Monoidal Structures on the Category of Sets

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This chapter contains some material on monoidal structures on Sets.

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5.1 The Monoidal Category of Sets and Products						
5.1	.1 Pr	oducts of Sets				
See Constructions With Sets, Section 4.1.3.						
5.1	.2 Th	e Internal Hom of Sets				

5.1.3 The Monoidal Unit

DEFINITION 5.1.3.1.1 ► THE MONOIDAL UNIT OF ×

See Constructions With Sets, Section 4.3.5.

The monoidal unit of the product of sets is the functor

$$\mathbb{1}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

defined by

$$1_{\mathsf{Sets}} \stackrel{\text{def}}{=} \mathsf{pt},$$

where pt is the terminal set of Constructions With Sets, Definition 4.1.1.1.1.

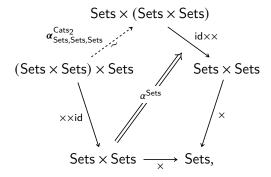
5.1.4 The Associator

DEFINITION 5.1.4.1.1 ► THE ASSOCIATOR OF ×

The associator of the product of sets is the natural isomorphism

$$\alpha^{\mathsf{Sets}} \colon \times \circ (\times \times \mathsf{id}_{\mathsf{Sets}}) \stackrel{\widetilde{-}}{\Longrightarrow} \times \circ (\mathsf{id}_{\mathsf{Sets}} \times \times) \circ \pmb{\alpha}^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}'}$$

as in the diagram



whose component

$$\alpha_{X,Y,Z}^{\mathsf{Sets}} \colon (X \times Y) \times Z \xrightarrow{\sim} X \times (Y \times Z)$$

at (X, Y, Z) is given by

$$\alpha_{X,Y,Z}^{\mathsf{Sets}}((x,y),z) \stackrel{\mathrm{def}}{=} (x,(y,z))$$

for each $((x, y), z) \in (X \times Y) \times Z$.

PROOF 5.1.4.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.1.4.1.1

Invertibility

The inverse of $\alpha_{X,Y,Z}^{\mathsf{Sets}}$ is the morphism

$$\alpha_{X,Y,Z}^{\mathsf{Sets},-1} \colon X \times (Y \times Z) \xrightarrow{\sim} (X \times Y) \times Z$$

defined by

$$\alpha_{X,Y,Z}^{\mathsf{Sets},-1}(x,(y,z)) \stackrel{\mathrm{def}}{=} ((x,y),z)$$

for each $(x, (y, z)) \in X \times (Y \times Z)$. Indeed:

• Invertibility I. We have

$$\begin{bmatrix} \alpha_{X,Y,Z}^{\mathsf{Sets},-1} \circ \alpha_{X,Y,Z}^{\mathsf{Sets}} \end{bmatrix} ((x,y),z) \stackrel{\text{def}}{=} \alpha_{X,Y,Z}^{\mathsf{Sets},-1} \Big(\alpha_{X,Y,Z}^{\mathsf{Sets}} ((x,y),z) \Big)$$

$$\stackrel{\text{def}}{=} \alpha_{X,Y,Z}^{\mathsf{Sets},-1} (x,(y,z))$$

$$\stackrel{\text{def}}{=} ((x, y), z)$$

$$\stackrel{\text{def}}{=} [\mathrm{id}_{(X \times Y) \times Z}] ((x, y), z)$$

for each $((x, y), z) \in (X \times Y) \times Z$, and therefore we have

$$\alpha_{X,Y,Z}^{\mathsf{Sets},-1} \circ \alpha_{X,Y,Z}^{\mathsf{Sets}} = \mathrm{id}_{(X \times Y) \times Z} \,.$$

• Invertibility II. We have

$$\begin{split} \left[\alpha_{X,Y,Z}^{\mathsf{Sets}} \circ \alpha_{X,Y,Z}^{\mathsf{Sets},-1}\right] &(x,(y,z)) \stackrel{\text{def}}{=} \alpha_{X,Y,Z}^{\mathsf{Sets}} \left(\alpha_{X,Y,Z}^{\mathsf{Sets},-1}(x,(y,z))\right) \\ \stackrel{\text{def}}{=} \alpha_{X,Y,Z}^{\mathsf{Sets}} ((x,y),z) \\ \stackrel{\text{def}}{=} (x,(y,z)) \\ \stackrel{\text{def}}{=} \left[\mathrm{id}_{(X \times Y) \times Z}\right] &(x,(y,z)) \end{split}$$

for each $(x, (y, z)) \in X \times (Y \times Z)$, and therefore we have $\alpha_{X,Y,Z}^{\mathsf{Sets},-1} \circ \alpha_{X,Y,Z}^{\mathsf{Sets}} = \mathrm{id}_{X \times (Y \times Z)} \ .$

Therefore $\alpha_{X,Y,Z}^{\mathsf{Sets}}$ is indeed an isomorphism.

Naturality

We need to show that, given functions

$$f: X \to X',$$

 $g: Y \to Y',$
 $h: Z \to Z'$

the diagram

$$\begin{array}{c|c} (X\times Y)\times Z & \xrightarrow{(f\times g)\times h} & (X'\times Y')\times Z' \\ \\ \alpha^{\mathsf{Sets}}_{X,Y,Z} & & & & & & \\ X\times (Y\times Z) & \xrightarrow{f\times (g\times h)} & X'\times (Y'\times Z') \end{array}$$

commutes. Indeed, this diagram acts on elements as

and hence indeed commutes, showing α^{Sets} to be a natural transformation.

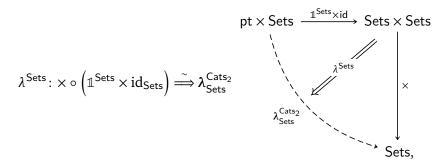
Being a Natural Isomorphism

Since α^{Sets} is natural and $\alpha^{\mathsf{Sets},-1}$ is a componentwise inverse to α^{Sets} , it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that $\alpha^{\mathsf{Sets},-1}$ is also natural. Thus α^{Sets} is a natural isomorphism.

5.1.5 The Left Unitor

DEFINITION 5.1.5.1.1 ► THE LEFT UNITOR OF ×

The left unitor of the product of sets is the natural isomorphism



whose component

$$\lambda_X^{\mathsf{Sets}} \colon \mathsf{pt} \times X \xrightarrow{\sim} X$$

at $X \in Obj(Sets)$ is given by

$$\lambda_X^{\mathsf{Sets}}(\star, x) \stackrel{\mathrm{def}}{=} x$$

for each $(\star, x) \in pt \times X$.

PROOF 5.1.5.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.1.5.1.1

Invertibility

The inverse of $\lambda_X^{\mathsf{Sets}}$ is the morphism

$$\lambda_X^{\mathsf{Sets},-1} \colon X \xrightarrow{\sim} \mathsf{pt} \times X$$

defined by

$$\lambda_X^{\mathsf{Sets},-1}(x) \stackrel{\mathrm{def}}{=} (\star, x)$$

for each $x \in X$. Indeed:

• Invertibility I. We have

$$\begin{split} \left[\lambda_X^{\mathsf{Sets},-1} \circ \lambda_X^{\mathsf{Sets}}\right] (\mathsf{pt},x) &= \lambda_X^{\mathsf{Sets},-1} \Big(\lambda_X^{\mathsf{Sets}} (\mathsf{pt},x)\Big) \\ &= \lambda_X^{\mathsf{Sets},-1} (x) \\ &= (\mathsf{pt},x) \\ &= \left[\mathrm{id}_{\mathsf{pt} \times X}\right] (\mathsf{pt},x) \end{split}$$

for each $(pt, x) \in pt \times X$, and therefore we have

$$\lambda_X^{\mathsf{Sets},-1} \circ \lambda_X^{\mathsf{Sets}} = \mathrm{id}_{\mathsf{pt} \times X} .$$

• Invertibility II. We have

$$\begin{split} \left[\lambda_X^{\mathsf{Sets}} \circ \lambda_X^{\mathsf{Sets},-1}\right](x) &= \lambda_X^{\mathsf{Sets}} \Big(\lambda_X^{\mathsf{Sets},-1}(x)\Big) \\ &= \lambda_X^{\mathsf{Sets},-1}(\mathsf{pt},x) \\ &= x \\ &= [\mathsf{id}_X](x) \end{split}$$

for each $x \in X$, and therefore we have

$$\lambda_X^{\mathsf{Sets}} \circ \lambda_X^{\mathsf{Sets},-1} = \mathrm{id}_X$$
.

Therefore $\lambda_X^{\mathsf{Sets}}$ is indeed an isomorphism.

Naturality

We need to show that, given a function $f: X \to Y$, the diagram

$$\begin{array}{ccc}
\operatorname{pt} \times X & \xrightarrow{\operatorname{id}_{\operatorname{pt}} \times f} & \operatorname{pt} \times Y \\
\lambda_X^{\operatorname{Sets}} & & & \downarrow \lambda_Y^{\operatorname{Sets}} \\
X & \xrightarrow{f} & Y
\end{array}$$

commutes. Indeed, this diagram acts on elements as

$$(\star, x) \qquad (\star, x) \longmapsto (\star, f(x))$$

$$\downarrow \qquad \qquad \downarrow$$

$$x \longmapsto f(x) \qquad f(x)$$

and hence indeed commutes. Therefore λ^{Sets} is a natural transformation.

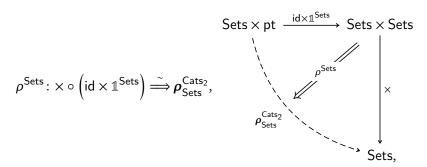
Being a Natural Isomorphism

Since λ^{Sets} is natural and $\lambda^{\text{Sets},-1}$ is a componentwise inverse to λ^{Sets} , it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that $\lambda^{\text{Sets},-1}$ is also natural. Thus λ^{Sets} is a natural isomorphism.

5.1.6 The Right Unitor

DEFINITION 5.1.6.1.1 ► THE RIGHT UNITOR OF ×

The **right unitor of the product of sets** is the natural isomorphism



whose component

$$\rho_X^{\mathsf{Sets}} \colon X \times \mathsf{pt} \xrightarrow{\sim} X$$

at $X \in \text{Obj}(\mathsf{Sets})$ is given by

$$\rho_X^{\mathsf{Sets}}(x,\star) \stackrel{\mathrm{def}}{=} x$$

for each $(x, \star) \in X \times pt$.

PROOF 5.1.6.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.1.6.1.1

Invertibility

The inverse of ρ_X^{Sets} is the morphism

$$\rho_X^{\mathsf{Sets},-1} \colon X \xrightarrow{\sim} X \times \mathsf{pt}$$

defined by

$$\rho_X^{\mathsf{Sets},-1}(x) \stackrel{\mathrm{def}}{=} (x,\star)$$

for each $x \in X$. Indeed:

• Invertibility I. We have

$$\begin{split} \left[\rho_X^{\mathsf{Sets},-1} \circ \rho_X^{\mathsf{Sets}} \right] (x, \star) &= \rho_X^{\mathsf{Sets},-1} \Big(\rho_X^{\mathsf{Sets}} (x, \star) \Big) \\ &= \rho_X^{\mathsf{Sets},-1} (x) \\ &= (x, \star) \\ &= \left[\mathrm{id}_{X \times \mathrm{pt}} \right] (x, \star) \end{split}$$

for each $(x, \star) \in X \times pt$, and therefore we have

$$\rho_X^{\mathsf{Sets},-1} \circ \rho_X^{\mathsf{Sets}} = \mathrm{id}_{X \times \mathrm{pt}}.$$

• Invertibility II. We have

$$\begin{split} \left[\rho_X^{\mathsf{Sets}} \circ \rho_X^{\mathsf{Sets},-1} \right] &(x) = \rho_X^{\mathsf{Sets}} \left(\rho_X^{\mathsf{Sets},-1}(x) \right) \\ &= \rho_X^{\mathsf{Sets},-1}(x, \star) \\ &= x \\ &= [\mathrm{id}_X](x) \end{split}$$

for each $x \in X$, and therefore we have

$$\rho_X^{\mathsf{Sets}} \circ \rho_X^{\mathsf{Sets},-1} = \mathrm{id}_X.$$

Therefore ρ_X^{Sets} is indeed an isomorphism.

Naturality

We need to show that, given a function $f: X \to Y$, the diagram

$$\begin{array}{ccc} X \times \operatorname{pt} & \xrightarrow{f \times \operatorname{id}_{\operatorname{pt}}} & Y \times \operatorname{pt} \\ \rho_X^{\operatorname{Sets}} & & & & \downarrow \rho_Y^{\operatorname{Sets}} \\ X & \xrightarrow{f} & & Y \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$(x, \star) \qquad (x, \star) \longmapsto (f(x), \star)$$

$$\downarrow \qquad \qquad \downarrow$$

$$x \longmapsto f(x) \qquad f(x)$$

and hence indeed commutes. Therefore ρ^{Sets} is a natural transformation.

