

Types of Morphisms in Bicategories

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In this chapter, we study special kinds of morphisms in bicategories:

1. *Monomorphisms and Epimorphisms in Bicategories* ([Sections 14.1 and 14.2](#)). There is a large number of different notions capturing the idea of a “monomorphism” or of an “epimorphism” in a bicategory.

Arguably, the notion that best captures these concepts is that of a *pseudomononic morphism* ([Definition 14.1.10.1.1](#)) and of a *pseudoepic morphism* ([Definition 14.2.10.1.1](#)), although the other notions introduced in [Sections 14.1](#) and [14.2](#) are also interesting on their own.

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14.1 Monomorphisms in Bicategories

14.1.1 Representably Faithful Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.1.1.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **representably faithful**¹ if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f_*: \text{Hom}_{\mathcal{C}}(X, A) \rightarrow \text{Hom}_{\mathcal{C}}(X, B)$$

given by postcomposition by f is faithful.

Remark 14.1.1.1.2. In detail, f is representably faithful if, for all diagrams in \mathcal{C} of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if we have

$$\text{id}_f \star \alpha = \text{id}_f \star \beta,$$

then $\alpha = \beta$.

Example 14.1.1.1.3. Here are some examples of representably faithful morphisms.

1. *Representably Faithful Morphisms in \mathbf{Cats}_2 .* The representably faithful morphisms in \mathbf{Cats}_2 are precisely the faithful functors; see [Categories, Item 2](#) of [Definition 11.6.1.1.2](#).

¹*Further Terminology:* Also called simply a **faithful morphism**, based on [Item 1](#) of

2. *Representably Faithful Morphisms in **Rel***. Every morphism of **Rel** is representably faithful; see [Relations](#), ?? of ??.

14.1.2 Representably Full Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.1.2.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **representably full**² if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f_*: \text{Hom}_{\mathcal{C}}(X, A) \rightarrow \text{Hom}_{\mathcal{C}}(X, B)$$

given by postcomposition by f is full.

Remark 14.1.2.1.2. In detail, f is representably full if, for each $X \in \text{Obj}(\mathcal{C})$ and each 2-morphism

$$\beta: f \circ \phi \Rightarrow f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of \mathcal{C} , there exists a 2-morphism

$$\alpha: \phi \Rightarrow \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of \mathcal{C} such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B = X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \text{id}_f \star \alpha.$$

Example 14.1.2.1.3. Here are some examples of representably full morphisms.

Definition 14.1.1.1.3.

²*Further Terminology:* Also called simply a **full morphism**, based on [Item 1](#) of

1. *Representably Full Morphisms in \mathbf{Cats}_2* . The representably full morphisms in \mathbf{Cats}_2 are precisely the full functors; see [Categories](#), ?? of [Definition 11.6.2.1.2](#).
2. *Representably Full Morphisms in \mathbf{Rel}* . The representably full morphisms in \mathbf{Rel} are characterised in [Relations](#), ?? of ??.

14.1.3 Representably Fully Faithful Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.1.3.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **representably fully faithful**³ if the following equivalent conditions are satisfied:

1. The 1-morphism f is representably faithful ([Definition 14.1.1.1.1](#)) and representably full ([Definition 14.1.2.1.1](#)).
2. For each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f_*: \text{Hom}_{\mathcal{C}}(X, A) \rightarrow \text{Hom}_{\mathcal{C}}(X, B)$$

given by postcomposition by f is fully faithful.

Remark 14.1.3.1.2. In detail, f is representably fully faithful if the conditions in [Definition 14.1.1.1.2](#) and [Definition 14.1.2.1.2](#) hold:

1. For all diagrams in \mathcal{C} of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if we have

$$\text{id}_f \star \alpha = \text{id}_f \star \beta,$$

then $\alpha = \beta$.

Definition 14.1.2.1.3.

³*Further Terminology:* Also called simply a **fully faithful morphism**, based on [Item 1](#) of [Definition 14.1.3.1.3](#).

2. For each $X \in \text{Obj}(\mathcal{C})$ and each 2-morphism

$$\beta: f \circ \phi \Rightarrow f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of \mathcal{C} , there exists a 2-morphism

$$\alpha: \phi \Rightarrow \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of \mathcal{C} such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B = X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \text{id}_f \star \alpha.$$

Example 14.1.3.1.3. Here are some examples of representably fully faithful morphisms.

