Monoidal Structures on the Category of Sets

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This chapter contains some material on monoidal structures on Sets.

Contents

5.1 The Monoidal Category of Sets and Products

5.1.1 Products of Sets

See Constructions With Sets, Section 4.1.3.

5.1.2 The Internal Hom of Sets

See Constructions With Sets, Section 4.3.5.

5.1.3 The Monoidal Unit

Definition 5.1.3.1.1. The **monoidal unit of the product of sets** is the functor

$$\mathbb{1}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

defined by

$$\mathbb{1}_{\mathsf{Sets}} \overset{\mathrm{def}}{=} \mathsf{pt},$$

where pt is the terminal set of Constructions With Sets, Definition 4.1.1.1.1.

5.1.4 The Associator

Definition 5.1.4.1.1. The **associator of the product of sets** is the natural isomorphism

$$\alpha^{\mathsf{Sets}} \colon \times \circ (\times \times \mathrm{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \times \circ (\mathrm{id}_{\mathsf{Sets}} \times \times) \circ \boldsymbol{\alpha}^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}},$$

as in the diagram

$$
\begin{array}{c}
\mathsf{Sets} \times (\mathsf{Sets} \times \mathsf{Sets}) \\[2em]
(\mathsf{Sets} \times \mathsf{Sets}) \times \mathsf{Sets} \qquad\qquad \mathsf{Sets} \times \mathsf{Sets} \\[2em]
\mathsf{Sets} \times \mathsf{Sets} \xrightarrow{\;\times\;} \mathsf{Sets},
\end{array}
$$

with arrows labeled $\alpha^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}$, $\mathsf{id}\times\times$, $\times\times\mathsf{id}$, α^{Sets}, $\times$.

whose component

$$
\alpha^{\mathsf{Sets}}_{X,Y,Z} \colon (X \times Y) \times Z \overset{\sim}{\dashrightarrow} X \times (Y \times Z)
$$

at (X, Y, Z) is given by

$$
\alpha^{\mathsf{Sets}}_{X,Y,Z}((x, y), z) \overset{\mathrm{def}}{=} (x, (y, z))
$$

for each $((x, y), z) \in (X \times Y) \times Z$.

Proof. *Invertibility*: The inverse of $\alpha^{\mathsf{Sets}}_{X,Y,Z}$ is the morphism

$$
\alpha^{\mathsf{Sets},-1}_{X,Y,Z} \colon X \times (Y \times Z) \overset{\sim}{\dashrightarrow} (X \times Y) \times Z
$$

defined by

$$
\alpha^{\mathsf{Sets},-1}_{X,Y,Z}(x, (y, z)) \overset{\mathrm{def}}{=} ((x, y), z)
$$

for each $(x, (y, z)) \in X \times (Y \times Z)$. Indeed:

· *Invertibility I.* We have

$$
\begin{aligned}
[\alpha^{\mathsf{Sets},-1}_{X,Y,Z} \circ \alpha^{\mathsf{Sets}}_{X,Y,Z}]((x, y), z) &\overset{\mathrm{def}}{=} \alpha^{\mathsf{Sets},-1}_{X,Y,Z}(\alpha^{\mathsf{Sets}}_{X,Y,Z}((x, y), z)) \\
&\overset{\mathrm{def}}{=} \alpha^{\mathsf{Sets},-1}_{X,Y,Z}(x, (y, z)) \\
&\overset{\mathrm{def}}{=} ((x, y), z) \\
&\overset{\mathrm{def}}{=} [\mathsf{id}_{(X\times Y)\times Z}]((x, y), z)
\end{aligned}
$$

for each $((x, y), z) \in (X \times Y) \times Z$, and therefore we have

$$
\alpha^{\mathsf{Sets},-1}_{X,Y,Z} \circ \alpha^{\mathsf{Sets}}_{X,Y,Z} = \mathsf{id}_{(X\times Y)\times Z} \, .
$$

· *Invertibility II.* We have

$$
\begin{aligned}
[\alpha_{X,Y,Z}^{\mathsf{Sets}} \circ \alpha_{X,Y,Z}^{\mathsf{Sets},-1}](x,(y,z)) &\overset{\mathrm{def}}{=} \alpha_{X,Y,Z}^{\mathsf{Sets}}(\alpha_{X,Y,Z}^{\mathsf{Sets},-1}(x,(y,z))) \\
&\overset{\mathrm{def}}{=} \alpha_{X,Y,Z}^{\mathsf{Sets}}((x,y),z) \\
&\overset{\mathrm{def}}{=} (x,(y,z)) \\
&\overset{\mathrm{def}}{=} [\mathrm{id}_{(X\times Y)\times Z}](x,(y,z))
\end{aligned}
$$

for each $(x,(y,z)) \in X \times (Y \times Z)$, and therefore we have

$$
\alpha_{X,Y,Z}^{\mathsf{Sets},-1} \circ \alpha_{X,Y,Z}^{\mathsf{Sets}} = \mathrm{id}_{X\times(Y\times Z)} \, .
$$

Therefore $\alpha_{X,Y,Z}^{\mathsf{Sets}}$ is indeed an isomorphism.
Naturality: We need to show that, given functions

$$
\begin{aligned}
f &: X \to X', \\
g &: Y \to Y', \\
h &: Z \to Z'
\end{aligned}
$$

the diagram

$$
\begin{array}{ccc}
(X \times Y) \times Z & \xrightarrow{\;(f\times g)\times h\;} & (X' \times Y') \times Z' \\
{\scriptstyle \alpha_{X,Y,Z}^{\mathsf{Sets}}}\Big\downarrow & & \Big\downarrow{\scriptstyle \alpha_{X',Y',Z'}^{\mathsf{Sets}}} \\
X \times (Y \times Z) & \xrightarrow[\;f\times(g\times h)\;]{} & X' \times (Y' \times Z')
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
((x,y),z) & & ((x,y),z) \longmapsto ((f(x),g(y)),h(z)) \\
\Big\downarrow & & \Big\downarrow \\
(x,(y,z)) \longmapsto (f(x),(g(y),h(z))) & & (f(x),(g(y),h(z)))
\end{array}
$$

and hence indeed commutes, showing α^{Sets} to be a natural transformation.
Being a Natural Isomorphism: Since α^{Sets} is natural and $\alpha^{\mathsf{Sets},-1}$ is a component-wise inverse to α^{Sets}, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\alpha^{\mathsf{Sets},-1}$ is also natural. Thus α^{Sets} is a natural isomorphism. □

5.1.5 The Left Unitor

Definition 5.1.5.1.1. The **left unitor of the product of sets** is the natural isomorphism

$$\lambda^{\mathsf{Sets}} \colon \times \circ (\mathbb{1}^{\mathsf{Sets}} \times \mathsf{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \lambda^{\mathsf{Cats2}}_{\mathsf{Sets}}$$

$$\begin{array}{ccc}
\mathsf{pt} \times \mathsf{Sets} & \xrightarrow{\ \mathbb{1}^{\mathsf{Sets}} \times \mathsf{id}\ } & \mathsf{Sets} \times \mathsf{Sets} \\
& \searrow\ {\scriptstyle \lambda^{\mathsf{Sets}}} & \Big\| \\
& {\scriptstyle \lambda^{\mathsf{Cats2}}_{\mathsf{Sets}}} & \times \\
& & \downarrow \\
& & \mathsf{Sets,}
\end{array}$$

whose component

$$\lambda^{\mathsf{Sets}}_X \colon \mathsf{pt} \times X \overset{\sim}{\dashrightarrow} X$$

at $X \in \mathrm{Obj}(\mathsf{Sets})$ is given by

$$\lambda^{\mathsf{Sets}}_X(\star, x) \overset{\mathrm{def}}{=} x$$

for each $(\star, x) \in \mathsf{pt} \times X$.

Proof. Invertibility: The inverse of $\lambda^{\mathsf{Sets}}_X$ is the morphism

$$\lambda^{\mathsf{Sets},-1}_X \colon X \overset{\sim}{\dashrightarrow} \mathsf{pt} \times X$$

defined by

$$\lambda^{\mathsf{Sets},-1}_X(x) \overset{\mathrm{def}}{=} (\star, x)$$

for each $x \in X$. Indeed:

· *Invertibility I.* We have

$$\begin{aligned}
[\lambda^{\mathsf{Sets},-1}_X \circ \lambda^{\mathsf{Sets}}_X](\mathsf{pt}, x) &= \lambda^{\mathsf{Sets},-1}_X(\lambda^{\mathsf{Sets}}_X(\mathsf{pt}, x)) \\
&= \lambda^{\mathsf{Sets},-1}_X(x) \\
&= (\mathsf{pt}, x) \\
&= [\mathsf{id}_{\mathsf{pt} \times X}](\mathsf{pt}, x)
\end{aligned}$$

for each $(\mathsf{pt}, x) \in \mathsf{pt} \times X$, and therefore we have

$$\lambda^{\mathsf{Sets},-1}_X \circ \lambda^{\mathsf{Sets}}_X = \mathsf{id}_{\mathsf{pt} \times X}.$$

- *Invertibility II.* We have

$$
\begin{aligned}
[\lambda_X^{\mathsf{Sets}} \circ \lambda_X^{\mathsf{Sets},-1}](x) &= \lambda_X^{\mathsf{Sets}}(\lambda_X^{\mathsf{Sets},-1}(x)) \\
&= \lambda_X^{\mathsf{Sets},-1}(\mathrm{pt}, x) \\
&= x \\
&= [\mathrm{id}_X](x)
\end{aligned}
$$

for each $x \in X$, and therefore we have

$$
\lambda_X^{\mathsf{Sets}} \circ \lambda_X^{\mathsf{Sets},-1} = \mathrm{id}_X \,.
$$

Therefore $\lambda_X^{\mathsf{Sets}}$ is indeed an isomorphism.

Naturality: We need to show that, given a function $f \colon X \to Y$, the diagram

$$
\begin{array}{ccc}
\mathrm{pt} \times X & \xrightarrow{\ \mathrm{id}_{\mathrm{pt}} \times f\ } & \mathrm{pt} \times Y \\
\downarrow{\scriptstyle \lambda_X^{\mathsf{Sets}}} & & \downarrow{\scriptstyle \lambda_Y^{\mathsf{Sets}}} \\
X & \xrightarrow[\ f\ ]{} & Y
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
(\star, x) & & (\star, x) \longmapsto (\star, f(x)) \\
\downarrow & & \downarrow \\
x \longmapsto f(x) & & f(x)
\end{array}
$$

and hence indeed commutes. Therefore λ^{Sets} is a natural transformation.

Being a Natural Isomorphism: Since λ^{Sets} is natural and $\lambda^{\mathsf{Sets},-1}$ is a componentwise inverse to λ^{Sets}, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\lambda^{\mathsf{Sets},-1}$ is also natural. Thus λ^{Sets} is a natural isomorphism. $\square$

5.1.6   The Right Unitor

Definition 5.1.6.1.1. The **right unitor of the product of sets** is the natural isomorphism

$$\rho^{\mathsf{Sets}} \colon \times \circ (\mathsf{id} \times \mathbb{1}^{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \boldsymbol{\rho}_{\mathsf{Sets}}^{\mathsf{Cats_2}},$$

whose component

$$\rho_X^{\mathsf{Sets}} \colon X \times \mathsf{pt} \overset{\sim}{\dashrightarrow} X$$

at $X \in \mathsf{Obj}(\mathsf{Sets})$ is given by

$$\rho_X^{\mathsf{Sets}}(x, \star) \overset{\mathrm{def}}{=} x$$

for each $(x, \star) \in X \times \mathsf{pt}$.

Proof. *Invertibility*: The inverse of ρ_X^{Sets} is the morphism

$$\rho_X^{\mathsf{Sets},-1} \colon X \overset{\sim}{\dashrightarrow} X \times \mathsf{pt}$$

defined by

$$\rho_X^{\mathsf{Sets},-1}(x) \overset{\mathrm{def}}{=} (x, \star)$$

for each $x \in X$. Indeed:

· *Invertibility I.* We have

$$
\begin{aligned}
[\rho_X^{\mathsf{Sets},-1} \circ \rho_X^{\mathsf{Sets}}](x, \star) &= \rho_X^{\mathsf{Sets},-1}(\rho_X^{\mathsf{Sets}}(x, \star)) \\
&= \rho_X^{\mathsf{Sets},-1}(x) \\
&= (x, \star) \\
&= [\mathsf{id}_{X \times \mathsf{pt}}](x, \star)
\end{aligned}
$$

for each $(x, \star) \in X \times \mathsf{pt}$, and therefore we have

$$\rho_X^{\mathsf{Sets},-1} \circ \rho_X^{\mathsf{Sets}} = \mathsf{id}_{X \times \mathsf{pt}}.$$

- *Invertibility II.* We have

$$
\begin{aligned}
[\rho_X^{\mathsf{Sets}} \circ \rho_X^{\mathsf{Sets},-1}](x) &= \rho_X^{\mathsf{Sets}}(\rho_X^{\mathsf{Sets},-1}(x)) \\
&= \rho_X^{\mathsf{Sets},-1}(x,\star) \\
&= x \\
&= [\mathrm{id}_X](x)
\end{aligned}
$$

for each $x \in X$, and therefore we have

$$
\rho_X^{\mathsf{Sets}} \circ \rho_X^{\mathsf{Sets},-1} = \mathrm{id}_X \,.
$$

Therefore ρ_X^{Sets} is indeed an isomorphism.

Naturality: We need to show that, given a function $f \colon X \to Y$, the diagram

$$
\begin{array}{ccc}
X \times \mathsf{pt} & \xrightarrow{\; f \times \mathrm{id}_{\mathsf{pt}} \;} & Y \times \mathsf{pt} \\
{\scriptstyle \rho_X^{\mathsf{Sets}}} \downarrow & & \downarrow {\scriptstyle \rho_Y^{\mathsf{Sets}}} \\
X & \xrightarrow{\quad f \quad} & Y
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
(x,\star) & & (x,\star) \longmapsto (f(x),\star) \\
\downarrow & & \downarrow \\
x \longmapsto f(x) & & f(x)
\end{array}
$$

and hence indeed commutes. Therefore ρ^{Sets} is a natural transformation.

Being a Natural Isomorphism: Since ρ^{Sets} is natural and $\rho^{\mathsf{Sets},-1}$ is a component-wise inverse to ρ^{Sets}, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\rho^{\mathsf{Sets},-1}$ is also natural. Thus ρ^{Sets} is a natural isomorphism.   □

5.1.7 The Symmetry

Definition 5.1.7.1.1. The **symmetry of the product of sets** is the natural isomorphism

$$\sigma^{\mathsf{Sets}} \colon \times \overset{\sim}{\Longrightarrow} \times \circ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}},$$

$$
\begin{array}{ccc}
\mathsf{Sets} \times \mathsf{Sets} & \xrightarrow{\ \times\ } & \mathsf{Sets}, \\
& & \| \\
\sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}} \searrow & \sigma^{\mathsf{Sets}} \Downarrow & \nearrow \times \\
& \mathsf{Sets} \times \mathsf{Sets} &
\end{array}
$$

whose component

$$\sigma^{\mathsf{Sets}}_{X,Y} \colon X \times Y \overset{\sim}{\dashrightarrow} Y \times X$$

at $X, Y \in \mathrm{Obj}(\mathsf{Sets})$ is defined by

$$\sigma^{\mathsf{Sets}}_{X,Y}(x, y) \overset{\mathrm{def}}{=} (y, x)$$

for each $(x, y) \in X \times Y$.