Being a Natural Isomorphism

Since ρ^{Sets} is natural and $\rho^{\text{Sets},-1}$ is a componentwise inverse to ρ^{Sets} , it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that $\rho^{\text{Sets},-1}$ is also natural. Thus ρ^{Sets} is a natural isomorphism.

5.1.7 The Symmetry

DEFINITION 5.1.7.1.1 ► THE SYMMETRY OF ×

The symmetry of the product of sets is the natural isomorphism

$$\sigma^{\mathsf{Sets}} : \times \stackrel{\sim}{\Longrightarrow} \times \circ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}}, \qquad \begin{array}{c} \mathsf{Sets} \times \mathsf{Sets} & \stackrel{\times}{\longrightarrow} \mathsf{Sets}, \\ & \downarrow \downarrow & \\ & \mathsf{Sets} \times \mathsf{Sets} & \\ & & \mathsf{Sets} \times \mathsf{Sets} & \\ & & \mathsf{Sets} \times \mathsf{Sets} & \\ \end{array}$$

whose component

$$\sigma_{X,Y}^{\mathsf{Sets}} : X \times Y \xrightarrow{\sim} Y \times X$$

at $X, Y \in Obj(Sets)$ is defined by

$$\sigma_{X,Y}^{\mathsf{Sets}}(x,y) \stackrel{\mathrm{def}}{=} (y,x)$$

for each $(x, y) \in X \times Y$.

PROOF 5.1.7.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.1.7.1.1

Invertibility

The inverse of $\sigma_{X,Y}^{\mathsf{Sets}}$ is the morphism

$$\sigma_{X,Y}^{\mathsf{Sets},-1} \colon Y \times X \xrightarrow{\sim} X \times Y$$

defined by

$$\sigma_{X,Y}^{\mathsf{Sets},-1}(y,x) \stackrel{\mathsf{def}}{=} (x,y)$$

for each $(y, x) \in Y \times X$. Indeed:

• Invertibility I. We have

$$\begin{split} \left[\sigma_{X,Y}^{\mathsf{Sets},-1} \circ \sigma_{X,Y}^{\mathsf{Sets}}\right] &(x,y) \stackrel{\text{def}}{=} \sigma_{X,Y}^{\mathsf{Sets},-1} \Big(\sigma_{X,Y}^{\mathsf{Sets}}(x,y)\Big) \\ &\stackrel{\text{def}}{=} \sigma_{X,Y}^{\mathsf{Sets},-1}(y,x) \\ &\stackrel{\text{def}}{=} (x,y) \\ &\stackrel{\text{def}}{=} [\mathrm{id}_{X\times Y}](x,y) \end{split}$$

for each $(x, y) \in X \times Y$, and therefore we have

$$\sigma_{X,Y}^{\mathsf{Sets},-1} \circ \sigma_{X,Y}^{\mathsf{Sets}} = \mathrm{id}_{X \times Y} \,.$$

• Invertibility II. We have

$$\begin{split} \left[\sigma_{X,Y}^{\mathsf{Sets}} \circ \sigma_{X,Y}^{\mathsf{Sets},-1}\right] (y,x) &\stackrel{\mathsf{def}}{=} \sigma_{X,Y}^{\mathsf{Sets},-1} \Big(\sigma_{X,Y}^{\mathsf{Sets}} (y,x) \Big) \\ &\stackrel{\mathsf{def}}{=} \sigma_{X,Y}^{\mathsf{Sets},-1} (x,y) \\ &\stackrel{\mathsf{def}}{=} (y,x) \\ &\stackrel{\mathsf{def}}{=} [\mathrm{id}_{Y \times X}] (y,x) \end{split}$$

for each $(y, x) \in Y \times X$, and therefore we have

$$\sigma_{X,Y}^{\mathsf{Sets}} \circ \sigma_{X,Y}^{\mathsf{Sets},-1} = \mathrm{id}_{Y \times X}$$
.

Therefore $\sigma_{X,Y}^{\mathsf{Sets}}$ is indeed an isomorphism.

Naturality

We need to show that, given functions

$$f: X \to A$$
, $g: Y \to B$

the diagram

$$\begin{array}{c|c} X \times Y & \xrightarrow{f \times g} & A \times B \\ \sigma_{X,Y}^{\mathsf{Sets}} & & & & & & \\ & & & & & & \\ & Y \times X & \xrightarrow{g \times f} & B \times A \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$(x,y) \qquad (x,y) \longmapsto (f(x),g(y))$$

$$\downarrow \qquad \qquad \downarrow$$

$$(y,x) \longmapsto (g(y),f(x)) \qquad (g(y),f(x))$$

and hence indeed commutes, showing σ^{Sets} to be a natural transformation.

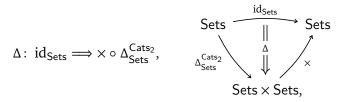
Being a Natural Isomorphism

Since σ^{Sets} is natural and $\sigma^{Sets,-1}$ is a componentwise inverse to σ^{Sets} , it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that $\sigma^{Sets,-1}$ is also natural. Thus σ^{Sets} is a natural isomorphism.

5.1.8 The Diagonal

DEFINITION 5.1.8.1.1 ► THE DIAGONAL OF ×

The **diagonal of the product of sets** is the natural transformation



whose component

$$\Delta_X : X \to X \times X$$

at $X \in \text{Obj}(\mathsf{Sets})$ is given by

$$\Delta_X(x) \stackrel{\text{def}}{=} (x, x)$$

for each $x \in X$.

PROOF 5.1.8.1.2 ➤ PROOF OF THE CLAIMS MADE IN DEFINITION 5.1.8.1.1

We need to show that, given a function $f: X \to Y$, the diagram

$$X \xrightarrow{f} Y$$

$$\downarrow^{\Delta_X} \qquad \qquad \downarrow^{\Delta_Y}$$

$$X \times X \xrightarrow{f \times f} Y \times Y$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
x & & x & \longrightarrow f(x) \\
\downarrow & & \downarrow \\
(x,x) & \longmapsto (f(x),f(x)) & & & (f(x),f(x))
\end{array}$$

and hence indeed commutes, showing Δ to be natural.

PROPOSITION 5.1.8.1.3 ► PROPERTIES OF THE DIAGONAL MAP

Let *X* be a set.

1. Monoidality. The diagonal map

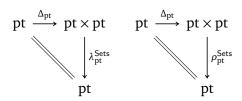
$$\Delta : id_{\mathsf{Sets}} \Longrightarrow \times \circ \Delta^{\mathsf{Cats}_2}_{\mathsf{Sets}},$$

is a monoidal natural transformation:

(a) Compatibility With Strong Monoidality Constraints. For each $X, Y \in Obj(Sets)$, the diagram

commutes.

(b) Compatibility With Strong Unitality Constraints. The diagrams



commute, i.e. we have

$$\begin{split} \Delta_{pt} &= \lambda_{pt}^{\mathsf{Sets},-1} \\ &= \rho_{pt}^{\mathsf{Sets},-1}, \end{split}$$

where we recall that the equalities

$$\begin{split} \lambda_{\text{pt}}^{\text{Sets}} &= \rho_{\text{pt}}^{\text{Sets}}, \\ \lambda_{\text{pt}}^{\text{Sets},-1} &= \rho_{\text{pt}}^{\text{Sets},-1} \end{split}$$

are always true in any monoidal category by Monoidal Categories, ?? of ??.

2. The Diagonal of the Unit. The component

$$\Delta_{pt} \colon pt \xrightarrow{\sim} pt \times pt$$

of Δ at pt is an isomorphism.

PROOF 5.1.8.1.4 ▶ PROOF OF PROPOSITION 5.1.8.1.3

Item 1: Monoidality

We claim that Δ is indeed monoidal:

1. Item 1a: Compatibility With Strong Monoidality Constraints: We need to show that the diagram

$$X \times Y \xrightarrow{\Delta_X \times \Delta_Y} (X \times X) \times (Y \times Y)$$

$$\downarrow \lambda \qquad \qquad \downarrow \lambda \qquad$$

commutes. Indeed, this diagram acts on elements as

and hence indeed commutes.

2. *Item 1b*: Compatibility With Strong Unitality Constraints: As shown in the proof of Definition 5.1.5.1.1, the inverse of the left unitor of Sets with respect to to the product at $X \in Obj(Sets)$ is given by

$$\lambda_X^{\mathsf{Sets},-1}(x) \stackrel{\mathrm{def}}{=} (\star, x)$$

for each $x \in X$, so when X = pt, we have

$$\lambda_{\mathrm{pt}}^{\mathrm{Sets},-1}(\star) \stackrel{\mathrm{def}}{=} (\star,\star),$$

and also

$$\Delta_{\mathrm{pt}}^{\mathrm{Sets}}(\star) \stackrel{\mathrm{def}}{=} (\star, \star)$$

 $\Delta_{pt}^{Sets}(\star)\stackrel{\text{def}}{=}(\star,\star),$ so we have $\Delta_{pt}=\lambda_{pt}^{Sets,-1}.$

This finishes the proof.

Item 2: The Diagonal of the Unit

This follows from <a>Item 1 and the invertibility of the left/right unitor of Sets with respect to ×, proved in the proof of Definition 5.1.5.1.1 for the left unitor or the proof of Definition 5.1.6.1.1 for the right unitor.

5.1.9 The Monoidal Category of Sets and Products

PROPOSITION 5.1.9.1.1 ➤ THE MONOIDAL STRUCTURE ON SETS ASSOCIATED TO THE PRODUCT

The category Sets admits a closed symmetric monoidal category with diagonals structure consisting of:

- *The Underlying Category*. The category Sets of pointed sets.
- The Monoidal Product. The product functor

$$\times$$
: Sets \times Sets \rightarrow Sets

of Constructions With Sets, Item 1 of Proposition 4.1.3.1.4.

• The Internal Hom. The internal Hom functor

$$\mathsf{Sets} \colon \mathsf{Sets}^\mathsf{op} \times \mathsf{Sets} \to \mathsf{Sets}$$

of Constructions With Sets, Item 1 of Proposition 4.3.5.1.2.

• The Monoidal Unit. The functor

$$\mathbb{1}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

of Definition 5.1.3.1.1.

• The Associators. The natural isomorphism

$$\alpha^{\mathsf{Sets}} \colon \times \circ (\times \times \mathsf{id}_{\mathsf{Sets}}) \stackrel{\widetilde{\longrightarrow}}{\Longrightarrow} \times \circ (\mathsf{id}_{\mathsf{Sets}} \times \times) \circ \pmb{\alpha}^{\mathsf{Cats}}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}$$

of Definition 5.1.4.1.1.

• The Left Unitors. The natural isomorphism

$$\lambda^{\mathsf{Sets}} : \times \circ \left(\mathbb{1}^{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}} \right) \stackrel{\sim}{\Longrightarrow} \lambda^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.1.5.1.1.

• The Right Unitors. The natural isomorphism

$$\rho^{\mathsf{Sets}} : \times \circ \left(\mathsf{id} \times \mathbb{1}^{\mathsf{Sets}} \right) \stackrel{\sim}{\Longrightarrow} \rho_{\mathsf{Sets}}^{\mathsf{Cats}_2}$$

of Definition 5.1.6.1.1.

• The Symmetry. The natural isomorphism

$$\sigma^{\mathsf{Sets}} \colon \times \stackrel{\widetilde{-}}{\Longrightarrow} \times \circ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}}$$

of Definition 5.1.7.1.1.

• The Diagonals. The monoidal natural transformation

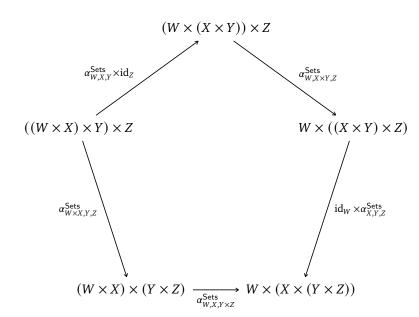
$$\Delta\colon\operatorname{id}_{\mathsf{Sets}}\Longrightarrow\times\circ\Delta^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.1.8.1.1.

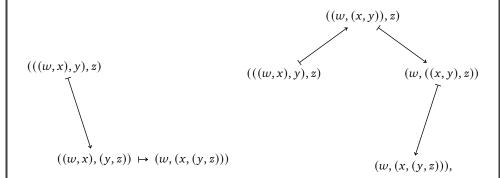
PROOF 5.1.9.1.2 ► PROOF OF PROPOSITION 5.1.9.1.1

The Pentagon Identity

Let W, X, Y and Z be sets. We have to show that the diagram



commutes. Indeed, this diagram acts on elements as



and thus the pentagon identity is satisfied.

