Constructions With Sets

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This chapter develops some material relating to constructions with sets with an eye towards its categorical and higher-categorical counterparts to be introduced later in this work. Of particular interest are perhaps the following:

- 1. Explicit descriptions of the major types of co/limits in Sets, including in particular explicit descriptions of pushouts and coequalisers (see Definitions 4.2.4.1.1, 4.2.4.1.3, 4.2.5.1.1 and 4.2.5.1.3).
- 2. A discussion of powersets as decategorifications of categories of presheaves, including in particular results such as:
 - (a) A discussion of the internal Hom of a powerset (Section 4.4.7).
 - (b) A o-categorical version of the Yoneda lemma (Presheaves and the Yoneda Lemma, Definition 12.1.5.1.1), which we term the Yoneda lemma for sets (Definition 4.5.5.1.1).
 - (c) A characterisation of powersets as free cocompletions (Section 4.4.5), mimicking the corresponding statement for categories of presheaves (??).
 - (d) A characterisation of powersets as free completions (Section 4.4.6), mimicking the corresponding statement for categories of copresheaves (??).
 - (e) A (-1)-categorical version of un/straightening (Item 2 of Definition 4.5.1.1.4 and Definition 4.5.1.1.5).
 - (f) A o-categorical form of Isbell duality internal to powersets (Section 4.4.8).
- 3. A lengthy discussion of the adjoint triple

$$f_! \dashv f^{-1} \dashv f_* \colon \mathcal{P}(A) \xrightarrow{\rightleftarrows} \mathcal{P}(B)$$

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of functors (i.e. morphisms of posets) between $\mathcal{P}(A)$ and $\mathcal{P}(B)$ induced by a map of sets $f:A\to B$, including in particular:

- (a) How f^{-1} can be described as a precomposition while $f_!$ and f_* can be described as Kan extensions (Definitions 4.6.1.1.4, 4.6.2.1.2 and 4.6.3.1.4).
- (b) An extensive list of the properties of $f_!$, f^{-1} , and f_* (Definitions 4.6.1.1.5, 4.6.1.1.6, 4.6.2.1.3, 4.6.2.1.4, 4.6.3.1.7 and 4.6.3.1.8).
- (c) How the functors $f_!$, f^{-1} , f_* , along with the functors

$$-_{1} \cap -_{2} \colon \mathcal{P}(X) \times \mathcal{P}(X) \to \mathcal{P}(X),$$
$$[-_{1}, -_{2}]_{X} \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

may be viewed as a six-functor formalism with the empty set \emptyset as the dualising object (Section 4.6.4).

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4.1 Limits of Sets

4.1.1 The Terminal Set

Definition 4.1.1.1.1. The **terminal set** is the terminal object of Sets as in Limits and Colimits, ??.

Construction 4.1.1.1.2. Concretely, the terminal set is the pair (pt, $\{!_A\}_{A \in \text{Obj}(\mathsf{Sets})}$) consisting of:

- 1. The Limit. The punctual set pt $\stackrel{\text{def}}{=} \{ \star \}$.
- 2. The Cone. The collection of maps

$$\{!_A : A \to \mathsf{pt}\}_{A \in \mathsf{Obj}(\mathsf{Sets})}$$

defined by

$$!_A(a) \stackrel{\text{def}}{=} \star$$

for each $a \in A$ and each $A \in Obj(Sets)$.

Proof. We claim that pt is the terminal object of Sets. Indeed, suppose we have a diagram of the form

$$A$$
 pt

in Sets. Then there exists a unique map $\phi:A\to \operatorname{pt}$ making the diagram

$$A - \frac{\phi}{\exists !} \rightarrow \mathsf{pt}$$

commute, namely $!_A$.

4.1.2 Products of Families of Sets

Let $\{A_i\}_{i\in I}$ be a family of sets.

Definition 4.1.2.1.1. The **product**¹ **of** $\{A_i\}_{i\in I}$ is the product of $\{A_i\}_{i\in I}$ in Sets as in Limits and Colimits, ??.

Construction 4.1.2.1.2. Concretely, the product of $\{A_i\}_{i\in I}$ is the pair $(\prod_{i\in I} A_i, \{pr_i\}_{i\in I})$ consisting of:

¹Further Terminology: Also called the **Cartesian product of** $\{A_i\}_{i\in I}$.

1. The Limit. The set $\prod_{i \in I} A_i$ defined by

$$\prod_{i \in I} A_i \stackrel{\text{def}}{=} \left\{ f \in \operatorname{Sets} \left(I, \bigcup_{i \in I} A_i \right) \middle| \begin{array}{l} \text{for each } i \in I, \text{ we} \\ \text{have } f(i) \in A_i \end{array} \right\}.$$

2. The Cone. The collection

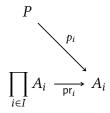
$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

of maps given by

$$\operatorname{pr}_i(f) \stackrel{\text{def}}{=} f(i)$$

for each $f \in \prod_{i \in I} A_i$ and each $i \in I$.

Proof. We claim that $\prod_{i \in I} A_i$ is the categorical product of $\{A_i\}_{i \in I}$ in Sets. Indeed, suppose we have, for each $i \in I$, a diagram of the form



in Sets. Then there exists a unique map $\phi\colon P\to\prod_{i\in I}A_i$ making the diagram

$$P$$

$$\phi \mid \exists ! \qquad p_i$$

$$\prod_{i \in I} A_i \xrightarrow{\mathsf{pr}_i} A_i$$

commute, being uniquely determined by the condition $\operatorname{pr}_i \circ \phi = p_i$ for each $i \in I$ via

$$\phi(x) = (p_i(x))_{i \in I}$$

for each $x \in P$.

Remark 4.1.2.1.3. Less formally, we may think of Cartesian products and projection maps as follows:

- 1. We think of $\prod_{i \in I} A_i$ as the set whose elements are I-indexed collections $(a_i)_{i \in I}$ with $a_i \in A_i$ for each $i \in I$.
- 2. We view the projection maps

$$\left\{ \operatorname{pr}_i \colon \prod_{i \in I} A_i \to A_i \right\}_{i \in I}$$

as being given by

$$\operatorname{pr}_i((a_j)_{j\in I})\stackrel{\text{def}}{=} a_i$$

for each $(a_i)_{i \in I} \in \prod_{i \in I} A_i$ and each $i \in I$.

Proposition 4.1.2.1.4. Let $\{A_i\}_{i\in I}$ be a family of sets.

1. Functoriality. The assignment $\{A_i\}_{i\in I}\mapsto \prod_{i\in I}A_i$ defines a functor

$$\prod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

· Action on Objects. For each $(A_i)_{i \in I} \in \mathsf{Obj}(\mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}))$, we have

$$\left[\prod_{i \in I} \right] ((A_i)_{i \in I}) \stackrel{\text{def}}{=} \prod_{i \in I} A_i$$

· Action on Morphisms. For each $(A_i)_{i \in I}$, $(B_i)_{i \in I} \in \mathsf{Obj}(\mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}))$, the action on Hom-sets

$$\left(\prod_{i\in I}\right)_{(A_i)_{i\in I},(B_i)_{i\in I}}: \operatorname{Nat}((A_i)_{i\in I},(B_i)_{i\in I}) \to \operatorname{Sets}\left(\prod_{i\in I}A_i,\prod_{i\in I}B_i\right)$$

of $\prod_{i \in I}$ at $((A_i)_{i \in I}, (B_i)_{i \in I})$ is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in Nat $((A_i)_{i \in I}, (B_i)_{i \in I})$ to the map of sets

$$\prod_{i\in I} f_i \colon \prod_{i\in I} A_i \to \prod_{i\in I} B_i$$

defined by

$$\left[\prod_{i\in I} f_i\right] ((a_i)_{i\in I}) \stackrel{\text{def}}{=} (f_i(a_i))_{i\in I}$$

for each $(a_i)_{i \in I} \in \prod_{i \in I} A_i$.

Proof. Item 1, *Functoriality*: This follows from Limits and Colimits, ?? of ??.

4.1.3 Binary Products of Sets

Let A and B be sets.

Definition 4.1.3.1.1. The **product of** A **and** B^2 is the product of A and B in Sets as in Limits and Colimits, ??.

Construction 4.1.3.1.2. Concretely, the product of A and B is the pair $(A \times B, \{pr_1, pr_2\})$ consisting of:

1. The Limit. The set $A \times B$ defined by

$$A \times B \stackrel{\text{def}}{=} \prod_{z \in \{A,B\}} z$$

$$\stackrel{\text{def}}{=} \{ f \in \mathsf{Sets}(\{0,1\}, A \cup B) \mid \mathsf{we have } f(0) \in A \, \mathsf{and} \, f(1) \in B \}$$

$$\cong \{ \{ \{a\}, \{a,b\} \} \in \mathcal{P}(\mathcal{P}(A \cup B)) \mid \mathsf{we have } a \in A \, \mathsf{and} \, b \in B \}$$

$$\cong \begin{cases} \mathsf{ordered pairs} \, (a,b) \, \mathsf{with} \\ a \in A \, \mathsf{and} \, b \in B \end{cases}.$$

2. The Cone. The maps

$$\operatorname{pr}_1 \colon A \times B \to A,$$

 $\operatorname{pr}_2 \colon A \times B \to B$

defined by

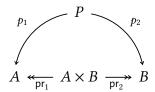
$$\operatorname{pr}_{1}(a, b) \stackrel{\text{def}}{=} a,$$

 $\operatorname{pr}_{2}(a, b) \stackrel{\text{def}}{=} b$

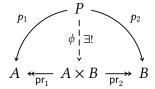
for each $(a, b) \in A \times B$.

² Further Terminology: Also called the **Cartesian product of** A **and** B.

Proof. We claim that $A \times B$ is the categorical product of A and B in the category of sets. Indeed, suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: P \to A \times B$ making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$

 $\operatorname{pr}_2 \circ \phi = p_2$

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each $x \in P$.

Proposition 4.1.3.1.3. Let *A*, *B*, *C*, and *X* be sets.

1. Functoriality. The assignments $A, B, (A, B) \mapsto A \times B$ define functors

$$A \times -:$$
 Sets \rightarrow Sets,
 $- \times B:$ Sets \rightarrow Sets,
 $-_1 \times -_2:$ Sets \times Sets \rightarrow Sets,

where -1×-2 is the functor where

· Action on Objects. For each $(A, B) \in \mathsf{Obj}(\mathsf{Sets} \times \mathsf{Sets})$, we have

$$[-_1 \times -_2](A, B) \stackrel{\text{def}}{=} A \times B.$$

· Action on Morphisms. For each (A, B), $(X, Y) \in Obj(Sets)$, the action on Hom-sets

$$\times_{(A,B),(X,Y)}$$
: Sets $(A,X) \times$ Sets $(B,Y) \to$ Sets $(A \times B, X \times Y)$
of \times at $((A,B),(X,Y))$ is defined by sending (f,g) to the function $f \times g \colon A \times B \to X \times Y$

defined by

$$[f \times g](a,b) \stackrel{\text{def}}{=} (f(a),g(b))$$

for each $(a, b) \in A \times B$.

and where $A \times -$ and $- \times B$ are the partial functors of $-_1 \times -_2$ at $A, B \in Obj(Sets)$.

2. *Adjointness I.* We have adjunctions

$$(A \times - + \operatorname{Sets}(A, -))$$
: Sets $\underbrace{A \times -}_{\operatorname{Sets}(A, -)}$ Sets, $\underbrace{- \times B}_{\operatorname{Sets}(B, -)}$ Sets, Sets, Sets,

witnessed by bijections

$$\mathsf{Sets}(A \times B, C) \cong \mathsf{Sets}(A, \mathsf{Sets}(B, C)),$$

 $\mathsf{Sets}(A \times B, C) \cong \mathsf{Sets}(B, \mathsf{Sets}(A, C)),$

natural in $A, B, C \in \mathsf{Obj}(\mathsf{Sets})$.

3. Adjointness II. We have an adjunction

$$(\Delta_{\mathsf{Sets}} \dashv -_1 \times -_2)$$
: $\mathsf{Sets} \underbrace{\perp}_{-_1 \times -_2} \mathsf{Sets} \times \mathsf{Sets},$

witnessed by a bijection

$$\mathsf{Hom}_{\mathsf{Sets} \times \mathsf{Sets}}((A, A), (B, C)) \cong \mathsf{Sets}(A, B \times C),$$

natural in $A \in \mathsf{Obj}(\mathsf{Sets})$ and in $(B, C) \in \mathsf{Obj}(\mathsf{Sets} \times \mathsf{Sets}).$

4. Associativity. We have an isomorphism of sets

$$\alpha_{ABC}^{\mathsf{Sets}} \colon (A \times B) \times C \xrightarrow{\sim} A \times (B \times C),$$

natural in $A, B, C \in \mathsf{Obj}(\mathsf{Sets})$.

5. Unitality. We have isomorphisms of sets

$$\lambda_A^{\mathsf{Sets}} \colon \mathsf{pt} \times A \xrightarrow{\sim} A,$$

 $\rho_A^{\mathsf{Sets}} \colon A \times \mathsf{pt} \xrightarrow{\sim} A,$

natural in $A \in Obj(Sets)$.

6. Commutativity. We have an isomorphism of sets

$$\sigma_{A,B}^{\mathsf{Sets}} : A \times B \xrightarrow{\sim} B \times A,$$

natural in $A, B \in Obj(Sets)$.

7. Distributivity Over Coproducts. We have isomorphisms of sets

$$\begin{array}{l} \delta_{\ell}^{\mathsf{Sets}} \colon A \times (B \coprod C) \xrightarrow{\sim} (A \times B) \coprod (A \times C), \\ \delta_{r}^{\mathsf{Sets}} \colon (A \coprod B) \times C \xrightarrow{\sim} (A \times C) \coprod (B \times C), \end{array}$$

natural in $A, B, C \in Obj(Sets)$.

8. Annihilation With the Empty Set. We have isomorphisms of sets

$$\zeta_{\ell}^{\mathsf{Sets}} \colon \emptyset \times A \xrightarrow{\sim} \emptyset,$$

 $\zeta_{r}^{\mathsf{Sets}} \colon A \times \emptyset \xrightarrow{\sim} \emptyset,$

natural in $A \in Obj(Sets)$.

9. Distributivity Over Unions. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \cup W) = (U \times V) \cup (U \times W),$$

$$(U \cup V) \times W = (U \times W) \cup (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

10. Distributivity Over Intersections. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \cap W) = (U \times V) \cap (U \times W),$$

$$(U \cap V) \times W = (U \times W) \cap (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

11. Distributivity Over Differences. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \setminus W) = (U \times V) \setminus (U \times W),$$

$$(U \setminus V) \times W = (U \times W) \setminus (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

12. Distributivity Over Symmetric Differences. Let X be a set. For each $U, V, W \in \mathcal{P}(X)$, we have equalities

$$U \times (V \triangle W) = (U \times V) \triangle (U \times W),$$

$$(U \triangle V) \times W = (U \times W) \triangle (V \times W)$$

of subsets of $\mathcal{P}(X \times X)$.

13. Middle-Four Exchange with Respect to Intersections. The diagram

$$(\mathcal{P}(X) \times \mathcal{P}(X)) \times (\mathcal{P}(X) \times \mathcal{P}(X)) \xrightarrow{\cap \times \cap} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow^{\mathcal{P}_{X,X}^{\times}} \times \mathcal{P}_{X,X}^{\times} \downarrow \qquad \qquad \downarrow^{\mathcal{P}_{X,X}^{\times}}$$

$$\mathcal{P}(X \times X) \times \mathcal{P}(X \times X) \xrightarrow{\cap} \mathcal{P}(X \times X)$$

commutes, i.e. we have

$$(U \times V) \cap (W \times T) = (U \cap V) \times (W \cap T).$$

for each $U, V, W, T \in \mathcal{P}(X)$.

14. Symmetric Monoidality. The 8-tuple (Sets, \times , pt, Sets(-1, -2), α^{Sets} , λ^{Sets} , ρ^{Sets} , σ^{Sets}) is a closed symmetric monoidal category.

15. Symmetric Bimonoidality. The 18-tuple

$$\begin{split} & \Big(\mathsf{Sets}, \coprod, \times, \emptyset, \mathsf{pt}, \mathsf{Sets}(-_1, -_2), \alpha^{\mathsf{Sets}}, \lambda^{\mathsf{Sets}}, \rho^{\mathsf{Sets}}, \sigma^{\mathsf{Sets}}, \\ & \alpha^{\mathsf{Sets}, \coprod}, \lambda^{\mathsf{Sets}, \coprod}, \rho^{\mathsf{Sets}, \coprod}, \sigma^{\mathsf{Sets}, \coprod}, \delta^{\mathsf{Sets}}_{\ell}, \delta^{\mathsf{Sets}}_{r}, \zeta^{\mathsf{Sets}}_{\ell}, \zeta^{\mathsf{Sets}}_{r} \Big), \end{split}$$

is a symmetric closed bimonoidal category, where $\alpha^{\text{Sets},\coprod}$, $\lambda^{\text{Sets},\coprod}$, $\rho^{\text{Sets},\coprod}$, and $\sigma^{\text{Sets},\coprod}$ are the natural transformations from Items 3 to 5 of Definition 4.2.3.1.3.

Proof. Item 1, Functoriality: This follows from Limits and Colimits, ?? of ??. Item 2, Adjointness: We prove only that there,s an adjunction $- \times B \dashv \mathsf{Sets}(B, -)$, witnessed by a bijection

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$

natural in $B, C \in \mathsf{Obj}(\mathsf{Sets})$, as the proof of the existence of the adjunction $A \times - \dashv \mathsf{Sets}(A, -)$ follows almost exactly in the same way.

· Map I. We define a map

$$\Phi_{BC}$$
: Sets $(A \times B, C) \rightarrow \text{Sets}(A, \text{Sets}(B, C))$.

by sending a function

$$\xi: A \times B \to C$$

to the function

$$\xi^{\dagger} : A \longrightarrow \mathsf{Sets}(B, C),$$

 $a \mapsto (\xi_a^{\dagger} : B \to C),$

where we define

$$\xi_a^{\dagger}(b) \stackrel{\text{def}}{=} \xi(a,b)$$

for each $b \in B$. In terms of the $[a \mapsto f(a)]$ notation of Sets, Definition 3.1.1.1.2, we have

$$\boldsymbol{\xi}^{\dagger} \stackrel{\mathrm{def}}{=} \llbracket \boldsymbol{a} \mapsto \llbracket \boldsymbol{b} \mapsto \boldsymbol{\xi}(\boldsymbol{a}, \boldsymbol{b}) \rrbracket \rrbracket.$$

· Map II. We define a map

$$\Psi_{BC}$$
: Sets(A, Sets(B, C)), \rightarrow Sets(A \times B, C)

given by sending a function

$$\xi: A \longrightarrow \mathsf{Sets}(B, C),$$

 $a \mapsto (\xi_a: B \to C),$

to the function

$$\xi^{\dagger}: A \times B \to C$$

defined by

$$\xi^{\dagger}(a,b) \stackrel{\text{def}}{=} \operatorname{ev}_b(\operatorname{ev}_a(\xi))$$

$$\stackrel{\text{def}}{=} \operatorname{ev}_b(\xi_a)$$

$$\stackrel{\text{def}}{=} \xi_a(b)$$

for each $(a, b) \in A \times B$.

· Invertibility I. We claim that

$$\Psi_{A,B} \circ \Phi_{A,B} = \mathrm{id}_{\mathsf{Sets}(A \times B,C)}$$
.

Indeed, given a function $\xi: A \times B \to C$, we have

$$\begin{split} \big[\Psi_{A,B} \circ \Phi_{A,B} \big] (\xi) &= \Psi_{A,B} (\Phi_{A,B} (\xi)) \\ &= \Psi_{A,B} (\Phi_{A,B} (\big[\big[(a,b) \mapsto \xi(a,b) \big] \big])) \\ &= \Psi_{A,B} (\big[\big[a \mapsto \big[\big[b \mapsto \xi(a,b) \big] \big] \big]) \\ &= \Psi_{A,B} (\big[\big[a' \mapsto \big[\big[b' \mapsto \xi(a',b') \big] \big] \big])) \\ &= \big[\big[(a,b) \mapsto \operatorname{ev}_b (\operatorname{ev}_a (\big[\big[a' \mapsto \big[b' \mapsto \xi(a',b') \big] \big] \big])) \big] \big] \\ &= \big[\big[(a,b) \mapsto \operatorname{ev}_b (\big[\big[b' \mapsto \xi(a,b') \big] \big] \big] \big] \big] \\ &= \big[\big[(a,b) \mapsto \xi(a,b) \big] \big] \\ &= \xi. \end{split}$$

· Invertibility II. We claim that

$$\Phi_{A,B} \circ \Psi_{A,B} = \mathsf{id}_{\mathsf{Sets}(A,\mathsf{Sets}(B,C))}$$
.

Indeed, given a function

$$\xi \colon A \longrightarrow \mathsf{Sets}(B,C),$$

 $a \mapsto (\xi_a \colon B \to C),$

we have

$$\begin{split} [\Phi_{A,B} \circ \Psi_{A,B}](\xi) &\stackrel{\text{def}}{=} \Phi_{A,B}(\Psi_{A,B}(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{A,B}(\llbracket(a,b) \mapsto \xi_a(b)\rrbracket) \\ &\stackrel{\text{def}}{=} \Phi_{A,B}(\llbracket(a',b') \mapsto \xi_{a'}(b')\rrbracket) \\ &\stackrel{\text{def}}{=} \llbracket a \mapsto \llbracket b \mapsto \operatorname{ev}_{(a,b)}(\llbracket(a',b') \mapsto \xi_{a'}(b')\rrbracket)\rrbracket \rrbracket \\ &\stackrel{\text{def}}{=} \llbracket a \mapsto \llbracket b \mapsto \xi_a(b)\rrbracket \rrbracket \rrbracket \\ &\stackrel{\text{def}}{=} \llbracket a \mapsto \xi_a \rrbracket \\ &\stackrel{\text{def}}{=} \xi. \end{split}$$

· Naturality for Φ , Part I. We need to show that, given a function $g\colon B\to B'$, the diagram

$$\begin{split} \mathsf{Sets}(A \times B', C) & \xrightarrow{\Phi_{B', C}} & \mathsf{Sets}(A, \mathsf{Sets}(B', C)), \\ & \mathsf{id}_A \times g^* \bigg| & & \bigg| (g^*)_! \\ & \mathsf{Sets}(A \times B, C) & \xrightarrow{\Phi_{B, C}} & \mathsf{Sets}(A, \mathsf{Sets}(B, C)) \end{split}$$

commutes. Indeed, given a function

$$\xi \colon A \times B' \to C$$

we have

$$\begin{aligned} [\Phi_{B,C} \circ (\mathsf{id}_A \times g^*)](\xi) &= \Phi_{B,C}([\mathsf{id}_A \times g^*](\xi)) \\ &= \Phi_{B,C}(\xi(-_1, g(-_2))) \\ &= [\xi(-_1, g(-_2))]^{\dagger} \\ &= \xi^{\dagger}_{-_1}(g(-_2)) \\ &= (g^*)_!(\xi^{\dagger}) \\ &= (g^*)_!(\Phi_{B',C}(\xi)) \end{aligned}$$

$$= [(q^*)_! \circ \Phi_{B',C}](\xi).$$

Alternatively, using the $[\![a\mapsto f(a)]\!]$ notation of Sets, Definition 3.1.1.1.2, we have

$$\begin{split} [\Phi_{B,C} \circ (\mathrm{id}_A \times g^*)](\xi) &= \Phi_{B,C}([\mathrm{id}_A \times g^*](\xi)) \\ &= \Phi_{B,C}([\mathrm{id}_A \times g^*]([\![(a,b') \mapsto \xi(a,b')]\!])) \\ &= \Phi_{B,C}([\![(a,b) \mapsto \xi(a,g(b))]\!]) \\ &= [\![a \mapsto [\![b \mapsto \xi(a,g(b))]\!]]] \\ &= [\![a \mapsto g^*([\![b' \mapsto \xi(a,b')]\!])]] \\ &= (g^*)_!([\![a \mapsto [\![b' \mapsto \xi(a,b')]\!]])) \\ &= (g^*)_!(\Phi_{B',C}([\![(a,b') \mapsto \xi(a,b')]\!])) \\ &= (g^*)_!(\Phi_{B',C}(\xi)) \\ &= [(g^*)_! \circ \Phi_{B',C}](\xi). \end{split}$$

· Naturality for Φ , Part II. We need to show that, given a function $h\colon C\to C'$, the diagram

$$\begin{split} \mathsf{Sets}(A \times B, C) & \xrightarrow{\Phi_{B,C}} & \mathsf{Sets}(A, \mathsf{Sets}(B, C)), \\ h_! & & \downarrow^{(h_!)_!} \\ \mathsf{Sets}(A \times B, C') & \xrightarrow{\Phi_{B,C'}} & \mathsf{Sets}(A, \mathsf{Sets}(B, C')) \end{split}$$

commutes. Indeed, given a function

$$\xi: A \times B \to C$$
,

we have

$$\begin{split} [\Phi_{B,C} \circ h_{!}](\xi) &= \Phi_{B,C}(h_{!}(\xi)) \\ &= \Phi_{B,C}(h_{!}([(a,b) \mapsto \xi(a,b)])) \\ &= \Phi_{B,C}([(a,b) \mapsto h(\xi(a,b))]) \\ &= [(a \mapsto [(b \mapsto h(\xi(a,b))]])] \\ &= [(a \mapsto h_{!}([(b \mapsto \xi(a,b))]])) \\ &= (h_{!})_{!}([(a \mapsto (b \mapsto \xi(a,b))]])) \end{split}$$

$$= (h_!)_! (\Phi_{B,C}(\llbracket (a,b) \mapsto \xi(a,b) \rrbracket))$$

= $(h_!)_! (\Phi_{B,C}(\xi))$
= $[(h_!)_! \circ \Phi_{B,C}](\xi)$.

• Naturality for Ψ . Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Ψ is also natural in each argument.

This finishes the proof.

Item 3, Adjointness II: This follows from the universal property of the product.

Item 4, Associativity: This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.1.4.1.1.

Item 5, Unitality: This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.1.5.1.1 and 5.1.6.1.1.

Item 6, Commutativity: This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.1.7.1.1.

Item 7, *Distributivity Over Coproducts*: This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.3.1.1.1 and 5.3.2.1.1.

Item 8, *Annihilation With the Empty Set*: This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.3.3.1.1 and 5.3.4.1.1.

Item 9, Distributivity Over Unions: See [Pro25c].

Item 10, *Distributivity Over Intersections*: See [Pro25d, Corollary 1].

Item 11, Distributivity Over Differences: See [Pro25a].

Item 12, Distributivity Over Symmetric Differences: See [Pro25b].

Item 13, Middle-Four Exchange With Respect to Intersections: See [Pro25d, Corollary 1].

Item 14, *Symmetric Monoidality*: This is a repetition of Monoidal Structures on the Category of Sets, Definition 5.1.9.1.1, and is proved there.

Item 15, Symmetric Bimonoidality: This is a repetition of Monoidal Structures on the Category of Sets, Definition 5.3.5.1.1, and is proved there.

Remark 4.1.3.1.4. As shown in Item 1 of Definition 4.1.3.1.3, the Cartesian product of sets defines a functor

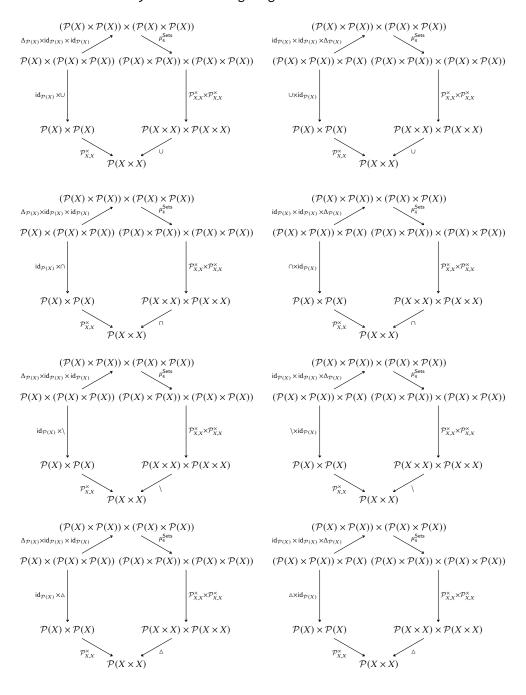
$$-1 \times -2$$
: Sets \times Sets \rightarrow Sets.

This functor is the $(k, \ell) = (-1, -1)$ case of a family of functors

$$\otimes_{k,\ell} \colon \mathsf{Mon}_{\mathbb{E}_k}(\mathsf{Sets}) \times \mathsf{Mon}_{\mathbb{E}_\ell}(\mathsf{Sets}) \to \mathsf{Mon}_{\mathbb{E}_{k+\ell}}(\mathsf{Sets})$$

of tensor products of \mathbb{E}_k -monoid objects on Sets with \mathbb{E}_ℓ -monoid objects on Sets; see ??.

Remark 4.1.3.1.5. We may state the equalities in Items 9 to 12 of Definition 4.1.3.1.3 as the commutativity of the following diagrams:



4.1.4 Pullbacks

Let A, B, and C be sets and let $f: A \to C$ and $g: B \to C$ be functions.

Definition 4.1.4.1.1. The **pullback of** A **and** B **over** C **along** f **and** g³ is the pullback of A and B over C along f and g in Sets as in Limits and Colimits, ??.

Construction 4.1.4.1.2. Concretely, the pullback of A and B over C along f and g is the pair $(A \times_C B, \{pr_1, pr_2\})$ consisting of:

1. The Limit. The set $A \times_C B$ defined by

$$A \times_C B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid f(a) = g(b)\}.$$

2. The Cone. The maps⁴

$$\operatorname{pr}_1 : A \times_C B \to A,$$

 $\operatorname{pr}_2 : A \times_C B \to B$

defined by

$$\operatorname{pr}_{1}(a, b) \stackrel{\text{def}}{=} a,$$

 $\operatorname{pr}_{2}(a, b) \stackrel{\text{def}}{=} b$

for each $(a, b) \in A \times_C B$.

Proof. We claim that $A \times_C B$ is the categorical pullback of A and B over C with respect to (f,g) in Sets. First we need to check that the relevant pullback diagram commutes, i.e. that we have

$$f \circ \operatorname{pr}_1 = g \circ \operatorname{pr}_2, \qquad A \times_C B \xrightarrow{\operatorname{pr}_2} B \\ \downarrow^g \\ A \xrightarrow{f} C.$$

Indeed, given $(a, b) \in A \times_C B$, we have

$$[f \circ \mathsf{pr}_1](a,b) = f(\mathsf{pr}_1(a,b))$$

³Further Terminology: Also called the **fibre product of** A **and** B **over** C **along** f **and** g.

⁴Further Notation: Also written $pr_1^{A \times_C B}$ and $pr_2^{A \times_C B}$.

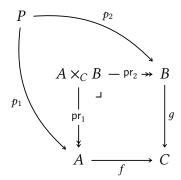
$$= f(a)$$

$$= g(b)$$

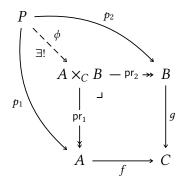
$$= g(pr_2(a, b))$$

$$= [g \circ pr_2](a, b),$$

where f(a) = g(b) since $(a, b) \in A \times_C B$. Next, we prove that $A \times_C B$ satisfies the universal property of the pullback. Suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi: P \to A \times_C B$ making the diagram



commute, being uniquely determined by the conditions

$$\operatorname{pr}_1 \circ \phi = p_1,$$

 $\operatorname{pr}_2 \circ \phi = p_2$

via

$$\phi(x) = (p_1(x), p_2(x))$$

for each $x \in P$, where we note that $(p_1(x), p_2(x)) \in A \times B$ indeed lies in $A \times_C B$ by the condition

$$f\circ p_1=g\circ p_2,$$

which gives

$$f(p_1(x)) = g(p_2(x))$$

for each $x \in P$, so that $(p_1(x), p_2(x)) \in A \times_C B$.

Remark 4.1.4.1.3. It is common practice to write $A \times_C B$ for the pullback of A and B over C along f and g, omitting the maps f and g from the notation and instead leaving them implicit, to be understood from the context.

However, the set $A \times_C B$ depends very much on the maps f and g, and sometimes it is necessary or useful to note this dependence explicitly. In such situations, we will write $A \times_{f,C,g} B$ or $A \times_C^{f,g} B$ for $A \times_C B$.

Example 4.1.4.1.4. Here are some examples of pullbacks of sets.

1. Unions via Intersections. Let X be a set. We have

$$A \cap B \cong A \times_{A \cup B} B, \qquad A \cap B \xrightarrow{J} B$$

$$\downarrow \qquad \downarrow \qquad \downarrow \iota_{B}$$

$$A \xrightarrow{\iota_{A}} A \cup B$$

for each $A, B \in \mathcal{P}(X)$.

Proof. Item 1, Unions via Intersections: Indeed, we have

$$A \times_{A \cup B} B \cong \{(x, y) \in A \times B \mid x = y\}$$

 $\cong A \cap B.$

This finishes the proof.

Proposition 4.1.4.1.5. Let *A*, *B*, *C*, and *X* be sets.

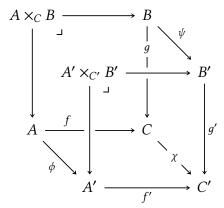
1. Functoriality. The assignment $(A, B, C, f, g) \mapsto A \times_{f,C,g} B$ defines a functor

$$-_1 \times_{-_3} -_1 : \operatorname{\mathsf{Fun}}(\mathcal{P}, \operatorname{\mathsf{Sets}}) \to \operatorname{\mathsf{Sets}},$$

where \mathcal{P} is the category that looks like this:



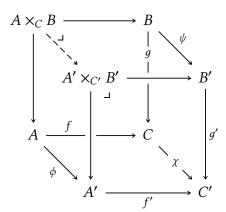
In particular, the action on morphisms of $-1 \times_{-3} -1$ is given by sending a morphism



in Fun(\mathcal{P} , Sets) to the map $\xi \colon A \times_C B \xrightarrow{\exists !} A' \times_{C'} B'$ given by

$$\xi(a,b) \stackrel{\text{def}}{=} (\phi(a), \psi(b))$$

for each $(a, b) \in A \times_C B$, which is the unique map making the diagram



commute.

2. *Adjointness I.* We have adjunctions

$$(A \times_X - \dashv \mathbf{Sets}_{/X}(A, -)) : \operatorname{Sets}_{/X} \xrightarrow{A \times_X -} \operatorname{Sets}_{/X},$$

$$(- \times_X B \dashv \mathbf{Sets}_{/X}(B, -)) : \operatorname{Sets}_{/X} \xrightarrow{\bot} \operatorname{Sets}_{/X},$$

$$\operatorname{Sets}_{/X}(B, -)$$

witnessed by bijections

$$\mathsf{Sets}_{/X}(A \times_X B, C) \cong \mathsf{Sets}_{/X}(A, \mathbf{Sets}_{/X}(B, C)),$$

 $\mathsf{Sets}_{/X}(A \times_X B, C) \cong \mathsf{Sets}_{/X}(B, \mathbf{Sets}_{/X}(A, C)),$

natural in (A, ϕ_A) , (B, ϕ_B) , $(C, \phi_C) \in \mathsf{Obj}(\mathsf{Sets}_{/X})$, where $\mathsf{Sets}_{/X}(A, B)$ is the object of $\mathsf{Sets}_{/X}$ consisting of (see Fibred Sets, ??):

· The Set. The set $\mathbf{Sets}_{/X}(A, B)$ defined by

$$\mathbf{Sets}_{/X}(A,B) \stackrel{\text{def}}{=} \coprod_{x \in X} \mathsf{Sets}(\phi_A^{-1}(x), \phi_Y^{-1}(x))$$

· The Map to X. The map

$$\phi_{\mathsf{Sets}_{/X}(A,B)} \colon \mathsf{Sets}_{/X}(A,B) \to X$$

defined by

$$\phi_{\mathbf{Sets}_{/X}(A,B)}(x,f) \stackrel{\text{def}}{=} x$$

for each
$$(x, f) \in \mathbf{Sets}_{/X}(A, B)$$
.

3. Adjointness II. We have an adjunction

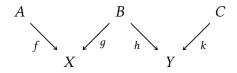
$$\left(\Delta_{\mathsf{Sets}_{/X}} \dashv -_1 \times -_2\right) \colon \mathsf{Sets}_{/X} \underbrace{\downarrow}_{-_1 \times -_2} \mathsf{Sets}_{/X} \times \mathsf{Sets}_{/X},$$

witnessed by a bijection

$$\mathsf{Hom}_{\mathsf{Sets}/_X} \times \mathsf{Sets}/_X ((A, A), (B, C)) \cong \mathsf{Sets}/_X (A, B \times_X C),$$

natural in $A \in \mathsf{Obj}(\mathsf{Sets}_{/X})$ and in $(B, C) \in \mathsf{Obj}(\mathsf{Sets}_{/X} \times \mathsf{Sets}_{/X})$.

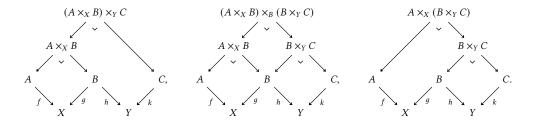
4. Associativity. Given a diagram



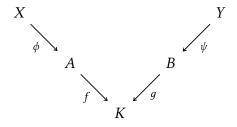
in Sets, we have isomorphisms of sets

$$(A \times_X B) \times_Y C \cong (A \times_X B) \times_B (B \times_Y C) \cong A \times_X (B \times_Y C),$$

where these pullbacks are built as in the diagrams



5. Interaction With Composition. Given a diagram



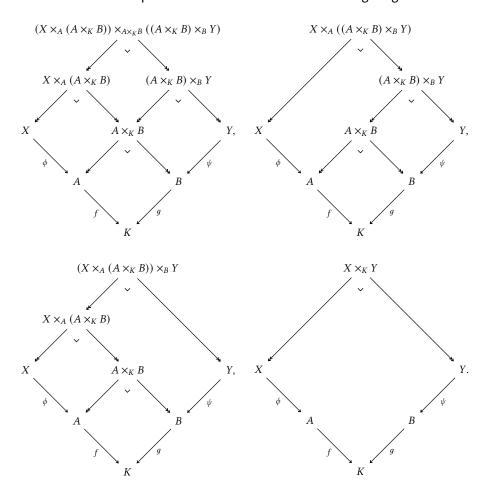
in Sets, we have isomorphisms of sets

$$\begin{split} X \times_K^{f \circ \phi, g \circ \psi} Y &\cong (X \times_A^{\phi, q_1} (A \times_K^{f, g} B)) \times_{A \times_K^{f, g} B}^{p_2, p_1} ((A \times_K^{f, g} B) \times_B^{q_2, \psi} Y) \\ &\cong X \times_A^{\phi, p} ((A \times_K^{f, g} B) \times_B^{q_2, \psi} Y) \\ &\cong (X \times_A^{\phi, q_1} (A \times_K^{f, g} B)) \times_B^{q, \psi} Y \end{split}$$

where

$$\begin{aligned} q_1 &= \mathsf{pr}_1^{A \times_K^{f,g} B}, & q_2 &= \mathsf{pr}_2^{A \times_K^{f,g} B}, \\ p_1 &= \mathsf{pr}_1^{(A \times_K^{f,g} B) \times_Y^{q_2, \psi}}, & X \times_{A \times_K^{f,g} B}^{\phi, q_1} (A \times_K^{f,g} B)} \\ p &= q_1 \circ \mathsf{pr}_1^{(A \times_K^{f,g} B) \times_B^{q_2, \psi} Y}, & q &= q_2 \circ \mathsf{pr}_2^{X \times_A^{\phi, q_1} (A \times_K^{f,g} B)}, \end{aligned}$$

and where these pullbacks are built as in the following diagrams:



6. Unitality. We have isomorphisms of sets

natural in $(A, f) \in \mathsf{Obj}(\mathsf{Sets}_{/X})$.

7. Commutativity. We have an isomorphism of sets

natural in (A, f), $(B, g) \in \mathsf{Obj}(\mathsf{Sets}_{/X})$.

8. Distributivity Over Coproducts. Let A, B, and C be sets and let $\phi_A \colon A \to X$, $\phi_B \colon B \to X$, and $\phi_C \colon C \to X$ be morphisms of sets. We have isomorphisms of sets

$$\delta_{\ell}^{\mathsf{Sets}_{/X}} : A \times_{X} (B \coprod C) \xrightarrow{\sim} (A \times_{X} B) \coprod (A \times_{X} C),$$

$$\delta_{r}^{\mathsf{Sets}_{/X}} : (A \coprod B) \times_{X} C \xrightarrow{\sim} (A \times_{X} C) \coprod (B \times_{X} C),$$

as in the diagrams

natural in $A, B, C \in \mathsf{Obj}(\mathsf{Sets}_{/X})$.

9. Annihilation With the Empty Set. We have isomorphisms of sets

$$\emptyset \longrightarrow \emptyset \qquad \qquad \emptyset \longrightarrow A$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \zeta_{\ell}^{\mathsf{Sets}/X} : A \times_{X} \emptyset \xrightarrow{\sim} \emptyset, \qquad \qquad \downarrow \qquad \downarrow f$$

$$A \longrightarrow f \qquad X, \qquad \qquad \emptyset \longrightarrow X,$$

natural in $(A, f) \in \mathsf{Obj}(\mathsf{Sets}_{/X})$.

10. Interaction With Products. We have an isomorphism of sets

$$A \times_{\mathsf{pt}} B \cong A \times B, \qquad A \times_{\mathsf{pt}} B \cong A \times B, \qquad A \xrightarrow{!_{A}} \mathsf{pt}.$$

Proof. Item 1, Functoriality: This is a special case of functoriality of co/limits, Limits and Colimits, ?? of ??, with the explicit expression for ξ following from the commutativity of the cube pullback diagram.

Item 2, Adjointness I: This is a repetition of Fibred Sets, ?? of ??, and is proved there. Item 3, Adjointness II: This follows from the universal property of the product (pullbacks are products in $\mathsf{Sets}_{/X}$).

Item 4, Associativity: We have

$$(A \times_X B) \times_Y C \cong \{((a,b),c) \in (A \times_X B) \times C \mid h(b) = k(c)\}$$

$$\cong \{((a,b),c) \in (A \times B) \times C \mid f(a) = g(b) \text{ and } h(b) = k(c)\}$$

$$\cong \{(a,(b,c)) \in A \times (B \times C) \mid f(a) = g(b) \text{ and } h(b) = k(c)\}$$

$$\cong \{(a,(b,c)) \in A \times (B \times_Y C) \mid f(a) = g(b)\}$$

$$\cong A \times_X (B \times_Y C)$$

and

$$(A \times_X B) \times_B (B \times_Y C) \cong \{((a,b),(b',c)) \in (A \times_X B) \times (B \times_Y C) \mid b = b'\}$$

$$\cong \left\{ ((a,b),(b',c)) \in (A \times B) \times (B \times C) \middle| \begin{array}{c} f(a) = g(b), b = b', \\ \text{and } h(b') = k(c) \end{array} \right\}$$

$$\cong \left\{ (a,(b,(b',c))) \in A \times (B \times (B \times C)) \middle| \begin{array}{c} f(a) = g(b), b = b', \\ \text{and } h(b') = k(c) \end{array} \right\}$$

$$\cong \left\{ (a,((b,b'),c)) \in A \times ((B \times B) \times C) \middle| \begin{array}{c} f(a) = g(b), b = b', \\ \text{and } h(b') = k(c) \end{array} \right\}$$

$$\cong \left\{ (a,((b,b'),c)) \in A \times ((B \times_B B) \times C) \middle| \begin{array}{c} f(a) = g(b) \text{ and } \\ h(b') = k(c) \end{array} \right\}$$

$$\cong \left\{ (a,((b,b'),c)) \in A \times ((B \times_B B) \times C) \middle| \begin{array}{c} f(a) = g(b) \text{ and } \\ h(b') = k(c) \end{array} \right\}$$

$$\cong \left\{ (a,(b,c)) \in A \times (B \times C) \middle| f(a) = g(b) \text{ and } h(b) = k(c) \right\}$$

$$\cong A \times_X (B \times_Y C),$$

where we have used Item 6 for the isomorphism $B \times_B B \cong B$.