Proof. *Invertibility*: The inverse of $\sigma^{\mathsf{Sets}}_{X,Y}$ is the morphism

$$\sigma^{\mathsf{Sets},-1}_{X,Y} \colon Y \times X \overset{\sim}{\dashrightarrow} X \times Y$$

defined by

$$\sigma^{\mathsf{Sets},-1}_{X,Y}(y, x) \overset{\mathrm{def}}{=} (x, y)$$

for each $(y, x) \in Y \times X$. Indeed:

- *Invertibility I.* We have

$$
\begin{aligned}
[\sigma^{\mathsf{Sets},-1}_{X,Y} \circ \sigma^{\mathsf{Sets}}_{X,Y}](x, y) &\overset{\mathrm{def}}{=} \sigma^{\mathsf{Sets},-1}_{X,Y}(\sigma^{\mathsf{Sets}}_{X,Y}(x, y)) \\
&\overset{\mathrm{def}}{=} \sigma^{\mathsf{Sets},-1}_{X,Y}(y, x) \\
&\overset{\mathrm{def}}{=} (x, y) \\
&\overset{\mathrm{def}}{=} [\mathrm{id}_{X \times Y}](x, y)
\end{aligned}
$$

for each $(x, y) \in X \times Y$, and therefore we have

$$\sigma^{\mathsf{Sets},-1}_{X,Y} \circ \sigma^{\mathsf{Sets}}_{X,Y} = \mathrm{id}_{X \times Y}.$$

· *Invertibility II*. We have

$$[\sigma_{X,Y}^{\mathsf{Sets}} \circ \sigma_{X,Y}^{\mathsf{Sets},-1}](y,x) \overset{\mathrm{def}}{=} \sigma_{X,Y}^{\mathsf{Sets},-1}(\sigma_{X,Y}^{\mathsf{Sets}}(y,x))$$
$$\overset{\mathrm{def}}{=} \sigma_{X,Y}^{\mathsf{Sets},-1}(x,y)$$
$$\overset{\mathrm{def}}{=} (y,x)$$
$$\overset{\mathrm{def}}{=} [\mathrm{id}_{Y \times X}](y,x)$$

for each $(y,x) \in Y \times X$, and therefore we have

$$\sigma_{X,Y}^{\mathsf{Sets}} \circ \sigma_{X,Y}^{\mathsf{Sets},-1} = \mathrm{id}_{Y \times X}\ .$$

Therefore $\sigma_{X,Y}^{\mathsf{Sets}}$ is indeed an isomorphism.
Naturality: We need to show that, given functions

$$f \colon X \to A,$$
$$g \colon Y \to B$$

the diagram

$$
\begin{array}{ccc}
X \times Y & \xrightarrow{\ f \times g\ } & A \times B \\
{\scriptstyle \sigma_{X,Y}^{\mathsf{Sets}}} \downarrow & & \downarrow {\scriptstyle \sigma_{A,B}^{\mathsf{Sets}}} \\
Y \times X & \xrightarrow[\ g \times f\ ]{} & B \times A
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
(x,y) & & (x,y) \longmapsto (f(x),g(y)) \\
\downarrow & & \downarrow \\
(y,x) \longmapsto (g(y),f(x)) & & (g(y),f(x))
\end{array}
$$

and hence indeed commutes, showing σ^{Sets} to be a natural transformation.
Being a Natural Isomorphism: Since σ^{Sets} is natural and $\sigma^{\mathsf{Sets},-1}$ is a component-wise inverse to σ^{Sets}, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\sigma^{\mathsf{Sets},-1}$ is also natural. Thus σ^{Sets} is a natural isomorphism.   □

5.1.8   The Diagonal

Definition 5.1.8.1.1. The **diagonal of the product of sets** is the natural transformation

$$\Delta\colon \mathrm{id}_{\mathsf{Sets}} \Longrightarrow \times \circ \Delta_{\mathsf{Sets}}^{\mathsf{Cats}_2},$$

whose component

$$\Delta_X\colon X \to X \times X$$

at $X \in \mathrm{Obj}(\mathsf{Sets})$ is given by

$$\Delta_X(x) \overset{\mathrm{def}}{=} (x, x)$$

for each $x \in X$.

Proof. We need to show that, given a function $f\colon X \to Y$, the diagram

$$
\begin{array}{ccc}
X & \xrightarrow{\;\;f\;\;} & Y \\
{\scriptstyle \Delta_X}\downarrow & & \downarrow{\scriptstyle \Delta_Y} \\
X \times X & \xrightarrow[f \times f]{} & Y \times Y
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
x & & x \longmapsto f(x) \\
\downarrow & & \downarrow \\
(x,x) \longmapsto (f(x), f(x)) & & (f(x), f(x))
\end{array}
$$

and hence indeed commutes, showing Δ to be natural. $\square$

Proposition 5.1.8.1.2. Let X be a set.

1. *Monoidality.* The diagonal map

$$\Delta\colon \mathrm{id}_{\mathsf{Sets}} \Longrightarrow \times \circ \Delta_{\mathsf{Sets}}^{\mathsf{Cats}_2},$$

   is a monoidal natural transformation:

(a) *Compatibility With Strong Monoidality Constraints.* For each $X, Y \in \mathrm{Obj}(\mathsf{Sets})$, the diagram

$$
\begin{array}{ccc}
X \times Y & \xrightarrow{\Delta_X \times \Delta_Y} & (X \times X) \times (Y \times Y) \\
& \searrow_{\Delta_{X \times Y}} & \Big\downarrow{\wr} \\
& & (X \times Y) \times (X \times Y)
\end{array}
$$

commutes.

(b) *Compatibility With Strong Unitality Constraints.* The diagrams

$$
\begin{array}{ccc}
\mathsf{pt} \xrightarrow{\Delta_{\mathsf{pt}}} \mathsf{pt} \times \mathsf{pt} & \qquad & \mathsf{pt} \xrightarrow{\Delta_{\mathsf{pt}}} \mathsf{pt} \times \mathsf{pt} \\
\searrow \Big\downarrow{\lambda_{\mathsf{pt}}^{\mathsf{Sets}}} & & \searrow \Big\downarrow{\rho_{\mathsf{pt}}^{\mathsf{Sets}}} \\
\mathsf{pt} & & \mathsf{pt}
\end{array}
$$

commute, i.e. we have

$$
\begin{aligned}
\Delta_{\mathsf{pt}} &= \lambda_{\mathsf{pt}}^{\mathsf{Sets},-1} \\
&= \rho_{\mathsf{pt}}^{\mathsf{Sets},-1},
\end{aligned}
$$

where we recall that the equalities

$$
\begin{aligned}
\lambda_{\mathsf{pt}}^{\mathsf{Sets}} &= \rho_{\mathsf{pt}}^{\mathsf{Sets}}, \\
\lambda_{\mathsf{pt}}^{\mathsf{Sets},-1} &= \rho_{\mathsf{pt}}^{\mathsf{Sets},-1}
\end{aligned}
$$

are always true in any monoidal category by Monoidal Categories, **??** of **??**.

2. *The Diagonal of the Unit.* The component

$$
\Delta_{\mathsf{pt}} \colon \mathsf{pt} \overset{\sim}{\dashrightarrow} \mathsf{pt} \times \mathsf{pt}
$$

of Δ at pt is an isomorphism.

Proof. *Item 1*, *Monoidality*: We claim that Δ is indeed monoidal:

1. *Item 1a: Compatibility With Strong Monoidality Constraints:* We need to show that the diagram

$$
X \times Y \xrightarrow{\;\Delta_X \times \Delta_Y\;} (X \times X) \times (Y \times Y)
$$

with $\Delta_{X \times Y}$ going down to $(X \times Y) \times (X \times Y)$ and a vertical arrow labeled $\wr$

commutes. Indeed, this diagram acts on elements as

$$(x, y) \longmapsto ((x, x), (y, y)) \qquad (x, y)$$

$$((x, y), (x, y)) \qquad\qquad\qquad ((x, y), (x, y))$$

   and hence indeed commutes.

2. *Item 1b: Compatibility With Strong Unitality Constraints:* As shown in the proof of Definition 5.1.5.1.1, the inverse of the left unitor of Sets with respect to to the product at $X \in \mathrm{Obj}(\mathsf{Sets})$ is given by

$$
\lambda_X^{\mathsf{Sets},-1}(x) \overset{\mathrm{def}}{=} (\star, x)
$$

   for each $x \in X$, so when $X = \mathrm{pt}$, we have

$$
\lambda_{\mathrm{pt}}^{\mathsf{Sets},-1}(\star) \overset{\mathrm{def}}{=} (\star, \star),
$$

   and also

$$
\Delta_{\mathrm{pt}}^{\mathsf{Sets}}(\star) \overset{\mathrm{def}}{=} (\star, \star),
$$

   so we have $\Delta_{\mathrm{pt}} = \lambda_{\mathrm{pt}}^{\mathsf{Sets},-1}$.

This finishes the proof.

Item 2, The Diagonal of the Unit: This follows from Item 1 and the invertibility of the left/right unitor of Sets with respect to $\times$, proved in the proof of Definition 5.1.5.1.1 for the left unitor or the proof of Definition 5.1.6.1.1 for the right unitor.   □

5.1.9   The Monoidal Category of Sets and Products

Proposition 5.1.9.1.1. The category Sets admits a closed symmetric monoidal category with diagonals structure consisting of:

· *The Underlying Category.* The category Sets of pointed sets.

· *The Monoidal Product.* The product functor

$$\times \colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}$$

of Constructions With Sets, Item 1 of Definition 4.1.3.1.3.

· *The Internal Hom.* The internal Hom functor

$$\mathsf{Sets} \colon \mathsf{Sets}^{\mathsf{op}} \times \mathsf{Sets} \to \mathsf{Sets}$$

of Constructions With Sets, Item 1 of Definition 4.3.5.1.2.

· *The Monoidal Unit.* The functor

$$\mathbb{1}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

of Definition 5.1.3.1.1.

· *The Associators.* The natural isomorphism

$$\alpha^{\mathsf{Sets}} \colon \times \circ (\times \times \mathsf{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \times \circ (\mathsf{id}_{\mathsf{Sets}} \times \times) \circ \boldsymbol{\alpha}^{\mathsf{Cats}}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}$$

of Definition 5.1.4.1.1.

· *The Left Unitors.* The natural isomorphism

$$\lambda^{\mathsf{Sets}} \colon \times \circ (\mathbb{1}^{\mathsf{Sets}} \times \mathsf{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \boldsymbol{\lambda}^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.1.5.1.1.

· *The Right Unitors.* The natural isomorphism

$$\rho^{\mathsf{Sets}} \colon \times \circ (\mathsf{id} \times \mathbb{1}^{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \boldsymbol{\rho}^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.1.6.1.1.

· *The Symmetry.* The natural isomorphism

$$\sigma^{\mathsf{Sets}} \colon \times \overset{\sim}{\Longrightarrow} \times \circ \boldsymbol{\sigma}^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}}$$

of Definition 5.1.7.1.1.

· *The Diagonals.* The monoidal natural transformation

$$\Delta \colon \mathsf{id}_{\mathsf{Sets}} \Longrightarrow \times \circ \Delta^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.1.8.1.1.

Proof. *The Pentagon Identity*: Let W, X, Y and Z be sets. We have to show that the diagram

$$
\begin{array}{c}
(W \times (X \times Y)) \times Z
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

and thus the pentagon identity is satisfied.

The Triangle Identity: Let X and Y be sets. We have to show that the diagram

$$
\begin{array}{ccc}
(X \times \mathrm{pt}) \times Y & \xrightarrow{\alpha^{\mathsf{Sets}}_{X,\mathrm{pt},Y}} & X \times (\mathrm{pt} \times Y) \\
& & \\
\rho^{\mathsf{Sets}}_X \times \mathrm{id}_Y \searrow & & \swarrow \mathrm{id}_X \times \lambda^{\mathsf{Sets}}_Y \\
& X \times Y &
\end{array}
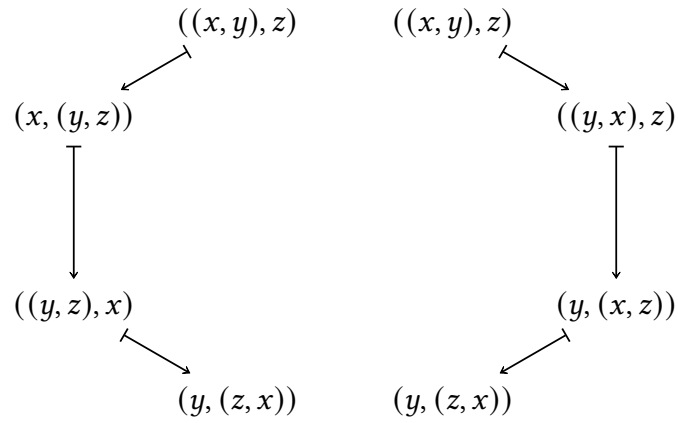$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
((x,\star),y) & \qquad\qquad ((x,\star),y) \longmapsto (x,(\star,y)) \\
\searrow & \qquad\qquad\qquad\qquad \swarrow \\
(x,y) & \qquad\qquad\qquad (x,y)
\end{array}
$$

and thus the triangle identity is satisfied.

The Left Hexagon Identity: Let X, Y, and Z be sets. We have to show that the diagram
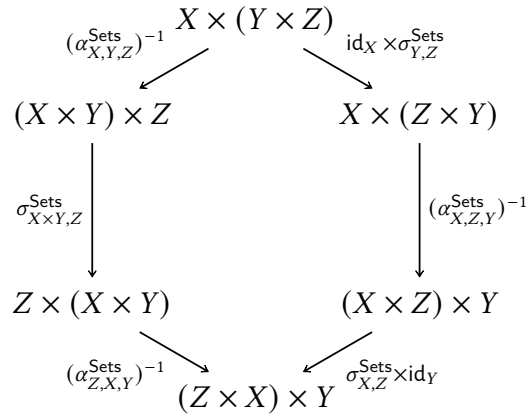
$$
\begin{array}{ccc}
& (X \times Y) \times Z & \\
\alpha^{\mathsf{Sets}}_{X,Y,Z} \swarrow & & \searrow \sigma^{\mathsf{Sets}}_{X,Y} \times \mathrm{id}_Z \\
X \times (Y \times Z) & & (Y \times X) \times Z \\
\sigma^{\mathsf{Sets}}_{X,Y \times Z} \downarrow & & \downarrow \alpha^{\mathsf{Sets}}_{Y,X,Z} \\
(Y \times Z) \times X & & Y \times (X \times Z) \\
\alpha^{\mathsf{Sets}}_{Y,Z,X} \searrow & & \swarrow \mathrm{id}_Y \times \sigma^{\mathsf{Sets}}_{X,Z} \\
& Y \times (Z \times X) &
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
& ((x,y),z) \qquad\qquad ((x,y),z) & \\
(x,(y,z)) & & ((y,x),z) \\
& & \\
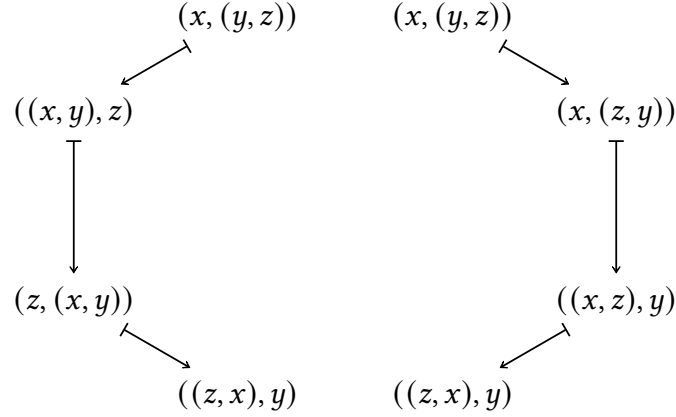((y,z),x) & & (y,(x,z)) \\
& (y,(z,x)) \qquad (y,(z,x)) &
\end{array}
$$

 and thus the left hexagon identity is satisfied.

The Right Hexagon Identity: Let X, Y, and Z be sets. We have to show that the diagram

$$
\begin{array}{ccc}
& X \times (Y \times Z) & \\
(\alpha^{\mathsf{Sets}}_{X,Y,Z})^{-1} \swarrow & & \searrow \mathrm{id}_X \times \sigma^{\mathsf{Sets}}_{Y,Z} \\
(X \times Y) \times Z & & X \times (Z \times Y) \\
\sigma^{\mathsf{Sets}}_{X \times Y,Z} \downarrow & & \downarrow (\alpha^{\mathsf{Sets}}_{X,Z,Y})^{-1} \\
Z \times (X \times Y) & & (X \times Z) \times Y \\
(\alpha^{\mathsf{Sets}}_{Z,X,Y})^{-1} \searrow & & \swarrow \sigma^{\mathsf{Sets}}_{X,Z} \times \mathrm{id}_Y \\
& (Z \times X) \times Y &
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
(x,(y,z)) & & (x,(y,z)) \\
& \searrow & & \nwarrow \\
((x,y),z) & & & (x,(z,y)) \\
\downarrow & & & \downarrow \\
(z,(x,y)) & & & ((x,z),y) \\
& \nwarrow & & \swarrow \\
& ((z,x),y) & ((z,x),y) &
\end{array}
$$

and thus the right hexagon identity is satisfied.

Monoidal Closedness: This follows from Constructions With Sets, Item 2 of Definition 4.3.5.1.2

Existence of Monoidal Diagonals: This follows from Items 1 and 2 of Definition 5.1.8.1.2.

□

5.1.10   The Universal Property of $(\mathsf{Sets}, \times, \mathbf{pt})$

Theorem 5.1.10.1.1. The symmetric monoidal structure on the category Sets of Definition 5.1.9.1.1 is uniquely determined by the following requirements:

1. *Existence of an Internal Hom*. The tensor product

    $$\otimes_{\mathsf{Sets}} \colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}$$

    of Sets admits an internal Hom $[-_1, -_2]_{\mathsf{Sets}}$.

2. *The Unit Object Is pt*. We have $\mathbb{1}_{\mathsf{Sets}} \cong \mathrm{pt}$.

More precisely, the full subcategory of the category $\mathcal{M}_{\mathbb{E}_\infty}^{\mathsf{cld}}(\mathsf{Sets})$ of **??** spanned by the closed symmetric monoidal categories $(\mathsf{Sets}, \otimes_{\mathsf{Sets}}, [-_1, -_2]_{\mathsf{Sets}}, \mathbb{1}_{\mathsf{Sets}}, \lambda^{\mathsf{Sets}}, \rho^{\mathsf{Sets}}, \sigma^{\mathsf{Sets}})$ satisfying Items 1 and 2 is contractible (i.e. equivalent to the punctual category).

