

# Categories

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This chapter contains some elementary material about categories, functors, and natural transformations. Notably, we discuss and explore:

1. Categories ([Section 11.1](#)).
2. Examples of categories ([Section 11.2](#)).
3. The quadruple adjunction  $\pi_0 \dashv (-)_{\text{disc}} \dashv \text{Obj} \dashv (-)_{\text{indisc}}$  between the category of categories and the category of sets ([Section 11.3](#)).
4. Groupoids, categories in which all morphisms admit inverses ([Section 11.4](#)).
5. Functors ([Section 11.5](#)).
6. The conditions one may impose on functors in decreasing order of importance:
  - (a) [Section 11.6](#) introduces the foundationally important conditions one may impose on functors, such as faithfulness, conservativity, essential surjectivity, etc.
  - (b) [Section 11.7](#) introduces more conditions one may impose on functors that are still important but less omni-present than those of [Section 11.6](#), such as being dominant, being a monomorphism, being pseudomononic, etc.
  - (c) [Section 11.8](#) introduces some rather rare or uncommon conditions one may impose on functors that are nevertheless still useful to explicit record in this chapter.

7. Natural transformations ([Section 11.9](#)).
8. The various categorical and 2-categorical structures formed by categories, functors, and natural transformations ([Section 11.10](#)).

This chapter is under active revision. TODO:

- Fix categories having an underlying set of objects by having them have an underlying setoid of objects (not necessarily by definition, as that'll likely be bothersome; at least [Section 11.3](#) should be fixed and several remarks should be added at several points). Related: [Definition 11.3.1.1.2](#)

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## II.I Categories

### II.I.I Foundations

**Definition II.I.I.I.** A **category**  $(C, \circ^C, \mathbb{1}^C)$  consists of:

- *Objects.* A class  $\text{Obj}(C)$  of **objects**.
- *Morphisms.* For each  $A, B \in \text{Obj}(C)$ , a class  $\text{Hom}_C(A, B)$ , called the **class of morphisms of  $C$  from  $A$  to  $B$** .
- *Identities.* For each  $A \in \text{Obj}(C)$ , a map of sets

$$\mathbb{1}_A^C: \text{pt} \rightarrow \text{Hom}_C(A, A),$$

called the **unit map of  $C$  at  $A$** , determining a morphism

$$\text{id}_A: A \rightarrow A$$

of  $C$ , called the **identity morphism of  $A$** .

- *Composition.* For each  $A, B, C \in \text{Obj}(C)$ , a map of sets

$$\circ_{A,B,C}^C: \text{Hom}_C(B, C) \times \text{Hom}_C(A, B) \rightarrow \text{Hom}_C(A, C),$$

called the **composition map of  $C$  at  $(A, B, C)$** .

such that the following conditions are satisfied:

1. *Associativity.* The diagram

$$\begin{array}{ccc}
 & \text{Hom}_C(C, D) \times (\text{Hom}_C(B, C) \times \text{Hom}_C(A, B)) & \\
 \alpha_{\text{Hom}_C(C,D), \text{Hom}_C(B,C), \text{Hom}_C(A,B)}^{\text{Sets}} \nearrow & & \searrow \text{id}_{\text{Hom}_C(C,D)} \times \circ_{A,B,C}^C \\
 (\text{Hom}_C(C, D) \times \text{Hom}_C(B, C)) \times \text{Hom}_C(A, B) & & \text{Hom}_C(C, D) \times \text{Hom}_C(A, C) \\
 \downarrow \circ_{B,C,D}^C \times \text{id}_{\text{Hom}_C(A,B)} & & \downarrow \circ_{A,C,D}^C \\
 \text{Hom}_C(B, D) \times \text{Hom}_C(A, B) & \xrightarrow{\circ_{A,B,D}^C} & \text{Hom}_C(A, D)
 \end{array}$$

commutes, i.e. for each composable triple  $(f, g, h)$  of morphisms of  $C$ , we have

$$(f \circ g) \circ h = f \circ (g \circ h).$$

2. *Left Unitality.* The diagram

$$\begin{array}{ccc}
 \text{pt} \times \text{Hom}_C(A, B) & & \\
 \downarrow \text{id}_B^C \times \text{id}_{\text{Hom}_C(A,B)} & \searrow \lambda_{\text{Hom}_C(A,B)}^{\text{Sets}} & \\
 \text{Hom}_C(B, B) \times \text{Hom}_C(A, B) & \xrightarrow{\circ_{A,B,B}^C} & \text{Hom}_C(A, B)
 \end{array}$$

commutes, i.e. for each morphism  $f: A \rightarrow B$  of  $C$ , we have

$$\text{id}_B \circ f = f.$$

3. *Right Unitality*. The diagram

$$\begin{array}{ccc}
 \text{Hom}_C(A, B) \times \text{pt} & & \\
 \text{id}_{\text{Hom}_C(A, B)} \times 1_A^C \downarrow & \nearrow \rho_{\text{Hom}_C(A, B)}^{\text{Sets}} & \\
 \text{Hom}_C(A, B) \times \text{Hom}_C(A, A) & \xrightarrow{\circ_{A, A, B}^C} & \text{Hom}_C(A, B)
 \end{array}$$

commutes, i.e. for each morphism  $f: A \rightarrow B$  of  $C$ , we have

$$f \circ \text{id}_A = f.$$

**Notation II.1.1.1.2.** Let  $C$  be a category.

1. We also write  $C(A, B)$  for  $\text{Hom}_C(A, B)$ .
2. We write  $\text{Mor}(C)$  for the class of all morphisms of  $C$ .

**Definition II.1.1.1.3.** Let  $\kappa$  be a regular cardinal. A category  $C$  is

1. **Locally small** if, for each  $A, B \in \text{Obj}(C)$ , the class  $\text{Hom}_C(A, B)$  is a set.
2. **Locally essentially small** if, for each  $A, B \in \text{Obj}(C)$ , the class

$$\text{Hom}_C(A, B) / \{\text{isomorphisms}\}$$

is a set.

3. **Small** if  $C$  is locally small and  $\text{Obj}(C)$  is a set.
4.  $\kappa$ -**Small** if  $C$  is locally small,  $\text{Obj}(C)$  is a set, and we have  $\#\text{Obj}(C) < \kappa$ .

## II.1.2 Subcategories

Let  $C$  be a category.

**Definition II.1.2.1.1.** A **subcategory** of  $C$  is a category  $\mathcal{A}$  satisfying the following conditions:

1. *Objects*. We have  $\text{Obj}(\mathcal{A}) \subset \text{Obj}(C)$ .

2. *Morphisms.* For each  $A, B \in \text{Obj}(\mathcal{A})$ , we have

$$\text{Hom}_{\mathcal{A}}(A, B) \subset \text{Hom}_C(A, B).$$

3. *Identities.* For each  $A \in \text{Obj}(\mathcal{A})$ , we have

$$1_A^{\mathcal{A}} = 1_A^C.$$

4. *Composition.* For each  $A, B, C \in \text{Obj}(\mathcal{A})$ , we have

$$\circ_{A,B,C}^{\mathcal{A}} = \circ_{A,B,C}^C.$$

**Definition II.1.2.1.2.** A subcategory  $\mathcal{A}$  of  $C$  is **full** if the canonical inclusion functor  $\mathcal{A} \rightarrow C$  is full, i.e. if, for each  $A, B \in \text{Obj}(\mathcal{A})$ , the inclusion

$$\iota_{A,B}: \text{Hom}_{\mathcal{A}}(A, B) \rightarrow \text{Hom}_C(A, B)$$

is surjective (and thus bijective).

**Definition II.1.2.1.3.** A subcategory  $\mathcal{A}$  of a category  $C$  is **strictly full** if it satisfies the following conditions:

1. *Fullness.* The subcategory  $\mathcal{A}$  is full.
2. *Closedness Under Isomorphisms.* The class  $\text{Obj}(\mathcal{A})$  is closed under isomorphisms.<sup>1</sup>

**Definition II.1.2.1.4.** A subcategory  $\mathcal{A}$  of  $C$  is **wide**<sup>2</sup> if  $\text{Obj}(\mathcal{A}) = \text{Obj}(C)$ .

### II.1.3 Skeletons of Categories

**Definition II.1.3.1.1.** A<sup>3</sup> **skeleton** of a category  $C$  is a full subcategory  $\text{Sk}(C)$  with one object from each isomorphism class of objects of  $C$ .

**Definition II.1.3.1.2.** A category  $C$  is **skeletal** if  $C \cong \text{Sk}(C)$ .<sup>4</sup>

<sup>1</sup>That is, given  $A \in \text{Obj}(\mathcal{A})$  and  $C \in \text{Obj}(C)$ , if  $C \cong A$ , then  $C \in \text{Obj}(\mathcal{A})$ .

<sup>2</sup>*Further Terminology:* Also called **lluf**.

<sup>3</sup>Due to **Item 3** of **Definition II.1.3.1.3**, which states that any two skeletons of a category are equivalent, we often refer to any such full subcategory  $\text{Sk}(C)$  of  $C$  as *the* skeleton of  $C$ .

<sup>4</sup>That is,  $C$  is **skeletal** if isomorphic objects of  $C$  are equal.

**Proposition II.1.3.1.3.** Let  $C$  be a category.

1. *Existence.* Assuming the axiom of choice,  $\text{Sk}(C)$  always exists.
2. *Pseudofunctoriality.* The assignment  $C \mapsto \text{Sk}(C)$  defines a pseudofunctor

$$\text{Sk}: \text{Cats}_2 \rightarrow \text{Cats}_2.$$

3. *Uniqueness Up to Equivalence.* Any two skeletons of  $C$  are equivalent.
4. *Inclusions of Skeletons Are Equivalences.* The inclusion

$$\iota_C: \text{Sk}(C) \rightarrow C$$

of a skeleton of  $C$  into  $C$  is an equivalence of categories.

*Proof.* **Item 1, Existence:** See [nLab23, Section “Existence of Skeletons of Categories”].

**Item 2, Pseudofunctoriality:** See [nLab23, Section “Skeletons as an Endo-Pseudofunctor on  $\mathbf{Cat}$ ”].

**Item 3, Uniqueness Up to Equivalence:** Clear.

**Item 4, Inclusions of Skeletons Are Equivalences:** Clear. □

## II.1.4 Precomposition and Postcomposition

Let  $C$  be a category and let  $A, B, C \in \text{Obj}(C)$ .

**Definition II.1.4.1.1.** Let  $f: A \rightarrow B$  and  $g: B \rightarrow C$  be morphisms of  $C$ .

1. The **precomposition function associated to  $f$**  is the function

$$f^*: \text{Hom}_C(B, C) \rightarrow \text{Hom}_C(A, C)$$

defined by

$$f^*(\phi) \stackrel{\text{def}}{=} \phi \circ f$$

for each  $\phi \in \text{Hom}_C(B, C)$ .

2. The **postcomposition function associated to  $g$**  is the function

$$g_*: \text{Hom}_C(A, B) \rightarrow \text{Hom}_C(A, C)$$

defined by

$$g_*(\phi) \stackrel{\text{def}}{=} g \circ \phi$$

for each  $\phi \in \text{Hom}_C(A, B)$ .



**Proposition II.1.4.1.2.** Let  $A, B, C, D \in \text{Obj}(C)$  and let  $f: A \rightarrow B$  and  $g: B \rightarrow C$  be morphisms of  $C$ .

1. *Interaction Between Precomposition and Postcomposition.* We have

$$\begin{array}{ccc}
 \text{Hom}_C(B, C) & \xrightarrow{g_*} & \text{Hom}_C(B, D) \\
 f^* \downarrow & & \downarrow f^* \\
 \text{Hom}_C(A, C) & \xrightarrow{g_*} & \text{Hom}_C(A, D).
 \end{array}$$

$g_* \circ f^* = f^* \circ g_*$

2. *Interaction With Composition I.* We have

$$\begin{array}{ccc}
 \text{Hom}_C(X, A) & \xrightarrow{f_*} & \text{Hom}_C(X, B) \\
 & \searrow (g \circ f)_* & \downarrow g_* \\
 & & \text{Hom}_C(X, C),
 \end{array}$$

$(g \circ f)^* = f^* \circ g^*$

$$\begin{array}{ccc}
 \text{Hom}_C(C, X) & \xrightarrow{g^*} & \text{Hom}_C(B, X) \\
 & \searrow (g \circ f)^* & \downarrow f^* \\
 & & \text{Hom}_C(A, X).
 \end{array}$$

$(g \circ f)_* = g_* \circ f_*$

3. *Interaction With Composition II.* We have

$$\begin{array}{ccc}
 \text{pt} \xrightarrow{[f]} \text{Hom}_C(A, B) & & \text{pt} \xrightarrow{[g]} \text{Hom}_C(B, C) \\
 \searrow [g \circ f] \quad \downarrow g_* & & \searrow [g \circ f] \quad \downarrow f^* \\
 \text{Hom}_C(A, C) & & \text{Hom}_C(A, C).
 \end{array}$$

$[g \circ f] = g_* \circ [f],$   
 $[g \circ f] = f^* \circ [g],$

4. *Interaction With Composition III.* We have

$$\begin{array}{ccc}
 \text{Hom}_C(B, C) \times \text{Hom}_C(A, B) & \xrightarrow{\circ_{A,B,C}^C} & \text{Hom}_C(A, C) \\
 \downarrow \text{id} \times f^* & & \downarrow f^* \\
 \text{Hom}_C(B, C) \times \text{Hom}_C(X, B) & \xrightarrow{\circ_{X,B,C}^C} & \text{Hom}_C(X, C), \\
 \\ 
 \text{Hom}_C(B, C) \times \text{Hom}_C(A, B) & \xrightarrow{\circ_{A,B,C}^C} & \text{Hom}_C(A, C) \\
 \downarrow g_* \times \text{id} & & \downarrow g_* \\
 \text{Hom}_C(B, D) \times \text{Hom}_C(A, B) & \xrightarrow{\circ_{A,B,D}^C} & \text{Hom}_C(A, D).
 \end{array}$$

$f^* \circ \circ_{A,B,C}^C = \circ_{X,B,C}^C \circ (f^* \times \text{id}),$   
 $g_* \circ \circ_{A,B,C}^C = \circ_{A,B,D}^C \circ (\text{id} \times g_*),$

5. *Interaction With Identities.* We have

$$\begin{aligned}
 (\text{id}_A)^* &= \text{id}_{\text{Hom}_C(A,B)}, \\
 (\text{id}_B)_* &= \text{id}_{\text{Hom}_C(A,B)}.
 \end{aligned}$$

*Proof.* **Item 1**, *Interaction Between Precomposition and Postcomposition*: Clear.

**Item 2**, *Interaction With Composition I*: Clear.

**Item 3**, *Interaction With Composition II*: Clear.

**Item 4**, *Interaction With Composition III*: Clear.

**Item 5**, *Interaction With Identities*: Clear. □

## II.2 Examples of Categories

### II.2.1 The Empty Category

**Example II.2.1.1.1.** The **empty category** is the category  $\emptyset_{\text{cat}}$  where

- *Objects.* We have

$$\text{Obj}(\emptyset_{\text{cat}}) \stackrel{\text{def}}{=} \emptyset.$$

- *Morphisms.* We have

$$\text{Mor}(\emptyset_{\text{cat}}) \stackrel{\text{def}}{=} \emptyset.$$

- *Identities and Composition.* Having no objects,  $\emptyset_{\text{cat}}$  has no unit nor composition maps.

## II.2.2 The Punctual Category

**Example II.2.2.1.1.** The **punctual category**<sup>5</sup> is the category  $\mathbf{pt}$  where

- *Objects.* We have

$$\mathrm{Obj}(\mathbf{pt}) \stackrel{\mathrm{def}}{=} \{\star\}.$$

- *Morphisms.* The unique Hom-set of  $\mathbf{pt}$  is defined by

$$\mathrm{Hom}_{\mathbf{pt}}(\star, \star) \stackrel{\mathrm{def}}{=} \{\mathrm{id}_{\star}\}.$$

- *Identities.* The unit map

$$1_{\star}^{\mathbf{pt}} : \mathbf{pt} \rightarrow \mathrm{Hom}_{\mathbf{pt}}(\star, \star)$$

of  $\mathbf{pt}$  at  $\star$  is defined by

$$\mathrm{id}_{\star}^{\mathbf{pt}} \stackrel{\mathrm{def}}{=} \mathrm{id}_{\star}.$$

- *Composition.* The composition map

$$\circ_{\star, \star, \star}^{\mathbf{pt}} : \mathrm{Hom}_{\mathbf{pt}}(\star, \star) \times \mathrm{Hom}_{\mathbf{pt}}(\star, \star) \rightarrow \mathrm{Hom}_{\mathbf{pt}}(\star, \star)$$

of  $\mathbf{pt}$  at  $(\star, \star, \star)$  is given by the bijection  $\mathbf{pt} \times \mathbf{pt} \cong \mathbf{pt}$ .

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<sup>5</sup>*Further Terminology:* Also called the **singleton category**.

### II.2.3 Monoids as One-Object Categories

**Example II.2.3.I.I.** We have an isomorphism of categories<sup>6</sup>

$$\text{Mon} \cong \text{pt} \times_{\text{Sets}} \text{Cats},$$

$$\begin{array}{ccc} \text{Mon} & \longrightarrow & \text{Cats} \\ \downarrow \lrcorner & & \downarrow \text{Obj} \\ \text{pt} & \xrightarrow{[\text{pt}]} & \text{Sets} \end{array}$$

via the delooping functor  $B: \text{Mon} \rightarrow \text{Cats}$  of ?? of ??, exhibiting monoids as exactly those categories having a single object.

*Proof.* Omitted. □

### II.2.4 Ordinal Categories

**Example II.2.4.I.I.** The  $n$ th ordinal category is the category  $\mathbf{n}$  where<sup>7</sup>

<sup>6</sup>This can be enhanced to an isomorphism of 2-categories

$$\text{Mon}_{2\text{disc}} \cong \text{pt}_{\text{bi}} \times_{\text{Sets}_{2\text{disc}}} \text{Cats}_{2,*},$$

$$\begin{array}{ccc} \text{Mon}_{2\text{disc}} & \longrightarrow & \text{Cats}_{2,*} \\ \downarrow \lrcorner & & \downarrow \text{Obj} \\ \text{pt}_{\text{bi}} & \xrightarrow{[\text{pt}]} & \text{Sets}_{2\text{disc}} \end{array}$$

between the discrete 2-category  $\text{Mon}_{2\text{disc}}$  on  $\text{Mon}$  and the 2-category of pointed categories with one object.

<sup>7</sup>In other words,  $\mathbf{n}$  is the category associated to the poset

$$[0] \rightarrow [1] \rightarrow \cdots \rightarrow [n-1] \rightarrow [n].$$

The category  $\mathbf{n}$  for  $n \geq 2$  may also be defined in terms of  $\mathbf{0}$  and joins (Constructions With Categories, ??): we have isomorphisms of categories

$$\begin{aligned} 1 &\cong \mathbf{0} \star \mathbf{0}, \\ 2 &\cong 1 \star \mathbf{0} \\ &\cong (\mathbf{0} \star \mathbf{0}) \star \mathbf{0}, \\ 3 &\cong 2 \star \mathbf{0} \\ &\cong (1 \star \mathbf{0}) \star \mathbf{0} \end{aligned}$$

- *Objects.* We have

$$\mathrm{Obj}(\mathfrak{n}) \stackrel{\mathrm{def}}{=} \{[0], \dots, [n]\}.$$

- *Morphisms.* For each  $[i], [j] \in \mathrm{Obj}(\mathfrak{n})$ , we have

$$\mathrm{Hom}_{\mathfrak{n}}([i], [j]) \stackrel{\mathrm{def}}{=} \begin{cases} \{\mathrm{id}_{[i]}\} & \text{if } [i] = [j], \\ \{[i] \rightarrow [j]\} & \text{if } [j] < [i], \\ \emptyset & \text{if } [j] > [i]. \end{cases}$$

- *Identities.* For each  $[i] \in \mathrm{Obj}(\mathfrak{n})$ , the unit map

$$1_{[i]}^{\mathfrak{n}} : \mathrm{pt} \rightarrow \mathrm{Hom}_{\mathfrak{n}}([i], [i])$$

of  $\mathfrak{n}$  at  $[i]$  is defined by

$$\mathrm{id}_{[i]}^{\mathfrak{n}} \stackrel{\mathrm{def}}{=} \mathrm{id}_{[i]}.$$

- *Composition.* For each  $[i], [j], [k] \in \mathrm{Obj}(\mathfrak{n})$ , the composition map

$$\circ_{[i],[j],[k]}^{\mathfrak{n}} : \mathrm{Hom}_{\mathfrak{n}}([j], [k]) \times \mathrm{Hom}_{\mathfrak{n}}([i], [j]) \rightarrow \mathrm{Hom}_{\mathfrak{n}}([i], [k])$$

of  $\mathfrak{n}$  at  $([i], [j], [k])$  is defined by

$$\begin{aligned} \mathrm{id}_{[i]} \circ \mathrm{id}_{[i]} &= \mathrm{id}_{[i]}, \\ ([j] \rightarrow [k]) \circ ([i] \rightarrow [j]) &= ([i] \rightarrow [k]). \end{aligned}$$

## II.2.5 The Walking Arrow

**Definition II.2.5.1.1.** The **walking arrow** is the category  $\mathbb{1}$  defined as the first ordinal category.

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$$\begin{aligned} &\cong ((0 \star 0) \star 0) \star 0, \\ \mathbb{4} &\cong 3 \star 0 \\ &\cong (2 \star 0) \star 0 \\ &\cong ((1 \star 0) \star 0) \star 0 \\ &\cong (((0 \star 0) \star 0) \star 0) \star 0, \end{aligned}$$

**Remark 11.2.5.1.2.** In detail, the walking arrow is the category  $\mathbf{1}$  where:

- *Objects.* We have  $\text{Obj}(\mathbf{1}) = \{0, 1\}$ .
- *Morphisms.* We have

$$\text{Hom}_{\mathbf{1}}(0, 0) = \{\text{id}_0\},$$

$$\text{Hom}_{\mathbf{1}}(1, 1) = \{\text{id}_1\},$$

$$\text{Hom}_{\mathbf{1}}(0, 1) = \{f_{01}\},$$

$$\text{Hom}_{\mathbf{1}}(1, 0) = \emptyset.$$

- *Identities and Composition.* The identities and composition of  $\mathbf{1}$  are completely determined by the unitality and associativity axioms for  $\mathbf{1}$ .

## 11.2.6 More Examples of Categories

**Example 11.2.6.1.1.** Here we list some of the other categories appearing throughout this work.

1. The category  $\text{Sets}_*$  of pointed sets of [Pointed Sets, Definition 6.1.3.1.1.](#)
2. The category  $\text{Rel}$  of sets and relations of [Relations, Definition 8.3.2.1.1.](#)
3. The category  $\text{Span}(\mathcal{A}, \mathcal{B})$  of spans from a set  $\mathcal{A}$  to a set  $\mathcal{B}$  of [??, ??](#).
4. The category  $\text{ISets}(K)$  of  $K$ -indexed sets of [Indexed Sets, ??](#).
5. The category  $\text{ISets}$  of indexed sets of [Indexed Sets, ??](#).
6. The category  $\text{FibSets}(K)$  of  $K$ -fibred sets of [Fibred Sets, ??](#).
7. The category  $\text{FibSets}$  of fibred sets of [Fibred Sets, ??](#).
8. Categories of functors  $\text{Fun}(\mathcal{C}, \mathcal{D})$  as in [Definition 11.10.1.1.1.](#)
9. The category of categories  $\text{Cats}$  of [Definition 11.10.2.1.1.](#)
10. The category of groupoids  $\text{Grpd}$  of [Definition 11.10.4.1.1.](#)

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and so on.

