mathrm{cp}}(U) = \mathbb{R}_{<0}$$

for any  $U \subset \mathbb{R}^2$ , and since

$$f^{-1}(r) = \begin{cases} \text{a circle of radius } r \text{ about the origin} & \text{if } r > 0, \\ \{(0,0)\} & \text{if } r = 0, \\ \emptyset & \text{if } r < 0, \end{cases}$$

we have e.g.:

$$f_{*,\text{im}}([-1,1] \times [-1,1]) = [0,1],$$
  
 $f_{*,\text{im}}(([-1,1] \times [-1,1]) \setminus [-1,1] \times \{0\}) = \emptyset.$ 

## PROPOSITION 4.6.3.1.7 ► PROPERTIES OF CODIRECT IMAGES I

Let  $f: X \to Y$  be a function.

1. Functoriality. The assignment  $U \mapsto f_*(U)$  defines a functor

$$f_* \colon (\mathcal{P}(X), \subset) \to (\mathcal{P}(Y), \subset).$$

In particular, for each  $U, V \in \mathcal{P}(X)$ , the following condition is satisfied:

- $(\star)$  If  $U \subset V$ , then  $f_*(U) \subset f_*(V)$ .
- 2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$id_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad id_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$
  
 $f_! \circ f^{-1} \hookrightarrow id_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow id_{\mathcal{P}(X)},$ 

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*\Big(f^{-1}(V)\Big),$$
  
 $f_!\Big(f^{-1}(V)\Big) \subset V, \qquad f^{-1}(f_*(U)) \subset U$ 

indexed by  $U \in \mathcal{P}(X)$  and  $V \in \mathcal{P}(Y)$ .

(b) Bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(Y)}(f_!(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f^{-1}(V)),$$
  
 $\operatorname{Hom}_{\mathcal{P}(X)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f_*(V)),$ 

natural in  $U \in \mathcal{P}(X)$  and  $V \in \mathcal{P}(Y)$  and (respectively)  $V \in \mathcal{P}(X)$  and  $U \in \mathcal{P}(Y)$ . In particular:

- i. The following conditions are equivalent:
  - A. We have  $f_!(U) \subset V$ .
  - B. We have  $U \subset f^{-1}(V)$ .
- ii. The following conditions are equivalent:
  - A. We have  $f^{-1}(U) \subset V$ .
  - B. We have  $U \subset f_*(V)$ .
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\cup \qquad \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_*(U) = \bigcup_{V \in f_*(\mathcal{U})} V$$

for each  $\mathcal{U} \in \mathcal{P}(X)$ , where  $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$ .

4. Interaction With Intersections of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcap_{U\in\mathcal{U}} f_*(U) = \bigcap_{V\in f_*(\mathcal{U})} V$$

for each  $\mathcal{U} \in \mathcal{P}(X)$ , where  $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$ .

5. *Interaction With Binary Unions.* Let  $f: X \to Y$  be a function. We have a natural transformation

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_* \times f_*} \mathcal{P}(Y) \times \mathcal{P}(Y) \\
\downarrow \qquad \qquad \qquad \qquad \downarrow \cup \\
\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

with components

$$f_*(U) \cup f_*(V) \subset f_*(U \cup V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

6. Interaction With Binary Intersections. The diagram

$$\begin{array}{cccc}
\mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_* \times f_*} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\
& & \downarrow & & \downarrow \\
\mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(U) \cap f_*(V) = f_*(U \cap V)$$

for each  $U, V \in \mathcal{P}(X)$ .

7. Interaction With Complements. The diagram

$$\mathcal{P}(X)^{\text{op}} \xrightarrow{f_!^{\text{op}}} \mathcal{P}(Y)^{\text{op}} \\
\stackrel{(-)^c}{\downarrow} \qquad \qquad \downarrow^{(-)^c} \\
\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U^{\mathsf{c}}) = f_!(U)^{\mathsf{c}}$$

for each  $U \in \mathcal{P}(X)$ .