Proof. Unwinding the Statement: Let $(\mathsf{Sets}, \otimes_{\mathsf{Sets}}, [-_1, -_2]_{\mathsf{Sets}}, \mathbb{1}_{\mathsf{Sets}}, \lambda', \rho', \sigma')$ be a closed symmetric monoidal category satisfying Items 1 and 2. We need to show that the identity functor

$$\mathrm{id}_{\mathsf{Sets}} \colon \mathsf{Sets} \to \mathsf{Sets}$$

admits a *unique* closed symmetric monoidal functor structure

$$
\begin{aligned}
\mathrm{id}^{\otimes}_{\mathsf{Sets}} &: & A \otimes_{\mathsf{Sets}} B & \xdashrightarrow{\sim} & A \times B, \\
\mathrm{id}^{\mathsf{Hom}}_{\mathsf{Sets}} &: & [A, B]_{\mathsf{Sets}} & \xdashrightarrow{\sim} & \mathsf{Sets}(A, B), \\
\mathrm{id}^{\otimes}_{\mathbb{1}|\mathsf{Sets}} &: & \mathbb{1}_{\mathsf{Sets}} & \xdashrightarrow{\sim} & \mathrm{pt},
\end{aligned}
$$

making it into a symmetric monoidal strongly closed isomorphism of categories from $(\mathsf{Sets}, \otimes_{\mathsf{Sets}}, [-_1, -_2]_{\mathsf{Sets}}, \mathbb{1}_{\mathsf{Sets}}, \lambda', \rho', \sigma')$ to the closed symmetric monoidal category $\left(\mathsf{Sets}, \times, \mathsf{Sets}(-_1, -_2), \mathbb{1}_{\mathsf{Sets}}, \lambda^{\mathsf{Sets}}, \rho^{\mathsf{Sets}}, \sigma^{\mathsf{Sets}}\right)$ of Definition 5.1.9.1.1.
Constructing an Isomorphism $[-_1, -_2]_{\mathsf{Sets}} \cong \mathsf{Sets}(-_1, -_2)$: By **??**, we have a natural isomorphism

$$\mathsf{Sets}(\mathrm{pt}, [-_1, -_2]_{\mathsf{Sets}}) \cong \mathsf{Sets}(-_1, -_2).$$

By Constructions With Sets, Item 3 of Definition 4.3.5.1.2, we also have a natural isomorphism

$$\mathsf{Sets}(\mathrm{pt}, [-_1, -_2]_{\mathsf{Sets}}) \cong [-_1, -_2]_{\mathsf{Sets}}.$$

Composing both natural isomorphisms, we obtain a natural isomorphism

$$\mathsf{Sets}(-_1, -_2) \cong [-_1, -_2]_{\mathsf{Sets}}.$$

Given $A, B \in \mathrm{Obj}(\mathsf{Sets})$, we will write

$$\mathrm{id}^{\mathsf{Hom}}_{A,B} : \mathsf{Sets}(A, B) \xdashrightarrow{\sim} [A, B]_{\mathsf{Sets}}$$

for the component of this isomorphism at (A, B).
Constructing an Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$: Since $\otimes_{\mathsf{Sets}}$ is adjoint in each variable to $[-_1, -_2]_{\mathsf{Sets}}$ by assumption and $\times$ is adjoint in each variable to $\mathsf{Sets}(-_1, -_2)$ by Constructions With Sets, Item 2 of Definition 4.3.5.1.2, uniqueness of adjoints (**??**) gives us natural isomorphisms

$$
\begin{aligned}
A \otimes_{\mathsf{Sets}} - &\cong A \times -, \\
- \otimes_{\mathsf{Sets}} B &\cong B \times -.
\end{aligned}
$$

By **??**, we then have $\otimes_{\mathsf{Sets}} \cong \times$. We will write

$$\mathrm{id}^{\otimes}_{\mathsf{Sets}|A,B} : A \otimes_{\mathsf{Sets}} B \xdashrightarrow{\sim} A \times B$$

for the component of this isomorphism at (A, B).
Alternative Construction of an Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$: Alternatively, we may construct a natural isomorphism $\otimes_{\mathsf{Sets}} \cong \times$ as follows:

1. Let $A \in \mathsf{Obj}(\mathsf{Sets})$.

2. Since $\otimes_{\mathsf{Sets}}$ is part of a closed monoidal structure, it preserves colimits in each variable by **??**.

3. Since $A \cong \coprod_{a \in A} \mathsf{pt}$ and $\otimes_{\mathsf{Sets}}$ preserves colimits in each variable, we have

$$A \otimes_{\mathsf{Sets}} B \cong (\coprod_{a \in A} \mathsf{pt}) \otimes_{\mathsf{Sets}} B$$
$$\cong \coprod_{a \in A} (\mathsf{pt} \otimes_{\mathsf{Sets}} B)$$
$$\cong \coprod_{a \in A} B$$
$$\cong A \times B,$$

   naturally in $B \in \mathsf{Obj}(\mathsf{Sets})$, where we have used that pt is the monoidal unit for $\otimes_{\mathsf{Sets}}$. Thus $A \otimes_{\mathsf{Sets}} - \cong A \times -$ for each $A \in \mathsf{Obj}(\mathsf{Sets})$.

4. Similarly, $- \otimes_{\mathsf{Sets}} B \cong - \times B$ for each $B \in \mathsf{Obj}(\mathsf{Sets})$.

5. By **??**, we then have $\otimes_{\mathsf{Sets}} \cong \times$.

Below, we,ll show that if a natural isomorphism $\otimes_{\mathsf{Sets}} \cong \times$ exists, then it must be unique. This will show that the isomorphism constructed above is equal to the isomorphism $\mathrm{id}^{\otimes}_{\mathsf{Sets}|A,B} : A \otimes_{\mathsf{Sets}} B \to A \times B$ from before.

Constructing an Isomorphism $\mathrm{id}^{\otimes}_{\mathbb{1}} : \mathbb{1}_{\mathsf{Sets}} \to pt$: We define an isomorphism $\mathrm{id}^{\otimes}_{\mathbb{1}} : \mathbb{1}_{\mathsf{Sets}} \to$ pt as the composition

$$\mathbb{1}_{\mathsf{Sets}} \overset{\rho^{\mathsf{Sets},-1}_{\mathbb{1}_{\mathsf{Sets}}}}{\underset{\sim}{\dashrightarrow}} \mathbb{1}_{\mathsf{Sets}} \times \mathsf{pt} \overset{\mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathbb{1}_{\mathsf{Sets}}}}{\underset{\sim}{\dashrightarrow}} \mathbb{1}_{\mathsf{Sets}} \otimes_{\mathsf{Sets}} \mathsf{pt} \overset{\lambda'_{\mathsf{pt}}}{\underset{\sim}{\dashrightarrow}} \mathsf{pt}$$

in Sets.

Monoidal Left Unity of the Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$: We have to show that the diagram

commutes. First, note that the diagram

$$
\begin{array}{ccc}
& \mathrm{pt} \otimes_{\mathsf{Sets}} \mathrm{pt} & \xrightarrow{\;\mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathrm{pt},\mathrm{pt}}\;} \mathrm{pt} \times \mathrm{pt} \\
\mathrm{id}^{\otimes}_{\mathbb{1}|\mathsf{Sets}} \otimes_{\mathsf{Sets}} \mathrm{id}_{\mathrm{pt}} \nearrow & & \searrow \lambda^{\mathsf{Sets}}_{\mathrm{pt}} \\
\mathbb{1}_{\mathsf{Sets}} \otimes_{\mathsf{Sets}} \mathrm{pt} & \xrightarrow[\lambda'_{\mathrm{pt}}]{\hspace{5cm}} & \mathrm{pt},
\end{array}
$$

corresponding to the case $A = \mathrm{pt}$, commutes by the terminality of pt (Constructions With Sets, Definition 4.1.1.1.2). Since this diagram commutes, so does the diagram

$$
\begin{array}{ccc}
\mathrm{pt} \times \mathrm{pt} & \xrightarrow{\;\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|\mathrm{pt},\mathrm{pt}}\;} & \mathrm{pt} \otimes_{\mathsf{Sets}} \mathrm{pt} \\
\lambda^{\mathsf{Sets},-1}_{\mathrm{pt}} \nearrow & (\dagger) & \searrow \mathrm{id}^{\otimes,-1}_{\mathbb{1}|\mathsf{Sets}} \otimes_{\mathsf{Sets}} \mathrm{id}_{\mathrm{pt}} \\
\mathrm{pt} & \xrightarrow[\lambda'^{,-1}_{\mathrm{pt}}]{} & \mathbb{1}_{\mathsf{Sets}} \otimes_{\mathsf{Sets}} \mathrm{pt}.
\end{array}
$$

Now, let $A \in \mathrm{Obj}(\mathsf{Sets})$, let $a \in A$, and consider the diagram

$$
\begin{array}{c}
\text{(diagram)}
\end{array}
$$

Since:

· Subdiagram (5) commutes by the naturality of $\lambda'^{,-1}$.

· Subdiagram (†) commutes, as proved above.

· Subdiagram (4) commutes by the naturality of $\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1}$.

· Subdiagram (1) commutes by the naturality of $\mathrm{id}_{\mathsf{Sets}}^{\otimes,-1}$.

· Subdiagram (3) commutes by the naturality of $\lambda^{\mathsf{Sets},-1}$.

it follows that the diagram

$$
\begin{array}{ccc}
& \mathrm{pt} \times A & \xrightarrow{\mathrm{id}_{\mathsf{Sets}|\mathrm{pt},A}^{\otimes,-1}} & \mathrm{pt} \otimes_{\mathsf{Sets}} A \\
\nearrow{\scriptstyle \lambda_A^{\mathsf{Sets},-1}} & & & \searrow{\scriptstyle \mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \otimes_{\mathsf{Sets}} \mathrm{id}_A} \\
\mathrm{pt} \xrightarrow{[a]} A & & \xrightarrow[\lambda_A'^{,-1}]{} & \mathbb{1}_{\mathsf{Sets}} \otimes_{\mathsf{Sets}} A
\end{array}
$$

Here,s a step-by-step showcase of this argument: [Link]. We then have

$$
\begin{aligned}
\lambda_A'^{-1}(a) &= [\lambda_A'^{-1} \circ [a]](\star) \\
&= [(\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \times \mathrm{id}_A) \circ \mathrm{id}_{\mathsf{Sets}|\mathrm{pt},A}^{\otimes,-1} \circ \lambda_A^{\mathsf{Sets},-1} \circ [a]](\star) \\
&= [(\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \times \mathrm{id}_A) \circ \mathrm{id}_{\mathsf{Sets}|\mathrm{pt},A}^{\otimes,-1} \circ \lambda_A^{\mathsf{Sets},-1}](a)
\end{aligned}
$$

for each $a \in A$, and thus we have

$$
\lambda_A'^{,-1} = (\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes,-1} \times \mathrm{id}_A) \circ \mathrm{id}_{\mathsf{Sets}|\mathrm{pt},A}^{\otimes,-1} \circ \lambda_A^{\mathsf{Sets},-1}.
$$

Taking inverses then gives

$$
\lambda_A' = \lambda_A^{\mathsf{Sets}} \circ \mathrm{id}_{\mathsf{Sets}|\mathrm{pt},A}^{\otimes} \circ (\mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes} \times \mathrm{id}_A),
$$

showing that the diagram

$$
\begin{array}{ccc}
& \mathrm{pt} \otimes_{\mathsf{Sets}} A & \xrightarrow{\mathrm{id}_{\mathsf{Sets}|\mathrm{pt},A}^{\otimes}} & \mathrm{pt} \times A \\
\nearrow{\scriptstyle \mathrm{id}_{\mathbb{1}|\mathsf{Sets}}^{\otimes} \otimes_{\mathsf{Sets}} \mathrm{id}_A} & & & \searrow{\scriptstyle \lambda_A^{\mathsf{Sets}}} \\
\mathbb{1}_{\mathsf{Sets}} \otimes_{\mathsf{Sets}} A & & \xrightarrow[\lambda_A']{} & A
\end{array}
$$

indeed commutes.

Monoidal Right Unity of the Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$: We can use the same argument we used to prove the monoidal left unity of the isomorphism $\otimes_{\mathsf{Sets}} \cong \times$ above. For completeness, we repeat it below.

We have to show that the diagram

$$
\begin{array}{ccc}
& A \otimes_{\mathsf{Sets}} \mathsf{pt} \xrightarrow{\ \mathrm{id}^{\otimes}_{\mathsf{Sets}|A,\mathsf{pt}}\ } A \times \mathsf{pt} & \\
{\scriptstyle \mathrm{id}_A \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes}_{\mathbb{1}|\mathsf{Sets}}} \nearrow & & \searrow {\scriptstyle \rho^{\mathsf{Sets}}_A} \\
A \otimes_{\mathsf{Sets}} \mathbb{1}_{\mathsf{Sets}} \xrightarrow[\ \rho'_A\ ]{} & & A
\end{array}
$$

commutes. First, note that the diagram

$$
\begin{array}{ccc}
& \mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt} \xrightarrow{\ \mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}}\ } \mathsf{pt} \times \mathsf{pt} & \\
{\scriptstyle \mathrm{id}_{\mathsf{pt}} \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes}_{\mathbb{1}|\mathsf{Sets}}} \nearrow & & \searrow {\scriptstyle \rho^{\mathsf{Sets}}_{\mathsf{pt}}} \\
\mathsf{pt} \otimes_{\mathsf{Sets}} \mathbb{1}_{\mathsf{Sets}} \xrightarrow[\ \rho'_{\mathsf{pt}}\ ]{} & & \mathsf{pt},
\end{array}
$$

corresponding to the case $A = \mathsf{pt}$, commutes by the terminality of pt (Constructions With Sets, Definition 4.1.1.1.2). Since this diagram commutes, so does the diagram

$$
\begin{array}{ccc}
& \mathsf{pt} \times \mathsf{pt} \xrightarrow{\ \mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}}\ } \mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt} & \\
{\scriptstyle \rho^{\mathsf{Sets},-1}_{\mathsf{pt}}} \nearrow & (\dagger) & \searrow {\scriptstyle \mathrm{id}_{\mathsf{pt}} \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes,-1}_{\mathbb{1}|\mathsf{Sets}}} \\
\mathsf{pt} \xrightarrow[\ \rho'^{,-1}_{\mathsf{pt}}\ ]{} & & \mathsf{pt} \otimes_{\mathsf{Sets}} \mathbb{1}_{\mathsf{Sets}}.
\end{array}
$$

Now, let $A \in \mathrm{Obj}(\mathsf{Sets})$, let $a \in A$, and consider the diagram

Since:

- · Subdiagram (5) commutes by the naturality of $\rho'^{,-1}$.

- · Subdiagram (†) commutes, as proved above.

- · Subdiagram (4) commutes by the naturality of $\mathrm{id}^{\otimes,-1}_{\mathbb{1}|\mathsf{Sets}}$.

- · Subdiagram (1) commutes by the naturality of $\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}}$.

- · Subdiagram (3) commutes by the naturality of $\rho^{\mathsf{Sets},-1}$.

it follows that the diagram

$$
\begin{array}{ccc}
& A \times \mathrm{pt} \xrightarrow{\ \mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|A,\mathrm{pt}}\ } A \otimes_{\mathsf{Sets}} \mathrm{pt} & \\
{\scriptstyle \rho^{\mathsf{Sets},-1}_A}\nearrow & & \searrow{\scriptstyle \mathrm{id}_A \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes,-1}_{\mathbb{1}|\mathsf{Sets}}} \\
\mathrm{pt} \xrightarrow{\ [a]\ } A & \xrightarrow{\qquad\quad \rho'^{,-1}_A \qquad\quad} & A \otimes_{\mathsf{Sets}} \mathbb{1}_{\mathsf{Sets}}
\end{array}
$$

Here,s a step-by-step showcase of this argument: [Link]. We then have

$$
\begin{aligned}
\rho'^{,-1}_A(a) &= [\rho'^{,-1}_A \circ [a]](\star) \\
&= [(\mathrm{id}_A \times \mathrm{id}^{\otimes,-1}_{\mathbb{1}|\mathsf{Sets}}) \circ \mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|\mathrm{pt},A} \circ \rho^{\mathsf{Sets},-1}_A \circ [a]](\star) \\
&= [(\mathrm{id}_A \times \mathrm{id}^{\otimes,-1}_{\mathbb{1}|\mathsf{Sets}}) \circ \mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|\mathrm{pt},A} \circ \rho^{\mathsf{Sets},-1}_A](a)
\end{aligned}
$$

for each $a \in A$, and thus we have

$$
\rho'^{,-1}_A = (\mathrm{id}_A \times \mathrm{id}^{\otimes,-1}_{\mathbb{1}|\mathsf{Sets}}) \circ \mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|\mathrm{pt},A} \circ \rho^{\mathsf{Sets},-1}_A.
$$

Taking inverses then gives

$$
\rho'_A = \rho^{\mathsf{Sets}}_A \circ \mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathrm{pt},A} \circ (\mathrm{id}_A \times \mathrm{id}^{\otimes}_{\mathbb{1}|\mathsf{Sets}}),
$$

showing that the diagram

$$
\begin{array}{ccc}
& A \otimes_{\mathsf{Sets}} \mathrm{pt} \xrightarrow{\ \mathrm{id}^{\otimes}_{\mathsf{Sets}|A,\mathrm{pt}}\ } A \times \mathrm{pt} & \\
{\scriptstyle \mathrm{id}_A \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes}_{\mathbb{1}|\mathsf{Sets}}}\nearrow & & \searrow{\scriptstyle \rho^{\mathsf{Sets}}_A} \\
A \otimes_{\mathsf{Sets}} \mathbb{1}_{\mathsf{Sets}} & \xrightarrow{\qquad\quad \rho'_A \qquad\quad} & A
\end{array}
$$

indeed commutes.