1. *Representably Fully Faithful Morphisms in \mathbf{Cats}_2 .* The representably fully faithful morphisms in \mathbf{Cats}_2 are precisely the fully faithful functors; see [Categories](#), [Item 6](#) of [Definition 11.6.3.1.2](#).
2. *Representably Fully Faithful Morphisms in \mathbf{Rel} .* The representably fully faithful morphisms of \mathbf{Rel} coincide ([Relations](#), ?? of ??) with the representably full morphisms in \mathbf{Rel} , which are characterised in [Relations](#), ?? of ??.

14.1.4 Morphisms Representably Faithful on Cores

Let \mathcal{C} be a bicategory.

Definition 14.1.4.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **representably faithful on cores** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f_*: \text{Core}(\text{Hom}_{\mathcal{C}}(X, A)) \rightarrow \text{Core}(\text{Hom}_{\mathcal{C}}(X, B))$$

given by postcomposition by f is faithful.

Remark 14.1.4.1.2. In detail, f is representably faithful on cores if, for all diagrams in \mathcal{C} of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if α and β are 2-isomorphisms and we have

$$\mathrm{id}_f \star \alpha = \mathrm{id}_f \star \beta,$$

then $\alpha = \beta$.

14.1.5 Morphisms Representably Full on Cores

Let \mathcal{C} be a bicategory.

Definition 14.1.5.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **representably full on cores** if, for each $X \in \mathrm{Obj}(\mathcal{C})$, the functor

$$f_*: \mathrm{Core}(\mathrm{Hom}_{\mathcal{C}}(X, A)) \rightarrow \mathrm{Core}(\mathrm{Hom}_{\mathcal{C}}(X, B))$$

given by postcomposition by f is full.

Remark 14.1.5.1.2. In detail, f is representably full on cores if, for each $X \in \mathrm{Obj}(\mathcal{C})$ and each 2-isomorphism

$$\beta: f \circ \phi \xRightarrow{\sim} f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of \mathcal{C} , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of \mathcal{C} such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B = X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \mathrm{id}_f \star \alpha.$$

14.1.6 Morphisms Representably Fully Faithful on Cores

Let \mathcal{C} be a bicategory.

Definition 14.1.6.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **representably fully faithful on cores** if the following equivalent conditions are satisfied:

1. The 1-morphism f is representably faithful on cores ([Definition 14.1.5.1.1](#)) and representably full on cores ([Definition 14.1.4.1.1](#)).
2. For each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f_*: \text{Core}(\text{Hom}_{\mathcal{C}}(X, A)) \rightarrow \text{Core}(\text{Hom}_{\mathcal{C}}(X, B))$$

given by postcomposition by f is fully faithful.

Remark 14.1.6.1.2. In detail, f is representably fully faithful on cores if the conditions in [Definition 14.1.4.1.2](#) and [Definition 14.1.5.1.2](#) hold:

1. For all diagrams in \mathcal{C} of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if α and β are 2-isomorphisms and we have

$$\text{id}_f \star \alpha = \text{id}_f \star \beta,$$

then $\alpha = \beta$.

2. For each $X \in \text{Obj}(\mathcal{C})$ and each 2-isomorphism

$$\beta: f \circ \phi \xrightarrow{\sim} f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of \mathcal{C} , there exists a 2-isomorphism

$$\alpha: \phi \xrightarrow{\sim} \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of \mathcal{C} such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B = X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \text{id}_f \star \alpha.$$

14.1.7 Representably Essentially Injective Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.1.7.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **representably essentially injective** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f_*: \text{Hom}_{\mathcal{C}}(X, A) \rightarrow \text{Hom}_{\mathcal{C}}(X, B)$$

given by postcomposition by f is essentially injective.

Remark 14.1.7.1.2. In detail, f is representably essentially injective if, for each pair of morphisms $\phi, \psi: X \rightrightarrows A$ of \mathcal{C} , the following condition is satisfied:

(\star) If $f \circ \phi \cong f \circ \psi$, then $\phi \cong \psi$.

14.1.8 Representably Conservative Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.1.8.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **representably conservative** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f_*: \text{Hom}_{\mathcal{C}}(X, A) \rightarrow \text{Hom}_{\mathcal{C}}(X, B)$$

given by postcomposition by f is conservative.