8. *Interaction With Symmetric Differences*. We have a natural transformation

with components

$$f_*(U \triangle V) \subset f_*(U) \triangle f_*(V)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

9. *Interaction With Internal Homs of Powersets.* We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_*} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \\
[-1,-2]_X \downarrow \qquad \qquad \downarrow [-1,-2]_Y \\
\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

with components

$$[f_!(U), f_*(V)]_V \subset f_*([U, V]_X)$$

indexed by  $U, V \in \mathcal{P}(X)$ .

10. Lax Preservation of Colimits. We have an inclusion of sets

$$\bigcup_{i\in I} f_*(U_i) \subset f_*\left(\bigcup_{i\in I} U_i\right),\,$$

natural in  $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{\times I}$ . In particular, we have inclusions

$$f_*(U) \cup f_*(V) \hookrightarrow f_*(U \cup V),$$
  
 $\emptyset \hookrightarrow f_*(\emptyset),$ 

natural in  $U, V \in \mathcal{P}(X)$ .

11. Preservation of Limits. We have an equality of sets

$$f_*\left(\bigcap_{i\in I}U_i\right)=\bigcap_{i\in I}f_*(U_i),$$

natural in  $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{\times I}$ . In particular, we have equalities

$$f^{-1}(U \cap V) = f_*(U) \cap f^{-1}(V),$$
  
 $f_*(X) = Y,$ 

natural in  $U, V \in \mathcal{P}(X)$ .

12. *Symmetric Lax Monoidality With Respect to Unions*. The codirect image function of Item 1 has a symmetric lax monoidal structure

$$\left(f_*, f_*^{\otimes}, f_{*|\mathbb{1}}^{\otimes}\right) \colon (\mathcal{P}(X), \cup, \emptyset) \to (\mathcal{P}(Y), \cup, \emptyset),$$

being equipped with inclusions

$$f_{*|U,V}^{\otimes} \colon f_*(U) \cup f_*(V) \hookrightarrow f_*(U \cup V),$$
  
 $f_{*|\mathfrak{1}}^{\otimes} \colon \emptyset \hookrightarrow f_*(\emptyset),$ 

natural in  $U, V \in \mathcal{P}(X)$ .

13. Symmetric Strict Monoidality With Respect to Intersections. The direct image function of Item 1 has a symmetric strict monoidal structure

$$(f_*, f_*^{\otimes}, f_{*|1}^{\otimes}) \colon (\mathcal{P}(X), \cap, X) \to (\mathcal{P}(Y), \cap, Y),$$

being equipped with equalities

$$f_{*|U,V}^{\otimes} \colon f_*(U \cap V) \xrightarrow{=} f_*(U) \cap f_*(V),$$
$$f_{*|1}^{\otimes} \colon f_*(X) \xrightarrow{=} Y,$$

natural in  $U, V \in \mathcal{P}(X)$ .

14. *Interaction With Coproducts.* Let  $f: X \to X'$  and  $g: Y \to Y'$  be maps of sets. We have

$$(f \coprod g)_*(U \coprod V) = f_*(U) \coprod g_*(V)$$

for each  $U \in \mathcal{P}(X)$  and each  $V \in \mathcal{P}(Y)$ .

15. *Interaction With Products.* Let  $f: X \to X'$  and  $g: Y \to Y'$  be maps of sets. We have

$$(f\boxtimes_{X\times Y}g)_*(U\boxtimes_{X\times Y}V)=f_*(U)\boxtimes_{X'\times Y'}g_*(V)$$

for each  $U \in \mathcal{P}(X)$  and each  $V \in \mathcal{P}(Y)$ .

16. Relation to Direct Images. We have

$$f_*(U) = f_!(U^c)^c$$
  
=  $Y \setminus f_!(X \setminus U)$ 

for each  $U \in \mathcal{P}(X)$ .