Item 5, Interaction With Composition: By Item 4, it suffices to construct only the isomorphism

$$X\times_K^{f\circ\phi,g\circ\psi}Y\cong (X\times_A^{\phi,q_1}(A\times_K^{f,g}B))\times_{A\times_K^{f,g}B}^{p_2,p_1}((A\times_K^{f,g}B)\times_B^{q_2,\psi}Y).$$

We have

$$(X \times_A^{\phi,q_1} (A \times_K^{f,g} B)) \stackrel{\text{def}}{=} \left\{ (x,(a,b)) \in X \times (A \times_K^{f,g} B) \,\middle|\, \phi(x) = q_1(a,b) \right\}$$

$$\stackrel{\text{def}}{=} \left\{ (x, (a, b)) \in X \times (A \times_K^{f, g} B) \,\middle|\, \phi(x) = a \right\}$$

$$\cong \left\{ (x, (a, b)) \in X \times (A \times B) \,\middle|\, \phi(x) = a \text{ and } f(a) = g(b) \right\},$$

$$((A \times_K^{f, g} B) \times_B^{q_2, \psi} Y) \stackrel{\text{def}}{=} \left\{ ((a, b), y) \in (A \times_K^{f, g} B) \times Y \,\middle|\, q_2(a, b) = \psi(y) \right\}$$

$$\stackrel{\text{def}}{=} \left\{ ((a, b), y) \in (A \times_K^{f, g} B) \times Y \,\middle|\, b = \psi(y) \right\}$$

$$\cong \left\{ ((a, b), y) \in (A \times B) \times Y \,\middle|\, b = \psi(y) \text{ and } f(a) = g(b) \right\},$$

so writing

$$S = (X \times_A^{\phi, q_1} (A \times_K^{f, g} B))$$

$$S' = ((A \times_K^{f, g} B) \times_B^{q_2, \psi} Y).$$

we have

$$\begin{split} S \times_{A \times_{K}^{f,g} B}^{p_{2},p_{1}} S' &\stackrel{\text{def}}{=} \{ ((x,(a,b)), ((a',b'),y)) \in S \times S' \mid p_{1}(x,(a,b)) = p_{2}((a',b'),y) \} \\ &\stackrel{\text{def}}{=} \{ ((x,(a,b)), ((a',b'),y)) \in S \times S' \mid (a,b) = (a',b') \} \\ &\cong \{ ((x,a,b,y)) \in X \times A \times B \times Y \mid \phi(x) = a, \psi(y) = b, \text{ and } f(a) = g(b) \} \\ &\stackrel{\text{def}}{=} \{ ((x,a,b,y)) \in X \times A \times B \times Y \mid f(\phi(x)) = g(\psi(y)) \} \\ &\stackrel{\text{def}}{=} X \times_{K} Y. \end{split}$$

This finishes the proof.

Item 6, Unitality: We have

$$X \times_X A \cong \{(x, a) \in X \times A \mid f(a) = x\},\$$

$$A \times_X X \cong \{(a, x) \in X \times A \mid f(a) = x\},\$$

which are isomorphic to A via the maps $(x, a) \mapsto a$ and $(a, x) \mapsto a$. The proof of the naturality of $\lambda^{\mathsf{Sets}_{/X}}$ and $\rho^{\mathsf{Sets}_{/X}}$ is omitted.

Item 7, Commutativity: We have

$$A \times_{C} B \stackrel{\text{def}}{=} \{(a, b) \in A \times B \mid f(a) = g(b)\}$$

$$= \{(a, b) \in A \times B \mid g(b) = f(a)\}$$

$$\cong \{(b, a) \in B \times A \mid g(b) = f(a)\}$$

$$\stackrel{\text{def}}{=} B \times_{C} A.$$

The proof of the naturality of $\sigma^{\text{Sets}/X}$ is omitted. Item 8, Distributivity Over Coproducts: We have

$$A \times_X (B \coprod C) \stackrel{\text{def}}{=} \left\{ (a, z) \in A \times (B \coprod C) \middle| \phi_A(a) = \phi_{B \coprod C}(z) \right\}$$

$$= \left\{ (a,z) \in A \times (B \coprod C) \,\middle|\, z = (0,b) \text{ and } \phi_A(a) = \phi_{B \coprod C}(z) \right\}$$

$$\cup \left\{ (a,z) \in A \times (B \coprod C) \,\middle|\, z = (1,c) \text{ and } \phi_A(a) = \phi_{B \coprod C}(z) \right\}$$

$$= \left\{ (a,z) \in A \times (B \coprod C) \,\middle|\, z = (0,b) \text{ and } \phi_A(a) = \phi_B(b) \right\}$$

$$\cup \left\{ (a,z) \in A \times (B \coprod C) \,\middle|\, z = (1,c) \text{ and } \phi_A(a) = \phi_C(c) \right\}$$

$$\cong \left\{ (a,b) \in A \times B \,\middle|\, \phi_A(a) = \phi_B(b) \right\}$$

$$\cup \left\{ (a,c) \in A \times C \,\middle|\, \phi_A(a) = \phi_C(c) \right\}$$

$$\stackrel{\text{def}}{=} (A \times_X B) \cup (A \times_X C)$$

$$\cong (A \times_X B) \coprod (A \times_X C),$$

with the construction of the isomorphism

$$\delta_r^{\mathsf{Sets}_{/X}} : (A \coprod B) \times_X C \xrightarrow{\sim} (A \times_X C) \coprod (B \times_X C)$$

being similar. The proof of the naturality of $\delta_\ell^{\mathsf{Sets}_{/X}}$ and $\delta_r^{\mathsf{Sets}_{/X}}$ is omitted. Item 9, Annihilation With the Empty Set: We have

$$A \times_X \emptyset \stackrel{\text{def}}{=} \{(a, b) \in A \times \emptyset \mid f(a) = g(b)\}$$
$$= \{k \in \emptyset \mid f(a) = g(b)\}$$
$$= \emptyset,$$

and similarly for $\emptyset \times_X A$, where we have used Item 8 of Definition 4.1.3.1.3. The proof of the naturality of $\zeta_\ell^{\mathsf{Sets}_{/X}}$ and $\zeta_r^{\mathsf{Sets}_{/X}}$ is omitted.

Item 10, Interaction With Products: We have

$$A \times_{\mathsf{pt}} B \stackrel{\text{def}}{=} \{ (a, b) \in A \times B \mid !_{A}(a) = !_{B}(b) \}$$

$$\stackrel{\text{def}}{=} \{ (a, b) \in A \times B \mid \star = \star \}$$

$$= \{ (a, b) \in A \times B \}$$

$$= A \times B.$$

Item 11, Symmetric Monoidality: Omitted.

4.1.5 Equalisers

Let *A* and *B* be sets and let $f, g: A \Rightarrow B$ be functions.

Definition 4.1.5.1.1. The **equaliser of** f **and** g is the equaliser of f and g in Sets as in Limits and Colimits, ??.

Construction 4.1.5.1.2. Concretely, the equaliser of f and g is the pair (Eq(f,g),eq(f,g)) consisting of:

1. The Limit. The set Eq(f, g) defined by

$$\mathsf{Eq}(f,g) \stackrel{\mathsf{def}}{=} \{ a \in A \mid f(a) = g(a) \}.$$

2. The Cone. The inclusion map

$$eq(f,g) : Eq(f,g) \hookrightarrow A.$$

Proof. We claim that Eq(f,g) is the categorical equaliser of f and g in Sets. First we need to check that the relevant equaliser diagram commutes, i.e. that we have

$$f \circ eq(f,g) = g \circ eq(f,g),$$

which indeed holds by the definition of the set ${\sf Eq}(f,g)$. Next, we prove that ${\sf Eq}(f,g)$ satisfies the universal property of the equaliser. Suppose we have a diagram of the form

$$\operatorname{Eq}(f,g) \xrightarrow{\operatorname{eq}(f,g)} A \xrightarrow{f} B$$

$$E$$

in Sets. Then there exists a unique map $\phi: E \to \text{Eq}(f,q)$ making the diagram

$$\mathsf{Eq}(f,g) \xrightarrow{\mathsf{eq}(f,g)} A \xrightarrow{f} B$$

$$\downarrow \phi \mid \exists ! \qquad e$$

$$E \qquad \qquad E$$

commute, being uniquely determined by the condition

$$eq(f, q) \circ \phi = e$$

via

$$\phi(x) = e(x)$$

for each $x \in E$, where we note that $e(x) \in A$ indeed lies in Eq(f,g) by the condition

$$f \circ e = g \circ e$$
,

which gives

$$f(e(x)) = g(e(x))$$

for each $x \in E$, so that $e(x) \in Eq(f, g)$.

Proposition 4.1.5.1.3. Let *A*, *B*, and *C* be sets.

1. Associativity. We have isomorphisms of sets⁵

$$\underbrace{ \mathsf{Eq}(f \circ \mathsf{eq}(g,h), g \circ \mathsf{eq}(g,h))}_{= \mathsf{Eq}(f \circ \mathsf{eq}(g,h), h \circ \mathsf{eq}(g,h))} \cong \underbrace{ \mathsf{Eq}(f,g,h) \cong \underbrace{ \mathsf{Eq}(f \circ \mathsf{eq}(f,g), h \circ \mathsf{eq}(f,g))}_{= \mathsf{Eq}(g \circ \mathsf{eq}(f,g), h \circ \mathsf{eq}(f,g))}$$

⁵That is, the following three ways of forming "the' equaliser of (f, g, h) agree:

1. Take the equaliser of (f, g, h), i.e. the limit of the diagram

$$A \xrightarrow{f} B$$

in Sets.

2. First take the equaliser of f and g, forming a diagram

$$\mathsf{Eq}(f,g) \overset{\mathsf{eq}(f,g)}{\hookrightarrow} A \overset{f}{\underset{q}{\Longrightarrow}} B$$

and then take the equaliser of the composition

$$\operatorname{Eq}(f,g) \stackrel{\operatorname{eq}(f,g)}{\hookrightarrow} A \stackrel{f}{\underset{h}{\Longrightarrow}} B,$$

obtaining a subset

$${\rm Eq}(f\circ {\rm eq}(f,g),h\circ {\rm eq}(f,g))={\rm Eq}(g\circ {\rm eq}(f,g),h\circ {\rm eq}(f,g))$$
 of ${\rm Eq}(f,g).$

3. First take the equaliser of g and h, forming a diagram

$$\mathsf{Eq}(g,h) \overset{\mathsf{eq}(g,h)}{\hookrightarrow} A \overset{g}{\underset{h}{\Longrightarrow}} B$$

and then take the equaliser of the composition

$$\mathsf{Eq}(g,h) \overset{\mathsf{eq}(g,h)}{\hookrightarrow} A \overset{f}{\underset{q}{\Longrightarrow}} B,$$

obtaining a subset

$${\rm Eq}(f\circ {\rm eq}(g,h),g\circ {\rm eq}(g,h))={\rm Eq}(f\circ {\rm eq}(g,h),h\circ {\rm eq}(g,h))$$
 of ${\rm Eq}(g,h).$

where Eq(f, g, h) is the limit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets, being explicitly given by

$$Eq(f, q, h) \cong \{a \in A \mid f(a) = g(a) = h(a)\}.$$

4. Unitality. We have an isomorphism of sets

$$\operatorname{Eq}(f, f) \cong A$$
.

5. Commutativity. We have an isomorphism of sets

$$\operatorname{Eq}(f, q) \cong \operatorname{Eq}(q, f)$$
.

6. Interaction With Composition. Let

$$A \stackrel{f}{\underset{q}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C$$

be functions. We have an inclusion of sets

$$\mathsf{Eq}(h \circ f \circ \mathsf{eq}(f, g), k \circ g \circ \mathsf{eq}(f, g)) \subset \mathsf{Eq}(h \circ f, k \circ g),$$

where Eq $(h \circ f \circ eq(f,g), k \circ g \circ eq(f,g))$ is the equaliser of the composition

$$\mathsf{Eq}(f,g) \overset{\mathsf{eq}(f,g)}{\hookrightarrow} A \overset{f}{\underset{q}{\Longrightarrow}} B \overset{h}{\underset{k}{\Longrightarrow}} C.$$

Proof. Item 1, Associativity: We first prove that Eq(f, g, h) is indeed given by

$$Eq(f, g, h) \cong \{a \in A \mid f(a) = g(a) = h(a)\}.$$

Indeed, suppose we have a diagram of the form

$$\mathsf{Eq}(f,g,h) \xrightarrow{\mathsf{eq}(f,g,h)} A \xrightarrow{f \atop h} B$$

in Sets. Then there exists a unique map $\phi \colon E \to \operatorname{Eq}(f,g,h)$, uniquely determined by the condition

$$eq(f,q) \circ \phi = e$$

being necessarily given by

$$\phi(x) = e(x)$$

for each $x \in E$, where we note that $e(x) \in A$ indeed lies in Eq(f, g, h) by the condition

$$f \circ e = g \circ e = h \circ e$$
,

which gives

$$f(e(x)) = g(e(x)) = h(e(x))$$

for each $x \in E$, so that $e(x) \in Eq(f, q, h)$.

We now check the equalities

$$\operatorname{Eq}(f \circ \operatorname{eq}(q, h), q \circ \operatorname{eq}(q, h)) \cong \operatorname{Eq}(f, q, h) \cong \operatorname{Eq}(f \circ \operatorname{eq}(f, q), h \circ \operatorname{eq}(f, q)).$$

Indeed, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(g,h), g \circ \operatorname{eq}(g,h)) &\cong \{x \in \operatorname{Eq}(g,h) \mid [f \circ \operatorname{eq}(g,h)](a) = [g \circ \operatorname{eq}(g,h)](a)\} \\ &\cong \{x \in \operatorname{Eq}(g,h) \mid f(a) = g(a)\} \\ &\cong \{x \in A \mid f(a) = g(a) \text{ and } g(a) = h(a)\} \\ &\cong \{x \in A \mid f(a) = g(a) = h(a)\} \\ &\cong \operatorname{Eq}(f,g,h). \end{split}$$

Similarly, we have

$$\begin{split} \operatorname{Eq}(f \circ \operatorname{eq}(f,g), h \circ \operatorname{eq}(f,g)) &\cong \{x \in \operatorname{Eq}(f,g) \mid [f \circ \operatorname{eq}(f,g)](a) = [h \circ \operatorname{eq}(f,g)](a) \} \\ &\cong \{x \in \operatorname{Eq}(f,g) \mid f(a) = h(a) \} \\ &\cong \{x \in A \mid f(a) = h(a) \text{ and } f(a) = g(a) \} \\ &\cong \{x \in A \mid f(a) = g(a) = h(a) \} \\ &\cong \operatorname{Eq}(f,g,h). \end{split}$$

Item 4, Unitality: Indeed, we have

$$\mathsf{Eq}(f, f) \stackrel{\mathsf{def}}{=} \{ a \in A \, | \, f(a) = f(a) \}$$

$$= A$$

Item 5, Commutativity: Indeed, we have

$$\mathsf{Eq}(f,g) \stackrel{\mathsf{def}}{=} \{ a \in A \, | \, f(a) = g(a) \}$$

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$$= \{a \in A \mid g(a) = f(a)\}$$

$$\stackrel{\text{def}}{=} \mathsf{Eq}(q, f).$$

Item 6, Interaction With Composition: Indeed, we have

$$\begin{split} \operatorname{Eq}(h \circ f \circ \operatorname{eq}(f,g), k \circ g \circ \operatorname{eq}(f,g)) & \cong \{a \in \operatorname{Eq}(f,g) \mid h(f(a)) = k(g(a))\} \\ & \cong \{a \in A \mid f(a) = g(a) \text{ and } h(f(a)) = k(g(a))\}. \end{split}$$

and

$$Eq(h \circ f, k \circ g) \cong \{a \in A \mid h(f(a)) = k(g(a))\},\$$

and thus there,s an inclusion from Eq $(h \circ f \circ eq(f,g), k \circ g \circ eq(f,g))$ to Eq $(h \circ f, k \circ g)$.

4.1.6 Inverse Limits

Let $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I} \colon (I, \preceq) \to \mathsf{Sets}$ be an inverse system of sets.

Definition 4.1.6.1.1. The **inverse limit of** $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the inverse limit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ in Sets as in Limits and Colimits, ??.

Construction 4.1.6.1.2. Concretely, the inverse limit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the pair $(\lim_{\longleftarrow} (X_{\alpha}), \{\operatorname{pr}_{\alpha}\}_{\alpha\in I})$ consisting of:

1. The Limit. The set $\lim_{\stackrel{\longleftarrow}{\alpha \in I}} (X_{\alpha})$ defined by

$$\lim_{\substack{\longleftarrow \\ \alpha \in I}} (X_{\alpha}) \stackrel{\text{def}}{=} \left\{ (x_{\alpha})_{\alpha \in I} \in \prod_{\alpha \in I} X_{\alpha} \middle| \begin{array}{l} \text{for each } \alpha, \beta \in I, \text{ if } \alpha \preceq \beta, \\ \text{then we have } x_{\alpha} = f_{\alpha\beta}(x_{\beta}) \end{array} \right\}.$$

2. The Cone. The collection

$$\left\{ \operatorname{pr}_{\gamma} \colon \lim_{\stackrel{\longleftarrow}{\alpha \in I}} (X_{\alpha}) \to X_{\gamma} \right\}_{\gamma \in I}$$

of maps of sets defined as the restriction of the maps

$$\left\{ \operatorname{pr}_{\gamma} \colon \prod_{\alpha \in I} X_{\alpha} \to X_{\gamma} \right\}_{\gamma \in I}$$

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of Item 2 of Definition 4.1.2.1.2 to $\lim_{\stackrel{\longleftarrow}{\alpha\in I}}(X_\alpha)$ and hence given by

$$\operatorname{pr}_{\gamma}((x_{\alpha})_{\alpha \in I}) \stackrel{\text{def}}{=} x_{\gamma}$$

for each $\gamma \in I$ and each $(x_{\alpha})_{\alpha \in I} \in \lim_{\stackrel{\longleftarrow}{\alpha \in I}} (X_{\alpha})$.

Proof. We claim that $\lim_{\leftarrow \alpha \in I} (X_{\alpha})$ is the limit of the inverse system of sets $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta \in I}$. First we need to check that the limit diagram defined by it commutes, i.e. that we have

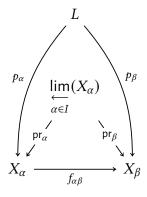
$$f_{\alpha\beta} \circ \operatorname{pr}_{\alpha} = \operatorname{pr}_{\beta}, \qquad \varprojlim_{\alpha \in I} (X_{\alpha})$$

$$X_{\alpha} \xrightarrow{f_{\alpha\beta}} X_{\beta}$$

for each $\alpha, \beta \in I$ with $\alpha \leq \beta$. Indeed, given $(x_{\gamma})_{\gamma \in I} \in \lim_{\leftarrow \gamma \in I} (X_{\gamma})$, we have

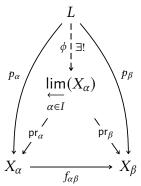
$$\begin{split} [f_{\alpha\beta} \circ \operatorname{pr}_{\alpha}]((x_{\gamma})_{\gamma \in I}) &\stackrel{\text{def}}{=} f_{\alpha\beta}(\operatorname{pr}_{\alpha}((x_{\gamma})_{\gamma \in I})) \\ &\stackrel{\text{def}}{=} f_{\alpha\beta}(x_{\alpha}) \\ &= x_{\beta} \\ &\stackrel{\text{def}}{=} \operatorname{pr}_{\beta}((x_{\gamma})_{\gamma \in I}), \end{split}$$

where the third equality comes from the definition of $\lim_{\stackrel{\longleftarrow}{\alpha}\in I}(X_{\alpha})$. Next, we prove that $\lim_{\stackrel{\longleftarrow}{\alpha}\in I}(X_{\alpha})$ satisfies the universal property of an inverse limit. Suppose that we have, for each $\alpha,\beta\in I$ with $\alpha\preceq\beta$, a diagram of the form



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in Sets. Then there indeed exists a unique map $\phi\colon L \xrightarrow{\exists !} \varprojlim_{\alpha \in I} (X_\alpha)$ making the diagram



commute, being uniquely determined by the family of conditions

$$\{p_{\alpha} = \operatorname{pr}_{\alpha} \circ \phi\}_{\alpha \in I}$$

via

$$\phi(\ell) = (p_{\alpha}(\ell))_{\alpha \in I}$$

for each $\ell \in L$, where we note that $(p_{\alpha}(\ell))_{\alpha \in I} \in \prod_{\alpha \in I} X_{\alpha}$ indeed lies in $\lim_{\longleftarrow \alpha \in I} (X_{\alpha})$, as we have

$$f_{\alpha\beta}(p_{\alpha}(\ell)) \stackrel{\text{def}}{=} [f_{\alpha\beta} \circ p_{\alpha}](\ell)$$

$$\stackrel{\text{def}}{=} p_{\beta}(\ell)$$

for each $\beta \in I$ with $\alpha \leq \beta$ by the commutativity of the diagram for $(L, \{p_{\alpha}\}_{\alpha \in I})$. \square **Example 4.1.6.1.3.** Here are some examples of inverse limits of sets.

1. The p-Adic Integers. The ring of p-adic integers \mathbb{Z}_p of $\mathbf{??}$ is the inverse limit

$$\mathbb{Z}_p \cong \lim_{\substack{\longleftarrow \\ n \in \mathbb{N}}} (\mathbb{Z}_{/p^n});$$

see ??.

2. Rings of Formal Power Series. The ring R[[t]] of formal power series in a variable t is the inverse limit

$$R[[t]] \cong \lim_{\substack{\longleftarrow \\ n \in \mathbb{N}}} (R[t]/t^n R[t]);$$

see ??.

3. Profinite Groups. Profinite groups are inverse limits of finite groups; see ??.

4.2 Colimits of Sets

4.2.1 The Initial Set

Definition 4.2.1.1.1. The **initial set** is the initial object of Sets as in Limits and Colimits, ??.

Construction 4.2.1.1.2. Concretely, the initial set is the pair $(\emptyset, \{\iota_A\}_{A \in \mathsf{Obj}(\mathsf{Sets})})$ consisting of:

- 1. The Colimit. The empty set Ø of Definition 4.3.1.1.1.
- 2. The Cocone. The collection of maps

$${\iota_A \colon \emptyset \to A}_{A \in \mathsf{Obj}(\mathsf{Sets})}$$

given by the inclusion maps from \emptyset to A.

Proof. We claim that \emptyset is the initial object of Sets. Indeed, suppose we have a diagram of the form

$$\emptyset$$
 A

in Sets. Then there exists a unique map $\phi: \emptyset \to A$ making the diagram

$$\emptyset \xrightarrow{\phi} A$$

commute, namely the inclusion map ι_A .

4.2.2 Coproducts of Families of Sets

Let $\{A_i\}_{i\in I}$ be a family of sets.

Definition 4.2.2.1.1. The **coproduct of** $\{A_i\}_{i\in I}^6$ is the coproduct of $\{A_i\}_{i\in I}$ in Sets as in Limits and Colimits, ??.

Construction 4.2.2.1.2. Concretely, the disjoint union of $\{A_i\}_{i\in I}$ is the pair $(\coprod_{i\in I} A_i, \{\text{inj}_i\}_{i\in I})$ consisting of:

⁶ Further Terminology: Also called the **disjoint union of the family** $\{A_i\}_{i\in I}$.

1. The Colimit. The set $\coprod_{i \in I} A_i$ defined by

$$\coprod_{i\in I} A_i \stackrel{\text{def}}{=} \left\{ (i, x) \in I \times \left(\bigcup_{i\in I} A_i \right) \middle| x \in A_i \right\}.$$

2. The Cocone. The collection

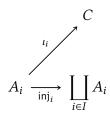
$$\left\{ \mathsf{inj}_i \colon A_i \to \coprod_{i \in I} A_i \right\}_{i \in I}$$

of maps given by

$$\operatorname{inj}_i(x) \stackrel{\text{def}}{=} (i, x)$$

for each $x \in A_i$ and each $i \in I$.

Proof. We claim that $\coprod_{i \in I} A_i$ is the categorical coproduct of $\{A_i\}_{i \in I}$ in Sets. Indeed, suppose we have, for each $i \in I$, a diagram of the form



in Sets. Then there exists a unique map $\phi:\coprod_{i\in I}A_i\to C$ making the diagram

$$A_i \xrightarrow[\inf_i]{l_i} C$$

$$\phi \downarrow \exists !$$

$$A_i \xrightarrow[\inf_i]{l_i} A$$

commute, being uniquely determined by the condition $\phi \circ \operatorname{inj}_i = \iota_i$ for each $i \in I$ via

$$\phi((i,x)) = \iota_i(x)$$

for each $(i, x) \in \coprod_{i \in I} A_i$.

Proposition 4.2.2.1.3. Let $\{A_i\}_{i\in I}$ be a family of sets.

1. Functoriality. The assignment $\{A_i\}_{i\in I}\mapsto \coprod_{i\in I}A_i$ defines a functor

$$\coprod_{i \in I} : \mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}) \to \mathsf{Sets}$$

where

· Action on Objects. For each $(A_i)_{i \in I} \in \mathsf{Obj}(\mathsf{Fun}(I_{\mathsf{disc}},\mathsf{Sets}))$, we have

$$\left[\bigsqcup_{i \in I} \right] ((A_i)_{i \in I}) \stackrel{\text{def}}{=} \bigsqcup_{i \in I} A_i$$

· Action on Morphisms. For each $(A_i)_{i \in I}$, $(B_i)_{i \in I} \in \mathsf{Obj}(\mathsf{Fun}(I_{\mathsf{disc}}, \mathsf{Sets}))$, the action on Hom-sets

$$\left(\bigsqcup_{i\in I}\right)_{(A_i)_{i\in I},(B_i)_{i\in I}}\colon\operatorname{Nat}((A_i)_{i\in I},(B_i)_{i\in I})\to\operatorname{Sets}\!\left(\bigsqcup_{i\in I}A_i,\bigsqcup_{i\in I}B_i\right)$$

of $\coprod_{i \in I}$ at $((A_i)_{i \in I}, (B_i)_{i \in I})$ is defined by sending a map

$$\{f_i\colon A_i\to B_i\}_{i\in I}$$

in Nat $((A_i)_{i \in I}, (B_i)_{i \in I})$ to the map of sets

$$\coprod_{i \in I} f_i \colon \coprod_{i \in I} A_i \to \coprod_{i \in I} B_i$$

defined by

$$\left[\bigsqcup_{i \in I} f_i \right] (i, a) \stackrel{\text{def}}{=} f_i(a)$$

for each $(i, a) \in \coprod_{i \in I} A_i$.

Proof. Item 1, Functoriality: This follows from Limits and Colimits, ?? of ??.

4.2.3 Binary Coproducts

Let A and B be sets.

Definition 4.2.3.1.1. The **coproduct of** A **and** B⁷ is the coproduct of A and B in Sets as in Limits and Colimits, ??.

Construction 4.2.3.1.2. Concretely, the coproduct of A and B is the pair $(A \coprod B, \{inj_1, inj_2\})$ consisting of:

1. The Colimit. The set $A \coprod B$ defined by

$$A \coprod B \stackrel{\text{def}}{=} \coprod_{z \in \{A,B\}} z$$

$$\stackrel{\text{def}}{=} \{(0,a) \in S \mid a \in A\} \cup \{(1,b) \in S \mid b \in B\},$$

where $S = \{0, 1\} \times (A \cup B)$.

2. The Cocone. The maps

$$inj_1: A \to A \coprod B,$$

 $inj_2: B \to A \coprod B,$

given by

$$\operatorname{inj}_{1}(a) \stackrel{\text{def}}{=} (0, a),$$

 $\operatorname{inj}_{2}(b) \stackrel{\text{def}}{=} (1, b),$

for each $a \in A$ and each $b \in B$.

Proof. We claim that $A \coprod B$ is the categorical coproduct of A and B in Sets. Indeed, suppose we have a diagram of the form

$$A \xrightarrow[\text{inj}_1]{C} A \coprod B \xrightarrow[\text{inj}_2]{l_2} B$$

in Sets. Then there exists a unique map $\phi:A\coprod B\to C$ making the diagram

$$A \underset{\text{inj}_{1}}{\overset{\iota_{1}}{\longrightarrow}} A \coprod B \underset{\text{inj}_{2}}{\overset{\iota_{2}}{\longrightarrow}} B$$

⁷ Further Terminology: Also called the **disjoint union of** A **and** B.

commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_A = \iota_A,$$
$$\phi \circ \operatorname{inj}_B = \iota_B$$

via

$$\phi(x) = \begin{cases} \iota_A(a) & \text{if } x = (0, a), \\ \iota_B(b) & \text{if } x = (1, b) \end{cases}$$

for each $x \in A \coprod B$.

Proposition 4.2.3.1.3. Let *A*, *B*, *C*, and *X* be sets.

1. Functoriality. The assignment $A, B, (A, B) \mapsto A \coprod B$ defines functors

$$A \coprod -:$$
 Sets \rightarrow Sets,
 $- \coprod B:$ Sets \rightarrow Sets,
 $-_1 \coprod -_2:$ Sets \times Sets \rightarrow Sets,

where $-1 \prod_{i=1}^{n} -2$ is the functor where

· Action on Objects. For each $(A, B) \in \mathsf{Obj}(\mathsf{Sets} \times \mathsf{Sets})$, we have

$$[-_1 \coprod -_2](A, B) \stackrel{\text{def}}{=} A \coprod B.$$

· Action on Morphisms. For each (A, B), $(X, Y) \in \mathsf{Obj}(\mathsf{Sets})$, the action on Hom-sets

$$\coprod_{(A,B),(X,Y)}$$
: $\mathsf{Sets}(A,X) \times \mathsf{Sets}(B,Y) \to \mathsf{Sets}(A \coprod B,X \coprod Y)$

of \coprod at ((A, B), (X, Y)) is defined by sending (f, g) to the function

$$f \coprod g : A \coprod B \to X \coprod Y$$

defined by

$$[f \coprod g](x) \stackrel{\text{def}}{=} \begin{cases} (0, f(a)) & \text{if } x = (0, a), \\ (1, g(b)) & \text{if } x = (1, b), \end{cases}$$

for each $x \in A \coprod B$.

and where $A \coprod -$ and $- \coprod B$ are the partial functors of $-_1 \coprod -_2$ at $A, B \in$ Obj(Sets).

2. Adjointness. We have an adjunction

$$(-_1 \coprod -_2 \dashv \Delta_{\mathsf{Sets}}) \colon \quad \mathsf{Sets} \times \mathsf{Sets} \underbrace{ \stackrel{-_1 \coprod -_2}{\Delta_{\mathsf{Sets}}}} \mathsf{Sets},$$

witnessed by a bijection

$$Sets(A \coprod B, C), \cong Hom_{Sets \times Sets}((A, B), (C, C))$$

natural in $(A, B) \in \mathsf{Obj}(\mathsf{Sets} \times \mathsf{Sets})$ and in $C \in \mathsf{Obj}(\mathsf{Sets})$.

3. Associativity. We have an isomorphism of sets

$$\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod} : (X \coprod Y) \coprod Z \xrightarrow{\sim} X \coprod (Y \coprod Z),$$

natural in $X, Y, Z \in \mathsf{Obj}(\mathsf{Sets})$.

4. Unitality. We have isomorphisms of sets

$$\lambda_X^{\mathsf{Sets},\coprod} : \emptyset \coprod X \xrightarrow{\sim} X,$$

$$\rho_X^{\mathsf{Sets},\coprod}\colon X \coprod \emptyset \stackrel{\sim}{\dashrightarrow} X,$$

natural in $X \in \mathsf{Obj}(\mathsf{Sets})$.

5. Commutativity. We have an isomorphism of sets

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod}: X \coprod Y \xrightarrow{\sim} Y \coprod X,$$

natural in $X, Y \in \mathsf{Obj}(\mathsf{Sets})$.

6. Symmetric Monoidality. The 7-tuple (Sets , \coprod , \emptyset , α^{Sets} , λ^{Sets} , ρ^{Sets} , σ^{Sets}) is a symmetric monoidal category.

Proof. Item 1, *Functoriality*: This follows from Limits and Colimits, ?? of ??.

Item 2, Adjointness: This follows from the universal property of the coproduct.

Item 3, *Associativity*: This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.2.3.1.1.

Item 4, Unitality: This is proved in the proof of Monoidal Structures on the Category of Sets, Definitions 5.2.4.1.1 and 5.2.5.1.1.

Item 5, Commutativity: This is proved in the proof of Monoidal Structures on the Category of Sets, Definition 5.2.6.1.1.

Item 6, Symmetric Monoidality: This is a repetition of Monoidal Structures on the Category of Sets, Definition 5.2.7.1.1, and is proved there.

4.2.4 Pushouts

Let A, B, and C be sets and let $f: C \to A$ and $g: C \to B$ be functions.

Definition 4.2.4.1.1. The **pushout of** A **and** B **over** C **along** f **and** g⁸ is the pushout of A and B over C along f and g in Sets as in Limits and Colimits, ??.

Construction 4.2.4.1.2. Concretely, the pushout of A and B over C along f and g is the pair $(A \coprod_C B, \{\text{inj}_1, \text{inj}_2\})$ consisting of:

1. The Colimit. The set $A \coprod_C B$ defined by

$$A \coprod_C B \stackrel{\text{def}}{=} A \coprod_C B / \sim_C$$

where \sim_C is the equivalence relation on $A \coprod B$ generated by $(0, f(c)) \sim_C (1, g(c))$.

2. The Cocone. The maps

$$\operatorname{inj}_1: A \to A \coprod_C B,$$

 $\operatorname{inj}_2: B \to A \coprod_C B$

given by

$$\operatorname{inj}_1(a) \stackrel{\text{def}}{=} [(0, a)]$$

 $\operatorname{inj}_2(b) \stackrel{\text{def}}{=} [(1, b)]$

for each $a \in A$ and each $b \in B$.

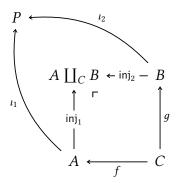
Proof. We claim that $A \coprod_C B$ is the categorical pushout of A and B over C with respect to (f,g) in Sets. First we need to check that the relevant pushout diagram commutes, i.e. that we have

⁸ Further Terminology: Also called the **fibre coproduct of** A **and** B **over** C **along** f **and** g.

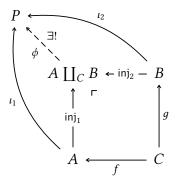
Indeed, given $c \in C$, we have

$$\begin{aligned} [\inf_1 \circ f](c) &= \inf_1(f(c)) \\ &= [(0, f(c))] \\ &= [(1, g(c))] \\ &= \inf_2(g(c)) \\ &= [\inf_2 \circ g](c), \end{aligned}$$

where [(0, f(c))] = [(1, g(c))] by the definition of the relation \sim on $A \coprod B$. Next, we prove that $A \coprod {}_{C}B$ satisfies the universal property of the pushout. Suppose we have a diagram of the form



in Sets. Then there exists a unique map $\phi \colon A \coprod_C B \to P$ making the diagram



commute, being uniquely determined by the conditions

$$\phi \circ \operatorname{inj}_1 = \iota_1,$$

 $\phi \circ \operatorname{inj}_2 = \iota_2$

via

$$\phi(x) = \begin{cases} \iota_1(a) & \text{if } x = [(0, a)], \\ \iota_2(b) & \text{if } x = [(1, b)] \end{cases}$$

for each $x \in A \coprod_C B$, where the well-definedness of ϕ is guaranteed by the equality $\iota_1 \circ f = \iota_2 \circ g$ and the definition of the relation \sim on $A \coprod B$ as follows:

1. Case 1: Suppose we have x = [(0, a)] = [(0, a')] for some $a, a' \in A$. Then, by Definition 4.2.4.1.3, we have a sequence

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0,a').$$

2. Case 2: Suppose we have x = [(1, b)] = [(1, b')] for some $b, b' \in B$. Then, by Definition 4.2.4.1.3, we have a sequence

$$(1,b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b').$$

3. Case 3: Suppose we have x = [(0, a)] = [(1, b)] for some $a \in A$ and $b \in B$. Then, by Definition 4.2.4.1.3, we have a sequence

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b).$$

In all these cases, we declare $x \sim' y$ iff there exists some $c \in C$ such that x = (0, f(c)) and y = (1, g(c)) or x = (1, g(c)) and y = (0, f(c)). Then, the equality $\iota_1 \circ f = \iota_2 \circ g$ gives

$$\phi([x]) = \phi([(0, f(c))])
\stackrel{\text{def}}{=} \iota_1(f(c))
= \iota_2(g(c))
\stackrel{\text{def}}{=} \phi([(1, g(c))])
= \phi([y]),$$

with the case where x=(1,g(c)) and y=(0,f(c)) similarly giving $\phi([x])=\phi([y])$. Thus, if $x\sim' y$, then $\phi([x])=\phi([y])$. Applying this equality pairwise to the sequences

$$(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (0,a'),$$

 $(1,b) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b'),$
 $(0,a) \sim' x_1 \sim' \cdots \sim' x_n \sim' (1,b)$

gives

$$\phi([(0,a)]) = \phi([(0,a')]),$$

$$\phi([(1,b)]) = \phi([(1,b')]),$$

$$\phi([(0,a)]) = \phi([(1,b)]),$$

showing ϕ to be well-defined.

Remark 4.2.4.1.3. In detail, by Conditions on Relations, Definition 10.5.2.1.2, the relation \sim of Definition 4.2.4.1.1 is given by declaring $a \sim b$ iff one of the following conditions is satisfied:

- 1. We have $a, b \in A$ and a = b.
- 2. We have $a, b \in B$ and a = b.
- 3. There exist $x_1, \ldots, x_n \in A \coprod B$ such that $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$, where we declare $x \sim' y$ if one of the following conditions is satisfied:
 - (a) There exists $c \in C$ such that x = (0, f(c)) and y = (1, g(c)).
 - (b) There exists $c \in C$ such that x = (1, g(c)) and y = (0, f(c)).

In other words, there exist $x_1, \ldots, x_n \in A \coprod B$ satisfying the following conditions:

- (c) There exists $c_0 \in C$ satisfying one of the following conditions:
 - i. We have $a = f(c_0)$ and $x_1 = g(c_0)$.
 - ii. We have $a = q(c_0)$ and $x_1 = f(c_0)$.
- (d) For each $1 \le i \le n-1$, there exists $c_i \in C$ satisfying one of the following conditions:
 - i. We have $x_i = f(c_i)$ and $x_{i+1} = q(c_i)$.
 - ii. We have $x_i = g(c_i)$ and $x_{i+1} = f(c_i)$.
- (e) There exists $c_n \in C$ satisfying one of the following conditions:
 - i. We have $x_n = f(c_n)$ and $b = g(c_n)$.
 - ii. We have $x_n = g(c_n)$ and $b = f(c_n)$.

Remark 4.2.4.1.4. It is common practice to write $A \coprod_C B$ for the pushout of A and B over C along f and g, omitting the maps f and g from the notation and instead

leaving them implicit, to be understood from the context.

However, the set $A \coprod_{C} B$ depends very much on the maps f and g, and sometimes it is necessary or useful to note this dependence explicitly. In such situations, we will write $A \coprod_{f,C,g} B$ or $A \coprod_{C} {f,g \choose G} B$ for $A \coprod_{C} B$.

Example 4.2.4.1.5. Here are some examples of pushouts of sets.

- 1. Wedge Sums of Pointed Sets. The wedge sum of two pointed sets of Pointed Sets, Definition 6.3.3.1.1 is an example of a pushout of sets.
- 2. Intersections via Unions. Let X be a set. We have

$$A \cup B \cong A \coprod_{A \cap B} B, \qquad A \longleftarrow B$$

$$A \longleftarrow A \cap B$$

for each $A, B \in \mathcal{P}(X)$.

Proof. Item 1, Wedge Sums of Pointed Sets: This follows by definition, as the wedge sum of two pointed sets is defined as a pushout.

Item 2, Intersections via Unions: Indeed, $A \coprod_{A \cap B} B$ is the quotient of $A \coprod B$ by the equivalence relation obtained by declaring $(0, a) \sim (1, b)$ iff $a = b \in A \cap B$, which is in bijection with $A \cup B$ via the map with $[(0, a)] \mapsto a$ and $[(1, b)] \mapsto b$.

Proposition 4.2.4.1.6. Let *A*, *B*, *C*, and *X* be sets.

1. Functoriality. The assignment $(A, B, C, f, g) \mapsto A \coprod_{f,C,g} B$ defines a functor

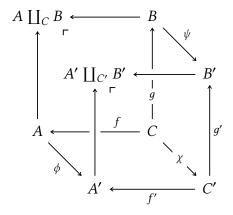
$$-_1 \coprod_{-_3} -_1 : \mathsf{Fun}(\mathcal{P},\mathsf{Sets}) \to \mathsf{Sets},$$

where \mathcal{P} is the category that looks like this:



In particular, the action on morphisms of $-1 \coprod_{-3} -1$ is given by sending a

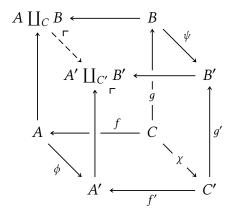
morphism



in Fun(\mathcal{P} , Sets) to the map $\xi \colon A \coprod_{C} B \xrightarrow{\exists !} A' \coprod_{C'} B'$ given by

$$\xi(x) \stackrel{\text{def}}{=} \begin{cases} \phi(a) & \text{if } x = [(0, a)], \\ \psi(b) & \text{if } x = [(1, b)] \end{cases}$$

for each $x \in A \coprod_C B$, which is the unique map making the diagram



commute.

2. Adjointness. We have an adjunction

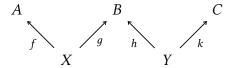
$$\left(-1 \coprod_{X} -_2 \dashv \Delta_{\mathsf{Sets}_{X/}}\right) : \quad \mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/} \xrightarrow{\perp} \mathsf{Sets}_{X/} \mathsf{Sets}_{X/}$$

witnessed by a bijection

$$\mathsf{Sets}_{X/}(A \coprod_X B, C), \cong \mathsf{Hom}_{\mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/}}((A, B), (C, C))$$

natural in $(A, B) \in \mathsf{Obj}(\mathsf{Sets}_{X/} \times \mathsf{Sets}_{X/})$ and in $C \in \mathsf{Obj}(\mathsf{Sets}_{X/})$.