Remark 14.1.8.1.2. In detail, f is representably conservative if, for each pair of morphisms $\phi, \psi: X \rightrightarrows A$ and each 2-morphism

$$\alpha: \phi \rightrightarrows \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of \mathcal{C} , if the 2-morphism

$$\mathrm{id}_f \star \alpha: f \circ \phi \Rightarrow f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \parallel \\ \mathrm{id}_f \star \alpha \\ \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

is a 2-isomorphism, then so is α .

14.1.9 Strict Monomorphisms

Let \mathcal{C} be a bicategory.

Definition 14.1.9.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is a **strict monomorphism** if, for each $X \in \mathrm{Obj}(\mathcal{C})$, the functor

$$f_*: \mathrm{Hom}_{\mathcal{C}}(X, A) \rightarrow \mathrm{Hom}_{\mathcal{C}}(X, B)$$

given by postcomposition by f is injective on objects, i.e. its action on objects

$$f_*: \mathrm{Obj}(\mathrm{Hom}_{\mathcal{C}}(X, A)) \rightarrow \mathrm{Obj}(\mathrm{Hom}_{\mathcal{C}}(X, B))$$

is injective.

Remark 14.1.9.1.2. In detail, f is a strict monomorphism in \mathcal{C} if, for each diagram in \mathcal{C} of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \Downarrow \psi \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if $f \circ \phi = f \circ \psi$, then $\phi = \psi$.

Example 14.1.9.1.3. Here are some examples of strict monomorphisms.

1. *Strict Monomorphisms in \mathbf{Cats}_2 .* The strict monomorphisms in \mathbf{Cats}_2 are precisely the functors which are injective on objects and injective on morphisms; see [Categories](#), [Item 1](#) of [Definition 11.7.2.1.2](#).
2. *Strict Monomorphisms in \mathbf{Rel} .* The strict monomorphisms in \mathbf{Rel} are characterised in [Relations](#), [??](#).

14.1.10 Pseudomonic Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.1.10.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **pseudomonic** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f_*: \text{Hom}_{\mathcal{C}}(X, A) \rightarrow \text{Hom}_{\mathcal{C}}(X, B)$$

given by postcomposition by f is pseudomonic.

Remark 14.1.10.1.2. In detail, a 1-morphism $f: A \rightarrow B$ of \mathcal{C} is pseudomonic if it satisfies the following conditions:

1. For all diagrams in \mathcal{C} of the form

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B,$$

if we have

$$\text{id}_f \star \alpha = \text{id}_f \star \beta,$$

then $\alpha = \beta$.

2. For each $X \in \text{Obj}(\mathcal{C})$ and each 2-isomorphism

$$\beta: f \circ \phi \xrightarrow{\sim} f \circ \psi, \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of \mathcal{C} , there exists a 2-isomorphism

$$\alpha: \phi \xrightarrow{\sim} \psi, \quad X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A$$

of \mathcal{C} such that we have an equality

$$X \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} A \xrightarrow{f} B \quad = \quad X \begin{array}{c} \xrightarrow{f \circ \phi} \\ \beta \Downarrow \\ \xrightarrow{f \circ \psi} \end{array} B$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \text{id}_f \star \alpha.$$

Proposition 14.1.10.1.3. Let $f: A \rightarrow B$ be a 1-morphism of \mathcal{C} .

1. *Characterisations.* The following conditions are equivalent:
 - (a) The morphism f is pseudomonic.
 - (b) The morphism f is representably full on cores and representably faithful.
 - (c) We have an isocomma square of the form

$$A \stackrel{\text{eq.}}{\cong} A \times_B^{\leftrightarrow} A, \quad \begin{array}{ccc} A & \xrightarrow{\text{id}_A} & A \\ \text{id}_A \downarrow & \nearrow \text{dashed} & \downarrow F \\ A & \xrightarrow{F} & B \end{array}$$

in \mathcal{C} up to equivalence.

2. *Interaction With Cotensors.* If \mathcal{C} has cotensors with $\mathbb{1}$, then the following conditions are equivalent:
 - (a) The morphism f is pseudomonic.
 - (b) We have an isocomma square of the form

$$A \stackrel{\text{eq.}}{\cong} A \times_{\mathbb{1} \pitchfork F}^{\leftrightarrow} B, \quad \begin{array}{ccc} A & \longrightarrow & \mathbb{1} \pitchfork A \\ F \downarrow & \nearrow \text{dashed} & \downarrow \mathbb{1} \pitchfork F \\ B & \longrightarrow & \mathbb{1} \pitchfork B \end{array}$$

in \mathcal{C} up to equivalence.