Monoidality of the Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$: We have to show that the diagram

$$
\begin{array}{ccc}
& (A \otimes_{\mathsf{Sets}} B) \otimes_{\mathsf{Sets}} C & \\
\mathrm{id}^{\otimes}_{\mathsf{Sets}|A,B} \otimes_{\mathsf{Sets}} \mathrm{id}_C \swarrow & & \searrow \alpha'_{A,B,C} \\
(A \times B) \otimes_{\mathsf{Sets}} C & & A \otimes_{\mathsf{Sets}} (B \otimes_{\mathsf{Sets}} C) \\
\mathrm{id}^{\otimes}_{\mathsf{Sets}|A\times B,C} \downarrow & & \downarrow \mathrm{id}_A \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes}_{\mathsf{Sets}|B,C} \\
(A \times B) \times C & & A \otimes_{\mathsf{Sets}} (B \times C) \\
\alpha^{\mathsf{Sets}}_{A,B,C} \searrow & & \swarrow \mathrm{id}^{\otimes}_{\mathsf{Sets}|A,B\times C} \\
& A \times (B \times C) &
\end{array}
$$

commutes. First, note that the diagram

$$
\begin{array}{ccc}
& (\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt}) \otimes_{\mathsf{Sets}} \mathsf{pt} & \\
\mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}} \otimes_{\mathsf{Sets}} \mathrm{id}_{\mathsf{pt}} \swarrow & & \searrow \alpha'_{\mathsf{pt},\mathsf{pt},\mathsf{pt}} \\
(\mathsf{pt} \times \mathsf{pt}) \otimes_{\mathsf{Sets}} \mathsf{pt} & & \mathsf{pt} \otimes_{\mathsf{Sets}} (\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt}) \\
\mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathsf{pt}\times\mathsf{pt},\mathsf{pt}} \downarrow & & \downarrow \mathrm{id}_{\mathsf{pt}} \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}} \\
(\mathsf{pt} \times \mathsf{pt}) \times \mathsf{pt} & & \mathsf{pt} \otimes_{\mathsf{Sets}} (\mathsf{pt} \times \mathsf{pt}) \\
\alpha^{\mathsf{Sets}}_{\mathsf{pt},\mathsf{pt},\mathsf{pt}} \searrow & & \swarrow \mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}\times\mathsf{pt}} \\
& \mathsf{pt} \times (\mathsf{pt} \times \mathsf{pt}) & \\
& \Big| {\scriptstyle !_{\mathsf{pt}\times(\mathsf{pt}\times\mathsf{pt})}} & \\
& \downarrow & \\
& \mathsf{pt} &
\end{array}
$$

commutes by the terminality of pt (Constructions With Sets, Definition 4.1.1.1.2).

Since the map $!_{\mathsf{pt}\times(\mathsf{pt}\times\mathsf{pt})} \colon \mathsf{pt} \times (\mathsf{pt} \times \mathsf{pt}) \to \mathsf{pt}$ is an isomorphism (e.g. having

inverse $\lambda_{\mathrm{pt}}^{\mathsf{Sets},-1} \circ \lambda_{\mathrm{pt}}^{\mathsf{Sets},-1}$), it follows that the diagram

$$
\begin{array}{c}
(\mathrm{pt} \otimes_{\mathsf{Sets}} \mathrm{pt}) \otimes_{\mathsf{Sets}} \mathrm{pt} \\[2em]
\end{array}
$$

also commutes. Taking inverses, we see that the diagram

commutes as well. Now, let $A, B, C \in \mathrm{Obj}(\mathsf{Sets})$, let $a \in A$, let $b \in B$, let $c \in C$,

and consider the diagram

$$\mathsf{pt} \times (\mathsf{pt} \times \mathsf{pt})$$

$$\alpha^{\mathsf{Sets},-1}_{\mathsf{pt},\mathsf{pt},\mathsf{pt}} \qquad \mathrm{id}^{\otimes,-1}_{\mathsf{pt},\mathsf{pt}\times\mathsf{pt}} \qquad [a]\times([b]\times[c])$$

$$(\mathsf{pt} \times \mathsf{pt}) \times \mathsf{pt} \qquad \mathsf{pt} \otimes_{\mathsf{Sets}} (\mathsf{pt} \times \mathsf{pt}) \qquad A \times (B \times C) \quad \mathrm{id}^{\otimes,-1}_{\mathsf{Sets}/A,B\times C}$$

$$([a]\times[b])\times[c] \qquad \alpha^{\mathsf{Sets},-1}_{A,B,C} \quad [a]\otimes_{\mathsf{Sets}}([b]\times[c]) \to$$

$$\mathrm{id}^{\otimes,-1}_{\mathsf{pt}\times\mathsf{pt},\mathsf{pt}} \qquad \mathrm{id}_{\mathsf{pt}} \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes,-1}_{\mathsf{pt},\mathsf{pt}} \qquad (A \times B) \times C \qquad A \otimes_{\mathsf{Sets}} (B \times C)$$

$$\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}/A\times B,C}$$

$$(\mathsf{pt} \times \mathsf{pt}) \otimes_{\mathsf{Sets}} \mathsf{pt} \qquad \mathsf{pt} \otimes_{\mathsf{Sets}} (\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt}) \qquad \mathrm{id}_A \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes,-1}_{\mathsf{Sets}/B,C}$$

$$\mathrm{id}^{\otimes,-1}_{\mathsf{pt},\mathsf{pt}} \otimes_{\mathsf{Sets}} \mathrm{id}_{\mathsf{pt}} \qquad ([a]\times[b])\otimes_{\mathsf{Sets}}[c] \qquad \alpha'^{-1}_{\mathsf{pt},\mathsf{pt},\mathsf{pt}} \quad \sim [a]\otimes_{\mathsf{Sets}}([b]\otimes_{\mathsf{Sets}}[c]) \to$$

$$(\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt}) \otimes_{\mathsf{Sets}} \mathsf{pt} \qquad (A \times B) \otimes_{\mathsf{Sets}} C \qquad A \otimes_{\mathsf{Sets}} (B \otimes_{\mathsf{Sets}} C)$$

$$\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}/A,B} \otimes_{\mathsf{Sets}} \mathrm{id}_C \qquad \alpha'^{-1}_{A,B,C}$$

$$([a]\otimes_{\mathsf{Sets}}[b])\otimes_{\mathsf{Sets}}[c] \qquad (A \otimes_{\mathsf{Sets}} B) \otimes_{\mathsf{Sets}} C$$

which we partition into subdiagrams as follows:

$$\mathsf{pt} \times (\mathsf{pt} \times \mathsf{pt})$$

$$\alpha^{\mathsf{Sets},-1}_{\mathsf{pt},\mathsf{pt},\mathsf{pt}} \qquad \mathrm{id}^{\otimes,-1}_{\mathsf{pt},\mathsf{pt}\times\mathsf{pt}}$$

$$(\mathsf{pt} \times \mathsf{pt}) \times \mathsf{pt} \qquad\qquad \mathsf{pt} \otimes_{\mathsf{Sets}} (\mathsf{pt} \times \mathsf{pt})$$

$$\mathrm{id}^{\otimes,-1}_{\mathsf{pt}\times\mathsf{pt},\mathsf{pt}} \qquad (\dagger) \qquad \mathrm{id}_{\mathsf{pt}} \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes,-1}_{\mathsf{pt},\mathsf{pt}}$$

$$(\mathsf{pt} \times \mathsf{pt}) \otimes_{\mathsf{Sets}} \mathsf{pt} \qquad\qquad \mathsf{pt} \otimes_{\mathsf{Sets}} (\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt})$$

$$\mathrm{id}^{\otimes,-1}_{\mathsf{pt},\mathsf{pt}} \otimes_{\mathsf{Sets}} \mathrm{id}_{\mathsf{pt}} \qquad \alpha'^{-1}_{\mathsf{pt},\mathsf{pt},\mathsf{pt}}$$

$$(\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt}) \otimes_{\mathsf{Sets}} \mathsf{pt}$$

Since:

- · Subdiagram (1) commutes by the naturality of $\alpha^{\mathsf{Sets},-1}$.

- · Subdiagram (2) commutes by the naturality of $\mathrm{id}_{\mathsf{Sets}}^{\otimes,-1}$.

- · Subdiagram (3) commutes by the naturality of $\mathrm{id}_{\mathsf{Sets}}^{\otimes,-1}$.

- · Subdiagram (†) commutes, as proved above.

· Subdiagram (4) commutes by the naturality of $\mathrm{id}_{\mathsf{Sets}}^{\otimes,-1}$.

· Subdiagram (5) commutes by the naturality of $\mathrm{id}_{\mathsf{Sets}}^{\otimes,-1}$.

· Subdiagram (6) commutes by the naturality of $\alpha'^{,-1}$.

it follows that the diagram

$$
\begin{array}{c}
\mathsf{pt} \times (\mathsf{pt} \times \mathsf{pt}) \\
| \\
[a]\times([b]\times[c]) \\
\downarrow \\
A \times (B \times C)
\end{array}
$$

with maps $\alpha_{A,B,C}^{\mathsf{Sets},-1}$ to $(A \times B) \times C$ and $\mathrm{id}_{\mathsf{Sets}|A,B\times C}^{\otimes,-1}$ to $A \otimes_{\mathsf{Sets}} (B \times C)$;

$\mathrm{id}_{\mathsf{Sets}|A\times B,C}^{\otimes,-1}: (A \times B) \times C \to (A \times B) \otimes_{\mathsf{Sets}} C$;

$\mathrm{id}_A \times \mathrm{id}_{\mathsf{Sets}|B,C}^{\otimes,-1}: A \otimes_{\mathsf{Sets}} (B \times C) \to A \otimes_{\mathsf{Sets}} (B \otimes_{\mathsf{Sets}} C)$;

$\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} \otimes_{\mathsf{Sets}} \mathrm{id}_C$ and $\alpha'^{,-1}_{A,B,C}$ to $(A \otimes_{\mathsf{Sets}} B) \otimes_{\mathsf{Sets}} C$

also commutes. We then have

$$
\begin{aligned}
\Big[(\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} & \otimes_{\mathsf{Sets}} \mathrm{id}_C) \circ \mathrm{id}_{\mathsf{Sets}|A\times B,C}^{\otimes,-1} \\
& \circ \alpha_{A,B,C}^{\mathsf{Sets},-1} \Big](a,(b,c)) = \Big[(\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} \otimes_{\mathsf{Sets}} \mathrm{id}_C) \circ \mathrm{id}_{\mathsf{Sets}|A\times B,C}^{\otimes,-1} \\
& \qquad \circ \alpha_{A,B,C}^{\mathsf{Sets},-1} \circ ([a] \times ([b] \times [c])) \Big](\star,(\star,\star)) \\
& = \Big[\alpha'^{,-1}_{A,B,C} \circ (\mathrm{id}_A \times \mathrm{id}_{\mathsf{Sets}|B,C}^{\otimes,-1}) \\
& \qquad \circ \mathrm{id}_{\mathsf{Sets}|A,B\times C}^{\otimes,-1} \circ ([a] \times ([b] \times [c])) \Big](\star,(\star,\star)) \\
& = [\alpha'^{,-1}_{A,B,C} \circ (\mathrm{id}_A \times \mathrm{id}_{\mathsf{Sets}|B,C}^{\otimes,-1}) \circ \mathrm{id}_{\mathsf{Sets}|A,B\times C}^{\otimes,-1}](a,(b,c))
\end{aligned}
$$

for each $(a,(b,c)) \in A \times (B \times C)$, and thus we have

$$
(\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes,-1} \otimes_{\mathsf{Sets}} \mathrm{id}_C) \circ \mathrm{id}_{\mathsf{Sets}|A\times B,C}^{\otimes,-1} \circ \alpha_{A,B,C}^{\mathsf{Sets},-1} = \alpha'^{,-1}_{A,B,C} \circ (\mathrm{id}_A \times \mathrm{id}_{\mathsf{Sets}|B,C}^{\otimes,-1}) \circ \mathrm{id}_{\mathsf{Sets}|A,B\times C}^{\otimes,-1} .
$$

Taking inverses then gives

$$
\alpha_{A,B,C}^{\mathsf{Sets}} \circ \mathrm{id}_{\mathsf{Sets}|A\times B,C}^{\otimes} \circ (\mathrm{id}_{\mathsf{Sets}|A,B}^{\otimes} \otimes_{\mathsf{Sets}} \mathrm{id}_C) = \mathrm{id}_{\mathsf{Sets}|A,B\times C}^{\otimes} \circ (\mathrm{id}_A \times \mathrm{id}_{\mathsf{Sets}|B,C}^{\otimes}) \circ \alpha'_{A,B,C},
$$

showing that the diagram

$$
\begin{array}{ccc}
& (A \otimes_{\mathsf{Sets}} B) \otimes_{\mathsf{Sets}} C & \\
\scriptstyle \mathrm{id}^{\otimes}_{\mathsf{Sets}|A,B} \otimes_{\mathsf{Sets}} \mathrm{id}_C \swarrow & & \searrow \scriptstyle \alpha'_{A,B,C} \\
(A \times B) \otimes_{\mathsf{Sets}} C & & A \otimes_{\mathsf{Sets}} (B \otimes_{\mathsf{Sets}} C) \\
\scriptstyle \mathrm{id}^{\otimes}_{\mathsf{Sets}|A\times B,C} \Big\downarrow & & \Big\downarrow \scriptstyle \mathrm{id}_A \otimes_{\mathsf{Sets}} \mathrm{id}^{\otimes}_{\mathsf{Sets}|B,C} \\
(A \times B) \times C & & A \otimes_{\mathsf{Sets}} (B \times C) \\
\scriptstyle \alpha^{\mathsf{Sets}}_{A,B,C} \searrow & & \swarrow \scriptstyle \mathrm{id}^{\otimes}_{\mathsf{Sets}|A,B\times C} \\
& A \times (B \times C) &
\end{array}
$$

indeed commutes.

Braidedness of the Isomorphism $\otimes_{\mathsf{Sets}} \cong \times$: We have to show that the diagram

$$
\begin{array}{ccc}
A \otimes_{\mathsf{Sets}} B & \xrightarrow{\ \mathrm{id}^{\otimes}_{\mathsf{Sets}|A,B}\ } & A \times B \\
\scriptstyle \sigma'_{A,B} \Big\downarrow & & \Big\downarrow \scriptstyle \sigma^{\mathsf{Sets}}_{A,B} \\
B \otimes_{\mathsf{Sets}} A & \xrightarrow[\ \mathrm{id}^{\otimes}_{\mathsf{Sets}|B,A}\ ]{} & B \times A
\end{array}
$$

commutes. First, note that the diagram

$$
\begin{array}{ccc}
\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt} & \xrightarrow{\ \mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}}\ } & \mathsf{pt} \times \mathsf{pt} \\
\scriptstyle \sigma'_{\mathsf{pt},\mathsf{pt}} \Big\downarrow & & \Big\downarrow \scriptstyle \sigma^{\mathsf{Sets}}_{\mathsf{pt},\mathsf{pt}} \\
\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt} & \xrightarrow[\ \mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}}\ ]{} & \mathsf{pt} \times \mathsf{pt} \\
& & \quad \searrow \scriptstyle !_{\mathsf{pt}\times\mathsf{pt}} \\
& & \qquad \mathsf{pt}
\end{array}
$$

commutes by the terminality of pt (Constructions With Sets, Definition 4.1.1.1.2).
Since the map $!_{\mathsf{pt}\times\mathsf{pt}}\colon \mathsf{pt} \times \mathsf{pt} \to \mathsf{pt}$ is invertible (e.g. with inverse $\lambda^{\mathsf{Sets},-1}_{\mathsf{pt}}$), the

diagram

$$
\begin{array}{ccc}
\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt} & \xrightarrow{\mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}}} & \mathsf{pt} \times \mathsf{pt} \\
{\scriptstyle\sigma'_{\mathsf{pt},\mathsf{pt}}}\downarrow & & \downarrow{\scriptstyle\sigma^{\mathsf{Sets}}_{\mathsf{pt},\mathsf{pt}}} \\
\mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt} & \xrightarrow[\mathrm{id}^{\otimes}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}}]{} & \mathsf{pt} \times \mathsf{pt}
\end{array}
$$

also commutes. Taking inverses, we see that the diagram

$$
\begin{array}{ccc}
\mathsf{pt} \times \mathsf{pt} & \xrightarrow{\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}}} & \mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt} \\
{\scriptstyle\sigma^{\mathsf{Sets},-1}_{\mathsf{pt},\mathsf{pt}}}\downarrow & (\dagger) & \downarrow{\scriptstyle\sigma'^{,-1}_{\mathsf{pt},\mathsf{pt}}} \\
\mathsf{pt} \times \mathsf{pt} & \xrightarrow[\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|\mathsf{pt},\mathsf{pt}}]{} & \mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt}
\end{array}
$$

commutes as well. Now, let $A, B \in \mathrm{Obj}(\mathsf{Sets})$, let $a \in A$, let $b \in B$, and consider the diagram

which we partition into subdiagrams as follows:

$$
\begin{array}{ccc}
\mathsf{pt} \times \mathsf{pt} & \xrightarrow{\mathrm{id}^{\otimes,-1}_{\mathsf{pt},\mathsf{pt}}} & \mathsf{pt} \otimes_{\mathsf{Sets}} \mathsf{pt} \\
\end{array}
$$

with subdiagrams (2), (3), and morphisms $\sigma^{\mathsf{Sets},-1}_{\mathsf{pt},\mathsf{pt}}$, $[b] \times [a]$, $\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|B,A}$, $\sigma^{\mathsf{Sets},-1}_{A,B}$, $[a]\times[b]$, $\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|A,B}$, $[b]\otimes_{\mathsf{Sets}}[a]$, $\sigma'^{,-1}_{A,B}$

and (†), (4), (5) diagram with $\mathrm{id}^{\otimes,-1}_{\mathsf{pt},\mathsf{pt}}$, $\sigma^{\mathsf{Sets},-1}_{\mathsf{pt},\mathsf{pt}}$, $\sigma'^{,-1}_{\mathsf{pt},\mathsf{pt}}$, $[a]\otimes_{\mathsf{Sets}}[b]$

Since:

- Subdiagram (2) commutes by the naturality of $\sigma^{\mathsf{Sets},-1}$.