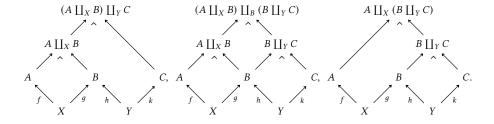
3. Associativity. Given a diagram



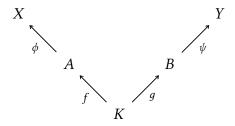
in Sets, we have isomorphisms of sets

$$(A \coprod_X B) \coprod_Y C \cong (A \coprod_X B) \coprod_B (B \coprod_Y C) \cong A \coprod_X (B \coprod_Y C)$$

where these pullbacks are built as in the diagrams



4. Interaction With Composition. Given a diagram



in Sets, we have isomorphisms of sets

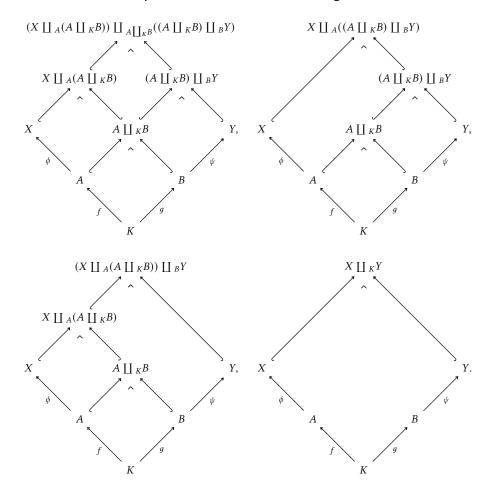
$$\begin{split} X \coprod_K^{\phi \circ f, \psi \circ g} Y &\cong (X \coprod_A^{\phi, j_1} (A \coprod_K^{f, g} B)) \coprod_{A \coprod_K^{f, g} B}^{i_2, i_1} ((A \coprod_K^{f, g} B) \coprod_B^{j_2, \psi} Y) \\ &\cong X \coprod_A^{\phi, i} ((A \coprod_K^{f, g} B) \coprod_B^{j_2, \psi} Y) \end{split}$$

$$\cong (X \coprod {}_A^{\phi,i_1}(A \coprod {}_K^{f,g}B)) \coprod {}_B^{j,\psi}Y$$

where

$$\begin{split} j_1 &= \mathsf{inj}_1^{A \times_K^{f,g} B}, & j_2 &= \mathsf{inj}_2^{A \times_K^{f,g} B}, \\ i_1 &= \mathsf{inj}_1^{(A \times_K^{f,g} B) \times_Y^{q_2,\psi}}, & X \times_{A \times_K^{f,g} B}^{\phi,q_1} (A \times_K^{f,g} B) \\ i_2 &= \mathsf{inj}_2^{(A \times_K^{f,g} B)}, & i_2 &= \mathsf{inj}_2^{X \times_A^{\phi,q_1}} (A \times_K^{f,g} B) \\ i &= j_1 \circ \mathsf{inj}_1^{(A \times_K^{f,g} B) \times_B^{q_2,\psi} Y}, & j &= j_2 \circ \mathsf{inj}_2^{X \times_A^{\phi,q_1}} (A \times_K^{f,g} B) \end{split}$$

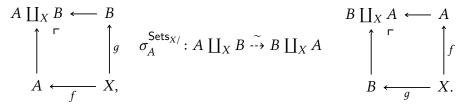
and where these pullbacks are built as in the diagrams



5. Unitality. We have isomorphisms of sets

natural in $(A, f) \in \mathsf{Obj}(\mathsf{Sets}_{X/})$.

6. Commutativity. We have an isomorphism of sets



natural in $(A, f), (B, g) \in \mathsf{Obj}(\mathsf{Sets}_{X/})$.

7. Interaction With Coproducts. We have

$$A \coprod_{\emptyset} B \cong A \coprod_{B}, \qquad A \coprod_{\iota_{A}} B \longleftarrow_{\iota_{B}} B$$

$$A \longleftarrow_{\iota_{A}} \emptyset.$$

8. Symmetric Monoidality. The triple (Sets $_{X/}$, \coprod_{X} , X) is a symmetric monoidal category.

Proof. Item 1, Functoriality: This is a special case of functoriality of co/limits, Limits and Colimits, ?? of ??, with the explicit expression for ξ following from the commutativity of the cube pushout diagram.

Item 2, : *Adjointness*: This follows from the universal property of the coproduct (pushouts are coproducts in $Sets_{X/}$).

Item 3, Associativity: Omitted.

Item 4, Interaction With Composition: Omitted.

Item 5, Unitality: Omitted.

Item 6, Commutativity: Omitted.

Item 7, Interaction With Coproducts: Omitted.

Item 8, Symmetric Monoidality: Omitted.

4.2.5 Coequalisers

Let *A* and *B* be sets and let $f, g: A \Rightarrow B$ be functions.

Definition 4.2.5.1.1. The **coequaliser of** f **and** g is the coequaliser of f and g in Sets as in Limits and Colimits, ??.

Construction 4.2.5.1.2. Concretely, the coequaliser of f and g is the pair (CoEq(f,g), coeq(f,g)) consisting of:

1. The Colimit. The set CoEq(f, g) defined by

$$CoEq(f, g) \stackrel{\text{def}}{=} B/\sim$$
,

where \sim is the equivalence relation on B generated by $f(a) \sim q(a)$.

2. The Cocone. The map

$$coeq(f, q) : B \rightarrow CoEq(f, q)$$

given by the quotient map $\pi \colon B \twoheadrightarrow B/\sim$ with respect to the equivalence relation generated by $f(a) \sim g(a)$.

Proof. We claim that CoEq(f,g) is the categorical coequaliser of f and g in Sets. First we need to check that the relevant coequaliser diagram commutes, i.e. that we have

$$coeq(f, g) \circ f = coeq(f, g) \circ g.$$

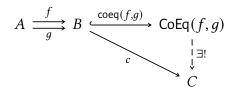
Indeed, we have

$$\begin{aligned} [\operatorname{coeq}(f,g) \circ f](a) &\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](f(a)) \\ &\stackrel{\text{def}}{=} [f(a)] \\ &= [g(a)] \\ &\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g)](g(a)) \\ &\stackrel{\text{def}}{=} [\operatorname{coeq}(f,g) \circ g](a) \end{aligned}$$

for each $a \in A$. Next, we prove that $\mathsf{CoEq}(f,g)$ satisfies the universal property of the coequaliser. Suppose we have a diagram of the form

$$A \xrightarrow{f} B \xrightarrow{\operatorname{coeq}(f,g)} \operatorname{CoEq}(f,g)$$

in Sets. Then, since c(f(a)) = c(g(a)) for each $a \in A$, it follows from Conditions on Relations, Items 4 and 5 of Definition 10.6.2.1.3 that there exists a unique map $CoEq(f,g) \xrightarrow{\exists !} C$ making the diagram



commute.

Remark 4.2.5.1.3. In detail, by Conditions on Relations, Definition 10.5.2.1.2, the relation \sim of Definition 4.2.5.1.1 is given by declaring $a \sim b$ iff one of the following conditions is satisfied:

- 1. We have a = b;
- 2. There exist $x_1, \ldots, x_n \in B$ such that $a \sim' x_1 \sim' \cdots \sim' x_n \sim' b$, where we declare $x \sim' y$ if one of the following conditions is satisfied:
 - (a) There exists $z \in A$ such that x = f(z) and y = g(z).
 - (b) There exists $z \in A$ such that x = g(z) and y = f(z).

In other words, there exist $x_1, \ldots, x_n \in B$ satisfying the following conditions:

- (a) There exists $z_0 \in A$ satisfying one of the following conditions:
 - i. We have $a = f(z_0)$ and $x_1 = g(z_0)$.
 - ii. We have $a = q(z_0)$ and $x_1 = f(z_0)$.
- (b) For each $1 \le i \le n-1$, there exists $z_i \in A$ satisfying one of the following conditions:
 - i. We have $x_i = f(z_i)$ and $x_{i+1} = g(z_i)$.
 - ii. We have $x_i = g(z_i)$ and $x_{i+1} = f(z_i)$.
- (c) There exists $z_n \in A$ satisfying one of the following conditions:
 - i. We have $x_n = f(z_n)$ and $b = g(z_n)$.
 - ii. We have $x_n = g(z_n)$ and $b = f(z_n)$.

Example 4.2.5.1.4. Here are some examples of coequalisers of sets.

1. Quotients by Equivalence Relations. Let R be an equivalence relation on a set X. We have a bijection of sets

$$X/{\sim_R} \cong \mathsf{CoEq}(R \hookrightarrow X \times X \overset{\mathsf{pr}_1}{\underset{\mathsf{pr}_2}{\Longrightarrow}} X).$$

Proof. Item 1, Quotients by Equivalence Relations: See [Pro25z].

Proposition 4.2.5.1.5. Let A, B, and C be sets.

1. Associativity. We have isomorphisms of sets⁹

$$\underbrace{\mathsf{CoEq}(\mathsf{coeq}(f,g) \circ f, \mathsf{coeq}(f,g) \circ h)}_{=\mathsf{CoEq}(\mathsf{coeq}(f,g) \circ g, \mathsf{coeq}(f,g) \circ h)} \cong \mathsf{CoEq}(f,g,h) \cong \underbrace{\mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ g)}_{=\mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ h)},$$

where CoEq(f, g, h) is the colimit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

 9 That is, the following three ways of forming "the' coequaliser of (f,g,h) agree:

1. Take the coequaliser of (f, g, h), i.e. the colimit of the diagram

$$A \xrightarrow{f \atop g \atop h} B$$

in Sets.

2. First take the coequaliser of f and g, forming a diagram

$$A \stackrel{f}{\underset{q}{\Longrightarrow}} B \stackrel{\mathsf{coeq}(f,g)}{\twoheadrightarrow} \mathsf{CoEq}(f,g)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\underset{h}{\Longrightarrow}} B \stackrel{\operatorname{coeq}(f,g)}{\twoheadrightarrow} \operatorname{CoEq}(f,g),$$

obtaining a quotient

$$\label{eq:coeq} \text{CoEq}(\text{coeq}(f,g)\circ f, \text{coeq}(f,g)\circ h) = \text{CoEq}(\text{coeq}(f,g)\circ g, \text{coeq}(f,g)\circ h)$$
 of
$$\text{CoEq}(f,g)$$

3. First take the coequaliser of g and h, forming a diagram

$$A \stackrel{g}{\underset{h}{\Longrightarrow}} B \stackrel{\mathsf{coeq}(g,h)}{\twoheadrightarrow} \mathsf{CoEq}(g,h)$$

and then take the coequaliser of the composition

$$A \stackrel{f}{\underset{q}{\Longrightarrow}} B \stackrel{\operatorname{coeq}(g,h)}{\Longrightarrow} \operatorname{CoEq}(g,h),$$

obtaining a quotient

$$\mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ g) = \mathsf{CoEq}(\mathsf{coeq}(g,h) \circ f, \mathsf{coeq}(g,h) \circ h)$$
 of
$$\mathsf{CoEq}(g,h).$$

in Sets.

4. Unitality. We have an isomorphism of sets

$$CoEq(f, f) \cong B$$
.

5. Commutativity. We have an isomorphism of sets

$$CoEq(f, g) \cong CoEq(g, f)$$
.

6. Interaction With Composition. Let

$$A \stackrel{f}{\underset{q}{\Longrightarrow}} B \stackrel{h}{\underset{k}{\Longrightarrow}} C$$

be functions. We have a surjection

$$CoEq(h \circ f, k \circ g) \twoheadrightarrow CoEq(coeq(h, k) \circ h \circ f, coeq(h, k) \circ k \circ g)$$

exhibiting CoEq(coeq(h,k) \circ h \circ f, coeq(h,k) \circ k \circ g) as a quotient of CoEq(h \circ f, k \circ g) by the relation generated by declaring h(y) \sim k(y) for each $y \in B$.

Proof. Item 1, Associativity: Omitted.

Item 4, Unitality: Omitted.

Item 5, Commutativity: Omitted.

Item 6, Interaction With Composition: Omitted.

4.2.6 Direct Colimits

Let $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I} \colon (I, \preceq) \to \mathbb{T}$ be a direct system of sets.

Definition 4.2.6.1.1. The **direct colimit of** $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the direct colimit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ in Sets as in Limits and Colimits, **??**.

Construction 4.2.6.1.2. Concretely, the direct colimit of $(X_{\alpha}, f_{\alpha\beta})_{\alpha,\beta\in I}$ is the pair $\left(\begin{array}{c} \operatorname{colim}(X_{\alpha}), \left\{ \operatorname{inj}_{\alpha} \right\}_{\alpha\in I} \end{array} \right)$ consisting of:

1. The Colimit. The set $\operatorname{colim}(X_{\alpha})$ defined by

$$\operatorname{colim}_{\underset{\alpha \in I}{\longrightarrow}} (X_{\alpha}) \stackrel{\text{def}}{=} \left(\left[\prod_{\alpha \in I} X_{\alpha} \right] \right) / \sim,$$

where \sim is the equivalence relation on $\coprod_{\alpha \in I} X_{\alpha}$ generated by declaring $(\alpha,x) \sim (\beta,y)$ iff there exists some $\gamma \in I$ satisfying the following conditions:

- (a) We have $\alpha \leq \gamma$.
- (b) We have $\beta \leq \gamma$.
- (c) We have $f_{\alpha \gamma}(x) = f_{\beta \gamma}(y)$.
- 2. The Cocone. The collection

$$\left\{\operatorname{inj}_{\gamma} \colon X_{\gamma} \to \operatornamewithlimits{colim}_{\underset{\alpha \in I}{\longrightarrow}} (X_{\alpha})\right\}_{\gamma \in I}$$

of maps of sets defined by

$$\operatorname{inj}_{\gamma}(x) \stackrel{\text{def}}{=} [(\gamma, x)]$$

for each $\gamma \in I$ and each $x \in X_{\gamma}$.

Proof. We will prove Definition 4.2.6.1.2 below in a bit, but first we need a lemma (which is interesting in its own right).

Lemma 4.2.6.1.3. For each $\alpha, \beta \in I$ and each $x \in X_{\alpha}$, if $\alpha \leq \beta$, then we have

$$(\alpha, x) \sim (\beta, f_{\alpha\beta}(x))$$

 $\inf \underset{\alpha \in I}{ \operatorname{colim}}(X_{\alpha}).$

Proof. Taking $\gamma=\beta$, we have $f_{\alpha\gamma}=f_{\alpha\beta}$, we have $f_{\beta\gamma}=f_{\beta\beta}\stackrel{\text{def}}{=} \mathrm{id}_{X_\beta}$, and we have

$$f_{\alpha\beta}(x) = f_{\beta\beta}(f_{\alpha\beta}(x))$$

$$\stackrel{\text{def}}{=} \mathrm{id}_{X_{\beta}}(f_{\alpha\beta}(x)),$$

$$= f_{\alpha\beta}(x).$$

As a result, since $\alpha \leq \beta$ and $\beta \leq \beta$ as well, Items 1a to 1c of Definition 4.2.6.1.2 are met. Thus we have $(\alpha, x) \sim (\beta, f_{\alpha\beta}(x))$.

4.2.6 Direct Colimits

We can now prove Definition 4.2.6.1.2:

Proof. We claim that $\operatornamewithlimits{colim}_{\alpha \in I}(X_\alpha)$ is the colimit of the direct system of sets $(X_\alpha, f_{\alpha\beta})_{\alpha,\beta \in I}$. Commutativity of the Colimit Diagram: First, we need to check that the colimit diagram defined by $\operatornamewithlimits{colim}_{\alpha \in I}(X_\alpha)$ commutes, i.e. that we have

$$\operatorname{inj}_{\alpha} = \operatorname{inj}_{\beta} \circ f_{\alpha\beta}, \qquad \underbrace{\bigcap_{\alpha \in I}^{\operatorname{colim}}(X_{\alpha})}_{\operatorname{inj}_{\alpha}}$$

$$X_{\alpha} \xrightarrow{f_{\alpha\beta}} X_{\beta}$$

for each $\alpha, \beta \in I$ with $\alpha \leq \beta$. Indeed, given $x \in X_{\alpha}$, we have

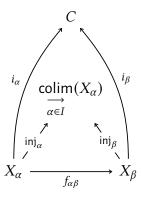
$$[\inf_{\beta} \circ f_{\alpha\beta}](x) \stackrel{\text{def}}{=} \inf_{\beta} (f_{\alpha\beta}(x))$$

$$\stackrel{\text{def}}{=} [(\beta, f_{\alpha\beta}(x))]$$

$$= [(\alpha, x)]$$

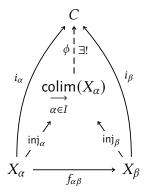
$$\stackrel{\text{def}}{=} \inf_{\alpha} (x),$$

where we have used Definition 4.2.6.1.3 for the third equality. Proof of the Universal Property of the Colimit: Next, we prove that $\operatornamewithlimits{colim}_{\alpha \in I}(X_{\alpha})$ as constructed in Definition 4.2.6.1.2 satisfies the universal property of a direct colimit. Suppose that we have, for each $\alpha, \beta \in I$ with $\alpha \preceq \beta$, a diagram of the form



in Sets. We claim that there exists a unique map $\phi \colon \operatornamewithlimits{colim}_{\alpha \in I}(X_\alpha) \overset{\exists !}{\longrightarrow} C$ making the

diagram



commute. To this end, first consider the diagram

$$\coprod_{\alpha \in I} X_{\alpha} \xrightarrow{\operatorname{pr}} \underset{\alpha \in I}{\operatorname{colim}}(X_{\alpha})$$

$$\coprod_{\alpha \in I} i_{\alpha}$$

$$C.$$

Lemma. If $(\alpha, x) \sim (\beta, y)$, then we have

$$\left[\bigsqcup_{\alpha \in I} i_{\alpha} \right](x) = \left[\bigsqcup_{\alpha \in I} i_{\alpha} \right](y).$$

Proof. Indeed, if $(\alpha, x) \sim (\beta, y)$, then there exists some $\gamma \in I$ satisfying the following conditions:

- 1. We have $\alpha \leq \gamma$.
- 2. We have $\beta \leq \gamma$.
- 3. We have $f_{\alpha\gamma}(x) = f_{\beta\gamma}(y)$.

We then have

$$\left[\bigsqcup_{\alpha \in I} i_{\alpha} \right](x) \stackrel{\text{def}}{=} i_{\alpha}(x)$$

$$\stackrel{\text{def}}{=} [i_{\gamma} \circ f_{\alpha\gamma}](x)$$

$$\stackrel{\text{def}}{=} i_{\gamma}(f_{\alpha\gamma}(x))$$

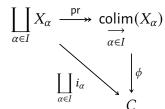
$$= i_{\gamma}(f_{\beta\gamma}(x))$$

$$\stackrel{\text{def}}{=} [i_{\gamma} \circ f_{\beta\gamma}](x)$$

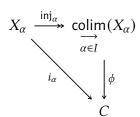
$$= i_{\beta}(y)$$

$$\stackrel{\text{def}}{=} \left[\prod_{\alpha \in I} i_{\alpha} \right](y).$$

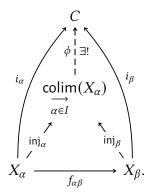
This finishes the proof of the lemma. Continuing, by Conditions on Relations, ?? of Definition 10.6.2.1.3, there then exists a map ϕ : colim $(X_{\alpha}) \stackrel{\exists !}{\longrightarrow} C$ making the diagram



commute. In particular, this implies that the diagram



also commutes, and thus so does the diagram



This finishes the proof.¹⁰

Example 4.2.6.1.4. Here are some examples of direct colimits of sets.

1. The Prüfer Group. The Prüfer group $\mathbb{Z}(p^{\infty})$ is defined as the direct colimit

$$\mathbb{Z}(p^{\infty}) \stackrel{\text{def}}{=} \underset{n \in \mathbb{N}}{\operatorname{colim}}(\mathbb{Z}_{/p^n});$$

see ??.

4.3 Operations With Sets

4.3.1 The Empty Set

Definition 4.3.1.1.1. The **empty set** is the set \emptyset defined by

$$\emptyset \stackrel{\text{def}}{=} \{ x \in X \mid x \neq x \},\$$

where X is the set in the set existence axiom, ?? of ??.

4.3.2 Singleton Sets

Let *X* be a set.

Definition 4.3.2.1.1. The **singleton set containing** X is the set $\{X\}$ defined by

$$\{X\} \stackrel{\text{def}}{=} \{X, X\},$$

where $\{X, X\}$ is the pairing of X with itself of Definition 4.3.3.1.1.

$$\{i_{\alpha} = \phi \circ \operatorname{inj}_{\alpha}\}_{\alpha \in I}$$

show that ϕ must be given by

$$\phi([(\alpha, x)]) = (i_{\alpha}(x))_{\alpha \in I}$$

for each $[(\alpha,x)] \in \underset{\alpha \in I}{\operatorname{colim}}(X_{\alpha})$, although we would need to show that this assignment is well-defined were we to prove Definition 4.2.6.1.2 in this way. Instead, invoking Conditions on Relations, ?? of Definition 10.6.2.1.3 gave us a way to avoid having to prove this, leading to a cleaner alternative proof.

¹⁰Incidentally, the conditions

4.3.3 Pairings of Sets

Let X and Y be sets.

Definition 4.3.3.1.1. The **pairing of** X **and** Y is the set $\{X, Y\}$ defined by

$${X, Y} \stackrel{\text{def}}{=} {x \in A \mid x = X \text{ or } x = Y},$$

where A is the set in the axiom of pairing, ?? of ??.

4.3.4 Ordered Pairs

Let A and B be sets.

Definition 4.3.4.1.1. The **ordered pair associated to** A **and** B is the set (A, B) defined by

$$(A, B) \stackrel{\text{def}}{=} \{ \{A\}, \{A, B\} \}.$$

Proposition 4.3.4.1.2. Let *A* and *B* be sets.

- 1. Uniqueness. Let A, B, C, and D be sets. The following conditions are equivalent:
 - (a) We have (A, B) = (C, D).
 - (b) We have A = C and B = D.

Proof. Item 1, Uniqueness: See [Cie97, Theorem 1.2.3].

4.3.5 Sets of Maps

Let A and B be sets.

Definition 4.3.5.1.1. The **set of maps from** A **to** B^{11} is the set $Sets(A, B)^{12}$ whose elements are the functions from A to B.

Proposition 4.3.5.1.2. Let *A* and *B* be sets.

1. Functoriality. The assignments $X, Y, (X, Y) \mapsto \mathsf{Hom}_{\mathsf{Sets}}(X, Y)$ define func-

¹¹ Further Terminology: Also called the **Hom set from** A **to** B.

¹² Further Notation: Also written $Hom_{Sets}(A, B)$.

tors

$$\begin{array}{lll} \mathsf{Sets}(X,-)\colon & \mathsf{Sets} & \to \mathsf{Sets}, \\ \mathsf{Sets}(-,Y)\colon & \mathsf{Sets}^\mathsf{op} & \to \mathsf{Sets}, \\ \mathsf{Sets}(-_1,-_2)\colon \mathsf{Sets}^\mathsf{op} \times \mathsf{Sets} \to \mathsf{Sets}. \end{array}$$

2. Adjointness. We have adjunctions

$$(A \times - + \mathsf{Sets}(A, -)): \quad \underbrace{\mathsf{Sets}}_{\mathsf{Sets}(A, -)}^{A \times -} \mathsf{Sets},$$

$$(- \times B + \mathsf{Sets}(B, -)): \quad \underbrace{\mathsf{Sets}}_{\mathsf{Sets}(B, -)}^{- \times B} \mathsf{Sets},$$

witnessed by bijections

$$Sets(A \times B, C) \cong Sets(A, Sets(B, C)),$$

 $Sets(A \times B, C) \cong Sets(B, Sets(A, C)),$

natural in $A, B, C \in \mathsf{Obj}(\mathsf{Sets})$.

3. Maps From the Punctual Set. We have a bijection

$$\mathsf{Sets}(\mathsf{pt},A) \cong A$$
,

natural in $A \in Obj(Sets)$.

4. Maps to the Punctual Set. We have a bijection

$$Sets(A, pt) \cong pt$$

natural in $A \in Obj(Sets)$.

Proof. Item 1, Functoriality: This follows from Categories, Items 2 and 5 of Definition 11.1.4.1.2.

Item 2, *Adjointness*: This is a repetition of Item 2 of Definition 4.1.3.1.3 and is proved there.

Item 3, Maps From the Punctual Set: The bijection

$$\Phi_A \colon \mathsf{Sets}(\mathsf{pt},A) \xrightarrow{\sim} A$$

is given by

$$\Phi_A(f) \stackrel{\text{def}}{=} f(\star)$$

for each $f \in Sets(pt, A)$, admitting an inverse

$$\Phi_A^{-1} \colon A \xrightarrow{\sim} \mathsf{Sets}(\mathsf{pt}, A)$$

given by

$$\Phi_A^{-1}(a) \stackrel{\text{def}}{=} \llbracket \star \mapsto a \rrbracket$$

for each $a \in A$. Indeed, we have

$$\begin{split} [\Phi_A^{-1} \circ \Phi_A](f) &\stackrel{\text{def}}{=} \Phi_A^{-1}(\Phi_A(f)) \\ &\stackrel{\text{def}}{=} \Phi_A^{-1}(f(\star)) \\ &\stackrel{\text{def}}{=} [\![\star \mapsto f(\star)]\!] \\ &\stackrel{\text{def}}{=} f \\ &\stackrel{\text{def}}{=} [\![\text{id}_{\mathsf{Sets}(\mathsf{pt},A)}]\!](f) \end{split}$$

for each $f \in \mathsf{Sets}(\mathsf{pt}, A)$ and

$$\begin{split} [\Phi_A \circ \Phi_A^{-1}](a) &\stackrel{\text{def}}{=} \Phi_A(\Phi_A^{-1}(a)) \\ &\stackrel{\text{def}}{=} \Phi_A([\![\star \mapsto a]\!]) \\ &\stackrel{\text{def}}{=} \operatorname{ev}_{\star}([\![\star \mapsto a]\!]) \\ &\stackrel{\text{def}}{=} a \\ &\stackrel{\text{def}}{=} [\operatorname{id}_A](a) \end{split}$$

for each $a \in A$, and thus we have

$$\begin{split} & \Phi_A^{-1} \circ \Phi_A = \mathrm{id}_{\mathsf{Sets}(\mathsf{pt},A)} \\ & \Phi_A \circ \Phi_A^{-1} = \mathrm{id}_A \,. \end{split}$$

To prove naturality, we need to show that the diagram

$$\begin{array}{ccc} \mathsf{Sets}(\mathsf{pt},A) & \xrightarrow{f_!} & \mathsf{Sets}(\mathsf{pt},B) \\ & & \downarrow & & \downarrow & \\ & A & \xrightarrow{f} & B \end{array}$$

commutes. Indeed, we have

$$[f \circ \Phi_A](\phi) \stackrel{\text{def}}{=} f(\Phi_A(\phi))$$

$$\stackrel{\text{def}}{=} f(\phi(\star))$$

$$\stackrel{\text{def}}{=} [f \circ \phi](\star)$$

$$\stackrel{\text{def}}{=} \Phi_B(f \circ \phi)$$

$$\stackrel{\text{def}}{=} \Phi_B(f_!(\phi))$$

$$\stackrel{\text{def}}{=} [\Phi_B \circ f_!](\phi)$$

for each $\phi \in \mathsf{Sets}(\mathsf{pt}, A)$. This finishes the proof.

Item 4, *Maps to the Punctual Set*: This follows from the universal property of pt as the terminal set, Definition 4.1.1.1.

4.3.6 Unions of Families of Subsets

Let X be a set and let $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

Definition 4.3.6.1.1. The **union of** \mathcal{U} is the set $\bigcup_{U \in \mathcal{U}} U$ defined by

$$\bigcup_{U \in \mathcal{U}} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \mathcal{U} \\ \text{such that we have } x \in U \end{array} \right\}.$$

Proposition 4.3.6.1.2. Let *X* be a set.

1. Functoriality. The assignment $\mathcal{U}\mapsto \bigcup_{U\in\mathcal{U}}U$ defines a functor

$$\bigcup : (\mathcal{P}(\mathcal{P}(X)), \subset) \to (\mathcal{P}(X), \subset).$$

In particular, for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$, the following condition is satisfied:

$$(\star) \ \ \mathsf{lf} \, \mathcal{U} \subset \mathcal{V}, \mathsf{then} \bigcup_{U \in \mathcal{U}} U \subset \bigcup_{V \in \mathcal{V}} V.$$

2. Associativity. The diagram

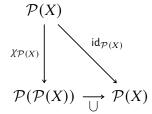
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(\mathcal{P}(X))) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \bigcup} & \mathcal{P}(\mathcal{P}(X)) \\
\cup \star \mathrm{id}_{\mathcal{P}(X)} & & & & & & & & & \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\qquad \qquad \downarrow} & & & & & & & & \\
\end{array}$$

commutes, i.e. we have

$$\bigcup_{U \in \bigcup_{A \in \mathcal{A}} A} U = \bigcup_{A \in \mathcal{A}} (\bigcup_{U \in A} U)$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$.

3. Left Unitality. The diagram

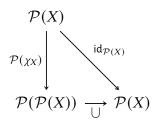


commutes, i.e. we have

$$\bigcup_{V \in \{U\}} V = U$$

for each $U \in \mathcal{P}(X)$.

4. Right Unitality. The diagram



commutes, i.e. we have

$$\bigcup_{\{u\}\in\chi_X(U)}\{u\}=U$$

for each $U \in \mathcal{P}(X)$.

5. Interaction With Unions I. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\cup} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\cup} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{W \in \mathcal{U} \cup \mathcal{V}} W = \left(\bigcup_{U \in \mathcal{U}} U\right) \cup \left(\bigcup_{V \in \mathcal{V}} V\right)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

6. Interaction With Unions II. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcup_{V \in \mathcal{V}} V\right) = \bigcup_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcup_{U \in \mathcal{U}} U\right) \cup V = \bigcup_{U \in \mathcal{U}} (U \cup V)$$

for each nonempty $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

7. Interaction With Intersections I. We have a natural transformation

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\cap} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \qquad \bigcup \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\cap} \mathcal{P}(X),$$

with components

$$\bigcup_{W \in \mathcal{U} \cap \mathcal{V}} W \subset \left(\bigcup_{U \in \mathcal{U}} U\right) \cap \left(\bigcup_{V \in \mathcal{V}} V\right)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

8. Interaction With Intersections II. The diagrams

commute, i.e. we have

$$U \cap \left(\bigcup_{V \in \mathcal{V}} V\right) = \bigcup_{V \in \mathcal{V}} (U \cap V),$$
$$\left(\bigcup_{U \in \mathcal{U}} U\right) \cap V = \bigcup_{U \in \mathcal{U}} (U \cap V)$$

for each $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

9. Interaction With Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\setminus} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\setminus} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in \mathcal{U} \setminus \mathcal{V}} W \neq \left(\bigcup_{U \in \mathcal{U}} U\right) \setminus \left(\bigcup_{V \in \mathcal{V}} V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

10. Interaction With Complements I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\operatorname{op}} \xrightarrow{(-)^{\operatorname{c}}} \mathcal{P}(\mathcal{P}(X))$$

$$\downarrow^{\operatorname{op}} \qquad \qquad \downarrow \cup$$

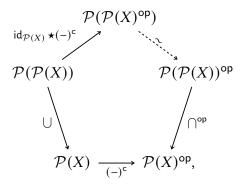
$$\mathcal{P}(X)^{\operatorname{op}} \xrightarrow{(-)^{\operatorname{c}}} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{U \in \mathcal{U}^{\mathsf{c}}} U \neq \bigcup_{U \in \mathcal{U}} U^{\mathsf{c}}$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

11. Interaction With Complements II. The diagram

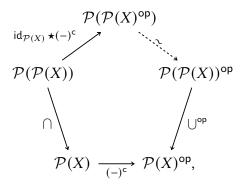


commutes, i.e. we have

$$\left(\bigcup_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

12. Interaction With Complements III. The diagram



commutes, i.e. we have

$$\left(\bigcap_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

13. Interaction With Symmetric Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\Delta} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \times \cup \downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\Delta} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in \mathcal{U} \triangle \mathcal{V}} W \neq \left(\bigcup_{U \in \mathcal{U}} U\right) \triangle \left(\bigcup_{V \in \mathcal{V}} V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

14. Interaction With Internal Homs I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\cup^{\mathsf{op}} \times \cup^{\mathsf{op}} \qquad \qquad \bigcup \cup$$

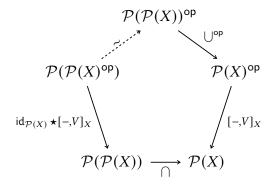
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcup_{U \in \mathcal{U}} U, \bigcup_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

15. Interaction With Internal Homs II. The diagram



commutes, i.e. we have

$$\left[\bigcup_{U\in\mathcal{U}}U,V\right]_{Y}=\bigcap_{U\in\mathcal{U}}[U,V]_{X}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

16. Interaction With Internal Homs III. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{\bigcup} \mathcal{P}(X)$$

$$id_{\mathcal{P}(X)} \star [U,-]_X \downarrow \qquad \qquad \downarrow [U,-]_X$$

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{\bigcup} \mathcal{P}(X)$$

commutes, i.e. we have

$$\left[U, \bigcup_{V \in \mathcal{V}} V\right]_X = \bigcup_{V \in \mathcal{V}} [U, V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

17. Interaction With Direct Images. Let $f\colon X\to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_!)_!} \mathcal{P}(\mathcal{P}(Y))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \cup$$

$$\mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_!(U) = \bigcup_{V \in f_!(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

18. Interaction With Inverse Images. Let $f\colon X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \qquad \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{V \in \mathcal{V}} f^{-1}(V) = \bigcup_{U \in f^{-1}(\mathcal{U})} U$$

for each $\mathcal{V} \in \mathcal{P}(Y)$, where $f^{-1}(\mathcal{V}) \stackrel{\text{def}}{=} (f^{-1})^{-1}(\mathcal{V})$.

19. Interaction With Codirect Images. Let $f: X \to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \bigcup$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_*(U) = \bigcup_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

20. Interaction With Intersections of Families I. The diagram

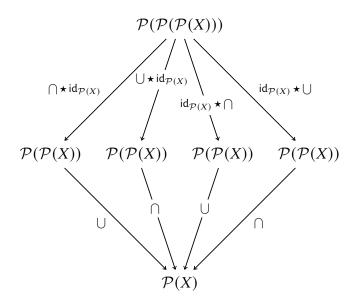
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \bigcap} & \mathcal{P}(\mathcal{P}(x)) \\
\downarrow \star \mathrm{id}_{\mathcal{P}(X)} & & & \downarrow \cap \\
\mathcal{P}(X) & \xrightarrow{\bigcap} & X
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U = \bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right)$$

for each $A \in \mathcal{P}(\mathcal{P}(X))$.

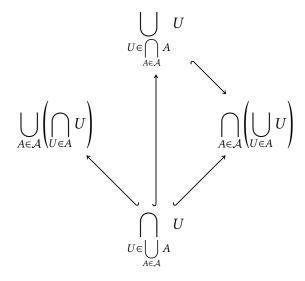
21. Interaction With Intersections of Families II. Let X be a set and consider the compositions



$$\mathcal{A} \mapsto \bigcup_{U \in \bigcap_{A \in \mathcal{A}} A} U, \quad \mathcal{A} \mapsto \bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U,$$

$$\mathcal{A} \mapsto \bigcup_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U\right), \quad \mathcal{A} \mapsto \bigcap_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U\right)$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$. We have the following inclusions:



All other possible inclusions fail to hold in general.

Proof. Item 1, Functoriality: Since $\mathcal{P}(X)$ is posetal, it suffices to prove the condition (\star) . So let $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ with $\mathcal{U} \subset \mathcal{V}$. We claim that

$$\bigcup_{U\in\mathcal{U}}U\subset\bigcup_{V\in\mathcal{V}}V.$$

Indeed, given $x \in \bigcup_{U \in \mathcal{U}} U$, there exists some $U \in \mathcal{U}$ such that $x \in U$, but since $\mathcal{U} \subset \mathcal{V}$, we have $U \in \mathcal{V}$ as well, and thus $x \in \bigcup_{V \in \mathcal{V}} V$, which gives our desired inclusion.

Item 2, Associativity: We have

$$\bigcup_{U \in \bigcup_{A \in A} A} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{c} \text{there exists some } U \in \bigcup_{A \in \mathcal{A}} A \\ \text{such that we have } x \in U \end{array} \right\}$$

$$= \begin{cases} x \in X & \text{there exists some } A \in \mathcal{A} \\ \text{and some } U \in A \text{ such that} \\ \text{we have } x \in U \end{cases}$$

$$= \begin{cases} x \in X & \text{there exists some } A \in \mathcal{A} \\ \text{such that we have } x \in \bigcup_{U \in A} U \end{cases}$$

$$\stackrel{\text{def}}{=} \bigcup_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U \right).$$

This finishes the proof.

Item 3, Left Unitality: We have

$$\bigcup_{V \in \{U\}} V \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } V \in \{U\} \\ \text{such that we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \mid x \in U \right\}$$

$$= U.$$

This finishes the proof.

Item 4, Right Unitality: We have

$$\bigcup_{\{u\} \in \chi_X(U)} \{u\} \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } \{u\} \in \chi_X(U) \right\}$$
 such that we have $x \in \{u\}$

$$= \left\{ x \in X \middle| \text{ there exists some } \{u\} \in \chi_X(U) \right\}$$
 such that we have $x = u$

$$= \left\{ x \in X \middle| \text{ there exists some } u \in U \right\}$$
 such that we have $x = u$

$$= \left\{ x \in X \middle| \text{ x } \in U \right\}$$

$$= \{ x \in X \mid x \in U \}$$

$$= U.$$

This finishes the proof.

Item 5, Interaction With Unions I: We have

$$\bigcup_{W \in \mathcal{U} \cup \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \text{ there exists some } W \in \mathcal{U} \cup \mathcal{V} \right\}$$
such that we have $x \in W$

$$= \begin{cases} x \in X & \text{there exists some } W \in \mathcal{U} \text{ or some} \\ W \in \mathcal{V} \text{ such that we have } x \in W \end{cases}$$

$$\stackrel{\text{def}}{=} \begin{cases} x \in X & \text{there exists some } W \in \mathcal{U} \\ \text{such that we have } x \in W \end{cases}$$

$$\cup \begin{cases} x \in X & \text{there exists some } W \in \mathcal{V} \\ \text{such that we have } x \in W \end{cases}$$

$$\stackrel{\text{def}}{=} \left(\bigcup_{W \in \mathcal{U}} W \right) \cup \left(\bigcup_{W \in \mathcal{V}} W \right)$$

$$= \left(\bigcup_{U \in \mathcal{U}} U \right) \cup \left(\bigcup_{W \in \mathcal{V}} V \right).$$

This finishes the proof.

Item 6, Interaction With Unions II: Assume \mathcal{V} is nonempty. We have

$$U \cup \bigcup_{V \in \mathcal{V}} V \stackrel{\text{def}}{=} \left\{ x \in X \middle| x \in U \text{ or } x \in \bigcup_{V \in \mathcal{V}} V \right\}$$

$$= \left\{ x \in X \middle| x \in U \text{ or there exists some} \right\}$$

$$= \left\{ x \in X \middle| \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \middle| \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \middle| \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \middle| \text{there exists some } V \in \mathcal{V} \right\}$$

$$\stackrel{\text{def}}{=} \bigcup_{V \in \mathcal{V}} U \cup V.$$

This concludes the proof of the first statement. For the second statement, use Item 4 of Definition 4.3.8.1.2 to rewrite

$$\left(\bigcup_{U \in \mathcal{U}} U\right) \cup V = V \cup \left(\bigcup_{U \in \mathcal{U}} U\right),$$

$$\bigcup_{U \in \mathcal{U}} (U \cup V) = \bigcup_{U \in \mathcal{U}} (V \cup U).$$

But these two sets are equal by the first statement.

Item 7, Interaction With Intersections I: We have

there exists some
$$W \in \mathcal{U} \cap \mathcal{V}$$
 such that we have $x \in W$

$$\subset \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \mathcal{U} \text{ and some } V \in \mathcal{V} \\ \text{such that we have } x \in U \text{ and } x \in V \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \mathcal{U} \text{ and } x \in V \\ \text{such that we have } x \in U \text{ and } x \in V \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } U \in \mathcal{U} \\ \text{such that we have } x \in U \end{array} \right\}$$

$$\cup \left\{ x \in X \middle| \begin{array}{l} \text{there exists some } V \in \mathcal{V} \\ \text{such that we have } x \in V \right\}$$

$$\stackrel{\text{def}}{=} \left(\bigcup_{U \in \mathcal{U}} U \right) \cap \left(\bigcup_{V \in \mathcal{V}} V \right).$$

This finishes the proof.

Item 8, Interaction With Intersections II: We have

$$U \cap \bigcup_{V \in \mathcal{V}} V \stackrel{\text{def}}{=} \left\{ x \in X \middle| x \in U \text{ and } x \in \bigcup_{V \in \mathcal{V}} V \right\}$$

$$= \left\{ x \in X \middle| x \in U \text{ and there exists some} \right\}$$

$$= \left\{ x \in X \middle| \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \middle| \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \middle| \text{there exists some } V \in \mathcal{V} \right\}$$

$$= \left\{ x \in X \middle| \text{there exists some } V \in \mathcal{V} \right\}$$

$$\stackrel{\text{def}}{=} \bigcup_{V \in \mathcal{V}} U \cap V.$$

This concludes the proof of the first statement. For the second statement, use Item 5 of Definition 4.3.9.1.2 to rewrite

$$\left(\bigcup_{U \in \mathcal{U}} U\right) \cap V = V \cap \left(\bigcup_{U \in \mathcal{U}} U\right),$$

$$\bigcup_{U\in\mathcal{U}}(U\cap V)=\bigcup_{U\in\mathcal{U}}(V\cap U).$$

But these two sets are equal by the first statement.

Item 9, Interaction With Differences: Let $X=\{0,1\}$, let $\mathcal{U}=\{\{0,1\}\}$, and let $\mathcal{V}=\{\{0\}\}$. We have

$$\bigcup_{W \in \mathcal{U} \setminus \mathcal{V}} U = \bigcup_{W \in \{\{0,1\}\}} W$$
$$= \{0,1\},$$

whereas

$$\left(\bigcup_{U \in \mathcal{U}} U\right) \setminus \left(\bigcup_{V \in \mathcal{V}} V\right) = \{0, 1\} \setminus \{0\}$$
$$= \{1\}.$$

Thus we have

$$\bigcup_{W\in\mathcal{U}\setminus\mathcal{V}}W=\left\{0,1\right\}\neq\left\{1\right\}=\left(\bigcup_{U\in\mathcal{U}}U\right)\setminus\left(\bigcup_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

Item 10, Interaction With Complements I: Let $X = \{0, 1\}$ and let $\mathcal{U} = \{0\}$. We have

$$\bigcup_{U \in \mathcal{U}^{c}} U = \bigcup_{U \in \{\emptyset, \{1\}, \{0,1\}\}} U$$
$$= \{0, 1\},$$

whereas

$$\bigcup_{U \in \mathcal{U}} U^{c} = \{0\}^{c}$$
$$= \{1\}.$$

Thus we have

$$\bigcup_{U\in\mathcal{U}^\mathsf{c}}U=\left\{0,1\right\}\neq\left\{1\right\}=\bigcup_{U\in\mathcal{U}}U^\mathsf{c}.$$

This finishes the proof.

Item 11, Interaction With Complements II: We have

$$\left(\bigcup_{U \in \mathcal{U}} U\right)^{\mathbf{c}} \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{there exists no } U \in \mathcal{U} \\ \text{such that we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for all } U \in \mathcal{U} \\ \text{we have } x \notin U \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for all } U \in \mathcal{U} \\ \text{we have } x \in U^{\mathbf{c}} \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{U \in \mathcal{U}} U^{\mathbf{c}}.$$

Item 12, *Interaction With Complements III*: By Item 11 Item 3 of Definition 4.3.11.1.2, we have

$$\left(\bigcap_{U \in \mathcal{U}} U\right)^{c} = \left(\bigcap_{U \in \mathcal{U}} (U^{c})^{c}\right)^{c}$$

$$= \left(\left(\bigcup_{U \in \mathcal{U}} U^{c}\right)^{c}\right)^{c}$$

$$= \bigcup_{U \in \mathcal{U}} U^{c}.$$

Item 13, Interaction With Symmetric Differences: Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}$, and let $\mathcal{V} = \{\{0\}, \{0, 1\}\}$. We have

$$\bigcup_{W \in \mathcal{U} \triangle \mathcal{V}} W = \bigcup_{W \in \{\{0\}\}} W$$
$$= \{0\},$$

whereas

$$\left(\bigcup_{U\in\mathcal{U}}U\right)\triangle\left(\bigcup_{V\in\mathcal{V}}V\right)=\left\{0,1\right\}\triangle\left\{0,1\right\}$$
$$=\emptyset,$$

Thus we have

$$\bigcup_{W\in\mathcal{U}\triangle\mathcal{V}}W=\left\{0\right\}\neq\emptyset=\left(\bigcup_{U\in\mathcal{U}}U\right)\triangle\left(\bigcup_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

Item 14, *Interaction With Internal Homs I*: This is a repetition of Item 7 of Definition 4.4.7.1.3 and is proved there.

Item 15, *Interaction With Internal Homs II*: This is a repetition of Item 8 of Definition 4.4.7.1.3 and is proved there.

Item 16, Interaction With Internal Homs III: This is a repetition of Item 9 of Definition 4.4.7.1.3 and is proved there.

Item 17, *Interaction With Direct Images*: This is a repetition of Item 3 of Definition 4.6.1.1.5 and is proved there.