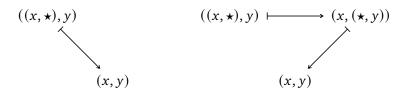
The Triangle Identity

Let *X* and *Y* be sets. We have to show that the diagram

$$(X \times \text{pt}) \times Y \xrightarrow{\alpha_{X,\text{pt},Y}^{\text{Sets}}} X \times (\text{pt} \times Y)$$

$$\rho_X^{\text{Sets}} \times \text{id}_Y \xrightarrow{Id_X \times \lambda_Y^{\text{Sets}}} X \times Y$$

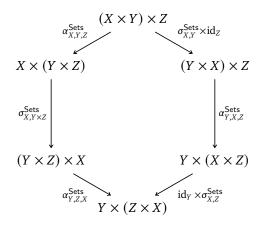
commutes. Indeed, this diagram acts on elements as



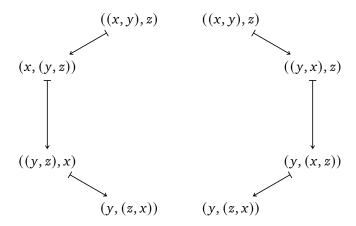
and thus the triangle identity is satisfied.

The Left Hexagon Identity

Let *X*, *Y*, and *Z* be sets. We have to show that the diagram



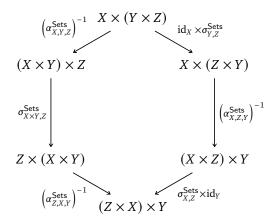
commutes. Indeed, this diagram acts on elements as



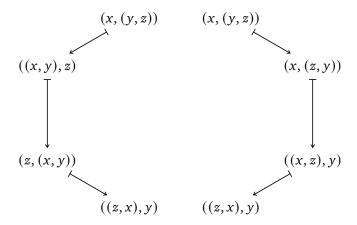
and thus the left hexagon identity is satisfied.

The Right Hexagon Identity

Let *X*, *Y*, and *Z* be sets. We have to show that the diagram



commutes. Indeed, this diagram acts on elements as



and thus the right hexagon identity is satisfied.

Monoidal Closedness

This follows from Constructions With Sets, Item 2 of Proposition 4.3.5.1.2

Existence of Monoidal Diagonals

This follows from Items 1 and 2 of Proposition 5.1.8.1.3.

5.1.10 The Universal Property of (Sets, \times , pt)

THEOREM 5.1.10.1.1 ► THE UNIVERSAL PROPERTY OF (Sets, ×, PT)

The symmetric monoidal structure on the category Sets of Proposition 5.1.9.1.1 is uniquely determined by the following requirements:

1. Existence of an Internal Hom. The tensor product

$$\otimes_{\mathsf{Sets}} \colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}$$

of Sets admits an internal Hom $[-1, -2]_{Sets}$

2. The Unit Object Is pt. We have $\mathbb{1}_{Sets} \cong pt$.

More precisely, the full subcategory of the category $\mathcal{M}^{cld}_{\mathbb{E}_{\infty}}(\mathsf{Sets})$ of ?? spanned by the closed symmetric monoidal categories (Sets, \otimes_{Sets} , $[-_1, -_2]_{\mathsf{Sets}}$, 1_{Sets} , $1_{\mathsf{Se$

PROOF 5.1.10.1.2 ▶ PROOF OF THEOREM 5.1.10.1.1

Unwinding the Statement

Let (Sets, \otimes_{Sets} , $[-_1, -_2]_{Sets}$, $\mathbb{1}_{Sets}$, λ' , ρ' , σ') be a closed symmetric monoidal category satisfying Items 1 and 2. We need to show that the identity functor

$$id_{Sets}$$
: Sets \rightarrow Sets

admits a unique closed symmetric monoidal functor structure

making it into a symmetric monoidal strongly closed isomorphism of categories from (Sets, \otimes_{Sets} , $[-1,-2]_{Sets}$, $\mathbb{1}_{Sets}$, λ' , ρ' , σ') to the closed symmetric monoidal category (Sets, \times , Sets(-1,-2), $\mathbb{1}_{Sets}$, λ^{Sets} , ρ^{Sets} , σ^{Sets}) of Proposition 5.1.9.1.1.

Constructing an Isomorphism $[-1, -2]_{Sets} \cong Sets(-1, -2)$

By ??, we have a natural isomorphism

$$Sets(pt, [-1, -2]_{Sets}) \cong Sets(-1, -2).$$

By Constructions With Sets, Item 3 of Proposition 4.3.5.1.2, we also have a natural isomorphism

$$\mathsf{Sets}(\mathsf{pt}, [-_1, -_2]_{\mathsf{Sets}}) \cong [-_1, -_2]_{\mathsf{Sets}}.$$

Composing both natural isomorphisms, we obtain a natural isomorphism

$$Sets(-1, -2) \cong [-1, -2]_{Sets}$$
.

Given $A, B \in Obj(Sets)$, we will write

$$id_{AB}^{Hom}$$
: Sets $(A, B) \xrightarrow{\sim} [A, B]_{Sets}$

for the component of this isomorphism at (A, B).

Constructing an Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$

Since \otimes_{Sets} is adjoint in each variable to $[-_1, -_2]_{\mathsf{Sets}}$ by assumption and \times is adjoint in each variable to $\mathsf{Sets}(-_1, -_2)$ by Constructions With Sets, Item 2 of Proposition 4.3.5.1.2, uniqueness of adjoints (??) gives us natural isomorphisms

$$A \otimes_{\mathsf{Sets}} - \cong A \times -,$$

 $- \otimes_{\mathsf{Sets}} B \cong B \times -.$

By ??, we then have $\otimes_{\mathsf{Sets}} \cong \times$. We will write

$$\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes} \colon A \otimes_{\mathsf{Sets}} B \xrightarrow{\sim} A \times B$$

for the component of this isomorphism at (A, B).

Alternative Construction of an Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$

Alternatively, we may construct a natural isomorphism $\otimes_{\mathsf{Sets}} \cong \times$ as follows:

- 1. Let $A \in Obj(Sets)$.
- 2. Since \otimes_{Sets} is part of a closed monoidal structure, it preserves colimits in each variable by $\ref{eq:sets}$?

3. Since $A \cong \coprod_{a \in A} pt$ and \otimes_{Sets} preserves colimits in each variable, we have

$$A \otimes_{\mathsf{Sets}} B \cong \left(\bigsqcup_{\alpha \in A} \mathsf{pt} \right) \otimes_{\mathsf{Sets}} B$$

$$\cong \bigsqcup_{\alpha \in A} (\mathsf{pt} \otimes_{\mathsf{Sets}} B)$$

$$\cong \bigsqcup_{\alpha \in A} B$$

$$\cong A \times B,$$

naturally in $B \in \text{Obj}(\mathsf{Sets})$, where we have used that pt is the monoidal unit for \otimes_{Sets} . Thus $A \otimes_{\mathsf{Sets}} - \cong A \times -$ for each $A \in \mathsf{Obj}(\mathsf{Sets})$.

- 4. Similarly, $-\otimes_{\mathsf{Sets}} B \cong -\times B$ for each $B \in \mathsf{Obj}(\mathsf{Sets})$.
- 5. By ??, we then have $\otimes_{\mathsf{Sets}} \cong \times$.

Below, we'll show that if a natural isomorphism $\otimes_{\mathsf{Sets}} \cong \times$ exists, then it must be unique. This will show that the isomorphism constructed above is equal to the isomorphism $\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes} \colon A \otimes_{\mathsf{Sets}} B \to A \times B$ from before.

Constructing an Isomorphism $id_1^{\otimes} : \mathbb{1}_{\mathsf{Sets}} \to \mathsf{pt}$

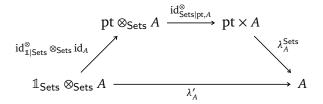
We define an isomorphism $id_1^{\otimes} : \mathbb{1}_{\mathsf{Sets}} \to \mathsf{pt}$ as the composition

$$\mathbb{1}_{\mathsf{Sets}} \overset{\rho^{\mathsf{Sets},-1}_{\mathbb{1}_{\mathsf{Sets}}}}{\overset{\circ}{\longrightarrow}} \mathbb{1}_{\mathsf{Sets}} \times \mathsf{pt} \overset{id^{\otimes}_{\mathsf{Sets} \mid \mathbb{1}_{\mathsf{Sets}}}}{\overset{\circ}{\longrightarrow}} \mathbb{1}_{\mathsf{Sets}} \otimes_{\mathsf{Sets}} \mathsf{pt} \overset{\lambda'_{\mathsf{pt}}}{\overset{\circ}{\longrightarrow}} \mathsf{pt}$$

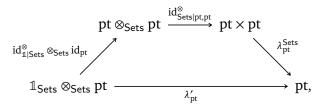
in Sets.

Monoidal Left Unity of the Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$

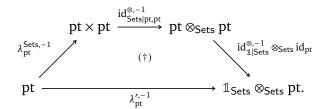
We have to show that the diagram



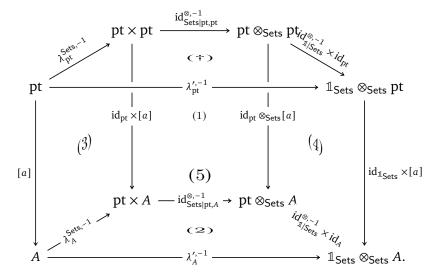
commutes. First, note that the diagram



corresponding to the case A = pt, commutes by the terminality of pt (Constructions With Sets, Construction 4.1.1.1.2). Since this diagram commutes, so does the diagram



Now, let $A \in \text{Obj}(\mathsf{Sets})$, let $a \in A$, and consider the diagram



Since:

• Subdiagram (5) commutes by the naturality of λ'^{-1} .

- Subdiagram (†) commutes, as proved above.
- Subdiagram (4) commutes by the naturality of $id_{1|Sets}^{\otimes,-1}$.
- Subdiagram (1) commutes by the naturality of $\mathrm{id}_{\mathsf{Sets}}^{\otimes,-1}$.
- Subdiagram (3) commutes by the naturality of $\lambda^{Sets,-1}$.

it follows that the diagram

$$\operatorname{pt} \times A \xrightarrow{\operatorname{id}_{\mathsf{Sets}|\mathsf{pt},A}^{\otimes,-1}} \operatorname{pt} \otimes_{\mathsf{Sets}} A$$

$$\downarrow^{\mathsf{Sets},-1} \wedge \qquad \qquad \downarrow^{\operatorname{id}_{\mathsf{Sets}}^{\otimes,-1}} \otimes_{\mathsf{Sets}} \operatorname{id}_{A}$$

$$\operatorname{pt} \xrightarrow{[a]} A \xrightarrow{\lambda_{A}^{\prime,-1}} \mathbb{1}_{\mathsf{Sets}} \otimes_{\mathsf{Sets}} A$$

Here's a step-by-step showcase of this argument: [Link]. We then have

$$\begin{split} \lambda_A^{\prime,-1}(a) &= \left[\lambda_A^{\prime,-1} \circ [a]\right](\star) \\ &= \left[\left(\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \times \mathrm{id}_A\right) \circ \mathrm{id}_{\mathsf{Sets}|\mathsf{pt},A}^{\otimes,-1} \circ \lambda_A^{\mathsf{Sets},-1} \circ [a]\right](\star) \\ &= \left[\left(\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \times \mathrm{id}_A\right) \circ \mathrm{id}_{\mathsf{Sets}|\mathsf{pt},A}^{\otimes,-1} \circ \lambda_A^{\mathsf{Sets},-1}\right](a) \end{split}$$

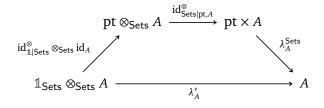
for each $a \in A$, and thus we have

$$\lambda_A^{\prime,-1} = \left(\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \times \mathrm{id}_A \right) \circ \mathrm{id}_{\mathsf{Sets}|\mathsf{pt},A}^{\otimes,-1} \circ \lambda_A^{\mathsf{Sets},-1}.$$

Taking inverses then gives

$$\lambda_A' = \lambda_A^{\mathsf{Sets}} \circ \mathsf{id}_{\mathsf{Sets}|\mathsf{pt},A}^{\otimes} \circ \left(\mathsf{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes} \times \mathsf{id}_A\right)$$

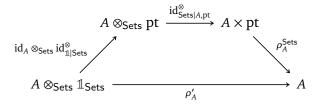
showing that the diagram



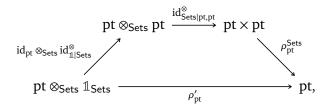
indeed commutes.