## II.2.7 Posetal Categories

**Definition II.2.7.1.1.** Let  $(X, \preceq_X)$  be a poset.

1. The **posetal category associated to**  $(X, \preceq_X)$  is the category  $X_{\text{pos}}$  where

- *Objects.* We have

$$\text{Obj}(X_{\text{pos}}) \stackrel{\text{def}}{=} X.$$

- *Morphisms.* For each  $a, b \in \text{Obj}(X_{\text{pos}})$ , we have

$$\text{Hom}_{X_{\text{pos}}}(a, b) \stackrel{\text{def}}{=} \begin{cases} \text{pt} & \text{if } a \preceq_X b, \\ \emptyset & \text{otherwise.} \end{cases}$$

- *Identities.* For each  $a \in \text{Obj}(X_{\text{pos}})$ , the unit map

$$1_a^{X_{\text{pos}}} : \text{pt} \rightarrow \text{Hom}_{X_{\text{pos}}}(a, a)$$

of  $X_{\text{pos}}$  at  $a$  is given by the identity map.

- *Composition.* For each  $a, b, c \in \text{Obj}(X_{\text{pos}})$ , the composition map

$$\circ_{a,b,c}^{X_{\text{pos}}} : \text{Hom}_{X_{\text{pos}}}(b, c) \times \text{Hom}_{X_{\text{pos}}}(a, b) \rightarrow \text{Hom}_{X_{\text{pos}}}(a, c)$$

of  $X_{\text{pos}}$  at  $(a, b, c)$  is defined as either the inclusion  $\emptyset \rightarrow \text{pt}$  or the identity map of  $\text{pt}$ , depending on whether we have  $a \preceq_X b$ ,  $b \preceq_X c$ , and  $a \preceq_X c$ .

2. A category  $C$  is **posetal**<sup>8</sup> if  $C$  is equivalent to  $X_{\text{pos}}$  for some poset  $(X, \preceq_X)$ .

**Proposition II.2.7.1.2.** Let  $(X, \preceq_X)$  be a poset and let  $C$  be a category.

1. *Functoriality.* The assignment  $(X, \preceq_X) \mapsto X_{\text{pos}}$  defines a functor

$$(-)_{\text{pos}} : \text{Pos} \rightarrow \text{Cats}.$$

2. *Fully Faithfulness.* The functor  $(-)_{\text{pos}}$  of **Item 1** is fully faithful.

---

<sup>8</sup>*Further Terminology:* Also called a **thin** category or a **(0, 1)-category**.

3. *Characterisations.* The following conditions are equivalent:
- (a) The category  $C$  is posetal.
  - (b) For each  $A, B \in \text{Obj}(C)$  and each  $f, g \in \text{Hom}_C(A, B)$ , we have  $f = g$ .
4. *Automatic Commutativity of Diagrams.* Every diagram in a posetal category commutes.

*Proof.* *Item 1, Functoriality:* Omitted.

*Item 2, Fully Faithfulness:* Omitted.

*Item 3, Characterisations:* Clear.

*Item 4, Automatic Commutativity of Diagrams:* This follows from the fact that if  $C$  is posetal, then there's at most one morphism between any two objects.  $\square$

## II.3 The Quadruple Adjunction With Sets

### II.3.1 Statement

Let  $C$  be a category.

**Proposition II.3.1.1.1.** We have a quadruple adjunction

$$(\pi_0 \dashv (-)_{\text{disc}} \dashv \text{Obj} \dashv (-)_{\text{indisc}}): \text{Sets} \begin{array}{c} \xrightarrow{\pi_0} \\ \dashv \uparrow \text{disc} \dashv \\ \text{Obj} \\ \dashv \downarrow \text{indisc} \dashv \\ \xrightarrow{(-)_{\text{indisc}}} \end{array} \text{Cats},$$

witnessed by bijections of sets

$$\begin{aligned} \text{Hom}_{\text{Sets}}(\pi_0(C), X) &\cong \text{Hom}_{\text{Cats}}(C, X_{\text{disc}}), \\ \text{Hom}_{\text{Cats}}(X_{\text{disc}}, C) &\cong \text{Hom}_{\text{Sets}}(X, \text{Obj}(C)), \\ \text{Hom}_{\text{Sets}}(\text{Obj}(C), X) &\cong \text{Hom}_{\text{Cats}}(C, X_{\text{indisc}}), \end{aligned}$$

natural in  $C \in \text{Obj}(\text{Cats})$  and  $X \in \text{Obj}(\text{Sets})$ , where



- The functor

$$\pi_0 : \mathbf{Cats} \rightarrow \mathbf{Sets},$$

the **connected components functor**, is the functor sending a category to its set of connected components of [Definition II.3.2.2.1](#).

- The functor

$$(-)_{\text{disc}} : \mathbf{Sets} \rightarrow \mathbf{Cats},$$

the **discrete category functor**, is the functor sending a set to its associated discrete category of [Item 1](#).

- The functor

$$\text{Obj} : \mathbf{Cats} \rightarrow \mathbf{Sets},$$

the **object functor**, is the functor sending a category to its set of objects.

- The functor

$$(-)_{\text{indisc}} : \mathbf{Sets} \rightarrow \mathbf{Cats},$$

the **indiscrete category functor**, is the functor sending a set to its associated indiscrete category of [Item 1](#).

*Proof.* Omitted. □

**Warning II.3.1.1.2.** (This is a stub, to be revised and expanded upon later.)

The discrete category functor of [Definition II.3.1.1.1](#) lifts to a 2-functor, but it fails to preserve 2-categorical colimits, and hence lacks a right 2-adjoint. For instance, the 2-pushout of  $\text{pt} \leftarrow S^0 \rightarrow \text{pt}$  in  $\mathbf{Sets}_{\text{ldisc}}$  is  $\text{pt}$ , but in  $\mathbf{Cats}_2$  it is given by  $\mathbf{B}\mathbb{Z}$ .

## II.3.2 Connected Components and Connected Categories

### II.3.2.1 Connected Components of Categories

Let  $C$  be a category.

**Definition II.3.2.1.1.** A **connected component** of  $C$  is a full subcategory  $\mathcal{I}$  of  $C$  satisfying the following conditions:<sup>9</sup>

<sup>9</sup>In other words, a **connected component** of  $C$  is an element of the set  $\text{Obj}(C)/\sim$  with  $\sim$  the equivalence relation generated by the relation  $\sim'$  obtained by declaring  $A \sim' B$  iff there exists a morphism of  $C$  from  $A$  to  $B$ .

1. *Non-Emptiness.* We have  $\text{Obj}(\mathcal{I}) \neq \emptyset$ .
2. *Connectedness.* There exists a zigzag of arrows between any two objects of  $\mathcal{I}$ .

### II.3.2.2 Sets of Connected Components of Categories

Let  $\mathcal{C}$  be a category.

**Definition II.3.2.2.1.** The **set of connected components of  $\mathcal{C}$**  is the set  $\pi_0(\mathcal{C})$  whose elements are the connected components of  $\mathcal{C}$ .

**Proposition II.3.2.2.2.** Let  $\mathcal{C}$  be a category.

1. *Functoriality.* The assignment  $\mathcal{C} \mapsto \pi_0(\mathcal{C})$  defines a functor

$$\pi_0: \text{Cats} \rightarrow \text{Sets}.$$

2. *Adjointness.* We have a quadruple adjunction

$$(\pi_0 \dashv (-)_{\text{disc}} \dashv \text{Obj} \dashv (-)_{\text{indisc}}): \text{Sets} \begin{array}{c} \xrightarrow{\pi_0} \\ \dashv \text{disc} \\ \dashv \text{Obj} \\ \dashv \text{indisc} \\ \xrightarrow{(-)_{\text{indisc}}} \end{array} \text{Cats}.$$

3. *Interaction With Groupoids.* If  $\mathcal{C}$  is a groupoid, then we have an isomorphism of categories

$$\pi_0(\mathcal{C}) \cong \text{K}(\mathcal{C}),$$

where  $\text{K}(\mathcal{C})$  is the set of isomorphism classes of  $\mathcal{C}$  of ??.

4. *Preservation of Colimits.* The functor  $\pi_0$  of **Item 1** preserves colimits. In particular, we have bijections of sets

$$\begin{aligned} \pi_0(\mathcal{C} \amalg \mathcal{D}) &\cong \pi_0(\mathcal{C}) \amalg \pi_0(\mathcal{D}), \\ \pi_0(\mathcal{C} \amalg_{\mathcal{E}} \mathcal{D}) &\cong \pi_0(\mathcal{C}) \amalg_{\pi_0(\mathcal{E})} \pi_0(\mathcal{D}), \\ \pi_0\left(\text{CoEq}\left(\mathcal{C} \begin{array}{c} \xrightarrow{F} \\ \xrightarrow{G} \end{array} \mathcal{D}\right)\right) &\cong \text{CoEq}\left(\pi_0(\mathcal{C}) \begin{array}{c} \xrightarrow{\pi_0(F)} \\ \xrightarrow{\pi_0(G)} \end{array} \pi_0(\mathcal{D})\right), \end{aligned}$$

natural in  $\mathcal{C}, \mathcal{D}, \mathcal{E} \in \text{Obj}(\text{Cats})$ .

5. *Symmetric Strong Monoidality With Respect to Coproducts.* The connected components functor of **Item 1** has a symmetric strong monoidal structure

$$\left(\pi_0, \pi_0^{\coprod}, \pi_{0|1}^{\coprod}\right): (\mathbf{Cats}, \coprod, \emptyset_{\text{cat}}) \rightarrow (\mathbf{Sets}, \coprod, \emptyset),$$

being equipped with isomorphisms

$$\begin{aligned} \pi_{0|C, \mathcal{D}}^{\coprod}: \pi_0(C) \coprod \pi_0(\mathcal{D}) &\xrightarrow{\sim} \pi_0(C \coprod \mathcal{D}), \\ \pi_{0|1}^{\coprod}: \emptyset &\xrightarrow{\sim} \pi_0(\emptyset_{\text{cat}}), \end{aligned}$$

natural in  $C, \mathcal{D} \in \text{Obj}(\mathbf{Cats})$ .

6. *Symmetric Strong Monoidality With Respect to Products.* The connected components functor of **Item 1** has a symmetric strong monoidal structure

$$\left(\pi_0, \pi_0^{\times}, \pi_{0|1}^{\times}\right): (\mathbf{Cats}, \times, \text{pt}) \rightarrow (\mathbf{Sets}, \times, \text{pt}),$$

being equipped with isomorphisms

$$\begin{aligned} \pi_{0|C, \mathcal{D}}^{\times}: \pi_0(C) \times \pi_0(\mathcal{D}) &\xrightarrow{\sim} \pi_0(C \times \mathcal{D}), \\ \pi_{0|1}^{\times}: \text{pt} &\xrightarrow{\sim} \pi_0(\text{pt}), \end{aligned}$$

natural in  $C, \mathcal{D} \in \text{Obj}(\mathbf{Cats})$ .

*Proof.* **Item 1**, *Functoriality*: Clear.

**Item 2**, *Adjointness*: This is proved in **Definition II.3.I.I.I.**

**Item 3**, *Interaction With Groupoids*: Clear.

**Item 4**, *Preservation of Colimits*: This follows from **Item 2** and ?? of ??.

**Item 5**, *Symmetric Strong Monoidality With Respect to Coproducts*: Clear.

**Item 6**, *Symmetric Strong Monoidality With Respect to Products*: Clear.  $\square$

### II.3.2.3 Connected Categories

**Definition II.3.2.3.I.** A category  $C$  is **connected** if  $\pi_0(C) \cong \text{pt}$ .<sup>10,11</sup>

<sup>10</sup>*Further Terminology*: A category is **disconnected** if it is not connected.

<sup>11</sup>*Example*: A groupoid is connected iff any two of its objects are isomorphic.

### II.3.3 Discrete Categories

**Definition II.3.3.1.1.** Let  $X$  be a set.

1. The **discrete category on  $X$**  is the category  $X_{\text{disc}}$  where

- *Objects.* We have

$$\text{Obj}(X_{\text{disc}}) \stackrel{\text{def}}{=} X.$$

- *Morphisms.* For each  $A, B \in \text{Obj}(X_{\text{disc}})$ , we have

$$\text{Hom}_{X_{\text{disc}}}(A, B) \stackrel{\text{def}}{=} \begin{cases} \text{id}_A & \text{if } A = B, \\ \emptyset & \text{if } A \neq B. \end{cases}$$

- *Identities.* For each  $A \in \text{Obj}(X_{\text{disc}})$ , the unit map

$$\mathbb{1}_A^{X_{\text{disc}}} : \text{pt} \rightarrow \text{Hom}_{X_{\text{disc}}}(A, A)$$

of  $X_{\text{disc}}$  at  $A$  is defined by

$$\text{id}_A^{X_{\text{disc}}} \stackrel{\text{def}}{=} \text{id}_A.$$

- *Composition.* For each  $A, B, C \in \text{Obj}(X_{\text{disc}})$ , the composition map

$$\circ_{A,B,C}^{X_{\text{disc}}} : \text{Hom}_{X_{\text{disc}}}(B, C) \times \text{Hom}_{X_{\text{disc}}}(A, B) \rightarrow \text{Hom}_{X_{\text{disc}}}(A, C)$$

of  $X_{\text{disc}}$  at  $(A, B, C)$  is defined by

$$\text{id}_A \circ \text{id}_A \stackrel{\text{def}}{=} \text{id}_A.$$

2. A category  $C$  is **discrete** if it is equivalent to  $X_{\text{disc}}$  for some set  $X$ .

**Proposition II.3.3.1.2.** Let  $X$  be a set.

1. *Functoriality.* The assignment  $X \mapsto X_{\text{disc}}$  defines a functor

$$(-)_{\text{disc}} : \text{Sets} \rightarrow \text{Cats}.$$

2. *Adjointness.* We have a quadruple adjunction

$$(\pi_0 \dashv (-)_{\text{disc}} \dashv \text{Obj} \dashv (-)_{\text{indisc}}): \text{Sets} \begin{array}{c} \xrightarrow{\pi_0} \\ \text{Cats.} \\ \xleftarrow{(-)_{\text{indisc}}} \end{array}$$

$\begin{array}{c} \perp \\ (-)_{\text{disc}} \\ \perp \\ \text{Obj} \\ \perp \\ (-)_{\text{indisc}} \end{array}$

3. *Symmetric Strong Monoidality With Respect to Coproducts.* The functor of **Item 1** has a symmetric strong monoidal structure

$$\left( (-)_{\text{disc}}, (-)_{\text{disc}}^{\amalg}, (-)_{\text{disc}|1}^{\amalg} \right): (\text{Sets}, \amalg, \emptyset) \rightarrow (\text{Cats}, \amalg, \emptyset_{\text{cat}}),$$

being equipped with isomorphisms

$$\begin{aligned} (-)_{\text{disc}|X,Y}^{\amalg}: X_{\text{disc}} \amalg Y_{\text{disc}} &\xrightarrow{\sim} (X \amalg Y)_{\text{disc}}, \\ (-)_{\text{disc}|1}^{\amalg}: \emptyset_{\text{cat}} &\xrightarrow{\sim} \emptyset_{\text{disc}}, \end{aligned}$$

natural in  $X, Y \in \text{Obj}(\text{Sets})$ .

4. *Symmetric Strong Monoidality With Respect to Products.* The functor of **Item 1** has a symmetric strong monoidal structure

$$\left( (-)_{\text{disc}}, (-)_{\text{disc}}^{\times}, (-)_{\text{disc}|1}^{\times} \right): (\text{Sets}, \times, \text{pt}) \rightarrow (\text{Cats}, \times, \text{pt}),$$

being equipped with isomorphisms

$$\begin{aligned} (-)_{\text{disc}|X,Y}^{\times}: X_{\text{disc}} \times Y_{\text{disc}} &\xrightarrow{\sim} (X \times Y)_{\text{disc}}, \\ (-)_{\text{disc}|1}^{\times}: \text{pt} &\xrightarrow{\sim} \text{pt}_{\text{disc}}, \end{aligned}$$

natural in  $X, Y \in \text{Obj}(\text{Sets})$ .

*Proof.* **Item 1, Functoriality:** Clear.

**Item 2, Adjointness:** This is proved in **Definition 11.3.1.1.1**.

**Item 3, Symmetric Strong Monoidality With Respect to Coproducts:** Clear.

**Item 4, Symmetric Strong Monoidality With Respect to Products:** Clear.  $\square$

## II.3.4 Indiscrete Categories

**Definition II.3.4.I.1.** Let  $X$  be a set.

1. The **indiscrete category on  $X$** <sup>12</sup> is the category  $X_{\text{indisc}}$  where

- *Objects.* We have

$$\text{Obj}(X_{\text{indisc}}) \stackrel{\text{def}}{=} X.$$

- *Morphisms.* For each  $A, B \in \text{Obj}(X_{\text{indisc}})$ , we have

$$\begin{aligned} \text{Hom}_{X_{\text{disc}}}(A, B) &\stackrel{\text{def}}{=} \{[A] \rightarrow [B]\} \\ &\cong \text{pt}. \end{aligned}$$

- *Identities.* For each  $A \in \text{Obj}(X_{\text{indisc}})$ , the unit map

$$\mathbb{1}_A^{X_{\text{indisc}}} : \text{pt} \rightarrow \text{Hom}_{X_{\text{indisc}}}(A, A)$$

of  $X_{\text{indisc}}$  at  $A$  is defined by

$$\text{id}_A^{X_{\text{indisc}}} \stackrel{\text{def}}{=} \{[A] \rightarrow [A]\}.$$

- *Composition.* For each  $A, B, C \in \text{Obj}(X_{\text{indisc}})$ , the composition map

$$\circ_{A,B,C}^{X_{\text{indisc}}} : \text{Hom}_{X_{\text{indisc}}}(B, C) \times \text{Hom}_{X_{\text{indisc}}}(A, B) \rightarrow \text{Hom}_{X_{\text{indisc}}}(A, C)$$

of  $X_{\text{disc}}$  at  $(A, B, C)$  is defined by

$$([B] \rightarrow [C]) \circ ([A] \rightarrow [B]) \stackrel{\text{def}}{=} ([A] \rightarrow [C]).$$

2. A category  $C$  is **indiscrete** if it is equivalent to  $X_{\text{indisc}}$  for some set  $X$ .

**Proposition II.3.4.I.2.** Let  $X$  be a set.

1. *Functoriality.* The assignment  $X \mapsto X_{\text{indisc}}$  defines a functor

$$(-)_{\text{indisc}} : \text{Sets} \rightarrow \text{Cats}.$$

---

<sup>12</sup>*Further Terminology:* Sometimes called the **chaotic category on  $X$** .

2. *Adjointness.* We have a quadruple adjunction

$$(\pi_0 \dashv (-)_{\text{disc}} \dashv \text{Obj} \dashv (-)_{\text{indisc}}): \text{Sets} \begin{array}{c} \xrightarrow{\pi_0} \\ \downarrow \perp \\ \xrightarrow{(-)_{\text{disc}}} \\ \downarrow \perp \\ \text{Obj} \\ \downarrow \perp \\ \xrightarrow{(-)_{\text{indisc}}} \end{array} \text{Cats.}$$

3. *Symmetric Strong Monoidality With Respect to Products.* The functor of **Item 1** has a symmetric strong monoidal structure

$$\left( (-)_{\text{indisc}}, (-)_{\text{indisc}}^{\times}, (-)_{\text{indisc}|\mathbb{1}}^{\times} \right): (\text{Sets}, \times, \text{pt}) \rightarrow (\text{Cats}, \times, \text{pt}),$$

being equipped with isomorphisms

$$\begin{aligned} (-)_{\text{indisc}|X,Y}^{\times}: X_{\text{indisc}} \times Y_{\text{indisc}} &\xrightarrow{\sim} (X \times Y)_{\text{indisc}}, \\ (-)_{\text{indisc}|\mathbb{1}}^{\times}: \text{pt} &\xrightarrow{\sim} \text{pt}_{\text{indisc}}, \end{aligned}$$

natural in  $X, Y \in \text{Obj}(\text{Sets})$ .

*Proof.* **Item 1, Functoriality:** Clear.

**Item 2, Adjointness:** This is proved in **Definition 11.3.1.1.1**.

**Item 3, Symmetric Strong Monoidality With Respect to Products:** Clear.  $\square$

## 11.4 Groupoids

### 11.4.1 Isomorphisms

Let  $C$  be a category.

**Definition 11.4.1.1.1.** A morphism  $f: A \rightarrow B$  of  $C$  is an **isomorphism** if there exists a morphism  $f^{-1}: B \rightarrow A$  of  $C$  such that

$$\begin{aligned} f \circ f^{-1} &= \text{id}_B, \\ f^{-1} \circ f &= \text{id}_A. \end{aligned}$$

**Notation 11.4.1.1.2.** We write  $\text{Iso}_C(A, B)$  for the set of all isomorphisms in  $C$  from  $A$  to  $B$ .

## II.4.2 Groupoids

**Definition II.4.2.1.1.** A **groupoid** is a category in which every morphism is an isomorphism.

**Example II.4.2.1.2.** The isomorphism of categories of [Definition II.2.3.1.1](#) restricts to an isomorphism

$$\begin{array}{ccc} \text{Grp} & \longrightarrow & \text{Grpd} \\ \downarrow & \lrcorner & \downarrow \text{Obj} \\ \text{pt} & \xrightarrow{[\text{pt}]} & \text{Sets} \end{array} \quad \text{Grp} \cong \text{pt}_{\text{Sets}} \times \text{Grpd},$$

where  $\text{Grpd}$  is the full subcategory of  $\text{Cats}$  spanned by the groupoids.

In other words, we have an identification

$$\{\text{Groups}\} \cong \{\text{One-object groupoids}\}.$$

## II.4.3 The Groupoid Completion of a Category

Let  $C$  be a category.

**Definition II.4.3.1.1.** The **groupoid completion** of  $C$ <sup>13</sup> is the pair  $(K_0(C), \iota_C)$  consisting of

- A groupoid  $K_0(C)$ ;
- A functor  $\iota_C : C \rightarrow K_0(C)$ ;

satisfying the following universal property:<sup>14</sup>

(UP) Given another such pair  $(\mathcal{G}, i)$ , there exists a unique functor  $K_0(C) \xrightarrow{\exists!} \mathcal{G}$  making the diagram

$$\begin{array}{ccc} & K_0(C) & \\ \iota_C \nearrow & \downarrow \exists! & \\ C & \xrightarrow{i} & \mathcal{G} \end{array}$$

<sup>13</sup>*Further Terminology:* Also called the **Grothendieck groupoid** of  $C$  or the **Grothendieck groupoid completion** of  $C$ .

<sup>14</sup>See [Item 5](#) of [Definition II.4.3.1.3](#) for an explicit construction.



commute.