17. Interaction With Injections. If f is injective, then we have

$$f_{*,\text{im}}(U) = f_!(U),$$
  
 $f_{*,\text{cp}}(U) = Y \setminus \text{Im}(f),$ 

and so

$$f_*(U) = f_{*,\text{im}}(U) \cup f_{*,\text{cp}}(U)$$
$$= f_!(U) \cup (Y \setminus \text{Im}(f))$$

for each  $U \in \mathcal{P}(X)$ .

18. *Interaction With Surjections*. If *f* is surjective, then we have

$$f_{*,\text{im}}(U) \subset f_!(U),$$
  
 $f_{*,\text{cp}}(U) = \emptyset,$ 

and so

$$f_*(U) \subset f_!(U)$$

for each  $U \in \mathcal{P}(X)$ .

#### PROOF 4.6.3.1.8 ► PROOF OF PROPOSITION 4.6.3.1.7

## Item 1: Functoriality

Omitted.

## Item 2: Triple Adjointness

This follows from Remark 4.6.1.1.4, Remark 4.6.2.1.2, Remark 4.6.3.1.4, and Kan Extensions, ?? of ??.

## Item 3: Interaction With Unions of Families of Subsets

We have

$$\bigcup_{V \in f_*(\mathcal{U})} V = \bigcup_{V \in \{f_*(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcup_{U \in \mathcal{U}} f_*(U).$$

This finishes the proof.

## Item 4: Interaction With Intersections of Families of Subsets

We have

$$\bigcap_{V \in f_*(\mathcal{U})} V = \bigcap_{V \in \{f_*(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$

$$= \bigcap_{U \in \mathcal{U}} f_*(U).$$

This finishes the proof.

## Item 5: Interaction With Binary Unions

We have

$$f_*(U) \cup f_*(V) = f_!(U^c)^c \cup f_!(V^c)^c$$

$$= (f_!(U^c) \cap f_!(V^c))^c$$

$$\subset (f_!(U^c \cap V^c))^c$$

$$= f_!((U \cup V)^c)^c$$

$$= f_*(U \cup V),$$

where:

1. We have used Item 16 for the first equality.

- 2. We have used Item 2 of Proposition 4.3.11.1.2 for the second equality.
- 3. We have used Item 6 of Proposition 4.6.1.1.5 for the third equality.
- 4. We have used Item 2 of Proposition 4.3.11.1.2 for the fourth equality.
- 5. We have used Item 16 for the last equality.

This finishes the proof.

## Item 6: Interaction With Binary Intersections

This follows from Item 11.

Item 7: Interaction With Complements

Omitted.

Item 8: Interaction With Symmetric Differences

Omitted.

## Item 9: Interaction With Internal Homs of Powersets

We have

$$\begin{split} \left[ f_!(U), f^!(V) \right]_X &\stackrel{\text{def}}{=} f_!(U)^{\mathsf{c}} \cup f_*(V) \\ &= f_*(U^{\mathsf{c}}) \cup f_*(V) \\ &\subset f_*(U^{\mathsf{c}} \cup V) \\ &\stackrel{\text{def}}{=} f_*([U, V]_X), \end{split}$$

where we have used:

- 1. Item 7 of Proposition 4.6.3.1.7 for the second equality.
- 2. Item 5 of Proposition 4.6.3.1.7 for the inclusion.

Since  $\mathcal{P}(X)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

Item 10: Lax Preservation of Colimits

Omitted.

## Item 11: Preservation of Limits

This follows from Item 2 and ??, ?? of ??.

Item 12: Symmetric Lax Monoidality With Respect to Unions

This follows from Item 10.

Item 13: Symmetric Strict Monoidality With Respect to Intersections

This follows from Item 11.

Item 14: Interaction With Coproducts

Omitted.

Item 15: Interaction With Products

Omitted.

Item 16: Relation to Direct Images

We claim that  $f_*(U) = Y \setminus f_!(X \setminus U)$ .