Proof. **Item 1, Characterisations:** Omitted.

Item 2, Interaction With Cotensors: Omitted. □

14.2 Epimorphisms in Bicategories

14.2.1 Corepresentably Faithful Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.2.1.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **corepresentably faithful** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f^*: \text{Hom}_{\mathcal{C}}(B, X) \rightarrow \text{Hom}_{\mathcal{C}}(A, X)$$

given by precomposition by f is faithful.

Remark 14.2.1.1.2. In detail, f is corepresentably faithful if, for all diagrams in \mathcal{C} of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

Example 14.2.1.1.3. Here are some examples of corepresentably faithful morphisms.

1. *Corepresentably Faithful Morphisms in \mathbf{Cats}_2 .* The corepresentably faithful morphisms in \mathbf{Cats}_2 are characterised in [Categories, Item 5 of Definition 11.6.1.1.2](#).
2. *Corepresentably Faithful Morphisms in \mathbf{Rel} .* Every morphism of \mathbf{Rel} is corepresentably faithful; see [Relations, ?? of ??](#).

14.2.2 Corepresentably Full Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.2.2.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **corepresentably full** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f^*: \text{Hom}_{\mathcal{C}}(B, X) \rightarrow \text{Hom}_{\mathcal{C}}(A, X)$$

given by precomposition by f is full.

Remark 14.2.2.1.2. In detail, f is corepresentably full if, for each $X \in$

$\text{Obj}(\mathcal{C})$ and each 2-morphism

$$\beta: \phi \circ f \Longrightarrow \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of \mathcal{C} , there exists a 2-morphism

$$\alpha: \phi \Longrightarrow \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of \mathcal{C} such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

Example 14.2.2.1.3. Here are some examples of corepresentably full morphisms.

1. *Corepresentably Full Morphisms in \mathbf{Cats}_2 .* The corepresentably full morphisms in \mathbf{Cats}_2 are characterised in [Categories](#), [Item 7](#) of [Definition 11.6.2.1.2](#).
2. *Corepresentably Full Morphisms in \mathbf{Rel} .* The corepresentably full morphisms in \mathbf{Rel} are characterised in [Relations](#), ?? of ??.

14.2.3 Corepresentably Fully Faithful Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.2.3.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **corepresentably fully faithful**⁴ if the following equivalent conditions are satisfied:

⁴*Further Terminology:* Corepresentably fully faithful morphisms have also been called **lax epimorphisms** in the literature (e.g. in [\[Ad +01\]](#)), though we will always use the name “corepresentably fully faithful morphism” instead in this work.

1. The 1-morphism f is corepresentably full ([Definition 14.2.2.1.1](#)) and corepresentably faithful ([Definition 14.2.1.1.1](#)).
2. For each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f^*: \text{Hom}_{\mathcal{C}}(B, X) \rightarrow \text{Hom}_{\mathcal{C}}(A, X)$$

given by precomposition by f is fully faithful.

Remark 14.2.3.1.2. In detail, f is corepresentably fully faithful if the conditions in [Definition 14.2.1.1.2](#) and [Definition 14.2.2.1.2](#) hold:

1. For all diagrams in \mathcal{C} of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

2. For each $X \in \text{Obj}(\mathcal{C})$ and each 2-morphism

$$\beta: \phi \circ f \Rightarrow \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of \mathcal{C} , there exists a 2-morphism

$$\alpha: \phi \Rightarrow \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of \mathcal{C} such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

Example 14.2.3.1.3. Here are some examples of corepresentably fully faithful morphisms.

1. *Corepresentably Fully Faithful Morphisms in \mathbf{Cats}_2 .* The fully faithful epimorphisms in \mathbf{Cats}_2 are characterised in [Categories, Item 10](#) of [Definition 11.6.3.1.2](#).
2. *Corepresentably Fully Faithful Morphisms in \mathbf{Rel} .* The corepresentably fully faithful morphisms of \mathbf{Rel} coincide ([Relations, ??](#) of [??](#)) with the corepresentably full morphisms in \mathbf{Rel} , which are characterised in [Relations, ??](#) of [??](#).

14.2.4 Morphisms Corepresentably Faithful on Cores

Let \mathcal{C} be a bicategory.