**Construction II.4.3.I.2.** Concretely, the groupoid completion of  $\mathcal{C}$  is the Gabriel–Zisman localisation  $\text{Mor}(\mathcal{C})^{-1}\mathcal{C}$  of  $\mathcal{C}$  at the set  $\text{Mor}(\mathcal{C})$  of all morphisms of  $\mathcal{C}$ ; see Constructions With Categories, ??.

(To be expanded upon later on.)

*Proof.* Omitted. □

**Proposition II.4.3.I.3.** Let  $\mathcal{C}$  be a category.

1. *Functoriality.* The assignment  $\mathcal{C} \mapsto K_0(\mathcal{C})$  defines a functor

$$K_0 : \text{Cats} \rightarrow \text{Grpd}.$$

2. *2-Functoriality.* The assignment  $\mathcal{C} \mapsto K_0(\mathcal{C})$  defines a 2-functor

$$K_0 : \text{Cats}_2 \rightarrow \text{Grpd}_2.$$

3. *Adjointness.* We have an adjunction

$$(K_0 \dashv \iota) : \text{Cats} \begin{array}{c} \xrightarrow{K_0} \\ \perp \\ \xleftarrow{\iota} \end{array} \text{Grpd},$$

witnessed by a bijection of sets

$$\text{Hom}_{\text{Grpd}}(K_0(\mathcal{C}), \mathcal{G}) \cong \text{Hom}_{\text{Cats}}(\mathcal{C}, \mathcal{G}),$$

natural in  $\mathcal{C} \in \text{Obj}(\text{Cats})$  and  $\mathcal{G} \in \text{Obj}(\text{Grpd})$ , forming, together with the functor  $\text{Core}$  of **Item 1** of **Definition II.4.4.I.4**, a triple adjunction

$$(K_0 \dashv \iota \dashv \text{Core}) : \text{Cats} \begin{array}{c} \xrightarrow{K_0} \\ \perp \\ \xleftarrow{\iota} \\ \perp \\ \xrightarrow{\text{Core}} \end{array} \text{Grpd},$$

witnessed by bijections of sets

$$\text{Hom}_{\text{Grpd}}(K_0(\mathcal{C}), \mathcal{G}) \cong \text{Hom}_{\text{Cats}}(\mathcal{C}, \mathcal{G}),$$

$$\text{Hom}_{\text{Cats}}(\mathcal{G}, \mathcal{D}) \cong \text{Hom}_{\text{Grpd}}(\mathcal{G}, \text{Core}(\mathcal{D})),$$

natural in  $\mathcal{C}, \mathcal{D} \in \text{Obj}(\text{Cats})$  and  $\mathcal{G} \in \text{Obj}(\text{Grpd})$ .

4. *2-Adjointness*. We have a 2-adjunction

$$(K_0 \dashv \iota): \text{Cats} \begin{array}{c} \xrightarrow{K_0} \\ \xleftarrow{\iota} \\ \xleftarrow{\perp_2} \end{array} \text{Grpd},$$

witnessed by an isomorphism of categories

$$\text{Fun}(K_0(C), \mathcal{G}) \cong \text{Fun}(C, \mathcal{G}),$$

natural in  $C \in \text{Obj}(\text{Cats})$  and  $\mathcal{G} \in \text{Obj}(\text{Grpd})$ , forming, together with the 2-functor  $\text{Core}$  of **Item 2** of **Definition II.4.4.1.4**, a triple 2-adjunction

$$(K_0 \dashv \iota \dashv \text{Core}): \text{Cats} \begin{array}{c} \xrightarrow{K_0} \\ \xleftarrow{\iota} \\ \xleftarrow{\perp_2} \\ \xleftarrow{\text{Core}} \end{array} \text{Grpd},$$

witnessed by isomorphisms of categories

$$\begin{aligned} \text{Fun}(K_0(C), \mathcal{G}) &\cong \text{Fun}(C, \mathcal{G}), \\ \text{Fun}(\mathcal{G}, \mathcal{D}) &\cong \text{Fun}(\mathcal{G}, \text{Core}(\mathcal{D})), \end{aligned}$$

natural in  $C, \mathcal{D} \in \text{Obj}(\text{Cats})$  and  $\mathcal{G} \in \text{Obj}(\text{Grpd})$ .

5. *Interaction With Classifying Spaces*. We have an isomorphism of groupoids

$$K_0(C) \cong \Pi_{\leq 1}(|N_\bullet(C)|),$$

natural in  $C \in \text{Obj}(\text{Cats})$ ; i.e. the diagram

$$\begin{array}{ccc} \text{Cats} & \xrightarrow{K_0} & \text{Grp} \\ \downarrow N_\bullet & \uparrow \wr & \uparrow \Pi_{\leq 1} \\ \text{sSets} & \xrightarrow{|\cdot|} & \Pi \end{array}$$

commutes up to natural isomorphism.

6. *Symmetric Strong Monoidality With Respect to Coproducts*. The groupoid

completion functor of **Item 1** has a symmetric strong monoidal structure

$$\left( K_0, K_0^{\coprod}, K_0^{\coprod|1} \right) : (\mathbf{Cats}, \coprod, \emptyset_{\text{cat}}) \rightarrow (\mathbf{Grpd}, \coprod, \emptyset_{\text{cat}})$$

being equipped with isomorphisms

$$\begin{aligned} K_0^{\coprod|C, \mathcal{D}} : K_0(C) \coprod K_0(\mathcal{D}) &\xrightarrow{\sim} K_0(C \coprod \mathcal{D}), \\ K_0^{\coprod|1} : \emptyset_{\text{cat}} &\xrightarrow{\sim} K_0(\emptyset_{\text{cat}}), \end{aligned}$$

natural in  $C, \mathcal{D} \in \text{Obj}(\mathbf{Cats})$ .

7. *Symmetric Strong Monoidality With Respect to Products.* The groupoid completion functor of **Item 1** has a symmetric strong monoidal structure

$$\left( K_0, K_0^{\times}, K_0^{\times|1} \right) : (\mathbf{Cats}, \times, \text{pt}) \rightarrow (\mathbf{Grpd}, \times, \text{pt})$$

being equipped with isomorphisms

$$\begin{aligned} K_0^{\times|C, \mathcal{D}} : K_0(C) \times K_0(\mathcal{D}) &\xrightarrow{\sim} K_0(C \times \mathcal{D}), \\ K_0^{\times|1} : \text{pt} &\xrightarrow{\sim} K_0(\text{pt}), \end{aligned}$$

natural in  $C, \mathcal{D} \in \text{Obj}(\mathbf{Cats})$ .

*Proof.* **Item 1, Functoriality:** Omitted.

**Item 2, 2-Functoriality:** Omitted.

**Item 3, Adjointness:** Omitted.

**Item 4, 2-Adjointness:** Omitted.

**Item 5, Interaction With Classifying Spaces:** See Corollary 18.33 of <https://web.ma.utexas.edu/users/dafr/M392C-2012/Notes/lecture18.pdf>.

**Item 6, Symmetric Strong Monoidality With Respect to Coproducts:** Omitted.

**Item 7, Symmetric Strong Monoidality With Respect to Products:** Omitted.  $\square$

## II.4.4 The Core of a Category

Let  $C$  be a category.

**Definition II.4.4.I.I.** The **core** of  $C$  is the pair  $(\text{Core}(C), \iota_C)$  consisting of

- A groupoid  $\text{Core}(C)$ ;
- A functor  $\iota_C : \text{Core}(C) \rightarrow C$ ;

satisfying the following universal property:

(UP) Given another such pair  $(\mathcal{G}, i)$ , there exists a unique functor  $\mathcal{G} \xrightarrow{\exists!} \text{Core}(C)$  making the diagram

$$\begin{array}{ccc} & \text{Core}(C) & \\ \exists! \nearrow & \downarrow \iota_C & \\ \mathcal{G} & \xrightarrow{i} & C \end{array}$$

commute.

**Notation II.4.4.I.2.** We also write  $C^\simeq$  for  $\text{Core}(C)$ .

**Construction II.4.4.I.3.** The core of  $C$  is the wide subcategory of  $C$  spanned by the isomorphisms of  $C$ , i.e. the category  $\text{Core}(C)$  where<sup>15</sup>

1. *Objects.* We have

$$\text{Obj}(\text{Core}(C)) \stackrel{\text{def}}{=} \text{Obj}(C).$$

2. *Morphisms.* The morphisms of  $\text{Core}(C)$  are the isomorphisms of  $C$ .

*Proof.* This follows from the fact that functors preserve isomorphisms (Item 1 of Definition II.5.I.1.6).  $\square$

**Proposition II.4.4.I.4.** Let  $C$  be a category.

1. *Functoriality.* The assignment  $C \mapsto \text{Core}(C)$  defines a functor

$$\text{Core}: \text{Cats} \rightarrow \text{Grpd}.$$

2. *2-Functoriality.* The assignment  $C \mapsto \text{Core}(C)$  defines a 2-functor

$$\text{Core}: \text{Cats}_2 \rightarrow \text{Grpd}_2.$$

---

<sup>15</sup>*Slogan:* The groupoid  $\text{Core}(C)$  is the maximal subgroupoid of  $C$ .

3. *Adjointness.* We have an adjunction

$$(\iota \dashv \text{Core}): \text{Grpd} \begin{array}{c} \xrightarrow{\iota} \\ \perp \\ \xleftarrow{\text{Core}} \end{array} \text{Cats},$$

witnessed by a bijection of sets

$$\text{Hom}_{\text{Cats}}(\mathcal{G}, \mathcal{D}) \cong \text{Hom}_{\text{Grpd}}(\mathcal{G}, \text{Core}(\mathcal{D})),$$

natural in  $\mathcal{G} \in \text{Obj}(\text{Grpd})$  and  $\mathcal{D} \in \text{Obj}(\text{Cats})$ , forming, together with the functor  $K_0$  of **Item 1** of **Definition II.4.3.1.3**, a triple adjunction

$$(K_0 \dashv \iota \dashv \text{Core}): \text{Cats} \begin{array}{c} \xrightarrow{K_0} \\ \perp \\ \xleftarrow{\iota} \\ \perp \\ \xrightarrow{\text{Core}} \end{array} \text{Grpd},$$

witnessed by bijections of sets

$$\begin{aligned} \text{Hom}_{\text{Grpd}}(K_0(C), \mathcal{G}) &\cong \text{Hom}_{\text{Cats}}(C, \mathcal{G}), \\ \text{Hom}_{\text{Cats}}(\mathcal{G}, \mathcal{D}) &\cong \text{Hom}_{\text{Grpd}}(\mathcal{G}, \text{Core}(\mathcal{D})), \end{aligned}$$

natural in  $C, \mathcal{D} \in \text{Obj}(\text{Cats})$  and  $\mathcal{G} \in \text{Obj}(\text{Grpd})$ .

4. *2-Adjointness.* We have an adjunction

$$(\iota \dashv \text{Core}): \text{Grpd} \begin{array}{c} \xrightarrow{\iota} \\ \perp_2 \\ \xleftarrow{\text{Core}} \end{array} \text{Cats},$$

witnessed by an isomorphism of categories

$$\text{Fun}(\mathcal{G}, \mathcal{D}) \cong \text{Fun}(\mathcal{G}, \text{Core}(\mathcal{D})),$$

natural in  $\mathcal{G} \in \text{Obj}(\text{Grpd})$  and  $\mathcal{D} \in \text{Obj}(\text{Cats})$ , forming, together with the 2-functor  $K_0$  of **Item 2** of **Definition II.4.3.1.3**, a triple 2-adjunction

$$(K_0 \dashv \iota \dashv \text{Core}): \text{Cats} \begin{array}{c} \xrightarrow{K_0} \\ \perp_2 \\ \xleftarrow{\iota} \\ \perp_2 \\ \xrightarrow{\text{Core}} \end{array} \text{Grpd},$$

witnessed by isomorphisms of categories

$$\begin{aligned}\mathrm{Fun}(\mathbf{K}_0(C), \mathcal{G}) &\cong \mathrm{Fun}(C, \mathcal{G}), \\ \mathrm{Fun}(\mathcal{G}, \mathcal{D}) &\cong \mathrm{Fun}(\mathcal{G}, \mathrm{Core}(\mathcal{D})),\end{aligned}$$

natural in  $C, \mathcal{D} \in \mathrm{Obj}(\mathrm{Cats})$  and  $\mathcal{G} \in \mathrm{Obj}(\mathrm{Grpd})$ .

5. *Symmetric Strong Monoidality With Respect to Products.* The core functor of **Item 1** has a symmetric strong monoidal structure

$$(\mathrm{Core}, \mathrm{Core}^\times, \mathrm{Core}_1^\times): (\mathrm{Cats}, \times, \mathrm{pt}) \rightarrow (\mathrm{Grpd}, \times, \mathrm{pt})$$

being equipped with isomorphisms

$$\begin{aligned}\mathrm{Core}_{C, \mathcal{D}}^\times: \mathrm{Core}(C) \times \mathrm{Core}(\mathcal{D}) &\xrightarrow{\sim} \mathrm{Core}(C \times \mathcal{D}), \\ \mathrm{Core}_1^\times: \mathrm{pt} &\xrightarrow{\sim} \mathrm{Core}(\mathrm{pt}),\end{aligned}$$

natural in  $C, \mathcal{D} \in \mathrm{Obj}(\mathrm{Cats})$ .

6. *Symmetric Strong Monoidality With Respect to Coproducts.* The core functor of **Item 1** has a symmetric strong monoidal structure

$$(\mathrm{Core}, \mathrm{Core}^\amalg, \mathrm{Core}_1^\amalg): (\mathrm{Cats}, \amalg, \emptyset_{\mathrm{cat}}) \rightarrow (\mathrm{Grpd}, \amalg, \emptyset_{\mathrm{cat}})$$

being equipped with isomorphisms

$$\begin{aligned}\mathrm{Core}_{C, \mathcal{D}}^\amalg: \mathrm{Core}(C) \amalg \mathrm{Core}(\mathcal{D}) &\xrightarrow{\sim} \mathrm{Core}(C \amalg \mathcal{D}), \\ \mathrm{Core}_1^\amalg: \emptyset_{\mathrm{cat}} &\xrightarrow{\sim} \mathrm{Core}(\emptyset_{\mathrm{cat}}),\end{aligned}$$

natural in  $C, \mathcal{D} \in \mathrm{Obj}(\mathrm{Cats})$ .

*Proof.* **Item 1, Functoriality:** Omitted.

**Item 2, 2-Functoriality:** Omitted.

**Item 3, Adjointness:** Omitted.

**Item 4, 2-Adjointness:** Omitted.

**Item 5, Symmetric Strong Monoidality With Respect to Products:** Omitted.

**Item 6, Symmetric Strong Monoidality With Respect to Coproducts:** Omitted.  $\square$

## II.5 Functors

### II.5.1 Foundations

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.5.1.1.1.** A **functor**  $F: \mathcal{C} \rightarrow \mathcal{D}$  **from  $\mathcal{C}$  to  $\mathcal{D}$** <sup>16</sup> consists of:

1. *Action on Objects.* A map of sets

$$F: \text{Obj}(\mathcal{C}) \rightarrow \text{Obj}(\mathcal{D}),$$

called the **action on objects of  $F$** .

2. *Action on Morphisms.* For each  $A, B \in \text{Obj}(\mathcal{C})$ , a map

$$F_{A,B}: \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(F(A), F(B)),$$

called the **action on morphisms of  $F$  at  $(A, B)$** <sup>17</sup>.

satisfying the following conditions:

1. *Preservation of Identities.* For each  $A \in \text{Obj}(\mathcal{C})$ , the diagram

$$\begin{array}{ccc} \text{pt} & & \\ \downarrow \scriptstyle 1_A^{\mathcal{C}} & \searrow \scriptstyle 1_{F(A)}^{\mathcal{D}} & \\ \text{Hom}_{\mathcal{C}}(A, A) & \xrightarrow{\scriptstyle F_{A,A}} & \text{Hom}_{\mathcal{D}}(F(A), F(A)) \end{array}$$

commutes, i.e. we have

$$F(\text{id}_A) = \text{id}_{F(A)}.$$

2. *Preservation of Composition.* For each  $A, B, C \in \text{Obj}(\mathcal{C})$ , the diagram

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}}(B, C) \times \text{Hom}_{\mathcal{C}}(A, B) & \xrightarrow{\scriptstyle \circ_{A,B,C}^{\mathcal{C}}} & \text{Hom}_{\mathcal{C}}(A, C) \\ \downarrow \scriptstyle F_{B,C} \times F_{A,B} & & \downarrow \scriptstyle F_{A,C} \\ \text{Hom}_{\mathcal{D}}(F(B), F(C)) \times \text{Hom}_{\mathcal{D}}(F(A), F(B)) & \xrightarrow{\scriptstyle \circ_{F(A),F(B),F(C)}^{\mathcal{D}}} & \text{Hom}_{\mathcal{D}}(F(A), F(C)) \end{array}$$

<sup>16</sup> *Further Terminology:* Also called a **covariant functor**.

<sup>17</sup> *Further Terminology:* Also called **action on Hom-sets of  $F$  at  $(A, B)$** .

commutes, i.e. for each composable pair  $(g, f)$  of morphisms of  $\mathcal{C}$ , we have

$$F(g \circ f) = F(g) \circ F(f).$$

**Notation II.5.1.1.2.** Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories, and write  $\mathcal{C}^{\text{op}}$  for the opposite category of  $\mathcal{C}$  of Constructions With Categories, ??.

1. Given a functor

$$F: \mathcal{C} \rightarrow \mathcal{D},$$

we also write  $F_A$  for  $F(A)$ .

2. Given a functor

$$F: \mathcal{C}^{\text{op}} \rightarrow \mathcal{D},$$

we also write  $F^A$  for  $F(A)$ .

3. Given a functor

$$F: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{D},$$

we also write  $F_{A,B}$  for  $F(A, B)$ .

4. Given a functor

$$F: \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{D},$$

we also write  $F_B^A$  for  $F(A, B)$ .

We employ a similar notation for morphisms, writing e.g.  $F_f$  for  $F(f)$  given a functor  $F: \mathcal{C} \rightarrow \mathcal{D}$ .

**Notation II.5.1.1.3.** Following the notation  $\llbracket x \mapsto f(x) \rrbracket$  for a function  $f: X \rightarrow Y$  introduced in [Sets, Definition 3.1.1.2](#), we will sometimes denote a functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  by

$$F \stackrel{\text{def}}{=} \llbracket A \mapsto F(A) \rrbracket,$$

specially when the action on morphisms of  $F$  is clear from its action on objects.

**Example II.5.1.1.4.** The **identity functor** of a category  $\mathcal{C}$  is the functor  $\text{id}_{\mathcal{C}}: \mathcal{C} \rightarrow \mathcal{C}$  where

1. *Action on Objects.* For each  $A \in \text{Obj}(\mathcal{C})$ , we have

$$\text{id}_{\mathcal{C}}(A) \stackrel{\text{def}}{=} A.$$



2. *Action on Morphisms.* For each  $A, B \in \text{Obj}(C)$ , the action on morphisms

$$(\text{id}_C)_{A,B}: \text{Hom}_C(A, B) \rightarrow \underbrace{\text{Hom}_C(\text{id}_C(A), \text{id}_C(B))}_{\stackrel{\text{def}}{=} \text{Hom}_C(A, B)}$$

of  $\text{id}_C$  at  $(A, B)$  is defined by

$$(\text{id}_C)_{A,B} \stackrel{\text{def}}{=} \text{id}_{\text{Hom}_C(A, B)}.$$

*Proof. Preservation of Identities:* We have  $\text{id}_C(\text{id}_A) \stackrel{\text{def}}{=} \text{id}_A$  for each  $A \in \text{Obj}(C)$  by definition.

*Preservation of Compositions:* For each composable pair  $A \xrightarrow{f} B \xrightarrow{g} C$  of morphisms of  $C$ , we have

$$\begin{aligned} \text{id}_C(g \circ f) &\stackrel{\text{def}}{=} g \circ f \\ &\stackrel{\text{def}}{=} \text{id}_C(g) \circ \text{id}_C(f). \end{aligned}$$

This finishes the proof.  $\square$

**Definition II.5.I.1.5.** The **composition** of two functors  $F: C \rightarrow D$  and  $G: D \rightarrow E$  is the functor  $G \circ F$  where

- *Action on Objects.* For each  $A \in \text{Obj}(C)$ , we have

$$[G \circ F](A) \stackrel{\text{def}}{=} G(F(A)).$$

- *Action on Morphisms.* For each  $A, B \in \text{Obj}(C)$ , the action on morphisms

$$(G \circ F)_{A,B}: \text{Hom}_C(A, B) \rightarrow \text{Hom}_E(G_{F_A}, G_{F_B})$$

of  $G \circ F$  at  $(A, B)$  is defined by

$$[G \circ F](f) \stackrel{\text{def}}{=} G(F(f)).$$

*Proof. Preservation of Identities:* For each  $A \in \text{Obj}(C)$ , we have

$$\begin{aligned} G_{F_{\text{id}_A}} &= G_{\text{id}_{F_A}} && \text{(functoriality of } F) \\ &= \text{id}_{G_{F_A}} && \text{(functoriality of } G) \end{aligned}$$

*Preservation of Composition:* For each composable pair  $(g, f)$  of morphisms of  $\mathcal{C}$ , we have

$$\begin{aligned} G_{F_g \circ f} &= G_{F_g \circ F_f} && \text{(functoriality of } F) \\ &= G_{F_g} \circ G_{F_f}. && \text{(functoriality of } G) \end{aligned}$$

This finishes the proof.  $\square$

**Proposition II.5.1.1.6.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor.

1. *Preservation of Isomorphisms.* If  $f$  is an isomorphism in  $\mathcal{C}$ , then  $F(f)$  is an isomorphism in  $\mathcal{D}$ .<sup>18</sup>

*Proof.* **Item 1, Preservation of Isomorphisms:** Indeed, we have

$$\begin{aligned} F(f)^{-1} \circ F(f) &= F(f^{-1} \circ f) \\ &= F(\text{id}_A) \\ &= \text{id}_{F(A)} \end{aligned}$$

and

$$\begin{aligned} F(f) \circ F(f)^{-1} &= F(f \circ f^{-1}) \\ &= F(\text{id}_B) \\ &= \text{id}_{F(B)}, \end{aligned}$$

showing  $F(f)$  to be an isomorphism.  $\square$

## II.5.2 Contravariant Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories, and let  $\mathcal{C}^{\text{op}}$  denote the opposite category of  $\mathcal{C}$  of Constructions With Categories, ??.

**Definition II.5.2.1.1.** A **contravariant functor** from  $\mathcal{C}$  to  $\mathcal{D}$  is a functor from  $\mathcal{C}^{\text{op}}$  to  $\mathcal{D}$ .

**Remark II.5.2.1.2.** In detail, a **contravariant functor** from  $\mathcal{C}$  to  $\mathcal{D}$  consists of:

<sup>18</sup>When the converse holds, we call  $F$  *conservative*, see [Definition II.6.4.1.1](#).

1. *Action on Objects.* A map of sets

$$F: \text{Obj}(C) \rightarrow \text{Obj}(\mathcal{D}),$$

called the **action on objects of  $F$** .