• The First Implication. We claim that

$$f_*(U) \subset Y \setminus f_!(X \setminus U).$$

Let  $y \in f_*(U)$ . We need to show that  $y \notin f_!(X \setminus U)$ , i.e. that there is no  $x \in X \setminus U$  such that f(x) = y.

This is indeed the case, as otherwise we would have  $x \in f^{-1}(y)$  and  $x \notin U$ , contradicting  $f^{-1}(y) \subset U$  (which holds since  $y \in f_*(U)$ ).

Thus  $y \in Y \setminus f_!(X \setminus U)$ .

• The Second Implication. We claim that

$$Y \setminus f_!(X \setminus U) \subset f_*(U)$$
.

Let  $y \in Y \setminus f_!(X \setminus U)$ . We need to show that  $y \in f_*(U)$ , i.e. that  $f^{-1}(y) \subset U$ .

Since  $y \notin f_!(X \setminus U)$ , there exists no  $x \in X \setminus U$  such that y = f(x), and hence  $f^{-1}(y) \subset U$ .

Thus  $y \in f_*(U)$ .

This finishes the proof of Item 16.

Item 17: Interaction With Injections

Omitted.

Item 18: Interaction With Surjections

Omitted.

#### PROPOSITION 4.6.3.1.9 ► PROPERTIES OF CODIRECT IMAGES II

Let  $f: X \to B$  be a function.

1. Functionality I. The assignment  $f \mapsto f_*$  defines a function

$$(-)_{Y|X|Y}: \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(X),\mathcal{P}(Y)).$$

2. Functionality II. The assignment  $f \mapsto f_*$  defines a function

$$(-)_{!|X,Y} : \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(X),\subset),(\mathcal{P}(Y),\subset)).$$

3. *Interaction With Identities.* For each  $X \in Obj(Sets)$ , we have

$$(\mathrm{id}_X)_* = \mathrm{id}_{\mathcal{P}(X)}$$
.

4. *Interaction With Composition*. For each pair of composable functions  $f: X \to Y$  and  $g: Y \to Z$ , we have

$$(g \circ f)_* = g_* \circ f_*,$$

$$(g \circ f)_* = g_* \circ f_*,$$

$$(g \circ f)_* \qquad g_*$$

$$\mathcal{P}(Z)$$

#### PROOF 4.6.3.1.10 ► PROOF OF PROPOSITION 4.6.3.1.9

## Item 1: Functionality I

There is nothing to prove.

## Item 2: Functionality II

This follows from Item 1 of Proposition 4.6.3.1.7.

## Item 3: Interaction With Identities

This follows from Remark 4.6.3.1.4 and Kan Extensions, ?? of ??.

## Item 4: Interaction With Composition

This follows from Remark 4.6.3.1.4 and Kan Extensions, ?? of ??.

#### 4.6.4 A Six-Functor Formalism for Sets

#### REMARK 4.6.4.1.1 ► A SIX-FUNCTOR FORMALISM FOR SETS

The assignment  $X \mapsto \mathcal{P}(X)$  together with the functors  $f_*$ ,  $f^{-1}$ , and  $f_!$  of Item 1 of Proposition 4.6.1.1.5, Item 1 of Proposition 4.6.2.1.3, and Item 1 of Proposition 4.6.3.1.7, and the functors

$$-_{1} \cap -_{2} \colon \mathcal{P}(X) \times \mathcal{P}(X) \to \mathcal{P}(X),$$
$$[-_{1}, -_{2}]_{X} \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

of Item 1 of Proposition 4.3.9.1.2 and Item 1 of Proposition 4.4.7.1.4 satisfy several properties reminiscent of a six functor formalism in the sense of ??. We collect these properties in Proposition 4.6.4.1.2 below.<sup>1</sup>

#### PROPOSITION 4.6.4.1.2 ► A SIX-FUNCTOR FORMALISM FOR SETS

Let *X* be a set.