- Subdiagram (5) commutes by the naturality of $\mathrm{id}^{\otimes,-1}$.

- Subdiagram (†) commutes, as proved above.

- Subdiagram (4) commutes by the naturality of $\sigma'^{,-1}$.

- Subdiagram (1) commutes by the naturality of $\mathrm{id}^{\otimes,-1}$.

it follows that the diagram

$$
\begin{array}{ccc}
\mathsf{pt} \times \mathsf{pt} & & \\
& \searrow^{[b]\times[a]} & \\
& B \times A & \xrightarrow{\mathrm{id}^{\otimes}_{\mathsf{Sets}|B,A}} B \otimes_{\mathsf{Sets}} A \\
& \sigma^{\mathsf{Sets}}_{A,B} \downarrow & \downarrow \sigma'_{A,B} \\
& A \times B & \xrightarrow{\mathrm{id}^{\otimes}_{\mathsf{Sets}|A,B}} A \otimes_{\mathsf{Sets}} B
\end{array}
$$

commutes. We then have

$$
[\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|A,B} \circ \sigma^{\mathsf{Sets},-1}_{A,B}](b,a) = [\mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|A,B} \circ \sigma^{\mathsf{Sets},-1}_{A,B} \circ ([b]\times[a])](\star,\star)
$$
$$
= [\sigma'^{,-1}_{A,B} \circ \mathrm{id}^{\otimes,-1}_{\mathsf{Sets}|B,A} \circ ([b]\times[a])](\star,\star)
$$

$$= [\sigma'^{,-1}_{A,B} \circ \text{id}^{\otimes,-1}_{\text{Sets}|B,A}](b,a)$$

for each $(b,a) \in B \times A$, and thus we have

$$\text{id}^{\otimes,-1}_{\text{Sets}|A,B} \circ \sigma^{\text{Sets},-1}_{A,B} = \sigma'^{,-1}_{A,B} \circ \text{id}^{\otimes,-1}_{\text{Sets}|B,A}.$$

Taking inverses then gives

$$\sigma^{\text{Sets}}_{A,B} \circ \text{id}^{\otimes}_{\text{Sets}|A,B} = \text{id}^{\otimes}_{\text{Sets}|B,A} \circ \sigma'_{A,B},$$

showing that the diagram

$$
\begin{array}{ccc}
A \otimes_{\text{Sets}} B & \xrightarrow{\text{id}^{\otimes}_{\text{Sets}|A,B}} & A \times B \\
\downarrow{\scriptstyle \sigma'_{A,B}} & & \downarrow{\scriptstyle \sigma^{\text{Sets}}_{A,B}} \\
B \otimes_{\text{Sets}} A & \xrightarrow[\text{id}^{\otimes}_{\text{Sets}|B,A}]{} & B \times A
\end{array}
$$

indeed commutes.

Uniqueness of the Isomorphism $\otimes_{\text{Sets}} \cong \times$: Let $\phi, \psi \colon -_1 \otimes_{\text{Sets}} -_2 \Rightarrow -_1 \times -_2$ be natural isomorphisms. Since these isomorphisms are compatible with the unitors of Sets with respect to $\times$ and $\otimes$ (as shown above), we have

$$\lambda'_B = \lambda^{\text{Sets}}_B \circ \phi_{\text{pt},B} \circ (\text{id}^{\otimes}_{\mathbb{1}|\text{Sets}} \otimes_{\text{Sets}} \text{id}_Y),$$

$$\lambda'_B = \lambda^{\text{Sets}}_B \circ \psi_{\text{pt},B} \circ (\text{id}^{\otimes}_{\mathbb{1}|\text{Sets}} \otimes_{\text{Sets}} \text{id}_Y).$$

Postcomposing both sides with $\lambda^{\text{Sets},-1}_B$ gives

$$\lambda^{\text{Sets},-1}_B \circ \lambda'_B \circ (\text{id}^{\otimes,-1}_{\mathbb{1}|\text{Sets}} \otimes_{\text{Sets}} \text{id}_Y) = \phi_{\text{pt},B},$$

$$\lambda^{\text{Sets},-1}_B \circ \lambda'_B \circ (\text{id}^{\otimes}_{\mathbb{1}|\text{Sets}} \otimes_{\text{Sets}} \text{id}_Y) = \psi_{\text{pt},B},$$

and thus we have

$$\phi_{\text{pt},B} = \psi_{\text{pt},B}$$

for each $B \in \text{Obj}(\text{Sets})$. Now, let $a \in A$ and consider the naturality diagrams

$$
\begin{array}{ccc}
\text{pt} \times B & \xrightarrow{[a]\times\text{id}_B} & A \times B \\
\downarrow{\scriptstyle \phi_{\text{pt},B}} & & \downarrow{\scriptstyle \phi_{A,B}} \\
\text{pt} \otimes_{\text{Sets}} B & \xrightarrow[{[a]\otimes_{\text{Sets}}\text{id}_B}]{} & A \otimes_{\text{Sets}} B
\end{array}
\qquad
\begin{array}{ccc}
\text{pt} \times B & \xrightarrow{[a]\times\text{id}_B} & A \times B \\
\downarrow{\scriptstyle \psi_{\text{pt},B}} & & \downarrow{\scriptstyle \psi_{A,B}} \\
\text{pt} \otimes_{\text{Sets}} B & \xrightarrow[{[a]\otimes_{\text{Sets}}\text{id}_B}]{} & A \otimes_{\text{Sets}} B
\end{array}
$$

for ϕ and ψ with respect to the morphisms $[a]$ and id_B. Having shown that $\phi_{\mathrm{pt},B} = \psi_{\mathrm{pt},B}$, we have

$$
\begin{aligned}
\phi_{A,B}(a,b) &= [\phi_{A,B} \circ ([a] \times \mathrm{id}_B)](\star, b) \\
&= [([a] \otimes_{\mathsf{Sets}} \mathrm{id}_B) \circ \phi_{\mathrm{pt},B}](\star, b) \\
&= [([a] \otimes_{\mathsf{Sets}} \mathrm{id}_B) \circ \psi_{\mathrm{pt},B}](\star, b) \\
&= [\psi_{A,B} \circ ([a] \times \mathrm{id}_B)](\star, b) \\
&= \psi_{A,B}(a,b)
\end{aligned}
$$

for each $(a,b) \in A \times B$. Therefore we have

$$
\phi_{A,B} = \psi_{A,B}
$$

for each $A, B \in \mathrm{Obj}(\mathsf{Sets})$ and thus $\phi = \psi$, showing the isomorphism $\otimes_{\mathsf{Sets}} \cong \times$ to be unique. $\qquad\square$

Corollary 5.1.10.1.2. The symmetric monoidal structure on the category Sets of Definition 5.1.9.1.1 is uniquely determined by the following requirements:

1. *Two-Sided Preservation of Colimits.* The tensor product

$$
\otimes_{\mathsf{Sets}} \colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}
$$

   of Sets preserves colimits separately in each variable.

2. *The Unit Object Is pt.* We have $\mathbb{1}_{\mathsf{Sets}} \cong \mathrm{pt}$.

More precisely, the full subcategory of the category $\mathcal{M}_{\mathbb{E}_\infty}(\mathsf{Sets})$ of **??** spanned by the symmetric monoidal categories $(\mathsf{Sets}, \otimes_{\mathsf{Sets}}, \mathbb{1}_{\mathsf{Sets}}, \lambda^{\mathsf{Sets}}, \rho^{\mathsf{Sets}}, \sigma^{\mathsf{Sets}})$ satisfying Items 1 and 2 is contractible.

Proof. Since Sets is locally presentable (**??**), it follows from **??** that Item 1 is equivalent to the existence of an internal Hom as in Item 1 of Definition 5.1.10.1.1. The result then follows from Definition 5.1.10.1.1. $\qquad\square$

5.2 The Monoidal Category of Sets and Coproducts

5.2.1 Coproducts of Sets

See Constructions With Sets, Section 4.2.3.

5.2.2   The Monoidal Unit

Definition 5.2.2.1.1. The **monoidal unit of the coproduct of sets** is the functor

$$\mathbb{0}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

defined by

$$\mathbb{0}_{\mathsf{Sets}} \overset{\mathrm{def}}{=} \emptyset,$$

where $\emptyset$ is the empty set of Constructions With Sets, Definition 4.3.1.1.1.

5.2.3   The Associator

Definition 5.2.3.1.1. The **associator of the coproduct of sets** is the natural isomorphism

$$\alpha^{\mathsf{Sets}, \coprod} \colon \coprod \circ (\coprod \times \mathrm{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \coprod \circ (\mathrm{id}_{\mathsf{Sets}} \times \coprod) \circ \boldsymbol{\alpha}^{\mathsf{Cats}}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}},$$

as in the diagram

whose component

$$\alpha^{\mathsf{Sets}, \coprod}_{X,Y,Z} \colon (X \coprod Y) \coprod Z \overset{\sim}{\dashrightarrow} X \coprod (Y \coprod Z)$$

at (X, Y, Z) is given by

$$\alpha^{\mathsf{Sets}, \coprod}_{X,Y,Z}(a) \overset{\mathrm{def}}{=} \begin{cases} (0, x) & \text{if } a = (0, (0, x)), \\ (1, (0, y)) & \text{if } a = (0, (1, y)), \\ (1, (1, a)) & \text{if } a = (1, z) \end{cases}$$

for each $a \in (X \coprod Y) \coprod Z$.

Proof. *Unwinding the Definitions of* $(X \coprod Y) \coprod Z$ *and* $X \coprod (Y \coprod Z)$: Firstly, we unwind the expressions for $(X \coprod Y) \coprod Z$ and $X \coprod (Y \coprod Z)$. We have

$$
\begin{aligned}
(X \coprod Y) \coprod Z &\overset{\text{def}}{=} \{(0, a) \in S \mid a \in X \coprod Y\} \cup \{(1, z) \in S \mid z \in Z\} \\
&= \{(0, (0, x)) \in S \mid x \in X\} \cup \{(0, (1, y)) \in S \mid y \in Y\} \\
&\quad \cup \{(1, z) \in S \mid z \in Z\},
\end{aligned}
$$

where $S = \{0, 1\} \times ((X \coprod Y) \cup Z)$ and

$$
\begin{aligned}
X \coprod (Y \coprod Z) &\overset{\text{def}}{=} \{(0, x) \in S' \mid x \in X\} \cup \{(1, a) \in S' \mid a \in Y \coprod Z\} \\
&= \{(0, x) \in S' \mid x \in X\} \cup \{(1, (0, y)) \in S' \mid y \in Y\} \\
&\quad \cup \{(1, (1, z)) \in S' \mid z \in Z\},
\end{aligned}
$$

where $S' = \{0, 1\} \times (X \cup (Y \coprod Z))$.

Invertibility: The inverse of $\alpha_{X,Y,Z}^{\mathsf{Sets}, \coprod}$ is the map

$$
\alpha_{X,Y,Z}^{\mathsf{Sets}, \coprod, -1} : X \coprod (Y \coprod Z) \to (X \coprod Y) \coprod Z
$$

given by

$$
\alpha_{X,Y,Z}^{\mathsf{Sets}, \coprod, -1}(a) \overset{\text{def}}{=} \begin{cases} (0, (0, x)) & \text{if } a = (0, x), \\ (0, (1, y)) & \text{if } a = (1, (0, y)), \\ (1, z) & \text{if } a = (1, (1, z)) \end{cases}
$$

for each $a \in X \coprod Y (\coprod Z)$. Indeed:

- *Invertibility I.* The map $\alpha_{X,Y,Z}^{\mathsf{Sets}, \coprod, -1} \circ \alpha_{X,Y,Z}^{\mathsf{Sets}, \coprod}$ acts on elements as

$$
\begin{aligned}
(0, (0, x)) &\mapsto (0, x) \mapsto (0, (0, x)), \\
(0, (0, y)) &\mapsto (1, (0, y)) \mapsto (0, (0, y)), \\
(1, z) &\mapsto (1, (1, z)) \mapsto (1, z)
\end{aligned}
$$

  and hence is equal to the identity map of $(X \coprod Y) \coprod Z$.

- *Invertibility II.* The map $\alpha_{X,Y,Z}^{\mathsf{Sets}, \coprod} \circ \alpha_{X,Y,Z}^{\mathsf{Sets}, \coprod, -1}$ acts on elements as

$$
\begin{aligned}
(0, x) &\mapsto (0, (0, x)) \mapsto (0, x), \\
(1, (0, y)) &\mapsto (0, (0, y)) \mapsto (1, (0, y)), \\
(1, (1, z)) &\mapsto (1, z) \mapsto (1, (1, z))
\end{aligned}
$$

  and hence is equal to the identity map of $X \coprod (Y \coprod Z)$.

Therefore $\alpha_{X,Y,Z}^{\mathsf{Sets},\coprod}$ is indeed an isomorphism.

Naturality: We need to show that, given functions

$$f\colon X \to X',$$
$$g\colon Y \to Y',$$
$$h\colon Z \to Z'$$

the diagram

$$
\begin{array}{ccc}
(X \coprod Y) \coprod Z & \xrightarrow{(f\coprod g)\coprod h} & (X' \coprod Y') \coprod Z' \\
\downarrow{\scriptstyle \alpha_{X,Y,Z}^{\mathsf{Sets},\coprod}} & & \downarrow{\scriptstyle \alpha_{X',Y',Z'}^{\mathsf{Sets},\coprod}} \\
X \coprod (Y \coprod Z) & \xrightarrow[f\coprod(g\coprod h)]{} & X' \coprod (Y' \coprod Z')
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{cc}
(0,(0,x)) & \qquad (0,(0,x)) \longmapsto (0,(0,f(x))) \\
\downarrow & \qquad\qquad \downarrow \\
(0,x) \longmapsto (0,f(x)) & \qquad (0,f(x)) \\[1em]
(0,(1,y)) & \qquad (0,(1,y)) \longmapsto (0,(1,g(y))) \\
\downarrow & \qquad\qquad \downarrow \\
(1,(0,y)) \longmapsto (1,(0,g(y))) & \qquad (1,(0,g(y))) \\[1em]
(1,z) & \qquad (1,z) \longmapsto (1,h(z)) \\
\downarrow & \qquad\qquad \downarrow \\
(1,(1,z)) \longmapsto (1,(1,h(z))) & \qquad (1,(1,h(z)))
\end{array}
$$

and hence indeed commutes, showing $\alpha^{\mathsf{Sets},\coprod}$ to be a natural transformation.

Being a Natural Isomorphism: Since $\alpha^{\mathsf{Sets},\coprod}$ is natural and $\alpha^{\mathsf{Sets},\coprod,-1}$ is a componentwise inverse to $\alpha^{\mathsf{Sets},\coprod}$, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\lambda^{\mathsf{Sets},-1}$ is also natural. Thus $\alpha^{\mathsf{Sets},\coprod}$ is a natural isomorphism. $\square$

5.2.4   The Left Unitor

Definition 5.2.4.1.1. The **left unitor of the coproduct of sets** is the natural isomorphism

$$\lambda^{\mathsf{Sets},\amalg} \colon \amalg \circ (\mathbb{0}^{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \lambda^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

whose component

$$\lambda^{\mathsf{Sets},\amalg}_X \colon \emptyset \amalg X \overset{\sim}{\dashrightarrow} X$$

at X is given by

$$\lambda^{\mathsf{Sets},\amalg}_X((1,x)) \overset{\mathrm{def}}{=} x$$

for each $(1, x) \in \emptyset \amalg X$.

Proof.   Unwinding the Definition of $\emptyset \amalg X$: Firstly, we unwind the expressions for $\emptyset \amalg X$. We have

$$
\begin{aligned}
\emptyset \amalg X &\overset{\mathrm{def}}{=} \{(0,z) \in S \mid z \in \emptyset\} \cup \{(1,x) \in S \mid x \in X\} \\
&= \emptyset \cup \{(1,x) \in S \mid x \in X\} \\
&= \{(1,x) \in S \mid x \in X\},
\end{aligned}
$$

where $S = \{0,1\} \times (\emptyset \cup X)$.

