Constructions With Sets

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This chapter develops some material relating to constructions with sets with an eye towards its categorical and higher-categorical counterparts to be introduced later in this work. Of particular interest are perhaps the following:

- 1. Explicit descriptions of the major types of co/limits in Sets, including in particular explicit descriptions of pushouts and coequalisers (see Definitions 4.2.4.1.1 and 4.2.5.1.1 and Remarks 4.2.4.1.4 and 4.2.5.1.4).
- 2. A discussion of powersets as decategorifications of categories of presheaves, including in particular results such as:
 - (a) A discussion of the internal Hom of a powerset (Section 4.4.7).
 - (b) A 0-categorical version of the Yoneda lemma (Presheaves and the Yoneda Lemma, Theorem 12.1.5.1.1), which we term the *Yoneda lemma for sets* (Proposition 4.5.5.1.1).
 - (c) A characterisation of powersets as free cocompletions (Section 4.4.5), mimicking the corresponding statement for categories of presheaves (??).
 - (d) A characterisation of powersets as free completions (Section 4.4.6), mimicking the corresponding statement for categories of copresheaves (??).
 - (e) A (-1)-categorical version of un/straightening (Item 2 of Proposition 4.5.1.1.4 and Remark 4.5.1.1.6).
 - (f) A 0-categorical form of Isbell duality internal to powersets (Section 4.4.8).
- 3. A lengthy discussion of the adjoint triple

$$f_! \dashv f^{-1} \dashv f_* \colon \mathcal{P}(A) \xrightarrow{\rightleftharpoons} \mathcal{P}(B)$$

of functors (i.e. morphisms of posets) between $\mathcal{P}(A)$ and $\mathcal{P}(B)$ induced by a map of sets $f: A \to B$, including in particular:

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(a) How f^{-1} can be described as a precomposition while $f_!$ and f_* can be described as Kan extensions (Remarks 4.6.1.1.4, 4.6.2.1.2 and 4.6.3.1.4).

- (b) An extensive list of the properties of f_1 , f^{-1} , and f_* (Propositions 4.6.1.1.5, 4.6.1.1.7, 4.6.2.1.3, 4.6.2.1.5, 4.6.3.1.7 and 4.6.3.1.9).
- (c) How the functors $f_!$, f^{-1} , f_* , along with the functors

$$-_{1} \cap -_{2} \colon \mathcal{P}(X) \times \mathcal{P}(X) \to \mathcal{P}(X),$$
$$[-_{1}, -_{2}]_{X} \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

may be viewed as a six-functor formalism with the empty set \emptyset as the dualising object (Section 4.6.4).

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4.1 Limits of Sets

4.1.1 The Terminal Set

DEFINITION 4.1.1.1.1 ► THE TERMINAL SET

The **terminal set** is the terminal object of Sets as in Limits and Colimits, ??.

CONSTRUCTION 4.1.1.1.2 ► CONSTRUCTION OF THE TERMINAL SET

Concretely, the terminal set is the pair (pt, $\{!_A\}_{A \in Obj(Sets)}$) consisting of:

- 1. *The Limit*. The punctual set pt $\stackrel{\text{def}}{=} \{ \star \}$.
- 2. The Cone. The collection of maps

$$\{!_A : A \to \mathsf{pt}\}_{A \in \mathsf{Obj}(\mathsf{Sets})}$$

defined by

$$!_A(a) \stackrel{\text{def}}{=} \star$$

for each $a \in A$ and each $A \in Obj(Sets)$.

PROOF 4.1.1.1.3 ► PROOF OF CONSTRUCTION 4.1.1.1.2

We claim that pt is the terminal object of Sets. Indeed, suppose we have a diagram of the form

$$A$$
 p

in Sets. Then there exists a unique map $\phi: A \to \operatorname{pt}$ making the diagram

$$A - \frac{\phi}{\exists !} \rightarrow pt$$

commute, namely $!_A$.

4.1.2 Products of Families of Sets

Let $\{A_i\}_{i\in I}$ be a family of sets.

DEFINITION 4.1.2.1.1 ► THE PRODUCT OF A FAMILY OF SETS

The **product**¹ **of** $\{A_i\}_{i\in I}$ is the product of $\{A_i\}_{i\in I}$ in Sets as in Limits and Colimits, **??**.

¹Further Terminology: Also called the **Cartesian product of** $\{A_i\}_{i\in I}$.

CONSTRUCTION 4.1.2.1.2 ► CONSTRUCTION OF THE PRODUCT OF A FAMILY OF SETS

Concretely, the product of $\{A_i\}_{i\in I}$ is the pair $(\prod_{i\in I} A_i, \{pr_i\}_{i\in I})$ consisting of:

1. The Limit. The set $\prod_{i \in I} A_i$ defined by

$$\prod_{i \in I} A_i \stackrel{\text{def}}{=} \left\{ f \in \mathsf{Sets} \left(I, \bigcup_{i \in I} A_i \right) \middle| \begin{array}{l} \text{for each } i \in I, \text{ we} \\ \text{have } f(i) \in A_i \end{array} \right\}.$$

2. *The Cone*. The collection

$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

of maps given by

$$\operatorname{pr}_{i}(f) \stackrel{\operatorname{def}}{=} f(i)$$

for each $f \in \prod_{i \in I} A_i$ and each $i \in I$.

PROOF 4.1.2.1.3 ► PROOF OF CONSTRUCTION 4.1.2.1.2

We claim that $\prod_{i \in I} A_i$ is the categorical product of $\{A_i\}_{i \in I}$ in Sets. Indeed, suppose we have, for each $i \in I$, a diagram of the form



in Sets. Then there exists a unique map $\phi: P \to \prod_{i \in I} A_i$ making the diagram

$$\begin{array}{c|c}
P \\
\phi \mid \exists ! & p_i \\
\downarrow & \\
\prod_{i \in I} A_i & \xrightarrow{\operatorname{pr}_i} A_i
\end{array}$$

commute, being uniquely determined by the condition $\operatorname{pr}_i \circ \phi = p_i$ for each $i \in I$ via

$$\phi(x) = (p_i(x))_{i \in I}$$

for each $x \in P$.

REMARK 4.1.2.1.4 ► Unwinding Construction 4.1.2.1.2

Less formally, we may think of Cartesian products and projection maps as follows:

- 1. We think of $\prod_{i \in I} A_i$ as the set whose elements are *I*-indexed collections $(a_i)_{i \in I}$ with $a_i \in A_i$ for each $i \in I$.
- 2. We view the projection maps

$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

as being given by

$$\operatorname{pr}_i((a_j)_{j\in I})\stackrel{\mathrm{def}}{=} a_i$$

for each $(a_j)_{j\in I}\in\prod_{i\in I}A_i$ and each $i\in I$.

PROPOSITION 4.1.2.1.5 ► PROPERTIES OF PRODUCTS OF FAMILIES OF SETS

Let $\{A_i\}_{i\in I}$ be a family of sets.

1. Functoriality. The assignment $\{A_i\}_{i\in I}\mapsto \prod_{i\in I}A_i$ defines a functor

$$\prod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

• *Action on Objects.* For each $(A_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$, we have

$$\left[\prod_{i\in I}\right]((A_i)_{i\in I})\stackrel{\text{def}}{=}\prod_{i\in I}A_i$$

• Action on Morphisms. For each $(A_i)_{i \in I}, (B_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$, the action on Hom-sets

$$\left(\prod_{i\in I}\right)_{(A_i)_{i\in I},(B_i)_{i\in I}}: \operatorname{Nat}((A_i)_{i\in I},(B_i)_{i\in I}) \to \operatorname{Sets}\left(\prod_{i\in I}A_i,\prod_{i\in I}B_i\right)$$

of $\prod_{i\in I}$ at $((A_i)_{i\in I},(B_i)_{i\in I})$ is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in Nat $((A_i)_{i\in I}, (B_i)_{i\in I})$ to the map of sets

$$\prod_{i \in I} f_i \colon \prod_{i \in I} A_i \to \prod_{i \in I} B_i$$

defined by

$$\left[\prod_{i\in I} f_i\right] ((a_i)_{i\in I}) \stackrel{\text{def}}{=} (f_i(a_i))_{i\in I}$$

for each $(a_i)_{i \in I} \in \prod_{i \in I} A_i$.

PROOF 4.1.2.1.6 ► PROOF OF PROPOSITION 4.1.2.1.5

Item 1: Functoriality

This follows from Limits and Colimits, ?? of ??.

4.1.3 Binary Products of Sets

Let *A* and *B* be sets.

DEFINITION 4.1.3.1.1 ► BINARY PRODUCTS OF SETS

The **product of** A and B^1 is the product of A and B in Sets as in Limits and Colimits, ??.

¹Further Terminology: Also called the **Cartesian product of** A **and** B.

CONSTRUCTION 4.1.3.1.2 ► CONSTRUCTION OF BINARY PRODUCTS OF SETS

Concretely, the product of *A* and *B* is the pair $(A \times B, \{pr_1, pr_2\})$ consisting of:

1. The Limit. The set $A \times B$ defined by

$$A \times B \stackrel{\text{def}}{=} \prod_{z \in \{A, B\}} z$$

$$\stackrel{\text{def}}{=} \{ f \in \text{Sets}(\{0, 1\}, A \cup B) \mid \text{we have } f(0) \in A \text{ and } f(1) \in B \}$$

$$\cong \{ \{ \{a\}, \{a, b\} \} \in \mathcal{P}(\mathcal{P}(A \cup B)) \mid \text{we have } a \in A \text{ and } b \in B \}$$

$$\cong \begin{cases} \text{ordered pairs } (a, b) \text{ with } \\ a \in A \text{ and } b \in B \end{cases}.$$

2. The Cone. The maps

$$\operatorname{pr}_1: A \times B \to A,$$

 $\operatorname{pr}_2: A \times B \to B$

defined by

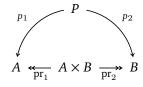
$$\operatorname{pr}_{1}(a,b) \stackrel{\text{def}}{=} a,$$

 $\operatorname{pr}_{2}(a,b) \stackrel{\text{def}}{=} b$

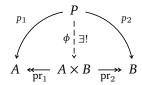
for each $(a, b) \in A \times B$.

PROOF 4.1.3.1.3 ► PROOF OF CONSTRUCTION 4.1.3.1.2

We claim that $A \times B$ is the categorical product of A and B in the category of sets. Indeed, suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: P \to A \times B$ making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$

 $\operatorname{pr}_2 \circ \phi = p_2$

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each $x \in P$.

PROPOSITION 4.1.3.1.4 ► PROPERTIES OF PRODUCTS OF SETS

Let A, B, C, and X be sets.

1. Functoriality. The assignments $A, B, (A, B) \mapsto A \times B$ define functors

$$A \times -:$$
 Sets \rightarrow Sets,
 $- \times B:$ Sets \rightarrow Sets,
 $-_1 \times -_2:$ Sets \times Sets \rightarrow Sets,

where -1×-2 is the functor where

• *Action on Objects.* For each $(A, B) \in \mathsf{Obj}(\mathsf{Sets} \times \mathsf{Sets})$, we have

$$[-1 \times -2](A, B) \stackrel{\text{def}}{=} A \times B.$$

• *Action on Morphisms.* For each $(A, B), (X, Y) \in Obj(Sets)$, the action on Hom-sets

$$\times_{(A,B),(X,Y)}$$
: $\mathsf{Sets}(A,X) \times \mathsf{Sets}(B,Y) \to \mathsf{Sets}(A \times B, X \times Y)$

of \times at ((A, B), (X, Y)) is defined by sending (f, g) to the function

$$f \times g \colon A \times B \to X \times Y$$

defined by

$$[f \times g](a,b) \stackrel{\text{def}}{=} (f(a),g(b))$$

for each $(a, b) \in A \times B$.

and where $A \times -$ and $- \times B$ are the partial functors of $-_1 \times -_2$ at $A, B \in \text{Obj}(\mathsf{Sets})$.

2. Adjointness I. We have adjunctions

$$(A \times - \dashv \operatorname{Sets}(A, -))$$
: Sets $\xrightarrow{A \times -}$ Sets, \subseteq Se

witnessed by bijections

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$

 $Sets(A \times B, C) \cong Sets(B, Sets(A, C)),$

natural in $A, B, C \in Obj(Sets)$.

3. Adjointness II. We have an adjunction

$$(\Delta_{\mathsf{Sets}} \dashv -_1 \times -_2)$$
: Sets $\stackrel{\Delta_{\mathsf{Sets}}}{=}$ Sets \times Sets,

witnessed by a bijection

$$Hom_{Sets \times Sets}((A, A), (B, C)) \cong Sets(A, B \times C),$$

natural in $A \in \text{Obj}(\mathsf{Sets})$ and in $(B, C) \in \text{Obj}(\mathsf{Sets} \times \mathsf{Sets})$.

4. Associativity. We have an isomorphism of sets

$$\alpha_{A.B.C}^{\mathsf{Sets}} : (A \times B) \times C \xrightarrow{\sim} A \times (B \times C),$$

natural in $A, B, C \in Obj(Sets)$.

5. Unitality. We have isomorphisms of sets

$$\lambda_A^{\mathsf{Sets}} \colon \mathsf{pt} \times A \xrightarrow{\sim} A,$$

 $\rho_A^{\mathsf{Sets}} \colon A \times \mathsf{pt} \xrightarrow{\sim} A,$

natural in $A \in Obj(Sets)$.

6. Commutativity. We have an isomorphism of sets

$$\sigma_{A,B}^{\mathsf{Sets}} \colon A \times B \xrightarrow{\sim} B \times A,$$

natural in $A, B \in Obj(Sets)$.

7. Distributivity Over Coproducts. We have isomorphisms of sets

$$\delta_{\ell}^{\mathsf{Sets}} : A \times (B \coprod C) \xrightarrow{\sim} (A \times B) \coprod (A \times C),$$

 $\delta_{r}^{\mathsf{Sets}} : (A \coprod B) \times C \xrightarrow{\sim} (A \times C) \coprod (B \times C),$

natural in $A, B, C \in Obj(Sets)$.

8. Annihilation With the Empty Set. We have isomorphisms of sets

$$\zeta_{\ell}^{\mathsf{Sets}} : \emptyset \times A \xrightarrow{\sim} \emptyset,$$

$$\zeta_{r}^{\mathsf{Sets}} : A \times \emptyset \xrightarrow{\sim} \emptyset,$$

natural in $A \in Obj(Sets)$.

9. *Distributivity Over Unions*. Let *X* be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \cup W) = (U \times V) \cup (U \times W),$$

$$(U \cup V) \times W = (U \times W) \cup (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

10. Distributivity Over Intersections. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \cap W) = (U \times V) \cap (U \times W),$$

$$(U \cap V) \times W = (U \times W) \cap (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

11. *Distributivity Over Differences*. Let *X* be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \setminus W) = (U \times V) \setminus (U \times W),$$

$$(U \setminus V) \times W = (U \times W) \setminus (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

12. Distributivity Over Symmetric Differences. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \triangle W) = (U \times V) \triangle (U \times W),$$

$$(U \triangle V) \times W = (U \times W) \triangle (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

13. Middle-Four Exchange with Respect to Intersections. The diagram

$$(\mathcal{P}(X) \times \mathcal{P}(X)) \times (\mathcal{P}(X) \times \mathcal{P}(X)) \xrightarrow{\cap \times \cap} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow^{\mathcal{P}_{X,X}^{\times}} \times \mathcal{P}_{X,X}^{\times}} \qquad \qquad \downarrow^{\mathcal{P}_{X,X}^{\times}}$$

$$\mathcal{P}(X \times X) \times \mathcal{P}(X \times X) \xrightarrow{\cap} \mathcal{P}(X \times X)$$

commutes, i.e. we have

$$(U \times V) \cap (W \times T) = (U \cap V) \times (W \cap T).$$

for each $U, V, W, T \in \mathcal{P}(X)$.

- 14. *Symmetric Monoidality*. The 8-tuple (Sets, \times , pt, Sets(-1, -2), α^{Sets} , λ^{Sets} , ρ^{Sets} , σ^{Sets}) is a closed symmetric monoidal category.
- 15. Symmetric Bimonoidality. The 18-tuple

$$\left(\mathsf{Sets}, \coprod, \times, \emptyset, \mathsf{pt}, \mathsf{Sets}(-_1, -_2), \alpha^{\mathsf{Sets}}, \lambda^{\mathsf{Sets}}, \rho^{\mathsf{Sets}}, \sigma^{\mathsf{Sets}}, \alpha^{\mathsf{Sets}}, \alpha^{\mathsf{Set$$

is a symmetric closed bimonoidal category, where $\alpha^{Sets, \coprod}$, $\lambda^{Sets, \coprod}$, $\rho^{Sets, \coprod}$, and $\sigma^{Sets, \coprod}$ are the natural transformations from Items 3 to 5 of Proposition 4.2.3.1.4.

PROOF 4.1.3.1.5 ► PROOF OF PROPOSITION 4.1.3.1.4

Item 1: Functoriality

This follows from Limits and Colimits, ?? of ??.

Item 2: Adjointness

We prove only that there's an adjunction $- \times B \dashv \mathsf{Sets}(B, -)$, witnessed by a bijection

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$

natural in $B, C \in \text{Obj}(\mathsf{Sets})$, as the proof of the existence of the adjunction $A \times - \exists \mathsf{Sets}(A, -)$ follows almost exactly in the same way.

• Map I. We define a map

$$\Phi_{B,C} : \mathsf{Sets}(A \times B, C) \to \mathsf{Sets}(A, \mathsf{Sets}(B, C)),$$

by sending a function

$$\xi: A \times B \to C$$

to the function

$$\xi^{\dagger} : A \longrightarrow \mathsf{Sets}(B, C),$$

 $a \mapsto (\xi_a^{\dagger} : B \to C),$

where we define

$$\xi_a^{\dagger}(b) \stackrel{\text{def}}{=} \xi(a,b)$$

for each $b \in B$. In terms of the $[a \mapsto f(a)]$ notation of Sets, Notation 3.1.1.1.2, we have

$$\xi^{\dagger} \stackrel{\text{def}}{=} \llbracket a \mapsto \llbracket b \mapsto \xi(a,b) \rrbracket \rrbracket.$$

• Map II. We define a map

$$\Psi_{B,C}$$
: Sets(A , Sets(B , C)), \rightarrow Sets($A \times B$, C)

given by sending a function

$$\xi: A \longrightarrow \mathsf{Sets}(B, C),$$

 $a \mapsto (\xi_a: B \to C),$

to the function

$$\xi^{\dagger}: A \times B \to C$$

defined by

$$\xi^{\dagger}(a,b) \stackrel{\text{def}}{=} \operatorname{ev}_b(\operatorname{ev}_a(\xi))$$

$$\stackrel{\text{def}}{=} \operatorname{ev}_b(\xi_a)$$

$$\stackrel{\text{def}}{=} \xi_a(b)$$

for each $(a, b) \in A \times B$.

• Invertibility I. We claim that

$$\Psi_{A,B} \circ \Phi_{A,B} = \mathrm{id}_{\mathsf{Sets}(A \times B,C)}$$
.

Indeed, given a function $\xi: A \times B \to C$, we have

$$\begin{split} & [\Psi_{A,B} \circ \Phi_{A,B}](\xi) = \Psi_{A,B}(\Phi_{A,B}(\xi)) \\ & = \Psi_{A,B}(\Phi_{A,B}([\![(a,b) \mapsto \xi(a,b)]\!])) \\ & = \Psi_{A,B}([\![a \mapsto [\![b \mapsto \xi(a,b)]\!]]\!]) \\ & = \Psi_{A,B}([\![a' \mapsto [\![b' \mapsto \xi(a',b')]\!]]\!]) \\ & = [\![(a,b) \mapsto \operatorname{ev}_b(\operatorname{ev}_a([\![a' \mapsto [\![b' \mapsto \xi(a',b')]\!]]\!]))]] \\ & = [\![(a,b) \mapsto \operatorname{ev}_b([\![b' \mapsto \xi(a,b')]\!])]\!] \\ & = [\![(a,b) \mapsto \xi(a,b)]\!] \\ & = \xi. \end{split}$$

• Invertibility II. We claim that

$$\Phi_{A,B} \circ \Psi_{A,B} = \mathrm{id}_{\mathsf{Sets}(A,\mathsf{Sets}(B,C))}$$
.

Indeed, given a function

$$\xi: A \longrightarrow \mathsf{Sets}(B, C),$$

 $a \mapsto (\xi_a: B \to C),$

we have

$$[\Phi_{A,B} \circ \Psi_{A,B}](\xi) \stackrel{\text{def}}{=} \Phi_{A,B}(\Psi_{A,B}(\xi))$$

$$\begin{array}{l} \overset{\mathrm{def}}{=} \Phi_{A,B}(\llbracket (a,b) \mapsto \xi_a(b) \rrbracket) \\ \overset{\mathrm{def}}{=} \Phi_{A,B}(\llbracket (a',b') \mapsto \xi_{a'}(b') \rrbracket) \\ \overset{\mathrm{def}}{=} \llbracket a \mapsto \llbracket b \mapsto \mathrm{ev}_{(a,b)}(\llbracket (a',b') \mapsto \xi_{a'}(b') \rrbracket) \rrbracket \rrbracket \\ \overset{\mathrm{def}}{=} \llbracket a \mapsto \llbracket b \mapsto \xi_a(b) \rrbracket \rrbracket \\ \overset{\mathrm{def}}{=} \llbracket a \mapsto \xi_a \rrbracket \\ \overset{\mathrm{def}}{=} \xi. \end{array}$$

• *Naturality for* Φ , *Part I.* We need to show that, given a function $g: B \to B'$, the diagram

$$\begin{split} \mathsf{Sets}(A \times B', C) & \xrightarrow{\Phi_{B', C}} \mathsf{Sets}(A, \mathsf{Sets}(B', C)), \\ & \mathsf{id}_A \times g^* \bigg| & & \bigg| & (g^*)_! \\ & \mathsf{Sets}(A \times B, C) & \xrightarrow{\Phi_{B, C}} \mathsf{Sets}(A, \mathsf{Sets}(B, C)) \end{split}$$

commutes. Indeed, given a function

$$\xi: A \times B' \to C$$

we have

$$\begin{split} [\Phi_{B,C} \circ (\mathrm{id}_A \times g^*)](\xi) &= \Phi_{B,C}([\mathrm{id}_A \times g^*](\xi)) \\ &= \Phi_{B,C}(\xi(-_1, g(-_2))) \\ &= [\xi(-_1, g(-_2))]^{\dagger} \\ &= \xi_{-_1}^{\dagger}(g(-_2)) \\ &= (g^*)_!(\xi^{\dagger}) \\ &= (g^*)_!(\Phi_{B',C}(\xi)) \\ &= [(g^*)_! \circ \Phi_{B',C}](\xi). \end{split}$$

Alternatively, using the $[a \mapsto f(a)]$ notation of Sets, Notation 3.1.1.2, we have

$$\begin{split} [\Phi_{B,C} \circ (\mathrm{id}_A \times g^*)](\xi) &= \Phi_{B,C}([\mathrm{id}_A \times g^*](\xi)) \\ &= \Phi_{B,C}([\mathrm{id}_A \times g^*]([\![(a,b') \mapsto \xi(a,b')]\!])) \end{split}$$

$$\begin{split} &= \Phi_{B,C}(\llbracket (a,b) \mapsto \xi(a,g(b)) \rrbracket) \\ &= \llbracket a \mapsto \llbracket b \mapsto \xi(a,g(b)) \rrbracket \rrbracket \\ &= \llbracket a \mapsto g^*(\llbracket b' \mapsto \xi(a,b') \rrbracket) \rrbracket \\ &= [g^*)_!(\llbracket a \mapsto \llbracket b' \mapsto \xi(a,b') \rrbracket] \rrbracket) \\ &= (g^*)_!(\Phi_{B',C}(\llbracket (a,b') \mapsto \xi(a,b') \rrbracket])) \\ &= (g^*)_!(\Phi_{B',C}(\xi)) \\ &= [(g^*)_! \circ \Phi_{B',C}](\xi). \end{split}$$

• *Naturality for* Φ , *Part II.* We need to show that, given a function $h: C \to C'$, the diagram

$$\begin{split} \mathsf{Sets}(A \times B, C) & \xrightarrow{\Phi_{B,C}} & \mathsf{Sets}(A, \mathsf{Sets}(B, C)), \\ h_! & & \downarrow^{(h_!)_!} \\ \mathsf{Sets}(A \times B, C') & \xrightarrow{\Phi_{B,C'}} & \mathsf{Sets}(A, \mathsf{Sets}(B, C')) \end{split}$$

commutes. Indeed, given a function

$$\xi: A \times B \to C$$

we have

$$\begin{split} [\Phi_{B,C} \circ h_!](\xi) &= \Phi_{B,C}(h_!(\xi)) \\ &= \Phi_{B,C}(h_!([(a,b) \mapsto \xi(a,b)])) \\ &= \Phi_{B,C}([(a,b) \mapsto h(\xi(a,b))]) \\ &= [(a \mapsto [(b \mapsto h(\xi(a,b))]]) \\ &= [(a \mapsto h_!([(b \mapsto \xi(a,b)]]))) \\ &= (h_!)_!([(a \mapsto [(b \mapsto \xi(a,b)]]))) \\ &= (h_!)_!(\Phi_{B,C}([((a,b) \mapsto \xi(a,b)]))) \\ &= (h_!)_!(\Phi_{B,C}(\xi)) \\ &= [(h_!)_! \circ \Phi_{B,C}](\xi). \end{split}$$

• Naturality for Ψ . Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from

Categories, Item 2 of Proposition 11.9.7.1.2 that Ψ is also natural in each argument.

This finishes the proof.

Item 3: Adjointness II

This follows from the universal property of the product.

Item 4: Associativity

This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.1.4.1.1.

Item 5: Unitality

This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.1.5.1.1 and 5.1.6.1.1.

Item 6: Commutativity

This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.1.7.1.1.

Item 7: Distributivity Over Coproducts

This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.3.1.1.1 and 5.3.2.1.1.

Item 8: Annihilation With the Empty Set

This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.3.3.1.1 and 5.3.4.1.1.

Item 9: Distributivity Over Unions

See [Pro25c].

Item 10: Distributivity Over Intersections

See [Pro25d, Corollary 1].

Item 11: Distributivity Over Differences

See [Pro25a].

Item 12: Distributivity Over Symmetric Differences

See [Pro25b].

Item 13: Middle-Four Exchange With Respect to Intersections

See [Pro25d, Corollary 1].

Item 14: Symmetric Monoidality

This is a repetition of Monoidal Structures on the Category of Sets, Proposition 5.1.9.1.1, and is proved there.

Item 15: Symmetric Bimonoidality

This is a repetition of Monoidal Structures on the Category of Sets, Proposition 5.3.5.1.1, and is proved there.

REMARK 4.1.3.1.6 \blacktriangleright The Cartesian Product of Sets as an $(\mathbb{E}_k, \mathbb{E}_\ell)$ -Tensor Product

As shown in Item 1 of Proposition 4.1.3.1.4, the Cartesian product of sets defines a functor

$$-1 \times -2$$
: Sets \times Sets \rightarrow Sets.

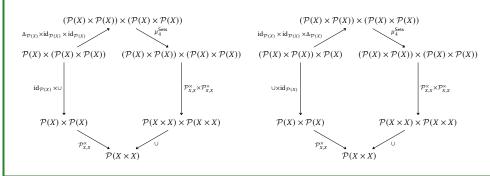
This functor is the $(k, \ell) = (-1, -1)$ case of a family of functors

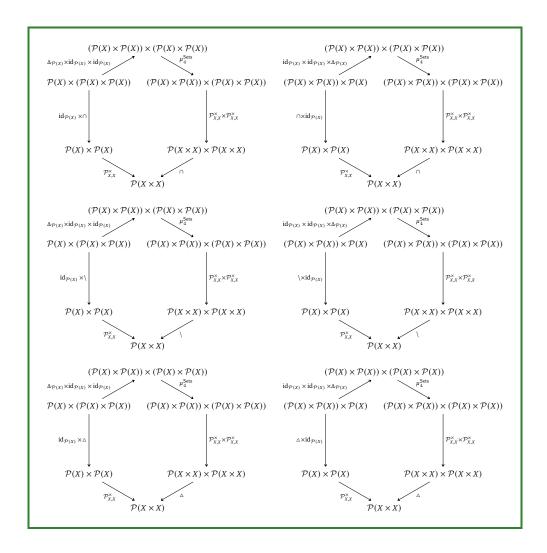
$$\otimes_{k,\ell} : \mathsf{Mon}_{\mathbb{E}_k}(\mathsf{Sets}) \times \mathsf{Mon}_{\mathbb{E}_\ell}(\mathsf{Sets}) \to \mathsf{Mon}_{\mathbb{E}_{k+\ell}}(\mathsf{Sets})$$

of tensor products of \mathbb{E}_k -monoid objects on Sets with \mathbb{E}_ℓ -monoid objects on Sets; see ??.

REMARK 4.1.3.1.7 ► DIAGRAMS FOR ITEMS 9 TO 12 OF PROPOSITION 4.1.3.1.4

We may state the equalities in Items 9 to 12 of Proposition 4.1.3.1.4 as the commutativity of the following diagrams:





4.1.4 Pullbacks

Let A, B, and C be sets and let $f: A \to C$ and $g: B \to C$ be functions.

DEFINITION 4.1.4.1.1 ► PULLBACKS OF SETS

The **pullback of** A **and** B **over** C **along** f **and** g^1 is the pullback of A and B over C along f and g in Sets as in Limits and Colimits, ??.

¹Further Terminology: Also called the **fibre product of** A **and** B **over** C **along** f **and** g.

CONSTRUCTION 4.1.4.1.2 ► CONSTRUCTION OF PULLBACKS OF SETS

Concretely, the pullback of *A* and *B* over *C* along *f* and *g* is the pair $(A \times_C B, \{pr_1, pr_2\})$ consisting of:

1. *The Limit*. The set $A \times_C B$ defined by

$$A \times_C B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid f(a) = g(b)\}.$$

2. The Cone. The maps¹

$$\operatorname{pr}_1: A \times_C B \to A,$$

 $\operatorname{pr}_2: A \times_C B \to B$

defined by

$$\operatorname{pr}_{1}(a,b) \stackrel{\text{def}}{=} a,$$

 $\operatorname{pr}_{2}(a,b) \stackrel{\text{def}}{=} b$

for each $(a, b) \in A \times_C B$.

PROOF 4.1.4.1.3 ► PROOF OF CONSTRUCTION 4.1.4.1.2

We claim that $A \times_C B$ is the categorical pullback of A and B over C with respect to (f,g) in Sets. First we need to check that the relevant pullback diagram commutes, i.e. that we have

$$f \circ \operatorname{pr}_1 = g \circ \operatorname{pr}_2, \qquad A \times_C B \xrightarrow{\operatorname{pr}_2} B$$

$$pr_1 \downarrow \qquad \qquad \downarrow g$$

$$A \xrightarrow{f} C.$$

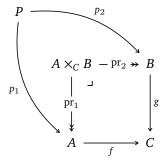
Indeed, given $(a, b) \in A \times_C B$, we have

$$[f \circ \operatorname{pr}_1](a, b) = f(\operatorname{pr}_1(a, b))$$
$$= f(a)$$
$$= g(b)$$

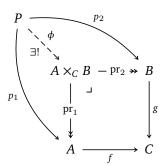
¹Further Notation: Also written $\operatorname{pr}_1^{A \times_C B}$ and $\operatorname{pr}_2^{A \times_C B}$.

$$= g(\operatorname{pr}_{2}(a, b))$$
$$= [g \circ \operatorname{pr}_{2}](a, b),$$

where f(a) = g(b) since $(a, b) \in A \times_C B$. Next, we prove that $A \times_C B$ satisfies the universal property of the pullback. Suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: P \to A \times_C B$ making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$

 $\operatorname{pr}_2 \circ \phi = p_2$

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each $x \in P$, where we note that $(p_1(x), p_2(x)) \in A \times B$ indeed lies in $A \times_C B$ by the condition

$$f \circ p_1 = g \circ p_2$$
,

which gives

$$f(p_1(x)) = g(p_2(x))$$

for each $x \in P$, so that $(p_1(x), p_2(x)) \in A \times_C B$.

REMARK 4.1.4.1.4 ▶ PULLBACKS OF SETS DEPEND ON THE MAPS

It is common practice to write $A \times_C B$ for the pullback of A and B over C along f and g, omitting the maps f and g from the notation and instead leaving them implicit, to be understood from the context.

However, the set $A \times_C B$ depends very much on the maps f and g, and sometimes it is necessary or useful to note this dependence explicitly. In such situations, we will write $A \times_{f,C,g} B$ or $A \times_C^{f,g} B$ for $A \times_C B$.

EXAMPLE 4.1.4.1.5 ► **EXAMPLES OF PULLBACKS OF SETS**

Here are some examples of pullbacks of sets.

1. *Unions via Intersections*. Let *X* be a set. We have

$$A \cap B \cong A \times_{A \cup B} B, \qquad A \cap B \xrightarrow{B} B$$

$$\downarrow \qquad \downarrow \qquad \downarrow \iota_{B}$$

$$A \hookrightarrow \iota_{A} \rightarrow A \cup B$$

for each $A, B \in \mathcal{P}(X)$.

PROOF 4.1.4.1.6 ► PROOF OF EXAMPLE 4.1.4.1.5

Item 1: Unions via Intersections

Indeed, we have

$$A \times_{A \cup B} B \cong \{(x, y) \in A \times B \mid x = y\}$$

\approx A \cap B.

This finishes the proof.

PROPOSITION 4.1.4.1.7 ► PROPERTIES OF PULLBACKS OF SETS

Let A, B, C, and X be sets.

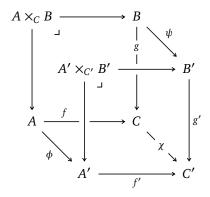
1. Functoriality. The assignment $(A, B, C, f, g) \mapsto A \times_{f,C,g} B$ defines a functor

$$-_1 \times_{-_3} -_1 \colon \mathsf{Fun}(\mathcal{P},\mathsf{Sets}) \to \mathsf{Sets},$$

where \mathcal{P} is the category that looks like this:



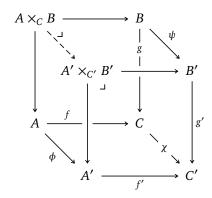
In particular, the action on morphisms of $-_1\times_{-_3}-_1$ is given by sending a morphism



in Fun(\mathcal{P} , Sets) to the map $\xi \colon A \times_C B \xrightarrow{\exists !} A' \times_{C'} B'$ given by

$$\xi(a,b) \stackrel{\text{def}}{=} (\phi(a), \psi(b))$$

for each $(a, b) \in A \times_C B$, which is the unique map making the diagram



commute.

2. Adjointness I. We have adjunctions

$$(A \times_X - \dashv \mathbf{Sets}_{/X}(A, -))$$
: $Sets_{/X} \xrightarrow{A \times_X -} Sets_{/X}$, $Sets_{/X}(A, -)$
 $(- \times_X B \dashv \mathbf{Sets}_{/X}(B, -))$: $Sets_{/X} \xrightarrow{A \times_X -} Sets_{/X}$, $Sets_{/X}(B, -)$

witnessed by bijections

$$\mathsf{Sets}_{/X}(A \times_X B, C) \cong \mathsf{Sets}_{/X}(A, \mathbf{Sets}_{/X}(B, C)),$$

 $\mathsf{Sets}_{/X}(A \times_X B, C) \cong \mathsf{Sets}_{/X}(B, \mathbf{Sets}_{/X}(A, C)),$

natural in (A, ϕ_A) , (B, ϕ_B) , $(C, \phi_C) \in \text{Obj}(\mathsf{Sets}_{/X})$, where $\mathsf{Sets}_{/X}(A, B)$ is the object of $\mathsf{Sets}_{/X}$ consisting of (see Fibred Sets, ??):

• *The Set*. The set $\mathbf{Sets}_{/X}(A, B)$ defined by

$$\mathbf{Sets}_{/X}(A,B) \stackrel{\text{def}}{=} \coprod_{x \in X} \mathsf{Sets}(\phi_A^{-1}(x), \phi_Y^{-1}(x))$$

• *The Map to X*. The map

$$\phi_{\mathsf{Sets}_{/X}(A,B)} : \mathsf{Sets}_{/X}(A,B) \to X$$

defined by

$$\phi_{\mathsf{Sets}_{/X}(A,B)}(x,f) \stackrel{\mathrm{def}}{=} x$$

for each $(x, f) \in \mathbf{Sets}_{/X}(A, B)$.

3. Adjointness II. We have an adjunction

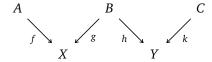
$$\left(\Delta_{\mathsf{Sets}_{/X}} \dashv -_1 \times -_2\right)$$
: $\mathsf{Sets}_{/X} \underbrace{\perp}_{-_1 \times -_2} \mathsf{Sets}_{/X} \times \mathsf{Sets}_{/X}$

witnessed by a bijection

$$\operatorname{Hom}_{\operatorname{\mathsf{Sets}}/_X\times\operatorname{\mathsf{Sets}}/_X}((A,A),(B,C))\cong\operatorname{\mathsf{Sets}}/_X(A,B\times_XC),$$

natural in $A \in \text{Obj}(\mathsf{Sets}_{/X})$ and in $(B, C) \in \text{Obj}(\mathsf{Sets}_{/X} \times \mathsf{Sets}_{/X})$.

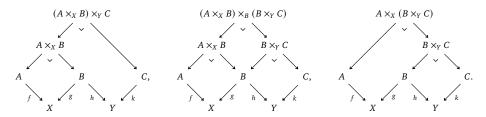
4. Associativity. Given a diagram



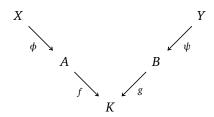
in Sets, we have isomorphisms of sets

$$(A \times_X B) \times_Y C \cong (A \times_X B) \times_B (B \times_Y C) \cong A \times_X (B \times_Y C),$$

where these pullbacks are built as in the diagrams



5. Interaction With Composition. Given a diagram



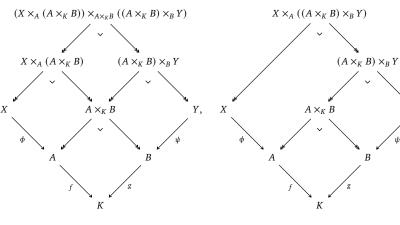
in Sets, we have isomorphisms of sets

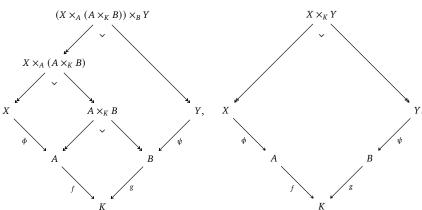
$$\begin{split} X \times_K^{f \circ \phi, g \circ \psi} Y &\cong (X \times_A^{\phi, q_1} (A \times_K^{f, g} B)) \times_{A \times_K^{f, g} B}^{p_2, p_1} ((A \times_K^{f, g} B) \times_B^{q_2, \psi} Y) \\ &\cong X \times_A^{\phi, p} ((A \times_K^{f, g} B) \times_B^{q_2, \psi} Y) \\ &\cong (X \times_A^{\phi, q_1} (A \times_K^{f, g} B)) \times_B^{q, \psi} Y \end{split}$$

where

$$\begin{aligned} q_1 &= \operatorname{pr}_1^{A \times_K^{f,g} B}, & q_2 &= \operatorname{pr}_2^{A \times_K^{f,g} B}, \\ p_1 &= \operatorname{pr}_1^{(A \times_K^{f,g} B) \times_Y^{q_2,\psi}}, & \underset{p_2}{X \times_K^{\phi,q_1}} (A \times_K^{f,g} B) \\ p &= q_1 \circ \operatorname{pr}_1^{(A \times_K^{f,g} B) \times_B^{q_2,\psi} Y}, & q &= q_2 \circ \operatorname{pr}_2^{X \times_A^{\phi,q_1}} (A \times_K^{f,g} B) \end{aligned},$$

and where these pullbacks are built as in the following diagrams:





6. Unitality. We have isomorphisms of sets

natural in $(A, f) \in \text{Obj}(\mathsf{Sets}_{/X})$.

7. Commutativity. We have an isomorphism of sets

natural in $(A, f), (B, g) \in \text{Obj}(\mathsf{Sets}_{/X})$.

8. *Distributivity Over Coproducts*. Let A, B, and C be sets and let $\phi_A \colon A \to X$, $\phi_B \colon B \to X$, and $\phi_C \colon C \to X$ be morphisms of sets. We have isomorphisms of sets

$$\delta_{\ell}^{\mathsf{Sets}_{/X}} : A \times_{X} (B \coprod C) \xrightarrow{\sim} (A \times_{X} B) \coprod (A \times_{X} C),$$

$$\delta_{r}^{\mathsf{Sets}_{/X}} : (A \coprod B) \times_{X} C \xrightarrow{\sim} (A \times_{X} C) \coprod (B \times_{X} C),$$

as in the diagrams

natural in $A, B, C \in \text{Obj}(\mathsf{Sets}_{/X})$.

9. Annihilation With the Empty Set. We have isomorphisms of sets

natural in $(A, f) \in \text{Obj}(\mathsf{Sets}_{/X})$.

10. Interaction With Products. We have an isomorphism of sets

$$A \times_{\text{pt}} B \cong A \times B,$$

$$A \times_{\text{pt}} B \cong A \times B,$$

$$A \xrightarrow{!_A} \text{pt.}$$

11. *Symmetric Monoidality*. The 8-tuple (Sets_{/X}, \times_X , X, **Sets**_{/X}, $\alpha^{\text{Sets}_{/X}}$, $\alpha^{\text{Sets}_{/X}}$, $\sigma^{\text{Sets}_{/X}}$) is a symmetric closed monoidal category.

PROOF 4.1.4.1.8 ► PROOF OF PROPOSITION 4.1.4.1.7

Item 1: Functoriality

This is a special case of functoriality of co/limits, Limits and Colimits, ?? of ??, with the explicit expression for ξ following from the commutativity of the cube pullback diagram.

Item 2: Adjointness I

This is a repetition of Fibred Sets, ?? of ??, and is proved there.

Item 3: Adjointness II

This follows from the universal property of the product (pullbacks are products in $\mathsf{Sets}_{/X}$).

Item 4: Associativity

We have

$$(A \times_X B) \times_Y C \cong \{((a,b),c) \in (A \times_X B) \times C \mid h(b) = k(c)\}$$

$$\cong \{((a,b),c) \in (A \times B) \times C \mid f(a) = g(b) \text{ and } h(b) = k(c)\}$$

$$\cong \{(a,(b,c)) \in A \times (B \times C) \mid f(a) = g(b) \text{ and } h(b) = k(c)\}$$

$$\cong \{(a,(b,c)) \in A \times (B \times_Y C) \mid f(a) = g(b)\}$$

$$\cong A \times_X (B \times_Y C)$$

and

$$(A \times_X B) \times_B (B \times_Y C) \cong \{((a,b),(b',c)) \in (A \times_X B) \times (B \times_Y C) \mid b = b'\}$$

$$\cong \left\{ ((a,b),(b',c)) \in (A \times B) \times (B \times C) \mid f(a) = g(b), b = b', \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,(b,(b',c))) \in A \times (B \times (B \times C)) \mid f(a) = g(b), b = b', \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,((b,b'),c)) \in A \times ((B \times B) \times C) \mid f(a) = g(b), b = b', \text{and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,((b,b'),c)) \in A \times ((B \times_B B) \times C) \mid f(a) = g(b) \text{ and } h(b') = k(c) \right\}$$

$$\cong \left\{ (a,(b,c)) \in A \times (B \times C) \mid f(a) = g(b) \text{ and } h(b) = k(c) \right\}$$

$$\cong A \times_X (B \times_Y C),$$

where we have used Item 6 for the isomorphism $B \times_B B \cong B$.