Monoidal Right Unity of the Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$

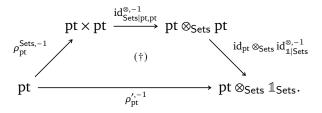
We can use the same argument we used to prove the monoidal left unity of the isomorphism $\otimes_{\mathsf{Sets}} \cong \times$ above. For completeness, we repeat it below. We have to show that the diagram

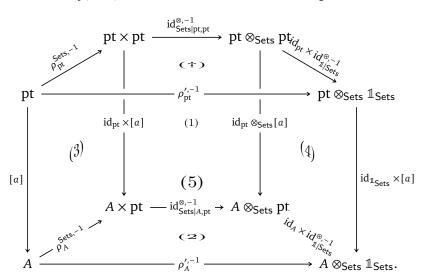


commutes. First, note that the diagram



corresponding to the case A = pt, commutes by the terminality of pt (Constructions With Sets, Construction 4.1.1.1.2). Since this diagram commutes, so does the diagram



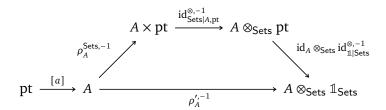


Now, let $A \in \text{Obj}(\mathsf{Sets})$, let $a \in A$, and consider the diagram

Since:

- Subdiagram (5) commutes by the naturality of ρ'^{-1} .
- Subdiagram (†) commutes, as proved above.
- Subdiagram (4) commutes by the naturality of $id_{1|Sets}^{\otimes,-1}$.
- Subdiagram (1) commutes by the naturality of $\mathrm{id}_{\mathsf{Sets}}^{\otimes,-1}$.
- Subdiagram (3) commutes by the naturality of $\rho^{\mathsf{Sets},-1}$.

it follows that the diagram



Here's a step-by-step showcase of this argument: [Link]. We then have

$$\rho_A^{\prime,-1}(a) = \left[\rho_A^{\prime,-1} \circ [a]\right](\star)$$

$$= \left[\left(id_A \times id_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \right) \circ id_{\mathsf{Sets}|\mathsf{pt},A}^{\otimes,-1} \circ \rho_A^{\mathsf{Sets},-1} \circ [a] \right] (\star)$$

$$= \left[\left(id_A \times id_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \right) \circ id_{\mathsf{Sets}|\mathsf{pt},A}^{\otimes,-1} \circ \rho_A^{\mathsf{Sets},-1} \right] (a)$$

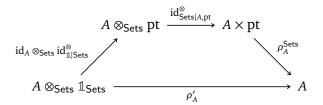
for each $a \in A$, and thus we have

$$\rho_A^{\prime,-1} = \left(\mathrm{id}_A \times \mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1}\right) \circ \mathrm{id}_{\mathsf{Sets}|\mathsf{pt},A}^{\otimes,-1} \circ \rho_A^{\mathsf{Sets},-1}.$$

Taking inverses then gives

$$\rho_A' = \rho_A^{\mathsf{Sets}} \circ \mathrm{id}_{\mathsf{Sets}|\mathsf{pt},A}^{\otimes} \circ (\mathrm{id}_A \times \mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes}),$$

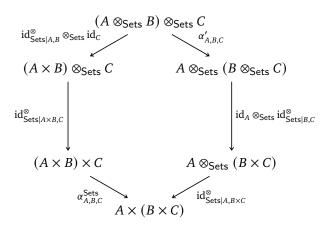
showing that the diagram

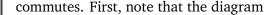


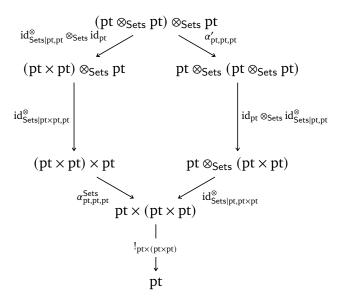
indeed commutes.

Monoidality of the Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$

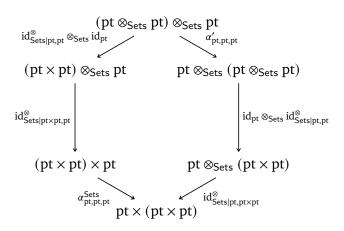
We have to show that the diagram



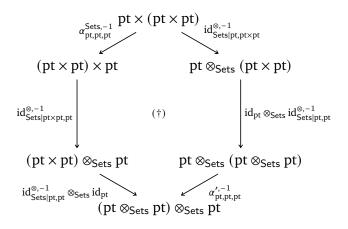




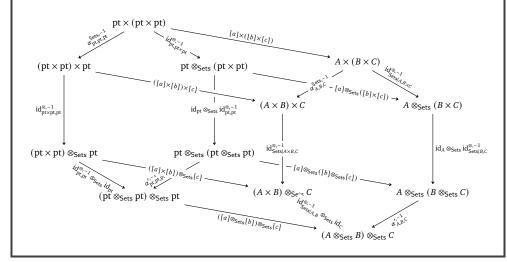
commutes by the terminality of pt (Constructions With Sets, Construction 4.1.1.2). Since the map $!_{pt \times (pt \times pt)} : pt \times (pt \times pt) \to pt$ is an isomorphism (e.g. having inverse $\lambda_{pt}^{\mathsf{Sets},-1} \circ \lambda_{pt}^{\mathsf{Sets},-1}$), it follows that the diagram

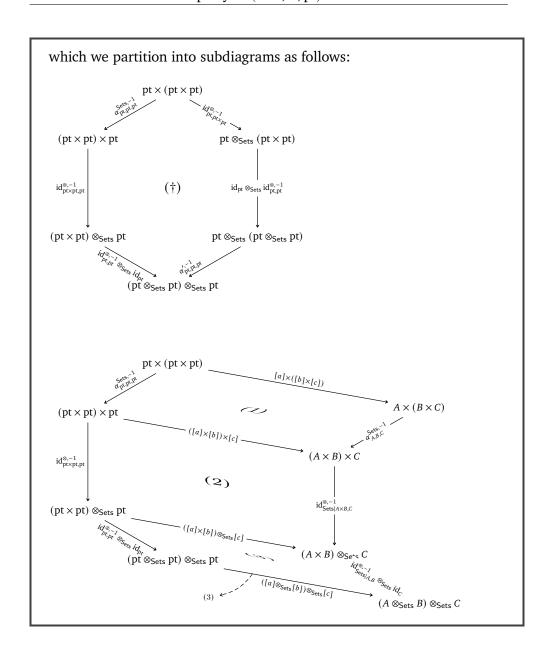


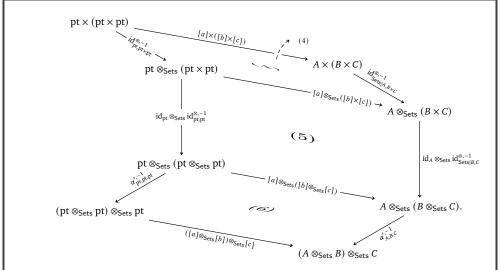
also commutes. Taking inverses, we see that the diagram



commutes as well. Now, let $A, B, C \in \text{Obj}(\mathsf{Sets})$, let $a \in A$, let $b \in B$, let $c \in C$, and consider the diagram



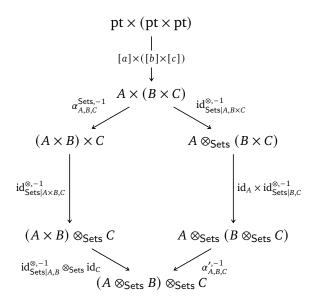




Since:

- Subdiagram (1) commutes by the naturality of $\alpha^{Sets,-1}$.
- Subdiagram (2) commutes by the naturality of $id_{Sets}^{\otimes,-1}$.
- Subdiagram (3) commutes by the naturality of $\mathrm{id}_{\mathsf{Sets}}^{\otimes,-1}$.
- Subdiagram (†) commutes, as proved above.
- Subdiagram (4) commutes by the naturality of $id_{\mathsf{Sets}}^{\otimes,-1}$.
- Subdiagram (5) commutes by the naturality of $\mathrm{id}_{\mathsf{Sets}}^{\otimes,-1}$
- Subdiagram (6) commutes by the naturality of α'^{-1} .

it follows that the diagram



also commutes. We then have

$$\begin{split} \left[\left(\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} \otimes_{\mathsf{Sets}} \mathrm{id}_{C} \right) \circ \mathrm{id}_{\mathsf{Sets}|A\times B,C}^{\otimes,-1} \\ & \circ \alpha_{A,B,C}^{\mathsf{Sets},-1} \right] (a,(b,c)) = \left[\left(\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} \otimes_{\mathsf{Sets}} \mathrm{id}_{C} \right) \circ \mathrm{id}_{\mathsf{Sets}|A\times B,C}^{\otimes,-1} \\ & \circ \alpha_{A,B,C}^{\mathsf{Sets},-1} \circ \left([a] \times ([b] \times [c]) \right) \right] (\star,(\star,\star)) \\ & = \left[\alpha_{A,B,C}^{\prime,-1} \circ \left(\mathrm{id}_{A} \times \mathrm{id}_{\mathsf{Sets}|B,C}^{\otimes,-1} \right) \\ & \circ \mathrm{id}_{\mathsf{Sets}|A,B\times C}^{\otimes,-1} \circ \left([a] \times ([b] \times [c]) \right) \right] (\star,(\star,\star)) \\ & = \left[\alpha_{A,B,C}^{\prime,-1} \circ \left(\mathrm{id}_{A} \times \mathrm{id}_{\mathsf{Sets}|B,C}^{\otimes,-1} \right) \circ \mathrm{id}_{\mathsf{Sets}|A,B\times C}^{\otimes,-1} \right] (a,(b,c)) \end{split}$$

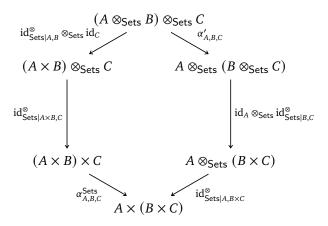
for each $(a, (b, c)) \in A \times (B \times C)$, and thus we have

$$\left(\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} \otimes_{\mathsf{Sets}} \mathrm{id}_{C}\right) \circ \mathrm{id}_{\mathsf{Sets}|A \times B,C}^{\otimes,-1} \circ \alpha_{A,B,C}^{\mathsf{Sets},-1} = \alpha_{A,B,C}'^{,-1} \circ \left(\mathrm{id}_{A} \times \mathrm{id}_{\mathsf{Sets}|B,C}^{\otimes,-1}\right) \circ \mathrm{id}_{\mathsf{Sets}|A,B \times C}^{\otimes,-1} \,.$$

Taking inverses then gives

$$\alpha_{A,B,C}^{\mathsf{Sets}} \circ \mathrm{id}_{\mathsf{Sets}|A \times B,C}^{\otimes} \circ \left(\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes} \otimes_{\mathsf{Sets}} \mathrm{id}_{C} \right) = \mathrm{id}_{\mathsf{Sets}|A,B \times C}^{\otimes} \circ \left(\mathrm{id}_{A} \times \mathrm{id}_{\mathsf{Sets}|B,C}^{\otimes} \right) \circ \alpha_{A,B,C}',$$

showing that the diagram



indeed commutes.

Braidedness of the Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$

We have to show that the diagram

$$\begin{array}{c|c} A \otimes_{\mathsf{Sets}} B & \xrightarrow{\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes}} & A \times B \\ & \sigma_{A,B}' & & & & \downarrow \sigma_{A,B}^{\mathsf{Sets}} \\ & B \otimes_{\mathsf{Sets}} A & \xrightarrow{\mathrm{id}_{\mathsf{Sets}|B,A}^{\otimes}} & B \times A \end{array}$$

commutes. First, note that the diagram

commutes by the terminality of pt (Constructions With Sets, Construction 4.1.1.2). Since the map $!_{pt \times pt} : pt \times pt \rightarrow pt$ is invertible (e.g. with

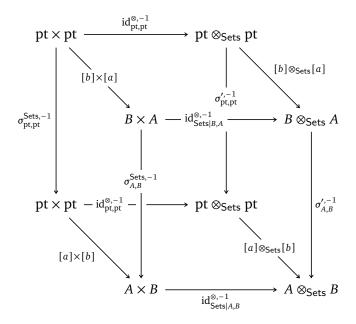
inverse $\lambda_{pt}^{\mathsf{Sets},-1}),$ the diagram

$$\begin{array}{ccc} pt \otimes_{\mathsf{Sets}} pt & \xrightarrow{id^{\otimes}_{\mathsf{Sets}|pt,pt}} & pt \times pt \\ & & & & \downarrow \sigma^{\mathsf{Sets}}_{pt,pt} \\ & & & \downarrow \sigma^{\mathsf{Sets}}_{pt,pt} \\ & & & \downarrow d^{\mathsf{Sets}}_{pt,pt} & pt \times pt \end{array}$$

