

Monoidal Structures on the Category of Sets

The Clowder Project Authors

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This chapter contains some material on monoidal structures on Sets.

Contents

5.1 The Monoidal Category of Sets and Products.....	2
5.1.1 Products of Sets	2
5.1.2 The Internal Hom of Sets	2
5.1.3 The Monoidal Unit	2
5.1.4 The Associator	2
5.1.5 The Left Unitor	5
5.1.6 The Right Unitor	7
5.1.7 The Symmetry	9
5.1.8 The Diagonal	11
5.1.9 The Monoidal Category of Sets and Products	14
5.1.10 The Universal Property of $(\text{Sets}, \times, \text{pt})$	18
5.2 The Monoidal Category of Sets and Coproducts	36
5.2.1 Coproducts of Sets	36
5.2.2 The Monoidal Unit	36
5.2.3 The Associator	36
5.2.4 The Left Unitor	39
5.2.5 The Right Unitor	42
5.2.6 The Symmetry	44
5.2.7 The Monoidal Category of Sets and Coproducts	47

5.3	The Bimonoidal Category of Sets, Products, and Coproducts ...	53
5.3.1	The Left Distributor	53
5.3.2	The Right Distributor	56
5.3.3	The Left Annihilator	58
5.3.4	The Right Annihilator	60
5.3.5	The Bimonoidal Category of Sets, Products, and Coproducts	61
A	Other Chapters.....	64

5.1 The Monoidal Category of Sets and Products

5.1.1 Products of Sets

See [Constructions With Sets, Section 4.1.3](#).

5.1.2 The Internal Hom of Sets

See [Constructions With Sets, Section 4.3.5](#).

5.1.3 The Monoidal Unit

Definition 5.1.3.1.1. The **monoidal unit of the product of sets** is the functor

$$1^{\text{Sets}}: \text{pt} \rightarrow \text{Sets}$$

defined by

$$1_{\text{Sets}} \stackrel{\text{def}}{=} \text{pt},$$

where pt is the terminal set of [Constructions With Sets, Definition 4.1.1.1.1](#).

5.1.4 The Associator

Definition 5.1.4.1.1. The **associator of the product of sets** is the natural isomorphism

$$\alpha^{\text{Sets}}: \times \circ (\times \times \text{id}_{\text{Sets}}) \xrightarrow{\sim} \times \circ (\text{id}_{\text{Sets}} \times \times) \circ \alpha_{\text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2},$$

as in the diagram

$$\begin{array}{ccccc}
 & & \text{Sets} \times (\text{Sets} \times \text{Sets}) & & \\
 & \nearrow \alpha_{\text{Sets}, \text{Sets}, \text{Sets}}^{\text{Sets}_2} & & \searrow \text{id} \times \times & \\
 (\text{Sets} \times \text{Sets}) \times \text{Sets} & & & & \text{Sets} \times \text{Sets} \\
 \downarrow \times \times \text{id} & \nearrow \alpha_{\text{Sets}}^{\text{Sets}} & & \searrow \times & \\
 \text{Sets} \times \text{Sets} & \xrightarrow{\times} & \text{Sets} & &
 \end{array}$$

whose component

$$\alpha_{X,Y,Z}^{\text{Sets}}: (X \times Y) \times Z \xrightarrow{\sim} X \times (Y \times Z)$$

at (X, Y, Z) is given by

$$\alpha_{X,Y,Z}^{\text{Sets}}((x, y), z) \stackrel{\text{def}}{=} (x, (y, z))$$

for each $((x, y), z) \in (X \times Y) \times Z$.

Proof. Invertibility: The inverse of $\alpha_{X,Y,Z}^{\text{Sets}}$ is the morphism

$$\alpha_{X,Y,Z}^{\text{Sets}, -1}: X \times (Y \times Z) \xrightarrow{\sim} (X \times Y) \times Z$$

defined by

$$\alpha_{X,Y,Z}^{\text{Sets}, -1}(x, (y, z)) \stackrel{\text{def}}{=} ((x, y), z)$$

for each $(x, (y, z)) \in X \times (Y \times Z)$. Indeed:

- *Invertibility I.* We have

$$\begin{aligned}
 [\alpha_{X,Y,Z}^{\text{Sets}, -1} \circ \alpha_{X,Y,Z}^{\text{Sets}}]((x, y), z) &\stackrel{\text{def}}{=} \alpha_{X,Y,Z}^{\text{Sets}, -1}(\alpha_{X,Y,Z}^{\text{Sets}}((x, y), z)) \\
 &\stackrel{\text{def}}{=} \alpha_{X,Y,Z}^{\text{Sets}, -1}(x, (y, z)) \\
 &\stackrel{\text{def}}{=} ((x, y), z) \\
 &\stackrel{\text{def}}{=} [\text{id}_{(X \times Y) \times Z}]((x, y), z)
 \end{aligned}$$

for each $((x, y), z) \in (X \times Y) \times Z$, and therefore we have

$$\alpha_{X,Y,Z}^{\text{Sets}, -1} \circ \alpha_{X,Y,Z}^{\text{Sets}} = \text{id}_{(X \times Y) \times Z}.$$

- *Invertibility II.* We have

$$\begin{aligned}
 [\alpha_{X,Y,Z}^{\text{Sets}} \circ \alpha_{X,Y,Z}^{\text{Sets},-1}](x, (y, z)) &\stackrel{\text{def}}{=} \alpha_{X,Y,Z}^{\text{Sets}}(\alpha_{X,Y,Z}^{\text{Sets},-1}(x, (y, z))) \\
 &\stackrel{\text{def}}{=} \alpha_{X,Y,Z}^{\text{Sets}}((x, y), z) \\
 &\stackrel{\text{def}}{=} (x, (y, z)) \\
 &\stackrel{\text{def}}{=} [\text{id}_{(X \times Y) \times Z}](x, (y, z))
 \end{aligned}$$

for each $(x, (y, z)) \in X \times (Y \times Z)$, and therefore we have

$$\alpha_{X,Y,Z}^{\text{Sets},-1} \circ \alpha_{X,Y,Z}^{\text{Sets}} = \text{id}_{X \times (Y \times Z)}.$$

Therefore $\alpha_{X,Y,Z}^{\text{Sets}}$ is indeed an isomorphism.

Naturality: We need to show that, given functions

$$\begin{aligned}
 f &: X \rightarrow X', \\
 g &: Y \rightarrow Y', \\
 h &: Z \rightarrow Z'
 \end{aligned}$$

the diagram

$$\begin{array}{ccc}
 (X \times Y) \times Z & \xrightarrow{(f \times g) \times h} & (X' \times Y') \times Z' \\
 \alpha_{X,Y,Z}^{\text{Sets}} \downarrow & & \downarrow \alpha_{X',Y',Z'}^{\text{Sets}} \\
 X \times (Y \times Z) & \xrightarrow{f \times (g \times h)} & X' \times (Y' \times Z')
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
 ((x, y), z) & \longmapsto & ((f(x), g(y)), h(z)) \\
 \downarrow & & \downarrow \\
 (x, (y, z)) & \longmapsto & (f(x), (g(y), h(z)))
 \end{array}$$

and hence indeed commutes, showing α^{Sets} to be a natural transformation.

Being a Natural Isomorphism: Since α^{Sets} is natural and $\alpha^{\text{Sets},-1}$ is a component-wise inverse to α^{Sets} , it follows from [Categories, Item 2](#) of [Definition 11.9.7.1.2](#) that $\alpha^{\text{Sets},-1}$ is also natural. Thus α^{Sets} is a natural isomorphism. \square

5.1.5 The Left Unitor

Definition 5.1.5.1.1. The **left unitor of the product of sets** is the natural isomorphism

$$\lambda^{\text{Sets}}: \times \circ (1^{\text{Sets}} \times \text{id}_{\text{Sets}}) \xRightarrow{\sim} \lambda_{\text{Sets}}^{\text{Cats}_2}$$

whose component

$$\lambda_X^{\text{Sets}}: \text{pt} \times X \xrightarrow{\sim} X$$

at $X \in \text{Obj}(\text{Sets})$ is given by

$$\lambda_X^{\text{Sets}}(\star, x) \stackrel{\text{def}}{=} x$$

for each $(\star, x) \in \text{pt} \times X$.

Proof. Invertibility: The inverse of λ_X^{Sets} is the morphism

$$\lambda_X^{\text{Sets}, -1}: X \xrightarrow{\sim} \text{pt} \times X$$

defined by

$$\lambda_X^{\text{Sets}, -1}(x) \stackrel{\text{def}}{=} (\star, x)$$

for each $x \in X$. Indeed:

- *Invertibility I.* We have

$$\begin{aligned} [\lambda_X^{\text{Sets}, -1} \circ \lambda_X^{\text{Sets}}](\text{pt}, x) &= \lambda_X^{\text{Sets}, -1}(\lambda_X^{\text{Sets}}(\text{pt}, x)) \\ &= \lambda_X^{\text{Sets}, -1}(x) \\ &= (\text{pt}, x) \\ &= [\text{id}_{\text{pt} \times X}](\text{pt}, x) \end{aligned}$$

for each $(\text{pt}, x) \in \text{pt} \times X$, and therefore we have

$$\lambda_X^{\text{Sets}, -1} \circ \lambda_X^{\text{Sets}} = \text{id}_{\text{pt} \times X}.$$

- *Invertibility II.* We have

$$\begin{aligned}
 [\lambda_X^{\text{Sets}} \circ \lambda_X^{\text{Sets}, -1}](x) &= \lambda_X^{\text{Sets}}(\lambda_X^{\text{Sets}, -1}(x)) \\
 &= \lambda_X^{\text{Sets}, -1}(\text{pt}, x) \\
 &= x \\
 &= [\text{id}_X](x)
 \end{aligned}$$

for each $x \in X$, and therefore we have

$$\lambda_X^{\text{Sets}} \circ \lambda_X^{\text{Sets}, -1} = \text{id}_X.$$

Therefore λ_X^{Sets} is indeed an isomorphism.

Naturality: We need to show that, given a function $f: X \rightarrow Y$, the diagram

$$\begin{array}{ccc}
 \text{pt} \times X & \xrightarrow{\text{id}_{\text{pt}} \times f} & \text{pt} \times Y \\
 \lambda_X^{\text{Sets}} \downarrow & & \downarrow \lambda_Y^{\text{Sets}} \\
 X & \xrightarrow{f} & Y
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
 (\star, x) & & (\star, x) \mapsto (\star, f(x)) \\
 \downarrow & & \downarrow \\
 x & \xrightarrow{\quad} & f(x)
 \end{array}$$

and hence indeed commutes. Therefore λ^{Sets} is a natural transformation.

Being a Natural Isomorphism: Since λ^{Sets} is natural and $\lambda^{\text{Sets}, -1}$ is a component-wise inverse to λ^{Sets} , it follows from **Categories, Item 2** of **Definition 11.9.7.1.2** that $\lambda^{\text{Sets}, -1}$ is also natural. Thus λ^{Sets} is a natural isomorphism. \square

5.1.6 The Right Unitor

Definition 5.1.6.1.1. The **right unitor of the product of sets** is the natural isomorphism

$$\rho^{\text{Sets}} : \times \circ (\text{id} \times \mathbb{1}^{\text{Sets}}) \xRightarrow{\sim} \rho_{\text{Sets}}^{\text{Cats}_2},$$

whose component

$$\rho_X^{\text{Sets}} : X \times \text{pt} \dashrightarrow X$$

at $X \in \text{Obj}(\text{Sets})$ is given by

$$\rho_X^{\text{Sets}}(x, \star) \stackrel{\text{def}}{=} x$$

for each $(x, \star) \in X \times \text{pt}$.

Proof. Invertibility: The inverse of ρ_X^{Sets} is the morphism

$$\rho_X^{\text{Sets}, -1} : X \dashrightarrow X \times \text{pt}$$

defined by

$$\rho_X^{\text{Sets}, -1}(x) \stackrel{\text{def}}{=} (x, \star)$$

for each $x \in X$. Indeed:

- *Invertibility I.* We have

$$\begin{aligned} [\rho_X^{\text{Sets}, -1} \circ \rho_X^{\text{Sets}}](x, \star) &= \rho_X^{\text{Sets}, -1}(\rho_X^{\text{Sets}}(x, \star)) \\ &= \rho_X^{\text{Sets}, -1}(x) \\ &= (x, \star) \\ &= [\text{id}_{X \times \text{pt}}](x, \star) \end{aligned}$$

for each $(x, \star) \in X \times \text{pt}$, and therefore we have

$$\rho_X^{\text{Sets}, -1} \circ \rho_X^{\text{Sets}} = \text{id}_{X \times \text{pt}}.$$

- *Invertibility II.* We have

$$\begin{aligned}
 [\rho_X^{\text{Sets}} \circ \rho_X^{\text{Sets}, -1}](x) &= \rho_X^{\text{Sets}}(\rho_X^{\text{Sets}, -1}(x)) \\
 &= \rho_X^{\text{Sets}, -1}(x, \star) \\
 &= x \\
 &= [\text{id}_X](x)
 \end{aligned}$$

for each $x \in X$, and therefore we have

$$\rho_X^{\text{Sets}} \circ \rho_X^{\text{Sets}, -1} = \text{id}_X.$$

Therefore ρ_X^{Sets} is indeed an isomorphism.

Naturality: We need to show that, given a function $f: X \rightarrow Y$, the diagram

$$\begin{array}{ccc}
 X \times \text{pt} & \xrightarrow{f \times \text{id}_{\text{pt}}} & Y \times \text{pt} \\
 \rho_X^{\text{Sets}} \downarrow & & \downarrow \rho_Y^{\text{Sets}} \\
 X & \xrightarrow{f} & Y
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
 (x, \star) & & (x, \star) \longmapsto (f(x), \star) \\
 \downarrow & & \downarrow \\
 x & \longmapsto & f(x)
 \end{array}$$

and hence indeed commutes. Therefore ρ^{Sets} is a natural transformation.

Being a Natural Isomorphism: Since ρ^{Sets} is natural and $\rho^{\text{Sets}, -1}$ is a component-wise inverse to ρ^{Sets} , it follows from **Categories, Item 2** of **Definition 11.9.7.1.2** that $\rho^{\text{Sets}, -1}$ is also natural. Thus ρ^{Sets} is a natural isomorphism. \square

5.1.7 The Symmetry

Definition 5.1.7.1.1. The **symmetry of the product of sets** is the natural isomorphism

$$\sigma^{\text{Sets}} : \times \xrightarrow{\sim} \times \circ \sigma_{\text{Sets}, \text{Sets}}^{\text{Cats}_2},$$

$$\begin{array}{ccc} \text{Sets} \times \text{Sets} & \xrightarrow{\times} & \text{Sets} \\ \sigma_{\text{Sets}, \text{Sets}}^{\text{Cats}_2} \searrow & \Downarrow \sigma^{\text{Sets}} & \nearrow \times \\ & \text{Sets} \times \text{Sets} & \end{array}$$

whose component

$$\sigma_{X,Y}^{\text{Sets}} : X \times Y \xrightarrow{\sim} Y \times X$$

at $X, Y \in \text{Obj}(\text{Sets})$ is defined by

$$\sigma_{X,Y}^{\text{Sets}}(x, y) \stackrel{\text{def}}{=} (y, x)$$

for each $(x, y) \in X \times Y$.

Proof. Invertibility: The inverse of $\sigma_{X,Y}^{\text{Sets}}$ is the morphism

$$\sigma_{X,Y}^{\text{Sets}, -1} : Y \times X \xrightarrow{\sim} X \times Y$$

defined by

$$\sigma_{X,Y}^{\text{Sets}, -1}(y, x) \stackrel{\text{def}}{=} (x, y)$$

for each $(y, x) \in Y \times X$. Indeed:

- *Invertibility I.* We have

$$\begin{aligned} [\sigma_{X,Y}^{\text{Sets}, -1} \circ \sigma_{X,Y}^{\text{Sets}}](x, y) &\stackrel{\text{def}}{=} \sigma_{X,Y}^{\text{Sets}, -1}(\sigma_{X,Y}^{\text{Sets}}(x, y)) \\ &\stackrel{\text{def}}{=} \sigma_{X,Y}^{\text{Sets}, -1}(y, x) \\ &\stackrel{\text{def}}{=} (x, y) \\ &\stackrel{\text{def}}{=} [\text{id}_{X \times Y}](x, y) \end{aligned}$$

for each $(x, y) \in X \times Y$, and therefore we have

$$\sigma_{X,Y}^{\text{Sets}, -1} \circ \sigma_{X,Y}^{\text{Sets}} = \text{id}_{X \times Y}.$$

- *Invertibility II.* We have

$$\begin{aligned}
 [\sigma_{X,Y}^{\text{Sets}} \circ \sigma_{X,Y}^{\text{Sets},-1}](y, x) &\stackrel{\text{def}}{=} \sigma_{X,Y}^{\text{Sets},-1}(\sigma_{X,Y}^{\text{Sets}}(y, x)) \\
 &\stackrel{\text{def}}{=} \sigma_{X,Y}^{\text{Sets},-1}(x, y) \\
 &\stackrel{\text{def}}{=} (y, x) \\
 &\stackrel{\text{def}}{=} [\text{id}_{Y \times X}](y, x)
 \end{aligned}$$

for each $(y, x) \in Y \times X$, and therefore we have

$$\sigma_{X,Y}^{\text{Sets}} \circ \sigma_{X,Y}^{\text{Sets},-1} = \text{id}_{Y \times X}.$$

Therefore $\sigma_{X,Y}^{\text{Sets}}$ is indeed an isomorphism.

Naturality: We need to show that, given functions

$$\begin{aligned}
 f &: X \rightarrow A, \\
 g &: Y \rightarrow B
 \end{aligned}$$

the diagram

$$\begin{array}{ccc}
 X \times Y & \xrightarrow{f \times g} & A \times B \\
 \sigma_{X,Y}^{\text{Sets}} \downarrow & & \downarrow \sigma_{A,B}^{\text{Sets}} \\
 Y \times X & \xrightarrow{g \times f} & B \times A
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
 (x, y) & & (x, y) \longmapsto (f(x), g(y)) \\
 \downarrow & & \downarrow \\
 (y, x) \longmapsto (g(y), f(x)) & & (g(y), f(x))
 \end{array}$$

and hence indeed commutes, showing σ^{Sets} to be a natural transformation.