Item 5: Interaction With Composition

By Item 4, it suffices to construct only the isomorphism

$$X\times_K^{f\circ\phi,g\circ\psi}Y\cong (X\times_A^{\phi,q_1}(A\times_K^{f,g}B))\times_{A\times_K^{f,g}B}^{p_2,p_1}((A\times_K^{f,g}B)\times_B^{q_2,\psi}Y).$$

We have

$$(X \times_{A}^{\phi,q_{1}} (A \times_{K}^{f,g} B)) \stackrel{\text{def}}{=} \left\{ (x, (a, b)) \in X \times (A \times_{K}^{f,g} B) \mid \phi(x) = q_{1}(a, b) \right\}$$

$$\stackrel{\text{def}}{=} \left\{ (x, (a, b)) \in X \times (A \times_{K}^{f,g} B) \mid \phi(x) = a \right\}$$

$$\cong \left\{ (x, (a, b)) \in X \times (A \times B) \mid \phi(x) = a \text{ and } f(a) = g(b) \right\},$$

$$((A \times_{K}^{f,g} B) \times_{B}^{q_{2}, \psi} Y) \stackrel{\text{def}}{=} \left\{ ((a, b), y) \in (A \times_{K}^{f,g} B) \times Y \mid q_{2}(a, b) = \psi(y) \right\}$$

$$\stackrel{\text{def}}{=} \left\{ ((a, b), y) \in (A \times_{K}^{f,g} B) \times Y \mid b = \psi(y) \text{ and } f(a) = g(b) \right\},$$

$$\cong \left\{ ((a, b), y) \in (A \times B) \times Y \mid b = \psi(y) \text{ and } f(a) = g(b) \right\},$$

so writing

$$S = (X \times_A^{\phi,q_1} (A \times_K^{f,g} B))$$

$$S' = ((A \times_K^{f,g} B) \times_B^{q_2,\psi} Y),$$

we have

$$\begin{split} S \times_{A \times_{K}^{f,S} B}^{P2,P1} S' &\stackrel{\text{def}}{=} \{ ((x,(a,b)), ((a',b'),y)) \in S \times S' \mid p_{1}(x,(a,b)) = p_{2}((a',b'),y) \} \\ &\stackrel{\text{def}}{=} \{ ((x,(a,b)), ((a',b'),y)) \in S \times S' \mid (a,b) = (a',b') \} \\ &\cong \{ ((x,a,b,y)) \in X \times A \times B \times Y \mid \phi(x) = a, \psi(y) = b, \text{ and } f(a) = g(b) \} \\ &\stackrel{\text{def}}{=} \{ ((x,a,b,y)) \in X \times A \times B \times Y \mid f(\phi(x)) = g(\psi(y)) \} \\ &\stackrel{\text{def}}{=} X \times_{K} Y. \end{split}$$

This finishes the proof.

Item 6: Unitality

We have

$$X \times_X A \cong \{(x, a) \in X \times A \mid f(a) = x\},\$$

$$A \times_X X \cong \{(a, x) \in X \times A \mid f(a) = x\},\$$

which are isomorphic to *A* via the maps $(x, a) \mapsto a$ and $(a, x) \mapsto a$. The proof of the naturality of $\lambda^{\text{Sets}_{/X}}$ and $\rho^{\text{Sets}_{/X}}$ is omitted.

Item 7: Commutativity

We have

$$A \times_C B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid f(a) = g(b)\}$$

$$= \{(a, b) \in A \times B \mid g(b) = f(a)\}$$

$$\cong \{(b, a) \in B \times A \mid g(b) = f(a)\}$$

$$\stackrel{\text{def}}{=} B \times_C A.$$

The proof of the naturality of $\sigma^{\text{Sets}_{/X}}$ is omitted.

Item 8: Distributivity Over Coproducts

We have

$$A \times_{X} (B \coprod C) \stackrel{\text{def}}{=} \left\{ (a, z) \in A \times (B \coprod C) \middle| \phi_{A}(a) = \phi_{B \coprod C}(z) \right\}$$

$$= \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (0, b) \text{ and } \phi_{A}(a) = \phi_{B \coprod C}(z) \right\}$$

$$\cup \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (1, c) \text{ and } \phi_{A}(a) = \phi_{B \coprod C}(z) \right\}$$

$$= \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (0, b) \text{ and } \phi_{A}(a) = \phi_{B}(b) \right\}$$

$$\cup \left\{ (a, z) \in A \times (B \coprod C) \middle| z = (1, c) \text{ and } \phi_{A}(a) = \phi_{C}(c) \right\}$$

$$\cong \left\{ (a, b) \in A \times B \middle| \phi_{A}(a) = \phi_{B}(b) \right\}$$

$$\cup \left\{ (a, c) \in A \times C \middle| \phi_{A}(a) = \phi_{C}(c) \right\}$$

$$\stackrel{\text{def}}{=} (A \times_{X} B) \cup (A \times_{X} C)$$

$$\cong (A \times_{X} B) \coprod (A \times_{X} C),$$

with the construction of the isomorphism

$$\delta_r^{\mathsf{Sets}_{/X}} : (A \coprod B) \times_X C \xrightarrow{\sim} (A \times_X C) \coprod (B \times_X C)$$

being similar. The proof of the naturality of $\delta_\ell^{\mathsf{Sets}_{/X}}$ and $\delta_r^{\mathsf{Sets}_{/X}}$ is omitted.

Item 9: Annihilation With the Empty Set

We have

$$A \times_X \emptyset \stackrel{\text{def}}{=} \{(a, b) \in A \times \emptyset \mid f(a) = g(b)\}$$
$$= \{k \in \emptyset \mid f(a) = g(b)\}$$
$$= \emptyset,$$

and similarly for $\emptyset \times_X A$, where we have used Item 8 of Proposition 4.1.3.1.4. The proof of the naturality of $\zeta_\ell^{\mathsf{Sets}_{/X}}$ and $\zeta_r^{\mathsf{Sets}_{/X}}$ is omitted.

Item 10: Interaction With Products

We have

$$A \times_{\text{pt}} B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid !_A(a) = !_B(b)\}$$

$$\stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid \star = \star\}$$

$$= \{(a, b) \in A \times B\}$$

$$= A \times B.$$

Item 11: Symmetric Monoidality

Omitted.

4.1.5 Equalisers

Let *A* and *B* be sets and let $f, g: A \Rightarrow B$ be functions.

DEFINITION 4.1.5.1.1 ► EQUALISERS OF SETS

The **equaliser of** f **and** g is the equaliser of f and g in Sets as in Limits and Colimits, ??.

CONSTRUCTION 4.1.5.1.2 ► CONSTRUCTION OF EQUALISERS OF SETS

Concretely, the equaliser of f and g is the pair (Eq(f,g),eq(f,g)) consisting of:

1. *The Limit*. The set Eq(f,g) defined by

$$\operatorname{Eq}(f,g) \stackrel{\text{def}}{=} \{ a \in A \mid f(a) = g(a) \}.$$

2. The Cone. The inclusion map

$$eq(f,g): Eq(f,g) \hookrightarrow A.$$

PROOF 4.1.5.1.3 ► PROOF OF CONSTRUCTION 4.1.5.1.2

We claim that Eq(f,g) is the categorical equaliser of f and g in Sets. First we need to check that the relevant equaliser diagram commutes, i.e. that we have

$$f \circ eq(f,g) = g \circ eq(f,g),$$

which indeed holds by the definition of the set Eq(f,g). Next, we prove that Eq(f,g) satisfies the universal property of the equaliser. Suppose we have a diagram of the form

$$\operatorname{Eq}(f,g) \xrightarrow{\operatorname{eq}(f,g)} A \xrightarrow{f \atop g} B$$

in Sets. Then there exists a unique map $\phi \colon E \to \operatorname{Eq}(f,g)$ making the diagram

$$Eq(f,g) \xrightarrow{eq(f,g)} A \xrightarrow{f} E$$

$$\downarrow f$$

$$\downarrow g$$

commute, being uniquely determined by the condition

$$eq(f,g) \circ \phi = e$$

via

$$\phi(x) = e(x)$$

for each $x \in E$, where we note that $e(x) \in A$ indeed lies in Eq(f,g) by the condition

$$f \circ e = g \circ e$$
,

which gives

$$f(e(x)) = g(e(x))$$

for each $x \in E$, so that $e(x) \in \text{Eq}(f, g)$.

PROPOSITION 4.1.5.1.4 ▶ PROPERTIES OF EQUALISERS OF SETS

Let A, B, and C be sets.

1. Associativity. We have isomorphisms of sets¹

$$\underbrace{\mathrm{Eq}(f \circ \mathrm{eq}(g,h), g \circ \mathrm{eq}(g,h))}_{=\mathrm{Eq}(f \circ \mathrm{eq}(g,h), h \circ \mathrm{eq}(g,h))} \cong \mathrm{Eq}(f,g,h) \cong \underbrace{\mathrm{Eq}(f \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))}_{=\mathrm{Eq}(g \circ \mathrm{eq}(f,g), h \circ \mathrm{eq}(f,g))}$$

where Eq(f, g, h) is the limit of the diagram

$$A \xrightarrow{\frac{f}{-g}} B$$

in Sets, being explicitly given by

$$\text{Eq}(f, g, h) \cong \{a \in A \mid f(a) = g(a) = h(a)\}.$$

2. Unitality. We have an isomorphism of sets

$$\operatorname{Eq}(f, f) \cong A$$
.

3. Commutativity. We have an isomorphism of sets

$$\operatorname{Eq}(f,g) \cong \operatorname{Eq}(g,f).$$

4. *Interaction With Composition*. Let

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C$$

be functions. We have an inclusion of sets

$$\operatorname{Eq}(h \circ f \circ \operatorname{eq}(f, g), k \circ g \circ \operatorname{eq}(f, g)) \subset \operatorname{Eq}(h \circ f, k \circ g),$$

where Eq($h \circ f \circ \text{eq}(f, g), k \circ g \circ \text{eq}(f, g)$) is the equaliser of the composition

$$\operatorname{Eq}(f,g) \stackrel{\operatorname{eq}(f,g)}{\hookrightarrow} A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C.$$

¹That is, the following three ways of forming "the" equaliser of (f, g, h) agree:

(a) Take the equaliser of (f, g, h), i.e. the limit of the diagram

$$A \xrightarrow{f \atop -g \xrightarrow{h}} B$$

in Sets.

(b) First take the equaliser of f and g, forming a diagram

$$\mathrm{Eq}(f,g) \overset{\mathrm{eq}(f,g)}{\hookrightarrow} A \overset{f}{\underset{g}{\Longrightarrow}} B$$

and then take the equaliser of the composition

$$\operatorname{Eq}(f,g) \stackrel{\operatorname{eq}(f,g)}{\hookrightarrow} A \stackrel{f}{\underset{h}{\Longrightarrow}} B,$$

obtaining a subset

$$Eq(f \circ eq(f,g), h \circ eq(f,g)) = Eq(g \circ eq(f,g), h \circ eq(f,g))$$

of Eq(f, g).

(c) First take the equaliser of g and h, forming a diagram

$$\operatorname{Eq}(g,h) \stackrel{\operatorname{eq}(g,h)}{\hookrightarrow} A \stackrel{g}{\underset{h}{\Longrightarrow}} B$$

and then take the equaliser of the composition

$$\mathrm{Eq}(g,h) \overset{\mathrm{eq}(g,h)}{\hookrightarrow} A \overset{f}{\underset{g}{\Longrightarrow}} B,$$

obtaining a subset

$$Eq(f \circ eq(g, h), g \circ eq(g, h)) = Eq(f \circ eq(g, h), h \circ eq(g, h))$$

of Eq(g, h).

PROOF 4.1.5.1.5 ► PROOF OF PROPOSITION 4.1.5.1.4

Item 1: Associativity

We first prove that Eq(f, g, h) is indeed given by

$$\operatorname{Eq}(f,g,h) \cong \{a \in A \mid f(a) = g(a) = h(a)\}.$$

Indeed, suppose we have a diagram of the form

$$\operatorname{Eq}(f,g,h) \xrightarrow{\operatorname{eq}(f,g,h)} A \xrightarrow{f \atop -g \xrightarrow{h}} B$$

in Sets. Then there exists a unique map $\phi \colon E \to \mathrm{Eq}(f,g,h)$, uniquely determined by the condition

$$eq(f,g) \circ \phi = e$$

being necessarily given by

$$\phi(x) = e(x)$$

for each $x \in E$, where we note that $e(x) \in A$ indeed lies in Eq(f, g, h) by the condition

$$f \circ e = g \circ e = h \circ e$$
,

which gives

$$f(e(x)) = g(e(x)) = h(e(x))$$

for each $x \in E$, so that $e(x) \in \text{Eq}(f, g, h)$.

We now check the equalities

$$\operatorname{Eq}(f \circ \operatorname{eq}(g, h), g \circ \operatorname{eq}(g, h)) \cong \operatorname{Eq}(f, g, h) \cong \operatorname{Eq}(f \circ \operatorname{eq}(f, g), h \circ \operatorname{eq}(f, g)).$$

Indeed, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(g,h), g \circ \operatorname{eq}(g,h)) &\cong \{x \in \operatorname{Eq}(g,h) \mid [f \circ \operatorname{eq}(g,h)](a) = [g \circ \operatorname{eq}(g,h)](a)\} \\ &\cong \{x \in \operatorname{Eq}(g,h) \mid f(a) = g(a)\} \\ &\cong \{x \in A \mid f(a) = g(a) \text{ and } g(a) = h(a)\} \\ &\cong \{x \in A \mid f(a) = g(a) = h(a)\} \\ &\cong \operatorname{Eq}(f,g,h). \end{split}$$

Similarly, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(f,g), h \circ \operatorname{eq}(f,g)) &\cong \{x \in \operatorname{Eq}(f,g) \mid [f \circ \operatorname{eq}(f,g)](a) = [h \circ \operatorname{eq}(f,g)](a)\} \\ &\cong \{x \in \operatorname{Eq}(f,g) \mid f(a) = h(a)\} \\ &\cong \{x \in A \mid f(a) = h(a) \text{ and } f(a) = g(a)\} \\ &\cong \{x \in A \mid f(a) = g(a) = h(a)\} \end{split}$$

$$\cong$$
 Eq (f, g, h) .

Item 2: Unitality

Indeed, we have

$$\operatorname{Eq}(f, f) \stackrel{\text{def}}{=} \{ a \in A \mid f(a) = f(a) \}$$
$$= A.$$

Item 3: Commutativity

Indeed, we have

$$\operatorname{Eq}(f,g) \stackrel{\text{def}}{=} \{ a \in A \mid f(a) = g(a) \}$$
$$= \{ a \in A \mid g(a) = f(a) \}$$
$$\stackrel{\text{def}}{=} \operatorname{Eq}(g,f).$$

Item 4: Interaction With Composition

Indeed, we have

$$\begin{split} \operatorname{Eq}(h \circ f \circ \operatorname{eq}(f,g), k \circ g \circ \operatorname{eq}(f,g)) & \cong \{a \in \operatorname{Eq}(f,g) \mid h(f(a)) = k(g(a))\} \\ & \cong \{a \in A \mid f(a) = g(a) \text{ and } h(f(a)) = k(g(a))\}. \end{split}$$

and

$$Eq(h \circ f, k \circ g) \cong \{a \in A \mid h(f(a)) = k(g(a))\},\$$

and thus there's an inclusion from Eq $(h \circ f \circ eq(f,g), k \circ g \circ eq(f,g))$ to Eq $(h \circ f, k \circ g)$.

4.1.6 Inverse Limits

Let $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I} \colon (I, \preceq) \to \mathsf{Sets}$ be an inverse system of sets.

DEFINITION 4.1.6.1.1 ► INVERSE LIMITS OF SETS

The **inverse limit of** $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the inverse limit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ in Sets as in Limits and Colimits, ??.

CONSTRUCTION 4.1.6.1.2 ► CONSTRUCTION OF INVERSE LIMITS OF SETS

Concretely, the inverse limit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the pair $\left(\lim_{\stackrel{\longleftarrow}{\alpha\in I}} (X_{\alpha}), \left\{ \operatorname{pr}_{\alpha} \right\}_{\alpha\in I} \right)$ consisting of:

1. *The Limit*. The set $\lim_{\substack{\longleftarrow \\ \alpha \in I}} (X_{\alpha})$ defined by

$$\lim_{\substack{\longleftarrow \\ \alpha \in I}} (X_{\alpha}) \stackrel{\text{def}}{=} \left\{ (x_{\alpha})_{\alpha \in I} \in \prod_{\alpha \in I} X_{\alpha} \mid \text{ for each } \alpha, \beta \in I, \text{ if } \alpha \preceq \beta, \\ \text{then we have } x_{\alpha} = f_{\alpha\beta}(x_{\beta}) \right\}.$$

2. The Cone. The collection

$$\left\{ \operatorname{pr}_{\gamma} \colon \lim_{\substack{\longleftarrow \\ \alpha \in I}} (X_{\alpha}) \to X_{\gamma} \right\}_{\gamma \in I}$$

of maps of sets defined as the restriction of the maps

$$\left\{ \operatorname{pr}_{\gamma} : \prod_{\alpha \in I} X_{\alpha} \to X_{\gamma} \right\}_{\gamma \in I}$$

of Item 2 of Construction 4.1.2.1.2 to $\lim_{\alpha \in I} (X_{\alpha})$ and hence given by

$$\operatorname{pr}_{\gamma}((x_{\alpha})_{\alpha\in I})\stackrel{\operatorname{def}}{=} x_{\gamma}$$

for each $\gamma \in I$ and each $(x_{\alpha})_{\alpha \in I} \in \lim_{\stackrel{\longleftarrow}{\alpha \in I}} (X_{\alpha})$.

PROOF 4.1.6.1.3 ► PROOF OF CONSTRUCTION 4.1.6.1.2

We claim that $\lim_{\leftarrow \alpha \in I} (X_{\alpha})$ is the limit of the inverse system of sets $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta \in I}$. First we need to check that the limit diagram defined by

it commutes, i.e. that we have

$$f_{\alpha\beta} \circ \operatorname{pr}_{\alpha} = \operatorname{pr}_{\beta}, \qquad \lim_{\substack{\longleftarrow \\ \operatorname{pr}_{\alpha} \\ X_{\alpha} \xrightarrow{f_{\alpha\beta}}}} X_{\beta}$$

for each $\alpha, \beta \in I$ with $\alpha \leq \beta$. Indeed, given $(x_{\gamma})_{\gamma \in I} \in \lim_{\longleftarrow \gamma \in I} (X_{\gamma})$, we have

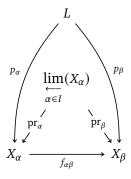
$$[f_{\alpha\beta} \circ \operatorname{pr}_{\alpha}]((x_{\gamma})_{\gamma \in I}) \stackrel{\text{def}}{=} f_{\alpha\beta}(\operatorname{pr}_{\alpha}((x_{\gamma})_{\gamma \in I}))$$

$$\stackrel{\text{def}}{=} f_{\alpha\beta}(x_{\alpha})$$

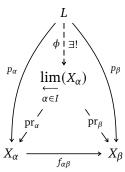
$$= x_{\beta}$$

$$\stackrel{\text{def}}{=} \operatorname{pr}_{\beta}((x_{\gamma})_{\gamma \in I}),$$

where the third equality comes from the definition of $\lim_{\stackrel{\longleftarrow}{\alpha}\in I}(X_{\alpha})$. Next, we prove that $\lim_{\stackrel{\longleftarrow}{\alpha}\in I}(X_{\alpha})$ satisfies the universal property of an inverse limit. Suppose that we have, for each $\alpha,\beta\in I$ with $\alpha\preceq\beta$, a diagram of the form



in Sets. Then there indeed exists a unique map $\phi: L \xrightarrow{\exists !} \varprojlim_{\alpha \in I} (X_{\alpha})$ making the diagram



commute, being uniquely determined by the family of conditions

$$\left\{p_{\alpha} = \operatorname{pr}_{\alpha} \circ \phi\right\}_{\alpha \in I}$$

via

$$\phi(\ell) = (p_{\alpha}(\ell))_{\alpha \in I}$$

for each $\ell \in L$, where we note that $(p_{\alpha}(\ell))_{\alpha \in I} \in \prod_{\alpha \in I} X_{\alpha}$ indeed lies in $\lim_{\leftarrow \alpha \in I} (X_{\alpha})$, as we have

$$f_{\alpha\beta}(p_{\alpha}(\ell)) \stackrel{\text{def}}{=} [f_{\alpha\beta} \circ p_{\alpha}](\ell)$$
$$\stackrel{\text{def}}{=} p_{\beta}(\ell)$$

for each $\beta \in I$ with $\alpha \leq \beta$ by the commutativity of the diagram for $(L, \{p_{\alpha}\}_{\alpha \in I})$.

EXAMPLE 4.1.6.1.4 ► EXAMPLES OF INVERSE LIMITS OF SETS

Here are some examples of inverse limits of sets.

1. *The p-Adic Integers*. The ring of *p*-adic integers \mathbb{Z}_p of $\ref{p-Adic}$ is the inverse limit

$$\mathbb{Z}_p \cong \lim_{\substack{\longleftarrow \\ n \in \mathbb{N}}} (\mathbb{Z}_{/p^n});$$

see ??.

2. Rings of Formal Power Series. The ring R[[t]] of formal power series in

a variable *t* is the inverse limit

$$R[[t]] \cong \lim_{\substack{\longleftarrow \\ n \in \mathbb{N}}} (R[t]/t^n R[t]);$$

see ??.

3. *Profinite Groups*. Profinite groups are inverse limits of finite groups; see ??.

4.2 Colimits of Sets

4.2.1 The Initial Set

DEFINITION 4.2.1.1.1 ► THE INITIAL SET

The initial set is the initial object of Sets as in Limits and Colimits, ??.

CONSTRUCTION 4.2.1.1.2 ► CONSTRUCTION OF THE INITIAL SET

Concretely, the initial set is the pair $(\emptyset, \{\iota_A\}_{A \in Obi(Sets)})$ consisting of:

- 1. *The Colimit*. The empty set Ø of Definition 4.3.1.1.1.
- 2. The Cocone. The collection of maps

$$\{\iota_A \colon \emptyset \to A\}_{A \in \mathrm{Obj}(\mathsf{Sets})}$$

given by the inclusion maps from \emptyset to A.

PROOF 4.2.1.1.3 ► PROOF OF CONSTRUCTION 4.2.1.1.2

We claim that \emptyset is the initial object of Sets. Indeed, suppose we have a diagram of the form

$$\emptyset$$
 A

in Sets. Then there exists a unique map $\phi: \emptyset \to A$ making the diagram

$$\emptyset - \frac{\phi}{\exists !} \rightarrow A$$

commute, namely the inclusion map ι_A .

4.2.2 Coproducts of Families of Sets

Let $\{A_i\}_{i\in I}$ be a family of sets.

DEFINITION 4.2.2.1.1 ► THE COPRODUCT OF A FAMILY OF SETS

The **coproduct of** $\{A_i\}_{i\in I}^1$ is the coproduct of $\{A_i\}_{i\in I}$ in Sets as in Limits and Colimits, ??.

CONSTRUCTION 4.2.2.1.2 ► CONSTRUCTION OF THE COPRODUCT OF A FAMILY OF SETS

Concretely, the disjoint union of $\{A_i\}_{i\in I}$ is the pair $(\coprod_{i\in I} A_i, \{\inf_i\}_{i\in I})$ consisting of:

1. *The Colimit.* The set $\coprod_{i \in I} A_i$ defined by

$$\coprod_{i \in I} A_i \stackrel{\text{def}}{=} \left\{ (i, x) \in I \times \left(\bigcup_{i \in I} A_i \right) \mid x \in A_i \right\}.$$

2. The Cocone. The collection

$$\left\{\inf_{i}\colon A_{i}\to \coprod_{i\in I}A_{i}\right\}_{i\in I}$$

of maps given by

$$inj_i(x) \stackrel{\text{def}}{=} (i, x)$$

for each $x \in A_i$ and each $i \in I$.

¹Further Terminology: Also called the **disjoint union of the family** $\{A_i\}_{i\in I}$.

PROOF 4.2.2.1.3 ► PROOF OF CONSTRUCTION 4.2.2.1.2

We claim that $\coprod_{i \in I} A_i$ is the categorical coproduct of $\{A_i\}_{i \in I}$ in Sets. Indeed, suppose we have, for each $i \in I$, a diagram of the form

$$A_i \xrightarrow[\text{inj}]{l} A_i$$

in Sets. Then there exists a unique map $\phi: \coprod_{i \in I} A_i \to C$ making the diagram

$$A_{i} \xrightarrow{\lim_{i \in I} A_{i}} A_{i}$$

commute, being uniquely determined by the condition $\phi \circ \text{inj}_i = \iota_i$ for each $i \in I$ via

$$\phi((i,x)) = \iota_i(x)$$

for each $(i, x) \in \coprod_{i \in I} A_i$.

PROPOSITION 4.2.2.1.4 ▶ PROPERTIES OF COPRODUCTS OF FAMILIES OF SETS

Let $\{A_i\}_{i\in I}$ be a family of sets.

1. Functoriality. The assignment $\{A_i\}_{i\in I}\mapsto \coprod_{i\in I}A_i$ defines a functor

$$\coprod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

• *Action on Objects.* For each $(A_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$, we have

$$\left[\bigsqcup_{i \in I} \right] ((A_i)_{i \in I}) \stackrel{\text{def}}{=} \bigsqcup_{i \in I} A_i$$

• Action on Morphisms. For each $(A_i)_{i \in I}, (B_i)_{i \in I} \in \text{Obj}(\text{Fun}(I_{\text{disc}}, \text{Sets}))$, the action on Hom-sets

$$\left(\bigsqcup_{i\in I}\right)_{(A_i)_{i\in I},(B_i)_{i\in I}}: \operatorname{Nat}((A_i)_{i\in I},(B_i)_{i\in I}) \to \operatorname{Sets}\left(\bigsqcup_{i\in I}A_i,\bigsqcup_{i\in I}B_i\right)$$

of $\coprod_{i\in I}$ at $((A_i)_{i\in I},(B_i)_{i\in I})$ is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in Nat($(A_i)_{i \in I}$, $(B_i)_{i \in I}$) to the map of sets

$$\coprod_{i \in I} f_i \colon \coprod_{i \in I} A_i \to \coprod_{i \in I} B_i$$

defined by

$$\left[\bigsqcup_{i \in I} f_i \right] (i, a) \stackrel{\text{def}}{=} f_i(a)$$

for each $(i, a) \in \coprod_{i \in I} A_i$.

PROOF 4.2.2.1.5 ► PROOF OF PROPOSITION 4.2.2.1.4

Item 1: Functoriality

This follows from Limits and Colimits, ?? of ??.

4.2.3 Binary Coproducts

Let A and B be sets.

DEFINITION 4.2.3.1.1 ► COPRODUCTS OF SETS

The **coproduct of** A **and** B^1 is the coproduct of A and B in Sets as in Limits and Colimits, ??.

¹Further Terminology: Also called the **disjoint union of** A **and** B.

CONSTRUCTION 4.2.3.1.2 ► CONSTRUCTION OF COPRODUCTS OF SETS

Concretely, the coproduct of *A* and *B* is the pair $(A \coprod B, \{\text{inj}_1, \text{inj}_2\})$ consisting of:

1. *The Colimit*. The set $A \coprod B$ defined by

$$A \coprod B \stackrel{\text{def}}{=} \coprod_{z \in \{A, B\}} z$$

$$\stackrel{\text{def}}{=} \{(0, a) \in S \mid a \in A\} \cup \{(1, b) \in S \mid b \in B\},$$

where $S = \{0, 1\} \times (A \cup B)$.

2. The Cocone. The maps

$$inj_1: A \to A \coprod B,$$

 $inj_2: B \to A \coprod B,$

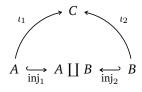
given by

$$\operatorname{inj}_{1}(a) \stackrel{\text{def}}{=} (0, a),$$
 $\operatorname{inj}_{2}(b) \stackrel{\text{def}}{=} (1, b),$

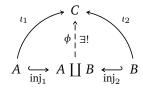
for each $a \in A$ and each $b \in B$.

PROOF 4.2.3.1.3 ► PROOF OF CONSTRUCTION 4.2.3.1.2

We claim that $A \coprod B$ is the categorical coproduct of A and B in Sets. Indeed, suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: A \coprod B \to C$ making the diagram



commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_A = \iota_A,$$
$$\phi \circ \operatorname{inj}_B = \iota_B$$

via

$$\phi(x) = \begin{cases} \iota_A(a) & \text{if } x = (0, a), \\ \iota_B(b) & \text{if } x = (1, b) \end{cases}$$

for each $x \in A \coprod B$.

PROPOSITION 4.2.3.1.4 ► PROPERTIES OF COPRODUCTS OF SETS

Let A, B, C, and X be sets.

1. Functoriality. The assignment $A, B, (A, B) \mapsto A \coprod B$ defines functors

$$\begin{array}{cccc} A \coprod -\colon & \mathsf{Sets} & \to \mathsf{Sets}, \\ - \coprod B\colon & \mathsf{Sets} & \to \mathsf{Sets}, \\ -_1 \coprod -_2\colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}, \end{array}$$

where $-1 \coprod -2$ is the functor where

• *Action on Objects.* For each $(A, B) \in Obj(Sets \times Sets)$, we have

$$[-_1 \coprod -_2](A, B) \stackrel{\text{def}}{=} A \coprod B.$$

• *Action on Morphisms*. For each $(A, B), (X, Y) \in Obj(Sets)$, the action on Hom-sets

$$\coprod_{(A,B),(X,Y)}$$
: Sets $(A,X) \times$ Sets $(B,Y) \rightarrow$ Sets $(A \coprod B,X \coprod Y)$ of \coprod at $((A,B),(X,Y))$ is defined by sending (f,g) to the function

$$f \mid \mid g: A \mid \mid B \rightarrow X \mid \mid Y$$

defined by

$$[f \coprod g](x) \stackrel{\text{def}}{=} \begin{cases} (0, f(a)) & \text{if } x = (0, a), \\ (1, g(b)) & \text{if } x = (1, b), \end{cases}$$

for each $x \in A \mid A \mid B$.

and where $A \coprod -$ and $- \coprod B$ are the partial functors of $-_1 \coprod -_2$ at $A, B \in \text{Obj}(\mathsf{Sets})$.

2. Adjointness. We have an adjunction

$$(-_1 \coprod -_2 \dashv \Delta_{\mathsf{Sets}})$$
: Sets \times Sets $\underbrace{\bot}_{\Delta_{\mathsf{Sets}}}$ Sets,

witnessed by a bijection

$$\mathsf{Sets}(A \coprod B, C), \cong \mathsf{Hom}_{\mathsf{Sets} \times \mathsf{Sets}}((A, B), (C, C))$$

natural in $(A, B) \in \text{Obj}(\mathsf{Sets} \times \mathsf{Sets})$ and in $C \in \text{Obj}(\mathsf{Sets})$.

3. Associativity. We have an isomorphism of sets

$$\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod} : (X \coprod Y) \coprod Z \xrightarrow{\sim} X \coprod (Y \coprod Z),$$

natural in $X, Y, Z \in Obj(Sets)$.

4. Unitality. We have isomorphisms of sets

$$\lambda_X^{\mathsf{Sets}, \coprod} : \emptyset \coprod X \xrightarrow{\sim} X,$$
$$\rho_X^{\mathsf{Sets}, \coprod} : X \coprod \emptyset \xrightarrow{\sim} X,$$

natural in $X \in \text{Obj}(\mathsf{Sets})$.

5. Commutativity. We have an isomorphism of sets

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod}: X\coprod Y\stackrel{\sim}{\dashrightarrow} Y\coprod X,$$

natural in $X, Y \in Obj(Sets)$.

6. *Symmetric Monoidality*. The 7-tuple (Sets, \coprod , \emptyset , α_{\coprod}^{Sets} , λ_{\coprod}^{Sets} , ρ_{\coprod}^{Sets} , σ_{\coprod}^{Sets}) is a symmetric monoidal category.

PROOF 4.2.3.1.5 ► PROOF OF PROPOSITION 4.2.3.1.4

Item 1: Functoriality

This follows from Limits and Colimits, ?? of ??.

Item 2: Adjointness

This follows from the universal property of the coproduct.

Item 3: Associativity

This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.2.3.1.1.

Item 4: Unitality

This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.2.4.1.1 and 5.2.5.1.1.

Item 5: Commutativity

This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.2.6.1.1.

Item 6: Symmetric Monoidality

This is a repetition of Monoidal Structures on the Category of Sets, Proposition 5.2.7.1.1, and is proved there.

4.2.4 Pushouts

Let A, B, and C be sets and let $f: C \to A$ and $g: C \to B$ be functions.

DEFINITION 4.2.4.1.1 ► PUSHOUTS OF SETS

The **pushout of** A **and** B **over** C **along** f **and** g^1 is the pushout of A and B over C along f and g in Sets as in Limits and Colimits, ??.

CONSTRUCTION 4.2.4.1.2 ➤ CONSTRUCTION OF PUSHOUTS OF SETS

Concretely, the pushout of *A* and *B* over *C* along *f* and *g* is the pair $(A \coprod_C B, \{inj_1, inj_2\})$ consisting of:

1. *The Colimit*. The set $A \coprod_C B$ defined by

$$A \coprod_C B \stackrel{\mathrm{def}}{=} A \coprod B/\sim_C,$$

¹Further Terminology: Also called the **fibre coproduct of** A **and** B **over** C **along** f **and** g.

where \sim_C is the equivalence relation on $A \coprod B$ generated by $(0, f(c)) \sim_C (1, g(c))$.

2. *The Cocone*. The maps

$$inj_1: A \to A \coprod_C B,$$

 $inj_2: B \to A \coprod_C B$

given by

$$inj_1(a) \stackrel{\text{def}}{=} [(0, a)]
inj_2(b) \stackrel{\text{def}}{=} [(1, b)]$$

for each $a \in A$ and each $b \in B$.

PROOF 4.2.4.1.3 ► PROOF OF CONSTRUCTION 4.2.4.1.2

We claim that $A \coprod_C B$ is the categorical pushout of A and B over C with respect to (f,g) in Sets. First we need to check that the relevant pushout diagram commutes, i.e. that we have

$$\begin{array}{ccc}
A \coprod_{C} B & \stackrel{\operatorname{inj}_{2}}{\longleftarrow} B \\
\operatorname{inj}_{1} \circ f & = \operatorname{inj}_{2} \circ g, & & \inf_{1} & f \\
A & \longleftarrow_{f} & C.
\end{array}$$

Indeed, given $c \in C$, we have

$$[inj_1 \circ f](c) = inj_1(f(c))$$

$$= [(0, f(c))]$$

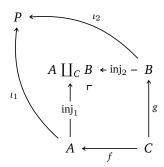
$$= [(1, g(c))]$$

$$= inj_2(g(c))$$

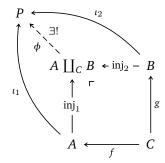
$$= [inj_2 \circ g](c),$$

where [(0, f(c))] = [(1, g(c))] by the definition of the relation \sim on $A \coprod B$. Next, we prove that $A \coprod {}_{C}B$ satisfies the universal property of the pushout.

Suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: A \coprod_C B \to P$ making the diagram



commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_1 = \iota_1,$$

$$\phi \circ \operatorname{inj}_2 = \iota_2$$

via

$$\phi(x) = \begin{cases} \iota_1(a) & \text{if } x = [(0, a)], \\ \iota_2(b) & \text{if } x = [(1, b)] \end{cases}$$

for each $x \in A \coprod_C B$, where the well-definedness of ϕ is guaranteed by the equality $\iota_1 \circ f = \iota_2 \circ g$ and the definition of the relation \sim on $A \coprod B$ as follows:

1. Case 1: Suppose we have x = [(0, a)] = [(0, a')] for some $a, a' \in A$. Then, by Remark 4.2.4.1.4, we have a sequence

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0,a').$$

2. Case 2: Suppose we have x = [(1,b)] = [(1,b')] for some $b,b' \in B$. Then, by Remark 4.2.4.1.4, we have a sequence

$$(1,b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b').$$

3. Case 3: Suppose we have x = [(0, a)] = [(1, b)] for some $a \in A$ and $b \in B$. Then, by Remark 4.2.4.1.4, we have a sequence

$$(0, a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1, b).$$

In all these cases, we declare $x \sim' y$ iff there exists some $c \in C$ such that x = (0, f(c)) and y = (1, g(c)) or x = (1, g(c)) and y = (0, f(c)). Then, the equality $\iota_1 \circ f = \iota_2 \circ g$ gives

$$\begin{aligned} \phi([x]) &= \phi([(0, f(c))]) \\ &\stackrel{\text{def}}{=} \iota_1(f(c)) \\ &= \iota_2(g(c)) \\ &\stackrel{\text{def}}{=} \phi([(1, g(c))]) \\ &= \phi([y]), \end{aligned}$$

with the case where x = (1, g(c)) and y = (0, f(c)) similarly giving $\phi([x]) = \phi([y])$. Thus, if $x \sim' y$, then $\phi([x]) = \phi([y])$. Applying this equality pairwise to the sequences

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0,a'),$$

 $(1,b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b'),$
 $(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b)$

gives

$$\begin{split} \phi([(0,a)]) &= \phi([(0,a')]), \\ \phi([(1,b)]) &= \phi([(1,b')]), \\ \phi([(0,a)]) &= \phi([(1,b)]), \end{split}$$

showing ϕ to be well-defined.

REMARK 4.2.4.1.4 ► Unwinding Definition 4.2.4.1.1

In detail, by Conditions on Relations, Construction 10.5.2.1.2, the relation \sim of Definition 4.2.4.1.1 is given by declaring $a \sim b$ iff one of the following conditions is satisfied:

- 1. We have $a, b \in A$ and a = b.
- 2. We have $a, b \in B$ and a = b.
- 3. There exist $x_1, \ldots, x_n \in A \coprod B$ such that $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$, where we declare $x \sim' y$ if one of the following conditions is satisfied:
 - (a) There exists $c \in C$ such that x = (0, f(c)) and y = (1, g(c)).
 - (b) There exists $c \in C$ such that x = (1, g(c)) and y = (0, f(c)).

In other words, there exist $x_1, \ldots, x_n \in A \coprod B$ satisfying the following conditions:

- (c) There exists $c_0 \in C$ satisfying one of the following conditions:
 - i. We have $a = f(c_0)$ and $x_1 = g(c_0)$.
 - ii. We have $a = g(c_0)$ and $x_1 = f(c_0)$.
- (d) For each $1 \le i \le n-1$, there exists $c_i \in C$ satisfying one of the following conditions:
 - i. We have $x_i = f(c_i)$ and $x_{i+1} = g(c_i)$.
 - ii. We have $x_i = g(c_i)$ and $x_{i+1} = f(c_i)$.
- (e) There exists $c_n \in C$ satisfying one of the following conditions:
 - i. We have $x_n = f(c_n)$ and $b = g(c_n)$.
 - ii. We have $x_n = g(c_n)$ and $b = f(c_n)$.

REMARK 4.2.4.1.5 ▶ PUSHOUTS OF SETS DEPEND ON THE MAPS

It is common practice to write $A \coprod_C B$ for the pushout of A and B over C along f and g, omitting the maps f and g from the notation and instead leaving them implicit, to be understood from the context.

However, the set $A \coprod_{C} B$ depends very much on the maps f and g, and sometimes it is necessary or useful to note this dependence explicitly. In such situations, we will write $A \coprod_{f,C,g} B$ or $A \coprod_{C} G B$ for $A \coprod_{C} B B$.

EXAMPLE 4.2.4.1.6 ► **EXAMPLES OF PUSHOUTS OF SETS**

Here are some examples of pushouts of sets.

- 1. Wedge Sums of Pointed Sets. The wedge sum of two pointed sets of Pointed Sets, Definition 6.3.3.1.1 is an example of a pushout of sets.
- 2. *Intersections via Unions*. Let *X* be a set. We have

$$A \cup B \cong A \coprod_{A \cap B} B, \qquad A \longleftarrow B$$

$$A \longleftarrow A \longrightarrow A \cap B$$

for each $A, B \in \mathcal{P}(X)$.

PROOF 4.2.4.1.7 ► PROOF OF EXAMPLE 4.2.4.1.6

Item 1: Wedge Sums of Pointed Sets

This follows by definition, as the wedge sum of two pointed sets is defined as a pushout.

Item 2: Intersections via Unions

Indeed, $A \coprod_{A \cap B} B$ is the quotient of $A \coprod B$ by the equivalence relation obtained by declaring $(0, a) \sim (1, b)$ iff $a = b \in A \cap B$, which is in bijection with $A \cup B$ via the map with $[(0, a)] \mapsto a$ and $[(1, b)] \mapsto b$.

PROPOSITION 4.2.4.1.8 ► PROPERTIES OF PUSHOUTS OF SETS

Let *A*, *B*, *C*, and *X* be sets.

1. Functoriality. The assignment $(A, B, C, f, g) \mapsto A \coprod_{f,C,g} B$ defines a functor

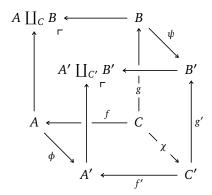
$$-_1 \coprod_{-_3} -_1 \colon \mathsf{Fun}(\mathcal{P},\mathsf{Sets}) \to \mathsf{Sets},$$

where \mathcal{P} is the category that looks like this:



In particular, the action on morphisms of $-1 \coprod_{-3} -1$ is given by sending

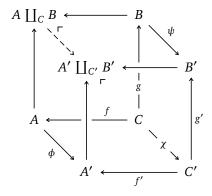
a morphism



in Fun(\mathcal{P} , Sets) to the map $\xi: A \coprod_C B \xrightarrow{\exists!} A' \coprod_{C'} B'$ given by

$$\xi(x) \stackrel{\text{def}}{=} \begin{cases} \phi(a) & \text{if } x = [(0, a)], \\ \psi(b) & \text{if } x = [(1, b)] \end{cases}$$

for each $x \in A \coprod_C B$, which is the unique map making the diagram



commute.

2. Adjointness. We have an adjunction

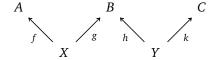
$$\left(-1 \coprod_{X} -_2 \dashv \Delta_{\mathsf{Sets}_{X/}}\right) \colon \quad \mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/} \underbrace{\bot_{\Delta_{\mathsf{Sets}_{X/}}}}_{\mathsf{Sets}_{X/}} \mathsf{Sets}_{X/},$$

witnessed by a bijection

$$\mathsf{Sets}_{X/}(A \coprod_X B, C), \cong \mathsf{Hom}_{\mathsf{Sets}_{X/}} \times \mathsf{Sets}_{X/}((A, B), (C, C))$$

natural in $(A, B) \in \mathsf{Obj}(\mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/})$ and in $C \in \mathsf{Obj}(\mathsf{Sets}_{X/})$.