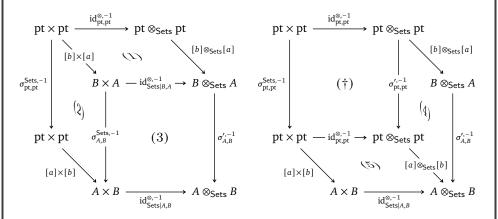
also commutes. Taking inverses, we see that the diagram

$$\begin{array}{ccc} pt \times pt & \xrightarrow{id_{\mathsf{Sets}|pt,pt}^{\otimes,-1}} & pt \otimes_{\mathsf{Sets}} pt \\ \sigma_{\mathsf{pt},\mathsf{pt}}^{\mathsf{Sets},-1} & & & & & & \\ \sigma_{\mathsf{pt},\mathsf{pt}}^{\mathsf{Sets},-1} & & & & & & \\ pt \times pt & \xrightarrow{id_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}}^{\otimes,-1}} & pt \otimes_{\mathsf{Sets}} pt \end{array}$$

commutes as well. Now, let $A, B \in \text{Obj}(\mathsf{Sets})$, let $a \in A$, let $b \in B$, and consider the diagram



which we partition into subdiagrams as follows:



Since:

- Subdiagram (2) commutes by the naturality of $\sigma^{\text{Sets},-1}$.
- Subdiagram (5) commutes by the naturality of $id^{\otimes,-1}$.
- Subdiagram (†) commutes, as proved above.
- Subdiagram (4) commutes by the naturality of σ'^{-1} .
- Subdiagram (1) commutes by the naturality of $id^{\otimes,-1}$.

it follows that the diagram

commutes. We then have

$$\left[\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} \circ \sigma_{A,B}^{\mathsf{Sets},-1}\right](b,a) = \left[\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} \circ \sigma_{A,B}^{\mathsf{Sets},-1} \circ ([b] \times [a])\right](\star,\star)$$

$$= \left[\sigma'_{A,B}^{,-1} \circ \mathrm{id}_{\mathsf{Sets}|B,A}^{\otimes,-1} \circ ([b] \times [a])\right](\star, \star)$$

$$= \left[\sigma'_{A,B}^{,-1} \circ \mathrm{id}_{\mathsf{Sets}|B,A}^{\otimes,-1}\right](b,a)$$

for each $(b, a) \in B \times A$, and thus we have

$$\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} \circ \sigma_{A,B}^{\mathsf{Sets},-1} = \sigma_{A,B}^{\prime,-1} \circ \mathrm{id}_{\mathsf{Sets}|B,A}^{\otimes,-1}$$

Taking inverses then gives

$$\sigma_{A,B}^{\mathsf{Sets}} \circ \mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes} = \mathrm{id}_{\mathsf{Sets}|B,A}^{\otimes} \circ \sigma_{A,B}'$$

showing that the diagram

$$A \otimes_{\mathsf{Sets}} B \xrightarrow{\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes}} A \times B$$

$$\sigma'_{A,B} \downarrow \qquad \qquad \qquad \downarrow \sigma^{\mathsf{Sets}}_{A,B}$$

$$B \otimes_{\mathsf{Sets}} A \xrightarrow{\mathrm{id}_{\mathsf{Sets}|B,A}^{\otimes}} B \times A$$

indeed commutes.

Uniqueness of the Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$

Let $\phi, \psi \colon \neg_1 \otimes_{\mathsf{Sets}} \neg_2 \Rightarrow \neg_1 \times \neg_2$ be natural isomorphisms. Since these isomorphisms are compatible with the unitors of Sets with respect to \times and \otimes (as shown above), we have

$$\lambda_B' = \lambda_B^{\mathsf{Sets}} \circ \phi_{\mathsf{pt},B} \circ \left(\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes} \otimes_{\mathsf{Sets}} \mathrm{id}_Y \right),$$
$$\lambda_B' = \lambda_B^{\mathsf{Sets}} \circ \psi_{\mathsf{pt},B} \circ \left(\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes} \otimes_{\mathsf{Sets}} \mathrm{id}_Y \right).$$

Postcomposing both sides with $\lambda_B^{\mathsf{Sets},-1}$ gives

$$\begin{split} & \lambda_{B}^{\mathsf{Sets},-1} \circ \lambda_{B}' \circ \left(\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \otimes_{\mathsf{Sets}} \mathrm{id}_{Y} \right) = \phi_{\mathsf{pt},B}, \\ & \lambda_{B}^{\mathsf{Sets},-1} \circ \lambda_{B}' \circ \left(\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes} \otimes_{\mathsf{Sets}} \mathrm{id}_{Y} \right) = \psi_{\mathsf{pt},B}, \end{split}$$

and thus we have

$$\phi_{\text{pt},B} = \psi_{\text{pt},B}$$

for each $B \in \text{Obj}(\mathsf{Sets})$. Now, let $a \in A$ and consider the naturality diagrams

for ϕ and ψ with respect to the morphisms [a] and id $_B$. Having shown that $\phi_{\text{pt},B} = \psi_{\text{pt},B}$, we have

$$\phi_{A,B}(a,b) = [\phi_{A,B} \circ ([a] \times id_B)](\star,b)$$

$$= [([a] \otimes_{\mathsf{Sets}} id_B) \circ \phi_{\mathsf{pt},B}](\star,b)$$

$$= [([a] \otimes_{\mathsf{Sets}} id_B) \circ \psi_{\mathsf{pt},B}](\star,b)$$

$$= [\psi_{A,B} \circ ([a] \times id_B)](\star,b)$$

$$= \psi_{A,B}(a,b)$$

for each $(a, b) \in A \times B$. Therefore we have

$$\phi_{A,B} = \psi_{A,B}$$

for each $A, B \in \text{Obj}(\mathsf{Sets})$ and thus $\phi = \psi$, showing the isomorphism $\otimes_{\mathsf{Sets}} \cong \mathsf{x}$ to be unique.

COROLLARY 5.1.10.1.3 ► A SECOND UNIVERSAL PROPERTY FOR (Sets, ×, PT)

The symmetric monoidal structure on the category Sets of Proposition 5.1.9.1.1 is uniquely determined by the following requirements:

1. Two-Sided Preservation of Colimits. The tensor product

$$\otimes_{\mathsf{Sets}} \colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}$$

of Sets preserves colimits separately in each variable.

2. The Unit Object Is pt. We have $\mathbb{1}_{Sets} \cong pt$.

More precisely, the full subcategory of the category $\mathcal{M}_{\mathbb{E}_{\infty}}(\mathsf{Sets})$ of $\ref{eq:subcategory}$ spanned by the symmetric monoidal categories (Sets, \otimes_{Sets} , $\mathbb{1}_{\mathsf{Sets}}$, λ^{Sets} , ρ^{Sets} , σ^{Sets}) satisfying Items 1 and 2 is contractible.

Since Sets is locally presentable (??), it follows from ?? that Item 1 is equivalent to the existence of an internal Hom as in Item 1 of Theorem 5.1.10.1.1.

The result then follows from Theorem 5.1.10.1.1.

5.2 The Monoidal Category of Sets and Coproducts

5.2.1 Coproducts of Sets

See Constructions With Sets, Section 4.2.3.

5.2.2 The Monoidal Unit

DEFINITION 5.2.2.1.1 ► THE MONOIDAL UNIT OF [

The monoidal unit of the coproduct of sets is the functor

$$\mathbb{O}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

defined by

$$\mathbb{O}_{\mathsf{Sets}} \stackrel{\mathrm{def}}{=} \emptyset$$
,

where Ø is the empty set of Constructions With Sets, Definition 4.3.1.1.1.

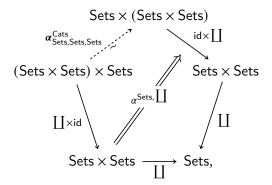
5.2.3 The Associator

DEFINITION 5.2.3.1.1 ► THE ASSOCIATOR OF ∐

The associator of the coproduct of sets is the natural isomorphism

$$\alpha^{\mathsf{Sets}, \coprod} : \coprod \circ (\coprod \times \mathsf{id}_{\mathsf{Sets}}) \stackrel{\widetilde{-}}{\Longrightarrow} \coprod \circ (\mathsf{id}_{\mathsf{Sets}} \times \coprod) \circ \alpha^{\mathsf{Cats}}_{\mathsf{Sets}, \mathsf{Sets}, \mathsf{Sets}},$$

as in the diagram



whose component

$$\alpha^{\mathsf{Sets}, \coprod}_{X,Y,Z} : (X \coprod Y) \coprod Z \xrightarrow{\sim} X \coprod (Y \coprod Z)$$

at (X, Y, Z) is given by

$$\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod}(a) \stackrel{\mathsf{def}}{=} \begin{cases} (0,x) & \text{if } a = (0,(0,x)), \\ (1,(0,y)) & \text{if } a = (0,(1,y)), \\ (1,(1,a)) & \text{if } a = (1,z) \end{cases}$$

for each $a \in (X \coprod Y) \coprod Z$.

PROOF 5.2.3.1.2 ► PROOF OF THE CLAIMS MADE IN DEFINITION 5.2.3.1.1

Unwinding the Definitions of $(X \coprod Y) \coprod Z$ and $X \coprod (Y \coprod Z)$

Firstly, we unwind the expressions for $(X \coprod Y) \coprod Z$ and $X \coprod (Y \coprod Z)$. We have

$$\begin{split} (X \coprod Y) \coprod Z &\stackrel{\text{def}}{=} \{ (0, a) \in S \mid a \in X \coprod Y \} \cup \{ (1, z) \in S \mid z \in Z \} \\ &= \{ (0, (0, x)) \in S \mid x \in X \} \cup \{ (0, (1, y)) \in S \mid y \in Y \} \\ & \cup \{ (1, z) \in S \mid z \in Z \}, \end{split}$$

where $S = \{0, 1\} \times ((X \coprod Y) \cup Z)$ and

$$X \coprod (Y \coprod Z) \stackrel{\mathsf{def}}{=} \{(0,x) \in S' \mid x \in X\} \cup \{(1,a) \in S' \mid a \in Y \coprod Z\}$$

$$= \{(0, x) \in S' \mid x \in X\} \cup \{(1, (0, y)) \in S' \mid y \in Y\}$$
$$\cup \{(1, (1, z)) \in S' \mid z \in Z\},$$

where $S' = \{0, 1\} \times (X \cup (Y \mid I \mid Z))$.

Invertibility

The inverse of $\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod}$ is the map

$$\alpha_{X,Y,Z}^{\mathsf{Sets}, \coprod, -1} \colon X \coprod (Y \coprod Z) \to (X \coprod Y) \coprod Z$$

given by

$$\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod,-1}(a) \stackrel{\mathsf{def}}{=} \begin{cases} (0,(0,x)) & \text{if } a = (0,x), \\ (0,(1,y)) & \text{if } a = (1,(0,y)), \\ (1,z) & \text{if } a = (1,(1,z)) \end{cases}$$

for each $a \in X \coprod Y(\coprod Z)$. Indeed:

• *Invertibility I*. The map $\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod,-1} \circ \alpha_{X,Y,Z}^{\mathsf{Sets},\coprod}$ acts on elements as

$$(0, (0, x)) \mapsto (0, x) \mapsto (0, (0, x)),$$

$$(0, (0, y)) \mapsto (1, (0, y)) \mapsto (0, (0, y)),$$

$$(1, z) \mapsto (1, (1, z)) \mapsto (1, z)$$

and hence is equal to the identity map of $(X \coprod Y) \coprod Z$.

• *Invertibility II.* The map $\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod} \circ \alpha_{X,Y,Z}^{\mathsf{Sets},\coprod,-1}$ acts on elements as

$$(0,x) \mapsto (0,(0,x)) \mapsto (0,x),$$

$$(1,(0,y)) \mapsto (0,(0,y)) \mapsto (1,(0,y)),$$

$$(1,(1,z)) \mapsto (1,z) \mapsto (1,(1,z))$$

and hence is equal to the identity map of $X \coprod (Y \coprod Z)$.

Therefore $\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod}$ is indeed an isomorphism.

Naturality

We need to show that, given functions

$$f: X \to X',$$

 $g: Y \to Y',$
 $h: Z \to Z'$

the diagram

commutes. Indeed, this diagram acts on elements as

$$(0, (0, x)) \qquad (0, (0, x)) \mapsto (0, (0, f(x)))$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad$$

and hence indeed commutes, showing $\alpha^{\mathsf{Sets}, \coprod}$ to be a natural transformation.

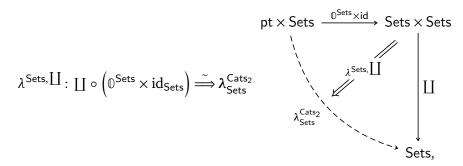
Being a Natural Isomorphism

Since $\alpha^{\mathsf{Sets}, \coprod}$ is natural and $\alpha^{\mathsf{Sets}, \coprod, -1}$ is a componentwise inverse to $\alpha^{\mathsf{Sets}, \coprod}$, it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that $\lambda^{\mathsf{Sets}, -1}$ is also natural. Thus $\alpha^{\mathsf{Sets}, \coprod}$ is a natural isomorphism.