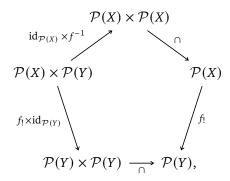
1. The Beck-Chevalley Condition. Let

$$\begin{array}{c|c} X \times_Z Y & \xrightarrow{\operatorname{pr}_2} Y \\ \operatorname{pr}_1 & & & \downarrow g \\ X & \xrightarrow{f} Z \end{array}$$

be a pullback diagram in Sets. We have

<sup>&</sup>lt;sup>1</sup>See also [nLa25].

2. The Projection Formula I. The diagram

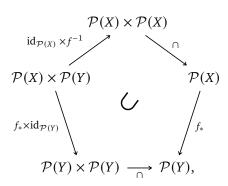


commutes, i.e. we have

$$f_!\Big(U\cap f^{-1}(V)\Big)=f_!(U)\cap V$$

for each  $U \in \mathcal{P}(X)$  and each  $V \in \mathcal{P}(Y)$ .

3. The Projection Formula II. We have a natural transformation



with components

$$f_*(U) \cap V \subset f_*\Big(U \cap f^{-1}(V)\Big)$$

indexed by  $U \in \mathcal{P}(X)$  and  $V \in \mathcal{P}(Y)$ .

4. Strong Closed Monoidality. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1}, \mathsf{op} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$[-1, -2]_{Y} \downarrow \qquad \qquad \downarrow [-1, -2]_{X}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_X,$$

natural in  $U, V \in \mathcal{P}(X)$ .

5. The External Tensor Product. We have an external tensor product

$$-1 \boxtimes_{X \times Y} -2: \mathcal{P}(X) \times \mathcal{P}(Y) \to \mathcal{P}(X \times Y)$$

given by

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \operatorname{pr}_{1}^{-1}(U) \cap \operatorname{pr}_{2}^{-1}(V)$$
$$= \{(u, v) \in X \times Y \mid u \in U \text{ and } v \in V\}.$$

This is the same map as the one in Item 5 of Proposition 4.4.1.1.4. Moreover, the following conditions are satisfied:

(a) *Interaction With Direct Images.* Let  $f: X \to X'$  and  $g: Y \to Y'$  be functions. The diagram

$$\begin{array}{c|c} \mathcal{P}(X) \times \mathcal{P}(Y) & \xrightarrow{f_1 \times g_1} & \mathcal{P}(X') \times \mathcal{P}(Y') \\ & \boxtimes_{X \times Y} & & & & & \boxtimes_{X' \times Y'} \\ & \mathcal{P}(X \times Y) & \xrightarrow{f_1 \times g_1} & \mathcal{P}(X' \times Y') & & \end{array}$$

commutes, i.e. we have

$$[f_! \times g_!](U \boxtimes_{X \times Y} V) = f_!(U) \boxtimes_{X' \times Y'} g_!(V)$$

for each  $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$ .

(b) *Interaction With Inverse Images.* Let  $f: X \to X'$  and  $g: Y \to Y'$  be functions. The diagram

$$\begin{array}{ccc} \mathcal{P}(X') \times \mathcal{P}(Y') & \xrightarrow{f^{-1} \times g^{-1}} & \mathcal{P}(X) \times \mathcal{P}(Y) \\ & \boxtimes_{X' \times Y'} & & & & & & & & & \\ & \boxtimes_{X' \times Y'} & & & & & & & & \\ & \mathcal{P}(X' \times Y') & \xrightarrow{f^{-1} \times g^{-1}} & \mathcal{P}(X \times Y) & & & & & \\ \end{array}$$

commutes, i.e. we have

$$[f^{-1} \times g^{-1}](U \boxtimes_{X' \times Y'} V) = f^{-1}(U) \boxtimes_{X \times Y} g^{-1}(V)$$

for each  $(U, V) \in \mathcal{P}(X') \times \mathcal{P}(Y')$ .