Being a Natural Isomorphism: Since σ^{Sets} is natural and $\sigma^{\text{Sets},-1}$ is a component-wise inverse to σ^{Sets} , it follows from [Categories, Item 2 of Definition 11.9.7.1.2](#) that $\sigma^{\text{Sets},-1}$ is also natural. Thus σ^{Sets} is a natural isomorphism. \square

5.1.8 The Diagonal

Definition 5.1.8.1.1. The **diagonal of the product of sets** is the natural transformation

$$\Delta: \text{id}_{\text{Sets}} \Rightarrow \times \circ \Delta_{\text{Sets}}^{\text{Cats}_2},$$

whose component

$$\Delta_X: X \rightarrow X \times X$$

at $X \in \text{Obj}(\text{Sets})$ is given by

$$\Delta_X(x) \stackrel{\text{def}}{=} (x, x)$$

for each $x \in X$.

Proof. We need to show that, given a function $f: X \rightarrow Y$, the diagram

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \Delta_X \downarrow & & \downarrow \Delta_Y \\ X \times X & \xrightarrow{f \times f} & Y \times Y \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc} x & \xrightarrow{\quad} & f(x) \\ \downarrow & & \downarrow \\ (x, x) & \xrightarrow{\quad} & (f(x), f(x)) \end{array}$$

and hence indeed commutes, showing Δ to be natural. \square

Proposition 5.1.8.1.2. Let X be a set.

1. *Monoidality.* The diagonal map

$$\Delta: \text{id}_{\text{Sets}} \Longrightarrow \times \circ \Delta_{\text{Sets}}^{\text{Cats}_2},$$

is a monoidal natural transformation:

(a) *Compatibility With Strong Monoidality Constraints.* For each $X, Y \in \text{Obj}(\text{Sets})$, the diagram

$$\begin{array}{ccc} X \times Y & \xrightarrow{\Delta_X \times \Delta_Y} & (X \times X) \times (Y \times Y) \\ & \searrow \Delta_{X \times Y} & \downarrow \wr \\ & & (X \times Y) \times (X \times Y) \end{array}$$

commutes.

(b) *Compatibility With Strong Unitality Constraints.* The diagrams

$$\begin{array}{ccc} \text{pt} & \xrightarrow{\Delta_{\text{pt}}} & \text{pt} \times \text{pt} \\ & \searrow & \downarrow \lambda_{\text{pt}}^{\text{Sets}} \\ & & \text{pt} \end{array} \quad \begin{array}{ccc} \text{pt} & \xrightarrow{\Delta_{\text{pt}}} & \text{pt} \times \text{pt} \\ & \searrow & \downarrow \rho_{\text{pt}}^{\text{Sets}} \\ & & \text{pt} \end{array}$$

commute, i.e. we have

$$\begin{aligned} \Delta_{\text{pt}} &= \lambda_{\text{pt}}^{\text{Sets}, -1} \\ &= \rho_{\text{pt}}^{\text{Sets}, -1}, \end{aligned}$$

where we recall that the equalities

$$\begin{aligned} \lambda_{\text{pt}}^{\text{Sets}} &= \rho_{\text{pt}}^{\text{Sets}}, \\ \lambda_{\text{pt}}^{\text{Sets}, -1} &= \rho_{\text{pt}}^{\text{Sets}, -1} \end{aligned}$$

are always true in any monoidal category by Monoidal Categories, ?? of ??.

2. *The Diagonal of the Unit.* The component

$$\Delta_{\text{pt}}: \text{pt} \xrightarrow{\sim} \text{pt} \times \text{pt}$$

of Δ at pt is an isomorphism.

Proof. **Item 1, Monoidality:** We claim that Δ is indeed monoidal:

1. **Item 1a: Compatibility With Strong Monoidality Constraints:** We need to show that the diagram

$$\begin{array}{ccc}
 X \times Y & \xrightarrow{\Delta_X \times \Delta_Y} & (X \times X) \times (Y \times Y) \\
 & \searrow \Delta_{X \times Y} & \downarrow \wr \\
 & & (X \times Y) \times (X \times Y)
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccccc}
 (x, y) & \longmapsto & ((x, x), (y, y)) & & (x, y) \\
 & & \downarrow & & \swarrow \\
 & & ((x, y), (x, y)) & & ((x, y), (x, y))
 \end{array}$$

and hence indeed commutes.

2. **Item 1b: Compatibility With Strong Unitality Constraints:** As shown in the proof of **Definition 5.1.5.1.1**, the inverse of the left unitor of **Sets** with respect to the product at $X \in \text{Obj}(\text{Sets})$ is given by

$$\lambda_X^{\text{Sets}, -1}(x) \stackrel{\text{def}}{=} (\star, x)$$

for each $x \in X$, so when $X = \text{pt}$, we have

$$\lambda_{\text{pt}}^{\text{Sets}, -1}(\star) \stackrel{\text{def}}{=} (\star, \star),$$

and also

$$\Delta_{\text{pt}}^{\text{Sets}}(\star) \stackrel{\text{def}}{=} (\star, \star),$$

so we have $\Delta_{\text{pt}} = \lambda_{\text{pt}}^{\text{Sets}, -1}$.

This finishes the proof.

Item 2, The Diagonal of the Unit: This follows from **Item 1** and the invertibility of the left/right unitor of **Sets** with respect to \times , proved in the proof of **Definition 5.1.5.1.1** for the left unitor or the proof of **Definition 5.1.6.1.1** for the right unitor. \square

5.1.9 The Monoidal Category of Sets and Products

Proposition 5.1.9.1.1. The category **Sets** admits a closed symmetric monoidal category with diagonals structure consisting of:

- *The Underlying Category.* The category **Sets** of pointed sets.
- *The Monoidal Product.* The product functor

$$\times: \mathbf{Sets} \times \mathbf{Sets} \rightarrow \mathbf{Sets}$$

of **Constructions With Sets, Item 1** of **Definition 4.1.3.1.3.**

- *The Internal Hom.* The internal Hom functor

$$\mathbf{Sets}: \mathbf{Sets}^{\text{op}} \times \mathbf{Sets} \rightarrow \mathbf{Sets}$$

of **Constructions With Sets, Item 1** of **Definition 4.3.5.1.2.**

- *The Monoidal Unit.* The functor

$$\mathbb{1}^{\mathbf{Sets}}: \text{pt} \rightarrow \mathbf{Sets}$$

of **Definition 5.1.3.1.1.**

- *The Associators.* The natural isomorphism

$$\alpha^{\mathbf{Sets}}: \times \circ (\times \times \text{id}_{\mathbf{Sets}}) \xrightarrow{\sim} \times \circ (\text{id}_{\mathbf{Sets}} \times \times) \circ \alpha_{\mathbf{Sets}, \mathbf{Sets}, \mathbf{Sets}}^{\mathbf{Cats}}$$

of **Definition 5.1.4.1.1.**

- *The Left Unitors.* The natural isomorphism

$$\lambda^{\mathbf{Sets}}: \times \circ (\mathbb{1}^{\mathbf{Sets}} \times \text{id}_{\mathbf{Sets}}) \xrightarrow{\sim} \lambda_{\mathbf{Sets}}^{\mathbf{Cats}_2}$$

of **Definition 5.1.5.1.1.**

- *The Right Unitors.* The natural isomorphism

$$\rho^{\mathbf{Sets}}: \times \circ (\text{id} \times \mathbb{1}^{\mathbf{Sets}}) \xrightarrow{\sim} \rho_{\mathbf{Sets}}^{\mathbf{Cats}_2}$$

of **Definition 5.1.6.1.1.**

- *The Symmetry.* The natural isomorphism

$$\sigma^{\text{Sets}}: \times \xrightarrow{\sim} \times \circ \sigma_{\text{Sets}, \text{Sets}}^{\text{Cats}_2}$$

of **Definition 5.1.7.1.1.**

- *The Diagonals.* The monoidal natural transformation

$$\Delta: \text{id}_{\text{Sets}} \Rightarrow \times \circ \Delta_{\text{Sets}}^{\text{Cats}_2}$$

of **Definition 5.1.8.1.1.**

Proof. The Pentagon Identity: Let W, X, Y and Z be sets. We have to show that the diagram

$$\begin{array}{ccccc}
 & & (W \times (X \times Y)) \times Z & & \\
 & \nearrow^{\alpha_{W,X,Y}^{\text{Sets}} \times \text{id}_Z} & & \searrow^{\alpha_{W,X \times Y,Z}^{\text{Sets}}} & \\
 ((W \times X) \times Y) \times Z & & & & W \times ((X \times Y) \times Z) \\
 \searrow^{\alpha_{W \times X,Y,Z}^{\text{Sets}}} & & & & \searrow^{\text{id}_W \times \alpha_{X,Y,Z}^{\text{Sets}}} \\
 (W \times X) \times (Y \times Z) & \xrightarrow{\alpha_{W,X,Y \times Z}^{\text{Sets}}} & W \times (X \times (Y \times Z)) & &
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccccc}
 & & ((w, (x, y)), z) & & \\
 & \swarrow & & \searrow & \\
 (((w, x), y), z) & & & & ((w, x), y), z) \quad (w, ((x, y), z)) \\
 \searrow & & & & \searrow \\
 ((w, x), (y, z)) \mapsto (w, (x, (y, z))) & & & & (w, (x, (y, z))),
 \end{array}$$

and thus the pentagon identity is satisfied.

The Triangle Identity: Let X and Y be sets. We have to show that the diagram

$$\begin{array}{ccc}
 (X \times \text{pt}) \times Y & \xrightarrow{\alpha_{X, \text{pt}, Y}^{\text{Sets}}} & X \times (\text{pt} \times Y) \\
 \searrow \rho_X^{\text{Sets} \times \text{id}_Y} & & \swarrow \text{id}_X \times \lambda_Y^{\text{Sets}} \\
 & X \times Y &
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

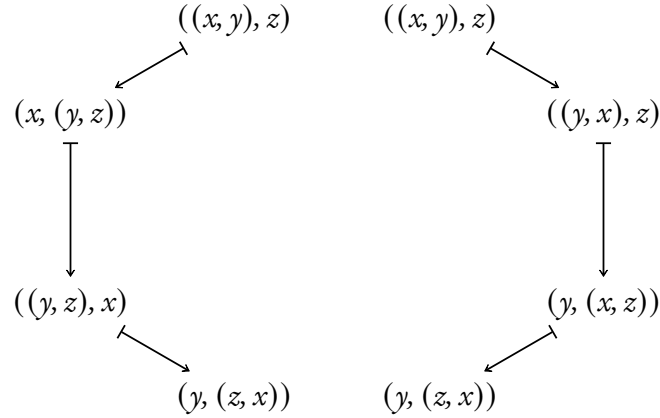
$$\begin{array}{ccc}
 ((x, \star), y) & \xrightarrow{\quad} & (x, (\star, y)) \\
 \searrow & & \swarrow \\
 & (x, y) &
 \end{array}$$

and thus the triangle identity is satisfied.

The Left Hexagon Identity: Let X , Y , and Z be sets. We have to show that the diagram

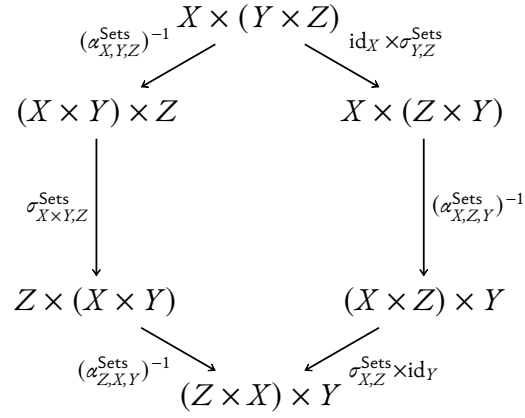
$$\begin{array}{ccc}
 & (X \times Y) \times Z & \\
 \alpha_{X, Y, Z}^{\text{Sets}} \swarrow & & \searrow \sigma_{X, Y}^{\text{Sets} \times \text{id}_Z} \\
 X \times (Y \times Z) & & (Y \times X) \times Z \\
 \downarrow \sigma_{X, Y \times Z}^{\text{Sets}} & & \downarrow \alpha_{Y, X, Z}^{\text{Sets}} \\
 (Y \times Z) \times X & & Y \times (X \times Z) \\
 \alpha_{Y, Z, X}^{\text{Sets}} \swarrow & & \swarrow \text{id}_Y \times \sigma_{X, Z}^{\text{Sets}} \\
 & Y \times (Z \times X) &
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

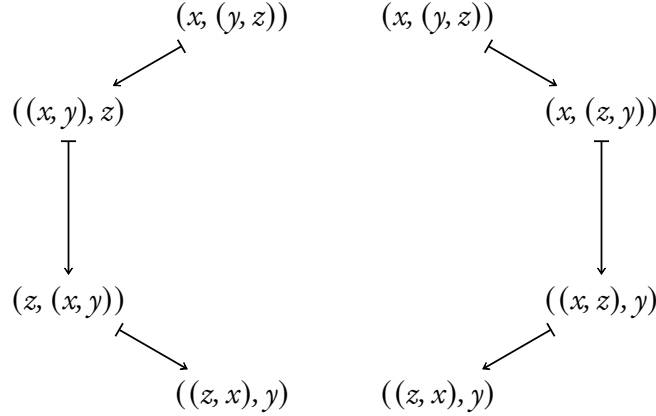


and thus the left hexagon identity is satisfied.

The Right Hexagon Identity: Let X , Y , and Z be sets. We have to show that the diagram



commutes. Indeed, this diagram acts on elements as



and thus the right hexagon identity is satisfied.

Monoidal Closedness: This follows from **Constructions With Sets, Item 2** of **Definition 4.3.5.1.2**

Existence of Monoidal Diagonals: This follows from **Items 1** and **2** of **Definition 5.1.8.1.2**. \square

5.1.10 The Universal Property of $(\mathbf{Sets}, \times, \text{pt})$

Theorem 5.1.10.1.1. The symmetric monoidal structure on the category \mathbf{Sets} of **Definition 5.1.9.1.1** is uniquely determined by the following requirements:

1. *Existence of an Internal Hom.* The tensor product

$$\otimes_{\mathbf{Sets}} : \mathbf{Sets} \times \mathbf{Sets} \rightarrow \mathbf{Sets}$$

of \mathbf{Sets} admits an internal Hom $[-1, -2]_{\mathbf{Sets}}$.

2. *The Unit Object Is pt.* We have $\mathbb{1}_{\mathbf{Sets}} \cong \text{pt}$.

More precisely, the full subcategory of the category $\mathcal{M}_{\mathbf{E}_{\infty}}^{\text{cld}}(\mathbf{Sets})$ of ?? spanned by the closed symmetric monoidal categories $(\mathbf{Sets}, \otimes_{\mathbf{Sets}}, [-1, -2]_{\mathbf{Sets}}, \mathbb{1}_{\mathbf{Sets}}, \lambda^{\mathbf{Sets}}, \rho^{\mathbf{Sets}}, \sigma^{\mathbf{Sets}})$ satisfying **Items 1** and **2** is contractible (i.e. equivalent to the punctual category).

Proof. Unwinding the Statement: Let $(\mathbf{Sets}, \otimes_{\mathbf{Sets}}, [-1, -2]_{\mathbf{Sets}}, \mathbb{1}_{\mathbf{Sets}}, \lambda', \rho', \sigma')$

be a closed symmetric monoidal category satisfying **Items 1** and **2**. We need to show that the identity functor

$$\text{id}_{\mathbf{Sets}} : \mathbf{Sets} \rightarrow \mathbf{Sets}$$

admits a *unique* closed symmetric monoidal functor structure

$$\begin{aligned} \text{id}_{\mathbf{Sets}}^{\otimes} : A \otimes_{\mathbf{Sets}} B &\xrightarrow{\sim} A \times B, \\ \text{id}_{\mathbf{Sets}}^{\text{Hom}} : [A, B]_{\mathbf{Sets}} &\xrightarrow{\sim} \mathbf{Sets}(A, B), \\ \text{id}_{1|\mathbf{Sets}}^{\otimes} : 1_{\mathbf{Sets}} &\xrightarrow{\sim} \text{pt}, \end{aligned}$$

making it into a symmetric monoidal strongly closed isomorphism of categories from $(\mathbf{Sets}, \otimes_{\mathbf{Sets}}, [-1, -2]_{\mathbf{Sets}}, 1_{\mathbf{Sets}}, \lambda', \rho', \sigma')$ to the closed symmetric monoidal category $(\mathbf{Sets}, \times, \mathbf{Sets}(-1, -2), 1_{\mathbf{Sets}}, \lambda^{\mathbf{Sets}}, \rho^{\mathbf{Sets}}, \sigma^{\mathbf{Sets}})$ of **Definition 5.1.9.1.1**. *Constructing an Isomorphism $[-1, -2]_{\mathbf{Sets}} \cong \mathbf{Sets}(-1, -2)$* : By ??, we have a natural isomorphism

$$\mathbf{Sets}(\text{pt}, [-1, -2]_{\mathbf{Sets}}) \cong \mathbf{Sets}(-1, -2).$$

By **Constructions With Sets, Item 3** of **Definition 4.3.5.1.2**, we also have a natural isomorphism

$$\mathbf{Sets}(\text{pt}, [-1, -2]_{\mathbf{Sets}}) \cong [-1, -2]_{\mathbf{Sets}}.$$

Composing both natural isomorphisms, we obtain a natural isomorphism

$$\mathbf{Sets}(-1, -2) \cong [-1, -2]_{\mathbf{Sets}}.$$

Given $A, B \in \text{Obj}(\mathbf{Sets})$, we will write

$$\text{id}_{A,B}^{\text{Hom}} : \mathbf{Sets}(A, B) \xrightarrow{\sim} [A, B]_{\mathbf{Sets}}$$

for the component of this isomorphism at (A, B) .

Constructing an Isomorphism $\otimes_{\mathbf{Sets}} \cong \times$: Since $\otimes_{\mathbf{Sets}}$ is adjoint in each variable to $[-1, -2]_{\mathbf{Sets}}$ by assumption and \times is adjoint in each variable to $\mathbf{Sets}(-1, -2)$ by **Constructions With Sets, Item 2** of **Definition 4.3.5.1.2**, uniqueness of adjoints (??) gives us natural isomorphisms

$$\begin{aligned} A \otimes_{\mathbf{Sets}} - &\cong A \times -, \\ - \otimes_{\mathbf{Sets}} B &\cong B \times -. \end{aligned}$$

By ??, we then have $\otimes_{\mathbf{Sets}} \cong \times$. We will write

$$\text{id}_{\mathbf{Sets}|A,B}^{\otimes}: A \otimes_{\mathbf{Sets}} B \xrightarrow{\sim} A \times B$$

for the component of this isomorphism at (A, B) .