2. *Action on Morphisms.* For each  $A, B \in \text{Obj}(C)$ , a map

$$F_{A,B}: \text{Hom}_C(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(F(B), F(A)),$$

called the **action on morphisms of  $F$  at  $(A, B)$** .

satisfying the following conditions:

1. *Preservation of Identities.* For each  $A \in \text{Obj}(C)$ , the diagram

$$\begin{array}{ccc} \text{pt} & & \\ \downarrow \mathbb{1}_A^C & \searrow \mathbb{1}_{F(A)}^{\mathcal{D}} & \\ \text{Hom}_C(A, A) & \xrightarrow{F_{A,A}} & \text{Hom}_{\mathcal{D}}(F(A), F(A)) \end{array}$$

commutes, i.e. we have

$$F(\text{id}_A) = \text{id}_{F(A)}.$$

2. *Preservation of Composition.* For each  $A, B, C \in \text{Obj}(C)$ , the diagram

$$\begin{array}{ccc} & \text{Hom}_{\mathcal{D}}(F(C), F(B)) \times \text{Hom}_{\mathcal{D}}(F(B), F(A)) & \\ & \nearrow F_{B,C} \times F_{A,B} & \searrow \sigma_{\text{Hom}_{\mathcal{D}}(F(C), F(B)), \text{Hom}_{\mathcal{D}}(F(B), F(A))}^{\text{Sets}} \\ \text{Hom}_C(B, C) \times \text{Hom}_C(A, B) & & \text{Hom}_{\mathcal{D}}(F(B), F(A)) \times \text{Hom}_{\mathcal{D}}(F(C), F(B)) \\ \downarrow \circ_{A,B,C}^C & & \downarrow \circ_{F(C), F(B), F(A)}^{\mathcal{D}} \\ \text{Hom}_C(A, C) & \xrightarrow{F_{A,C}} & \text{Hom}_{\mathcal{D}}(F(C), F(A)) \end{array}$$

commutes, i.e. for each composable pair  $(g, f)$  of morphisms of  $C$ , we have

$$F(g \circ f) = F(f) \circ F(g).$$

**Remark 11.5.2.1.3.** Throughout this work we will not use the term “contravariant” functor, speaking instead simply of functors  $F: C^{\text{op}} \rightarrow \mathcal{D}$ . We will usually, however, write

$$F_{A,B}: \text{Hom}_C(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(F(B), F(A))$$

for the action on morphisms

$$F_{A,B}: \text{Hom}_{C^{\text{op}}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(F(A), F(B))$$

of  $F$ , as well as write  $F(g \circ f) = F(f) \circ F(g)$ .

### 11.5.3 Forgetful Functors

**Definition 11.5.3.1.1.** There isn’t a precise definition of a **forgetful functor**.

**Remark 11.5.3.1.2.** Despite there not being a formal or precise definition of a forgetful functor, the term is often very useful in practice, similarly to the word “canonical”. The idea is that a “forgetful functor” is a functor that forgets structure or properties, and is best explained through examples, such as the ones below (see [Definitions 11.5.3.1.3](#) and [11.5.3.1.4](#)).

**Example 11.5.3.1.3.** Examples of forgetful functors that forget structure include:

1. *Forgetting Group Structures.* The functor  $\text{Grp} \rightarrow \text{Sets}$  sending a group  $(G, \mu_G, \eta_G)$  to its underlying set  $G$ , forgetting the multiplication and unit maps  $\mu_G$  and  $\eta_G$  of  $G$ .
2. *Forgetting Topologies.* The functor  $\Pi \rightarrow \text{Sets}$  sending a topological space  $(X, T_X)$  to its underlying set  $X$ , forgetting the topology  $T_X$ .
3. *Forgetting Fibrations.* The functor  $\text{FibSets}(K) \rightarrow \text{Sets}$  sending a  $K$ -fibred set  $\phi_X: X \rightarrow K$  to the set  $X$ , forgetting the map  $\phi_X$  and the base set  $K$ .

**Example 11.5.3.1.4.** Examples of forgetful functors that forget properties include:

1. *Forgetting Commutativity.* The inclusion functor  $\iota: \text{CMon} \rightarrow \text{Mon}$  which forgets the property of being commutative.
2. *Forgetting Inverses.* The inclusion functor  $\iota: \text{Grp} \rightarrow \text{Mon}$  which forgets the property of having inverses.

**Notation 11.5.3.1.5.** Throughout this work, we will denote forgetful functors that forget structure by  $\text{忘}$ , e.g. as in

$$\text{忘}: \text{Grp} \rightarrow \text{Sets}.$$

The symbol  $\text{忘}$ , pronounced *wasureru* (see [Item 1](#) of [Definition 11.5.3.1.6](#) below), means *to forget*, and is a kanji found in the following words in Japanese and Chinese:

1. 忘れる, transcribed as *wasureru*, meaning *to forget*.
2. 忘却関手, transcribed as *boukyaku kanshu*, meaning *forgetful functor*.
3. 忘记 or 忘記, transcribed as *wàngjì*, meaning *to forget*.
4. 遗忘函子 or 遺忘函子, transcribed as *yíwàng hánzǐ*, meaning *forgetful functor*.

**Remark 11.5.3.1.6.** Here we collect the pronunciation of the words in [Definition 11.5.3.1.5](#) for accuracy and completeness.

1. Pronunciation of 忘れる:
  - See [here](#).
  - IPA broad transcription: [wäsureru].
  - IPA narrow transcription: [ʷäsi̯r̥ɐɹ̥u̯].
2. Pronunciation of 忘却関手: Pronunciation:
  - See [here](#).
  - IPA broad transcription: [bɔ:käku kãũɕu].
  - IPA narrow transcription: [bɔ:käku̯ kãũɕu̯].
3. Pronunciation of 忘记:
  - See [here](#).
  - Broad IPA transcription: [waŋtɕi].
  - Sinological IPA transcription: [waŋ<sup>51-53</sup>tɕi<sup>51</sup>].
4. Pronunciation of 遗忘函子:

- See [here](#).
- Broad IPA transcription: [iwaŋ xänfʂzi].
- Sinological IPA transcription: [i<sup>35</sup>waŋ<sup>51</sup> xän<sup>35</sup>fʂz<sup>214-21(4)</sup>].

## II.5.4 The Natural Transformation Associated to a Functor

**Definition II.5.4.1.1.** Every functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  defines a natural transformation<sup>19</sup>

$$F^\dagger: \text{Hom}_{\mathcal{C}} \Rightarrow \text{Hom}_{\mathcal{D}} \circ (F^{\text{op}} \times F),$$

$$\begin{array}{ccc} \mathcal{C}^{\text{op}} \times \mathcal{C} & \xrightarrow{F^{\text{op}} \times F} & \mathcal{D}^{\text{op}} \times \mathcal{D} \\ & \searrow \text{Hom}_{\mathcal{C}} \quad \xRightarrow{F^\dagger} \quad \swarrow \text{Hom}_{\mathcal{D}} & \\ & \text{Sets,} & \end{array}$$

called the **natural transformation associated to  $F$** , consisting of the collection

$$\left\{ F_{A,B}^\dagger: \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(F_A, F_B) \right\}_{(A,B) \in \text{Obj}(\mathcal{C}^{\text{op}} \times \mathcal{C})}$$

with

$$F_{A,B}^\dagger \stackrel{\text{def}}{=} F_{A,B}.$$

*Proof.* The naturality condition for  $F^\dagger$  is the requirement that for each morphism

$$(\phi, \psi): (X, Y) \rightarrow (A, B)$$

of  $\mathcal{C}^{\text{op}} \times \mathcal{C}$ , the diagram

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}}(X, Y) & \xrightarrow{\phi^* \circ \psi_* = \psi_* \circ \phi^*} & \text{Hom}_{\mathcal{C}}(A, B) \\ F_{X,Y} \downarrow & & \downarrow F_{A,B} \\ \text{Hom}_{\mathcal{D}}(F_X, F_Y) & \xrightarrow{F(\phi)^* \circ F(\psi)_* = F(\psi)_* \circ F(\phi)^*} & \text{Hom}_{\mathcal{D}}(F_A, F_B), \end{array}$$

<sup>19</sup>This is the 1-categorical version of **Constructions With Sets**, ?? of ??.

acting on elements as

$$\begin{array}{ccc} f & \xrightarrow{\quad} & \psi \circ f \circ \phi \\ \downarrow & & \downarrow \\ F(f) & \xrightarrow{\quad} & F(\psi) \circ F(f) \circ F(\psi) = F(\psi \circ f \circ \phi) \end{array}$$

commutes, which follows from the functoriality of  $F$ .  $\square$

**Proposition II.5.4.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  and  $G: \mathcal{D} \rightarrow \mathcal{E}$  be functors.

1. *Interaction With Natural Isomorphisms.* The following conditions are equivalent:
  - (a) The natural transformation  $F^\dagger: \text{Hom}_{\mathcal{C}} \Rightarrow \text{Hom}_{\mathcal{D}} \circ (F^{\text{op}} \times F)$  associated to  $F$  is a natural isomorphism.
  - (b) The functor  $F$  is fully faithful.
2. *Interaction With Composition.* We have an equality of pasting diagrams

$$\begin{array}{ccc} \mathcal{C}^{\text{op}} \times \mathcal{C} & \xrightarrow{F^{\text{op}} \times F} & \mathcal{D}^{\text{op}} \times \mathcal{D} \xrightarrow{G^{\text{op}} \times G} \mathcal{E}^{\text{op}} \times \mathcal{E} \\ \downarrow \text{Hom}_{\mathcal{C}} & \nearrow F^\dagger & \downarrow \text{Hom}_{\mathcal{D}} \xRightarrow{G^\dagger} \downarrow \text{Hom}_{\mathcal{E}} \\ & \text{Sets} & \end{array} = \begin{array}{ccc} \mathcal{C}^{\text{op}} \times \mathcal{C} & \xrightarrow{(G \circ F)^{\text{op}} \times (G \circ F)} & \mathcal{E}^{\text{op}} \times \mathcal{E} \\ \downarrow \text{Hom}_{\mathcal{C}} & \nearrow (G \circ F)^\dagger & \downarrow \text{Hom}_{\mathcal{E}} \\ & \text{Sets} & \end{array}$$

in  $\text{Cats}_2$ , i.e. we have

$$(G \circ F)^\dagger = (G^\dagger \star \text{id}_{F^{\text{op}} \times F}) \circ F^\dagger.$$

3. *Interaction With Identities.* We have

$$\text{id}_{\mathcal{C}}^\dagger = \text{id}_{\text{Hom}_{\mathcal{C}}(-1, -2)},$$

i.e. the natural transformation associated to  $\text{id}_{\mathcal{C}}$  is the identity natural transformation of the functor  $\text{Hom}_{\mathcal{C}}(-1, -2)$ .

*Proof.* **Item 1, Interaction With Natural Isomorphisms:** Clear.

**Item 2, Interaction With Composition:** Clear.

**Item 3, Interaction With Identities:** Clear.  $\square$

## II.6 Conditions on Functors

### II.6.1 Faithful Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.6.1.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **faithful** if, for each  $A, B \in \text{Obj}(\mathcal{C})$ , the action on morphisms

$$F_{A,B}: \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(F_A, F_B)$$

of  $F$  at  $(A, B)$  is injective.

**Proposition II.6.1.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  and  $G: \mathcal{D} \rightarrow \mathcal{E}$  be functors.

1. *Interaction With Composition.* If  $F$  and  $G$  are faithful, then so is  $G \circ F$ .
2. *Interaction With Postcomposition.* The following conditions are equivalent:

- (a) The functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is faithful.
- (b) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_*: \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is faithful.

- (c) The functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is a representably faithful morphism in  $\text{Cats}_2$  in the sense of **Types of Morphisms in Bicategories, Definition 14.1.1.1.1.**

3. *Interaction With Precomposition I.* Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor.

- (a) If  $F$  is faithful, then the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

**can fail** to be faithful.

- (b) Conversely, if the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is faithful, then  $F$  **can fail** to be faithful.



4. *Interaction With Precomposition II.* If  $F$  is essentially surjective, then the precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is faithful.

5. *Interaction With Precomposition III.* The following conditions are equivalent:

- (a) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is faithful.

- (b) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is conservative.

- (c) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is monadic.

- (d) The functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  is a corepresentably faithful morphism in  $\text{Cats}_2$  in the sense of [Types of Morphisms in Bicategories, Definition 14.2.1.1.1.](#)

- (e) The components

$$\eta_G : G \Longrightarrow \text{Ran}_F(G \circ F)$$

of the unit

$$\eta : \text{id}_{\text{Fun}(\mathcal{D}, \mathcal{X})} \Longrightarrow \text{Ran}_F \circ F^*$$

of the adjunction  $F^* \dashv \text{Ran}_F$  are all monomorphisms.

(f) The components

$$\varepsilon_G: \text{Lan}_F(G \circ F) \Longrightarrow G$$

of the counit

$$\varepsilon: \text{Lan}_F \circ F^* \Longrightarrow \text{id}_{\text{Fun}(\mathcal{D}, \mathcal{X})}$$

of the adjunction  $\text{Lan}_F \dashv F^*$  are all epimorphisms.

(g) The functor  $F$  is dominant (**Definition 11.7.1.1.1**), i.e. every object of  $\mathcal{D}$  is a retract of some object in  $\text{Im}(F)$ :

(★) For each  $B \in \text{Obj}(\mathcal{D})$ , there exist:

- An object  $A$  of  $\mathcal{C}$ ;
- A morphism  $s: B \rightarrow F(A)$  of  $\mathcal{D}$ ;
- A morphism  $r: F(A) \rightarrow B$  of  $\mathcal{D}$ ;

such that  $r \circ s = \text{id}_B$ .

*Proof.* **Item 1, Interaction With Composition:** Since the map

$$(G \circ F)_{A,B}: \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(G_{F_A}, G_{F_B}),$$

defined as the composition

$$\text{Hom}_{\mathcal{C}}(A, B) \xrightarrow{F_{A,B}} \text{Hom}_{\mathcal{D}}(F_A, F_B) \xrightarrow{G_{F(A), F(B)}} \text{Hom}_{\mathcal{D}}(G_{F_A}, G_{F_B}),$$

is a composition of injective functions, it follows from ?? that it is also injective. Therefore  $G \circ F$  is faithful.

**Item 2, Interaction With Postcomposition:** Omitted.

**Item 3, Interaction With Precomposition I:** See [MSE 733163] for **Item 3a**. **Item 3b** follows from **Item 4** and the fact that there are essentially surjective functors that are not faithful.

**Item 4, Interaction With Precomposition II:** Omitted, but see [https://unimath.github.io/doc/UniMath/d4de26f//UniMath.CategoryTheory.precursive\\_fully\\_faithful.html](https://unimath.github.io/doc/UniMath/d4de26f//UniMath.CategoryTheory.precursive_fully_faithful.html) for a formalised proof.

**Item 5, Interaction With Precomposition III:** We claim **Items 5a** to **5g** are equivalent:

- *Items  $\mathfrak{sa}$  and  $\mathfrak{sd}$  Are Equivalent:* This is true by the definition of corepresentably faithful morphism; see [Types of Morphisms in Bicategories](#), [Definition 14.2.1.1.1](#).
- *Items  $\mathfrak{sa}$  to  $\mathfrak{sc}$  and  $\mathfrak{sg}$  Are Equivalent:* See [Adá+01, Proposition 4.1] or alternatively [Fre09, Lemmas 3.1 and 3.2] for the equivalence between [Items  \$\mathfrak{sa}\$](#)  and [Items  \$\mathfrak{sg}\$](#) .
- *Items  $\mathfrak{sa}$ ,  $\mathfrak{se}$  and  $\mathfrak{sf}$  Are Equivalent:* See ??, ?? of ??.

This finishes the proof.  $\square$

## 11.6.2 Full Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.6.2.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **full** if, for each  $A, B \in \text{Obj}(\mathcal{C})$ , the action on morphisms

$$F_{A,B}: \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(F_A, F_B)$$

of  $F$  at  $(A, B)$  is surjective.

**Proposition 11.6.2.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  and  $G: \mathcal{D} \rightarrow \mathcal{E}$  be functors.

1. *Interaction With Composition.* If  $F$  and  $G$  are full, then so is  $G \circ F$ .
2. *Interaction With Postcomposition I.* If  $F$  is full, then the postcomposition functor

$$F_*: \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

**can fail** to be full.

3. *Interaction With Postcomposition II.* If, for each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_*: \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is full, then  $F$  is also full.

4. *Interaction With Precomposition I.* If  $F$  is full, then the precomposition

functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

*can fail* to be full.

5. *Interaction With Precomposition II.* If, for each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is full, then  $F$  *can fail* to be full.

6. *Interaction With Precomposition III.* If  $F$  is essentially surjective and full, then the precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is full (and also faithful by [Item 4](#) of [Definition 11.6.1.1.2](#)).

7. *Interaction With Precomposition IV.* The following conditions are equivalent:

- (a) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is full.

- (b) The functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  is a corepresentably full morphism in  $\text{Cats}_2$  in the sense of [Types of Morphisms in Bicategories, Definition 14.2.1.1.1](#).

- (c) The components

$$\eta_G : G \Longrightarrow \text{Ran}_F(G \circ F)$$

of the unit

$$\eta : \text{id}_{\text{Fun}(\mathcal{D}, \mathcal{X})} \Longrightarrow \text{Ran}_F \circ F^*$$

of the adjunction  $F^* \dashv \text{Ran}_F$  are all retractions/split epimorphisms.

(d) The components

$$\varepsilon_G: \text{Lan}_F(G \circ F) \Longrightarrow G$$

of the counit

$$\varepsilon: \text{Lan}_F \circ F^* \Longrightarrow \text{id}_{\text{Fun}(\mathcal{D}, \mathcal{X})}$$

of the adjunction  $\text{Lan}_F \dashv F^*$  are all sections/split monomorphisms.

(e) For each  $B \in \text{Obj}(\mathcal{D})$ , there exist:

- An object  $A_B$  of  $\mathcal{C}$ ;
- A morphism  $s_B: B \rightarrow F(A_B)$  of  $\mathcal{D}$ ;
- A morphism  $r_B: F(A_B) \rightarrow B$  of  $\mathcal{D}$ ;

satisfying the following condition:

(★) For each  $A \in \text{Obj}(\mathcal{C})$  and each pair of morphisms

$$r: F(A) \rightarrow B,$$

$$s: B \rightarrow F(A)$$

of  $\mathcal{D}$ , we have

$$[(A_B, s_B, r_B)] = [(A, s, r \circ s_B \circ r_B)]$$

$$\text{in } \int^{A \in \mathcal{C}} b_{F_A}^{B'} \times b_B^{F_A}.$$

*Proof.* **Item 1, Interaction With Composition:** Since the map

$$(G \circ F)_{A,B}: \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(G_{F_A}, G_{F_B}),$$

defined as the composition

$$\text{Hom}_{\mathcal{C}}(A, B) \xrightarrow{F_{A,B}} \text{Hom}_{\mathcal{D}}(F_A, F_B) \xrightarrow{G_{F(A), F(B)}} \text{Hom}_{\mathcal{D}}(G_{F_A}, G_{F_B}),$$

is a composition of surjective functions, it follows from ?? that it is also surjective.

Therefore  $G \circ F$  is full.

**Item 2, Interaction With Postcomposition I:** We follow the proof (completely formalised in cubical Agda!) given by Naïm Camille Favier in [**favier:postcompose-not-full**].

Let  $\mathcal{C}$  be the category where:

- *Objects.* We have  $\text{Obj}(C) = \{A, B\}$ .
- *Morphisms.* We have

$$\begin{aligned}\text{Hom}_C(A, A) &= \{e_A, \text{id}_A\}, \\ \text{Hom}_C(B, B) &= \{e_B, \text{id}_B\}, \\ \text{Hom}_C(A, B) &= \{f, g\}, \\ \text{Hom}_C(B, A) &= \emptyset.\end{aligned}$$

- *Composition.* The nontrivial compositions in  $C$  are the following:

$$\begin{aligned}e_A \circ e_A &= \text{id}_A, & f \circ e_A &= g, & e_B \circ f &= f, \\ e_B \circ e_B &= \text{id}_B, & g \circ e_A &= f, & e_B \circ g &= g.\end{aligned}$$

We may picture  $C$  as follows:

$$e_A \circlearrowleft A \begin{matrix} \xrightarrow{f} \\ \xleftarrow{g} \end{matrix} B \circlearrowright e_B.$$

Next, let  $\mathcal{D}$  be the walking arrow category **1** of [Definition 11.2.5.1.1](#) and let  $F: C \rightarrow \mathbf{1}$  be the functor given on objects by

$$\begin{aligned}F(A) &= 0, \\ F(B) &= 1\end{aligned}$$

and on non-identity morphisms by

$$\begin{aligned}F(f) &= f_{01}, & F(e_A) &= \text{id}_0, \\ F(g) &= f_{01}, & F(e_B) &= \text{id}_1.\end{aligned}$$

Finally, let  $\mathcal{X} = \mathbf{B}\mathbb{Z}/2$  be the walking involution and let  $\iota_A, \iota_B: \mathbf{B}\mathbb{Z}/2 \rightrightarrows C$  be the inclusion functors from  $\mathbf{B}\mathbb{Z}/2$  to  $C$  with

$$\begin{aligned}\iota_A(\bullet) &= A, \\ \iota_B(\bullet) &= B.\end{aligned}$$

Since every morphism in  $\mathbb{1}$  has a preimage in  $\mathcal{C}$  by  $F$ , the functor  $F$  is full. Now, for  $F_*$  to be full, the map

$$\begin{aligned} F_{*|\iota_A, \iota_B}: \text{Nat}(\iota_A, \iota_B) &\longrightarrow \text{Nat}(F \circ \iota_A, F \circ \iota_B) \\ \alpha &\longmapsto \text{id}_F \star \alpha \end{aligned}$$

would need to be surjective. However, as we will show next, we have

$$\begin{aligned} \text{Nat}(\iota_A, \iota_B) &= \emptyset, \\ \text{Nat}(F \circ \iota_A, F \circ \iota_B) &\cong \text{pt}, \end{aligned}$$

so this is impossible:

- *Proof of  $\text{Nat}(\iota_A, \iota_B) = \emptyset$ :* A natural transformation  $\alpha: \iota_A \Rightarrow \iota_B$  consists of a morphism

$$\alpha: \underbrace{\iota_A(\bullet)}_{=A} \rightarrow \underbrace{\iota_B(\bullet)}_{=B}$$

in  $\mathcal{C}$  making the diagram

$$\begin{array}{ccc} \iota_A(\bullet) & \xrightarrow{\iota_A(e)} & \iota_A(\bullet) \\ \alpha \downarrow & & \downarrow \alpha \\ \iota_B(\bullet) & \xrightarrow{\iota_B(e)} & \iota_B(\bullet) \end{array}$$

commute for each  $e \in \text{Hom}_{\mathbb{B}\mathbb{Z}/2}(\bullet, \bullet) \cong \mathbb{Z}/2$ . We have two cases:

- If  $\alpha = f$ , the naturality diagram for the unique nonidentity element of  $\mathbb{Z}/2$  is given by

$$\begin{array}{ccc} A & \xrightarrow{e_A} & A \\ f \downarrow & & \downarrow f \\ B & \xrightarrow{e_B} & B. \end{array}$$

However,  $e_B \circ f = f$  and  $f \circ e_A = g$ , so this diagram does not commute.