(c) Interaction With Codirect Images. Let  $f: X \to X'$  and  $g: Y \to Y'$  be functions. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(Y) \xrightarrow{f_* \times g_*} \mathcal{P}(X') \times \mathcal{P}(Y')$$

$$\boxtimes_{X \times Y} \qquad \qquad \qquad \qquad \qquad \downarrow^{\boxtimes_{X' \times Y'}}$$

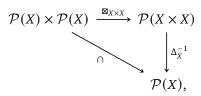
$$\mathcal{P}(X \times Y) \xrightarrow{f_* \times g_*} \mathcal{P}(X' \times Y')$$

commutes, i.e. we have

$$[f_* \times g_*](U \boxtimes_{X \times Y} V) = f_*(U) \boxtimes_{X' \times Y'} g_*(V)$$

for each  $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$ .

(d) Interaction With Diagonals. The diagram



i.e. we have

$$U\cap V=\Delta_X^{-1}(U\boxtimes_{X\times X}V)$$

for each  $U, V \in \mathcal{P}(X)$ .

6. The Dualisation Functor. We have a functor

$$D_X \colon \mathcal{P}(X)^{\mathsf{op}} \to \mathcal{P}(X)$$

given by

$$D_X(U) \stackrel{\text{def}}{=} [U, \emptyset]_X$$
$$\stackrel{\text{def}}{=} U^{\mathsf{c}}$$

for each  $U \in \mathcal{P}(X)$ , as in Item 5 of Proposition 4.4.7.1.4, satisfying the following conditions:

(a) Duality. We have

$$D_X(D_X(U)) = U, \qquad D_X \xrightarrow{\mathrm{id}_{\mathcal{P}(X)}} \mathcal{P}(X)$$

$$\mathcal{P}(X)$$

(b) Duality. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{\cap^{\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}} \\
\mathrm{id}_{\mathcal{P}(X)^{\mathsf{op}}} \times \mathcal{D}_{X} & & & & \\
\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) & & & & & \\
\end{array}$$

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) & & & & \\
& & & & & \\
\hline
\end{array}$$

commutes, i.e. we have

$$\underbrace{D_X(U\cap D_X(V))}_{\stackrel{\mathrm{def}}{=}[U\cap[V,\emptyset]_X,\emptyset]_X}=[U,V]_X$$

for each  $U, V \in \mathcal{P}(X)$ .

(c) Interaction With Direct Images. The diagram

$$\begin{array}{ccc} \mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_*^{\mathsf{op}}} & \mathcal{P}(Y)^{\mathsf{op}} \\ & & \downarrow & & \downarrow \\ \mathcal{P}(X) & \xrightarrow{f_!} & \mathcal{P}(Y) \end{array}$$

commutes, i.e. we have

$$f_!(D_X(U)) = D_Y(f_*(U))$$

for each  $U \in \mathcal{P}(X)$ .

(d) Interaction With Inverse Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(Y)^{\mathsf{op}} & \xrightarrow{f^{-1},\mathsf{op}} & \mathcal{P}(X)^{\mathsf{op}} \\
\downarrow^{D_{Y}} & & \downarrow^{D_{X}} \\
\mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$f^{-1}(D_Y(U)) = D_X(f^{-1}(U))$$

for each  $U \in \mathcal{P}(X)$ .

(e) Interaction With Codirect Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_!^{\mathsf{op}}} & \mathcal{P}(Y)^{\mathsf{op}} \\
\downarrow^{D_X} & & \downarrow^{D_Y} \\
\mathcal{P}(X) & \xrightarrow{f} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(D_X(U)) = D_Y(f_!(U))$$

for each  $U \in \mathcal{P}(X)$ .