Alternative Construction of an Isomorphism $\otimes_{\mathbf{Sets}} \cong \times$: Alternatively, we may construct a natural isomorphism $\otimes_{\mathbf{Sets}} \cong \times$ as follows:

1. Let $A \in \text{Obj}(\mathbf{Sets})$.
2. Since $\otimes_{\mathbf{Sets}}$ is part of a closed monoidal structure, it preserves colimits in each variable by ??.
3. Since $A \cong \coprod_{a \in A} \text{pt}$ and $\otimes_{\mathbf{Sets}}$ preserves colimits in each variable, we have

$$\begin{aligned} A \otimes_{\mathbf{Sets}} B &\cong \left(\coprod_{a \in A} \text{pt} \right) \otimes_{\mathbf{Sets}} B \\ &\cong \coprod_{a \in A} (\text{pt} \otimes_{\mathbf{Sets}} B) \\ &\cong \coprod_{a \in A} B \\ &\cong A \times B, \end{aligned}$$

naturally in $B \in \text{Obj}(\mathbf{Sets})$, where we have used that pt is the monoidal unit for $\otimes_{\mathbf{Sets}}$. Thus $A \otimes_{\mathbf{Sets}} - \cong A \times -$ for each $A \in \text{Obj}(\mathbf{Sets})$.

4. Similarly, $- \otimes_{\mathbf{Sets}} B \cong - \times B$ for each $B \in \text{Obj}(\mathbf{Sets})$.
5. By ??, we then have $\otimes_{\mathbf{Sets}} \cong \times$.

Below, we'll show that if a natural isomorphism $\otimes_{\mathbf{Sets}} \cong \times$ exists, then it must be unique. This will show that the isomorphism constructed above is equal to the isomorphism $\text{id}_{\mathbf{Sets}|A,B}^{\otimes}: A \otimes_{\mathbf{Sets}} B \rightarrow A \times B$ from before.

Constructing an Isomorphism $\text{id}_1^{\otimes}: \mathbf{1}_{\mathbf{Sets}} \rightarrow \text{pt}$: We define an isomorphism $\text{id}_1^{\otimes}: \mathbf{1}_{\mathbf{Sets}} \rightarrow \text{pt}$ as the composition

$$\mathbf{1}_{\mathbf{Sets}} \xrightarrow[\sim]{\rho_{\mathbf{1}_{\mathbf{Sets}}}^{\mathbf{Sets}, -1}} \mathbf{1}_{\mathbf{Sets}} \times \text{pt} \xrightarrow[\sim]{\text{id}_{\mathbf{Sets}}^{\otimes} \mathbf{1}_{\mathbf{Sets}}} \mathbf{1}_{\mathbf{Sets}} \otimes_{\mathbf{Sets}} \text{pt} \xrightarrow[\sim]{\lambda'_{\text{pt}}} \text{pt}$$

in \mathbf{Sets} .

Monoidal Left Unity of the Isomorphism $\otimes_{\mathbf{Sets}} \cong \times$: We have to show that the diagram

$$\begin{array}{ccc}
 & \text{pt} \otimes_{\mathbf{Sets}} A & \xrightarrow{\text{id}_{\mathbf{Sets}}^{\otimes} |_{\text{pt}, A}} \text{pt} \times A \\
 \text{id}_{1|\mathbf{Sets}}^{\otimes} \otimes_{\mathbf{Sets}} \text{id}_A \nearrow & & \searrow \lambda_A^{\mathbf{Sets}} \\
 1_{\mathbf{Sets}} \otimes_{\mathbf{Sets}} A & \xrightarrow{\lambda'_A} & A
 \end{array}$$

commutes. First, note that the diagram

$$\begin{array}{ccc}
 & \text{pt} \otimes_{\mathbf{Sets}} \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}}^{\otimes} |_{\text{pt}, \text{pt}}} \text{pt} \times \text{pt} \\
 \text{id}_{1|\mathbf{Sets}}^{\otimes} \otimes_{\mathbf{Sets}} \text{id}_{\text{pt}} \nearrow & & \searrow \lambda_{\text{pt}}^{\mathbf{Sets}} \\
 1_{\mathbf{Sets}} \otimes_{\mathbf{Sets}} \text{pt} & \xrightarrow{\lambda'_{\text{pt}}} & \text{pt},
 \end{array}$$

corresponding to the case $A = \text{pt}$, commutes by the terminality of pt (**Constructions With Sets, Definition 4.1.1.2**). Since this diagram commutes, so does the diagram

$$\begin{array}{ccc}
 & \text{pt} \times \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}}^{\otimes, -1} |_{\text{pt}, \text{pt}}} \text{pt} \otimes_{\mathbf{Sets}} \text{pt} \\
 \lambda_{\text{pt}}^{\mathbf{Sets}, -1} \nearrow & & \searrow \text{id}_{1|\mathbf{Sets}}^{\otimes, -1} \otimes_{\mathbf{Sets}} \text{id}_{\text{pt}} \\
 \text{pt} & \xrightarrow{\lambda'_{\text{pt}}{}^{-1}} & 1_{\mathbf{Sets}} \otimes_{\mathbf{Sets}} \text{pt}.
 \end{array}
 \quad (\dagger)$$

Now, let $A \in \text{Obj}(\mathbf{Sets})$, let $a \in A$, and consider the diagram

$$\begin{array}{ccccc}
 & & \text{pt} \times \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},\text{pt}}^{\otimes,-1}} & \text{pt} \otimes_{\mathbf{Sets}} \text{pt} & \xrightarrow{\text{id}_{\mathbf{1}|\mathbf{Sets}}^{\otimes,-1} \times \text{id}_{\text{pt}}} & \mathbf{1}_{\mathbf{Sets}} \otimes_{\mathbf{Sets}} \text{pt} \\
 & \nearrow \lambda_{\text{pt}}^{\mathbf{Sets},-1} & \downarrow & \text{(\dagger)} & \downarrow & \searrow & \downarrow \\
 \text{pt} & \xrightarrow{\lambda_{\text{pt}}'^{-1}} & \text{pt} \times A & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},A}^{\otimes,-1}} & \text{pt} \otimes_{\mathbf{Sets}} A & \xrightarrow{\text{id}_{\mathbf{1}|\mathbf{Sets}}^{\otimes,-1} \times \text{id}_A} & \mathbf{1}_{\mathbf{Sets}} \otimes_{\mathbf{Sets}} A \\
 & \downarrow \text{id}_{\text{pt}} \times [a] & \downarrow \text{id}_{\text{pt}} \times [a] & \text{(1)} & \downarrow \text{id}_{\text{pt}} \otimes_{\mathbf{Sets}} [a] & \downarrow \text{id}_{\mathbf{1}_{\mathbf{Sets}}} \times [a] & \downarrow \\
 & \text{(3)} & \text{(5)} & & \text{(4)} & & \\
 & \downarrow & \downarrow & & \downarrow & & \downarrow \\
 & \text{pt} \times A & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},A}^{\otimes,-1}} & \text{pt} \otimes_{\mathbf{Sets}} A & \xrightarrow{\text{id}_{\mathbf{1}|\mathbf{Sets}}^{\otimes,-1} \times \text{id}_A} & \mathbf{1}_{\mathbf{Sets}} \otimes_{\mathbf{Sets}} A & \\
 & \nearrow \lambda_A^{\mathbf{Sets},-1} & \downarrow & \text{(2)} & \downarrow & \searrow & \downarrow \\
 A & \xrightarrow{\lambda_A'^{-1}} & A & \xrightarrow{\lambda_A'^{-1}} & A & \xrightarrow{\lambda_A'^{-1}} & A
 \end{array}$$

Since:

- Subdiagram (5) commutes by the naturality of λ'^{-1} .
- Subdiagram (†) commutes, as proved above.
- Subdiagram (4) commutes by the naturality of $\text{id}_{\mathbf{1}|\mathbf{Sets}}^{\otimes,-1}$.
- Subdiagram (1) commutes by the naturality of $\text{id}_{\mathbf{Sets}}^{\otimes,-1}$.
- Subdiagram (3) commutes by the naturality of $\lambda^{\mathbf{Sets},-1}$.

it follows that the diagram

$$\begin{array}{ccccc}
 & & \text{pt} \times A & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},A}^{\otimes,-1}} & \text{pt} \otimes_{\mathbf{Sets}} A \\
 & \nearrow \lambda_A^{\mathbf{Sets},-1} & \downarrow & & \downarrow \text{id}_{\mathbf{1}|\mathbf{Sets}}^{\otimes,-1} \otimes_{\mathbf{Sets}} \text{id}_A \\
 \text{pt} & \xrightarrow{[a]} & A & \xrightarrow{\lambda_A'^{-1}} & \mathbf{1}_{\mathbf{Sets}} \otimes_{\mathbf{Sets}} A
 \end{array}$$

Here's a step-by-step showcase of this argument: [\[Link\]](#). We then have

$$\lambda_A'^{-1}(a) = [\lambda_A'^{-1} \circ [a]](\star)$$

$$\begin{aligned}
&= [(\text{id}_{1|\text{Sets}}^{\otimes, -1} \times \text{id}_A) \circ \text{id}_{\text{Sets}|\text{pt}, A}^{\otimes, -1} \circ \lambda_A^{\text{Sets}, -1} \circ [a]](\star) \\
&= [(\text{id}_{1|\text{Sets}}^{\otimes, -1} \times \text{id}_A) \circ \text{id}_{\text{Sets}|\text{pt}, A}^{\otimes, -1} \circ \lambda_A^{\text{Sets}, -1}](a)
\end{aligned}$$

for each $a \in A$, and thus we have

$$\lambda'_A = (\text{id}_{1|\text{Sets}}^{\otimes, -1} \times \text{id}_A) \circ \text{id}_{\text{Sets}|\text{pt}, A}^{\otimes, -1} \circ \lambda_A^{\text{Sets}, -1}.$$

Taking inverses then gives

$$\lambda'_A = \lambda_A^{\text{Sets}} \circ \text{id}_{\text{Sets}|\text{pt}, A}^{\otimes} \circ (\text{id}_{1|\text{Sets}}^{\otimes} \times \text{id}_A),$$

showing that the diagram

$$\begin{array}{ccc}
& \text{pt} \otimes_{\text{Sets}} A & \xrightarrow{\text{id}_{\text{Sets}|\text{pt}, A}^{\otimes}} \text{pt} \times A \\
\text{id}_{1|\text{Sets}}^{\otimes} \otimes_{\text{Sets}} \text{id}_A \nearrow & & \searrow \lambda_A^{\text{Sets}} \\
1_{\text{Sets}} \otimes_{\text{Sets}} A & \xrightarrow{\lambda'_A} & A
\end{array}$$

indeed commutes.

Monoidal Right Unity of the Isomorphism $\otimes_{\text{Sets}} \cong \times$: We can use the same argument we used to prove the monoidal left unity of the isomorphism $\otimes_{\text{Sets}} \cong \times$ above. For completeness, we repeat it below.

We have to show that the diagram

$$\begin{array}{ccc}
& A \otimes_{\text{Sets}} \text{pt} & \xrightarrow{\text{id}_{\text{Sets}|A, \text{pt}}^{\otimes}} A \times \text{pt} \\
\text{id}_A \otimes_{\text{Sets}} \text{id}_{1|\text{Sets}}^{\otimes} \nearrow & & \searrow \rho_A^{\text{Sets}} \\
A \otimes_{\text{Sets}} 1_{\text{Sets}} & \xrightarrow{\rho'_A} & A
\end{array}$$

commutes. First, note that the diagram

$$\begin{array}{ccc}
& \text{pt} \otimes_{\text{Sets}} \text{pt} & \xrightarrow{\text{id}_{\text{Sets}|\text{pt}, \text{pt}}^{\otimes}} \text{pt} \times \text{pt} \\
\text{id}_{\text{pt}} \otimes_{\text{Sets}} \text{id}_{1|\text{Sets}}^{\otimes} \nearrow & & \searrow \rho_{\text{pt}}^{\text{Sets}} \\
\text{pt} \otimes_{\text{Sets}} 1_{\text{Sets}} & \xrightarrow{\rho'_{\text{pt}}} & \text{pt}
\end{array}$$

corresponding to the case $A = \text{pt}$, commutes by the terminality of pt (**Constructions With Sets, Definition 4.1.1.1.2**). Since this diagram commutes, so does the diagram

$$\begin{array}{ccc}
 & \text{pt} \times \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},\text{pt}}^{\otimes,-1}} \text{pt} \otimes_{\mathbf{Sets}} \text{pt} \\
 \nearrow \rho_{\text{pt}}^{\mathbf{Sets},-1} & & \searrow \text{id}_{\text{pt}} \otimes_{\mathbf{Sets}} \text{id}_{1|\mathbf{Sets}}^{\otimes,-1} \\
 \text{pt} & \xrightarrow{\rho_{\text{pt}}'^{-1}} & \text{pt} \otimes_{\mathbf{Sets}} 1_{\mathbf{Sets}}
 \end{array}
 \quad (\dagger)$$

Now, let $A \in \text{Obj}(\mathbf{Sets})$, let $a \in A$, and consider the diagram

$$\begin{array}{ccccc}
 & \text{pt} \times \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},\text{pt}}^{\otimes,-1}} & \text{pt} \otimes_{\mathbf{Sets}} \text{pt} & \\
 \nearrow \rho_{\text{pt}}^{\mathbf{Sets},-1} & & & \searrow \text{id}_{\text{pt}} \times \text{id}_{1|\mathbf{Sets}}^{\otimes,-1} & \\
 \text{pt} & & \xrightarrow{\rho_{\text{pt}}'^{-1}} & & \text{pt} \otimes_{\mathbf{Sets}} 1_{\mathbf{Sets}} \\
 \downarrow [a] & \downarrow \text{id}_{\text{pt}} \times [a] & (1) & \downarrow \text{id}_{\text{pt}} \otimes_{\mathbf{Sets}} [a] & \downarrow \text{id}_{1_{\mathbf{Sets}}} \times [a] \\
 & & (3) & & (4) \\
 & & \downarrow & & \downarrow \\
 & & A \times \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|A,\text{pt}}^{\otimes,-1}} & A \otimes_{\mathbf{Sets}} \text{pt} \\
 \nearrow \rho_A^{\mathbf{Sets},-1} & & & \searrow \text{id}_A \times \text{id}_{1|\mathbf{Sets}}^{\otimes,-1} & \\
 A & & \xrightarrow{\rho_A'^{-1}} & & A \otimes_{\mathbf{Sets}} 1_{\mathbf{Sets}} \\
 & & (2) & &
 \end{array}
 \quad (5)$$

Since:

- Subdiagram (5) commutes by the naturality of ρ'^{-1} .
- Subdiagram (\dagger) commutes, as proved above.
- Subdiagram (4) commutes by the naturality of $\text{id}_{1|\mathbf{Sets}}^{\otimes,-1}$.
- Subdiagram (1) commutes by the naturality of $\text{id}_{\mathbf{Sets}}^{\otimes,-1}$.
- Subdiagram (3) commutes by the naturality of $\rho^{\mathbf{Sets},-1}$.

it follows that the diagram

$$\begin{array}{ccccc}
 & & A \times \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|A, \text{pt}}^{\otimes, -1}} & A \otimes_{\mathbf{Sets}} \text{pt} \\
 & \nearrow \rho_A^{\mathbf{Sets}, -1} & & & \searrow \text{id}_A \otimes_{\mathbf{Sets}} \text{id}_{1|\mathbf{Sets}}^{\otimes, -1} \\
 \text{pt} & \xrightarrow{[a]} & A & \xrightarrow{\rho_A'^{-1}} & A \otimes_{\mathbf{Sets}} 1_{\mathbf{Sets}}
 \end{array}$$

Here's a step-by-step showcase of this argument: [\[Link\]](#). We then have

$$\begin{aligned}
 \rho_A'^{-1}(a) &= [\rho_A'^{-1} \circ [a]](\star) \\
 &= [(\text{id}_A \times \text{id}_{1|\mathbf{Sets}}^{\otimes, -1}) \circ \text{id}_{\mathbf{Sets}|\text{pt}, A}^{\otimes, -1} \circ \rho_A^{\mathbf{Sets}, -1} \circ [a]](\star) \\
 &= [(\text{id}_A \times \text{id}_{1|\mathbf{Sets}}^{\otimes, -1}) \circ \text{id}_{\mathbf{Sets}|\text{pt}, A}^{\otimes, -1} \circ \rho_A^{\mathbf{Sets}, -1}](a)
 \end{aligned}$$

for each $a \in A$, and thus we have

$$\rho_A'^{-1} = (\text{id}_A \times \text{id}_{1|\mathbf{Sets}}^{\otimes, -1}) \circ \text{id}_{\mathbf{Sets}|\text{pt}, A}^{\otimes, -1} \circ \rho_A^{\mathbf{Sets}, -1}.$$

Taking inverses then gives

$$\rho_A' = \rho_A^{\mathbf{Sets}} \circ \text{id}_{\mathbf{Sets}|\text{pt}, A}^{\otimes} \circ (\text{id}_A \times \text{id}_{1|\mathbf{Sets}}^{\otimes}),$$

showing that the diagram

$$\begin{array}{ccccc}
 & & A \otimes_{\mathbf{Sets}} \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|A, \text{pt}}^{\otimes}} & A \times \text{pt} \\
 & \nearrow \text{id}_A \otimes_{\mathbf{Sets}} \text{id}_{1|\mathbf{Sets}}^{\otimes} & & & \searrow \rho_A^{\mathbf{Sets}} \\
 A \otimes_{\mathbf{Sets}} 1_{\mathbf{Sets}} & \xrightarrow{\rho_A'} & & & A
 \end{array}$$

indeed commutes.