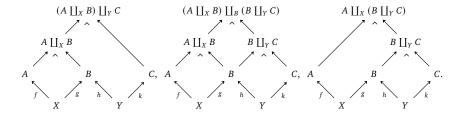
3. Associativity. Given a diagram



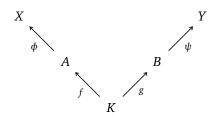
in Sets, we have isomorphisms of sets

$$(A \coprod_X B) \coprod_Y C \cong (A \coprod_X B) \coprod_B (B \coprod_Y C) \cong A \coprod_X (B \coprod_Y C)$$

where these pullbacks are built as in the diagrams



4. Interaction With Composition. Given a diagram



in Sets, we have isomorphisms of sets

$$\begin{split} X \coprod_K^{\phi \circ f, \psi \circ g} Y &\cong (X \coprod_A^{\phi, j_1} (A \coprod_K^{f, g} B)) \coprod_{A \coprod_K^{f, g} B}^{i_2, i_1} ((A \coprod_K^{f, g} B) \coprod_B^{j_2, \psi} Y) \\ &\cong X \coprod_A^{\phi, i} ((A \coprod_K^{f, g} B) \coprod_B^{j_2, \psi} Y) \\ &\cong (X \coprod_A^{\phi, i_1} (A \coprod_K^{f, g} B)) \coprod_B^{j, \psi} Y \end{split}$$

where

$$j_{1} = \text{inj}_{1}^{A \times_{K}^{f,g} B}, \qquad j_{2} = \text{inj}_{2}^{A \times_{K}^{f,g} B},$$

$$i_{1} = \text{inj}_{1}^{(A \times_{K}^{f,g} B) \times_{Y}^{q_{2},\psi}}, \qquad X \times_{K}^{\phi,q_{1}} (A \times_{K}^{f,g} B)},$$

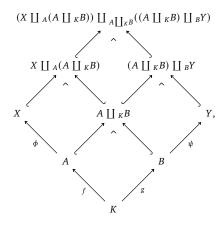
$$i_{2} = \text{inj}_{2}^{A \times_{K}^{f,g} B} (A \times_{K}^{f,g} B),$$

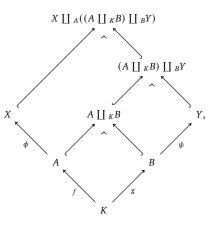
$$i_{2} = \text{inj}_{2}^{A \times_{K}^{f,g} B},$$

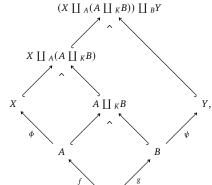
$$j_{3} = j_{2} \circ \text{inj}_{2}^{X \times_{K}^{\phi,q_{1}} (A \times_{K}^{f,g} B)},$$

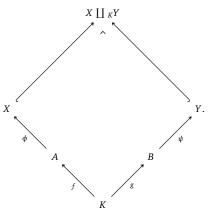
$$j_{3} = j_{4} \circ \text{inj}_{2}^{X \times_{K}^{\phi,q_{1}} (A \times_{K}^{f,g} B)},$$

and where these pullbacks are built as in the diagrams









5. Unitality. We have isomorphisms of sets

$$\begin{array}{ccc}
A & \longrightarrow & A \\
f & & \uparrow \\
X & \longrightarrow & X
\end{array}$$

$$\lambda_A^{\mathsf{Sets}_{X/}} : X \coprod_X A \xrightarrow{\sim} A,$$
 $\rho_A^{\mathsf{Sets}_{X/}} : A \coprod_X X \xrightarrow{\sim} A,$

$$\begin{array}{cccc}
A & \stackrel{f}{\longleftarrow} & X \\
\parallel & & \parallel \\
X & \stackrel{f}{\longleftarrow} & X,
\end{array}$$

natural in $(A, f) \in \mathsf{Obj}(\mathsf{Sets}_{X/})$.

6. Commutativity. We have an isomorphism of sets

natural in $(A, f), (B, g) \in \text{Obj}(\mathsf{Sets}_{X/})$.

7. Interaction With Coproducts. We have

$$A \coprod_{\emptyset} B \cong A \coprod_{B}, \qquad \uparrow \qquad \uparrow_{\iota_{B}}$$

$$A \longleftarrow_{\iota_{A}} \emptyset.$$

8. *Symmetric Monoidality*. The triple $(\mathsf{Sets}_{X/}, \coprod_X, X)$ is a symmetric monoidal category.

PROOF 4.2.4.1.9 ► PROOF OF PROPOSITION 4.2.4.1.8

Item 1: Functoriality

This is a special case of functoriality of co/limits, Limits and Colimits, ?? of ??, with the explicit expression for ξ following from the commutativity of the cube pushout diagram.

Item 2: : Adjointness

This follows from the universal property of the coproduct (pushouts are coproducts in $Sets_{X/}$).

Item 3: Associativity

Omitted.

Item 4: Interaction With Composition

Omitted.

Item 5: Unitality

Omitted.

Item 6: Commutativity
Omitted.

Item 7: Interaction With Coproducts
Omitted.

Item 8: Symmetric Monoidality
Omitted.

4.2.5 Coequalisers

Let *A* and *B* be sets and let $f, g: A \Rightarrow B$ be functions.

DEFINITION 4.2.5.1.1 ► COEQUALISERS OF SETS

The **coequaliser of** f **and** g is the coequaliser of f and g in Sets as in Limits and Colimits, ??.

CONSTRUCTION 4.2.5.1.2 ► CONSTRUCTION OF COEQUALISERS OF SETS

Concretely, the coequaliser of f and g is the pair (CoEq(f, g), coeq(f, g)) consisting of:

1. *The Colimit.* The set CoEq(f, g) defined by

$$CoEq(f, g) \stackrel{\text{def}}{=} B/\sim$$
,

where \sim is the equivalence relation on *B* generated by $f(a) \sim g(a)$.

2. The Cocone. The map

$$coeq(f,g): B \rightarrow CoEq(f,g)$$

given by the quotient map $\pi: B \twoheadrightarrow B/\sim$ with respect to the equivalence relation generated by $f(a) \sim g(a)$.

PROOF 4.2.5.1.3 ► PROOF OF CONSTRUCTION 4.2.5.1.2

We claim that CoEq(f,g) is the categorical coequaliser of f and g in Sets. First we need to check that the relevant coequaliser diagram commutes, i.e. that we have

$$coeq(f, g) \circ f = coeq(f, g) \circ g$$
.

Indeed, we have

$$[\operatorname{coeq}(f,g) \circ f](a) \stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](f(a))$$

$$\stackrel{\text{def}}{=} [f(a)]$$

$$= [g(a)]$$

$$\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](g(a))$$

$$\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g) \circ g](a)$$

for each $a \in A$. Next, we prove that CoEq(f,g) satisfies the universal property of the coequaliser. Suppose we have a diagram of the form

$$A \xrightarrow{f \atop g} B \xrightarrow{\operatorname{coeq}(f,g)} \operatorname{CoEq}(f,g)$$

in Sets. Then, since c(f(a)) = c(g(a)) for each $a \in A$, it follows from Conditions on Relations, Items 4 and 5 of Proposition 10.6.2.1.3 that there exists a unique map $CoEq(f,g) \xrightarrow{\exists !} C$ making the diagram

$$A \xrightarrow{f} B \xrightarrow{\operatorname{coeq}(f,g)} \operatorname{CoEq}(f,g)$$

$$\downarrow g$$

$$\downarrow \exists !$$

$$C$$

commute.

REMARK 4.2.5.1.4 ► Unwinding Definition 4.2.5.1.1

In detail, by Conditions on Relations, Construction 10.5.2.1.2, the relation \sim of Definition 4.2.5.1.1 is given by declaring $a \sim b$ iff one of the following conditions is satisfied:

1. We have a = b;

- 2. There exist $x_1, \ldots, x_n \in B$ such that $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$, where we declare $x \sim' y$ if one of the following conditions is satisfied:
 - (a) There exists $z \in A$ such that x = f(z) and y = g(z).
 - (b) There exists $z \in A$ such that x = g(z) and y = f(z).

In other words, there exist $x_1, \ldots, x_n \in B$ satisfying the following conditions:

- (a) There exists $z_0 \in A$ satisfying one of the following conditions:
 - i. We have $a = f(z_0)$ and $x_1 = g(z_0)$.
 - ii. We have $a = g(z_0)$ and $x_1 = f(z_0)$.
- (b) For each $1 \le i \le n-1$, there exists $z_i \in A$ satisfying one of the following conditions:
 - i. We have $x_i = f(z_i)$ and $x_{i+1} = g(z_i)$.
 - ii. We have $x_i = g(z_i)$ and $x_{i+1} = f(z_i)$.
- (c) There exists $z_n \in A$ satisfying one of the following conditions:
 - i. We have $x_n = f(z_n)$ and $b = g(z_n)$.
 - ii. We have $x_n = g(z_n)$ and $b = f(z_n)$.

EXAMPLE 4.2.5.1.5 ► **EXAMPLES OF COEQUALISERS OF SETS**

Here are some examples of coequalisers of sets.

1. *Quotients by Equivalence Relations*. Let *R* be an equivalence relation on a set *X*. We have a bijection of sets

$$X/\sim_R \cong \mathsf{CoEq}(R \hookrightarrow X \times X \overset{\mathsf{pr}_1}{\underset{\mathsf{pr}_2}{\Longrightarrow}} X).$$

PROOF 4.2.5.1.6 ► PROOF OF EXAMPLE 4.2.5.1.5

Item 1: Quotients by Equivalence Relations

See [Pro25z].

PROPOSITION 4.2.5.1.7 ▶ PROPERTIES OF COEQUALISERS OF SETS

Let A, B, and C be sets.

1. Associativity. We have isomorphisms of sets¹

$$\underbrace{\mathsf{CoEq}(\mathsf{coeq}(f,g) \circ f, \mathsf{coeq}(f,g) \circ h)}_{=\mathsf{CoEq}(\mathsf{coeq}(f,g) \circ g, \mathsf{coeq}(f,g) \circ h)} \cong \mathsf{CoEq}(f,g,h) \cong \underbrace{\mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ g)}_{=\mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ h)}$$

where CoEq(f, g, h) is the colimit of the diagram

$$A \xrightarrow{f \atop -g \xrightarrow{h}} B$$

in Sets.

2. Unitality. We have an isomorphism of sets

$$CoEq(f, f) \cong B$$
.

3. Commutativity. We have an isomorphism of sets

$$CoEq(f, g) \cong CoEq(g, f)$$
.

4. Interaction With Composition. Let

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C$$

be functions. We have a surjection

$$CoEq(h \circ f, k \circ g) \twoheadrightarrow CoEq(coeq(h, k) \circ h \circ f, coeq(h, k) \circ k \circ g)$$

exhibiting $CoEq(coeq(h, k) \circ h \circ f, coeq(h, k) \circ k \circ g)$ as a quotient of $CoEq(h \circ f, k \circ g)$ by the relation generated by declaring $h(y) \sim k(y)$ for each $y \in B$.

(a) Take the coequaliser of (f, g, h), i.e. the colimit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets.

¹That is, the following three ways of forming "the" coequaliser of (f, g, h) agree:

(b) First take the coequaliser of f and g, forming a diagram

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{\operatorname{coeq}(f,g)}{\Longrightarrow} \operatorname{CoEq}(f,g)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\underset{h}{\Longrightarrow}} B \stackrel{\text{coeq}(f,g)}{\Longrightarrow} \text{CoEq}(f,g),$$

obtaining a quotient

$$\label{eq:coeq} {\sf CoEq}({\sf coeq}(f,g)\circ f, {\sf coeq}(f,g)\circ h) = {\sf CoEq}({\sf coeq}(f,g)\circ g, {\sf coeq}(f,g)\circ h)$$
 of ${\sf CoEq}(f,g)$

(c) First take the coequaliser of g and h, forming a diagram

$$A \stackrel{g}{\underset{h}{\Longrightarrow}} B \stackrel{\operatorname{coeq}(g,h)}{\Longrightarrow} \operatorname{CoEq}(g,h)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\underset{g}{\Longrightarrow}} B \stackrel{\text{coeq}(g,h)}{\Longrightarrow} \text{CoEq}(g,h),$$

obtaining a quotient

$$\label{eq:coeq} \begin{split} \mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ g) &= \mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ h) \\ \mathsf{of} \ \mathsf{CoEq}(g,h). \end{split}$$

PROOF 4.2.5.1.8 ► PROOF OF PROPOSITION 4.2.5.1.7

Item 1: Associativity

Omitted.

Item 2: Unitality

Omitted.

Item 3: Commutativity

Omitted.

Item 4: Interaction With Composition

Omitted.

4.2.6 Direct Colimits

Let $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I} \colon (I, \preceq) \to \pi$ be a direct system of sets.

DEFINITION 4.2.6.1.1 ► **DIRECT COLIMITS OF SETS**

The **direct colimit of** $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the direct colimit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ in Sets as in Limits and Colimits, ??.

CONSTRUCTION 4.2.6.1.2 ➤ CONSTRUCTION OF DIRECT COLIMITS OF SETS

Concretely, the direct colimit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the pair $\left(\begin{array}{c} \operatorname{colim}(X_{\alpha}), \\ \overrightarrow{\inf_{\alpha\in I}} \end{array} \right)$ consisting of:

1. *The Colimit*. The set $\underset{\longrightarrow}{\operatorname{colim}}(X_{\alpha})$ defined by

$$\operatorname{colim}_{\underset{\alpha \in I}{\longrightarrow}} (X_{\alpha}) \stackrel{\text{def}}{=} \left(\left[\prod_{\alpha \in I} X_{\alpha} \right] \right) / \sim,$$

where \sim is the equivalence relation on $\coprod_{\alpha \in I} X_{\alpha}$ generated by declaring $(\alpha, x) \sim (\beta, y)$ iff there exists some $\gamma \in I$ satisfying the following conditions:

- (a) We have $\alpha \leq \gamma$.
- (b) We have $\beta \leq \gamma$.
- (c) We have $f_{\alpha \gamma}(x) = f_{\beta \gamma}(y)$.
- 2. The Cocone. The collection

$$\left\{\inf_{\gamma}\colon X_{\gamma}\to \underset{\alpha\in I}{\operatorname{colim}}(X_{\alpha})\right\}_{\gamma\in I}$$

of maps of sets defined by

$$\operatorname{inj}_{\gamma}(x) \stackrel{\text{def}}{=} [(\gamma, x)]$$

for each $y \in I$ and each $x \in X_y$.

PROOF 4.2.6.1.3 ► PROOF OF CONSTRUCTION 4.2.6.1.2

We will prove Construction 4.2.6.1.2 below in a bit, but first we need a lemma (which is interesting in its own right).

LEMMA 4.2.6.1.4 \blacktriangleright Identification of x with $f_{\alpha\beta}(x)$ in Direct Colimits

For each $\alpha, \beta \in I$ and each $x \in X_{\alpha}$, if $\alpha \leq \beta$, then we have

$$(\alpha, x) \sim (\beta, f_{\alpha\beta}(x))$$

in $\underset{\alpha \in I}{\overset{\longrightarrow}{colim}}(X_{\alpha})$.

PROOF 4.2.6.1.5 ► PROOF OF LEMMA 4.2.6.1.4

Taking $\gamma = \beta$, we have $f_{\alpha\gamma} = f_{\alpha\beta}$, we have $f_{\beta\gamma} = f_{\beta\beta} \stackrel{\text{def}}{=} \mathrm{id}_{X_{\beta}}$, and we have

$$f_{\alpha\beta}(x) = f_{\beta\beta}(f_{\alpha\beta}(x))$$

$$\stackrel{\text{def}}{=} \mathrm{id}_{X_{\beta}}(f_{\alpha\beta}(x)),$$

$$= f_{\alpha\beta}(x).$$

As a result, since $\alpha \leq \beta$ and $\beta \leq \beta$ as well, Items 1a to 1c of Construction 4.2.6.1.2 are met. Thus we have $(\alpha, x) \sim (\beta, f_{\alpha\beta}(x))$.

We can now prove Construction 4.2.6.1.2:

PROOF 4.2.6.1.6 ► PROOF OF CONSTRUCTION 4.2.6.1.2

We claim that $\operatornamewithlimits{colim}_{\alpha \in I}(X_{\alpha})$ is the colimit of the direct system of sets $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta \in I}$.

Commutativity of the Colimit Diagram

First, we need to check that the colimit diagram defined by $\underset{\alpha}{\text{colim}}(X_{\alpha})$

commutes, i.e. that we have

for each $\alpha, \beta \in I$ with $\alpha \leq \beta$. Indeed, given $x \in X_{\alpha}$, we have

$$[\inf_{\beta} \circ f_{\alpha\beta}](x) \stackrel{\text{def}}{=} \inf_{\beta} (f_{\alpha\beta}(x))$$

$$\stackrel{\text{def}}{=} [(\beta, f_{\alpha\beta}(x))]$$

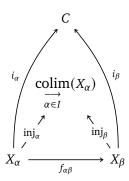
$$= [(\alpha, x)]$$

$$\stackrel{\text{def}}{=} \inf_{\alpha} (x),$$

where we have used Lemma 4.2.6.1.4 for the third equality.

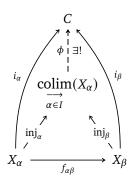
Proof of the Universal Property of the Colimit

Next, we prove that $\operatornamewithlimits{colim}_{\alpha \in I}(X_{\alpha})$ as constructed in Construction 4.2.6.1.2 satisfies the universal property of a direct colimit. Suppose that we have, for each $\alpha, \beta \in I$ with $\alpha \leq \beta$, a diagram of the form



in Sets. We claim that there exists a unique map ϕ : $\operatornamewithlimits{colim}_{\longrightarrow}(X_\alpha) \stackrel{\exists !}{\longrightarrow} C$ making

the diagram



commute. To this end, first consider the diagram

$$\coprod_{\alpha \in I} X_{\alpha} \xrightarrow{\operatorname{pr}} \operatorname{colim}_{\alpha \in I} (X_{\alpha})$$

$$\coprod_{\alpha \in I} i_{\alpha} \qquad \qquad C.$$

Lemma. If $(\alpha, x) \sim (\beta, y)$, then we have

$$\left[\bigsqcup_{\alpha \in I} i_{\alpha} \right](x) = \left[\bigsqcup_{\alpha \in I} i_{\alpha} \right](y).$$

Proof. Indeed, if $(\alpha, x) \sim (\beta, y)$, then there exists some $\gamma \in I$ satisfying the following conditions:

- 1. We have $\alpha \leq \gamma$.
- 2. We have $\beta \leq \gamma$.
- 3. We have $f_{\alpha\gamma}(x) = f_{\beta\gamma}(y)$.

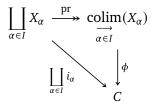
We then have

$$\left[\bigsqcup_{\alpha \in I} i_{\alpha} \right](x) \stackrel{\text{def}}{=} i_{\alpha}(x)$$

$$\stackrel{\text{def}}{=} [i_{\gamma} \circ f_{\alpha\gamma}](x)$$

$$\stackrel{\text{def}}{=} i_{\gamma}(f_{\alpha\gamma}(x))
= i_{\gamma}(f_{\beta\gamma}(x))
\stackrel{\text{def}}{=} [i_{\gamma} \circ f_{\beta\gamma}](x)
= i_{\beta}(y)
\stackrel{\text{def}}{=} \left[\prod_{\alpha \in I} i_{\alpha} \right](y).$$

This finishes the proof of the lemma. Continuing, by Conditions on Relations, ?? of Proposition 10.6.2.1.3, there then exists a map $\phi: \underset{\alpha \in I}{\operatorname{colim}}(X_{\alpha}) \xrightarrow{\exists !} C$ making the diagram

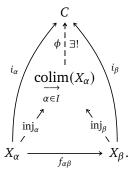


commute. In particular, this implies that the diagram

$$X_{\alpha} \xrightarrow{\operatorname{inj}_{\alpha}} \operatorname{colim}_{\alpha \in I}(X_{\alpha})$$

$$\downarrow_{\alpha} \qquad \downarrow_{\phi}$$

also commutes, and thus so does the diagram



This finishes the proof.¹



¹Incidentally, the conditions

$$\left\{i_{\alpha} = \phi \circ \operatorname{inj}_{\alpha}\right\}_{\alpha \in I}$$

show that ϕ must be given by

$$\phi([(\alpha, x)]) = (i_{\alpha}(x))_{\alpha \in I}$$

for each $[(\alpha,x)] \in \operatornamewithlimits{colim}_{\alpha \in I}(X_\alpha)$, although we would need to show that this assignment is well-defined were we to prove Construction 4.2.6.1.2 in this way. Instead, invoking Conditions on Relations, ?? of Proposition 10.6.2.1.3 gave us a way to avoid having to prove this, leading to a cleaner alternative proof.

EXAMPLE 4.2.6.1.7 ► **EXAMPLES OF DIRECT COLIMITS OF SETS**

Here are some examples of direct colimits of sets.

1. *The Prüfer Group*. The Prüfer group $\mathbb{Z}(p^{\infty})$ is defined as the direct colimit

$$\mathbb{Z}(p^{\infty}) \stackrel{\text{def}}{=} \underset{n \in \mathbb{N}}{\operatorname{colim}}(\mathbb{Z}_{/p^n});$$

see ??.

4.3 Operations With Sets

4.3.1 The Empty Set

DEFINITION 4.3.1.1.1 ► THE EMPTY SET

The **empty set** is the set Ø defined by

$$\emptyset \stackrel{\text{def}}{=} \{ x \in X \mid x \neq x \},\$$

where X is the set in the set existence axiom, ?? of ??.

4.3.2 Singleton Sets

Let *X* be a set.

DEFINITION 4.3.2.1.1 ► SINGLETON SETS

The **singleton set containing** X is the set $\{X\}$ defined by

$$\{X\} \stackrel{\mathrm{def}}{=} \{X, X\},\,$$

where $\{X, X\}$ is the pairing of X with itself of Definition 4.3.3.1.1.

4.3.3 Pairings of Sets

Let *X* and *Y* be sets.

DEFINITION 4.3.3.1.1 ► PAIRINGS OF SETS

The **pairing of** X **and** Y is the set $\{X,Y\}$ defined by

$${X,Y} \stackrel{\text{def}}{=} {x \in A \mid x = X \text{ or } x = Y},$$

where *A* is the set in the axiom of pairing, ?? of ??.

4.3.4 Ordered Pairs

Let *A* and *B* be sets.

DEFINITION 4.3.4.1.1 ► ORDERED PAIRS

The **ordered pair associated to** A **and** B is the set (A, B) defined by

$$(A, B) \stackrel{\text{def}}{=} \{ \{A\}, \{A, B\} \}.$$

PROPOSITION 4.3.4.1.2 ► PROPERTIES OF ORDERED PAIRS

Let A and B be sets.

- 1. *Uniqueness*. Let *A*, *B*, *C*, and *D* be sets. The following conditions are equivalent:
 - (a) We have (A, B) = (C, D).
 - (b) We have A = C and B = D.

PROOF 4.3.4.1.3 ► PROOF OF PROPOSITION 4.3.4.1.2

Item 1: Uniqueness

See [Cie97, Theorem 1.2.3].

4.3.5 Sets of Maps

Let *A* and *B* be sets.

DEFINITION 4.3.5.1.1 ► **SETS OF MAPS**

The **set of maps from** A **to** B^1 is the set $Sets(A, B)^2$ whose elements are the functions from A to B.

PROPOSITION 4.3.5.1.2 ► PROPERTIES OF SETS OF MAPS

Let A and B be sets.

1. *Functoriality*. The assignments $X, Y, (X, Y) \mapsto \text{Hom}_{\mathsf{Sets}}(X, Y)$ define functors

$$\begin{array}{lll} \mathsf{Sets}(X,-)\colon & \mathsf{Sets} & \to \mathsf{Sets}, \\ \mathsf{Sets}(-,Y)\colon & \mathsf{Sets}^\mathsf{op} & \to \mathsf{Sets}, \\ \mathsf{Sets}(-_1,-_2)\colon & \mathsf{Sets}^\mathsf{op} \times \mathsf{Sets} \to \mathsf{Sets}. \end{array}$$

2. Adjointness. We have adjunctions

$$(A \times - \dashv \mathsf{Sets}(A, -))$$
: Sets $\underbrace{ \xrightarrow{A \times -}}_{\mathsf{Sets}(A, -)} \mathsf{Sets},$

$$(- \times B + \operatorname{Sets}(B, -))$$
: Sets $\underbrace{- \times B}_{\text{Sets}(B, -)}$ Sets,

witnessed by bijections

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$

 $Sets(A \times B, C) \cong Sets(B, Sets(A, C)),$

natural in $A, B, C \in Obj(Sets)$.

¹Further Terminology: Also called the **Hom set from** A **to** B.

²Further Notation: Also written $Hom_{Sets}(A, B)$.

3. Maps From the Punctual Set. We have a bijection

$$\mathsf{Sets}(\mathsf{pt},A) \cong A$$
,

natural in $A \in Obj(Sets)$.

4. Maps to the Punctual Set. We have a bijection

$$Sets(A, pt) \cong pt$$
,

natural in $A \in Obj(Sets)$.

PROOF 4.3.5.1.3 ► PROOF OF PROPOSITION 4.3.5.1.2

Item 1: Functoriality

This follows from Categories, Items 2 and 5 of Proposition 11.1.4.1.2.

Item 2: Adjointness

This is a repetition of Item 2 of Proposition 4.1.3.1.4 and is proved there.

Item 3: Maps From the Punctual Set

The bijection

$$\Phi_A \colon \mathsf{Sets}(\mathsf{pt},A) \xrightarrow{\sim} A$$

is given by

$$\Phi_A(f) \stackrel{\text{def}}{=} f(\star)$$

for each $f \in Sets(pt, A)$, admitting an inverse

$$\Phi_A^{-1} : A \xrightarrow{\sim} \mathsf{Sets}(\mathsf{pt}, A)$$

given by

$$\Phi_A^{-1}(a) \stackrel{\mathrm{def}}{=} \llbracket \star \mapsto a \rrbracket$$

for each $a \in A$. Indeed, we have

$$\begin{split} [\Phi_A^{-1} \circ \Phi_A](f) &\stackrel{\text{def}}{=} \Phi_A^{-1}(\Phi_A(f)) \\ &\stackrel{\text{def}}{=} \Phi_A^{-1}(f(\star)) \\ &\stackrel{\text{def}}{=} \llbracket \star \mapsto f(\star) \rrbracket \\ &\stackrel{\text{def}}{=} f \end{split}$$

$$\stackrel{\text{def}}{=} [\mathrm{id}_{\mathsf{Sets}(\mathsf{pt},A)}](f)$$

for each $f \in Sets(pt, A)$ and

$$\begin{split} [\Phi_A \circ \Phi_A^{-1}](a) &\stackrel{\text{def}}{=} \Phi_A(\Phi_A^{-1}(a)) \\ &\stackrel{\text{def}}{=} \Phi_A([\![\star \mapsto a]\!]) \\ &\stackrel{\text{def}}{=} \operatorname{ev}_{\star}([\![\star \mapsto a]\!]) \\ &\stackrel{\text{def}}{=} a \\ &\stackrel{\text{def}}{=} [\operatorname{id}_A](a) \end{split}$$

for each $a \in A$, and thus we have

$$\Phi_A^{-1} \circ \Phi_A = \mathrm{id}_{\mathsf{Sets}(\mathsf{pt},A)}$$

$$\Phi_A \circ \Phi_A^{-1} = \mathrm{id}_A.$$

To prove naturality, we need to show that the diagram

$$\begin{array}{ccc}
\operatorname{Sets}(\operatorname{pt},A) & \xrightarrow{f_!} & \operatorname{Sets}(\operatorname{pt},B) \\
& & \downarrow \\
\Phi_A & \downarrow \downarrow \\
A & \xrightarrow{f} & B
\end{array}$$

commutes. Indeed, we have

$$[f \circ \Phi_A](\phi) \stackrel{\text{def}}{=} f(\Phi_A(\phi))$$

$$\stackrel{\text{def}}{=} f(\phi(\star))$$

$$\stackrel{\text{def}}{=} [f \circ \phi](\star)$$

$$\stackrel{\text{def}}{=} \Phi_B(f \circ \phi)$$

$$\stackrel{\text{def}}{=} \Phi_B(f_!(\phi))$$

$$\stackrel{\text{def}}{=} [\Phi_B \circ f_!](\phi)$$

for each $\phi \in Sets(pt, A)$. This finishes the proof.

Item 4: Maps to the Punctual Set

This follows from the universal property of pt as the terminal set, Definition 4.1.1.1.

4.3.6 Unions of Families of Subsets

Let *X* be a set and let $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

DEFINITION 4.3.6.1.1 ► Unions of Families of Subsets

The **union of** \mathcal{U} is the set $\bigcup_{U \in \mathcal{U}} U$ defined by

$$\bigcup_{U \in \mathcal{U}} U \stackrel{\text{def}}{=} \left\{ x \in X \;\middle|\; \text{there exists some } U \in \mathcal{U} \right\}.$$

PROPOSITION 4.3.6.1.2 ► PROPERTIES OF UNIONS OF FAMILIES OF SUBSETS

Let *X* be a set.

1. Functoriality. The assignment $\mathcal{U} \mapsto \bigcup_{U \in \mathcal{U}} U$ defines a functor

$$\bigcup : (\mathcal{P}(\mathcal{P}(X)), \subset) \to (\mathcal{P}(X), \subset).$$

In particular, for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$, the following condition is satisfied:

$$(\star) \text{ If } \mathcal{U} \subset \mathcal{V} \text{, then } \bigcup_{U \in \mathcal{U}} U \subset \bigcup_{V \in \mathcal{V}} V.$$

2. Associativity. The diagram

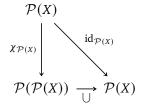
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(\mathcal{P}(X))) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \bigcup} & \mathcal{P}(\mathcal{P}(X)) \\
\cup \star \mathrm{id}_{\mathcal{P}(X)} & & & & & & & & & \\
\mathcal{P}(\mathcal{P}(X)) & & & & & & & & & & \\
\end{array}$$

commutes, i.e. we have

$$\bigcup_{U \in \bigcup_{A \in \mathcal{A}} A} U = \bigcup_{A \in \mathcal{A}} (\bigcup_{U \in A} U)$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$.

3. Left Unitality. The diagram

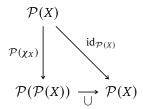


commutes, i.e. we have

$$\bigcup_{V \in \{U\}} V = U$$

for each $U \in \mathcal{P}(X)$.

4. Right Unitality. The diagram



commutes, i.e. we have

$$\bigcup_{\{u\}\in\chi_X(U)}\{u\}=U$$

for each $U \in \mathcal{P}(X)$.

5. Interaction With Unions I. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\cup} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\cup} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{W\in\mathcal{U}\cup\mathcal{V}}W=\left(\bigcup_{U\in\mathcal{U}}U\right)\cup\left(\bigcup_{V\in\mathcal{V}}V\right)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

6. Interaction With Unions II. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcup_{V \in \mathcal{V}} V\right) = \bigcup_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcup_{U \in \mathcal{U}} U\right) \cup V = \bigcup_{U \in \mathcal{U}} (U \cup V)$$

for each nonempty $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

7. *Interaction With Intersections I.* We have a natural transformation

with components

$$\bigcup_{W \in \mathcal{U} \cap \mathcal{V}} W \subset \left(\bigcup_{U \in \mathcal{U}} U\right) \cap \left(\bigcup_{V \in \mathcal{V}} V\right)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

8. Interaction With Intersections II. The diagrams

commute, i.e. we have

$$U \cap \left(\bigcup_{V \in \mathcal{V}} V\right) = \bigcup_{V \in \mathcal{V}} (U \cap V),$$

$$\left(\bigcup_{U\in\mathcal{U}}U\right)\cap V=\bigcup_{U\in\mathcal{U}}(U\cap V)$$

for each $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

9. Interaction With Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\setminus} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\setminus} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in \mathcal{U} \setminus \mathcal{V}} W \neq \left(\bigcup_{U \in \mathcal{U}} U\right) \setminus \left(\bigcup_{V \in \mathcal{V}} V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

10. Interaction With Complements I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\operatorname{op}} \xrightarrow{(-)^{c}} \mathcal{P}(\mathcal{P}(X))$$

$$\cup^{\operatorname{op}} \bigvee \qquad \qquad \bigcup \cup$$

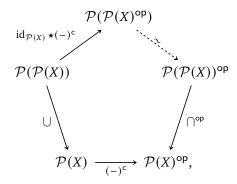
$$\mathcal{P}(X)^{\operatorname{op}} \xrightarrow{(-)^{c}} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{U\in\mathcal{U}^\mathsf{c}}U\neq\bigcup_{U\in\mathcal{U}}U^\mathsf{c}$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

11. Interaction With Complements II. The diagram

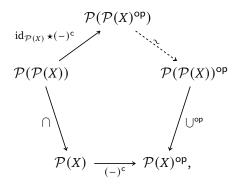


commutes, i.e. we have

$$\left(\bigcup_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

12. Interaction With Complements III. The diagram



commutes, i.e. we have

$$\left(\bigcap_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

13. Interaction With Symmetric Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\Delta} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\Delta} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W\in\mathcal{U}\triangle\mathcal{V}}W\neq\left(\bigcup_{U\in\mathcal{U}}U\right)\triangle\left(\bigcup_{V\in\mathcal{V}}V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

14. Interaction With Internal Homs I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\downarrow^{\mathsf{op}} \times \cup^{\mathsf{op}} \qquad \qquad \downarrow \cup$$

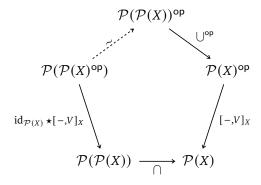
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcup_{U \in \mathcal{U}} U, \bigcup_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

15. Interaction With Internal Homs II. The diagram



commutes, i.e. we have

$$\left[\bigcup_{U\in\mathcal{U}}U,V\right]_X=\bigcap_{U\in\mathcal{U}}[U,V]_X$$

for each $U \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

16. Interaction With Internal Homs III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcup}{\longrightarrow} \mathcal{P}(X) \\
id_{\mathcal{P}(X)} \star [U,-]_X & & \downarrow [U,-]_X \\
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcup}{\longrightarrow} \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U, \bigcup_{V \in \mathcal{V}} V\right]_X = \bigcup_{V \in \mathcal{V}} [U, V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

17. *Interaction With Direct Images.* Let $f: X \to Y$ be a map of sets. The diagram

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_!(U) = \bigcup_{V \in f_!(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

18. Interaction With Inverse Images. Let $f: X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \qquad \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{V \in \mathcal{V}} f^{-1}(V) = \bigcup_{U \in f^{-1}(\mathcal{U})} U$$

for each $V \in \mathcal{P}(Y)$, where $f^{-1}(V) \stackrel{\text{def}}{=} (f^{-1})^{-1}(V)$.

19. *Interaction With Codirect Images.* Let $f: X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \bigcup$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_*(U) = \bigcup_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

20. Interaction With Intersections of Families I. The diagram

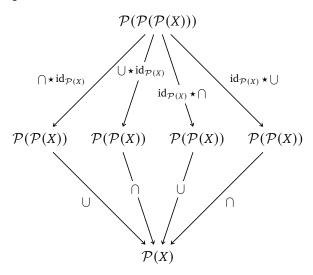
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \cap} & \mathcal{P}(\mathcal{P}(X)) \\
\downarrow \star \mathrm{id}_{\mathcal{P}(X)} & & & \downarrow \cap \\
\mathcal{P}(X) & \xrightarrow{\qquad \cap} & X
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U = \bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right)$$

for each $A \in \mathcal{P}(\mathcal{P}(X))$.

21. *Interaction With Intersections of Families II.* Let *X* be a set and consider the compositions

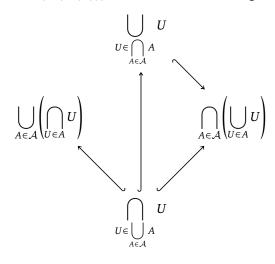


given by

$$A \mapsto \bigcup_{U \in \bigcap_{A \in \mathcal{A}} A} U, \quad A \mapsto \bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U,$$

$$A \mapsto \bigcup_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U\right), \quad A \mapsto \bigcap_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U\right)$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$. We have the following inclusions:



All other possible inclusions fail to hold in general.

PROOF 4.3.6.1.3 ► PROOF OF PROPOSITION 4.3.6.1.2

Item 1: Functoriality

Since $\mathcal{P}(X)$ is posetal, it suffices to prove the condition (\star) . So let $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ with $\mathcal{U} \subset \mathcal{V}$. We claim that

$$\bigcup_{U\in\mathcal{U}}U\subset\bigcup_{V\in\mathcal{V}}V.$$

Indeed, given $x \in \bigcup_{U \in \mathcal{U}} U$, there exists some $U \in \mathcal{U}$ such that $x \in U$, but since $\mathcal{U} \subset \mathcal{V}$, we have $U \in \mathcal{V}$ as well, and thus $x \in \bigcup_{V \in \mathcal{V}} V$, which gives our desired inclusion.

Item 2: Associativity

We have

$$\bigcup_{U \in \bigcup_{A \in \mathcal{A}} A} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } U \in \bigcup_{A \in \mathcal{A}} A \right\}$$

$$= \left\{ x \in X \middle| \text{ there exists some } A \in \mathcal{A} \right\}$$

$$= \left\{ x \in X \middle| \text{ there exists some } A \in \mathcal{A} \right\}$$

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$$= \left\{ x \in X \middle| \text{ there exists some } A \in \mathcal{A} \right\}$$

This finishes the proof.

Item 3: Left Unitality

$$\bigcup_{V \in \{U\}} V \stackrel{\text{def}}{=} \left\{ x \in X \;\middle|\; \text{there exists some } V \in \{U\} \right\}$$
 such that we have $x \in U$

$$= \{x \in X \mid x \in U\}$$
$$= U.$$

Item 4: Right Unitality

We have

$$\bigcup_{\{u\} \in \chi_X(U)} \{u\} \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } \{u\} \in \chi_X(U) \right\}$$
 such that we have $x \in \{u\}$

$$= \left\{ x \in X \middle| \text{ there exists some } \{u\} \in \chi_X(U) \right\}$$
 such that we have $x = u$

$$= \left\{ x \in X \middle| \text{ there exists some } u \in U \right\}$$
 such that we have $x = u$

$$= \left\{ x \in X \middle| \text{ there exists some } u \in U \right\}$$
 such that we have $x = u$

$$= \left\{ x \in X \middle| x \in U \right\}$$

$$= U.$$

This finishes the proof.

Item 5: Interaction With Unions I

$$\bigcup_{W \in \mathcal{U} \cup \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{U} \cup \mathcal{V} \right\} \\
= \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{U} \text{ or some} \right\} \\
\stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{U} \text{ or some} \right\} \\
\downarrow \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{U} \right\} \\
\downarrow \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{V} \right\} \\
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Item 6: Interaction With Unions II

Assume V is nonempty. We have

$$U \cup \bigcup_{V \in \mathcal{V}} V \stackrel{\text{def}}{=} \left\{ x \in X \mid x \in U \text{ or } x \in \bigcup_{V \in \mathcal{V}} V \right\}$$

$$= \left\{ x \in X \mid x \in U \text{ or there exists some } V \in \mathcal{V} \text{ such that } x \in V \right\}$$

$$= \left\{ x \in X \mid \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \mid \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \mid \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \mid \text{there exists some } V \in \mathcal{V} \right\}$$

$$\stackrel{\text{def}}{=} \bigcup_{V \in \mathcal{V}} U \cup V.$$

This concludes the proof of the first statement. For the second statement, use Item 4 of Proposition 4.3.8.1.2 to rewrite

$$\left(\bigcup_{U\in\mathcal{U}}U\right)\cup V=V\cup\left(\bigcup_{U\in\mathcal{U}}U\right),$$

$$\bigcup_{U\in\mathcal{U}}(U\cup V)=\bigcup_{U\in\mathcal{U}}(V\cup U).$$

But these two sets are equal by the first statement.

Item 7: Interaction With Intersections I

$$\bigcup_{W \in \mathcal{U} \cap \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } W \in \mathcal{U} \cap \mathcal{V} \\ \text{such that we have } x \in W \end{array} \right\}$$

$$\subset \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \mathcal{U} \text{ and some } V \in \mathcal{V} \\ \text{such that we have } x \in U \text{ and } x \in V \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \mathcal{U} \\ \text{such that we have } x \in U \end{array} \right\}$$

$$\bigcup \left\{ x \in X \middle| \text{ there exists some } V \in \mathcal{V} \right\} \\
\text{such that we have } x \in V \right\} \\
\stackrel{\text{def}}{=} \left(\bigcup_{U \in \mathcal{U}} U \right) \cap \left(\bigcup_{V \in \mathcal{V}} V \right).$$

Item 8: Interaction With Intersections II

We have

$$U \cap \bigcup_{V \in \mathcal{V}} V \stackrel{\text{def}}{=} \left\{ x \in X \mid x \in U \text{ and } x \in \bigcup_{V \in \mathcal{V}} V \right\}$$

$$= \left\{ x \in X \mid x \in U \text{ and there exists some} \right\}$$

$$= \left\{ x \in X \mid \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \mid \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \mid \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \mid \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ u \in X \mid \text{such that } u \in U \cap V \right\}$$

$$= \left\{ u \in X \mid \text{there exists some } v \in \mathcal{V} \right\}$$

This concludes the proof of the first statement. For the second statement, use Item 5 of Proposition 4.3.9.1.2 to rewrite

$$\left(\bigcup_{U\in\mathcal{U}}U\right)\cap V=V\cap\left(\bigcup_{U\in\mathcal{U}}U\right),$$
$$\bigcup_{U\in\mathcal{U}}(U\cap V)=\bigcup_{U\in\mathcal{U}}(V\cap U).$$

But these two sets are equal by the first statement.

Item 9: Interaction With Differences

Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}\$, and let $\mathcal{V} = \{\{0\}\}\$. We have

$$\bigcup_{W\in\mathcal{U}\backslash\mathcal{V}}U=\bigcup_{W\in\{\{0,1\}\}}W$$

$$= \{0, 1\},\$$

whereas

$$\left(\bigcup_{U \in \mathcal{U}} U\right) \setminus \left(\bigcup_{V \in \mathcal{V}} V\right) = \{0, 1\} \setminus \{0\}$$

$$= \{1\}$$

Thus we have

$$\bigcup_{W\in\mathcal{U}\backslash\mathcal{V}}W=\left\{0,1\right\}\neq\left\{1\right\}=\left(\bigcup_{U\in\mathcal{U}}U\right)\backslash\left(\bigcup_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

Item 10: Interaction With Complements I

Let $X = \{0, 1\}$ and let $\mathcal{U} = \{0\}$. We have

$$\bigcup_{U \in \mathcal{U}^{c}} U = \bigcup_{U \in \{\emptyset, \{1\}, \{0,1\}\}} U$$
$$= \{0, 1\},$$

whereas

$$\bigcup_{U \in \mathcal{U}} U^{c} = \{0\}^{c}$$
$$= \{1\}.$$

Thus we have

$$\bigcup_{U\in\mathcal{U}^\mathsf{c}}U=\left\{0,1\right\}\neq\left\{1\right\}=\bigcup_{U\in\mathcal{U}}U^\mathsf{c}.$$

This finishes the proof.