#### PROOF 4.6.4.1.3 ► PROOF OF PROPOSITION 4.6.4.1.2

## Item 1: The Beck–Chevalley Condition

We have

$$[g^{-1} \circ f_{!}](U) \stackrel{\text{def}}{=} g^{-1}(f_{!}(U))$$

$$\stackrel{\text{def}}{=} \{y \in Y \mid g(y) \in f_{!}(U)\}$$

$$= \begin{cases} y \in Y \mid \text{there exists some } x \in U \\ \text{such that } f(x) = g(y) \end{cases}$$

$$= \begin{cases} y \in Y \mid \text{there exists some} \\ (x, y) \in \{(x, y) \in X \times_{Z} Y \mid x \in U\} \end{cases}$$

$$= \begin{cases} \text{there exists some} \\ (x, y) \in \{(x, y) \in X \times_{Z} Y \mid x \in U\} \\ \text{such that } y = y \end{cases}$$

$$= \begin{cases} y \in Y & \text{there exists some} \\ (x, y) \in \{(x, y) \in X \times_Z Y \mid x \in U\} \\ \text{such that } \operatorname{pr}_2(x, y) = y \end{cases}$$

$$\stackrel{\text{def}}{=} (\operatorname{pr}_2)_! (\{(x, y) \in X \times_Z Y \mid x \in U\})$$

$$= (\operatorname{pr}_2)_! (\{(x, y) \in X \times_Z Y \mid \operatorname{pr}_1(x, y) \in U\})$$

$$\stackrel{\text{def}}{=} (\operatorname{pr}_2)_! (\operatorname{pr}_1^{-1}(U))$$

$$\stackrel{\text{def}}{=} [(\operatorname{pr}_2)_! \circ \operatorname{pr}_1^{-1}](U)$$

for each  $U \in \mathcal{P}(X)$ . Therefore, we have

$$g^{-1} \circ f_! = (pr_2)_! \circ pr_1^{-1}$$
.

For the second equality, we have

$$[f^{-1} \circ g_!](U) \stackrel{\text{def}}{=} f^{-1}(g_!(U))$$

$$\stackrel{\text{def}}{=} \{x \in X \mid f(x) \in g_!(V)\}$$

$$= \left\{x \in X \mid \text{there exists some } y \in V \right\}$$

$$\text{such that } f(x) = g(y)$$

$$= \left\{x \in X \mid \text{there exists some } (x, y) \in \{(x, y) \in X \times_Z Y \mid y \in V\}\right\}$$

$$= \left\{x \in X \mid \text{there exists some } (x, y) \in \{(x, y) \in X \times_Z Y \mid y \in V\}\right\}$$

$$\text{such that } x = x$$

$$= \left\{x \in X \mid \text{there exists some } (x, y) \in \{(x, y) \in X \times_Z Y \mid y \in V\}\right\}$$

$$\text{such that } \text{pr}_1(x, y) = x$$

$$\stackrel{\text{def}}{=} (\text{pr}_1)_!(\{(x, y) \in X \times_Z Y \mid y \in V\})$$

$$= (\text{pr}_1)_!(\{(x, y) \in X \times_Z Y \mid \text{pr}_2(x, y) \in V\})$$

$$\stackrel{\text{def}}{=} (\text{pr}_1)_!(\text{pr}_2^{-1}(V))$$

$$\stackrel{\text{def}}{=} [(\text{pr}_1)_! \circ \text{pr}_2^{-1}](V)$$

for each  $V \in \mathcal{P}(Y)$ . Therefore, we have

$$f^{-1} \circ g_! = (pr_1)_! \circ pr_2^{-1}$$
.

This finishes the proof.

## Item 2: The Projection Formula I

We claim that

$$f_!(U) \cap V \subset f_!(U \cap f^{-1}(V)).$$

Indeed, we have

$$f_!(U) \cap V \subset f_!(U) \cap f_!(f^{-1}(V))$$
  
=  $f_!(U \cap f^{-1}(V)),$ 

where we have used:

- 1. Item 2 of Proposition 4.6.1.1.5 for the inclusion.
- 2. Item 6 of Proposition 4.6.1.1.5 for the equality.