5.2.4 The Left Unitor

DEFINITION 5.2.4.1.1 ► THE LEFT UNITOR OF [

The left unitor of the coproduct of sets is the natural isomorphism



whose component

$$\lambda_X^{\mathsf{Sets},\coprod} : \emptyset \coprod X \xrightarrow{\sim} X$$

at X is given by

$$\lambda_v^{\text{Sets},\coprod}((1,x)) \stackrel{\text{def}}{=} x$$

for each $(1, x) \in \emptyset \coprod X$.

PROOF 5.2.4.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.2.4.1.1

Unwinding the Definition of $\emptyset \coprod X$

Firstly, we unwind the expressions for $\emptyset \coprod X$. We have

where $S = \{0, 1\} \times (\emptyset \cup X)$.

Invertibility

The inverse of $\lambda_X^{\mathsf{Sets}, \coprod}$ is the map

$$\lambda_X^{\mathsf{Sets}, \coprod, -1} \colon X \to \emptyset \coprod X$$

given by

$$\lambda_X^{\mathsf{Sets},\coprod,-1}(x) \stackrel{\mathrm{def}}{=} (1,x)$$

for each $x \in X$. Indeed:

• Invertibility I. We have

$$\begin{split} \left[\lambda_X^{\mathsf{Sets}, \coprod, -1} \circ \lambda_X^{\mathsf{Sets}, \coprod}\right] (1, x) &= \lambda_X^{\mathsf{Sets}, \coprod, -1} \left(\lambda_X^{\mathsf{Sets}, \coprod} (1, x)\right) \\ &= \lambda_X^{\mathsf{Sets}, \coprod, -1} (x) \\ &= (1, x) \\ &= \left[\mathrm{id}_{\emptyset \coprod X}\right] (1, x) \end{split}$$

for each $(1, x) \in \emptyset \coprod X$, and therefore we have

$$\lambda_X^{\mathsf{Sets}, \coprod, -1} \circ \lambda_X^{\mathsf{Sets}, \coprod} = \mathrm{id}_{\emptyset \coprod X} \,.$$

• Invertibility II. We have

$$\begin{split} \left[\lambda_X^{\mathsf{Sets}, \coprod} \circ \lambda_X^{\mathsf{Sets}, \coprod, -1}\right](x) &= \lambda_X^{\mathsf{Sets}, \coprod} \left(\lambda_X^{\mathsf{Sets}, \coprod, -1}(x)\right) \\ &= \lambda_X^{\mathsf{Sets}, \coprod, -1}(1, x) \\ &= x \\ &= [\mathrm{id}_X](x) \end{split}$$

for each $x \in X$, and therefore we have

$$\lambda_X^{\mathsf{Sets}, \coprod} \, \circ \, \lambda_X^{\mathsf{Sets}, \coprod, -1} = \mathrm{id}_X \, .$$

Therefore $\lambda_X^{\text{Sets}, \coprod}$ is indeed an isomorphism.

Naturality

We need to show that, given a function $f: X \to Y$, the diagram

$$\begin{array}{ccc}
\emptyset \coprod X & \xrightarrow{\mathrm{id}_{\emptyset} \coprod f} \emptyset \coprod Y \\
\downarrow^{\lambda_{X}^{\mathsf{Sets}, \coprod}} & & \downarrow^{\lambda_{Y}^{\mathsf{Sets}, \coprod}} \\
X & \xrightarrow{f} & Y
\end{array}$$

commutes. Indeed, this diagram acts on elements as

$$(1,x) \qquad (1,x) \longmapsto (1,f(x))$$

$$\downarrow \qquad \qquad \qquad \downarrow$$

$$x \longmapsto f(x) \qquad \qquad f(x)$$

and hence indeed commutes. Therefore $\lambda^{\mathsf{Sets}, \coprod}$ is a natural transformation.

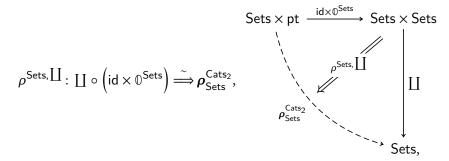
Being a Natural Isomorphism

Since $\lambda^{Sets, \coprod}$ is natural and $\lambda^{Sets, -1}$ is a componentwise inverse to $\lambda^{Sets, \coprod}$, it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that $\lambda^{Sets, -1}$ is also natural. Thus $\lambda^{Sets, \coprod}$ is a natural isomorphism.

5.2.5 The Right Unitor

DEFINITION 5.2.5.1.1 ► THE RIGHT UNITOR OF []

The **right unitor of the coproduct of sets** is the natural isomorphism



whose component

$$\rho_X^{\mathsf{Sets},\coprod}: X \coprod \overset{\sim}{ \mathbb{Q}} \overset{\sim}{\longrightarrow} X$$

at *X* is given by

$$\rho_X^{\mathsf{Sets},\coprod}((0,x)) \stackrel{\mathsf{def}}{=} x$$

for each $(0, x) \in X \coprod \emptyset$.

PROOF 5.2.5.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.2.5.1.1

Unwinding the Definition of $X \mid I \emptyset$

Firstly, we unwind the expression for $X \coprod \emptyset$. We have

$$X \coprod \emptyset \stackrel{\text{def}}{=} \{(0, x) \in S \mid x \in X\} \cup \{(1, z) \in S \mid z \in \emptyset\}$$
$$= \{(0, x) \in S \mid x \in X\} \cup \emptyset$$
$$= \{(0, x) \in S \mid x \in X\},$$

where $S = \{0, 1\} \times (X \cup \emptyset) = \{0, 1\} \times (\emptyset \cup X) = S$.

Invertibility

The inverse of $\rho_X^{\mathsf{Sets}, \coprod}$ is the map

$$\rho_X^{\mathsf{Sets}, \coprod, -1} \colon X \to X \coprod \emptyset$$

given by

$$\rho_X^{\mathsf{Sets},\coprod,-1}(x)\stackrel{\mathrm{def}}{=}(0,x)$$

for each $x \in X$. Indeed:

• Invertibility I. We have

$$\begin{split} \left[\rho_X^{\mathsf{Sets}, \coprod, -1} \circ \rho_X^{\mathsf{Sets}, \coprod} \right] (0, x) &= \rho_X^{\mathsf{Sets}, \coprod, -1} \Big(\rho_X^{\mathsf{Sets}, \coprod} (0, x) \Big) \\ &= \rho_X^{\mathsf{Sets}, \coprod, -1} (x) \\ &= (0, x) \\ &= \left[\mathrm{id}_{X \coprod \emptyset} \right] (0, x) \end{split}$$

for each $(0, x) \in \emptyset \coprod X$, and therefore we have

$$\rho_X^{\mathsf{Sets}, \coprod, -1} \circ \rho_X^{\mathsf{Sets}, \coprod} = \mathrm{id}_{\emptyset \coprod X}.$$

• Invertibility II. We have

$$\begin{split} \left[\rho_X^{\mathsf{Sets}, \coprod} \circ \rho_X^{\mathsf{Sets}, \coprod, -1} \right] (x) &= \rho_X^{\mathsf{Sets}, \coprod} \left(\rho_X^{\mathsf{Sets}, \coprod, -1} (x) \right) \\ &= \rho_X^{\mathsf{Sets}, \coprod, -1} (0, x) \\ &= x \\ &= [\mathrm{id}_X] (x) \end{split}$$

for each $x \in X$, and therefore we have

$$\rho_X^{\mathsf{Sets}, \coprod} \circ \rho_X^{\mathsf{Sets}, \coprod, -1} = \mathrm{id}_X.$$

Therefore $\rho_X^{\mathsf{Sets}, \coprod}$ is indeed an isomorphism.

Naturality

We need to show that, given a function $f: X \to Y$, the diagram

$$\begin{array}{cccc} X \coprod \varnothing & \xrightarrow{f \coprod \mathrm{id}_{\varnothing}} & Y \coprod \varnothing \\ & & & & \downarrow \rho_{Y}^{\mathsf{Sets}, \coprod} & & & \downarrow \rho_{Y}^{\mathsf{Sets}, \coprod} \\ & & & X & \xrightarrow{f} & Y \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$(0,x) \qquad (0,x) \longmapsto (1,f(x))$$

$$\downarrow \qquad \qquad \qquad \downarrow$$

$$x \longmapsto f(x) \qquad f(x)$$

and hence indeed commutes. Therefore $\rho^{\mathsf{Sets}, \coprod}$ is a natural transformation.

Being a Natural Isomorphism

Since $\rho^{\mathsf{Sets}, \coprod}$ is natural and $\rho^{\mathsf{Sets}, -1}$ is a componentwise inverse to $\rho^{\mathsf{Sets}, \coprod}$, it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that $\rho^{\mathsf{Sets}, -1}$ is also natural. Thus $\rho^{\mathsf{Sets}, \coprod}$ is a natural isomorphism.

5.2.6 The Symmetry

DEFINITION 5.2.6.1.1 ► THE SYMMETRY OF []

The symmetry of the coproduct of sets is the natural isomorphism

$$\sigma^{\mathsf{Sets}, \coprod} : \coprod \overset{\sim}{\Longrightarrow} \coprod \circ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets}, \mathsf{Sets}}, \qquad \begin{array}{c} \mathsf{Sets} \times \mathsf{Sets} & \overset{\coprod}{\longleftrightarrow} \mathsf{Sets}, \\ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets}, \mathsf{Sets}} & & \downarrow & \downarrow \\ \mathsf{Sets} \times \mathsf{Sets} & & \mathsf{Sets} \end{array}$$

whose component

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod}:X\coprod Y\stackrel{\sim}{\dashrightarrow} Y\coprod X$$

at $X, Y \in Obj(Sets)$ is defined by

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod}(x,y)\stackrel{\mathrm{def}}{=}(y,x)$$

for each $(x, y) \in X \times Y$.

PROOF 5.2.6.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.2.6.1.1

Unwinding the Definitions of $X \coprod Y$ and $Y \coprod X$

Firstly, we unwind the expressions for $X \coprod Y$ and $Y \coprod X$. We have

$$X \coprod Y \stackrel{\text{def}}{=} \{(0, x) \in S \mid x \in X\} \cup \{(1, y) \in S \mid y \in Y\},\$$

where $S = \{0, 1\} \times (X \cup Y)$ and

$$Y \coprod X \stackrel{\text{def}}{=} \{(0, y) \in S' \mid y \in Y\} \cup \{(1, x) \in S' \mid x \in X\},\$$

where
$$S' = \{0, 1\} \times (Y \cup X) = \{0, 1\} \times (X \cup Y) = S$$
.

Invertibility

The inverse of $\sigma_{X,Y}^{\mathsf{Sets},\coprod}$ is the map

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1} \colon Y \coprod X \to X \coprod Y$$

defined by

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1} \stackrel{\mathrm{def}}{=} \sigma_{Y,X}^{\mathsf{Sets},\coprod}$$

and hence given by

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1}(z) \stackrel{\text{def}}{=} \begin{cases} (0,x) & \text{if } z = (1,x), \\ (1,y) & \text{if } z = (0,y) \end{cases}$$

for each $z \in Y \mid X$. Indeed:

• Invertibility I. We have

$$\begin{split} \left[\sigma_{X,Y}^{\mathsf{Sets}, \coprod, -1} \circ \sigma_{X,Y}^{\mathsf{Sets}, \coprod}\right] (0, x) &= \sigma_{X}^{\mathsf{Sets}, \coprod, -1} \left(\sigma_{X}^{\mathsf{Sets}, \coprod} (0, x)\right) \\ &= \sigma_{X}^{\mathsf{Sets}, \coprod, -1} (1, x) \\ &= (0, x) \\ &= \left[\mathrm{id}_{X \coprod Y}\right] (0, x) \end{split}$$

for each $(0, x) \in X \coprod Y$ and

$$\left[\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1} \circ \sigma_{X,Y}^{\mathsf{Sets},\coprod}\right](1,y) = \sigma_{X}^{\mathsf{Sets},\coprod,-1} \left(\sigma_{X}^{\mathsf{Sets},\coprod}(1,y)\right)$$

$$= \sigma_X^{\mathsf{Sets}, \coprod, -1}(0, y)$$

$$= (1, y)$$

$$= \left[\mathsf{id}_{X \coprod Y} \right] (1, y)$$

for each $(1, y) \in X \coprod Y$, and therefore we have

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1} \circ \sigma_{X,Y}^{\mathsf{Sets},\coprod} = \mathrm{id}_{X\coprod Y} \,.$$

• Invertibility II. We have

$$\begin{split} \left[\sigma_{X,Y}^{\mathsf{Sets}, \coprod} \circ \sigma_{X,Y}^{\mathsf{Sets}, \coprod, -1}\right] (0, y) &= \sigma_{X}^{\mathsf{Sets}, \coprod} \left(\sigma_{X}^{\mathsf{Sets}, \coprod, -1} (0, y)\right) \\ &= \sigma_{X}^{\mathsf{Sets}, \coprod, -1} (1, y) \\ &= (0, y) \\ &= \left[\mathrm{id}_{Y \coprod X}\right] (0, y) \end{split}$$

for each $(0, y) \in Y \coprod X$ and

$$\begin{split} \left[\sigma_{X,Y}^{\mathsf{Sets}, \coprod} \circ \sigma_{X,Y}^{\mathsf{Sets}, \coprod, -1}\right] (1, x) &= \sigma_{X}^{\mathsf{Sets}, \coprod} \left(\sigma_{X}^{\mathsf{Sets}, \coprod, -1} (1, x)\right) \\ &= \sigma_{X}^{\mathsf{Sets}, \coprod, -1} (0, x) \\ &= (1, x) \\ &= \left[\mathrm{id}_{Y \coprod X}\right] (1, x) \end{split}$$

for each $(1, x) \in Y \coprod X$, and therefore we have

$$\sigma_X^{\mathsf{Sets}, \coprod} \circ \sigma_X^{\mathsf{Sets}, \coprod, -1} = \mathrm{id}_{Y \coprod X} \,.$$

Therefore $\sigma_{X,Y}^{\mathsf{Sets},\coprod}$ is indeed an isomorphism.