Item 11: Interaction With Complements II

$$\left(\bigcup_{U \in \mathcal{U}} U\right)^{\mathsf{c}} \stackrel{\text{def}}{=} \left\{ x \in X \mid \text{ there exists no } U \in \mathcal{U} \\ \text{ such that we have } x \in U \right\}$$

$$= \begin{cases} x \in X & \text{for all } U \in \mathcal{U} \\ \text{we have } x \notin U \end{cases}$$

$$\stackrel{\text{def}}{=} \begin{cases} x \in X & \text{for all } U \in \mathcal{U} \\ \text{we have } x \in U^{c} \end{cases}$$

$$\stackrel{\text{def}}{=} \bigcap_{U \in \mathcal{U}} U^{c}.$$

Item 12: Interaction With Complements III

By Item 11 Item 3 of Proposition 4.3.11.1.2, we have

$$\left(\bigcap_{U\in\mathcal{U}}U\right)^{c} = \left(\bigcap_{U\in\mathcal{U}}(U^{c})^{c}\right)^{c}$$

$$= \left(\left(\bigcup_{U\in\mathcal{U}}U^{c}\right)^{c}\right)^{c}$$

$$= \bigcup_{U\in\mathcal{U}}U^{c}.$$

Item 13: Interaction With Symmetric Differences

Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}\$, and let $\mathcal{V} = \{\{0\}, \{0, 1\}\}\$. We have

$$\bigcup_{W \in \mathcal{U} \triangle \mathcal{V}} W = \bigcup_{W \in \{\{0\}\}} W$$
$$= \{0\},$$

whereas

$$\left(\bigcup_{U \in \mathcal{U}} U\right) \triangle \left(\bigcup_{V \in \mathcal{V}} V\right) = \{0, 1\} \triangle \{0, 1\}$$
$$= \emptyset,$$

Thus we have

$$\bigcup_{W\in\mathcal{U}\triangle\mathcal{V}}W=\left\{0\right\}\neq\emptyset=\left(\bigcup_{U\in\mathcal{U}}U\right)\triangle\left(\bigcup_{V\in\mathcal{V}}V\right).$$

Item 14: Interaction With Internal Homs I

This is a repetition of Item 7 of Proposition 4.4.7.1.4 and is proved there.

Item 15: Interaction With Internal Homs II

This is a repetition of Item 8 of Proposition 4.4.7.1.4 and is proved there.

Item 16: Interaction With Internal Homs III

This is a repetition of Item 9 of Proposition 4.4.7.1.4 and is proved there.

Item 17: Interaction With Direct Images

This is a repetition of Item 3 of Proposition 4.6.1.1.5 and is proved there.

Item 18: Interaction With Inverse Images

This is a repetition of Item 3 of Proposition 4.6.2.1.3 and is proved there.

Item 19: Interaction With Codirect Images

This is a repetition of Item 3 of Proposition 4.6.3.1.7 and is proved there.

Item 20: Interaction With Intersections of Families I

We have

$$\bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \bigcup_{A \in \mathcal{A}} A, \\ \text{we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } A \in \mathcal{A} \text{ and each } \\ U \in A, \text{ we have } x \in U \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right).$$

This finishes the proof.

Item 21: Interaction With Intersections of Families II

Omitted.

4.3.7 Intersections of Families of Subsets

Let *X* be a set and let $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

DEFINITION 4.3.7.1.1 ► INTERSECTIONS OF FAMILIES OF SUBSETS

The **intersection of** \mathcal{U} is the set $\bigcap_{U \in \mathcal{U}} U$ defined by

$$\bigcap_{U \in \mathcal{U}} U \stackrel{\text{def}}{=} \left\{ x \in X \, \middle| \, \text{for each } U \in \mathcal{U}, \right\}.$$
we have $x \in U$

PROPOSITION 4.3.7.1.2 ► PROPERTIES OF INTERSECTIONS OF FAMILIES OF SUBSETS

Let *X* be a set.

1. Functoriality. The assignment $\mathcal{U} \mapsto \bigcap_{U \in \mathcal{U}} U$ defines a functor

$$\bigcap \colon (\mathcal{P}(\mathcal{P}(X)),\supset) \to (\mathcal{P}(X),\subset).$$

In particular, for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$, the following condition is satisfied:

$$(\star) \ \ \text{If} \ \mathcal{U} \subset \mathcal{V} \text{, then} \bigcap_{V \in \mathcal{V}} V \subset \bigcap_{U \in \mathcal{U}} U.$$

2. Oplax Associativity. We have a natural transformation

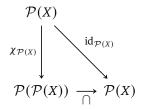
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(\mathcal{P}(X))) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \cap} & \mathcal{P}(\mathcal{P}(X)) \\
\cap \star \mathrm{id}_{\mathcal{P}(X)} & & & & & & & & & \\
\mathcal{P}(\mathcal{P}(X)) & & & & & & & & & \\
\end{array}$$

with components

$$\bigcap_{A\in\mathcal{A}} \left(\bigcap_{U\in A} U\right) \subset \bigcap_{U\in \bigcap_{A\in\mathcal{A}} A} U$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$.

3. Left Unitality. The diagram



commutes, i.e. we have

$$\bigcap_{V\in\{U\}}V=U.$$

for each $U \in \mathcal{P}(X)$.

4. Oplax Right Unitality. The diagram

$$\begin{array}{c|c}
\mathcal{P}(X) & \operatorname{id}_{\mathcal{P}(X)} \\
 & \times & \\
\mathcal{P}(\mathcal{P}(X)) & \longrightarrow & \mathcal{P}(X)
\end{array}$$

does not commute in general, i.e. we may have

$$\bigcap_{\{x\}\in\chi_X(U)}\{x\}\neq U$$

in general, where $U \in \mathcal{P}(X)$. However, when U is nonempty, we have

$$\bigcap_{\{x\}\in\chi_X(U)}\{x\}\subset U.$$

5. Interaction With Unions I. The diagram

$$\begin{array}{cccc} \mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) & \stackrel{\cup}{\longrightarrow} & \mathcal{P}(\mathcal{P}(X)) \\ & & & \downarrow & & \downarrow \\ & \mathcal{P}(X) \times \mathcal{P}(X) & \stackrel{\cap}{\longrightarrow} & \mathcal{P}(X) \end{array}$$

commutes, i.e. we have

$$\bigcap_{W\in\mathcal{U}\cup\mathcal{V}}W=\left(\bigcap_{U\in\mathcal{U}}U\right)\cap\left(\bigcap_{V\in\mathcal{V}}V\right)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

6. Interaction With Unions II. The diagram

commute, i.e. we have

$$U \cup \left(\bigcap_{V \in \mathcal{V}} V\right) = \bigcap_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cup V = \bigcap_{U \in \mathcal{U}} (U \cup V)$$

for each $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

7. Interaction With Intersections I. We have a natural transformation

with components

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\cap\left(\bigcap_{V\in\mathcal{V}}V\right)\subset\bigcap_{W\in\mathcal{U}\cap\mathcal{V}}W$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

8. Interaction With Intersections II. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcap_{V \in \mathcal{V}} V\right) = \bigcap_{V \in \mathcal{V}} (U \cup V),$$

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\cup V=\bigcap_{U\in\mathcal{U}}(U\cup V)$$

for each $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

9. Interaction With Differences. The diagram

$$\begin{array}{cccc} \mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) & \stackrel{\backslash}{\longrightarrow} & \mathcal{P}(\mathcal{P}(X)) \\ & & & \swarrow & & \downarrow \cap \\ & \mathcal{P}(X) \times \mathcal{P}(X) & \stackrel{\backslash}{\longrightarrow} & \mathcal{P}(X), \end{array}$$

does not commute in general, i.e. we may have

$$\bigcap_{W\in\mathcal{U}\setminus\mathcal{V}}W\neq\left(\bigcap_{U\in\mathcal{U}}U\right)\setminus\left(\bigcap_{V\in\mathcal{V}}V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

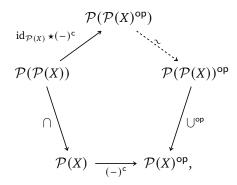
10. Interaction With Complements I. The diagram

does not commute in general, i.e. we may have

$$\bigcap_{W\in\mathcal{U}^{\mathsf{c}}}W\neq\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

11. Interaction With Complements II. The diagram

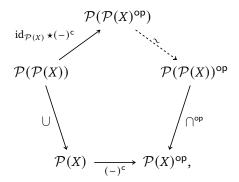


commutes, i.e. we have

$$\left(\bigcap_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

12. Interaction With Complements III. The diagram



commutes, i.e. we have

$$\left(\bigcup_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

13. Interaction With Symmetric Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\triangle} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \times \cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\triangle} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W \in \mathcal{U} \triangle \mathcal{V}} W \neq \left(\bigcap_{U \in \mathcal{U}} U\right) \triangle \left(\bigcap_{V \in \mathcal{V}} V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

14. Interaction With Internal Homs I. The diagram

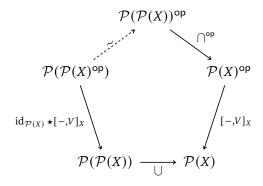
$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X)) \\
\cap^{\mathsf{op}} \times \cap^{\mathsf{op}} \downarrow \qquad \qquad \downarrow \cap \\
\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W\in [\mathcal{U},\mathcal{V}]_{\mathcal{P}(X)}}W\neq\left[\bigcap_{U\in\mathcal{U}}U,\bigcap_{V\in\mathcal{V}}V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

15. Interaction With Internal Homs II. The diagram



commutes, i.e. we have

$$\left[\bigcap_{U\in\mathcal{U}}U,V\right]_X=\bigcup_{U\in\mathcal{U}}[U,V]_X$$

for each $U \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

16. Interaction With Internal Homs III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcap}{\longrightarrow} & \mathcal{P}(X) \\
id_{\mathcal{P}(X)} \star [U,-]_X & & \downarrow [U,-]_X \\
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcap}{\longrightarrow} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U,\bigcap_{V\in\mathcal{V}}V\right]_X=\bigcap_{V\in\mathcal{V}}[U,V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

17. *Interaction With Direct Images.* Let $f: X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_!)_!} \mathcal{P}(\mathcal{P}(Y))$$

$$\cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcap_{U\in\mathcal{U}}f_!(U)=\bigcap_{V\in f_!(\mathcal{U})}V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

18. *Interaction With Inverse Images.* Let $f: X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \bigcup_{f^{-1}} \bigcup_{f^{-1}} \bigcap_{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcap_{V \in \mathcal{V}} f^{-1}(V) = \bigcap_{U \in f^{-1}(\mathcal{U})} U$$

for each $V \in \mathcal{P}(Y)$, where $f^{-1}(V) \stackrel{\text{def}}{=} (f^{-1})^{-1}(V)$.

19. *Interaction With Codirect Images.* Let $f: X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcap_{U\in\mathcal{U}} f_*(U) = \bigcap_{V\in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

20. Interaction With Unions of Families I. The diagram

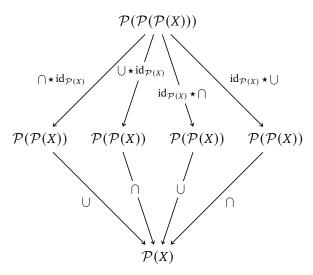
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \cap} & \mathcal{P}(\mathcal{P}(X)) \\
\downarrow \star \mathrm{id}_{\mathcal{P}(X)} & & & \downarrow \cap \\
\mathcal{P}(X) & \xrightarrow{\qquad \cap} & X
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \bigcup_{A \in A} A} U = \bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right)$$

for each $A \in \mathcal{P}(\mathcal{P}(X))$.

21. *Interaction With Unions of Families II.* Let *X* be a set and consider the compositions

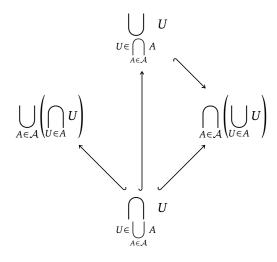


given by

$$A \mapsto \bigcup_{U \in \bigcap_{A \in \mathcal{A}} A} U, \quad A \mapsto \bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U,$$

$$A \mapsto \bigcup_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U\right), \quad A \mapsto \bigcap_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U\right)$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$. We have the following inclusions:



All other possible inclusions fail to hold in general.

PROOF 4.3.7.1.3 ► PROOF OF PROPOSITION 4.3.6.1.2

Item 1: Functoriality

Since $\mathcal{P}(X)$ is posetal, it suffices to prove the condition (\star) . So let $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ with $\mathcal{U} \subset \mathcal{V}$. We claim that

$$\bigcap_{V\in\mathcal{V}}V\subset\bigcap_{U\in\mathcal{U}}U.$$

Indeed, if $x \in \bigcap_{V \in \mathcal{V}} V$, then $x \in V$ for all $V \in \mathcal{V}$. But since $\mathcal{U} \subset \mathcal{V}$, it follows that $x \in U$ for all $U \in \mathcal{U}$ as well. Thus $x \in \bigcap_{U \in \mathcal{U}} U$, which gives our desired inclusion.

Item 2: Oplax Associativity

We have

$$\bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right) \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ for each } A \in \mathcal{A}, \\ \text{we have } x \in \bigcap_{U \in A} U \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ for each } A \in \mathcal{A} \text{ and each } \\ U \in A, \text{ we have } x \in U \right\}$$

$$= \left\{ x \in X \middle| \text{ for each } U \in \bigcup_{A \in \mathcal{A}} A, \\ \text{ we have } x \in U \right\}$$

$$\subset \left\{ x \in X \middle| \text{ for each } U \in \bigcap_{A \in \mathcal{A}} A, \\ \text{ we have } x \in U \right\}$$

$$\stackrel{\text{def}}{=} \bigcup_{A \in \mathcal{A}} U.$$

$$U \in \bigcap_{A \in \mathcal{A}} A$$

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

Item 3: Left Unitality

We have

$$\bigcap_{V \in \{U\}} V \stackrel{\text{def}}{=} \left\{ x \in X \mid \text{for each } V \in \{U\}, \right\}$$

$$= \left\{ x \in X \mid x \in U \right\}$$

$$= U.$$

This finishes the proof.

Item 4: Oplax Right Unitality

If $U = \emptyset$, then we have

$$\bigcap_{\{u\} \in \chi_X(U)} \{u\} = \bigcap_{\{u\} \in \emptyset} \{u\}$$
$$= X,$$

so $\bigcap_{\{u\} \in \chi_X(U)} \{u\} = X \neq \emptyset = U$. When U is nonempty, we have two cases:

1. If *U* is a singleton, say $U = \{u\}$, we have

$$\bigcap_{\{u\}\in\chi_X(U)}\{u\}=\{u\}$$

$$\stackrel{\text{def}}{=}U.$$

2. If *U* contains at least two elements, we have

$$\bigcap_{\{u\} \in \chi_X(U)} \{u\} = \emptyset$$

$$\subset U.$$

This finishes the proof.

Item 5: Interaction With Unions I

$$\bigcap_{W \in \mathcal{U} \cup \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cup \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$
$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \text{ and each} \\ W \in \mathcal{V}, \text{ we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ for each } W \in \mathcal{U}, \right\} \\
\text{we have } x \in W \right\} \\
\cap \left\{ x \in X \middle| \text{ for each } W \in \mathcal{V}, \right\} \\
\stackrel{\text{def}}{=} \left(\bigcap_{W \in \mathcal{U}} W \right) \cap \left(\bigcap_{W \in \mathcal{V}} W \right) \\
= \left(\bigcap_{U \in \mathcal{U}} U \right) \cap \left(\bigcap_{V \in \mathcal{V}} V \right).$$

Item 6: Interaction With Unions II

Omitted.

Item 7: Interaction With Intersections 1

We have

$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cap \left(\bigcap_{V \in \mathcal{V}} V\right) \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \mathcal{U}, \\ \text{we have } x \in U \end{array} \right\}$$

$$\cup \left\{ x \in X \middle| \begin{array}{l} \text{for each } V \in \mathcal{V}, \\ \text{we have } x \in V \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cap \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\subset \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cap \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{W \in \mathcal{U} \cap \mathcal{V}} W.$$

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

Item 8: Interaction With Intersections II

Omitted.

Item 9: Interaction With Differences

Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0\}, \{0, 1\}\}\$, and let $\mathcal{V} = \{\{0\}\}\$. We have

$$\bigcap_{W \in \mathcal{U} \setminus \mathcal{V}} U = \bigcap_{W \in \{\{0,1\}\}} W$$
$$= \{0,1\},$$

whereas

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\setminus\left(\bigcap_{V\in\mathcal{V}}V\right)=\left\{0\right\}\setminus\left\{0\right\}$$
$$=\emptyset.$$

Thus we have

$$\bigcap_{W\in\mathcal{U}\backslash\mathcal{V}}W=\left\{0,1\right\}\neq\emptyset=\left(\bigcap_{U\in\mathcal{U}}U\right)\backslash\left(\bigcap_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

Item 10: Interaction With Complements I

Let $X = \{0, 1\}$ and let $\mathcal{U} = \{\{0\}\}$. We have

$$\bigcap_{W \in \mathcal{U}^{c}} U = \bigcap_{W \in \{\emptyset, \{1\}, \{0,1\}\}} W$$
$$= \emptyset,$$

whereas

$$\bigcap_{U \in \mathcal{U}} U^{c} = \{0\}^{c}$$
$$= \{1\}.$$

Thus we have

$$\bigcap_{W\in\mathcal{U}^\mathsf{c}}U=\emptyset\neq\{1\}=\bigcap_{U\in\mathcal{U}}U^\mathsf{c}.$$

This finishes the proof.

Item 11: Interaction With Complements II

This is a repetition of Item 12 of Proposition 4.3.6.1.2 and is proved there.

Item 12: Interaction With Complements III

This is a repetition of Item 11 of Proposition 4.3.6.1.2 and is proved there.

Item 13: Interaction With Symmetric Differences

Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}\$, and let $\mathcal{V} = \{\{0\}, \{0, 1\}\}\$. We have

$$\bigcap_{W \in \mathcal{U} \triangle \mathcal{V}} W = \bigcap_{W \in \{\{0\}\}} W$$
$$= \{0\},$$

whereas

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\triangle\left(\bigcap_{V\in\mathcal{V}}V\right)=\left\{0,1\right\}\triangle\left\{0\right\}$$
$$=\emptyset,$$

Thus we have

$$\bigcap_{W \in \mathcal{U} \triangle \mathcal{V}} W = \{0\} \neq \emptyset = \left(\bigcap_{U \in \mathcal{U}} U\right) \triangle \left(\bigcap_{V \in \mathcal{V}} V\right).$$

This finishes the proof.

Item 14: Interaction With Internal Homs I

This is a repetition of Item 10 of Proposition 4.4.7.1.4 and is proved there.

Item 15: Interaction With Internal Homs II

This is a repetition of Item 11 of Proposition 4.4.7.1.4 and is proved there.

Item 16: Interaction With Internal Homs III

This is a repetition of Item 12 of Proposition 4.4.7.1.4 and is proved there.

Item 17: Interaction With Direct Images

This is a repetition of Item 4 of Proposition 4.6.1.1.5 and is proved there.

Item 18: Interaction With Inverse Images

This is a repetition of Item 4 of Proposition 4.6.2.1.3 and is proved there.

Item 19: Interaction With Codirect Images

This is a repetition of Item 4 of Proposition 4.6.3.1.7 and is proved there.

Item 20: Interaction With Unions of Families I

This is a repetition of Item 20 of Proposition 4.3.6.1.2 and is proved there.

Item 21: Interaction With Unions of Families II

This is a repetition of Item 21 of Proposition 4.3.6.1.2 and is proved there.

4.3.8 Binary Unions

Let *X* be a set and let $U, V \in \mathcal{P}(X)$.

DEFINITION 4.3.8.1.1 ► **BINARY UNIONS**

The **union of** *U* **and** *V* is the set $U \cup V$ defined by

$$U \cup V \stackrel{\text{def}}{=} \bigcup_{z \in \{U, V\}} z$$

$$\stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in V\}.$$

PROPOSITION 4.3.8.1.2 ► PROPERTIES OF BINARY UNIONS

Let *X* be a set.

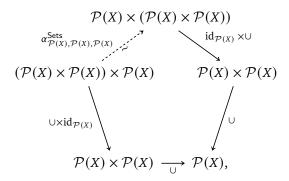
1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cup V$ define functors

$$\begin{array}{ll} U \cup -\colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ - \cup V \colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ -_1 \cup -_2 \colon (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset). \end{array}$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $U \cup V \subset A \cup V$.
- (b) If $V \subset B$, then $U \cup V \subset U \cup B$.
- (c) If $U \subset A$ and $V \subset B$, then $U \cup V \subset A \cup B$.

2. Associativity. The diagram

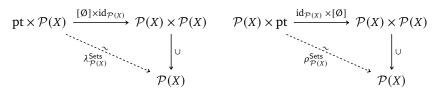


commutes, i.e. we have an equality of sets

$$(U \cup V) \cup W = U \cup (V \cup W)$$

for each $U, V, W \in \mathcal{P}(X)$.

3. Unitality. The diagrams



commute, i.e. we have equalities of sets

$$\emptyset \cup U = U$$
, $U \cup \emptyset = U$

for each $U \in \mathcal{P}(X)$.

4. Commutativity. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\sigma_{\mathcal{P}(X),\mathcal{P}(X)}^{\mathsf{Sets}}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \cup \qquad \qquad \downarrow \cup$$

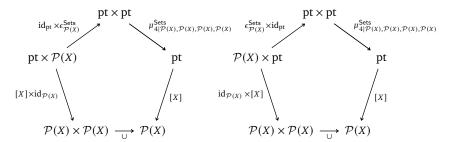
$$\mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cup V = V \cup U$$

for each $U, V \in \mathcal{P}(X)$.

5. Annihilation With X. The diagrams

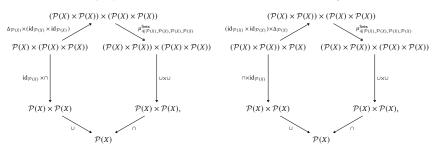


commute, i.e. we have equalities of sets

$$U \cup X = X,$$
$$X \cup V = X$$

for each $U, V \in \mathcal{P}(X)$.

6. Distributivity of Unions Over Intersections. The diagrams



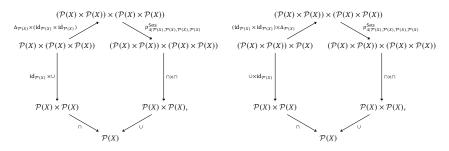
commute, i.e. we have equalities of sets

$$U \cup (V \cap W) = (U \cup V) \cap (U \cup W),$$

$$(U \cap V) \cup W = (U \cup W) \cap (V \cup W)$$

for each $U, V, W \in \mathcal{P}(X)$.

7. Distributivity of Intersections Over Unions. The diagrams

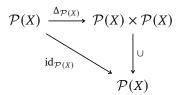


commute, i.e. we have equalities of sets

$$U \cap (V \cup W) = (U \cap V) \cup (U \cap W),$$
$$(U \cup V) \cap W = (U \cap W) \cup (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

8. *Idempotency*. The diagram



commutes, i.e. we have an equality of sets

$$U \cup U = U$$

for each $U \in \mathcal{P}(X)$.

9. Via Intersections and Symmetric Differences. The diagram

$$(\mathcal{P}(X) \times \mathcal{P}(X)) \times (\mathcal{P}(X) \times \mathcal{P}(X)) \xrightarrow{\Delta \times \cap} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\Delta_{\mathcal{P}(X) \times \mathcal{P}(X)} / \Delta$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \longrightarrow \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cup V = (U \triangle V) \triangle (U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

10. Interaction With Characteristic Functions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

11. Interaction With Characteristic Functions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each $U, V \in \mathcal{P}(X)$.

12. *Interaction With Direct Images.* Let $f: X \to Y$ be a function. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_! \times f_!} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\
\downarrow & & \downarrow & & \downarrow \\
\mathcal{P}(X) & \xrightarrow{f_!} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$

for each $U, V \in \mathcal{P}(X)$.

13. *Interaction With Inverse Images.* Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

14. *Interaction With Codirect Images.* Let $f: X \to Y$ be a function. We have a natural transformation

with components

$$f_*(U) \cup f_*(V) \subset f_*(U \cup V)$$

indexed by $U, V \in \mathcal{P}(X)$.

15. *Interaction With Powersets and Semirings.* The quintuple $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$ is an idempotent commutative semiring.

PROOF 4.3.8.1.3 ► PROOF OF PROPOSITION 4.3.8.1.2

Item 1: Functoriality

See [Pro25an].

Item 2: Associativity

See [Pro25ba].

Item 3: Unitality

This follows from [Pro25bd] and Item 4.

Item 4: Commutativity

See [Pro25bb].

Item 5: Annihilation With *X*

We have

$$U \cup X \stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in X\}$$
$$= \{x \in X \mid x \in X\},$$
$$= X$$

and

$$X \cup V \stackrel{\text{def}}{=} \{x \in X \mid x \in X \text{ or } x \in V\}$$
$$= \{x \in X \mid x \in X\}$$
$$= X.$$

This finishes the proof.

Item 6: Distributivity of Unions Over Intersections

See [Pro25az].

Item 7: Distributivity of Intersections Over Unions

See [Pro25aj].

Item 8: Idempotency

See [Pro25am].

Item 9: Via Intersections and Symmetric Differences

See [Pro25ay].

Item 10: Interaction With Characteristic Functions I

See [Pro25h].

Item 11: Interaction With Characteristic Functions II

See [Pro25h].

Item 12: Interaction With Direct Images

See [Pro25p].

Item 13: Interaction With Inverse Images

See [Pro25y].

Item 14: Interaction With Codirect Images

This is a repetition of Item 5 of Proposition 4.6.3.1.7 and is proved there.

Item 15: Interaction With Powersets and Semirings

This follows from Items 2 to 4 and 8 of this proposition and Items 3 to 6 and 8 of Proposition 4.3.9.1.2.

4.3.9 Binary Intersections

Let *X* be a set and let $U, V \in \mathcal{P}(X)$.

DEFINITION 4.3.9.1.1 ► BINARY INTERSECTIONS

The **intersection of** *U* **and** *V* is the set $U \cap V$ defined by

$$U \cap V \stackrel{\text{def}}{=} \bigcap_{z \in \{U, V\}} z$$

$$\stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in V\}.$$

PROPOSITION 4.3.9.1.2 ► PROPERTIES OF BINARY INTERSECTIONS

Let *X* be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cap V$ define functors

$$U \cap -: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$- \cap V: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$-_1 \cap -_2: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset).$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $U \cap V \subset A \cap V$.
- (b) If $V \subset B$, then $U \cap V \subset U \cap B$.

- (c) If $U \subset A$ and $V \subset B$, then $U \cap V \subset A \cap B$.
- 2. Adjointness. We have adjunctions

$$(U \cap - \dashv [U, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{U \cap -}{\downarrow}} \mathcal{P}(X),$$

$$(- \cap V \dashv [V, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{- \cap V}{\downarrow}} \mathcal{P}(X),$$

witnessed by bijections

$$\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, [V, W]_X),$$

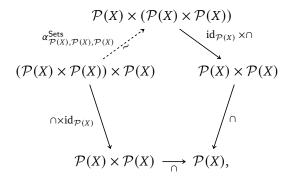
 $\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(V, [U, W]_X),$

natural in $U, V, W \in \mathcal{P}(X)$, where

$$[-1,-2]_X \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

is the bifunctor of Section 4.4.7. In particular, the following statements hold for each $U, V, W \in \mathcal{P}(X)$:

- (a) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $U \subset [V, W]_X$.
- (b) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $V \subset [U, W]_X$.
- 3. Associativity. The diagram

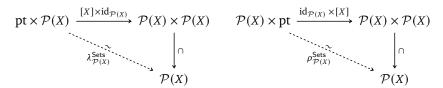


commutes, i.e. we have an equality of sets

$$(U \cap V) \cap W = U \cap (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

4. Unitality. The diagrams



commute, i.e. we have equalities of sets

$$X \cap U = U,$$
$$U \cap X = U$$

for each $U \in \mathcal{P}(X)$.

5. Commutativity. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\sigma_{\mathcal{P}(X),\mathcal{P}(X)}^{\mathsf{Sets}}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\uparrow \cap \qquad \qquad \downarrow \cap$$

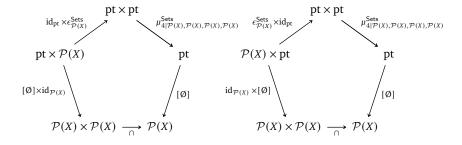
$$\mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cap V = V \cap U$$

for each $U, V \in \mathcal{P}(X)$.

6. Annihilation With the Empty Set. The diagrams

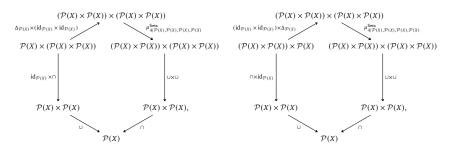


commute, i.e. we have equalities of sets

$$\emptyset \cap X = \emptyset$$
, $X \cap \emptyset = \emptyset$

for each $U \in \mathcal{P}(X)$.

7. Distributivity of Unions Over Intersections. The diagrams



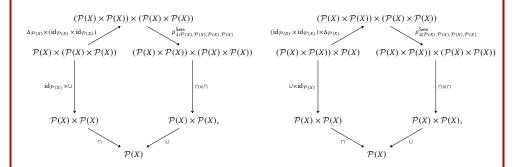
commute, i.e. we have equalities of sets

$$U \cup (V \cap W) = (U \cup V) \cap (U \cup W),$$

$$(U \cap V) \cup W = (U \cup W) \cap (V \cup W)$$

for each $U, V, W \in \mathcal{P}(X)$.

8. Distributivity of Intersections Over Unions. The diagrams



commute, i.e. we have equalities of sets

$$U \cap (V \cup W) = (U \cap V) \cup (U \cap W),$$
$$(U \cup V) \cap W = (U \cap W) \cup (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

9. *Idempotency*. The diagram

$$\mathcal{P}(X) \xrightarrow{\Delta_{\mathcal{P}(X)}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \cap$$

$$\mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cap U = U$$

for each $U \in \mathcal{P}(X)$.

10. Interaction With Characteristic Functions I. We have

$$\chi_{U\cap V}=\chi_U\chi_V$$

for each $U, V \in \mathcal{P}(X)$.

11. Interaction With Characteristic Functions II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

12. *Interaction With Direct Images.* Let $f: X \to Y$ be a function. We have a natural transformation

with components

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

13. Interaction With Inverse Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\uparrow \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U \cap V) = f^{-1}(U) \cap f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

14. *Interaction With Codirect Images.* Let $f: X \to Y$ be a function. The diagram

$$\begin{array}{ccccc} \mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_* \times f_*} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\ & & & & \downarrow \cap \\ & & & \downarrow \cap \\ & \mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y) \end{array}$$

commutes, i.e. we have

$$f_*(U) \cap f_*(V) = f_*(U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

- 15. *Interaction With Powersets and Monoids With Zero.* The quadruple $((\mathcal{P}(X), \emptyset), \cap, X)$ is a commutative monoid with zero.
- 16. *Interaction With Powersets and Semirings.* The quintuple $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$ is an idempotent commutative semiring.

PROOF 4.3.9.1.3 ► PROOF OF PROPOSITION 4.3.9.1.2 Item 1: Functoriality See [Pro25al]. Item 2: Adjointness See [MSE 267469]. Item 3: Associativity See [Pro25r]. Item 4: Unitality This follows from [Pro25v] and Item 5. Item 5: Commutativity See [Pro25s]. Item 6: Annihilation With the Empty Set This follows from [Pro25t] and Item 5. Item 7: Distributivity of Unions Over Intersections See [Pro25az]. Item 8: Distributivity of Intersections Over Unions See [Pro25aj]. Item 9: Idempotency See [Pro25ak]. Item 10: Interaction With Characteristic Functions I See [Pro25e]. Item 11: Interaction With Characteristic Functions II See [Pro25e]. Item 12: Interaction With Direct Images See [Pro25n]. Item 13: Interaction With Inverse Images See [Pro25w]. Item 14: Interaction With Codirect Images

This is a repetition of Item 6 of Proposition 4.6.3.1.7 and is proved there.

Item 15: Interaction With Powersets and Monoids With Zero

This follows from Items 3 to 6.

Item 16: Interaction With Powersets and Semirings

This follows from Items 2 to 4 and 8 and Items 3 to 6 and 8 of Proposition 4.3.9.1.2.

4.3.10 Differences

Let *X* and *Y* be sets.

DEFINITION 4.3.10.1.1 ► **DIFFERENCES**

The **difference of** X **and** Y is the set $X \setminus Y$ defined by

$$X \setminus Y \stackrel{\mathrm{def}}{=} \{ a \in X \mid a \notin Y \}.$$

PROPOSITION 4.3.10.1.2 ► PROPERTIES OF DIFFERENCES

Let *X* be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cap V$ define functors

$$\begin{array}{ccc} U \setminus -\colon & (\mathcal{P}(X), \supset) & \to (\mathcal{P}(X), \subset), \\ - \setminus V \colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ -_1 \setminus -_2 \colon (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \supset) \to (\mathcal{P}(X), \subset). \end{array}$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $U \setminus V \subset A \setminus V$.
- (b) If $V \subset B$, then $U \setminus B \subset U \setminus V$.
- (c) If $U \subset A$ and $V \subset B$, then $U \setminus B \subset A \setminus V$.
- 2. De Morgan's Laws. We have equalities of sets

$$X \setminus (U \cup V) = (X \setminus U) \cap (X \setminus V),$$

$$X \setminus (U \cap V) = (X \setminus U) \cup (X \setminus V)$$

for each $U, V \in \mathcal{P}(X)$.

3. Interaction With Unions I. We have equalities of sets

$$U \setminus (V \cup W) = (U \setminus V) \cap (U \setminus W)$$

for each $U, V, W \in \mathcal{P}(X)$.

4. Interaction With Unions II. We have equalities of sets

$$(U \setminus V) \cup W = (U \cup W) \setminus (V \setminus W)$$

for each $U, V, W \in \mathcal{P}(X)$.

5. Interaction With Unions III. We have equalities of sets

$$U \setminus (V \cup W) = (U \cup W) \setminus (V \cup W)$$
$$= (U \setminus V) \setminus W$$
$$= (U \setminus W) \setminus V$$

for each $U, V, W \in \mathcal{P}(X)$.

6. Interaction With Unions IV. We have equalities of sets

$$(U \cup V) \setminus W = (U \setminus W) \cup (V \setminus W)$$

for each $U, V, W \in \mathcal{P}(X)$.

7. Interaction With Intersections. We have equalities of sets

$$(U \setminus V) \cap W = (U \cap W) \setminus V$$
$$= U \cap (W \setminus V)$$

for each $U, V, W \in \mathcal{P}(X)$.

8. *Interaction With Complements*. We have an equality of sets

$$U\setminus V=U\cap V^\mathsf{c}$$

for each $U, V \in \mathcal{P}(X)$.

9. Interaction With Symmetric Differences. We have an equality of sets

$$U\setminus V=U\bigtriangleup (U\cap V)$$

for each $U, V \in \mathcal{P}(X)$.

10. Triple Differences. We have

$$U \setminus (V \setminus W) = (U \cap W) \cup (U \setminus V)$$

for each $U, V, W \in \mathcal{P}(X)$.

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11. Left Annihilation. We have

$$\emptyset \setminus U = \emptyset$$

for each $U \in \mathcal{P}(X)$.

12. Right Unitality. We have

$$U \setminus \emptyset = U$$

for each $U \in \mathcal{P}(X)$.

13. Right Annihilation. We have

$$U \setminus X = \emptyset$$

for each $U \in \mathcal{P}(X)$.

14. *Invertibility*. We have

$$U \setminus U = \emptyset$$

for each $U \in \mathcal{P}(X)$.

- 15. *Interaction With Containment*. The following conditions are equivalent:
 - (a) We have $V \setminus U \subset W$.
 - (b) We have $V \setminus W \subset U$.
- 16. Interaction With Characteristic Functions. We have

$$\chi_{U\setminus V} = \chi_U - \chi_{U\cap V}$$

for each $U, V \in \mathcal{P}(X)$.

17. Interaction With Direct Images. We have a natural transformation

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by $U, V \in \mathcal{P}(X)$.

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18. Interaction With Inverse Images. The diagram

commutes, i.e. we have

$$f^{-1}(U \setminus V) = f^{-1}(U) \setminus f^{-1}(V)$$

for each $U, V \in \mathcal{P}(X)$.

19. Interaction With Codirect Images. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_{!}^{\mathsf{op}} \times f_{!}} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \\
\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \\
\mathcal{P}(X) \xrightarrow{f_{!}} \mathcal{P}(Y)$$

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by $U, V \in \mathcal{P}(X)$.

PROOF 4.3.10.1.3 ► PROOF OF PROPOSITION 4.3.10.1.2

Item 1: Functoriality

See [Pro25ad] and [Pro25ah].

Item 2: De Morgan's Laws

See [Pro25k].

Item 3: Interaction With Unions I

See [Pro251].

Item 4: Interaction With Unions II

We have

```
(U \setminus V) \cup W \stackrel{\text{def}}{=} \{x \in X \mid (x \in U \text{ and } x \notin V) \text{ or } x \in W\}
= \{x \in X \mid (x \in U \text{ or } x \in W) \text{ and } (x \notin V \text{ or } x \in W)\}
= \{x \in X \mid (x \in U \cup W) \text{ and not } (x \in V \text{ and } x \notin W)\}
= \{x \in X \mid (x \in U \cup W) \text{ and not } (x \in V \setminus W)\}
= \{x \in X \mid (x \in (U \cup W) \setminus (V \setminus W))\}
= (U \cup W) \setminus (V \setminus W).
```

Item 5: Interaction With Unions III

See [Pro25ai].

Item 6: Interaction With Unions IV

See [Pro25ac].

Item 7: Interaction With Intersections

See [Pro25u].

Item 8: Interaction With Complements

See [Pro25aa].

Item 9: Interaction With Symmetric Differences

See [Pro25ab].

Item 10: Triple Differences

See [Pro25ag].

Item 11: Left Annihilation

The direction $\emptyset \subset \emptyset \setminus U$ always holds. Now assume $x \in \emptyset \setminus U$. Then, $x \in \emptyset$ and $x \notin U$. Hence $\emptyset \setminus U \subset \emptyset$ must hold and the sets are equal.

Item 12: Right Unitality

See [Pro25ae].

Item 13: Right Annihilation

It suffices to show that no $x \in X$ can be an element of $U \setminus X$. Assume $x \in U \setminus X$. Then $x \notin X$, contradicting $x \in X$. This completes the proof.

Item 14: Invertibility

See [Pro25af].

Item 15: Interaction With Containment

The conditions are symmetric in U, W, hence it suffices to show that $V \setminus U \subset W$ implies $V \setminus W \subset U$. So assume $V \setminus U \subset W$, $x \in V \setminus W$. Then $x \in V$, $x \notin W$. So by contraposition, $x \notin V \setminus U$. But $x \in V$, so we must have $x \in U$, completing the proof.

Item 16: Interaction With Characteristic Functions

See [Pro25f].

Item 17: Interaction With Direct Images

See [Pro25o].

Item 18: Interaction With Inverse Images

See [Pro25x].

4.3.11 Complements

Let *X* be a set and let $U \in \mathcal{P}(X)$.

DEFINITION 4.3.11.1.1 ► **COMPLEMENTS**

The **complement of** U is the set U^{c} defined by

$$U^{c} \stackrel{\text{def}}{=} X \setminus U$$

$$\stackrel{\text{def}}{=} \{ a \in X \mid a \notin U \}.$$

PROPOSITION 4.3.11.1.2 ► PROPERTIES OF COMPLEMENTS

Let *X* be a set.

1. Functoriality. The assignment $U \mapsto U^{c}$ defines a functor

$$(-)^{c} : \mathcal{P}(X)^{op} \to \mathcal{P}(X).$$

In particular, the following statements hold for each $U, V \in \mathcal{P}(X)$:

 (\star) If $U \subset V$, then $V^{c} \subset U^{c}$.

2. De Morgan's Laws. The diagrams

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)^{\mathsf{op}} \xrightarrow{\cup^{\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}} \qquad \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)^{\mathsf{op}} \xrightarrow{\cap^{\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}}$$

$$(-)^{\mathsf{c}} \times (-)^{\mathsf{c}} \downarrow \qquad \qquad (-)^{\mathsf{c}} \times (-)^{\mathsf{c}} \downarrow \qquad \qquad \downarrow (-)^{\mathsf{c}}$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\quad \cap \quad} \mathcal{P}(X) \qquad \qquad \mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\quad \cup \quad} \mathcal{P}(X)$$

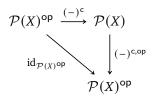
commute, i.e. we have equalities of sets

$$(U \cup V)^{c} = U^{c} \cap V^{c},$$

$$(U \cap V)^{c} = U^{c} \cup V^{c}$$

for each $U, V \in \mathcal{P}(X)$.

3. Involutority. The diagram



commutes, i.e. we have

$$(U^{\mathsf{c}})^{\mathsf{c}} = U$$

for each $U \in \mathcal{P}(X)$.

4. Interaction With Characteristic Functions. We have

$$\chi_{U^c} = 1 - \chi_U$$

for each $U \in \mathcal{P}(X)$.

5. *Interaction With Direct Images.* Let $f: X \to Y$ be a function. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\text{op}} & \xrightarrow{f_{*}^{\text{op}}} \mathcal{P}(Y)^{\text{op}} \\
 & \downarrow^{(-)^{c}} & \downarrow^{(-)^{c}} \\
\mathcal{P}(X) & \xrightarrow{f_{*}} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_!(U^{\mathsf{c}}) = f_*(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

6. *Interaction With Inverse Images.* Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \xrightarrow{f^{-1,\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}} \\
\xrightarrow{(-)^{\mathsf{c}}} \qquad \qquad \downarrow^{(-)^{\mathsf{c}}} \\
\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U^{c}) = f^{-1}(U)^{c}$$

for each $U \in \mathcal{P}(X)$.

7. *Interaction With Codirect Images.* Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(X)^{\text{op}} \xrightarrow{f_!^{\text{op}}} \mathcal{P}(Y)^{\text{op}} \\
\stackrel{(-)^c}{\downarrow} \qquad \qquad \downarrow^{(-)^c} \\
\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U^{\mathsf{c}}) = f_!(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

PROOF 4.3.11.1.3 ► PROOF OF PROPOSITION 4.3.11.1.2

Item 1: Functoriality

This follows from Item 1 of Proposition 4.3.10.1.2.

Item 2: De Morgan's Laws

See [Pro25k].

Item 3: Involutority

See [Pro25i].

Item 4: Interaction With Characteristic Functions

We consider the two cases $x \in U, x \notin U$.

1. If $x \in U$, then $x \notin U^{c}$. So $\chi_{U}(x) = 1$ and

$$\chi_{U^{c}}(x) = 0$$
$$= 1 - \chi_{U}(x).$$

2. If $x \notin U$, then $x \in U^{c}$. So $\chi_{U}(x) = 0$ and

$$\chi_{U^{c}}(x) = 1$$
$$= 1 - \chi_{U}(x).$$

Hence, the equation holds for all $x \in X$.

Item 5: Interaction With Direct Images

This is a repetition of Item 8 of Proposition 4.6.1.1.5 and is proved there.

Item 6: Interaction With Inverse Images

This is a repetition of Item 8 of Proposition 4.6.2.1.3 and is proved there.

Item 7: Interaction With Codirect Images

This is a repetition of Item 7 of Proposition 4.6.3.1.7 and is proved there.

4.3.12 Symmetric Differences

Let *X* be a set and let $U, V \in \mathcal{P}(X)$.