Conversely, we claim that

$$f_!\Big(U\cap f^{-1}(V)\Big)\subset f_!(U)\cap V.$$

Indeed:

- 1. Let  $y \in f_!(U \cap f^{-1}(V))$ .
- 2. Since  $y \in f_!(U \cap f^{-1}(V))$ , there exists some  $x \in U \cap f^{-1}(V)$  such that f(x) = y.
- 3. Since  $x \in U \cap f^{-1}(V)$ , we have  $x \in U$ , and thus  $f(x) \in f_!(U)$ .
- 4. Since  $x \in U \cap f^{-1}(V)$ , we have  $x \in f^{-1}(V)$ , and thus  $f(x) \in V$ .
- 5. Since  $f(x) \in f_!(U)$  and  $f(x) \in V$ , we have  $f(x) \in f_!(U) \cap V$ .
- 6. But y = f(x), so  $y \in f_{!}(U) \cap V$ .
- 7. Thus  $f_!(U \cap f^{-1}(V)) \subset f_!(U) \cap V$ .

This finishes the proof.

Item 3: The Projection Formula II

We have

$$f_*(U) \cap V \subset f_*(U) \cap f_*(f^{-1}(V))$$
$$= f_*(U \cap f^{-1}(V)),$$

where we have used:

- 1. Item 2 of Proposition 4.6.3.1.7 for the inclusion.
- 2. Item 6 of Proposition 4.6.3.1.7 for the equality.

Since  $\mathcal{P}(Y)$  is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2).

## Item 4: Strong Closed Monoidality

This is a repetition of Item 19 of Proposition 4.4.7.1.4 and is proved there.

## Item 5: The External Tensor Product

We have

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \operatorname{pr}_{1}^{-1}(U) \cap \operatorname{pr}_{2}^{-1}(V)$$

$$\stackrel{\text{def}}{=} \left\{ (x, y) \in X \times Y \mid \operatorname{pr}_{1}(x, y) \in U \right\}$$

$$\cup \left\{ (x, y) \in X \times Y \mid \operatorname{pr}_{2}(x, y) \in V \right\}$$

$$= \left\{ (x, y) \in X \times Y \mid x \in U \right\}$$

$$\cup \left\{ (x, y) \in X \times Y \mid x \in U \text{ and } y \in V \right\}$$

$$\stackrel{\text{def}}{=} U \times V.$$

Next, we claim that Items 5a to 5d are indeed true:

- 1. *Proof of Item 5a*: This is a repetition of Item 16 of Proposition 4.6.1.1.5 and is proved there.
- 2. *Proof of Item 5b*: This is a repetition of Item 16 of Proposition 4.6.2.1.3 and is proved there.
- 3. *Proof of Item 5c:* This is a repetition of Item 15 of Proposition 4.6.3.1.7 and is proved there.

4. Proof of Item 5d: We have

$$\Delta_X^{-1}(U \boxtimes_{X \times X} V) \stackrel{\text{def}}{=} \{ x \in X \mid (x, x) \in U \boxtimes_{X \times X} V \}$$

$$= \{ x \in X \mid (x, x) \in \{ (u, v) \in X \times X \mid u \in U \text{ and } v \in V \} \}$$

$$= U \cap V.$$

This finishes the proof.

## Item 6: The Dualisation Functor

This is a repetition of Items 5 and 6 of Proposition 4.4.7.1.4 and is proved there.

# **Appendices**

## A Other Chapters

#### **Preliminaries**

- 1. Introduction
- 2. A Guide to the Literature

#### Sets

- 3. Sets
- 4. Constructions With Sets
- 5. Monoidal Structures on the Category of Sets
- 6. Pointed Sets
- 7. Tensor Products of Pointed Sets

#### Relations

- 8. Relations
- 9. Constructions With Relations

10. Conditions on Relations

## **Categories**

- 11. Categories
- 12. Presheaves and the Yoneda Lemma

## **Monoidal Categories**

13. Constructions With Monoidal Categories

#### **Bicategories**

14. Types of Morphisms in Bicategories

#### Extra Part

15. Notes

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