Item 18, *Interaction With Inverse Images*: This is a repetition of Item 3 of Definition 4.6.2.1.3 and is proved there.

Item 19, *Interaction With Codirect Images*: This is a repetition of Item 3 of Definition 4.6.3.1.7 and is proved there.

Item 20, Interaction With Intersections of Families I: We have

$$\bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \bigcup_{A \in \mathcal{A}} A, \\ \text{we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } A \in \mathcal{A} \text{ and each} \\ U \in A, \text{ we have } x \in U \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right).$$

This finishes the proof.

Item 21, Interaction With Intersections of Families II: Omitted.

4.3.7 Intersections of Families of Subsets

Let X be a set and let $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

Definition 4.3.7.1.1. The **intersection of** \mathcal{U} is the set $\bigcap_{U \in \mathcal{U}} U$ defined by

$$\bigcap_{U \in \mathcal{U}} U \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \mathcal{U}, \\ \text{we have } x \in U \end{array} \right\}.$$

Proposition 4.3.7.1.2. Let *X* be a set.

1. Functoriality. The assignment $\mathcal{U} \mapsto \bigcap_{U \in \mathcal{U}} U$ defines a functor

$$\bigcap \colon (\mathcal{P}(\mathcal{P}(X)), \supset) \to (\mathcal{P}(X), \subset).$$

In particular, for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$, the following condition is satisfied:

$$(\star) \ \ \mathsf{lf} \, \mathcal{U} \subset \mathcal{V}, \mathsf{then} \bigcap_{V \in \mathcal{V}} V \subset \bigcap_{U \in \mathcal{U}} U.$$

2. Oplax Associativity. We have a natural transformation

$$\begin{array}{c|c}
\mathcal{P}(\mathcal{P}(\mathcal{P}(X))) & \xrightarrow{\mathrm{id}_{\mathcal{P}(X)} \star \bigcap} & \mathcal{P}(\mathcal{P}(X)) \\
\cap \star \mathrm{id}_{\mathcal{P}(X)} & & & & & & & & \\
\mathcal{P}(\mathcal{P}(X)) & & & & & & & & \\
\end{array}$$

with components

$$\bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right) \subset \bigcap_{U \in \bigcap_{A \in A} A} U$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$.

3. Left Unitality. The diagram

$$\begin{array}{c|c} \mathcal{P}(X) & & \\ \downarrow^{\chi_{\mathcal{P}(X)}} & & \text{id}_{\mathcal{P}(X)} \\ & & \mathcal{P}(\mathcal{P}(X)) & \longrightarrow \mathcal{P}(X) \end{array}$$

commutes, i.e. we have

$$\bigcap_{V \in \{U\}} V = U.$$

for each $U \in \mathcal{P}(X)$.

4. Oplax Right Unitality. The diagram

$$\begin{array}{c|c}
\mathcal{P}(X) & & \text{id}_{\mathcal{P}(X)} \\
\downarrow & & \times \\
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{} \mathcal{P}(X)
\end{array}$$

does not commute in general, i.e. we may have

$$\bigcap_{\{x\}\in\chi_X(U)}\{x\}\neq U$$

in general, where $U \in \mathcal{P}(X)$. However, when U is nonempty, we have

$$\bigcap_{\{x\}\in\chi_X(U)}\{x\}\subset U.$$

5. Interaction With Unions I. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\cup} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \times \cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\cap} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcap_{W\in\mathcal{U}\cup\mathcal{V}}W=\left(\bigcap_{U\in\mathcal{U}}U\right)\cap\left(\bigcap_{V\in\mathcal{V}}V\right)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

6. Interaction With Unions II. The diagram

commute, i.e. we have

$$U \cup \left(\bigcap_{V \in \mathcal{V}} V\right) = \bigcap_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cup V = \bigcap_{U \in \mathcal{U}} (U \cup V)$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ and each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(X)$.

7. Interaction With Intersections I. We have a natural transformation

with components

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\cap\left(\bigcap_{V\in\mathcal{V}}V\right)\subset\bigcap_{W\in\mathcal{U}\cap\mathcal{V}}W$$

for each $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

8. Interaction With Intersections II. The diagrams

commute, i.e. we have

$$U \cup \left(\bigcap_{V \in \mathcal{V}} V\right) = \bigcap_{V \in \mathcal{V}} (U \cup V),$$
$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cup V = \bigcap_{U \in \mathcal{U}} (U \cup V)$$

for each $U, V \in \mathcal{P}(\mathcal{P}(X))$ and each $U, V \in \mathcal{P}(X)$.

9. Interaction With Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\setminus} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \times \cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\setminus} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W \in \mathcal{U} \setminus \mathcal{V}} W \neq \left(\bigcap_{U \in \mathcal{U}} U\right) \setminus \left(\bigcap_{V \in \mathcal{V}} V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

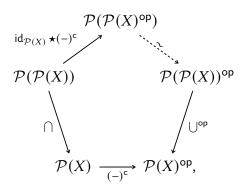
10. Interaction With Complements I. The diagram

does not commute in general, i.e. we may have

$$\bigcap_{W \in \mathcal{U}^{\mathsf{c}}} W \neq \bigcap_{U \in \mathcal{U}} U^{\mathsf{c}}$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

11. Interaction With Complements II. The diagram

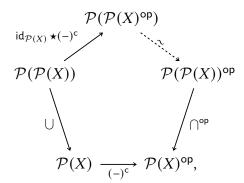


commutes, i.e. we have

$$\left(\bigcap_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

12. Interaction With Complements III. The diagram



commutes, i.e. we have

$$\left(\bigcup_{U\in\mathcal{U}}U\right)^{\mathsf{c}}=\bigcap_{U\in\mathcal{U}}U^{\mathsf{c}}$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

13. Interaction With Symmetric Differences. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{\Delta} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \times \cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\wedge} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W\in\mathcal{U}\triangle\mathcal{V}}W\neq\left(\bigcap_{U\in\mathcal{U}}U\right)\triangle\left(\bigcap_{V\in\mathcal{V}}V\right)$$

in general, where $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

14. Interaction With Internal Homs I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-_{1}, -_{2}]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\bigcap^{\mathsf{op}} \times \bigcap^{\mathsf{op}} \bigvee \qquad \bigvee \qquad \qquad \bigcup \bigcap$$

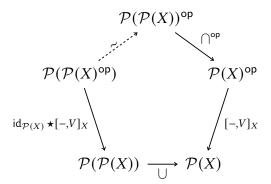
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-_{1}, -_{2}]_{X}} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcap_{U \in \mathcal{U}} U, \bigcap_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

15. Interaction With Internal Homs II. The diagram



commutes, i.e. we have

$$\left[\bigcap_{U\in\mathcal{U}}U,V\right]_X=\bigcup_{U\in\mathcal{U}}[U,V]_X$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

16. Interaction With Internal Homs III. The diagram

$$\begin{array}{ccc} \mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcap}{\longrightarrow} \mathcal{P}(X) \\ \operatorname{id}_{\mathcal{P}(X)} \star [U,-]_X & & & & |_{[U,-]_X} \\ \mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcap}{\longrightarrow} \mathcal{P}(X) & & & \end{array}$$

commutes, i.e. we have

$$\left[U, \bigcap_{V \in \mathcal{V}} V\right]_{X} = \bigcap_{V \in \mathcal{V}} [U, V]_{X}$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

17. Interaction With Direct Images. Let $f\colon X\to Y$ be a map of sets. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_{1})_{1}} & \mathcal{P}(\mathcal{P}(Y)) \\
& & & \downarrow \cap \\
\mathcal{P}(X) & \xrightarrow{f_{1}} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \mathcal{U}} f_!(U) = \bigcap_{V \in f_!(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

18. Interaction With Inverse Images. Let $f\colon X\to Y$ be a map of sets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\cap \downarrow \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcap_{V \in \mathcal{V}} f^{-1}(V) = \bigcap_{U \in f^{-1}(\mathcal{U})} U$$

for each $\mathcal{V} \in \mathcal{P}(Y)$, where $f^{-1}(\mathcal{V}) \stackrel{\text{def}}{=} (f^{-1})^{-1}(\mathcal{V})$.

19. Interaction With Codirect Images. Let $f\colon X\to Y$ be a map of sets. The diagram

commutes, i.e. we have

$$\bigcap_{U \in \mathcal{U}} f_*(U) = \bigcap_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

20. Interaction With Unions of Families I. The diagram

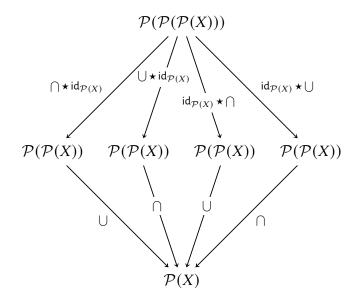
$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{\operatorname{id}_{\mathcal{P}(X)} \star \bigcap} & \mathcal{P}(\mathcal{P}(x)) \\
\downarrow \star \operatorname{id}_{\mathcal{P}(X)} & & & \downarrow \cap \\
\mathcal{P}(X) & \xrightarrow{\bigcap} & X
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U = \bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right)$$

for each $A \in \mathcal{P}(\mathcal{P}(X))$.

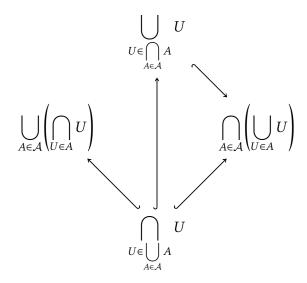
21. Interaction With Unions of Families II. Let X be a set and consider the compositions



$$\mathcal{A} \mapsto \bigcup_{U \in \bigcap_{A \in \mathcal{A}} A} U, \quad \mathcal{A} \mapsto \bigcap_{U \in \bigcup_{A \in \mathcal{A}} A} U,$$

$$\mathcal{A} \mapsto \bigcup_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U\right), \quad \mathcal{A} \mapsto \bigcap_{A \in \mathcal{A}} \left(\bigcup_{U \in A} U\right)$$

for each $A \in \mathcal{P}(\mathcal{P}(\mathcal{P}(X)))$. We have the following inclusions:



All other possible inclusions fail to hold in general.

Proof. Item 1, Functoriality: Since $\mathcal{P}(X)$ is posetal, it suffices to prove the condition (\star) . So let $\mathcal{U}, \mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$ with $\mathcal{U} \subset \mathcal{V}$. We claim that

$$\bigcap_{V\in\mathcal{V}}V\subset\bigcap_{U\in\mathcal{U}}U.$$

Indeed, if $x \in \bigcap_{V \in \mathcal{V}} V$, then $x \in V$ for all $V \in \mathcal{V}$. But since $\mathcal{U} \subset \mathcal{V}$, it follows that $x \in U$ for all $U \in \mathcal{U}$ as well. Thus $x \in \bigcap_{U \in \mathcal{U}} U$, which gives our desired inclusion. *Item* 2, *Oplax Associativity*: We have

$$\bigcap_{A \in \mathcal{A}} \left(\bigcap_{U \in A} U \right) \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } A \in \mathcal{A}, \\ \text{we have } x \in \bigcap_{U \in A} U \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } A \in \mathcal{A} \text{ and each} \\ U \in A, \text{ we have } x \in U \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \bigcup_{A \in \mathcal{A}} A, \\ \text{we have } x \in U \end{array} \right\}$$

$$\subset \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \bigcap_{A \in \mathcal{A}} A, \\ \text{we have } x \in U \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{A \in \mathcal{A}} U.$$

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 3, Left Unitality: We have

$$\bigcap_{V \in \{U\}} V \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } V \in \{U\}, \\ \text{we have } x \in U \end{array} \right\}$$
$$= \left\{ x \in X \middle| x \in U \right\}$$
$$= U$$

This finishes the proof.

Item 4, Oplax Right Unitality: If $U = \emptyset$, then we have

$$\bigcap_{\{u\}\in\chi_X(U)} \{u\} = \bigcap_{\{u\}\in\emptyset} \{u\}$$
$$= X,$$

so $\bigcap_{\{u\}\in\chi_X(U)}\{u\}=X
eq\emptyset=U$. When U is nonempty, we have two cases:

1. If U is a singleton, say $U = \{u\}$, we have

$$\bigcap_{\{u\}\in\chi_X(U)} \{u\} = \{u\}$$

$$\stackrel{\text{def}}{=} U.$$

2. If U contains at least two elements, we have

$$\bigcap_{\{u\}\in\chi_X(U)}\{u\}=\emptyset$$

$$\subset U.$$

This finishes the proof.

Item 5, Interaction With Unions I: We have

$$\bigcap_{W \in \mathcal{U} \cup \mathcal{V}} W \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cup \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \text{ and each } \\ W \in \mathcal{V}, \text{ we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\cap \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \left(\bigcap_{W \in \mathcal{U}} W \right) \cap \left(\bigcap_{W \in \mathcal{V}} W \right)$$

$$= \left(\bigcap_{W \in \mathcal{U}} U \right) \cap \left(\bigcap_{W \in \mathcal{V}} V \right).$$

This finishes the proof.

Item 6, Interaction With Unions II: Omitted.

Item 7, Interaction With Intersections I: We have

$$\left(\bigcap_{U \in \mathcal{U}} U\right) \cap \left(\bigcap_{V \in \mathcal{V}} V\right) \stackrel{\text{def}}{=} \left\{ x \in X \middle| \begin{array}{l} \text{for each } U \in \mathcal{U}, \\ \text{we have } x \in U \end{array} \right\}$$

$$\cup \left\{ x \in X \middle| \begin{array}{l} \text{for each } V \in \mathcal{V}, \\ \text{we have } x \in V \end{array} \right\}$$

$$= \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cap \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\subset \left\{ x \in X \middle| \begin{array}{l} \text{for each } W \in \mathcal{U} \cup \mathcal{V}, \\ \text{we have } x \in W \end{array} \right\}$$

$$\stackrel{\text{def}}{=} \bigcap_{W \in \mathcal{U} \cap \mathcal{V}} W.$$

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 8, Interaction With Intersections II: Omitted.

Item 9, Interaction With Differences: Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0\}, \{0, 1\}\}$, and let $\mathcal{V} = \{\{0\}\}$. We have

$$\bigcap_{W \in \mathcal{U} \setminus \mathcal{V}} U = \bigcap_{W \in \{\{0,1\}\}} W$$
$$= \{0,1\},$$

whereas

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\setminus\left(\bigcap_{V\in\mathcal{V}}V\right)=\{0\}\setminus\{0\}$$
$$=\emptyset.$$

Thus we have

$$\bigcap_{W\in\mathcal{U}\backslash\mathcal{V}}W=\left\{0,1\right\}\neq\emptyset=\left(\bigcap_{U\in\mathcal{U}}U\right)\backslash\left(\bigcap_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

Item 10, Interaction With Complements I: Let $X = \{0, 1\}$ and let $\mathcal{U} = \{\{0\}\}$. We have

$$\bigcap_{W \in \mathcal{U}^{c}} U = \bigcap_{W \in \{\emptyset, \{1\}, \{0,1\}\}} W$$

$$= \emptyset.$$

whereas

$$\bigcap_{U \in \mathcal{U}} U^{c} = \{0\}^{c}$$
$$= \{1\}.$$

Thus we have

$$\bigcap_{W\in\mathcal{U}^\mathsf{c}}U=\emptyset\neq\{1\}=\bigcap_{U\in\mathcal{U}}U^\mathsf{c}.$$

This finishes the proof.

Item 11, *Interaction With Complements II*: This is a repetition of Item 12 of Definition 4.3.6.1.2 and is proved there.

Item 12, *Interaction With Complements III*: This is a repetition of Item 11 of Definition 4.3.6.1.2 and is proved there.

Item 13, Interaction With Symmetric Differences: Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}$, and let $\mathcal{V} = \{\{0\}, \{0, 1\}\}$. We have

$$\bigcap_{W \in \mathcal{U} \triangle \mathcal{V}} W = \bigcap_{W \in \{\{0\}\}} W$$
$$= \{0\},$$

whereas

$$\left(\bigcap_{U\in\mathcal{U}}U\right)\triangle\left(\bigcap_{V\in\mathcal{V}}V\right)=\{0,1\}\triangle\{0\}$$
$$=\emptyset.$$

Thus we have

$$\bigcap_{W\in\mathcal{U}\triangle\mathcal{V}}W=\left\{0\right\}\neq\emptyset=\left(\bigcap_{U\in\mathcal{U}}U\right)\triangle\left(\bigcap_{V\in\mathcal{V}}V\right).$$

This finishes the proof.

Item 14, Interaction With Internal Homs I: This is a repetition of Item 10 of Definition 4.4.7.1.3 and is proved there.

Item 15, *Interaction With Internal Homs II*: This is a repetition of Item 11 of Definition 4.4.7.1.3 and is proved there.

Item 16, Interaction With Internal Homs III: This is a repetition of Item 12 of Definition 4.4.7.1.3 and is proved there.

Item 17, *Interaction With Direct Images*: This is a repetition of Item 4 of Definition 4.6.1.1.5 and is proved there.

Item 18, *Interaction With Inverse Images*: This is a repetition of Item 4 of Definition 4.6.2.1.3 and is proved there.

Item 19, *Interaction With Codirect Images*: This is a repetition of Item 4 of Definition 4.6.3.1.7 and is proved there.

Item 20, *Interaction With Unions of Families I*: This is a repetition of Item 20 of Definition 4.3.6.1.2 and is proved there.

Item 21, Interaction With Unions of Families II: This is a repetition of Item 21 of Definition 4.3.6.1.2 and is proved there.

4.3.8 Binary Unions

Let X be a set and let $U, V \in \mathcal{P}(X)$.

Definition 4.3.8.1.1. The union of U and V is the set $U \cup V$ defined by

$$\begin{split} U \cup V &\stackrel{\text{def}}{=} \bigcup_{z \in \{U,V\}} z \\ &\stackrel{\text{def}}{=} \{x \in X \,|\, x \in U \text{ or } x \in V\}. \end{split}$$

Proposition 4.3.8.1.2. Let *X* be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cup V$ define functors

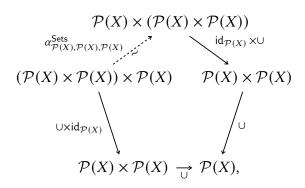
$$U \cup -: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$- \cup V: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$-_1 \cup -_2: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset).$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $U \cup V \subset A \cup V$.
- (b) If $V \subset B$, then $U \cup V \subset U \cup B$.
- (c) If $U \subset A$ and $V \subset B$, then $U \cup V \subset A \cup B$.
- 2. Associativity. The diagram

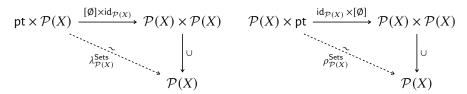


commutes, i.e. we have an equality of sets

$$(U \cup V) \cup W = U \cup (V \cup W)$$

for each $U, V, W \in \mathcal{P}(X)$.

3. Unitality. The diagrams



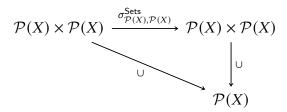
commute, i.e. we have equalities of sets

$$\emptyset \cup U = U,$$

 $U \cup \emptyset = U$

for each $U \in \mathcal{P}(X)$.

4. Commutativity. The diagram

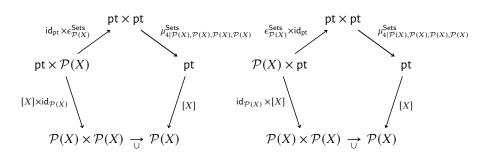


commutes, i.e. we have an equality of sets

$$U \cup V = V \cup U$$

for each $U, V \in \mathcal{P}(X)$.

5. Annihilation With X. The diagrams

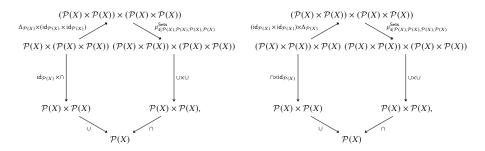


commute, i.e. we have equalities of sets

$$U \cup X = X,$$
$$X \cup V = X$$

for each $U, V \in \mathcal{P}(X)$.

6. Distributivity of Unions Over Intersections. The diagrams



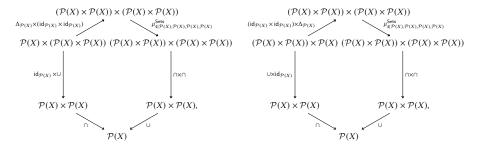
commute, i.e. we have equalities of sets

$$U \cup (V \cap W) = (U \cup V) \cap (U \cup W),$$

$$(U \cap V) \cup W = (U \cup W) \cap (V \cup W)$$

for each $U, V, W \in \mathcal{P}(X)$.

7. Distributivity of Intersections Over Unions. The diagrams



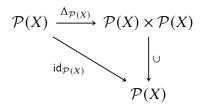
commute, i.e. we have equalities of sets

$$U \cap (V \cup W) = (U \cap V) \cup (U \cap W),$$

$$(U \cup V) \cap W = (U \cap W) \cup (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

8. Idempotency. The diagram



commutes, i.e. we have an equality of sets

$$U \cup U = U$$

for each $U \in \mathcal{P}(X)$.

9. Via Intersections and Symmetric Differences. The diagram

commutes, i.e. we have an equality of sets

$$U \cup V = (U \triangle V) \triangle (U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

10. Interaction With Characteristic Functions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

11. Interaction With Characteristic Functions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each $U, V \in \mathcal{P}(X)$.

12. Interaction With Direct Images. Let $f: X \to Y$ be a function. The diagram

$$\begin{array}{ccc} \mathcal{P}(X) \times \mathcal{P}(X) & \xrightarrow{f_! \times f_!} & \mathcal{P}(Y) \times \mathcal{P}(Y) \\ & & & \downarrow & & \downarrow \\ \mathcal{P}(X) & \xrightarrow{f_!} & \mathcal{P}(Y) \end{array}$$

commutes, i.e. we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$

for each $U, V \in \mathcal{P}(X)$.

13. Interaction With Inverse Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

14. Interaction With Codirect Images. Let $f\colon X\to Y$ be a function. We have a natural transformation

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_* \times f_*} \mathcal{P}(Y) \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

with components

$$f_*(U) \cup f_*(V) \subset f_*(U \cup V)$$

indexed by $U, V \in \mathcal{P}(X)$.

15. Interaction With Powersets and Semirings. The quintuple $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$ is an idempotent commutative semiring.

Proof. Item 1, *Functoriality*: See [Pro25an].

Item 2, Associativity: See [Pro25ba].

Item 3, Unitality: This follows from [Pro25bd] and Item 4.

Item 4, Commutativity: See [Pro25bb].

Item 5, Annihilation With X: We have

$$U \cup X \stackrel{\text{def}}{=} \{x \in X \mid x \in U \text{ or } x \in X\}$$
$$= \{x \in X \mid x \in X\},$$
$$= X$$

and

$$X \cup V \stackrel{\text{def}}{=} \{x \in X \mid x \in X \text{ or } x \in V\}$$
$$= \{x \in X \mid x \in X\}$$
$$= X.$$

This finishes the proof.

Item 6, Distributivity of Unions Over Intersections: See [Pro25az].

Item 7, *Distributivity of Intersections Over Unions*: See [Pro25aj].

Item 8, Idempotency: See [Pro25am].

Item 9, Via Intersections and Symmetric Differences: See [Pro25ay].

Item 10, *Interaction With Characteristic Functions I*: See [Pro25h].

Item 11, Interaction With Characteristic Functions II: See [Pro25h].

Item 12, Interaction With Direct Images: See [Pro25p].

Item 13, Interaction With Inverse Images: See [Pro25y].

Item 14, Interaction With Codirect Images: This is a repetition of Item 5 of Definition 4.6.3.1.7 and is proved there.

Item 15, Interaction With Powersets and Semirings: This follows from Items 2 to 4 and 8 of this proposition and Items 3 to 6 and 8 of Definition 4.3.9.1.2.

4.3.9 Binary Intersections

Let X be a set and let $U, V \in \mathcal{P}(X)$.

Definition 4.3.9.1.1. The **intersection of** U **and** V is the set $U \cap V$ defined by

$$\begin{split} U \cap V &\stackrel{\mathrm{def}}{=} \bigcap_{z \in \{U,V\}} z \\ &\stackrel{\mathrm{def}}{=} \{x \in X \,|\, x \in U \text{ or } x \in V\}. \end{split}$$

Proposition 4.3.9.1.2. Let *X* be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cap V$ define functors

$$U \cap -: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$- \cap V: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$-_{1} \cap -_{2}: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset).$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $U \cap V \subset A \cap V$.
- (b) If $V \subset B$, then $U \cap V \subset U \cap B$.
- (c) If $U \subset A$ and $V \subset B$, then $U \cap V \subset A \cap B$.
- 2. Adjointness. We have adjunctions

$$(U \cap - + [U, -]_X): \quad \mathcal{P}(X) \xrightarrow{\downarrow} \mathcal{P}(X),$$

$$(- \cap V + [V, -]_X): \quad \mathcal{P}(X) \xrightarrow{\downarrow} \mathcal{P}(X),$$

witnessed by bijections

$$\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U\cap V,W)\cong\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U,[V,W]_X),$$

 $\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U\cap V,W)\cong\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(V,[U,W]_X),$

natural in $U, V, W \in \mathcal{P}(X)$, where

$$[-1,-2]_X \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

is the bifunctor of Section 4.4.7. In particular, the following statements hold for each $U, V, W \in \mathcal{P}(X)$:

- (a) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $U \subset [V, W]_X$.
- (b) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $V \subset [U, W]_X$.
- 3. Associativity. The diagram

$$\begin{array}{c} \mathcal{P}(X)\times(\mathcal{P}(X)\times\mathcal{P}(X)) \\ \alpha^{\mathsf{Sets}}_{\mathcal{P}(X),\mathcal{P}(X),\mathcal{P}(X)} & \mathsf{id}_{\mathcal{P}(X)}\times\cap \\ (\mathcal{P}(X)\times\mathcal{P}(X))\times\mathcal{P}(X) & \mathcal{P}(X)\times\mathcal{P}(X) \\ & & & & & & & & & \\ \cap \times \mathsf{id}_{\mathcal{P}(X)} & & & & & & & \\ & & & & & & & & \\ \mathcal{P}(X)\times\mathcal{P}(X) & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & & \\ & & & \\ & & & & \\ & & & \\ & & & & \\ & & & \\ & & & & \\ & & & \\ & & & & \\ &$$

commutes, i.e. we have an equality of sets

$$(U \cap V) \cap W = U \cap (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

4. Unitality. The diagrams

$$\operatorname{pt} \times \mathcal{P}(X) \xrightarrow{[X] \times \operatorname{id}_{\mathcal{P}(X)}} \mathcal{P}(X) \times \mathcal{P}(X) \qquad \mathcal{P}(X) \times \operatorname{pt} \xrightarrow{\operatorname{id}_{\mathcal{P}(X)} \times [X]} \mathcal{P}(X) \times \mathcal{P}(X)$$

commute, i.e. we have equalities of sets

$$X \cap U = U,$$

$$U \cap X = U$$

for each $U \in \mathcal{P}(X)$.

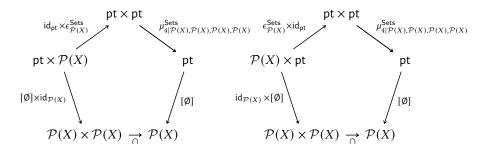
5. Commutativity. The diagram

commutes, i.e. we have an equality of sets

$$U \cap V = V \cap U$$

for each $U, V \in \mathcal{P}(X)$.

6. Annihilation With the Empty Set. The diagrams

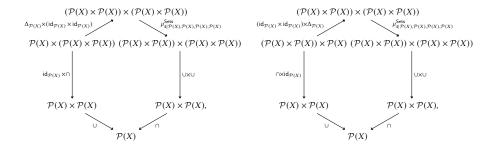


commute, i.e. we have equalities of sets

$$\emptyset \cap X = \emptyset$$
, $X \cap \emptyset = \emptyset$

for each $U \in \mathcal{P}(X)$.

7. Distributivity of Unions Over Intersections. The diagrams



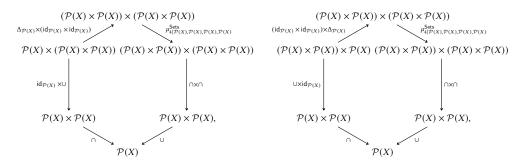
commute, i.e. we have equalities of sets

$$U \cup (V \cap W) = (U \cup V) \cap (U \cup W),$$

$$(U \cap V) \cup W = (U \cup W) \cap (V \cup W)$$

for each $U, V, W \in \mathcal{P}(X)$.

8. Distributivity of Intersections Over Unions. The diagrams



commute, i.e. we have equalities of sets

$$U \cap (V \cup W) = (U \cap V) \cup (U \cap W),$$

$$(U \cup V) \cap W = (U \cap W) \cup (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

9. Idempotency. The diagram

$$\mathcal{P}(X) \xrightarrow{\Delta_{\mathcal{P}(X)}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \cap$$

$$\mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$U \cap U = U$$

for each $U \in \mathcal{P}(X)$.

10. Interaction With Characteristic Functions I. We have

$$\chi_{U\cap V} = \chi_U \chi_V$$

for each $U, V \in \mathcal{P}(X)$.

11. Interaction With Characteristic Functions II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

12. Interaction With Direct Images. Let $f: X \to Y$ be a function. We have a natural transformation

with components

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

13. Interaction With Inverse Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\uparrow \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U\cap V)=f^{-1}(U)\cap f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

14. Interaction With Codirect Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_* \times f_*} \mathcal{P}(Y) \times \mathcal{P}(Y)$$

$$\uparrow \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U) \cap f_*(V) = f_*(U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

- 15. Interaction With Powersets and Monoids With Zero. The quadruple $((\mathcal{P}(X), \emptyset), \cap, X)$ is a commutative monoid with zero.
- 16. Interaction With Powersets and Semirings. The quintuple $(\mathcal{P}(X), \cup, \cap, \emptyset, X)$ is an idempotent commutative semiring.

Proof. Item 1, Functoriality: See [Pro25al].

Item 2, Adjointness: See [MSE 267469].

Item 3, *Associativity*: See [Pro25r].

Item 4, *Unitality*: This follows from [Pro25v] and Item 5.

Item 5, Commutativity: See [Pro25s].

Item 6, Annihilation With the Empty Set: This follows from [Pro25t] and Item 5.

Item 7, Distributivity of Unions Over Intersections: See [Pro25az].

Item 8, Distributivity of Intersections Over Unions: See [Pro25aj].

Item 9, Idempotency: See [Pro25ak].

Item 10, *Interaction With Characteristic Functions I*: See [Pro25e].

Item 11, Interaction With Characteristic Functions II: See [Pro25e].

Item 12, Interaction With Direct Images: See [Pro25n].

Item 13, Interaction With Inverse Images: See [Pro25w].

Item 14, Interaction With Codirect Images: This is a repetition of Item 6 of Definition 4.6.3.1.7 and is proved there.

Item 15, *Interaction With Powersets and Monoids With Zero*: This follows from Items 3 to 6.

Item 16, *Interaction With Powersets and Semirings*: This follows from Items 2 to 4 and 8 and Items 3 to 6 and 8 of Definition 4.3.9.1.2.

4.3.10 Differences

Let X and Y be sets.

Definition 4.3.10.1.1. The **difference of** X **and** Y is the set $X \setminus Y$ defined by

$$X \setminus Y \stackrel{\text{def}}{=} \{ a \in X \mid a \notin Y \}.$$

Proposition 4.3.10.1.2. Let *X* be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto U \cap V$ define functors

$$U \setminus -: \qquad (\mathcal{P}(X), \supset) \qquad \to (\mathcal{P}(X), \subset),$$

$$- \setminus V: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$-_1 \setminus -_2: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \supset) \to (\mathcal{P}(X), \subset).$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $U \setminus V \subset A \setminus V$.
- (b) If $V \subset B$, then $U \setminus B \subset U \setminus V$.
- (c) If $U \subset A$ and $V \subset B$, then $U \setminus B \subset A \setminus V$.
- 2. De Morgan's Laws. We have equalities of sets

$$X \setminus (U \cup V) = (X \setminus U) \cap (X \setminus V),$$

$$X \setminus (U \cap V) = (X \setminus U) \cup (X \setminus V)$$

for each $U, V \in \mathcal{P}(X)$.

3. *Interaction With Unions I.* We have equalities of sets

$$U\setminus (V\cup W)=(U\setminus V)\cap (U\setminus W)$$

for each $U, V, W \in \mathcal{P}(X)$.

4. Interaction With Unions II. We have equalities of sets

$$(U \setminus V) \cup W = (U \cup W) \setminus (V \setminus W)$$

for each $U, V, W \in \mathcal{P}(X)$.

5. Interaction With Unions III. We have equalities of sets

$$U \setminus (V \cup W) = (U \cup W) \setminus (V \cup W)$$
$$= (U \setminus V) \setminus W$$
$$= (U \setminus W) \setminus V$$

for each $U, V, W \in \mathcal{P}(X)$.

6. Interaction With Unions IV. We have equalities of sets

$$(U \cup V) \setminus W = (U \setminus W) \cup (V \setminus W)$$

for each $U, V, W \in \mathcal{P}(X)$.

7. Interaction With Intersections. We have equalities of sets

$$(U \setminus V) \cap W = (U \cap W) \setminus V$$
$$= U \cap (W \setminus V)$$

for each $U, V, W \in \mathcal{P}(X)$.

8. Interaction With Complements. We have an equality of sets

$$U \setminus V = U \cap V^{\mathsf{c}}$$

for each $U, V \in \mathcal{P}(X)$.

9. Interaction With Symmetric Differences. We have an equality of sets

$$U \setminus V = U \triangle (U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

10. Triple Differences. We have

$$U \setminus (V \setminus W) = (U \cap W) \cup (U \setminus V)$$

for each $U, V, W \in \mathcal{P}(X)$.

11. Left Annihilation. We have

$$\emptyset \setminus U = \emptyset$$

for each $U \in \mathcal{P}(X)$.

12. Right Unitality. We have

$$U \setminus \emptyset = U$$

for each $U \in \mathcal{P}(X)$.

13. Right Annihilation. We have

$$U \setminus X = \emptyset$$

for each $U \in \mathcal{P}(X)$.

14. Invertibility. We have

$$U \setminus U = \emptyset$$

for each $U \in \mathcal{P}(X)$.

- 15. Interaction With Containment. The following conditions are equivalent:
 - (a) We have $V \setminus U \subset W$.
 - (b) We have $V \setminus W \subset U$.
- 16. Interaction With Characteristic Functions. We have

$$\chi_{U\setminus V} = \chi_U - \chi_{U\cap V}$$

for each $U, V \in \mathcal{P}(X)$.

17. Interaction With Direct Images. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow$$

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by $U, V \in \mathcal{P}(X)$.

18. Interaction With Inverse Images. The diagram

$$\mathcal{P}(Y)^{\text{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\text{op},-1} \times f^{-1}} \mathcal{P}(X)^{\text{op}} \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad$$

commutes, i.e. we have

$$f^{-1}(U \setminus V) = f^{-1}(U) \setminus f^{-1}(V)$$

for each $U, V \in \mathcal{P}(X)$.

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19. Interaction With Codirect Images. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad$$

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by $U, V \in \mathcal{P}(X)$.

Proof. Item 1, Functoriality: See [Pro25ad] and [Pro25ah].

Item 2, De Morgan's Laws: See [Pro25k].

Item 3, *Interaction With Unions I*: See [Pro25].

Item 4, Interaction With Unions II: We have

$$(U \setminus V) \cup W \stackrel{\text{def}}{=} \{x \in X \mid (x \in U \text{ and } x \notin V) \text{ or } x \in W\}$$

$$= \{x \in X \mid (x \in U \text{ or } x \in W) \text{ and } (x \notin V \text{ or } x \in W)\}$$

$$= \{x \in X \mid (x \in U \cup W) \text{ and not } (x \in V \text{ and } x \notin W)\}$$

$$= \{x \in X \mid (x \in U \cup W) \text{ and not } (x \in V \setminus W)\}$$

$$= \{x \in X \mid (x \in (U \cup W) \setminus (V \setminus W))\}$$

$$= (U \cup W) \setminus (V \setminus W).$$

Item 5, Interaction With Unions III: See [Pro25ai].

Item 6, Interaction With Unions IV: See [Pro25ac].

Item 7, *Interaction With Intersections*: See [Pro25u].

Item 8, Interaction With Complements: See [Pro25aa].

Item 9, Interaction With Symmetric Differences: See [Pro25ab].

Item 10, Triple Differences: See [Pro25ag].

Item 11, Left Annihilation: The direction $\emptyset \subset \emptyset \setminus U$ always holds. Now assume $x \in \emptyset \setminus U$. Then, $x \in \emptyset$ and $x \notin U$. Hence $\emptyset \setminus U \subset \emptyset$ must hold and the sets are equal.

Item 12, Right Unitality: See [Pro25ae].

Item 13, Right Annihilation: It suffices to show that no $x \in X$ can be an element of $U \setminus X$. Assume $x \in U \setminus X$. Then $x \notin X$, contradicting $x \in X$. This completes the proof.

Item 14, Invertibility: See [Pro25af].

Item 15, Interaction With Containment: The conditions are symmetric in U, W, hence it suffices to show that $V \setminus U \subset W$ implies $V \setminus W \subset U$. So assume $V \setminus U \subset W, x \in V \setminus W$. Then $x \in V, x \notin W$. So by contraposition, $x \notin V \setminus U$. But $x \in V$, so we must have $x \in U$, completing the proof.

Item 16, Interaction With Characteristic Functions: See [Pro25f].

Item 17, Interaction With Direct Images: See [Pro250].

Item 18, Interaction With Inverse Images: See [Pro25x].

4.3.11 Complements

Let X be a set and let $U \in \mathcal{P}(X)$.

Definition 4.3.11.1.1. The **complement of** U is the set U^{c} defined by

$$U^{c} \stackrel{\text{def}}{=} X \setminus U$$

$$\stackrel{\text{def}}{=} \{ a \in X \mid a \notin U \}.$$

Proposition 4.3.11.1.2. Let *X* be a set.

1. Functoriality. The assignment $U \mapsto U^{c}$ defines a functor

$$(-)^{\mathsf{c}} \colon \mathcal{P}(X)^{\mathsf{op}} \to \mathcal{P}(X).$$

In particular, the following statements hold for each $U, V \in \mathcal{P}(X)$:

$$(\star)$$
 If $U \subset V$, then $V^{c} \subset U^{c}$.

2. De Morgan's Laws. The diagrams

$$\mathcal{P}(X)^{\mathrm{op}} \times \mathcal{P}(X)^{\mathrm{op}} \xrightarrow{\cup^{\mathrm{op}}} \mathcal{P}(X)^{\mathrm{op}} \qquad \mathcal{P}(X)^{\mathrm{op}} \times \mathcal{P}(X)^{\mathrm{op}} \xrightarrow{\cap^{\mathrm{op}}} \mathcal{P}(X)^{\mathrm{op}}$$

$$(-)^{\mathrm{c}} \times (-)^{\mathrm{c}} \downarrow \qquad \qquad (-)^{\mathrm{c}} \times (-)^{\mathrm{c}} \downarrow \qquad \qquad \downarrow (-)^{\mathrm{c}}$$

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\quad \cap \quad} \mathcal{P}(X) \qquad \qquad \mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\quad \cup \quad} \mathcal{P}(X)$$

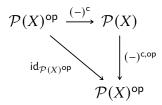
commute, i.e. we have equalities of sets

$$(U \cup V)^{c} = U^{c} \cap V^{c},$$

$$(U \cap V)^{c} = U^{c} \cup V^{c}$$

for each $U, V \in \mathcal{P}(X)$.

3. Involutority. The diagram



commutes, i.e. we have

$$(U^{\mathsf{c}})^{\mathsf{c}} = U$$

for each $U \in \mathcal{P}(X)$.

4. Interaction With Characteristic Functions. We have

$$\chi_{U^c} = 1 - \chi_U$$

for each $U \in \mathcal{P}(X)$.

5. Interaction With Direct Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(X)^{\operatorname{op}} \xrightarrow{f_{*}^{\operatorname{op}}} \mathcal{P}(Y)^{\operatorname{op}}$$

$$(-)^{\operatorname{c}} \qquad \qquad \downarrow (-)^{\operatorname{c}}$$

$$\mathcal{P}(X) \xrightarrow{f} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_!(U^\mathsf{c}) = f_*(U)^\mathsf{c}$$

for each $U \in \mathcal{P}(X)$.

6. Interaction With Inverse Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \xrightarrow{f^{-1,\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}}$$

$$(-)^{\mathsf{c}} \qquad \qquad \downarrow (-)^{\mathsf{c}}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U^{\mathsf{c}}) = f^{-1}(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

7. Interaction With Codirect Images. Let $f: X \to Y$ be a function. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_!^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
 & \downarrow & \downarrow \\
 & \downarrow \\
 & \mathcal{P}(X) & \xrightarrow{f_*} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(U^{\mathsf{c}}) = f_!(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

Proof. Item 1, *Functoriality*: This follows from Item 1 of Definition 4.3.10.1.2.

Item 2, De Morgan's Laws: See [Pro25k].

Item 3, Involutority: See [Pro25i].

Item 4, Interaction With Characteristic Functions: We consider the two cases $x \in U, x \notin U$.

1. If $x \in U$, then $x \notin U^c$. So $\chi_U(x) = 1$ and

$$\chi_{U^{c}}(x) = 0$$
$$= 1 - \chi_{U}(x).$$

2. If $x \notin U$, then $x \in U^{c}$. So $\chi_{U}(x) = 0$ and

$$\chi_{U^{c}}(x) = 1$$
$$= 1 - \chi_{U}(x).$$

Hence, the equation holds for all $x \in X$.

Item 5, Interaction With Direct Images: This is a repetition of Item 8 of Definition 4.6.1.1.5 and is proved there.

Item 6, *Interaction With Inverse Images*: This is a repetition of *Item 8* of *Definition 4.6.2.1.3* and is proved there.

Item 7, *Interaction With Codirect Images*: This is a repetition of Item 7 of Definition 4.6.3.1.7 and is proved there. □

4.3.12 Symmetric Differences

Let X be a set and let $U, V \in \mathcal{P}(X)$.

Definition 4.3.12.1.1. The **symmetric difference of** U **and** V is the set $U \triangle V$ defined by 13

$$U \triangle V \stackrel{\text{def}}{=} (U \setminus V) \cup (V \setminus U).$$

Proposition 4.3.12.1.2. Let X be a set.

1. Lack of Functoriality. The assignment $(U,V)\mapsto U\vartriangle V$ does not in general define functors

$$U \triangle -: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$- \triangle V: \qquad (\mathcal{P}(X), \subset) \qquad \to (\mathcal{P}(X), \subset),$$

$$-_{1} \triangle -_{2}: (\mathcal{P}(X) \times \mathcal{P}(X), \subset \times \subset) \to (\mathcal{P}(X), \subset).$$

2. Via Unions and Intersections. We have

$$U \triangle V = (U \cup V) \setminus (U \cap V)$$

for each $U, V \in \mathcal{P}(X)$, as in the Venn diagram

$$\boxed{\bigcup_{U \triangle V}} = \boxed{\bigcup_{U \cup V}} \setminus \boxed{\bigcup_{U \cap V}}$$

3. Symmetric Differences of Disjoint Sets. If U and V are disjoint, then we have

$$U \triangle V = U \cup V$$
.