DEFINITION 4.3.12.1.1 ► SYMMETRIC DIFFERENCES

The **symmetric difference of** *U* **and** *V* is the set $U \triangle V$ defined by $^{\mathbf{1}}$

$$U \triangle V \stackrel{\mathrm{def}}{=} (U \setminus V) \cup (V \setminus U).$$

¹*Illustration:*



PROPOSITION 4.3.12.1.2 ► PROPERTIES OF SYMMETRIC DIFFERENCES

Let *X* be a set.

1. *Lack of Functoriality.* The assignment $(U,V) \mapsto U \triangle V$ *does not* in general define functors

$$\begin{array}{lll} U \bigtriangleup - \colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ - \bigtriangleup V \colon & (\mathcal{P}(X), \subset) & \to (\mathcal{P}(X), \subset), \\ -_1 \bigtriangleup -_2 \colon (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset). \end{array}$$

2. Via Unions and Intersections. We have

$$U \triangle V = (U \cup V) \setminus (U \cap V)$$

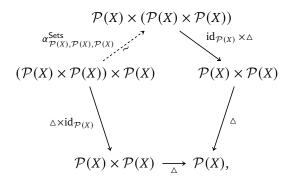
for each $U, V \in \mathcal{P}(X)$, as in the Venn diagram

$$\boxed{\bigcup_{U \triangle V}} = \boxed{\bigcup_{U \cup V}} \setminus \boxed{\bigcup_{U \cap V}}.$$

3. *Symmetric Differences of Disjoint Sets.* If *U* and *V* are disjoint, then we have

$$U \triangle V = U \cup V$$
.

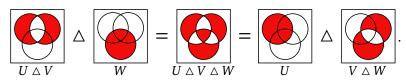
4. Associativity. The diagram



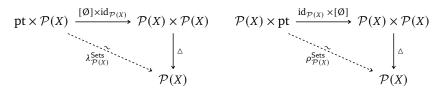
commutes, i.e. we have

$$(U \triangle V) \triangle W = U \triangle (V \triangle W)$$

for each $U, V, W \in \mathcal{P}(X)$, as in the Venn diagram



5. *Unitality*. The diagrams



commute, i.e. we have

$$U \triangle \emptyset = U,$$

$$\emptyset \triangle U = U$$

for each $U \in \mathcal{P}(X)$.

6. Commutativity. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\sigma_{\mathcal{P}(X),\mathcal{P}(X)}^{\mathsf{Sets}}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow^{\triangle}$$

$$\mathcal{P}(X)$$

commutes, i.e. we have

$$U \triangle V = V \triangle U$$

for each $U, V \in \mathcal{P}(X)$.

7. Invertibility. We have

$$U \triangle U = \emptyset$$

for each $U \in \mathcal{P}(X)$.

8. Interaction With Unions. We have

$$(U \triangle V) \cup (V \triangle T) = (U \cup V \cup W) \setminus (U \cap V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

9. Interaction With Complements I. We have

$$U \triangle U^{c} = X$$

for each $U \in \mathcal{P}(X)$.

10. Interaction With Complements II. We have

$$U \triangle X = U^{c},$$

 $X \triangle U = U^{c}$

for each $U \in \mathcal{P}(X)$.

11. Interaction With Complements III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X) \times \mathcal{P}(X) & \stackrel{\triangle}{\longrightarrow} & \mathcal{P}(X) \\
 & \downarrow^{(-)^{c}} \downarrow & & \downarrow^{(-)^{c}} \\
\mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{\triangle} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$U^{c} \triangle V^{c} = U \triangle V$$

for each $U, V \in \mathcal{P}(X)$.

12. "Transitivity". We have

$$(U \triangle V) \triangle (V \triangle W) = U \triangle W$$

for each $U, V, W \in \mathcal{P}(X)$.

13. The Triangle Inequality for Symmetric Differences. We have

$$U \triangle W \subset U \triangle V \cup V \triangle W$$

for each $U, V, W \in \mathcal{P}(X)$.

14. Distributivity Over Intersections. We have

$$U \cap (V \triangle W) = (U \cap V) \triangle (U \cap W),$$

$$(U \triangle V) \cap W = (U \cap W) \triangle (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

15. Interaction With Characteristic Functions. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \pmod{2}$$

for each $U, V \in \mathcal{P}(X)$.

16. *Bijectivity*. Given $U, V \in \mathcal{P}(X)$, the maps

$$U \triangle -: \mathcal{P}(X) \to \mathcal{P}(X),$$
$$- \triangle V: \mathcal{P}(X) \to \mathcal{P}(X)$$

are self-inverse bijections. Moreover, the map

$$\mathcal{P}(X) \longrightarrow \mathcal{P}(X)$$

$$C \longmapsto C \triangle (U \triangle V)$$

is a bijection of $\mathcal{P}(X)$ onto itself sending U to V and V to U.

- 17. *Interaction With Powersets and Groups.* Let *X* be a set.
 - (a) The quadruple $(\mathcal{P}(X), \triangle, \emptyset, \mathrm{id}_{\mathcal{P}(X)})$ is an abelian group.¹
 - (b) Every element of $\mathcal{P}(X)$ has order 2 with respect to \triangle , and thus $\mathcal{P}(X)$ is a *Boolean group* (i.e. an abelian 2-group).
- 18. *Interaction With Powersets and Vector Spaces I.* The pair $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$ consisting of
 - The group $\mathcal{P}(X)$ of Item 17;
 - The map $\alpha_{\mathcal{P}(X)} : \mathbb{F}_2 \times \mathcal{P}(X) \to \mathcal{P}(X)$ defined by

$$0 \cdot U \stackrel{\text{def}}{=} \emptyset,$$
$$1 \cdot U \stackrel{\text{def}}{=} U;$$

is an \mathbb{F}_2 -vector space.

- 19. *Interaction With Powersets and Vector Spaces II.* If *X* is finite, then:
 - (a) The set of singletons sets on the elements of X forms a basis for the \mathbb{F}_2 -vector space $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$ of Item 18.
 - (b) We have

$$\dim(\mathcal{P}(X)) = \#X.$$

- 20. *Interaction With Powersets and Rings*. The quintuple $(\mathcal{P}(X), \triangle, \cap, \emptyset, X)$ is a commutative ring.²
- 21. Interaction With Direct Images. We have a natural transformation

with components

$$f_1(U) \triangle f_1(V) \subset f_1(U \triangle V)$$

indexed by $U, V \in \mathcal{P}(X)$.

22. Interaction With Inverse Images. The diagram

$$\begin{array}{cccc} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) & \xrightarrow{f^{\mathsf{op},-1} \times f^{-1}} & \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \\ & & & & \downarrow^{\triangle} \\ & & & \mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X) \end{array}$$

i.e. we have

$$f^{-1}(U) \triangle f^{-1}(V) = f^{-1}(U \triangle V)$$

for each $U, V \in \mathcal{P}(Y)$.

23. Interaction With Codirect Images. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_{*}^{\mathsf{op}} \times f_{*}} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(X) \xrightarrow{f_{*}} \mathcal{P}(Y)$$

with components

$$f_*(U \triangle V) \subset f_*(U) \triangle f_*(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

i. When $X = \emptyset$, we have an isomorphism of groups between $\mathcal{P}(\emptyset)$ and the trivial group:

$$(\mathcal{P}(\emptyset), \triangle, \emptyset, id_{\mathcal{P}(\emptyset)}) \cong pt.$$

ii. When X = pt, we have an isomorphism of groups between $\mathcal{P}(pt)$ and $\mathbb{Z}_{/2}$:

$$(\mathcal{P}(\mathsf{pt}), \triangle, \emptyset, \mathsf{id}_{\mathcal{P}(\mathsf{pt})}) \cong \mathbb{Z}_{/2}.$$

iii. When $X = \{0, 1\}$, we have an isomorphism of groups between $\mathcal{P}(\{0, 1\})$ and $\mathbb{Z}_{/2} \times \mathbb{Z}_{/2}$:

$$(\mathcal{P}(\{0,1\}), \triangle, \emptyset, id_{\mathcal{P}(\{0,1\})}) \cong \mathbb{Z}_{/2} \times \mathbb{Z}_{/2}.$$

Warning: The analogous statement replacing intersections by unions (i.e. that the quintuple $(\mathcal{P}(X), \Delta, \cup, \emptyset, X)$ is a ring) is false, however. See [Pro25aw] for a proof.

PROOF 4.3.12.1.3 ► PROOF OF PROPOSITION 4.3.12.1.2

Item 1: Lack of Functoriality

Let $X = \{0, 1\}, U = \{0\}$. Then $\emptyset \subset U$, but $U \triangle \emptyset = U \not\subset \emptyset = U \triangle U$ from Item 5 and Item 7. This gives a counterexample to the first statement. By using Item 6, we can adapt it to the second and third statement.

Item 2: Via Unions and Intersections

See [Pro25m].

Item 3: Symmetric Differences of Disjoint Sets

¹Here are some examples:

Since *U* and *V* are disjoint, we have $U \cap V = \emptyset$, and therefore we have

$$U \triangle V = (U \cup V) \setminus (U \cap V)$$
$$= (U \cup V) \setminus \emptyset$$
$$= U \cup V,$$

where we've used Item 2 and Item 12 of Proposition 4.3.10.1.2.

Item 4: Associativity

See [Pro25ao].

Item 5: Unitality

This follows from Item 6 and [Pro25at].

Item 6: Commutativity

See [Pro25ap].

Item 7: Invertibility

See [Pro25av].

Item 8: Interaction With Unions

See [Pro25bc].

Item 9: Interaction With Complements I

See [Pro25as].

Item 10: Interaction With Complements II

This follows from Item 6 and [Pro25ax].

Item 11: Interaction With Complements III

See [Pro25aq].

Item 12: "Transitivity"

We have

```
(U \triangle V) \triangle (V \triangle W) = U \triangle (V \triangle (V \triangle W)) \quad \text{(by Item 4)}
= U \triangle ((V \triangle V) \triangle W) \quad \text{(by Item 4)}
= U \triangle (\emptyset \triangle W) \quad \text{(by Item 7)}
= U \triangle W. \quad \text{(by Item 5)}
```

This finishes the proof.

Item 13: The Triangle Inequality for Symmetric Differences

This follows from Items 2 and 12.

Item 14: Distributivity Over Intersections

See [Pro25q].

Item 15: Interaction With Characteristic Functions

See [Pro25g].

Item 16: Bijectivity

We show that

$$(U \triangle -): \mathcal{P}(X) \to \mathcal{P}(X)$$

is self-inverse.

Let $W \in \mathcal{P}(X)$. Then,

$$U \triangle (U \triangle W) = (U \triangle U) \triangle W$$
 (by Item 4)
= $\emptyset \triangle W$ (by Item 7)
= W . (by Item 5)

- By Item 6, $(-\triangle V) = (V \triangle -)$, hence the former is also self-inverse by the first point.
- The map $-\triangle (U \triangle V)$ is a bijection as a special case of the second point. From the first two points and Item 6, we get

$$U \triangle (U \triangle V) = V$$
, $V \triangle (U \triangle V) = V \triangle (V \triangle U) = U$.

Hence the function maps U to V and V to U.

Item 17: Interaction With Powersets and Groups

Item 17a follows from Items 4 to 7, while Item 17b follows from Item 7.1

Item 18: Interaction With Powersets and Vector Spaces I

See [MSE 2719059].

Item 19: Interaction With Powersets and Vector Spaces II

See [MSE 2719059].

Item 20: Interaction With Powersets and Rings

This follows from Items 6 and 15 of Proposition 4.3.9.1.2 and Items 14 and 17.2

Item 21: Interaction With Direct Images

This is a repetition of Item 9 of Proposition 4.6.1.1.5 and is proved there.

Item 22: Interaction With Inverse Images

This is a repetition of Item 9 of Proposition 4.6.2.1.3 and is proved there.

Item 23: Interaction With Codirect Images

This is a repetition of Item 8 of Proposition 4.6.3.1.7 and is proved there.

¹Reference: [Pro25ar]. ²Reference: [Pro25au].

4.4 Powersets

4.4.1 Foundations

Let *X* be a set.

DEFINITION 4.4.1.1.1 ▶ POWERSETS

The **powerset of** *X* is the set $\mathcal{P}(X)$ defined by

$$\mathcal{P}(X) \stackrel{\text{def}}{=} \{ U \in P \mid U \subset X \},\,$$

where *P* is the set in the axiom of powerset, ?? of ??.

REMARK 4.4.1.1.2 ▶ POWERSETS AS DECATEGORIFICATIONS OF CO/PRESHEAF CATEGORIES

Under the analogy that $\{t, f\}$ should be the (-1)-categorical analogue of Sets, we may view the powerset of a set as a decategorification of the category of presheaves of a category (or of the category of copresheaves):

• The powerset of a set *X* is equivalently (Item 2 of Proposition 4.5.1.1.4) the set

$$Sets(X, \{t, f\})$$

of functions from *X* to the set {t, f} of classical truth values.

• The category of presheaves on a category *C* is the category

$$\operatorname{Fun}(C^{\operatorname{op}},\operatorname{Sets})$$

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of functors from C^{op} to the category Sets of sets.

NOTATION 4.4.1.1.3 ► FURTHER NOTATION FOR POWERSETS

Let *X* be a set.

- 1. We write $\mathcal{P}_0(X)$ for the set of nonempty subsets of X.
- 2. We write $\mathcal{P}_{fin}(X)$ for the set of finite subsets of X.

PROPOSITION 4.4.1.1.4 ► ELEMENTARY PROPERTIES OF POWERSETS

Let *X* be a set.

- 1. *Co/Completeness*. The (posetal) category (associated to) $(\mathcal{P}(X), \subset)$ is complete and cocomplete:
 - (a) *Products*. The products in $\mathcal{P}(X)$ are given by intersection of subsets.
 - (b) *Coproducts*. The coproducts in $\mathcal{P}(X)$ are given by union of subsets.
 - (c) *Co/Equalisers*. Being a posetal category, $\mathcal{P}(X)$ only has at most one morphisms between any two objects, so co/equalisers are trivial.
- 2. *Cartesian Closedness*. The category $\mathcal{P}(X)$ is Cartesian closed.
- 3. Powersets as Sets of Relations. We have bijections

$$\mathcal{P}(X) \cong \text{Rel}(\text{pt}, X),$$

 $\mathcal{P}(X) \cong \text{Rel}(X, \text{pt}),$

natural in $X \in \text{Obj}(\mathsf{Sets})$.

4. Interaction With Products I. The map

$$\mathcal{P}(X) \times \mathcal{P}(Y) \longrightarrow \mathcal{P}(X \coprod Y)$$
$$(U, V) \longmapsto U \cup V$$

is an isomorphism of sets, natural in $X,Y \in \text{Obj}(\mathsf{Sets})$ with respect to each of the functor structures $\mathcal{P}_!$, \mathcal{P}^{-1} , and \mathcal{P}_* on \mathcal{P} of Proposition 4.4.2.1.1. Moreover, this makes each of $\mathcal{P}_!$, \mathcal{P}^{-1} , and \mathcal{P}_* into a symmetric monoidal functor.

5. Interaction With Products II. The map

$$\mathcal{P}(X) \times \mathcal{P}(Y) \longrightarrow \mathcal{P}(X \coprod Y)$$
$$(U, V) \longmapsto U \boxtimes_{X \times Y} V,$$

where¹

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \{(u, v) \in X \times Y \mid u \in U \text{ and } v \in V\}$$

is an inclusion of sets, natural in $X,Y \in \text{Obj}(\mathsf{Sets})$ with respect to each of the functor structures $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* on \mathcal{P} of Proposition 4.4.2.1.1. Moreover, this makes each of $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* into a symmetric monoidal functor.

6. Interaction With Products III. We have an isomorphism

$$\mathcal{P}(X) \otimes \mathcal{P}(Y) \cong \mathcal{P}(X \times Y),$$

natural in $X, Y \in \text{Obj}(\mathsf{Sets})$ with respect to each of the functor structures $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* on \mathcal{P} of Proposition 4.4.2.1.1, where \otimes denotes the tensor product of suplattices of ??. Moreover, this makes each of $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* into a symmetric monoidal functor.

PROOF 4.4.1.1.5 ► PROOF OF PROPOSITION 4.4.1.1.4

Item 1: Co/Completeness

Omitted.

Item 2: Cartesian Closedness

See Section 4.4.7.

Item 3: Powersets as Sets of Relations

Indeed, we have

$$Rel(pt, X) \stackrel{\text{def}}{=} \mathcal{P}(pt \times X)$$
$$\cong \mathcal{P}(X)$$

¹The set $U \boxtimes_{X \times Y} V$ is usually denoted simply $U \times V$. Here we denote it in this somewhat weird way to highlight the similarity to external tensor products in six-functor formalisms (see also Section 4.6.4).

and

$$Rel(X, pt) \stackrel{\text{def}}{=} \mathcal{P}(X \times pt)$$

$$\cong \mathcal{P}(X),$$

where we have used Item 5 of Proposition 4.1.3.1.4.

Item 4: Interaction With Products I

The inverse of the map in the statement is the map

$$\Phi \colon \mathcal{P}(X \coprod Y) \to \mathcal{P}(X) \times \mathcal{P}(Y)$$

defined by

$$\Phi(S) \stackrel{\mathrm{def}}{=} (S_X, S_Y)$$

for each $S \in \mathcal{P}(X \coprod Y)$, where

$$S_X \stackrel{\text{def}}{=} \{ x \in X \mid (0, x) \in S \}$$

 $S_Y \stackrel{\text{def}}{=} \{ y \in Y \mid (1, y) \in S \}.$

The rest of the proof is omitted.

Item 5: Interaction With Products II

Omitted.

Item 6: Interaction With Products III

Omitted.

4.4.2 Functoriality of Powersets

PROPOSITION 4.4.2.1.1 ► FUNCTORIALITY OF POWERSETS

Let *X* be a set.

1. Functoriality I. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}_1 \colon \mathsf{Sets} \to \mathsf{Sets}$$
,

where

• Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}_!(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

• *Action on Morphisms*. For each $A, B \in Obj(Sets)$, the action on morphisms

$$\mathcal{P}_{*|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of $\mathcal{P}_!$ at (A, B) is the map defined by sending a map of sets $f: A \to B$ to the map

$$\mathcal{P}_!(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_!(f) \stackrel{\text{def}}{=} f_!,$$

as in Definition 4.6.1.1.1.

2. Functoriality II. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}^{-1}$$
: Sets^{op} \rightarrow Sets,

where

• Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}^{-1}(A) \stackrel{\text{def}}{=} \mathcal{P}(A)$$
.

• *Action on Morphisms*. For each $A, B \in Obj(Sets)$, the action on morphisms

$$\mathcal{P}_{A,B}^{-1} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(B),\mathcal{P}(A))$$

of \mathcal{P}^{-1} at (A,B) is the map defined by sending a map of sets $f\colon A\to B$ to the map

$$\mathcal{P}^{-1}(f) \colon \mathcal{P}(B) \to \mathcal{P}(A)$$

defined by

$$\mathcal{P}^{-1}(f) \stackrel{\text{def}}{=} f^{-1},$$

as in Definition 4.6.2.1.1.

3. Functoriality III. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}_*$$
: Sets \rightarrow Sets,

where

• *Action on Objects.* For each $A \in Obj(Sets)$, we have

$$\mathcal{P}_*(A) \stackrel{\text{def}}{=} \mathcal{P}(A)$$
.

• *Action on Morphisms*. For each $A, B \in Obj(Sets)$, the action on morphisms

$$\mathcal{P}_{!|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of \mathcal{P}_* at (A, B) is the map defined by sending a map of sets $f: A \to B$ to the map

$$\mathcal{P}_*(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_*(f) \stackrel{\text{def}}{=} f_*,$$

as in Definition 4.6.3.1.1.

PROOF 4.4.2.1.2 ▶ PROOF OF PROPOSITION 4.4.2.1.1

Item 1: Functoriality I

This follows from Items 3 and 4 of Proposition 4.6.1.1.7.

Item 2: Functoriality II

This follows from Items 3 and 4 of Proposition 4.6.2.1.5.

Item 3: Functoriality III

This follows from Items 3 and 4 of Proposition 4.6.3.1.9.

4.4.3 Adjointness of Powersets I

PROPOSITION 4.4.3.1.1 ► ADJOINTNESS OF POWERSETS I

We have an adjunction

$$\left(\mathcal{P}^{-1}\dashv\mathcal{P}^{-1,\mathsf{op}}\right)$$
: Sets $\overset{\mathcal{P}^{-1}}{\underset{\mathcal{P}^{-1,\mathsf{op}}}{\smile}}\mathsf{Sets},$

witnessed by a bijection

$$\underbrace{\mathsf{Sets}^{\mathsf{op}}(\mathcal{P}(X),Y)}_{\substack{\text{def}\\ = \\ \mathsf{Sets}(Y,\mathcal{P}(X))}} \cong \mathsf{Sets}(X,\mathcal{P}(Y)),$$

natural in $X \in \text{Obj}(\mathsf{Sets})$ and $Y \in \text{Obj}(\mathsf{Sets}^{\mathsf{op}})$.

PROOF 4.4.3.1.2 ► PROOF OF PROPOSITION 4.4.3.1.1

We have

```
Sets<sup>op</sup>(\mathcal{P}(A), B) \stackrel{\text{def}}{=} Sets(B, \mathcal{P}(A))
\cong \text{Sets}(B, \text{Sets}(A, \{t, f\})) \text{ (by Item 2 of Proposition 4.5.1 1.4)}
\cong \text{Sets}(A \times B, \{t, f\}) \text{ (by Item 2 of Proposition 4.1.3 1.4)}
\cong \text{Sets}(A, \text{Sets}(B, \{t, f\})) \text{ (by Item 2 of Proposition 4.1.3 1.4)}
\cong \text{Sets}(A, \mathcal{P}(B)), \text{ (by Item 2 of Proposition 4.5.1 1.4)}
```

where all bijections are natural in A and B.

4.4.4 Adjointness of Powersets II

PROPOSITION 4.4.4.1.1 ► ADJOINTNESS OF POWERSETS II

We have an adjunction

$$(Gr\dashv \mathcal{P}_!)\colon \quad \mathsf{Sets}\underbrace{\overset{Gr}{\underset{\mathcal{P}_!}{\bot}}}_{\mathsf{Rel}} \mathsf{Rel},$$

witnessed by a bijection of sets

$$Rel(Gr(X), Y) \cong Sets(X, \mathcal{P}(Y))$$

natural in $X \in \text{Obj}(\mathsf{Sets})$ and $Y \in \text{Obj}(\mathsf{Rel})$, where Gr is the graph functor of Relations, Item 1 of Proposition 8.2.2.1.2 and $\mathcal{P}_!$ is the functor of Relations, Proposition 8.7.5.1.1.

¹Here we are using Item 3 of Proposition 4.5.1.1.4.

PROOF 4.4.4.1.2 ▶ PROOF OF PROPOSITION 4.4.4.1.1

We have

$$Rel(Gr(A), B) \cong \mathcal{P}(A \times B)$$

$$\cong Sets(A \times B, \{t, f\}) \qquad \text{(by Item 2 of Proposition 4.5.1.1.4)}$$

$$\cong Sets(A, Sets(B, \{t, f\})) \qquad \text{(by Item 2 of Proposition 4.1.3.1.4)}$$

$$\cong Sets(A, \mathcal{P}(B)), \qquad \text{(by Item 2 of Proposition 4.5.1.1.4)}$$

where all bijections are natural in A, (where we are using Item 3 of Proposition 4.5.1.1.4). Explicitly, this isomorphism is given by sending a relation $R: Gr(A) \to B$ to the map $R^{\dagger}: A \to \mathcal{P}(B)$ sending a to the subset R(a) of B, as in Relations, Definition 8.1.1.1.1.

Naturality in *B* is then the statement that given a relation $R: B \to B'$, the diagram

commutes, which follows from Relations, Remark 8.7.1.1.3.

4.4.5 Powersets as Free Cocompletions

Let *X* be a set.

PROPOSITION 4.4.5.1.1 ➤ POWERSETS AS FREE COCOMPLETIONS: UNIVERSAL PROPERTY

The pair $(\mathcal{P}(X), \chi_{(-)})$ consisting of

- The powerset $(\mathcal{P}(X), \subset)$ of X of Definition 4.4.1.1.1;
- The characteristic embedding $\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$ of X into $\mathcal{P}(X)$ of Definition 4.5.4.1.1;

satisfies the following universal property:

- (\star) Given another pair (Y, f) consisting of
 - A suplattice (Y, ≤);
 - A function $f: X \to Y$;

there exists a unique morphism of suplattices

$$(\mathcal{P}(X),\subset) \xrightarrow{\exists !} (Y,\preceq)$$

making the diagram



commute.

PROOF 4.4.5.1.2 ► PROOF OF PROPOSITION 4.4.5.1.1

This is a rephrasing of Proposition 4.4.5.1.3, which we prove below.

¹Here we only remark that the unique morphism of suplattices in the statement is given by the left Kan extension $\text{Lan}_{\chi_X}(f)$ of f along χ_X .

PROPOSITION 4.4.5.1.3 ► POWERSETS AS FREE COCOMPLETIONS: ADJOINTNESS

We have an adjunction

$$(\mathcal{P} \dashv \overline{\mathbb{R}})$$
: Sets $\underbrace{\overset{\mathcal{P}}{}}_{\Xi}$ SupLat,

witnessed by a bijection

$$\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))\cong\mathsf{Sets}(X,Y),$$

natural in $X \in \text{Obj}(\mathsf{Sets})$ and $(Y, \preceq) \in \text{Obj}(\mathsf{SupLat})$, where:

- The category SupLat is the category of suplattices of ??.
- The map

$$\chi_X^* : \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

witnessing the above bijection is defined by

$$\chi_X^*(f) \stackrel{\text{def}}{=} f \circ \chi_X,$$

i.e. by sending a morphism of suplattices $f \colon \mathcal{P}(X) \to Y$ to the composition

$$X \stackrel{\chi_X}{\longleftrightarrow} \mathcal{P}(X) \stackrel{f}{\longrightarrow} Y.$$

• The map

$$\operatorname{Lan}_{\gamma_X} : \operatorname{\mathsf{Sets}}(X,Y) \to \operatorname{\mathsf{SupLat}}((\mathcal{P}(X),\subset),(Y,\preceq))$$

witnessing the above bijection is given by sending a function $f: X \to Y$ to its left Kan extension along χ_X ,

$$\operatorname{Lan}_{\chi_X}(f) \colon \mathcal{P}(X) \to Y, \qquad X \xrightarrow{\chi_X} \downarrow \operatorname{Lan}_{\chi_X}(f)$$

$$X \xrightarrow{f} Y.$$

Moreover, invoking the bijection $\mathcal{P}(X) \cong \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ of Item 2 of Proposition 4.5.1.1.4, $\mathsf{Lan}_{\mathsf{Y}_X}(f)$ can be explicitly computed by

$$[\operatorname{Lan}_{\chi_X}(f)](U) = \int_{x \in X}^{x \in X} \chi_{\mathcal{D}(X)}(\chi_x, U) \odot f(x)$$

$$= \int_{x \in X}^{x \in X} \chi_{\mathcal{U}}(x) \odot f(x)$$

$$= \bigvee_{x \in X} (\chi_{\mathcal{U}}(x) \odot f(x))$$

$$= \left(\bigvee_{x \in U} (\chi_{\mathcal{U}}(x) \odot f(x))\right) \vee \left(\bigvee_{x \in U^c} (\chi_{\mathcal{U}}(x) \odot f(x))\right)$$

$$= \left(\bigvee_{x \in U} f(x)\right) \vee \left(\bigvee_{x \in U^c} \varnothing_Y\right)$$

$$= \bigvee_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$, where:

- We have used ?? for the first equality.
- We have used Proposition 4.5.5.1.1 for the second equality.
- We have used ?? for the third equality.
- The symbol \vee denotes the join in (Y, \preceq) .

– The symbol \odot denotes the tensor of an element of Y by a truth value as in $\ref{eq:total_point}$. In particular, we have

true
$$\odot f(x) \stackrel{\text{def}}{=} f(x)$$
,
false $\odot f(x) \stackrel{\text{def}}{=} \varnothing_Y$,

where \emptyset_Y is the bottom element of (Y, \preceq) .

In particular, when $(Y, \preceq_Y) = (\mathcal{P}(B), \subset)$ for some set B, the Kan extension $\text{Lan}_{\chi_X}(f)$ is given by

$$[\operatorname{Lan}_{\chi_X}(f)](U) = \bigvee_{x \in U} f(x)$$
$$= \bigcup_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$.

PROOF 4.4.5.1.4 ► PROOF OF PROPOSITION 4.4.5.1.3

Map I

We define a map

$$\Phi_{X,Y} \colon \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

as in the statement, i.e. by

$$\Phi_{X,Y}(f) \stackrel{\mathrm{def}}{=} f \circ \chi_X$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$.

Map II

We define a map

$$\Psi_{X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$$

as in the statement, i.e. by

$$\Psi_{X,Y}(f) \stackrel{\text{def}}{=} \operatorname{Lan}_{\chi_X}(f), \qquad X \xrightarrow{\chi_X} \downarrow \operatorname{Lan}_{\chi_X}(f)$$

$$X \xrightarrow{f} Y,$$

for each $f \in Sets(X, Y)$.

Invertibility I

We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = id_{\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))}$$
.

We have

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) \stackrel{\text{def}}{=} \Psi_{X,Y}(\Phi_{X,Y}(f))$$

$$\stackrel{\text{def}}{=} \Psi_{X,Y}(f \circ \chi_X)$$

$$\stackrel{\text{def}}{=} \operatorname{Lan}_{\chi_X}(f \circ \chi_X)$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$. We now claim that

$$\operatorname{Lan}_{\chi_X}(f \circ \chi_X) = f$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$. Indeed, we have

$$\begin{aligned} \left[\operatorname{Lan}_{\chi_X} (f \circ \chi_X) \right] (U) &= \bigvee_{x \in U} f(\chi_X(x)) \\ &= f \left(\bigvee_{x \in U} \chi_X(x) \right) \\ &= f \left(\bigcup_{x \in U} \{x\} \right) \\ &= f(U) \end{aligned}$$

for each $U \in \mathcal{P}(X)$, where we have used that f is a morphism of suplattices and hence preserves joins for the second equality. This proves our claim. Since we have shown that

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) = f$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$, it follows that $\Psi_{X,Y} \circ \Phi_{X,Y}$ must be equal to the identity map $\mathsf{id}_{\mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))}$ of $\mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$.

Invertibility II

We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathsf{Sets}(X,Y)} \; .$$

We have

$$\begin{aligned} [\Phi_{X,Y} \circ \Psi_{X,Y}](f) &\stackrel{\text{def}}{=} \Phi_{X,Y}(\Psi_{X,Y}(f)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}(\text{Lan}_{\chi_X}(f)) \\ &\stackrel{\text{def}}{=} \text{Lan}_{\chi_X}(f) \circ \chi_X \end{aligned}$$

for each $f \in Sets(X, Y)$. We now claim that

$$\operatorname{Lan}_{\chi_X}(f)\circ\chi_X=f$$

for each $f \in Sets(X, Y)$. Indeed, we have

$$[\operatorname{Lan}_{\chi_X}(f) \circ \chi_X](x) = \bigvee_{y \in \{x\}} f(y)$$
$$= f(x)$$

for each $x \in X$. This proves our claim. Since we have shown that

$$[\Phi_{X,Y} \circ \Psi_{X,Y}](f) = f$$

for each $f \in \mathsf{Sets}(X,Y)$, it follows that $\Phi_{X,Y} \circ \Psi_{X,Y}$ must be equal to the identity map $\mathsf{id}_{\mathsf{Sets}(X,Y)}$ of $\mathsf{Sets}(X,Y)$.

Naturality for Φ, Part I

We need to show that, given a function $f: X \to X'$, the diagram

$$\begin{split} \mathsf{SupLat}((\mathcal{P}(X'),\subset),(Y,\preceq)) &\xrightarrow{\Phi_{X',Y}} \mathsf{Sets}(X',Y) \\ & & \downarrow^{f^*} \\ \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq)) &\xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y) \end{split}$$

commutes. Indeed, we have

$$\begin{split} [\Phi_{X,Y} \circ \mathcal{P}_{!}(f)^{*}](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y}(\mathcal{P}_{!}(f)^{*}(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}(\xi \circ f_{!}) \\ &\stackrel{\text{def}}{=} (\xi \circ f_{!}) \circ \chi_{X} \\ &= \xi \circ (f_{!} \circ \chi_{X}) \\ &\stackrel{(\uparrow)}{=} \xi \circ (\chi_{X'} \circ f) \\ &= (\xi \circ \chi_{X'}) \circ f \\ &\stackrel{\text{def}}{=} \Phi_{X',Y}(\xi) \circ f \\ &\stackrel{\text{def}}{=} f^{*}(\Phi_{X',Y}(\xi)) \\ &\stackrel{\text{def}}{=} [f^{*} \circ \Phi_{X',Y}](\xi), \end{split}$$

for each $\xi \in \text{SupLat}((\mathcal{P}(X'), \subset), (Y, \preceq))$, where we have used Item 1 of Proposition 4.5.4.1.3 for the fifth equality above.

Naturality for Φ, Part II

We need to show that, given a morphism of suplattices

$$g: (Y, \preceq_Y) \rightarrow (Y', \preceq_{Y'}),$$

the diagram

$$\begin{split} \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y) \\ & & & \downarrow^{g_!} \\ \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y',\preceq)) & \xrightarrow{\Phi_{X,Y'}} \mathsf{Sets}(X,Y') \end{split}$$

commutes. Indeed, we have

$$\begin{split} \big[\Phi_{X,Y'} \circ g_!\big](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g_!(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g \circ \xi) \\ &\stackrel{\text{def}}{=} (g \circ \xi) \circ \chi_X \\ &= g \circ (\xi \circ \chi_X) \\ &\stackrel{\text{def}}{=} g \circ (\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} g_!(\Phi_{X,Y}(\xi)) \end{split}$$

$$\stackrel{\text{def}}{=} [g_! \circ \Phi_{X,Y}](\xi).$$

for each $\xi \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$.

Naturality for Ψ

Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that Ψ is also natural in each argument.

WARNING 4.4.5.1.5 ► FREE COCOMPLETION IS NOT AN IDEMPOTENT OPERATION



Although the assignment $X \mapsto \mathcal{P}(X)$ is called the *free cocompletion of* X, it is not an idempotent operation, i.e. we have $\mathcal{P}(\mathcal{P}(X)) \neq \mathcal{P}(X)$.

4.4.6 Powersets as Free Completions

Let *X* be a set.

PROPOSITION 4.4.6.1.1 ➤ POWERSETS AS FREE COMPLETIONS: UNIVERSAL PROPERTY

The pair $(\mathcal{P}(X), \chi_{(-)})$ consisting of

- The powerset of *X* together with reverse inclusion $\mathcal{P}(X)^{\mathsf{op}} = (\mathcal{P}(X), \supset)$ of Definition 4.4.1.1.1;
- The characteristic embedding $\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$ of X into $\mathcal{P}(X)$ of Definition 4.5.4.1.1;

satisfies the following universal property:

- (\star) Given another pair (Y, f) consisting of
 - An inflattice (Y, \preceq) ;
 - A function $f: X \to Y$;

there exists a unique morphism of inflattices

$$(\mathcal{P}(X),\supset) \xrightarrow{\exists!} (Y,\preceq)$$

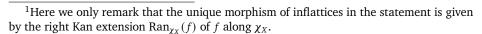
making the diagram



commute.

PROOF 4.4.6.1.2 ► PROOF OF PROPOSITION 4.4.6.1.1

This is a rephrasing of Proposition 4.4.6.1.3, which we prove below.¹



PROPOSITION 4.4.6.1.3 ► POWERSETS AS FREE COMPLETIONS: ADJOINTNESS

We have an adjunction

$$(\mathcal{P} \dashv \overline{\mathbb{k}})$$
: Sets $\underbrace{\hspace{1em}}_{\overline{\mathbb{k}}}$ InfLat,

witnessed by a bijection

$$InfLat((\mathcal{P}(X),\supset),(Y,\preceq)) \cong Sets(X,Y),$$

natural in $X \in \text{Obj}(\mathsf{Sets})$ and $(Y, \preceq) \in \text{Obj}(\mathsf{InfLat})$, where:

- The category InfLat is the category of inflattices of ??.
- The map

$$\chi_X^* : \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

witnessing the above bijection is defined by

$$\chi_X^*(f) \stackrel{\text{def}}{=} f \circ \chi_X,$$

i.e. by sending a morphism of inflattices $f \colon \mathcal{P}(X)^{\mathsf{op}} \to Y$ to the composition

$$X \stackrel{\chi_X}{\hookrightarrow} \mathcal{P}(X)^{\mathsf{op}} \stackrel{f}{\longrightarrow} Y.$$

The map

$$\operatorname{Ran}_{\gamma_X} \colon \operatorname{\mathsf{Sets}}(X,Y) \to \operatorname{\mathsf{InfLat}}((\mathcal{P}(X),\supset),(Y,\preceq))$$

witnessing the above bijection is given by sending a function $f: X \to Y$ to its right Kan extension along χ_X ,

$$\operatorname{Ran}_{\chi_X}(f) \colon \mathcal{P}(X)^{\operatorname{op}} \to Y, \qquad \begin{array}{c} \mathcal{P}(X)^{\operatorname{op}} \\ \chi_X & \downarrow \\$$

Moreover, invoking the bijection $\mathcal{P}(X) \cong \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ of Item 2 of Proposition 4.5.1.1.4, $\mathsf{Ran}_{\chi_X}(f)$ can be explicitly computed by

$$[\operatorname{Ran}_{\chi_X}(f)](U) = \int_{x \in X} \chi_{\mathcal{P}(X)^{\operatorname{op}}}(\chi_x, U) \, \, \, \, \, f(x)$$

$$= \int_{x \in X} \chi_{\mathcal{P}(X)}(U, \chi_x) \, \, \, \, \, f(x)$$

$$= \int_{x \in X} \chi_U(x) \, \, \, \, \, f(x)$$

$$= \left(\bigwedge_{x \in U} \chi_U(x) \, \, \, \, \, f(x) \right) \wedge \left(\bigwedge_{x \in U^c} \chi_U(x) \, \, \, \, \, f(x) \right)$$

$$= \left(\bigwedge_{x \in U} f(x) \right) \wedge \left(\bigwedge_{x \in U^c} \omega_Y \right)$$

$$= \left(\bigwedge_{x \in U} f(x) \right) \wedge \infty_Y$$

$$= \bigwedge_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$, where:

- We have used ?? for the first equality.
- We have used Proposition 4.5.5.1.1 for the second equality.
- We have used ?? for the third equality.

- The symbol \land denotes the meet in (Y, ≤).
- The symbol \pitchfork denotes the cotensor of an element of *Y* by a truth value as in **??**. In particular, we have

true
$$\pitchfork f(x) \stackrel{\text{def}}{=} f(x)$$
, false $\pitchfork f(x) \stackrel{\text{def}}{=} \infty_Y$,

where ∞_Y is the top element of (Y, \preceq) .

In particular, when $(Y, \preceq_Y) = (\mathcal{P}(B), \subset)$ for some set B, the Kan extension $\operatorname{Ran}_{\chi_X}(f)$ is given by

$$[\operatorname{Ran}_{\chi_X}(f)](U) = \bigwedge_{x \in U} f(x)$$
$$= \bigcap_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$.

PROOF 4.4.6.1.4 ► PROOF OF PROPOSITION 4.4.5.1.3

Map I

We define a map

$$\Phi_{X,Y} : \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

as in the statement, i.e. by

$$\Phi_{X,Y}(f)\stackrel{\mathrm{def}}{=} f\circ\chi_X$$

for each $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$.

Map II

We define a map

$$\Psi_{X,Y} : \mathsf{Sets}(X,Y) \to \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq))$$

as in the statement, i.e. by

$$\Psi_{X,Y}(f) \stackrel{\text{def}}{=} \operatorname{Ran}_{\chi_X}(f), \qquad X \xrightarrow{\chi_X} \bigvee_{f} \operatorname{Ran}_{\chi_X}(f)$$

$$X \xrightarrow{f} Y,$$

for each $f \in Sets(X, Y)$.

Invertibility I

We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = id_{\mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq))}$$
.

We have

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) \stackrel{\text{def}}{=} \Psi_{X,Y}(\Phi_{X,Y}(f))$$

$$\stackrel{\text{def}}{=} \Psi_{X,Y}(f \circ \chi_X)$$

$$\stackrel{\text{def}}{=} \operatorname{Ran}_{\chi_X}(f \circ \chi_X)$$

for each $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$. We now claim that

$$\operatorname{Ran}_{\chi_X}(f \circ \chi_X) = f$$

for each $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$. Indeed, we have

$$[\operatorname{Ran}_{\chi_X}(f \circ \chi_X)](U) = \bigwedge_{x \in U} f(\chi_X(x))$$
$$= f\left(\bigwedge_{x \in U} \chi_X(x)\right)$$
$$= f\left(\bigcup_{x \in U} \{x\}\right)$$
$$= f(U)$$

for each $U \in \mathcal{P}(X)$, where we have used that f is a morphism of inflattices and hence preserves meets in $(\mathcal{P}(X), \supset)$ (i.e. joins in $(\mathcal{P}(X), \subset)$) for the second equality. This proves our claim. Since we have shown that

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) = f$$

for each $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$, it follows that $\Psi_{X,Y} \circ \Phi_{X,Y}$ must be equal to the identity map $id_{InfLat((\mathcal{P}(X), \supset), (Y, \preceq))}$ of $InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$.

Invertibility II

We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathsf{Sets}(X,Y)}$$
.

We have

$$[\Phi_{X,Y} \circ \Psi_{X,Y}](f) \stackrel{\text{def}}{=} \Phi_{X,Y}(\Psi_{X,Y}(f))$$
$$\stackrel{\text{def}}{=} \Phi_{X,Y}(\text{Ran}_{\chi_X}(f))$$
$$\stackrel{\text{def}}{=} \text{Ran}_{\gamma_Y}(f) \circ \chi_X$$

for each $f \in Sets(X, Y)$. We now claim that

$$\operatorname{Ran}_{\chi_X}(f) \circ \chi_X = f$$

for each $f \in Sets(X, Y)$. Indeed, we have

$$[\operatorname{Ran}_{\chi_X}(f) \circ \chi_X](x) = \bigwedge_{y \in \{x\}} f(y)$$
$$= f(x)$$

for each $x \in X$. This proves our claim. Since we have shown that

$$[\Phi_{X,Y} \circ \Psi_{X,Y}](f) = f$$

for each $f \in \mathsf{Sets}(X,Y)$, it follows that $\Phi_{X,Y} \circ \Psi_{X,Y}$ must be equal to the identity map $\mathsf{id}_{\mathsf{Sets}(X,Y)}$ of $\mathsf{Sets}(X,Y)$.

Naturality for Φ, Part I

We need to show that, given a function $f: X \to X'$, the diagram

$$\begin{array}{c|c} \mathsf{InfLat}((\mathcal{P}(X'),\supset),(Y,\preceq)) & \xrightarrow{\Phi_{X',Y}} & \mathsf{Sets}(X',Y) \\ & & & \downarrow f^* \\ \\ \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} & \mathsf{Sets}(X,Y) \end{array}$$

commutes. Indeed, we have

$$\begin{split} [\Phi_{X,Y} \circ \mathcal{P}_{!}(f)^{*}](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y}(\mathcal{P}_{!}(f)^{*}(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}(\xi \circ f_{!}) \\ &\stackrel{\text{def}}{=} (\xi \circ f_{!}) \circ \chi_{X} \\ &= \xi \circ (f_{!} \circ \chi_{X}) \\ &\stackrel{(\uparrow)}{=} \xi \circ (\chi_{X'} \circ f) \\ &= (\xi \circ \chi_{X'}) \circ f \\ &\stackrel{\text{def}}{=} \Phi_{X',Y}(\xi) \circ f \\ &\stackrel{\text{def}}{=} f^{*}(\Phi_{X',Y}(\xi)) \\ &\stackrel{\text{def}}{=} [f^{*} \circ \Phi_{X',Y}](\xi), \end{split}$$

for each $\xi \in InfLat((\mathcal{P}(X'), \supset), (Y, \preceq))$, where we have used Item 1 of Proposition 4.5.4.1.3 for the fifth equality above.