Monoidality of the Isomorphism $\otimes_{\text{Sets}} \cong \times$: We have to show that the diagram

$$\begin{array}{ccc}
 & (A \otimes_{\text{Sets}} B) \otimes_{\text{Sets}} C & \\
 \text{id}_{\text{Sets}|A,B}^{\otimes} \otimes_{\text{Sets}} \text{id}_C \swarrow & & \searrow \alpha'_{A,B,C} \\
 (A \times B) \otimes_{\text{Sets}} C & & A \otimes_{\text{Sets}} (B \otimes_{\text{Sets}} C) \\
 \downarrow \text{id}_{\text{Sets}|A \times B, C}^{\otimes} & & \downarrow \text{id}_A \otimes_{\text{Sets}} \text{id}_{\text{Sets}|B,C}^{\otimes} \\
 (A \times B) \times C & & A \otimes_{\text{Sets}} (B \times C) \\
 \searrow \alpha_{A,B,C}^{\text{Sets}} & & \swarrow \text{id}_{\text{Sets}|A, B \times C}^{\otimes} \\
 & A \times (B \times C) &
 \end{array}$$

commutes. First, note that the diagram

$$\begin{array}{ccc}
 & (\text{pt} \otimes_{\text{Sets}} \text{pt}) \otimes_{\text{Sets}} \text{pt} & \\
 \text{id}_{\text{Sets}|\text{pt},\text{pt}}^{\otimes} \otimes_{\text{Sets}} \text{id}_{\text{pt}} \swarrow & & \searrow \alpha'_{\text{pt},\text{pt},\text{pt}} \\
 (\text{pt} \times \text{pt}) \otimes_{\text{Sets}} \text{pt} & & \text{pt} \otimes_{\text{Sets}} (\text{pt} \otimes_{\text{Sets}} \text{pt}) \\
 \downarrow \text{id}_{\text{Sets}|\text{pt} \times \text{pt}, \text{pt}}^{\otimes} & & \downarrow \text{id}_{\text{pt}} \otimes_{\text{Sets}} \text{id}_{\text{Sets}|\text{pt}, \text{pt}}^{\otimes} \\
 (\text{pt} \times \text{pt}) \times \text{pt} & & \text{pt} \otimes_{\text{Sets}} (\text{pt} \times \text{pt}) \\
 \searrow \alpha_{\text{pt},\text{pt},\text{pt}}^{\text{Sets}} & & \swarrow \text{id}_{\text{Sets}|\text{pt}, \text{pt} \times \text{pt}}^{\otimes} \\
 & \text{pt} \times (\text{pt} \times \text{pt}) & \\
 \downarrow !_{\text{pt} \times (\text{pt} \times \text{pt})} & & \\
 & \text{pt} &
 \end{array}$$

commutes by the terminality of pt ([Constructions With Sets, Definition 4.1.1.2](#)). Since the map $!_{\text{pt} \times (\text{pt} \times \text{pt})} : \text{pt} \times (\text{pt} \times \text{pt}) \rightarrow \text{pt}$ is an isomorphism (e.g. having

inverse $\lambda_{\text{pt}}^{\mathbf{Sets}, -1} \circ \lambda_{\text{pt}}^{\mathbf{Sets}, -1}$), it follows that the diagram

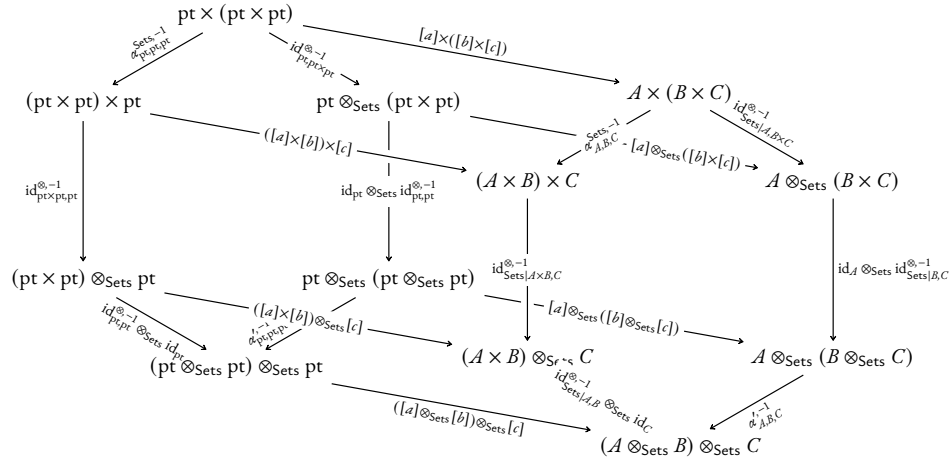
$$\begin{array}{ccc}
 & (\text{pt} \otimes_{\mathbf{Sets}} \text{pt}) \otimes_{\mathbf{Sets}} \text{pt} & \\
 \text{id}_{\mathbf{Sets}|\text{pt}, \text{pt}}^{\otimes} \otimes_{\mathbf{Sets}} \text{id}_{\text{pt}} \swarrow & & \searrow \alpha'_{\text{pt}, \text{pt}, \text{pt}} \\
 (\text{pt} \times \text{pt}) \otimes_{\mathbf{Sets}} \text{pt} & & \text{pt} \otimes_{\mathbf{Sets}} (\text{pt} \otimes_{\mathbf{Sets}} \text{pt}) \\
 \downarrow \text{id}_{\mathbf{Sets}|\text{pt} \times \text{pt}, \text{pt}}^{\otimes} & & \downarrow \text{id}_{\text{pt}} \otimes_{\mathbf{Sets}} \text{id}_{\mathbf{Sets}|\text{pt}, \text{pt}}^{\otimes} \\
 (\text{pt} \times \text{pt}) \times \text{pt} & & \text{pt} \otimes_{\mathbf{Sets}} (\text{pt} \times \text{pt}) \\
 \searrow \alpha_{\text{pt}, \text{pt}, \text{pt}}^{\mathbf{Sets}} & & \swarrow \text{id}_{\mathbf{Sets}|\text{pt}, \text{pt} \times \text{pt}}^{\otimes} \\
 & \text{pt} \times (\text{pt} \times \text{pt}) &
 \end{array}$$

also commutes. Taking inverses, we see that the diagram

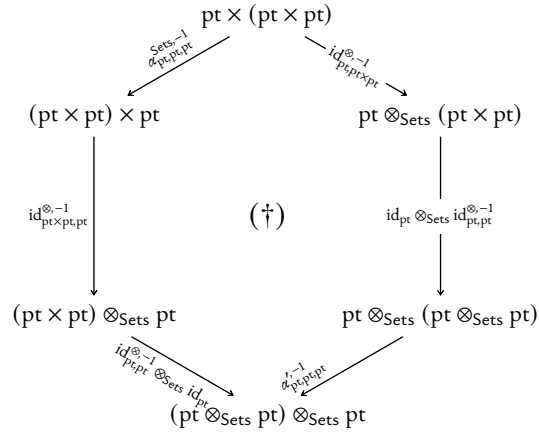
$$\begin{array}{ccc}
 & \text{pt} \times (\text{pt} \times \text{pt}) & \\
 \alpha_{\text{pt}, \text{pt}, \text{pt}}^{\mathbf{Sets}, -1} \swarrow & & \searrow \text{id}_{\mathbf{Sets}|\text{pt}, \text{pt} \times \text{pt}}^{\otimes, -1} \\
 (\text{pt} \times \text{pt}) \times \text{pt} & & \text{pt} \otimes_{\mathbf{Sets}} (\text{pt} \times \text{pt}) \\
 \downarrow \text{id}_{\mathbf{Sets}|\text{pt} \times \text{pt}, \text{pt}}^{\otimes, -1} & (+) & \downarrow \text{id}_{\text{pt}} \otimes_{\mathbf{Sets}} \text{id}_{\mathbf{Sets}|\text{pt}, \text{pt}}^{\otimes, -1} \\
 (\text{pt} \times \text{pt}) \otimes_{\mathbf{Sets}} \text{pt} & & \text{pt} \otimes_{\mathbf{Sets}} (\text{pt} \otimes_{\mathbf{Sets}} \text{pt}) \\
 \text{id}_{\mathbf{Sets}|\text{pt}, \text{pt}}^{\otimes, -1} \otimes_{\mathbf{Sets}} \text{id}_{\text{pt}} \swarrow & & \swarrow \alpha'_{\text{pt}, \text{pt}, \text{pt}}{}^{\otimes, -1} \\
 & (\text{pt} \otimes_{\mathbf{Sets}} \text{pt}) \otimes_{\mathbf{Sets}} \text{pt} &
 \end{array}$$

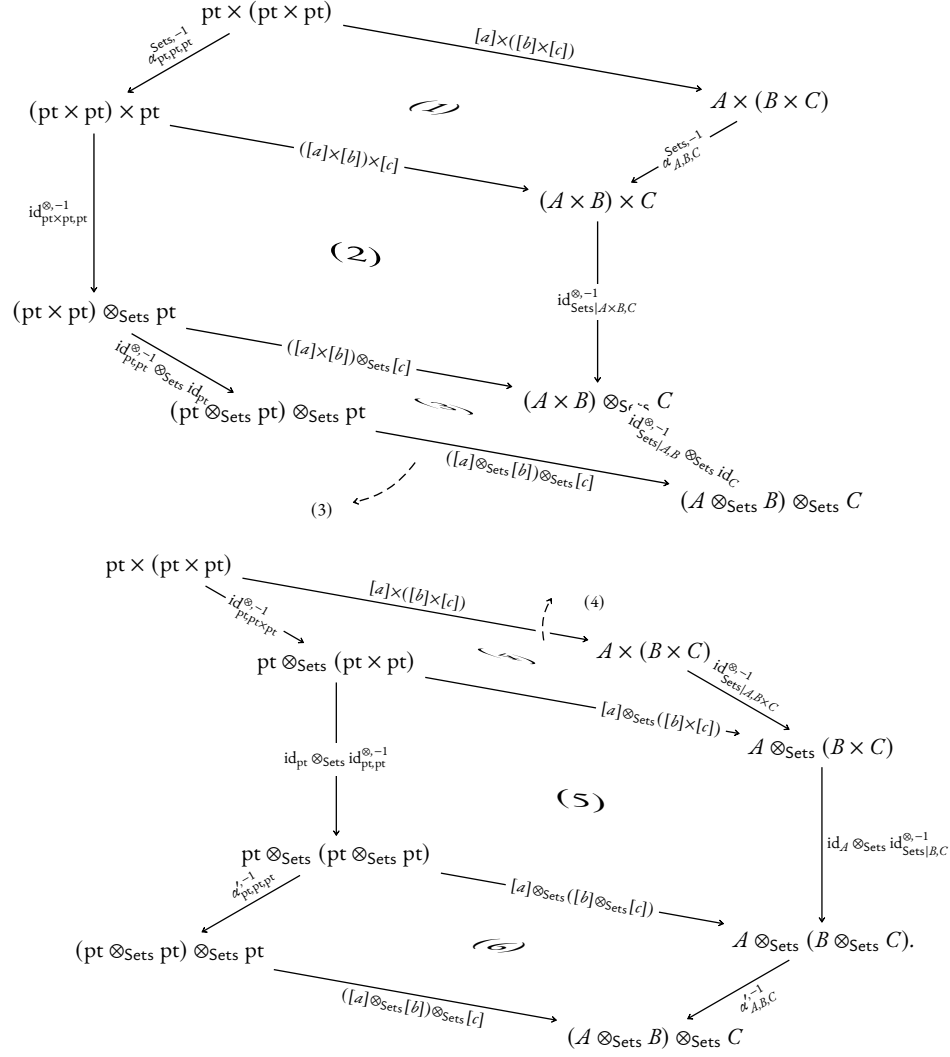
commutes as well. Now, let $A, B, C \in \text{Obj}(\mathbf{Sets})$, let $a \in A$, let $b \in B$, let $c \in C$,

and consider the diagram



which we partition into subdiagrams as follows:





Since:

- Subdiagram (1) commutes by the naturality of $\alpha^{\text{Sets}, -1}$.
- Subdiagram (2) commutes by the naturality of $\text{id}_{\text{Sets}}^{\otimes, -1}$.
- Subdiagram (3) commutes by the naturality of $\text{id}_{\text{Sets}}^{\otimes, -1}$.
- Subdiagram (\dagger) commutes, as proved above.

- Subdiagram (4) commutes by the naturality of $\text{id}_{\text{Sets}}^{\otimes, -1}$.
- Subdiagram (5) commutes by the naturality of $\text{id}_{\text{Sets}}^{\otimes, -1}$.
- Subdiagram (6) commutes by the naturality of α'^{-1} .

it follows that the diagram

$$\begin{array}{ccc}
 & \text{pt} \times (\text{pt} \times \text{pt}) & \\
 & \downarrow [a] \times ([b] \times [c]) & \\
 & A \times (B \times C) & \\
 \alpha_{A,B,C}^{\text{Sets}, -1} \swarrow & & \searrow \text{id}_{\text{Sets}|A, B \times C}^{\otimes, -1} \\
 (A \times B) \times C & & A \otimes_{\text{Sets}} (B \times C) \\
 \downarrow \text{id}_{\text{Sets}|A \times B, C}^{\otimes, -1} & & \downarrow \text{id}_A \times \text{id}_{\text{Sets}|B, C}^{\otimes, -1} \\
 (A \times B) \otimes_{\text{Sets}} C & & A \otimes_{\text{Sets}} (B \otimes_{\text{Sets}} C) \\
 \swarrow \text{id}_{\text{Sets}|A, B}^{\otimes, -1} \otimes_{\text{Sets}} \text{id}_C & & \nwarrow \alpha'_{A,B,C}{}^{-1} \\
 & (A \otimes_{\text{Sets}} B) \otimes_{\text{Sets}} C &
 \end{array}$$

also commutes. We then have

$$\begin{aligned}
 & \left[(\text{id}_{\text{Sets}|A, B}^{\otimes, -1} \otimes_{\text{Sets}} \text{id}_C) \circ \text{id}_{\text{Sets}|A \times B, C}^{\otimes, -1} \right. \\
 & \quad \left. \circ \alpha_{A,B,C}^{\text{Sets}, -1} \right] (a, (b, c)) = \left[(\text{id}_{\text{Sets}|A, B}^{\otimes, -1} \otimes_{\text{Sets}} \text{id}_C) \circ \text{id}_{\text{Sets}|A \times B, C}^{\otimes, -1} \right. \\
 & \quad \left. \circ \alpha_{A,B,C}^{\text{Sets}, -1} \circ ([a] \times ([b] \times [c])) \right] (\star, (\star, \star)) \\
 & = \left[\alpha_{A,B,C}'^{-1} \circ (\text{id}_A \times \text{id}_{\text{Sets}|B, C}^{\otimes, -1}) \right. \\
 & \quad \left. \circ \text{id}_{\text{Sets}|A, B \times C}^{\otimes, -1} \circ ([a] \times ([b] \times [c])) \right] (\star, (\star, \star)) \\
 & = [\alpha_{A,B,C}'^{-1} \circ (\text{id}_A \times \text{id}_{\text{Sets}|B, C}^{\otimes, -1}) \circ \text{id}_{\text{Sets}|A, B \times C}^{\otimes, -1}] (a, (b, c))
 \end{aligned}$$

for each $(a, (b, c)) \in A \times (B \times C)$, and thus we have

$$(\text{id}_{\text{Sets}|A, B}^{\otimes, -1} \otimes_{\text{Sets}} \text{id}_C) \circ \text{id}_{\text{Sets}|A \times B, C}^{\otimes, -1} \circ \alpha_{A,B,C}^{\text{Sets}, -1} = \alpha_{A,B,C}'^{-1} \circ (\text{id}_A \times \text{id}_{\text{Sets}|B, C}^{\otimes, -1}) \circ \text{id}_{\text{Sets}|A, B \times C}^{\otimes, -1}.$$

Taking inverses then gives

$$\alpha_{A,B,C}^{\text{Sets}} \circ \text{id}_{\text{Sets}|A \times B, C}^{\otimes} \circ (\text{id}_{\text{Sets}|A, B}^{\otimes} \otimes_{\text{Sets}} \text{id}_C) = \text{id}_{\text{Sets}|A, B \times C}^{\otimes} \circ (\text{id}_A \times \text{id}_{\text{Sets}|B, C}^{\otimes}) \circ \alpha_{A,B,C}',$$

showing that the diagram

$$\begin{array}{ccc}
 (A \otimes_{\text{Sets}} B) \otimes_{\text{Sets}} C & & \\
 \text{id}_{\text{Sets}|A,B}^{\otimes} \otimes_{\text{Sets}} \text{id}_C \swarrow & & \searrow \alpha'_{A,B,C} \\
 (A \times B) \otimes_{\text{Sets}} C & & A \otimes_{\text{Sets}} (B \otimes_{\text{Sets}} C) \\
 \downarrow \text{id}_{\text{Sets}|A \times B, C}^{\otimes} & & \downarrow \text{id}_A \otimes_{\text{Sets}} \text{id}_{\text{Sets}|B, C}^{\otimes} \\
 (A \times B) \times C & & A \otimes_{\text{Sets}} (B \times C) \\
 \searrow \alpha_{A,B,C}^{\text{Sets}} & & \swarrow \text{id}_{\text{Sets}|A, B \times C}^{\otimes} \\
 & A \times (B \times C) &
 \end{array}$$

indeed commutes.

Braidedness of the Isomorphism $\otimes_{\text{Sets}} \cong \times$: We have to show that the diagram

$$\begin{array}{ccc}
 A \otimes_{\text{Sets}} B & \xrightarrow{\text{id}_{\text{Sets}|A,B}^{\otimes}} & A \times B \\
 \downarrow \sigma'_{A,B} & & \downarrow \sigma_{A,B}^{\text{Sets}} \\
 B \otimes_{\text{Sets}} A & \xrightarrow{\text{id}_{\text{Sets}|B,A}^{\otimes}} & B \times A
 \end{array}$$

commutes. First, note that the diagram

$$\begin{array}{ccc}
 \text{pt} \otimes_{\text{Sets}} \text{pt} & \xrightarrow{\text{id}_{\text{Sets}|\text{pt}, \text{pt}}^{\otimes}} & \text{pt} \times \text{pt} \\
 \downarrow \sigma'_{\text{pt}, \text{pt}} & & \downarrow \sigma_{\text{pt}, \text{pt}}^{\text{Sets}} \\
 \text{pt} \otimes_{\text{Sets}} \text{pt} & \xrightarrow{\text{id}_{\text{Sets}|\text{pt}, \text{pt}}^{\otimes}} & \text{pt} \times \text{pt} \\
 & & \searrow !_{\text{pt} \times \text{pt}} \\
 & & \text{pt}
 \end{array}$$

commutes by the terminality of pt (**Constructions With Sets, Definition 4.1.1.2**).