4. Associativity. The diagram

$$\begin{array}{c|c} \mathcal{P}(X)\times(\mathcal{P}(X)\times\mathcal{P}(X)) \\ \alpha^{\mathsf{Sets}}_{\mathcal{P}(X),\mathcal{P}(X),\mathcal{P}(X)} & \mathsf{id}_{\mathcal{P}(X)}\times\triangle \\ \\ (\mathcal{P}(X)\times\mathcal{P}(X))\times\mathcal{P}(X) & \mathcal{P}(X)\times\mathcal{P}(X) \\ & & & & & & & & \\ \triangle\times\mathsf{id}_{\mathcal{P}(X)} & & & & & & \\ \mathcal{P}(X)\times\mathcal{P}(X) & \xrightarrow{\Delta} & \mathcal{P}(X), \end{array}$$

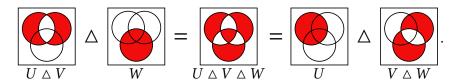
$$\boxed{\bigcup_{U \, \triangle \, V}} = \boxed{\bigcup_{U \, \backslash \, V}} \cup \boxed{\bigcup_{V \, \backslash \, U}}$$

¹³Illustration:

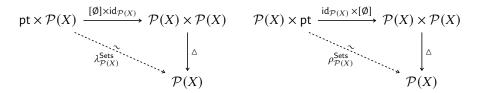
commutes, i.e. we have

$$(U \triangle V) \triangle W = U \triangle (V \triangle W)$$

for each $U, V, W \in \mathcal{P}(X)$, as in the Venn diagram



5. Unitality. The diagrams



commute, i.e. we have

$$U \triangle \emptyset = U,$$
$$\emptyset \triangle U = U$$

for each $U \in \mathcal{P}(X)$.

6. Commutativity. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{\sigma_{\mathcal{P}(X),\mathcal{P}(X)}^{\mathsf{Sets}}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow^{\triangle}$$

$$\mathcal{P}(X)$$

commutes, i.e. we have

$$U \triangle V = V \triangle U$$

for each $U, V \in \mathcal{P}(X)$.

7. Invertibility. We have

$$U \triangle U = \emptyset$$

for each $U \in \mathcal{P}(X)$.

8. Interaction With Unions. We have

$$(U \triangle V) \cup (V \triangle T) = (U \cup V \cup W) \setminus (U \cap V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

9. Interaction With Complements I. We have

$$U \triangle U^{c} = X$$

for each $U \in \mathcal{P}(X)$.

10. Interaction With Complements II. We have

$$U \triangle X = U^{\mathsf{c}},$$

$$X \triangle U = U^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

11. Interaction With Complements III. The diagram

commutes, i.e. we have

$$U^{c} \triangle V^{c} = U \triangle V$$

for each $U, V \in \mathcal{P}(X)$.

12. "Transitivity". We have

$$(U \triangle V) \triangle (V \triangle W) = U \triangle W$$

for each $U, V, W \in \mathcal{P}(X)$.

13. The Triangle Inequality for Symmetric Differences. We have

$$U \triangle W \subset U \triangle V \cup V \triangle W$$

for each $U, V, W \in \mathcal{P}(X)$.

14. Distributivity Over Intersections. We have

$$U \cap (V \triangle W) = (U \cap V) \triangle (U \cap W),$$

$$(U \triangle V) \cap W = (U \cap W) \triangle (V \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

15. Interaction With Characteristic Functions. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \pmod{2}$$

for each $U, V \in \mathcal{P}(X)$.

16. Bijectivity. Given $U, V \in \mathcal{P}(X)$, the maps

$$U \triangle -: \mathcal{P}(X) \to \mathcal{P}(X),$$

- $\triangle V: \mathcal{P}(X) \to \mathcal{P}(X)$

are self-inverse bijections. Moreover, the map

$$\mathcal{P}(X) \longrightarrow \mathcal{P}(X)$$

$$C \longmapsto C \triangle (U \triangle V)$$

is a bijection of $\mathcal{P}(X)$ onto itself sending U to V and V to U.

- 17. Interaction With Powersets and Groups. Let X be a set.
 - (a) The quadruple $(\mathcal{P}(X), \Delta, \emptyset, \mathrm{id}_{\mathcal{P}(X)})$ is an abelian group.¹⁴

¹⁴Here are some examples:

- (b) Every element of $\mathcal{P}(X)$ has order 2 with respect to \triangle , and thus $\mathcal{P}(X)$ is a *Boolean group* (i.e. an abelian 2-group).
- 4. Interaction With Powersets and Vector Spaces I. The pair $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$ consisting of
 - · The group $\mathcal{P}(X)$ of Item 17;
 - · The map $\alpha_{\mathcal{P}(X)} \colon \mathbb{F}_2 \times \mathcal{P}(X) \to \mathcal{P}(X)$ defined by

$$0 \cdot U \stackrel{\text{def}}{=} \emptyset,$$
$$1 \cdot U \stackrel{\text{def}}{=} U;$$

is an \mathbb{F}_2 -vector space.

- 5. Interaction With Powersets and Vector Spaces II. If X is finite, then:
 - (a) The set of singletons sets on the elements of X forms a basis for the \mathbb{F}_2 -vector space $(\mathcal{P}(X), \alpha_{\mathcal{P}(X)})$ of Item 4.
 - (b) We have

$$\dim(\mathcal{P}(X)) = \#X.$$

- 6. Interaction With Powersets and Rings. The quintuple $(\mathcal{P}(X), \triangle, \cap, \emptyset, X)$ is a commutative ring.¹⁵
- 1. When $X = \emptyset$, we have an isomorphism of groups between $\mathcal{P}(\emptyset)$ and the trivial group:

$$(\mathcal{P}(\emptyset), \triangle, \emptyset, id_{\mathcal{P}(\emptyset)}) \cong pt.$$

2. When $X = \operatorname{pt}$, we have an isomorphism of groups between $\mathcal{P}(\operatorname{pt})$ and $\mathbb{Z}_{/2}$:

$$(\mathcal{P}(\mathsf{pt}), \Delta, \emptyset, \mathsf{id}_{\mathcal{P}(\mathsf{pt})}) \cong \mathbb{Z}_{/2}.$$

3. When $X=\{0,1\}$, we have an isomorphism of groups between $\mathcal{P}(\{0,1\})$ and $\mathbb{Z}_{/2}\times\mathbb{Z}_{/2}$:

$$(\mathcal{P}(\{0,1\}), \triangle, \emptyset, \mathsf{id}_{\mathcal{P}(\{0,1\})}) \cong \mathbb{Z}_{/2} \times \mathbb{Z}_{/2}.$$

Warning: The analogous statement replacing intersections by unions (i.e. that the quintuple $(\mathcal{P}(X), \Delta, \cup, \emptyset, X)$ is a ring) is false, however. See [Pro25aw] for a proof.

7. Interaction With Direct Images. We have a natural transformation

with components

$$f_!(U) \triangle f_!(V) \subset f_!(U \triangle V)$$

indexed by $U, V \in \mathcal{P}(X)$.

8. Interaction With Inverse Images. The diagram

i.e. we have

$$f^{-1}(U) \triangle f^{-1}(V) = f^{-1}(U \triangle V)$$

for each $U, V \in \mathcal{P}(Y)$.

9. Interaction With Codirect Images. We have a natural transformation

with components

$$f_*(U \triangle V) \subset f_*(U) \triangle f_*(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

Proof. Item 1, Lack of Functoriality: Let $X = \{0,1\}, U = \{0\}$. Then $\emptyset \subset U$, but $U \triangle \emptyset = U \not\subset \emptyset = U \triangle U$ from Item 5 and Item 7. This gives a counterexample to the first statement. By using Item 6, we can adapt it to the second and third statement.

Item 2, *Via Unions and Intersections*: See [Pro25m].

Item 3, Symmetric Differences of Disjoint Sets: Since U and V are disjoint, we have $U \cap V = \emptyset$, and therefore we have

$$U \triangle V = (U \cup V) \setminus (U \cap V)$$
$$= (U \cup V) \setminus \emptyset$$
$$= U \cup V.$$

where we,ve used Item 2 and Item 12 of Definition 4.3.10.1.2.

Item 4, Associativity: See [Pro25ao].

Item 5, *Unitality*: This follows from Item 6 and [Pro25at].

Item 6, *Commutativity*: See [Pro25ap].

Item 7, *Invertibility*: See [Pro25av].

Item 8, *Interaction With Unions*: See [Pro25bc].

Item 9, Interaction With Complements I: See [Pro25as].

Item 10, *Interaction With Complements II*: This follows from Item 6 and [Pro25ax].

Item 11, Interaction With Complements III: See [Pro25aq].

Item 12, "Transitivity": We have

$$\begin{array}{rcl} (U \bigtriangleup V) \bigtriangleup (V \bigtriangleup W) & = & U \bigtriangleup (V \bigtriangleup (V \bigtriangleup W)) & \text{(by Item 4)} \\ & = & U \bigtriangleup ((V \bigtriangleup V) \bigtriangleup W) & \text{(by Item 4)} \\ & = & U \bigtriangleup (\emptyset \bigtriangleup W) & \text{(by Item 7)} \\ & = & U \bigtriangleup W. & \text{(by Item 5)} \end{array}$$

This finishes the proof.

Item 13, *The Triangle Inequality for Symmetric Differences*: This follows from Items 2 and 12.

Item 14, *Distributivity Over Intersections*: See [Pro25q].

Item 15, *Interaction With Characteristic Functions*: See [Pro25g].

Item 16, Bijectivity:

· We show that

$$(U \triangle -) \colon \mathcal{P}(X) \to \mathcal{P}(X)$$

is self-inverse.

Let $W \in \mathcal{P}(X)$. Then,

$$U \triangle (U \triangle W) = (U \triangle U) \triangle W$$
 (by Item 4)
= $\emptyset \triangle W$ (by Item 7)
= W . (by Item 5)

- By Item 6, $(- \triangle V) = (V \triangle -)$, hence the former is also self-inverse by the first point.
- The map $\triangle (U \triangle V)$ is a bijection as a special case of the second point. From the first two points and Item 6, we get

$$U \triangle (U \triangle V) = V$$
, $V \triangle (U \triangle V) = V \triangle (V \triangle U) = U$.

Hence the function maps U to V and V to U.

Item 17, *Interaction With Powersets and Groups*: Item 17a follows from Items 4 to 7, while Item 3b follows from Item 7.¹⁶

Item 4, Interaction With Powersets and Vector Spaces I: See [MSE 2719059].

Item 5, Interaction With Powersets and Vector Spaces II: See [MSE 2719059].

Item 6, Interaction With Powersets and Rings: This follows from Items 6 and 15 of Definition 4.3.9.1.2 and Items 14 and 17.¹⁷

Item 7, *Interaction With Direct Images*: This is a repetition of Item 9 of Definition 4.6.1.1.5 and is proved there.

Item 8, Interaction With Inverse Images: This is a repetition of Item 9 of Definition 4.6.2.1.3 and is proved there.

Item 9, Interaction With Codirect Images: This is a repetition of Item 8 of Definition 4.6.3.1.7 and is proved there. □

4.4 Powersets

4.4.1 Foundations

Let *X* be a set.

Definition 4.4.1.1.1. The **powerset of** X is the set $\mathcal{P}(X)$ defined by

$$\mathcal{P}(X) \stackrel{\text{def}}{=} \{ U \in P \mid U \subset X \},\$$

where P is the set in the axiom of powerset, ?? of ??.

¹⁶ Reference: [Pro25ar].

¹⁷ Reference: [Pro25au].

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Remark 4.4.1.1.2. Under the analogy that $\{t, f\}$ should be the (-1)-categorical analogue of Sets, we may view the powerset of a set as a decategorification of the category of presheaves of a category (or of the category of copresheaves):

• The powerset of a set X is equivalently (Item 2 of Definition 4.5.1.1.4) the set

$$Sets(X, \{t, f\})$$

of functions from X to the set $\{t, f\}$ of classical truth values.

· The category of presheaves on a category C is the category

$$\operatorname{Fun}(C^{\operatorname{op}},\operatorname{Sets})$$

of functors from C^{op} to the category Sets of sets.

Notation 4.4.1.1.3. Let *X* be a set.

- 1. We write $\mathcal{P}_0(X)$ for the set of nonempty subsets of X.
- 2. We write $\mathcal{P}_{fin}(X)$ for the set of finite subsets of X.

Proposition 4.4.1.1.4. Let *X* be a set.

- 1. *Co/Completeness*. The (posetal) category (associated to) $(\mathcal{P}(X), \subset)$ is complete and cocomplete:
 - (a) *Products*. The products in $\mathcal{P}(X)$ are given by intersection of subsets.
 - (b) Coproducts. The coproducts in $\mathcal{P}(X)$ are given by union of subsets.
 - (c) Co/Equalisers. Being a posetal category, $\mathcal{P}(X)$ only has at most one morphisms between any two objects, so co/equalisers are trivial.
- 2. Cartesian Closedness. The category $\mathcal{P}(X)$ is Cartesian closed.
- 3. Powersets as Sets of Relations. We have bijections

$$\mathcal{P}(X) \cong \text{Rel}(\text{pt}, X),$$

 $\mathcal{P}(X) \cong \text{Rel}(X, \text{pt}),$

natural in $X \in Obj(Sets)$.

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4. Interaction With Products I. The map

$$\mathcal{P}(X) \times \mathcal{P}(Y) \longrightarrow \mathcal{P}(X \coprod Y)$$
$$(U, V) \longmapsto U \cup V$$

is an isomorphism of sets, natural in $X, Y \in \mathsf{Obj}(\mathsf{Sets})$ with respect to each of the functor structures $\mathcal{P}_1, \mathcal{P}^{-1}$, and \mathcal{P}_* on \mathcal{P} of Definition 4.4.2.1.1. Moreover, this makes each of $\mathcal{P}_1, \mathcal{P}^{-1}$, and \mathcal{P}_* into a symmetric monoidal functor.

5. Interaction With Products II. The map

$$\mathcal{P}(X) \times \mathcal{P}(Y) \longrightarrow \mathcal{P}(X \coprod Y)$$
$$(U, V) \longmapsto U \boxtimes_{X \times Y} V,$$

where 18

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \{(u, v) \in X \times Y \mid u \in U \text{ and } v \in V\}$$

is an inclusion of sets, natural in $X,Y\in Obj(Sets)$ with respect to each of the functor structures $\mathcal{P}_1,\mathcal{P}^{-1}$, and \mathcal{P}_* on \mathcal{P} of Definition 4.4.2.1.1. Moreover, this makes each of $\mathcal{P}_1,\mathcal{P}^{-1}$, and \mathcal{P}_* into a symmetric monoidal functor.

6. Interaction With Products III. We have an isomorphism

$$\mathcal{P}(X) \otimes \mathcal{P}(Y) \cong \mathcal{P}(X \times Y),$$

natural in $X, Y \in \text{Obj}(\mathsf{Sets})$ with respect to each of the functor structures $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* on \mathcal{P} of Definition 4.4.2.1.1, where \otimes denotes the tensor product of suplattices of ??. Moreover, this makes each of $\mathcal{P}_!, \mathcal{P}^{-1}$, and \mathcal{P}_* into a symmetric monoidal functor.

Proof. Item 1, *Co/Completeness*: Omitted.

Item 2, Cartesian Closedness: See Section 4.4.7.

Item 3, Powersets as Sets of Relations: Indeed, we have

$$\mathsf{Rel}(\mathsf{pt},X) \stackrel{\mathsf{def}}{=} \mathcal{P}(\mathsf{pt} \times X)$$

¹⁸The set $U \boxtimes_{X \times Y} V$ is usually denoted simply $U \times V$. Here we denote it in this somewhat weird way to highlight the similarity to external tensor products in six-functor formalisms (see also Section 4.6.4).

$$\cong \mathcal{P}(X)$$

and

$$\mathsf{Rel}(X,\mathsf{pt}) \stackrel{\mathsf{def}}{=} \mathcal{P}(X \times \mathsf{pt})$$

 $\cong \mathcal{P}(X),$

where we have used Item 5 of Definition 4.1.3.1.3.

Item 4, Interaction With Products I: The inverse of the map in the statement is the map

$$\Phi \colon \mathcal{P}(X \mid \mid Y) \to \mathcal{P}(X) \times \mathcal{P}(Y)$$

defined by

$$\Phi(S) \stackrel{\mathsf{def}}{=} (S_X, S_Y)$$

for each $S \in \mathcal{P}(X \mid Y)$, where

$$S_X \stackrel{\text{def}}{=} \{ x \in X \mid (0, x) \in S \}$$

 $S_Y \stackrel{\text{def}}{=} \{ y \in Y \mid (1, y) \in S \}.$

The rest of the proof is omitted.

Item 5, Interaction With Products II: Omitted.

Item 6, Interaction With Products III: Omitted.

4.4.2 Functoriality of Powersets

Proposition 4.4.2.1.1. Let *X* be a set.

1. Functoriality I. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}_1 \colon \mathsf{Sets} \to \mathsf{Sets}$$

where

· Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}_!(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

· Action on Morphisms. For each $A, B \in \mathsf{Obj}(\mathsf{Sets})$, the action on morphisms

$$\mathcal{P}_{*|A|B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of $\mathcal{P}_!$ at (A,B) is the map defined by sending a map of sets $f\colon A\to B$ to the map

$$\mathcal{P}_!(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_!(f) \stackrel{\text{def}}{=} f_!,$$

as in Definition 4.6.1.1.1.

2. Functoriality II. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}^{-1}$$
: Sets^{op} \rightarrow Sets,

where

· Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}^{-1}(A) \stackrel{\text{def}}{=} \mathcal{P}(A).$$

· Action on Morphisms. For each $A, B \in \mathsf{Obj}(\mathsf{Sets})$, the action on morphisms

$$\mathcal{P}_{AB}^{-1}$$
: Sets $(A, B) \to \text{Sets}(\mathcal{P}(B), \mathcal{P}(A))$

of \mathcal{P}^{-1} at (A,B) is the map defined by by sending a map of sets $f\colon A\to B$ to the map

$$\mathcal{P}^{-1}(f) \colon \mathcal{P}(B) \to \mathcal{P}(A)$$

defined by

$$\mathcal{P}^{-1}(f) \stackrel{\text{def}}{=} f^{-1},$$

as in Definition 4.6.2.1.1.

3. Functoriality III. The assignment $X \mapsto \mathcal{P}(X)$ defines a functor

$$\mathcal{P}_* \colon \mathsf{Sets} \to \mathsf{Sets},$$

where

· Action on Objects. For each $A \in Obj(Sets)$, we have

$$\mathcal{P}_*(A) \stackrel{\mathsf{def}}{=} \mathcal{P}(A).$$

· Action on Morphisms. For each $A, B \in \mathsf{Obj}(\mathsf{Sets})$, the action on morphisms

$$\mathcal{P}_{!|A,B} \colon \mathsf{Sets}(A,B) \to \mathsf{Sets}(\mathcal{P}(A),\mathcal{P}(B))$$

of \mathcal{P}_* at (A,B) is the map defined by sending a map of sets $f\colon A\to B$ to the map

$$\mathcal{P}_*(f) \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

defined by

$$\mathcal{P}_*(f) \stackrel{\text{def}}{=} f_*,$$

as in Definition 4.6.3.1.1.

Proof. Item 1, Functoriality I: This follows from Items 3 and 4 of Definition 4.6.1.1.6.

Item 2, Functoriality II: This follows from Items 3 and 4 of Definition 4.6.2.1.4.

Item 3, Functoriality III: This follows from Items 3 and 4 of Definition 4.6.3.1.8. □

4.4.3 Adjointness of Powersets I

Proposition 4.4.3.1.1. We have an adjunction

$$(\mathcal{P}^{-1} \dashv \mathcal{P}^{-1,op})$$
: Sets^{op} $\underset{\mathcal{P}^{-1,op}}{\stackrel{\mathcal{P}^{-1}}{\smile}}$ Sets,

witnessed by a bijection

$$\underbrace{\mathsf{Sets}^{\mathsf{op}}(\mathcal{P}(X),Y)}_{\overset{\mathsf{def}}{=}\mathsf{Sets}(Y,\mathcal{P}(X))} \cong \mathsf{Sets}(X,\mathcal{P}(Y)),$$

natural in $X \in Obj(Sets)$ and $Y \in Obj(Sets^{op})$.

Proof. We have

$$\begin{aligned} \mathsf{Sets}^\mathsf{op}(\mathcal{P}(A),B) &\stackrel{\mathsf{def}}{=} & \mathsf{Sets}(B,\mathcal{P}(A)) \\ &\cong & \mathsf{Sets}(B,\mathsf{Sets}(A,\{\mathsf{t},\mathsf{f}\})) & (\mathsf{by}\,\mathsf{ltem}\,\mathsf{2}\,\mathsf{of}\,\mathsf{Definition}\,\mathsf{4.5.1.1.4}) \\ &\cong & \mathsf{Sets}(A\times B,\{\mathsf{t},\mathsf{f}\}) & (\mathsf{by}\,\mathsf{ltem}\,\mathsf{2}\,\mathsf{of}\,\mathsf{Definition}\,\mathsf{4.1.3.1.3}) \\ &\cong & \mathsf{Sets}(A,\mathsf{Sets}(B,\{\mathsf{t},\mathsf{f}\})) & (\mathsf{by}\,\mathsf{ltem}\,\mathsf{2}\,\mathsf{of}\,\mathsf{Definition}\,\mathsf{4.1.3.1.3}) \\ &\cong & \mathsf{Sets}(A,\mathcal{P}(B)), & (\mathsf{by}\,\mathsf{ltem}\,\mathsf{2}\,\mathsf{of}\,\mathsf{Definition}\,\mathsf{4.5.1.1.4}) \end{aligned}$$

where all bijections are natural in A and B.¹⁹

¹⁹Here we are using Item 3 of Definition 4.5.1.1.4.

4.4.4 Adjointness of Powersets II

Proposition 4.4.4.1.1. We have an adjunction

$$(\operatorname{Gr} \dashv \mathcal{P}_!)$$
: Sets $\stackrel{\operatorname{Gr}}{\underset{\mathcal{P}_!}{\longleftarrow}} \operatorname{Rel},$

witnessed by a bijection of sets

$$Rel(Gr(X), Y) \cong Sets(X, \mathcal{P}(Y))$$

natural in $X \in \text{Obj}(\mathsf{Sets})$ and $Y \in \text{Obj}(\mathsf{Rel})$, where Gr is the graph functor of Relations, Item 1 of Definition 8.2.2.1.2 and $\mathcal{P}_!$ is the functor of Relations, Definition 8.7.5.1.1.

Proof. We have

where all bijections are natural in A, (where we are using Item 3 of Definition 4.5.1.1.4). Explicitly, this isomorphism is given by sending a relation R: $Gr(A) \to B$ to the map $R^{\dagger} : A \to \mathcal{P}(B)$ sending a to the subset R(a) of B, as in Relations, Definition 8.1.1.1.1.

Naturality in B is then the statement that given a relation $R \colon B \to B'$, the diagram

commutes, which follows from Relations, Definition 8.7.1.1.3.

4.4.5 Powersets as Free Cocompletions

Let X be a set.

Proposition 4.4.5.1.1. The pair $(\mathcal{P}(X), \chi_{(-)})$ consisting of

- · The powerset $(\mathcal{P}(X), \subset)$ of X of Definition 4.4.1.1.1;
- The characteristic embedding $\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$ of X into $\mathcal{P}(X)$ of Definition 4.5.4.1.1;

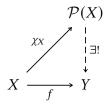
satisfies the following universal property:

- (\star) Given another pair (Y, f) consisting of
 - A suplattice (Y, ≤);
 - A function $f: X \to Y$;

there exists a unique morphism of suplattices

$$(\mathcal{P}(X),\subset) \xrightarrow{\exists !} (Y,\preceq)$$

making the diagram



commute.

Proof. This is a rephrasing of Definition 4.4.5.1.2, which we prove below.²⁰

Proposition 4.4.5.1.2. We have an adjunction

$$(\mathcal{P} \dashv \overline{\Xi})$$
: Sets $\stackrel{\mathcal{P}}{=}$ SupLat,

witnessed by a bijection

$$\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))\cong\mathsf{Sets}(X,Y),$$

natural in $X \in \mathsf{Obj}(\mathsf{Sets})$ and $(Y, \preceq) \in \mathsf{Obj}(\mathsf{SupLat})$, where:

²⁰Here we only remark that the unique morphism of suplattices in the statement is given by the

- · The category SupLat is the category of suplattices of ??.
- · The map

$$\chi_X^* : \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

witnessing the above bijection is defined by

$$\chi_X^*(f) \stackrel{\text{def}}{=} f \circ \chi_X,$$

i.e. by sending a morphism of suplattices $f: \mathcal{P}(X) \to Y$ to the composition

$$X \stackrel{\chi_X}{\hookrightarrow} \mathcal{P}(X) \stackrel{f}{\longrightarrow} Y.$$

· The map

$$\mathsf{Lan}_{\chi_X} \colon \mathsf{Sets}(X,Y) \to \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$$

witnessing the above bijection is given by sending a function $f: X \to Y$ to its left Kan extension along χ_X ,

$$\operatorname{Lan}_{\chi_X}(f) \colon \mathcal{P}(X) \to Y, \qquad \begin{array}{c|c} \mathcal{P}(X) \\ \chi_X & \downarrow \\ \chi_X & \downarrow \\ X & \xrightarrow{f} Y. \end{array}$$

Moreover, invoking the bijection $\mathcal{P}(X) \cong \operatorname{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ of Item 2 of Definition 4.5.1.1.4, $\operatorname{Lan}_{\chi_X}(f)$ can be explicitly computed by

$$[\operatorname{Lan}_{\chi_X}(f)](U) = \int_{x \in X}^{x \in X} \chi_{\mathcal{P}(X)}(\chi_x, U) \odot f(x)$$

$$= \int_{x \in X}^{x \in X} \chi_U(x) \odot f(x)$$

$$= \bigvee_{x \in X} (\chi_U(x) \odot f(x))$$

$$= \left(\bigvee_{x \in U} (\chi_U(x) \odot f(x))\right) \vee \left(\bigvee_{x \in U^c} (\chi_U(x) \odot f(x))\right)$$

$$= \left(\bigvee_{x \in U} f(x)\right) \vee \left(\bigvee_{x \in U^{c}} \varnothing_{Y}\right)$$
$$= \bigvee_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$, where:

- We have used ?? for the first equality.
- We have used Definition 4.5.5.1.1 for the second equality.
- We have used ?? for the third equality.
- The symbol \bigvee denotes the join in (Y, \preceq) .
- The symbol \odot denotes the tensor of an element of Y by a truth value as in $\ref{eq: 1}$. In particular, we have

true
$$\odot f(x) \stackrel{\text{def}}{=} f(x)$$
,
false $\odot f(x) \stackrel{\text{def}}{=} \varnothing_Y$,

where \emptyset_Y is the bottom element of (Y, \preceq) .

In particular, when $(Y, \preceq_Y) = (\mathcal{P}(B), \subset)$ for some set B, the Kan extension $\operatorname{Lan}_{\chi_X}(f)$ is given by

$$[\operatorname{Lan}_{\chi_X}(f)](U) = \bigvee_{x \in U} f(x)$$
$$= \bigcup_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$.

Proof. Map I: We define a map

$$\Phi_{XY} : \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

as in the statement, i.e. by

$$\Phi_{X,Y}(f) \stackrel{\mathsf{def}}{=} f \circ \chi_X$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$.

left Kan extension $\operatorname{Lan}_{\chi_X}(f)$ of f along χ_X .

Map II: We define a map

$$\Psi_{X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$$

as in the statement, i.e. by

$$\Psi_{X,Y}(f) \stackrel{\text{def}}{=} \mathsf{Lan}_{\chi_X}(f), \qquad \qquad \downarrow^{\chi_X} \qquad \downarrow^{\mathsf{Lan}_{\chi_X}(f)} \\ X \xrightarrow{f} Y,$$

for each $f \in Sets(X, Y)$. Invertibility I: We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathsf{id}_{\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))}$$
.

We have

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) \stackrel{\text{def}}{=} \Psi_{X,Y}(\Phi_{X,Y}(f))$$

$$\stackrel{\text{def}}{=} \Psi_{X,Y}(f \circ \chi_X)$$

$$\stackrel{\text{def}}{=} \text{Lan}_{\chi_X}(f \circ \chi_X)$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$. We now claim that

$$\mathsf{Lan}_{\chi_X}(f\circ\chi_X)=f$$

for each $f \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$. Indeed, we have

$$\begin{aligned} \left[\mathsf{Lan}_{\chi_X} (f \circ \chi_X) \right] (U) &= \bigvee_{x \in U} f(\chi_X(x)) \\ &= f \bigg(\bigvee_{x \in U} \chi_X(x) \bigg) \\ &= f \bigg(\bigcup_{x \in U} \{x\} \bigg) \\ &= f(U) \end{aligned}$$

for each $U \in \mathcal{P}(X)$, where we have used that f is a morphism of suplattices and hence preserves joins for the second equality. This proves our claim. Since we have shown that

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) = f$$

for each $f \in \operatorname{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$, it follows that $\Psi_{X,Y} \circ \Phi_{X,Y}$ must be equal to the identity map $\operatorname{id}_{\operatorname{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))}$ of $\operatorname{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq))$. Invertibility II: We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathsf{id}_{\mathsf{Sets}(X,Y)}$$
.

We have

$$\begin{split} [\Phi_{X,Y} \circ \Psi_{X,Y}](f) &\stackrel{\text{def}}{=} \Phi_{X,Y}(\Psi_{X,Y}(f)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}(\mathsf{Lan}_{\chi_X}(f)) \\ &\stackrel{\text{def}}{=} \mathsf{Lan}_{\chi_X}(f) \circ \chi_X \end{split}$$

for each $f \in Sets(X, Y)$. We now claim that

$$\mathsf{Lan}_{\chi_X}(f) \circ \chi_X = f$$

for each $f \in Sets(X, Y)$. Indeed, we have

$$[\operatorname{Lan}_{\chi_X}(f) \circ \chi_X](x) = \bigvee_{y \in \{x\}} f(y)$$
$$= f(x)$$

for each $x \in X$. This proves our claim. Since we have shown that

$$[\Phi_{X,Y} \circ \Psi_{X,Y}](f) = f$$

for each $f \in \text{Sets}(X, Y)$, it follows that $\Phi_{X,Y} \circ \Psi_{X,Y}$ must be equal to the identity map $\text{id}_{\text{Sets}(X,Y)}$ of Sets(X,Y).

Naturality for Φ , Part I: We need to show that, given a function $f\colon X\to X'$, the diagram

$$\begin{split} \mathsf{SupLat}((\mathcal{P}(X'),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X',Y}} \mathsf{Sets}(X',Y) \\ & & \downarrow f^* \\ \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y) \end{split}$$

commutes. Indeed, we have

$$\begin{split} [\Phi_{X,Y} \circ \mathcal{P}_!(f)^*](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y}(\mathcal{P}_!(f)^*(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}(\xi \circ f_!) \\ &\stackrel{\text{def}}{=} (\xi \circ f_!) \circ \chi_X \\ &= \xi \circ (f_! \circ \chi_X) \\ &\stackrel{(\dagger)}{=} \xi \circ (\chi_{X'} \circ f) \\ &= (\xi \circ \chi_{X'}) \circ f \\ &\stackrel{\text{def}}{=} \Phi_{X',Y}(\xi) \circ f \\ &\stackrel{\text{def}}{=} f^*(\Phi_{X',Y}(\xi)) \\ &\stackrel{\text{def}}{=} [f^* \circ \Phi_{X',Y}](\xi), \end{split}$$

for each $\xi \in \text{SupLat}((\mathcal{P}(X'), \subset), (Y, \preceq))$, where we have used Item 1 of Definition 4.5.4.1.3 for the fifth equality above.

Naturality for Φ , Part II: We need to show that, given a morphism of suplattices

$$g: (Y, \preceq_Y) \to (Y', \preceq_{Y'}),$$

the diagram

$$\begin{array}{ccc} \mathsf{SupLat}((\mathcal{P}(X),\subset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} & \mathsf{Sets}(X,Y) \\ & & & & \downarrow g_! \\ & & & \downarrow g_! \end{array}$$

$$\mathsf{SupLat}((\mathcal{P}(X),\subset),(Y',\preceq)) \xrightarrow{\Phi_{X,Y'}} & \mathsf{Sets}(X,Y')$$

commutes. Indeed, we have

$$\begin{split} [\Phi_{X,Y'} \circ g_!](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g_!(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g \circ \xi) \\ &\stackrel{\text{def}}{=} (g \circ \xi) \circ \chi_X \\ &= g \circ (\xi \circ \chi_X) \\ &\stackrel{\text{def}}{=} g \circ (\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} g_!(\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} [g_! \circ \Phi_{X,Y}](\xi). \end{split}$$

for each $\xi \in \mathsf{SupLat}((\mathcal{P}(X), \subset), (Y, \preceq))$.

Naturality for Ψ : Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Ψ is also natural in each argument.

Warning 4.4.5.1.3. Although the assignment $X \mapsto \mathcal{P}(X)$ is called the *free cocompletion of X*, it is not an idempotent operation, i.e. we have $\mathcal{P}(\mathcal{P}(X)) \neq \mathcal{P}(X)$.

4.4.6 Powersets as Free Completions

Let X be a set.

Proposition 4.4.6.1.1. The pair $(\mathcal{P}(X), \chi_{(-)})$ consisting of

- The powerset of X together with reverse inclusion $\mathcal{P}(X)^{\mathsf{op}} = (\mathcal{P}(X), \supset)$ of Definition 4.4.1.1.1;
- The characteristic embedding $\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$ of X into $\mathcal{P}(X)$ of Definition 4.5.4.1.1;

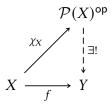
satisfies the following universal property:

- (\star) Given another pair (Y, f) consisting of
 - An inflattice (Y, \preceq) ;
 - A function $f: X \to Y$;

there exists a unique morphism of inflattices

$$(\mathcal{P}(X),\supset) \xrightarrow{\exists!} (Y,\preceq)$$

making the diagram



commute.

Proof. This is a rephrasing of Definition 4.4.6.1.2, which we prove below.²¹

Proposition 4.4.6.1.2. We have an adjunction

$$(\mathcal{P} \dashv \overline{\Xi})$$
: Sets $\stackrel{\mathcal{P}}{\underset{\overline{\Xi}}{\longleftarrow}}$ InfLat,

witnessed by a bijection

$$InfLat((\mathcal{P}(X),\supset),(Y,\preceq)) \cong Sets(X,Y),$$

natural in $X \in \mathsf{Obj}(\mathsf{Sets})$ and $(Y, \preceq) \in \mathsf{Obj}(\mathsf{InfLat})$, where:

- · The category InfLat is the category of inflattices of ??.
- · The map

$$\chi_X^* : \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq)) \to \mathsf{Sets}(X, Y)$$

witnessing the above bijection is defined by

$$\chi_X^*(f) \stackrel{\text{def}}{=} f \circ \chi_X,$$

i.e. by sending a morphism of inflattices $f: \mathcal{P}(X)^{op} \to Y$ to the composition

$$X \stackrel{\chi_X}{\hookrightarrow} \mathcal{P}(X)^{\mathsf{op}} \stackrel{f}{\longrightarrow} Y.$$

· The map

$$\operatorname{\mathsf{Ran}}_{\chi_X} \colon \operatorname{\mathsf{Sets}}(X,Y) \to \operatorname{\mathsf{InfLat}}((\mathcal{P}(X),\supset),(Y,\preceq))$$

witnessing the above bijection is given by sending a function $f\colon X\to Y$ to its right Kan extension along χ_X ,

$$\operatorname{Ran}_{\chi_X}(f) \colon \mathcal{P}(X)^{\operatorname{op}} \to Y, \qquad \chi_X / \underset{f}{ \downarrow_{\operatorname{Ran}_{\chi_X}(f)}} X \xrightarrow{f} Y.$$

²¹Here we only remark that the unique morphism of inflattices in the statement is given by the

Moreover, invoking the bijection $\mathcal{P}(X) \cong \operatorname{Sets}(X, \{t, f\})$ of Item 2 of Definition 4.5.1.1.4, $\operatorname{Ran}_{\chi_X}(f)$ can be explicitly computed by

for each $U \in \mathcal{P}(X)$, where:

- We have used ?? for the first equality.
- We have used Definition 4.5.5.1.1 for the second equality.
- We have used ?? for the third equality.
- **–** The symbol \land denotes the meet in (Y, \preceq) .
- The symbol \pitchfork denotes the cotensor of an element of Y by a truth value as in $\ref{eq: Y}$. In particular, we have

true
$$\pitchfork f(x) \stackrel{\text{def}}{=} f(x)$$
, false $\pitchfork f(x) \stackrel{\text{def}}{=} \infty_Y$,

where ∞_Y is the top element of (Y, \preceq) .

right Kan extension $\operatorname{Ran}_{\chi_X}(f)$ of f along χ_X .

In particular, when $(Y, \preceq_Y) = (\mathcal{P}(B), \subset)$ for some set B, the Kan extension $\operatorname{Ran}_{Y_X}(f)$ is given by

$$[\operatorname{Ran}_{\chi_X}(f)](U) = \bigwedge_{x \in U} f(x)$$
$$= \bigcap_{x \in U} f(x)$$

for each $U \in \mathcal{P}(X)$.

Proof. Map I: We define a map

$$\Phi_{X,Y} \colon \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) \to \mathsf{Sets}(X,Y)$$

as in the statement, i.e. by

$$\Phi_{X,Y}(f) \stackrel{\text{def}}{=} f \circ \chi_X$$

for each $f \in \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq))$. Map II: We define a map

$$\Psi_{X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq))$$

as in the statement, i.e. by

$$\Psi_{X,Y}(f) \stackrel{\text{def}}{=} \operatorname{Ran}_{\chi_X}(f), \qquad X \xrightarrow{\chi_X} \bigvee_{f} \operatorname{Ran}_{\chi_X}(f)$$

$$X \xrightarrow{f} Y,$$

for each $f \in Sets(X, Y)$. Invertibility I: We claim that

$$\Psi_{X,Y} \circ \Phi_{X,Y} = \mathsf{id}_{\mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq))}$$
.

We have

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) \stackrel{\text{def}}{=} \Psi_{X,Y}(\Phi_{X,Y}(f))$$

$$\stackrel{\text{def}}{=} \Psi_{X,Y}(f \circ \chi_X)$$

$$\stackrel{\text{def}}{=} \operatorname{Ran}_{\chi_X}(f \circ \chi_X)$$

for each $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$. We now claim that

$$\mathsf{Ran}_{\chi_X}(f\circ\chi_X)=f$$

for each $f \in InfLat((\mathcal{P}(X), \supset), (Y, \preceq))$. Indeed, we have

$$\begin{aligned} \left[\operatorname{Ran}_{\chi_X} (f \circ \chi_X) \right] (U) &= \bigwedge_{x \in U} f(\chi_X(x)) \\ &= f \left(\bigwedge_{x \in U} \chi_X(x) \right) \\ &= f \left(\bigcup_{x \in U} \{x\} \right) \\ &= f(U) \end{aligned}$$

for each $U \in \mathcal{P}(X)$, where we have used that f is a morphism of inflattices and hence preserves meets in $(\mathcal{P}(X), \supset)$ (i.e. joins in $(\mathcal{P}(X), \subset)$) for the second equality. This proves our claim. Since we have shown that

$$[\Psi_{X,Y} \circ \Phi_{X,Y}](f) = f$$

for each $f \in \mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq))$, it follows that $\Psi_{X,Y} \circ \Phi_{X,Y}$ must be equal to the identity map $\mathsf{id}_{\mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq))}$ of $\mathsf{InfLat}((\mathcal{P}(X), \supset), (Y, \preceq))$. Invertibility II: We claim that

$$\Phi_{X,Y} \circ \Psi_{X,Y} = \mathrm{id}_{\mathsf{Sets}(X,Y)}$$
.

We have

$$\begin{split} [\Phi_{X,Y} \circ \Psi_{X,Y}](f) &\stackrel{\text{def}}{=} \Phi_{X,Y}(\Psi_{X,Y}(f)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}(\mathsf{Ran}_{\chi_X}(f)) \\ &\stackrel{\text{def}}{=} \mathsf{Ran}_{\chi_X}(f) \circ \chi_X \end{split}$$

for each $f \in Sets(X, Y)$. We now claim that

$$\mathsf{Ran}_{\chi_X}(f) \circ \chi_X = f$$

for each $f \in Sets(X, Y)$. Indeed, we have

$$[\operatorname{Ran}_{\chi_X}(f) \circ \chi_X](x) = \bigwedge_{y \in \{x\}} f(y)$$

$$= f(x)$$

for each $x \in X$. This proves our claim. Since we have shown that

$$[\Phi_{X,Y} \circ \Psi_{X,Y}](f) = f$$

for each $f \in \mathsf{Sets}(X,Y)$, it follows that $\Phi_{X,Y} \circ \Psi_{X,Y}$ must be equal to the identity map $\mathsf{id}_{\mathsf{Sets}(X,Y)}$ of $\mathsf{Sets}(X,Y)$.

Naturality for Φ , Part I: We need to show that, given a function $f\colon X\to X'$, the diagram

$$\begin{split} \mathsf{InfLat}((\mathcal{P}(X'),\supset),(Y,\preceq)) & \xrightarrow{\Phi_{X',Y}} \mathsf{Sets}(X',Y) \\ \mathcal{P}_{!}(f)^{*} & & \downarrow f^{*} \\ \mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) & \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y) \end{split}$$

commutes. Indeed, we have

$$\begin{split} [\Phi_{X,Y} \circ \mathcal{P}_!(f)^*](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y}(\mathcal{P}_!(f)^*(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y}(\xi \circ f_!) \\ &\stackrel{\text{def}}{=} (\xi \circ f_!) \circ \chi_X \\ &= \xi \circ (f_! \circ \chi_X) \\ &\stackrel{(\dagger)}{=} \xi \circ (\chi_{X'} \circ f) \\ &= (\xi \circ \chi_{X'}) \circ f \\ &\stackrel{\text{def}}{=} \Phi_{X',Y}(\xi) \circ f \\ &\stackrel{\text{def}}{=} f^*(\Phi_{X',Y}(\xi)) \\ &\stackrel{\text{def}}{=} [f^* \circ \Phi_{X',Y}](\xi), \end{split}$$

for each $\xi \in InfLat((\mathcal{P}(X'), \supset), (Y, \preceq))$, where we have used Item 1 of Definition 4.5.4.1.3 for the fifth equality above.

Naturality for Φ , Part II: We need to show that, given a cocontinuous morphism of posets

$$g: (Y, \preceq_Y) \rightarrow (Y', \preceq_{Y'}),$$

the diagram

$$\mathsf{InfLat}((\mathcal{P}(X),\supset),(Y,\preceq)) \xrightarrow{\Phi_{X,Y}} \mathsf{Sets}(X,Y)$$

$$\downarrow^{g_!} \qquad \qquad \downarrow^{g_!}$$

$$\mathsf{InfLat}((\mathcal{P}(X),\supset),(Y',\preceq)) \xrightarrow{\Phi_{X,Y'}} \mathsf{Sets}(X,Y')$$

commutes. Indeed, we have

$$\begin{split} [\Phi_{X,Y'} \circ g_!](\xi) &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g_!(\xi)) \\ &\stackrel{\text{def}}{=} \Phi_{X,Y'}(g \circ \xi) \\ &\stackrel{\text{def}}{=} (g \circ \xi) \circ \chi_X \\ &= g \circ (\xi \circ \chi_X) \\ &\stackrel{\text{def}}{=} g \circ (\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} g_!(\Phi_{X,Y}(\xi)) \\ &\stackrel{\text{def}}{=} [g_! \circ \Phi_{X,Y}](\xi). \end{split}$$

for each $\xi \in InfLat((\mathcal{P}(X),\supset),(Y,\preceq))$.

Naturality for Ψ : Since Φ is natural in each argument and Φ is a componentwise inverse to Ψ in each argument, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Ψ is also natural in each argument.

Warning 4.4.6.1.3. Although the assignment $X \mapsto \mathcal{P}(X)^{\text{op}}$ is called the *free completion of X*, it is not an idempotent operation, i.e. we have $\mathcal{P}(\mathcal{P}(X)^{\text{op}})^{\text{op}} \neq \mathcal{P}(X)^{\text{op}}$.

4.4.7 The Internal Hom of a Powerset

Let X be a set and let $U, V \in \mathcal{P}(X)$.

Proposition 4.4.7.1.1. The internal Hom of $\mathcal{P}(X)$ from U to V is the subset $[U,V]_X^{22}$ of X given by

$$[U, V]_X = U^{\mathsf{c}} \cup V$$
$$= (U \setminus V)^{\mathsf{c}}$$

where U^{c} is the complement of U of Definition 4.3.11.1.1.