Naturality for Φ, Part II

We need to show that, given a cocontinuous morphism of posets

$$g: (Y, \preceq_Y) \rightarrow (Y', \preceq_{Y'}),$$

the diagram

$$\begin{array}{c|c} \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} & \mathsf{Sets}(X,Y) \\ & & & & & \downarrow^{g_!} \\ \\ \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y',\preceq)) & \xrightarrow{\Phi_{X,Y'}} & \mathsf{Sets}(X,Y') \end{array}$$

commutes. Indeed, we have

$$\begin{split} \big[\Phi_{X,Y'} \circ g_!\big](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g_!(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g \circ \xi) \\ &\stackrel{\text{def}}{=} (g \circ \xi) \circ \chi_X \\ &= g \circ (\xi \circ \chi_X) \\ &\stackrel{\text{def}}{=} g \circ (\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} g_!(\Phi_{X,Y}(\xi)) \end{split}$$

$$\stackrel{\text{def}}{=} [g_! \circ \Phi_{X,Y}](\xi).$$

for each $\xi \in InfLat((\mathcal{P}(X),\supset),(Y,\preceq))$.

Naturality for Ψ

Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that Ψ is also natural in each argument.

WARNING 4.4.6.1.5 ► FREE COMPLETION IS NOT AN IDEMPOTENT OPERATION



Although the assignment $X \mapsto \mathcal{P}(X)^{\text{op}}$ is called the *free completion of* X, it is not an idempotent operation, i.e. we have $\mathcal{P}(\mathcal{P}(X)^{\text{op}})^{\text{op}} \neq \mathcal{P}(X)^{\text{op}}$.

4.4.7 The Internal Hom of a Powerset

Let *X* be a set and let $U, V \in \mathcal{P}(X)$.

PROPOSITION 4.4.7.1.1 ► THE INTERNAL HOM OF A POWERSET

The **internal Hom of** $\mathcal{P}(X)$ **from** U **to** V is the subset $[U,V]_X^{-1}$ of X given by

$$[U, V]_X = U^{c} \cup V$$
$$= (U \setminus V)^{c}$$

where U^{c} is the complement of U of Definition 4.3.11.1.1.

¹Further Notation: Also written $\mathbf{Hom}_{\mathcal{P}(X)}(U,V)$.

PROOF 4.4.7.1.2 ► PROOF OF PROPOSITION 4.4.7.1.1

Proof of the Equality $U^{c} \cup V = (U \setminus V)^{c}$

We have

$$(U \setminus V)^{c} \stackrel{\text{def}}{=} X \setminus (U \setminus V)$$
$$= (X \cap V) \cup (X \setminus U)$$

$$= V \cup (X \setminus U)$$

$$\stackrel{\text{def}}{=} V \cup U^{c}$$

$$= U^{c} \cup V,$$

where we have used:

- 1. Item 10 of Proposition 4.3.10.1.2 for the second equality.
- 2. Item 4 of Proposition 4.3.9.1.2 for the third equality.
- 3. Item 4 of Proposition 4.3.8.1.2 for the last equality.

This finishes the proof.

Proof that $U^c \cup V$ Is Indeed the Internal Hom

This follows from Item 2 of Proposition 4.3.9.1.2.

REMARK 4.4.7.1.3 \blacktriangleright Intuition for the Internal Hom of $\mathcal{P}(X)$

Henning Makholm suggests the following heuristic intuition for the internal Hom of $\mathcal{P}(X)$ from U to V ([MSE 267365]):

- 1. Since products in $\mathcal{P}(X)$ are given by binary intersections (Item 1 of Proposition 4.4.1.1.4), the right adjoint $\mathbf{Hom}_{\mathcal{P}(X)}(U, -)$ of $U \cap -$ may be thought of as a function type [U, V].
- 2. Under the Curry–Howard correspondence (??), the function type [U,V] corresponds to implication $U \Rightarrow V$.
- 3. Implication $U \Rightarrow V$ is logically equivalent to $\neg U \lor V$.
- 4. The expression $\neg U \lor V$ then corresponds to the set $U^c \cup V$ in $\mathcal{P}(X)$.
- 5. The set $U^{c} \vee V$ turns out to indeed be the internal Hom of $\mathcal{P}(X)$.

PROPOSITION 4.4.7.1.4 ▶ PROPERTIES OF INTERNAL HOMS OF POWERSETS

Let X be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto \mathbf{Hom}_{\mathcal{P}(X)}$ define func-

tors

$$[U, -]_X: \qquad (\mathcal{P}(X), \supset) \qquad \to (\mathcal{P}(X), \subset),$$

$$[-, V]_X: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$[-_1, -_2]_X: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \supset) \to (\mathcal{P}(X), \subset).$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $[A, V]_X \subset [U, V]_X$.
- (b) If $V \subset B$, then $[U, V]_X \subset [U, B]_X$.
- (c) If $U \subset A$ and $V \subset B$, then $[A, V]_X \subset [U, B]_X$.
- 2. Adjointness. We have adjunctions

$$(U \cap - + [U, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{U \cap -}{\bot}} \mathcal{P}(X),$$

$$(- \cap V + [V, -]_X): \quad \mathcal{P}(X) \xrightarrow{\stackrel{- \cap V}{\bot}} \mathcal{P}(X),$$

witnessed by bijections

$$\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, [V, W]_X),$$

 $\operatorname{Hom}_{\mathcal{P}(X)}(U \cap V, W) \cong \operatorname{Hom}_{\mathcal{P}(X)}(V, [U, W]_X).$

In particular, the following statements hold for each $U, V, W \in \mathcal{P}(X)$:

- (a) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $U \subset [V, W]_X$.
- (b) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $V \subset [U, W]_X$.
- 3. Interaction With the Empty Set I. We have

$$[U, \emptyset]_X = U^{\mathsf{c}},$$

 $[\emptyset, V]_X = X,$

natural in $U, V \in \mathcal{P}(X)$.

4. Interaction With X. We have

$$[U, X]_X = X,$$
$$[X, V]_X = V,$$

natural in $U, V \in \mathcal{P}(X)$.

5. Interaction With the Empty Set II. The functor

$$D_X \colon \mathcal{P}(X)^{\mathsf{op}} \to \mathcal{P}(X)$$

defined by

$$D_X \stackrel{\text{def}}{=} [-, \emptyset]_X$$
$$= (-)^{\mathsf{c}}$$

is an involutory isomorphism of categories, making \emptyset into a dualising object for $(\mathcal{P}(X), \cap, X, [-, -]_X)$ in the sense of **??**. In particular:

(a) The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{D_X} \mathcal{P}(X)$$

$$\mathsf{id}_{\mathcal{P}(X)^{\mathsf{op}}} \qquad \qquad \qquad \downarrow^{D_X}$$

$$\mathcal{P}(X)^{\mathsf{op}}$$

commutes, i.e. we have

$$\underbrace{D_X(D_X(U))}_{\substack{\text{def}\\=[[U,\emptyset]_X,\emptyset]_X}} = U$$

for each $U \in \mathcal{P}(X)$.

(b) The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{\cap^{\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}} \\
\downarrow^{\mathsf{id}_{\mathcal{P}(X)^{\mathsf{op}}} \times D_X} & & & \downarrow^{D_X} \\
\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) & & & & \downarrow^{D_X}
\end{array}$$

commutes, i.e. we have

$$\underbrace{D_X(U\cap D_X(V))}_{\stackrel{\mathrm{def}}{=}[U\cap [V,\emptyset]_X,\emptyset]_X} = [U,V]_X$$

for each $U, V \in \mathcal{P}(X)$.

- 6. *Interaction With the Empty Set III.* Let $f: X \to Y$ be a function.
 - (a) Interaction With Direct Images. The diagram

$$\begin{array}{ccc} \mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_{\ast}^{\mathsf{op}}} & \mathcal{P}(Y)^{\mathsf{op}} \\ & & \downarrow^{D_{X}} & & \downarrow^{D_{Y}} \\ \mathcal{P}(X) & \xrightarrow{f_{1}} & \mathcal{P}(Y) \end{array}$$

commutes, i.e. we have

$$f_!(D_X(U)) = D_Y(f_*(U))$$

for each $U \in \mathcal{P}(X)$.

(b) Interaction With Inverse Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(Y)^{\mathsf{op}} & \xrightarrow{f^{-1},\mathsf{op}} & \mathcal{P}(X)^{\mathsf{op}} \\
\downarrow^{D_{Y}} & & \downarrow^{D_{X}} \\
\mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$f^{-1}(D_Y(U)) = D_X(f^{-1}(U))$$

for each $U \in \mathcal{P}(X)$.

(c) Interaction With Codirect Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\text{op}} & \xrightarrow{f_{!}^{\text{op}}} & \mathcal{P}(Y)^{\text{op}} \\
\downarrow D_{X} & & \downarrow D_{Y} \\
\mathcal{P}(X) & \xrightarrow{f_{*}} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(D_X(U)) = D_Y(f_!(U))$$

for each $U \in \mathcal{P}(X)$.

7. Interaction With Unions of Families of Subsets I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\cup^{\mathsf{op}} \times \cup^{\mathsf{op}} \bigvee \qquad \qquad \bigcup \cup$$

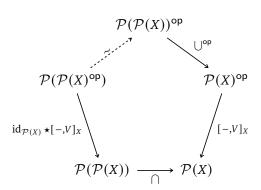
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcup_{U \in \mathcal{U}} U, \bigcup_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

8. Interaction With Unions of Families of Subsets II. The diagram



commutes, i.e. we have

$$\left[\bigcup_{U\in\mathcal{U}}U,V\right]_X=\bigcap_{U\in\mathcal{U}}[U,V]_X$$

for each $U \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

9. Interaction With Unions of Families of Subsets III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcup}{\longrightarrow} & \mathcal{P}(X) \\
id_{\mathcal{P}(X)} \star [U,-]_X & & & \downarrow [U,-]_X \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\downarrow \downarrow} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U, \bigcup_{V \in \mathcal{V}} V\right]_X = \bigcup_{V \in \mathcal{V}} [U, V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

10. Interaction With Intersections of Families of Subsets I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\cap^{\mathsf{op}} \times \cap^{\mathsf{op}} \downarrow \qquad \qquad \downarrow \cap$$

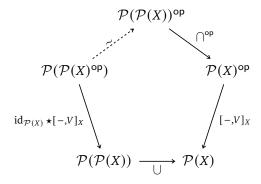
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcap_{U \in \mathcal{U}} U, \bigcap_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

11. Interaction With Intersections of Families of Subsets II. The diagram



commutes, i.e. we have

$$\left[\bigcap_{U\in\mathcal{U}}U,V\right]_X=\bigcup_{U\in\mathcal{U}}[U,V]_X$$

for each $U \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

12. Interaction With Intersections of Families of Subsets III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcap} & \mathcal{P}(X) \\
id_{\mathcal{P}(X)} \star [U,-]_X & & & \downarrow [U,-]_X \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\bigcap} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U,\bigcap_{V\in\mathcal{V}}V\right]_X=\bigcap_{V\in\mathcal{V}}[U,V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

13. Interaction With Binary Unions. We have equalities of sets

$$[U \cap V, W]_X = [U, W]_X \cup [V, W]_X,$$

 $[U, V \cap W]_X = [U, V]_X \cap [U, W]_X$

for each $U, V, W \in \mathcal{P}(X)$.

14. Interaction With Binary Intersections. We have equalities of sets

$$[U \cup V, W]_X = [U, W]_X \cap [V, W]_X,$$

 $[U, V \cup W]_X = [U, V]_X \cup [U, W]_X$

for each $U, V, W \in \mathcal{P}(X)$.

15. Interaction With Differences. We have equalities of sets

$$[U \setminus V, W]_X = [U, W]_X \cup [V^{c}, W]_X$$
$$= [U, W]_X \cup [U, V]_X,$$
$$[U, V \setminus W]_X = [U, V]_X \setminus (U \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

16. Interaction With Complements. We have equalities of sets

$$[U^{c}, V]_{X} = U \cup V,$$

$$[U, V^{c}]_{X} = U \cap V,$$

$$[U, V]_{X}^{c} = U \setminus V$$

for each $U, V \in \mathcal{P}(X)$.

17. Interaction With Characteristic Functions. We have

$$\chi_{[U,V]_{\mathcal{P}(X)}}(x) = \max(1 - \chi_U \pmod{2}, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

18. *Interaction With Direct Images.* Let $f: X \to Y$ be a function. The diagram

commutes, i.e. we have an equality of sets

$$f_!([U,V]_X) = [f_*(U), f_!(V)]_Y,$$

natural in $U, V \in \mathcal{P}(X)$.

19. *Interaction With Inverse Images.* Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1}, \mathsf{op} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$[-1, -2]_Y \downarrow \qquad \qquad \downarrow [-1, -2]_X$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_X,$$

natural in $U, V \in \mathcal{P}(X)$.

20. *Interaction With Codirect Images.* Let $f: X \to Y$ be a function. We have a natural transformation

with components

$$[f_!(U), f_*(V)]_Y \subset f_*([U, V]_X)$$

indexed by $U, V \in \mathcal{P}(X)$.

PROOF 4.4.7.1.5 ► PROOF OF PROPOSITION 4.4.7.1.4

Item 1: Functoriality

Since $\mathcal{P}(X)$ is posetal, it suffices to prove Items 1a to 1c.

1. Proof of Item 1a: We have

$$[A, V]_X \stackrel{\text{def}}{=} A^{c} \cup V$$

$$\subset U^{c} \cup V$$

$$\stackrel{\text{def}}{=} [U, V]_X,$$

where we have used:

- (a) Item 1 of Proposition 4.3.11.1.2, which states that if $U \subset A$, then $A^{c} \subset U^{c}$.
- (b) Item 1a of Item 1 of Proposition 4.3.11.1.2, which states that if $A^c \subset U^c$, then $A^c \cup K \subset U^c \cup K$ for any $K \in \mathcal{P}(X)$.
- 2. Proof of Item 1b: We have

$$\begin{split} [U,V]_X &\stackrel{\text{def}}{=} U^\mathsf{c} \cup V \\ &\subset U^\mathsf{c} \cup B \\ &\stackrel{\text{def}}{=} [U,B]_X, \end{split}$$

where we have used Item 1b of Item 1 of Proposition 4.3.11.1.2, which states that if $V \subset B$, then $K \cup V \subset K \cup B$ for any $K \in \mathcal{P}(X)$.

3. Proof of Item 1c: We have

$$[A,V]_X \subset [U,V]_X$$
$$\subset [U,B]_X,$$

where we have used Items 1a and 1b.

This finishes the proof.

Item 2: Adjointness

This is a repetition of Item 2 of Proposition 4.3.9.1.2 and is proved there.

Item 3: Interaction With the Empty Set I

We have

$$[U,\emptyset]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup \emptyset$$
$$= U^{\mathsf{c}},$$

where we have used Item 3 of Proposition 4.3.8.1.2, and we have

$$[\emptyset, V]_X \stackrel{\text{def}}{=} \emptyset^{c} \cup V$$

$$\stackrel{\text{def}}{=} (X \setminus \emptyset) \cup V$$

$$= X \cup V$$

$$= X,$$

where we have used:

- 1. Item 12 of Proposition 4.3.10.1.2 for the first equality.
- 2. Item 5 of Proposition 4.3.8.1.2 for the last equality.

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2).

Item 4: Interaction With *X*

We have

$$[U,X]_X \stackrel{\mathrm{def}}{=} U^{\mathsf{c}} \cup X$$

$$=X$$
,

where we have used Item 5 of Proposition 4.3.8.1.2, and we have

$$[X, V]_X \stackrel{\text{def}}{=} X^{c} \cup V$$

$$\stackrel{\text{def}}{=} (X \setminus X) \cup V$$

$$= \emptyset \cup V$$

$$= V,$$

where we have used Item 3 of Proposition 4.3.8.1.2 for the last equality. Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2).

Item 5: Interaction With the Empty Set II

We have

$$D_X(D_X(U)) \stackrel{\text{def}}{=} [[U, \emptyset]_X, \emptyset]_X$$
$$= [U^c, \emptyset]_X$$
$$= (U^c)^c$$
$$= U,$$

where we have used:

- 1. Item 3 for the second and third equalities.
- 2. Item 3 of Proposition 4.3.11.1.2 for the fourth equality.

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2), and thus we have

$$[[-,\emptyset]_X,\emptyset]_X \cong \mathrm{id}_{\mathcal{P}(X)}$$

This finishes the proof.

Item 6: Interaction With the Empty Set III

Since $D_X = (-)^c$, this is essentially a repetition of the corresponding results for $(-)^c$, namely Items 5 to 7 of Proposition 4.3.11.1.2.

Item 7: Interaction With Unions of Families of Subsets I

By Item 3 of Proposition 4.4.7.1.4, we have

$$[\mathcal{U}, \emptyset]_{\mathcal{P}(X)} = \mathcal{U}^{\mathsf{c}},$$

 $[\mathcal{U}, \emptyset]_X = \mathcal{U}^{\mathsf{c}}.$

With this, the counterexample given in the proof of Item 10 of Proposition 4.3.6.1.2 then applies.

Item 8: Interaction With Unions of Families of Subsets II

We have

$$\left[\bigcup_{U \in \mathcal{U}} U, V\right]_{X}^{\text{def}} \left(\bigcup_{U \in \mathcal{U}} U\right)^{c} \cup V$$

$$= \left(\bigcap_{U \in \mathcal{U}} U^{c}\right) \cup V$$

$$= \bigcap_{U \in \mathcal{U}} (U^{c} \cup V)$$

$$\stackrel{\text{def}}{=} \bigcap_{U \in \mathcal{U}} [U, V]_{X},$$

where we have used:

- 1. Item 11 of Proposition 4.3.6.1.2 for the second equality.
- 2. Item 6 of Proposition 4.3.7.1.2 for the third equality.

This finishes the proof.

Item 9: Interaction With Unions of Families of Subsets III

We have

$$\bigcup_{V \in \mathcal{V}} [U, V]_X \stackrel{\text{def}}{=} \bigcup_{V \in \mathcal{V}} (U^c \cup V)$$
$$= U^c \cup \left(\bigcup_{V \in \mathcal{V}} V\right)$$
$$\stackrel{\text{def}}{=} \left[U, \bigcup_{V \in \mathcal{V}} V\right]_X.$$

where we have used Item 6. This finishes the proof.

Item 10: Interaction With Intersections of Families of Subsets I

Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}\$, and let $\mathcal{V} = \{\{0\}, \{0, 1\}\}\$. We have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W = \bigcap_{W \in \mathcal{P}(X)} W$$
$$= \{0, 1\},$$

whereas

$$\left[\bigcap_{U\in\mathcal{U}}U,\bigcap_{V\in\mathcal{V}}V\right]_X=\left[\{0,1\},\{0\}\right]$$
$$=\{0\},$$

Thus we have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W = \{0, 1\} \neq \{0\} = \left[\bigcap_{U \in \mathcal{U}} U, \bigcap_{V \in \mathcal{V}} V\right]_X.$$

This finishes the proof.

Item 11: Interaction With Intersections of Families of Subsets II

We have

$$\left[\bigcap_{U \in \mathcal{U}} U, V\right]_{X} \stackrel{\text{def}}{=} \left(\bigcap_{U \in \mathcal{U}} U\right)^{c} \cup V$$

$$= \left(\bigcup_{U \in \mathcal{U}} U^{c}\right) \cup V$$

$$= \bigcup_{U \in \mathcal{U}} (U^{c} \cup V)$$

$$\stackrel{\text{def}}{=} \bigcup_{U \in \mathcal{U}} [U, V]_{X},$$

where we have used:

- 1. Item 12 of Proposition 4.3.6.1.2 for the second equality.
- 2. Item 6 of Proposition 4.3.7.1.2 for the third equality.

This finishes the proof.

Item 12: Interaction With Intersections of Families of Subsets III

We have

$$\bigcap_{V \in \mathcal{V}} [U, V]_X \stackrel{\text{def}}{=} \bigcap_{V \in \mathcal{V}} (U^c \cup V)$$

$$= U^c \cup \left(\bigcap_{V \in \mathcal{V}} V\right)$$

$$\stackrel{\text{def}}{=} \left[U, \bigcap_{V \in \mathcal{V}} V\right]_Y$$

where we have used Item 6. This finishes the proof.

Item 13: Interaction With Binary Unions

We have

$$[U \cap V, W]_X \stackrel{\text{def}}{=} (U \cap V)^c \cup W$$

$$= (U^c \cup V^c) \cup W$$

$$= (U^c \cup V^c) \cup (W \cup W)$$

$$= (U^c \cup W) \cup (V^c \cup W)$$

$$\stackrel{\text{def}}{=} [U, W]_X \cup [V, W]_X,$$

where we have used:

- 1. Item 2 of Proposition 4.3.11.1.2 for the second equality.
- 2. Item 8 of Proposition 4.3.8.1.2 for the third equality.
- 3. Several applications of Items 2 and 4 of Proposition 4.3.8.1.2 and for the fourth equality.

For the second equality in the statement, we have

$$[U, V \cap W]_X \stackrel{\text{def}}{=} U^{c} \cup (V \cap W)$$
$$= (U^{c} \cup V) \cap (U^{c} \cap W)$$
$$\stackrel{\text{def}}{=} [U, V]_X \cap [U, W]_X,$$

where we have used Item 6 of Proposition 4.3.8.1.2 for the second equality.

Item 14: Interaction With Binary Intersections

We have

$$[U \cup V, W]_X \stackrel{\text{def}}{=} (U \cup V)^c \cup W$$
$$= (U^c \cap V^c) \cup W$$
$$= (U^c \cup W) \cap (V^c \cup W)$$
$$\stackrel{\text{def}}{=} [U, W]_X \cap [V, W]_X,$$

where we have used:

- 1. Item 2 of Proposition 4.3.11.1.2 for the second equality.
- 2. Item 6 of Proposition 4.3.8.1.2 for the third equality.

Now, for the second equality in the statement, we have

$$\begin{split} [U, V \cup W]_X &\stackrel{\text{def}}{=} U^{\mathsf{c}} \cup (V \cup W) \\ &= (U^{\mathsf{c}} \cup U^{\mathsf{c}}) \cup (V \cup W) \\ &= (U^{\mathsf{c}} \cup V) \cup (U^{\mathsf{c}} \cup W) \\ &\stackrel{\text{def}}{=} [U, V]_X \cup [U, W]_X, \end{split}$$

where we have used:

- 1. Item 8 of Proposition 4.3.8.1.2 for the second equality.
- 2. Several applications of Items 2 and 4 of Proposition 4.3.8.1.2 and for the third equality.

This finishes the proof.

Item 15: Interaction With Differences

We have

$$[U \setminus V, W]_X \stackrel{\text{def}}{=} (U \setminus V)^{c} \cup W$$

$$\stackrel{\text{def}}{=} (X \setminus (U \setminus V)) \cup W$$

$$= ((X \cap V) \cup (X \setminus U)) \cup W$$

$$= (V \cup (X \setminus U)) \cup W$$

$$\stackrel{\text{def}}{=} (V \cup U^{c}) \cup W$$

$$= (V \cup (U^{c} \cup U^{c})) \cup W$$
$$= (U^{c} \cup W) \cup (U^{c} \cup V)$$
$$\stackrel{\text{def}}{=} [U, W]_{X} \cup [U, V]_{X},$$

where we have used:

- 1. Item 10 of Proposition 4.3.10.1.2 for the third equality.
- 2. Item 4 of Proposition 4.3.9.1.2 for the fourth equality.
- 3. Item 8 of Proposition 4.3.8.1.2 for the sixth equality.
- 4. Several applications of Items 2 and 4 of Proposition 4.3.8.1.2 and for the seventh equality.

We also have

$$\begin{split} [U \setminus V, W]_X &\stackrel{\text{def}}{=} (U \setminus V)^c \cup W \\ &\stackrel{\text{def}}{=} (X \setminus (U \setminus V)) \cup W \\ &= ((X \cap V) \cup (X \setminus U)) \cup W \\ &= (V \cup (X \setminus U)) \cup W \\ &\stackrel{\text{def}}{=} (V \cup U^c) \cup W \\ &= (V \cup U^c) \cup (W \cup W) \\ &= (U^c \cup W) \cup ((V^c)^c \cup W) \\ &\stackrel{\text{def}}{=} [U, W]_X \cup [V^c, W]_X, \end{split}$$

where we have used:

- 1. Item 10 of Proposition 4.3.10.1.2 for the third equality.
- 2. Item 4 of Proposition 4.3.9.1.2 for the fourth equality.
- 3. Item 8 of Proposition 4.3.8.1.2 for the sixth equality.
- 4. Several applications of Items 2 and 4 of Proposition 4.3.8.1.2 and for the seventh equality.
- 5. Item 3 of Proposition 4.3.11.1.2 for the eighth equality.

Now, for the second equality in the statement, we have

$$\begin{split} [U, V \setminus W]_X &\stackrel{\text{def}}{=} U^c \cup (V \setminus W) \\ &= (V \setminus W) \cup U^c \\ &= (V \cup U^c) \setminus (W \setminus U^c) \\ &\stackrel{\text{def}}{=} (V \cup U^c) \setminus (W \setminus (X \setminus U)) \\ &= (V \cup U^c) \setminus ((W \cap U) \cup (W \setminus X)) \\ &= (V \cup U^c) \setminus ((W \cap U) \cup \emptyset) \\ &= (V \cup U^c) \setminus (W \cap U) \\ &= (V \cup U^c) \setminus (U \cap W) \\ &\stackrel{\text{def}}{=} [U, V]_X \setminus (U \cap W) \end{split}$$

where we have used:

- 1. Item 4 of Proposition 4.3.8.1.2 for the second equality.
- 2. Item 4 of Proposition 4.3.10.1.2 for the third equality.
- 3. Item 10 of Proposition 4.3.10.1.2 for the fifth equality.
- 4. Item 13 of Proposition 4.3.10.1.2 for the sixth equality.
- 5. Item 3 of Proposition 4.3.8.1.2 for the seventh equality.
- 6. Item 5 of Proposition 4.3.9.1.2 for the eighth equality.

This finishes the proof.

Item 16: Interaction With Complements

We have

$$[U^{c}, V]_{X} \stackrel{\text{def}}{=} (U^{c})^{c} \cup V,$$
$$= U \cup V,$$

where we have used Item 3 of Proposition 4.3.11.1.2. We also have

$$[U, V^{c}]_{X} \stackrel{\text{def}}{=} U^{c} \cup V^{c}$$
$$= U \cap V$$

where we have used Item 2 of Proposition 4.3.11.1.2. Finally, we have

$$[U, V]_X^{c} = ((U \setminus V)^{c})^{c}$$
$$= U \setminus V,$$

where we have used Item 2 of Proposition 4.3.11.1.2.

Item 17: Interaction With Characteristic Functions

We have

$$\chi_{[U,V]_{\mathcal{P}(X)}}(x) \stackrel{\text{def}}{=} \chi_{U^{c} \cup V}(x)$$

$$= \max(\chi_{U^{c}}, \chi_{V})$$

$$= \max(1 - \chi_{U} \pmod{2}, \chi_{V}),$$

where we have used:

- 1. Item 10 of Proposition 4.3.8.1.2 for the second equality.
- 2. Item 4 of Proposition 4.3.11.1.2 for the third equality.

This finishes the proof.

Item 18: Interaction With Direct Images

This is a repetition of Item 10 of Proposition 4.6.1.1.5 and is proved there.

Item 19: Interaction With Inverse Images

This is a repetition of Item 10 of Proposition 4.6.2.1.3 and is proved there.

Item 20: Interaction With Codirect Images

This is a repetition of Item 9 of Proposition 4.6.3.1.7 and is proved there.

4.4.8 Isbell Duality for Sets

Let *X* be a set.

DEFINITION 4.4.8.1.1 ► THE ISBELL FUNCTION

The **Isbell function** of *X* is the map

$$I: \mathcal{P}(X) \to \mathsf{Sets}(X, \mathcal{P}(X))$$

defined by

$$\mathsf{I}(U) \stackrel{\mathrm{def}}{=} \llbracket x \mapsto [U, \{x\}]_X \rrbracket$$

for each $U \in \mathcal{P}(X)$.

REMARK 4.4.8.1.2 ► MOTIVATION FOR THE ISBELL FUNCTION

Recall from Remark 4.4.1.1.2 that we may view the powerset $\mathcal{P}(X)$ of a set X as the decategorification of the category of presheaves $\mathsf{PSh}(C)$ of a category C. Building upon this analogy, we want to mimic the definition of the Isbell Spec functor, which is given on objects by

$$\mathsf{Spec}(\mathcal{F}) \stackrel{\mathsf{def}}{=} \mathsf{Nat}(\mathcal{F}, h_{(-)})$$

for each $\mathcal{F} \in \mathsf{Obj}(\mathsf{PSh}(C))$. To this end, we could define

$$I(U) \stackrel{\text{def}}{=} [U, \chi_{(-)}]_X,$$

replacing:

- The Yoneda embedding $X \mapsto h_X$ of C into PSh(C) with the characteristic embedding $x \mapsto \chi_x$ of X into $\mathcal{P}(X)$ of Definition 4.5.4.1.1.
- The internal Hom Nat of PSh(C) with the internal Hom $[-,-]_X$ of $\mathcal{P}(X)$ of Proposition 4.4.7.1.1.

However, since $[U, \chi_x]_X$ is a subset of U instead of a truth value, we get a function

$$I: \mathcal{P}(X) \to \mathsf{Sets}(X, \mathcal{P}(X))$$

instead of a function

$$I: \mathcal{P}(X) \to \mathcal{P}(X)$$
.

This makes some of the properties involving I a bit more cumbersome to state, although we still have an analogue of Isbell duality in that I! \circ I evaluates to $id_{\mathcal{P}(X)}$ in the sense of Proposition 4.4.8.1.3.

PROPOSITION 4.4.8.1.3 ► ISBELL DUALITY FOR SETS

The diagram

$$\mathcal{P}(X) \xrightarrow{\mathsf{I}} \mathsf{Sets}(X, \mathcal{P}(X))$$

$$\Delta_{\Delta_{\mathrm{id}_{\mathcal{P}(X)}}} \qquad \qquad \mathsf{I}_{!}$$

$$\mathsf{Sets}(X, \mathsf{Sets}(X, \mathcal{P}(X)))$$

commutes, i.e. we have

$$l_!(\mathsf{I}(U)) = \llbracket x \mapsto \llbracket y \mapsto U \rrbracket \rrbracket$$

for each $U \in \mathcal{P}(X)$.

PROOF 4.4.8.1.4 ► PROOF OF PROPOSITION 4.4.8.1.3

We have

$$I_{!}(I(U)) \stackrel{\text{def}}{=} I_{!}(\llbracket x \mapsto U^{c} \cup \{x\} \rrbracket)$$

$$\stackrel{\text{def}}{=} \llbracket x \mapsto I(U^{c} \cup \{x\}) \rrbracket$$

$$\stackrel{\text{def}}{=} \llbracket x \mapsto \llbracket y \mapsto (U^{c} \cup \{x\})^{c} \cup \{x\} \rrbracket \rrbracket$$

$$= \llbracket x \mapsto \llbracket y \mapsto (U \cap (X \setminus \{x\})) \cup \{x\} \rrbracket \rrbracket$$

$$= \llbracket x \mapsto \llbracket y \mapsto (U \setminus \{x\}) \cup \{x\} \rrbracket \rrbracket$$

$$= \llbracket x \mapsto \llbracket y \mapsto U \setminus \{x\} \end{bmatrix}$$

where we have used Item 2 of Proposition 4.3.11.1.2 for the fourth equality above.

4.5 Characteristic Functions

4.5.1 The Characteristic Function of a Subset

Let *X* be a set and let $U \in \mathcal{P}(X)$.

DEFINITION 4.5.1.1.1 ► THE CHARACTERISTIC FUNCTION OF A SUBSET

The **characteristic function of** U^1 is the function $\chi_U: X \to \{t, f\}^2$ defined by

$$\chi_U(x) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x \in U, \\ \text{false} & \text{if } x \notin U \end{cases}$$

for each $x \in X$.

¹Further Terminology: Also called the **indicator function of** *U*.

²Further Notation: Also written $\chi_X(U, -)$ or $\chi_X(-, U)$.

REMARK 4.5.1.1.2 ► CHARACTERISTIC FUNCTIONS OF SUBSETS AS DECATEGORIFICATIONS OF PRESHEAVES

Under the analogy that $\{t, f\}$ should be the (-1)-categorical analogue of Sets, we may view a function

$$f: X \to \{\mathsf{t},\mathsf{f}\}$$

as a decategorification of presheaves and copresheaves

$$\mathcal{G}: C^{\mathsf{op}} \to \mathsf{Sets},$$

 $F: C \to \mathsf{Sets}.$

The characteristic functions χ_U of the subsets of X are then the primordial examples of such functions (and, in fact, all of them).

NOTATION 4.5.1.1.3 ► FURTHER NOTATION FOR CHARACTERISTIC FUNCTIONS

We will often employ the bijection $\{t, f\} \cong \{0, 1\}$ to make use of the arithmetical operations defined on $\{0, 1\}$ when disucssing characteristic functions. Examples of this include Items 4 to 11 of Proposition 4.5.1.1.4 below.

PROPOSITION 4.5.1.1.4 ► PROPERTIES OF CHARACTERISTIC FUNCTIONS OF SUBSETS

Let *X* be a set.

1. Functionality. The assignment $U \mapsto \chi_U$ defines a function

$$\chi_{(-)} \colon \mathcal{P}(X) \to \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\}).$$

- 2. *Bijectivity*. The function $\chi_{(-)}$ from Item 1 is bijective.
- 3. Naturality. The collection

$$\left\{\chi_{(-)}\colon \mathcal{P}(X)\to \mathsf{Sets}(X,\{\mathsf{t},\mathsf{f}\})\right\}_{X\in \mathsf{Obj}(\mathsf{Sets})}$$

defines a natural isomorphism between \mathcal{P}^{-1} and Sets(-, {t, f}). In particular, given a function $f: X \to Y$, the diagram

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

$$\chi_{(-)} \downarrow \downarrow \qquad \qquad \downarrow \chi_{(-)}$$

$$\mathsf{Sets}(Y, \{\mathsf{t}, \mathsf{f}\}) \xrightarrow{f^*} \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

commutes, i.e. we have

$$\chi_V \circ f = \chi_{f^{-1}(V)}$$

for each $V \in \mathcal{P}(Y)$.

4. Interaction With Unions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

5. Interaction With Unions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each $U, V \in \mathcal{P}(X)$.

6. Interaction With Intersections I. We have

$$\chi_{U\cap V} = \chi_U \chi_V$$

for each $U, V \in \mathcal{P}(X)$.

7. Interaction With Intersections II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

8. Interaction With Differences. We have

$$\chi_{U\setminus V} = \chi_U - \chi_{U\cap V}$$

for each $U, V \in \mathcal{P}(X)$.

9. Interaction With Complements. We have

$$\chi_{U^{c}} = 1 - \chi_{U}$$

for each $U \in \mathcal{P}(X)$.

10. Interaction With Symmetric Differences. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \pmod{2}$$

for each $U, V \in \mathcal{P}(X)$.

11. Interaction With Internal Homs. We have

$$\chi_{[U,V]_{\mathcal{P}(X)}} = \max(1 - \chi_U \pmod{2}, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

PROOF 4.5.1.1.5 ► PROOF OF PROPOSITION 4.5.1.1.4

Item 1: Functionality

There is nothing to prove.

Item 2: Bijectivity

We proceed in three steps:

1. The Inverse of $\chi_{(-)}$. The inverse of $\chi_{(-)}$ is the map

$$\Phi \colon \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\}) \xrightarrow{\sim} \mathcal{P}(X),$$

defined by

$$\begin{split} \Phi(f) &\stackrel{\text{def}}{=} U_f \\ &\stackrel{\text{def}}{=} f^{-1}(\mathsf{true}) \\ &\stackrel{\text{def}}{=} \{x \in X \mid f(x) = \mathsf{true}\} \end{split}$$

for each $f \in Sets(X, \{t, f\})$.

2. Invertibility I. We have

$$[\Phi \circ \chi_{(-)}](U) \stackrel{\text{def}}{=} \Phi(\chi_U)$$

$$\stackrel{\text{def}}{=} \chi_U^{-1}(\text{true})$$

$$\stackrel{\text{def}}{=} \{x \in X \mid \chi_U(x) = \text{true}\}$$

$$\stackrel{\text{def}}{=} \{x \in X \mid x \in U\}$$

$$= U$$

$$\stackrel{\text{def}}{=} [\text{id}_{\mathcal{P}(X)}](U)$$

for each $U \in \mathcal{P}(X)$. Thus, we have

$$\Phi \circ \chi_{(-)} = \mathrm{id}_{\mathcal{P}(X)} .$$

3. Invertibility II. We have

$$\begin{split} & [\chi_{(-)} \circ \Phi](U) \overset{\text{def}}{=} \chi_{\Phi(f)} \\ & \overset{\text{def}}{=} \chi_{f^{-1}(\mathsf{true})} \\ & \overset{\text{def}}{=} [\![x \mapsto \begin{cases} \mathsf{true} & \text{if } x \in f^{-1}(\mathsf{true}) \\ \mathsf{false} & \text{otherwise} \end{cases}] \\ & = [\![x \mapsto f(x)]\!] \\ & = f \\ & \overset{\text{def}}{=} [id_{\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})}](f) \end{split}$$

for each $f \in Sets(X, \{t, f\})$. Thus, we have

$$\chi_{(-)} \circ \Phi = \mathrm{id}_{\mathsf{Sets}(X, \{\mathsf{t},\mathsf{f}\})}$$
.

This finishes the proof.

Item 3: Naturality

We proceed in two steps:

1. *Naturality of* $\chi_{(-)}$. We have

$$[\chi_V \circ f](v) \stackrel{\text{def}}{=} \chi_V(f(v))$$

$$= \begin{cases} \text{true} & \text{if } f(v) \in V, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } v \in f^{-1}(V), \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_{f^{-1}(V)}(v)$$

for each $v \in V$.

2. Naturality of Φ . Since $\chi_{(-)}$ is natural and a componentwise inverse to Φ , it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that Φ is also natural in each argument.

This finishes the proof.

Item 4: Interaction With Unions I

This is a repetition of Item 10 of Proposition 4.3.8.1.2 and is proved there.

Item 5: Interaction With Unions II

This is a repetition of Item 11 of Proposition 4.3.8.1.2 and is proved there.

Item 6: Interaction With Intersections I

This is a repetition of Item 10 of Proposition 4.3.9.1.2 and is proved there.

Item 7: Interaction With Intersections II

This is a repetition of Item 11 of Proposition 4.3.9.1.2 and is proved there.

Item 8: Interaction With Differences

This is a repetition of Item 16 of Proposition 4.3.10.1.2 and is proved there.

Item 9: Interaction With Complements

This is a repetition of Item 4 of Proposition 4.3.11.1.2 and is proved there.

Item 10: Interaction With Symmetric Differences

This is a repetition of Item 15 of Proposition 4.3.12.1.2 and is proved there.

Item 11: Interaction With Internal Homs

This is a repetition of Item 17 of Proposition 4.4.7.1.4 and is proved there.

REMARK 4.5.1.1.6 ▶ POWERSETS AS SETS OF FUNCTIONS AND UN/STRAIGHTENING

The bijection

$$\mathcal{P}(X) \cong \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

of Item 2 of Proposition 4.5.1.1.4, which

- Takes a subset $U \hookrightarrow X$ of X and *straightens* it to a function $\chi_U \colon X \to \{\text{true}, \text{false}\};$
- Takes a function $f: X \to \{\text{true}, \text{false}\}\$ and *unstraightens* it to a subset $f^{-1}(\text{true}) \hookrightarrow X \text{ of } X;$

may be viewed as the (-1)-categorical version of the 0-categorical un/straightening isomorphism between indexed and fibred sets

$$\underbrace{\mathsf{FibSets}_X}_{\substack{\mathsf{def}\\ = \mathsf{Sets}_{/X}}} \cong \underbrace{\mathsf{ISets}_X}_{\substack{\mathsf{def}\\ = \mathsf{Fun}(X_{\mathsf{disc}}, \mathsf{Sets})}}$$

of Un/Straightening for Indexed and Fibred Sets, ??. Here we view:

• Subsets $U \hookrightarrow X$ as being analogous to X-fibred sets $\phi_X : A \to X$.

• Functions $f: X \to \{t, f\}$ as being analogous to *X*-indexed sets $A: X_{disc} \to Sets$.

4.5.2 The Characteristic Function of a Point

Let *X* be a set and let $x \in X$.

DEFINITION 4.5.2.1.1 ► THE CHARACTERISTIC FUNCTION OF A POINT

The **characteristic function of** x is the function¹

$$\chi_X : X \to \{\mathsf{t},\mathsf{f}\}$$

defined by

$$\chi_x \stackrel{\text{def}}{=} \chi_{\{x\}},$$

i.e. by

$$\chi_x(y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each $y \in X$.

REMARK 4.5.2.1.2 ► CHARACTERISTIC FUNCTIONS OF POINTS AS DECATEGORIFICATIONS OF REPRESENTABLE PRESHEAVES

Expanding upon Remark 4.5.1.1.2, we may think of the characteristic function

$$\chi_x \colon X \to \{\mathsf{t},\mathsf{f}\}$$

of an *element* x of X as a decategorification of the representable presheaf and of the representable copresheaf

$$h_X \colon C^{\mathsf{op}} \to \mathsf{Sets},$$

 $h^X \colon C \to \mathsf{Sets}$

associated of an object X of a category C.

4.5.3 The Characteristic Relation of a Set

Let *X* be a set.

¹Further Notation: Also written χ^x , $\chi_X(x, -)$, or $\chi_X(-, x)$.

The **characteristic relation on** X^1 is the relation²

$$\chi_X(-1,-2): X \times X \to \{\mathsf{t},\mathsf{f}\}$$

on X defined by³

$$\chi_X(x,y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each $x, y \in X$.

REMARK 4.5.3.1.2 ► THE CHARACTERISTIC RELATION OF A SET AS A DECATEGORIFICATION OF THE HOM PROFUNCTOR

Expanding upon Remarks 4.5.1.1.2 and 4.5.2.1.2, we may view the characteristic relation

$$\chi_X(-1,-2): X \times X \to \{t,f\}$$

of *X* as a decategorification of the Hom profunctor

$$\operatorname{Hom}_{\mathcal{C}}(-1,-2): \mathcal{C}^{\operatorname{op}} \times \mathcal{C} \to \operatorname{Sets}$$

of a category C.

PROPOSITION 4.5.3.1.3 ► PROPERTIES OF CHARACTERISTIC RELATIONS

Let $f: X \to Y$ be a function.

1. The Inclusion of Characteristic Relations Associated to a Function. Let $f: A \to B$ be a function. We have an inclusion 1

$$\chi_B \circ (f \times f) \subset \chi_A, \qquad A \times A \xrightarrow{f \times f} B \times B$$

$$\chi_A \searrow \searrow \chi_B$$

$$\{t, f\}.$$

¹Further Terminology: Also called the **identity relation on** X.