Naturality

We need to show that, given functions $f: A \to X$ and $g: B \to Y$, the diagram

$$A \coprod B \xrightarrow{f \coprod g} X \coprod Y$$

$$\sigma_{A,B}^{\mathsf{Sets}, \coprod} \downarrow \qquad \qquad \qquad \downarrow \sigma_{X,Y}^{\mathsf{Sets}, \coprod}$$

$$B \coprod A \xrightarrow{g \coprod f} Y \coprod X$$

commutes. Indeed, this diagram acts on elements as

$$(0,a) \qquad (0,a) \longmapsto (0,f(a))$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow$$

and hence indeed commutes. Therefore $\sigma^{Sets, \coprod}$ is a natural transformation.

Being a Natural Isomorphism

Since $\sigma^{Sets, \coprod}$ is natural and $\sigma^{Sets, -1}$ is a componentwise inverse to $\sigma^{Sets, \coprod}$, it follows from Categories, Item 2 of Proposition 11.9.7.1.2 that $\sigma^{Sets, -1}$ is also natural. Thus $\sigma^{Sets, \coprod}$ is a natural isomorphism.

5.2.7 The Monoidal Category of Sets and Coproducts

PROPOSITION 5.2.7.1.1 ➤ THE MONOIDAL STRUCTURE ON SETS ASSOCIATED TO []

The category Sets admits a closed symmetric monoidal category structure consisting of:

- *The Underlying Category*. The category Sets of pointed sets.
- The Monoidal Product. The coproduct functor

II: Sets
$$\times$$
 Sets \rightarrow Sets

of Constructions With Sets, Item 1 of Proposition 4.2.3.1.4.

• *The Monoidal Unit*. The functor

$$\mathbb{0}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

of Definition 5.2.2.1.1.

• The Associators. The natural isomorphism

$$\alpha^{\mathsf{Sets}, \coprod} : \coprod \circ (\coprod \times \mathsf{id}_{\mathsf{Sets}}) \stackrel{^{\sim}}{\Longrightarrow} \coprod \circ (\mathsf{id}_{\mathsf{Sets}} \times \coprod) \circ \pmb{\alpha}_{\mathsf{Sets}, \mathsf{Sets}, \mathsf{Sets}}^{\mathsf{Cats}}$$

of Definition 5.2.3.1.1.

• The Left Unitors. The natural isomorphism

$$\lambda^{\mathsf{Sets}, \coprod} : \coprod \circ \left(\mathbb{O}^{\mathsf{Sets}} \times id_{\mathsf{Sets}} \right) \stackrel{\widetilde{}}{\Longrightarrow} \lambda^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.2.4.1.1.

• The Right Unitors. The natural isomorphism

$$\rho^{\mathsf{Sets}, \coprod} : \coprod \circ \left(\mathsf{id} \times \mathbb{O}^{\mathsf{Sets}}\right) \stackrel{\sim}{\Longrightarrow} \rho_{\mathsf{Sets}}^{\mathsf{Cats}_2}$$

of Definition 5.2.5.1.1.

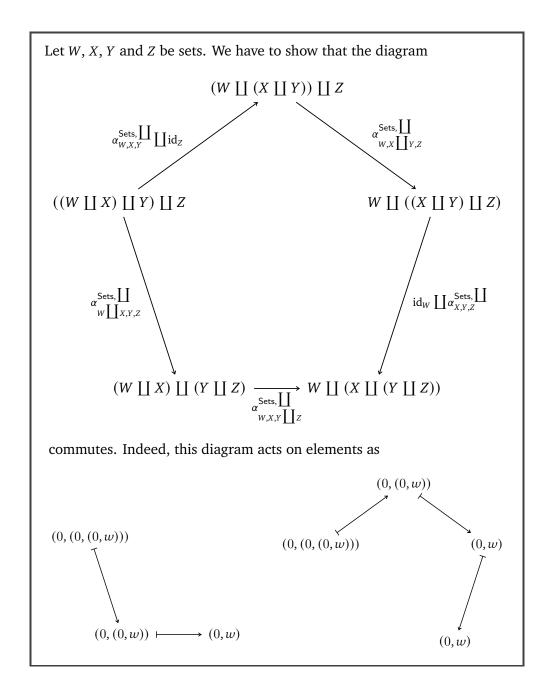
• The Symmetry. The natural isomorphism

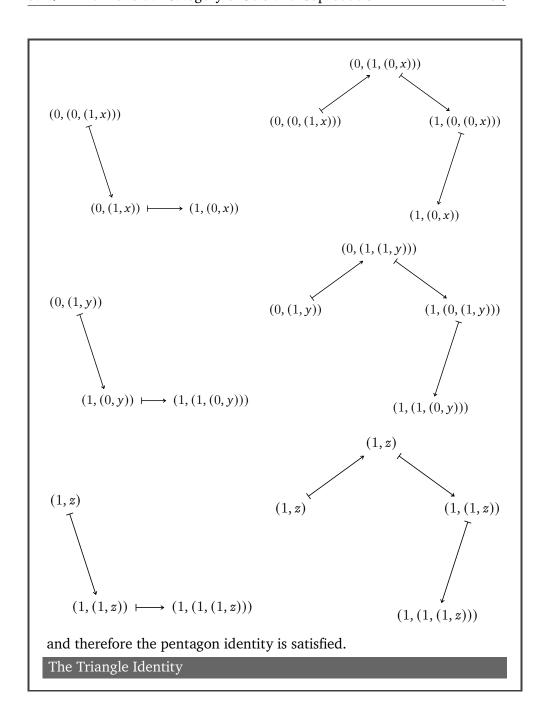
$$\sigma^{\mathsf{Sets}, \coprod} : \times \stackrel{\widetilde{}}{\Longrightarrow} \times \circ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets}, \mathsf{Sets}}$$

of Definition 5.2.6.1.1.

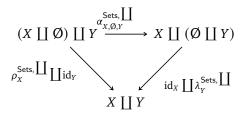
PROOF 5.2.7.1.2 ▶ PROOF OF PROPOSITION 5.2.7.1.1

The Pentagon Identity

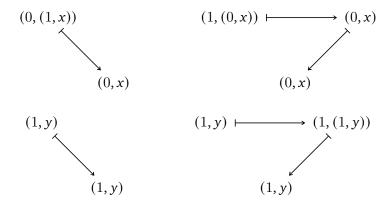




Let *X* and *Y* be sets. We have to show that the diagram



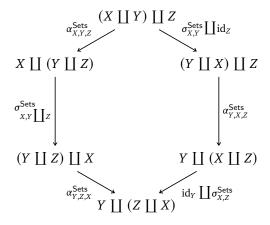
commutes. Indeed, this diagram acts on elements as

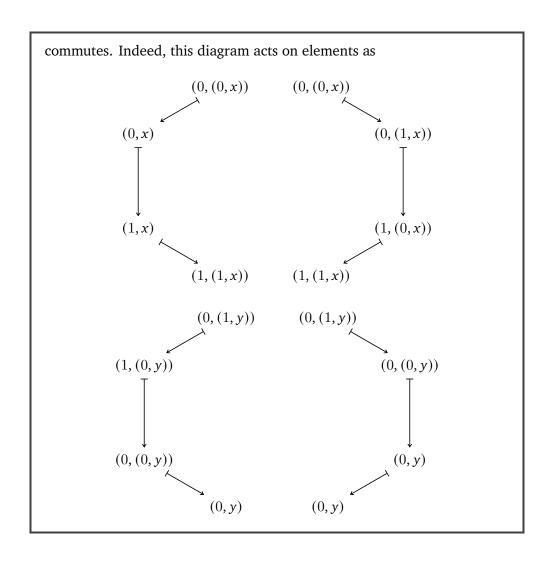


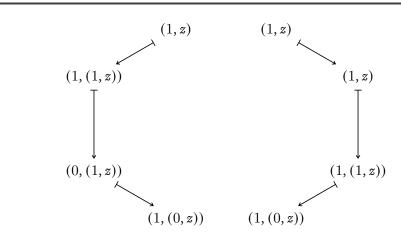
and therefore the triangle identity is satisfied.

The Left Hexagon Identity

Let X, Y, and Z be sets. We have to show that the diagram



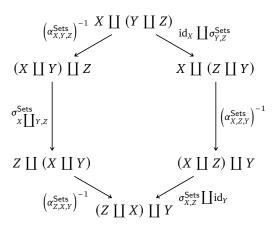


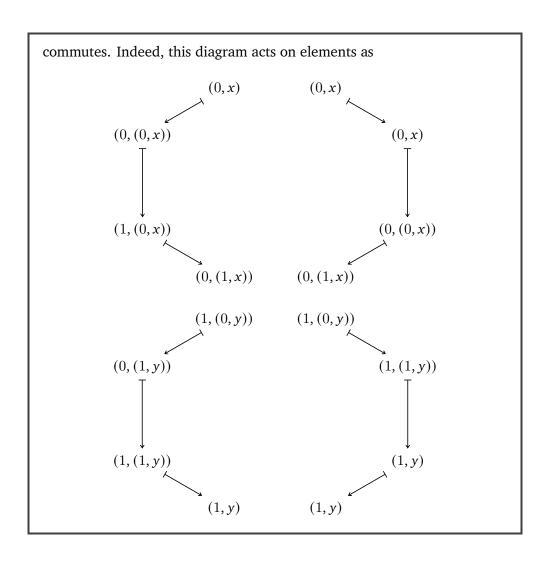


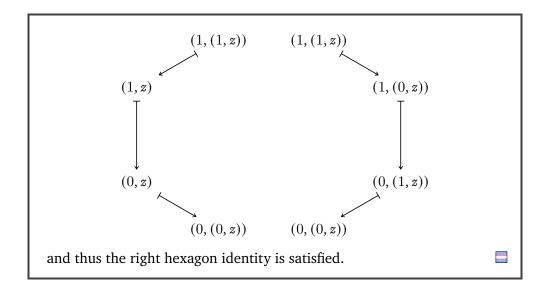
and thus the left hexagon identity is satisfied.

The Right Hexagon Identity

Let *X*, *Y*, and *Z* be sets. We have to show that the diagram







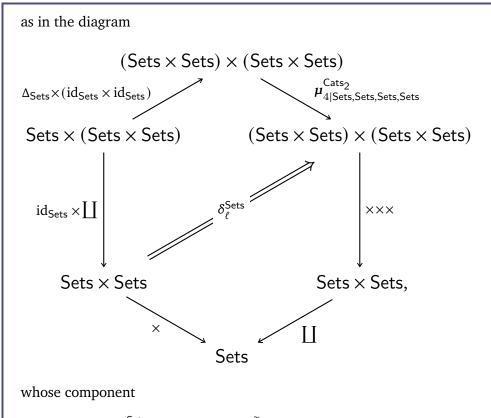
5.3 The Bimonoidal Category of Sets, Products, and Coproducts

5.3.1 The Left Distributor

DEFINITION 5.3.1.1.1 ► THE LEFT DISTRIBUTOR OF × OVER ∐

The left distributor of the product of sets over the coproduct of sets is the natural isomorphism

$$\delta_{\ell}^{\mathsf{Sets}} \colon \times \circ (\mathsf{id}_{\mathsf{Sets}} \times \coprod) \overset{\widetilde{}}{\Longrightarrow} \coprod \circ (\times \times \times) \circ \boldsymbol{\mu}_{4 \mid \mathsf{Sets}, \mathsf{Sets}, \mathsf{Sets}, \mathsf{Sets}}^{\mathsf{Cats}_2} \circ (\Delta_{\mathsf{Sets}} \times (\mathsf{id}_{\mathsf{Sets}} \times \mathsf{id}_{\mathsf{Sets}}))$$



$$\mathcal{S}^{\mathsf{Sets}}_{\ell\mid X,Y,Z} \colon X \times (Y \coprod Z) \stackrel{\sim}{\dashrightarrow} (X \times Y) \coprod (X \times Z)$$

at (X, Y, Z) is defined by

$$\delta^{\mathsf{Sets}}_{\ell|X,Y,Z}(x,a) \stackrel{\text{def}}{=} \begin{cases} (0,(x,y)) & \text{if } a = (0,y), \\ (1,(x,z)) & \text{if } a = (1,z) \end{cases}$$

for each $(x, a) \in X \times (Y \coprod Z)$.