Proof. Proof of the Equality $U^c \cup V = (U \setminus V)^c$: We have

$$(U \setminus V)^{c} \stackrel{\text{def}}{=} X \setminus (U \setminus V)$$

$$= (X \cap V) \cup (X \setminus U)$$

$$= V \cup (X \setminus U)$$

$$\stackrel{\text{def}}{=} V \cup U^{c}$$

$$= U^{c} \cup V,$$

where we have used:

²² Further Notation: Also written $\mathbf{Hom}_{\mathcal{P}(X)}(U,V)$.

- 1. Item 10 of Definition 4.3.10.1.2 for the second equality.
- 2. Item 4 of Definition 4.3.9.1.2 for the third equality.
- 3. Item 4 of Definition 4.3.8.1.2 for the last equality.

This finishes the proof.

Proof that $U^c \cup V$ *Is Indeed the Internal Hom*: This follows from Item 2 of Definition 4.3.9.1.2.

Remark 4.4.7.1.2. Henning Makholm suggests the following heuristic intuition for the internal Hom of $\mathcal{P}(X)$ from U to V ([MSE 267365]):

- 1. Since products in $\mathcal{P}(X)$ are given by binary intersections (Item 1 of Definition 4.4.1.1.4), the right adjoint $\mathbf{Hom}_{\mathcal{P}(X)}(U,-)$ of $U\cap-$ may be thought of as a function type [U,V].
- 2. Under the Curry–Howard correspondence (??), the function type [U, V] corresponds to implication $U \Rightarrow V$.
- 3. Implication $U \Rightarrow V$ is logically equivalent to $\neg U \lor V$.
- 4. The expression $\neg U \lor V$ then corresponds to the set $U^{c} \cup V$ in $\mathcal{P}(X)$.
- 5. The set $U^{c} \vee V$ turns out to indeed be the internal Hom of $\mathcal{P}(X)$.

Proposition 4.4.7.1.3. Let *X* be a set.

1. Functoriality. The assignments $U, V, (U, V) \mapsto \operatorname{Hom}_{\mathcal{P}(X)}$ define functors

$$[U,-]_X \colon (\mathcal{P}(X),\supset) \to (\mathcal{P}(X),\subset),$$

$$[-,V]_X \colon (\mathcal{P}(X),\subset) \to (\mathcal{P}(X),\subset),$$

$$[-_1,-_2]_X \colon (\mathcal{P}(X) \times \mathcal{P}(X),\subset \times \supset) \to (\mathcal{P}(X),\subset).$$

In particular, the following statements hold for each $U, V, A, B \in \mathcal{P}(X)$:

- (a) If $U \subset A$, then $[A, V]_X \subset [U, V]_X$.
- (b) If $V \subset B$, then $[U, V]_X \subset [U, B]_X$.
- (c) If $U \subset A$ and $V \subset B$, then $[A, V]_X \subset [U, B]_X$.

2. Adjointness. We have adjunctions

$$(U \cap - + [U, -]_X): \quad \mathcal{P}(X) \xrightarrow{\bot} \mathcal{P}(X),$$

$$(- \cap V + [V, -]_X): \quad \mathcal{P}(X) \xrightarrow{\bot} \mathcal{P}(X),$$

witnessed by bijections

$$\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U\cap V,W)\cong\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U,[V,W]_X),$$

 $\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U\cap V,W)\cong\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(V,[U,W]_X).$

In particular, the following statements hold for each $U, V, W \in \mathcal{P}(X)$:

- (a) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $U \subset [V, W]_X$.
- (b) The following conditions are equivalent:
 - i. We have $U \cap V \subset W$.
 - ii. We have $V \subset [U, W]_X$.
- 3. Interaction With the Empty Set I. We have

$$[U, \emptyset]_X = U^{\mathsf{c}},$$

 $[\emptyset, V]_X = X,$

natural in $U, V \in \mathcal{P}(X)$.

4. Interaction With X. We have

$$[U, X]_X = X,$$
$$[X, V]_X = V,$$

natural in $U, V \in \mathcal{P}(X)$.

5. Interaction With the Empty Set II. The functor

$$D_X \colon \mathcal{P}(X)^{\mathsf{op}} \to \mathcal{P}(X)$$

defined by

$$D_X \stackrel{\text{def}}{=} [-, \emptyset]_X$$
$$= (-)^{\mathsf{c}}$$

is an involutory isomorphism of categories, making \emptyset into a dualising object for $(\mathcal{P}(X), \cap, X, [-, -]_X)$ in the sense of $\ref{eq:property}$. In particular:

(a) The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{D_X} \mathcal{P}(X)$$

$$\downarrow^{D_X}$$

$$\mathcal{P}(X)^{\mathsf{op}}$$

commutes, i.e. we have

$$\underbrace{D_X(D_X(U))}_{\stackrel{\text{def}}{=}[[U,\emptyset]_X,\emptyset]_X} = U$$

for each $U \in \mathcal{P}(X)$.

(b) The diagram

$$\mathcal{P}(X)^{\mathrm{op}} \times \mathcal{P}(X)^{\mathrm{op}} \stackrel{\cap^{\mathrm{op}}}{\to} \mathcal{P}(X)^{\mathrm{op}}$$

$$\downarrow^{D_X}$$

$$\mathcal{P}(X)^{\mathrm{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X)$$

commutes, i.e. we have

$$\underbrace{D_X(U \cap D_X(V))}_{\substack{\text{def}\\=[U \cap [V,\emptyset]_X,\emptyset]_X}} = [U,V]_X$$

for each $U, V \in \mathcal{P}(X)$.

- 6. Interaction With the Empty Set III. Let $f: X \to Y$ be a function.
 - (a) Interaction With Direct Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_*^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
D_X & & \downarrow D_Y \\
\mathcal{P}(X) & \xrightarrow{f} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_!(D_X(U)) = D_Y(f_*(U))$$

for each $U \in \mathcal{P}(X)$.

(b) Interaction With Inverse Images. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \xrightarrow{f^{-1,\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}}
\downarrow^{D_X}
\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(D_Y(U)) = D_X(f^{-1}(U))$$

for each $U \in \mathcal{P}(X)$.

(c) Interaction With Codirect Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_!^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
D_X & & \downarrow \\
D_Y & & \downarrow \\
\mathcal{P}(X) & \xrightarrow{f_*} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(D_X(U)) = D_Y(f_!(U))$$

for each $U \in \mathcal{P}(X)$.

7. Interaction With Unions of Families of Subsets I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-_{1}, -_{2}]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\cup^{\mathsf{op}} \times \cup^{\mathsf{op}} \qquad \qquad \bigcup \cup$$

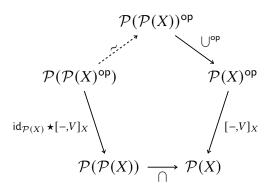
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-_{1}, -_{2}]_{X}} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcup_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcup_{U \in \mathcal{U}} U, \bigcup_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

8. Interaction With Unions of Families of Subsets II. The diagram



commutes, i.e. we have

$$\left[\bigcup_{U\in\mathcal{U}}U,V\right]_X=\bigcap_{U\in\mathcal{U}}[U,V]_X$$

for each $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

9. Interaction With Unions of Families of Subsets III. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcup}{\longrightarrow} \mathcal{P}(X) \\
\downarrow^{[U,-]_X} & & \downarrow^{[U,-]_X} \\
\mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcup}{\longrightarrow} \mathcal{P}(X)
\end{array}$$

commutes, i.e. we have

$$\left[U, \bigcup_{V \in \mathcal{V}} V\right]_X = \bigcup_{V \in \mathcal{V}} [U, V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

10. Interaction With Intersections of Families of Subsets I. The diagram

$$\mathcal{P}(\mathcal{P}(X))^{\mathsf{op}} \times \mathcal{P}(\mathcal{P}(X)) \xrightarrow{[-1,-2]_{\mathcal{P}(X)}} \mathcal{P}(\mathcal{P}(X))$$

$$\uparrow^{\mathsf{op}} \times \uparrow^{\mathsf{op}} \qquad \qquad \downarrow \cap$$

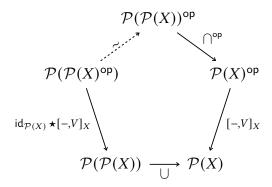
$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_X} \mathcal{P}(X),$$

does not commute in general, i.e. we may have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W \neq \left[\bigcap_{U \in \mathcal{U}} U, \bigcap_{V \in \mathcal{V}} V\right]_X$$

in general, where $\mathcal{U} \in \mathcal{P}(\mathcal{P}(X))$.

11. Interaction With Intersections of Families of Subsets II. The diagram



commutes, i.e. we have

$$\left[\bigcap_{U\in\mathcal{U}}U,V\right]_X=\bigcup_{U\in\mathcal{U}}[U,V]_X$$

for each $U \in \mathcal{P}(\mathcal{P}(X))$ and each $V \in \mathcal{P}(X)$.

12. Interaction With Intersections of Families of Subsets III. The diagram

$$\begin{array}{ccc} \mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcap}{\longrightarrow} \mathcal{P}(X) \\ \operatorname{id}_{\mathcal{P}(X)} \star [U,-]_X & & & \downarrow [U,-]_X \\ & \mathcal{P}(\mathcal{P}(X)) & \stackrel{\bigcap}{\longrightarrow} \mathcal{P}(X) \end{array}$$

commutes, i.e. we have

$$\left[U, \bigcap_{V \in \mathcal{V}} V\right]_X = \bigcap_{V \in \mathcal{V}} [U, V]_X$$

for each $U \in \mathcal{P}(X)$ and each $\mathcal{V} \in \mathcal{P}(\mathcal{P}(X))$.

13. Interaction With Binary Unions. We have equalities of sets

$$[U \cap V, W]_X = [U, W]_X \cup [V, W]_X,$$

 $[U, V \cap W]_X = [U, V]_X \cap [U, W]_X$

for each $U, V, W \in \mathcal{P}(X)$.

14. *Interaction With Binary Intersections*. We have equalities of sets

$$[U \cup V, W]_X = [U, W]_X \cap [V, W]_X,$$

 $[U, V \cup W]_X = [U, V]_X \cup [U, W]_X$

for each $U, V, W \in \mathcal{P}(X)$.

15. Interaction With Differences. We have equalities of sets

$$[U \setminus V, W]_X = [U, W]_X \cup [V^{c}, W]_X$$
$$= [U, W]_X \cup [U, V]_X,$$
$$[U, V \setminus W]_X = [U, V]_X \setminus (U \cap W)$$

for each $U, V, W \in \mathcal{P}(X)$.

16. Interaction With Complements. We have equalities of sets

$$[U^{c}, V]_{X} = U \cup V,$$

$$[U, V^{c}]_{X} = U \cap V,$$

$$[U, V]_{X}^{c} = U \setminus V$$

for each $U, V \in \mathcal{P}(X)$.

17. Interaction With Characteristic Functions. We have

$$\chi_{[U,V]_{\mathcal{P}(X)}}(x) = \max(1 - \chi_U \pmod{2}, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

18. Interaction With Direct Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_*^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \\
 \downarrow [-_1, -_2]_X \downarrow \qquad \qquad \downarrow [-_1, -_2]_Y \\
 \mathcal{P}(X) \xrightarrow{f_!} \mathcal{P}(Y)$$

commutes, i.e. we have an equality of sets

$$f_!([U,V]_X) = [f_*(U), f_!(V)]_Y,$$

natural in $U, V \in \mathcal{P}(X)$.

19. Interaction With Inverse Images. Let $f: X \to Y$ be a function. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1,\mathsf{op}} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow [-_{1},-_{2}]_{Y} \qquad \qquad \downarrow [-_{1},-_{2}]_{X}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_X,$$

natural in $U, V \in \mathcal{P}(X)$.

20. Interaction With Codirect Images. Let $f: X \to Y$ be a function. We have a natural transformation

with components

$$[f_!(U), f_*(V)]_Y \subset f_*([U, V]_X)$$

indexed by $U, V \in \mathcal{P}(X)$.

Proof. Item 1, Functoriality: Since $\mathcal{P}(X)$ is posetal, it suffices to prove Items 1a to 1c.

1. Proof of Item 1a: We have

$$[A, V]_X \stackrel{\text{def}}{=} A^{c} \cup V$$

$$\subset U^{c} \cup V$$

$$\stackrel{\text{def}}{=} [U, V]_X,$$

where we have used:

- (a) Item 1 of Definition 4.3.11.1.2, which states that if $U \subset A$, then $A^c \subset U^c$.
- (b) Item 1a of Item 1 of Definition 4.3.11.1.2, which states that if $A^c \subset U^c$, then $A^c \cup K \subset U^c \cup K$ for any $K \in \mathcal{P}(X)$.
- 2. Proof of Item 1b: We have

$$[U, V]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup V$$
$$\subset U^{\mathsf{c}} \cup B$$
$$\stackrel{\text{def}}{=} [U, B]_X,$$

where we have used Item 1b of Item 1 of Definition 4.3.11.1.2, which states that if $V \subset B$, then $K \cup V \subset K \cup B$ for any $K \in \mathcal{P}(X)$.

3. Proof of Item 1c: We have

$$[A, V]_X \subset [U, V]_X$$

$$\subset [U, B]_X,$$

where we have used Items 1a and 1b.

This finishes the proof.

Item 2, *Adjointness*: This is a repetition of Item 2 of Definition 4.3.9.1.2 and is proved there.

Item 3, Interaction With the Empty Set I: We have

$$[U, \emptyset]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup \emptyset$$
$$= U^{\mathsf{c}},$$

where we have used Item 3 of Definition 4.3.8.1.2, and we have

$$\begin{split} [\emptyset, V]_X &\stackrel{\text{def}}{=} \emptyset^{\mathsf{c}} \cup V \\ &\stackrel{\text{def}}{=} (X \setminus \emptyset) \cup V \\ &= X \cup V \\ &= X, \end{split}$$

where we have used:

- 1. Item 12 of Definition 4.3.10.1.2 for the first equality.
- 2. Item 5 of Definition 4.3.8.1.2 for the last equality.

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). *Item* 4, *Interaction With* X: We have

$$[U,X]_X \stackrel{\text{def}}{=} U^{\mathsf{c}} \cup X$$
$$= X.$$

where we have used Item 5 of Definition 4.3.8.1.2, and we have

$$[X, V]_X \stackrel{\text{def}}{=} X^{c} \cup V$$

$$\stackrel{\text{def}}{=} (X \setminus X) \cup V$$

$$= \emptyset \cup V$$

$$= V,$$

where we have used Item 3 of Definition 4.3.8.1.2 for the last equality. Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). Item 5, Interaction With the Empty Set II: We have

$$D_X(D_X(U)) \stackrel{\text{def}}{=} [[U, \emptyset]_X, \emptyset]_X$$
$$= [U^c, \emptyset]_X$$
$$= (U^c)^c$$
$$= U,$$

where we have used:

- 1. Item 3 for the second and third equalities.
- 2. Item 3 of Definition 4.3.11.1.2 for the fourth equality.

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2), and thus we have

$$[[-,\emptyset]_X,\emptyset]_X \cong id_{\mathcal{P}(X)}$$

This finishes the proof.

Item 6, Interaction With the Empty Set III: Since $D_X = (-)^c$, this is essentially a repetition of the corresponding results for $(-)^c$, namely Items 5 to 7 of Definition 4.3.11.1.2. Item 7, Interaction With Unions of Families of Subsets I: By Item 3 of Definition 4.4.7.1.3, we have

$$[\mathcal{U}, \emptyset]_{\mathcal{P}(X)} = \mathcal{U}^{\mathsf{c}},$$
$$[\mathcal{U}, \emptyset]_{X} = \mathcal{U}^{\mathsf{c}}.$$

With this, the counterexample given in the proof of Item 10 of Definition 4.3.6.1.2 then applies.

Item 8, Interaction With Unions of Families of Subsets II: We have

$$\left[\bigcup_{U \in \mathcal{U}} U, V\right]_X \stackrel{\text{def}}{=} \left(\bigcup_{U \in \mathcal{U}} U\right)^{\mathsf{c}} \cup V$$

$$= \left(\bigcap_{U \in \mathcal{U}} U^{\mathsf{c}}\right) \cup V$$

$$= \bigcap_{U \in \mathcal{U}} (U^{\mathsf{c}} \cup V)$$

$$\stackrel{\text{def}}{=} \bigcap_{U \in \mathcal{U}} [U, V]_X,$$

where we have used:

- 1. Item 11 of Definition 4.3.6.1.2 for the second equality.
- 2. Item 6 of Definition 4.3.7.1.2 for the third equality.

This finishes the proof.

Item 9, Interaction With Unions of Families of Subsets III: We have

$$\bigcup_{V \in \mathcal{V}} [U, V]_X \stackrel{\text{def}}{=} \bigcup_{V \in \mathcal{V}} (U^{\mathsf{c}} \cup V)$$

$$= U^{c} \cup \left(\bigcup_{V \in \mathcal{V}} V\right)$$

$$\stackrel{\text{def}}{=} \left[U, \bigcup_{V \in \mathcal{V}} V\right]_{X}.$$

where we have used Item 6. This finishes the proof.

Item 10, Interaction With Intersections of Families of Subsets I: Let $X = \{0, 1\}$, let $\mathcal{U} = \{\{0, 1\}\}$, and let $\mathcal{V} = \{\{0\}, \{0, 1\}\}$. We have

$$\bigcap_{W \in [\mathcal{U}, \mathcal{V}]_{\mathcal{P}(X)}} W = \bigcap_{W \in \mathcal{P}(X)} W$$
$$= \{0, 1\},$$

whereas

$$\left[\bigcap_{U\in\mathcal{U}}U,\bigcap_{V\in\mathcal{V}}V\right]_X = \left[\{0,1\},\{0\}\right]$$
$$= \{0\},$$

Thus we have

$$\bigcap_{W\in [\mathcal{U},\mathcal{V}]_{\mathcal{P}(X)}}W=\left\{0,1\right\}\neq\left\{0\right\}=\left[\bigcap_{U\in\mathcal{U}}U,\bigcap_{V\in\mathcal{V}}V\right]_X.$$

This finishes the proof.

Item 11, Interaction With Intersections of Families of Subsets II: We have

$$\left[\bigcap_{U\in\mathcal{U}}U,V\right]_X \stackrel{\text{def}}{=} \left(\bigcap_{U\in\mathcal{U}}U\right)^{\mathsf{c}} \cup V$$

$$= \left(\bigcup_{U\in\mathcal{U}}U^{\mathsf{c}}\right) \cup V$$

$$= \bigcup_{U\in\mathcal{U}}(U^{\mathsf{c}} \cup V)$$

$$\stackrel{\text{def}}{=} \bigcup_{U\in\mathcal{U}}[U,V]_X,$$

where we have used:

- 1. Item 12 of Definition 4.3.6.1.2 for the second equality.
- 2. Item 6 of Definition 4.3.7.1.2 for the third equality.

This finishes the proof.

Item 12, Interaction With Intersections of Families of Subsets III: We have

$$\bigcap_{V \in \mathcal{V}} [U, V]_X \stackrel{\text{def}}{=} \bigcap_{V \in \mathcal{V}} (U^{\mathsf{c}} \cup V)$$

$$= U^{\mathsf{c}} \cup \left(\bigcap_{V \in \mathcal{V}} V\right)$$

$$\stackrel{\text{def}}{=} \left[U, \bigcap_{V \in \mathcal{V}} V\right]_Y$$

where we have used Item 6. This finishes the proof. *Item* 13, *Interaction With Binary Unions*: We have

$$[U \cap V, W]_X \stackrel{\text{def}}{=} (U \cap V)^c \cup W$$
$$= (U^c \cup V^c) \cup W$$
$$= (U^c \cup V^c) \cup (W \cup W)$$
$$= (U^c \cup W) \cup (V^c \cup W)$$
$$\stackrel{\text{def}}{=} [U, W]_X \cup [V, W]_X,$$

where we have used:

- 1. Item 2 of Definition 4.3.11.1.2 for the second equality.
- 2. Item 8 of Definition 4.3.8.1.2 for the third equality.
- 3. Several applications of Items 2 and 4 of Definition 4.3.8.1.2 and for the fourth equality.

For the second equality in the statement, we have

$$[U, V \cap W]_X \stackrel{\text{def}}{=} U^{c} \cup (V \cap W)$$
$$= (U^{c} \cup V) \cap (U^{c} \cap W)$$
$$\stackrel{\text{def}}{=} [U, V]_X \cap [U, W]_X,$$

where we have used Item 6 of Definition 4.3.8.1.2 for the second equality.

Item 14, Interaction With Binary Intersections: We have

$$[U \cup V, W]_X \stackrel{\text{def}}{=} (U \cup V)^{c} \cup W$$
$$= (U^{c} \cap V^{c}) \cup W$$
$$= (U^{c} \cup W) \cap (V^{c} \cup W)$$
$$\stackrel{\text{def}}{=} [U, W]_X \cap [V, W]_X,$$

where we have used:

- 1. Item 2 of Definition 4.3.11.1.2 for the second equality.
- 2. Item 6 of Definition 4.3.8.1.2 for the third equality.

Now, for the second equality in the statement, we have

$$[U, V \cup W]_X \stackrel{\text{def}}{=} U^{c} \cup (V \cup W)$$
$$= (U^{c} \cup U^{c}) \cup (V \cup W)$$
$$= (U^{c} \cup V) \cup (U^{c} \cup W)$$
$$\stackrel{\text{def}}{=} [U, V]_X \cup [U, W]_X,$$

where we have used:

- 1. Item 8 of Definition 4.3.8.1.2 for the second equality.
- 2. Several applications of Items 2 and 4 of Definition 4.3.8.1.2 and for the third equality.

This finishes the proof.

Item 15, Interaction With Differences: We have

$$\begin{split} [U \setminus V, W]_X &\stackrel{\text{def}}{=} (U \setminus V)^{\mathsf{c}} \cup W \\ &\stackrel{\text{def}}{=} (X \setminus (U \setminus V)) \cup W \\ &= ((X \cap V) \cup (X \setminus U)) \cup W \\ &= (V \cup (X \setminus U)) \cup W \\ &\stackrel{\text{def}}{=} (V \cup U^{\mathsf{c}}) \cup W \\ &= (V \cup (U^{\mathsf{c}} \cup U^{\mathsf{c}})) \cup W \\ &= (U^{\mathsf{c}} \cup W) \cup (U^{\mathsf{c}} \cup V) \\ &\stackrel{\text{def}}{=} [U, W]_X \cup [U, V]_X, \end{split}$$

where we have used:

- 1. Item 10 of Definition 4.3.10.1.2 for the third equality.
- 2. Item 4 of Definition 4.3.9.1.2 for the fourth equality.
- 3. Item 8 of Definition 4.3.8.1.2 for the sixth equality.
- 4. Several applications of Items 2 and 4 of Definition 4.3.8.1.2 and for the seventh equality.

We also have

$$\begin{split} [U \setminus V, W]_X &\stackrel{\mathrm{def}}{=} (U \setminus V)^{\mathsf{c}} \cup W \\ &\stackrel{\mathrm{def}}{=} (X \setminus (U \setminus V)) \cup W \\ &= ((X \cap V) \cup (X \setminus U)) \cup W \\ &= (V \cup (X \setminus U)) \cup W \\ &\stackrel{\mathrm{def}}{=} (V \cup U^{\mathsf{c}}) \cup W \\ &= (V \cup U^{\mathsf{c}}) \cup (W \cup W) \\ &= (U^{\mathsf{c}} \cup W) \cup (V \cup W) \\ &= (U^{\mathsf{c}} \cup W) \cup ((V^{\mathsf{c}})^{\mathsf{c}} \cup W) \\ &\stackrel{\mathrm{def}}{=} [U, W]_X \cup [V^{\mathsf{c}}, W]_X, \end{split}$$

where we have used:

- 1. Item 10 of Definition 4.3.10.1.2 for the third equality.
- 2. Item 4 of Definition 4.3.9.1.2 for the fourth equality.
- 3. Item 8 of Definition 4.3.8.1.2 for the sixth equality.
- 4. Several applications of Items 2 and 4 of Definition 4.3.8.1.2 and for the seventh equality.
- 5. Item 3 of Definition 4.3.11.1.2 for the eighth equality.

Now, for the second equality in the statement, we have

$$[U, V \setminus W]_X \stackrel{\text{def}}{=} U^{c} \cup (V \setminus W)$$

$$= (V \setminus W) \cup U^{c}$$

$$= (V \cup U^{c}) \setminus (W \setminus U^{c})$$

$$\stackrel{\text{def}}{=} (V \cup U^{c}) \setminus (W \setminus (X \setminus U))$$

$$= (V \cup U^{c}) \setminus ((W \cap U) \cup (W \setminus X))$$

$$= (V \cup U^{c}) \setminus ((W \cap U) \cup \emptyset)$$

$$= (V \cup U^{c}) \setminus (W \cap U)$$

$$= (V \cup U^{c}) \setminus (U \cap W)$$

$$\stackrel{\text{def}}{=} [U, V]_{X} \setminus (U \cap W)$$

where we have used:

- 1. Item 4 of Definition 4.3.8.1.2 for the second equality.
- 2. Item 4 of Definition 4.3.10.1.2 for the third equality.
- 3. Item 10 of Definition 4.3.10.1.2 for the fifth equality.
- 4. Item 13 of Definition 4.3.10.1.2 for the sixth equality.
- 5. Item 3 of Definition 4.3.8.1.2 for the seventh equality.
- 6. Item 5 of Definition 4.3.9.1.2 for the eighth equality.

This finishes the proof.

Item 16, Interaction With Complements: We have

$$[U^{c}, V]_{X} \stackrel{\text{def}}{=} (U^{c})^{c} \cup V,$$
$$= U \cup V,$$

where we have used Item 3 of Definition 4.3.11.1.2. We also have

$$[U, V^{\mathsf{c}}]_X \stackrel{\mathsf{def}}{=} U^{\mathsf{c}} \cup V^{\mathsf{c}}$$
$$= U \cap V$$

where we have used Item 2 of Definition 4.3.11.1.2. Finally, we have

$$[U, V]_X^{c} = ((U \setminus V)^{c})^{c}$$
$$= U \setminus V,$$

where we have used Item 2 of Definition 4.3.11.1.2.

Item 17, Interaction With Characteristic Functions: We have

$$\chi_{[U,V]_{\mathcal{P}(X)}}(x) \stackrel{\text{def}}{=} \chi_{U^{\mathsf{c}} \cup V}(x)$$

$$= \max(\chi_{U^c}, \chi_V)$$

= \text{max}(1 - \chi_U \text{(mod 2)}, \chi_V),

where we have used:

- 1. Item 10 of Definition 4.3.8.1.2 for the second equality.
- 2. Item 4 of Definition 4.3.11.1.2 for the third equality.

This finishes the proof.

Item 18, *Interaction With Direct Images*: This is a repetition of Item 10 of Definition 4.6.1.1.5 and is proved there.

Item 19, *Interaction With Inverse Images*: This is a repetition of Item 10 of Definition 4.6.2.1.3 and is proved there.

Item 20, Interaction With Codirect Images: This is a repetition of Item 9 of Definition 4.6.3.1.7 and is proved there.

4.4.8 Isbell Duality for Sets

Let X be a set.

Definition 4.4.8.1.1. The **Isbell function** of X is the map

$$I: \mathcal{P}(X) \to \mathsf{Sets}(X, \mathcal{P}(X))$$

defined by

$$I(U) \stackrel{\text{def}}{=} \llbracket x \mapsto [U, \{x\}]_X \rrbracket$$

for each $U \in \mathcal{P}(X)$.

Remark 4.4.8.1.2. Recall from Definition 4.4.1.1.2 that we may view the powerset $\mathcal{P}(X)$ of a set X as the decategorification of the category of presheaves $\mathsf{PSh}(C)$ of a category C. Building upon this analogy, we want to mimic the definition of the Isbell Spec functor, which is given on objects by

$$\mathsf{Spec}(\mathcal{F}) \stackrel{\mathsf{def}}{=} \mathsf{Nat}(\mathcal{F}, h_{(-)})$$

for each $\mathcal{F} \in \mathsf{Obj}(\mathsf{PSh}(C))$. To this end, we could define

$$I(U) \stackrel{\text{def}}{=} [U, \chi_{(-)}]_X$$

replacing:

- The Yoneda embedding $X \mapsto h_X$ of C into $\mathsf{PSh}(C)$ with the characteristic embedding $x \mapsto \chi_X$ of X into $\mathcal{P}(X)$ of Definition 4.5.4.1.1.
- The internal Hom Nat of PSh(C) with the internal Hom $[-, -]_X$ of $\mathcal{P}(X)$ of Definition 4.4.7.1.1.

However, since $[U, \chi_x]_X$ is a subset of U instead of a truth value, we get a function

$$I: \mathcal{P}(X) \to \mathsf{Sets}(X, \mathcal{P}(X))$$

instead of a function

$$I: \mathcal{P}(X) \to \mathcal{P}(X).$$

This makes some of the properties involving I a bit more cumbersome to state, although we still have an analogue of Isbell duality in that $I_! \circ I$ evaluates to $id_{\mathcal{P}(X)}$ in the sense of Definition 4.4.8.1.3.

Proposition 4.4.8.1.3. The diagram

$$\mathcal{P}(X) \xrightarrow{\mathsf{I}} \mathsf{Sets}(X, \mathcal{P}(X))$$

$$\Delta_{\Delta_{\mathsf{id}_{\mathcal{P}(X)}}} \qquad \qquad \mathsf{I}_{!}$$

$$\mathsf{Sets}(X, \mathsf{Sets}(X, \mathcal{P}(X)))$$

commutes, i.e. we have

$$I_!(I(U)) = \llbracket x \mapsto \llbracket y \mapsto U \rrbracket \rrbracket$$

for each $U \in \mathcal{P}(X)$.

Proof. We have

$$I_{!}(I(U)) \stackrel{\text{def}}{=} I_{!}([\![x \mapsto U^{c} \cup \{x\}]\!])$$

$$\stackrel{\text{def}}{=} [\![x \mapsto I(U^{c} \cup \{x\})]\!]$$

$$\stackrel{\text{def}}{=} [\![x \mapsto [\![y \mapsto (U^{c} \cup \{x\})^{c} \cup \{x\}]\!]]\!]$$

$$= [\![x \mapsto [\![y \mapsto (U \cap (X \setminus \{x\})) \cup \{x\}]\!]]\!]$$

$$= [\![x \mapsto [\![y \mapsto (U \setminus \{x\}) \cup \{x\}]\!]]\!]$$

$$= [\![x \mapsto [\![y \mapsto U]\!]]\!],$$

where we have used Item 2 of Definition 4.3.11.1.2 for the fourth equality above. \Box

4.5 Characteristic Functions

4.5.1 The Characteristic Function of a Subset

Let X be a set and let $U \in \mathcal{P}(X)$.

Definition 4.5.1.1.1. The **characteristic function of** U^{23} is the function $\chi_U \colon X \to \{t, f\}^{24}$ defined by

$$\chi_U(x) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x \in U, \\ \text{false} & \text{if } x \notin U \end{cases}$$

for each $x \in X$.

Remark 4.5.1.1.2. Under the analogy that $\{t, f\}$ should be the (-1)-categorical analogue of Sets, we may view a function

$$f: X \to \{\mathsf{t},\mathsf{f}\}$$

as a decategorification of presheaves and copresheaves

$$\mathcal{F} \colon C^{\mathsf{op}} \to \mathsf{Sets},$$

 $F \colon C \to \mathsf{Sets}.$

The characteristic functions χ_U of the subsets of X are then the primordial examples of such functions (and, in fact, all of them).

Notation 4.5.1.1.3. We will often employ the bijection $\{t, f\} \cong \{0, 1\}$ to make use of the arithmetical operations defined on $\{0, 1\}$ when disucssing characteristic functions.

Examples of this include Items 4 to 11 of Definition 4.5.1.1.4 below.

Proposition 4.5.1.1.4. Let *X* be a set.

1. Functionality. The assignment $U \mapsto \chi_U$ defines a function

$$\chi_{(-)} \colon \mathcal{P}(X) \to \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\}).$$

2. *Bijectivity*. The function $\chi_{(-)}$ from Item 1 is bijective.

²³ Further Terminology: Also called the **indicator function of** U.

²⁴ Further Notation: Also written $\chi_X(U, -)$ or $\chi_X(-, U)$.

3. Naturality. The collection

$$\left\{\chi_{(-)} \colon \mathcal{P}(X) \to \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})\right\}_{X \in \mathsf{Obi}(\mathsf{Sets})}$$

defines a natural isomorphism between \mathcal{P}^{-1} and $\mathsf{Sets}(-, \{\mathsf{t}, \mathsf{f}\})$. In particular, given a function $f \colon X \to Y$, the diagram

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

$$\chi_{(-)} \downarrow \chi \qquad \qquad \downarrow \chi_{(-)}$$

$$\text{Sets}(Y, \{t, f\}) \xrightarrow{f^*} \text{Sets}(X, \{t, f\})$$

commutes, i.e. we have

$$\chi_V \circ f = \chi_{f^{-1}(V)}$$

for each $V \in \mathcal{P}(Y)$.

4. Interaction With Unions I. We have

$$\chi_{U \cup V} = \max(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

5. Interaction With Unions II. We have

$$\chi_{U \cup V} = \chi_U + \chi_V - \chi_{U \cap V}$$

for each $U, V \in \mathcal{P}(X)$.

6. Interaction With Intersections I. We have

$$\chi_{U\cap V} = \chi_U \chi_V$$

for each $U, V \in \mathcal{P}(X)$.

7. Interaction With Intersections II. We have

$$\chi_{U\cap V} = \min(\chi_U, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

8. Interaction With Differences. We have

$$\chi_{U\setminus V} = \chi_U - \chi_{U\cap V}$$

for each $U, V \in \mathcal{P}(X)$.

9. Interaction With Complements. We have

$$\chi_{U^c} = 1 - \chi_U$$

for each $U \in \mathcal{P}(X)$.

10. Interaction With Symmetric Differences. We have

$$\chi_{U \triangle V} = \chi_U + \chi_V - 2\chi_{U \cap V}$$

and thus, in particular, we have

$$\chi_{U \triangle V} \equiv \chi_U + \chi_V \pmod{2}$$

for each $U, V \in \mathcal{P}(X)$.

11. Interaction With Internal Homs. We have

$$\chi_{[U,V]_{\mathcal{P}(X)}} = \max(1 - \chi_U \pmod{2}, \chi_V)$$

for each $U, V \in \mathcal{P}(X)$.

Proof. Item 1, *Functionality*: There is nothing to prove. *Item* 2, *Bijectivity*: We proceed in three steps:

1. The Inverse of $\chi_{(-)}$. The inverse of $\chi_{(-)}$ is the map

$$\Phi \colon \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\}) \xrightarrow{\sim} \mathcal{P}(X),$$

defined by

$$\Phi(f) \stackrel{\text{def}}{=} U_f$$

$$\stackrel{\text{def}}{=} f^{-1}(\text{true})$$

$$\stackrel{\text{def}}{=} \{x \in X \mid f(x) = \text{true}\}$$

for each $f \in Sets(X, \{t, f\})$.

2. Invertibility I. We have

$$\begin{split} [\Phi \circ \chi_{(-)}](U) &\stackrel{\text{def}}{=} \Phi(\chi_U) \\ &\stackrel{\text{def}}{=} \chi_U^{-1}(\mathsf{true}) \\ &\stackrel{\text{def}}{=} \{x \in X \,|\, \chi_U(x) = \mathsf{true}\} \\ &\stackrel{\text{def}}{=} \{x \in X \,|\, x \in U\} \\ &= U \\ &\stackrel{\text{def}}{=} [\mathsf{id}_{\mathcal{P}(X)}](U) \end{split}$$

for each $U \in \mathcal{P}(X)$. Thus, we have

$$\Phi \circ \chi_{(-)} = id_{\mathcal{P}(X)}$$
.

3. Invertibility II. We have

$$\begin{split} [\chi_{(-)} \circ \Phi](U) &\stackrel{\text{def}}{=} \chi_{\Phi(f)} \\ &\stackrel{\text{def}}{=} \chi_{f^{-1}(\mathsf{true})} \\ &\stackrel{\text{def}}{=} [\![x \mapsto \begin{cases} \mathsf{true} & \mathsf{if} \, x \in f^{-1}(\mathsf{true}) \\ \mathsf{false} & \mathsf{otherwise} \end{cases}]\!] \\ &= [\![x \mapsto f(x)]\!] \\ &= f \\ &\stackrel{\text{def}}{=} [\mathsf{id}_{\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})}](f) \end{split}$$

for each $f \in Sets(X, \{t, f\})$. Thus, we have

$$\chi_{(-)} \circ \Phi = \mathsf{id}_{\mathsf{Sets}(X, \{\mathsf{t},\mathsf{f}\})}$$
 .

This finishes the proof.

Item 3, *Naturality*: We proceed in two steps:

1. Naturality of $\chi_{(-)}$. We have

$$\begin{split} [\chi_V \circ f](v) &\stackrel{\text{def}}{=} \chi_V(f(v)) \\ &= \begin{cases} \text{true} & \text{if } f(v) \in V, \\ \text{false} & \text{otherwise} \end{cases} \\ &= \begin{cases} \text{true} & \text{if } v \in f^{-1}(V), \\ \text{false} & \text{otherwise} \end{cases} \end{split}$$

$$\stackrel{\text{def}}{=} \chi_{f^{-1}(V)}(v)$$

for each $v \in V$.

2. Naturality of Φ . Since $\chi_{(-)}$ is natural and a componentwise inverse to Φ , it follows from Categories, Item 2 of Definition 11.9.7.1.2 that Φ is also natural in each argument.

This finishes the proof.

Item 4, Interaction With Unions I: This is a repetition of Item 10 of Definition 4.3.8.1.2 and is proved there.

Item 5, Interaction With Unions II: This is a repetition of Item 11 of Definition 4.3.8.1.2 and is proved there.

Item 6, *Interaction With Intersections I*: This is a repetition of Item 10 of Definition 4.3.9.1.2 and is proved there.

Item 7, *Interaction With Intersections II*: This is a repetition of Item 11 of Definition 4.3.9.1.2 and is proved there.

Item 8, Interaction With Differences: This is a repetition of Item 16 of Definition 4.3.10.1.2 and is proved there.

Item 9, Interaction With Complements: This is a repetition of Item 4 of Definition 4.3.11.1.2 and is proved there.

Item 10, *Interaction With Symmetric Differences*: This is a repetition of Item 15 of Definition 4.3.12.1.2 and is proved there.

Item 11, *Interaction With Internal Homs*: This is a repetition of Item 17 of Definition 4.4.7.1.3 and is proved there. □

Remark 4.5.1.1.5. The bijection

$$\mathcal{P}(X) \cong \mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$$

of Item 2 of Definition 4.5.1.1.4, which

- · Takes a subset $U \hookrightarrow X$ of X and straightens it to a function $\chi_U \colon X \to \{\text{true, false}\};$
- · Takes a function $f: X \to \{\text{true}, \text{false}\}$ and unstraightens it to a subset $f^{-1}(\text{true}) \hookrightarrow X$ of X:

may be viewed as the (-1)-categorical version of the o-categorical un/straightening isomorphism between indexed and fibred sets

$$\underbrace{\mathsf{FibSets}_X}_{\overset{\mathsf{def}}{=}\mathsf{Sets}_{/X}} \cong \underbrace{\mathsf{ISets}_X}_{\overset{\mathsf{def}}{=}\mathsf{Fun}(X_{\mathsf{disc}},\mathsf{Sets})}$$

of Un/Straightening for Indexed and Fibred Sets, ??. Here we view:

- · Subsets $U \hookrightarrow X$ as being analogous to X-fibred sets $\phi_X \colon A \to X$.
- · Functions $f: X \to \{\mathsf{t}, \mathsf{f}\}$ as being analogous to X-indexed sets $A: X_{\mathsf{disc}} \to \mathsf{Sets}$.

4.5.2 The Characteristic Function of a Point

Let X be a set and let $x \in X$.

Definition 4.5.2.1.1. The **characteristic function of** x is the function x

$$\chi_x \colon X \to \{\mathsf{t},\mathsf{f}\}$$

defined by

$$\chi_x \stackrel{\text{def}}{=} \chi_{\{x\}},$$

i.e. by

$$\chi_x(y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each $y \in X$.

Remark 4.5.2.1.2. Expanding upon Definition 4.5.1.1.2, we may think of the characteristic function

$$\chi_x \colon X \to \{\mathsf{t},\mathsf{f}\}$$

of an element x of X as a decategorification of the representable presheaf and of the representable copresheaf

$$h_X \colon C^{\mathsf{op}} \to \mathsf{Sets},$$

 $h^X \colon C \to \mathsf{Sets}$

associated of an object X of a category C.

4.5.3 The Characteristic Relation of a Set

Let *X* be a set.

²⁵ Further Notation: Also written $\chi^{\overline{x}}$, $\chi_X(x,-)$, or $\chi_X(-,x)$.

Definition 4.5.3.1.1. The **characteristic relation on** X^{26} is the relation²⁷

$$\gamma_X(-1,-2): X \times X \to \{\mathsf{t},\mathsf{f}\}$$

on X defined by 28

$$\chi_X(x,y) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x = y, \\ \text{false} & \text{if } x \neq y \end{cases}$$

for each $x, y \in X$.

Remark 4.5.3.1.2. Expanding upon Definitions 4.5.1.1.2 and 4.5.2.1.2, we may view the characteristic relation

$$\gamma_X(-1,-2): X \times X \to \{\mathsf{t},\mathsf{f}\}$$

of X as a decategorification of the Hom profunctor

$$\operatorname{Hom}_C(-1,-2)\colon C^{\operatorname{op}}\times C\to\operatorname{Sets}$$

of a category C.

Proposition 4.5.3.1.3. Let $f: X \to Y$ be a function.

1. The Inclusion of Characteristic Relations Associated to a Function. Let $f:A\to B$ be a function. We have an inclusion²⁹

$$\chi_{B} \circ (f \times f) \subset \chi_{A}, \qquad A \times A \xrightarrow{f \times f} B \times B$$

$$\chi_{A} \searrow \chi_{A} \searrow \chi_{B}$$

$$\{t, f\}.$$

Proof. Item 1, The Inclusion of Characteristic Relations Associated to a Function: The inclusion $\chi_B(f(a), f(b)) \subset \chi_A(a, b)$ is equivalent to the statement "if a = b, then f(a) = f(b)', which is true.

²⁶ Further Terminology: Also called the **identity relation on** X.

²⁷ Further Notation: Also written χ_{-2}^{-1} , or \sim_{id} in the context of relations.

²⁸Under the bijection Sets($X \times X$, $\{t, f\}$) $\cong \mathcal{P}(X \times X)$ of Item 2 of Definition 4.5.1.1.4, the relation χ_X corresponds to the diagonal $\Delta_X \subset X \times X$ of X.

²⁹ Note: This is the 0-categorical version of Categories, Definition 11.5.4.1.1.

4.5.4 The Characteristic Embedding of a Set

Let *X* be a set.

Definition 4.5.4.1.1. The **characteristic embedding**³⁰ **of** X **into** $\mathcal{P}(X)$ is the function

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

defined by31

$$\chi_{(-)}(x) \stackrel{\text{def}}{=} \chi_x$$
$$= \{x\}$$

for each $x \in X$.

Remark 4.5.4.1.2. Expanding upon Definitions 4.5.1.1.2, 4.5.2.1.2 and 4.5.3.1.2, we may view the characteristic embedding

$$\chi_{(-)}: X \hookrightarrow \mathcal{P}(X)$$

of X into $\mathcal{P}(X)$ as a decategorification of the Yoneda embedding

$$\mbox{\sharp}: \mbox{\mathcal{C}}^{\sf op} \hookrightarrow \mbox{\sf PSh}(\mbox{\mathcal{C}})$$

of a category C into PSh(C).

Proposition 4.5.4.1.3. Let $f: X \to Y$ be a map of sets.

1. Interaction With Functions. We have

$$f_! \circ \chi_X = \chi_Y \circ f, \qquad \chi_X \qquad \downarrow \chi_Y$$

$$\mathcal{P}(X) \xrightarrow{f} \mathcal{P}(Y).$$

Proof. Item 1, Interaction With Functions: Indeed, we have

$$[f_! \circ \chi_X](x) \stackrel{\text{def}}{=} f_!(\chi_X(x))$$

³⁰The name "characteristic *embedding*" is justified by Definition 4.5.5.1.2, which gives an analogue of fully faithfulness for $\chi_{(-)}$.