²Further Notation: Also written χ_{-2}^{-1} , or \sim_{id} in the context of relations.

³Under the bijection Sets($X \times X$, {t, f}) $\cong \mathcal{P}(X \times X)$ of Item 2 of Proposition 4.5.1.1.4, the relation χ_X corresponds to the diagonal $\Delta_X \subset X \times X$ of X.

¹Note: This is the 0-categorical version of Categories, Definition 11.5.4.1.1.

PROOF 4.5.3.1.4 ► PROOF OF PROPOSITION 4.5.3.1.3

Item 1: The Inclusion of Characteristic Relations Associated to a Function

The inclusion $\chi_B(f(a), f(b)) \subset \chi_A(a, b)$ is equivalent to the statement "if a = b, then f(a) = f(b)", which is true.

4.5.4 The Characteristic Embedding of a Set

Let *X* be a set.

DEFINITION 4.5.4.1.1 ► THE CHARACTERISTIC EMBEDDING OF A SET

The **characteristic embedding**¹ **of** *X* **into** $\mathcal{P}(X)$ is the function

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

defined by²

$$\chi_{(-)}(x) \stackrel{\text{def}}{=} \chi_x$$
$$= \{x$$

for each $x \in X$.

REMARK 4.5.4.1.2 ► THE CHARACTERISTIC EMBEDDING OF A SET AS A DECATEGORIFICATION OF THE YONEDA EMBEDDING

Expanding upon Remarks 4.5.1.1.2, 4.5.2.1.2 and 4.5.3.1.2, we may view the characteristic embedding

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

of *X* into $\mathcal{P}(X)$ as a decategorification of the Yoneda embedding

$$\sharp : C^{\mathsf{op}} \hookrightarrow \mathsf{PSh}(C)$$

of a category C into PSh(C).

¹The name "characteristic *embedding*" is justified by Corollary 4.5.5.1.3, which gives an analogue of fully faithfulness for $\chi_{(-)}$.

²Here we are identifying $\mathcal{P}(X)$ with Sets(X, {t, f}) as per Item 2 of Proposition 4.5.1.1.4.

PROPOSITION 4.5.4.1.3 ► PROPERTIES OF CHARACTERISTIC EMBEDDINGS

Let $f: X \to Y$ be a map of sets.

1. Interaction With Functions. We have

$$f_! \circ \chi_X = \chi_Y \circ f, \qquad \chi_X \qquad \qquad \downarrow \chi_Y$$

$$\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y).$$

PROOF 4.5.4.1.4 ► PROOF OF PROPOSITION 4.5.4.1.3

Item 1: Interaction With Functions

Indeed, we have

$$[f! \circ \chi_X](x) \stackrel{\text{def}}{=} f!(\chi_X(x))$$

$$\stackrel{\text{def}}{=} f!(\{x\})$$

$$= \{f(x)\}$$

$$\stackrel{\text{def}}{=} \chi_{X'}(f(x))$$

$$\stackrel{\text{def}}{=} [\chi_{X'} \circ f](x),$$

for each $x \in X$, showing the desired equality.

4.5.5 The Yoneda Lemma for Sets

Let *X* be a set and let $U \subset X$ be a subset of *X*.

PROPOSITION 4.5.5.1.1 ► THE YONEDA LEMMA FOR SETS

We have

$$\chi_{\mathcal{P}(X)}(\chi_x,\chi_U)=\chi_U(x)$$

for each $x \in X$, giving an equality of functions

$$\chi_{\mathcal{P}(X)}(\chi_{(-)},\chi_U)=\chi_U,$$

where

$$\chi_{\mathcal{P}(X)}(U,V) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } U \subset V, \\ \text{false} & \text{otherwise.} \end{cases}$$

PROOF 4.5.5.1.2 ▶ PROOF OF PROPOSITION 4.5.5.1.1

We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_U) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } \{x\} \subset U, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } x \in U \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_U(x).$$

This finishes the proof.

COROLLARY 4.5.5.1.3 ► THE CHARACTERISTIC EMBEDDING IS FULLY FAITHFUL

The characteristic embedding is fully faithful, i.e., we have

$$\chi_{\mathcal{P}(X)}(\chi_x,\chi_y)\cong\chi_X(x,y)$$

for each $x, y \in X$.

PROOF 4.5.5.1.4 ► PROOF OF COROLLARY 4.5.5.1.3

We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_y) = \chi_y(x)$$

$$\stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x \in \{y\} \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } x = y \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_X(x, y).$$

where we have used Proposition 4.5.5.1.1 for the first equality.

4.6 The Adjoint Triple $f_! \dashv f^{-1} \dashv f_*$

4.6.1 Direct Images

Let $f: X \to Y$ be a function.

DEFINITION 4.6.1.1.1 ► **DIRECT IMAGES**

The **direct image function associated to** f is the function¹

$$f_! \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by²

$$f_!(U) \stackrel{\text{def}}{=} \left\{ y \in Y \mid \text{there exists some } x \in U \right\}$$

= $\left\{ f(x) \in Y \mid x \in U \right\}$

for each $U \in \mathcal{P}(X)$.

NOTATION 4.6.1.1.2 ► FURTHER NOTATION FOR DIRECT IMAGES

Sometimes one finds the notation

$$\exists_f \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

for $f_!$. This notation comes from the fact that the following statements are equivalent, where $y \in Y$ and $U \in \mathcal{P}(X)$:

- We have $y \in \exists_f(U)$.
- There exists some $x \in U$ such that f(x) = y.

We will not make use of this notation elsewhere in Clowder.

¹Further Notation: Also written simply $f: \mathcal{P}(X) \to \mathcal{P}(Y)$.

²Further Terminology: The set f(U) is called the **direct image of** U **by** f.

WARNING 4.6.1.1.3 ► NOTATION FOR DIRECT IMAGES IS CONFUSING

Notation for direct images between powersets is tricky:

- 1. Direct images for powersets and presheaves are both adjoint to their corresponding inverse image functors. However, the direct image functor for powersets is a *left* adjoint, while the direct image functor for presheaves is a *right* adjoint:
 - (a) *Powersets*. Given a function $f: X \to Y$, we have an inverse image functor

$$f^{-1} \colon \mathcal{P}(Y) \to \mathcal{P}(X)$$
.

The *left* adjoint of this functor is the usual direct image, defined above in Definition 4.6.1.1.1.

(b) *Presheaves*. Given a morphism of topological spaces $f: X \to Y$, we have an inverse image functor

$$f^{-1} \colon \mathsf{PSh}(Y) \to \mathsf{PSh}(X).$$

The *right* adjoint of this functor is the direct image functor of presheaves, defined in ??.

- 2. The presheaf direct image functor is denoted f_* , but the direct image functor for powersets is denoted $f_!$ (as it's a left adjoint).
- 3. Adding to the confusion, it's somewhat common for $f_!: \mathcal{P}(X) \to \mathcal{P}(Y)$ to be denoted f_* .

We chose to write $f_!$ for the direct image to keep the notation aligned with the following similar adjoint situations:

SITUATION	Adjoint String
Functoriality of Powersets	$(f_! \dashv f^{-1} \dashv f_*) \colon \mathcal{P}(X) \xrightarrow{\rightleftarrows} \mathcal{P}(Y)$
Functoriality of Presheaf Categories	$(f_! \dashv f^{-1} \dashv f_*) \colon PSh(X) \xrightarrow{\rightleftarrows} PSh(Y)$
Base Change	$(f_!\dashv f^*\dashv f_*)\colon C_{/X}\overset{\rightleftarrows}{ o} C_{/Y}$
Kan Extensions	$(F_! \dashv F^* \dashv F_*) \colon \operatorname{Fun}(C, \mathcal{E}) \xrightarrow{\rightleftarrows} \operatorname{Fun}(\mathcal{D}, \mathcal{E})$



REMARK 4.6.1.1.4 ► Unwinding Definition 4.6.1.1.1

Identifying $\mathcal{P}(X)$ with $\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ via Item 2 of Proposition 4.5.1.1.4, we see that the direct image function associated to f is equivalently the function

$$f_! \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by

$$f_{!}(\chi_{U}) \stackrel{\text{def}}{=} \operatorname{Lan}_{f}(\chi_{U})$$

$$= \operatorname{colim}((f \times (-1)) \stackrel{\operatorname{pr}}{\twoheadrightarrow} A \xrightarrow{\chi_{U}} \{t, f\})$$

$$= \operatorname{colim}_{x \in X} (\chi_{U}(x))$$

$$= \bigvee_{\substack{x \in X \\ f(x) = -1}} (\chi_{U}(x)),$$

where we have used ?? for the second equality. In other words, we have

$$[f_!(\chi_U)](y) = \bigvee_{\substack{x \in X \\ f(x) = y}} (\chi_U(x))$$

$$= \begin{cases} \text{true} & \text{if there exists some } x \in X \text{ such that } f(x) = y \text{ and } x \in U, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if there exists some } x \in U \\ & \text{such that } f(x) = y, \\ \text{false} & \text{otherwise} \end{cases}$$

for each $y \in Y$.

PROPOSITION 4.6.1.1.5 ► PROPERTIES OF DIRECT IMAGES I

Let $f: X \to Y$ be a function.

1. Functoriality. The assignment $U \mapsto f_!(U)$ defines a functor

$$f_! : (\mathcal{P}(X), \subset) \to (\mathcal{P}(Y), \subset).$$

In particular, for each $U, V \in \mathcal{P}(X)$, the following condition is satisfied:

(\star) If $U \subset V$, then $f_!(U) \subset f_!(V)$.

2. Triple Adjointness. We have a triple adjunction

$$(f_! + f^{-1} + f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$id_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad id_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$

 $f_! \circ f^{-1} \hookrightarrow id_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow id_{\mathcal{P}(X)},$

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*(f^{-1}(V)),$$

 $f_!(f^{-1}(V)) \subset V, \qquad f^{-1}(f_*(U)) \subset U$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

(b) Bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(Y)}(f_!(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f^{-1}(V)),$$

 $\operatorname{Hom}_{\mathcal{P}(X)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f_*(V)),$

natural in $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$ and (respectively) $V \in \mathcal{P}(X)$ and $U \in \mathcal{P}(Y)$. In particular:

- i. The following conditions are equivalent:
 - A. We have $f_!(U) \subset V$.
 - B. We have $U \subset f^{-1}(V)$.
- ii. The following conditions are equivalent:
 - A. We have $f^{-1}(U) \subset V$.
 - B. We have $U \subset f_*(V)$.
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_!)_!} \mathcal{P}(\mathcal{P}(Y))$$

$$\cup \qquad \qquad \qquad \bigcup \qquad \qquad \bigcup \cup$$

$$\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U\in\mathcal{U}}f_!(U)=\bigcup_{V\in f_!(\mathcal{U})}V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

4. Interaction With Intersections of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_!)_!} \mathcal{P}(\mathcal{P}(Y))$$

$$\cap \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcap_{U\in\mathcal{U}}f_!(U)=\bigcap_{V\in f_!(\mathcal{U})}V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

5. Interaction With Binary Unions. The diagram

$$\begin{array}{ccc} \mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_! \times f_!} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\ & & & \downarrow \cup & & \downarrow \cup \\ & \mathcal{P}(X) & \xrightarrow{f_!} & \mathcal{P}(Y) & \end{array}$$

commutes, i.e. we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$

for each $U, V \in \mathcal{P}(X)$.

6. *Interaction With Binary Intersections*. We have a natural transformation

with components

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

7. Interaction With Differences. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad$$

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by $U, V \in \mathcal{P}(X)$.

8. Interaction With Complements. The diagram

$$\mathcal{P}(X)^{\text{op}} \xrightarrow{f_*^{\text{op}}} \mathcal{P}(Y)^{\text{op}} \\
\xrightarrow{(-)^c} \qquad \qquad \downarrow^{(-)^c} \\
\mathcal{P}(X) \xrightarrow{f} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_!(U^\mathsf{c}) = f_*(U)^\mathsf{c}$$

for each $U \in \mathcal{P}(X)$.

9. *Interaction With Symmetric Differences*. We have a natural transformation

with components

$$f_!(U) \triangle f_!(V) \subset f_!(U \triangle V)$$

indexed by $U, V \in \mathcal{P}(X)$.

10. Interaction With Internal Homs of Powersets. The diagram

commutes, i.e. we have an equality of sets

$$f_!([U,V]_X) = [f_*(U), f_!(V)]_Y,$$

natural in $U, V \in \mathcal{P}(X)$.

11. Preservation of Colimits. We have an equality of sets

$$f!\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f_!(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{\times I}$. In particular, we have equalities

$$f_!(U) \cup f_!(V) = f_!(U \cup V),$$

 $f_!(\emptyset) = \emptyset,$

natural in $U, V \in \mathcal{P}(X)$.

12. Oplax Preservation of Limits. We have an inclusion of sets

$$f!\left(\bigcap_{i\in I}U_i\right)\subset\bigcap_{i\in I}f_!(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{\times I}$. In particular, we have inclusions

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V),$$

 $f_!(X) \subset Y,$

natural in $U, V \in \mathcal{P}(X)$.

13. Symmetric Strict Monoidality With Respect to Unions. The direct image function of Item 1 has a symmetric strict monoidal structure

$$(f_!, f_!^{\otimes}, f_{!\mid 1}^{\otimes}) \colon (\mathcal{P}(X), \cup, \emptyset) \to (\mathcal{P}(Y), \cup, \emptyset),$$

being equipped with equalities

$$f_{!|U,V}^{\otimes} \colon f_{!}(U) \cup f_{!}(V) \xrightarrow{=} f_{!}(U \cup V),$$
$$f_{!|\mathfrak{A}}^{\otimes} \colon \emptyset \xrightarrow{=} \emptyset,$$

natural in $U, V \in \mathcal{P}(X)$.

14. *Symmetric Oplax Monoidality With Respect to Intersections.* The direct image function of <u>Item 1</u> has a symmetric oplax monoidal structure

$$(f_!, f_!^{\otimes}, f_{!\mid 1}^{\otimes}) \colon (\mathcal{P}(X), \cap, X) \to (\mathcal{P}(Y), \cap, Y),$$

being equipped with inclusions

$$f_{!|U,V}^{\otimes} : f_{!}(U \cap V) \hookrightarrow f_{!}(U) \cap f_{!}(V),$$

 $f_{!|1}^{\otimes} : f_{!}(X) \hookrightarrow Y,$

natural in $U, V \in \mathcal{P}(X)$.

15. *Interaction With Coproducts.* Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \coprod g)_!(U \coprod V) = f_!(U) \coprod g_!(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

16. *Interaction With Products.* Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \boxtimes_{X \times Y} g)_!(U \boxtimes_{X \times Y} V) = f_!(U) \boxtimes_{X' \times Y'} g_!(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

17. Relation to Codirect Images. We have

$$f_!(U) = f_*(U^{\mathsf{c}})^{\mathsf{c}}$$

$$\stackrel{\text{def}}{=} Y \setminus f_*(X \setminus U)$$

for each $U \in \mathcal{P}(X)$.

PROOF 4.6.1.1.6 ► PROOF OF PROPOSITION 4.6.1.1.5

Item 1: Functoriality

Omitted.

Item 2: Triple Adjointness

This follows from Remark 4.6.1.1.4, Remark 4.6.2.1.2, Remark 4.6.3.1.4, and Kan Extensions, ?? of ??.

Item 3: Interaction With Unions of Families of Subsets

We have

$$\bigcup_{V \in f_{!}(\mathcal{U})} V = \bigcup_{V \in \{f_{!}(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcup_{U \in \mathcal{U}} f_{!}(U).$$

This finishes the proof.

Item 4: Interaction With Intersections of Families of Subsets

We have

$$\bigcap_{V \in f_{!}(\mathcal{U})} V = \bigcap_{V \in \{f_{!}(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$

$$= \bigcap_{U \in \mathcal{U}} f_{!}(U).$$

This finishes the proof.

Item 5: Interaction With Binary Unions

See [Pro25p].

Item 6: Interaction With Binary Intersections

See [Pro25n].

Item 7: Interaction With Differences

See [Pro25o].

Item 8: Interaction With Complements

Applying Item 17 to $X \setminus U$, we have

$$f_!(U^c) = f_!(X \setminus U)$$

$$= Y \setminus f_*(X \setminus (X \setminus U))$$

= $Y \setminus f_*(U)$
= $f_*(U)^{c}$.

This finishes the proof.

Item 9: Interaction With Symmetric Differences

We have

$$f_{!}(U) \triangle f_{!}(V) = (f_{!}(U) \cup f_{!}(V)) \setminus (f_{!}(U) \cap f_{!}(V))$$

$$\subset (f_{!}(U) \cup f_{!}(V)) \setminus (f_{!}(U \cap V))$$

$$= (f_{!}(U \cup V)) \setminus (f_{!}(U \cap V))$$

$$\subset f_{!}((U \cup V) \setminus (U \cap V))$$

$$= f_{!}(U \triangle V),$$

where we have used:

- 1. Item 2 of Proposition 4.3.12.1.2 for the first equality.
- 2. Item 6 of this proposition together with Item 1 of Proposition 4.3.10.1.2 for the first inclusion.
- 3. Item 5 for the second equality.
- 4. Item 7 for the second inclusion.
- 5. Item 2 of Proposition 4.3.12.1.2 for the tchird equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

Item 10: Interaction With Internal Homs of Powersets

We have

$$f_{!}([U,V]_{X}) \stackrel{\text{def}}{=} f_{!}(U^{c} \cup V)$$

$$= f_{!}(U^{c}) \cup f_{!}(V)$$

$$= f_{*}(U)^{c} \cup f_{!}(V)$$

$$\stackrel{\text{def}}{=} [f_{*}(U), f_{!}(V)]_{Y},$$

where we have used:

- 1. Item 5 for the second equality.
- 2. Item 17 for the third equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

Item 11: Preservation of Colimits

This follows from Item 2 and ??, ?? of ??.¹

Item 12: Oplax Preservation of Limits

The inclusion $f_!(X) \subset Y$ is automatic. See [Pro25n] for the other inclusions.

Item 13: Symmetric Strict Monoidality With Respect to Unions

This follows from Item 11.

Item 14: Symmetric Oplax Monoidality With Respect to Intersections

The inclusions in the statement follow from Item 12. Since $\mathcal{P}(Y)$ is posetal, the commutativity of the diagrams in the definition of a symmetric oplax monoidal functor is automatic (Categories, Item 4 of Proposition 11.2.7.1.2).

Item 15: Interaction With Coproducts

Omitted.

Item 16: Interaction With Products

Omitted.

Item 17: Relation to Codirect Images

Applying Item 16 of Proposition 4.6.3.1.7 to $X \setminus U$, we have

$$f_*(X \setminus U) = B \setminus f_!(X \setminus (X \setminus U))$$
$$= B \setminus f_!(U).$$

Taking complements, we then obtain

$$f_!(U) = B \setminus (B \setminus f_!(U)),$$

= $B \setminus f_*(X \setminus U),$

which finishes the proof.

¹Reference: [Pro25p].

PROPOSITION 4.6.1.1.7 ▶ PROPERTIES OF DIRECT IMAGES II

Let $f: X \to Y$ be a function.

1. Functionality I. The assignment $f \mapsto f_!$ defines a function

$$(-)_{*|X,Y} : \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(X),\mathcal{P}(Y)).$$

2. Functionality II. The assignment $f \mapsto f_!$ defines a function

$$(-)_{*|X,Y}: \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(X),\subset),(\mathcal{P}(Y),\subset)).$$

3. *Interaction With Identities.* For each $X \in \text{Obj}(\mathsf{Sets})$, we have

$$(id_X)_! = id_{\mathcal{P}(X)}$$
.

4. *Interaction With Composition*. For each pair of composable functions $f: X \to Y$ and $g: Y \to Z$, we have

$$(g \circ f)_! = g_! \circ f_!,$$

$$\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

$$\downarrow^{g_!}$$

$$\mathcal{P}(Z).$$

PROOF 4.6.1.1.8 ► PROOF OF PROPOSITION 4.6.1.1.7

Item 1: Functionality I

There is nothing to prove.

Item 2: Functionality II

This follows from Item 1 of Proposition 4.6.1.1.5.

Item 3: Interaction With Identities

This follows from Remark 4.6.1.1.4 and Kan Extensions, ?? of ??.

Item 4: Interaction With Composition

This follows from Remark 4.6.1.1.4 and Kan Extensions, ?? of ??.

4.6.2 Inverse Images

Let $f: X \to Y$ be a function.

The **inverse image function associated to** f is the function¹

$$f^{-1} \colon \mathcal{P}(Y) \to \mathcal{P}(X)$$

defined by²

$$f^{-1}(V) \stackrel{\text{def}}{=} \{x \in X \mid \text{we have } f(x) \in V\}$$

for each $V \in \mathcal{P}(Y)$.

REMARK 4.6.2.1.2 ► UNWINDING DEFINITION 4.6.2.1.1

Identifying $\mathcal{P}(Y)$ with $\mathsf{Sets}(Y, \{\mathsf{t}, \mathsf{f}\})$ via Item 2 of Proposition 4.5.1.1.4, we see that the inverse image function associated to f is equivalently the function

$$f^* \colon \mathcal{P}(Y) \to \mathcal{P}(X)$$

defined by

$$f^*(\chi_V) \stackrel{\mathrm{def}}{=} \chi_V \circ f$$

for each $\chi_V \in \mathcal{P}(Y)$, where $\chi_V \circ f$ is the composition

$$X \xrightarrow{f} Y \xrightarrow{\chi_V} \{\text{true, false}\}$$

in Sets.

PROPOSITION 4.6.2.1.3 ► PROPERTIES OF INVERSE IMAGES I

Let $f: X \to Y$ be a function.

1. Functoriality. The assignment $V \mapsto f^{-1}(V)$ defines a functor

$$f^{-1}: (\mathcal{P}(Y), \subset) \to (\mathcal{P}(X), \subset).$$

In particular, for each $U, V \in \mathcal{P}(Y)$, the following condition is satisfied:

(*) If
$$U \subset V$$
, then $f^{-1}(U) \subset f^{-1}(V)$.

¹Further Notation: Also written $f^* : \mathcal{P}(Y) \to \mathcal{P}(X)$.

²Further Terminology: The set $f^{-1}(V)$ is called the **inverse image of** V **by** f.

2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$id_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad id_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$

 $f_! \circ f^{-1} \hookrightarrow id_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow id_{\mathcal{P}(X)},$

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*(f^{-1}(V)),$$

 $f_!(f^{-1}(V)) \subset V, \qquad f^{-1}(f_*(U)) \subset U$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

(b) Bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(Y)}(f_!(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f^{-1}(V)),$$

 $\operatorname{Hom}_{\mathcal{P}(X)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f_*(V)),$

natural in $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$ and (respectively) $V \in \mathcal{P}(X)$ and $U \in \mathcal{P}(Y)$. In particular:

- i. The following conditions are equivalent:
 - A. We have $f_1(U) \subset V$.
 - B. We have $U \subset f^{-1}(V)$.
- ii. The following conditions are equivalent:
 - A. We have $f^{-1}(U) \subset V$.
 - B. We have $U \subset f_*(V)$.
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{V \in \mathcal{V}} f^{-1}(V) = \bigcup_{U \in f^{-1}(\mathcal{U})} U$$

for each $V \in \mathcal{P}(Y)$, where $f^{-1}(V) \stackrel{\text{def}}{=} (f^{-1})^{-1}(V)$.

4. Interaction With Intersections of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcap_{V \in \mathcal{V}} f^{-1}(V) = \bigcap_{U \in f^{-1}(\mathcal{U})} U$$

for each $V \in \mathcal{P}(Y)$, where $f^{-1}(V) \stackrel{\text{def}}{=} (f^{-1})^{-1}(V)$.

5. Interaction With Binary Unions. The diagram

$$\begin{array}{cccc} \mathcal{P}(Y) \times \mathcal{P}(Y) & \xrightarrow{f^{-1} \times f^{-1}} & \mathcal{P}(X) \times \mathcal{P}(X) \\ & & & \downarrow \cup \\ & & & \downarrow \cup \\ & \mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X) \end{array}$$

commutes, i.e. we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

6. Interaction With Binary Intersections. The diagram

commutes, i.e. we have

$$f^{-1}(U \cap V) = f^{-1}(U) \cap f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

7. Interaction With Differences. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\mathsf{op},-1} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad$$

commutes, i.e. we have

$$f^{-1}(U \setminus V) = f^{-1}(U) \setminus f^{-1}(V)$$

for each $U, V \in \mathcal{P}(X)$.

8. Interaction With Complements. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \xrightarrow{f^{-1,\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}}$$

$$(-)^{\mathsf{c}} \qquad \qquad \downarrow (-)^{\mathsf{c}}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U^{\mathsf{c}}) = f^{-1}(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

9. Interaction With Symmetric Differences. The diagram

i.e. we have

$$f^{-1}(U) \triangle f^{-1}(V) = f^{-1}(U \triangle V)$$

for each $U, V \in \mathcal{P}(Y)$.

10. Interaction With Internal Homs of Powersets. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1}, \mathsf{op} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$[-1, -2]_Y \downarrow \qquad \qquad \downarrow [-1, -2]_X$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_X,$$

natural in $U, V \in \mathcal{P}(X)$.

11. Preservation of Colimits. We have an equality of sets

$$f^{-1}\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f^{-1}(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(Y)^{\times I}$. In particular, we have equalities

$$f^{-1}(U) \cup f^{-1}(V) = f^{-1}(U \cup V),$$

 $f^{-1}(\emptyset) = \emptyset,$

natural in $U, V \in \mathcal{P}(Y)$.

12. Preservation of Limits. We have an equality of sets

$$f^{-1}\left(\bigcap_{i\in I}U_i\right)=\bigcap_{i\in I}f^{-1}(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(Y)^{\times I}$. In particular, we have equalities

$$f^{-1}(U) \cap f^{-1}(V) = f^{-1}(U \cap V),$$

 $f^{-1}(Y) = X,$

natural in $U, V \in \mathcal{P}(Y)$.

13. Symmetric Strict Monoidality With Respect to Unions. The inverse image function of Item 1 has a symmetric strict monoidal structure

$$(f^{-1}, f^{-1, \otimes}, f_{\mathbb{I}}^{-1, \otimes}) \colon (\mathcal{P}(Y), \cup, \emptyset) \to (\mathcal{P}(X), \cup, \emptyset),$$

being equipped with equalities

$$f_{U,V}^{-1,\otimes} \colon f^{-1}(U) \cup f^{-1}(V) \xrightarrow{=} f^{-1}(U \cup V),$$
$$f_{1}^{-1,\otimes} \colon \emptyset \xrightarrow{=} f^{-1}(\emptyset),$$

natural in $U, V \in \mathcal{P}(Y)$.

14. *Symmetric Strict Monoidality With Respect to Intersections*. The inverse image function of <u>Item 1</u> has a symmetric strict monoidal structure

$$(f^{-1}, f^{-1, \otimes}, f_{\parallel}^{-1, \otimes}) \colon (\mathcal{P}(Y), \cap, Y) \to (\mathcal{P}(X), \cap, X),$$

being equipped with equalities

$$f_{U,V}^{-1,\otimes} \colon f^{-1}(U) \cap f^{-1}(V) \xrightarrow{=} f^{-1}(U \cap V),$$

$$f_{1}^{-1,\otimes} \colon X \xrightarrow{=} f^{-1}(Y),$$

natural in $U, V \in \mathcal{P}(Y)$.

15. *Interaction With Coproducts.* Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \coprod g)^{-1}(U' \coprod V') = f^{-1}(U') \coprod g^{-1}(V')$$

for each $U' \in \mathcal{P}(X')$ and each $V' \in \mathcal{P}(Y')$.

16. *Interaction With Products.* Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f\boxtimes_{X'\times Y'}g)^{-1}(U'\boxtimes_{X'\times Y'}V')=f^{-1}(U')\boxtimes_{X\times Y}g^{-1}(V')$$

for each $U' \in \mathcal{P}(X')$ and each $V' \in \mathcal{P}(Y')$.

PROOF 4.6.2.1.4 ► PROOF OF PROPOSITION 4.6.2.1.3

Item 1: Functoriality

Omitted.

Item 2: Triple Adjointness

This follows from Remark 4.6.1.1.4, Remark 4.6.2.1.2, Remark 4.6.3.1.4, and Kan Extensions, ?? of ??.

Item 3: Interaction With Unions of Families of Subsets

We have

$$\bigcup_{U \in f^{-1}(\mathcal{V})} U = \bigcup_{U \in \left\{ f^{-1}(V) \in \mathcal{P}(X) \mid V \in \mathcal{V} \right\}} U$$
$$= \bigcup_{V \in \mathcal{V}} f^{-1}(V).$$

This finishes the proof.

Item 4: Interaction With Intersections of Families of Subsets

We have

$$\bigcap_{U \in f^{-1}(\mathcal{V})} U = \bigcap_{U \in \{f^{-1}(V) \in \mathcal{P}(X) \mid V \in \mathcal{V}\}} U$$
$$= \bigcap_{V \in \mathcal{V}} f^{-1}(V).$$

This finishes the proof.

Item 5: Interaction With Binary Unions

See [Pro25y].

Item 6: Interaction With Binary Intersections

See [Pro25w].

Item 7: Interaction With Differences

See [Pro25x].

Item 8: Interaction With Complements

See [Pro25j].

Item 9: Interaction With Symmetric Differences

We have

$$\begin{split} f^{-1}(U \triangle V) &= f^{-1}((U \cup V) \setminus (U \cap V)) \\ &= f^{-1}(U \cup V) \setminus f^{-1}(U \cap V) \\ &= f^{-1}(U) \cup f^{-1}(V) \setminus f^{-1}(U \cap V) \\ &= f^{-1}(U) \cup f^{-1}(V) \setminus f^{-1}(U) \cap f^{-1}(V) \\ &= f^{-1}(U) \triangle f^{-1}(V), \end{split}$$

where we have used:

- 1. Item 2 of Proposition 4.3.12.1.2 for the first equality.
- 2. Item 7 for the second equality.
- 3. Item 5 for the third equality.
- 4. Item 6 for the fourth equality.
- 5. Item 2 of Proposition 4.3.12.1.2 for the fifth equality.

This finishes the proof.

Item 10: Interaction With Internal Homs of Powersets

We have

$$f^{-1}([U,V]_Y) \stackrel{\text{def}}{=} f^{-1}(U^c \cup V)$$

$$= f^{-1}(U^c) \cup f^{-1}(V)$$

$$= f^{-1}(U)^c \cup f^{-1}(V)$$

$$\stackrel{\text{def}}{=} [f^{-1}(U), f^{-1}(V)]_X,$$

where we have used:

- 1. Item 8 for the second equality.
- 2. Item 5 for the third equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

Item 11: Preservation of Colimits

This follows from Item 2 and ??, ?? of ??.¹

Item 12: Preservation of Limits

This follows from Item 2 and ??, ?? of ??.²

Item 13: Symmetric Strict Monoidality With Respect to Unions

This follows from Item 11.

Item 14: Symmetric Strict Monoidality With Respect to Intersections

This follows from Item 12.

Item 15: Interaction With Coproducts

Omitted.

Item 16: Interaction With Products

Omitted.

¹Reference: [Pro25y]. ²Reference: [Pro25w].

PROPOSITION 4.6.2.1.5 ► PROPERTIES OF INVERSE IMAGES II

Let $f: X \to Y$ be a function.

1. Functionality I. The assignment $f \mapsto f^{-1}$ defines a function

$$(-)_{XY}^{-1} : \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(Y),\mathcal{P}(X)).$$

2. Functionality II. The assignment $f \mapsto f^{-1}$ defines a function

$$(-)_{XY}^{-1} \colon \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(Y),\subset),(\mathcal{P}(X),\subset)).$$

3. *Interaction With Identities.* For each $X \in Obj(Sets)$, we have

$$\mathrm{id}_X^{-1}=\mathrm{id}_{\mathcal{P}(X)}\;.$$

4. *Interaction With Composition*. For each pair of composable functions $f: X \to Y$ and $g: Y \to Z$, we have

$$(g \circ f)^{-1} = f^{-1} \circ g^{-1},$$

$$\mathcal{P}(Z) \xrightarrow{g^{-1}} \mathcal{P}(Y)$$

$$\downarrow^{f^{-1}}$$

$$\mathcal{P}(X).$$

Item 1: Functionality I

There is nothing to prove.

Item 2: Functionality II

This follows from Item 1 of Proposition 4.6.2.1.3.

Item 3: Interaction With Identities

This follows from Remark 4.6.2.1.2 and Categories, Item 5 of Proposition 11.1.4.1.2.

Item 4: Interaction With Composition

This follows from Remark 4.6.2.1.2 and Categories, Item 2 of Proposition 11.1.4.1.2.

4.6.3 Codirect Images

Let $f: X \to Y$ be a function.

DEFINITION 4.6.3.1.1 ► CODIRECT IMAGES

The **codirect image function associated to** f is the function

$$f_* \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by^{1,2}

$$f_*(U) \stackrel{\text{def}}{=} \left\{ y \in Y \middle| \begin{array}{l} \text{for each } x \in X, \text{ if we have} \\ f(x) = y, \text{ then } x \in U \end{array} \right\}$$
$$= \left\{ y \in Y \middle| \text{ we have } f^{-1}(y) \subset U \right\}$$

for each $U \in \mathcal{P}(X)$.

$$f_*(U) = f_!(U^{\mathsf{c}})^{\mathsf{c}}$$

$$\stackrel{\text{def}}{=} Y \setminus f_!(X \setminus U);$$

see Item 16 of Proposition 4.6.3.1.7.

¹Further Terminology: The set $f_*(U)$ is called the **codirect image of** U **by** f.

²We also have

NOTATION 4.6.3.1.2 ► FURTHER NOTATION FOR CODIRECT IMAGES

Sometimes one finds the notation

$$\forall_f \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

for $f_!$. This notation comes from the fact that the following statements are equivalent, where $y \in Y$ and $U \in \mathcal{P}(X)$:

- We have $y \in \forall_f(U)$.
- For each $x \in X$, if y = f(x), then $x \in U$.

We will not make use of this notation elsewhere in Clowder.

WARNING 4.6.3.1.3 ► NOTATION FOR CODIRECT IMAGES IS CONFUSING



See Warning 4.6.1.1.3.

REMARK 4.6.3.1.4 ► Unwinding Definition 4.6.3.1.1

Identifying $\mathcal{P}(X)$ with $\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ via Item 2 of Proposition 4.5.1.1.4, we see that the codirect image function associated to f is equivalently the function

$$f_* \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by

$$f_*(\chi_U) \stackrel{\text{def}}{=} \operatorname{Ran}_f(\chi_U)$$

$$= \lim ((\underbrace{(-_1)}_{x \in X} \xrightarrow{f}) \xrightarrow{\operatorname{pr}} X \xrightarrow{\chi_U} \{ \text{true, false} \})$$

$$= \lim_{\substack{x \in X \\ f(x) = -_1}} (\chi_U(x))$$

$$= \bigwedge_{\substack{x \in X \\ f(x) = -_1}} (\chi_U(x)).$$

where we have used ?? for the second equality. In other words, we have

$$[f_*(\chi_U)](y) = \bigwedge_{\substack{x \in X \\ f(x) = y}} (\chi_U(x))$$

$$= \begin{cases} \mathsf{true} & \text{if, for each } x \in X \text{ such that} \\ & f(x) = y, \text{ we have } x \in U, \\ \mathsf{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \mathsf{true} & \text{if } f^{-1}(y) \subset U \\ \mathsf{false} & \text{otherwise} \end{cases}$$

for each $y \in Y$.

DEFINITION 4.6.3.1.5 \blacktriangleright The Image and Complement Parts of f_*

Let *U* be a subset of X.^{1,2}

1. The **image part of the codirect image** $f_*(U)$ **of** U is the set $f_{*,im}(U)$ defined by

$$f_{*,\text{im}}(U) \stackrel{\text{def}}{=} f_*(U) \cap \text{Im}(f)$$

$$= \left\{ y \in Y \middle| \begin{array}{l} \text{we have } f^{-1}(y) \subset U \\ \text{and } f^{-1}(y) \neq \emptyset. \end{array} \right\}.$$

2. The complement part of the codirect image $f_*(U)$ of U is the set $f_{*,cp}(U)$ defined by

$$f_{*,cp}(U) \stackrel{\text{def}}{=} f_*(U) \cap (Y \setminus \text{Im}(f))$$

$$= Y \setminus \text{Im}(f)$$

$$= \left\{ y \in Y \middle| \text{ we have } f^{-1}(y) \subset U \right\}$$

$$= \left\{ y \in Y \middle| f^{-1}(y) = \emptyset \right\}.$$

$$f_*(U) = f_{*,im}(U) \cup f_{*,cp}(U),$$

as

$$\begin{split} f_*(U) &= f_*(U) \cap Y \\ &= f_*(U) \cap (\operatorname{Im}(f) \cup (Y \setminus \operatorname{Im}(f))) \\ &= (f_*(U) \cap \operatorname{Im}(f)) \cup (f_*(U) \cap (Y \setminus \operatorname{Im}(f))) \\ &\stackrel{\text{def}}{=} f_{*,\operatorname{im}}(U) \cup f_{*,\operatorname{cp}}(U). \end{split}$$

¹Note that we have

²In terms of the meet computation of $f_*(U)$ of Remark 4.6.3.1.4, namely

$$f_*(\chi_U) = \bigwedge_{\substack{x \in X \\ f(x) = -1}} (\chi_U(x)),$$

we see that $f_{*,im}$ corresponds to meets indexed over nonempty sets, while $f_{*,cp}$ corresponds to meets indexed over the empty set.

EXAMPLE 4.6.3.1.6 ► **EXAMPLES OF CODIRECT IMAGES**

Here are some examples of codirect images.

1. *Multiplication by Two*. Consider the function $f: \mathbb{N} \to \mathbb{N}$ given by

$$f(n) \stackrel{\text{def}}{=} 2n$$

for each $n \in \mathbb{N}$. Since f is injective, we have

$$f_{*,\text{im}}(U) = f_!(U)$$

 $f_{*,\text{cp}}(U) = \{\text{odd natural numbers}\}$

for any $U \subset \mathbb{N}$. In particular, we have

 $f_*(\{\text{even natural numbers}\}) = \mathbb{N}.$

2. *Parabolas*. Consider the function $f: \mathbb{R} \to \mathbb{R}$ given by

$$f(x) \stackrel{\text{def}}{=} x^2$$

for each $x \in \mathbb{R}$. We have

$$f_{*,cp}(U) = \mathbb{R}_{<0}$$

for any $U \subset \mathbb{R}$. Moreover, since $f^{-1}(x) = \{-\sqrt{x}, \sqrt{x}\}$, we have e.g.:

$$f_{*,\text{im}}([0,1]) = \{0\},$$

$$f_{*,\text{im}}([-1,1]) = [0,1],$$

$$f_{*,\text{im}}([1,2]) = \emptyset,$$

$$f_{*,\text{im}}([-2,-1] \cup [1,2]) = [1,4].$$

3. *Circles*. Consider the function $f: \mathbb{R}^2 \to \mathbb{R}$ given by

$$f(x,y) \stackrel{\text{def}}{=} x^2 + y^2$$

for each $(x, y) \in \mathbb{R}^2$. We have

$$f_{*,\mathrm{cp}}(U) = \mathbb{R}_{<0}$$

for any $U \subset \mathbb{R}^2$, and since

$$f^{-1}(r) = \begin{cases} \text{a circle of radius } r \text{ about the origin} & \text{if } r > 0, \\ \{(0,0)\} & \text{if } r = 0, \\ \emptyset & \text{if } r < 0, \end{cases}$$

we have e.g.:

$$f_{*,\text{im}}([-1,1] \times [-1,1]) = [0,1],$$

$$f_{*,\text{im}}(([-1,1] \times [-1,1]) \setminus [-1,1] \times \{0\}) = \emptyset.$$

PROPOSITION 4.6.3.1.7 ► PROPERTIES OF CODIRECT IMAGES I

Let $f: X \to Y$ be a function.

1. Functoriality. The assignment $U \mapsto f_*(U)$ defines a functor

$$f_* \colon (\mathcal{P}(X), \subset) \to (\mathcal{P}(Y), \subset).$$

In particular, for each $U, V \in \mathcal{P}(X)$, the following condition is satisfied:

- (\star) If $U \subset V$, then $f_*(U) \subset f_*(V)$.
- 2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$id_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad id_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$

 $f_! \circ f^{-1} \hookrightarrow id_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow id_{\mathcal{P}(X)},$

having components of the form

$$U \subset f^{-1}(f_{!}(U)), \qquad V \subset f_{*}(f^{-1}(V)),$$

 $f_{!}(f^{-1}(V)) \subset V, \qquad f^{-1}(f_{*}(U)) \subset U$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

(b) Bijections of sets

$$\operatorname{Hom}_{\mathcal{P}(Y)}(f_!(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f^{-1}(V)),$$

 $\operatorname{Hom}_{\mathcal{P}(X)}(f^{-1}(U), V) \cong \operatorname{Hom}_{\mathcal{P}(X)}(U, f_*(V)),$

natural in $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$ and (respectively) $V \in \mathcal{P}(X)$ and $U \in \mathcal{P}(Y)$. In particular:

- i. The following conditions are equivalent:
 - A. We have $f_1(U) \subset V$.
 - B. We have $U \subset f^{-1}(V)$.
- ii. The following conditions are equivalent:
 - A. We have $f^{-1}(U) \subset V$.
 - B. We have $U \subset f_*(V)$.
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\begin{array}{ccc} \mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_*)_*} & \mathcal{P}(\mathcal{P}(Y)) \\ & & & \downarrow \cup \\ & & & \downarrow \cup \\ & \mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y) \end{array}$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_*(U) = \bigcup_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

4. Interaction With Intersections of Families of Subsets. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_*)_*} & \mathcal{P}(\mathcal{P}(Y)) \\
& & \downarrow & & \downarrow \\
\mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U\in\mathcal{U}} f_*(U) = \bigcap_{V\in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

5. *Interaction With Binary Unions.* Let $f: X \to Y$ be a function. We have a natural transformation

$$\begin{array}{cccc}
\mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_* \times f_*} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\
\downarrow & & & \downarrow & & \downarrow \\
\mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y)
\end{array}$$

with components

$$f_*(U) \cup f_*(V) \subset f_*(U \cup V)$$

indexed by $U, V \in \mathcal{P}(X)$.

6. Interaction With Binary Intersections. The diagram

$$\begin{array}{cccc}
\mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_* \times f_*} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\
& & \downarrow & & \downarrow \\
\mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(U) \cap f_*(V) = f_*(U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

7. Interaction With Complements. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{f_!^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
\stackrel{(-)^{\mathsf{c}}}{\downarrow} \qquad \qquad \downarrow^{(-)^{\mathsf{c}}} \\
\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U^{\mathsf{c}}) = f_!(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

8. *Interaction With Symmetric Differences*. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_{*}^{\mathsf{op}} \times f_{*}} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(X) \xrightarrow{f_{*}} \mathcal{P}(Y)$$

with components

$$f_*(U \triangle V) \subset f_*(U) \triangle f_*(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

9. *Interaction With Internal Homs of Powersets*. We have a natural transformation

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) & \xrightarrow{f_{!}^{\mathsf{op}} \times f_{*}} & \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \\
[-1,-2]_{X} & & & \downarrow [-1,-2]_{Y} \\
\mathcal{P}(X) & \xrightarrow{f_{*}} & & \mathcal{P}(Y)
\end{array}$$

with components

$$[f_!(U), f_*(V)]_Y \subset f_*([U, V]_X)$$

indexed by $U, V \in \mathcal{P}(X)$.