Invertibility: The inverse of $\lambda^{\mathsf{Sets},\amalg}_X$ is the map

$$\lambda^{\mathsf{Sets},\amalg,-1}_X \colon X \to \emptyset \amalg X$$

given by

$$\lambda^{\mathsf{Sets},\amalg,-1}_X(x) \overset{\mathrm{def}}{=} (1,x)$$

for each $x \in X$. Indeed:

· *Invertibility I.* We have

$$[\lambda_X^{\mathsf{Sets},\coprod,-1} \circ \lambda_X^{\mathsf{Sets},\coprod}](1,x) = \lambda_X^{\mathsf{Sets},\coprod,-1}(\lambda_X^{\mathsf{Sets},\coprod}(1,x))$$
$$= \lambda_X^{\mathsf{Sets},\coprod,-1}(x)$$
$$= (1,x)$$
$$= [\mathrm{id}_{\emptyset\coprod X}](1,x)$$

for each $(1,x) \in \emptyset \coprod X$, and therefore we have

$$\lambda_X^{\mathsf{Sets},\coprod,-1} \circ \lambda_X^{\mathsf{Sets},\coprod} = \mathrm{id}_{\emptyset\coprod X}\,.$$

· *Invertibility II.* We have

$$[\lambda_X^{\mathsf{Sets},\coprod} \circ \lambda_X^{\mathsf{Sets},\coprod,-1}](x) = \lambda_X^{\mathsf{Sets},\coprod}(\lambda_X^{\mathsf{Sets},\coprod,-1}(x))$$
$$= \lambda_X^{\mathsf{Sets},\coprod,-1}(1,x)$$
$$= x$$
$$= [\mathrm{id}_X](x)$$

for each $x \in X$, and therefore we have

$$\lambda_X^{\mathsf{Sets},\coprod} \circ \lambda_X^{\mathsf{Sets},\coprod,-1} = \mathrm{id}_X\,.$$

Therefore $\lambda_X^{\mathsf{Sets},\coprod}$ is indeed an isomorphism.
Naturality: We need to show that, given a function $f\colon X \to Y$, the diagram

$$
\begin{array}{ccc}
\emptyset \coprod X & \xrightarrow{\mathrm{id}_\emptyset \coprod f} & \emptyset \coprod Y \\
{\scriptstyle \lambda_X^{\mathsf{Sets},\coprod}}\Big\downarrow & & \Big\downarrow{\scriptstyle \lambda_Y^{\mathsf{Sets},\coprod}} \\
X & \xrightarrow{\ f\ } & Y
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
(1,x) & & (1,x) \longmapsto (1,f(x)) \\
\Big\downarrow & & \Big\downarrow \\
x \longmapsto f(x) & & f(x)
\end{array}
$$

and hence indeed commutes. Therefore $\lambda^{\mathsf{Sets},\amalg}$ is a natural transformation.
Being a Natural Isomorphism: Since $\lambda^{\mathsf{Sets},\amalg}$ is natural and $\lambda^{\mathsf{Sets},-1}$ is a component-wise inverse to $\lambda^{\mathsf{Sets},\amalg}$, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\lambda^{\mathsf{Sets},-1}$ is also natural. Thus $\lambda^{\mathsf{Sets},\amalg}$ is a natural isomorphism.   □

5.2.5   The Right Unitor

Definition 5.2.5.1.1. The **right unitor of the coproduct of sets** is the natural isomorphism

$$\rho^{\mathsf{Sets},\amalg} : \amalg \circ (\mathrm{id} \times \emptyset^{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \boldsymbol{\rho}^{\mathsf{Cats}_2}_{\mathsf{Sets}},$$

whose component

$$\rho^{\mathsf{Sets},\amalg}_X : X \amalg \emptyset \overset{\sim}{\dashrightarrow} X$$

at X is given by

$$\rho^{\mathsf{Sets},\amalg}_X((0,x)) \overset{\mathrm{def}}{=} x$$

for each $(0,x) \in X \amalg \emptyset$.

Proof. *Unwinding the Definition of $X \amalg \emptyset$:* Firstly, we unwind the expression for $X \amalg \emptyset$. We have

$$
\begin{aligned}
X \amalg \emptyset &\overset{\mathrm{def}}{=} \{(0,x) \in S \mid x \in X\} \cup \{(1,z) \in S \mid z \in \emptyset\} \\
&= \{(0,x) \in S \mid x \in X\} \cup \emptyset \\
&= \{(0,x) \in S \mid x \in X\},
\end{aligned}
$$

where $S = \{0,1\} \times (X \cup \emptyset) = \{0,1\} \times (\emptyset \cup X) = S$.
Invertibility: The inverse of $\rho^{\mathsf{Sets},\amalg}_X$ is the map

$$\rho^{\mathsf{Sets},\amalg,-1}_X : X \to X \amalg \emptyset$$

given by

$$\rho_X^{\mathsf{Sets},\coprod,-1}(x) \overset{\text{def}}{=} (0, x)$$

for each $x \in X$. Indeed:

- *Invertibility I.* We have

$$
\begin{aligned}
[\rho_X^{\mathsf{Sets},\coprod,-1} \circ \rho_X^{\mathsf{Sets},\coprod}](0, x) &= \rho_X^{\mathsf{Sets},\coprod,-1}(\rho_X^{\mathsf{Sets},\coprod}(0, x)) \\
&= \rho_X^{\mathsf{Sets},\coprod,-1}(x) \\
&= (0, x) \\
&= [\mathrm{id}_{X\coprod\emptyset}](0, x)
\end{aligned}
$$

  for each $(0, x) \in \emptyset \coprod X$, and therefore we have

$$\rho_X^{\mathsf{Sets},\coprod,-1} \circ \rho_X^{\mathsf{Sets},\coprod} = \mathrm{id}_{\emptyset\coprod X}\,.$$

- *Invertibility II.* We have

$$
\begin{aligned}
[\rho_X^{\mathsf{Sets},\coprod} \circ \rho_X^{\mathsf{Sets},\coprod,-1}](x) &= \rho_X^{\mathsf{Sets},\coprod}(\rho_X^{\mathsf{Sets},\coprod,-1}(x)) \\
&= \rho_X^{\mathsf{Sets},\coprod,-1}(0, x) \\
&= x \\
&= [\mathrm{id}_X](x)
\end{aligned}
$$

  for each $x \in X$, and therefore we have

$$\rho_X^{\mathsf{Sets},\coprod} \circ \rho_X^{\mathsf{Sets},\coprod,-1} = \mathrm{id}_X\,.$$

Therefore $\rho_X^{\mathsf{Sets},\coprod}$ is indeed an isomorphism.

Naturality: We need to show that, given a function $f\colon X \to Y$, the diagram

$$
\begin{CD}
X \coprod \emptyset @>{f\coprod \mathrm{id}_\emptyset}>> Y \coprod \emptyset \\
@V{\rho_X^{\mathsf{Sets},\coprod}}VV @VV{\rho_Y^{\mathsf{Sets},\coprod}}V \\
X @>>{f}> Y
\end{CD}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
(0,x) & & (0,x) \longmapsto (1,f(x)) \\
\Big\downarrow & & \Big\downarrow \\
x \longmapsto f(x) & & f(x)
\end{array}
$$

and hence indeed commutes. Therefore $\rho^{\mathsf{Sets},\coprod}$ is a natural transformation.
Being a Natural Isomorphism: Since $\rho^{\mathsf{Sets},\coprod}$ is natural and $\rho^{\mathsf{Sets},-1}$ is a component-wise inverse to $\rho^{\mathsf{Sets},\coprod}$, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\rho^{\mathsf{Sets},-1}$ is also natural. Thus $\rho^{\mathsf{Sets},\coprod}$ is a natural isomorphism. $\square$

5.2.6   The Symmetry

Definition 5.2.6.1.1. The **symmetry of the coproduct of sets** is the natural isomorphism

$$
\sigma^{\mathsf{Sets},\coprod} : \coprod \xRightarrow{\sim} \coprod \circ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}},
$$

$$
\begin{array}{ccc}
\mathsf{Sets} \times \mathsf{Sets} & \xrightarrow{\coprod} & \mathsf{Sets}, \\
{\scriptstyle \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}}} \searrow & {\overset{\parallel}{\underset{\Downarrow}{\sigma^{\mathsf{Sets},\coprod}}}} & \nearrow {\scriptstyle \coprod} \\
& \mathsf{Sets} \times \mathsf{Sets} &
\end{array}
$$

whose component

$$
\sigma^{\mathsf{Sets},\coprod}_{X,Y} : X \coprod Y \xdashrightarrow{\sim} Y \coprod X
$$

at $X, Y \in \mathrm{Obj}(\mathsf{Sets})$ is defined by

$$
\sigma^{\mathsf{Sets},\coprod}_{X,Y}(x,y) \overset{\mathrm{def}}{=} (y,x)
$$

for each $(x,y) \in X \times Y$.

Proof.   Unwinding the Definitions of $X \coprod Y$ and $Y \coprod X$: Firstly, we unwind the expressions for $X \coprod Y$ and $Y \coprod X$. We have

$$
X \coprod Y \overset{\mathrm{def}}{=} \{(0,x) \in S \mid x \in X\} \cup \{(1,y) \in S \mid y \in Y\},
$$

where $S = \{0,1\} \times (X \cup Y)$ and

$$
Y \coprod X \overset{\mathrm{def}}{=} \{(0,y) \in S' \mid y \in Y\} \cup \{(1,x) \in S' \mid x \in X\},
$$

where $S' = \{0, 1\} \times (Y \cup X) = \{0, 1\} \times (X \cup Y) = S$.

Invertibility: The inverse of $\sigma_{X,Y}^{\mathsf{Sets},\coprod}$ is the map

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1} \colon Y \coprod X \to X \coprod Y$$

defined by

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1} \overset{\text{def}}{=} \sigma_{Y,X}^{\mathsf{Sets},\coprod}$$

and hence given by

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1}(z) \overset{\text{def}}{=} \begin{cases} (0, x) & \text{if } z = (1, x), \\ (1, y) & \text{if } z = (0, y) \end{cases}$$

for each $z \in Y \coprod X$. Indeed:

- *Invertibility I.* We have

$$\begin{aligned}
[\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1} \circ \sigma_{X,Y}^{\mathsf{Sets},\coprod}](0, x) &= \sigma_X^{\mathsf{Sets},\coprod,-1}(\sigma_X^{\mathsf{Sets},\coprod}(0, x)) \\
&= \sigma_X^{\mathsf{Sets},\coprod,-1}(1, x) \\
&= (0, x) \\
&= [\mathsf{id}_{X\coprod Y}](0, x)
\end{aligned}$$

  for each $(0, x) \in X \coprod Y$ and

$$\begin{aligned}
[\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1} \circ \sigma_{X,Y}^{\mathsf{Sets},\coprod}](1, y) &= \sigma_X^{\mathsf{Sets},\coprod,-1}(\sigma_X^{\mathsf{Sets},\coprod}(1, y)) \\
&= \sigma_X^{\mathsf{Sets},\coprod,-1}(0, y) \\
&= (1, y) \\
&= [\mathsf{id}_{X\coprod Y}](1, y)
\end{aligned}$$

  for each $(1, y) \in X \coprod Y$, and therefore we have

$$\sigma_{X,Y}^{\mathsf{Sets},\coprod,-1} \circ \sigma_{X,Y}^{\mathsf{Sets},\coprod} = \mathsf{id}_{X\coprod Y}.$$

- *Invertibility II.* We have

$$[\sigma_{X,Y}^{\mathsf{Sets},\coprod} \circ \sigma_{X,Y}^{\mathsf{Sets},\coprod,-1}](0, y) = \sigma_X^{\mathsf{Sets},\coprod}(\sigma_X^{\mathsf{Sets},\coprod,-1}(0, y))$$

$$= \sigma_X^{\mathsf{Sets},\coprod,-1}(1, y)$$
$$= (0, y)$$
$$= [\mathsf{id}_{Y \coprod X}](0, y)$$

for each $(0, y) \in Y \coprod X$ and

$$[\sigma_{X,Y}^{\mathsf{Sets},\coprod} \circ \sigma_{X,Y}^{\mathsf{Sets},\coprod,-1}](1, x) = \sigma_X^{\mathsf{Sets},\coprod}(\sigma_X^{\mathsf{Sets},\coprod,-1}(1, x))$$
$$= \sigma_X^{\mathsf{Sets},\coprod,-1}(0, x)$$
$$= (1, x)$$
$$= [\mathsf{id}_{Y \coprod X}](1, x)$$

for each $(1, x) \in Y \coprod X$, and therefore we have

$$\sigma_X^{\mathsf{Sets},\coprod} \circ \sigma_X^{\mathsf{Sets},\coprod,-1} = \mathsf{id}_{Y \coprod X} \, .$$

Therefore $\sigma_{X,Y}^{\mathsf{Sets},\coprod}$ is indeed an isomorphism.
Naturality: We need to show that, given functions $f : A \to X$ and $g : B \to Y$, the diagram

$$
\begin{array}{ccc}
A \coprod B & \xrightarrow{f \coprod g} & X \coprod Y \\
\sigma_{A,B}^{\mathsf{Sets},\coprod} \downarrow & & \downarrow \sigma_{X,Y}^{\mathsf{Sets},\coprod} \\
B \coprod A & \xrightarrow[g \coprod f]{} & Y \coprod X
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
(0, a) & & (0, a) \longmapsto (0, f(a)) \\
\downarrow & & \downarrow \\
(1, a) \longmapsto (1, f(a)) & & (1, f(a))
\end{array}
$$

$$
\begin{array}{ccc}
(1, b) & & (1, b) \longmapsto (1, g(b)) \\
\downarrow & & \downarrow \\
(0, b) \longmapsto (0, g(b)) & & (0, g(b))
\end{array}
$$

and hence indeed commutes. Therefore $\sigma^{\mathsf{Sets},\coprod}$ is a natural transformation.
Being a Natural Isomorphism: Since $\sigma^{\mathsf{Sets},\coprod}$ is natural and $\sigma^{\mathsf{Sets},-1}$ is a component-wise inverse to $\sigma^{\mathsf{Sets},\coprod}$, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\sigma^{\mathsf{Sets},-1}$ is also natural. Thus $\sigma^{\mathsf{Sets},\coprod}$ is a natural isomorphism. □

5.2.7   The Monoidal Category of Sets and Coproducts

Proposition 5.2.7.1.1. The category Sets admits a closed symmetric monoidal category structure consisting of:

  · *The Underlying Category.* The category Sets of pointed sets.

  · *The Monoidal Product.* The coproduct functor

$$\coprod \colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}$$

  of Constructions With Sets, Item 1 of Definition 4.2.3.1.3.

  · *The Monoidal Unit.* The functor

$$\mathbb{0}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

  of Definition 5.2.2.1.1.

  · *The Associators.* The natural isomorphism

$$\alpha^{\mathsf{Sets},\coprod} \colon \coprod \circ (\coprod \times \mathrm{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \coprod \circ (\mathrm{id}_{\mathsf{Sets}} \times \coprod) \circ \boldsymbol{\alpha}^{\mathsf{Cats}}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}$$

  of Definition 5.2.3.1.1.

  · *The Left Unitors.* The natural isomorphism

$$\lambda^{\mathsf{Sets},\coprod} \colon \coprod \circ (\mathbb{0}^{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \lambda^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

  of Definition 5.2.4.1.1.

  · *The Right Unitors.* The natural isomorphism

$$\rho^{\mathsf{Sets},\coprod} \colon \coprod \circ (\mathrm{id} \times \mathbb{0}^{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \boldsymbol{\rho}^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

  of Definition 5.2.5.1.1.

· *The Symmetry.* The natural isomorphism

$$\sigma^{\mathsf{Sets},\coprod} : \times \overset{\sim}{\implies} \times \circ \, \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}}$$

of Definition 5.2.6.1.1.

Proof.   *The Pentagon Identity*: Let W, X, Y and Z be sets. We have to show that the diagram

$$(W \coprod (X \coprod Y)) \coprod Z$$

with maps $\alpha^{\mathsf{Sets},\coprod}_{W,X,Y} \coprod \mathrm{id}_Z$ and $\alpha^{\mathsf{Sets},\coprod}_{W,X\coprod Y,Z}$

$$((W \coprod X) \coprod Y) \coprod Z \qquad\qquad W \coprod ((X \coprod Y) \coprod Z)$$

with maps $\alpha^{\mathsf{Sets},\coprod}_{W\coprod X,Y,Z}$ and $\mathrm{id}_W \coprod \alpha^{\mathsf{Sets},\coprod}_{X,Y,Z}$

$$(W \coprod X) \coprod (Y \coprod Z) \xrightarrow{\;\;\alpha^{\mathsf{Sets},\coprod}_{W,X,Y\coprod Z}\;\;} W \coprod (X \coprod (Y \coprod Z))$$

commutes. Indeed, this diagram acts on elements as

$$(0, (0, w))$$

$$(0, (0, (0, w)))\qquad\qquad (0, (0, (0, w)))\qquad\qquad (0, w)$$

$$(0, (0, w)) \longmapsto (0, w)\qquad\qquad\qquad\qquad (0, w)$$

$$(0, (1, (0, x)))$$

$$(0, (0, (1, x)))$$

$$(0, (0, (1, x)))\qquad\qquad (1, (0, (0, x)))$$

$$(0, (1, x)) \longmapsto (1, (0, x))$$

$$(1, (0, x))$$

$$(0, (1, (1, y)))$$

$$(0, (1, y))$$

$$(0, (1, y))\qquad\qquad (1, (0, (1, y)))$$

$$(1, (0, y)) \longmapsto (1, (1, (0, y)))$$

$$(1, (1, (0, y)))$$

$$(1, z)$$

$$(1, z)$$

$$(1, z)\qquad\qquad (1, (1, z))$$

$$(1, (1, z)) \longmapsto (1, (1, (1, z)))$$

$$(1, (1, (1, z)))$$

and therefore the pentagon identity is satisfied.