Since the map $!_{\text{pt} \times \text{pt}} : \text{pt} \times \text{pt} \rightarrow \text{pt}$ is invertible (e.g. with inverse $\lambda_{\text{pt}}^{\text{Sets}, -1}$), the

diagram

$$\begin{array}{ccc}
 \text{pt} \otimes_{\mathbf{Sets}} \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},\text{pt}}^{\otimes}} & \text{pt} \times \text{pt} \\
 \sigma'_{\text{pt},\text{pt}} \downarrow & & \downarrow \sigma_{\text{pt},\text{pt}}^{\mathbf{Sets}} \\
 \text{pt} \otimes_{\mathbf{Sets}} \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},\text{pt}}^{\otimes}} & \text{pt} \times \text{pt}
 \end{array}$$

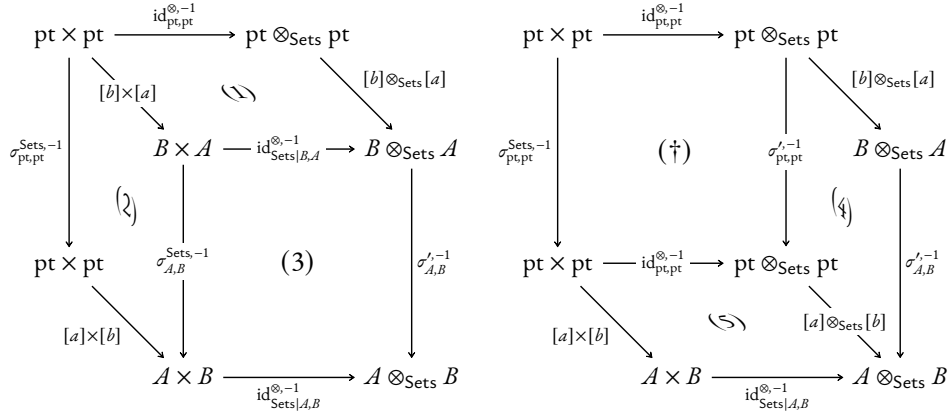
also commutes. Taking inverses, we see that the diagram

$$\begin{array}{ccc}
 \text{pt} \times \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},\text{pt}}^{\otimes, -1}} & \text{pt} \otimes_{\mathbf{Sets}} \text{pt} \\
 \sigma_{\text{pt},\text{pt}}^{\mathbf{Sets}, -1} \downarrow & (\dagger) & \downarrow \sigma'_{\text{pt},\text{pt}}{}^{-1} \\
 \text{pt} \times \text{pt} & \xrightarrow{\text{id}_{\mathbf{Sets}|\text{pt},\text{pt}}^{\otimes, -1}} & \text{pt} \otimes_{\mathbf{Sets}} \text{pt}
 \end{array}$$

commutes as well. Now, let $A, B \in \text{Obj}(\mathbf{Sets})$, let $a \in A$, let $b \in B$, and consider the diagram

$$\begin{array}{ccccc}
 \text{pt} \times \text{pt} & \xrightarrow{\text{id}_{\text{pt},\text{pt}}^{\otimes, -1}} & \text{pt} \otimes_{\mathbf{Sets}} \text{pt} & & \\
 \downarrow \sigma_{\text{pt},\text{pt}}^{\mathbf{Sets}, -1} & \searrow [b] \times [a] & \downarrow \sigma'_{\text{pt},\text{pt}}{}^{-1} & \searrow [b] \otimes_{\mathbf{Sets}} [a] & \\
 & B \times A & \xrightarrow{\text{id}_{\mathbf{Sets}|B,A}^{\otimes, -1}} & B \otimes_{\mathbf{Sets}} A & \\
 & \downarrow \sigma_{A,B}^{\mathbf{Sets}, -1} & \downarrow \sigma'_{A,B}{}^{-1} & & \\
 \text{pt} \times \text{pt} & \xrightarrow{\text{id}_{\text{pt},\text{pt}}^{\otimes, -1}} & \text{pt} \otimes_{\mathbf{Sets}} \text{pt} & & \\
 \searrow [a] \times [b] & \downarrow \sigma_{A,B}^{\mathbf{Sets}, -1} & \searrow [a] \otimes_{\mathbf{Sets}} [b] & \downarrow \sigma'_{A,B}{}^{-1} & \\
 & A \times B & \xrightarrow{\text{id}_{\mathbf{Sets}|A,B}^{\otimes, -1}} & A \otimes_{\mathbf{Sets}} B &
 \end{array}$$

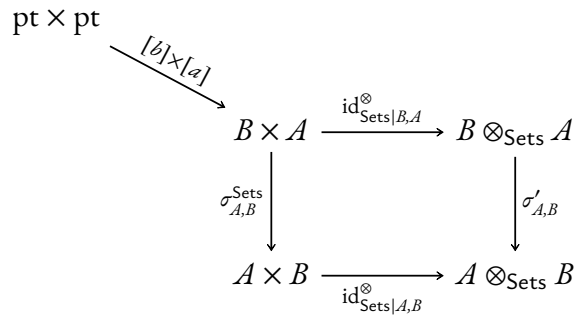
which we partition into subdiagrams as follows:



Since:

- Subdiagram (2) commutes by the naturality of $\sigma^{\text{Sets}, -1}$.
- Subdiagram (5) commutes by the naturality of $\text{id}^{\otimes, -1}$.
- Subdiagram (\dagger) commutes, as proved above.
- Subdiagram (4) commutes by the naturality of σ'^{-1} .
- Subdiagram (1) commutes by the naturality of $\text{id}^{\otimes, -1}$.

it follows that the diagram



commutes. We then have

$$[\text{id}_{\text{Sets}|A, B}^{\otimes, -1} \circ \sigma_{A, B}^{\text{Sets}, -1}](b, a) = [\text{id}_{\text{Sets}|A, B}^{\otimes, -1} \circ \sigma_{A, B}^{\text{Sets}, -1} \circ ([b] \times [a])](\star, \star)$$

$$\begin{aligned}
&= [\sigma'_{A,B}{}^{-1} \circ \text{id}_{\mathbf{Sets}|B,A}^{\otimes,-1} \circ ([b] \times [a])](\star, \star) \\
&= [\sigma'_{A,B}{}^{-1} \circ \text{id}_{\mathbf{Sets}|B,A}^{\otimes,-1}](b, a)
\end{aligned}$$

for each $(b, a) \in B \times A$, and thus we have

$$\text{id}_{\mathbf{Sets}|A,B}^{\otimes,-1} \circ \sigma_{A,B}^{\mathbf{Sets},-1} = \sigma'_{A,B}{}^{-1} \circ \text{id}_{\mathbf{Sets}|B,A}^{\otimes,-1}.$$

Taking inverses then gives

$$\sigma_{A,B}^{\mathbf{Sets}} \circ \text{id}_{\mathbf{Sets}|A,B}^{\otimes} = \text{id}_{\mathbf{Sets}|B,A}^{\otimes} \circ \sigma'_{A,B},$$

showing that the diagram

$$\begin{array}{ccc}
A \otimes_{\mathbf{Sets}} B & \xrightarrow{\text{id}_{\mathbf{Sets}|A,B}^{\otimes}} & A \times B \\
\sigma'_{A,B} \downarrow & & \downarrow \sigma_{A,B}^{\mathbf{Sets}} \\
B \otimes_{\mathbf{Sets}} A & \xrightarrow{\text{id}_{\mathbf{Sets}|B,A}^{\otimes}} & B \times A
\end{array}$$

indeed commutes.

Uniqueness of the Isomorphism $\otimes_{\mathbf{Sets}} \cong \times$: Let $\phi, \psi: -_1 \otimes_{\mathbf{Sets}} -_2 \Rightarrow -_1 \times -_2$ be natural isomorphisms. Since these isomorphisms are compatible with the unitors of \mathbf{Sets} with respect to \times and \otimes (as shown above), we have

$$\begin{aligned}
\lambda'_B &= \lambda_B^{\mathbf{Sets}} \circ \phi_{\text{pt},B} \circ (\text{id}_{1|\mathbf{Sets}}^{\otimes} \otimes_{\mathbf{Sets}} \text{id}_Y), \\
\lambda'_B &= \lambda_B^{\mathbf{Sets}} \circ \psi_{\text{pt},B} \circ (\text{id}_{1|\mathbf{Sets}}^{\otimes} \otimes_{\mathbf{Sets}} \text{id}_Y).
\end{aligned}$$

Postcomposing both sides with $\lambda_B^{\mathbf{Sets},-1}$ gives

$$\begin{aligned}
\lambda_B^{\mathbf{Sets},-1} \circ \lambda'_B \circ (\text{id}_{1|\mathbf{Sets}}^{\otimes,-1} \otimes_{\mathbf{Sets}} \text{id}_Y) &= \phi_{\text{pt},B}, \\
\lambda_B^{\mathbf{Sets},-1} \circ \lambda'_B \circ (\text{id}_{1|\mathbf{Sets}}^{\otimes} \otimes_{\mathbf{Sets}} \text{id}_Y) &= \psi_{\text{pt},B},
\end{aligned}$$

and thus we have

$$\phi_{\text{pt},B} = \psi_{\text{pt},B}$$

for each $B \in \text{Obj}(\text{Sets})$. Now, let $a \in A$ and consider the naturality diagrams

$$\begin{array}{ccc} \text{pt} \times B & \xrightarrow{[a] \times \text{id}_B} & A \times B \\ \phi_{\text{pt}, B} \downarrow & & \downarrow \phi_{A, B} \\ \text{pt} \otimes_{\text{Sets}} B & \xrightarrow{[a] \otimes_{\text{Sets}} \text{id}_B} & A \otimes_{\text{Sets}} B \end{array} \quad \begin{array}{ccc} \text{pt} \times B & \xrightarrow{[a] \times \text{id}_B} & A \times B \\ \psi_{\text{pt}, B} \downarrow & & \downarrow \psi_{A, B} \\ \text{pt} \otimes_{\text{Sets}} B & \xrightarrow{[a] \otimes_{\text{Sets}} \text{id}_B} & A \otimes_{\text{Sets}} B \end{array}$$

for ϕ and ψ with respect to the morphisms $[a]$ and id_B . Having shown that $\phi_{\text{pt}, B} = \psi_{\text{pt}, B}$, we have

$$\begin{aligned} \phi_{A, B}(a, b) &= [\phi_{A, B} \circ ([a] \times \text{id}_B)](\star, b) \\ &= [([a] \otimes_{\text{Sets}} \text{id}_B) \circ \phi_{\text{pt}, B}](\star, b) \\ &= [([a] \otimes_{\text{Sets}} \text{id}_B) \circ \psi_{\text{pt}, B}](\star, b) \\ &= [\psi_{A, B} \circ ([a] \times \text{id}_B)](\star, b) \\ &= \psi_{A, B}(a, b) \end{aligned}$$

for each $(a, b) \in A \times B$. Therefore we have

$$\phi_{A, B} = \psi_{A, B}$$

for each $A, B \in \text{Obj}(\text{Sets})$ and thus $\phi = \psi$, showing the isomorphism $\otimes_{\text{Sets}} \cong \times$ to be unique. \square

Corollary 5.1.10.1.2. The symmetric monoidal structure on the category Sets of [Definition 5.1.9.1.1](#) is uniquely determined by the following requirements:

1. *Two-Sided Preservation of Colimits.* The tensor product

$$\otimes_{\text{Sets}} : \text{Sets} \times \text{Sets} \rightarrow \text{Sets}$$

of Sets preserves colimits separately in each variable.

2. *The Unit Object Is pt.* We have $1_{\text{Sets}} \cong \text{pt}$.

More precisely, the full subcategory of the category $\mathcal{M}_{\text{E}_\infty}(\text{Sets})$ of ?? spanned by the symmetric monoidal categories $(\text{Sets}, \otimes_{\text{Sets}}, 1_{\text{Sets}}, \lambda^{\text{Sets}}, \rho^{\text{Sets}}, \sigma^{\text{Sets}})$ satisfying [Items 1](#) and [2](#) is contractible.

Proof. Since Sets is locally presentable (??), it follows from ?? that [Item 1](#) is equivalent to the existence of an internal Hom as in [Item 1](#) of [Definition 5.1.10.1.1](#). The result then follows from [Definition 5.1.10.1.1](#). \square

5.2 The Monoidal Category of Sets and Coproducts

5.2.1 Coproducts of Sets

See [Constructions With Sets, Section 4.2.3](#).

5.2.2 The Monoidal Unit

Definition 5.2.2.1.1. The **monoidal unit of the coproduct of sets** is the functor

$$0^{\text{Sets}}: \text{pt} \rightarrow \text{Sets}$$

defined by

$$0_{\text{Sets}} \stackrel{\text{def}}{=} \emptyset,$$

where \emptyset is the empty set of [Constructions With Sets, Definition 4.3.1.1.1](#).

5.2.3 The Associator

Definition 5.2.3.1.1. The **associator of the coproduct of sets** is the natural isomorphism

$$\alpha^{\text{Sets}, \amalg}: \amalg \circ (\amalg \times \text{id}_{\text{Sets}}) \xrightarrow{\sim} \amalg \circ (\text{id}_{\text{Sets}} \times \amalg) \circ \alpha_{\text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}},$$

as in the diagram

$$\begin{array}{ccccc}
 & & \text{Sets} \times (\text{Sets} \times \text{Sets}) & & \\
 & \nearrow \alpha_{\text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}} & & \searrow \text{id} \times \amalg & \\
 (\text{Sets} \times \text{Sets}) \times \text{Sets} & & & & \text{Sets} \times \text{Sets} \\
 \downarrow \amalg \times \text{id} & \nearrow \alpha_{\text{Sets}, \amalg} & & \searrow \amalg & \\
 \text{Sets} \times \text{Sets} & & \text{Sets} \times \text{Sets} & & \text{Sets} \\
 & \xrightarrow{\amalg} & & &
 \end{array}$$

whose component

$$\alpha_{X,Y,Z}^{\text{Sets}, \amalg}: (X \amalg Y) \amalg Z \xrightarrow{\sim} X \amalg (Y \amalg Z)$$

at (X, Y, Z) is given by

$$\alpha_{X,Y,Z}^{\text{Sets}, \amalg}(a) \stackrel{\text{def}}{=} \begin{cases} (0, x) & \text{if } a = (0, (0, x)), \\ (1, (0, y)) & \text{if } a = (0, (1, y)), \\ (1, (1, z)) & \text{if } a = (1, z) \end{cases}$$

for each $a \in (X \amalg Y) \amalg Z$.

Proof. Unwinding the Definitions of $(X \amalg Y) \amalg Z$ and $X \amalg (Y \amalg Z)$: Firstly, we unwind the expressions for $(X \amalg Y) \amalg Z$ and $X \amalg (Y \amalg Z)$. We have

$$\begin{aligned} (X \amalg Y) \amalg Z &\stackrel{\text{def}}{=} \{(0, a) \in S \mid a \in X \amalg Y\} \cup \{(1, z) \in S \mid z \in Z\} \\ &= \{(0, (0, x)) \in S \mid x \in X\} \cup \{(0, (1, y)) \in S \mid y \in Y\} \\ &\quad \cup \{(1, z) \in S \mid z \in Z\}, \end{aligned}$$

where $S = \{0, 1\} \times ((X \amalg Y) \cup Z)$ and

$$\begin{aligned} X \amalg (Y \amalg Z) &\stackrel{\text{def}}{=} \{(0, x) \in S' \mid x \in X\} \cup \{(1, a) \in S' \mid a \in Y \amalg Z\} \\ &= \{(0, x) \in S' \mid x \in X\} \cup \{(1, (0, y)) \in S' \mid y \in Y\} \\ &\quad \cup \{(1, (1, z)) \in S' \mid z \in Z\}, \end{aligned}$$

where $S' = \{0, 1\} \times (X \cup (Y \amalg Z))$.

Invertibility: The inverse of $\alpha_{X,Y,Z}^{\text{Sets}, \amalg}$ is the map

$$\alpha_{X,Y,Z}^{\text{Sets}, \amalg, -1} : X \amalg (Y \amalg Z) \rightarrow (X \amalg Y) \amalg Z$$

given by

$$\alpha_{X,Y,Z}^{\text{Sets}, \amalg, -1}(a) \stackrel{\text{def}}{=} \begin{cases} (0, (0, x)) & \text{if } a = (0, x), \\ (0, (1, y)) & \text{if } a = (1, (0, y)), \\ (1, z) & \text{if } a = (1, (1, z)) \end{cases}$$

for each $a \in X \amalg Y(\amalg Z)$. Indeed:

- *Invertibility I.* The map $\alpha_{X,Y,Z}^{\text{Sets}, \amalg, -1} \circ \alpha_{X,Y,Z}^{\text{Sets}, \amalg}$ acts on elements as

$$\begin{aligned} (0, (0, x)) &\mapsto (0, x) \mapsto (0, (0, x)), \\ (0, (0, y)) &\mapsto (1, (0, y)) \mapsto (0, (0, y)), \\ (1, z) &\mapsto (1, (1, z)) \mapsto (1, z) \end{aligned}$$

and hence is equal to the identity map of $(X \amalg Y) \amalg Z$.

- *Invertibility II.* The map $\alpha_{X,Y,Z}^{\text{Sets}, \coprod} \circ \alpha_{X,Y,Z}^{\text{Sets}, \coprod, -1}$ acts on elements as

$$\begin{aligned} (0, x) &\mapsto (0, (0, x)) \mapsto (0, x), \\ (1, (0, y)) &\mapsto (0, (0, y)) \mapsto (1, (0, y)), \\ (1, (1, z)) &\mapsto (1, z) \mapsto (1, (1, z)) \end{aligned}$$

and hence is equal to the identity map of $X \coprod (Y \coprod Z)$.

Therefore $\alpha_{X,Y,Z}^{\text{Sets}, \coprod}$ is indeed an isomorphism.

Naturality: We need to show that, given functions

$$\begin{aligned} f &: X \rightarrow X', \\ g &: Y \rightarrow Y', \\ h &: Z \rightarrow Z' \end{aligned}$$

the diagram

$$\begin{array}{ccc} (X \coprod Y) \coprod Z & \xrightarrow{(f \coprod g) \coprod h} & (X' \coprod Y') \coprod Z' \\ \downarrow \alpha_{X,Y,Z}^{\text{Sets}, \coprod} & & \downarrow \alpha_{X',Y',Z'}^{\text{Sets}, \coprod} \\ X \coprod (Y \coprod Z) & \xrightarrow{f \coprod (g \coprod h)} & X' \coprod (Y' \coprod Z') \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
 (0, (0, x)) & \xrightarrow{\quad} & (0, (0, f(x))) \\
 \downarrow & & \downarrow \\
 (0, x) \xrightarrow{\quad} (0, f(x)) & & (0, f(x)) \\
 \\
 (0, (1, y)) & \xrightarrow{\quad} & (0, (1, g(y))) \\
 \downarrow & & \downarrow \\
 (1, (0, y)) \xrightarrow{\quad} (1, (0, g(y))) & & (1, (0, g(y))) \\
 \\
 (1, z) & \xrightarrow{\quad} & (1, b(z)) \\
 \downarrow & & \downarrow \\
 (1, (1, z)) \xrightarrow{\quad} (1, (1, b(z))) & & (1, (1, b(z)))
 \end{array}$$

and hence indeed commutes, showing $\alpha^{\text{Sets}, \amalg}$ to be a natural transformation. *Being a Natural Isomorphism:* Since $\alpha^{\text{Sets}, \amalg}$ is natural and $\alpha^{\text{Sets}, \amalg, -1}$ is a componentwise inverse to $\alpha^{\text{Sets}, \amalg}$, it follows from **Categories, Item 2 of Definition II.9.7.1.2** that $\lambda^{\text{Sets}, -1}$ is also natural. Thus $\alpha^{\text{Sets}, \amalg}$ is a natural isomorphism. \square

5.2.4 The Left Unitor

Definition 5.2.4.1.1. The **left unitor of the coproduct of sets** is the natural isomorphism

$$\lambda^{\text{Sets}, \amalg} : \amalg \circ (\amalg^{\text{Sets}} \times \text{id}_{\text{Sets}}) \xrightarrow{\sim} \lambda_{\text{Sets}}^{\text{Cats}_2}$$

whose component

$$\lambda_X^{\text{Sets}, \amalg} : \emptyset \amalg X \xrightarrow{\sim} X$$

at X is given by

$$\lambda_X^{\text{Sets}, \amalg}((1, x)) \stackrel{\text{def}}{=} x$$

for each $(1, x) \in \emptyset \amalg X$.