³¹Here we are identifying $\mathcal{P}(X)$ with Sets $(X, \{t, f\})$ as per Item 2 of Definition 4.5.1.1.4.

$$\stackrel{\text{def}}{=} f_!(\{x\})$$

$$= \{f(x)\}$$

$$\stackrel{\text{def}}{=} \chi_{X'}(f(x))$$

$$\stackrel{\text{def}}{=} [\chi_{X'} \circ f](x),$$

for each $x \in X$, showing the desired equality.

4.5.5 The Yoneda Lemma for Sets

Let X be a set and let $U \subset X$ be a subset of X.

Proposition 4.5.5.1.1. We have

$$\chi_{\mathcal{P}(X)}(\chi_x,\chi_U)=\chi_U(x)$$

for each $x \in X$, giving an equality of functions

$$\chi_{\mathcal{P}(X)}(\chi_{(-)},\chi_U)=\chi_U,$$

where

$$\chi_{\mathcal{P}(X)}(U,V) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } U \subset V, \\ \text{false} & \text{otherwise.} \end{cases}$$

Proof. We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_U) \stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } \{x\} \subset U, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } x \in U \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_U(x).$$

This finishes the proof.

Corollary 4.5.5.1.2. The characteristic embedding is fully faithful, i.e., we have

$$\chi_{\mathcal{P}(X)}(\chi_x,\chi_y)\cong\chi_X(x,y)$$

for each $x, y \in X$.

Proof. We have

$$\chi_{\mathcal{P}(X)}(\chi_x, \chi_y) = \chi_y(x)$$

$$\stackrel{\text{def}}{=} \begin{cases} \text{true} & \text{if } x \in \{y\} \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if } x = y \\ \text{false} & \text{otherwise} \end{cases}$$

$$\stackrel{\text{def}}{=} \chi_X(x, y).$$

where we have used Definition 4.5.5.1.1 for the first equality.

4.6 The Adjoint Triple $f_! \dashv f^{-1} \dashv f_*$

4.6.1 Direct Images

Let $f: X \to Y$ be a function.

Definition 4.6.1.1.1. The direct image function associated to f is the function 32

$$f_i \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by³³

$$f_{!}(U) \stackrel{\text{def}}{=} \left\{ y \in Y \middle| \begin{array}{l} \text{there exists some } x \in U \\ \text{such that } y = f(x) \end{array} \right\}$$
$$= \left\{ f(x) \in Y \middle| x \in U \right\}$$

for each $U \in \mathcal{P}(X)$.

Notation 4.6.1.1.2. Sometimes one finds the notation

$$\exists_f \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

for $f_!$. This notation comes from the fact that the following statements are equivalent, where $y \in Y$ and $U \in \mathcal{P}(X)$:

³² Further Notation: Also written simply $f: \mathcal{P}(X) \to \mathcal{P}(Y)$.

³³ Further Terminology: The set f(U) is called the **direct image of** U **by** f.

- · We have $y \in \exists_f(U)$.
- · There exists some $x \in U$ such that f(x) = y.

We will not make use of this notation elsewhere in Clowder.

Warning 4.6.1.1.3. Notation for direct images between powersets is tricky:

- Direct images for powersets and presheaves are both adjoint to their corresponding inverse image functors. However, the direct image functor for powersets is a *left* adjoint, while the direct image functor for presheaves is a *right* adjoint:
 - (a) *Powersets*. Given a function $f: X \to Y$, we have an inverse image functor

$$f^{-1} \colon \mathcal{P}(Y) \to \mathcal{P}(X).$$

The *left* adjoint of this functor is the usual direct image, defined above in Definition 4.6.1.1.1.

(b) *Presheaves*. Given a morphism of topological spaces $f: X \to Y$, we have an inverse image functor

$$f^{-1} \colon \mathsf{PSh}(Y) \to \mathsf{PSh}(X).$$

The *right* adjoint of this functor is the direct image functor of presheaves, defined in ??.

- 2. The presheaf direct image functor is denoted f_* , but the direct image functor for powersets is denoted $f_!$ (as it,s a left adjoint).
- 3. Adding to the confusion, it,s somewhat common for $f_! : \mathcal{P}(X) \to \mathcal{P}(Y)$ to be denoted f_* .

We chose to write $f_!$ for the direct image to keep the notation aligned with the following similar adjoint situations:

SITUATION	Adjoint String
Functoriality of Powersets	$(f_! \dashv f^{-1} \dashv f_*) \colon \mathcal{P}(X) \xrightarrow{\rightleftarrows} \mathcal{P}(Y)$
Functoriality of Presheaf Categories	$(f_! \dashv f^{-1} \dashv f_*) \colon PSh(X) \xrightarrow{\rightleftarrows} PSh(Y)$
Base Change	$(f_! \dashv f^* \dashv f_*) \colon C_{/X} \stackrel{\rightleftarrows}{\to} C_{/Y}$
Kan Extensions	$(F_! \dashv F^* \dashv F_*) \colon Fun(\mathcal{C}, \mathcal{E}) \xrightarrow{\rightleftarrows} Fun(\mathcal{D}, \mathcal{E})$

Remark 4.6.1.1.4. Identifying $\mathcal{P}(X)$ with $\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ via $\mathsf{Item 2}$ of $\mathsf{Definition 4.5.1.1.4}$, we see that the direct image function associated to f is equivalently the function

$$f_! : \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by

$$f_{!}(\chi_{U}) \stackrel{\text{def}}{=} \operatorname{Lan}_{f}(\chi_{U})$$

$$= \operatorname{colim}((f \overset{\rightarrow}{\times} (\underline{-_{1}})) \stackrel{\text{pr}}{\twoheadrightarrow} A \overset{\chi_{U}}{\longrightarrow} \{\mathsf{t},\mathsf{f}\})$$

$$= \operatorname{colim}_{x \in X} (\chi_{U}(x))$$

$$f(x) = -_{1}$$

$$= \bigvee_{x \in X} (\chi_{U}(x)),$$

$$f(x) = -_{1}$$

where we have used ?? for the second equality. In other words, we have

$$[f!(\chi_U)](y) = \bigvee_{\substack{x \in X \\ f(x) = y}} (\chi_U(x))$$

$$= \begin{cases} \text{true} & \text{if there exists some } x \in X \text{ such} \\ & \text{that } f(x) = y \text{ and } x \in U, \\ \text{false} & \text{otherwise} \end{cases}$$

$$= \begin{cases} \text{true} & \text{if there exists some } x \in U \\ & \text{such that } f(x) = y, \\ \text{false} & \text{otherwise} \end{cases}$$

for each $y \in Y$.

Proposition 4.6.1.1.5. Let $f: X \to Y$ be a function.

1. Functoriality. The assignment $U\mapsto f_!(U)$ defines a functor

$$f_i : (\mathcal{P}(X), \subset) \to (\mathcal{P}(Y), \subset).$$

In particular, for each $U, V \in \mathcal{P}(X)$, the following condition is satisfied:

$$(\star)$$
 If $U \subset V$, then $f_1(U) \subset f_1(V)$.

2. Triple Adjointness. We have a triple adjunction

$$(f_! + f^{-1} + f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$\mathrm{id}_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad \mathrm{id}_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$

 $f_! \circ f^{-1} \hookrightarrow \mathrm{id}_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow \mathrm{id}_{\mathcal{P}(X)},$

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*(f^{-1}(V)),$$

 $f_!(f^{-1}(V)) \subset V, \qquad f^{-1}(f_*(U)) \subset U$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

(b) Bijections of sets

$$\operatorname{\mathsf{Hom}}_{\mathcal{P}(Y)}(f_!(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U,f^{-1}(V)),$$

 $\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(f^{-1}(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U,f_*(V)),$

natural in $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$ and (respectively) $V \in \mathcal{P}(X)$ and $U \in \mathcal{P}(Y)$. In particular:

- i. The following conditions are equivalent:
 - A. We have $f_!(U) \subset V$.
 - B. We have $U \subset f^{-1}(V)$.
- ii. The following conditions are equivalent:
 - A. We have $f^{-1}(U) \subset V$.
 - B. We have $U \subset f_*(V)$.
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\begin{array}{ccc} \mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_{i})_{!}} & \mathcal{P}(\mathcal{P}(Y)) \\ & & & \downarrow \cup \\ & \mathcal{P}(X) & \xrightarrow{f_{i}} & \mathcal{P}(Y) \end{array}$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_!(U) = \bigcup_{V \in f_!(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

4. Interaction With Intersections of Families of Subsets. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_{1})_{!}} & \mathcal{P}(\mathcal{P}(Y)) \\
& & & \downarrow & \\
\mathcal{P}(X) & \xrightarrow{f_{1}} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \mathcal{U}} f_!(U) = \bigcap_{V \in f_!(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_!(\mathcal{U}) \stackrel{\text{def}}{=} (f_!)_!(\mathcal{U})$.

5. Interaction With Binary Unions. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_{!} \times f_{!}} \mathcal{P}(Y) \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(X) \xrightarrow{f_{!}} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$

for each $U, V \in \mathcal{P}(X)$.

6. Interaction With Binary Intersections. We have a natural transformation

with components

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

7. Interaction With Differences. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad$$

with components

$$f_!(U) \setminus f_!(V) \subset f_!(U \setminus V)$$

indexed by $U, V \in \mathcal{P}(X)$.

8. Interaction With Complements. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{f_*^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
\xrightarrow{(-)^c} \qquad \qquad \downarrow^{(-)^c} \\
\mathcal{P}(X) \xrightarrow{f} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_!(U^\mathsf{c}) = f_*(U)^\mathsf{c}$$

for each $U \in \mathcal{P}(X)$.

9. Interaction With Symmetric Differences. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_!^{\mathsf{op}} \times f_!} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(X) \xrightarrow{f_!} \qquad \mathcal{P}(Y)$$

with components

$$f_{!}(U) \triangle f_{!}(V) \subset f_{!}(U \triangle V)$$

indexed by $U, V \in \mathcal{P}(X)$.

10. Interaction With Internal Homs of Powersets. The diagram

commutes, i.e. we have an equality of sets

$$f_!([U,V]_X) = [f_*(U), f_!(V)]_Y,$$

natural in $U, V \in \mathcal{P}(X)$.

11. Preservation of Colimits. We have an equality of sets

$$f!\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f_!(U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{ imes I}$. In particular, we have equalities

$$f_!(U) \cup f_!(V) = f_!(U \cup V),$$

 $f_!(\emptyset) = \emptyset,$

natural in $U, V \in \mathcal{P}(X)$.

12. Oplax Preservation of Limits. We have an inclusion of sets

$$f_! \left(\bigcap_{i \in I} U_i\right) \subset \bigcap_{i \in I} f_! (U_i),$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{ imes I}$. In particular, we have inclusions

$$f_!(U \cap V) \subset f_!(U) \cap f_!(V),$$

 $f_!(X) \subset Y.$

natural in $U, V \in \mathcal{P}(X)$.

13. Symmetric Strict Monoidality With Respect to Unions. The direct image function of Item 1 has a symmetric strict monoidal structure

$$(f_!, f_!^{\otimes}, f_{!|\mathbb{1}}^{\otimes}) \colon (\mathcal{P}(X), \cup, \emptyset) \to (\mathcal{P}(Y), \cup, \emptyset),$$

being equipped with equalities

$$\begin{split} f^{\otimes}_{!|U,V} \colon f_!(U) \cup f_!(V) &\stackrel{=}{\to} f_!(U \cup V), \\ f^{\otimes}_{!|\mathbb{1}} \colon \emptyset &\stackrel{=}{\to} \emptyset, \end{split}$$

natural in $U, V \in \mathcal{P}(X)$.

14. Symmetric Oplax Monoidality With Respect to Intersections. The direct image function of Item 1 has a symmetric oplax monoidal structure

$$(f_!, f_!^{\otimes}, f_{!|\mathbb{1}}^{\otimes}) \colon (\mathcal{P}(X), \cap, X) \to (\mathcal{P}(Y), \cap, Y),$$

being equipped with inclusions

$$f_{!|U,V}^{\otimes} \colon f_{!}(U \cap V) \hookrightarrow f_{!}(U) \cap f_{!}(V),$$
$$f_{!|\mathcal{X}}^{\otimes} \colon f_{!}(X) \hookrightarrow Y,$$

natural in $U, V \in \mathcal{P}(X)$.

15. Interaction With Coproducts. Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \coprod g)_!(U \coprod V) = f_!(U) \coprod g_!(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

16. Interaction With Products. Let $f\colon X\to X'$ and $g\colon Y\to Y'$ be maps of sets. We have

$$(f \boxtimes_{X \times Y} q)_!(U \boxtimes_{X \times Y} V) = f_!(U) \boxtimes_{X' \times Y'} q_!(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

17. Relation to Codirect Images. We have

$$f_!(U) = f_*(U^{\mathsf{c}})^{\mathsf{c}}$$

$$\stackrel{\text{def}}{=} Y \setminus f_*(X \setminus U)$$

for each $U \in \mathcal{P}(X)$.

Proof. Item 1, Functoriality: Omitted.

Item 2, Triple Adjointness: This follows from Definition 4.6.1.1.4, Definition 4.6.2.1.2,

Definition 4.6.3.1.4, and Kan Extensions, ?? of ??.

Item 3, Interaction With Unions of Families of Subsets: We have

$$\bigcup_{V \in f_!(\mathcal{U})} V = \bigcup_{V \in \{f_!(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcup_{U \in \mathcal{U}} f_!(U).$$

This finishes the proof.

Item 4, Interaction With Intersections of Families of Subsets: We have

$$\bigcap_{V \in f_!(\mathcal{U})} V = \bigcap_{V \in \{f_!(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcap_{U \in \mathcal{U}} f_!(U).$$

This finishes the proof.

Item 5, *Interaction With Binary Unions*: See [Pro25p].

Item 6, *Interaction With Binary Intersections*: See [Pro25n].

Item 7, Interaction With Differences: See [Pro250].

Item 8, Interaction With Complements: Applying Item 17 to $X \setminus U$, we have

$$f_{!}(U^{c}) = f_{!}(X \setminus U)$$

$$= Y \setminus f_{*}(X \setminus (X \setminus U))$$

$$= Y \setminus f_{*}(U)$$

$$= f_{*}(U)^{c}.$$

This finishes the proof.

Item 9, Interaction With Symmetric Differences: We have

$$f_{!}(U) \triangle f_{!}(V) = (f_{!}(U) \cup f_{!}(V)) \setminus (f_{!}(U) \cap f_{!}(V))$$

$$\subset (f_{!}(U) \cup f_{!}(V)) \setminus (f_{!}(U \cap V))$$

$$= (f_{!}(U \cup V)) \setminus (f_{!}(U \cap V))$$

$$\subset f_{!}((U \cup V) \setminus (U \cap V))$$

$$= f_{!}(U \triangle V),$$

where we have used:

- 1. Item 2 of Definition 4.3.12.1.2 for the first equality.
- 2. Item 6 of this proposition together with Item 1 of Definition 4.3.10.1.2 for the first inclusion.
- 3. Item 5 for the second equality.
- 4. Item 7 for the second inclusion.
- 5. Item 2 of Definition 4.3.12.1.2 for the tchird equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 10, Interaction With Internal Homs of Powersets: We have

$$f_!([U, V]_X) \stackrel{\text{def}}{=} f_!(U^c \cup V)$$

$$= f_!(U^c) \cup f_!(V)$$

$$= f_*(U)^c \cup f_!(V)$$

$$\stackrel{\text{def}}{=} [f_*(U), f_!(V)]_Y,$$

where we have used:

- 1. Item 5 for the second equality.
- 2. Item 17 for the third equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 11, Preservation of Colimits: This follows from Item 2 and ??, ?? of ??.³⁴

Item 12, *Oplax Preservation of Limits*: The inclusion $f_!(X) \subset Y$ is automatic. See [Pro25n] for the other inclusions.

Item 13, Symmetric Strict Monoidality With Respect to Unions: This follows from Item 11. Item 14, Symmetric Oplax Monoidality With Respect to Intersections: The inclusions in the statement follow from Item 12. Since $\mathcal{P}(Y)$ is posetal, the commutativity of the diagrams in the definition of a symmetric oplax monoidal functor is automatic (Categories, Item 4 of Definition 11.2.7.1.2).

Item 15, Interaction With Coproducts: Omitted.

Item 16, Interaction With Products: Omitted.

Item 17, Relation to Codirect Images: Applying Item 16 of Definition 4.6.3.1.7 to $X \setminus U$,

³⁴ Reference: [Pro25p].

we have

$$f_*(X \setminus U) = B \setminus f_!(X \setminus (X \setminus U))$$
$$= B \setminus f_!(U).$$

Taking complements, we then obtain

$$f_!(U) = B \setminus (B \setminus f_!(U)),$$

= $B \setminus f_*(X \setminus U),$

which finishes the proof.

Proposition 4.6.1.1.6. Let $f: X \to Y$ be a function.

1. Functionality I. The assignment $f \mapsto f_!$ defines a function

$$(-)_{*|X,Y} : \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(X),\mathcal{P}(Y)).$$

2. Functionality II. The assignment $f \mapsto f_!$ defines a function

$$(-)_{*|X,Y} \colon \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(X),\subset),(\mathcal{P}(Y),\subset)).$$

3. Interaction With Identities. For each $X \in Obj(Sets)$, we have

$$(id_X)_1 = id_{\mathcal{D}(X)}$$
.

4. *Interaction With Composition*. For each pair of composable functions $f: X \to Y$ and $g: Y \to Z$, we have

$$(g \circ f)_{!} = g_{!} \circ f_{!}, \qquad \begin{matrix} \mathcal{P}(X) & \xrightarrow{f_{!}} & \mathcal{P}(Y) \\ & \downarrow g_{!} & & \downarrow g_{!} \\ & \mathcal{P}(Z). \end{matrix}$$

Proof. Item 1, *Functionality I*: There is nothing to prove.

Item 2, Functionality II: This follows from Item 1 of Definition 4.6.1.1.5.

Item 3, Interaction With Identities: This follows from Definition 4.6.1.1.4 and Kan Extensions, ?? of ??.

Item 4, Interaction With Composition: This follows from Definition 4.6.1.1.4 and Kan Extensions, ?? of ??.

4.6.2 Inverse Images

Let $f: X \to Y$ be a function.

Definition 4.6.2.1.1. The inverse image function associated to f is the function³⁵

$$f^{-1} \colon \mathcal{P}(Y) \to \mathcal{P}(X)$$

defined by³⁶

$$f^{-1}(V) \stackrel{\text{def}}{=} \{x \in X \mid \text{we have } f(x) \in V\}$$

for each $V \in \mathcal{P}(Y)$.

Remark 4.6.2.1.2. Identifying $\mathcal{P}(Y)$ with $\mathsf{Sets}(Y, \{\mathsf{t}, \mathsf{f}\})$ via $\mathsf{Item\,2}$ of $\mathsf{Definition\,4.5.1.1.4}$, we see that the inverse image function associated to f is equivalently the function

$$f^* \colon \mathcal{P}(Y) \to \mathcal{P}(X)$$

defined by

$$f^*(\chi_V)\stackrel{\mathrm{def}}{=} \chi_V\circ f$$

for each $\chi_V \in \mathcal{P}(Y)$, where $\chi_V \circ f$ is the composition

$$X \xrightarrow{f} Y \xrightarrow{\chi_V} \{\text{true, false}\}$$

in Sets.

Proposition 4.6.2.1.3. Let $f: X \to Y$ be a function.

1. Functoriality. The assignment $V\mapsto f^{-1}(V)$ defines a functor

$$f^{-1} \colon (\mathcal{P}(Y), \subset) \to (\mathcal{P}(X), \subset).$$

In particular, for each $U, V \in \mathcal{P}(Y)$, the following condition is satisfied:

$$(\star)$$
 If $U \subset V$, then $f^{-1}(U) \subset f^{-1}(V)$.

³⁵ Further Notation: Also written $f^* : \mathcal{P}(Y) \to \mathcal{P}(X)$.

³⁶ Further Terminology: The set $f^{-1}(V)$ is called the **inverse image of** V by f.

2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$\mathrm{id}_{\mathcal{P}(X)} \hookrightarrow f^{-1} \circ f_!, \qquad \mathrm{id}_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1},$$

 $f_! \circ f^{-1} \hookrightarrow \mathrm{id}_{\mathcal{P}(Y)}, \qquad f^{-1} \circ f_* \hookrightarrow \mathrm{id}_{\mathcal{P}(X)},$

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*(f^{-1}(V)),$$

 $f_!(f^{-1}(V)) \subset V, \qquad f^{-1}(f_*(U)) \subset U$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

(b) Bijections of sets

$$\operatorname{\mathsf{Hom}}_{\mathcal{P}(Y)}(f_!(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U,f^{-1}(V)),$$

 $\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(f^{-1}(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U,f_*(V)),$

natural in $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$ and (respectively) $V \in \mathcal{P}(X)$ and $U \in \mathcal{P}(Y)$. In particular:

- i. The following conditions are equivalent:
 - A. We have $f_!(U) \subset V$.
 - B. We have $U \subset f^{-1}(V)$.
- ii. The following conditions are equivalent:
 - A. We have $f^{-1}(U) \subset V$.
 - B. We have $U \subset f_*(V)$.
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\cup \qquad \qquad \qquad \bigcup \qquad \qquad \qquad \bigcup \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcup_{V \in \mathcal{V}} f^{-1}(V) = \bigcup_{U \in f^{-1}(\mathcal{U})} U$$

for each $\mathcal{V} \in \mathcal{P}(Y)$, where $f^{-1}(\mathcal{V}) \stackrel{\text{def}}{=} (f^{-1})^{-1}(\mathcal{V})$.

4. Interaction With Intersections of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(Y)) \xrightarrow{(f^{-1})^{-1}} \mathcal{P}(\mathcal{P}(X))$$

$$\uparrow \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\bigcap_{V \in \mathcal{V}} f^{-1}(V) = \bigcap_{U \in f^{-1}(\mathcal{U})} U$$

for each $\mathcal{V} \in \mathcal{P}(Y)$, where $f^{-1}(\mathcal{V}) \stackrel{\text{def}}{=} (f^{-1})^{-1}(\mathcal{V})$

5. Interaction With Binary Unions. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \downarrow \cup$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

6. Interaction With Binary Intersections. The diagram

$$\mathcal{P}(Y) \times \mathcal{P}(Y) \xrightarrow{f^{-1} \times f^{-1}} \mathcal{P}(X) \times \mathcal{P}(X)$$

$$\uparrow \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U \cap V) = f^{-1}(U) \cap f^{-1}(V)$$

for each $U, V \in \mathcal{P}(Y)$.

7. Interaction With Differences. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{\mathsf{op},-1} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad$$

commutes, i.e. we have

$$f^{-1}(U \setminus V) = f^{-1}(U) \setminus f^{-1}(V)$$

for each $U, V \in \mathcal{P}(X)$.

8. Interaction With Complements. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \xrightarrow{f^{-1,\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}} \\
\xrightarrow{(-)^{\mathsf{c}}} \qquad \qquad \downarrow^{(-)^{\mathsf{c}}} \\
\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(U^{\mathsf{c}}) = f^{-1}(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

9. Interaction With Symmetric Differences. The diagram

i.e. we have

$$f^{-1}(U) \triangle f^{-1}(V) = f^{-1}(U \triangle V)$$

for each $U, V \in \mathcal{P}(Y)$.

10. Interaction With Internal Homs of Powersets. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1,\mathsf{op}} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$[-1,-2]_Y \downarrow \qquad \qquad \downarrow [-1,-2]_X$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_X,$$

natural in $U, V \in \mathcal{P}(X)$.

11. Preservation of Colimits. We have an equality of sets

$$f^{-1}\left(\bigcup_{i\in I}U_i\right)=\bigcup_{i\in I}f^{-1}(U_i),$$

natural in $\{U_i\}_{i\in I}\in\mathcal{P}(Y)^{\times I}$. In particular, we have equalities

$$f^{-1}(U) \cup f^{-1}(V) = f^{-1}(U \cup V),$$

 $f^{-1}(\emptyset) = \emptyset,$

natural in $U, V \in \mathcal{P}(Y)$.

12. Preservation of Limits. We have an equality of sets

$$f^{-1}\left(\bigcap_{i\in I}U_i\right)=\bigcap_{i\in I}f^{-1}(U_i),$$

natural in $\{U_i\}_{i\in I}\in\mathcal{P}(Y)^{\times I}$. In particular, we have equalities

$$f^{-1}(U) \cap f^{-1}(V) = f^{-1}(U \cap V),$$

 $f^{-1}(Y) = X,$

natural in $U, V \in \mathcal{P}(Y)$.

13. Symmetric Strict Monoidality With Respect to Unions. The inverse image function of Item 1 has a symmetric strict monoidal structure

$$(f^{-1}, f^{-1, \otimes}, f_{\mathbb{I}}^{-1, \otimes}) \colon (\mathcal{P}(Y), \cup, \emptyset) \to (\mathcal{P}(X), \cup, \emptyset),$$

being equipped with equalities

$$f_{U,V}^{-1,\otimes} : f^{-1}(U) \cup f^{-1}(V) \xrightarrow{=} f^{-1}(U \cup V),$$

$$f_{1}^{-1,\otimes} : \emptyset \xrightarrow{=} f^{-1}(\emptyset),$$

natural in $U, V \in \mathcal{P}(Y)$.

14. Symmetric Strict Monoidality With Respect to Intersections. The inverse image function of Item1 has a symmetric strict monoidal structure

$$(f^{-1}, f^{-1, \otimes}, f_{\mathbb{1}}^{-1, \otimes}) \colon (\mathcal{P}(Y), \cap, Y) \to (\mathcal{P}(X), \cap, X),$$

being equipped with equalities

$$f_{U,V}^{-1,\otimes} \colon f^{-1}(U) \cap f^{-1}(V) \xrightarrow{=} f^{-1}(U \cap V),$$

$$f_{1}^{-1,\otimes} \colon X \xrightarrow{=} f^{-1}(Y),$$

natural in $U, V \in \mathcal{P}(Y)$.

15. Interaction With Coproducts. Let $f: X \to X'$ and $g: Y \to Y'$ be maps of sets. We have

$$(f \coprod g)^{-1}(U' \coprod V') = f^{-1}(U') \coprod g^{-1}(V')$$

for each $U' \in \mathcal{P}(X')$ and each $V' \in \mathcal{P}(Y')$.

16. Interaction With Products. Let $f\colon X\to X'$ and $g\colon Y\to Y'$ be maps of sets. We have

$$(f \boxtimes_{X' \times Y'} g)^{-1} (U' \boxtimes_{X' \times Y'} V') = f^{-1} (U') \boxtimes_{X \times Y} g^{-1} (V')$$

for each $U' \in \mathcal{P}(X')$ and each $V' \in \mathcal{P}(Y')$.

Proof. Item 1, *Functoriality*: Omitted.

Item 2, *Triple Adjointness*: This follows from Definition 4.6.1.1.4, Definition 4.6.2.1.2, Definition 4.6.3.1.4, and Kan Extensions, ?? of ??.

Item 3, Interaction With Unions of Families of Subsets: We have

$$\bigcup_{U \in f^{-1}(\mathcal{V})} U = \bigcup_{U \in \{f^{-1}(V) \in \mathcal{P}(X) \mid V \in \mathcal{V}\}} U$$
$$= \bigcup_{V \in \mathcal{V}} f^{-1}(V).$$

This finishes the proof.

Item 4, Interaction With Intersections of Families of Subsets: We have

$$\bigcap_{U \in f^{-1}(\mathcal{V})} U = \bigcap_{U \in \{f^{-1}(V) \in \mathcal{P}(X) \mid V \in \mathcal{V}\}} U$$
$$= \bigcap_{V \in \mathcal{V}} f^{-1}(V).$$

This finishes the proof.

Item 5, Interaction With Binary Unions: See [Pro25y].

Item 6, *Interaction With Binary Intersections*: See [Pro25w].

Item 7, *Interaction With Differences*: See [Pro25x].

Item 8, Interaction With Complements: See [Pro25j].

Item 9, Interaction With Symmetric Differences: We have

$$f^{-1}(U \triangle V) = f^{-1}((U \cup V) \setminus (U \cap V))$$

$$= f^{-1}(U \cup V) \setminus f^{-1}(U \cap V)$$

$$= f^{-1}(U) \cup f^{-1}(V) \setminus f^{-1}(U \cap V)$$

$$= f^{-1}(U) \cup f^{-1}(V) \setminus f^{-1}(U) \cap f^{-1}(V)$$

$$= f^{-1}(U) \triangle f^{-1}(V),$$

where we have used:

- 1. Item 2 of Definition 4.3.12.1.2 for the first equality.
- 2. Item 7 for the second equality.
- 3. Item 5 for the third equality.
- 4. Item 6 for the fourth equality.
- 5. Item 2 of Definition 4.3.12.1.2 for the fifth equality.

This finishes the proof.

Item 10, Interaction With Internal Homs of Powersets: We have

$$f^{-1}([U, V]_Y) \stackrel{\text{def}}{=} f^{-1}(U^{c} \cup V)$$

$$= f^{-1}(U^{c}) \cup f^{-1}(V)$$

$$= f^{-1}(U)^{c} \cup f^{-1}(V)$$

$$\stackrel{\text{def}}{=} [f^{-1}(U), f^{-1}(V)]_X,$$

where we have used:

- 1. Item 8 for the second equality.
- 2. Item 5 for the third equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 11, Preservation of Colimits: This follows from Item 2 and ??, ?? of ??. 37

Item 12, Preservation of Limits: This follows from Item 2 and ??, ?? of ??. 38

Item 13, Symmetric Strict Monoidality With Respect to Unions: This follows from Item 11.

Item 14, *Symmetric Strict Monoidality With Respect to Intersections*: This follows from Item 12.

Item 15, Interaction With Coproducts: Omitted.

Item 16, Interaction With Products: Omitted.

Proposition 4.6.2.1.4. Let $f: X \to Y$ be a function.

1. Functionality I. The assignment $f \mapsto f^{-1}$ defines a function

$$(-)_{X,Y}^{-1} \colon \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(Y),\mathcal{P}(X)).$$

2. Functionality II. The assignment $f\mapsto f^{-1}$ defines a function

$$(-)^{-1}_{XY} : \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(Y),\subset),(\mathcal{P}(X),\subset)).$$

3. *Interaction With Identities.* For each $X \in Obj(Sets)$, we have

$$\operatorname{id}_X^{-1} = \operatorname{id}_{\mathcal{P}(X)}$$
.

³⁷ Reference: [Pro25v].

³⁸ Reference: [Pro25w].

4. *Interaction With Composition*. For each pair of composable functions $f: X \to Y$ and $g: Y \to Z$, we have

$$(g \circ f)^{-1} = f^{-1} \circ g^{-1},$$

$$\mathcal{P}(Z) \xrightarrow{g^{-1}} \mathcal{P}(Y)$$

$$\downarrow^{f^{-1}}$$

$$\mathcal{P}(X).$$

Proof. Item 1, *Functionality I:* There is nothing to prove.

Item 2, Functionality II: This follows from Item 1 of Definition 4.6.2.1.3.

Item 3, *Interaction With Identities*: This follows from Definition 4.6.2.1.2 and Categories, Item 5 of Definition 11.1.4.1.2.

Item 4, Interaction With Composition: This follows from Definition 4.6.2.1.2 and Categories, Item 2 of Definition 11.1.4.1.2.

4.6.3 Codirect Images

Let $f: X \to Y$ be a function.

Definition 4.6.3.1.1. The **codirect image function associated to** f is the function

$$f_* \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by^{39,40}

$$f_*(U) \stackrel{\text{def}}{=} \left\{ y \in Y \middle| \begin{array}{l} \text{for each } x \in X, \text{ if we have} \\ f(x) = y, \text{ then } x \in U \end{array} \right\}$$
$$= \left\{ y \in Y \middle| \text{ we have } f^{-1}(y) \subset U \right\}$$

for each $U \in \mathcal{P}(X)$.

$$f_*(U) = f_!(U^{\mathsf{c}})^{\mathsf{c}}$$

$$\stackrel{\text{def}}{=} Y \setminus f_!(X \setminus U);$$

see Item 16 of Definition 4.6.3.1.7.

³⁹ Further Terminology: The set $f_*(U)$ is called the **codirect image of** U by f.

⁴⁰We also have

Notation 4.6.3.1.2. Sometimes one finds the notation

$$\forall_f \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

for $f_!$. This notation comes from the fact that the following statements are equivalent, where $y \in Y$ and $U \in \mathcal{P}(X)$:

- · We have $y \in \forall_f(U)$.
- · For each $x \in X$, if y = f(x), then $x \in U$.

We will not make use of this notation elsewhere in Clowder.

Warning 4.6.3.1.3. See Definition 4.6.1.1.3.

Remark 4.6.3.1.4. Identifying $\mathcal{P}(X)$ with $\mathsf{Sets}(X, \{\mathsf{t}, \mathsf{f}\})$ via $\mathsf{Item 2}$ of $\mathsf{Definition 4.5.1.1.4}$, we see that the codirect image function associated to f is equivalently the function

$$f_* \colon \mathcal{P}(X) \to \mathcal{P}(Y)$$

defined by

$$f_*(\chi_U) \stackrel{\text{def}}{=} \operatorname{Ran}_f(\chi_U)$$

$$= \lim ((\underbrace{(-_1)}_{x \in X} \xrightarrow{x} f) \stackrel{\text{pr}}{\to} X \xrightarrow{\chi_U} \{\text{true, false}\})$$

$$= \lim_{\substack{x \in X \\ f(x) = -_1}} (\chi_U(x))$$

$$= \bigwedge_{\substack{x \in X \\ f(x) = -_1}} (\chi_U(x)).$$

where we have used ?? for the second equality. In other words, we have

$$\begin{split} [f_*(\chi_U)](y) &= \bigwedge_{\substack{x \in X \\ f(x) = y}} (\chi_U(x)) \\ &= \begin{cases} \text{true} & \text{if, for each } x \in X \text{ such that} \\ & f(x) = y, \text{ we have } x \in U, \\ \text{false} & \text{otherwise} \end{cases} \\ &= \begin{cases} \text{true} & \text{if } f^{-1}(y) \subset U \\ \text{false} & \text{otherwise} \end{cases} \end{split}$$

for each $y \in Y$.

Definition 4.6.3.1.5. Let U be a subset of X.^{41,42}

1. The **image part of the codirect image** $f_*(U)$ **of** U is the set $f_{*, \text{im}}(U)$ defined by

$$f_{*,\mathsf{im}}(U) \stackrel{\mathsf{def}}{=} f_*(U) \cap \mathsf{Im}(f)$$

$$= \left\{ y \in Y \middle| \begin{array}{l} \mathsf{we have} \ f^{-1}(y) \subset U \\ \mathsf{and} \ f^{-1}(y) \neq \emptyset. \end{array} \right\}.$$

2. The complement part of the codirect image $f_*(U)$ of U is the set $f_{*,\mathrm{cp}}(U)$ defined by

$$\begin{split} f_{*,\mathsf{cp}}(U) &\stackrel{\mathsf{def}}{=} f_*(U) \cap (Y \setminus \mathsf{Im}(f)) \\ &= Y \setminus \mathsf{Im}(f) \\ &= \left\{ y \in Y \middle| \begin{aligned} \mathsf{we have} \, f^{-1}(y) \subset U \\ \mathsf{and} \, f^{-1}(y) = \emptyset. \end{aligned} \right\} \\ &= \left\{ y \in Y \middle| f^{-1}(y) = \emptyset \right\}. \end{split}$$

Example 4.6.3.1.6. Here are some examples of codirect images.

1. Multiplication by Two. Consider the function $f: \mathbb{N} \to \mathbb{N}$ given by

$$f(n) \stackrel{\text{def}}{=} 2n$$

$$f_*(U) = f_{*,im}(U) \cup f_{*,cp}(U),$$

as

$$\begin{split} f_*(U) &= f_*(U) \cap Y \\ &= f_*(U) \cap (\operatorname{Im}(f) \cup (Y \setminus \operatorname{Im}(f))) \\ &= (f_*(U) \cap \operatorname{Im}(f)) \cup (f_*(U) \cap (Y \setminus \operatorname{Im}(f))) \\ &\stackrel{\text{def}}{=} f_{*,\operatorname{im}}(U) \cup f_{*,\operatorname{cp}}(U). \end{split}$$

$$f_*(\chi_U) = \bigwedge_{\substack{x \in X \\ f(x) = -1}} (\chi_U(x)),$$

⁴¹Note that we have

⁴²In terms of the meet computation of $f_*(U)$ of Definition 4.6.3.1.4, namely

for each $n \in \mathbb{N}$. Since f is injective, we have

$$f_{*,im}(U) = f_!(U)$$

 $f_{*,cp}(U) = \{ \text{odd natural numbers} \}$

for any $U \subset \mathbb{N}$. In particular, we have

 $f_*(\{\text{even natural numbers}\}) = \mathbb{N}.$

2. Parabolas. Consider the function $f: \mathbb{R} \to \mathbb{R}$ given by

$$f(x) \stackrel{\text{def}}{=} x^2$$

for each $x \in \mathbb{R}$. We have

$$f_{*,\mathsf{cp}}(U) = \mathbb{R}_{<0}$$

for any $U \subset \mathbb{R}$. Moreover, since $f^{-1}(x) = \{-\sqrt{x}, \sqrt{x}\}$, we have e.g.:

$$f_{*,\text{im}}([0,1]) = \{0\},$$

$$f_{*,\text{im}}([-1,1]) = [0,1],$$

$$f_{*,\text{im}}([1,2]) = \emptyset,$$

$$f_{*,\text{im}}([-2,-1] \cup [1,2]) = [1,4].$$

3. *Circles*. Consider the function $f: \mathbb{R}^2 \to \mathbb{R}$ given by

$$f(x,y) \stackrel{\text{def}}{=} x^2 + y^2$$

for each $(x, y) \in \mathbb{R}^2$. We have

$$f_{*,\mathsf{cp}}(U) = \mathbb{R}_{<0}$$

for any $U \subset \mathbb{R}^2$, and since

$$f^{-1}(r) = \begin{cases} \text{a circle of radius } r \text{ about the origin} & \text{if } r > 0, \\ \{(0,0)\} & \text{if } r = 0, \\ \emptyset & \text{if } r < 0, \end{cases}$$

we have e.g.:

$$f_{*,\text{im}}([-1,1] \times [-1,1]) = [0,1],$$

 $f_{*,\text{im}}(([-1,1] \times [-1,1]) \setminus [-1,1] \times \{0\}) = \emptyset.$

Proposition 4.6.3.1.7. Let $f: X \to Y$ be a function.

1. Functoriality. The assignment $U \mapsto f_*(U)$ defines a functor

$$f_* : (\mathcal{P}(X), \subset) \to (\mathcal{P}(Y), \subset).$$

In particular, for each $U, V \in \mathcal{P}(X)$, the following condition is satisfied:

- (\star) If $U \subset V$, then $f_*(U) \subset f_*(V)$.
- 2. Triple Adjointness. We have a triple adjunction

$$(f_! \dashv f^{-1} \dashv f_*): \mathcal{P}(X) \leftarrow f^{-1} - \mathcal{P}(Y),$$

witnessed by:

(a) Units and counits of the form

$$\begin{split} \mathrm{id}_{\mathcal{P}(X)} &\hookrightarrow f^{-1} \circ f_!, & \mathrm{id}_{\mathcal{P}(Y)} \hookrightarrow f_* \circ f^{-1}, \\ f_! \circ f^{-1} &\hookrightarrow \mathrm{id}_{\mathcal{P}(Y)}, & f^{-1} \circ f_* \hookrightarrow \mathrm{id}_{\mathcal{P}(X)}, \end{split}$$

having components of the form

$$U \subset f^{-1}(f_!(U)), \qquad V \subset f_*(f^{-1}(V)),$$

 $f_!(f^{-1}(V)) \subset V, \qquad f^{-1}(f_*(U)) \subset U$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

(b) Bijections of sets

$$\operatorname{\mathsf{Hom}}_{\mathcal{P}(Y)}(f_!(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U,f^{-1}(V)),$$

 $\operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(f^{-1}(U),V) \cong \operatorname{\mathsf{Hom}}_{\mathcal{P}(X)}(U,f_*(V)),$

natural in $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$ and (respectively) $V \in \mathcal{P}(X)$ and $U \in \mathcal{P}(Y)$. In particular:

i. The following conditions are equivalent:

- A. We have $f_1(U) \subset V$.
- B. We have $U \subset f^{-1}(V)$.
- ii. The following conditions are equivalent:
 - A. We have $f^{-1}(U) \subset V$.
 - B. We have $U \subset f_*(V)$.
- 3. Interaction With Unions of Families of Subsets. The diagram

$$\mathcal{P}(\mathcal{P}(X)) \xrightarrow{(f_*)_*} \mathcal{P}(\mathcal{P}(Y))$$

$$\bigcup \qquad \qquad \bigcup \qquad \qquad \bigcup \bigcup$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$\bigcup_{U \in \mathcal{U}} f_*(U) = \bigcup_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

4. Interaction With Intersections of Families of Subsets. The diagram

$$\begin{array}{ccc}
\mathcal{P}(\mathcal{P}(X)) & \xrightarrow{(f_*)_*} & \mathcal{P}(\mathcal{P}(Y)) \\
& & & & \downarrow \\
& & & \downarrow \\
\mathcal{P}(X) & \xrightarrow{f_*} & \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$\bigcap_{U \in \mathcal{U}} f_*(U) = \bigcap_{V \in f_*(\mathcal{U})} V$$

for each $\mathcal{U} \in \mathcal{P}(X)$, where $f_*(\mathcal{U}) \stackrel{\text{def}}{=} (f_*)_*(\mathcal{U})$.

we see that $f_{*,im}$ corresponds to meets indexed over nonempty sets, while $f_{*,cp}$ corresponds to meets indexed over the empty set.

5. Interaction With Binary Unions. Let $f\colon X\to Y$ be a function. We have a natural transformation

with components

$$f_*(U) \cup f_*(V) \subset f_*(U \cup V)$$

indexed by $U, V \in \mathcal{P}(X)$.

6. Interaction With Binary Intersections. The diagram

$$\mathcal{P}(X) \times \mathcal{P}(X) \xrightarrow{f_* \times f_*} \mathcal{P}(Y) \times \mathcal{P}(Y)$$

$$\uparrow \qquad \qquad \qquad \downarrow \cap$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U) \cap f_*(V) = f_*(U \cap V)$$

for each $U, V \in \mathcal{P}(X)$.

7. Interaction With Complements. The diagram

$$\mathcal{P}(X)^{\mathsf{op}} \xrightarrow{f_!^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
\xrightarrow{(-)^{\mathsf{c}}} \qquad \qquad \downarrow^{(-)^{\mathsf{c}}} \\
\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

commutes, i.e. we have

$$f_*(U^{\mathsf{c}}) = f_!(U)^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$.

8. Interaction With Symmetric Differences. We have a natural transformation

$$\mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \xrightarrow{f_*^{\mathsf{op}} \times f_*} \mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y)$$

$$\downarrow^{\triangle} \qquad \qquad \downarrow^{\triangle}$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

with components

$$f_*(U \triangle V) \subset f_*(U) \triangle f_*(V)$$

indexed by $U, V \in \mathcal{P}(X)$.

9. Interaction With Internal Homs of Powersets. We have a natural transformation

with components

$$[f_!(U), f_*(V)]_Y \subset f_*([U, V]_X)$$

indexed by $U, V \in \mathcal{P}(X)$.

10. Lax Preservation of Colimits. We have an inclusion of sets

$$\bigcup_{i\in I} f_*(U_i) \subset f_*\left(\bigcup_{i\in I} U_i\right),\,$$

natural in $\{U_i\}_{i\in I}\in \mathcal{P}(X)^{ imes I}$. In particular, we have inclusions

$$f_*(U) \cup f_*(V) \hookrightarrow f_*(U \cup V),$$

 $\emptyset \hookrightarrow f_*(\emptyset),$

natural in $U, V \in \mathcal{P}(X)$.