2. If  $\alpha = g$ , the naturality diagram for the unique nonidentity element of  $\mathbb{Z}/2$  is given by

$$\begin{array}{ccc} A & \xrightarrow{e_A} & A \\ \downarrow g & & \downarrow g \\ B & \xrightarrow{e_B} & B. \end{array}$$

However,  $e_B \circ g = g$  and  $g \circ e_A = f$ , so this diagram does not commute.

As a result, there are no natural transformations from  $\iota_A$  to  $\iota_B$ .

- *Proof of*  $\text{Nat}(F \circ \iota_A, F \circ \iota_B) \cong \text{pt}$ : A natural transformation

$$\beta: F \circ \iota_A \Rightarrow F \circ \iota_B$$

consists of a morphism

$$\beta: \underbrace{[F \circ \iota_A](\bullet)}_{=0} \rightarrow \underbrace{[F \circ \iota_B](\bullet)}_{=1}$$

in  $\mathbb{1}$  making the diagram

$$\begin{array}{ccc} [F \circ \iota_A](\bullet) & \xrightarrow{[F \circ \iota_A](e)} & [F \circ \iota_A](\bullet) \\ \downarrow \beta & & \downarrow \beta \\ [F \circ \iota_B](\bullet) & \xrightarrow{[F \circ \iota_B](e)} & [F \circ \iota_B](\bullet) \end{array}$$

commute for each  $e \in \text{Hom}_{\mathbb{B}\mathbb{Z}/2}(\bullet, \bullet) \cong \mathbb{Z}/2$ . Since the only morphism from 0 to 1 in  $\mathbb{1}$  is  $f_{01}$ , we must have  $\beta = f_{01}$  if such a transformation were to exist, and in fact it indeed does, as in this case the naturality diagram above becomes

$$\begin{array}{ccc} 0 & \xrightarrow{\text{id}_0} & 0 \\ \downarrow f_{01} & & \downarrow f_{01} \\ 1 & \xrightarrow{\text{id}_1} & 1 \end{array}$$



for each  $e \in \mathbb{Z}/2$ , and this diagram indeed commutes, making  $\beta$  into a natural transformation.

This finishes the proof.

*Item 3, Interaction With Postcomposition II:* Taking  $\mathcal{X} = \text{pt}$ , it follows by assumption that the functor

$$F_* : \text{Fun}(\text{pt}, C) \rightarrow \text{Fun}(\text{pt}, \mathcal{D})$$

is full. However, by *Item 5* of *Definition 11.10.1.1.2*, we have isomorphisms of categories

$$\begin{aligned} \text{Fun}(\text{pt}, C) &\cong C, \\ \text{Fun}(\text{pt}, \mathcal{D}) &\cong \mathcal{D} \end{aligned}$$

and the diagram

$$\begin{array}{ccc} \text{Fun}(\text{pt}, C) & \xrightarrow{F_*} & \text{Fun}(\text{pt}, \mathcal{D}) \\ \downarrow \wr & & \downarrow \wr \\ C & \xrightarrow{F} & \mathcal{D} \end{array}$$

commutes. It then follows from *Item 1* that  $F$  is full.

*Item 4, Interaction With Precomposition I:* Omitted.

*Item 5, Interaction With Precomposition II:* See [BS10, p. 47].

*Item 6, Interaction With Precomposition III:* Omitted, but see [https://unimath.github.io/doc/UniMath/d4de26f//UniMath.CategoryTheory.precomp\\_fully\\_faithful.html](https://unimath.github.io/doc/UniMath/d4de26f//UniMath.CategoryTheory.precomp_fully_faithful.html) for a formalised proof.

*Item 7, Interaction With Precomposition IV:* We claim *Items 7a* to *7e* are equivalent:

- *Items 7a and 7b Are Equivalent:* This is true by the definition of corepresentably full morphism; see *Types of Morphisms in Bicategories*, *Definition 14.2.2.1.1*.
- *Items 7a, 7c and 7d Are Equivalent:* See ??, ?? of ??.
- *Items 7a and 7e Are Equivalent:* See [Adá+01, Item (b) of Remark 4.3].

This finishes the proof.  $\square$

**Question 11.6.2.1.3.** Item 7 of Definition 11.6.2.1.2 gives a characterisation of the functors  $F$  for which  $F^*$  is full, but the characterisations given there are really messy. Are there better ones?

This question also appears as [MO 468121b].

### 11.6.3 Fully Faithful Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.6.3.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **fully faithful** if  $F$  is full and faithful, i.e. if, for each  $A, B \in \text{Obj}(\mathcal{C})$ , the action on morphisms

$$F_{A,B}: \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(F_A, F_B)$$

of  $F$  at  $(A, B)$  is bijective.

**Proposition 11.6.3.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  and  $G: \mathcal{D} \rightarrow \mathcal{E}$  be functors.

1. *Characterisations.* The following conditions are equivalent:

- (a) The functor  $F$  is fully faithful.
- (b) We have a pullback square

$$\begin{array}{ccc} \text{Arr}(\mathcal{C}) & \xrightarrow{\text{Arr}(F)} & \text{Arr}(\mathcal{D}) \\ \text{src} \times \text{tgt} \downarrow & \lrcorner & \downarrow \text{src} \times \text{tgt} \\ \mathcal{C} \times \mathcal{C} & \xrightarrow{F \times F} & \mathcal{D} \times \mathcal{D} \end{array}$$

in  $\mathbf{Cats}$ .

- 2. *Interaction With Composition.* If  $F$  and  $G$  are fully faithful, then so is  $G \circ F$ .
- 3. *Conservativity.* If  $F$  is fully faithful, then  $F$  is conservative.
- 4. *Essential Injectivity.* If  $F$  is fully faithful, then  $F$  is essentially injective.
- 5. *Interaction With Co/Limits.* If  $F$  is fully faithful, then  $F$  reflects co/limits.

6. *Interaction With Postcomposition.* The following conditions are equivalent:

- (a) The functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is fully faithful.
- (b) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_*: \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is fully faithful.

- (c) The functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is a representably fully faithful morphism in  $\text{Cats}_2$  in the sense of [Types of Morphisms in Bicategories, Definition 14.1.3.1.1.](#)

7. *Interaction With Precomposition I.* If  $F$  is fully faithful, then the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

**can fail** to be fully faithful.

8. *Interaction With Precomposition II.* If the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is fully faithful, then  $F$  **can fail** to be fully faithful (and in fact it can also fail to be either full or faithful).

9. *Interaction With Precomposition III.* If  $F$  is essentially surjective and full, then the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is fully faithful.

10. *Interaction With Precomposition IV.* The following conditions are equivalent:

- (a) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is fully faithful.

(b) The precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \text{Sets}) \rightarrow \text{Fun}(\mathcal{C}, \text{Sets})$$

is fully faithful.

(c) The functor

$$\text{Lan}_F : \text{Fun}(\mathcal{C}, \text{Sets}) \rightarrow \text{Fun}(\mathcal{D}, \text{Sets})$$

is fully faithful.

(d) The functor  $F$  is a corepresentably fully faithful morphism in  $\text{Cats}_2$  in the sense of [Types of Morphisms in Bicategories, Definition 14.2.3.1.1.](#)

(e) The functor  $F$  is absolutely dense.

(f) The components

$$\eta_G : G \Longrightarrow \text{Ran}_F(G \circ F)$$

of the unit

$$\eta : \text{id}_{\text{Fun}(\mathcal{D}, \mathcal{X})} \Longrightarrow \text{Ran}_F \circ F^*$$

of the adjunction  $F^* \dashv \text{Ran}_F$  are all isomorphisms.

(g) The components

$$\varepsilon_G : \text{Lan}_F(G \circ F) \Longrightarrow G$$

of the counit

$$\varepsilon : \text{Lan}_F \circ F^* \Longrightarrow \text{id}_{\text{Fun}(\mathcal{D}, \mathcal{X})}$$

of the adjunction  $\text{Lan}_F \dashv F^*$  are all isomorphisms.

(h) The natural transformation

$$\alpha : \text{Lan}_{h_F}(h^F) \Longrightarrow h$$

with components

$$\alpha_{B', B} : \int^{A \in \mathcal{C}} h_{F_A}^{B'} \times h_B^{F_A} \rightarrow h_B^{B'}$$

given by

$$\alpha_{B', B}([\langle \phi, \psi \rangle]) = \psi \circ \phi$$

is a natural isomorphism.

(i) For each  $B \in \text{Obj}(\mathcal{D})$ , there exist:

- An object  $A_B$  of  $\mathcal{C}$ ;
- A morphism  $s_B: B \rightarrow F(A_B)$  of  $\mathcal{D}$ ;
- A morphism  $r_B: F(A_B) \rightarrow B$  of  $\mathcal{D}$ ;

satisfying the following conditions:

- i. The triple  $(F(A_B), r_B, s_B)$  is a retract of  $B$ , i.e. we have  $r_B \circ s_B = \text{id}_B$ .
- ii. For each morphism  $f: B' \rightarrow B$  of  $\mathcal{D}$ , we have

$$[(A_B, s_{B'}, f \circ r_{B'})] = [(A_B, s_B \circ f, r_B)]$$

$$\text{in } \int^{A \in \mathcal{C}} h_{F_A}^{B'} \times h_B^{F_A}.$$

*Proof.* **Item 1, Characterisations:** Omitted.

**Item 2, Interaction With Composition:** Since the map

$$(G \circ F)_{A,B}: \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(G_{F_A}, G_{F_B}),$$

defined as the composition

$$\text{Hom}_{\mathcal{C}}(A, B) \xrightarrow{F_{A,B}} \text{Hom}_{\mathcal{D}}(F_A, F_B) \xrightarrow{G_{F(A), F(B)}} \text{Hom}_{\mathcal{D}}(G_{F_A}, G_{F_B}),$$

is a composition of bijective functions, it follows from ?? that it is also bijective. Therefore  $G \circ F$  is fully faithful.

**Item 3, Conservativity:** This is a repetition of **Item 2** of **Definition 11.6.4.1.2**, and is proved there.

**Item 4, Essential Injectivity:** Omitted.

**Item 5, Interaction With Co/Limits:** Omitted.

**Item 6, Interaction With Postcomposition:** This follows from **Item 2** of **Definition 11.6.1.1.2** and ?? of **Definition 11.6.2.1.2**.

**Item 7, Interaction With Precomposition I:** See [MSE 733161] for an example of a fully faithful functor whose precomposition with which fails to be full.

**Item 8, Interaction With Precomposition II:** See [MSE 749304, Item 3].

**Item 9, Interaction With Precomposition III:** Omitted, but see [https://unimath.github.io/doc/UniMath/d4de26f//UniMath.CategoryTheory.precomp\\_fully\\_faithful.html](https://unimath.github.io/doc/UniMath/d4de26f//UniMath.CategoryTheory.precomp_fully_faithful.html) for a formalised proof.

**Item 10, Interaction With Precomposition IV:** We claim **Items 10a** to **10i** are equivalent:

- *Items 10a and 10d Are Equivalent:* This is true by the definition of corepresentably fully faithful morphism; see [Types of Morphisms in Bicategories](#), [Definition 14.2.3.1.1](#).
- *Items 10a, 10f and 10g Are Equivalent:* See ??, ?? of ??.
- *Items 10a to 10c Are Equivalent:* This follows from [Low15, Proposition A.1.5].
- *Items 10a, 10e, 10h and 10i Are Equivalent:* See [Fre09, Theorem 4.1] and [Adá+oi, Theorem 1.1].

This finishes the proof.  $\square$

## 11.6.4 Conservative Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.6.4.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **conservative** if it satisfies the following condition:<sup>20</sup>

- (★) For each  $f \in \text{Mor}(\mathcal{C})$ , if  $F(f)$  is an isomorphism in  $\mathcal{D}$ , then  $f$  is an isomorphism in  $\mathcal{C}$ .

**Proposition 11.6.4.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor.

1. *Characterisations.* The following conditions are equivalent:
  - (a) The functor  $F$  is conservative.
  - (b) For each  $f \in \text{Mor}(\mathcal{C})$ , the morphism  $F(f)$  is an isomorphism in  $\mathcal{D}$  iff  $f$  is an isomorphism in  $\mathcal{C}$ .
2. *Interaction With Fully Faithfulness.* Every fully faithful functor is conservative.
3. *Interaction With Precomposition.* The following conditions are equivalent:
  - (a) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is conservative.

---

<sup>20</sup>*Slogan:* A functor  $F$  is **conservative** if it reflects isomorphisms.

- (b) The equivalent conditions of **Item 5** of **Definition 11.6.1.1.2** are satisfied.

*Proof.* **Item 1, Characterisations:** This follows from **Item 1** of **Definition 11.5.1.1.6**. **Item 2, Interaction With Fully Faithfulness:** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a fully faithful functor, let  $f: A \rightarrow B$  be a morphism of  $\mathcal{C}$ , and suppose that  $F_f$  is an isomorphism. We have

$$\begin{aligned} F(\mathrm{id}_B) &= \mathrm{id}_{F(B)} \\ &= F(f) \circ F(f)^{-1} \\ &= F(f \circ f^{-1}). \end{aligned}$$

Similarly,  $F(\mathrm{id}_A) = F(f^{-1} \circ f)$ . But since  $F$  is fully faithful, we must have

$$\begin{aligned} f \circ f^{-1} &= \mathrm{id}_B, \\ f^{-1} \circ f &= \mathrm{id}_A, \end{aligned}$$

showing  $f$  to be an isomorphism. Thus  $F$  is conservative.  $\square$

**Question 11.6.4.1.3.** Is there a characterisation of functors  $F: \mathcal{C} \rightarrow \mathcal{D}$  satisfying the following condition:

- ( $\star$ ) For each  $\mathcal{X} \in \mathrm{Obj}(\mathrm{Cats})$ , the postcomposition functor

$$F_*: \mathrm{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \mathrm{Fun}(\mathcal{X}, \mathcal{D})$$

is conservative?

This question also appears as [MO 468121a].

## 11.6.5 Essentially Injective Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.6.5.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **essentially injective** if it satisfies the following condition:

- ( $\star$ ) For each  $A, B \in \mathrm{Obj}(\mathcal{C})$ , if  $F(A) \cong F(B)$ , then  $A \cong B$ .

**Question 11.6.5.1.2.** Is there a characterisation of functors  $F: \mathcal{C} \rightarrow \mathcal{D}$  such that:

1. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is essentially injective, i.e. if  $\phi \circ F \cong \psi \circ F$ , then  $\phi \cong \psi$  for all functors  $\phi$  and  $\psi$ ?

2. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_* : \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is essentially injective, i.e. if  $F \circ \phi \cong F \circ \psi$ , then  $\phi \cong \psi$ ?

This question also appears as [MO 468121a].

## II.6.6 Essentially Surjective Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.6.6.I.1.** A functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  is **essentially surjective**<sup>21</sup> if it satisfies the following condition:

- (★) For each  $D \in \text{Obj}(\mathcal{D})$ , there exists some object  $A$  of  $\mathcal{C}$  such that  $F(A) \cong D$ .

**Question II.6.6.I.2.** Is there a characterisation of functors  $F : \mathcal{C} \rightarrow \mathcal{D}$  such that:

1. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is essentially surjective?

2. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_* : \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is essentially surjective?

This question also appears as [MO 468121a].

<sup>21</sup>*Further Terminology:* Also called an **eso** functor, meaning *essentially surjective on objects*.



## II.6.7 Equivalences of Categories

**Definition II.6.7.1.1.** Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

1. An **equivalence of categories** between  $\mathcal{C}$  and  $\mathcal{D}$  consists of a pair of functors

$$\begin{aligned} F: \mathcal{C} &\rightarrow \mathcal{D}, \\ G: \mathcal{D} &\rightarrow \mathcal{C} \end{aligned}$$

together with natural isomorphisms

$$\begin{aligned} \eta: \text{id}_{\mathcal{C}} &\xrightarrow{\sim} G \circ F, \\ \varepsilon: F \circ G &\xrightarrow{\sim} \text{id}_{\mathcal{D}}. \end{aligned}$$

2. An **adjoint equivalence of categories** between  $\mathcal{C}$  and  $\mathcal{D}$  is an equivalence  $(F, G, \eta, \varepsilon)$  between  $\mathcal{C}$  and  $\mathcal{D}$  which is also an adjunction.

**Proposition II.6.7.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor.

1. *Characterisations.* If  $\mathcal{C}$  and  $\mathcal{D}$  are small<sup>22</sup>, then the following conditions are equivalent:<sup>23</sup>
  - (a) The functor  $F$  is an equivalence of categories.
  - (b) The functor  $F$  is fully faithful and essentially surjective.
  - (c) The induced functor

$$\uparrow F\text{Sk}(\mathcal{C}): \text{Sk}(\mathcal{C}) \rightarrow \text{Sk}(\mathcal{D})$$

is an *isomorphism* of categories.

- (d) For each  $X \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, X) \rightarrow \text{Fun}(\mathcal{C}, X)$$

is an equivalence of categories.

<sup>22</sup>Otherwise there will be size issues. One can also work with large categories and universes, or require  $F$  to be *constructively* essentially surjective; see [MSE 1465107].

<sup>23</sup>In ZFC, the equivalence between **Item 1a** and **Item 1b** is equivalent to the axiom of choice;

(e) For each  $X \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_* : \text{Fun}(X, C) \rightarrow \text{Fun}(X, D)$$

is an equivalence of categories.

2. *Two-Out-of-Three*. Let

$$\begin{array}{ccc} C & \xrightarrow{G \circ F} & E \\ & \searrow F \quad \nearrow G & \\ & D & \end{array}$$

be a diagram in  $\text{Cats}$ . If two out of the three functors among  $F$ ,  $G$ , and  $G \circ F$  are equivalences of categories, then so is the third.

3. *Stability Under Composition*. Let

$$C \begin{array}{c} \xrightarrow{F} \\ \xleftarrow{G} \end{array} D \begin{array}{c} \xrightarrow{F'} \\ \xleftarrow{G'} \end{array} E$$

be a diagram in  $\text{Cats}$ . If  $(F, G)$  and  $(F', G')$  are equivalences of categories, then so is their composite  $(F' \circ F, G' \circ G)$ .

4. *Equivalences vs. Adjoint Equivalences*. Every equivalence of categories can be promoted to an adjoint equivalence.<sup>24</sup>

5. *Interaction With Groupoids*. If  $C$  and  $D$  are groupoids, then the following conditions are equivalent:

(a) The functor  $F$  is an equivalence of groupoids.

(b) The following conditions are satisfied:

i. The functor  $F$  induces a bijection

$$\pi_0(F) : \pi_0(C) \rightarrow \pi_0(D)$$

of sets.

---

see [MO 119454].

In Univalent Foundations, this is true without requiring neither the axiom of choice nor the law of excluded middle.

<sup>24</sup>More precisely, we can promote an equivalence of categories  $(F, G, \eta, \varepsilon)$  to adjoint equiva-

ii. For each  $A \in \text{Obj}(C)$ , the induced map

$$F_{x,x}: \text{Aut}_C(A) \rightarrow \text{Aut}_{\mathcal{D}}(F_A)$$

is an isomorphism of groups.

*Proof.* **Item 1, Characterisations:** We claim that **Items 1a** to **1e** are indeed equivalent:

1. **Item 1a**  $\implies$  **Item 1b**: Clear.
2. **Item 1b**  $\implies$  **Item 1a**: Since  $F$  is essentially surjective and  $C$  and  $\mathcal{D}$  are small, we can choose, using the axiom of choice, for each  $B \in \text{Obj}(\mathcal{D})$ , an object  $j_B$  of  $C$  and an isomorphism  $i_B: B \rightarrow F_{j_B}$  of  $\mathcal{D}$ .  
Since  $F$  is fully faithful, we can extend the assignment  $B \mapsto j_B$  to a *unique* functor  $j: \mathcal{D} \rightarrow C$  such that the isomorphisms  $i_B: B \rightarrow F_{j_B}$  assemble into a natural isomorphism  $\eta: \text{id}_{\mathcal{D}} \xrightarrow{\sim} F \circ j$ , with a similar natural isomorphism  $\varepsilon: \text{id}_C \xrightarrow{\sim} j \circ F$ . Hence  $F$  is an equivalence.
3. **Item 1a**  $\implies$  **Item 1c**: This follows from **Item 4** of **Definition II.1.3.1.3**.
4. **Item 1c**  $\implies$  **Item 1a**: Omitted.
5. **Items 1a, 1d and 1e Are Equivalent**: This follows from ??.

This finishes the proof of **Item 1**.

**Item 2, Two-Out-of-Three**: Omitted.

**Item 3, Stability Under Composition**: Clear.

**Item 4, Equivalences vs. Adjoint Equivalences**: See [Rier6, Proposition 4.4.5].

**Item 5, Interaction With Groupoids**: See [nLa25, Proposition 4.4].  $\square$

## II.6.8 Isomorphisms of Categories

**Definition II.6.8.1.1.** An **isomorphism of categories** is a pair of functors

$$\begin{aligned} F: C &\rightarrow \mathcal{D}, \\ G: \mathcal{D} &\rightarrow C \end{aligned}$$

---

lences  $(F, G, \eta', \varepsilon)$  and  $(F, G, \eta, \varepsilon')$ .

such that we have

$$G \circ F = \text{id}_C,$$

$$F \circ G = \text{id}_D.$$

**Example 11.6.8.1.2.** Categories can be equivalent but non-isomorphic. For example, the category consisting of two isomorphic objects is equivalent to  $\text{pt}$ , but not isomorphic to it.

**Proposition 11.6.8.1.3.** Let  $F: C \rightarrow D$  be a functor.

1. *Characterisations.* If  $C$  and  $D$  are small, then the following conditions are equivalent:

- (a) The functor  $F$  is an isomorphism of categories.
- (b) The functor  $F$  is fully faithful and bijective on objects.
- (c) For each  $X \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^*: \text{Fun}(D, X) \rightarrow \text{Fun}(C, X)$$

is an isomorphism of categories.

- (d) For each  $X \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_*: \text{Fun}(X, C) \rightarrow \text{Fun}(X, D)$$

is an isomorphism of categories.

*Proof.* **Item 1, Characterisations:** We claim that **Items 1a** to **1d** are indeed equivalent:

1. **Items 1a and 1b Are Equivalent:** Omitted, but similar to **Item 1** of **Definition 11.6.7.1.2**.
2. **Items 1a, 1c and 1d Are Equivalent:** This follows from ??.

This finishes the proof. □

## II.7 More Conditions on Functors

### II.7.1 Dominant Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.7.1.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **dominant** if every object of  $\mathcal{D}$  is a retract of some object in  $\text{Im}(F)$ , i.e.:

(★) For each  $B \in \text{Obj}(\mathcal{D})$ , there exist:

- An object  $A$  of  $\mathcal{C}$ ;
- A morphism  $r: F(A) \rightarrow B$  of  $\mathcal{D}$ ;
- A morphism  $s: B \rightarrow F(A)$  of  $\mathcal{D}$ ;

such that we have

$$r \circ s = \text{id}_B,$$

$$\begin{array}{ccc} B & \xrightarrow{s} & F(A) \\ & \searrow \text{id}_B & \downarrow r \\ & & B. \end{array}$$

**Proposition II.7.1.1.2.** Let  $F, G: \mathcal{C} \rightarrow \mathcal{D}$  be functors and let  $I: \mathcal{X} \rightarrow \mathcal{C}$  be a functor.

1. *Interaction With Right Whiskering.* If  $I$  is full and dominant, then the map

$$- \star \text{id}_I: \text{Nat}(F, G) \rightarrow \text{Nat}(F \circ I, G \circ I)$$

is a bijection.