Proof. Unwinding the Definition of $\emptyset \amalg X$: Firstly, we unwind the expressions for $\emptyset \amalg X$. We have

$$\begin{aligned} \emptyset \amalg X &\stackrel{\text{def}}{=} \{(0, z) \in S \mid z \in \emptyset\} \cup \{(1, x) \in S \mid x \in X\} \\ &= \emptyset \cup \{(1, x) \in S \mid x \in X\} \\ &= \{(1, x) \in S \mid x \in X\}, \end{aligned}$$

where $S = \{0, 1\} \times (\emptyset \cup X)$.

Invertibility: The inverse of $\lambda_X^{\text{Sets}, \amalg}$ is the map

$$\lambda_X^{\text{Sets}, \amalg, -1} : X \rightarrow \emptyset \amalg X$$

given by

$$\lambda_X^{\text{Sets}, \amalg, -1}(x) \stackrel{\text{def}}{=} (1, x)$$

for each $x \in X$. Indeed:

- *Invertibility I.* We have

$$\begin{aligned} [\lambda_X^{\text{Sets}, \amalg, -1} \circ \lambda_X^{\text{Sets}, \amalg}](1, x) &= \lambda_X^{\text{Sets}, \amalg, -1}(\lambda_X^{\text{Sets}, \amalg}(1, x)) \\ &= \lambda_X^{\text{Sets}, \amalg, -1}(x) \\ &= (1, x) \\ &= [\text{id}_{\emptyset \amalg X}](1, x) \end{aligned}$$

for each $(1, x) \in \emptyset \amalg X$, and therefore we have

$$\lambda_X^{\text{Sets}, \amalg, -1} \circ \lambda_X^{\text{Sets}, \amalg} = \text{id}_{\emptyset \amalg X}.$$

- *Invertibility II.* We have

$$\begin{aligned}
 [\lambda_X^{\text{Sets}, \sqcup} \circ \lambda_X^{\text{Sets}, \sqcup, -1}](x) &= \lambda_X^{\text{Sets}, \sqcup}(\lambda_X^{\text{Sets}, \sqcup, -1}(x)) \\
 &= \lambda_X^{\text{Sets}, \sqcup, -1}(1, x) \\
 &= x \\
 &= [\text{id}_X](x)
 \end{aligned}$$

for each $x \in X$, and therefore we have

$$\lambda_X^{\text{Sets}, \sqcup} \circ \lambda_X^{\text{Sets}, \sqcup, -1} = \text{id}_X.$$

Therefore $\lambda_X^{\text{Sets}, \sqcup}$ is indeed an isomorphism.

Naturality: We need to show that, given a function $f: X \rightarrow Y$, the diagram

$$\begin{array}{ccc}
 \emptyset \sqcup X & \xrightarrow{\text{id}_\emptyset \sqcup f} & \emptyset \sqcup Y \\
 \lambda_X^{\text{Sets}, \sqcup} \downarrow & & \downarrow \lambda_Y^{\text{Sets}, \sqcup} \\
 X & \xrightarrow{f} & Y
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
 (1, x) & & (1, x) \mapsto (1, f(x)) \\
 \downarrow & & \downarrow \\
 x & \mapsto & f(x)
 \end{array}$$

and hence indeed commutes. Therefore $\lambda^{\text{Sets}, \sqcup}$ is a natural transformation.

Being a Natural Isomorphism: Since $\lambda^{\text{Sets}, \sqcup}$ is natural and $\lambda^{\text{Sets}, -1}$ is a componentwise inverse to $\lambda^{\text{Sets}, \sqcup}$, it follows from **Categories, Item 2** of **Definition 11.9.7.1.2** that $\lambda^{\text{Sets}, -1}$ is also natural. Thus $\lambda^{\text{Sets}, \sqcup}$ is a natural isomorphism. \square

5.2.5 The Right Unitor

Definition 5.2.5.1.1. The **right unitor of the coproduct of sets** is the natural isomorphism

$$\rho^{\text{Sets}, \amalg} : \amalg \circ (\text{id} \times 0^{\text{Sets}}) \xrightarrow{\sim} \rho_{\text{Sets}}^{\text{Cats}_2},$$

whose component

$$\rho_X^{\text{Sets}, \amalg} : X \amalg \emptyset \xrightarrow{\sim} X$$

at X is given by

$$\rho_X^{\text{Sets}, \amalg}((0, x)) \stackrel{\text{def}}{=} x$$

for each $(0, x) \in X \amalg \emptyset$.

Proof. Unwinding the Definition of $X \amalg \emptyset$: Firstly, we unwind the expression for $X \amalg \emptyset$. We have

$$\begin{aligned} X \amalg \emptyset &\stackrel{\text{def}}{=} \{(0, x) \in S \mid x \in X\} \cup \{(1, z) \in S \mid z \in \emptyset\} \\ &= \{(0, x) \in S \mid x \in X\} \cup \emptyset \\ &= \{(0, x) \in S \mid x \in X\}, \end{aligned}$$

where $S = \{0, 1\} \times (X \cup \emptyset) = \{0, 1\} \times (\emptyset \cup X) = S$.

Invertibility: The inverse of $\rho_X^{\text{Sets}, \amalg}$ is the map

$$\rho_X^{\text{Sets}, \amalg, -1} : X \rightarrow X \amalg \emptyset$$

given by

$$\rho_X^{\text{Sets}, \amalg, -1}(x) \stackrel{\text{def}}{=} (0, x)$$

for each $x \in X$. Indeed:

- *Invertibility I.* We have

$$\begin{aligned}
 [\rho_X^{\text{Sets}, \amalg, -1} \circ \rho_X^{\text{Sets}, \amalg}](0, x) &= \rho_X^{\text{Sets}, \amalg, -1}(\rho_X^{\text{Sets}, \amalg}(0, x)) \\
 &= \rho_X^{\text{Sets}, \amalg, -1}(x) \\
 &= (0, x) \\
 &= [\text{id}_X \amalg \emptyset](0, x)
 \end{aligned}$$

for each $(0, x) \in \emptyset \amalg X$, and therefore we have

$$\rho_X^{\text{Sets}, \amalg, -1} \circ \rho_X^{\text{Sets}, \amalg} = \text{id}_{\emptyset \amalg X}.$$

- *Invertibility II.* We have

$$\begin{aligned}
 [\rho_X^{\text{Sets}, \amalg} \circ \rho_X^{\text{Sets}, \amalg, -1}](x) &= \rho_X^{\text{Sets}, \amalg}(\rho_X^{\text{Sets}, \amalg, -1}(x)) \\
 &= \rho_X^{\text{Sets}, \amalg, -1}(0, x) \\
 &= x \\
 &= [\text{id}_X](x)
 \end{aligned}$$

for each $x \in X$, and therefore we have

$$\rho_X^{\text{Sets}, \amalg} \circ \rho_X^{\text{Sets}, \amalg, -1} = \text{id}_X.$$

Therefore $\rho_X^{\text{Sets}, \amalg}$ is indeed an isomorphism.

Naturality: We need to show that, given a function $f: X \rightarrow Y$, the diagram

$$\begin{array}{ccc}
 X \amalg \emptyset & \xrightarrow{f \amalg \text{id}_{\emptyset}} & Y \amalg \emptyset \\
 \rho_X^{\text{Sets}, \amalg} \downarrow & & \downarrow \rho_Y^{\text{Sets}, \amalg} \\
 X & \xrightarrow{f} & Y
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
 (0, x) & & (0, x) \longmapsto (1, f(x)) \\
 \downarrow & & \downarrow \\
 x & \longmapsto & f(x)
 \end{array}$$

and hence indeed commutes. Therefore $\rho^{\text{Sets}, \amalg}$ is a natural transformation.

Being a Natural Isomorphism: Since $\rho^{\text{Sets}, \amalg}$ is natural and $\rho^{\text{Sets}, -1}$ is a component-wise inverse to $\rho^{\text{Sets}, \amalg}$, it follows from **Categories, Item 2** of **Definition 11.9.7.1.2** that $\rho^{\text{Sets}, -1}$ is also natural. Thus $\rho^{\text{Sets}, \amalg}$ is a natural isomorphism. \square

5.2.6 The Symmetry

Definition 5.2.6.1.1. The **symmetry of the coproduct of sets** is the natural isomorphism

$$\sigma^{\text{Sets}, \amalg} : \amalg \xrightarrow{\sim} \amalg \circ \sigma_{\text{Sets}, \text{Sets}}^{\text{Cats}_2}, \quad \begin{array}{ccc} \text{Sets} \times \text{Sets} & \xrightarrow{\amalg} & \text{Sets} \\ \sigma_{\text{Sets}, \text{Sets}}^{\text{Cats}_2} \searrow & \parallel \sigma^{\text{Sets}, \amalg} & \nearrow \amalg \\ & \text{Sets} \times \text{Sets} & \end{array}$$

whose component

$$\sigma_{X,Y}^{\text{Sets}, \amalg} : X \amalg Y \xrightarrow{\sim} Y \amalg X$$

at $X, Y \in \text{Obj}(\text{Sets})$ is defined by

$$\sigma_{X,Y}^{\text{Sets}, \amalg}(x, y) \stackrel{\text{def}}{=} (y, x)$$

for each $(x, y) \in X \times Y$.

Proof. Unwinding the Definitions of $X \amalg Y$ and $Y \amalg X$: Firstly, we unwind the expressions for $X \amalg Y$ and $Y \amalg X$. We have

$$X \amalg Y \stackrel{\text{def}}{=} \{(0, x) \in S \mid x \in X\} \cup \{(1, y) \in S \mid y \in Y\},$$

where $S = \{0, 1\} \times (X \cup Y)$ and

$$Y \amalg X \stackrel{\text{def}}{=} \{(0, y) \in S' \mid y \in Y\} \cup \{(1, x) \in S' \mid x \in X\},$$

where $S' = \{0, 1\} \times (Y \cup X) = \{0, 1\} \times (X \cup Y) = S$.

Invertibility: The inverse of $\sigma_{X,Y}^{\text{Sets}, \amalg}$ is the map

$$\sigma_{X,Y}^{\text{Sets}, \amalg, -1} : Y \amalg X \rightarrow X \amalg Y$$

defined by

$$\sigma_{X,Y}^{\text{Sets}, \coprod, -1} \stackrel{\text{def}}{=} \sigma_{Y,X}^{\text{Sets}, \coprod}$$

and hence given by

$$\sigma_{X,Y}^{\text{Sets}, \coprod, -1}(z) \stackrel{\text{def}}{=} \begin{cases} (0, x) & \text{if } z = (1, x), \\ (1, y) & \text{if } z = (0, y) \end{cases}$$

for each $z \in Y \coprod X$. Indeed:

- *Invertibility I.* We have

$$\begin{aligned} [\sigma_{X,Y}^{\text{Sets}, \coprod, -1} \circ \sigma_{X,Y}^{\text{Sets}, \coprod}](0, x) &= \sigma_X^{\text{Sets}, \coprod, -1}(\sigma_X^{\text{Sets}, \coprod}(0, x)) \\ &= \sigma_X^{\text{Sets}, \coprod, -1}(1, x) \\ &= (0, x) \\ &= [\text{id}_X \coprod \text{id}_Y](0, x) \end{aligned}$$

for each $(0, x) \in X \coprod Y$ and

$$\begin{aligned} [\sigma_{X,Y}^{\text{Sets}, \coprod, -1} \circ \sigma_{X,Y}^{\text{Sets}, \coprod}](1, y) &= \sigma_X^{\text{Sets}, \coprod, -1}(\sigma_X^{\text{Sets}, \coprod}(1, y)) \\ &= \sigma_X^{\text{Sets}, \coprod, -1}(0, y) \\ &= (1, y) \\ &= [\text{id}_X \coprod \text{id}_Y](1, y) \end{aligned}$$

for each $(1, y) \in X \coprod Y$, and therefore we have

$$\sigma_{X,Y}^{\text{Sets}, \coprod, -1} \circ \sigma_{X,Y}^{\text{Sets}, \coprod} = \text{id}_{X \coprod Y}.$$

- *Invertibility II.* We have

$$\begin{aligned} [\sigma_{X,Y}^{\text{Sets}, \coprod} \circ \sigma_{X,Y}^{\text{Sets}, \coprod, -1}](0, y) &= \sigma_X^{\text{Sets}, \coprod}(\sigma_X^{\text{Sets}, \coprod, -1}(0, y)) \\ &= \sigma_X^{\text{Sets}, \coprod, -1}(1, y) \\ &= (0, y) \\ &= [\text{id}_Y \coprod \text{id}_X](0, y) \end{aligned}$$

for each $(0, y) \in Y \amalg X$ and

$$\begin{aligned} [\sigma_{X,Y}^{\text{Sets}, \amalg} \circ \sigma_{X,Y}^{\text{Sets}, \amalg, -1}](1, x) &= \sigma_X^{\text{Sets}, \amalg}(\sigma_X^{\text{Sets}, \amalg, -1}(1, x)) \\ &= \sigma_X^{\text{Sets}, \amalg, -1}(0, x) \\ &= (1, x) \\ &= [\text{id}_Y \amalg \text{id}_X](1, x) \end{aligned}$$

for each $(1, x) \in Y \amalg X$, and therefore we have

$$\sigma_X^{\text{Sets}, \amalg} \circ \sigma_X^{\text{Sets}, \amalg, -1} = \text{id}_{Y \amalg X}.$$

Therefore $\sigma_{X,Y}^{\text{Sets}, \amalg}$ is indeed an isomorphism.

Naturality: We need to show that, given functions $f: A \rightarrow X$ and $g: B \rightarrow Y$, the diagram

$$\begin{array}{ccc} A \amalg B & \xrightarrow{f \amalg g} & X \amalg Y \\ \sigma_{A,B}^{\text{Sets}, \amalg} \downarrow & & \downarrow \sigma_{X,Y}^{\text{Sets}, \amalg} \\ B \amalg A & \xrightarrow{g \amalg f} & Y \amalg X \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc} (0, a) & \xrightarrow{\quad} & (0, f(a)) \\ \downarrow & & \downarrow \\ (1, a) & \xrightarrow{\quad} & (1, f(a)) \\ \\ (1, b) & \xrightarrow{\quad} & (1, g(b)) \\ \downarrow & & \downarrow \\ (0, b) & \xrightarrow{\quad} & (0, g(b)) \end{array}$$

and hence indeed commutes. Therefore $\sigma^{\text{Sets}, \amalg}$ is a natural transformation.

Being a Natural Isomorphism: Since $\sigma^{\text{Sets}, \amalg}$ is natural and $\sigma^{\text{Sets}, -1}$ is a componentwise inverse to $\sigma^{\text{Sets}, \amalg}$, it follows from **Categories, Item 2** of **Definition 11.9.7.1.2** that $\sigma^{\text{Sets}, -1}$ is also natural. Thus $\sigma^{\text{Sets}, \amalg}$ is a natural isomorphism. \square

5.2.7 The Monoidal Category of Sets and Coproducts

Proposition 5.2.7.1.1. The category **Sets** admits a closed symmetric monoidal category structure consisting of:

- *The Underlying Category.* The category **Sets** of pointed sets.
- *The Monoidal Product.* The coproduct functor

$$\amalg : \mathbf{Sets} \times \mathbf{Sets} \rightarrow \mathbf{Sets}$$

of **Constructions With Sets**, **Item 1** of **Definition 4.2.3.1.3**.

- *The Monoidal Unit.* The functor

$$0^{\mathbf{Sets}} : \mathbf{pt} \rightarrow \mathbf{Sets}$$

of **Definition 5.2.2.1.1**.

- *The Associators.* The natural isomorphism

$$\alpha^{\mathbf{Sets}, \amalg} : \amalg \circ (\amalg \times \mathrm{id}_{\mathbf{Sets}}) \xrightarrow{\sim} \amalg \circ (\mathrm{id}_{\mathbf{Sets}} \times \amalg) \circ \alpha_{\mathbf{Sets}, \mathbf{Sets}, \mathbf{Sets}}^{\mathbf{Cats}}$$

of **Definition 5.2.3.1.1**.

- *The Left Unitors.* The natural isomorphism

$$\lambda^{\mathbf{Sets}, \amalg} : \amalg \circ (0^{\mathbf{Sets}} \times \mathrm{id}_{\mathbf{Sets}}) \xrightarrow{\sim} \lambda_{\mathbf{Sets}}^{\mathbf{Cats}_2}$$

of **Definition 5.2.4.1.1**.

- *The Right Unitors.* The natural isomorphism

$$\rho^{\mathbf{Sets}, \amalg} : \amalg \circ (\mathrm{id} \times 0^{\mathbf{Sets}}) \xrightarrow{\sim} \rho_{\mathbf{Sets}}^{\mathbf{Cats}_2}$$

of **Definition 5.2.5.1.1**.