The Triangle Identity: Let X and Y be sets. We have to show that the diagram

$$
(X \coprod \emptyset) \coprod Y \xrightarrow{\alpha^{\mathsf{Sets},\coprod}_{X,\emptyset,Y}} X \coprod (\emptyset \coprod Y)
$$

$$
\rho^{\mathsf{Sets},\coprod}_{X} \coprod \mathrm{id}_Y \qquad\qquad \mathrm{id}_X \coprod \lambda^{\mathsf{Sets},\coprod}_{Y}
$$

$$
X \coprod Y
$$

commutes. Indeed, this diagram acts on elements as

$$
(0,(1,x)) \qquad\qquad\qquad (1,(0,x)) \longmapsto (0,x)
$$

$$
(0,x) \qquad\qquad\qquad\qquad (0,x)
$$

$$
(1,y) \qquad\qquad\qquad\qquad (1,y) \longmapsto (1,(1,y))
$$

$$
(1,y) \qquad\qquad\qquad\qquad (1,y)
$$

and therefore the triangle identity is satisfied.

The Left Hexagon Identity: Let X, Y, and Z be sets. We have to show that the diagram

$$
(X \coprod Y) \coprod Z
$$

$$
\alpha^{\mathsf{Sets}}_{X,Y,Z} \qquad\qquad \sigma^{\mathsf{Sets}}_{X,Y} \coprod \mathrm{id}_Z
$$

$$
X \coprod (Y \coprod Z) \qquad\qquad\qquad (Y \coprod X) \coprod Z
$$

$$
\sigma^{\mathsf{Sets}}_{X,Y} \coprod Z \qquad\qquad\qquad\qquad \alpha^{\mathsf{Sets}}_{Y,X,Z}
$$

$$
(Y \coprod Z) \coprod X \qquad\qquad\qquad Y \coprod (X \coprod Z)
$$

$$
\alpha^{\mathsf{Sets}}_{Y,Z,X} \qquad\qquad \mathrm{id}_Y \coprod \sigma^{\mathsf{Sets}}_{X,Z}
$$

$$
Y \coprod (Z \coprod X)
$$

commutes. Indeed, this diagram acts on elements as

$$(0, (0, x)) \qquad (0, (0, x))$$

$$(0, x) \qquad\qquad (0, (1, x))$$

$$(1, x) \qquad\qquad (1, (0, x))$$

$$(1, (1, x)) \quad (1, (1, x))$$

$$(0, (1, y)) \qquad (0, (1, y))$$

$$(1, (0, y)) \qquad\qquad (0, (0, y))$$

$$(0, (0, y)) \qquad\qquad (0, y)$$

$$(0, y) \qquad (0, y)$$

$$(1, z) \qquad (1, z)$$

$$(1, (1, z)) \qquad\qquad (1, z)$$

$$(0, (1, z)) \qquad\qquad (1, (1, z))$$
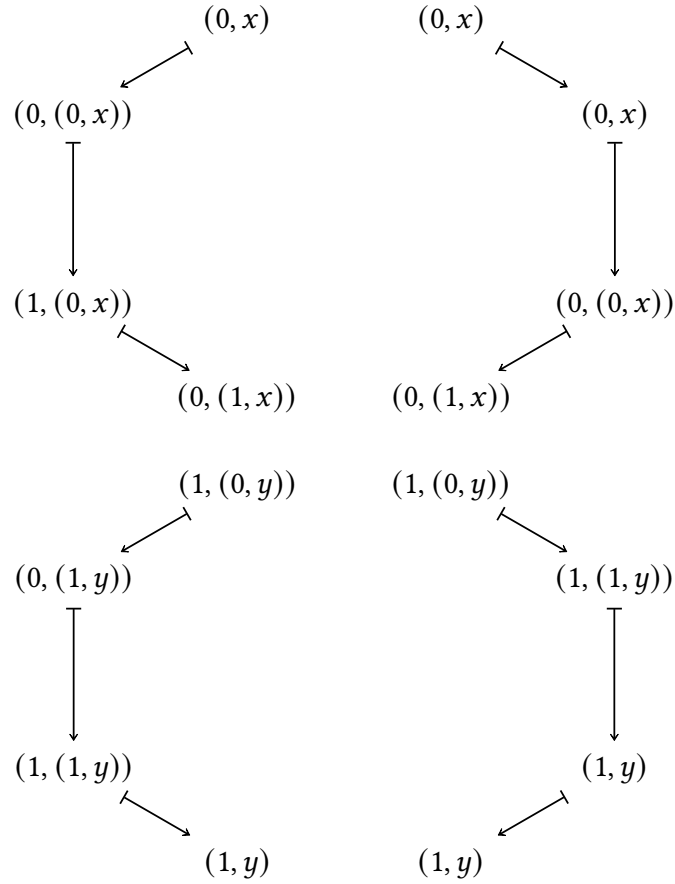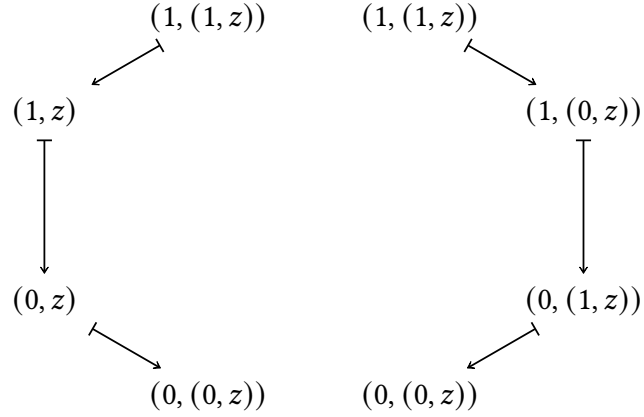
$$(1, (0, z)) \quad (1, (0, z))$$

and thus the left hexagon identity is satisfied.

The Right Hexagon Identity: Let X, Y, and Z be sets. We have to show that the diagram

$$
\begin{array}{ccc}
 & X \coprod (Y \coprod Z) & \\
{}^{(\alpha^{\mathsf{Sets}}_{X,Y,Z})^{-1}} \swarrow & & \searrow {}^{\mathrm{id}_X \coprod \sigma^{\mathsf{Sets}}_{Y,Z}} \\
(X \coprod Y) \coprod Z & & X \coprod (Z \coprod Y) \\
\Big\downarrow {}^{\sigma^{\mathsf{Sets}}_{X \coprod Y,Z}} & & \Big\downarrow {}^{(\alpha^{\mathsf{Sets}}_{X,Z,Y})^{-1}} \\
Z \coprod (X \coprod Y) & & (X \coprod Z) \coprod Y \\
{}^{(\alpha^{\mathsf{Sets}}_{Z,X,Y})^{-1}} \searrow & & \swarrow {}^{\sigma^{\mathsf{Sets}}_{X,Z} \coprod \mathrm{id}_Y} \\
 & (Z \coprod X) \coprod Y &
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
 (0,x) & \qquad & (0,x) \\
 \swarrow & & \nwarrow \\
(0,(0,x)) & & (0,x) \\
\big\downarrow & & \big\downarrow \\
(1,(0,x)) & & (0,(0,x)) \\
 \nwarrow & & \nearrow \\
(0,(1,x)) & & (0,(1,x)) \\
 & & \\
(1,(0,y)) & & (1,(0,y)) \\
 \nearrow & & \nwarrow \\
(0,(1,y)) & & (1,(1,y)) \\
\big\downarrow & & \big\downarrow \\
(1,(1,y)) & & (1,y) \\
 \searrow & & \nwarrow \\
(1,y) & & (1,y)
\end{array}
$$

$$(1, (1, z)) \qquad (1, (1, z))$$

$$(1, z) \qquad\qquad\qquad\qquad (1, (0, z))$$

$$(0, z) \qquad\qquad\qquad\qquad (0, (1, z))$$

$$(0, (0, z)) \qquad (0, (0, z))$$

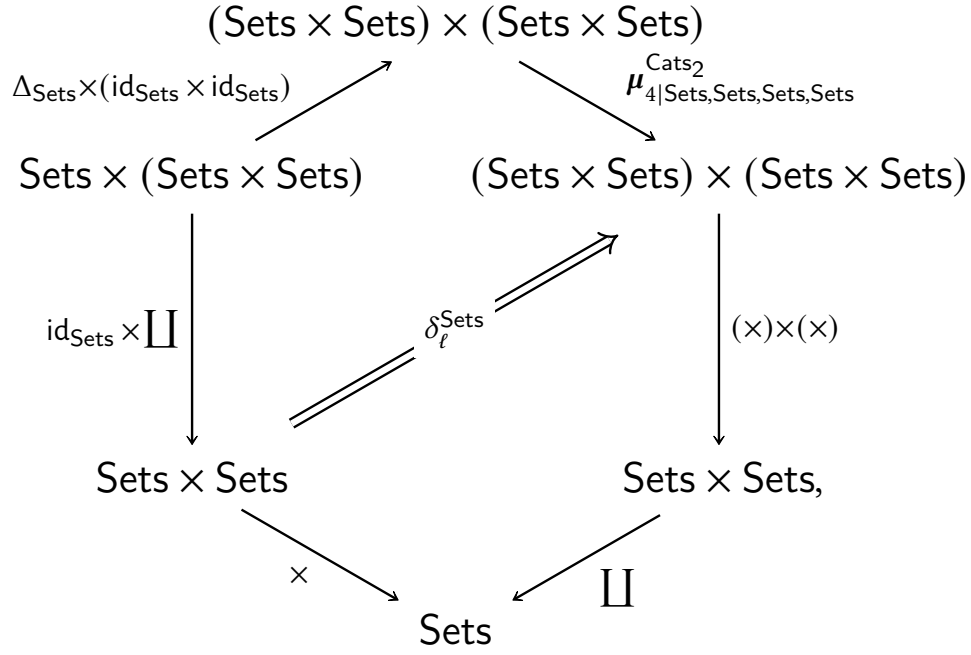and thus the right hexagon identity is satisfied. $\qquad\qquad\qquad\square$

5.3 The Bimonoidal Category of Sets, Products, and Coproducts

5.3.1 The Left Distributor

Definition 5.3.1.1.1. The **left distributor of the product of sets over the coproduct of sets** is the natural isomorphism

$$\delta_\ell^{\mathsf{Sets}} : \times \circ (\mathrm{id}_{\mathsf{Sets}} \times \coprod) \overset{\sim}{\Longrightarrow} \coprod \circ ((\times) \times (\times)) \circ \mu_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}^{\mathsf{Cats}_2} \circ (\Delta_{\mathsf{Sets}} \times (\mathrm{id}_{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}))$$

as in the diagram

$$(\mathsf{Sets} \times \mathsf{Sets}) \times (\mathsf{Sets} \times \mathsf{Sets})$$

$$\Delta_{\mathsf{Sets}} \times (\mathrm{id}_{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}) \qquad\qquad \mu^{\mathsf{Cats}_2}_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}$$

$$\mathsf{Sets} \times (\mathsf{Sets} \times \mathsf{Sets}) \qquad\qquad (\mathsf{Sets} \times \mathsf{Sets}) \times (\mathsf{Sets} \times \mathsf{Sets})$$

$$\mathrm{id}_{\mathsf{Sets}} \times \coprod \qquad\qquad \delta^{\mathsf{Sets}}_{\ell} \qquad\qquad (\times) \times (\times)$$

$$\mathsf{Sets} \times \mathsf{Sets} \qquad\qquad\qquad \mathsf{Sets} \times \mathsf{Sets},$$

$$\times \qquad\qquad \coprod$$

$$\mathsf{Sets}$$

whose component

$$\delta^{\mathsf{Sets}}_{\ell|X,Y,Z} \colon X \times (Y \textstyle\coprod Z) \xdashrightarrow{\sim} (X \times Y) \coprod (X \times Z)$$

at (X, Y, Z) is defined by

$$\delta^{\mathsf{Sets}}_{\ell|X,Y,Z}(x, a) \overset{\text{def}}{=} \begin{cases} (0, (x, y)) & \text{if } a = (0, y), \\ (1, (x, z)) & \text{if } a = (1, z) \end{cases}$$

for each $(x, a) \in X \times (Y \coprod Z)$.

Proof. *Invertibility*: The inverse of $\delta^{\mathsf{Sets}}_{\ell|X,Y,Z}$ is the map

$$\delta^{\mathsf{Sets},-1}_{\ell|X,Y,Z} \colon (X \times Y) \coprod (X \times Z) \xdashrightarrow{\sim} X \times (Y \textstyle\coprod Z)$$

given by

$$\delta^{\mathsf{Sets},-1}_{\ell|X,Y,Z}(a) \overset{\text{def}}{=} \begin{cases} (x, (0, y)) & \text{if } a = (0, (x, y)), \\ (x, (1, z)) & \text{if } a = (1, (x, z)) \end{cases}$$

for $a \in (X \times Y) \coprod (X \times Z)$. Indeed:

· *Invertibility I.* The map $\delta^{\mathsf{Sets},-1}_{\ell|X,Y,Z} \circ \delta^{\mathsf{Sets}}_{\ell|X,Y,Z}$ acts on elements as

$$(x,(0,y)) \;\mapsto\; (0,(x,y)) \;\mapsto\; (x,(0,y)),$$
$$(x,(1,z)) \;\mapsto\; (1,(x,z)) \;\mapsto\; (x,(1,z)),$$

but these are the two possible cases for elements of $X \times (Y \coprod Z)$. Hence the map is equal to the identity.

· *Invertibility II.* The map $\delta^{\mathsf{Sets}}_{\ell|X,Y,Z} \circ \delta^{\mathsf{Sets},-1}_{\ell|X,Y,Z}$ acts on elements as

$$(0,(x,y)) \;\mapsto\; (x,(0,y)) \;\mapsto\; (0,(x,y)),$$
$$(1,(x,z)) \;\mapsto\; (x,(1,z)) \;\mapsto\; (1,(x,z)),$$

but these are the two possible cases for elements of $(X \times Y) \coprod (X \times Z)$. Hence the map is equal to the identity.

Thus $\delta^{\mathsf{Sets}}_{\ell|X,Y,Z}$ is an isomorphism for all X, Y, Z.
Naturality: We need to show that, given functions

$$f\colon X \to X',$$
$$g\colon Y \to Y',$$
$$h\colon Z \to Z'$$

the diagram

$$
\begin{array}{ccc}
X \times (Y \coprod Z) & \xrightarrow{\;f \times \left(g \coprod h\right)\;} & X' \times (Y' \coprod Z') \\
\Big\downarrow{\scriptstyle \delta^{\mathsf{Sets}}_{\ell|X,Y,Z}} & & \Big\downarrow{\scriptstyle \delta^{\mathsf{Sets}}_{\ell|X',Y',Z'}} \\
(X \times Y) \coprod (X \times Z) & \xrightarrow[\;(f \times g) \coprod (f \times h)\;]{} & (X' \times Y') \coprod (X' \times Z')
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{ccc}
(x,(0,y)) & & (x,(0,y)) \longmapsto (f(x),(0,f(y))) \\
\Big\downarrow & & \Big\downarrow \\
(0,(x,y)) \longmapsto (0,(f(x),g(y))) & & (0,(f(x),g(y)))
\end{array}
$$

$$(x, (1, z))$$

$$(x, (1, z)) \longmapsto (f(x), (1, h(z)))$$

$$(1, (x, z)) \longmapsto (1, (f(x), h(z)))$$

$$(1, (f(x), h(z))),$$

so it commutes, showing $\delta_\ell^{\mathsf{Sets}}$ to be a natural transformation.

Being a Natural Isomorphism: Since $\delta_\ell^{\mathsf{Sets}}$ is natural and $\delta_\ell^{\mathsf{Sets},-1}$ is a componentwise inverse to $\delta_\ell^{\mathsf{Sets}}$, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\delta_\ell^{\mathsf{Sets},-1}$ is also natural. Thus $\delta_\ell^{\mathsf{Sets}}$ is a natural isomorphism. □

5.3.2 The Right Distributor

Definition 5.3.2.1.1. The **right distributor of the product of sets over the coproduct of sets** is the natural isomorphism

$$\delta_r^{\mathsf{Sets}} \colon \times \circ (\coprod \times \mathrm{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \coprod \circ ((\times) \times (\times)) \circ \mu_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}^{\mathsf{Cats}_2} \circ ((\mathrm{id}_{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}) \times \Delta_{\mathsf{Sets}})$$

as in the diagram

whose component

$$\delta_{r|X,Y,Z}^{\mathsf{Sets}} \colon (X \coprod Y) \times Z \overset{\sim}{\dashrightarrow} (X \times Z) \coprod (Y \times Z)$$

at (X, Y, Z) is defined by

$$\delta^{\mathsf{Sets}}_{r|X,Y,Z}(a, z) \overset{\text{def}}{=} \begin{cases} (0, (x, z)) & \text{if } a = (0, x), \\ (1, (y, z)) & \text{if } a = (1, y) \end{cases}$$

for each $(a, z) \in (X \amalg Y) \times Z$.