Definition 14.2.4.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **corepresentably faithful on cores** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f^*: \text{Core}(\text{Hom}_{\mathcal{C}}(B, X)) \rightarrow \text{Core}(\text{Hom}_{\mathcal{C}}(A, X))$$

given by precomposition by f is faithful.

Remark 14.2.4.1.2. In detail, f is corepresentably faithful on cores if, for all diagrams in \mathcal{C} of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if α and β are 2-isomorphisms and we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

14.2.5 Morphisms Corepresentably Full on Cores

Let \mathcal{C} be a bicategory.

Definition 14.2.5.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **corepresentably full on cores** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f^*: \text{Core}(\text{Hom}_{\mathcal{C}}(B, X)) \rightarrow \text{Core}(\text{Hom}_{\mathcal{C}}(A, X))$$

given by precomposition by f is full.

Remark 14.2.5.1.2. In detail, f is corepresentably full on cores if, for each $X \in \text{Obj}(\mathcal{C})$ and each 2-isomorphism

$$\beta: \phi \circ f \xRightarrow{\sim} \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \Downarrow \beta \\ \xrightarrow{\psi \circ f} \end{array} X$$

of \mathcal{C} , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \Downarrow \alpha \\ \xrightarrow{\psi} \end{array} X$$

of \mathcal{C} such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \Downarrow \alpha \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \Downarrow \beta \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

14.2.6 Morphisms Corepresentably Fully Faithful on Cores

Let \mathcal{C} be a bicategory.

Definition 14.2.6.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **corepresentably fully faithful on cores** if the following equivalent conditions are satisfied:

1. The 1-morphism f is corepresentably full on cores ([Definition 14.2.5.1.1](#)) and corepresentably faithful on cores ([Definition 14.2.1.1.1](#)).
2. For each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f^*: \text{Core}(\text{Hom}_{\mathcal{C}}(B, X)) \rightarrow \text{Core}(\text{Hom}_{\mathcal{C}}(A, X))$$

given by precomposition by f is fully faithful.

Remark 14.2.6.1.2. In detail, f is corepresentably fully faithful on cores if the conditions in [Definition 14.2.4.1.2](#) and [Definition 14.2.5.1.2](#) hold:

1. For all diagrams in \mathcal{C} of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xrightarrow{\psi} \end{array} X,$$

if α and β are 2-isomorphisms and we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

2. For each $X \in \text{Obj}(\mathcal{C})$ and each 2-isomorphism

$$\beta: \phi \circ f \xRightarrow{\sim} \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of \mathcal{C} , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of \mathcal{C} such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

14.2.7 Corepresentably Essentially Injective Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.2.7.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **corepresentably essentially injective** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f^*: \text{Hom}_{\mathcal{C}}(B, X) \rightarrow \text{Hom}_{\mathcal{C}}(A, X)$$

given by precomposition by f is essentially injective.

Remark 14.2.7.1.2. In detail, f is corepresentably essentially injective if, for each pair of morphisms $\phi, \psi: B \rightrightarrows X$ of \mathcal{C} , the following condition is satisfied:

(\star) If $\phi \circ f \cong \psi \circ f$, then $\phi \cong \psi$.

14.2.8 Corepresentably Conservative Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.2.8.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **corepresentably conservative** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f^*: \text{Hom}_{\mathcal{C}}(B, X) \rightarrow \text{Hom}_{\mathcal{C}}(A, X)$$

given by precomposition by f is conservative.

Remark 14.2.8.1.2. In detail, f is corepresentably conservative if, for each pair of morphisms $\phi, \psi: B \rightrightarrows X$ and each 2-morphism

$$\alpha: \phi \rightrightarrows \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of \mathcal{C} , if the 2-morphism

$$\alpha \star \text{id}_f: \phi \circ f \rightrightarrows \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \alpha \star \text{id}_f \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

is a 2-isomorphism, then so is α .

14.2.9 Strict Epimorphisms

Let \mathcal{C} be a bicategory.

Definition 14.2.9.1.1. A 1-morphism $f: A \rightarrow B$ is a **strict epimorphism in \mathcal{C}** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f_*: \text{Hom}_{\mathcal{C}}(B, X) \rightarrow \text{Hom}_{\mathcal{C}}(A, X)$$

given by precomposition by f is injective on objects, i.e. its action on objects

$$f_*: \text{Obj}(\text{Hom}_{\mathcal{C}}(B, X)) \rightarrow \text{Obj}(\text{Hom}_{\mathcal{C}}(A, X))$$

is injective.