- *The Symmetry.* The natural isomorphism

$$\sigma^{\mathbf{Sets}, \amalg} : \times \xrightarrow{\sim} \times \circ \sigma_{\mathbf{Sets}, \mathbf{Sets}}^{\mathbf{Cats}_2}$$

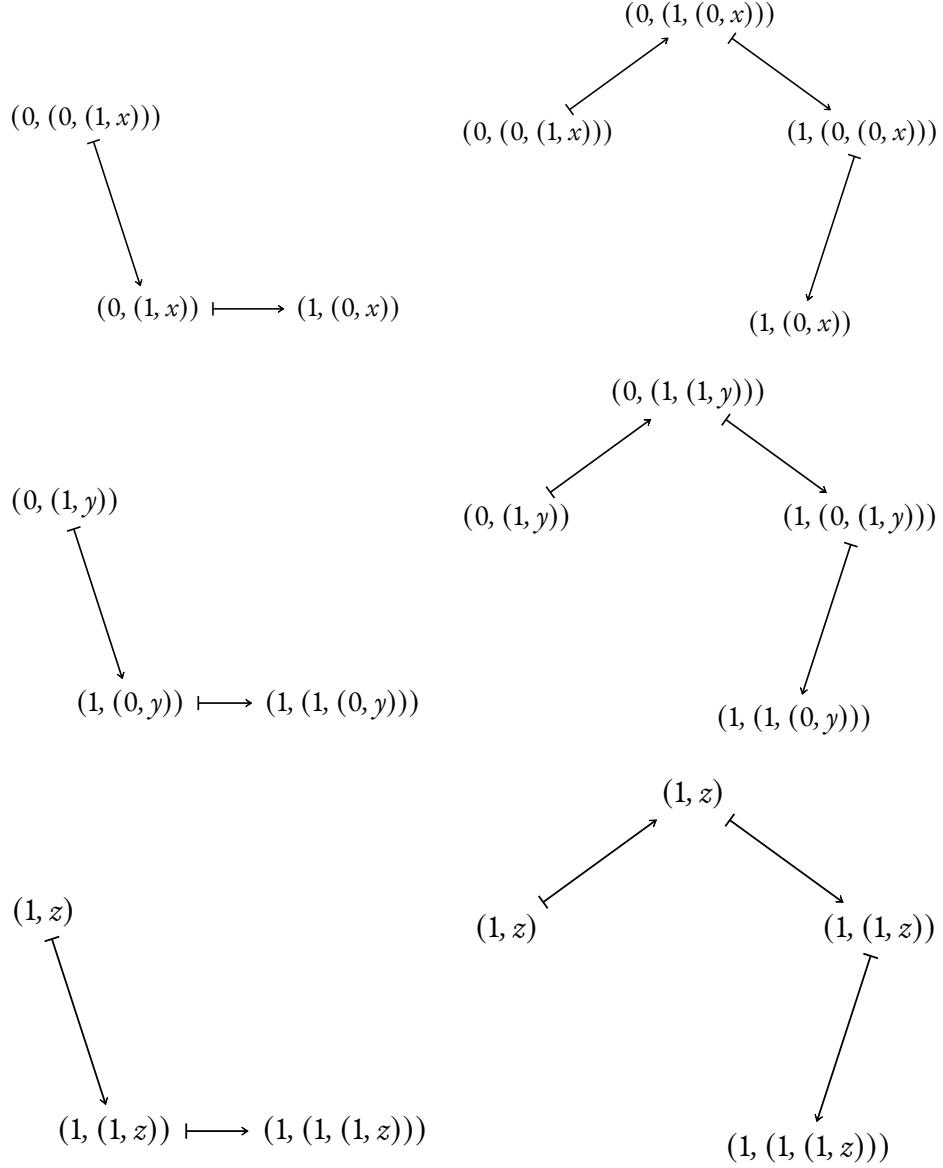
of **Definition 5.2.6.1.1**.

Proof. The Pentagon Identity: Let W, X, Y and Z be sets. We have to show that the diagram

$$\begin{array}{ccc}
 & (W \amalg (X \amalg Y)) \amalg Z & \\
 \alpha_{W,X,Y}^{\text{Sets}, \amalg} \amalg \text{id}_Z \nearrow & & \searrow \alpha_{W,X}^{\text{Sets}, \amalg} \amalg \alpha_{Y,Z}^{\text{Sets}, \amalg} \\
 ((W \amalg X) \amalg Y) \amalg Z & & W \amalg ((X \amalg Y) \amalg Z) \\
 \alpha_{W,X,Y,Z}^{\text{Sets}, \amalg} \searrow & & \nearrow \text{id}_W \amalg \alpha_{X,Y,Z}^{\text{Sets}, \amalg} \\
 (W \amalg X) \amalg (Y \amalg Z) & \xrightarrow{\alpha_{W,X,Y}^{\text{Sets}, \amalg} \amalg \alpha_{Z}^{\text{Sets}, \amalg}} & W \amalg (X \amalg (Y \amalg Z))
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
 & (0, (0, w)) & \\
 & \swarrow \quad \searrow & \\
 (0, (0, (0, w))) & & (0, (0, (0, w))) \quad (0, w) \\
 \searrow & & \searrow \\
 (0, (0, w)) & \longmapsto & (0, w)
 \end{array}$$



and therefore the pentagon identity is satisfied.

The Triangle Identity: Let X and Y be sets. We have to show that the diagram

$$\begin{array}{ccc}
 (X \amalg \emptyset) \amalg Y & \xrightarrow{\alpha_{X,\emptyset,Y}^{\text{Sets}, \amalg}} & X \amalg (\emptyset \amalg Y) \\
 \searrow \rho_X^{\text{Sets}, \amalg} \amalg \text{id}_Y & & \swarrow \text{id}_X \amalg \lambda_Y^{\text{Sets}, \amalg} \\
 & X \amalg Y &
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

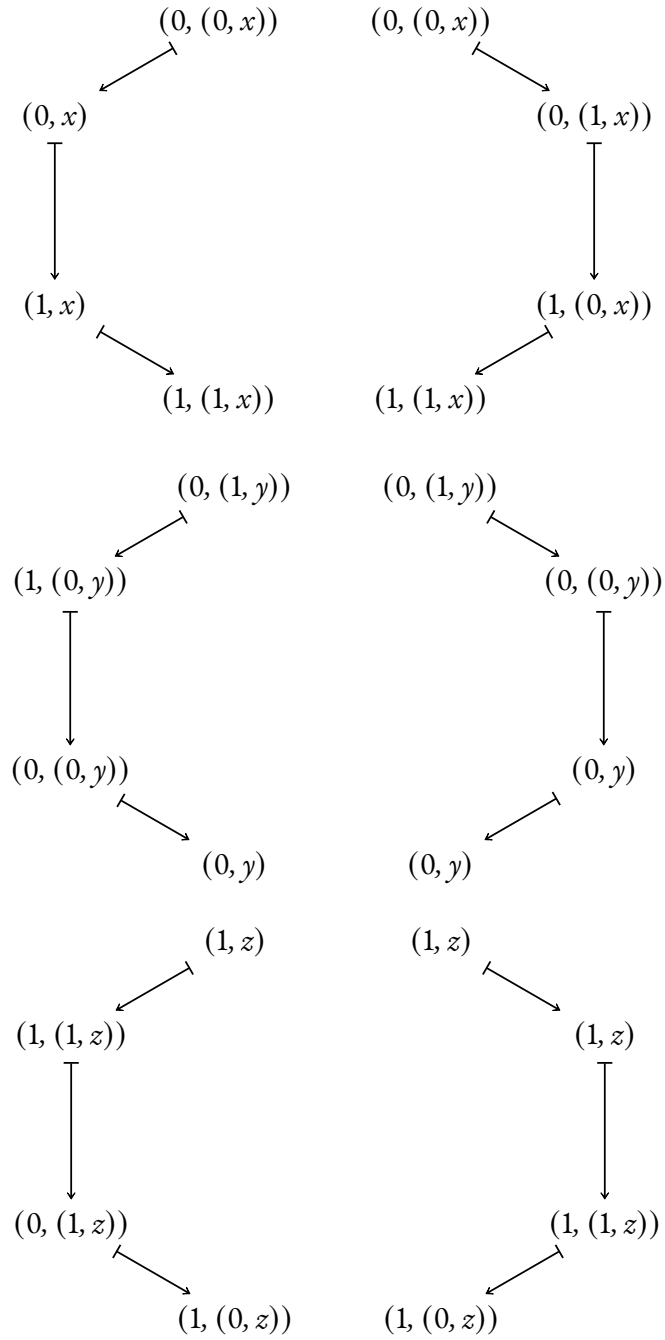
$$\begin{array}{ccc}
 (0, (1, x)) & & (1, (0, x)) \xrightarrow{\quad} (0, x) \\
 \searrow & & \swarrow \\
 & (0, x) & \\
 \\
 (1, y) & & (1, y) \xrightarrow{\quad} (1, (1, y)) \\
 \searrow & & \swarrow \\
 & (1, y) &
 \end{array}$$

and therefore the triangle identity is satisfied.

The Left Hexagon Identity: Let X , Y , and Z be sets. We have to show that the diagram

$$\begin{array}{ccc}
 & (X \amalg Y) \amalg Z & \\
 \alpha_{X,Y,Z}^{\text{Sets}} \swarrow & & \searrow \sigma_{X,Y}^{\text{Sets}} \amalg \text{id}_Z \\
 X \amalg (Y \amalg Z) & & (Y \amalg X) \amalg Z \\
 \downarrow \sigma_{X,Y}^{\text{Sets}} \amalg \text{id}_Z & & \downarrow \alpha_{Y,X,Z}^{\text{Sets}} \\
 (Y \amalg Z) \amalg X & & Y \amalg (X \amalg Z) \\
 \alpha_{Y,Z,X}^{\text{Sets}} \swarrow & & \swarrow \text{id}_Y \amalg \sigma_{X,Z}^{\text{Sets}} \\
 & Y \amalg (Z \amalg X) &
 \end{array}$$

commutes. Indeed, this diagram acts on elements as



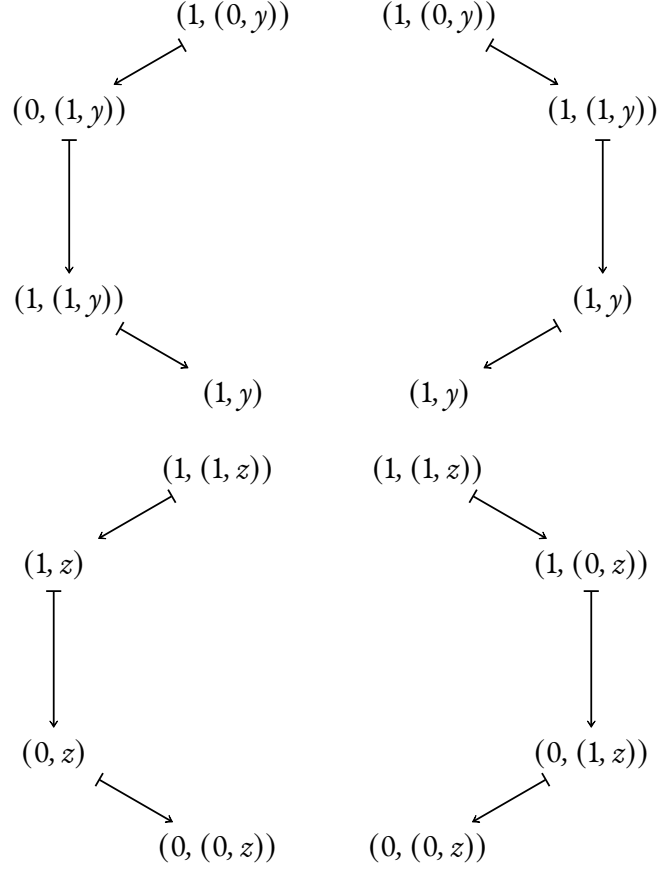
and thus the left hexagon identity is satisfied.

The Right Hexagon Identity: Let X , Y , and Z be sets. We have to show that the diagram

$$\begin{array}{ccc}
 & X \amalg (Y \amalg Z) & \\
 (\alpha_{X,Y,Z}^{\text{Sets}})^{-1} \swarrow & & \searrow \text{id}_X \amalg \sigma_{Y,Z}^{\text{Sets}} \\
 (X \amalg Y) \amalg Z & & X \amalg (Z \amalg Y) \\
 \downarrow \sigma_{X \amalg Y, Z}^{\text{Sets}} & & \downarrow (\alpha_{X,Z,Y}^{\text{Sets}})^{-1} \\
 Z \amalg (X \amalg Y) & & (X \amalg Z) \amalg Y \\
 (\alpha_{Z,X,Y}^{\text{Sets}})^{-1} \swarrow & & \searrow \sigma_{X,Z}^{\text{Sets}} \amalg \text{id}_Y \\
 & (Z \amalg X) \amalg Y &
 \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc}
 & (0, x) & \\
 \swarrow & & \searrow \\
 (0, (0, x)) & & (0, x) \\
 \downarrow & & \downarrow \\
 (1, (0, x)) & & (0, (0, x)) \\
 \swarrow & & \swarrow \\
 & (0, (1, x)) &
 \end{array}$$



and thus the right hexagon identity is satisfied. \square

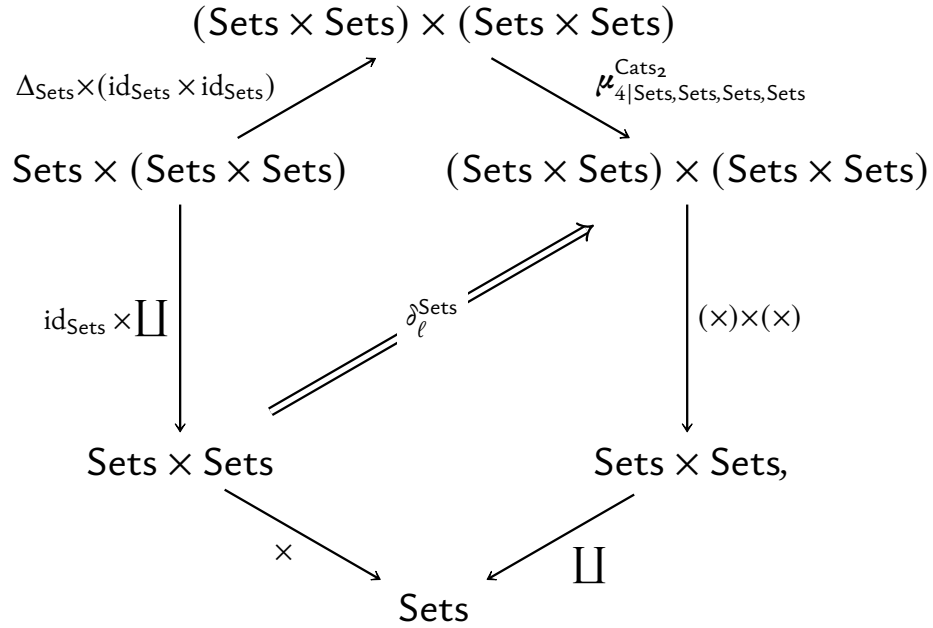
5.3 The Bimonoidal Category of Sets, Products, and Coproducts

5.3.1 The Left Distributor

Definition 5.3.1.1.1. The **left distributor of the product of sets over the coproduct of sets** is the natural isomorphism

$$\delta_\ell^{\text{Sets}} : \times \circ (\text{id}_{\text{Sets}} \times \amalg) \xrightarrow{\sim} \amalg \circ ((\times) \times (\times)) \circ \mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2} \circ (\Delta_{\text{Sets}} \times (\text{id}_{\text{Sets}} \times \text{id}_{\text{Sets}}))$$

as in the diagram



whose component

$$\delta_{\ell|X,Y,Z}^{\text{Sets}}: X \times (Y \amalg Z) \xrightarrow{\sim} (X \times Y) \amalg (X \times Z)$$

at (X, Y, Z) is defined by

$$\delta_{\ell|X,Y,Z}^{\text{Sets}}(x, a) \stackrel{\text{def}}{=} \begin{cases} (0, (x, y)) & \text{if } a = (0, y), \\ (1, (x, z)) & \text{if } a = (1, z) \end{cases}$$

for each $(x, a) \in X \times (Y \amalg Z)$.

Proof. Invertibility: The inverse of $\delta_{\ell|X,Y,Z}^{\text{Sets}}$ is the map

$$\delta_{\ell|X,Y,Z}^{\text{Sets},-1}: (X \times Y) \amalg (X \times Z) \xrightarrow{\sim} X \times (Y \amalg Z)$$

given by

$$\delta_{\ell|X,Y,Z}^{\text{Sets},-1}(a) \stackrel{\text{def}}{=} \begin{cases} (x, (0, y)) & \text{if } a = (0, (x, y)), \\ (x, (1, z)) & \text{if } a = (1, (x, z)) \end{cases}$$

for $a \in (X \times Y) \amalg (X \times Z)$. Indeed:

- *Invertibility I.* The map $\delta_{\ell|X,Y,Z}^{\text{Sets},-1} \circ \delta_{\ell|X,Y,Z}^{\text{Sets}}$ acts on elements as

$$\begin{aligned} (x, (0, y)) &\mapsto (0, (x, y)) \mapsto (x, (0, y)), \\ (x, (1, z)) &\mapsto (1, (x, z)) \mapsto (x, (1, z)), \end{aligned}$$

but these are the two possible cases for elements of $X \times (Y \amalg Z)$. Hence the map is equal to the identity.

- *Invertibility II.* The map $\delta_{\ell|X,Y,Z}^{\text{Sets}} \circ \delta_{\ell|X,Y,Z}^{\text{Sets},-1}$ acts on elements as

$$\begin{aligned} (0, (x, y)) &\mapsto (x, (0, y)) \mapsto (0, (x, y)), \\ (1, (x, z)) &\mapsto (x, (1, z)) \mapsto (1, (x, z)), \end{aligned}$$

but these are the two possible cases for elements of $(X \times Y) \amalg (X \times Z)$. Hence the map is equal to the identity.

Thus $\delta_{\ell|X,Y,Z}^{\text{Sets}}$ is an isomorphism for all X, Y, Z .

Naturality: We need to show that, given functions

$$\begin{aligned} f &: X \rightarrow X', \\ g &: Y \rightarrow Y', \\ h &: Z \rightarrow Z' \end{aligned}$$

the diagram

$$\begin{array}{ccc} X \times (Y \amalg Z) & \xrightarrow{f \times (g \amalg h)} & X' \times (Y' \amalg Z') \\ \delta_{\ell|X,Y,Z}^{\text{Sets}} \downarrow & & \downarrow \delta_{\ell|X',Y',Z'}^{\text{Sets}} \\ (X \times Y) \amalg (X \times Z) & \xrightarrow{(f \times g) \amalg (f \times h)} & (X' \times Y') \amalg (X' \times Z') \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc} (x, (0, y)) & & (x, (0, y)) \mapsto (f(x), (0, f(y))) \\ \downarrow & & \downarrow \\ (0, (x, y)) \mapsto (0, (f(x), g(y))) & & (0, (f(x), g(y))) \end{array}$$

$$\begin{array}{ccc}
 (x, (1, z)) & & (x, (1, z)) \longmapsto (f(x), (1, b(z))) \\
 \downarrow & & \downarrow \\
 (1, (x, z)) \longmapsto (1, (f(x), b(z))) & & (1, (f(x), b(z))),
 \end{array}$$

so it commutes, showing $\delta_\ell^{\text{Sets}}$ to be a natural transformation.