Proof. *Invertibility*: The inverse of $\delta^{\mathsf{Sets}}_{r|X,Y,Z}$ is the map

$$\delta^{\mathsf{Sets},-1}_{r|X,Y,Z} \colon (X \times Z) \amalg (Y \times Z) \overset{\sim}{\dashrightarrow} (X \amalg Y) \times Z$$

given by

$$\delta^{\mathsf{Sets},-1}_{r|X,Y,Z}(a) \overset{\text{def}}{=} \begin{cases} ((0, x), z) & \text{if } a = (0, (x, z)), \\ ((1, y), z) & \text{if } a = (1, (y, z)) \end{cases}$$

for $a \in (X \times Z) \amalg (Y \times Z)$. Indeed:

- *Invertibility I.* The map $\delta^{\mathsf{Sets},-1}_{r|X,Y,Z} \circ \delta^{\mathsf{Sets}}_{r|X,Y,Z}$ acts on elements as

$$\begin{aligned} ((0, x), z) &\mapsto (0, (x, z)) \mapsto (0, (x, z)), \\ ((1, y), z) &\mapsto (1, (y, z)) \mapsto (1, (y, z)), \end{aligned}$$

  but these are the two possible cases for elements of $(X \amalg Y) \times Z$. Hence the map is equal to the identity.

- *Invertibility II.* The map $\delta^{\mathsf{Sets}}_{r|X,Y,Z} \circ \delta^{\mathsf{Sets},-1}_{r|X,Y,Z}$ acts on elements as

$$\begin{aligned} (0, (x, z)) &\mapsto ((0, x), z) \mapsto (0, (x, z)), \\ (1, (y, z)) &\mapsto ((1, y), z) \mapsto (1, (y, z)), \end{aligned}$$

  but these are the two possible cases for elements of $(X \times Z) \amalg (Y \times Z)$. Hence the map is equal to the identity.

So $\delta^{\mathsf{Sets}}_{r|X,Y,Z}$ is an isomorphism for all X, Y, Z.
Naturality: We need to show that, given functions

$$\begin{aligned} f &\colon X \to X', \\ g &\colon Y \to Y', \\ h &\colon Z \to Z' \end{aligned}$$

the diagram

$$
\begin{array}{ccc}
(X \amalg Y) \times Z' & \xrightarrow{\;(f \amalg g) \times h\;} & (X' \amalg Y') \times Z' \\
{\scriptstyle \delta^{\mathsf{Sets}}_{r|X,Y,Z}} \Big\downarrow & & \Big\downarrow {\scriptstyle \delta^{\mathsf{Sets}}_{r|X',Y',Z'}} \\
(X \times Z) \amalg (Y \times Z) & \xrightarrow[\;(f \times h) \amalg (g \times h)\;]{} & (X' \times Z') \amalg (Y' \times Z')
\end{array}
$$

commutes. Indeed, this diagram acts on elements as

$$
\begin{array}{lll}
((0,x),z) & & ((0,x),z) \longmapsto ((0,f(x)),h(z)) \\
\Big\downarrow & & \qquad\qquad\qquad\quad \Big\downarrow \\
(0,(x,z)) \longmapsto (0,(f(x),h(z))) & & (0,(f(x),h(z)))
\end{array}
$$

$$
\begin{array}{lll}
((1,y),z) & & ((1,y),z) \longmapsto ((1,g(y)),h(z)) \\
\Big\downarrow & & \qquad\qquad\qquad\quad \Big\downarrow \\
(1,(y,z)) \longmapsto (1,(g(y),h(z))) & & (1,(g(y),h(z)))
\end{array}
$$

so it commutes and δ^{Sets}_r is a natural transformation.

Being a Natural Isomorphism: Since δ^{Sets}_r is natural and $\delta^{\mathsf{Sets},-1}_r$ is a componentwise inverse to δ^{Sets}_r, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\delta^{\mathsf{Sets},-1}_r$ is also natural. Thus δ^{Sets}_r is a natural isomorphism. $\qquad\square$

5.3.3   The Left Annihilator

Definition 5.3.3.1.1. The **left annihilator of the product of sets** is the natural isomorphism

$$
\zeta^{\mathsf{Sets}}_\ell \colon \mathbb{0}^{\mathsf{Sets}} \circ \mu^{\mathsf{Cats}_2}_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}} \circ (\mathrm{id}_{\mathsf{pt}} \times \epsilon^{\mathsf{Cats}_2}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \times \circ (\mathbb{0}^{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}})
$$

as in the diagram

$$
\begin{array}{ccc}
& \text{pt} \times \text{pt} & \\
{\scriptstyle \text{id}_{\text{pt}} \times \epsilon_{\text{Sets}}^{\text{Cats}_2}} \nearrow & & \searrow {\scriptstyle \mu_{4|\text{Sets},\text{Sets},\text{Sets},\text{Sets}}^{\text{Cats}_2}} \\
\text{pt} \times \text{Sets} & \zeta_\ell^{\text{Sets}} & \text{pt} \\
{\scriptstyle \mathbb{0}^{\text{Sets}} \times \text{id}_{\text{Sets}}} \downarrow & \Downarrow & \downarrow {\scriptstyle \mathbb{0}^{\text{Sets}}} \\
\text{Sets} \times \text{Sets} & \xrightarrow[\times]{} & \text{Sets}
\end{array}
$$
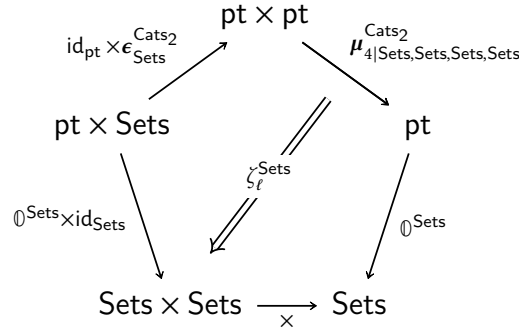
with components

$$
\zeta_{\ell|A}^{\text{Sets}} : \varnothing \times A \overset{\sim}{\dashrightarrow} \varnothing
$$

given by $\zeta_{\ell|A}^{\text{Sets}} \overset{\text{def}}{=} \text{pr}_1$.

Proof. *Invertibility*: The inverse of $\zeta_{\ell|A}^{\text{Sets}}$ is the map

$$
\zeta_{\ell|A}^{\text{Sets},-1} : \varnothing \overset{\sim}{\dashrightarrow} \varnothing \times A
$$

given by

$$
\zeta_{\ell|A}^{\text{Sets},-1} \overset{\text{def}}{=} \iota_A,
$$

where ι_A is as defined in Constructions With Sets, Definition 4.2.1.1.2:

· *Invertibility I*. The map $\zeta_{\ell|A}^{\text{Sets}} \circ \iota_A : \varnothing \to \varnothing$ is equal to $\text{id}_\varnothing$, as $\varnothing$ is the initial object of Sets.

· *Invertibility II*. The map $\iota_A \circ \zeta_{\ell|A}^{\text{Sets}}$ is equal to the identity on every $(x, a) \in \varnothing \times A$, of which there are none.

Hence $\zeta_{\ell|A}^{\text{Sets}}$ is an isomorphism.
Naturality: We need to show that given a function $f : A \to B$, the diagram

$$
\begin{array}{ccc}
\varnothing \times A & \xrightarrow{\text{id}_\varnothing \times f} & \varnothing \times B \\
{\scriptstyle \zeta_{\ell|A}^{\text{Sets}}} \downarrow & & \downarrow {\scriptstyle \zeta_{\ell|B}^{\text{Sets}}} \\
\varnothing & \xrightarrow[\text{id}_\varnothing]{} & \varnothing
\end{array}
$$

commutes. But since $\emptyset \times A$ has no elements, this is trivially true.

Being a Natural Isomorphism: Since $\zeta_\ell^{\mathsf{Sets}}$ is natural and $\zeta_\ell^{\mathsf{Sets},-1}$ is a componentwise inverse to $\zeta_\ell^{\mathsf{Sets}}$, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\zeta_\ell^{\mathsf{Sets},-1}$ is also natural. Thus $\zeta_\ell^{\mathsf{Sets}}$ is a natural isomorphism. $\qquad\square$

5.3.4   The Right Annihilator

Definition 5.3.4.1.1. The **right annihilator of the product of sets** is the natural isomorphism

$$\zeta_r^{\mathsf{Sets}} \colon \mathbb{0}^{\mathsf{Sets}} \circ \mu_{4|\mathsf{Sets,Sets,Sets,Sets}}^{\mathsf{Cats}_2} \circ (\epsilon_{\mathsf{Sets}}^{\mathsf{Cats}_2} \times \mathrm{id}_{\mathsf{pt}}) \overset{\sim}{\dashrightarrow} \times \circ (\mathrm{id}_{\mathsf{Sets}} \times \mathbb{0}^{\mathsf{Sets}})$$

as in the diagram
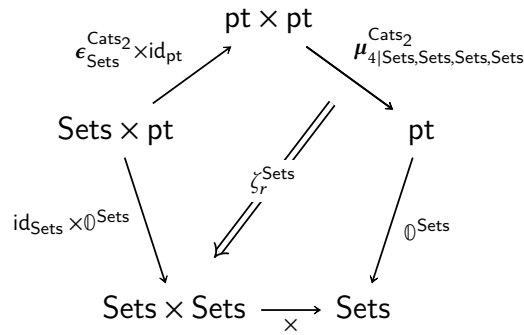
with components

$$\zeta_{r|A}^{\mathsf{Sets}} \colon A \times \emptyset \overset{\sim}{\dashrightarrow} \emptyset$$

given by $\zeta_{r|A}^{\mathsf{Sets}} \overset{\mathrm{def}}{=} \mathrm{pr}_2$.

Proof. *Invertibility*: The inverse of $\zeta_{r|A}^{\mathsf{Sets}}$ is the map

$$\zeta_{r|A}^{\mathsf{Sets},-1} \colon \emptyset \overset{\sim}{\dashrightarrow} A \times \emptyset$$

given by

$$\zeta_{r|A}^{\mathsf{Sets},-1} \overset{\mathrm{def}}{=} \iota_A,$$

where ι_A is as defined in Constructions With Sets, Definition 4.2.1.1.2:

- *Invertibility I*. The map $\zeta_{r|A}^{\mathsf{Sets}} \circ \iota_A \colon \emptyset \to \emptyset$ is equal to $\mathrm{id}_\emptyset$, as $\emptyset$ is the initial object of Sets.

- *Invertibility II*. The map $\iota_A \circ \zeta_{r|A}^{\mathsf{Sets}}$ is equal to the identity on every $(a, x) \in A \times \emptyset$, of which there are none.

Hence $\zeta_{r|A}^{\mathsf{Sets}}$ is an isomorphism.
Naturality: We need to show that given a function $f \colon A \to B$, the diagram

$$
\begin{array}{ccc}
A \times \emptyset & \xrightarrow{f \times \mathrm{id}_\emptyset} & B \times \emptyset \\
\downarrow{\scriptstyle \zeta_{r|A}^{\mathsf{Sets}}} & & \downarrow{\scriptstyle \zeta_{r|B}^{\mathsf{Sets}}} \\
\emptyset & \xrightarrow[\mathrm{id}_\emptyset]{} & \emptyset
\end{array}
$$

commutes. But since $A \times \emptyset$ has no elements, this is trivially true.
Being a Natural Isomorphism: Since ζ_r^{Sets} is natural and $\zeta_r^{\mathsf{Sets},-1}$ is a componentwise inverse to ζ_r^{Sets}, it follows from Categories, Item 2 of Definition 11.9.7.1.2 that $\zeta_r^{\mathsf{Sets},-1}$ is also natural. Thus ζ_r^{Sets} is a natural isomorphism.   □

5.3.5   The Bimonoidal Category of Sets, Products, and Coproducts

Proposition 5.3.5.1.1. The category Sets admits a closed symmetric bimonoidal category structure consisting of:

- *The Underlying Category.* The category Sets of pointed sets.

- *The Additive Monoidal Product.* The coproduct functor

$$
\coprod \colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}
$$

   of Constructions With Sets, Item 1 of Definition 4.2.3.1.3.

- *The Multiplicative Monoidal Product.* The product functor

$$
\times \colon \mathsf{Sets} \times \mathsf{Sets} \to \mathsf{Sets}
$$

   of Constructions With Sets, Item 1 of Definition 4.1.3.1.3.

- *The Monoidal Unit.* The functor

$$
\mathbb{1}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}
$$

   of Definition 5.1.3.1.1.

- *The Monoidal Zero.* The functor

$$\mathbb{0}^{\mathsf{Sets}} \colon \mathsf{pt} \to \mathsf{Sets}$$

  of Definition 5.1.3.1.1.

- *The Internal Hom.* The internal Hom functor

$$\mathsf{Sets} \colon \mathsf{Sets}^{\mathsf{op}} \times \mathsf{Sets} \to \mathsf{Sets}$$

  of Constructions With Sets, **??** of **??**.

- *The Additive Associators.* The natural isomorphism

$$\alpha^{\mathsf{Sets},\coprod} \colon \coprod \circ (\coprod \times \mathsf{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \coprod \circ (\mathsf{id}_{\mathsf{Sets}} \times \coprod) \circ \alpha^{\mathsf{Cats}}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}$$

  of Definition 5.2.3.1.1.

- *The Additive Left Unitors.* The natural isomorphism

$$\lambda^{\mathsf{Sets},\coprod} \colon \coprod \circ (\mathbb{0}^{\mathsf{Sets}} \times \mathsf{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \lambda^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

  of Definition 5.2.4.1.1.

- *The Additive Right Unitors.* The natural isomorphism

$$\rho^{\mathsf{Sets},\coprod} \colon \coprod \circ (\mathsf{id} \times \mathbb{0}^{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \rho^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

  of Definition 5.2.5.1.1.

- *The Additive Symmetry.* The natural isomorphism

$$\sigma^{\mathsf{Sets},\coprod} \colon \coprod \overset{\sim}{\Longrightarrow} \coprod \circ \sigma^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}}$$

  of Definition 5.2.6.1.1.

- *The Multiplicative Associators.* The natural isomorphism

$$\alpha^{\mathsf{Sets}} \colon \times \circ (\times \times \mathsf{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \times \circ (\mathsf{id}_{\mathsf{Sets}} \times \times) \circ \alpha^{\mathsf{Cats}}_{\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}}$$

  of Definition 5.1.4.1.1.

· *The Multiplicative Left Unitors.* The natural isomorphism

$$\lambda^{\mathsf{Sets}} \colon \times \circ (\mathbb{1}^{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \boldsymbol{\lambda}^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.1.5.1.1.

· *The Multiplicative Right Unitors.* The natural isomorphism

$$\rho^{\mathsf{Sets}} \colon \times \circ (\mathrm{id} \times \mathbb{1}^{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \boldsymbol{\rho}^{\mathsf{Cats}_2}_{\mathsf{Sets}}$$

of Definition 5.1.6.1.1.

· *The Multiplicative Symmetry.* The natural isomorphism

$$\sigma^{\mathsf{Sets}} \colon \times \overset{\sim}{\Longrightarrow} \times \circ \boldsymbol{\sigma}^{\mathsf{Cats}_2}_{\mathsf{Sets},\mathsf{Sets}}$$

of Definition 5.1.7.1.1.

· *The Left Distributor.* The natural isomorphism

$$\delta^{\mathsf{Sets}}_{\ell} \colon \times \circ (\mathrm{id}_{\mathsf{Sets}} \times \textstyle\coprod) \overset{\sim}{\Longrightarrow} \textstyle\coprod \circ ((\times) \times (\times)) \circ \boldsymbol{\mu}^{\mathsf{Cats}_2}_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}} \circ (\Delta_{\mathsf{Sets}} \times (\mathrm{id}_{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}))$$

of Definition 5.3.1.1.1.

· *The Right Distributor.* The natural isomorphism

$$\delta^{\mathsf{Sets}}_{r} \colon \times \circ (\textstyle\coprod \times \mathrm{id}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \textstyle\coprod \circ ((\times) \times (\times)) \circ \boldsymbol{\mu}^{\mathsf{Cats}_2}_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}} \circ ((\mathrm{id}_{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}}) \times \Delta_{\mathsf{Sets}})$$

of Definition 5.3.2.1.1.

· *The Left Annihilator.* The natural isomorphism

$$\zeta^{\mathsf{Sets}}_{\ell} \colon \mathbb{0}^{\mathsf{Sets}} \circ \boldsymbol{\mu}^{\mathsf{Cats}_2}_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}} \circ (\mathrm{id}_{\mathsf{pt}} \times \boldsymbol{\epsilon}^{\mathsf{Cats}_2}_{\mathsf{Sets}}) \overset{\sim}{\Longrightarrow} \times \circ (\mathbb{0}^{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{Sets}})$$

of Definition 5.3.3.1.1.

· *The Right Annihilator.* The natural isomorphism

$$\zeta^{\mathsf{Sets}}_{r} \colon \mathbb{0}^{\mathsf{Sets}} \circ \boldsymbol{\mu}^{\mathsf{Cats}_2}_{4|\mathsf{Sets},\mathsf{Sets},\mathsf{Sets},\mathsf{Sets}} \circ (\boldsymbol{\epsilon}^{\mathsf{Cats}_2}_{\mathsf{Sets}} \times \mathrm{id}_{\mathsf{pt}}) \overset{\sim}{\dashrightarrow} \times \circ (\mathrm{id}_{\mathsf{Sets}} \times \mathbb{0}^{\mathsf{Sets}})$$

of Definition 5.3.4.1.1.

Proof. Omitted. $\qquad\qquad\square$

Appendices

A   Other Chapters