10. Lax Preservation of Colimits. We have an inclusion of sets

$$\bigcup_{i\in I} f_*(U_i) \subset f_*\left(\bigcup_{i\in I} U_i\right),\,$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{\times I}$. In particular, we have inclusions

$$f_*(U) \cup f_*(V) \hookrightarrow f_*(U \cup V),$$

 $\emptyset \hookrightarrow f_*(\emptyset),$

natural in $U, V \in \mathcal{P}(X)$.

11. Preservation of Limits. We have an equality of sets

$$f_*\left(\bigcap_{i\in I}U_i\right)=\bigcap_{i\in I}f_*(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{\times I}$. In particular, we have equalities

$$f^{-1}(U \cap V) = f_*(U) \cap f^{-1}(V),$$

 $f_*(X) = Y,$

natural in $U, V \in \mathcal{P}(X)$.

12. *Symmetric Lax Monoidality With Respect to Unions*. The codirect image function of Item 1 has a symmetric lax monoidal structure

$$(f_*, f_*^{\otimes}, f_{*|1}^{\otimes}) \colon (\mathcal{P}(X), \cup, \emptyset) \to (\mathcal{P}(Y), \cup, \emptyset),$$

being equipped with inclusions

$$f_{*|U,V}^{\otimes} \colon f_*(U) \cup f_*(V) \hookrightarrow f_*(U \cup V),$$

 $f_{*|\mathfrak{1}}^{\otimes} \colon \emptyset \hookrightarrow f_*(\emptyset),$

natural in $U, V \in \mathcal{P}(X)$.

13. *Symmetric Strict Monoidality With Respect to Intersections*. The direct image function of Item 1 has a symmetric strict monoidal structure

$$(f_*, f_*^{\otimes}, f_{*|1}^{\otimes}) \colon (\mathcal{P}(X), \cap, X) \to (\mathcal{P}(Y), \cap, Y),$$

being equipped with equalities

$$f_{*|U,V}^{\otimes} : f_*(U \cap V) \xrightarrow{=} f_*(U) \cap f_*(V),$$
$$f_{*|1}^{\otimes} : f_*(X) \xrightarrow{=} Y,$$

natural in $U, V \in \mathcal{P}(X)$.

14. *Interaction With Coproducts.* Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \coprod g)_*(U \coprod V) = f_*(U) \coprod g_*(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

15. *Interaction With Products.* Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \boxtimes_{X \times Y} g)_*(U \boxtimes_{X \times Y} V) = f_*(U) \boxtimes_{X' \times Y'} g_*(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

16. Relation to Direct Images. We have

$$f_*(U) = f_!(U^{\mathsf{c}})^{\mathsf{c}}$$
$$= Y \setminus f_!(X \setminus U)$$

for each $U \in \mathcal{P}(X)$.

17. Interaction With Injections. If f is injective, then we have

$$f_{*,\text{im}}(U) = f_!(U),$$

 $f_{*,\text{cp}}(U) = Y \setminus \text{Im}(f),$

and so

$$f_*(U) = f_{*,\text{im}}(U) \cup f_{*,\text{cp}}(U)$$
$$= f_!(U) \cup (Y \setminus \text{Im}(f))$$

for each $U \in \mathcal{P}(X)$.

18. *Interaction With Surjections*. If *f* is surjective, then we have

$$f_{*,\text{im}}(U) \subset f_!(U),$$

 $f_{*,\text{cp}}(U) = \emptyset,$

and so

$$f_*(U) \subset f_!(U)$$

for each $U \in \mathcal{P}(X)$.

PROOF 4.6.3.1.8 ► PROOF OF PROPOSITION 4.6.3.1.7

Item 1: Functoriality

Omitted.

Item 2: Triple Adjointness

This follows from Remark 4.6.1.1.4, Remark 4.6.2.1.2, Remark 4.6.3.1.4, and Kan Extensions, ?? of ??.

Item 3: Interaction With Unions of Families of Subsets

We have

$$\bigcup_{V \in f_*(\mathcal{U})} V = \bigcup_{V \in \{f_*(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcup_{U \in \mathcal{U}} f_*(U).$$

This finishes the proof.

Item 4: Interaction With Intersections of Families of Subsets

We have

$$\bigcap_{V \in f_*(\mathcal{U})} V = \bigcap_{V \in \{f_*(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcap_{U \in \mathcal{U}} f_*(U).$$

This finishes the proof.

Item 5: Interaction With Binary Unions

We have

$$f_*(U) \cup f_*(V) = f_!(U^c)^c \cup f_!(V^c)^c$$

$$= (f_!(U^c) \cap f_!(V^c))^c$$

$$\subset (f_!(U^c \cap V^c))^c$$

$$= f_!((U \cup V)^c)^c$$

$$= f_*(U \cup V),$$

where:

1. We have used Item 16 for the first equality.

- 2. We have used Item 2 of Proposition 4.3.11.1.2 for the second equality.
- 3. We have used Item 6 of Proposition 4.6.1.1.5 for the third equality.
- 4. We have used Item 2 of Proposition 4.3.11.1.2 for the fourth equality.
- 5. We have used Item 16 for the last equality.

This finishes the proof.

Item 6: Interaction With Binary Intersections

This follows from Item 11.

Item 7: Interaction With Complements

Omitted.

Item 8: Interaction With Symmetric Differences

Omitted.

Item 9: Interaction With Internal Homs of Powersets

We have

$$\begin{split} \left[f_!(U), f^!(V) \right]_X &\stackrel{\text{def}}{=} f_!(U)^\mathsf{c} \cup f_*(V) \\ &= f_*(U^\mathsf{c}) \cup f_*(V) \\ &\subset f_*(U^\mathsf{c} \cup V) \\ &\stackrel{\text{def}}{=} f_*([U, V]_X), \end{split}$$

where we have used:

- 1. Item 7 of Proposition 4.6.3.1.7 for the second equality.
- 2. Item 5 of Proposition 4.6.3.1.7 for the inclusion.

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2). This finishes the proof.

Item 10: Lax Preservation of Colimits

Omitted.

Item 11: Preservation of Limits

This follows from Item 2 and ??, ?? of ??.

Item 12: Symmetric Lax Monoidality With Respect to Unions

This follows from Item 10.

Item 13: Symmetric Strict Monoidality With Respect to Intersections

This follows from Item 11.

Item 14: Interaction With Coproducts

Omitted.

Item 15: Interaction With Products

Omitted.

Item 16: Relation to Direct Images

We claim that $f_*(U) = Y \setminus f_!(X \setminus U)$.

• The First Implication. We claim that

$$f_*(U) \subset Y \setminus f_!(X \setminus U).$$

Let $y \in f_*(U)$. We need to show that $y \notin f_!(X \setminus U)$, i.e. that there is no $x \in X \setminus U$ such that f(x) = y.

This is indeed the case, as otherwise we would have $x \in f^{-1}(y)$ and $x \notin U$, contradicting $f^{-1}(y) \subset U$ (which holds since $y \in f_*(U)$).

Thus $y \in Y \setminus f_!(X \setminus U)$.

• The Second Implication. We claim that

$$Y \setminus f_!(X \setminus U) \subset f_*(U)$$
.

Let $y \in Y \setminus f_!(X \setminus U)$. We need to show that $y \in f_*(U)$, i.e. that $f^{-1}(y) \subset U$.

Since $y \notin f_!(X \setminus U)$, there exists no $x \in X \setminus U$ such that y = f(x), and hence $f^{-1}(y) \subset U$.

Thus $y \in f_*(U)$.

This finishes the proof of Item 16.

Item 17: Interaction With Injections

Omitted.

Item 18: Interaction With Surjections

Omitted.

PROPOSITION 4.6.3.1.9 ► PROPERTIES OF CODIRECT IMAGES II

Let $f: X \to B$ be a function.

1. Functionality I. The assignment $f \mapsto f_*$ defines a function

$$(-)_{Y|X|Y}: \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(X),\mathcal{P}(Y)).$$

2. Functionality II. The assignment $f \mapsto f_*$ defines a function

$$(-)_{!|X,Y}: \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(X),\subset),(\mathcal{P}(Y),\subset)).$$

3. *Interaction With Identities.* For each $X \in Obj(Sets)$, we have

$$(\mathrm{id}_X)_* = \mathrm{id}_{\mathcal{P}(X)}$$
.

4. *Interaction With Composition*. For each pair of composable functions $f: X \to Y$ and $g: Y \to Z$, we have

$$(g \circ f)_* = g_* \circ f_*,$$

$$(g \circ f)_* = g_* \circ f_*,$$

$$(g \circ f)_* \qquad g_*$$

$$\mathcal{P}(Z)$$

PROOF 4.6.3.1.10 ► PROOF OF PROPOSITION 4.6.3.1.9

Item 1: Functionality I

There is nothing to prove.

Item 2: Functionality II

This follows from Item 1 of Proposition 4.6.3.1.7.

Item 3: Interaction With Identities

This follows from Remark 4.6.3.1.4 and Kan Extensions, ?? of ??.

Item 4: Interaction With Composition

This follows from Remark 4.6.3.1.4 and Kan Extensions, ?? of ??.

4.6.4 A Six-Functor Formalism for Sets

REMARK 4.6.4.1.1 ► A SIX-FUNCTOR FORMALISM FOR SETS

The assignment $X \mapsto \mathcal{P}(X)$ together with the functors f_* , f^{-1} , and $f_!$ of Item 1 of Proposition 4.6.1.1.5, Item 1 of Proposition 4.6.2.1.3, and Item 1 of Proposition 4.6.3.1.7, and the functors

$$-_{1} \cap -_{2} \colon \mathcal{P}(X) \times \mathcal{P}(X) \to \mathcal{P}(X),$$
$$[-_{1}, -_{2}]_{X} \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

of Item 1 of Proposition 4.3.9.1.2 and Item 1 of Proposition 4.4.7.1.4 satisfy several properties reminiscent of a six functor formalism in the sense of ??. We collect these properties in Proposition 4.6.4.1.2 below.¹

PROPOSITION 4.6.4.1.2 ► A SIX-FUNCTOR FORMALISM FOR SETS

Let *X* be a set.

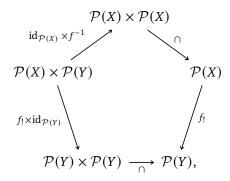
1. The Beck-Chevalley Condition. Let

$$\begin{array}{c|c}
X \times_Z Y & \xrightarrow{\operatorname{pr}_2} Y \\
 \operatorname{pr}_1 & & \downarrow g \\
 X & \xrightarrow{f} Z
\end{array}$$

be a pullback diagram in Sets. We have

¹See also [nLa25].

2. The Projection Formula I. The diagram

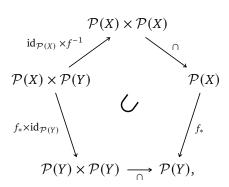


commutes, i.e. we have

$$f_!(U \cap f^{-1}(V)) = f_!(U) \cap V$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

3. The Projection Formula II. We have a natural transformation



with components

$$f_*(U) \cap V \subset f_*(U \cap f^{-1}(V))$$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

4. Strong Closed Monoidality. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1}, \mathsf{op} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$[-1, -2]_Y \downarrow \qquad \qquad \downarrow [-1, -2]_X$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_X,$$

natural in $U, V \in \mathcal{P}(X)$.

5. The External Tensor Product. We have an external tensor product

$$-1 \boxtimes_{X \times Y} -2: \mathcal{P}(X) \times \mathcal{P}(Y) \to \mathcal{P}(X \times Y)$$

given by

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \operatorname{pr}_{1}^{-1}(U) \cap \operatorname{pr}_{2}^{-1}(V)$$
$$= \{(u, v) \in X \times Y \mid u \in U \text{ and } v \in V\}.$$

This is the same map as the one in Item 5 of Proposition 4.4.1.1.4. Moreover, the following conditions are satisfied:

(a) *Interaction With Direct Images.* Let $f: X \to X'$ and $g: Y \to Y'$ be functions. The diagram

$$\begin{array}{ccc} \mathcal{P}(X) \times \mathcal{P}(Y) & \xrightarrow{f_1 \times g_1} & \mathcal{P}(X') \times \mathcal{P}(Y') \\ & \boxtimes_{X \times Y} & & & & & & \boxtimes_{X' \times Y'} \\ & \mathcal{P}(X \times Y) & \xrightarrow{f_1 \times g_1} & \mathcal{P}(X' \times Y') & & & & \end{array}$$

commutes, i.e. we have

$$[f_! \times g_!](U \boxtimes_{X \times Y} V) = f_!(U) \boxtimes_{X' \times Y'} g_!(V)$$

for each $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$.

(b) *Interaction With Inverse Images.* Let $f: X \to X'$ and $g: Y \to Y'$ be functions. The diagram

$$\begin{array}{ccc} \mathcal{P}(X') \times \mathcal{P}(Y') & \xrightarrow{f^{-1} \times g^{-1}} & \mathcal{P}(X) \times \mathcal{P}(Y) \\ & \boxtimes_{X' \times Y'} & & & & & & & & & \\ & \boxtimes_{X' \times Y'} & & & & & & & & \\ & \mathcal{P}(X' \times Y') & \xrightarrow{f^{-1} \times g^{-1}} & \mathcal{P}(X \times Y) & & & & & \\ \end{array}$$

commutes, i.e. we have

$$[f^{-1} \times g^{-1}](U \boxtimes_{X' \times Y'} V) = f^{-1}(U) \boxtimes_{X \times Y} g^{-1}(V)$$

for each $(U, V) \in \mathcal{P}(X') \times \mathcal{P}(Y')$.

(c) Interaction With Codirect Images. Let $f: X \to X'$ and $g: Y \to Y'$ be functions. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(Y) \xrightarrow{f_* \times g_*} \mathcal{P}(X') \times \mathcal{P}(Y')$$

$$\boxtimes_{X \times Y} \qquad \qquad \qquad \qquad \downarrow \boxtimes_{X' \times Y'}$$

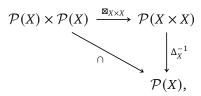
$$\mathcal{P}(X \times Y) \xrightarrow{f_* \times g_*} \mathcal{P}(X' \times Y')$$

commutes, i.e. we have

$$[f_* \times g_*](U \boxtimes_{X \times Y} V) = f_*(U) \boxtimes_{X' \times Y'} g_*(V)$$

for each $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$.

(d) Interaction With Diagonals. The diagram



i.e. we have

$$U\cap V=\Delta_X^{-1}(U\boxtimes_{X\times X}V)$$

for each $U, V \in \mathcal{P}(X)$.

6. The Dualisation Functor. We have a functor

$$D_X \colon \mathcal{P}(X)^{\mathsf{op}} \to \mathcal{P}(X)$$

given by

$$D_X(U) \stackrel{\text{def}}{=} [U, \emptyset]_X$$
$$\stackrel{\text{def}}{=} U^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$, as in Item 5 of Proposition 4.4.7.1.4, satisfying the following conditions:

(a) Duality. We have

$$D_X(D_X(U)) = U, \qquad D_X \xrightarrow{\mathrm{id}_{\mathcal{P}(X)}} \mathcal{P}(X)$$

$$\mathcal{P}(X)$$

(b) Duality. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)^{\mathsf{op}} \xrightarrow{\cap^{\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}}$$

$$\mathrm{id}_{\mathcal{P}(X)^{\mathsf{op}}} \times \mathcal{D}_{X} \xrightarrow{D_{X}} \mathcal{P}(X)$$

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_{X}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\underbrace{D_X(U\cap D_X(V))}_{\stackrel{\mathrm{def}}{=}[U\cap[V,\emptyset]_X,\emptyset]_X}=[U,V]_X$$

for each $U, V \in \mathcal{P}(X)$.

(c) Interaction With Direct Images. The diagram

$$\begin{array}{ccc} \mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_*^{\mathsf{op}}} & \mathcal{P}(Y)^{\mathsf{op}} \\ & & & \downarrow D_Y & \\ \mathcal{P}(X) & \xrightarrow{f_!} & \mathcal{P}(Y) & \end{array}$$

commutes, i.e. we have

$$f_!(D_X(U)) = D_Y(f_*(U))$$

for each $U \in \mathcal{P}(X)$.

(d) Interaction With Inverse Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(Y)^{\mathsf{op}} & \xrightarrow{f^{-1,\mathsf{op}}} & \mathcal{P}(X)^{\mathsf{op}} \\
\downarrow^{D_{Y}} & & \downarrow^{D_{X}} \\
\mathcal{P}(Y) & \xrightarrow{f^{-1}} & \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$f^{-1}(D_Y(U)) = D_X(f^{-1}(U))$$

for each $U \in \mathcal{P}(X)$.

(e) Interaction With Codirect Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\text{op}} & \xrightarrow{f_!^{\text{op}}} \mathcal{P}(Y)^{\text{op}} \\
\downarrow^{D_X} & & \downarrow^{D_Y} \\
\mathcal{P}(X) & \xrightarrow{f_*} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(D_X(U)) = D_Y(f_!(U))$$

for each $U \in \mathcal{P}(X)$.

PROOF 4.6.4.1.3 ► PROOF OF PROPOSITION 4.6.4.1.2

Item 1: The Beck–Chevalley Condition

We have

$$[g^{-1} \circ f_!](U) \stackrel{\text{def}}{=} g^{-1}(f_!(U))$$

$$\stackrel{\text{def}}{=} \{ y \in Y \mid g(y) \in f_!(U) \}$$

$$= \left\{ y \in Y \mid \text{there exists some } x \in U \right\}$$

$$= \left\{ y \in Y \mid \text{there exists some} \right\}$$

$$= \left\{ y \in Y \mid \text{there exists some} \right\}$$

$$= \left\{ y \in Y \mid \text{there exists some} \right\}$$

$$= \left\{ y \in Y \mid \text{there exists some} \right\}$$

$$= \left\{ y \in Y \mid (x, y) \in \{(x, y) \in X \times_Z Y \mid x \in U\} \right\}$$

$$= \left\{ \text{such that } y = y \right\}$$

$$= \begin{cases} y \in Y & \text{there exists some} \\ (x, y) \in \{(x, y) \in X \times_Z Y \mid x \in U\} \\ \text{such that } \mathrm{pr}_2(x, y) = y \end{cases}$$

$$\stackrel{\mathrm{def}}{=} (\mathrm{pr}_2)_! (\{(x, y) \in X \times_Z Y \mid x \in U\})$$

$$= (\mathrm{pr}_2)_! (\{(x, y) \in X \times_Z Y \mid \mathrm{pr}_1(x, y) \in U\})$$

$$\stackrel{\mathrm{def}}{=} (\mathrm{pr}_2)_! (\mathrm{pr}_1^{-1}(U))$$

$$\stackrel{\mathrm{def}}{=} [(\mathrm{pr}_2)_! \circ \mathrm{pr}_1^{-1}](U)$$

for each $U \in \mathcal{P}(X)$. Therefore, we have

$$g^{-1} \circ f_! = (pr_2)_! \circ pr_1^{-1}$$
.

For the second equality, we have

$$[f^{-1} \circ g_{!}](U) \stackrel{\text{def}}{=} f^{-1}(g_{!}(U))$$

$$\stackrel{\text{def}}{=} \{x \in X \mid f(x) \in g_{!}(V)\}$$

$$= \begin{cases} x \in X \mid \text{there exists some } y \in V \\ \text{such that } f(x) = g(y) \end{cases}$$

$$= \begin{cases} x \in X \mid \text{there exists some} \\ (x,y) \in \{(x,y) \in X \times_{Z} Y \mid y \in V\} \end{cases}$$

$$= \begin{cases} x \in X \mid \text{there exists some} \\ (x,y) \in \{(x,y) \in X \times_{Z} Y \mid y \in V\} \end{cases}$$

$$= \begin{cases} x \in X \mid \text{there exists some} \\ (x,y) \in \{(x,y) \in X \times_{Z} Y \mid y \in V\} \end{cases}$$

$$= \begin{cases} x \in X \mid \text{there exists some} \\ (x,y) \in \{(x,y) \in X \times_{Z} Y \mid y \in V\} \end{cases}$$

$$= (pr_{1})_{!}(\{(x,y) \in X \times_{Z} Y \mid pr_{2}(x,y) \in V\})$$

$$\stackrel{\text{def}}{=} (pr_{1})_{!}(pr_{2}^{-1}(V))$$

$$\stackrel{\text{def}}{=} [(pr_{1})_{!} \circ pr_{2}^{-1}](V)$$

for each $V \in \mathcal{P}(Y)$. Therefore, we have

$$f^{-1} \circ g_! = (pr_1)_! \circ pr_2^{-1}$$
.

This finishes the proof.

Item 2: The Projection Formula I

We claim that

$$f_!(U) \cap V \subset f_!(U \cap f^{-1}(V)).$$

Indeed, we have

$$f_!(U) \cap V \subset f_!(U) \cap f_!(f^{-1}(V))$$
$$= f_!(U \cap f^{-1}(V)),$$

where we have used:

- 1. Item 2 of Proposition 4.6.1.1.5 for the inclusion.
- 2. Item 6 of Proposition 4.6.1.1.5 for the equality.

Conversely, we claim that

$$f_!(U \cap f^{-1}(V)) \subset f_!(U) \cap V.$$

Indeed:

- 1. Let $y \in f_!(U \cap f^{-1}(V))$.
- 2. Since $y \in f_!(U \cap f^{-1}(V))$, there exists some $x \in U \cap f^{-1}(V)$ such that f(x) = y.
- 3. Since $x \in U \cap f^{-1}(V)$, we have $x \in U$, and thus $f(x) \in f_!(U)$.
- 4. Since $x \in U \cap f^{-1}(V)$, we have $x \in f^{-1}(V)$, and thus $f(x) \in V$.
- 5. Since $f(x) \in f_!(U)$ and $f(x) \in V$, we have $f(x) \in f_!(U) \cap V$.
- 6. But y = f(x), so $y \in f_!(U) \cap V$.
- 7. Thus $f_!(U \cap f^{-1}(V)) \subset f_!(U) \cap V$.

This finishes the proof.

Item 3: The Projection Formula II

We have

$$f_*(U) \cap V \subset f_*(U) \cap f_*(f^{-1}(V))$$

= $f_*(U \cap f^{-1}(V)),$

where we have used:

- 1. Item 2 of Proposition 4.6.3.1.7 for the inclusion.
- 2. Item 6 of Proposition 4.6.3.1.7 for the equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Proposition 11.2.7.1.2).

Item 4: Strong Closed Monoidality

This is a repetition of Item 19 of Proposition 4.4.7.1.4 and is proved there.

Item 5: The External Tensor Product

We have

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \operatorname{pr}_{1}^{-1}(U) \cap \operatorname{pr}_{2}^{-1}(V)$$

$$\stackrel{\text{def}}{=} \left\{ (x, y) \in X \times Y \mid \operatorname{pr}_{1}(x, y) \in U \right\}$$

$$\cup \left\{ (x, y) \in X \times Y \mid \operatorname{pr}_{2}(x, y) \in V \right\}$$

$$= \left\{ (x, y) \in X \times Y \mid x \in U \right\}$$

$$\cup \left\{ (x, y) \in X \times Y \mid x \in U \right\}$$

$$= \left\{ (x, y) \in X \times Y \mid x \in U \text{ and } y \in V \right\}$$

$$\stackrel{\text{def}}{=} U \times V.$$

Next, we claim that Items 5a to 5d are indeed true:

- 1. *Proof of Item 5a*: This is a repetition of Item 16 of Proposition 4.6.1.1.5 and is proved there.
- 2. *Proof of Item 5b*: This is a repetition of Item 16 of Proposition 4.6.2.1.3 and is proved there.
- 3. *Proof of Item 5c:* This is a repetition of Item 15 of Proposition 4.6.3.1.7 and is proved there.
- 4. Proof of Item 5d: We have

```
\begin{split} \Delta_X^{-1}(U \boxtimes_{X \times X} V) &\stackrel{\text{def}}{=} \{ x \in X \mid (x, x) \in U \boxtimes_{X \times X} V \} \\ &= \{ x \in X \mid (x, x) \in \{ (u, v) \in X \times X \mid u \in U \text{ and } v \in V \} \} \\ &= U \cap V. \end{split}
```

This finishes the proof.

Item 6: The Dualisation Functor

This is a repetition of Items 5 and 6 of Proposition 4.4.7.1.4 and is proved there.

4.7 Miscellany

4.7.1 Injective Functions

Let *A* and *B* be sets.

DEFINITION 4.7.1.1.1 ► INJECTIVE FUNCTIONS

A function $f: A \to B$ is **injective** if it satisfies the following condition:

 (\star) For each $a, a' \in A$, if f(a) = f(a'), then a = a'.

PROPOSITION 4.7.1.1.2 ► PROPERTIES OF INJECTIVE FUNCTIONS

Let $f: A \rightarrow B$ be a function.

- 1. Characterisations. The following conditions are equivalent:
 - (a) The function f is injective.
 - (b) The function f is a monomorphism in Sets.
 - (c) The direct image function

$$f_! : \mathcal{P}(A) \to \mathcal{P}(B)$$

associated to f is injective.

(d) The codirect image function

$$f_* \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

associated to f is injective.

(e) The direct image functor

$$f_! : (\mathcal{P}(A), \subset) \to (\mathcal{P}(B), \subset)$$

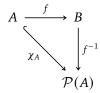
associated to f is full.

(f) The codirect image function

$$f_* \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

associated to f is full.

(g) The diagram



commutes. That is, we have

$$f^{-1}(f(a)) = \{a\}$$

for each $a \in A$.

(h) We have

$$f^{-1} \circ f_! = \mathrm{id}_{\mathcal{P}(A)} \qquad \qquad \int_{f^{-1}}^{f_!} \mathcal{P}(B)$$

$$\mathcal{P}(A).$$

In other words, we have

$${a \in A \mid f(a) \in f(U)} = U$$

for each $U \in \mathcal{P}(A)$.

(i) We have

$$f^{-1} \circ f_* = \mathrm{id}_{\mathcal{P}(A)} \qquad \qquad \downarrow^{f_*} \mathcal{P}(B)$$

$$\mathcal{P}(A).$$

In other words, we have

$$\left\{a \in A \mid f^{-1}(f(a)) \subset U\right\} = U$$

for each $U \in \mathcal{P}(A)$.

¹Items 1c to 1f unwind respectively to the following statements:

- For each $U, V \in \mathcal{P}(A)$, if $f_!(U) = f_!(V)$, then U = V.
- For each $U, V \in \mathcal{P}(A)$, if $f_*(U) = f_*(V)$, then U = V.
- For each $U, V \in \mathcal{P}(A)$, if $f_!(U) \subset f_!(V)$, then $U \subset V$.
- For each $U, V \in \mathcal{P}(A)$, if $f_*(U) \subset f_*(V)$, then $U \subset V$.

PROOF 4.7.1.1.3 ► PROOF OF PROPOSITION 4.7.1.1.2

Item 1: Characterisations

We will proceed by showing:

- Step 1: Item $1a \iff Item 1b$.
- Step 2: Item $1a \iff Item 1c$.
- Step 3: Item $1a \iff Item 1d$.
- Step 4: Item $1c \iff$ Item 1e.
- Step 5: Item $1e \iff Item 1f$.
- Step 6: Item $1a \iff Item 1g$.
- Step 7: Item 1g \iff Item 1h.
- Step 8: Item $1a \iff Item 1i$.

Step 1: Item $1a \iff Item \ 1b$

We claim that Items 1a and 1b are equivalent:

- Item $1a \Longrightarrow Item \ 1b$: We proceed in a few steps:
 - Proceeding by contrapositive, we claim that given a pair of maps $g, h: C \Rightarrow A$ such that $g \neq h$, we have $f \circ g \neq f \circ h$.
 - Indeed, as g and h are different maps, there must exist at least one element $x \in C$ such that $g(x) \neq h(x)$.
 - But then we have $f(g(x)) \neq f(h(x))$, as f is injective.
 - Thus $f \circ g \neq f \circ h$, and we are done.
- Item $1b \Longrightarrow Item 1a$: We proceed in a few steps:
 - Consider the diagram

$$pt \xrightarrow{[y]}^{[x]} A \xrightarrow{f} B,$$

where [x] and [y] are the morphisms picking the elements x and y of A.

- Note that we have f(x) = f(y) iff $f \circ [x] = f \circ [y]$.
- Since f is assumed to be a monomorphism, if f(x) = f(y), then $f \circ [x] = f \circ [y]$ and therefore [x] = [y].
- This shows that if f(x) = f(y), then x = y, so f is injective.

Step 2: Item $1a \iff Item 1c$

We claim that Items 1a and 1c are indeed equivalent:

- Item $1a \Longrightarrow Item 1c$: We proceed in a few steps:
 - Assume that f is injective and let $U, V \in \mathcal{P}(A)$ such that $f_!(U) = f_!(V)$. We wish to show that U = V.
 - **–** To show that $U \subset V$, let $u \in U$.
 - − By the definition of the direct image, we have $f(u) \in f_!(U)$.
 - Since $f_!(U) = f_!(V)$, it follows that $f(u) ∈ f_!(V)$.
 - Thus, there exists some $v \in V$ such that f(v) = f(u).
 - Since f is injective, the equality f(v) = f(u) implies that v = u.
 - **–** Thus u ∈ V and U ⊂ V.
 - A symmetric argument shows that $V \subset U$.
 - Therefore U = V, showing $f_!$ to be injective.
- Item $1c \Longrightarrow Item 1a$: We proceed in a few steps:
 - Assume that the direct image function $f_!$ is injective and let $a, a' \in A$ such that f(a) = f(a'). We wish to show that a = a'.
 - Since

$$f_!(\{a\}) = \{f(a)\}\$$

$$= \{f(a')\}\$$

$$= f_!(\{a'\}),\$$

we must have $\{a\} = \{a'\}$, as $f_!$ is injective, so a = a', showing f to be injective.

Step 3: Item $1c \iff Item 1d$

This follows from Item 17 of Proposition 4.6.1.1.5.

Step 4: Item $1c \iff Item 1e$

We claim that Items 1c and 1e are equivalent:

- *Item 1c* \Longrightarrow *Item 1e*: We proceed in a few steps:
 - Let $U, V \in \mathcal{P}(A)$ such that $f_!(U) \subset f_!(V)$, assume $f_!$ to be injective, and consider the set $U \cup V$.
 - Since $f_!(U)$ ⊂ $f_!(V)$, we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$
$$= f_!(V),$$

where we have used Item 5 of Proposition 4.6.1.1.5 for the first equality.

- Since f_1 is injective, this implies $U \cup V = V$.
- **–** Thus U ⊂ V, as we wished to show.
- Item $1c \Longrightarrow Item 1e$: We proceed in a few steps:
 - Suppose Item 1e holds, and let $U, V ∈ \mathcal{P}(A)$ such that $f_!(U) = f_!(V)$.
 - Since $f_!(U) = f_!(V)$, we have $f_!(U) \subset f_!(V)$ and $f_!(V) \subset f_!(U)$.
 - **–** By assumption, this implies U ⊂ V and V ⊂ U.
 - Thus U = V, showing $f_!$ to be injective.

Step 5: Item $1e \iff Item 1f$

This follows from Item 17 of Proposition 4.6.1.1.5.

Step 6: Item 1a ← Item 1g

We have

$$f^{-1}(f(a)) = \{ a' \in A \mid f(a') = f(a) \}$$

so the condition $f^{-1}(f(a)) = \{a\}$ states precisely that if f(a') = f(a), then a' = a.

Step 7: Item 1g ← Item 1h

We claim that Items 1g and 1h are indeed equivalent:

• *Item 1g* \Longrightarrow *Item 1h*: We have

$$[f^{-1} \circ f_!](U) \stackrel{\text{def}}{=} f^{-1}(f_!(U))$$

$$= f^{-1} \left(\int_{u \in U} \{u\} \right)$$

$$= f^{-1} \left(\bigcup_{u \in U} f_!(\{u\}) \right)$$

$$= \bigcup_{u \in U} f^{-1}(f_!(\{u\}))$$

$$= \bigcup_{u \in U} \{u\}$$

$$= \bigcup_{u \in U} \{u\}$$

for each $U \in \mathcal{P}(A)$, where we have used Item 5 of Proposition 4.6.1.1.5 for the third equality and Item 5 of Proposition 4.6.2.1.3 for the fourth equality.

• *Item 1h* \Longrightarrow *Item 1g*: Applying the condition $f^{-1} \circ f_! = \mathrm{id}_{\mathcal{P}(A)}$ to $U = \{a\}$ gives

$$f^{-1}(f_!(\{a\})) = \{a\}.$$

Step 8: Item $1a \iff Item 1i$

We claim that Items 1a and 1i are equivalent:

• Item 1a \Longrightarrow Item 1i: If f is injective, then $f^{-1}(f(a)) = \{a\}$, so we have

$$f^{-1}(f_*(a)) = \{ a \in A \mid \{a\} \subset U \}$$

= U.

• *Item 1i* \Longrightarrow *Item 1a*: For $U = \{a\}$, the condition $f^{-1}(f_*(U)) = U$ becomes

$$\{a' \in A \mid f^{-1}(f(a')) \subset \{a\}\} = \{a\}.$$

Since the set $f^{-1}(f(a'))$ is given by

$${a \in A \mid f(a) = f(a')},$$

it follows that f is injective.

This finishes the proof.

4.7.2 Surjective Functions

Let *A* and *B* be sets.

DEFINITION 4.7.2.1.1 ► **SURJECTIVE FUNCTIONS**

A function $f: A \to B$ is **surjective** if it satisfies the following condition:

 (\star) For each $b \in B$, there exists some $a \in A$ such that f(a) = b.

PROPOSITION 4.7.2.1.2 ► PROPERTIES OF SURJECTIVE FUNCTIONS

Let $f: A \rightarrow B$ be a function.

- 1. Characterisations. The following conditions are equivalent:
 - (a) The function f is surjective.
 - (b) The function *f* is an epimorphism in Sets.
 - (c) The inverse image function

$$f^{-1} \colon \mathcal{P}(B) \to \mathcal{P}(A)$$

associated to *f* is injective.

(d) The inverse image functor

$$f^{-1}: (\mathcal{P}(B), \subset) \to (\mathcal{P}(A), \subset)$$

associated to f is full.

(e) The diagram

$$B \xrightarrow{f^{-1}} \mathcal{P}(A)$$

$$\downarrow_{\chi_B} \qquad \downarrow_{f_!}$$

$$\mathcal{P}(B)$$

commutes. That is, we have

$$f_!(f^{-1}(b)) = \{b\}$$

for each $b \in B$.

(f) We have

$$f_! \circ f^{-1} = \mathrm{id}_{\mathcal{P}(B)} \qquad \qquad \downarrow^{f^{-1}} \mathcal{P}(A)$$

$$\mathcal{P}(B) \xrightarrow{f} \mathcal{P}(B)$$

In other words, we have

$$\left\{b \in B \mid \text{ there exists some } a \in f^{-1}(U) \\ \text{ such that } f(a) = b \right\} = U$$

for each $U \in \mathcal{P}(A)$.

(g) We have

$$\mathcal{P}(B) \xrightarrow{f^{-1}} \mathcal{P}(A)$$

$$f_* \circ f^{-1} = \mathrm{id}_{\mathcal{P}(B)}$$

$$\mathcal{P}(B).$$

In other words, we have

$$\{b \in B \mid f^{-1}(b) \subset f^{-1}(U)\} = U$$

for each $U \in \mathcal{P}(B)$.

PROOF 4.7.2.1.3 ► PROOF OF PROPOSITION 4.7.2.1.2

Item 1: Characterisations

We will proceed by showing:

- Step 1: Item $1a \iff Item \ 1b$.
- Step 2: Item $1a \iff Item 1c$.
- Step 3: Item $1c \iff Item 1d$.
- Step 4: Item 1a ← Item 1e.
- Step 5: Item $1e \iff Item 1f$.
- Step 6: Item $1a \iff Item 1g$.

Step 1: Item $1a \iff Item 1b$

We claim Items 1a and 1b are indeed equivalent:

- *Item 1a* \Longrightarrow *Item 1b*: We proceed in a few steps:
 - Let g, h: B \Rightarrow C be morphisms such that $g \circ f = h \circ f$.
 - **–** For each a ∈ A, we have

$$g(f(a)) = h(f(a)).$$

- However, this implies that

$$g(b) = h(b)$$

for each $b \in B$, as f is surjective.

- Thus g = h and f is an epimorphism.
- *Item 1b* ⇒ *Item 1a*: We proceed by contrapositive. Consider the diagram

$$A \xrightarrow{f} B \xrightarrow{g} C,$$

where h is the map defined by h(b) = 0 for each $b \in B$ and g is the map defined by

$$g(b) = \begin{cases} 1 & \text{if } b \in \text{Im}(f), \\ 0 & \text{otherwise.} \end{cases}$$

Then $h \circ f = g \circ f$, as h(f(a)) = 1 = g(f(a)) for each $a \in A$. However, for any $b \in B \setminus \text{Im}(f)$, we have

$$g(b) = 0 \neq 1 = h(b)$$
.

Therefore $g \neq h$ and f is not an epimorphism.

Step 2: Item $1a \iff Item 1c$

We claim Items 1a and 1c are indeed equivalent:

- Item $1a \Longrightarrow Item 1c$: We proceed in a few steps:
 - Assume that f is surjective. Let $U, V \in \mathcal{P}(B)$ such that $f^{-1}(U) = f^{-1}(V)$. We wish to show that U = V.
 - **–** To show that $U \subset V$, let $b \in U$.
 - Since f is surjective, there must exist some a ∈ A such that f(a) = b.
 - By the definition of the inverse image, since f(a) = b and $b \in U$, we have $a \in f^{-1}(U)$.
 - By our initial assumption, $f^{-1}(U) = f^{-1}(V)$, so it follows that $a \in f^{-1}(V)$.
 - Again, by the definition of the inverse image, $a \in f^{-1}(V)$ means that $f(a) \in V$.
 - Since f(a) = b, we have shown that $b \in V$.
 - − This establishes that $U \subset V$. A symmetric argument shows that $V \subset U$.
 - Thus U = V, proving that f^{-1} is injective.
 - Item $1c \Longrightarrow Item 1a$: We proceed in a few steps:

- Assume that the inverse image function f^{-1} is injective. Suppose, for the sake of contradiction, that f is not surjective.
- The assumption that f is not surjective means there exists some $b_0 ∈ B$ such that for all a ∈ A, we have $f(a) ≠ b_0$.
- By the definition of the inverse image, this is equivalent to stating that $f^{-1}(\{b_0\}) = \emptyset$.
- Since $f^{-1}(\emptyset) = \emptyset$, we have $f^{-1}(\{b_0\}) = f^{-1}(\emptyset)$.
- Since f^{-1} is injective, this implies that $\{b_0\} = \emptyset$.
- This is a contradiction, as the singleton set $\{b_0\}$ is non-empty.
- Therefore, *f* is surjective.

Step 3: Item $1c \iff Item 1d$

We claim that Items 1c and 1d are equivalent:

- Item $1c \Longrightarrow Item 1d$: We proceed in a few steps:
 - Let $U, V \in \mathcal{P}(B)$ such that $f^{-1}(U) \subset f^{-1}(V)$, assume f^{-1} to be injective, and consider the set $U \cup V$.
 - Since $f^{-1}(U) \subset f^{-1}(V)$, we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$

= $f^{-1}(V)$,

where we have used Item 5 of Proposition 4.6.2.1.3 for the first equality.

- Since f^{-1} is injective, this implies $U \cup V = V$.
- **–** Thus U ⊂ V, as we wished to show.
- *Item* $1d \Longrightarrow Item$ 1c: We proceed in a few steps:
 - Suppose Item 1d holds, and let $U, V \in \mathcal{P}(B)$ such that $f^{-1}(U) = f^{-1}(V)$.
 - Since $f^{-1}(U) = f^{-1}(V)$, we have $f^{-1}(U) \subset f^{-1}(V)$ and $f^{-1}(V) \subset f^{-1}(U)$.
 - **–** By assumption, this implies U ⊂ V and V ⊂ U.

- Thus U = V, showing f^{-1} to be injective.

Step 4: Item $1a \iff Item 1e$

We have

$$f_!(f^{-1}(b)) = \left\{ b \in B \mid \text{ there exists some } a \in f^{-1}(b) \right\},$$

so the condition $f_!(f^{-1}(b)) = \{b\}$ holds iff f is surjective.

Step 5: Item $1e \iff Item 1f$

We claim that Items 1e and 1f are indeed equivalent:

• Item $1e \Longrightarrow Item 1f$: We have

$$[f! \circ f^{-1}](U) \stackrel{\text{def}}{=} f!(f^{-1}(U))$$

$$= f! \left(\int_{u \in U} f^{-1} \left(\bigcup_{u \in U} \{u\} \right) \right)$$

$$= f! \left(\bigcup_{u \in U} f^{-1} (\{u\}) \right)$$

$$= \bigcup_{u \in U} f!(f^{-1}(u))$$

$$= \bigcup_{u \in U} f!(f^{-1}(u))$$

$$= \bigcup_{u \in U} f!$$

for each $U \in \mathcal{P}(B)$, where we have used Item 5 of Proposition 4.6.1.1.5 for the third equality and Item 5 of Proposition 4.6.2.1.3 for the fourth equality.

• *Item 1f* \Longrightarrow *Item 1e*: Applying the condition $f_! \circ f^{-1} = \mathrm{id}_{\mathcal{P}(B)}$ to $U = \{b\}$ gives

$$f_!(f^{-1}(\{b\}))=\{b\}.$$

Step 6: Item 1a ← Item 1g

First, note that for the condition $f^{-1}(b) \subset f^{-1}(U)$ to hold, we must have $b \in U$ or $f^{-1}(b) = \emptyset$. Thus

$$f_*(f^{-1}(U)) = (U \cap \operatorname{Im}(f)) \cup (B \setminus \operatorname{Im}(f)).$$

We now claim that Items 1a and 1g are indeed equivalent:

• Item $1a \Longrightarrow Item 1g$: If f is surjective, we have

$$(U \cap \operatorname{Im}(f)) \cup (B \setminus \operatorname{Im}(f)) = U \cup \emptyset$$
$$= U,$$

so
$$f_* \circ f^{-1} = \mathrm{id}_{\mathcal{P}(B)}$$
.

• Item $1g \Longrightarrow Item 1a$: Taking $U = \emptyset$ gives

$$f_*(f^{-1}(\emptyset)) = (\emptyset \cap \operatorname{Im}(f)) \cup (B \setminus \operatorname{Im}(f))$$

= $B \setminus \operatorname{Im}(f)$,

so the condition $f_*(f^{-1}(\emptyset)) = \emptyset$ implies $B \setminus \text{Im}(f) = \emptyset$. Thus Im(f) = B and f is surjective.

This finishes the proof.

Appendices

A Other Chapters

Preliminaries

- 1. Introduction
- 2. A Guide to the Literature

Sets

3. Sets

- 4. Constructions With Sets
- 5. Monoidal Structures on the Category of Sets
- 6. Pointed Sets
- 7. Tensor Products of Pointed Sets

Relations

8. Relations		Monoidal Categories
9. Construct	ions With Relations	13. Constructions With Monoidal Categories
10. Condition	s on Relations	Bicategories
Categories 11. Categorie	s	14. Types of Morphisms in Bicategories
		Extra Part
12. Presheave Lemma	es and the Yoneda	15. Notes
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