Being a Natural Isomorphism: Since $\delta_\ell^{\text{Sets}}$ is natural and $\delta_\ell^{\text{Sets}, -1}$ is a component-wise inverse to $\delta_\ell^{\text{Sets}}$, it follows from **Categories, Item 2** of **Definition II.9.7.1.2** that $\delta_\ell^{\text{Sets}, -1}$ is also natural. Thus $\delta_\ell^{\text{Sets}}$ is a natural isomorphism. \square

5.3.2 The Right Distributor

Definition 5.3.2.1.1. The **right distributor of the product of sets over the coproduct of sets** is the natural isomorphism

$$\delta_r^{\text{Sets}}: \times \circ (\coprod \times \text{id}_{\text{Sets}}) \xrightarrow{\sim} \coprod \circ ((\times) \times (\times)) \circ \mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2} \circ ((\text{id}_{\text{Sets}} \times \text{id}_{\text{Sets}}) \times \Delta_{\text{Sets}})$$

as in the diagram

$$\begin{array}{ccccc}
 & & (\text{Sets} \times \text{Sets}) \times (\text{Sets} \times \text{Sets}) & & \\
 & \nearrow^{(\text{id}_{\text{Sets}} \times \text{id}_{\text{Sets}}) \times \Delta_{\text{Sets}}} & & \searrow^{\mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2}} & \\
 (\text{Sets} \times \text{Sets}) \times \text{Sets} & & & & (\text{Sets} \times \text{Sets}) \times (\text{Sets} \times \text{Sets}) \\
 \downarrow \coprod \times \text{id}_{\text{Sets}} & & \nearrow \delta_r^{\text{Sets}} & & \downarrow (\times) \times (\times) \\
 \text{Sets} \times \text{Sets} & & & & \text{Sets} \times \text{Sets} \\
 \searrow \times & & & & \swarrow \coprod \\
 & \text{Sets} & & &
 \end{array}$$

whose component

$$\partial_{r|X,Y,Z}^{\text{Sets}}: (X \amalg Y) \times Z \xrightarrow{\sim} (X \times Z) \amalg (Y \times Z)$$

at (X, Y, Z) is defined by

$$\partial_{r|X,Y,Z}^{\text{Sets}}(a, z) \stackrel{\text{def}}{=} \begin{cases} (0, (x, z)) & \text{if } a = (0, x), \\ (1, (y, z)) & \text{if } a = (1, y) \end{cases}$$

for each $(a, z) \in (X \amalg Y) \times Z$.

Proof. Invertibility: The inverse of $\partial_{r|X,Y,Z}^{\text{Sets}}$ is the map

$$\partial_{r|X,Y,Z}^{\text{Sets}, -1}: (X \times Z) \amalg (Y \times Z) \xrightarrow{\sim} (X \amalg Y) \times Z$$

given by

$$\partial_{r|X,Y,Z}^{\text{Sets}, -1}(a) \stackrel{\text{def}}{=} \begin{cases} ((0, x), z) & \text{if } a = (0, (x, z)), \\ ((1, y), z) & \text{if } a = (1, (y, z)) \end{cases}$$

for $a \in (X \times Z) \amalg (Y \times Z)$. Indeed:

- *Invertibility I.* The map $\partial_{r|X,Y,Z}^{\text{Sets}, -1} \circ \partial_{r|X,Y,Z}^{\text{Sets}}$ acts on elements as

$$\begin{aligned} ((0, x), z) &\mapsto (0, (x, z)) \mapsto (0, (x, z)), \\ ((1, y), z) &\mapsto (1, (y, z)) \mapsto (1, (y, z)), \end{aligned}$$

but these are the two possible cases for elements of $(X \amalg Y) \times Z$. Hence the map is equal to the identity.

- *Invertibility II.* The map $\partial_{r|X,Y,Z}^{\text{Sets}} \circ \partial_{r|X,Y,Z}^{\text{Sets}, -1}$ acts on elements as

$$\begin{aligned} (0, (x, z)) &\mapsto ((0, x), z) \mapsto (0, (x, z)), \\ (1, (y, z)) &\mapsto ((1, y), z) \mapsto (1, (y, z)), \end{aligned}$$

but these are the two possible cases for elements of $(X \times Z) \amalg (Y \times Z)$. Hence the map is equal to the identity.

So $\partial_{r|X,Y,Z}^{\text{Sets}}$ is an isomorphism for all X, Y, Z .

Naturality: We need to show that, given functions

$$\begin{aligned} f &: X \rightarrow X', \\ g &: Y \rightarrow Y', \\ h &: Z \rightarrow Z' \end{aligned}$$

the diagram

$$\begin{array}{ccc} (X \amalg Y) \times Z' & \xrightarrow{(f \amalg g) \times h} & (X' \amalg Y') \times Z' \\ \downarrow \delta_{r|X,Y,Z}^{\text{Sets}} & & \downarrow \delta_{r|X',Y',Z'}^{\text{Sets}} \\ (X \times Z) \amalg (Y \times Z) & \xrightarrow{(f \times h) \amalg (g \times b)} & (X' \times Z') \amalg (Y' \times Z') \end{array}$$

commutes. Indeed, this diagram acts on elements as

$$\begin{array}{ccc} ((0, x), z) & \longmapsto & ((0, f(x)), b(z)) \\ \downarrow & & \downarrow \\ (0, (x, z)) & \longmapsto & (0, (f(x), b(z))) \\ \\ ((1, y), z) & \longmapsto & ((1, g(y)), b(z)) \\ \downarrow & & \downarrow \\ (1, (y, z)) & \longmapsto & (1, (g(y), b(z))) \end{array}$$

so it commutes and δ_r^{Sets} is a natural transformation.

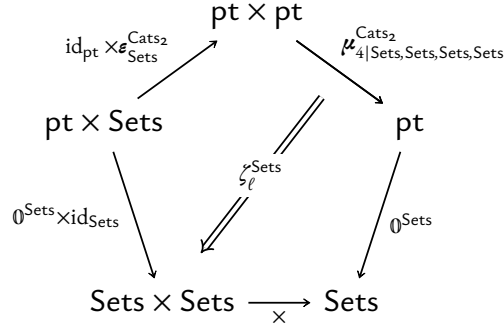
Being a Natural Isomorphism: Since δ_r^{Sets} is natural and $\delta_r^{\text{Sets}, -1}$ is a component-wise inverse to δ_r^{Sets} , it follows from **Categories, Item 2** of **Definition II.9.7.1.2** that $\delta_r^{\text{Sets}, -1}$ is also natural. Thus δ_r^{Sets} is a natural isomorphism. \square

5.3.3 The Left Annihilator

Definition 5.3.3.1.1. The **left annihilator of the product of sets** is the natural isomorphism

$$\zeta_\ell^{\text{Sets}} : \mathbb{0}^{\text{Sets}} \circ \mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2} \circ (\text{id}_{\text{pt}} \times \varepsilon_{\text{Sets}}^{\text{Cats}_2}) \xrightarrow{\sim} \times \circ (\mathbb{0}^{\text{Sets}} \times \text{id}_{\text{Sets}})$$

as in the diagram



with components

$$\zeta_{\ell|A}^{\text{Sets}} : \emptyset \times A \xrightarrow{\sim} \emptyset$$

given by $\zeta_{\ell|A}^{\text{Sets}} \stackrel{\text{def}}{=} \text{pr}_1$.

Proof. Invertibility: The inverse of $\zeta_{\ell|A}^{\text{Sets}}$ is the map

$$\zeta_{\ell|A}^{\text{Sets}, -1} : \emptyset \xrightarrow{\sim} \emptyset \times A$$

given by

$$\zeta_{\ell|A}^{\text{Sets}, -1} \stackrel{\text{def}}{=} \iota_A,$$

where ι_A is as defined in [Constructions With Sets, Definition 4.2.I.I.2](#):

- *Invertibility I.* The map $\zeta_{\ell|A}^{\text{Sets}} \circ \iota_A : \emptyset \rightarrow \emptyset$ is equal to id_{\emptyset} , as \emptyset is the initial object of **Sets**.
- *Invertibility II.* The map $\iota_A \circ \zeta_{\ell|A}^{\text{Sets}}$ is equal to the identity on every $(x, a) \in \emptyset \times A$, of which there are none.

Hence $\zeta_{\ell|A}^{\text{Sets}}$ is an isomorphism.

Naturality: We need to show that given a function $f : A \rightarrow B$, the diagram

$$\begin{array}{ccc} \emptyset \times A & \xrightarrow{\text{id}_{\emptyset} \times f} & \emptyset \times B \\ \zeta_{\ell|A}^{\text{Sets}} \downarrow & & \downarrow \zeta_{\ell|B}^{\text{Sets}} \\ \emptyset & \xrightarrow{\text{id}_{\emptyset}} & \emptyset \end{array}$$

commutes. But since $\emptyset \times A$ has no elements, this is trivially true.

Being a Natural Isomorphism: Since ζ_ℓ^{Sets} is natural and $\zeta_\ell^{\text{Sets}, -1}$ is a component-wise inverse to ζ_ℓ^{Sets} , it follows from **Categories, Item 2 of Definition 11.9.7.1.2** that $\zeta_\ell^{\text{Sets}, -1}$ is also natural. Thus ζ_ℓ^{Sets} is a natural isomorphism. \square

5.3.4 The Right Annihilator

Definition 5.3.4.1.1. The **right annihilator of the product of sets** is the natural isomorphism

$$\zeta_r^{\text{Sets}} : \mathbf{0}^{\text{Sets}} \circ \mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2} \circ (\epsilon_{\text{Sets}}^{\text{Cats}_2} \times \text{id}_{\text{pt}}) \xrightarrow{\sim} \times \circ (\text{id}_{\text{Sets}} \times \mathbf{0}^{\text{Sets}})$$

as in the diagram

$$\begin{array}{ccccc}
 & & \text{pt} \times \text{pt} & & \\
 \epsilon_{\text{Sets}}^{\text{Cats}_2} \times \text{id}_{\text{pt}} \nearrow & & & \nwarrow \mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2} & \\
 \text{Sets} \times \text{pt} & & & & \text{pt} \\
 \text{id}_{\text{Sets}} \times \mathbf{0}^{\text{Sets}} \searrow & \zeta_r^{\text{Sets}} \swarrow \! \! \! \swarrow & & \searrow \mathbf{0}^{\text{Sets}} & \\
 \text{Sets} \times \text{Sets} & \xrightarrow{\times} & \text{Sets} & &
 \end{array}$$

with components

$$\zeta_{r|A}^{\text{Sets}} : A \times \emptyset \xrightarrow{\sim} \emptyset$$

given by $\zeta_{r|A}^{\text{Sets}} \stackrel{\text{def}}{=} \text{pr}_2$.

Proof. Invertibility: The inverse of $\zeta_{r|A}^{\text{Sets}}$ is the map

$$\zeta_{r|A}^{\text{Sets}, -1} : \emptyset \xrightarrow{\sim} A \times \emptyset$$

given by

$$\zeta_{r|A}^{\text{Sets}, -1} \stackrel{\text{def}}{=} \iota_A,$$

where ι_A is as defined in **Constructions With Sets, Definition 4.2.1.1.2**:

- *Invertibility I.* The map $\zeta_{r|A}^{\text{Sets}} \circ \iota_A : \emptyset \rightarrow \emptyset$ is equal to id_\emptyset , as \emptyset is the initial object of **Sets**.

- *Invertibility II.* The map $\iota_A \circ \zeta_{r|A}^{\text{Sets}}$ is equal to the identity on every $(a, x) \in A \times \emptyset$, of which there are none.

Hence $\zeta_{r|A}^{\text{Sets}}$ is an isomorphism.

Naturality: We need to show that given a function $f: A \rightarrow B$, the diagram

$$\begin{array}{ccc} A \times \emptyset & \xrightarrow{f \times \text{id}_\emptyset} & B \times \emptyset \\ \zeta_{r|A}^{\text{Sets}} \downarrow & & \downarrow \zeta_{r|B}^{\text{Sets}} \\ \emptyset & \xrightarrow{\text{id}_\emptyset} & \emptyset \end{array}$$

commutes. But since $A \times \emptyset$ has no elements, this is trivially true.

Being a Natural Isomorphism: Since ζ_r^{Sets} is natural and $\zeta_r^{\text{Sets}, -1}$ is a component-wise inverse to ζ_r^{Sets} , it follows from [Categories, Item 2](#) of [Definition II.9.7.1.2](#) that $\zeta_r^{\text{Sets}, -1}$ is also natural. Thus ζ_r^{Sets} is a natural isomorphism. \square

5.3.5 The Bimonoidal Category of Sets, Products, and Coproducts

Proposition 5.3.5.1.1. The category **Sets** admits a closed symmetric bimonoidal category structure consisting of:

- *The Underlying Category.* The category **Sets** of pointed sets.
- *The Additive Monoidal Product.* The coproduct functor

$$\amalg: \mathbf{Sets} \times \mathbf{Sets} \rightarrow \mathbf{Sets}$$

of [Constructions With Sets, Item 1](#) of [Definition 4.2.3.1.3](#).

- *The Multiplicative Monoidal Product.* The product functor

$$\times: \mathbf{Sets} \times \mathbf{Sets} \rightarrow \mathbf{Sets}$$

of [Constructions With Sets, Item 1](#) of [Definition 4.1.3.1.3](#).

- *The Monoidal Unit.* The functor

$$\mathbf{1}^{\text{Sets}}: \text{pt} \rightarrow \mathbf{Sets}$$

of [Definition 5.1.3.1.1](#).

- *The Monoidal Zero.* The functor

$$0^{\text{Sets}}: \text{pt} \rightarrow \text{Sets}$$

of [Definition 5.1.3.1.1.](#)

- *The Internal Hom.* The internal Hom functor

$$\text{Sets}: \text{Sets}^{\text{op}} \times \text{Sets} \rightarrow \text{Sets}$$

of [Constructions With Sets](#), ?? of ??.

- *The Additive Associators.* The natural isomorphism

$$\alpha^{\text{Sets}, \amalg}: \amalg \circ (\amalg \times \text{id}_{\text{Sets}}) \xrightarrow{\sim} \amalg \circ (\text{id}_{\text{Sets}} \times \amalg) \circ \alpha_{\text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}}$$

of [Definition 5.2.3.1.1.](#)

- *The Additive Left Unitors.* The natural isomorphism

$$\lambda^{\text{Sets}, \amalg}: \amalg \circ (0^{\text{Sets}} \times \text{id}_{\text{Sets}}) \xrightarrow{\sim} \lambda_{\text{Sets}}^{\text{Cats}_2}$$

of [Definition 5.2.4.1.1.](#)

- *The Additive Right Unitors.* The natural isomorphism

$$\rho^{\text{Sets}, \amalg}: \amalg \circ (\text{id} \times 0^{\text{Sets}}) \xrightarrow{\sim} \rho_{\text{Sets}}^{\text{Cats}_2}$$

of [Definition 5.2.5.1.1.](#)

- *The Additive Symmetry.* The natural isomorphism

$$\sigma^{\text{Sets}, \amalg}: \amalg \xrightarrow{\sim} \amalg \circ \sigma_{\text{Sets}, \text{Sets}}^{\text{Cats}_2}$$

of [Definition 5.2.6.1.1.](#)

- *The Multiplicative Associators.* The natural isomorphism

$$\alpha^{\text{Sets}}: \times \circ (\times \times \text{id}_{\text{Sets}}) \xrightarrow{\sim} \times \circ (\text{id}_{\text{Sets}} \times \times) \circ \alpha_{\text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}}$$

of [Definition 5.1.4.1.1.](#)

- *The Multiplicative Left Unitors.* The natural isomorphism

$$\lambda^{\text{Sets}} : \times \circ (\mathbb{1}^{\text{Sets}} \times \text{id}_{\text{Sets}}) \xRightarrow{\sim} \lambda_{\text{Sets}}^{\text{Cats}_2}$$

of **Definition 5.1.5.1.1.**

- *The Multiplicative Right Unitors.* The natural isomorphism

$$\rho^{\text{Sets}} : \times \circ (\text{id} \times \mathbb{1}^{\text{Sets}}) \xRightarrow{\sim} \rho_{\text{Sets}}^{\text{Cats}_2}$$

of **Definition 5.1.6.1.1.**

- *The Multiplicative Symmetry.* The natural isomorphism

$$\sigma^{\text{Sets}} : \times \xRightarrow{\sim} \times \circ \sigma_{\text{Sets}, \text{Sets}}^{\text{Cats}_2}$$

of **Definition 5.1.7.1.1.**

- *The Left Distributor.* The natural isomorphism

$$\delta_\ell^{\text{Sets}} : \times \circ (\text{id}_{\text{Sets}} \times \amalg) \xRightarrow{\sim} \amalg \circ ((\times) \times (\times)) \circ \mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2} \circ (\Delta_{\text{Sets}} \times (\text{id}_{\text{Sets}} \times \text{id}_{\text{Sets}}))$$

of **Definition 5.3.1.1.1.**

- *The Right Distributor.* The natural isomorphism

$$\delta_r^{\text{Sets}} : \times \circ (\amalg \times \text{id}_{\text{Sets}}) \xRightarrow{\sim} \amalg \circ ((\times) \times (\times)) \circ \mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2} \circ ((\text{id}_{\text{Sets}} \times \text{id}_{\text{Sets}}) \times \Delta_{\text{Sets}})$$

of **Definition 5.3.2.1.1.**

- *The Left Annihilator.* The natural isomorphism

$$\zeta_\ell^{\text{Sets}} : \mathbb{0}^{\text{Sets}} \circ \mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2} \circ (\text{id}_{\text{pt}} \times \varepsilon_{\text{Sets}}^{\text{Cats}_2}) \xRightarrow{\sim} \times \circ (\mathbb{0}^{\text{Sets}} \times \text{id}_{\text{Sets}})$$

of **Definition 5.3.3.1.1.**

- *The Right Annihilator.* The natural isomorphism

$$\zeta_r^{\text{Sets}} : \mathbb{0}^{\text{Sets}} \circ \mu_{4|\text{Sets}, \text{Sets}, \text{Sets}, \text{Sets}}^{\text{Cats}_2} \circ (\varepsilon_{\text{Sets}}^{\text{Cats}_2} \times \text{id}_{\text{pt}}) \xrightarrow{\sim} \times \circ (\text{id}_{\text{Sets}} \times \mathbb{0}^{\text{Sets}})$$

of **Definition 5.3.4.1.1.**

Proof. Omitted. □

Appendices

A Other Chapters

Preliminaries

1. Introduction
2. A Guide to the Literature

Sets

3. Sets
4. Constructions With Sets
5. Monoidal Structures on the Category of Sets
6. Pointed Sets
7. Tensor Products of Pointed Sets

Relations

8. Relations
9. Constructions With Relations

10. Conditions on Relations

Categories

11. Categories
12. Presheaves and the Yoneda Lemma

Monoidal Categories

13. Constructions With Monoidal Categories

Bicategories

14. Types of Morphisms in Bicategories

Extra Part

15. Notes