2. *Interaction With Adjunctions.* Let  $(F, G): \mathcal{C} \rightleftarrows \mathcal{D}$  be an adjunction.

- (a) If  $F$  is dominant, then  $G$  is faithful.
- (b) The following conditions are equivalent:
  - i. The functor  $G$  is full.
  - ii. The restriction

$$\upharpoonright G\text{Im}_F: \text{Im}(F) \rightarrow \mathcal{C}$$

of  $G$  to  $\text{Im}(F)$  is full.

*Proof.* **Item 1, Interaction With Right Whiskering:** See [DFH75, Proposition 1.4].

**Item 2, Interaction With Adjunctions:** See [DFH75, Proposition 1.7].  $\square$

**Question II.7.1.1.3.** Is there a characterisation of functors  $F: \mathcal{C} \rightarrow \mathcal{D}$  such that:

1. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is dominant?

2. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_*: \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is dominant?

This question also appears as [MO 468121a].

## II.7.2 Monomorphisms of Categories

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.7.2.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is a **monomorphism of categories** if it is a monomorphism in  $\text{Cats}$  (see ??, ??).

**Proposition II.7.2.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor.

1. *Characterisations.* The following conditions are equivalent:
  - (a) The functor  $F$  is a monomorphism of categories.
  - (b) The functor  $F$  is injective on objects and morphisms, i.e.  $F$  is injective on objects and the map

$$F: \text{Mor}(\mathcal{C}) \rightarrow \text{Mor}(\mathcal{D})$$

is injective.

*Proof.* **Item 1, Characterisations:** Omitted.  $\square$

**Question II.7.2.1.3.** Is there a characterisation of functors  $F: \mathcal{C} \rightarrow \mathcal{D}$  such that:

1. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is a monomorphism of categories?

2. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_*: \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is a monomorphism of categories?

This question also appears as [MO 468121a].

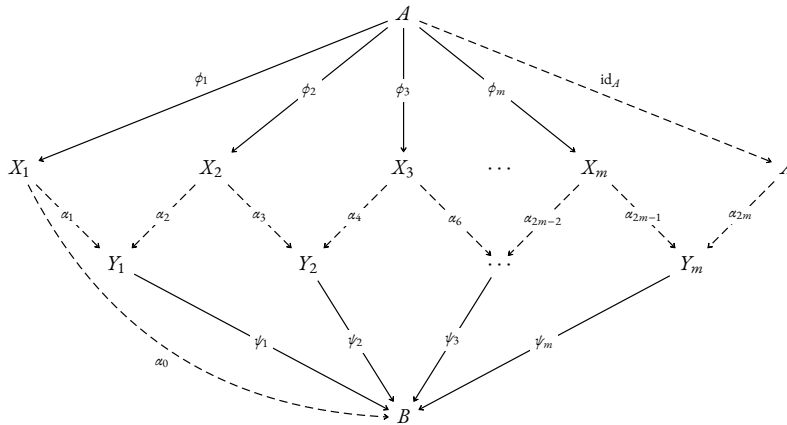
### II.7.3 Epimorphisms of Categories

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.7.3.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is a **epimorphism of categories** if it is an epimorphism in  $\text{Cats}$  (see ??, ??).

**Proposition II.7.3.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor.

1. *Characterisations.* The following conditions are equivalent:<sup>25</sup>
  - (a) The functor  $F$  is an epimorphism of categories.
  - (b) For each morphism  $f: A \rightarrow B$  of  $\mathcal{D}$ , we have a diagram



<sup>25</sup>*Further Terminology:* This statement is known as **Isbell's zigzag theorem**.

in  $\mathcal{D}$  satisfying the following conditions:

- i. We have  $f = \alpha_0 \circ \phi_1$ .
  - ii. We have  $f = \psi_m \circ \alpha_{2m}$ .
  - iii. For each  $0 \leq i \leq 2m$ , we have  $\alpha_i \in \text{Mor}(\text{Im}(F))$ .
2. *Surjectivity on Objects.* If  $F$  is an epimorphism of categories, then  $F$  is surjective on objects.

*Proof.* **Item 1, Characterisations:** See [Isb68].

**Item 2, Surjectivity on Objects:** Omitted.  $\square$

**Question II.7.3.I.3.** Is there a characterisation of functors  $F: \mathcal{C} \rightarrow \mathcal{D}$  such that:

1. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is an epimorphism of categories?

2. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_*: \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is an epimorphism of categories?

This question also appears as [MO 468121a].

## II.7.4 Pseudomononic Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.7.4.I.I.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **pseudomononic** if it satisfies the following conditions:

1. For all diagrams of the form

$$\mathcal{X} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} \mathcal{C} \xrightarrow{F} \mathcal{D},$$

if we have

$$\text{id}_F \star \alpha = \text{id}_F \star \beta,$$

then  $\alpha = \beta$ .



2. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$  and each natural isomorphism

$$\beta: F \circ \phi \xrightarrow{\sim} F \circ \psi, \quad \mathcal{X} \begin{array}{c} \xrightarrow{F \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{F \circ \psi} \end{array} \mathcal{D},$$

there exists a natural isomorphism

$$\alpha: \phi \xrightarrow{\sim} \psi, \quad \mathcal{X} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} \mathcal{C}$$

such that we have an equality

$$\mathcal{X} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} \mathcal{C} \xrightarrow{F} \mathcal{D} = \mathcal{X} \begin{array}{c} \xrightarrow{F \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{F \circ \psi} \end{array} \mathcal{D}$$

of pasting diagrams, i.e. such that we have

$$\beta = \text{id}_F \star \alpha.$$

**Proposition II.7.4.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor.

I. *Characterisations.* The following conditions are equivalent:

(a) The functor  $F$  is pseudomonoid.

(b) The functor  $F$  satisfies the following conditions:

i. The functor  $F$  is faithful, i.e. for each  $A, B \in \text{Obj}(\mathcal{C})$ , the action on morphisms

$$F_{A,B}: \text{Hom}_{\mathcal{C}}(A, B) \rightarrow \text{Hom}_{\mathcal{D}}(F_A, F_B)$$

of  $F$  at  $(A, B)$  is injective.

ii. For each  $A, B \in \text{Obj}(\mathcal{C})$ , the restriction

$$F_{A,B}^{\text{iso}}: \text{Iso}_{\mathcal{C}}(A, B) \rightarrow \text{Iso}_{\mathcal{D}}(F_A, F_B)$$

of the action on morphisms of  $F$  at  $(A, B)$  to isomorphisms is surjective.

(c) We have an isocomma square of the form

$$C \stackrel{\text{eq.}}{\cong} C \times_{\mathcal{D}} C, \quad \begin{array}{ccc} C & \xrightarrow{\text{id}_C} & C \\ \text{id}_C \downarrow & \nearrow \text{dashed} & \downarrow F \\ C & \xrightarrow{F} & \mathcal{D} \end{array}$$

in  $\text{Cats}_2$  up to equivalence.

(d) We have an isocomma square of the form

$$C \stackrel{\text{eq.}}{\cong} C \times_{\text{Arr}(\mathcal{D})} \mathcal{D}, \quad \begin{array}{ccc} C & \hookrightarrow & \text{Arr}(C) \\ F \downarrow & \nearrow \text{dashed} & \downarrow \text{Arr}(F) \\ \mathcal{D} & \hookrightarrow & \text{Arr}(\mathcal{D}) \end{array}$$

in  $\text{Cats}_2$  up to equivalence.

(e) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition<sup>26</sup> functor

$$F_* : \text{Fun}(\mathcal{X}, C) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is pseudomononic.

2. *Conservativity.* If  $F$  is pseudomononic, then  $F$  is conservative.
3. *Essential Injectivity.* If  $F$  is pseudomononic, then  $F$  is essentially injective.

*Proof.* **Item 1, Characterisations:** Omitted.

**Item 2, Conservativity:** Omitted.

**Item 3, Essential Injectivity:** Omitted. □

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<sup>26</sup>Asking the precomposition functors

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(C, \mathcal{X})$$

to be pseudomononic leads to pseudoepic functors; see **Item 1b** of **Item 1** of **Definition 11.7.5.1.2**.

## II.7.5 Pseudoepic Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.7.5.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **pseudoepic** if it satisfies the following conditions:

1. For all diagrams of the form

$$\mathcal{C} \xrightarrow{F} \mathcal{D} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \beta \\ \xrightarrow{\psi} \end{array} \mathcal{X},$$

if we have

$$\alpha \star \text{id}_F = \beta \star \text{id}_F,$$

then  $\alpha = \beta$ .

2. For each  $X \in \text{Obj}(\mathcal{C})$  and each 2-isomorphism

$$\beta: \phi \circ F \xRightarrow{\sim} \psi \circ F, \quad \mathcal{C} \begin{array}{c} \xrightarrow{\phi \circ F} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ F} \end{array} \mathcal{X}$$

of  $\mathcal{C}$ , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad \mathcal{D} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} \mathcal{X}$$

of  $\mathcal{C}$  such that we have an equality

$$\mathcal{C} \xrightarrow{F} \mathcal{D} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} \mathcal{X} = \mathcal{C} \begin{array}{c} \xrightarrow{\phi \circ F} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ F} \end{array} \mathcal{X}$$

of pasting diagrams in  $\mathcal{C}$ , i.e. such that we have

$$\beta = \alpha \star \text{id}_F.$$

**Proposition II.7.5.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor.

1. *Characterisations.* The following conditions are equivalent:

- (a) The functor  $F$  is pseudoepic.
- (b) For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the functor

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(C, \mathcal{X})$$

given by precomposition by  $F$  is pseudomonic.

- (c) We have an isococoma square of the form

$$\mathcal{D} \xrightarrow{\text{eq.}} \mathcal{D} \xleftrightarrow{\quad} \coprod_C \mathcal{D}, \quad \begin{array}{ccc} \mathcal{D} & \xleftarrow{\text{id}_{\mathcal{D}}} & \mathcal{D} \\ \uparrow \text{id}_{\mathcal{D}} & \nearrow & \uparrow F \\ \mathcal{D} & \xleftarrow{F} & C \end{array}$$

in  $\text{Cats}_2$  up to equivalence.

- 2. *Dominance.* If  $F$  is pseudoepic, then  $F$  is dominant ([Definition 11.7.1.1.1](#)).

*Proof.* [Item 1, Characterisations](#): Omitted.

[Item 2, Dominance](#): If  $F$  is pseudoepic, then

$$F^* : \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(C, \mathcal{X})$$

is pseudomonic for all  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , and thus in particular faithful. By [Item 5g of Item 5 of Definition 11.6.1.1.2](#), this is equivalent to requiring  $F$  to be dominant.  $\square$

**Question 11.7.5.1.3.** Is there a nice characterisation of the pseudoepic functors, similarly to the characterisation of pseudomonic functors given in [Item 1b of Item 1 of Definition 11.7.4.1.2](#)?

This question also appears as [\[MO 321971\]](#).

**Question 11.7.5.1.4.** A pseudomonic and pseudoepic functor is dominant, faithful, essentially injective, and full on isomorphisms. Is it necessarily an equivalence of categories? If not, how bad can this fail, i.e. how far can a pseudomonic and pseudoepic functor be from an equivalence of categories?

This question also appears as [\[MO 468334\]](#).

**Question II.7.5.1.5.** Is there a characterisation of functors  $F: \mathcal{C} \rightarrow \mathcal{D}$  such that:

1. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the precomposition functor

$$F^*: \text{Fun}(\mathcal{D}, \mathcal{X}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{X})$$

is pseudoepic?

2. For each  $\mathcal{X} \in \text{Obj}(\text{Cats})$ , the postcomposition functor

$$F_*: \text{Fun}(\mathcal{X}, \mathcal{C}) \rightarrow \text{Fun}(\mathcal{X}, \mathcal{D})$$

is pseudoepic?

This question also appears as [MO 468121a].

## II.8 Even More Conditions on Functors

### II.8.1 Injective on Objects Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.8.1.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **injective on objects** if the action on objects

$$F: \text{Obj}(\mathcal{C}) \rightarrow \text{Obj}(\mathcal{D})$$

of  $F$  is injective.

**Proposition II.8.1.1.2.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor.

1. *Characterisations.* The following conditions are equivalent:
  - (a) The functor  $F$  is injective on objects.
  - (b) The functor  $F$  is an isofibration in  $\text{Cats}_2$ .

*Proof.* **Item 1, Characterisations:** Omitted. □

## 11.8.2 Surjective on Objects Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.8.2.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **surjective on objects** if the action on objects

$$F: \text{Obj}(\mathcal{C}) \rightarrow \text{Obj}(\mathcal{D})$$

of  $F$  is surjective.

## 11.8.3 Bijective on Objects Functors

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.8.3.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **bijective on objects**<sup>27</sup> if the action on objects

$$F: \text{Obj}(\mathcal{C}) \rightarrow \text{Obj}(\mathcal{D})$$

of  $F$  is a bijection.

## 11.8.4 Functors Representably Faithful on Cores

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.8.4.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **representably faithful on cores** if, for each  $X \in \text{Obj}(\text{Cats})$ , the postcomposition by  $F$  functor

$$F_*: \text{Core}(\text{Fun}(X, \mathcal{C})) \rightarrow \text{Core}(\text{Fun}(X, \mathcal{D}))$$

is faithful.

**Remark 11.8.4.1.2.** In detail, a functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **representably faithful on cores** if, given a diagram of the form

$$\begin{array}{ccc} & \phi & \\ \alpha \downarrow & \xrightarrow{\quad} & C \xrightarrow{F} \mathcal{D} \\ \beta \downarrow & \xleftarrow{\quad} & \\ & \psi & \end{array}$$

if  $\alpha$  and  $\beta$  are natural isomorphisms and we have

$$\text{id}_F \star \alpha = \text{id}_F \star \beta,$$

then  $\alpha = \beta$ .

<sup>27</sup>*Further Terminology:* Also called a **bo** functor.

**Question 11.8.4.1.3.** Is there a characterisation of functors representably faithful on cores?

## 11.8.5 Functors Representably Full on Cores

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.8.5.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **representably full on cores** if, for each  $X \in \text{Obj}(\text{Cats})$ , the postcomposition by  $F$  functor

$$F_*: \text{Core}(\text{Fun}(X, \mathcal{C})) \rightarrow \text{Core}(\text{Fun}(X, \mathcal{D}))$$

is full.

**Remark 11.8.5.1.2.** In detail, a functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **representably full on cores** if, for each  $X \in \text{Obj}(\text{Cats})$  and each natural isomorphism

$$\beta: F \circ \phi \xRightarrow{\sim} F \circ \psi, \quad X \begin{array}{c} \xrightarrow{F \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{F \circ \psi} \end{array} \mathcal{D},$$

there exists a natural isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} \mathcal{C}$$

such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} \mathcal{C} \xrightarrow{F} \mathcal{D} = X \begin{array}{c} \xrightarrow{F \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{F \circ \psi} \end{array} \mathcal{D}$$

of pasting diagrams in  $\text{Cats}_2$ , i.e. such that we have

$$\beta = \text{id}_F \star \alpha.$$

**Question 11.8.5.1.3.** Is there a characterisation of functors representably full on cores?

This question also appears as [MO 468121a].

## II.8.6 Functors Representably Fully Faithful on Cores

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition II.8.6.I.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **representably fully faithful on cores** if, for each  $X \in \text{Obj}(\text{Cats})$ , the postcomposition by  $F$  functor

$$F_*: \text{Core}(\text{Fun}(X, \mathcal{C})) \rightarrow \text{Core}(\text{Fun}(X, \mathcal{D}))$$

is fully faithful.

**Remark II.8.6.I.2.** In detail, a functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **representably fully faithful on cores** if it satisfies the conditions in [Definitions II.8.4.I.2](#) and [II.8.5.I.2](#), i.e.:

1. For all diagrams of the form

$$\mathcal{X} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} \mathcal{C} \xrightarrow{F} \mathcal{D},$$

with  $\alpha$  and  $\beta$  natural isomorphisms, if we have  $\text{id}_F \star \alpha = \text{id}_F \star \beta$ , then  $\alpha = \beta$ .

2. For each  $X \in \text{Obj}(\text{Cats})$  and each natural isomorphism

$$\beta: F \circ \phi \xrightarrow{\sim} F \circ \psi, \quad \mathcal{X} \begin{array}{c} \xrightarrow{F \circ \phi} \\ \beta \Downarrow \Downarrow \\ \xrightarrow{F \circ \psi} \end{array} \mathcal{D}$$

of  $\mathcal{C}$ , there exists a natural isomorphism

$$\alpha: \phi \xrightarrow{\sim} \psi, \quad \mathcal{X} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \\ \xrightarrow{\psi} \end{array} \mathcal{C}$$

of  $\mathcal{C}$  such that we have an equality

$$\mathcal{X} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \\ \xrightarrow{\psi} \end{array} \mathcal{C} \xrightarrow{F} \mathcal{D} = \mathcal{X} \begin{array}{c} \xrightarrow{F \circ \phi} \\ \beta \Downarrow \Downarrow \\ \xrightarrow{F \circ \psi} \end{array} \mathcal{D}$$

of pasting diagrams in  $\text{Cats}_2$ , i.e. such that we have

$$\beta = \text{id}_F \star \alpha.$$



**Question 11.8.6.1.3.** Is there a characterisation of functors representably fully faithful on cores?

## 11.8.7 Functors Corepresentably Faithful on Cores

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.8.7.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **corepresentably faithful on cores** if, for each  $X \in \text{Obj}(\text{Cats})$ , the postcomposition by  $F$  functor

$$F_*: \text{Core}(\text{Fun}(X, \mathcal{C})) \rightarrow \text{Core}(\text{Fun}(X, \mathcal{D}))$$

is faithful.

**Remark 11.8.7.1.2.** In detail, a functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **corepresentably faithful on cores** if, given a diagram of the form

$$\mathcal{C} \xrightarrow{F} \mathcal{D} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if  $\alpha$  and  $\beta$  are natural isomorphisms and we have

$$\alpha \star \text{id}_F = \beta \star \text{id}_F,$$

then  $\alpha = \beta$ .

**Question 11.8.7.1.3.** Is there a characterisation of functors corepresentably faithful on cores?

## 11.8.8 Functors Corepresentably Full on Cores

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories.

**Definition 11.8.8.1.1.** A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is **corepresentably full on cores** if, for each  $X \in \text{Obj}(\text{Cats})$ , the postcomposition by  $F$  functor

$$F_*: \text{Core}(\text{Fun}(X, \mathcal{C})) \rightarrow \text{Core}(\text{Fun}(X, \mathcal{D}))$$

is full.

**Remark 11.8.8.1.2.** In detail, a functor  $F: C \rightarrow D$  is **corepresentably full on cores** if, for each  $X \in \text{Obj}(\text{Cats})$  and each natural isomorphism

$$\beta: \phi \circ F \xrightarrow{\sim} \psi \circ F, \quad C \begin{array}{c} \xrightarrow{\phi \circ F} \\ \Downarrow \beta \\ \xrightarrow{\psi \circ F} \end{array} X,$$

there exists a natural isomorphism

$$\alpha: \phi \xrightarrow{\sim} \psi, \quad D \begin{array}{c} \xrightarrow{\phi} \\ \Downarrow \alpha \\ \xrightarrow{\psi} \end{array} X$$

such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \Downarrow \alpha \\ \xrightarrow{\psi} \end{array} C \xrightarrow{F} D = X \begin{array}{c} \xrightarrow{F \circ \phi} \\ \Downarrow \beta \\ \xrightarrow{F \circ \psi} \end{array} D$$

of pasting diagrams in  $\text{Cats}_2$ , i.e. such that we have

$$\beta = \alpha \star \text{id}_F.$$

**Question 11.8.8.1.3.** Is there a characterisation of functors corepresentably full on cores?

This question also appears as [MO 468121a].

## 11.8.9 Functors Corepresentably Fully Faithful on Cores

Let  $C$  and  $D$  be categories.

**Definition 11.8.9.1.1.** A functor  $F: C \rightarrow D$  is **corepresentably fully faithful on cores** if, for each  $X \in \text{Obj}(\text{Cats})$ , the postcomposition by  $F$  functor

$$F_*: \text{Core}(\text{Fun}(X, C)) \rightarrow \text{Core}(\text{Fun}(X, D))$$

is fully faithful.

**Remark 11.8.9.1.2.** In detail, a functor  $F: C \rightarrow D$  is **corepresentably fully faithful on cores** if it satisfies the conditions in [Definitions 11.8.7.1.2](#) and [11.8.8.1.2](#), i.e.:

1. For all diagrams of the form

$$C \xrightarrow{F} \mathcal{D} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if  $\alpha$  and  $\beta$  are natural isomorphisms and we have

$$\alpha \star \text{id}_F = \beta \star \text{id}_F,$$

then  $\alpha = \beta$ .

2. For each  $X \in \text{Obj}(\text{Cats})$  and each natural isomorphism

$$\beta: \phi \circ F \xrightarrow{\sim} \psi \circ F, \quad C \begin{array}{c} \xrightarrow{\phi \circ F} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ F} \end{array} X,$$

there exists a natural isomorphism

$$\alpha: \phi \xrightarrow{\sim} \psi, \quad \mathcal{D} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} C \xrightarrow{F} \mathcal{D} = X \begin{array}{c} \xrightarrow{F \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{F \circ \psi} \end{array} \mathcal{D}$$

of pasting diagrams in  $\text{Cats}_2$ , i.e. such that we have

$$\beta = \alpha \star \text{id}_F.$$

**Question 11.8.9.1.3.** Is there a characterisation of functors corepresentably fully faithful on cores?

## II.9 Natural Transformations

### II.9.1 Transformations

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories and let  $F, G: \mathcal{C} \Rightarrow \mathcal{D}$  be functors.

**Definition II.9.1.1.1.** A **transformation**<sup>28</sup>  $\alpha: F \Rightarrow G$  **from  $F$  to  $G$**  is a collection

$$\{\alpha_A: F(A) \rightarrow G(A)\}_{A \in \text{Obj}(\mathcal{C})}$$

of morphisms of  $\mathcal{D}$ .

**Notation II.9.1.1.2.** We write  $\text{Trans}(F, G)$  for the set of transformations from  $F$  to  $G$ .

**Remark II.9.1.1.3.** We have an isomorphism

$$\text{Trans}(F, G) \cong \prod_{A \in \mathcal{C}} \text{Hom}_{\mathcal{D}}(F_A, G_A).$$

*Proof.* Clear. □

### II.9.2 Natural Transformations

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories and  $F, G: \mathcal{C} \Rightarrow \mathcal{D}$  be functors.

**Definition II.9.2.1.1.** A **natural transformation**  $\alpha: F \Rightarrow G$  **from  $F$  to  $G$**  is a transformation

$$\{\alpha_A: F(A) \rightarrow G(A)\}_{A \in \text{Obj}(\mathcal{C})}$$

from  $F$  to  $G$  such that, for each morphism  $f: A \rightarrow B$  of  $\mathcal{C}$ , the diagram

$$\begin{array}{ccc} F(A) & \xrightarrow{F(f)} & F(B) \\ \alpha_A \downarrow & & \downarrow \alpha_B \\ G(A) & \xrightarrow{G(f)} & G(B) \end{array}$$

commutes.