11. Preservation of Limits. We have an equality of sets

$$f_*\left(\bigcap_{i\in I}U_i\right)=\bigcap_{i\in I}f_*(U_i),$$

natural in $\{U_i\}_{i\in I}\in\mathcal{P}(X)^{\times I}$. In particular, we have equalities

$$f^{-1}(U \cap V) = f_*(U) \cap f^{-1}(V),$$

$$f_*(X) = Y,$$

natural in $U, V \in \mathcal{P}(X)$.

12. Symmetric Lax Monoidality With Respect to Unions. The codirect image function of Item1 has a symmetric lax monoidal structure

$$(f_*, f_*^{\otimes}, f_{*|1}^{\otimes}) \colon (\mathcal{P}(X), \cup, \emptyset) \to (\mathcal{P}(Y), \cup, \emptyset),$$

being equipped with inclusions

$$\begin{split} f^{\otimes}_{*|U,V} \colon f_{*}(U) \cup f_{*}(V) & \hookrightarrow f_{*}(U \cup V), \\ f^{\otimes}_{*|\mathfrak{I}} \colon \emptyset & \hookrightarrow f_{*}(\emptyset), \end{split}$$

natural in $U, V \in \mathcal{P}(X)$.

13. Symmetric Strict Monoidality With Respect to Intersections. The direct image function of Item1 has a symmetric strict monoidal structure

$$(f_*, f_*^{\otimes}, f_{*|\mathbb{1}}^{\otimes}) \colon (\mathcal{P}(X), \cap, X) \to (\mathcal{P}(Y), \cap, Y),$$

being equipped with equalities

$$f_{*|U,V}^{\otimes} \colon f_{*}(U \cap V) \xrightarrow{=} f_{*}(U) \cap f_{*}(V),$$
$$f_{*|\mathbb{1}}^{\otimes} \colon f_{*}(X) \xrightarrow{=} Y,$$

natural in $U, V \in \mathcal{P}(X)$.

14. Interaction With Coproducts. Let $f\colon X\to X'$ and $g\colon Y\to Y'$ be maps of sets. We have

$$(f \coprod g)_*(U \coprod V) = f_*(U) \coprod g_*(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

15. Interaction With Products. Let $f\colon X\to X'$ and $g\colon Y\to Y'$ be maps of sets. We have

$$(f \boxtimes_{X \times Y} q)_*(U \boxtimes_{X \times Y} V) = f_*(U) \boxtimes_{X' \times Y'} q_*(V)$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

16. Relation to Direct Images. We have

$$f_*(U) = f_!(U^{c})^{c}$$
$$= Y \setminus f_!(X \setminus U)$$

for each $U \in \mathcal{P}(X)$.

17. Interaction With Injections. If f is injective, then we have

$$f_{*,\text{im}}(U) = f_!(U),$$

 $f_{*,\text{cp}}(U) = Y \setminus \text{Im}(f),$

and so

$$f_*(U) = f_{*,\text{im}}(U) \cup f_{*,\text{cp}}(U)$$
$$= f_!(U) \cup (Y \setminus \text{Im}(f))$$

for each $U \in \mathcal{P}(X)$.

18. Interaction With Surjections. If f is surjective, then we have

$$\begin{split} f_{*,\mathrm{im}}(U) &\subset f_!(U), \\ f_{*,\mathrm{cp}}(U) &= \emptyset, \end{split}$$

and so

$$f_*(U) \subset f_!(U)$$

for each $U \in \mathcal{P}(X)$.

Proof. Item 1, Functoriality: Omitted.

Item 2, *Triple Adjointness*: This follows from Definition 4.6.1.1.4, Definition 4.6.2.1.2, Definition 4.6.3.1.4, and Kan Extensions, ?? of ??.

Item 3, Interaction With Unions of Families of Subsets: We have

$$\bigcup_{V \in f_*(\mathcal{U})} V = \bigcup_{V \in \{f_*(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcup_{U \in \mathcal{U}} f_*(U).$$

This finishes the proof.

Item 4, Interaction With Intersections of Families of Subsets: We have

$$\bigcap_{V \in f_*(\mathcal{U})} V = \bigcap_{V \in \{f_*(U) \in \mathcal{P}(X) \mid U \in \mathcal{U}\}} V$$
$$= \bigcap_{U \in \mathcal{U}} f_*(U).$$

This finishes the proof.

Item 5, Interaction With Binary Unions: We have

$$f_*(U) \cup f_*(V) = f_!(U^c)^c \cup f_!(V^c)^c$$

$$= (f_!(U^c) \cap f_!(V^c))^c$$

$$\subset (f_!(U^c \cap V^c))^c$$

$$= f_!((U \cup V)^c)^c$$

$$= f_*(U \cup V),$$

where:

- 1. We have used Item 16 for the first equality.
- 2. We have used Item 2 of Definition 4.3.11.1.2 for the second equality.
- 3. We have used Item 6 of Definition 4.6.1.1.5 for the third equality.
- 4. We have used Item 2 of Definition 4.3.11.1.2 for the fourth equality.
- 5. We have used Item 16 for the last equality.

This finishes the proof.

Item 6, *Interaction With Binary Intersections*: This follows from *Item 11*.

Item 7, Interaction With Complements: Omitted.

Item 8, Interaction With Symmetric Differences: Omitted.

Item 9, Interaction With Internal Homs of Powersets: We have

$$\begin{split} \left[f_!(U), f^!(V) \right]_X &\stackrel{\text{def}}{=} f_!(U)^{\mathsf{c}} \cup f_*(V) \\ &= f_*(U^{\mathsf{c}}) \cup f_*(V) \\ &\subset f_*(U^{\mathsf{c}} \cup V) \\ &\stackrel{\text{def}}{=} f_*([U, V]_X), \end{split}$$

where we have used:

- 1. Item 7 of Definition 4.6.3.1.7 for the second equality.
- 2. Item 5 of Definition 4.6.3.1.7 for the inclusion.

Since $\mathcal{P}(X)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). This finishes the proof.

Item 10, Lax Preservation of Colimits: Omitted.

Item 11, Preservation of Limits: This follows from Item 2 and ??, ?? of ??.

Item 12, Symmetric Lax Monoidality With Respect to Unions: This follows from Item 10.

Item 13, *Symmetric Strict Monoidality With Respect to Intersections*: This follows from Item 11.

Item 14, Interaction With Coproducts: Omitted.

Item 15, Interaction With Products: Omitted.

Item 16, *Relation to Direct Images*: We claim that $f_*(U) = Y \setminus f_!(X \setminus U)$.

· The First Implication. We claim that

$$f_*(U) \subset Y \setminus f_!(X \setminus U).$$

Let $y \in f_*(U)$. We need to show that $y \notin f_!(X \setminus U)$, i.e. that there is no $x \in X \setminus U$ such that f(x) = y.

This is indeed the case, as otherwise we would have $x \in f^{-1}(y)$ and $x \notin U$, contradicting $f^{-1}(y) \subset U$ (which holds since $y \in f_*(U)$).

Thus $y \in Y \setminus f_!(X \setminus U)$.

· The Second Implication. We claim that

$$Y \setminus f_!(X \setminus U) \subset f_*(U)$$
.

Let $y \in Y \setminus f_!(X \setminus U)$. We need to show that $y \in f_*(U)$, i.e. that $f^{-1}(y) \subset U$.

Since $y \notin f_!(X \setminus U)$, there exists no $x \in X \setminus U$ such that y = f(x), and hence $f^{-1}(y) \subset U$.

Thus $y \in f_*(U)$.

This finishes the proof of Item 16.

Item 17, Interaction With Injections: Omitted.

Item 18, Interaction With Surjections: Omitted.

Proposition 4.6.3.1.8. Let $f: X \to B$ be a function.

1. Functionality I. The assignment $f \mapsto f_*$ defines a function

$$(-)_{!|X,Y} : \mathsf{Sets}(X,Y) \to \mathsf{Sets}(\mathcal{P}(X),\mathcal{P}(Y)).$$

2. Functionality II. The assignment $f \mapsto f_*$ defines a function

$$(-)_{!|X,Y}: \mathsf{Sets}(X,Y) \to \mathsf{Pos}((\mathcal{P}(X),\subset),(\mathcal{P}(Y),\subset)).$$

3. Interaction With Identities. For each $X \in \mathsf{Obj}(\mathsf{Sets})$, we have

$$(id_X)_* = id_{\mathcal{P}(X)}$$
.

4. *Interaction With Composition*. For each pair of composable functions $f: X \to Y$ and $g: Y \to Z$, we have

$$(g \circ f)_* = g_* \circ f_*,$$

$$\mathcal{P}(X) \xrightarrow{f_*} \mathcal{P}(Y)$$

$$\downarrow^{g_*}$$

$$\mathcal{P}(Z).$$

Proof. Item 1, *Functionality I*: There is nothing to prove.

Item 2, Functionality II: This follows from Item 1 of Definition 4.6.3.1.7.

Item 3, *Interaction With Identities*: This follows from Definition 4.6.3.1.4 and Kan Extensions, ?? of ??.

Item 4, Interaction With Composition: This follows from Definition 4.6.3.1.4 and Kan Extensions, **??** of **??**. □

4.6.4 A Six-Functor Formalism for Sets

Remark 4.6.4.1.1. The assignment $X \mapsto \mathcal{P}(X)$ together with the functors f_* , f^{-1} , and $f_!$ of Item 1 of Definition 4.6.1.1.5, Item 1 of Definition 4.6.2.1.3, and Item 1 of Definition 4.6.3.1.7, and the functors

$$-_1 \cap -_2 \colon \mathcal{P}(X) \times \mathcal{P}(X) \to \mathcal{P}(X),$$

$$[-1,-2]_X \colon \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X) \to \mathcal{P}(X)$$

of Item 1 of Definition 4.3.9.1.2 and Item 1 of Definition 4.4.7.1.3 satisfy several properties reminiscent of a six functor formalism in the sense of ??.

We collect these properties in Definition 4.6.4.1.2 below.⁴³

Proposition 4.6.4.1.2. Let *X* be a set.

1. The Beck-Chevalley Condition. Let

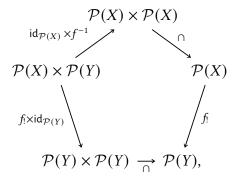
$$X \times_{Z} Y \xrightarrow{\operatorname{pr}_{2}} Y$$

$$\downarrow g$$

$$X \xrightarrow{f} Z$$

be a pullback diagram in Sets. We have

2. The Projection Formula I. The diagram



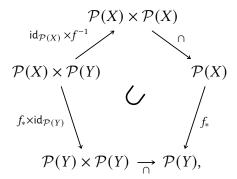
commutes, i.e. we have

$$f_!(U \cap f^{-1}(V)) = f_!(U) \cap V$$

for each $U \in \mathcal{P}(X)$ and each $V \in \mathcal{P}(Y)$.

⁴³See also [nLa25].

3. The Projection Formula II. We have a natural transformation



with components

$$f_*(U) \cap V \subset f_*(U \cap f^{-1}(V))$$

indexed by $U \in \mathcal{P}(X)$ and $V \in \mathcal{P}(Y)$.

4. Strong Closed Monoidality. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \times \mathcal{P}(Y) \xrightarrow{f^{-1,\mathsf{op}} \times f^{-1}} \mathcal{P}(X)^{\mathsf{op}} \times \mathcal{P}(X)$$

$$\downarrow [-1,-2]_{Y} \qquad \qquad \downarrow [-1,-2]_{X}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have an equality of sets

$$f^{-1}([U,V]_Y) = [f^{-1}(U), f^{-1}(V)]_X,$$

natural in $U, V \in \mathcal{P}(X)$.

5. The External Tensor Product. We have an external tensor product

$$-_1 \boxtimes_{X \times Y} -_2 : \mathcal{P}(X) \times \mathcal{P}(Y) \to \mathcal{P}(X \times Y)$$

given by

$$U \boxtimes_{X \times Y} V \stackrel{\text{def}}{=} \operatorname{pr}_{1}^{-1}(U) \cap \operatorname{pr}_{2}^{-1}(V)$$
$$= \{(u, v) \in X \times Y \mid u \in U \text{ and } v \in V\}.$$

This is the same map as the one in Item 5 of Definition 4.4.1.1.4. Moreover, the following conditions are satisfied:

(a) Interaction With Direct Images. Let $f\colon X\to X'$ and $g\colon Y\to Y'$ be functions. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X) \times \mathcal{P}(Y) & \xrightarrow{f_! \times g_!} & \mathcal{P}(X') \times \mathcal{P}(Y') \\
\boxtimes_{X \times Y} & & & & & & & \\
\mathbb{Z}_{X \times Y} & & & & & & \\
\mathcal{P}(X \times Y) & \xrightarrow{f_! \times g_!} & \mathcal{P}(X' \times Y')
\end{array}$$

commutes, i.e. we have

$$[f_! \times g_!](U \boxtimes_{X \times Y} V) = f_!(U) \boxtimes_{X' \times Y'} g_!(V)$$

for each $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$.

(b) Interaction With Inverse Images. Let $f\colon X\to X'$ and $g\colon Y\to Y'$ be functions. The diagram

commutes, i.e. we have

$$[f^{-1}\times g^{-1}](U\boxtimes_{X'\times Y'}V)=f^{-1}(U)\boxtimes_{X\times Y}g^{-1}(V)$$

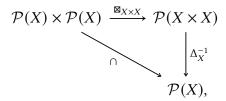
for each $(U, V) \in \mathcal{P}(X') \times \mathcal{P}(Y')$.

(c) Interaction With Codirect Images. Let $f: X \to X'$ and $g: Y \to Y'$ be functions. The diagram

commutes, i.e. we have

$$[f_* \times g_*](U \boxtimes_{X \times Y} V) = f_*(U) \boxtimes_{X' \times Y'} g_*(V)$$
 for each $(U, V) \in \mathcal{P}(X) \times \mathcal{P}(Y)$.

(d) Interaction With Diagonals. The diagram



i.e. we have

$$U \cap V = \Delta_X^{-1}(U \boxtimes_{X \times X} V)$$

for each $U, V \in \mathcal{P}(X)$.

6. The Dualisation Functor. We have a functor

$$D_X \colon \mathcal{P}(X)^{\mathsf{op}} \to \mathcal{P}(X)$$

given by

$$D_X(U) \stackrel{\text{def}}{=} [U, \emptyset]_X$$
$$\stackrel{\text{def}}{=} U^{\mathsf{c}}$$

for each $U \in \mathcal{P}(X)$, as in Item 5 of Definition 4.4.7.1.3, satisfying the following conditions:

(a) Duality. We have

$$D_X(D_X(U)) = U, \qquad D_X \xrightarrow{\operatorname{id}_{\mathcal{P}(X)}} \mathcal{P}(X)$$

$$\mathcal{P}(X) \xrightarrow{D_X} \mathcal{P}(X)$$

$$\mathcal{P}(X).$$

(b) Duality. The diagram

$$\mathcal{P}(X)^{\mathrm{op}} \times \mathcal{P}(X)^{\mathrm{op}} \stackrel{\cap^{\mathrm{op}}}{\to} \mathcal{P}(X)^{\mathrm{op}}$$

$$id_{\mathcal{P}(X)^{\mathrm{op}}} \times D_{X} / D_{X}$$

$$\mathcal{P}(X)^{\mathrm{op}} \times \mathcal{P}(X) \xrightarrow{[-1,-2]_{X}} \mathcal{P}(X)$$

commutes, i.e. we have

$$\underbrace{D_X(U\cap D_X(V))}_{\stackrel{\text{def}}{=}[U\cap [V,\emptyset]_X,\emptyset]_X}=[U,V]_X$$

for each $U, V \in \mathcal{P}(X)$.

(c) Interaction With Direct Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_*^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
\downarrow^{D_X} & & \downarrow^{D_Y} \\
\mathcal{P}(X) & \xrightarrow{f_*} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_!(D_X(U)) = D_Y(f_*(U))$$

for each $U \in \mathcal{P}(X)$.

(d) Interaction With Inverse Images. The diagram

$$\mathcal{P}(Y)^{\mathsf{op}} \xrightarrow{f^{-1,\mathsf{op}}} \mathcal{P}(X)^{\mathsf{op}}$$

$$D_{Y} \downarrow \qquad \qquad \downarrow D_{X}$$

$$\mathcal{P}(Y) \xrightarrow{f^{-1}} \mathcal{P}(X)$$

commutes, i.e. we have

$$f^{-1}(D_Y(U)) = D_X(f^{-1}(U))$$

for each $U \in \mathcal{P}(X)$.

(e) Interaction With Codirect Images. The diagram

$$\begin{array}{ccc}
\mathcal{P}(X)^{\mathsf{op}} & \xrightarrow{f_!^{\mathsf{op}}} \mathcal{P}(Y)^{\mathsf{op}} \\
\downarrow^{D_X} & & \downarrow^{D_Y} \\
\mathcal{P}(X) & \xrightarrow{f_*} \mathcal{P}(Y)
\end{array}$$

commutes, i.e. we have

$$f_*(D_X(U)) = D_Y(f_!(U))$$

for each $U \in \mathcal{P}(X)$.

Proof. Item 1, The Beck-Chevalley Condition: We have

$$[g^{-1} \circ f_!](U) \stackrel{\text{def}}{=} g^{-1}(f_!(U))$$

$$\stackrel{\text{def}}{=} \{y \in Y \mid g(y) \in f_!(U)\}$$

$$= \left\{y \in Y \mid \text{there exists some } x \in U \right\}$$

$$\text{such that } f(x) = g(y)$$

$$= \left\{y \in Y \mid \text{there exists some} \\ (x, y) \in \{(x, y) \in X \times_Z Y \mid x \in U\}\right\}$$

$$= \left\{y \in Y \mid \text{there exists some} \\ (x, y) \in \{(x, y) \in X \times_Z Y \mid x \in U\}\right\}$$

$$\text{such that } y = y$$

$$= \left\{y \in Y \mid \text{there exists some} \\ (x, y) \in \{(x, y) \in X \times_Z Y \mid x \in U\}\right\}$$

$$\text{such that pr}_2(x, y) = y$$

$$\stackrel{\text{def}}{=} (\text{pr}_2)_!(\{(x, y) \in X \times_Z Y \mid \text{pr}_1(x, y) \in U\})$$

$$\stackrel{\text{def}}{=} (\text{pr}_2)_!(\text{pr}_1^{-1}(U))$$

$$\stackrel{\text{def}}{=} [(\text{pr}_2)_! \circ \text{pr}_1^{-1}](U)$$

for each $U \in \mathcal{P}(X)$. Therefore, we have

$$g^{-1} \circ f_! = (\mathsf{pr}_2)_! \circ \mathsf{pr}_1^{-1}$$
.

For the second equality, we have

$$[f^{-1} \circ g_!](U) \stackrel{\text{def}}{=} f^{-1}(g_!(U))$$

$$\stackrel{\text{def}}{=} \{x \in X \mid f(x) \in g_!(V)\}$$

$$= \left\{x \in X \middle| \text{ there exists some } y \in V \right\}$$
such that $f(x) = g(y)$

$$= \begin{cases} x \in X & \text{there exists some} \\ (x,y) \in \{(x,y) \in X \times_Z Y \mid y \in V\} \end{cases}$$

$$= \begin{cases} x \in X & \text{there exists some} \\ (x,y) \in \{(x,y) \in X \times_Z Y \mid y \in V\} \end{cases}$$

$$= \begin{cases} x \in X & \text{there exists some} \\ (x,y) \in \{(x,y) \in X \times_Z Y \mid y \in V\} \end{cases}$$

$$= \begin{cases} x \in X & \text{there exists some} \\ (x,y) \in \{(x,y) \in X \times_Z Y \mid y \in V\} \end{cases}$$

$$= (pr_1)!(\{(x,y) \in X \times_Z Y \mid y \in V\})$$

$$= (pr_1)!(\{(x,y) \in X \times_Z Y \mid pr_2(x,y) \in V\})$$

$$\stackrel{\text{def}}{=} (pr_1)!(pr_2^{-1}(V))$$

$$\stackrel{\text{def}}{=} (pr_1)!(pr_2^{-1}(V))$$

$$\stackrel{\text{def}}{=} (pr_1)!(pr_2^{-1}(V))$$

for each $V \in \mathcal{P}(Y)$. Therefore, we have

$$f^{-1} \circ g_! = (\mathsf{pr}_1)_! \circ \mathsf{pr}_2^{-1}$$
.

This finishes the proof.

Item 2, The Projection Formula I: We claim that

$$f_!(U) \cap V \subset f_!(U \cap f^{-1}(V)).$$

Indeed, we have

$$f_!(U) \cap V \subset f_!(U) \cap f_!(f^{-1}(V))$$
$$= f_!(U \cap f^{-1}(V)),$$

where we have used:

- 1. Item 2 of Definition 4.6.1.1.5 for the inclusion.
- 2. Item 6 of Definition 4.6.1.1.5 for the equality.

Conversely, we claim that

$$f_!(U \cap f^{-1}(V)) \subset f_!(U) \cap V.$$

Indeed:

- 1. Let $y \in f_!(U \cap f^{-1}(V))$.
- 2. Since $y \in f_!(U \cap f^{-1}(V))$, there exists some $x \in U \cap f^{-1}(V)$ such that f(x) = y.
- 3. Since $x \in U \cap f^{-1}(V)$, we have $x \in U$, and thus $f(x) \in f_!(U)$.
- 4. Since $x \in U \cap f^{-1}(V)$, we have $x \in f^{-1}(V)$, and thus $f(x) \in V$.
- 5. Since $f(x) \in f_!(U)$ and $f(x) \in V$, we have $f(x) \in f_!(U) \cap V$.
- 6. But y = f(x), so $y \in f(U) \cap V$.
- 7. Thus $f_!(U \cap f^{-1}(V)) \subset f_!(U) \cap V$.

This finishes the proof.

Item 3, The Projection Formula II: We have

$$f_*(U) \cap V \subset f_*(U) \cap f_*(f^{-1}(V))$$

= $f_*(U \cap f^{-1}(V))$,

where we have used:

- 1. Item 2 of Definition 4.6.3.1.7 for the inclusion.
- 2. Item 6 of Definition 4.6.3.1.7 for the equality.

Since $\mathcal{P}(Y)$ is posetal, naturality is automatic (Categories, Item 4 of Definition 11.2.7.1.2). *Item 4, Strong Closed Monoidality*: This is a repetition of Item 19 of Definition 4.4.7.1.3 and is proved there.

Item 5, The External Tensor Product: We have

$$\begin{split} U \boxtimes_{X \times Y} V &\stackrel{\text{def}}{=} \operatorname{pr}_{1}^{-1}(U) \cap \operatorname{pr}_{2}^{-1}(V) \\ &\stackrel{\text{def}}{=} \{(x,y) \in X \times Y \mid \operatorname{pr}_{1}(x,y) \in U\} \\ & \cup \{(x,y) \in X \times Y \mid \operatorname{pr}_{2}(x,y) \in V\} \\ &= \{(x,y) \in X \times Y \mid x \in U\} \\ & \cup \{(x,y) \in X \times Y \mid y \in V\} \\ &= \{(x,y) \in X \times Y \mid x \in U \text{ and } y \in V\} \\ &\stackrel{\text{def}}{=} U \times V. \end{split}$$

Next, we claim that Items 5a to 5d are indeed true:

- 1. *Proof of Item 5a*: This is a repetition of Item 16 of Definition 4.6.1.1.5 and is proved there.
- 2. *Proof of Item 5b*: This is a repetition of Item 16 of Definition 4.6.2.1.3 and is proved there.
- 3. *Proof of Item 5c*: This is a repetition of Item 15 of Definition 4.6.3.1.7 and is proved there.
- 4. Proof of Item 5d: We have

$$\begin{split} \Delta_X^{-1}(U\boxtimes_{X\times X}V) &\stackrel{\text{def}}{=} \{x\in X \mid (x,x)\in U\boxtimes_{X\times X}V\}\\ &= \{x\in X \mid (x,x)\in \{(u,v)\in X\times X \mid u\in U \text{ and } v\in V\}\}\\ &= U\cap V. \end{split}$$

This finishes the proof.

Item 6, *The Dualisation Functor*: This is a repetition of Items 5 and 6 of Definition 4.4.7.1.3 and is proved there.

4.7 Miscellany

4.7.1 Injective Functions

Let *A* and *B* be sets.

Definition 4.7.1.1.1. A function $f: A \to B$ is **injective** if it satisfies the following condition:

 (\star) For each $a, a' \in A$, if f(a) = f(a'), then a = a'.

Proposition 4.7.1.1.2. Let $f: A \rightarrow B$ be a function.

1. Characterisations. The following conditions are equivalent: 44

⁴⁴Items 1c to 1f unwind respectively to the following statements:

[·] For each $U, V \in \mathcal{P}(A)$, if $f_!(U) = f_!(V)$, then U = V.

[·] For each $U, V \in \mathcal{P}(A)$, if $f_*(U) = f_*(V)$, then U = V.

[·] For each $U, V \in \mathcal{P}(A)$, if $f_!(U) \subset f_!(V)$, then $U \subset V$.

[·] For each $U, V \in \mathcal{P}(A)$, if $f_*(U) \subset f_*(V)$, then $U \subset V$.

- (a) The function *f* is injective.
- (b) The function f is a monomorphism in Sets.
- (c) The direct image function

$$f_! \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

associated to f is injective.

(d) The codirect image function

$$f_* \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

associated to f is injective.

(e) The direct image functor

$$f_! \colon (\mathcal{P}(A), \subset) \to (\mathcal{P}(B), \subset)$$

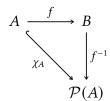
associated to f is full.

(f) The codirect image function

$$f_* \colon \mathcal{P}(A) \to \mathcal{P}(B)$$

associated to f is full.

(g) The diagram

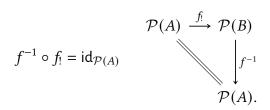


commutes. That is, we have

$$f^{-1}(f(a)) = \{a\}$$

for each $a \in A$.

(h) We have

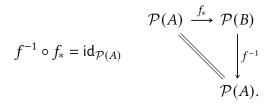


In other words, we have

$${a \in A \mid f(a) \in f(U)} = U$$

for each $U \in \mathcal{P}(A)$.

(i) We have



In other words, we have

$$\left\{a \in A \,\middle|\, f^{-1}(f(a)) \subset U\right\} = U$$

for each $U \in \mathcal{P}(A)$.

Proof. Item 1, *Characterisations*: We will proceed by showing:

- · Step 1: Item 1a \iff Item 1b.
- · Step 2: Item 1a \iff Item 1c.
- · Step 3: Item 1a \iff Item 1d.
- · Step 4: Item 1c \iff Item 1e.
- · Step 5: Item 1e \iff Item 1f.
- · Step 6: Item 1a \iff Item 1g.
- · Step 7: Item 1g \iff Item 1h.
- · Step 8: Item 1a \iff Item 1i.

Step 1: Item 1a \iff **Item 1b.** We claim that Items 1a and 1b are equivalent:

- · $ltem 1a \implies ltem 1b$: We proceed in a few steps:
 - Proceeding by contrapositive, we claim that given a pair of maps $g, h \colon C \Rightarrow A$ such that $g \neq h$, we have $f \circ g \neq f \circ h$.

- Indeed, as g and h are different maps, there must exist at least one element $x \in C$ such that $g(x) \neq h(x)$.
- But then we have $f(g(x)) \neq f(h(x))$, as f is injective.
- Thus $f \circ q \neq f \circ h$, and we are done.
- · $ltem 1b \implies ltem 1a$: We proceed in a few steps:
 - Consider the diagram

$$\mathsf{pt} \xrightarrow{[y]}^{[x]} A \xrightarrow{f} B,$$

where [x] and [y] are the morphisms picking the elements x and y of A.

- Note that we have f(x) = f(y) iff $f \circ [x] = f \circ [y]$.
- Since f is assumed to be a monomorphism, if f(x) = f(y), then $f \circ [x] = f \circ [y]$ and therefore [x] = [y].
- This shows that if f(x) = f(y), then x = y, so f is injective.

Step 2: Item 1a \iff Item 1c. We claim that Items 1a and 1c are indeed equivalent:

- · Item 1a \Longrightarrow Item 1c: We proceed in a few steps:
 - Assume that f is injective and let $U, V \in \mathcal{P}(A)$ such that $f_!(U) = f_!(V)$. We wish to show that U = V.
 - To show that $U \subset V$, let $u \in U$.
 - By the definition of the direct image, we have $f(u) \in f(U)$.
 - Since $f_!(U) = f_!(V)$, it follows that $f(u) \in f_!(V)$.
 - Thus, there exists some $v \in V$ such that f(v) = f(u).
 - Since f is injective, the equality f(v) = f(u) implies that v = u.
 - Thus $u \in V$ and $U \subset V$.
 - A symmetric argument shows that $V \subset U$.
 - Therefore U = V, showing $f_!$ to be injective.
- · $ltem1c \implies ltem1a$: We proceed in a few steps:

- Assume that the direct image function $f_!$ is injective and let $a, a' \in A$ such that f(a) = f(a'). We wish to show that a = a'.
- Since

$$f_!(\{a\}) = \{f(a)\}\$$

$$= \{f(a')\}\$$

$$= f_!(\{a'\}),\$$

we must have $\{a\} = \{a'\}$, as $f_!$ is injective, so a = a', showing f to be injective.

Step 3: Item 1c \iff Item 1d. This follows from Item 17 of Definition 4.6.1.1.5. Step 4: Item 1c \iff Item 1e. We claim that Items 1c and 1e are equivalent:

- · $ltem1c \implies ltem1e$: We proceed in a few steps:
 - Let $U, V \in \mathcal{P}(A)$ such that $f_!(U) \subset f_!(V)$, assume $f_!$ to be injective, and consider the set $U \cup V$.
 - Since $f_!(U) \subset f_!(V)$, we have

$$f_!(U \cup V) = f_!(U) \cup f_!(V)$$
$$= f_!(V),$$

where we have used Item 5 of Definition 4.6.1.1.5 for the first equality.

- Since $f_!$ is injective, this implies $U \cup V = V$.
- Thus $U \subset V$, as we wished to show.
- · $ltem 1c \implies ltem 1e$: We proceed in a few steps:
 - Suppose Item 1e holds, and let $U, V \in \mathcal{P}(A)$ such that $f_!(U) = f_!(V)$.
 - Since $f_!(U) = f_!(V)$, we have $f_!(U) \subset f_!(V)$ and $f_!(V) \subset f_!(U)$.
 - By assumption, this implies $U \subset V$ and $V \subset U$.
 - Thus U = V, showing $f_!$ to be injective.

Step 5: Item 1e \iff **Item 1f.** This follows from Item 17 of Definition 4.6.1.1.5.

Step 6: Item 1a \iff Item 1g. We have

$$f^{-1}(f(a)) = \{ a' \in A \mid f(a') = f(a) \}$$

so the condition $f^{-1}(f(a)) = \{a\}$ states precisely that if f(a') = f(a), then a' = a. **Step 7: Item 1g** \iff **Item 1h.** We claim that **Items 1g** and **1h** are indeed equivalent:

· Item 1 $g \Longrightarrow Item 1h$: We have

$$[f^{-1} \circ f_!](U) \stackrel{\text{def}}{=} f^{-1}(f_!(U))$$

$$= f^{-1} \left(\int_{u \in U} \{u\} \right)$$

$$= f^{-1} \left(\bigcup_{u \in U} f_!(\{u\}) \right)$$

$$= \bigcup_{u \in U} f^{-1}(f_!(\{u\}))$$

$$= \bigcup_{u \in U} f^{-1}(f_!(u))$$

$$= \bigcup_{u \in U} \{u\}$$

$$= U$$

for each $U \in \mathcal{P}(A)$, where we have used Item 5 of Definition 4.6.1.1.5 for the third equality and Item 5 of Definition 4.6.2.1.3 for the fourth equality.

· Item 1h \Longrightarrow Item 1g: Applying the condition $f^{-1} \circ f_! = \mathrm{id}_{\mathcal{P}(A)}$ to $U = \{a\}$ gives

$$f^{-1}(f_!(\{a\})) = \{a\}.$$

Step 8: Item 1a \iff Item 1i. We claim that Items 1a and 1i are equivalent:

· Item 1a \Longrightarrow Item 1i: If f is injective, then $f^{-1}(f(a)) = \{a\}$, so we have

$$f^{-1}(f_*(a)) = \{ a \in A \mid \{a\} \subset U \}$$

= U.

· Item 1i \Longrightarrow Item 1a: For $U=\{a\}$, the condition $f^{-1}(f_*(U))=U$ becomes

$$\left\{ a' \in A \,\middle|\, f^{-1}(f(a')) \subset \{a\} \right\} = \{a\}.$$

Since the set $f^{-1}(f(a'))$ is given by

$${a \in A \mid f(a) = f(a')},$$

it follows that f is injective.

This finishes the proof.

4.7.2 Surjective Functions

Let A and B be sets.

Definition 4.7.2.1.1. A function $f: A \rightarrow B$ is **surjective** if it satisfies the following condition:

 (\star) For each $b \in B$, there exists some $a \in A$ such that f(a) = b.

Proposition 4.7.2.1.2. Let $f: A \rightarrow B$ be a function.

- 1. Characterisations. The following conditions are equivalent:
 - (a) The function *f* is surjective.
 - (b) The function f is an epimorphism in Sets.
 - (c) The inverse image function

$$f^{-1} \colon \mathcal{P}(B) \to \mathcal{P}(A)$$

associated to f is injective.

(d) The inverse image functor

$$f^{-1} \colon (\mathcal{P}(B), \subset) \to (\mathcal{P}(A), \subset)$$

associated to f is full.

(e) The diagram

$$B \xrightarrow{f^{-1}} \mathcal{P}(A)$$

$$\downarrow_{f}$$

$$\mathcal{P}(B)$$

commutes. That is, we have

$$f_!(f^{-1}(b)) = \{b\}$$

for each $b \in B$.

(f) We have

$$f_! \circ f^{-1} = \mathrm{id}_{\mathcal{P}(B)} \qquad \mathcal{P}(A)$$

$$f_! \circ f^{-1} = \mathrm{id}_{\mathcal{P}(B)} \qquad f_!$$

In other words, we have

$$\left\{b \in B \middle| \begin{array}{l} \text{there exists some } a \in f^{-1}(U) \\ \text{such that } f(a) = b \end{array} \right\} = U$$

for each $U \in \mathcal{P}(A)$.

(g) We have

$$f_* \circ f^{-1} = \mathrm{id}_{\mathcal{P}(B)} \qquad \qquad \int_{f_*}^{f^{-1}} \mathcal{P}(A) \\ \mathcal{P}(B).$$

In other words, we have

$$\left\{b\in B\left|f^{-1}(b)\subset f^{-1}(U)\right\}=U$$

for each $U \in \mathcal{P}(B)$.

Proof. Item 1, *Characterisations*: We will proceed by showing:

- · Step 1: Item 1a \iff Item 1b.
- · Step 2: Item 1a \iff Item 1c.
- · Step 3: Item 1c \iff Item 1d.
- · Step 4: Item 1a \iff Item 1e.
- · Step 5: Item 1e \iff Item 1f.
- · Step 6: Item 1a \iff Item 1g.

Step 1: Item 1a ← Item 1b. We claim Items 1a and 1b are indeed equivalent:

- · $ltem 1a \implies ltem 1b$: We proceed in a few steps:
 - Let $g, h: B \Rightarrow C$ be morphisms such that $g \circ f = h \circ f$.
 - For each $a \in A$, we have

$$g(f(a)) = h(f(a)).$$

- However, this implies that

$$g(b) = h(b)$$

for each $b \in B$, as f is surjective.

- Thus q = h and f is an epimorphism.
- · $ltem 1b \implies ltem 1a$: We proceed by contrapositive. Consider the diagram

$$A \xrightarrow{f} B \xrightarrow{g} C$$
,

where h is the map defined by h(b) = 0 for each $b \in B$ and g is the map defined by

$$g(b) = \begin{cases} 1 & \text{if } b \in \text{Im}(f), \\ 0 & \text{otherwise.} \end{cases}$$

Then $h\circ f=g\circ f$, as h(f(a))=1=g(f(a)) for each $a\in A$. However, for any $b\in B\setminus {\rm Im}(f)$, we have

$$q(b) = 0 \neq 1 = h(b)$$
.

Therefore $g \neq h$ and f is not an epimorphism.

Step 2: Item 1a ← Item 1c. We claim Items 1a and 1c are indeed equivalent:

- · $ltem 1a \implies ltem 1c$: We proceed in a few steps:
 - Assume that f is surjective. Let $U, V \in \mathcal{P}(B)$ such that $f^{-1}(U) = f^{-1}(V)$. We wish to show that U = V.
 - To show that $U \subset V$, let $b \in U$.
 - Since f is surjective, there must exist some $a \in A$ such that f(a) = b.

- By the definition of the inverse image, since f(a) = b and $b \in U$, we have $a \in f^{-1}(U)$.
- By our initial assumption, $f^{-1}(U) = f^{-1}(V)$, so it follows that $a \in f^{-1}(V)$.
- Again, by the definition of the inverse image, $a \in f^{-1}(V)$ means that $f(a) \in V$.
- Since f(a) = b, we have shown that $b \in V$.
- This establishes that $U \subset V$. A symmetric argument shows that $V \subset U$.
- Thus U = V, proving that f^{-1} is injective.
- · $ltem1c \implies ltem1a$: We proceed in a few steps:
 - Assume that the inverse image function f^{-1} is injective. Suppose, for the sake of contradiction, that f is not surjective.
 - The assumption that f is not surjective means there exists some $b_0 \in B$ such that for all $a \in A$, we have $f(a) \neq b_0$.
 - By the definition of the inverse image, this is equivalent to stating that $f^{-1}(\{b_0\}) = \emptyset$.
 - Since $f^{-1}(\emptyset) = \emptyset$, we have $f^{-1}(\{b_0\}) = f^{-1}(\emptyset)$.
 - Since f^{-1} is injective, this implies that $\{b_0\} = \emptyset$.
 - This is a contradiction, as the singleton set $\{b_0\}$ is non-empty.
 - Therefore, *f* is surjective.

Step 3: Item 1c \iff Item 1d. We claim that Items 1c and 1d are equivalent:

- · $ltem 1c \implies ltem 1d$: We proceed in a few steps:
 - Let $U, V \in \mathcal{P}(B)$ such that $f^{-1}(U) \subset f^{-1}(V)$, assume f^{-1} to be injective, and consider the set $U \cup V$.
 - Since $f^{-1}(U) \subset f^{-1}(V)$, we have

$$f^{-1}(U \cup V) = f^{-1}(U) \cup f^{-1}(V)$$
$$= f^{-1}(V),$$

where we have used Item 5 of Definition 4.6.2.1.3 for the first equality.

- Since f^{-1} is injective, this implies $U \cup V = V$.
- Thus $U \subset V$, as we wished to show.
- · Item 1d \Longrightarrow Item 1c: We proceed in a few steps:
 - Suppose Item 1d holds, and let $U,V\in\mathcal{P}(B)$ such that $f^{-1}(U)=f^{-1}(V)$.
 - Since $f^{-1}(U)=f^{-1}(V)$, we have $f^{-1}(U)\subset f^{-1}(V)$ and $f^{-1}(V)\subset f^{-1}(U)$.
 - By assumption, this implies $U \subset V$ and $V \subset U$.
 - Thus U = V, showing f^{-1} to be injective.

Step 4: Item 1a \iff Item 1e. We have

$$f_!(f^{-1}(b)) = \left\{ b \in B \middle| \begin{array}{l} \text{there exists some } a \in f^{-1}(b) \\ \text{such that } f(a) = b \end{array} \right\},$$

so the condition $f_!(f^{-1}(b)) = \{b\}$ holds iff f is surjective.

Step 5: Item 1e \in Item 1f. We claim that Items 1e and 1f are indeed equivalent:

· Item 1e \Longrightarrow Item 1f: We have

$$[f! \circ f^{-1}](U) \stackrel{\text{def}}{=} f!(f^{-1}(U))$$

$$= f! \left(\bigcup_{u \in U} \{u\} \right) \right)$$

$$= f! \left(\bigcup_{u \in U} f^{-1}(\{u\}) \right)$$

$$= \bigcup_{u \in U} f!(f^{-1}(\{u\}))$$

$$= \bigcup_{u \in U} \{u\}$$

$$= U$$

for each $U \in \mathcal{P}(B)$, where we have used Item 5 of Definition 4.6.1.1.5 for the third equality and Item 5 of Definition 4.6.2.1.3 for the fourth equality.

· Item 1f \Longrightarrow Item 1e: Applying the condition $f_! \circ f^{-1} = \mathrm{id}_{\mathcal{P}(B)}$ to $U = \{b\}$ gives

$$f_!(f^{-1}(\{b\})) = \{b\}.$$

Step 6: Item 1a \iff Item 1g. First, note that for the condition $f^{-1}(b) \subset f^{-1}(U)$ to hold, we must have $b \in U$ or $f^{-1}(b) = \emptyset$. Thus

$$f_*(f^{-1}(U)) = (U \cap \operatorname{Im}(f)) \cup (B \setminus \operatorname{Im}(f)).$$

We now claim that Items 1a and 1g are indeed equivalent:

· Item 1a \Longrightarrow Item 1g: If f is surjective, we have

$$(U \cap \mathsf{Im}(f)) \cup (B \setminus \mathsf{Im}(f)) = U \cup \emptyset$$
$$= U,$$

so
$$f_* \circ f^{-1} = \mathrm{id}_{\mathcal{P}(B)}$$
.

· Item 1g \Longrightarrow Item 1a: Taking $U = \emptyset$ gives

$$f_*(f^{-1}(\emptyset)) = (\emptyset \cap \operatorname{Im}(f)) \cup (B \setminus \operatorname{Im}(f))$$
$$= B \setminus \operatorname{Im}(f),$$

so the condition $f_*(f^{-1}(\emptyset)) = \emptyset$ implies $B \setminus \text{Im}(f) = \emptyset$. Thus Im(f) = B and f is surjective.

This finishes the proof.

Appendices

A Other Chapters

- 1. Introduction
- 2. A Guide to the Literature

Sets

3. Sets

Preliminaries

- 4. Constructions With Sets
- Monoidal Structures on the Category of Sets
- 6. Pointed Sets
- 7. Tensor Products of Pointed Sets

Relations

- 8. Relations
- 9. Constructions With Relations

10. Conditions on Relations

Categories

- 11. Categories
- 12. Presheaves and the Yoneda Lemma

Monoidal Categories

13. Constructions With Monoidal Categories

Bicategories

14. Types of Morphisms in Bicategories

Extra Part

15. Notes

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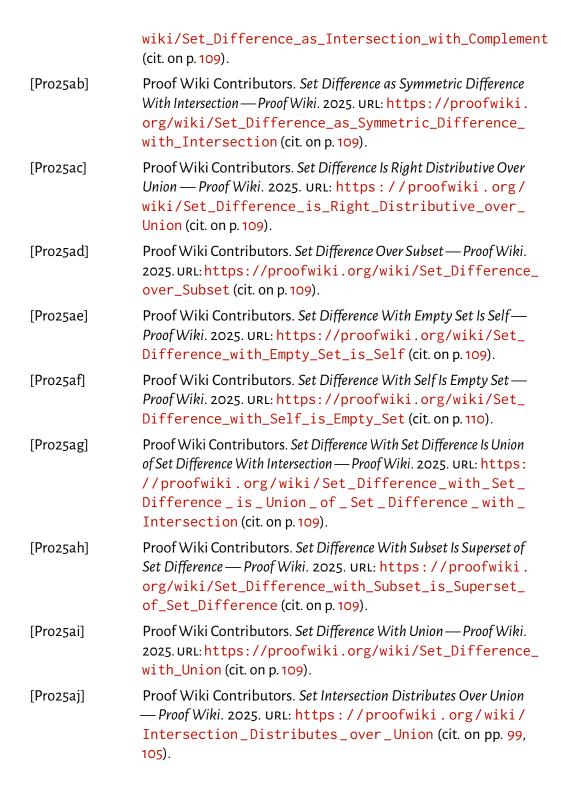
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