Remark 14.2.9.1.2. In detail, f is a strict epimorphism if, for each diagram in \mathcal{C} of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \xRightarrow{\psi} \end{array} X,$$

if $\phi \circ f = \psi \circ f$, then $\phi = \psi$.

Example 14.2.9.1.3. Here are some examples of strict epimorphisms.

1. *Strict Epimorphisms in \mathbf{Cats}_2 .* The strict epimorphisms in \mathbf{Cats}_2 are characterised in [Categories](#), [Item 1](#) of [Definition 11.7.3.1.2](#).
2. *Strict Epimorphisms in \mathbf{Rel} .* The strict epimorphisms in \mathbf{Rel} are characterised in [Relations](#), [??](#).

14.2.10 Pseudoepic Morphisms

Let \mathcal{C} be a bicategory.

Definition 14.2.10.1.1. A 1-morphism $f: A \rightarrow B$ of \mathcal{C} is **pseudoepic** if, for each $X \in \text{Obj}(\mathcal{C})$, the functor

$$f^*: \text{Hom}_{\mathcal{C}}(B, X) \rightarrow \text{Hom}_{\mathcal{C}}(A, X)$$

given by precomposition by f is pseudomonic.

Remark 14.2.10.1.2. In detail, a 1-morphism $f: A \rightarrow B$ of \mathcal{C} is pseudoepic if it satisfies the following conditions:

1. For all diagrams in \mathcal{C} of the form

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \Downarrow \beta \\ \xRightarrow{\psi} \end{array} X,$$

if we have

$$\alpha \star \text{id}_f = \beta \star \text{id}_f,$$

then $\alpha = \beta$.

2. For each $X \in \text{Obj}(\mathcal{C})$ and each 2-isomorphism

$$\beta: \phi \circ f \xRightarrow{\sim} \psi \circ f, \quad A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xRightarrow{\psi \circ f} \end{array} X$$

of \mathcal{C} , there exists a 2-isomorphism

$$\alpha: \phi \xRightarrow{\sim} \psi, \quad B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X$$

of \mathcal{C} such that we have an equality

$$A \xrightarrow{f} B \begin{array}{c} \xrightarrow{\phi} \\ \alpha \Downarrow \\ \xrightarrow{\psi} \end{array} X = A \begin{array}{c} \xrightarrow{\phi \circ f} \\ \beta \Downarrow \\ \xrightarrow{\psi \circ f} \end{array} X$$

of pasting diagrams in \mathcal{C} , i.e. such that we have

$$\beta = \alpha \star \text{id}_f.$$

Proposition 14.2.10.1.3. Let $f: A \rightarrow B$ be a 1-morphism of \mathcal{C} .

1. *Characterisations.* The following conditions are equivalent:
 - (a) The morphism f is pseudoepic.
 - (b) The morphism f is corepresentably full on cores and corepresentably faithful.
 - (c) We have an isococcomma square of the form

$$B \stackrel{\text{eq.}}{\cong} B \coprod_A^{\leftrightarrow} B, \quad \begin{array}{ccc} B & \xleftarrow{\text{id}_B} & B \\ \text{id}_B \uparrow & \nearrow \text{dashed} & \uparrow F \\ B & \xleftarrow{F} & A \end{array}$$

in \mathcal{C} up to equivalence.

Proof. *Item 1, Characterisations:* Omitted. □

Appendices

A Other Chapters

Preliminaries

1. [Introduction](#)
2. [A Guide to the Literature](#)

Sets

3. [Sets](#)
4. [Constructions With Sets](#)
5. [Monoidal Structures on the Category of Sets](#)
6. [Pointed Sets](#)
7. [Tensor Products of Pointed Sets](#)

Relations

8. [Relations](#)
9. [Constructions With Relations](#)

10. [Conditions on Relations](#)

Categories

11. [Categories](#)
12. [Presheaves and the Yoneda Lemma](#)

Monoidal Categories

13. [Constructions With Monoidal Categories](#)

Bicategories

14. [Types of Morphisms in Bicategories](#)

Extra Part

15. [Notes](#)

References

- [Adá+01] Jiří Adámek, Robert El Bashir, Manuela Sobral, and Jiří Velebil. “On Functors Which Are Lax Epimorphisms”. In: *Theory Appl. Categ.* 8 (2001), pp. 509–521. ISSN: 1201-561X (cit. on p. [13](#)).