---

<sup>28</sup>*Further Terminology:* Also called an **unnatural transformation** for emphasis.

**Remark II.9.2.1.2.** Let  $\alpha: F \Rightarrow G$  be a natural transformation.

1. For each  $A \in \text{Obj}(C)$ , the morphism  $\alpha_A: F_A \rightarrow G_A$  is called the **component of  $\alpha$  at  $A$** .
2. We denote natural transformations such as  $\alpha$  in diagrams as

$$C \begin{array}{c} \xrightarrow{F} \\ \alpha \Downarrow \\ \xrightarrow{G} \end{array} \mathcal{D}.$$

**Notation II.9.2.1.3.** We write  $\text{Nat}(F, G)$  for the set of natural transformations from  $F$  to  $G$ .

**Definition II.9.2.1.4.** Two natural transformations  $\alpha, \beta: F \Rightarrow G$  are **equal** if we have

$$\alpha_A = \beta_A$$

for each  $A \in \text{Obj}(C)$ .

### II.9.3 Examples of Natural Transformations

**Example II.9.3.1.1.** The **identity natural transformation**  $\text{id}_F: F \Rightarrow F$  of  $F$  is the natural transformation consisting of the collection

$$\{(\text{id}_F)_A: F(A) \rightarrow F(A)\}_{A \in \text{Obj}(C)}$$

defined by

$$(\text{id}_F)_A \stackrel{\text{def}}{=} \text{id}_{F(A)}$$

for each  $A \in \text{Obj}(C)$ .

*Proof.* The naturality condition for  $\text{id}_F$  is the requirement that, for each morphism  $f: A \rightarrow B$  of  $C$ , the diagram

$$\begin{array}{ccc} F(A) & \xrightarrow{F(f)} & F(B) \\ \text{id}_{F(A)} \downarrow & & \downarrow \text{id}_{F(B)} \\ F(A) & \xrightarrow{F(f)} & F(B) \end{array}$$

commutes. This follows from unitality of the composition of  $\mathcal{D}$ , as we have

$$\begin{aligned} F(f) \circ \text{id}_{F(A)} &= F(f) \\ &= \text{id}_{F(B)} \circ F(f), \end{aligned}$$

where we have applied unitality twice.  $\square$

**Example 11.9.3.1.2.** Let  $A$  and  $B$  be monoids and let  $f, g: A \rightrightarrows B$  be morphisms of monoids. Applying the delooping construction of ??, we obtain functors  $Bf, Bg: BA \rightrightarrows BB$ . We then have

$$\text{Nat}(Bf, Bg) \cong \left\{ b \in B \left| \begin{array}{l} \text{for each } a \in A, \text{ we} \\ \text{have } bf(a) = g(a)b \end{array} \right. \right\}.$$

*Proof.* Unwinding the definitions in this case, we see that a transformation  $\alpha$  from  $Bf$  to  $Bg$  consists of a collection

$$\{\alpha_{\bullet}: \bullet \rightarrow \bullet\}_{\bullet \in \text{Obj}(BA)}$$

of morphisms of  $BB$  indexed by  $\text{Obj}(BA)$ . Since  $\text{Obj}(BA) = \text{pt}$  and the morphisms of  $BB$  are precisely the elements of  $B$ , it follows that  $\alpha$  corresponds precisely to the data of an element  $b \in B$ . Now, a transformation  $[b]: Bf \Rightarrow Bg$  is natural precisely if, for each  $a \in \text{Hom}_{BA}(\bullet, \bullet) \stackrel{\text{def}}{=} A$ , the diagram

$$\begin{array}{ccc} Bf(\bullet) & \xrightarrow{Bf(a)} & Bf(\bullet) \\ [b]_{\bullet} \downarrow & & \downarrow [b]_{\bullet} \\ Bg(\bullet) & \xrightarrow{Bg(a)} & Bg(\bullet) \end{array}$$

commutes. Unwinding the definitions, we see that this diagram is given by

$$\begin{array}{ccc} \bullet & \xrightarrow{f(a)} & \bullet \\ b \downarrow & & \downarrow b \\ \bullet & \xrightarrow{g(a)} & \bullet, \end{array}$$

and hence corresponds precisely to the condition  $g(a)b = bf(a)$ .  $\square$

## II.9.4 Vertical Composition of Natural Transformations

**Definition II.9.4.1.1.** The **vertical composition** of two natural transformations  $\alpha: F \Rightarrow G$  and  $\beta: G \Rightarrow H$  as in the diagram

$$\begin{array}{ccc} & F & \\ \curvearrowright & \downarrow \alpha & \curvearrowright \\ C & \xrightarrow{G} & \mathcal{D} \\ \curvearrowleft & \downarrow \beta & \curvearrowleft \\ & H & \end{array}$$

is the natural transformation  $\beta \circ \alpha: F \Rightarrow H$  consisting of the collection

$$\{(\beta \circ \alpha)_A: F(A) \rightarrow H(A)\}_{A \in \text{Obj}(C)}$$

with

$$(\beta \circ \alpha)_A \stackrel{\text{def}}{=} \beta_A \circ \alpha_A$$

for each  $A \in \text{Obj}(C)$ .

*Proof.* The naturality condition for  $\beta \circ \alpha$  is the requirement that the boundary of the diagram

$$\begin{array}{ccc} F(A) & \xrightarrow{F(f)} & F(B) \\ \alpha_A \downarrow & (1) & \downarrow \alpha_B \\ G(A) & \xrightarrow{G(f)} & G(B) \\ \beta_A \downarrow & (2) & \downarrow \beta_B \\ H(A) & \xrightarrow{H(f)} & H(B) \end{array}$$

commutes. Since

- Subdiagram (1) commutes by the naturality of  $\alpha$ .
- Subdiagram (2) commutes by the naturality of  $\beta$ .

so does the boundary diagram. Hence  $\beta \circ \alpha$  is a natural transformation.  $\square$

**Proposition II.9.4.1.2.** Let  $\mathcal{C}$ ,  $\mathcal{D}$ , and  $\mathcal{E}$  be categories.

1. *Functionality.* The assignment  $(\beta, \alpha) \mapsto \beta \circ \alpha$  defines a function

$$\circ_{F,G,H}: \text{Nat}(G, H) \times \text{Nat}(F, G) \rightarrow \text{Nat}(F, H).$$

2. *Associativity.* Let  $F, G, H, K: \mathcal{C} \rightrightarrows \mathcal{D}$  be functors. The diagram

$$\begin{array}{ccc}
 & \text{Nat}(H, K) \times (\text{Nat}(G, H) \times \text{Nat}(F, G)) & \\
 \alpha_{\text{Nat}(H,K), \text{Nat}(G,H), \text{Nat}(F,G)}^{\text{Sets}} \nearrow & & \searrow \text{id}_{\text{Nat}(H,K)} \times \circ_{F,G,H} \\
 (\text{Nat}(H, K) \times \text{Nat}(G, H)) \times \text{Nat}(F, G) & & \text{Nat}(H, K) \times \text{Nat}(F, H) \\
 \downarrow \circ_{G,H,K} \times \text{id}_{\text{Nat}(F,G)} & & \downarrow \circ_{F,H,K} \\
 \text{Nat}(G, K) \times \text{Nat}(F, G) & \xrightarrow{\circ_{F,G,K}} & \text{Nat}(F, K)
 \end{array}$$

commutes, i.e. given natural transformations

$$F \xRightarrow{\alpha} G \xRightarrow{\beta} H \xRightarrow{\gamma} K,$$

we have

$$(\gamma \circ \beta) \circ \alpha = \gamma \circ (\beta \circ \alpha).$$

3. *Unitality.* Let  $F, G: \mathcal{C} \rightrightarrows \mathcal{D}$  be functors.

(a) *Left Unitality.* The diagram

$$\begin{array}{ccc}
 \text{pt} \times \text{Nat}(F, G) & & \\
 \downarrow [\text{id}_G] \times \text{id}_{\text{Nat}(F,G)} & \nearrow \lambda_{\text{Nat}(F,G)}^{\text{Sets}} & \\
 \text{Nat}(G, G) \times \text{Nat}(F, G) & \xrightarrow{\circ_{F,G,G}} & \text{Nat}(F, G)
 \end{array}$$

commutes, i.e. given a natural transformation  $\alpha: F \rightrightarrows G$ , we have

$$\text{id}_G \circ \alpha = \alpha.$$



(b) *Right Unitality*. The diagram

$$\begin{array}{ccc}
 \text{Nat}(F, G) \times \text{pt} & & \\
 \downarrow \text{id}_{\text{Nat}(F, G)} \times [\text{id}_F] & \searrow \beta_{\text{Nat}(F, G)}^{\text{Sets}} & \\
 \text{Nat}(F, G) \times \text{Nat}(F, F) & \xrightarrow{\circ_{F, F, G}^C} & \text{Nat}(F, G)
 \end{array}$$

commutes, i.e. given a natural transformation  $\alpha: F \Rightarrow G$ , we have

$$\alpha \circ \text{id}_F = \alpha.$$

4. *Middle Four Exchange*. Let  $F_1, F_2, F_3: \mathcal{C} \rightarrow \mathcal{D}$  and  $G_1, G_2, G_3: \mathcal{D} \rightarrow \mathcal{E}$  be functors. The diagram

$$\begin{array}{ccc}
 (\text{Nat}(G_2, G_3) \times \text{Nat}(G_1, G_2)) \times (\text{Nat}(F_2, F_3) \times \text{Nat}(F_1, F_2)) & \xleftarrow{\beta_4} & (\text{Nat}(G_2, G_3) \times \text{Nat}(F_2, F_3)) \times (\text{Nat}(G_1, G_2) \times \text{Nat}(F_1, F_2)) \\
 \downarrow \circ_{G_1, G_2, G_3} \times \circ_{F_1, F_2, F_3} & & \downarrow \star_{F_2, F_3, G_2, G_3} \times \star_{F_1, F_2, G_1, G_2} \\
 \text{Nat}(G_1, G_3) \times \text{Nat}(F_1, F_3) & & \text{Nat}(G_2 \circ F_2, G_3 \circ F_3) \times \text{Nat}(G_1 \circ F_1, G_2 \circ F_2) \\
 \searrow \star_{F_1, F_3, G_1, G_3} & & \swarrow \circ_{G_1 \circ F_1, G_2 \circ F_2, G_3 \circ F_3} \\
 & \text{Nat}(G_1 \circ F_1, G_3 \circ F_3) &
 \end{array}$$

commutes, i.e. given a diagram

$$\begin{array}{ccccc}
 & F_1 & & G_1 & \\
 & \downarrow \alpha & & \downarrow \beta & \\
 \mathcal{C} & \xrightarrow{F_2} & \mathcal{D} & \xrightarrow{G_2} & \mathcal{E} \\
 & \downarrow \alpha' & & \downarrow \beta' & \\
 & F_3 & & G_3 &
 \end{array}$$

in  $\text{Cats}_2$ , we have

$$(\beta' \star \alpha') \circ (\beta \star \alpha) = (\beta' \circ \beta) \star (\alpha' \circ \alpha).$$

*Proof.* **Item 1, Functionality:** Clear.

**Item 2, Associativity:** Indeed, we have

$$((\gamma \circ \beta) \circ \alpha)_A \stackrel{\text{def}}{=} (\gamma \circ \beta)_A \circ \alpha_A$$

$$\begin{aligned}
 &\stackrel{\text{def}}{=} (\gamma_A \circ \beta_A) \circ \alpha_A \\
 &= \gamma_A \circ (\beta_A \circ \alpha_A) \\
 &\stackrel{\text{def}}{=} \gamma_A \circ (\beta \circ \alpha)_A \\
 &\stackrel{\text{def}}{=} (\gamma \circ (\beta \circ \alpha))_A
 \end{aligned}$$

for each  $A \in \text{Obj}(C)$ , showing the desired equality.

*Item 3, Unitality:* We have

$$\begin{aligned}
 (\text{id}_G \circ \alpha)_A &= \text{id}_G \circ \alpha_A \\
 &= \alpha_A, \\
 (\alpha \circ \text{id}_F)_A &= \alpha_A \circ \text{id}_F \\
 &= \alpha_A
 \end{aligned}$$

for each  $A \in \text{Obj}(C)$ , showing the desired equality.

*Item 4, Middle Four Exchange:* This is proved in *Item 4* of *Definition 11.9.5.1.3*.  $\square$

## 11.9.5 Horizontal Composition of Natural Transformations

**Definition 11.9.5.1.1.** The **horizontal composition**<sup>29,30</sup> of two natural transformations  $\alpha: F \Rightarrow G$  and  $\beta: H \Rightarrow K$  as in the diagram

$$\begin{array}{ccccc}
 & F & & H & \\
 C & \xrightarrow{\quad} & \mathcal{D} & \xrightarrow{\quad} & \mathcal{E} \\
 & \alpha \Downarrow & & \beta \Downarrow & \\
 & G & & K &
 \end{array}$$

of  $\alpha$  and  $\beta$  is the natural transformation

$$\beta \star \alpha: (H \circ F) \Rightarrow (K \circ G),$$

as in the diagram

$$\begin{array}{ccc}
 & H \circ F & \\
 C & \xrightarrow{\quad} & \mathcal{E}, \\
 & \beta \star \alpha \Downarrow & \\
 & K \circ G &
 \end{array}$$

<sup>29</sup>*Further Terminology:* Also called the **Godement product** of  $\alpha$  and  $\beta$ .

<sup>30</sup>Horizontal composition forms a map

$$\star_{(F,H),(G,K)}: \text{Nat}(H, K) \times \text{Nat}(F, G) \rightarrow \text{Nat}(H \circ F, K \circ G).$$

consisting of the collection

$$\{(\beta \star \alpha)_A: H(F(A)) \rightarrow K(G(A))\}_{A \in \text{Obj}(\mathcal{C})},$$

of morphisms of  $\mathcal{E}$  with

$$\begin{aligned} (\beta \star \alpha)_A &\stackrel{\text{def}}{=} \beta_{G(A)} \circ H(\alpha_A) \\ &= K(\alpha_A) \circ \beta_{F(A)}, \end{aligned} \quad \begin{array}{ccc} H(F(A)) & \xrightarrow{H(\alpha_A)} & H(G(A)) \\ \beta_{F(A)} \downarrow & & \downarrow \beta_{G(A)} \\ K(F(A)) & \xrightarrow{K(\alpha_A)} & K(G(A)). \end{array}$$

*Proof.* First, we claim that we indeed have

$$\beta_{G(A)} \circ H(\alpha_A) = K(\alpha_A) \circ \beta_{F(A)}, \quad \begin{array}{ccc} H(F(A)) & \xrightarrow{H(\alpha_A)} & H(G(A)) \\ \beta_{F(A)} \downarrow & & \downarrow \beta_{G(A)} \\ K(F(A)) & \xrightarrow{K(\alpha_A)} & K(G(A)). \end{array}$$

This is, however, simply the naturality square for  $\beta$  applied to the morphism  $\alpha_A: F(A) \rightarrow G(A)$ . Next, we check the naturality condition for  $\beta \star \alpha$ , which is the requirement that the boundary of the diagram

$$\begin{array}{ccc} H(F(A)) & \xrightarrow{H(F(f))} & H(F(B)) \\ \downarrow H(\alpha_A) & (1) & \downarrow H(\alpha_B) \\ H(G(A)) & \xrightarrow{H(G(f))} & H(G(B)) \\ \downarrow \beta_{G(A)} & (2) & \downarrow \beta_{G(B)} \\ K(G(A)) & \xrightarrow{K(G(f))} & K(G(B)) \end{array}$$

commutes. Since

- Subdiagram (1) commutes by the naturality of  $\alpha$ .
- Subdiagram (2) commutes by the naturality of  $\beta$ .

so does the boundary diagram. Hence  $\beta \circ \alpha$  is a natural transformation.<sup>31</sup>  $\square$

**Definition II.9.5.1.2.** Let

$$\mathcal{X} \xrightarrow{F} \mathcal{C} \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} \mathcal{D} \xrightarrow{G} \mathcal{Y}$$

be a diagram in  $\mathbf{Cats}_2$ .

1. The **left whiskering of  $\alpha$  with  $G$**  is the natural transformation<sup>32</sup>

$$\mathrm{id}_G \star \alpha: G \circ \phi \Longrightarrow G \circ \psi.$$

2. The **right whiskering of  $\alpha$  with  $F$**  is the natural transformation<sup>33</sup>

$$\alpha \star \mathrm{id}_F: \phi \circ F \Longrightarrow \psi \circ F.$$

**Proposition II.9.5.1.3.** Let  $\mathcal{C}$ ,  $\mathcal{D}$ , and  $\mathcal{E}$  be categories.

1. *Functionality.* The assignment  $(\beta, \alpha) \mapsto \beta \star \alpha$  defines a function

$$\star_{(F,G),(H,K)}: \mathrm{Nat}(H, K) \times \mathrm{Nat}(F, G) \rightarrow \mathrm{Nat}(H \circ F, K \circ G).$$

2. *Associativity.* Let

$$\mathcal{C} \begin{array}{c} \xrightarrow{F_1} \\ \xrightarrow{G_1} \end{array} \mathcal{D} \begin{array}{c} \xrightarrow{F_2} \\ \xrightarrow{G_2} \end{array} \mathcal{E} \begin{array}{c} \xrightarrow{F_3} \\ \xrightarrow{G_3} \end{array} \mathcal{F}$$

<sup>31</sup>Reference: [Bor94, Proposition 1.3.4].

<sup>32</sup>Further Notation: Also written  $G\alpha$  or  $G \star \alpha$ , although we won't use either of these notations in this work.

<sup>33</sup>Further Notation: Also written  $\alpha F$  or  $\alpha \star F$ , although we won't use either of these notations in this work.

be a diagram in  $\mathbf{Cats}_2$ . The diagram

$$\begin{array}{ccc}
 \text{Nat}(F_3, G_3) \times \text{Nat}(F_2, G_2) \times \text{Nat}(F_1, G_1) & \xrightarrow{\star_{(F_2, G_2), (F_3, G_3)} \times \text{id}} & \text{Nat}(F_3 \circ F_2, G_3 \circ G_2) \times \text{Nat}(F_1, G_1) \\
 \downarrow \text{id} \times \star_{(F_1, G_1), (F_2, G_2)} & & \downarrow \star_{(F_3 \circ F_2), (G_3 \circ G_2, F_1, G_1)} \\
 \text{Nat}(F_3, G_3) \times \text{Nat}(F_2 \circ F_1, G_2 \circ G_1) & \xrightarrow{\star_{(F_2 \circ F_1), (G_2 \circ G_1, F_3, G_3)}} & \text{Nat}(F_3 \circ F_2 \circ F_1, G_3 \circ G_2 \circ G_1)
 \end{array}$$

commutes, i.e. given natural transformations

$$\begin{array}{ccccc}
 C & \xrightleftharpoons[\alpha \downarrow]{F_1} & D & \xrightleftharpoons[\beta \downarrow]{F_2} & E & \xrightleftharpoons[\gamma \downarrow]{F_3} & F, \\
 & G_1 & & G_2 & & G_3 &
 \end{array}$$

we have

$$(\gamma \star \beta) \star \alpha = \gamma \star (\beta \star \alpha).$$

3. *Interaction With Identities.* Let  $F: C \rightarrow D$  and  $G: D \rightarrow E$  be functors. The diagram

$$\begin{array}{ccc}
 \text{pt} \times \text{pt} & \xrightarrow{[\text{id}_G] \times [\text{id}_F]} & \text{Nat}(G, G) \times \text{Nat}(F, F) \\
 \uparrow \wr & & \downarrow \star_{(F, F), (G, G)} \\
 \text{pt} & \xrightarrow{[\text{id}_{G \circ F}]} & \text{Nat}(G \circ F, G \circ F)
 \end{array}$$

commutes, i.e. we have

$$\text{id}_G \star \text{id}_F = \text{id}_{G \circ F}.$$

4. *Middle Four Exchange.* Let  $F_1, F_2, F_3: C \rightarrow D$  and  $G_1, G_2, G_3: D \rightarrow E$  be functors. The diagram

$$\begin{array}{ccc}
 (\text{Nat}(G_2, G_3) \times \text{Nat}(G_1, G_2)) \times (\text{Nat}(F_2, F_3) \times \text{Nat}(F_1, F_2)) & \xleftarrow{\mu_4} & (\text{Nat}(G_2, G_3) \times \text{Nat}(F_2, F_3)) \times (\text{Nat}(G_1, G_2) \times \text{Nat}(F_1, F_2)) \\
 \downarrow \circ_{G_1, G_2, G_3} \times \circ_{F_1, F_2, F_3} & & \downarrow \star_{F_2, F_3, G_2, G_3} \times \star_{F_1, F_2, G_1, G_2} \\
 \text{Nat}(G_1, G_3) \times \text{Nat}(F_1, F_3) & & \text{Nat}(G_2 \circ F_2, G_3 \circ F_3) \times \text{Nat}(G_1 \circ F_1, G_2 \circ F_2) \\
 \searrow \star_{F_1, F_3, G_1, G_3} & & \swarrow \circ_{G_1 \circ F_1, G_2 \circ F_2, G_3 \circ F_3} \\
 & \text{Nat}(G_1 \circ F_1, G_3 \circ F_3) &
 \end{array}$$

commutes, i.e. given a diagram

$$\begin{array}{ccccc}
 & F_1 & & G_1 & \\
 & \alpha \Downarrow & & \beta \Downarrow & \\
 C & \xrightarrow{F_2} & \mathcal{D} & \xrightarrow{G_2} & \mathcal{E} \\
 & \alpha' \Downarrow & & \beta' \Downarrow & \\
 & F_3 & & G_3 &
 \end{array}$$

in  $\mathbf{Cats}_2$ , we have

$$(\beta' \star \alpha') \circ (\beta \star \alpha) = (\beta' \circ \beta) \star (\alpha' \circ \alpha).$$

*Proof.* **Item 1, Functionality:** Clear.

**Item 2, Associativity:** Omitted.

**Item 3, Interaction With Identities:** We have

$$\begin{aligned}
 (\mathrm{id}_G \star \mathrm{id}_F)_A &\stackrel{\mathrm{def}}{=} (\mathrm{id}_G)_{F_A} \circ G_{(\mathrm{id}_F)_A} \\
 &\stackrel{\mathrm{def}}{=} \mathrm{id}_{G_{F_A}} \circ G_{\mathrm{id}_{F_A}} \\
 &= \mathrm{id}_{G_{F_A}} \circ \mathrm{id}_{G_{F_A}} \\
 &= \mathrm{id}_{G_{F_A}} \\
 &\stackrel{\mathrm{def}}{=} (\mathrm{id}_{G \circ F})_A
 \end{aligned}$$

for each  $A \in \mathrm{Obj}(C)$ , showing the desired equality.