PROOF 5.3.1.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.3.1.1.1

Omitted.

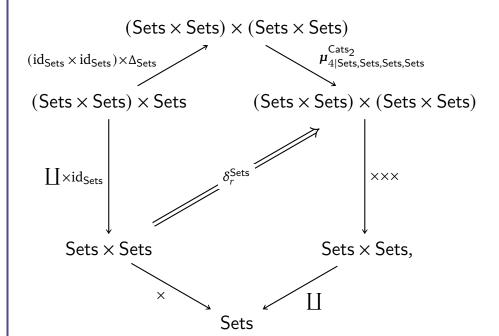
5.3.2 The Right Distributor

DEFINITION 5.3.2.1.1 ► THE RIGHT DISTRIBUTOR OF × OVER ∐

The **right distributor of the product of sets over the coproduct of sets** is the natural isomorphism

$$\delta_r^{\mathsf{Sets}} \colon \times \circ (\coprod \times \mathsf{id}_{\mathsf{Sets}}) \stackrel{\scriptstyle \sim}{\Longrightarrow} \coprod \circ (\times \times \times) \circ \boldsymbol{\mu}_{\mathsf{4}|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}^{\mathsf{Cats}_2} \circ ((\mathsf{id}_{\mathsf{Sets}} \times \mathsf{id}_{\mathsf{Sets}}) \times \Delta_{\mathsf{Sets}})$$

as in the diagram



whose component

$$\mathcal{S}^{\mathsf{Sets}}_{r|X,Y,Z} \colon (X \coprod Y) \times Z \xrightarrow{\sim} (X \times Z) \coprod (Y \times Z)$$

at (X, Y, Z) is defined by

$$\delta^{\mathsf{Sets}}_{r|X,Y,Z}(a,z) \stackrel{\text{def}}{=} \begin{cases} (0,(x,z)) & \text{if } a = (0,x), \\ (1,(y,z)) & \text{if } a = (1,y) \end{cases}$$

for each $(a, z) \in (X \coprod Y) \times Z$.

PROOF 5.3.2.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.3.2.1.1

Omitted.

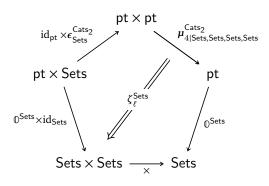
5.3.3 The Left Annihilator

DEFINITION 5.3.3.1.1 ► THE LEFT ANNIHILATOR OF ×

The **left annihilator of the product of sets** is the natural isomorphism

$$\boldsymbol{\zeta}^{\mathsf{Sets}}_{\ell} \colon \mathbb{O}^{\mathsf{Sets}} \circ \boldsymbol{\mu}^{\mathsf{Cats}_2}_{4 \mid \mathsf{Sets}, \mathsf{Sets}, \mathsf{Sets}, \mathsf{Sets}} \circ \left(\mathrm{id}_{\mathsf{pt}} \times \boldsymbol{\epsilon}^{\mathsf{Cats}_2}_{\mathsf{Sets}} \right) \stackrel{\sim}{\Longrightarrow} \times \circ \left(\mathbb{O}^{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}} \right)$$

as in the diagram



with components

$$\zeta_{\ell|A}^{\mathsf{Sets}} \colon \emptyset \times A \xrightarrow{\sim} \emptyset.$$

PROOF 5.3.3.1.2 ▶ PROOF OF THE CLAIMS MADE IN DEFINITION 5.3.3.1.1

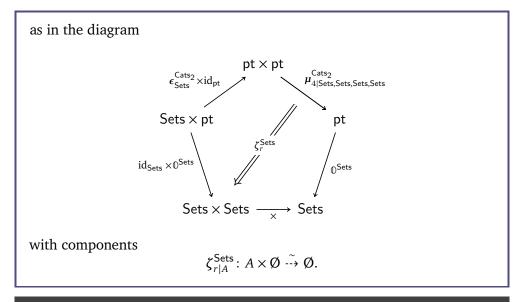
Omitted. For a partial proof, see [Pro25].

5.3.4 The Right Annihilator

DEFINITION 5.3.4.1.1 ► THE RIGHT ANNIHILATOR OF ×

The right annihilator of the product of sets is the natural isomorphism

$$\zeta_{\mathit{r}}^{\mathsf{Sets}} \colon \mathbb{O}^{\mathsf{Sets}} \circ \boldsymbol{\mu}_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}^{\mathsf{Cats}_2} \circ \left(\boldsymbol{\epsilon}_{\mathsf{Sets}}^{\mathsf{Cats}_2} \times \mathrm{id}_{\mathsf{pt}}\right) \overset{\sim}{\dashrightarrow} \times \circ \left(\mathrm{id}_{\mathsf{Sets}} \times \mathbb{O}^{\mathsf{Sets}}\right)$$



PROOF 5.3.4.1.2 ➤ PROOF OF THE CLAIMS MADE IN DEFINITION 5.3.4.1.1

Omitted. For a partial proof, see [Pro25].

5.3.5 The Bimonoidal Category of Sets, Products, and Coproducts

PROPOSITION 5.3.5.1.1 ► THE BIMONOIDAL STRUCTURE ON SETS ASSOCIATED TO × AND \[\]

The category Sets admits a closed symmetric bimonoidal category structure consisting of:

- *The Underlying Category*. The category Sets of pointed sets.
- The Additive Monoidal Product. The coproduct functor

$$II: Sets \times Sets \rightarrow Sets$$

of Constructions With Sets, Item 1 of Proposition 4.2.3.1.4.

• The Multiplicative Monoidal Product. The product functor

$$\times$$
: Sets \times Sets \rightarrow Sets

of Constructions With Sets, Item 1 of Proposition 4.1.3.1.4.

• The Monoidal Unit. The functor

$$\mathbb{1}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

of Definition 5.1.3.1.1.

• The Monoidal Zero. The functor

$$\mathbb{O}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

of Definition 5.1.3.1.1.

• The Internal Hom. The internal Hom functor

Sets:
$$Sets^{op} \times Sets \rightarrow Sets$$

of Constructions With Sets, ?? of ??.

• The Additive Associators. The natural isomorphism

$$\alpha^{\mathsf{Sets}, \coprod} : \coprod \circ (\coprod \times \mathsf{id}_{\mathsf{Sets}}) \stackrel{\widetilde{-}}{\Longrightarrow} \coprod \circ (\mathsf{id}_{\mathsf{Sets}} \times \coprod) \circ \pmb{\alpha}^{\mathsf{Cats}}_{\mathsf{Sets}, \mathsf{Sets}, \mathsf{Sets}}$$

of Definition 5.2.3.1.1.

• The Additive Left Unitors. The natural isomorphism

$$\lambda^{\mathsf{Sets}, \coprod} : \coprod \circ \left(\mathbb{O}^{\mathsf{Sets}} \times id_{\mathsf{Sets}} \right) \stackrel{^{\sim}}{\Longrightarrow} \lambda^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.2.4.1.1.

• The Additive Right Unitors. The natural isomorphism

$$\rho^{\mathsf{Sets},\coprod} : \coprod \circ \left(\mathsf{id} \times \mathbb{O}^{\mathsf{Sets}}\right) \stackrel{\sim}{\Longrightarrow} \rho_{\mathsf{Sets}}^{\mathsf{Cats}_2}$$

of Definition 5.2.5.1.1.

• The Additive Symmetry. The natural isomorphism

$$\sigma^{\mathsf{Sets},\coprod} : \coprod \stackrel{\widetilde{}}{\Longrightarrow} \coprod \circ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}}$$

of Definition 5.2.6.1.1.

• The Multiplicative Associators. The natural isomorphism

$$\alpha^{\mathsf{Sets}} : \times \circ (\times \times \mathsf{id}_{\mathsf{Sets}}) \stackrel{\widetilde{-}}{\Longrightarrow} \times \circ (\mathsf{id}_{\mathsf{Sets}} \times \times) \circ \pmb{\alpha}^{\mathsf{Cats}}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}$$

of Definition 5.1.4.1.1.

• The Multiplicative Left Unitors. The natural isomorphism

$$\lambda^{\mathsf{Sets}} : \times \circ \left(\mathbb{1}^{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}} \right) \stackrel{\sim}{\Longrightarrow} \lambda^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.1.5.1.1.

• The Multiplicative Right Unitors. The natural isomorphism

$$\rho^{\mathsf{Sets}} : \times \circ \left(\mathsf{id} \times \mathbb{1}^{\mathsf{Sets}} \right) \stackrel{\sim}{\Longrightarrow} \rho_{\mathsf{Sets}}^{\mathsf{Cats}_2}$$

of Definition 5.1.6.1.1.

• The Multiplicative Symmetry. The natural isomorphism

$$\sigma^{\mathsf{Sets}} : \times \stackrel{\sim}{\Longrightarrow} \times \circ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}}$$

of Definition 5.1.7.1.1.

• The Left Distributor. The natural isomorphism

$$\delta_{\ell}^{\mathsf{Sets}} : \times \circ (\mathrm{id}_{\mathsf{Sets}} \times \coprod) \overset{\sim}{\Longrightarrow} \coprod \circ (\times \times \times) \circ \mu_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}^{\mathsf{Cats}_2} \circ (\Delta_{\mathsf{Sets}} \times (\mathrm{id}_{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}))$$
of Definition 5.3.1.1.1.

• The Right Distributor. The natural isomorphism

$$\delta_r^{\mathsf{Sets}} : \times \circ (\coprod \times \mathsf{id}_{\mathsf{Sets}}) \xrightarrow{\sim} \coprod \circ (\times \times \times) \circ \boldsymbol{\mu}_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}^{\mathsf{Cats}_2} \circ ((\mathsf{id}_{\mathsf{Sets}} \times \mathsf{id}_{\mathsf{Sets}}) \times \Delta_{\mathsf{Sets}})$$
of Definition 5.3.2.1.1.

• The Left Annihilator. The natural isomorphism

$$\zeta_{\ell}^{\mathsf{Sets}} \colon \mathbb{O}^{\mathsf{Sets}} \circ \boldsymbol{\mu}_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}^{\mathsf{Cats}_2} \circ \left(\mathrm{id}_{\mathsf{pt}} \times \boldsymbol{\epsilon}_{\mathsf{Sets}}^{\mathsf{Cats}_2}\right) \stackrel{\sim}{\Longrightarrow} \times \circ \left(\mathbb{O}^{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}\right)$$
of Definition 5.3.3.1.1.

• The Right Annihilator. The natural isomorphism

$$\zeta_r^{\mathsf{Sets}} \colon \mathbb{O}^{\mathsf{Sets}} \circ \pmb{\mu}_{4 \mid \mathsf{Sets}, \mathsf{Sets}, \mathsf{Sets}, \mathsf{Sets}}^{\mathsf{Cats}_2} \circ \left(\pmb{\epsilon}_{\mathsf{Sets}}^{\mathsf{Cats}_2} \times \mathsf{id}_{\mathsf{pt}}\right) \overset{\sim}{\dashrightarrow} \times \circ \left(\mathsf{id}_{\mathsf{Sets}} \times \mathbb{O}^{\mathsf{Sets}}\right)$$

of Definition 5.3.4.1.1.



Appendices

A Other Chapters

Preliminaries	10. Conditions on Relations
1. Introduction	Categories
2. A Guide to the Literature	11. Categories
Sets	12. Presheaves and the Yoneda
3. Sets	Lemma
4. Constructions With Sets	Monoidal Categories
5. Monoidal Structures on the Category of Sets	13. Constructions With Monoidal Categories
6. Pointed Sets	Bicategories
7. Tensor Products of Pointed Sets	14. Types of Morphisms in Bicate-
Relations	gories
8. Relations	Extra Part

References

9. Constructions With Relations

[Pro25] Proof Wiki Contributors. Cartesian Product Is Empty Iff Factor Is Empty — Proof Wiki. 2025. URL: https://proofwiki.org/wiki/Cartesian_ Product_is_Empty_iff_Factor_is_Empty (cit. on pp. 62, 63).

15. Notes