**Item 4, Middle Four Exchange:** Let  $A \in \mathrm{Obj}(C)$  and consider the diagram

$$\begin{array}{ccccccc}
 & & G_1(F_3(A)) & & & & \\
 & G_1(\alpha'_A) \nearrow & & \searrow \beta_{F_3(A)} & & & \\
 G_1(F_1(A)) & \xrightarrow{G_1(\alpha_A)} & G_1(F_2(A)) & (1) & G_2(F_3(A)) & \xrightarrow{\beta'_{F_3(A)}} & G_3(F_3(A)). \\
 & \searrow \beta_{F_2(A)} & & \nearrow G_2(\alpha'_A) & & & \\
 & & G_2(F_2(A)) & & & &
 \end{array}$$

The top composition

$$\begin{array}{ccccc}
 & & G_1(F_3(A)) & & \\
 & \nearrow^{G_1(\alpha'_A)} & & \searrow_{\beta_{F_3(A)}} & \\
 G_1(F_1(A)) & \xrightarrow{G_1(\alpha_A)} & G_1(F_2(A)) & \quad (1) \quad & G_2(F_3(A)) \xrightarrow{\beta'_{F_3(A)}} G_3(F_3(A)). \\
 & \searrow_{\beta_{F_2(A)}} & & \nearrow_{G_2(\alpha'_A)} & \\
 & & G_2(F_2(A)) & & 
 \end{array}$$

is given by  $((\beta' \circ \beta) \star (\alpha' \circ \alpha))_A$ , while the bottom composition

$$\begin{array}{ccccc}
 & & G_1(F_3(A)) & & \\
 & \nearrow^{G_1(\alpha'_A)} & & \searrow_{\beta_{F_3(A)}} & \\
 G_1(F_1(A)) & \xrightarrow{G_1(\alpha_A)} & G_1(F_2(A)) & \quad (1) \quad & G_2(F_3(A)) \xrightarrow{\beta'_{F_3(A)}} G_3(F_3(A)). \\
 & \searrow_{\beta_{F_2(A)}} & & \nearrow_{G_2(\alpha'_A)} & \\
 & & G_2(F_2(A)) & & 
 \end{array}$$

is given by  $((\beta' \star \alpha') \circ (\beta \star \alpha))_A$ . Now, Subdiagram (1) corresponds to the naturality condition

$$\begin{array}{ccc}
 G_1(F_2(A)) & \xrightarrow{G_1(\alpha'_A)} & G_1(F_3(A)) \\
 \downarrow \beta_{F_2(A)} & & \downarrow \beta_{F_3(A)} \\
 G_2(F_2(A)) & \xrightarrow{G_2(\alpha'_A)} & G_2(F_3(A))
 \end{array}$$

$G_2(\alpha'_A) \circ \beta_{F_2(A)} = \beta_{F_3(A)} \circ G_1(\alpha'_A),$

for  $\beta: G_1 \Rightarrow G_2$  at  $\alpha'_A: F_2(A) \rightarrow F_3(A)$ , and thus commutes. Thus we have

$$((\beta' \circ \beta) \star (\alpha' \circ \alpha))_A = ((\beta' \star \alpha') \circ (\beta \star \alpha))_A$$

for each  $A \in \text{Obj}(C)$  and therefore

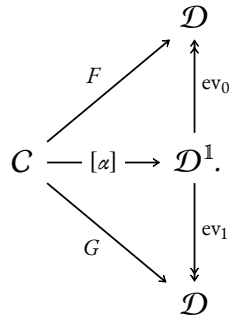
$$(\beta' \star \alpha') \circ (\beta \star \alpha) = (\beta' \circ \beta) \star (\alpha' \circ \alpha).$$

This finishes the proof.  $\square$

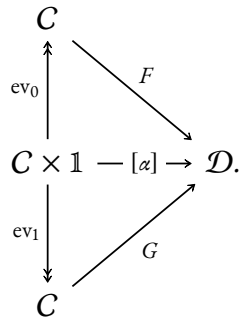
## II.9.6 Properties of Natural Transformations

**Proposition II.9.6.I.I.** Let  $F, G: C \rightrightarrows \mathcal{D}$  be functors. The following data are equivalent:<sup>34</sup>

1. A natural transformation  $\alpha: F \Rightarrow G$ .
2. A functor  $[\alpha]: C \rightarrow \mathcal{D}^1$  filling the diagram



3. A functor  $[\alpha]: C \times 1 \rightarrow \mathcal{D}$  filling the diagram



<sup>34</sup>Taken from [MO 64365].



*Proof. From Item 1 to Item 2 and Back:* We may identify  $\mathcal{D}^1$  with  $\text{Arr}(\mathcal{D})$ . Given a natural transformation  $\alpha: F \Rightarrow G$ , we have a functor

$$\begin{aligned} [\alpha]: C &\longrightarrow \mathcal{D}^1 \\ A &\longmapsto \alpha_A \end{aligned}$$

$$(f: A \rightarrow B) \longmapsto \left( \begin{array}{ccc} F_A & \xrightarrow{F_f} & F_B \\ \alpha_A \downarrow & & \downarrow \alpha_B \\ G_A & \xrightarrow{G_f} & G_B \end{array} \right)$$

making the diagram in **Item 2** commute. Conversely, every such functor gives rise to a natural transformation from  $F$  to  $G$ , and these constructions are inverse to each other.

*From Item 2 to Item 3 and Back:* This follows from **Item 3** of **Definition II.10.1.1.2**.

□

## II.9.7 Natural Isomorphisms

Let  $C$  and  $\mathcal{D}$  be categories and let  $F, G: C \Rightarrow \mathcal{D}$  be functors.

**Definition II.9.7.1.1.** A natural transformation  $\alpha: F \Rightarrow G$  is a **natural isomorphism** if there exists a natural transformation  $\alpha^{-1}: G \Rightarrow F$  such that

$$\begin{aligned} \alpha^{-1} \circ \alpha &= \text{id}_F, \\ \alpha \circ \alpha^{-1} &= \text{id}_G. \end{aligned}$$

**Proposition II.9.7.1.2.** Let  $\alpha: F \Rightarrow G$  be a natural transformation.

1. *Characterisations.* The following conditions are equivalent:
  - (a) The natural transformation  $\alpha$  is a natural isomorphism.
  - (b) For each  $A \in \text{Obj}(C)$ , the morphism  $\alpha_A: F_A \rightarrow G_A$  is an isomorphism.
2. *Componentwise Inverses of Natural Transformations Assemble Into Natural*

*Transformations.* Let  $\alpha^{-1}: G \Rightarrow F$  be a transformation such that, for each  $A \in \text{Obj}(C)$ , we have

$$\begin{aligned}\alpha_A^{-1} \circ \alpha_A &= \text{id}_{F(A)}, \\ \alpha_A \circ \alpha_A^{-1} &= \text{id}_{G(A)}.\end{aligned}$$

Then  $\alpha^{-1}$  is a natural transformation.

*Proof.* **Item 1, Characterisations:** The implication **Item 1a**  $\Rightarrow$  **Item 1b** is clear, whereas the implication **Item 1b**  $\Rightarrow$  **Item 1a** follows from **Item 2**.

**Item 2, Componentwise Inverses of Natural Transformations Assemble Into Natural Transformations:** The naturality condition for  $\alpha^{-1}$  corresponds to the commutativity of the diagram

$$\begin{array}{ccc} G(A) & \xrightarrow{G(f)} & G(B) \\ \alpha_A^{-1} \downarrow & & \downarrow \alpha_B^{-1} \\ F(A) & \xrightarrow{F(f)} & F(B) \end{array}$$

for each  $A, B \in \text{Obj}(C)$  and each  $f \in \text{Hom}_C(A, B)$ . Considering the diagram

$$\begin{array}{ccccc} G(A) & \xrightarrow{G(f)} & G(B) & & \\ \alpha_A^{-1} \downarrow & & \downarrow \alpha_B^{-1} & & \\ F(A) & \xrightarrow{F(f)} & F(B) & & \\ \alpha_A \downarrow & & \downarrow \alpha_B & & \\ G(A) & \xrightarrow{G(f)} & G(B), & & \end{array} \quad \begin{array}{c} (1) \\ (2) \end{array}$$

where the boundary diagram as well as Subdiagram (2) commute, we have

$$\begin{aligned}G(f) &= G(f) \circ \text{id}_{G(A)} \\ &= G(f) \circ \alpha_A \circ \alpha_A^{-1} \\ &= \alpha_B \circ F(f) \circ \alpha_A^{-1}.\end{aligned}$$

Postcomposing both sides with  $\alpha_B^{-1}$ , we get

$$\begin{aligned}\alpha_B^{-1} \circ G(f) &= \alpha_B^{-1} \circ \alpha_B \circ F(f) \circ \alpha_A^{-1} \\ &= \text{id}_{F(B)} \circ F(f) \circ \alpha_A^{-1} \\ &= F(f) \circ \alpha_A^{-1},\end{aligned}$$

which is the naturality condition we wanted to show. Thus  $\alpha^{-1}$  is a natural transformation.  $\square$

## II.10 Categories of Categories

### II.10.1 Functor Categories

Let  $\mathcal{C}$  be a category and  $\mathcal{D}$  be a small category.

**Definition II.10.1.1.1.** The **category of functors from  $\mathcal{C}$  to  $\mathcal{D}$** <sup>35</sup> is the category  $\text{Fun}(\mathcal{C}, \mathcal{D})$ <sup>36</sup> where

- *Objects.* The objects of  $\text{Fun}(\mathcal{C}, \mathcal{D})$  are functors from  $\mathcal{C}$  to  $\mathcal{D}$ .
- *Morphisms.* For each  $F, G \in \text{Obj}(\text{Fun}(\mathcal{C}, \mathcal{D}))$ , we have

$$\text{Hom}_{\text{Fun}(\mathcal{C}, \mathcal{D})}(F, G) \stackrel{\text{def}}{=} \text{Nat}(F, G).$$

- *Identities.* For each  $F \in \text{Obj}(\text{Fun}(\mathcal{C}, \mathcal{D}))$ , the unit map

$$\mathbb{1}_F^{\text{Fun}(\mathcal{C}, \mathcal{D})} : \text{pt} \rightarrow \text{Nat}(F, F)$$

of  $\text{Fun}(\mathcal{C}, \mathcal{D})$  at  $F$  is given by

$$\text{id}_F^{\text{Fun}(\mathcal{C}, \mathcal{D})} \stackrel{\text{def}}{=} \text{id}_F,$$

where  $\text{id}_F : F \implies F$  is the identity natural transformation of  $F$  of **Definition II.9.3.1.1**.

<sup>35</sup>*Further Terminology:* Also called the **functor category**  $\text{Fun}(\mathcal{C}, \mathcal{D})$ .

<sup>36</sup>*Further Notation:* Also written  $\mathcal{D}^{\mathcal{C}}$  and  $[\mathcal{C}, \mathcal{D}]$ .

- *Composition.* For each  $F, G, H \in \text{Obj}(\text{Fun}(C, \mathcal{D}))$ , the composition map

$$\circ_{F,G,H}^{\text{Fun}(C,\mathcal{D})} : \text{Nat}(G, H) \times \text{Nat}(F, G) \rightarrow \text{Nat}(F, H)$$

of  $\text{Fun}(C, \mathcal{D})$  at  $(F, G, H)$  is given by

$$\beta \circ_{F,G,H}^{\text{Fun}(C,\mathcal{D})} \alpha \stackrel{\text{def}}{=} \beta \circ \alpha,$$

where  $\beta \circ \alpha$  is the vertical composition of  $\alpha$  and  $\beta$  of [Item 1](#) of [Definition II.9.4.1.2](#).

**Proposition II.10.1.1.2.** Let  $C$  and  $\mathcal{D}$  be categories and let  $F: C \rightarrow \mathcal{D}$  be a functor.

1. *Functoriality.* The assignments  $C, \mathcal{D}, (C, \mathcal{D}) \mapsto \text{Fun}(C, \mathcal{D})$  define functors

$$\begin{aligned} \text{Fun}(C, -) &: \text{Cats} && \rightarrow \text{Cats}, \\ \text{Fun}(-, \mathcal{D}) &: \text{Cats}^{\text{op}} && \rightarrow \text{Cats}, \\ \text{Fun}(-_1, -_2) &: \text{Cats}^{\text{op}} \times \text{Cats} && \rightarrow \text{Cats}. \end{aligned}$$

2. *2-Functoriality.* The assignments  $C, \mathcal{D}, (C, \mathcal{D}) \mapsto \text{Fun}(C, \mathcal{D})$  define 2-functors

$$\begin{aligned} \text{Fun}(C, -) &: \text{Cats}_2 && \rightarrow \text{Cats}_2, \\ \text{Fun}(-, \mathcal{D}) &: \text{Cats}_2^{\text{op}} && \rightarrow \text{Cats}_2, \\ \text{Fun}(-_1, -_2) &: \text{Cats}_2^{\text{op}} \times \text{Cats}_2 && \rightarrow \text{Cats}_2. \end{aligned}$$

3. *Adjointness.* We have adjunctions

$$\begin{aligned} (C \times - \dashv \text{Fun}(C, -)) &: \text{Cats} \begin{array}{c} \xrightarrow{C \times -} \\ \perp \\ \xleftarrow{\text{Fun}(C, -)} \end{array} \text{Cats}, \\ (- \times \mathcal{D} \dashv \text{Fun}(\mathcal{D}, -)) &: \text{Cats} \begin{array}{c} \xrightarrow{- \times \mathcal{D}} \\ \perp \\ \xleftarrow{\text{Fun}(\mathcal{D}, -)} \end{array} \text{Cats}, \end{aligned}$$

witnessed by bijections of sets

$$\begin{aligned} \text{Hom}_{\text{Cats}}(C \times \mathcal{D}, \mathcal{E}) &\cong \text{Hom}_{\text{Cats}}(\mathcal{D}, \text{Fun}(C, \mathcal{E})), \\ \text{Hom}_{\text{Cats}}(C \times \mathcal{D}, \mathcal{E}) &\cong \text{Hom}_{\text{Cats}}(C, \text{Fun}(\mathcal{D}, \mathcal{E})), \end{aligned}$$

natural in  $C, \mathcal{D}, \mathcal{E} \in \text{Obj}(\text{Cats})$ .

4. *2-Adjointness*. We have 2-adjunctions

$$\begin{aligned} (C \times - \dashv \text{Fun}(C, -)) : \quad & \text{Cats}_2 \begin{array}{c} \xrightarrow{C \times -} \\ \perp_2 \\ \xleftarrow{\text{Fun}(C, -)} \end{array} \text{Cats}_2, \\ (- \times \mathcal{D} \dashv \text{Fun}(\mathcal{D}, -)) : \quad & \text{Cats}_2 \begin{array}{c} \xrightarrow{- \times \mathcal{D}} \\ \perp_2 \\ \xleftarrow{\text{Fun}(\mathcal{D}, -)} \end{array} \text{Cats}_2, \end{aligned}$$

witnessed by isomorphisms of categories

$$\begin{aligned} \text{Fun}(C \times \mathcal{D}, \mathcal{E}) &\cong \text{Fun}(\mathcal{D}, \text{Fun}(C, \mathcal{E})), \\ \text{Fun}(C \times \mathcal{D}, \mathcal{E}) &\cong \text{Fun}(C, \text{Fun}(\mathcal{D}, \mathcal{E})), \end{aligned}$$

natural in  $C, \mathcal{D}, \mathcal{E} \in \text{Obj}(\text{Cats}_2)$ .

5. *Interaction With Punctual Categories*. We have a canonical isomorphism of categories

$$\text{Fun}(\text{pt}, C) \cong C,$$

natural in  $C \in \text{Obj}(\text{Cats})$ .

6. *Objectwise Computation of Co/Limits*. Let

$$D: I \rightarrow \text{Fun}(C, \mathcal{D})$$

be a diagram in  $\text{Fun}(C, \mathcal{D})$ . We have isomorphisms

$$\begin{aligned} \lim(D)_A &\cong \lim_{i \in I} (D_i(A)), \\ \text{colim}(D)_A &\cong \text{colim}_{i \in I} (D_i(A)), \end{aligned}$$

naturally in  $A \in \text{Obj}(C)$ .

7. *Interaction With Co/Completeness*. If  $\mathcal{E}$  is co/complete, then so is  $\text{Fun}(C, \mathcal{E})$ .

8. *Monomorphisms and Epimorphisms*. Let  $\alpha: F \Rightarrow G$  be a morphism of  $\text{Fun}(C, \mathcal{D})$ . The following conditions are equivalent:

(a) The natural transformation

$$\alpha: F \Rightarrow G$$

is a monomorphism (resp. epimorphism) in  $\text{Fun}(C, \mathcal{D})$ .

(b) For each  $A \in \text{Obj}(C)$ , the morphism

$$\alpha_A: F_A \rightarrow G_A$$

is a monomorphism (resp. epimorphism) in  $\mathcal{D}$ .

*Proof.* **Item 1, Functoriality:** Omitted.

**Item 2, 2-Functoriality:** Omitted.

**Item 3, Adjointness:** Omitted.

**Item 4, 2-Adjointness:** Omitted.

**Item 5, Interaction With Punctual Categories:** Omitted.

**Item 6, Objectwise Computation of Co/Limits:** Omitted.

**Item 7, Interaction With Co/Completeness:** This follows from ??.

**Item 8, Monomorphisms and Epimorphisms:** Omitted.  $\square$

## II.10.2 The Category of Categories and Functors

**Definition II.10.2.1.1.** The **category of (small) categories and functors** is the category  $\text{Cats}$  where

- *Objects.* The objects of  $\text{Cats}$  are small categories.
- *Morphisms.* For each  $C, \mathcal{D} \in \text{Obj}(\text{Cats})$ , we have

$$\text{Hom}_{\text{Cats}}(C, \mathcal{D}) \stackrel{\text{def}}{=} \text{Obj}(\text{Fun}(C, \mathcal{D})).$$

- *Identities.* For each  $C \in \text{Obj}(\text{Cats})$ , the unit map

$$\mathbb{1}_C^{\text{Cats}}: \text{pt} \rightarrow \text{Hom}_{\text{Cats}}(C, C)$$

of  $\text{Cats}$  at  $C$  is defined by

$$\text{id}_C^{\text{Cats}} \stackrel{\text{def}}{=} \text{id}_C,$$

where  $\text{id}_C: C \rightarrow C$  is the identity functor of  $C$  of **Definition II.5.1.1.4**.

- *Composition.* For each  $C, \mathcal{D}, \mathcal{E} \in \text{Obj}(\text{Cats})$ , the composition map

$$\circ_{C, \mathcal{D}, \mathcal{E}}^{\text{Cats}}: \text{Hom}_{\text{Cats}}(\mathcal{D}, \mathcal{E}) \times \text{Hom}_{\text{Cats}}(C, \mathcal{D}) \rightarrow \text{Hom}_{\text{Cats}}(C, \mathcal{E})$$

of  $\text{Cats}$  at  $(C, \mathcal{D}, \mathcal{E})$  is given by

$$G \circ_{C, \mathcal{D}, \mathcal{E}}^{\text{Cats}} F \stackrel{\text{def}}{=} G \circ F,$$

where  $G \circ F: C \rightarrow \mathcal{E}$  is the composition of  $F$  and  $G$  of **Definition II.5.1.1.5**.

**Proposition II.10.2.1.2.** Let  $C$  be a category.

1. *Co/Completeness.* The category  $\mathbf{Cats}$  is complete and cocomplete.
2. *Cartesian Monoidal Structure.* The quadruple  $(\mathbf{Cats}, \times, \text{pt}, \text{Fun})$  is a Cartesian closed monoidal category.

*Proof.* **Item 1, Co/Completeness:** Omitted.

**Item 2, Cartesian Monoidal Structure:** Omitted.  $\square$

### II.10.3 The 2-Category of Categories, Functors, and Natural Transformations

**Definition II.10.3.1.1.** The 2-category of (small) categories, functors, and natural transformations is the 2-category  $\mathbf{Cats}_2$  where

- *Objects.* The objects of  $\mathbf{Cats}_2$  are small categories.
- *Hom-Categories.* For each  $C, \mathcal{D} \in \text{Obj}(\mathbf{Cats}_2)$ , we have

$$\text{Hom}_{\mathbf{Cats}_2}(C, \mathcal{D}) \stackrel{\text{def}}{=} \text{Fun}(C, \mathcal{D}).$$

- *Identities.* For each  $C \in \text{Obj}(\mathbf{Cats}_2)$ , the unit functor

$$1_C^{\mathbf{Cats}_2}: \text{pt} \rightarrow \text{Fun}(C, C)$$

of  $\mathbf{Cats}_2$  at  $C$  is the functor picking the identity functor  $\text{id}_C: C \rightarrow C$  of  $C$ .

- *Composition.* For each  $C, \mathcal{D}, \mathcal{E} \in \text{Obj}(\mathbf{Cats}_2)$ , the composition bifunctor

$$\circ_{C, \mathcal{D}, \mathcal{E}}^{\mathbf{Cats}_2}: \text{Hom}_{\mathbf{Cats}_2}(\mathcal{D}, \mathcal{E}) \times \text{Hom}_{\mathbf{Cats}_2}(C, \mathcal{D}) \rightarrow \text{Hom}_{\mathbf{Cats}_2}(C, \mathcal{E})$$

of  $\mathbf{Cats}_2$  at  $(C, \mathcal{D}, \mathcal{E})$  is the functor where

- *Action on Objects.* For each object  $(G, F) \in \text{Obj}(\text{Hom}_{\mathbf{Cats}_2}(\mathcal{D}, \mathcal{E}) \times \text{Hom}_{\mathbf{Cats}_2}(C, \mathcal{D}))$ , we have

$$\circ_{C, \mathcal{D}, \mathcal{E}}^{\mathbf{Cats}_2}(G, F) \stackrel{\text{def}}{=} G \circ F.$$

- *Action on Morphisms.* For each morphism  $(\beta, \alpha): (K, H) \Longrightarrow (G, F)$

of  $\text{Hom}_{\text{Cats}_2}(\mathcal{D}, \mathcal{E}) \times \text{Hom}_{\text{Cats}_2}(C, \mathcal{D})$ , we have

$$\circ_{C, \mathcal{D}, \mathcal{E}}^{\text{Cats}_2}(\beta, \alpha) \stackrel{\text{def}}{=} \beta \star \alpha,$$

where  $\beta \star \alpha$  is the horizontal composition of  $\alpha$  and  $\beta$  of [Definition II.9.5.I.I.](#)

**Proposition II.10.3.I.2.** Let  $C$  be a category.

1. *2-Categorical Co/Completeness.* The 2-category  $\text{Cats}_2$  is complete and co-complete as a 2-category, having all 2-categorical and bicategorical co/limits.

*Proof.* [Item I](#), *Co/Completeness*: Omitted. □

## II.10.4 The Category of Groupoids

**Definition II.10.4.I.I.** The **category of (small) groupoids** is the full subcategory  $\text{Grpd}$  of  $\text{Cats}$  spanned by the groupoids.

## II.10.5 The 2-Category of Groupoids

**Definition II.10.5.I.I.** The **2-category of (small) groupoids** is the full sub-2-category  $\text{Grpd}_2$  of  $\text{Cats}_2$  spanned by the groupoids.

# Appendices

## A Other Chapters

### Preliminaries

1. [Introduction](#)
2. [A Guide to the Literature](#)

### Sets

3. [Sets](#)

4. [Constructions With Sets](#)

5. [Monoidal Structures on the Category of Sets](#)

6. [Pointed Sets](#)

7. [Tensor Products of Pointed Sets](#)



**Relations**

- 8. Relations
- 9. Constructions With Relations
- 10. Conditions on Relations

**Categories**

- 11. Categories
- 12. Presheaves and the Yoneda Lemma

**Monoidal Categories**

- 13. Constructions With Monoidal Categories

**Bicategories**

- 14. Types of Morphisms in Bicategories

**Extra Part**

- 15. Notes

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